

# ASSIGNMENT – 11.2

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Batch - 03

Task 1:-

```
assignment11.py > Stack > __init__
19
20 class Stack:
21     """
22     A Stack data structure that follows Last-In-First-Out (LIFO) principle.
23
24     This implementation uses a Python list as the underlying data structure.
25     Supports basic stack operations: push, pop, peek, and empty check.
26
27     Attributes:
28     | items (list): Internal list to store stack elements
29     """
30
31     def __init__(self):
32         """Initialize an empty stack."""
33         self.items = []
34
35     def push(self, item):
36         """
37         Add an item to the top of the stack.
38
39         Args:
40         | item: The item to be added to the stack
41         """
42         self.items.append(item)
43
44     def pop(self):
45         """
46         Remove and return the top item from the stack.
47
48         Returns:
49         | The top item from the stack
50
51         Raises:
52         | IndexError: If the stack is empty
53         """
54         if self.is_empty():
55             raise IndexError("Pop from empty stack")
56         return self.items.pop()
57
58     def peek(self):
59         """
60         Return the top item from the stack without removing it.
61
62         Returns:
63         | The top item from the stack
64
65         Raises:
66         | IndexError: If the stack is empty
67         """
68         if self.is_empty():
69             raise IndexError("Peek from empty stack")
```

Output:-

```
Data Structures Implementation - Lab 11
=====

1. Stack Operations:
Stack after pushes: size = 3
Peek: 30
Pop: 30
Pop: 20
Is empty: False
```

## Task 1: Stack Implementation

The Stack was implemented to demonstrate the **Last-In-First-Out (LIFO)** principle, which is fundamental in computer science. Stacks are widely used in:

- Expression evaluation

- Undo/Redo operations
- Function call management (call stack)
- Backtracking algorithms

This implementation uses a Python list for simplicity and efficiency. Proper error handling (raising `IndexError`) ensures robustness. The object-oriented design improves modularity, readability, and reusability of the code.

## Task 2:

```
class Queue:
    """
    A Queue data structure that follows First-In-First-Out (FIFO) principle.

    This implementation uses a Python list as the underlying data structure.
    Supports basic queue operations: enqueue, dequeue, front access, and size.

    Attributes:
    | items (list): Internal list to store queue elements
    """

    def __init__(self):
        """Initialize an empty queue."""
        self.items = []

    def enqueue(self, item):
        """
        Add an item to the rear of the queue.

        Args:
        | item: The item to be added to the queue
        """
        self.items.append(item)

    def dequeue(self):
        """
        Remove and return the front item from the queue.

        Returns:
        | The front item from the queue

        Raises:
        | IndexError: If the queue is empty
        """
        if self.is_empty():
            raise IndexError("Dequeue from empty queue")
        return self.items.pop(0)
```

```

def front(self):
    """
    Return the front item from the queue without removing it.

    Returns:
        The front item from the queue

    Raises:
        IndexError: If the queue is empty
    """
    if self.is_empty():
        raise IndexError("Front from empty queue")
    return self.items[0]

def is_empty(self):
    """
    Check if the queue is empty.

    Returns:
        bool: True if queue is empty, False otherwise
    """
    return len(self.items) == 0

def size(self):
    """
    Return the number of items in the queue.

    Returns:
        int: Number of items in the queue
    """
    return len(self.items)

```

Output:-

```

2. Queue Operations:
Queue after enqueues: size = 3
Front: 10
Dequeue: 10
Dequeue: 20
Is empty: False

```

## Task 2: Queue Implementation

The Queue was implemented to demonstrate the **First-In-First-Out (FIFO)** principle. Queues are essential in:

- CPU scheduling
- Breadth-First Search (BFS)
- Printer task scheduling
- Real-time systems

Using a Python list makes the implementation easy to understand. The structure includes enqueue, dequeue, front, and size operations with proper boundary checks to prevent runtime errors.

### Task 3:-

```
class Node:
    """
    A node in a singly linked list.

    Attributes:
        data: The data stored in the node
        next: Reference to the next node in the list
    """

    def __init__(self, data):
        """
        Initialize a node with given data.

        Args:
            data: The data to store in the node
        """
        self.data = data
        self.next = None

class SinglyLinkedList:
    """
    A singly linked list data structure.

    Supports insertion at the beginning and traversal/display operations.

    Attributes:
        head: Reference to the first node in the list
    """

    def __init__(self):
        """Initialize an empty singly linked list."""
        self.head = None

    def insert_at_beginning(self, data):
        """
        Insert a new node with given data at the beginning of the list.

        Args:
            data: The data to insert
        """
        new_node = Node(data)
        new_node.next = self.head
        self.head = new_node
```

```
def display(self):
    """
    Display all elements in the linked list.

    Prints the elements separated by ' -> ' and ends with 'None'
    """
    current = self.head
    while current:
        print(current.data, end=" -> ")
        current = current.next
    print("None")
```

Output:-

```
Is empty: False  
  
3. Singly Linked List Operations:  
Linked List after insertions:  
10 -> 20 -> 30 -> None
```

## Task 3: Singly Linked List Implementation

The Singly Linked List demonstrates dynamic memory allocation and non-contiguous data storage. Unlike arrays, linked lists allow efficient insertion and deletion operations.

This implementation helps in understanding:

- Pointer/reference-based data structures
- Dynamic memory usage
- Node-based traversal

It provides insertion at the beginning and traversal functionality, forming the foundation for more complex structures like doubly linked lists and circular linked lists.

Task 4:

```

class BSTNode:
    """
    Attributes:
        data: The data stored in the node
        left: Reference to the left child node
        right: Reference to the right child node
    """

    def __init__(self, data):
        """
        Initialize a BST node with given data.

        Args:
            data: The data to store in the node (must be comparable)
        """
        self.data = data
        self.left = None
        self.right = None

class BinarySearchTree:
    """
    A Binary Search Tree data structure.

    Supports insertion and in-order traversal operations.

    Attributes:
        root: Reference to the root node of the tree
    """

    def __init__(self):
        """Initialize an empty binary search tree."""
        self.root = None

    def insert(self, data):
        """
        Insert a new node with given data into the BST.

        Args:
            data: The data to insert (must be comparable)
        """
        if self.root is None:

```

```

            self.root = BSTNode(data)
        else:
            if data < node.data:
                if node.left is None:
                    node.left = BSTNode(data)
                else:
                    self._insert_recursive(node.left, data)
            else:
                if node.right is None:
                    node.right = BSTNode(data)
                else:
                    self._insert_recursive(node.right, data)

    def in_order_traversal(self):
        """
        Perform in-order traversal of the BST and print the elements.

        In-order traversal visits left subtree, current node, right subtree.
        """
        self._in_order_recursive(self.root)
        print() # New line after traversal

    def _in_order_recursive(self, node):
        """
        Helper method for recursive in-order traversal.

        Args:
            node: Current node in the recursion
        """
        if node:
            self._in_order_recursive(node.left)
            print(node.data, end=" ")
            self._in_order_recursive(node.right)

```

Output:

```
4. Binary Search Tree Operations:  
BST after insertions - In-order traversal:  
20 30 40 50 60 70 80
```

## Task 4: Binary Search Tree (BST) Implementation

The Binary Search Tree was implemented to demonstrate hierarchical data storage and efficient searching.

BST provides:

- Average-case search complexity:  **$O(\log n)$**
- Ordered data storage
- Efficient insertion and traversal

In-order traversal prints elements in sorted order, demonstrating the correctness of BST properties. Recursive implementation improves clarity and reflects real-world tree algorithms.

## Task 5:-

```
class HashTable:
    """
    A Hash Table data structure with collision handling using chaining.

    Uses separate chaining (linked lists) to handle collisions.
    Supports insert, search, and delete operations.

    Attributes:
        size (int): The size of the hash table
        table (list): Internal list of buckets (each bucket is a list)
    """

    def __init__(self, size=10):
        """
        Initialize a hash table with given size.

        Args:
            size (int): The number of buckets in the hash table
        """
        self.size = size
        self.table = [[] for _ in range(size)]

    def _hash(self, key):
        """
        Compute the hash value for a given key.

        Args:
            key: The key to hash

        Returns:
            int: The hash value (bucket index)
        """
        return hash(key) % self.size

    def insert(self, key, value):
        """
        Insert a key-value pair into the hash table.

        If the key already exists, updates the value.

        Args:
            key: The key to insert
            value: The value to insert
        """
```



```

        # Check if key already exists
        for i, (k, v) in enumerate(bucket):
            if k == key:
                bucket[i] = (key, value) # Update value
                return

        # Key doesn't exist, add new entry
        bucket.append((key, value))

def search(self, key):
    """
    Search for a value by its key in the hash table.

    Args:
        key: The key to search for

    Returns:
        The value associated with the key, or None if not found
    """
    bucket_index = self._hash(key)
    bucket = self.table[bucket_index]

    for k, v in bucket:
        if k == key:
            return v

    return None

def delete(self, key):
    """
    Delete a key-value pair from the hash table.

    Args:
        key: The key to delete

    Returns:
        bool: True if key was found and deleted, False otherwise
    """
    bucket_index = self._hash(key)
    bucket = self.table[bucket_index]

    for i, (k, v) in enumerate(bucket):

```

Output:-

```
5. Hash Table Operations:
Hash Table after insertions:
Bucket 0: [('cherry', 30)]
Bucket 1: []
Bucket 2: [('apple', 10), ('banana', 20)]
Bucket 3: []
Bucket 4: []
Bucket 5: []
Bucket 6: [('date', 40)]
Bucket 7: []
Bucket 8: []
Bucket 9: []
Search 'apple': 10
Search 'grape': None
Delete 'banana': True
Hash Table after deletion:
Bucket 0: [('cherry', 30)]
Bucket 1: []
Bucket 2: [('apple', 10)]
Bucket 3: []
Bucket 4: []
Bucket 5: []
Bucket 6: [('date', 40)]
Bucket 7: []
Bucket 8: []
Bucket 9: []
PS C:\Users\hp\OneDrive\Desktop\ai>
```

## Task 5: Hash Table Implementation

The Hash Table demonstrates fast data retrieval using hashing.

Key advantages:

- Average-case time complexity:  $O(1)$  for insert, search, delete
- Efficient key-value mapping
- Collision handling using separate chaining

This implementation introduces hashing concepts and collision resolution techniques, which are fundamental in database indexing, caching systems, and symbol tables.

Extra question:

```
5
6 SIZE = 5
7 stack = []
8 top = -1
9
10 def push(value):
11     global top
12     if top == SIZE - 1:
13         print("Stack Overflow")
14     else:
15         stack.append(value)
16         top += 1
17         print(value, "inserted")
18
19 def pop():
20     global top
21     if top == -1:
22         print("Stack Underflow")
23     else:
24         removed = stack.pop()
25         print(removed, "removed")
26         top -= 1
27
28 # Testing the additional stack implementation
29 print("\n6. Additional Stack Implementation (Array-based):")
30
31 # Inserting elements
32 push('a')
33 push('b')
34 push('c')
35 push('d')
36 push('e')
37 push('f') # This should show Overflow
38
39 print()
40
41 # Removing elements
42 pop()
43 pop()
44 pop()
45 pop()
46 pop()
47 pop() # This should show Underflow
```

Output:

#### 6. Additional Stack Implementation (Array-based):

```
a inserted  
b inserted  
c inserted  
d inserted  
e inserted  
Stack Overflow  
  
e removed  
d removed  
c removed  
b removed  
a removed  
Stack Underflow
```

This version demonstrates stack implementation using a fixed-size array and global variables. It highlights:

- Stack overflow condition
- Stack underflow condition
- Manual top pointer management

This helps compare structured (OOP-based) and procedural approaches, improving conceptual clarity.