

Compiling and Running of Program

Dr. Zheng Xiaojuan Professor

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Main Content of Chapter 6

§ 6. Semantic Analysis

- 6.1 Overview of Semantic Analysis
- 6.2 Symbol Table
- 6.3 Semantic Analysis of Types
- 6.4 Semantic Analysis of Declaration
- 6.5 Semantic Analysis of Body
- 6.6 Attribute Grammar and Action Grammar



Semantic Analysis (review)

• What?

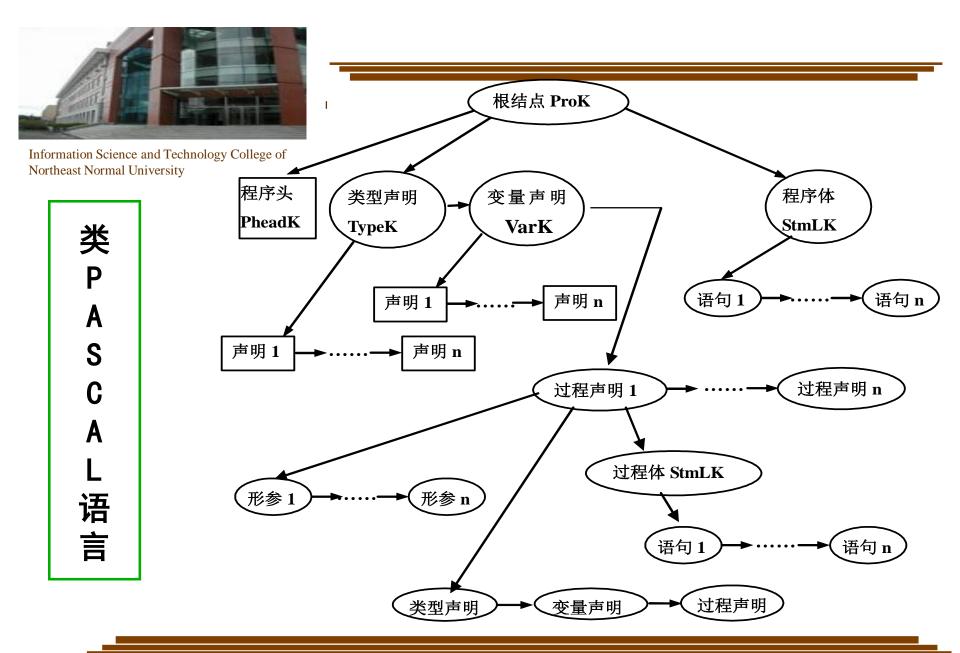
- Establish symbol table;
- Check semantic errors;

Where?

- Declaration --- establish symbol table; check "repeat declaration" error;
- Body -- search symbol table for attributes, check "use without declaration" error and type-related errors;

How?

- Attribute Grammar
- Action Grammar
- Parse Tree based semantic analysis;





6.3 Semantic Analysis of Types

- Some Concepts
 - 类型分析
 - _ 类型等价性
 - 类型相容性
- What?
- Where?
- How?



6.3 Semantic Analysis of Types

• 类型分析:

分解出类型表达式内部的各种信息。构造类型的内部 表示。这些信息要直接从类型表达式文本结构直接得 到是不容易的。

定义性出现时,把标识符与某个类型属性相关联,使用性出现时,取到与标识符相关联的类型属性。



6.3 Semantic Analysis of Types

• 类型等价性

- 按名等价(比较指针值)
- 按结构等价(复杂:需要定义一个 判定两个类型是否等价的子程序)

• 类型相容性

- 具体的编译器定义不同;
- 常见的三种相容性
 - 操作分量相容
 - 赋值相容性
 - 形参和值参相容性

```
typedef int T1[10];
typedef int T2[10];
```

```
T1 a, b;
T2 c, d;
```



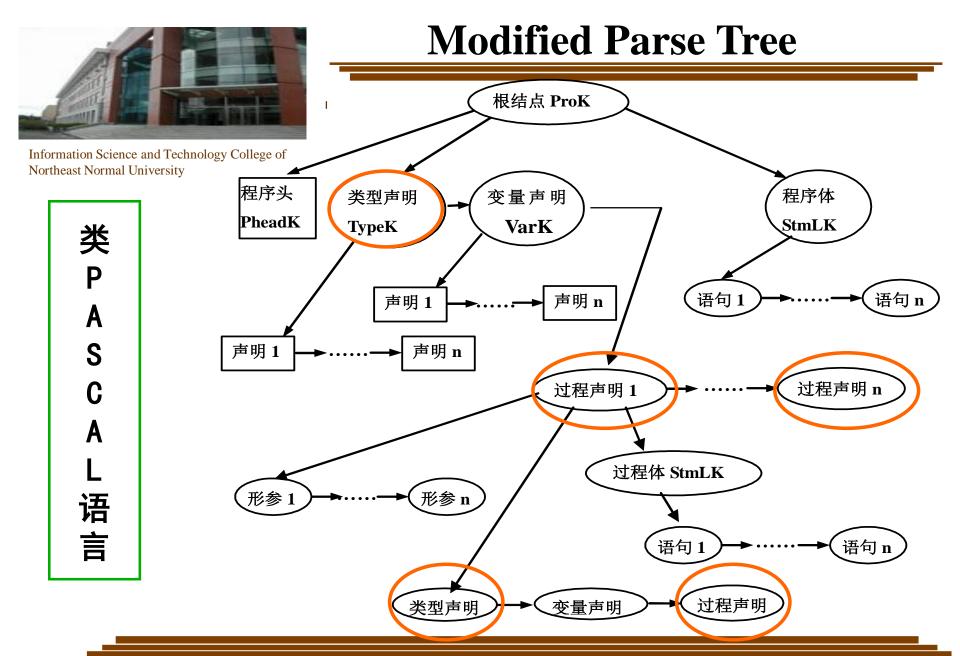
6.3 Semantic Analysis of Types

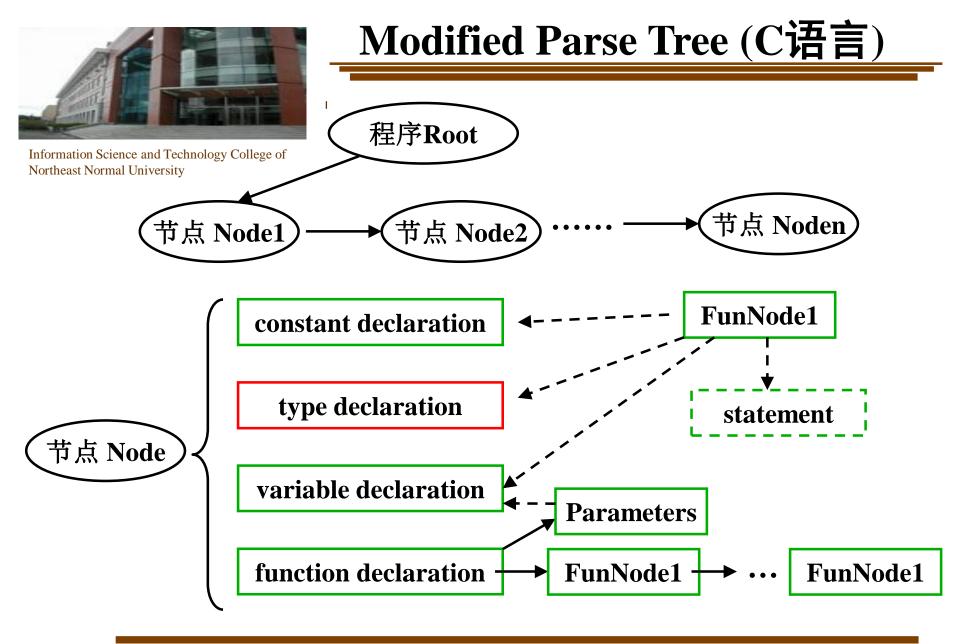
What?

- Establish internal representation of types;
- Check semantic errors in type definition;

• Where?

- Type Declaration
- Variable declaration
- Parameter declaration







Semantic Analysis of Types

• How?

- According to different types, establish their internal representations;
 - 类型名;
 - 基本类型;
 - 枚举类型:
 - 数组类型;
 - 结构类型;
 - 联合类型;
 - 指针类型;
 - 递归类型



6.3 Semantic Analysis of Types

- 类型名
 - search symbol table for pointer to the internal representation;
- 基本类型
 - predefined & stored in symbol table

Name	Kind	TypePtr	Size		Kind
int	typeKind	intPtr —	┝━┥	intSize	intTy
bool	typeKind	boolPtr —		boolSize	boolTy
char	typeKind	charPtr _	├	charSize	charTy
real	typeKind	realPtr —	 [realSize	realTy

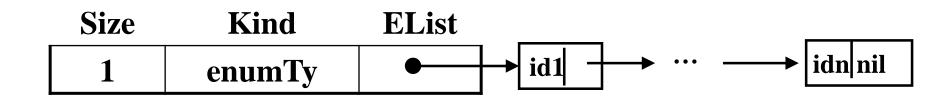


6.3 Semantic Analysis of Types

- 仅千天王
 - General form: (id1, ..., idn)
- Constant identifiers:(id1, intType, constKind, 0)

(idn, intType, constKind, n-1)

- Establish list of constant identifiers;
- Establish internal representation of enumeration;





6.3 Semantic Analysis of Types

• 数组类型

- Get low and up;
- Get or establish an pointer for internal representation of element type;
- Calculate size: size = sizeof(ElemTyp)× (Up-Low+1)
- Establish internal representation of array;

Size	Kind	Low	$\mathbf{U}\mathbf{p}$	ElemTy
	arrayTy			



6.3 Semantic Analysis of Types

Structure

- For each field identifier, establish the internal representation of its type;
- Link pointers of all field to form a chain for Body;
- Set the size as the sum of each field's size;
- Establish the internal representation of structure;

Size	Kind	Body
	structTy	



6.3 Semantic Analysis of Types

Union

- For each field identifier, establish the internal representation of its type (the offset of each field is 0);
- Link pointers of all field to form a chain for Body;
- Set the size as the maximum of all fields' size;
- Establish the internal representation of structure;

Size	Kind	Body
	unionTy	



6.3 Semantic Analysis of Types

- 指针类型
 - Get the pointer to the base type;
 - Establish the internal representation of pointer;
- 递归类型

Typedef struct treenode

{ int val;

struct treenode * left;

struct treenode * right;} TreeNode;

- Solution: 回填;



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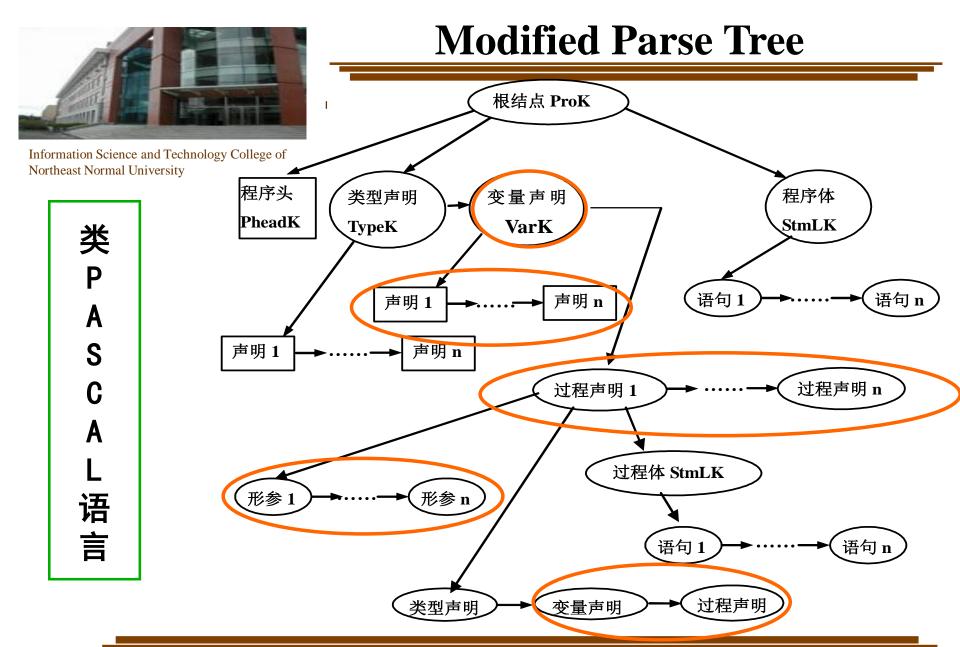
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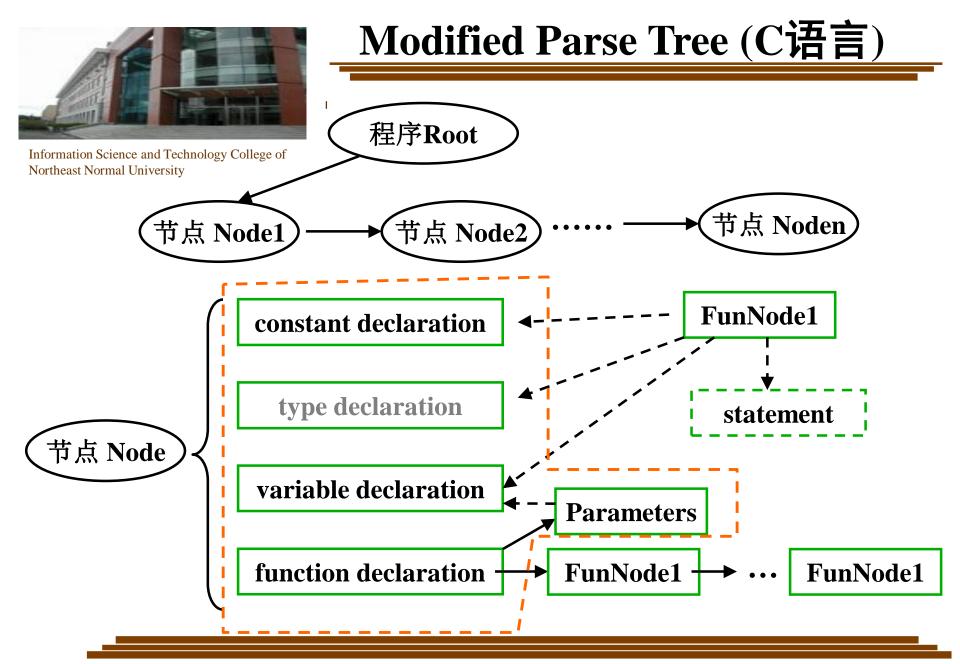
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- 分析任务:扫描声明部分的程序段,并且做语法分析的同时建立定 义性标识符的符号表,其间通过调用类型处理子程序检查类型的正 确性。
- General Process
 - <u>Collect</u> attributes of identifiers declared;
 - Establish internal representation of identifiers declared;
 - <u>Check</u> "repeat declaration" error through searching symbol table;
 - *Insert* into symbol table;
- Declaration
 - Constant declaration
 - Type declaration
 - Variable declaration
 - Function declaration
 - Label declaration (optional)







- Constant declaration
 - General form: id = number;
 - Internal representation:
 - Establish the internal representation of number;
 - Establish the internal representation of the type of number;
 - Establish the internal representation of id;

TypePtr	Kind	Value		
	constKind			



- Type declaration
 - General form: id = type;
 - Internal Representation:
 - Establish the internal representation of *type*;
 - Establish the internal representation of *id*;

TypePtr	Kind
	typeKind



6.4 Semantic Analysis of Declaration

Variable declaration

- General form: type id;
- Internal Representation :
 - Establish the internal representation of *type*;
 - <CurrentLevel, CurrentOffset>
 - Access: dir/indir
 - Establish the internal representation of *id*;
 - CurrentOffset += sizeof(type)

TypePtr	Kind	Access	Level	Offset
	varKind			



- Function/Procedure declaration
 - General form: type id (params) { body }
 - Process of analysis:
 - Pointer to internal representation of type;
 - 保存CurrentOffset , CurrentOffset = 0;
 - Establish internal representation of parameters (as variables, whose level is CurrentLevel+1)
 - Link all parameters;
 - Establish internal representation of *id* (whose level is CurrentLevel);
 - 进入函数体, CurrentLevel ++;
 - <u>Code</u> will be set after codes have been created;
 - Size will be set after all variables declared in this function are handled;
 - 恢复CurrentOffset;

TypePtr	Kind	Class	Level	Param	Code	Size	Forward
	routKind	actual					



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- General checking
 - Whether identifiers in the statement are already declared;
 - Type-related semantic errors;
- Different elements in the Body
 - Assignment
 - Conditional statement
 - Loop statement
 - Function call
 - Expressions
 - Variables



- Assignment
 - General form: $\underline{var} = \underline{Expression}$
 - General process:
 - Get the type of <u>var</u>;
 - Get the result type of *Expression*;
 - Check whether the type of <u>var</u> and the type of <u>Expression</u> are compatible(赋值相容);



- Conditional statement
 - General form: if Exp S1 [else S2]
 - General Process:
 - Check whether the result type of Exp is boolean;



- Loop statement
 - Different loops
 - While Exp {S}
 - Check whether the result type of Exp is boolean;
 - For i=E1 to E2 {S}
 - Check whether the types of i, E1 and E2 are compatible;



6.5 Semantic Analysis of Body

Function call

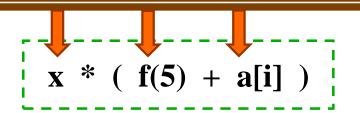
- General form: f(E1, ..., En)
- General Process:
 - Search symbol table to get the attributes of f;
 - Check whether f is "routeKind";
 - Get the <u>result type</u> of Ei (i=1, ..., n);
 - Check whether the number of actual parameters and the number of formal parameter of f is the same;
 - Check whether the type of ith formal parameter and the ith actual parameter is compatible; (类型和种类)



6.5 Semantic Analysis of Body

Expressions

- Check "运算分量类型相容";
 - 求表达式的结果类型;[
 - 求表达式的种类
 - 数据
 - 地址(指针)
 - 过程/函数



- (1) Search x in Symbol Table
- (2) x: (t1, varKind, L1, off1, dir)
- (3) Search f in symbol table
- (4) **f**: (**t2**, routKind, L2, paramL,...)
- (5) Search a in Symbol Table
- (6) a: (t3, varKind, L3, Off3, indir)
- (7) Check whether t3 is arrayType;
- (8) Check whether t2 and t3.elemTy are "+" compatible; if yes, get the (result type t', 种类);
- (9) Check whether t1 and t' are "*" compatible; get (结果类型t,种类c)
- (10) **Return** (t,c)



6.5 Semantic Analysis of Body

Variable

- General form: $V \rightarrow id \mid V[Exp] \mid V.id \mid *V$
- 标识符情形(id):
 - search symbol table to get its attributes, from which return (type, 种类);
- ─ 数组下标变量V[Exp]:
 - check whether the result type of V is arrayKind,
 - whether the result type of Exp is correct,
 - returns (ElemType, 种类);



6.5 Semantic Analysis of Body

Variable

- 结构/联合的域名V.id:
 - check whether the result type of V is structure or union,
 - whether id is one of its field,
 - return (fieldType, 种类)
- *V:
 - check whether the result type of V is pointTy,
 - return (BaseType, 种类)



Some Issues

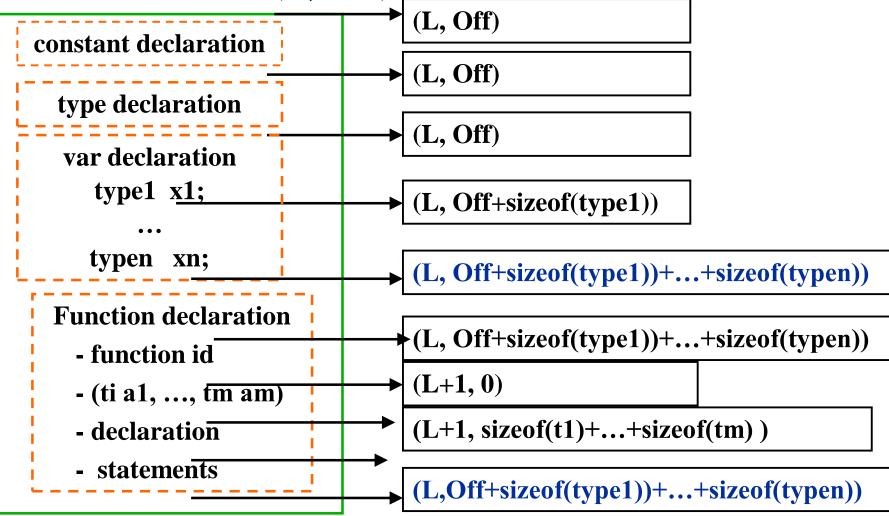
Symbol Table

constant declaration Add constant identifier to symbol table Add enum constant id to symbol table; type declaration Add enum type id to symbol table; Add type id to symbol table; other var declaration Add variable id to symbol table; **Function declaration** Add function id to symbol table; - function id Add parameters(var id) to symbol table; - parameters - declaration Add id declared to symbol table; - statements Search for symbol table;

me Table

Some Issues

• Level & Offset (L, Off)





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Exercise

• Please write down (level,offset) at different red points of following program. If at point 0 it is (L, Off).



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6.6 Attribute Grammar & Action Grammar

- How to formally describe static semantics and implement static semantic analysis?
- · Syntax directed methods(语法制导的方法)
 - Attribute Grammar = CFG + attributes, which was introduced by Knuth and used widely;
 - Action Grammar = CFG + actions, which was introduced in Prof. JIN's book;
- Attribute grammar and Action Grammar can be implemented automatically, the resulting program can be semantic analyzer;



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Attribute Grammar(属性文法)

Overview

- 在CFG的基础上,描述静态语义(相关属性);
- 一种规范定义(specification);
- 广泛应用静态语义描述、程序转换 、程序分析、数据流 分析等领域,是一个非常有用的抽象描述方式;
- 在编译器中的应用
 - 语义检查
 - 代码生成
- 一个例子



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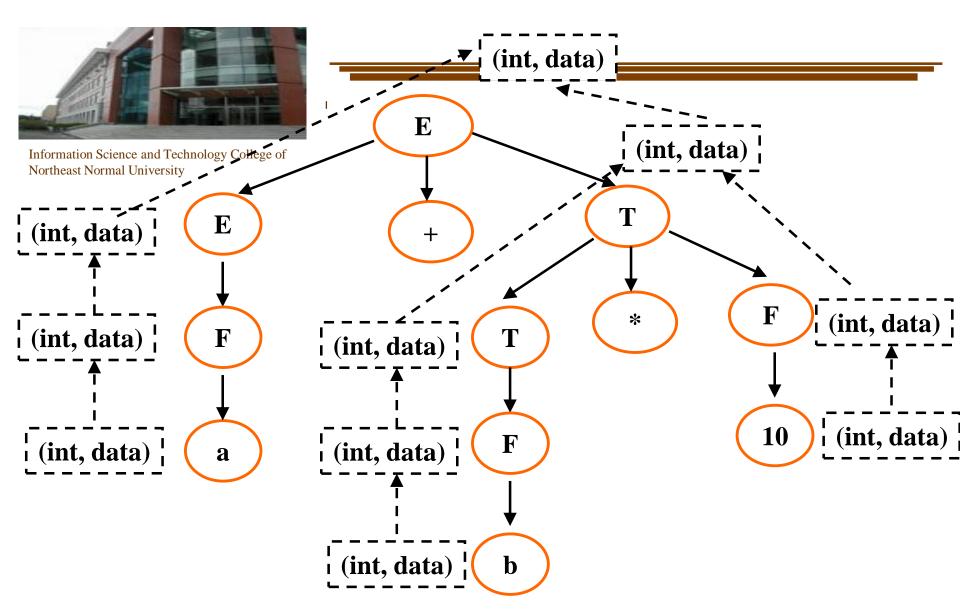
Attribute Grammar(属性文法)

Main idea

- A given context free grammar G
- Define attributes for symbols (non-terminal or terminal)
 - Attributes could be of complex structure, representing value, type, parse tree, code,
- For each production, define a rule for how to calculate the attributes of symbols in the production;

Example of Attribute Grammar

符号	属性定义	产生式	属性计算规则
E	type, class	$E^1 \to E^2 + T$	if (E ² .type=T.type) then {E ¹ .type= E ² .type; E ¹ .class = Data} else error;
T	type, class	$E \rightarrow T$	E.type= T.type; E.class = T.class
\mathbf{F}	type, class	$T^1 \rightarrow T^2 * F$	if (T ² .type=F.type) then {T ¹ .type= T ² .type; T ¹ .class = Data} else error;
n	type, class	$T \rightarrow F$	T.type= F.type; T.class = F.class
id	type, class	$F \rightarrow id$	F.type = id.type; F.class = id.class
+ *		$F \rightarrow n$	F.type = n.type; F.class = n.class
()		$\mathbf{F} \rightarrow (\mathbf{E})$	F.type = E.type; F.class = E.class





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Action Grammar(动作文法)

- Action Grammar
 - A CFG with action symbols appearing the right parts of productions;
 - each action symbol corresponds to a function or procedure;
- The definition of an action grammar depends on the parsing methods;
- Example

Action Grammar for LL(1) Parsing

An action grammar for calculating the number of 1 and 0 in the input string

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Northeast Normal University $V_T = \{0, 1, a\}$ $\mathbf{V_N} = \{\mathbf{S, L}\}$ S = S**P**: $S \rightarrow L$ $L \rightarrow 0 L$ $L \rightarrow 1 L$ $L \rightarrow a$

```
V_T = \{0, 1, a\}
                                           init()
\mathbf{V_N} = \{\mathbf{S, L}\}
                                           \{sum0 = 0; sum1 = 0;\}
S = S
                                           print()
                                           {printf(sum0, sum1);}
  S \rightarrow \langle init \rangle L \langle print \rangle
                                           add0()
  L \rightarrow 0 <add0> L
                                           {sum0++;}
  L \rightarrow 1 <add1> L
                                           add1()
  L \rightarrow a
                                           {sum1++;}
```



Working Process of Action Grammar

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$V_T = \{0, 1, a\}$	
$\mathbf{V_N} = \{\mathbf{S, L}\}$	
S = S	
P:	
{	
$(1)S \rightarrow \langle init \rangle L \langle print \rangle$	{0,1}
$(2)L \rightarrow 0 < add0 > L$	{0}
$(3)L \rightarrow 1 < add1 > L$	{1 }
$(4)L \rightarrow a$	{a}
}	

0		
分析栈	输入串	动作
S	01a#	推导(1)
<init> L <print></print></init>	01a#	sum0=0;
		sum1=0;
L <print></print>	01a#	推导(2)
0 <add0> L</add0>	01a#	匹配
<print></print>		
<add0> L <print></print></add0>	1a#	sum0=1;
L <print></print>	1a#	推导(3)
1 <add1> L</add1>	1a#	匹配
<print></print>		
<add1> L <print></print></add1>	a#	sum1=1
L <print></print>	a#	推导(4)
a <print></print>	a#	匹配
<pri><print></print></pri>	#	"1,1"
	#	接受

Compiling and Running of Program



Action Grammar for LR Parsing

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```
V_T = \{0, 1, a\}
\mathbf{V_N} = \{\mathbf{S, L}\}\
S = S
P:
  S \rightarrow L
  L \rightarrow 0 L
  L \rightarrow 1 L
  L \rightarrow a
```

```
V_T = \{0, 1, a\}
                                      init()
V_N = \{S, L\}
                                      \{sum0 = 0; sum1 = 1;\}
S = S
                                      print()
                                      {printf(sum0, sum1);}
 S \rightarrow L
                 <print>
                                      add0()
 L \rightarrow D L
                                      {sum0++;}
 D \rightarrow 1
                 <add1>
 D \rightarrow 0
                 <add0>
                                      add1()
 L \rightarrow a
                                      {sum1++;}
```



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$$\begin{split} V_T &= \{0, 1, a\} \\ V_N &= \{S, L\} \\ S &= S \\ P; \\ \{ & (1) \ S \to L \quad <print> \\ (2) \ L \to D \quad L \\ (3) \ D \to 1 \quad <add1> \\ (4) \ D \to 0 \quad <add0> \\ (5) \ L \to a \end{split}$$

Working Process of Action Grammar

符号栈	输入串	动作
<init></init>		sum0=0; sum1=0
	01a#	shift
0	1a#	Reduce (4)
D <add0></add0>	1a#	sum0=1
D	1a#	Shift
D1	a#	Reduce (3)
DD <add1></add1>	a#	sum1=1
DD	a#	Shift
DDa	#	Reduce(5)
DDL	#	Reduce(2)
DL	#	Reduce(2)
L <print></print>	#	Reduce(1) print
S	#	Accept



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Summary of Chapter 6

- Main function of Semantic Analysis?
- About Symbol table
 - What ?
 - The purpose of establishing symbol table;
 - Internal representation of identifiers, types & values;
 - Understanding (level, offset) for variables;
 - Understanding the scope for identifiers;
 - The organization of symbol table;
- How to establish symbol table during semantic analysis?
 - Add new item;
 - Check repeat declaration error;
 - Get attributes for identifiers;
 - Especially (level, offset)



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Summary of Chapter 6

- The main task of semantic analysis for different syntactic structures of a program
 - Declaration
 - Constant
 - Type
 - Variable
 - Function/procedure
 - Body
 - Statements
 - Expressions
 - variables



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Summary of Chapter 6

- The main idea of Attribute Grammar?
- The main idea of Action Grammar?