

# COL 351 Quiz 6B

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TOTAL POINTS

10 / 10

QUESTION 1

1 Q1 10 / 10

$n = k$  in  $G$

+ 0 pts No submission / Incorrect

✓ + 2 pts We first show that this problem is in NP.

The verifier takes as input a graph  $G$  and a subset  $S$  of size  $k$  vertices and checks if  $S$  is an independent set or not

✓ + 2 pts We now reduce the independent set problem to the LargeIndSet problem. Let  $(G, k)$  be an input to the independent set problem.

We map it to an input  $(G', k')$  of the LargeIndSet problem. Let  $n$  be the number of vertices in  $G$ . We obtain  $G'$  as follows: we add a set  $W$  of  $n$  new vertices to  $G$  and there are no edges incident with any vertex in  $W$ .

✓ + 1 pts The parameter  $k' = k + n$ .

✓ + 1 pts Note that if  $n'$  is the number of vertices in  $G'$ , then

$n' = 2n$  and so,  $k' \geq n'/2$

✓ + 2 pts Argue that  $G$  has an independent set of size  $k$  iff  $G'$  has a large independent set of size  $k'$ .

Suppose  $G$  has an independent  $S$  of size  $k$ . Then  $S \cup W$  is an independent set in  $G'$  of size  $n + k = k'$

✓ + 2 pts Conversely, suppose  $G'$  has an independent set  $S$  of size  $k'$ . Now at most  $n$  of the vertices in  $S$  can belong to  $W$ . The remaining vertices  $S \setminus W$  belong to  $G$  and form an independent set in  $G$ . Thus, we have an independent set of size at least  $k' -$

Quiz 6B (COL 351)

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Give precise arguments. You can use the fact that the following problems are NP-complete: 3-Satisfiability, Clique, Vertex Cover, Independent Set, Subset Sum.

Given an undirected graph  $G$  on  $n$  vertices, we say that a subset  $S$  of vertices in  $G$  is a large independent set if  $S$  is an independent set AND  $|S| \geq n/2$ . Prove that the following problem, called LargeIndSet, is NP-complete: given a graph  $G$  and a value  $k \geq n/2$ , does  $G$  have a large independent set of size at least  $k$ ? Recall that a subset  $S$  of vertices is said to be an independent set if there is no edge between any pair of vertices in  $S$ .

We will first show that LargeIndSet is in NP.

Given any subset  $S$  which is a solution / LargeIndSet, we ~~show~~ can check in polynomial time that  $|S| \geq n/2$  AND  $|S|$  is independent (check every pair of vertices in  $O(|S|^2)$  time)

We will now show that

IS  $\leq_p$  LargeIndSet.

$G, k$

$(G, k)$

Decision version.

$G', k'$

(any general instance of IS  $(G, k)$  may be converted to a specific instance of LargeIndSet, i.e. it is reducible to LargeIndSet).

If  $k \leq n/2$

Construct  $G'$  by adding  $n-2k$  disconnected vertices to  $G$ , and choose  $k' = n-k$ . independent  $n=|G|$

Else if  $k > n/2$ ,

Choose  $G' = G$  and  $k' = k$ .  $G' = G \cup S'$  where  $S'$  has  $n-2k$  vertices

Note that this conversion is done in polynomial time (adding  $n-2k$  vertices  $\rightarrow$  polynomial)



We will now show that corresponding to every solved instance of  $IS(G, k)$ ,  $\exists$  solved instance of  $Large\ IndSet(G', k')$  i.e.  $k$  vice-versa i.e.

$$IS(G, k) \rightarrow Large\ IndSet(G', k') \quad (1)$$

$$Large\ IndSet(G', k') \rightarrow IS(G, k) \quad (2)$$

(also note that if  $k \geq n/2$ ,  $Large\ IndSet(G', k')$  can end as  $|G'| \leq k$  if  $k < n/2$ , " as  $|G'| = \frac{2n-2k}{2} \leq k' = \frac{n-k}{2}$ )

① Now, if there exists a solution to  $IS(G, k)$

$\hookrightarrow$  if  $k \geq n/2 \rightarrow$  it exists in  $Large\ IndSet(G', k')$  as  $|G'| \leq k$   
if  $k < n/2 \rightarrow$   $|G'| = n$   $\frac{k'}{2}$

On adding  $n-2k$  disconnected vertices to  $G$ , we now have  $k + n-2k = n-k$  disconnected vertices in  $G'$  (disconnected = independent) (independent).

Since  $k' = n-k$ , we would be able to find a sol<sup>n</sup> for  $Large\ IndSet(G', k')$  as  $k' = n-k = |G'|$   
 $k' \geq \frac{|G'|}{2}$

② Now, if  $\exists$  a sol<sup>n</sup> to  $Large\ IndSet(G', k')$

$\hookrightarrow$  if  $k \geq n/2$  same argument as above ( $G'=G$  &  $k'=k$ )  
if  $k < n/2$  Corresponding to sol<sup>n</sup> for  $Large\ IndSet(G', k')$   
 $\exists$  a sol<sup>n</sup> to  $IS(G, k)$

Since we added only  $n-2k$  independent vertices

&  $k' = n-k$ ,  $\therefore$  there must exist  $k$  independent vertices in the original graph  $G$ , &  $k' \geq |G'|/2$

Corresponding to those  $k$  vertices, we have a sol<sup>n</sup> to  $IS(G, k)$ .

reduces to

$\therefore IS(G, k) \sim Large\ IndSet(G, k)$  reduces in polynomial time