

菁英班作业第5课

环境

真机：华为nova7 HarmonyOS 3.0 未root

IDA pro

项目目录

Crackme2_unupx: Crackme2脱upx壳

assets: 说明文档图片目录

说明文档.pdf

一、CrackMe1分析

1、实验环境搭建

adb连接手机。

```
(base) PS C:\Users\22057\Tools\adb> .\adb.exe devices -l
List of devices attached
E6E4C20721006850      device product:JEF-AN00 model:JEF_AN00 device:HWJEF transport_id:1
```

将可执行文件push进手机。

```
(base) PS C:\Users\22057\Tools\adb> .\adb.exe push .\CrackMe1 /data/local/tmp/
.\CrackMe1: 1 file pushed, 0 skipped. 1.4 MB/s (34596 bytes in 0.024s)
```

进入adb shell中，给CrackMe1可执行权限。

```
(base) PS C:\Users\22057\Tools\adb> .\adb.exe shell
HWJEF:/ $ cd /data/local/tmp
HWJEF:/data/local/tmp $ chmod +x Cra
CrackMe1 CrackMe2
HWJEF:/data/local/tmp $ chmod +x CrackMe1
```

初步运行，输入1，显示错误答案。

```
HWJEF:/data/local/tmp $ ./CrackMe1
Input Your Answer
1
Wrong Answer
```

分析其可能采用字符串比较的方式进行跳转。

2、使用IDA进行静态分析

```

; int __cdecl main(int argc, const char **argv, const char **envp)
main

var_C= -0xC

; __unwind {
PUSH      {R4,R10,R11,LR}
ADD       R11, SP, #8
SUB       SP, SP, #0x3F0
LDR       R0, (__stack_chk_guard_ptr - 0x105C)
MOV       R4, SP
MOV       R1, #0x3E8
LDR       R0, [PC,R0] ; __stack_chk_guard
LDR       R0, [R0]
STR       R0, [R11,#var_C]
MOV       R0, R4
BL        __aeabi_memclr8
LDR       R0, =(byte_9070 - 0x1074)
ADD       R0, PC, R0 ; byte_9070 ; format
BL        printf
LDR       R0, =(byte_9083 - 0x1084)
MOV       R1, R4
ADD       R0, PC, R0 ; byte_9083 ; format
BL        scanf
MOV       R0, R4
BL        sub_E08
LDR       R0, (__stack_chk_guard_ptr - 0x109C)
LDR       R1, [R11,#var_C]
LDR       R0, [PC,R0] ; __stack_chk_guard
LDR       R0, [R0]
SUBS      R0, R0, R1
MOVEQ     R0, #0
SUBEQ     SP, R11, #8
POPEQ     {R4,R10,R11,PC}

```

BL __stack_chk_fail
; End of function main

对main函数进行反编译

```

int __cdecl main(int argc, const char **argv, const char **envp)
{
    _BYTE v4[1004]; // [sp+0h] [bp-3F8h] BYREF

    memset(v4, 0, 0x3E8u);
    printf(&byte_9070);
    scanf(&byte_9083, v4);
    sub_E08(v4);
    return 0;
}

```

发现其先输出提示字符串"Input Your Answer".

再读入字符串，保存在v4中。

将v4传入函数sub_E08。

进入函数sub_E08

```

void *v5; // r5
void *v6; // r7
size_t v7; // r0
size_t v8; // r0
int v9; // r0
int v11; // [sp+10h] [bp-1B8h]
__int16 v12[50]; // [sp+14h] [bp-1B4h] BYREF
char dest[100]; // [sp+78h] [bp-150h] BYREF
char v14[100]; // [sp+DCh] [bp-ECh] BYREF
char v15[100]; // [sp+140h] [bp-88h] BYREF
int v16; // [sp+1A4h] [bp-24h]

memset(v15, 0, sizeof(v15));
memset(v14, 0, sizeof(v14));
memset(dest, 0, sizeof(dest));
memset(v12, 0, sizeof(v12));
qmemcpy(v12, "123", 3);
if ( sub_B04(51, v2, v3, v4) == 1 )
{
    printf("TOEDCTF+!");
    sleep(1u);
    abort();
}
strcat(dest, &byte_9030);
v5 = malloc(0x64u);
memset(v5, 0, 0x64u);
v6 = malloc(0x64u);
memset(v6, 0, 0x64u);
v7 = strlen(dest);
sub_1664(dest, v7, v6);
sub_2550(v5, v12, v6, 16, v15, v14, dest, v12);
v8 = strlen((const char *)v5);
v11 = strncmp((const char *)v5, a1, v8);
v9 = 334548934;
do
{
    if ( v9 == -1636053901 )
    {
        printf(&byte_9056);
        goto LABEL_3;
    }
    v9 = -1636053901;
    if ( !v11 )

```

前面一部分内容暂且不管，发现函数调用strncmp

`v11 = strncmp((const char *)v5, a1, v8);`

其中的a1参数为传入的输入字符串地址，v8是v5的字符串长度。判断正确字符串保存在v5中。

而v5通过函数sub_2550得出。

进入函数sub_2550

```

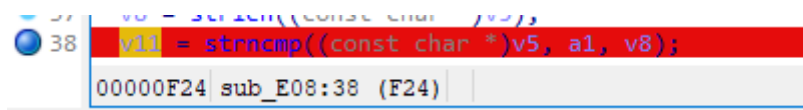
unsigned int v8; // r4
int v9; // r1
int v10; // r6
unsigned int v11; // r8
int v12; // r5
int v15; // [sp+Ch] [bp-49Ch] BYREF
int v16; // [sp+10h] [bp-498h] BYREF
int v17; // [sp+14h] [bp-494h]
int v18; // [sp+18h] [bp-490h]
unsigned int v19; // [sp+3DCh] [bp-CCh]

v18 = 0;
v17 = 0;
v16 = 0;
v15 = 0;
sub_1888(&v16, a2, &v15, 16, 16);
v6 = (a4 + ((unsigned int)(a4 >> 31) >> 28)) & 0xFFFFFFFF0;
v7 = a4 % 16;
if ( (unsigned int)(a4 + 15) >= 0x1F )
{
    if ( (_BYTE)v17 )
    {
        v8 = v19;
        sub_2884(v6, v19);
        if ( v8 <= v6 && !v9 )
        {
            v10 = a3;
            v11 = 0;
            v12 = a1;
            do
            {
                sub_21DC(&v16, v10, v12);
                v12 += v19;
                v10 += v19;
                ++v11;
            }
            while ( v11 < sub_27DC(v6) );
        }
    }
}
if ( v7 >= 1 )
    qmemcpy((void *)(a1 + v6), (const void *)(a3 + v6), v7);
return 1;
}

```

发现其函数实现较为复杂，故放弃静态分析正确字符串，转向动态分析。

记录其关键比较函数strncmp地址及其参数，在此位置设置断点



其参数位于R0,R1寄存器内

text:00000F14	00 20 A0 E1	MOV	R2, R0	; n
text:00000F18	05 00 A0 E1	MOV	R0, R5	; s1
text:00000F1C	04 10 A0 E1	MOV	R1, R4	; s2
text:00000F20	91 FE FF EB	BL	strncmp	
text:00000F20				
text:00000F24	10 00 8D E5	STR	R0, [SP,#0x1C8+var_1B8]	

3、IDA动态调试环境搭建

将IDA自带的android_server传入手机中

```
(base) PS C:\Users\22057\Tools\adb> .\adb.exe push .\android_server /data/local/tmp/
.\android_server: 1 file pushed, 0 skipped. 84.5 MB/s (803256 bytes in 0.009s)
(base) PS C:\Users\22057\Tools\adb>
```

启动android_server

```
HWJEF:/data/local/tmp $ ./android_server
IDA Android 32-bit remote debug server(ST) v7.7.27. Hex-Rays (c) 2004-2022
Listening on 0.0.0.0:23946...
```

进行端口映射

```
.\adb.exe forward tcp:23946 tcp:23946
```

4、IDA动态调试

运行CrackMe1

```
HWJEF:/data/local/tmp $ ./CrackMe1
Input Your Answer
```

在strcmp处打下断点

```
32  v19 = ((int (fastcall *) (int, int, int))unk_6D8196C)(v14, v13, v19);
33  v22 = ((int (fastcall *) (int, int, int))unk_6D8196C)(v14, v13, v19);
34  v20 = 334548934;
```

输入测试数据1，继续运行

```
HWJEF:/data/local/tmp $ ./CrackMe1
Input Your Answer
1|
```

产生中断，错误，判断其可能存在反调试策略。

```
HWJEF:/data/local/tmp $ ./CrackMe1
Input Your Answer
1
undebug
Aborted
```

5、绕过反调试策略

静态分析发现此处存在printf与abort函数，在此处打下断点尝试

```
23  if ( sub_B04(51, v2, v3, v4) == 1 )
24  {
25      printf("TOEDCTF+!");
26      sleep(1u);
27      abort();
28  }
```

CrackMe1:0421FE98		
CrackMe1:0421FE9C	CMP	R0, #1
CrackMe1:0421FEA0	BEQ	loc_4220008
CrackMe1:0421FEA0		
CrackMe1:0421FEA4	LDR	R0, =(a0wpbyeqgdjfope - 0x421FEB4) ; "0wPBYEQGDjFOpeiKFxHqEQ=="
CrackMe1:0421FEA8	ADD	R6, SP, #0x78 ; 'x'
CrackMe1:0421FEAC	ADD	R1, PC, R0 ; "0wPBYEQGDjFOpeiKFxHqEQ=="
CrackMe1:0421FEB0	MOV	R0, R6
CrackMe1:0421FEB4	BL	unk_421F9B4
CrackMe1:0421FEB4		
CrackMe1:0421FEB8	MOV	R0, #0x64 ; 'd'
CrackMe1:0421FEBc	BL	unk_421F9C0
CrackMe1:0421FEBc		

运行到此处，发现R0值为1

FE9C	CMP	R0, #1
FEA0	BEQ	loc_4220008
FEA0		
FEA4	LDR	R0, =(a0wpbyeqgdjfope - 0x421FEB4) ; "0wPBYEQGDjFOpeiKFxHqEQ=="

二者相等，则跳转至目的地址，输出undebug。

8	loc_4220008		; CODE XREF: CrackMe1:0421FEA0↑j
8	LDR	R0, =(aUndebug - 0x4220014) ; "undebug\n"	
C	ADD	R0, PC, R0 ; "undebug\n"	
0	BL	unk_421F984	
0			
4	MOV	R0, #1	
8	BL	unk_421F990	
8			
C	MOV	LR, PC	
0	B	sub_421F99C	
0			
4	ANDEQ	R8, R0, R12, LSR R1	
4			

此处修改R0的值为0，让其不进行跳转。

CrackMe1:0421FE9C	CMP	R0, #1
CrackMe1:0421FEA0	BEQ	loc_4220008
CrackMe1:0421FEA0		
CrackMe1:0421FEA4	LDR	R0, =(a0wpbyeqgdjfope - 0x421FEB4) ; "0wPBYEQGDjFOpeiKFxHqEQ=="
CrackMe1:0421FEA8	ADD	R6, SP, #0x78 ; 'x'
CrackMe1:0421FEAC	ADD	R1, PC, R0 ; "0wPBYEQGDjFOpeiKFxHqEQ=="
CrackMe1:0421FEB0	MOV	R0, R6
CrackMe1:0421FEB4	BL	unk_421F9B4
CrackMe1:0421FEB4		
CrackMe1:0421FEB8	MOV	R0, #0x64 ; 'd'
CrackMe1:0421FEBc	BL	unk_421F9C0
CrackMe1:0421FEBc		
CrackMe1:0421FEC0	MOV	R1, #0x64 ; 'd'

6、获取最终结果

运行至strncmp函数处。

CrackMe1:0421FF10		
CrackMe1:0421FF14	MOV	R2, R0
CrackMe1:0421FF18	MOV	R0, R5
CrackMe1:0421FF1C	MOV	R1, R4
CrackMe1:0421FF20	BL	unk_421F96C
CrackMe1:0421FF20		
CrackMe1:0421FF24	STR	R0, [SP, #0x10]

查看R0, R1寄存器内容。

R0: 为正确字符串GameSecurity

```

[anon:libc_malloc]:F3B3D000 DCB 0x47 ; G
[anon:libc_malloc]:F3B3D001 DCB 0x61 ; a
[anon:libc_malloc]:F3B3D002 DCB 0x6D ; m
[anon:libc_malloc]:F3B3D003 DCB 0x65 ; e
[anon:libc_malloc]:F3B3D004 DCB 0x53 ; S
[anon:libc_malloc]:F3B3D005 DCB 0x65 ; e
[anon:libc_malloc]:F3B3D006 DCB 0x63 ; c
[anon:libc_malloc]:F3B3D007 DCB 0x75 ; u
[anon:libc_malloc]:F3B3D008 DCB 0x72 ; r
[anon:libc_malloc]:F3B3D009 DCB 0x69 ; i
[anon:libc_malloc]:F3B3D00A DCB 0x74 ; t
[anon:libc_malloc]:F3B3D00B DCB 0x79 ; y
[anon:libc_malloc]:F3B3D00C DCB 0

```

R1: 为测试内容1

```

[stack]:FFE8B8EF DCB 4
[stack]:FFE8B8F0 DCB 0x31 ; 1
[stack]:FFE8B8F1 DCB 0

```

最终结果应为GameSecurity

7、测试获取的结果

```

HWJEF:/data/local/tmp $ ./CrackMe1
Input Your Answer
GameSecurity
True Answer
HWJEF:/data/local/tmp $

```

结果正确

二、CrackMe2分析

1、实验环境搭建

略

2、使用IDA进行静态分析

对start函数进行反编译

```

IDA View-A Pseudocode-A
1 void __noreturn start()
2 {
3     _BYTE *v0; // r0
4     char v1; // t1
5     char *v2; // r0
6     _BYTE *v3; // r2
7     char v4; // cf
8
9     v0 = (_BYTE *)((__int64 (*)(void))loc_6008)();
0     while ( 1 )
1     {
2         v2 = (char *)sub_5DDC(v0);
3         if ( !v4 )
4             break;
5         v1 = *v2;
6         v0 = v2 + 1;
7         *v3 = v1;
8     }
9     sub_5DF4(v2);

```

均为数据加载代码，判断其为加壳应用。

通过查看信息，判断其为upx加壳

```
... ..#INFO..
·This·file·is·pa
cked·with·the·UP
X·executable·pac
ker·http://upx.s
f.net/$..$Id:·UP
X·3.95·Copyright
·(C)·1996-2018·t
he·UPX·Team·All
·Rights·Reserved
·$.....0.....
```

3、upx脱壳

使用upx工具进行脱壳

```
upx -d CrackMe2
```

脱壳成功

```
CrackMe1 CrackMe2
HWJEF:/data/local/tmp $ ./upx -d CrackMe2
                          Ultimate Packer for eXecutables
                          Copyright (C) 1996 - 2022
UPX 4.0.1      Markus Oberhumer, Laszlo Molnar & John Reiser   Nov 16th 2022

   File size      Ratio      Format      Name
   -----
   40536 <-      27308      67.37%     linux/arm      CrackMe2

Unpacked 1 file.
```

4、使用IDA进行静态分析脱壳后可执行文件

main函数与CrackMe1差不多

```
int __cdecl main(int argc, const char **argv, const char **envp)
{
    _BYTE v4[1004]; // [sp+0h] [bp-3F8h] BYREF

    memset(v4, 0, 0x3E8u);
    printf(&byte_A0A0);
    scanf(&byte_A0B3, v4);
    sub_1470(v4);
    return 0;
}
```

sub_1420函数

```
sub_1420(54, 49);
v2 = -889797365;
```

进入后发现可能与反调试有关


```

int sub_1420()
{
    int result; // r0

    if ( sub_B18() == 1 || (result = sub_1324(), result == 1) )
    {
        printf(&byte_A03A);
        sleep(1u);
        abort();
    }
    return result;
}

```

以下三个函数可能与字符串比较有关，动态调试时在此打下断点

```

132 ;
133 sub_20EC(v37, v15 - v37, v14);
134 sub_2FD8(v13, v36, v14, 16);
135 v17 = v30;
136 while ( *(unsigned __int8 *)++v17 )
137 ;
138 sub_1F10(v31, v17 - v31, v38);
139 v19 = v13 - 1;

```

5、IDA动态调试

IDA尚未附加便检测到调试退出

```

HWJEF:/data/local/tmp $ ./CrackMe2
undebug
Aborted

```

判断其存在检测android_server端口代码

修改android_server端口为22222

```

1|HWJEF:/data/local/tmp $ ./android_server -p22222
IDA Android 32-bit remote debug server(ST) v7.7.27. Hex-Rays (c) 2004-2022
Listening on 0.0.0.0:22222...

```

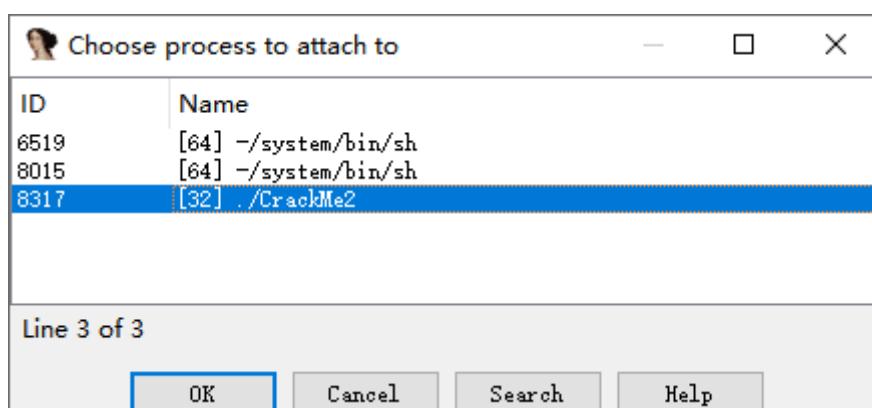
正常启动说明其绕过了检测

```

134|HWJEF:/data/local/tmp $
134|HWJEF:/data/local/tmp $ ./CrackMe2
Input Your Answer

```

附加到CrackMe2上



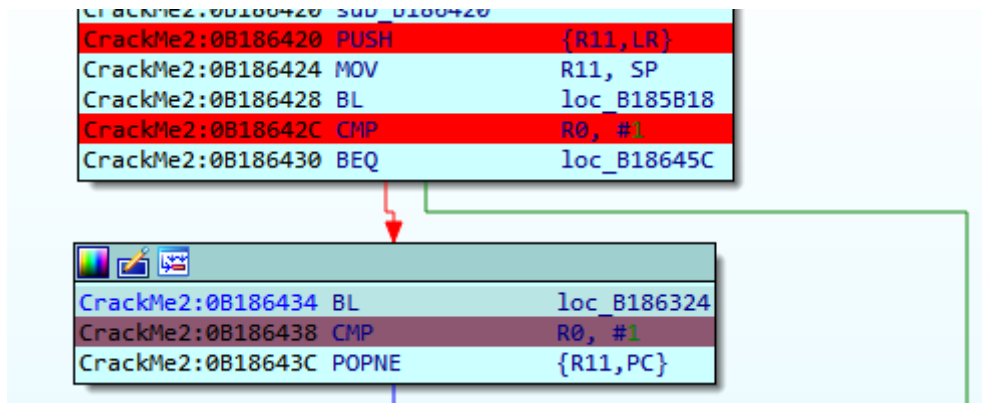
sub1420函数中存在输出undebug信息

```

1 int sub_4DE7420()
2 {
3     int result; // r0
4     int v1; // r0
5
6     if ( ((int (*)(void))unk_4DE6B18)() == 1 || (result = ((int (*)(void))unk_4DE7324)(), result == 1) )
7     {
8         ((void (__fastcall *)(char *))unk_4DE69BC)(aUndebug);
9         v1 = ((int (__fastcall *)(int))unk_4DE69C8)(1);
10        return sub_4DE69D4(v1);
11    }
12    return result;
13}

```

判断其为反调试



其中loc_B185B18函数为检测IDA android_server的tcp端口，loc_B186324为检测是否调试状态。

由于此前已经修改了端口，只需要修改第二次CMP R0, #1的指令即可。

将R0的值赋值为0即可绕过调试

接下来对于三个可疑函数调用进行分析

函数sub20EC:

```

CrackMe2:0B1868B0 MOVW    R3, #0xB4BF
CrackMe2:0B1868B4 LDR     R0, [R11, #-0xF0]
CrackMe2:0B1868B8 MOVT   R3, #0xA563
CrackMe2:0B1868BC ADD     R2, R2, R3
CrackMe2:0B1868C0 SUB     R1, R2, R1
CrackMe2:0B1868C4 MOV     R2, R7
CrackMe2:0B1868C8 SUB     R1, R1, R3
CrackMe2:0B1868C8 ; END OF FUNCTION CHUNK FOR sub_B1868CC
CrackMe2:0B1868CC ; ===== S U B R O U T I N E =====
CrackMe2:0B1868CC ; void __fastcall sub_B1868CC(int, int, int, int)
CrackMe2:0B1868CC sub_B1868CC
CrackMe2:0B1868CC ; FUNCTION CHUNK AT CrackMe2:0B1864FC SIZE 0006
CrackMe2:0B1868CC BL      unk_B1870EC

```

R0值为

```

[stack]:FFB2D917 DCB    0
[stack]:FFB2D918 aWqvZkhagjpu3da DCB "wQV/zkhagjpu3dAo4YVPFoOXM5l0lDN+99FNfjx8Cs=",0
[stack]:FFB2D945 DCB    0

```

判断其应为base64加密后内容，但解码为乱码

R1为2c，R3地址处为空值，判断其应当不是字符串比较函数

函数sub_1FD8

```

CrackMe2:0B1868C0
CrackMe2:0B1868D0 LDR      R1, [R11,#-0xF4]
CrackMe2:0B1868D4 MOV      R0, R5
CrackMe2:0B1868D8 MOV      R2, R7
CrackMe2:0B1868DC MOV      R3, #0x10
CrackMe2:0B1868E0 BL       loc_B187FD8
CrackMe2:0B1868E0
CrackMe2:0B1868E4 LDR      R2, [R11,#-0x10C]
CrackMe2:0B1868F4

```

R0为R5的值

在函数执行完成后R5的位置发现一字符串

```

[anon:libc_malloc]:EA0EB000 aVgvuy2vudedhbw DCB "VGVuY2VudEdhbwU=",0
[anon:libc_malloc]:EA0EB011 DCB 0
[anon:libc_malloc]:EA0EB012 DCB 0

```

判断其为base64加密后内容

请输入要进行 Base64 编码或解码的字符

VGVuY2VudEdhbwU=

编码 (Encode)

解码 (Decode)

↑ 交换

(编码快捷键: **Ct**

Enter)

Base64 编码或解码的结果:

☐ 编/解码后自动

TencentGame

解码结果为TencentGame，初步猜测为结果

对于函数1F10:

其参数R0为输入字符串1，R1位字符串长度1。

在其函数内部

```

dMe2:0B186F14 ADD      R11, SP, #0x1C
dMe2:0B186F18 SUB      SP, SP, #0x14
dMe2:0B186F1C STR      R2, [SP,#4]
dMe2:0B186F20 CMP      R1, #0
dMe2:0B186F24 LDR      R2, =(off_B18EF60 - 0xB186F30)
dMe2:0B186F28 LDR      R2, [PC,R2] ; dword_E9F2C274
dMe2:0B186F2C LDR      R2, [R2]
dMe2:0B186F30 STR      R2, [SP,#0x10]
dMe2:0B186F34 LDRB     LR, [SP,#0xE]
dMe2:0B186F38 LDRB     R12, [SP,#0xF]
dMe2:0B186F3C BEQ      loc_B186FC8
dMe2:0B186F3C
dMe2:0B186F40 LDR      R2, =(aBcDeFgHiJkLmN - 0xB186F50) ; "ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/",0
dMe2:0B186F44 MOV      R9, #0
dMe2:0B186F48 ADD      R10, PC, R2 ; "ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/",0
dMe2:0B186F4C MOV      R2, #0 ; DATA XREF: CrackMe2:0B186F40to
dMe2:0B186F4C ; CrackMe2:0B186F48to
dMe2:0B186F50 ; CrackMe2:0B187020to
dMe2:0B186F50 loc_B186F50 ; CrackMe2:0B187044to
dMe2:0B186F50 LDRB     R4, [R0,R2] ; CrackMe2:0B187064to
dMe2:0B186F54 ADD      R3, R2, #1 ; CrackMe2:0B18706Cto
dMe2:0B186F58 MOV      R6, #1 ; CrackMe2:off_B1870DCto
dMe2:0B186F5C CMP      R3, R1 ; CrackMe2:off_B1870E0to
dMe2:0B186F60 BCS     loc_B186FE0 ; CrackMe2:off_B1870E4to
loc_B186FE0
; CODE XREF: CrackMe2:0B186F50
; 0B186F10: CrackMe2:loc_B186F10 (Synchronized with PC. Pseudocode-A)

```

存在ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/

其应为base64加密函数

6、测试获取的结果

使用TencentGame进行测试，结果正确。

```
139|HWJEF:/data/local/tmp $ ./CrackMe2
Input Your Answer
TencentGame
True Answer
HWJEF:/data/local/tmp $ |
```