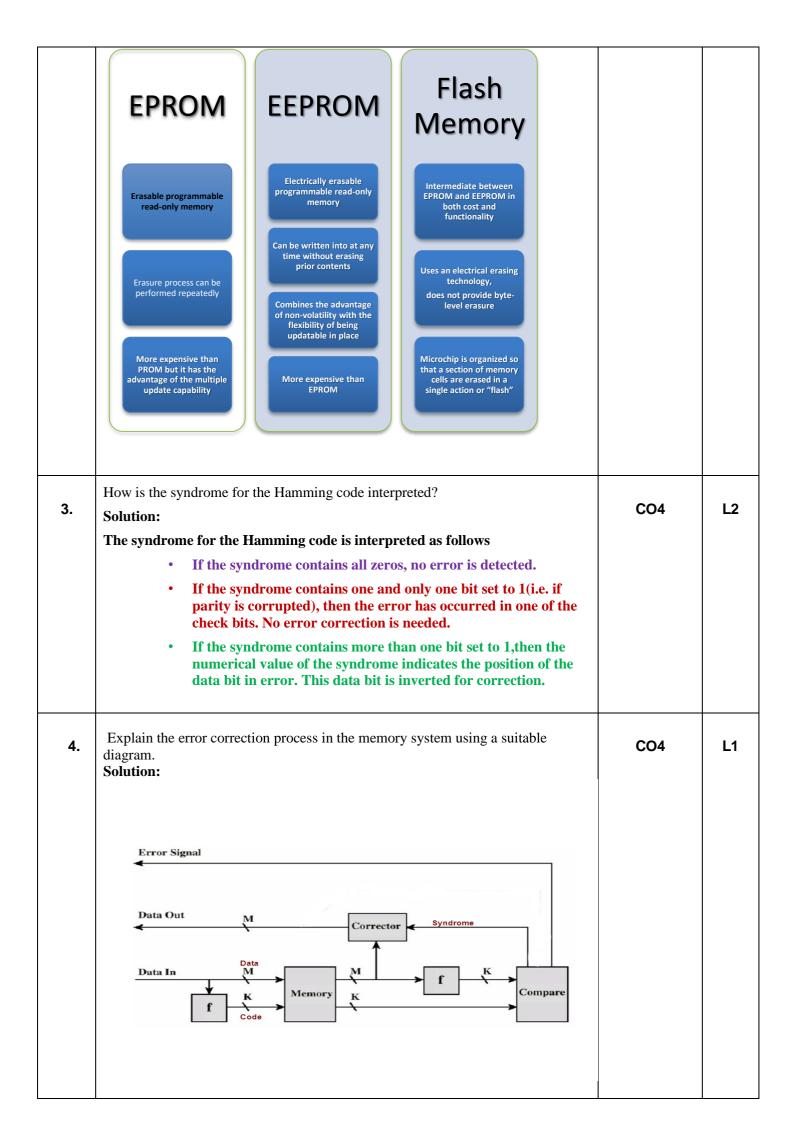
O. A. O. O. A. O. A. O. O. A. O. O. A. O. A. O. O. A. O. O. A. O. O. A.	ANUS AND TAN	ITER, SIKSHA 'C	O' ANUSANDHAN University)	l (Deeme	ed to	be	Assig	gnment
Branch		Computer Science and E	ingineering	Program	me		В.	Tech
Course	Name	Computer Organization a	Computer Organization and Architecture Semester					
Course Code		EET 2211		Academi	c Year		3/Even	
Solutio Assign	n of ment-4	Topic-Cache memory a				GP-1		
	g Level	L1: Remembering	L3: Applying	•	<b>L5</b> : E	valuating		
(LL)	J	L2: Understanding	L4: Analyzing		<b>L6</b> : C	reating		
Q's		Qu	estion S			COs		LL
1	What is Solution	the difference between DRAM n:	and SRAM in terms of ap	pplication?		CO4	4	L1
	DRAM							
	_	esence or absence of charge pacitor is interpreted as a 1 or 0	sed					
	Essenti	ally an analog device	Basically, digital device					
	Cells an	e smaller and simpler	Cells are not simpler					
	More d		Less Dense  More expensive					
	Less ex	pensive						
	Used fo	Used for main memory  Used for cache memory (both on and off-chip)						
	circuitr	es the supporting refresh ry which Tend to be ed for large memory	Refresh circuitry is not					
2	Explain different types of ROM. Solution:						4	L1



	<ul> <li>Prior to storing data, a code is generated from the bits in the word.</li> <li>Code is stored along with the data word in memory.</li> <li>Code used to identify and correct errors.</li> <li>When the word is fetched a new code is generated and compared to the stored code. <ul> <li>No errors detected.</li> <li>An error is detected and it is possible to correct the error.</li> <li>An error is detected, but it is not possible to correct it.</li> </ul> </li> </ul>		
5.	If there are m input lines n output lines for a decoder that is used to uniquely address a byte addressable 1 KB RAM, then the minimum value of m+n is		
	(A) 18 (B) 1034 (C) 10 (D) 1024 [GATE 2020]	CO4	L2
	Solution: A Decoder is represented as $m \times n(2^m)$ used to identify a byte address of 1KB RAM. 1KB memory is represented as, 1 KB = 1K $\times$ B = $2^{10} \times$ 8 = $2^m \times$ 8. So $m = 10$ $n = 2^m = 2^{10} = 1024$ so minimum value of $m + n = 10 + 1024 = 1034$ , option (B)		
6.	Suppose an 8-bit data word stored in memory is 11000010. Using the Hamming algorithm, determine what check bits would be stored in memory with the data word. Show how you got your answer.	CO4	L3
	Ans: 1. Length of data word =m=8-bits According to inequality condition 2 <sup>k</sup> -1 ≥ m+k, no of check bits required is 4 So, k=4-bit (C8, C4, C2, C1) Length of codeword=m+k=12-bit  2. Arrange the data bits and check bits in the cord word  Cord word 12 11 10 9 8 7 6 5 4 3 2 1		
	Position		
	Check bits         C8         C4         C2         C1           Data bits         D8         D7         D6         D5         D         D3         D2         D1         D1		
	1 1 0 0 C8 0 0 1 C4 0 C2 C1		

Cord word	12	1.1	10	0	0	7		- E	1	2	1	1	1		
word	12	11	10	9	8	7	6	5	4	3	2	1			
Position															
Check					С				C4		С	С			
bits					8						2	1			
Data bits	D8	<b>D7</b>	<b>D6</b>	<b>D5</b>		D4	D3	D2		D1					
	0	0	1	1	С	1	0	0	<b>C4</b>	1	С	С	-		
					8						2	1			
$K_{\text{old}} = 0111$ $K_{\text{new}} = 1101$ So, Syndrome Since, Syndrome decimal equiverror.  Correct data voletine and com	ome ≠ valent of word is	0, so end of the sign of the s	rror h yndro	as bee	en de e. 10	. The	10 <sup>th</sup> p	ositio	on of the	he cod			in	CO4	
RAID1: Mirro RAID 2: Redu	ındant	•							_				ted		
the corresponding the corresponding the corresponding the corresponding to the corresponding the corre	onding interle	bit po aved p	sitio arity	ns on ; sim	mul ilar t	tiple o lev	nd the parity el 2 l	e bits y disl out in	ks. istead	of an	erro	stoi or-	ed		
n the corresponding the correcting code	onding interle e, a sin	bit po aved p mple p	sition arity arity	ns on ; sim bit is	mul ilar t s con	tiple o lev npute	nd the parity el 2 l	e bits y disl out in	ks. istead	of an	erro	stoi or-	ed		
n the corresponding the correcting code in the same position.	onding interle e, a sin ion on	bit po aved p mple p all of	osition parity parity the d	ns on ; sim bit is lata d	mulilar t s con lisks.	tiple o lev npute	nd the parity el 2 le ed for	e bits y disl out in the	ks. istead set of	of an indiv	erro idua	stor or- l bits	red s in		
n the corresponding the correcting code he same positions and the corresponding the	interle e, a sin ion on ek-inte strips	aved posterior above a bit posterior above a	osition parity parity the d d par ch da	ns on; sim; bit is lata dity; a	mulilar to see continuation in the continuatio	tiple o lev npute by-bi	nd the parity el 2 le ed for	e bits y disl out in the s	ks.  stead  set of  rip is	of an indivi	erro idua	stor or- l bits	red s in		
cross corresponding corresponding corresponding	interle e, a sin ion on ek-inte strips	aved posterior above a bit posterior above a	osition parity parity the d d par ch da	ns on; sim; bit is lata dity; a	mulilar to see continuation in the continuatio	tiple o lev npute by-bi	nd the parity el 2 le ed for	e bits y disl out in the s	ks.  stead  set of  rip is	of an indivi	erro idua	stor or- l bits	red s in		
RAID 3: Bit-in orrecting code he same positions are same positions or responding or responding or responding or responding to the same positions are same positions.	onding interle e, a sin ion on ck-inte strips strip (	g bit poor aved proper parall of crleave on each on the Data Disks	parity the dad parity character da parity	ns on; sim bit is lata d ity; a ta disty disl	mulilar to see combined to see	tiple o lev npute by-bi nd the	nd the parity el 2 le ed for it parite pari	e bits y disl out in the s	ks.  stead  set of  rip is	of an indivi	erro idua	stor or- l bits	red s in		
RAID 3: Bit-in correcting code in e same positions and positions are positionally as a same position or responding corresponding corresponding are positionally as a same position or responding corresponding are positionally as a same position or responding are positionally as a same position of the same position	onding interle e, a sin ion on ck-inte strips strip (DBk1 BB)	g bit poor aved proper proper all of crieave on each on the Data Disks	parity the dad parity character da parity Disk 3 BB BB BB Strippi	ns on ; sim bit is lata d ity; a ta dis y disl	mulilar to see contisks.  a bit- sk, and k.	tiple o lev npute by-bi nd the	nd the parity el 2 le de for el te parity	e bits y disl out in the s ity st ity bi	ks. astead set of rip is ts are	of an individual calculations	erro idua lated d in	estor or- l bits l acro	red in oss		
RAID 3: Bit-in orrecting code he same position or esponding or responding or responding or responding or responding RAID 5: Block	onding interle e, a sin ion on ek-inte strips strip (  D-4: Block-inter	g bit poor aved proper parall of crieave on each on the Data Disks B2 B3 B6 Ck-Level rleaved	parity the d parity ch da parity barity the da parity ch da parity barity barit	ns on; sim bit is lata d ity; a ta dis y disl	mulilar to see contisks.  a bit- sk, and k.	tiple o lev npute by-bi nd the	nd the parity el 2 le de for el te parity	e bits y disl out in the s ity st ity bi	ks. astead set of rip is ts are	of an individual calculations	erro idua lated d in	estor or- l bits l acro	red in oss		
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RAID 3: Bit-incorrecting code in earner position in the corresponding corresponding corresponding corresponding corresponding call 5: Block parity strip	onding interle e, a sin ion on ck-inter strips strip (DBk1 Block).  D-4: Block-inter strips ad B7 Block acro	g bit poor aved proper proper all of crieave on each on the Data Disks    Disk 2   B5   B6	parity the dad parity the dad parity barity the dad parity barity barity barity the dad parity barity barit	ns on ; sim bit is lata d ity; a ta dis y disl Party D Disk 4 Disk 3 Disk 4 Disk 3 Disk 3 Disk 4 Disk 3 Disk 3 Disk 4 Disk 3 Disk 4 Disk 3 Disk 4 Disk 3 Disk 4 Disk 4 Disk 3 Disk 4 Disk 4 Disk 4 Disk 3 Disk 4 Disk 4 Disk 4 Disk 4 Disk 4 Disk 5 Disk 4 Disk 5 Disk 4 Disk 5 Disk 6 Dis	mulilar to see contisks.  a bit- sk, and k.  bit- sk, and the Dediced pa	tiple o lev npute by-bi nd the	nd the parity el 2 le parity simil	e bits y disl but in the s ity st ity bi	ks. astead set of rip is ts are	of an individual of a store of the store of	erro	stor stor l bits acre the	red in oss		
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RAID 3: Bit-in orrecting code he same positions are positional to the same position or responding or responding or responding to the same parity strip in the parity strip	onding interle e, a sin ion on one k-inter strips strip of the strip of the strips strip of the s	g bit poor aved proper proper all of crieave on each on the Data Disks beck-Level crieave crieave ried out	parity the d d parity the d d parity the d d parity the d distance of the dist	ns on ; sim bit is lata d ity; a ta dis y disl  Party D  Disk 4  Disk 3  Disk 4  Disk 3  Disk 4  Disk 3  Disk 3  Disk 4  Disk 5  Disk 4  Disk 4  Disk 5  Disk 4  Disk 6  Disk 7  Disk 6  Disk 6  Disk 6  Disk 6  Disk 6  Disk 7  Disk 8  Disk 6  Disk 8  Disk	mulilar to see contisks.  a bit- sk, and k.  bit- sk and bit- sk a	tiple o lev npute by-bi nd the cated P arity;	nd the parity el 2 le Parity arity	e bits y disl but in the s ity st ity bi	ks. astead set of rip is ts are	of an individual of a store of the store of	erro	e stor	red in oss		
RAID 3: Bit-icorrecting code ne same positive and positiv	onding interle e, a sin ion on onek-inters strips strip (D-4: Block-inters acro	g bit poor aved property all of erleaved on the Data Disks  Ck-Level rleaved ss all of the cock-Level criteaved series are criteaved to the cock-Level criteaved to the criteaved to the criteaved to the cock-Level criteaved to the criteaved to th	parity the d d parity the d d parity the d d parity the d distance of the dist	ns on ; sim bit is lata d ity; a ta dis y disl parky D parky D pisk 4 point with ribute  D point with ribute  All dis storece k3	mulilar to see condisks.  The bis bit	tiple o lev npute by-bi nd the cated P arity;	nd the parity el 2 le Parity arity	e bits y disl but in the s ity st ity bi	ks. astead set of rip is ts are	of an individual of a store of the store of	erro	e stor	red in oss		

8.

9.	Consider a 4-drive, 200 GB-per-drive RAID array. What is the available data storage capacity for each of the RAID levels 0, 1, 3, 4, 5, and 6?  Solution: No of disks available =N=4 Capacity of each disk =200GB For RAID level 0, all the disks are used to store data. So no of data disk=N= 4 Data storage capacity = no of data disks * capacity of each disk  =4*200GB =800GB	CO4	L2
	For RAID level 1, no of data disk=N/2= 2  Data storage capacity = no of data disks * capacity of each disk = 2*200GB = 400GB  For RAID level 3, no of data disk=N-1= 3  Data storage capacity = no of data disks * capacity of each disk = 3*200GB = 600GB  For RAID level 4, no of data disk=N-1= 3  Data storage capacity = no of data disks * capacity of each disk = 3*200GB = 600GB  For RAID level 5, no of data disk=N-1= 3  Data storage capacity = no of data disks * capacity of each disk = 3*200GB = 600GB  For RAID level 4, no of data disk=N-2= 2  Data storage capacity = no of data disks * capacity of each disk = 2*200GB = 400GB		
10	What are the major functions of an I/O module?  Solution:  The major functions of an I/O module are  Control and timing Coordinates the flow of traffic between internal resources and external devices  Processor communication Involves command decoding, data, status reporting, address recognition  Device communication Involves commands, status information, and data  Data buffering Performs the needed buffering operation to balance device and memory speeds  Error detection Detects and reports transmission errors	CO1	L1

11.	List three broad classifications	of external, or peri	pheral, device	es		CO1	L1
	Solution:						
	External device categories						
	<ul><li>Human readable: commun</li><li>Machine-readable: comm</li></ul>		•		a driva		
	Communication: commun			_			
	readable or machine-readab		devices may	oc namar			
12.	Explain the different modes of 8255A.	operation of				CO1	L1
	Solution:						
	Mode 0: This is the basic I/O m						
	designated as input or output podefined as output, and data may	• •		•			
	defined as output, and data may	only be read from	i a port ii uie	port is set	io input.		
	■ Mode 1: In this mode, ports and lines from port C serve as	control lines for	<b>A and B.</b> The	control sig	gnals serve		
	two principal purposes: "handsl simple timing mechanism. The						
	indicate when the data are prese	ent on the I/O data	lines, ACKN	OWLED	GE,		
	indicating that the data have been INTERRUPT REQUEST line						
	■ Mode 2: This is a bidirectional either the input or output line for Port A only.						
13.	Suppose that the 8255A is confand all the bits of port C as out this configuration.	•	CO1	L2			
	Solution:						
	Control word Format of 825						
	D7 D6 D5 1 0 0	D4 D3	B D2 0	D1 0	D0 0		
		1 0					
	The 8255 Programmable Per	d below					
	(i) An A/D converter is into						
	The conversion is initiated by	nal on Port					
	C causes data to be strobe into (ii) Two computers exchange	vorks as a					
14.	bidirectional data port supporte	VOIRS as a	CO1	L2			
	The appropriate modes of opera						
	(A) Mode 0 for (i) and Mode (B) Mode 1 for (i) and Mode	* *					
	(C) Mode for (i) and Mode (	for (ii)					
	( <b>D</b> ) Mode 2 for (i) and Mode	e 1 for (ii)		[GA	TE 2004]		
	Solution:						
	Option (B)						

Explain the use of 82C59A interrupt handler to handle up to 64- I/O request using a suitable diagram.	CO1	L1
Solution: Figure 7.8 of text book.		
	suitable diagram. Solution:	suitable diagram.  Solution: