## - MODULE P2PBroadcastSpec

Spec for a reliable p2p broadcast

EXTENDS Naturals, Sequences, FiniteSets

CONSTANT

Proc, Set of processes

Data

VARIABLES

sent, All messages sent by all processors

 $received\_by$  All messges received. Function from message to receiving processors

 $vars \triangleq \langle sent, received\_by \rangle$ 

Message is a record including the sending proc and a data

 $Message \triangleq [from : Proc, data : Data]$ 

 $Init \stackrel{\triangle}{=} \land sent = \{\} \\ \land received\_by = [m \in Message \mapsto \{\}]$ 

 $\begin{array}{ccc} \textit{TypeOK} & \triangleq & \land \textit{sent} \in \textit{SUBSET Message} \\ & \land \textit{received\_by} \in [\textit{Message} \rightarrow \textit{SUBSET Proc}] \end{array}$ 

Send message m

 $Send(m) \stackrel{\triangle}{=} \land m \notin sent$  Message is sent only once by the original sender  $\land sent' = sent \cup \{m\}$   $\land$  UNCHANGED  $\land received\_by \land$ 

Receive a message m at proc p

 $\begin{array}{lll} Recv(m,\,p) & \triangleq & \wedge \, m \in sent & \text{receive only if } \, m \text{ was sent first} \\ & \wedge \, p \notin received\_by[m] & \text{receive only once} \\ & \wedge \, received\_by' = [received\_by \text{ EXCEPT } ![m] = @ \cup \{p\}] \\ & \wedge \, \text{UNCHANGED } \langle sent \rangle \\ \end{array}$ 

 $Next \triangleq \exists m \in Message, p \in Proc : Send(m) \lor Recv(m, p)$ 

 $Spec \stackrel{\Delta}{=} Init \wedge \Box [Next]_{vars}$ 

FairSpec is Spec with the addition requirement that it keeps taking steps.

 $FairSpec \triangleq Spec \wedge WF_{vars}(Next)$ 

Liveness is a temporal property that captures the property that if a message is sent, it is eventually received.

The WF Conjuction rule is used here ref: Specifying Systems p 105

 $Liveness \triangleq \exists m \in Message, p \in Proc : \mathrm{WF}_{vars}(Send(m) \vee Recv(m, p))$