- Module BlockGeneration -

Block generation specifies when and how braidpool miners generate blocks. The protocol to build current pool key and threshold signatures is assumed

EXTENDS

Sequences, Integers, DAG, FiniteSets

CONSTANT

Miner, Set of miners

Share Seq No, Share seq numbers each miner generates BlockReward, Block reward in a difficulty period

MaxPathLen Max path length to check for message stability.

This helps constrain the search space in the DAG.

VARIABLES

TODO: Replace these last_. * variables with operators on DAG

last_sent, Function mapping miner to last sent share seq_no share_dag, A DAG with valid shares for now implemented as a set stable, Set of shares that are stable in the DAG, i.e. received

by all other miners

unpaid_coinbases, coinbases for braidpool blocks that

haven t been paid yet

uhpo, Function mapping miner to unpaid balance

pool_key Current public key for TS

Share is a record of miner and sequence number. All shares are assumed to be mined at same difficulty

 $Share \stackrel{\triangle}{=} [miner : Miner, seq_no : ShareSeqNo]$

Acks are the implicit acknowledgements sent with each share. These are the sequence number of shares received from each miner.

 $Acks \triangleq \langle Share \rangle$

PublicKey is defined as sequence of miner identifier for now. The miners in a key sequence are the miners contributing to the key generated using DKG.

 $PublicKey \triangleq Seq(Miner)$

Coinbase is a payment to a DKG public key with an value.

 $Coinbase \stackrel{\Delta}{=} [key : PublicKey, value : BlockReward]$

 $No Val \triangleq 0$

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Init \triangleq
         \land last\_sent = [m \in Miner \mapsto NoVal]
         \land \mathit{share\_dag} = [\mathit{node} \mapsto \{\}, \mathit{edge} \mapsto \{\}]
         \land stable = \{\}
         \land \ unpaid\_coinbases = \{\}
         \land uhpo = [m \in Miner \mapsto NoVal]
         \land pool\_key = \langle \rangle
TypeInvariant \triangleq
          \land last\_sent \in [Miner \rightarrow Int \cup \{NoVal\}]
          \land share\_dag.node \in subset Share
          \land share_dag.edge \in SUBSET (Share \times Share)
          \land stable \in SUBSET Share
          \land unpaid\_coinbases \in subset Coinbase
          \land uhpo \in [Miner \rightarrow Int \cup \{NoVal\}]
          \land pool\_key \in Seq(Miner)
vars \triangleq \langle last\_sent, share\_daq, stable, unpaid\_coinbases, uhpo, pool\_key \rangle
Send a share from a miner with a seqno = last share sent +1 and in ShareSeqNo. The share is
assumed to be successfully broadcast to all miners.
SendShare \triangleq \exists m \in Miner, sno \in ShareSeqNo:
                \land sno = last\_sent[m] + 1
                \land last\_sent' = [last\_sent \ EXCEPT \ ![m] = @ + 1]
                \land share\_dag' = [share\_dag \ EXCEPT]
                                          Add share to node list of graph
                                         !.node = @ \cup \{[miner \mapsto m, seq\_no \mapsto sno]\},
                                          Add edge from share to all non NoVal last_sent
                                          This can be replaced by last share in DAG from others
                                         !.edge = @ \cup
                                             \{[miner \mapsto m, seq\_no \mapsto sno]\}
                                             \{[miner \mapsto mo, seq\_no \mapsto last\_sent[mo]]:
                                                    mo \in \{mm \in Miner : last\_sent[mm] \neq NoVal\}\}
                \land UNCHANGED \langle stable, unpaid\_coinbases, uhpo, pool\_key <math>\rangle
Stabilise a share if there is a path from the share to any share from all other miners. How do we
know all other miners? This comes from a separate protocol where a miner is dropped from the
set of all other miners. Miners are dropped from the list if they have not sent shares since the last
bitcoin block was found. For now, we assume the list of to the group of miners is known.
StabiliseShare \triangleq \exists s \in share\_dag.node :
                          \land s \not\in stable
                          \land \forall m \in Miner \setminus \{s.miner\}:
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 $\land stable' = stable \cup \{s\}$

 $\exists p \in Paths(MaxPathLen, share_dag) : \land p[1].miner = m$

 $\land p[2].miner = s.miner$