## MODULE ShamirSecretSharing

Sepcification for simple Shamir Secret Sharing. This is not a veriable secret sharing scheme.

We specify that dealer first sends shares to all players, and once all players have received their shares the can eventually reconstruct the secret.

We do not deal with the communication protocol between players to send their shares to each other before reconstructing the secret.

## EXTENDS Integers, Sequences

#### CONSTANT

Dealer, The dealer sharing the secret with the players

Players, Set of all players

Coefficients The coefficient of the polynomial. These are provided by the model

#### VARIABLES

shares,
shares\_sent,
shares\_received,
reconstructed
Function mapping Player to shares received shares
Function mapping Player to received shares
Function mapping Player to flag if secret
has been successfully constructed

 $vars \triangleq \langle shares, shares\_sent, shares\_received, reconstructed \rangle$ 

```
No Value \triangleq -1
```

# $Init \triangleq$

```
Compute shares as a + bx + cx^2
```

 $\land shares = [p \in Players \mapsto Coefficients[1] + Coefficients[2] * p + Coefficients[3] * p^2]$ 

 $\land shares\_sent = [p \in Players \mapsto NoValue]$ 

 $\land shares\_received = [p \in Players \mapsto NoValue]$ 

 $\land reconstructed = [p \in Players \mapsto FALSE]$ 

## The type invariant for all variables.

# $TypeOK \triangleq$

 $\land shares \in [Players \rightarrow Int]$ 

 $\land shares\_sent \in [Players \rightarrow Int]$ 

 $\land \mathit{shares\_received} \in [\mathit{Players} \rightarrow \mathit{Int}]$ 

 $\land reconstructed \in [Players \rightarrow BOOLEAN]$ 

## Send the share to Player p.

$$SendShare(p) \triangleq$$

 $\land shares\_sent[p] = NoValue$ 

Send a share that has not been sent to anyone

 $\land shares\_sent' = [shares\_sent \ EXCEPT \ ![p] = shares[p]]$ 

∧ UNCHANGED ⟨shares, shares\_received, reconstructed⟩

Receive the share at Player p. It should have been sent before.

Reconstruct secret with Players p and q. The payers should have received share.

The next step either sends shares, receieves them or reconstructs the secret.

```
 \begin{array}{ll} Next & \triangleq \\ & \lor \exists \: p \in Players : \\ & SendShare(p) \lor ReceiveShare(p) \\ & \lor \exists \: p \in Players, \: q \in Players : Reconstruct(p, \: q) \\ Spec & \triangleq \\ & \land \: Init \\ & \land \: \Box [Next]_{vars} \\ \end{array}
```

Liveness states that eventually all players reconstruct the secret.

```
Liveness \stackrel{\triangle}{=} \forall p \in Players, q \in Players : WF_{vars}(Reconstruct(p, q))
```

For a fair specification, we assure the spec takes next steps and liveness is guaranteed.

 $FairSpec \stackrel{\Delta}{=} Spec \wedge Liveness$ 

```
\ * Modification History
```

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