Homework 3 Fundamental Algorithms, Fall 2020, Professor Yap

Due: Wed Oct 21, by 11pm in GradeScope.

INSTRUCTIONS:

- Remember that we have a "no late homework" policy. Special permission must be obtained in advance if you have a valid reason.
- Only hand in answers to the questions with positive points attached to them. Do the others at your leisure.
- 1. (6 Points) Recall that a rotation can be implemented with 6 pointer assignments. Suppose a BST maintains successor and predecessor links (denoted u.succ and u.pred in the ¶III.22, p.23). How many of pointer assignments does a rotation need in such a BST? First specify an integer, then justify your answer. (But be careful!)
- 2. (12 Points) Please do Exercise III.6.9 (p. 52) on AVL height.
- 3. (12 Points)
 - (a) Consider a linear recurrence of the form

$$T(n) = C + \sum_{i=1}^{k} a_i T(n-i)$$

where C, a_1, \ldots, a_k are constants. We say the recurrence is **homogeneous** if C = 0, else **non-homogeneous**. If T(n) is non-homogeneous, show how you can transform it into a homogeneous linear recurrence,

$$t(n) = \sum_{i=1}^{k} a_i t(n-i).$$

What is the connection between T(n) and t(n)?

(b) Recall that the min-size AVL tree of height h satisfies the non-homogeneous recurrence

$$\mu(h) = \begin{cases} 1+h & \text{if } h = 0, 1, \\ 1+\mu(h-1) + \mu(h-2) & \text{if } h \ge 2. \end{cases}$$

Give an exact solution for $\mu(h)$ of the form

$$\mu(h) = A\phi^h + B\widehat{\phi}^h$$

for all $h \ge 0$. You must determine the constants A, B, using the initial conditions for $\mu(0)$ and $\mu(1)$. Recall that $\phi = \frac{1+\sqrt{5}}{2} = 1.6180\ldots$ and $\widehat{\phi} = 1 - \phi = -0.6180\ldots$

4. (20 Points) (This is a essentially Ex.III.6.10, p. 52.)

We know that a single insertion can always be fixed by a single rebalancing act. By a **rebalancing** act, we mean either a single rotation or a double rotation as performed during our Rebalance Phase. But how many rebalancing acts can be caused by a deletion?

Let m(k) be the minimum size AVL trees such that a single deletion will cause k rebalancing acts.

- (a) Determine m(1), m(2), m(3). You must draw the AVL tree for each m(k), and mark a node whose deletion will cause k rebalancing acts.
- (b) Prove that $m(k) = \mu(2k)$.
- (c) Can you design the AVL tree for m(k) in part(b) to be any combination of k single (S) or double
- (D) rotations? E.g., k = 2 has four combinations: SS, SD, DS, DD.
- (d) In an AVL tree of size n, what is the most number of rebalancing acts you can get after a deletion. Note that $O(\log n)$ is a trivial answer, so we really what a non-asymptotic bound.

- 5. (16 Points) (a) Give a non-recursive routine called isBST(u) which returns true if the node u represents (the root of) a BST; otherwise it returns false. Although isBST(u) is non-recursive, it calls a recursive routine which we denote by R(u). HINT: The subroutine R(u) returns a pair of values.
 - (b) Please specify clearly what is returned by R(u), and give a brief proof that your routine for R(u) is correct.

Notes. Assume that a node u is either nil, or else it has three fields: u.key, u.left, u.right. The keys k in a BST are, by definition finite, i.e., $-\infty < k < +\infty$.

- 6. (16 Points) The ratio bound ρ in the ratio balanced class $RB[\rho]$ is normally restricted to $(0, \frac{1}{2})$. Let us explore why we don't extend to $\rho \in [\frac{1}{2}, 1)$.
 - (a) What is wrong with the class $RB\left[\frac{2}{3}\right]$?
 - (b) Which other values of $\rho \in [\frac{1}{2}, 1)$ pose problem similar to part(a)?
 - (c) Suppose we really want to allow values of ρ that lie in the range $(\frac{1}{2}, 1)$. How can we overcome the above objections?
- 7. (24 Points) Relaxed AVL Trees

Let us define **AVL[2]** balance condition to mean that at each node u in the binary tree, $|balance(u)| \le 2$.

- (a) Draw the smallest size AVL[2] trees of heights h = 0, 1, 2, 3, 4, 5.
- (b) Derive an upper bound on the height of a AVL[2] tree on n nodes of the form. You may imitate the arguments in our Lectures Notes for the usual AVL trees. Give the best bound you can.
- (c) Give an insertion algorithm that preserves AVL[2] trees. Try to follow the original AVL insertion as much as possible; but point out differences from the original insertion.
- (d) Give the deletion algorithm for AVL[2] trees.
- 8. (0 Points) This is for your own practice (don't submit). Insert into an initially empty AVL tree the following sequence of keys: 1, 2, 3, ..., 14, 15.
 - (a) Draw the AVL after inserting the key 15.
 - (b) Prove the following: if we continue in this manner, we will have a complete binary tree at the end of inserting key $2^n 1$ for all $n \ge 1$.