

Object

1- Yes

2- No

3- No

1 - $g(n) = n^3$

$$f(n) = 3n^3 + n^2 + n$$

$$f(n) \leq C_1 \cdot g(n)$$

$$3n^3 + n^2 + n \leq n^3 \cdot C_1$$

به الصدا كثر المهمه
 $C_1 = 5$

$$f(n) \geq C_2 \cdot g(n)$$

$$3n^3 + n^2 + n \geq C_2 \cdot n^3$$

$$3n^3 + n^2 + n \geq 2n^3 \cdot C_2$$

$$C_2 = 2$$

$$5n^3 \geq 3n^3 + n^2 + n \geq 2n^3$$

$$-2 \geq 0 \geq 0$$

$$\Theta(n^3) = f(n)$$

$$\therefore \Theta(g(n)) = f(n)$$

2- $g(n) = 2^n$, $f(n) = 2^{n+1}$

$$f(n) \leq C_1 \cdot g(n)$$

$$2^{n+1} = 2^n \times 2^1 \leq 2^n \cdot C_1$$

$$\therefore C_1 = 2$$

$$f(n) \geq C_2 \cdot g(n)$$

$$2^n \times 2 \geq 2^n \cdot C_2$$

$$C_2 = 2$$

$$-2 \geq 0 \geq 0$$

$$2 \cdot 2^n \geq 2^{n+1} \geq 2 \cdot 2^n$$

$$\therefore \Theta(2^n) = f(n)$$

Object

Q₃

1- $f(n) = n^3$, $g(n) = n^2$

$\Omega \rightarrow f(n) \geq C_1 \cdot g(n)$

$$n^3 \geq C_1 \cdot n^2$$

$$C_1 \leq n$$

$\therefore n^3$ grows faster than n^2

$$f(n) = \Omega(n^2)$$

$$n^3 = \Omega(n^2)$$

$O \rightarrow f(n) \leq C_2 \cdot g(n)$

$$n^3 \leq C_2 \cdot g(n)$$

$$n^3 \leq C_2 \cdot n^2 \quad \times$$

2- $f(n) = \log(n)$, $g(n) = \log^2(n)$

$\Omega \rightarrow f(n) \geq C_1 \cdot g(n)$

$\log(n) \geq C_1 \cdot \log^2(n)$

$$\log(n) \geq C_1 \frac{\log(n)}{\log(2)}$$

$$C_1 \leq \log(2)$$

$$\therefore \Omega(g(n)) = f(n)$$

$$\therefore \log(n) = \Omega(\log^2(n))$$

$O \rightarrow f(n) \leq C_2 \cdot g(n)$

$$\log(n) \leq C_2 \frac{\log(n)}{\log(2)}$$

$$\log(2) \leq C_2$$

$$\therefore O(g(n)) = f(n)$$

$$\log(n) = O(\log^2(n))$$

\times



Object

Q4.

$$\begin{array}{ccc} \text{best case} & & \text{worst case} \\ \square n & \leftarrow F(n) \leftarrow & \square n \\ \sim n(n) & & O(n) \\ C_1 = & & C_2 = \\ n_0 = & & n_0 = \end{array}$$

حيث

حيث

best case

worst case

$$F(n) \geq C_1 \cdot g(n)$$

$$F(n) \leq C_2 \cdot g(n)$$

$$n \geq n_0$$

$$n \geq n_0$$

Q5.

$$O(g(n))$$

$$F(n) \leq C_1 \cdot g(n)$$

$$, n \geq n_0$$

$$W(g(n))$$

$$F(n) \geq C_2 \cdot g(n)$$

$$, n \geq n_0$$

$$\therefore C_1 \cdot g(n) \geq F(n) \geq C_2 \cdot g(n)$$

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سريعة

بطيئة

$$= O(g(n)) \cap W(g(n)) = \emptyset$$