

## **Lab 12 – Algorithms with AI Assistance**

AI Assisted Coding

Assignment Number: 12.4

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### **Task 1: Bubble Sort for Ranking Exam Scores**

#### **Question**

You are working on a college result processing system where a small list of student exam scores must be sorted after each internal assessment.

Implement Bubble Sort in Python to sort a list of student scores.  
Explain comparisons, swaps, early termination, and time complexity.

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#### **Python Code**

```
def bubble_sort(scores):
    n = len(scores)
    for i in range(n):
        swapped = False
        for j in range(0, n - i - 1):
            if scores[j] > scores[j + 1]:
                scores[j], scores[j + 1] = scores[j + 1], scores[j]
                swapped = True
        if not swapped:
            break
    return scores
```

# Testing

```
marks = [45, 78, 62, 90, 55]
print(bubble_sort(marks))
```

## Output

[45, 55, 62, 78, 90]

## Code Explanation

- Bubble Sort repeatedly compares adjacent elements and swaps them if they are in the wrong order.
- The swapped flag allows early termination when the list becomes sorted.
- Best Case:  $O(n)$  when the list is already sorted.
- Average Case:  $O(n^2)$ .
- Worst Case:  $O(n^2)$  when the list is reverse sorted.
- Bubble Sort is suitable only for small datasets.

## Task 2: Improving Sorting for Nearly Sorted Attendance Records

### Question

Student roll numbers are almost sorted with a few late updates.

Implement Bubble Sort and Insertion Sort, then explain why Insertion Sort is better for nearly sorted data.

### Python Code (Bubble Sort)

```
def bubble_sort(data):  
    for i in range(len(data)):  
        for j in range(0, len(data) - i - 1):  
            if data[j] > data[j + 1]:  
                data[j], data[j + 1] = data[j + 1], data[j]  
  
    return data
```

---

## Output

Bubble Sort executed

## Code Explanation

- Bubble Sort always performs comparisons even if the list is nearly sorted.

- This leads to unnecessary iterations.

### **Python Code (Insertion Sort)**

```
def insertion_sort(data):
    for i in range(1, len(data)):
        key = data[i]
        j = i - 1
        while j >= 0 and data[j] > key:
            data[j + 1] = data[j]
            j -= 1
        data[j + 1] = key
    return data
```

# Testing

```
roll_numbers = [1, 2, 3, 5, 4, 6]
print(insertion_sort(roll_numbers))
```

### **Output**

[1, 2, 3, 4, 5, 6]

### **Code Explanation**

- Insertion Sort shifts only misplaced elements.
- It performs efficiently on nearly sorted data.
- Time Complexity:
  - Best Case:  $O(n)$
  - Worst Case:  $O(n^2)$
- Insertion Sort is preferred for small or partially sorted datasets.

### **Task 3: Searching Student Records in a Database**

#### **Question**

Implement Linear Search for unsorted data and Binary Search for sorted data.  
Explain use cases and performance differences.

### Python Code (Linear Search)

```
def linear_search(data, target):
```

```
    """
```

Searches for a target value in an unsorted list.

Parameters:

data (list): List of student roll numbers

target (int): Roll number to search

Returns:

int: Index if found, else -1

```
    """
```

```
for i in range(len(data)):
```

```
    if data[i] == target:
```

```
        return i
```

```
return -1
```

```
# Testing
```

```
students = [104, 101, 109, 102]
```

```
print(linear_search(students, 109))
```

### Output

2

### Code Explanation

- Linear Search checks each element sequentially.
- Works on both sorted and unsorted lists.

- Time Complexity: O(n).

### Python Code (Binary Search)

```
def binary_search(data, target):
```

```
    """
```

Searches for a target value in a sorted list.

Parameters:

data (list): Sorted list of roll numbers

target (int): Roll number to search

Returns:

int: Index if found, else -1

```
    """
```

```
low = 0
```

```
high = len(data) - 1
```

```
while low <= high:
```

```
    mid = (low + high) // 2
```

```
    if data[mid] == target:
```

```
        return mid
```

```
    elif data[mid] < target:
```

```
        low = mid + 1
```

```
    else:
```

```
        high = mid - 1
```

```
return -1
```

```
# Testing
```

```
sorted_students = [101, 102, 104, 109]
```

```
print(binary_search(sorted_students, 109))
```

## **Output**

3

## **Code Explanation**

- Binary Search requires sorted data.
- It reduces the search space by half each iteration.
- Time Complexity:  $O(\log n)$ .
- Binary Search is much faster than Linear Search for large datasets.

## **Task 4: Choosing Between Quick Sort and Merge Sort**

### **Question**

Implement Quick Sort and Merge Sort using recursion.  
Explain recursion and compare their performance.

### **Python Code (Quick Sort)**

```
def quick_sort(arr):  
    """  
    Sorts a list using Quick Sort algorithm.  
    """  
  
    if len(arr) <= 1:  
        return arr  
  
    pivot = arr[len(arr) // 2]  
    left = [x for x in arr if x < pivot]  
    middle = [x for x in arr if x == pivot]  
    right = [x for x in arr if x > pivot]  
  
    return quick_sort(left) + middle + quick_sort(right)
```

```
# Testing  
data = [5, 3, 8, 6, 2]  
print(quick_sort(data))
```

## Output

```
[2, 3, 5, 6, 8]
```

## Code Explanation

- Quick Sort divides the list using a pivot.
- Recursion sorts left and right partitions.
- Average Time Complexity:  $O(n \log n)$ .
- Worst Case:  $O(n^2)$ .

## Python Code (Merge Sort)

```
def merge_sort(arr):  
    """  
    Sorts a list using Merge Sort algorithm.  
    """  
  
    if len(arr) <= 1:  
        return arr  
  
    mid = len(arr) // 2  
    left = merge_sort(arr[:mid])  
    right = merge_sort(arr[mid:])  
  
    return merge(left, right)
```

## def merge(left, right):

```
    result = []  
    i = j = 0
```

```
while i < len(left) and j < len(right):  
    if left[i] < right[j]:  
        result.append(left[i])  
        i += 1  
    else:  
        result.append(right[j])  
        j += 1  
  
    result.extend(left[i:])  
    result.extend(right[j:])  
return result
```

```
# Testing  
print(merge_sort(data))
```

## Output

```
[2, 3, 5, 6, 8]
```

## Code Explanation

- Merge Sort always divides data into halves.
- Recursion continues until single elements remain.
- Time Complexity is always  $O(n \log n)$ .
- Merge Sort is preferred when stable performance is required.

## Task 5: Optimizing a Duplicate Detection Algorithm

### Question

Detect duplicate user IDs using a naive approach and an optimized approach.

### Python Code (Brute Force)

```
def has_duplicates_bruteforce(data):
    for i in range(len(data)):
        for j in range(i + 1, len(data)):
            if data[i] == data[j]:
                return True
    return False

# Testing
ids = [101, 203, 405, 101]
print(has_duplicates_bruteforce(ids))
```

## Output

True

## Code Explanation

- Uses nested loops.
- Time Complexity:  $O(n^2)$ .
- Inefficient for large datasets.

## Python Code (Optimized Using Set)

```
def has_duplicates_optimized(data):
    seen = set()
    for item in data:
        if item in seen:
            return True
        seen.add(item)
    return False
```

# Testing

```
print(has_duplicates_optimized(ids))
```

## Output

True

## Code Explanation

- Uses a set for constant-time lookups.
- Time Complexity:  $O(n)$ .
- Space-Time tradeoff improves performance significantly.