

Example Breakdown: Lambda Calculus Reductions

Problem 1: Reducing $(\lambda x. x x) (\lambda y. y)$

Step 1: Understanding the Expression

- We have a function $\lambda x. x x$ applied to another function $\lambda y. y$.
- In lambda calculus, function application means substituting the argument into the function.

Step 2: Applying Beta Reduction

- The function $\lambda x. x x$ expects an input and applies it to itself.
- We substitute $(\lambda y. y)$ for x in $x x$:

$$(\lambda x. x x) (\lambda y. y) \rightarrow (\lambda y. y) (\lambda y. y)$$

Step 3: Evaluating the New Expression

- Now we have $(\lambda y. y) (\lambda y. y)$.
- The function $\lambda y. y$ is the identity function, meaning it returns whatever is given to it.
- Applying it to $\lambda y. y$ results in:

$$(\lambda y. y) \rightarrow \lambda y. y$$

- The expression reduces to $\lambda y. y$, which is the simplest form (normal form).

Key Concepts in This Reduction:

- **Beta Reduction:** Replacing a bound variable with its argument.
 - **Identity Function:** $\lambda y. y$ always returns its input.
 - **Normal Form:** When no further reductions can be applied.
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Problem 2: Reducing $(\lambda x. x (\lambda z. z)) (\lambda y. y y)$

Step 1: Understanding the Expression

- The function $\lambda x. x (\lambda z. z)$ is applied to $\lambda y. y y$.
- This means x will be replaced with $(\lambda y. y y)$ in the expression $x (\lambda z. z)$.

Step 2: Applying Beta Reduction

- Substitute $(\lambda y. y y)$ for x in $x (\lambda z. z)$:

$(\lambda x. x (\lambda z. z)) (\lambda y. y y) \rightarrow (\lambda y. y y) (\lambda z. z)$

Step 3: Evaluating the New Expression

- Now we have $(\lambda y. y y) (\lambda z. z)$.
- $\lambda y. y y$ applies itself to its argument, which is $\lambda z. z$.
- This results in:

$(\lambda z. z) (\lambda z. z)$

- The function $\lambda z. z$ is the identity function, so it just returns $\lambda z. z$.
- The final result is $\lambda z. z$, which is the normal form.

Key Concepts in This Reduction:

- **Function Substitution:** Replacing x with its argument.
- **Self-Application:** $\lambda y. y y$ applies itself to an argument.
- **Normal Form:** The simplest form after reduction.

General Notes for Understanding Lambda Calculus:

- **Alpha Conversion:** Renaming bound variables to avoid confusion.
- **Beta Reduction:** The main step of computation in lambda calculus.
- **Eta Reduction:** Simplifying functions when possible (e.g., $\lambda x. f x \rightarrow f$ if x does not appear in f).

This document should help you practice breaking down lambda expressions step by step!