

# In-EVM Solana State Verification

## Technical Reference

Alisa Cherniaeva

[a.cherniaeva@nil.foundation](mailto:a.cherniaeva@nil.foundation)

=nil; Crypto3 (<https://crypto3.nil.foundation>)

Ilia Shirobokov

[i.shirobokov@nil.foundation](mailto:i.shirobokov@nil.foundation)

=nil; Crypto3 (<https://crypto3.nil.foundation>)

Mikhail Komarov

[nemo@nil.foundation](mailto:nemo@nil.foundation)

=nil; Foundation (<https://nil.foundation>)

November 25, 2021

# Contents

<b>1</b>	<b>Introduction</b>	<b>2</b>
1.1	Overview . . . . .	2
<b>2</b>	<b>State Proof Generator</b>	<b>3</b>
2.1	'Light-Client' State . . . . .	3
2.2	Proof System . . . . .	4
2.3	Optimizations . . . . .	4
2.3.1	Batched FRI . . . . .	4
2.3.2	Hash By Column . . . . .	4
2.3.3	Hash By Subset . . . . .	4
2.4	RedShift Protocol . . . . .	5
2.4.1	Prover View . . . . .	5
2.4.2	Verifier View . . . . .	7
2.5	Circuit Definition . . . . .	8
2.5.1	Verification Circuit Overview . . . . .	8
2.5.2	SHA-256 Circuit . . . . .	8
2.5.3	SHA2-512 Circuit . . . . .	11
2.5.4	Poseidon Circuit . . . . .	15
2.5.5	Merkle Tree Circuit . . . . .	15
2.5.6	Ed25519 Circuit . . . . .	16
2.5.7	Elliptic Curves Arithmetics . . . . .	16
2.5.8	Validator Set Proof Circuit . . . . .	18
<b>3</b>	<b>In-EVM State Proof Verifier</b>	<b>19</b>
3.1	Verification Logic Architecture . . . . .	19
3.2	Verification Logic API Reference . . . . .	19
3.3	Input Data Structures . . . . .	19
	<b>Bibliography</b>	<b>19</b>

# Chapter 1

## Introduction

This document is a technical reference to the in-EVM Solana's 'Light-Client' state verification project.

### 1.1 Overview

The project's purpose is to provide Ethereum users with reliable Solana's cluster state and necessary transactions proof.

The project UX consists of several steps:

1. Retrieve Solana's 'Light-Client' state.
2. Generate a proof for it.
3. Submit the proof to EVM-enabled cluster.
4. Verify the proof with EVM.

Such a UX defines projects parts:

1. Solana's 'Light-Client' state retriever.
2. State proof generator.
3. Ethereum RPC proof submitter.
4. EVM-based proof verifier.

Each of these parts will be considered independently.

## Chapter 2

# State Proof Generator

This introduces a description for Solana's 'Light-Client' state proof generator. Crucial components which define this part design and performance are:

1. Input data format ('Light-Client' state data structure).
2. Proof system used for the proof generation.
3. Circuit definition used for the proof system.

### 2.1 'Light-Client' State

Block Information  $\bar{B}_k$  is defined as follows:

- $k$  - the number of the block
- $B_k = H(B_{k-1}||\text{account\_hash}||\text{signature\_count\_buf}||b_k||\text{validators\_state})$  - bank hash of the block<sup>1</sup>
- $b_k$  Merkle Block
- $B_{k-1}$  - the previous block's bank hash
- $\text{validators\_state}$  is not implemented for now.

Proof algorithm input is defined as follows:

- $n_1$  - current confirmed block number
- $n_2$  - new confirmed block number
- $\{\bar{B}_{n_1}, \dots, \bar{B}_{n_2}, \dots, \bar{B}_{n_2+32}\}$  - block information for blocks from  $n_1$  to  $n_2 + 32$ .
- $\sigma_0, \dots, \sigma_N$  - signatures for  $B_{n_2+32}$

Approximate code representation of such a state data structure is as follows:

```
template<typename Hash>
struct block_data {
    typedef typename Hash::digest_type digest_type;

    std::size_t block_number;
    digest_type bank_hash;
    digest_type merkle_hash;
    digest_type previous_bank_hash;
    // std::vector<vote_state> votes;
};

template<typename Hash, typename SignatureSchemeType>
struct state_type {
    typedef Hash hash_type;
    typedef SignatureSchemeType signature_scheme_type;
    typedef typename signature_scheme_type::signature_type signature_type;
```

<sup>1</sup>See <https://docs.solana.com/proposals/simple-payment-and-state-verification#block-headers>

```

std::size_t n_1 confirmed;
std::size_t n_2 new_confirmed;
std::vector<block_data<hash_type>> repl_data;
std::vector<signature_type> signatures;
};

```

Validator state-representing data structure (`vote_state`) supposes such a state to begin being handled by Solana replication protocol (or its implementation) for handling the tracking of votes state being unchanged 'till the end of epoch.

## 2.2 Proof System

### WIP

The proof system used for proving Solana's 'Light-Client' state on EVM is Redshift SNARK[1]. RedShift is a transparent SNARK that uses PLONK[2] proof system but replaces the commitment scheme. Initial paper proposal is to employ FRI[3] protocol to obtain transparency for the PLONK system.

However, FRI cannot be straightforwardly used with the PLONK system. To achieve the required security level without huge overheads, the authors introduce *list polynomial commitment* scheme as a part of the protocol. For more details, the reader gets referred to [1].

The original RedShift protocol utilizes the classic PLONK[2] system. To provide better performance, the original protocol is generalized to be used with PLONK with custom gates [4], [5] and lookup arguments [6], [7].

## 2.3 Optimizations

### WIP

### 2.3.1 Batched FRI

Instead of check each commitment individually, we can aggregate them for FRI. For polynomials  $f_0, \dots, f_k$ :

1. Get  $\theta$  from transcript
2.  $f = f_0 \cdot \theta^{k-1} + \dots + f_k$
3. Run FRI over  $f$ , using oracles to  $f_0, \dots, f_k$

Thus, we can run only one FRI instance for all committed polynomials.  
See [1] for details.

### 2.3.2 Hash By Column

Instead of committing each of the polynomials, we can use the same Merkle tree for several polynomials. It decreases the number of Merkle tree paths that need to be provided by the prover.

See [8], [1] for details.

### 2.3.3 Hash By Subset

On the each  $i + 1$  FRI round, the prover should send all elements from a coset  $H \in D^{(i)}$ . Each Merkle leaf is able to contain the whole coset instead of separate values.

See [8] for details. Similar approach is described in [1]. However, the authors of [1] use more values per leaf, that leads to better performance.

## 2.4 RedShift Protocol

### WIP

Notations:

$N_{\text{wires}}$	Number of wires ('advice columns')
$N_{\text{perm}}$	Number of wires that are included in the permutation argument
$N_{\text{sel}}$	Number of selectors used in the circuit
$N_{\text{const}}$	Number of constant columns
$N_{\text{lookups}}$	Number of lookups
$\mathbf{f}_i$	Witness polynomials, $0 \leq i < N_{\text{wires}}$
$\mathbf{f}_{c_i}$	Constant-related polynomials, $0 \leq i < N_{\text{const}}$
$\mathbf{gate}_i$	Gate polynomials, $0 \leq i < N_{\text{sel}}$
$\sigma(\text{col} : i, \text{row} : j) = (\text{col} : i', \text{row} : j')$	Permutation over the table

For details on polynomial commitment scheme and polynomial evaluation scheme, we refer the reader to [1].

- 
1.  $\mathcal{L}' = (\mathbf{q}_0, \dots, \mathbf{q}_{N_{\text{sel}}})$
  2. Let  $\omega$  be a  $2^k$  root of unity
  3. Let  $\delta$  be a  $T$  root of unity, where  $T \cdot 2^S + 1 = p$  with  $T$  odd and  $k \leq S$
  4. Compute  $N_{\text{perm}}$  permutation polynomials  $S_{\sigma_i}(X)$  such that  $S_{\sigma_i}(\omega^j) = \delta^{i'} \cdot \omega^{j'}$
  5. Compute  $N_{\text{perm}}$  identity permutation polynomials:  $S_{id_i}(X)$  such that  $S_{id_i}(\omega^j) = \delta^i \cdot \omega^j$
  6. Let  $H = \{\omega^0, \dots, \omega^n\}$  be a cyclic subgroup of  $\mathbb{F}^*$
  7. Let  $Z(X) = \prod_{a \in H^*} (X - a)$
  8. Let  $A_i$  be a witness lookup columns and  $S_i$  be a table columns,  $i = 0, \dots, m$ .
- 

### Preprocessing:

#### 2.4.1 Prover View

1. Choose masking polynomials:

$$h_i(X) \leftarrow \mathbb{F}_{<k}[X] \text{ for } 0 \leq i < N_{\text{wires}}$$

**Remark:** For details on choice of  $k$ , we refer the reader to [1].

2. Define new witness polynomials:

$$f_i(X) = \mathbf{f}_i(X) + h_i(X)Z(X) \text{ for } 0 \leq i < N_{\text{wires}}$$

3. Add commitments to  $f_i$  to transcript
4. Get  $\theta \in \mathbb{F}$  from  $\text{hash}(\text{transcript})$
5. Construct the witness lookup compression and table compression  $S(\theta)$  and  $A(\theta)$ :

$$\begin{aligned} A(\theta) &= \theta^{m-1}A_0 + \theta^{m-2}A_1 + \dots + \theta A_{m-2} + A_{m-1} \\ S(\theta) &= \theta^{m-1}S_0 + \theta^{m-2}S_1 + \dots + \theta S_{m-2} + S_{m-1} \end{aligned}$$

6. Produce the permutation polynomials  $S'(X)$  and  $A'(X)$  such that:

6.1 All the cells of column  $A'$  are arranged so that like-valued cells are vertically adjacent to each other.

6.2 The first row in a sequence of values in  $A'$  is the row that has the corresponding value in  $S'$ .

7. Compute and add commitments to  $A'$  and  $S'$  to transcript

8. Get  $\beta, \gamma \in \mathbb{F}$  from  $\text{hash}(\text{transcript})$

9. For  $0 \leq i < N_{\text{perm}}$

$$\begin{aligned} p_i &= f_i + \beta \cdot S_{id_i} + \gamma \\ q_i &= f_i + \beta \cdot S_{\sigma_i} + \gamma \end{aligned}$$

10. Define:

$$\begin{aligned} p'(X) &= \prod_{0 \leq i < N_{\text{perm}}} p_i(X) \in \mathbb{F}_{<N_{\text{perm}} \cdot n}[X] \\ q'(X) &= \prod_{0 \leq i < N_{\text{perm}}} q_i(X) \in \mathbb{F}_{<N_{\text{perm}} \cdot n}[X] \end{aligned}$$

11. Compute  $P(X), Q(X) \in \mathbb{F}_{<n+1}[X]$ , such that:

$$\begin{aligned} P(\omega) &= Q(\omega) = 1 \\ P(\omega^i) &= \prod_{1 \leq j < i} p'(\omega^j) \text{ for } i \in 2, \dots, n+1 \\ Q(\omega^i) &= \prod_{1 \leq j < i} q'(\omega^j) \text{ for } i \in 2, \dots, n+1 \end{aligned}$$

12. Compute and add commitments to  $P$  and  $Q$  to transcript

13. Compute permutation product column:

$$\begin{aligned} V(\omega^i) &= \frac{(\theta^{m-1}A_0(\omega^i) + \theta^{m-2}A_1(\omega^i) + \dots + \theta A_{m-2}(\omega^i) + A_{m-1}(\omega^i) + \beta) \cdot (\theta^{m-1}S_0(\omega^i) + \theta^{m-2}S_1(\omega^i) + \dots + \theta S_{m-2}(\omega^i) + S_{m-1}(\omega^i) + \gamma)}{(A'(\omega^i) + \beta)(S'(\omega^i) + \gamma)} \\ V(1) &= V(\omega^{N_{\text{lookups}}}) = 1 \end{aligned}$$

14. Compute and add commitments to  $V$  to transcript

15. Get  $\alpha_0, \dots, \alpha_5 \in \mathbb{F}$  from  $\text{hash}(\text{transcript})$

16. Get  $\tau$  from  $\text{hash}(\text{transcript})$

17. Define polynomials ( $F_0, \dots, F_4$  - copy-satisfiability,  $\text{gate}_0$  is  $PI$ -constraining gate):

$$\begin{aligned} F_0(X) &= L_1(X)(P(X) - 1) \\ F_1(X) &= L_1(X)(Q(X) - 1) \\ F_2(X) &= P(X)p'(X) - P(X\omega) \\ F_3(X) &= Q(X)q'(X) - Q(X\omega) \\ F_4(X) &= L_n(X)(P(X\omega) - Q(X\omega)) \\ F_5(X) &= \sum_{0 \leq i < N_{\text{sel}}} (\tau^i \cdot \mathbf{q}_i(X) \cdot \text{gate}_i(X)) + PI(X) \end{aligned}$$

18. For the lookup:

18.1 Two selectors  $q_{last}$  and  $q_{blind}$  are used, where  $q_{last} = 1$  for  $t$  last blinding rows and  $q_{blind} = 1$  on the row in between the usable rows and the blinding rows.

18.2  $F_6(X) = L_0(X)(1 - V(X))$

18.3  $F_7(X) = q_{last} \cdot (V(X)^2 - V(X))$

18.4  $F_8(X) = (1 - (q_{last} + q_{blind})) \cdot (V(\omega X)(A'(X) + \beta)(S'(X) + \gamma) - V(X)(\theta^{m-1}A_0(X) + \dots + A_{m-1}(X) + \beta)(\theta^{m-1}S_0(X) + \dots + S_{m-1}(X) + \gamma))$

18.5  $F_9(X) = L_0(X) \cdot (A'(X) - S'(X))$

18.6  $F_{10}(X) = (1 - (q_{last} + q_{blind})) \cdot (A'(X) - S'(X)) \cdot (A'(X) - A'(\omega^{-1}X))$

19. Compute:

$$F(X) = \sum_{i=0}^{10} \alpha_i F_i(X)$$

$$T(X) = \frac{F(X)}{Z(X)}$$

20.  $N_T := \max(N_{\text{perm}}, \text{deg}_{\text{gates}} - 1)$ , where  $\text{deg}_{\text{gates}}$  is the highest degree of the degrees of gate polynomials.

21. Split  $T(X)$  into separate polynomials  $T_0(X), \dots, T_{N_T-1}(X)$ <sup>2</sup>

22. Add commitments to  $T_0(X), \dots, T_{N_T-1}(X)$  to transcript

23. Get  $y \in \mathbb{F}/H$  from  $\text{hash}_{\mathbb{F}/H}(\text{transcript})$

24. Run evaluation scheme with the committed polynomials and  $y$

**Remark:** Depending on the circuit, evaluation can be done also on  $y\omega, y\omega^{-1}$ .

25. The proof is  $\pi_{\text{comm}}$  and  $\pi_{\text{eval}}$ , where:

- $\pi_{\text{comm}} = \{f_{0,\text{comm}}, \dots, f_{N_{\text{wires}}-1,\text{comm}}, P_{\text{comm}}, Q_{\text{comm}}, T_{0,\text{comm}}, \dots, T_{N_T-1,\text{comm}}, A'_{\text{comm}}, S'_{\text{comm}}, V_{\text{comm}}\}$
- $\pi_{\text{eval}}$  is evaluation proofs for  $f_0(y), \dots, f_{N_{\text{wires}}}(y), P(y), P(y\omega), Q(y), Q(y\omega), T_0(y), \dots, T_{N_T-1}(y), A'(y), A'(y\omega^{-1}), S'(y), V(y), V(y\omega)$

### 2.4.2 Verifier View

1. Let  $f_{0,\text{comm}}, \dots, f_{N_{\text{wires}}-1,\text{comm}}$  be commitments to  $f_0(X), \dots, f_{N_{\text{wires}}-1}(X)$
2.  $\text{transcript} = \text{setup\_values} || f_{0,\text{comm}} || \dots || f_{N_{\text{wires}}-1,\text{comm}}$
3.  $\theta = \text{hash}(\text{transcript})$
4. Let  $A'_{\text{comm}}, S'_{\text{comm}}$  be commitments to  $A'(X), S'(X)$ .
5.  $\text{transcript} = \text{transcript} || A'_{\text{comm}} || S'_{\text{comm}}$
6.  $\beta, \gamma = \text{hash}(\text{transcript})$
7. Let  $P_{\text{comm}}, Q_{\text{comm}}, V_{i,\text{comm}}$  be commitments to  $P(X), Q(X), V(X)$ .
8.  $\text{transcript} = \text{transcript} || P_{\text{comm}} || Q_{\text{comm}} || V_{\text{comm}}$
9.  $\alpha_0, \dots, \alpha_5 = \text{hash}(\text{transcript})$
10.  $\tau = \text{hash}(\text{transcript})$
11.  $N_T := \max(N_{\text{perm}}, \text{deg}_{\text{gates}} - 1)$ , where  $\text{deg}_{\text{gates}}$  is the highest degree of the degrees of gate polynomials.
12. Let  $T_{0,\text{comm}}, \dots, T_{N_T-1,\text{comm}}$  be commitments to  $T_0(X), \dots, T_{N_T-1}(X)$
13.  $\text{transcript} = \text{transcript} || T_{0,\text{comm}} || \dots || T_{N_T-1,\text{comm}}$
14.  $y = \text{hash}_{\mathbb{F}/H}(\text{transcript})$
15. Run evaluation scheme verification with the committed polynomials and  $y$  to get values  $f_i(y), P(y), P(y\omega), Q(y), Q(y\omega), T_j(y), A'(y), S'(y), V(y), A'(y\omega^{-1}), V(y\omega)$ .  
**Remark:** Depending on the circuit, evaluation can be done also on  $f_i(y\omega), f_i(y\omega^{-1})$  for some  $i$ .
16. Calculate:

$$F_0(y) = L_1(y)(P(y) - 1)$$

$$F_1(y) = L_1(y)(Q(y) - 1)$$

$$p'(y) = \prod p_i(y) = \prod f_i(y) + \beta \cdot S_{id_i}(y) + \gamma$$

$$F_2(y) = P(y)p'(y) - P(y\omega)$$

$$q'(y) = \prod q_i(y) = \prod f_i(y) + \beta \cdot S_{\sigma_i}(y) + \gamma$$

$$F_3(y) = Q(y)q'(y) - Q(y\omega)$$

---

<sup>2</sup>Commit scheme supposes that polynomials should be degree  $\leq n$



$$\begin{aligned}
F_4(y) &= L_n(y)(P(y\omega) - Q(y\omega)) \\
F_5(y) &= \sum_{0 \leq i < N_{sel}} (\tau^i \cdot \mathbf{q}_i(y) \cdot \mathbf{gate}_i(y)) + PI(y) \\
T(y) &= \sum_{0 \leq j < N_T} y^{n \cdot j} T_j(y) \quad F_6(y) = L_0(y)(1 - V(y)) \\
F_7(y) &= q_{last} \cdot (V(y)^2 - V(y)) \\
F_8(y) &= (1 - (q_{last} + q_{blind})) \cdot (V(\omega y)(A'(y) + \beta)(S'(y) + \gamma) - V(y)(\theta^{m-1}A_0(y) + \dots + A_{m-1}(y) + \\
&\quad \beta)(\theta^{m-1}S_{i,0}(y) + \dots + S_{m-1}(y) + \gamma)) \\
F_9(y) &= L_0(y) \cdot (A'(y) - S'(y)) \\
F_{10}(y) &= (1 - (q_{last} + q_{blind})) \cdot (A'(y) - S'(y)) \cdot (A'(y) - A'(\omega^{-1}y))
\end{aligned}$$

17. Check the identity:

$$\sum_{i=0}^{10} \alpha_i F_i(y) = Z(y)T(y)$$

## 2.5 Circuit Definition

This section contains a description of PLONK-style circuits for In-EVM Solana's "Light Client" state verification<sup>3</sup>.

This section provides a high-level overview of the circuit used for proof generation and verification. Following sections provide sub-circuits details.

### 2.5.1 Verification Circuit Overview

Let bank-hashes of proving block set be  $\{H_{B_{n_1}}, \dots, H_{B_{n_2}}\}$ . The last confirmed block is  $H_{B_L}$ . Each positively confirmed block is signed by  $M$  validators.

Denote by `block_data` the data that is included in the bank hash other than the bank hash of the parent block.

1.  $H_{B_{n_1}} = H_{B_L} // H_{B_L}$  is a public input
2. Validator set constraints. // see Section 2.5.8
3. for  $i$  from  $n_1 + 1$  to  $n_2 + 32$ :

$$3.1 \quad H_{B_i} = \text{sha256}(\text{block\_data} || H_{B_{i-1}}) // \text{ see Section 2.5.2}$$

4. for  $j$  from 0 to  $M$ :

$$4.1 \quad \text{Ed25519 constraints for } H_{B_{n_2+32}} // \text{ see Section 2.5.6}$$

5. Merkle tree constraints for the set  $\{H_{B_{n_1}}, \dots, H_{B_{n_2}}\}$  // see Section 2.5.5

### 2.5.2 SHA-256 Circuit

Suppose that input data is in the 32-bits form, which is already padded to the required size. We suppose that the checking that chunked input data corresponds to the original data out of the circuit. However, we do not need to range constrain these chunks as we get them for free from the SHA-256 circuit.

Thus, the preprocessing constraints for the SHA-256 circuit is a decomposition of  $k$  message blocks to 32 bits chunks without range proofs. For 'Solana-EVM' circuit,  $k = 3$ .

**Lookup tables** We use the following lookup tables:

1. **SHA-256 NORMALIZE4** with 2 columns and  $2^{14}$  rows. The first column contains all possible 14-bits words. The second column contains corresponding sparse representations. The constraints can be used for the range check and sparse representation simultaneously.

<sup>3</sup><https://blog.nil.foundation/2021/10/14/solana-ethereum-bridge.html>

**The function  $\sigma_0$**  contains sparse mapping with base 4. Let  $a$  be divided to chunks  $a_0, a_1, a_2, a_3$  which equals to 3, 4, 11, 14 bits respectively. The values  $a'_0, a'_1, a'_2, a'_3$  are in sparse form, and  $a'$  is a sparse  $a$ . **SHA-256 NORMALIZE4** lookup table is used for mapping to sparse representation and range-constraining for each chunk  $a_i$ , where bit-length of  $a_i > 3$ . If a chunk is 14 bits long, then it is constrained for free. Else the prover has to calculate the sparse representation  $\hat{a}_i$  for  $2^j \cdot a_i$ , where  $j + \text{len}(a_i) = 14$  and  $\text{len}(a_i)$  is bit-length of  $a_i$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$\hat{a}_1$	$\hat{a}_2$	$a_3$	$a$
$j + 1$	$a'_0$	$a_1$	$a_2$	$a'_3$	—
$j + 2$	$s_0$	$a'_1$	$a'_2$	$s_3$	$\sigma_0$
$j + 3$	$b$	$s_1$	$s_2$	$c$	$w$

Constraints:

$$\begin{aligned}
w_{o,j} &= w_{1,j} + w_{2,j+1} \cdot 2^3 + w_{3,j+1} \cdot 2^7 + w_{4,j} \cdot 2^{18} \\
w_{o,j+3} &= w_{1,j+2} + w_{2,j+3} \cdot 2^3 + w_{3,j+3} \cdot 2^7 + w_{4,j+2} \cdot 2^{18} \\
(w_{1,j} - 7) \cdot (w_{1,j} - 6) \cdot \dots \cdot w_{1,j} &= 0 \\
10 \text{ plookup constraints: } &(w_{1,j}, w_{1,j+1}), (2^{10} \cdot w_{2,j+1}, w_{2,j}), (w_{2,j+1}, w_{2,j+2}), (2^3 \cdot \\
&w_{3,j+1}, w_{3,j}), (w_{3,j+1}, w_{3,j+2}), (w_{4,j}, w_{4,j+1}), (w_{1,j+2}, (w_{3,j+2} + w_{4,j+1} + w_{2,j+2})), (w_{2,j+3}, (w_{4,j+1} + \\
&w_{1,j+1} + w_{3,j+2}), (w_{3,j+3}, (w_{1,j+1} + w_{2,j+2} + w_{4,j+1})), (w_{4,j+2}, (w_{2,j+2} + w_{3,j+2}))
\end{aligned}$$

**The function  $\sigma_1$**  contains sparse mapping subcircuit with base 4. Let  $a$  be divided to chunks  $a_0, a_1, a_2, a_3$  which equals to 10, 7, 2, 13 bits respectively. The values  $a'_0, a'_1, a'_2, a'_3$  are in sparse form and  $a'$  is a sparse  $a$ . **SHA-256 NORMALIZE4** lookup table is used for mapping to sparse representation and range-constraining in the same way as for  $\sigma_0$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$s_0$	$s_1$	—	$s_3$	$\sigma_1$
$j + 1$	$a'_0$	$a'_1$	$s_2$	$a'_3$	—
$j + 2$	$a_0$	$a_1$	$a'_2$	$a_3$	—
$j + 3$	$\hat{a}_0$	$\hat{a}_1$	$a_2$	$\hat{a}_3$	$a$

Constraints:

$$\begin{aligned}
w_{o,j} &= w_{1,j} + w_{2,j} \cdot 2^{10} + w_{3,j+1} \cdot 2^{17} + w_{4,j} \cdot 2^{19} \\
w_{o,j+3} &= w_{1,j+2} + w_{2,j+2} \cdot 2^{10} + w_{3,j+3} \cdot 2^{17} + w_{4,j+2} \cdot 2^{19} \\
(w_{3,j+3} - 3) \cdot (w_{3,j+3} - 2) \cdot \dots \cdot w_{3,j+3} &= 0 \\
11 \text{ plookup constraints: } &(2^4 \cdot (w_{1,j+2}, w_{1,j+1})), (2^7 \cdot w_{2,j+2}, w_{2,j+3}), (2 \cdot \\
&w_{3,j+2}, w_{3,j+3}), (w_{1,j+2}, w_{1,j+1}), (w_{2,j+2}, w_{2,j+1}), (w_{3,j+3}, w_{3,j+2}), (w_{4,j+2}, w_{4,j+1}), (w_{1,j}, (w_{3,j+2} + \\
&w_{4,j+1} + w_{2,j+1})), (w_{2,j}, (w_{4,j+1} + w_{1,j+1} + w_{3,j+2})), (w_{3,j+1}, (w_{1,j+1} + w_{3,j+2} + w_{4,j+1})), (w_{4,j}, (w_{2,j+1} + w_{3,j+2}))
\end{aligned}$$

**Message scheduling** For each block of 512 bits of the padded message the 64 words are constructed in the following way:

- The first 16 words are obtained by splitting the message.
- The last 48 words are obtained by using the functions  $\sigma_0, \sigma_1$ :

$$W_i = \sigma_1(W_{i-2}) \oplus W_{i-7} \oplus \sigma_0(W_{i-15}) \oplus W_{i-16}$$

The first 16 words require a range check. We get it for free from range-constraining chunks inside functions  $\sigma_0$  and  $\sigma_1$ . Thus, for  $i$  from 16 to 63:

1. Apply  $\sigma_0$  to  $W_{i-15}$ . It costs 4 rows.
2. Add the following constraint for  $W_i$ :

$$w_{o,3} = w_{1,3} + w_{4,3} + w_{o,2} + w_{o,4},$$

where rows 2 and 3 belong to the  $\sigma_0$  table and row 4 is the first row of the  $\sigma_1$  table. This constraint does not require additional rows.

3. Apply  $\sigma_1$  to  $W_{i-2}$ . It costs 4 rows.

Thus, the message schedule takes  $8 \cdot 48 = 384$  rows.

**The  $\Sigma_0$  function** contain sparse mapping subcircuit with base 2. Let  $a$  be divided to 8 bits-chunks  $a_0, a_1, a_2, a_3$ . The values  $a'_0, a'_1, a'_2, a'_3$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-256 NORMALIZE4**: Read  $a_i$  to  $a'_i$
2. **SHA-256 8ROT2 32**: Read  $a'_0$  to  $r_1$
3. **SHA-256 8ROT5 32**: Read  $a'_1$  to  $r_2$
4. **SHA-256 8ROT6 32**: Read  $a'_2$  to  $r_3$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a'$
$j + 2$	$r_1$	$r_2$	$r_3$		$\Sigma_0$

Sparse map gate constraints:

$$\begin{aligned}
w_{o,j} &= w_{1,j} + w_{2,j} \cdot 2^8 + w_{3,j} \cdot 2^{8 \cdot 2} + w_{4,j} \cdot 2^{8 \cdot 3} \\
w_{o,j+1} &= w_{1,j+1} + w_{2,j+1} \cdot 4^8 + w_{3,j+1} \cdot 4^{8 \cdot 2} + w_{4,j+1} \cdot 4^{8 \cdot 3} \\
w_{o,j+2} &= w_{2,j+1} \cdot 4^{8-2} + w_{3,j+1} \cdot 4^{8 \cdot 2-2} + w_{4,j+1} \cdot 4^{8 \cdot 3-2} + w_{1,j+1} \cdot 4^{8 \cdot 3-5} + w_{3,j+1} \cdot 4^{8-5} + w_{4,j+1} \cdot 4^{8 \cdot 2-5} + \\
&\quad w_{1,j+1} \cdot 4^{8 \cdot 2-6} + w_{2,j+1} \cdot 4^{8 \cdot 3-6} + w_{4,j+1} \cdot 4^{8-6} + w_{1,j+2} + w_{2,j+2} + w_{3,j+2}
\end{aligned}$$

7 plookup constraints

**The  $\Sigma_1$  function** contain sparse mapping subcircuit with base 2. Let  $a$  be divided to 8 bits-chunks  $a_0, a_1, a_2, a_3$ . The values  $a'_0, a'_1, a'_2, a'_3$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-256 NORMALIZE7**: Read  $a_i$  to  $a'_i$
2. **SHA-256 8ROT6 32**: Read  $a'_0$  to  $r_1$
3. **SHA-256 8ROT3 32**: Read  $a'_1$  to  $r_2$
4. **SHA-256 8ROT1 32**: Read  $a'_3$  to  $r_3$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a'$
$j + 2$	$r_1$	$r_2$	$r_3$		$\Sigma_1$

Sparse map gate constraints:

$$\begin{aligned}
w_{o,j} &= w_{1,j} + w_{2,j} \cdot 2^8 + w_{3,j} \cdot 2^{8 \cdot 2} + w_{4,j} \cdot 2^{8 \cdot 3} \\
w_{o,j+1} &= w_{1,j+1} + w_{2,j+1} \cdot 7^8 + w_{3,j+1} \cdot 7^{8 \cdot 2} + w_{4,j+1} \cdot 7^{8 \cdot 3} \\
w_{o,j+2} &= w_{2,j+1} \cdot 7^{8-6} + w_{3,j+1} \cdot 7^{8 \cdot 2-6} + w_{7,j+1} \cdot 4^{8 \cdot 3-6} + w_{1,j+1} \cdot 7^{8 \cdot 3-3} + w_{3,j+1} \cdot 7^{8-3} + w_{4,j+1} \cdot 7^{8 \cdot 2-3} + \\
&\quad w_{1,j+1} \cdot 7^{8-1} + w_{2,j+1} \cdot 7^{8 \cdot 2-1} + w_{3,j+1} \cdot 7^{8 \cdot 3-1} + w_{1,j+2} + w_{2,j+2} + w_{3,j+2}
\end{aligned}$$

7 plookup constraints

The sparse values  $\Sigma_0$  and  $\Sigma_1$  have to be normalized. We use **SHA256 NORMALIZE4** and **SHA256 NORMALIZE7**.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	
$j + 1$	$a_0$	$a_1$	$a_2$	$a_3$	$\Sigma_i$

Normalize gate constraints:

$$\begin{aligned}
w_{o,j-1} &= w_{4,j} \cdot 4^8 \cdot 3 + w_{3,j} \cdot 4^8 \cdot 2 + w_{2,j} \cdot 4^8 + w_{1,j} \text{ for } \Sigma_1 \text{ replace 4 with 7} \\
w_{o,i} &= w_{4,i} \cdot 256^3 + w_{3,i} \cdot 256^2 + w_{2,i} \cdot 256 + w_{1,i}
\end{aligned}$$

7 plookup constraints

**The Maj function** contain sparse mapping subcircuit with base 2 for  $a, b, c$ . Let  $a; b; c$  be divided to 8 bits-chunks  $a_0, a_1, a_2, a_3; b_0, b_1, b_2, b_3; c_0, c_1, c_2, c_3$ . The values  $a'_0, a'_1, a'_2, a'_3$  are in sparse form, and  $a'$  is a sparse  $a$ . Similarly for  $b$  and  $c$ . Note, that  $a$  we already have in the sparse from  $\Sigma_0$  in the circuit. The variables  $b$  and  $c$  were represented in sparse form in the previous rounds or it is public inputs.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
j - k	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a'$
...					
j - l	$b'_0$	$b'_1$	$b'_2$	$b'_3$	$b'$
...					
j - t	$c'_0$	$c'_1$	$c'_2$	$c'_3$	$c'$
...					
j + 0	$a'$	$b'$	$c'$		maj

Sparse map gate constraints:

$$w_{o,j} = w_{1,j} + w_{2,j} + w_{3,j}$$

The sparse values  $maj$  have to be normalized. We use **SHA256 MAJ NORMALIZE4**

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
j + 0	$a'_0$	$a'_1$	$a'_2$	$a'_3$	
j + 1	$a_0$	$a_1$	$a_2$	$a_3$	$maj$

Normalize gate constraints:

$$w_{o,i} = w_{4,i} \cdot 256^3 + w_{3,i} \cdot 256^2 + w_{2,i} \cdot 256 + w_{1,i}$$

The final addition requires one add gate.

**The Ch function** contain sparse mapping subcircuit with base 2 for  $e, f, g$ . Let  $e; f; g$  be divided to 8 bits-chunks  $e_0, e_1, e_2, e_3; f_0, f_1, f_2, f_3; g_0, g_1, g_2, g_3$ . The values  $e'_0, e'_1, e'_2, e'_3$  are in sparse form, and  $e'$  is a sparse  $e$ . Similarly for  $b$  and  $c$ . Note, that  $e$  we already have in the sparse from  $\Sigma_1$  in the circuit. The variables  $f$  and  $g$  were represented in sparse form in the previous rounds or it is public inputs.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
j - k	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a'$
...					
j - l	$b'_0$	$b'_1$	$b'_2$	$b'_3$	$b'$
...					
j - t	$c'_0$	$c'_1$	$c'_2$	$c'_3$	$c'$
...					
j + 0	$a'$	$b'$	$c'$		ch

Sparse map gate constraints:

$$w_{o,j} = w_{1,j} + 2 * w_{2,j} + 3 * w_{3,j}$$

The sparse values  $ch$  have to be normalized. We use **SHA256 CH NORMALIZE7**

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
j + 0	$a'_0$	$a'_1$	$a'_2$	$a'_3$	
j + 1	$a_0$	$a_1$	$a_2$	$a_3$	$ch$

Normalize gate constraints:

$$w_{o,i} = w_{4,i} \cdot 256^3 + w_{3,i} \cdot 256^2 + w_{2,i} \cdot 256 + w_{1,i}$$

The final addition requires one add gate.

The updating of variables for new rounds costs 10 add gates.

Producing the final hash value costs two add gates.

### 2.5.3 SHA2-512 Circuit

SHA-512 uses the similar logical functions as in 2.5.2 which operates on 64-bits words. Thus each input uses the same range proof which extended to 64-bits.

**Range proof that  $a < 2^{64}$**  Let  $a = \{a_0, \dots, a_{32}\}$ , where  $a_i$  is two bits.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_{29}$	$a_{30}$	$a_{31}$	$a_{32}$	acc
$j + 1$	$a_{25}$	$a_{26}$	$a_{27}$	$a_{28}$	acc
...					
$j + 6$	$a_4$	$a_5$	$a_6$	$a_7$	acc
$j + 7$	$a_0$	$a_1$	$a_2$	$a_3$	a

Range gate constraints:

$$w_{1,i}(w_{1,i} - 1)(w_{1,i} - 2)(w_{1,i} - 3) + w_{2,i}(w_{2,i} - 1)(w_{2,i} - 2)(w_{2,i} - 3) + w_{3,i}(w_{3,i} - 1)(w_{3,i} - 2)(w_{3,i} - 3) + w_{4,i}(w_{4,i} - 1)(w_{4,i} - 2)(w_{4,i} - 3) \\ w_{o,i} = w_{o,i-1} \cdot 4^4 + w_{4,i} \cdot 4^3 + w_{3,i} \cdot 4^2 + w_{2,i} \cdot 4 + w_{1,i}$$

The range proofs are included for each input data block.

**The function  $\sigma_0$**  contain sparse mapping subcircuit with base 4. Let  $a$  be divided to 8 bits-chunks  $a_0, a_1, a_2, \dots, a_7$ . The values  $a'_0, a'_1, a'_2, \dots, a'_7$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-256 NORMALIZE4**: Read  $a_i$  to  $a'_i$
2. **SHA-512 8ROT1 64**: Read  $a'_0$  to  $r_1$
3. **SHA-512 8SHR7 64**: Read  $a'_0$  to  $r_3$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a_4$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	a
$j + 2$	$a_5$	$a_6$	$a_7$	$a'_4$	$\sigma_0$
$j + 3$	$a'_5$	$a'_6$	$a'_7$	$r_1$	$r_2$

Sparse map gate constraints:

$$w_{o,j+1} = w_{1,j} + w_{2,j} \cdot 2^8 + w_{3,j} \cdot 2^{8 \cdot 2} + w_{4,j} \cdot 2^{8 \cdot 3} + w_{o,j} \cdot 2^{8 \cdot 4} + w_{1,j+2} \cdot 2^{8 \cdot 5} + w_{2,j+2} \cdot 2^{8 \cdot 6} + w_{3,j+2} \cdot 2^{8 \cdot 7} \\ w_{o,j+2} = w_{2,j+1} \cdot 4^{8-1} + w_{3,j+1} \cdot 4^{8 \cdot 2-1} + w_{4,j+1} \cdot 4^{8 \cdot 3-1} + w_{4,j+2} \cdot 4^{8 \cdot 4-1} + w_{1,j+3} \cdot 4^{8 \cdot 5-1} + w_{2,j+3} \cdot 4^{8 \cdot 6-1} + w_{3,j+3} \cdot 4^{8 \cdot 7-1} + w_{1,j+1} \cdot 4^{8 \cdot 7} + w_{2,j+1} + w_{3,j+1} \cdot 4^8 + w_{4,j+1} \cdot 4^{8 \cdot 2} + w_{4,j+2} \cdot 4^{8 \cdot 3} + w_{1,j+3} \cdot 4^{8 \cdot 4} + w_{2,j+3} \cdot 4^{8 \cdot 5} + w_{3,j+3} \cdot 4^{8 \cdot 6} + w_{2,j+1} \cdot 4^{8-7} + w_{3,j+1} \cdot 4^{8 \cdot 2-7} + w_{4,j+1} \cdot 4^{8 \cdot 3-7} + w_{4,j+2} \cdot 4^{8 \cdot 4-7} + w_{1,j+3} \cdot 4^{8 \cdot 5-7} + w_{2,j+3} \cdot 4^{8 \cdot 6-7} + w_{3,j+3} \cdot 4^{8 \cdot 7-7} + w_{4,j+3} + w_{o,j+3} \\ 10 \text{ lookup constraints}$$

**The function  $\sigma_1$**  contain sparse mapping subcircuit with base 4. Let  $a$  be divided to 8 bits-chunks  $a_0, a_1, a_2, \dots, a_7$ . The values  $a'_0, a'_1, a'_2, \dots, a'_7$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-256 NORMALIZE4**: Read  $a_i$  to  $a'_i$
2. **SHA-512 8ROT3 64**: Read  $a'_2$  to  $r_1$
3. **SHA-512 8ROT5 SHR6 64**: Read  $a'_7 + a'_0$  to  $r_2$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a_4$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	a
$j + 2$	$a_5$	$a_6$	$a_7$	$a'_4$	$\sigma_1$
$j + 3$	$a'_5$	$a'_6$	$a'_7$	$r_1$	$r_2$

Sparse map gate constraints:

$$w_{o,j+1} = w_{1,j} + w_{2,j} \cdot 2^8 + w_{3,j} \cdot 2^{8 \cdot 2} + w_{4,j} \cdot 2^{8 \cdot 3} + w_{o,j} \cdot 2^{8 \cdot 4} + w_{1,j+2} \cdot 2^{8 \cdot 5} + w_{2,j+2} \cdot 2^{8 \cdot 6} + w_{3,j+2} \cdot 2^{8 \cdot 7} \\ w_{o,j+2} = w_{1,j+1} \cdot 4^{64-19} + w_{2,j+1} \cdot 4^{64+(8-19)} + w_{4,j+1} \cdot 4^{8 \cdot 3-19} + w_{4,j+2} \cdot 4^{8 \cdot 4-19} + w_{1,j+3} \cdot 4^{8 \cdot 5-19} + w_{2,j+3} \cdot 4^{8 \cdot 6-19} + w_{3,j+3} \cdot 4^{8 \cdot 7-19} + w_{1,j+1} \cdot 4^{64-61} + w_{2,j+1} \cdot 4^{64+(8-61)} + w_{3,j+1} \cdot 4^{64+(8 \cdot 2-61)} + w_{4,j+1} \cdot 4^{64+(8 \cdot 3-61)} + w_{4,j+2} \cdot 4^{64+(8 \cdot 4-61)} + w_{1,j+3} \cdot 4^{64+(8 \cdot 5-61)} + w_{2,j+3} \cdot 4^{64+(8 \cdot 6-61)} + w_{2,j+1} \cdot 4^{8-6} + w_{3,j+1} \cdot 4^{8 \cdot 2-6} + w_{4,j+1} \cdot 4^{8 \cdot 3-6} + w_{4,j+2} \cdot 4^{8 \cdot 4-6} + w_{1,j+3} \cdot 4^{8 \cdot 5-6} + w_{2,j+3} \cdot 4^{8 \cdot 6-6} + w_{3,j+3} \cdot 4^{8 \cdot 7-6} + w_{4,j+3} + w_{o,j+3} \\ 10 \text{ lookup constraints}$$

The sparse values  $\sigma_0$  and  $\sigma_1$  have to be normalized. The final addition requires one add gate. Note, that  $a'$  already initialized in the row  $j - 2$ . We use **SHA256 NORMALIZE4**

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$acc$
$j + 1$	$a_0$	$a_1$	$a_2$	$a_3$	0
$j + 2$	$a'_4$	$a'_5$	$a'_6$	$a'_7$	$\sigma_i$
$j + 3$	$a_4$	$a_5$	$a_6$	$a_7$	

Normalize gate constraints:

$$\begin{aligned}
w_{o,j+2} &= w_{4,j+1} \cdot 256^3 + w_{3,j+1} \cdot 256^2 + w_{2,j+1} \cdot 256 + w_{1,j+1} + w_{1,j+3} \cdot 256^4 \\
&\quad + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7 \\
w_{o,j} &= w_{o,j-2} - (w_{4,j} \cdot 256^3 + w_{3,j} \cdot 256^2 + w_{2,j} \cdot 256 + w_{1,j}) \\
w_{o,j+1} &= w_{o,j} - (w_{1,j+3} \cdot 256^4 + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7)
\end{aligned}$$

8 plookup constraints

**The  $\Sigma_0$  function** contain sparse mapping subcircuit with base 4. Let  $a$  be divided to 7-bits chunks  $a_0, a_1, a_2, a_3$  and 9 bits-chunks  $a_4, a_5, a_6, a_7$ . The values  $a'_0, a'_1, a'_2, \dots, a'_7$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-512 9NORMALIZE4**: Read  $a_i$  to  $a'_i$
2. **SHA-512 7NORMALIZE4**: Read  $a_i$  to  $a'_i$
3. **SHA-512 9ROT6 64**: Read  $a'_4$  to  $r_2$
4. **SHA-512 9ROT2 64**: Read  $a'_5$  to  $r_3$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a_4$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a$
$j + 2$	$a_5$	$a_6$	$a_7$	$a'_4$	$\Sigma_0$
$j + 3$	$a'_5$	$a'_6$	$a'_7$	$r_1$	$r_2$

Sparse map gate constraints:

$$\begin{aligned}
w_{o,j+1} &= w_{1,j} + w_{2,j} \cdot 2^7 + w_{3,j} \cdot 2^{7 \cdot 2} + w_{4,j} \cdot 2^{7 \cdot 3} + w_{o,j} \cdot 2^{7 \cdot 4} + w_{1,j+2} \cdot 2^{7 \cdot 4+9} + w_{2,j+2} \cdot 2^{7 \cdot 4+9 \cdot 2} + w_{3,j+2} \cdot 2^{7 \cdot 4+9 \cdot 3} \\
&\quad + w_{4,j+2} \cdot 2^{7 \cdot 4+9 \cdot 4} + w_{1,j+3} \cdot 4^9 + w_{2,j+3} \cdot 4^{9 \cdot 2} + w_{3,j+3} \cdot 4^{9 \cdot 3} + w_{1,j+1} \cdot 4^{9 \cdot 4} + w_{2,j+1} \cdot 4^{9 \cdot 4+7} \\
&\quad + w_{3,j+1} \cdot 4^{9 \cdot 4+7 \cdot 2} + w_{4,j+1} \cdot 4^{9 \cdot 4+7 \cdot 3} + w_{1,j+1} \cdot 4^{64-34} + w_{2,j+1} \cdot 4^{64+(7-34)} + w_{3,j+1} \cdot 4^{64+(7 \cdot 2-34)} + \\
&\quad + w_{4,j+1} \cdot 4^{64+(7 \cdot 3-34)} + w_{1,j+3} \cdot 4^{7 \cdot 4+9-34} + w_{2,j+3} \cdot 4^{7 \cdot 4+9 \cdot 2-34} + w_{3,j+3} \cdot 4^{7 \cdot 4+9 \cdot 3-34} + w_{1,j+1} \cdot 4^{64-39} + \\
&\quad + w_{2,j+1} \cdot 4^{64+(7-39)} + w_{3,j+1} \cdot 4^{64+(7 \cdot 2-39)} + w_{4,j+1} \cdot 4^{64+(7 \cdot 3-39)} + w_{4,j+2} \cdot 4^{64+(7 \cdot 4-39)} + w_{2,j+3} \cdot \\
&\quad + 4^{7 \cdot 4+9 \cdot 2-39} + w_{3,j+3} \cdot 4^{7 \cdot 4+9 \cdot 3-39} + w_{4,j+3} + w_{o,j+3}
\end{aligned}$$

10 plookup constraints

**The  $\Sigma_1$  function** contain sparse mapping subcircuit with base 7. Let  $a$  be divided to 7-bits chunks  $a_0, a_1, a_2, a_3$  and 9 bits-chunks  $a_4, a_5, a_6, a_7$ . The values  $a'_0, a'_1, a'_2, \dots, a'_7$  are in sparse form, and  $a'$  is a sparse  $a$ . We need the following lookup tables:

1. **SHA-512 9NORMALIZE7**: Read  $a_i$  to  $a'_i$
2. **SHA-512 7NORMALIZE7**: Read  $a_i$  to  $a'_i$
3. **SHA-512 7ROT4 32**: Read  $a'_2$  to  $r_2$
4. **SHA-512 9ROT4 32**: Read  $a'_5$  to  $r_3$

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a_0$	$a_1$	$a_2$	$a_3$	$a_4$
$j + 1$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a$
$j + 2$	$a_5$	$a_6$	$a_7$	$a'_4$	$\Sigma_1$
$j + 3$	$a'_5$	$a'_6$	$a'_7$	$r_1$	$r_2$

Sparse map gate constraints:

$$\begin{aligned}
w_{o,j+1} &= w_{1,j} + w_{2,j} \cdot 2^7 + w_{3,j} \cdot 2^{7 \cdot 2} + w_{4,j} \cdot 2^{7 \cdot 3} + w_{o,j} \cdot 2^{7 \cdot 4} + w_{1,j+2} \cdot 2^{7 \cdot 4+9} + w_{2,j+2} \cdot 2^{7 \cdot 4+9 \cdot 2} + w_{3,j+2} \cdot 2^{7 \cdot 4+9 \cdot 3} \\
w_{o,j+2} &= w_{3,j+1} + w_{4,j+1} \cdot 7^7 + w_{4,j+2} \cdot 7^{7 \cdot 2} + w_{1,j+3} \cdot 7^{7 \cdot 2+9} + w_{2,j+3} \cdot 7^{7 \cdot 2+9 \cdot 2} + w_{3,j+3} \cdot 7^{9 \cdot 3+7 \cdot 2} + w_{1,j+1} \cdot 7^{9 \cdot 4+7 \cdot 2} + \\
&w_{2,j+1} \cdot 7^{9 \cdot 4+7 \cdot 3} + w_{1,j+1} \cdot 7^{64-18} + w_{2,j+1} \cdot 7^{64+(7-18)} + w_{4,j+1} \cdot 7^{7 \cdot 3-18} + w_{4,j+2} \cdot 7^{7 \cdot 4-18} + w_{1,j+3} \cdot 7^{7 \cdot 4+9-18} + \\
&w_{2,j+3} \cdot 7^{7 \cdot 4+9 \cdot 2-18} + w_{3,j+3} \cdot 7^{7 \cdot 4+9 \cdot 3-18} + w_{1,j+1} \cdot 7^{64-41} + w_{2,j+1} \cdot 7^{64+(7-41)} + w_{3,j+1} \cdot 7^{64+(7 \cdot 2-41)} + \\
&w_{4,j+1} \cdot 7^{64+(7 \cdot 3-41)} + w_{4,j+2} \cdot 7^{64+(7 \cdot 3+9-41)} + w_{2,j+3} \cdot 7^{64+(7 \cdot 3+9 \cdot 2-41)} + w_{3,j+3} \cdot 7^{7 \cdot 3+9 \cdot 3-41} + w_{4,j+3} + w_{o,j+3}
\end{aligned}$$

10 plookup constraints

The sparse values  $\Sigma_0$  and  $\Sigma_1$  have to be normalized. We use **SHA256 NORMALIZE4** and **SHA256 NORMALIZE7**. Note, that  $a'$  already initialized in the row  $j - 2$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$a''$
$j + 1$	$a_0$	$a_1$	$a_2$	$a_3$	0
$j + 2$	$a'_4$	$a'_5$	$a'_6$	$a'_7$	$\Sigma_i$
$j + 3$	$a_4$	$a_5$	$a_6$	$a_7$	

Normalize gate constraints:

$$\begin{aligned}
w_{o,j+2} &= w_{4,j+1} \cdot 256^3 + w_{3,j+1} \cdot 256^2 + w_{2,j+1} \cdot 256 + w_{1,j+1} + w_{1,j+3} \cdot 256^4 \\
&\quad + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7 \\
w_{o,j} &= w_{1,j-3} + w_{2,j-3} \cdot 4^7 + w_{3,j-3} \cdot 4^{7 \cdot 2} + w_{4,j-3} \cdot 4^{7 \cdot 3} + w_{4,j-2} \cdot 4^{7 \cdot 4} + w_{1,j-1} \cdot 7^{7 \cdot 4+9} \\
&\quad + w_{2,j-1} \cdot 7^{7 \cdot 4+9 \cdot 2} + w_{2,j-1} \cdot 7^{7 \cdot 4+9 \cdot 3} \text{ for maj or ch function. For } \Sigma_1 \text{ replace 4 with 7} \\
&\quad w_{o,j+1} = \\
w_{o,j-2} &- (w_{4,j} \cdot 256^3 + w_{3,j} \cdot 256^2 + w_{2,j} \cdot 256 + w_{1,j} + w_{1,j+3} \cdot 256^4 + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7)
\end{aligned}$$

8 plookup constraints

**The Maj function** contain sparse mapping subcircuit with base 4 for  $a, b, c$ . Note, that the sparse chunks of  $a$  we already have in  $\Sigma_0$  in the circuit. The variables  $b$  and  $c$  were represented in sparse chunks in the previous rounds or it is public inputs.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j$	$a'$	$b'$	$c'$		maj

Sparse map gate constraints:

$$w_{o,j} = w_{1,j} + w_{2,j} + w_{3,j}$$

The sparse values  $maj$  have to be normalized. We use **SHA256 MAJ NORMALIZE4** Note, that the sparse  $maj$  already initialized in the row  $j - 1$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$acc$
$j + 1$	$a_0$	$a_1$	$a_2$	$a_3$	0
$j + 2$	$a'_4$	$a'_5$	$a'_6$	$a'_7$	$maj$
$j + 3$	$a_4$	$a_5$	$a_6$	$a_7$	

Normalize gate constraints:

$$\begin{aligned}
w_{o,j+2} &= w_{4,j+1} \cdot 256^3 + w_{3,j+1} \cdot 256^2 + w_{2,j+1} \cdot 256 + w_{1,j+1} + w_{1,j+3} \cdot 256^4 \\
&\quad + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7 \\
w_{o,j} &= w_{o,j-1} - (w_{4,j} \cdot 256^3 + w_{3,j} \cdot 256^2 + w_{2,j} \cdot 256 + w_{1,j}) \\
w_{o,j+1} &= w_{o,j} - (w_{1,j+3} \cdot 256^4 + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7)
\end{aligned}$$

8 plookup constraints

The final addition requires one add gate.

**The Ch function** contain sparse mapping subcircuit with base 7 for  $e, f, g$ . Note, that  $e$  we already have in the sparse from  $\Sigma_1$  in the circuit. The variables  $f$  and  $g$  were represented in sparse form in the previous rounds or it is public inputs.

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$e'$	$f'$	$g'$		ch

Sparse map gate constraints:

$$w_{o,j} = w_{1,j} + 2 \cdot w_{2,j} + 3 \cdot w_{3,j}$$

The sparse values  $ch$  have to be normalized. Note, that  $ch$  already initialized in the row  $j - 1$ . We use **SHA256 CH NORMALIZE7**

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$a'_0$	$a'_1$	$a'_2$	$a'_3$	$acc$
$j + 1$	$a_0$	$a_1$	$a_2$	$a_3$	$0$
$j + 2$	$a'_4$	$a'_5$	$a'_6$	$a'_7$	$ch$
$j + 3$	$a_4$	$a_5$	$a_6$	$a_7$	

Normalize gate constraints:

$$\begin{aligned}
w_{o,j+2} &= w_{4,j+1} \cdot 256^3 + w_{3,j+1} \cdot 256^2 + w_{2,j+1} \cdot 256 + w_{1,j+1} + w_{1,j+3} \cdot 256^4 + w_{2,j+3} \cdot 256^5 \\
&\quad + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7 \\
w_{o,j} &= w_{o,j-1} - (w_{4,j} \cdot 256^3 + w_{3,j} \cdot 256^2 + w_{2,j} \cdot 256 + w_{1,j}) \\
w_{o,j+1} &= w_{o,j} - (w_{1,j+3} \cdot 256^4 + w_{2,j+3} \cdot 256^5 + w_{3,j+3} \cdot 256^6 + w_{4,j+4} \cdot 256^7) \\
&\quad 8 \text{ plookup constraints}
\end{aligned}$$

The final addition requires one add gate.

The updating of variables for new rounds costs 10 add gates.

Producing the final hash value costs two add gates.

#### 2.5.4 Poseidon Circuit

Consider a poseidon permutation  $F : [0_{\mathbb{F}}, I[2], I[3]] \rightarrow [O[1], H, O[3]]$  of width 3 and  $\alpha = 5$ . The 1-call sponge function is used:

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$0_{\mathbb{F}}$	$I[2]$	$I[3]$	$T_{1,0}$	$T_{1,1}$
$j + 1$	$T_{1,2}$	$T_{2,0}$	$T_{2,1}$	$T_{2,2}$	$T_{3,0}$
$\dots$					
$j + 39$	$O[1]$	$H$	$O[3]$	$-$	$-$

Constraints:

$$\begin{aligned}
&\text{For 4 rounds:} \\
&[w_{4,j}, w_{o,j}, w_{1,j+1}] = [w_{1,j}^5, w_{2,j}^5, w_{3,j}^5] \times M + RC \\
&\text{For 57 rounds:} \\
&[w_{1,j+4}, w_{2,j+4}, w_{3,j+4}] = [w_{3,j+3}, w_{4,j+3}, w_{o,j+3}^5] \times M + RC \\
&\text{For 4 rounds:} \\
&[w_{2,j+37}, w_{3,j+37}, w_{4,j+37}] = [w_{4,j+36}^5, w_{o,j+36}^5, w_{1,j+37}^5] \times M + RC
\end{aligned}$$

#### 2.5.5 Merkle Tree Circuit

Merkle Tree generation for set  $\{H_{B_{n_1}}, \dots, H_{B_{n_2}}\}$ . Let  $k = \lceil \log(n_2 - n_1) \rceil$

1.  $n = n_2 - n_1$
2.  $2^k = n$
3. for  $i$  from 0 to  $n - 1$ :
  - 3.1  $T_i := H_i$  // just notation for simplicity, not a real part of the circuit
4. for  $i$  from 0 to  $k - 1$ :
  - 4.1 for  $j$  from 0 to  $(n - 1)/2$ :
    - 4.1.1  $T'_i = \text{hash}(T_{2 \cdot i}, T_{2 \cdot i + 1})$ . // see Section 2.5.4
  - 4.2  $n = \frac{n}{2}$
  - 4.3 for  $j$  from 0 to  $n - 1$ :
    - 4.3.1  $T_i := T'_i$ . // just notation for simplicity, not a real part of the circuit



### 2.5.6 Ed25519 Circuit

To verify a signature  $(R, s)$  on a message  $M$  using public key  $A$  and a generator  $B$  do:

1. Prove that  $s$  in the range  $L = 2^{252} + 27742317777372353535851937790883648493$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$s$	$z_0$	$z_1$	$z_2$	$z_3$
$\dots$					
$j + 5$	$z_{25}$			$--$	

Constraints:

$$w_{2,j} = w_{1,j} + 2^{253} - L$$

Each  $w_{i,k} - 2^{10} \cdot w_{next}$ , where  $i = 2, \dots, o$  for  $k = 0$  and  $i = 1, \dots, o$  for  $k = 1, \dots, 4$  is range-constrained by 10-bits plookup table.

$w_{1,j+5} \cdot 2^7$  is range-constrained by 10-bits plookup table.

2.  $k == \text{SHA-512}(\text{data} || R || A || M) // \text{See section ??}$

3.  $sB = ?R + kA$ :

3.1 Fixed-base scalar multiplication circuit is used for  $sB = S$

3.2 One addition is used for  $S + (-R)$ . The coordinates of  $R$  and  $T = S + (-R)$  are placed on the last row of fixed-base scalar multiplication circuit. In total, three constraints are used for addition:

$$\begin{aligned} x_t \cdot (1 + dx_s \cdot (-x_r) \cdot y_s \cdot y_r) &= x_s \cdot y_r + (-x_r) \cdot y_s \\ y_t \cdot (1 - dx_s \cdot (-x_r) \cdot y_s \cdot y_r) &= x_s \cdot (-x_r) + y_r \cdot y_s \\ -x_r^2 + y_r^2 &= 1 - d \cdot x_r^2 \cdot y_r^2 \end{aligned}$$

3.3 Variable-base scalar multiplication circuit has to be used in reversed order, where  $(x_n, y_n) = (x_t, y_t)$ .

### 2.5.7 Elliptic Curves Arithmetics

#### WIP

This section instantiates the arithmetic of edwards25519 curve:

$$-x^2 + y^2 = 1 - (121665/121666) \cdot x^2 \cdot y^2$$

Affine coordinates are used for points. Let  $d$  be equal to  $121665/121666$ .

**Fixed-base scalar multiplication circuit** : We precompute all values  $w(B, s, k) = k_i \cdot 8^s B$ , where  $k_i \in \{0, \dots, 7\}$ ,  $s \in \{0, \dots, 84\}$ .

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$b_{n-1}$	$b_{n-2}$	$b_{n-3}$	$u_1$	acc
$j + 1$	$x_2$	$y_2$	$b_{n-6}$	$u_2$	$v_1$
$j + 2$	$b_{n-4}$	$b_{n-5}$	$v_2$	$b_{n-7}$	acc
$j + 3$	$x_3$	$y_3$	$b_{n-8}$	$b_{n-9}$	$u_3$
$j + 4$	$x_4$	$y_4$	$v_3$	$--$	acc
$\dots$					
$j + 84$	$--$	$--$	$v_{85}$	$--$	$--$

Define the following functions:

1.  $\phi_1 : (x_1, x_2, x_3, x_4) \mapsto$   

$$x_3 \cdot (-u'_0 \cdot x_2 \cdot x_1 + u'_0 \cdot x_1 + u'_0 \cdot x_2 - u'_0 + u'_2 \cdot x_1 \cdot x_2 - u'_2 \cdot x_2 + u'_4 \cdot x_1 \cdot x_2 - u'_4 \cdot x_2 - u'_6 \cdot x_1 \cdot x_2 + u'_1 \cdot x_2 \cdot x_1 - u'_1 \cdot x_1 - u'_1 \cdot x_2 + u'_1 - u'_3 \cdot x_1 \cdot x_2 + u'_3 \cdot x_2 - u'_5 \cdot x_1 \cdot x_2 + u'_5 \cdot x_2 + u'_7 \cdot x_1 \cdot x_2) - (x_4 - u'_0 \cdot x_2 \cdot x_1 + u'_0 \cdot x_1 + u'_0 \cdot x_2 - u'_0 + u'_2 \cdot x_1 \cdot x_2 - u'_2 \cdot x_2 + u'_4 \cdot x_1 \cdot x_2 - u'_4 \cdot x_2 - u'_6 \cdot x_1 \cdot x_2)$$

2.  $\phi_2 : (x_1, x_2, x_3, x_4) \mapsto$   
 $x_3 \cdot (-v'_0 \cdot x_2 \cdot x_1 + v'_0 \cdot x_1 + v'_0 \cdot x_2 - v'_0 \cdot v'_2 \cdot x_1 \cdot x_2 - v'_2 \cdot x_2 + v'_4 \cdot x_1 \cdot x_2 - v'_4 \cdot x_2 - v'_6 \cdot x_1 \cdot x_2 + v'_1 \cdot$   
 $x_2 \cdot x_1 - v'_1 \cdot x_1 - v'_1 \cdot x_2 + v'_1 - v'_3 \cdot x_1 \cdot x_2 + v'_3 \cdot x_2 - v'_5 \cdot x_1 \cdot x_2 + v'_5 \cdot x_2 + v'_7 \cdot x_1 \cdot x_2) - (x_4 - v'_0 \cdot$   
 $x_2 \cdot x_1 + v'_0 \cdot x_1 + v'_0 \cdot x_2 - v'_0 \cdot v'_2 \cdot x_1 \cdot x_2 - v'_2 \cdot x_2 + v'_4 \cdot x_1 \cdot x_2 - v'_4 \cdot x_2 - v'_6 \cdot x_1 \cdot x_2)$
3.  $\phi_3 : (x_1, x_2, x_3, x_4, x_5, x_6) \mapsto$   
 $x_1 \cdot (1 + d \cdot x_3 \cdot x_4 \cdot x_5 \cdot x_6) - (x_3 \cdot x_6 + x_4 \cdot x_5)$
4.  $\phi_4 : (x_1, x_2, x_3, x_4, x_5, x_6) \mapsto$   
 $x_2 \cdot (1 - d \cdot x_3 \cdot x_4 \cdot x_5 \cdot x_6) - (x_3 \cdot x_5 + x_4 \cdot x_6)$

Constraints:

- For  $j + 0$ :
  - $w_{o,j} = w_{1,j} \cdot 2^2 + w_{2,j} \cdot 2 + w_{3,j}$
  - $\phi_3(w_{1,j+1}, w_{2,j+1}, w_{4,j}, w_{o,j+1}, w_{4,j+1}, w_{3,j+2}) = 0$
  - $\phi_4(w_{1,j+1}, w_{2,j+1}, w_{4,j}, w_{o,j+1}, w_{4,j+1}, w_{3,j+2}) = 0$
- For  $j + z, z \equiv 0 \pmod{5}, z \neq 0$ :
  - $w_{o,j+z} = w_{1,j+z} \cdot 2^2 + w_{2,j+z} \cdot 2 + w_{3,j+z} + w_{o,j+z-1} \cdot 2^3$
  - $\phi_1(w_{1,j+z}, w_{2,j+z}, w_{3,j+z}, w_{4,j+z}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z}{5}), i)$
  - $\phi_2(w_{1,j+z}, w_{2,j+z}, w_{3,j+z}, w_{o,j+z+1}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z}{5}), i)$
  - $\phi_3(w_{1,j+z+1}, w_{2,j+z+1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{4,j+z+1}, w_{3,j+z+2}) = 0$
  - $\phi_4(w_{1,j+z+1}, w_{2,j+z+1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{4,j+z+1}, w_{3,j+z+2}) = 0$
- For  $j + z, z \equiv 2 \pmod{5}$ :
  - $w_{o,j+z} = w_{1,j+z} \cdot 2^2 + w_{2,j+z} \cdot 2 + w_{3,j+z-1} + w_{o,j+z-2} \cdot 2^3$
  - $\phi_1(w_{1,j+z}, w_{2,j+z}, w_{3,j+z-1}, w_{4,j+z-1}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z-2}{5}) + 1, i)$
  - $\phi_2(w_{1,j+z}, w_{2,j+z}, w_{3,j+z-1}, w_{3,j+z}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z-2}{5}) + 1, i)$
  - $\phi_3(w_{1,j+z+1}, w_{2,j+z+1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{o,j+z+1}, w_{3,j+z+2}) = 0$
  - $\phi_4(w_{1,j+z+1}, w_{2,j+z+1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{o,j+z+1}, w_{3,j+z+2}) = 0$
- For  $j + z, z \equiv 3 \pmod{5}$ :
  - $\phi_1(w_{4,j+z-1}, w_{3,j+z}, w_{4,j+z}, w_{o,j+z}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z-3}{5}) + 2, i)$
  - $\phi_2(w_{4,j+z-1}, w_{3,j+z}, w_{4,j+z}, w_{3,j+z+1}) = 0$ , where  $(u'_i, v'_i) = w(B, 3 \cdot (\frac{z-3}{5}) + 2, i)$
- For  $j + z, z \equiv 4 \pmod{5}$ :
  - $w_{o,j+z} = w_{4,j+z-2} \cdot 2^2 + w_{3,j+z-3} \cdot 2 + w_{4,j+z-3} + w_{o,j+z-2} \cdot 2^3$
  - $\phi_3(w_{1,j+z-2}, w_{2,j+z}, w_{1,j+z-1}, w_{2,j+z-1}, w_{4,j+z+1}, w_{o,j+z+2}) = 0$
  - $\phi_4(w_{1,j+z-2}, w_{2,j+z}, w_{1,j+z-1}, w_{2,j+z-1}, w_{4,j+z+1}, w_{o,j+z+2}) = 0$

Variable-base scalar multiplication circuit :

	$w_1$	$w_2$	$w_3$	$w_4$	$w_o$
$j + 0$	$b_{n-1}$	$x_2$	$y_2$	$b_{n-2}$	acc
$j + 1$	$x_3$	$y_3$	$x_4$	$b_{n-3}$	acc
$j + 2$	$x_1$	$y_1$	$y_4$	$b_{n-4}$	acc
$j + 3$	$x_5$	$y_5$	$x_6$	$b_{n-5}$	acc
$j + 4$	$y_6$	$x_7$	$y_7$	$b_{n-6}$	acc
...					
$j + 210$	...	$x_{n-3}$	$y_{n-3}$	$b_2$	b
$j + 211$	$x_{n-2}$	$y_{n-2}$	$b_1$	$b_0$	$x_{n-1}$
$j + 212$	$x_1$	$y_1$	$y_{n-1}$	$x_n$	$y_n$

Define the following functions:

1.  $\phi_1 : (b, x_1, y_1, x_2, y_2, x_3) \mapsto$   
 $x_3 \cdot ((y_1^2 - x_1^2) \cdot (2 - y_1^2 + x_1^2) + 2dx_1y_1(y_1^2 + x_1^2) \cdot x_2y_2b) - (2x_1y_1 \cdot (2 - y_1^2 + x_1^2) \cdot (y_2b + (1 - b)) +$   
 $(y_1^2 + x_1^2) \cdot (y_1^2 - x_1^2) \cdot x_2b)$
2.  $\phi_2 : (b, x_1, y_1, x_2, y_2, y_3) \mapsto$   
 $y_3 \cdot ((y_1^2 - x_1^2) \cdot (2 - y_1^2 + x_1^2) - 2dx_1y_1(y_1^2 + x_1^2) \cdot x_2y_2b) - (2x_1y_1 \cdot (2 - y_1^2 + x_1^2) \cdot x_2b + (y_1^2 + x_1^2) \cdot$   
 $(y_1^2 - x_1^2) \cdot (y_2b + (1 - b)))$

Constraints:

- For  $j + 0$ :
  - $w_{o,j} = w_{1,j} \cdot 2 + w_{4,j}$
  - $\phi_1(w_{1,j+0}, w_{1,j+2}, w_{2,j+2}, w_{1,j+2}, w_{2,j+2}, w_{2,j+0})$
  - $\phi_2(w_{1,j+0}, w_{1,j+2}, w_{2,j+2}, w_{1,j+2}, w_{2,j+2}, w_{3,j+0})$
- For  $j + z, z \equiv 0 \pmod{5}, z \neq 0$ :
  - $w_{o,j+z} = w_{1,j+z} \cdot 2 + w_{4,j+z} + w_{o,j+z-1}$
  - $\phi_1(w_{4,j+z}, w_{2,j+z-1}, w_{3,j+z-1}, w_{1,j+z+2}, w_{2,j+z+2}, w_{2,j+z})$
  - $\phi_2(w_{4,j+z}, w_{2,j+z-1}, w_{3,j+z-1}, w_{1,j+z+2}, w_{2,j+z+2}, w_{3,j+z})$
- For  $j + z, z \equiv 1 \pmod{5}$ :
  - $w_{o,j+z} = 2 \cdot w_{o,j+z-1} + w_{4,j+z}$
  - $\phi_1(w_{4,j+z-1}, w_{2,j+z-1}, w_{3,j+z-1}, w_{1,j+z+1}, w_{2,j+z+1}, w_{1,j+z})$
  - $\phi_2(w_{4,j+z-1}, w_{2,j+z-1}, w_{3,j+z-1}, w_{1,j+z+1}, w_{2,j+z+1}, w_{2,j+z})$
  - $\phi_1(w_{4,j+z}, w_{1,j+z}, w_{2,j+z}, w_{1,j+z+1}, w_{2,j+z+1}, w_{3,j+z})$
- For  $j + z, z \equiv 2 \pmod{5}$ :
  - $w_{o,j+z} = 2 \cdot w_{o,j+z-1} + w_{4,j+z}$
  - $\phi_2(w_{4,j+z-1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{1,j+z}, w_{2,j+z}, w_{3,j+z})$
- For  $j + z, z \equiv 3 \pmod{5}$ :
  - $w_{o,j+z} = 2 \cdot w_{o,j+z-1} + w_{4,j+z}$
  - $w_{o,j+z} = 2 \cdot w_{o,j+z-1} + w_{4,j+z}$
  - $\phi_1(w_{4,j+z-1}, w_{3,j+z-2}, w_{3,j+z-1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{1,j+z})$
  - $\phi_2(w_{4,j+z-1}, w_{3,j+z-2}, w_{3,j+z-1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{2,j+z})$
  - $\phi_1(w_{4,j+z}, w_{1,j+z}, w_{2,j+z}, w_{1,j+z-1}, w_{2,j+z-1}, w_{3,j+z})$
- For  $j + z, z \equiv 4 \pmod{5}$ :
  - $w_{o,j+z} = 2 \cdot w_{o,j+z-1} + w_{4,j+z}$
  - $\phi_2(w_{4,j+z-1}, w_{1,j+z-1}, w_{2,j+z-1}, w_{1,j+z-2}, w_{2,j+z-2}, w_{1,j+z})$
  - $\phi_1(w_{4,j+z}, w_{3,j+z-1}, w_{1,j+z}, w_{1,j+z-2}, w_{2,j+z-2}, w_{2,j+z})$
  - $\phi_2(w_{4,j+z}, w_{3,j+z-1}, w_{1,j+z}, w_{1,j+z-2}, w_{2,j+z-2}, w_{3,j+z})$

### 2.5.8 Validator Set Proof Circuit

WIP

## Chapter 3

# In-EVM State Proof Verifier

This introduces a description for Solana's 'Light-Client' state proof in-EVM verifier. Crucial components which define this part design are:

1. Verification architecture description.
2. Verification logic API reference.
3. Input data structures description.

### 3.1 Verification Logic Architecture

### 3.2 Verification Logic API Reference

### 3.3 Input Data Structures

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