```
1 定个大目标: (春招)吊打(HR)
2020年11月13日晚
```

栈和队列刷题

232. Implement Queue using Stacks

1、一入一出: 双栈法

①、代码

```
1 class MyQueue {
2
      //一、算法思想
          //通过:双栈实现队列,一个栈(inStack),另一个栈
   (outStack)
      //二、实现细节
5
          //1、创建两个栈:用LinkedList,构造方法中进行初始化
          //2、写一个outStack为空,则inStack全出栈,入outStack
7
   的方法: inPutOut
          //3、入栈时:直接入inStack
9
          //4、出栈时:调用 inPutOut 方法
         //5、判断是否为空: instack && outStack 都为空是才为空
10
11
         //6、获取栈顶元素:调用inPutOut方法,然后outstak.peek
   即可
12
      //三、测试用例
13
          //1, [1, 2, 3] --->> push, push, pop, peek
14
   [null, null, 2 , 1]
15
      //四、复杂度
16
          //1 Time Complexity: O(n)
17
18
          //2 Space Complexity: O(n)
19
      private Deque<Integer> inStack;
20
21
      private Deque<Integer> outStack;
22
23
      public MyQueue() {
24
          inStack = new LinkedList<>();
```

```
outStack = new LinkedList<>();
25
        }
26
27
        /** Push element x to the back of queue. */
28
29
        public void push(int x) {
30
            inStack.push(x);
31
        }
32
33
        /** Removes the element from in front of queue and
    returns that element. */
        public int pop() {
34
35
            inPutOut(inStack, outStack);
36
37
            return outStack.pop();
38
        }
39
        /** Get the front element. */
40
        public int peek() {
41
            inPutOut(inStack, outStack);
42
43
            return outStack.peek();
44
        }
45
        /** Returns whether the queue is empty. */
46
        public boolean empty() {
47
48
            return inStack.isEmpty() && outStack.isEmpty();
49
        }
50
        /**inStack入栈 outStack的方法*/
51
        public void inPutOut(Deque inStack, Deque outStack){
52
            if(outStack.isEmpty()){
53
54
                while(!inStack.isEmpty()){
55
                    outStack.push(inStack.pop());
56
                }
57
            }
        }
58
59 }
```

②、代码图

```
private Deque<Integer> inStack;
public MyQueue() {
    inStack = new LinkedList<>();
public int pop() {
     inPutOut(inStack, outStack);
     return outStack.pop();
     inPutOut(inStack, outStack);
     return outStack.peek();
 public boolean empty() {
     return inStack.isEmpty() && outStack.isEmpty();
/**inStack入栈 outStack的方法*/
public void inPutOut(Deque inStack, Deque outStack){
  if(outStack.isEmpty()){
        while(!inStack.isEmpty()){
             outStack.push(inStack.pop());
             核心所在: 当outStack为空时, inStack全部入栈outStack
```

739、Daily Temperatures (Medium)

1、暴力法:

①、(强烈不推荐)代码如下

```
class Solution {
       //一、算法思路:
 2
 3
            //通过暴力法:两个for循环
 4
       //二、实现细节
 5
           //1、a[i] , j = i+1, a[i]与a[j]依次比较
 6
7
           //2、如果a[i] >= a[i] 则直接相减: 存入a[i]
           //3、break跳出内层循环
9
       //三、测试用例:
10
11
           //1、[70, 92, 66]
           //2、[73, 74, 75, 71]
12
13
14
       //四、Complexity
15
           //1, Time Complexity: O(n\2)
16
           //2 Space Complexity: O(1)
17
       public int[] dailyTemperatures(int[] T) {
18
           for(int i = 0; i < T.length; i++){
19
               boolean flag = true;
                                                  //true为
   T[i] < T[i], 作为标志
20
               for(int j = i+1; j < T.length; j++){
21
                   if(T[i] > T[i]){
22
                       T[i] = j-i;
23
                       flag = false;
                                                  //此时
   T[i] > T[i]
24
                       break;
25
                   }
26
```

```
27
28
29
               if(flag){ //此时 T[i], 后面没有比它: 气温高的日子
                   T[i] = 0;
30
               }
31
32
           }
33
34
           return T;
35
      }
36 }
```

②、代码图 (便于分析)

③、复杂度 (LeetCode)

```
Success Details >
```

Runtime: $922\ ms$, faster than 11.40% of Java online submissions for Daily Temperatures.

Memory Usage: 47.4~MB, less than 5.12% of Java online submissions for Daily Temperatures.

Next challenges:

Next Greater Element I

Show off your acceptance:







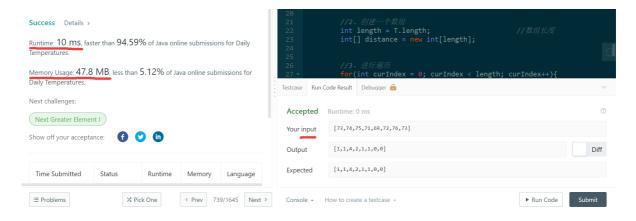
2、栈+数组

①、(推荐)代码如下

```
class Solution {
1
      //一、算法思路
2
          //通过利用: 栈(存放遍历温度下标) + 数组(存放升温距离天
 3
   数)
4
5
      //二、实现细节
6
          //1、栈采用: Deque<Integer> stack = new
   LinkedList<>() --->>> 双端链表模拟,速度优于Stack
          //2、数组采用: int[] distance;
7
8
      //三、测试用例
9
10
          //1、[73, 56, 88]
          //2、[74, 75, 71, 69, 72, 76, 73]
11
12
13
      //四、complexity
          //1、Time Complexity: O(n) 进行了for循环 n次 +
14
   栈中元素出栈次数
15
          //2、Space Complexity: O(n) 栈n + 数组n
16
17
       public int[] dailyTemperatures(int[] T) {
          //1、创建一个栈
18
19
          Deque<Integer> stack = new LinkedList<>();
20
          //2、创建一个数组
21
```

```
int length = T.length;
22
                                                     //数组
   长度
23
           int[] distance = new int[length];
24
25
26
           //3、进行遍历
27
           for(int curIndex = 0; curIndex < length;</pre>
   curIndex++){
28
                //distance[curIndex] = 0;
                                                       //此
   处不用进行初始化
29
                                                       //因
   为java中:数组没有赋初值时:为0
30
31
               while(!stack.isEmpty() && T[curIndex] >
   T[stack.peek()]){
                  int preIndex = stack.pop();
                                                     //栈顶
32
   元素的下标
                  distance[preIndex] = curIndex - preIndex;
33
34
               }
35
               stack.push(curIndex);
36
           }
37
38
39
           return distance;
40
       }
41 }
```

②、代码图 (便于分析)



503、Next Greater Element II

- 一、栈+数组
- ①、代码图

```
class Solution {
1
      //一、算法思想
2
          //1、利用栈 + 返回数组: 初始化-1
3
4
      //二、算法细节
5
         //1、将数组: 通过取余 2 * length 可以,模拟循环数组
6
7
          //2、遍历需要两遍, 才能最终确认下来,每一个元素的下一
   个:最大值
          //3、提前将:返回数组赋值为 -1, 避免了找不到:赋值-1操
8
   作
9
         //4、核心: if(!stack.isEmpty && num > stack.peek())
   则 next[stack.pop()] = num
10
          //5、核心: i < length 则 stack.push(i)
11
      //三、测试用例
12
13
         //1、[1, 2, 1]
14
15
      //四、复杂度
```

```
//1、Time Complexity: O(N) -->>因为一个
16
   for (2*length次) + 小于等于n次出栈 + i < length入栈
           //2、Space Complexity: O(N) -->>一个栈(入栈n次)
17
   + 一个数组长n
18
       public int[] nextGreaterElements(int[] nums) {
19
           //1、获取nums的长度
20
           int length = nums.length;
21
22
23
           //2、前置预判断: length == 0 则直接返回
           if(length == 0){
24
25
               return nums;
26
           }
27
28
           //3、创建一个栈
29
           Deque<Integer> stack = new LinkedList<>();
30
31
           //4、创建一个:存放nums[]数组的(下一个大元素的数组),
   并赋值为: -1
32
           int maxNext[] = new int[length];
33
           Arrays.fill(maxNext, -1);
34
35
           //5、模拟循环数组:遍历
           for(int i = 0; i < 2*length; <math>i++){
36
37
               //5.1、记录当前数组元素的: value
38
               int curvalue = nums[i%length];
39
40
               //5.2、进行判断
41
               while(!stack.isEmpty() && curValue >
   nums[stack.peek()]){
42
                   maxNext[stack.pop()] = curValue;
43
               }
44
               if(i < length){</pre>
45
46
                   stack.push(i);
47
               }
48
           }
49
50
           return maxNext;
51
52
       }
53 }
```

Success Details >

Runtime: 6~ms, faster than 94.75% of Java online submissions for Next Greater Element II.

Memory Usage: 41 MB, less than 6.70% of Java online submissions for Next Greater Element II.

Next challenges:

Next Greater Element I

Next Greater Element III

Show off your acceptance:







225. Implement Stack using Queues

- 一、双端队列法
- ①、代码图

```
//2、求出: 例列长度
int length = queue.size();
   /** Removes the element on top of the stack and returns that element. */
public int pop() {
    return queue.remove();
```

```
1
   class MyStack {
2
3
      //一、算法思想
         //采用:双端队列Deque
4
5
      //二、算法细节
6
7
         //1、add时: 判断
             //第一步: (元素)直接入队,求出队列长度 length
8
             //第二步: 通过while(length-- > 1)的元素全部出队,
9
   再入队,实现逆转,相当于将(新元素)前的所有元素(都重新入队)
10
11
         //2、remove时: 直接出栈, 因为 (All the calls to pop
   and top are valid.)
```

```
12
          //3、peek时: 直接查看
13
14
          //4、empty: 直接判断
15
16
17
          //注意:切记poll()和remove()效果是一样的,别乱用
18
       //三、测试用例
19
20
          //push(1) , push(2) --->>> pop() --->>> 2
21
22
       //四、复杂度
           //1、Time Complexity: O(n) --->>> 主要在 push()
23
   处累加和用到了n, 其余操作为0(1)
24
           //2、Space Complexity: O(n) --->>> 双端队列
25
26
       /** Initialize your data structure here. */
27
28
       private Deque<Integer> queue;
29
30
       public MyStack() {
31
           queue = new LinkedList<>();
32
       }
33
       /** Push element x onto stack. */
34
35
       public void push(int x) {
36
          //1、直接入队
           queue.add(x);
37
38
          //2、求出:队列长度
39
40
          int length = queue.size();
41
42
          //3、将入队(新元素)前的(所有元素)重新入队, 实现逆
   转,变成栈的效果
          while(length-- > 1){
43
              queue.add(queue.remove());
44
45
           }
46
       }
47
       /** Removes the element on top of the stack and
48
   returns that element. */
49
       public int pop() {
50
           return queue.remove();
51
       }
```

```
52
        /** Get the top element. */
53
        public int top() {
54
55
            return queue.peek();
56
        }
57
        /** Returns whether the stack is empty. */
58
        public boolean empty() {
59
            return queue.isEmpty();
60
61
        }
62
   }
63
```

Success Details >

Runtime: 0 ms, faster than 100.00% of Java online submissions for Implement Stack using Queues.

Memory Usage: 36.8 MB, less than 8.40% of Java online submissions for Implement Stack using Queues.

Next challenges:



155、Min Stack (Easy)

一、栈 + Node结点法

①、代码图

```
public void push(int x) {
    if(!stack.isEmpty() && x > stack.peek().minValue){ //此时模不为空,(当前栈顶)存放的最小值:依然是最小值    int minValue = stack.peek().minValue;    stack.push(new Node(x, minValue));
                                                                            此处为精髓:
1、如果: 栈不为空 && x > 栈顶.minValue 也就是最小值时,最小值如旧。
2、否则: x就是最小值, 直接入栈,成为新的(最小值)
public void pop() {
    stack.pop();
}
public int top() {
    return stack.peek().value;
}
public int getMin() {
    return stack.peek().minValue;
}
private class Node{
private int value; //存放当前: 入稅值x
private int minValue; //存放找中: 最小值
      public Node(int value, int minValue){
   this.value = value;
   this.minValue = minValue;
```

```
class MinStack {
1
2
      //一、算法思想
3
         //利用:内部结点类Node + 模拟栈Deque 实现, 始终保持栈
   顶元素(结点.minValue属性)存放着(栈最小值)
4
5
      //二、算法细节
         //1、创建一个内部类: Node
6
7
             //成员变量(value): 存放当前:入栈值x
             //成员变量(minvalue): 存放栈中:最小值
8
9
10
         //2、创建一个LinkedList模拟栈
11
         //3、在push: 进行判断
12
```

```
//1、当stack为空时: 直接Node入栈,
13
               //2、当stack不为空时:将(Node.minValue值)与
14
    (栈顶元素.minValue) 进行比较
                   //若Node.minValue < 栈顶元素.minValue 则直
15
   接入栈
                   //否则: 将Node.minvalue = 栈顶元素.minvalue
16
   再入栈
              //3、stack.getMin() --->>> 直接获取(栈
17
   顶.minvalue)即可
              //4、其余: pop, peek, isEmpty 不用更改
18
19
20
      //三、测试用例
21
              //push 1 -->>> push -2 -->>> push 0
   \RightarrowgetMin = -2, pop = 0, getMin = -2,
22
23
       //四、复杂度
              //Time Complexity: 0(1) --->>> push.
24
   pop、top、getMin: 这几个方法(都是常数阶),所以为O(1)
25
             //Space Complexity: O(n) --->> 用到了: 栈
   + Node, 当最坏情况是,全部入栈, 每次x都比栈顶的(minvalue)要
   小,0(n)跑不掉了
26
27
28
       //1、创建一个栈
29
       private Deque<Node> stack;
30
       /** initialize your data structure here. */
31
32
       public MinStack() {
          stack = new LinkedList<>();
33
34
       }
35
36
       public void push(int x) {
37
          if(!stack.isEmpty() && x > stack.peek().minValue){
   //此时栈不为空, (当前栈顶) 存放的最小值: 依然是最小值
              int minvalue = stack.peek().minvalue;
38
39
              stack.push(new Node(x, minValue));
40
          }else{
              stack.push(new Node(x, x));
41
          }
42
43
44
       }
45
       public void pop() {
46
```

```
47
           stack.pop();
       }
48
49
50
       public int top() {
51
           return stack.peek().value;
52
       }
53
       public int getMin() {
54
           return stack.peek().minvalue;
55
       }
56
57
       private class Node{
58
           private int value; //存放当前: 入栈值x
59
           private int minvalue;
                                      //存放栈中:最小值
60
61
           public Node(int value, int minValue){
62
               this.value = value;
63
               this.minValue = minValue;
64
65
           }
66
       }
67
   }
68
69
   /**
   * Your MinStack object will be instantiated and called as
70
   such:
    * MinStack obj = new MinStack();
71
   * obj.push(x);
72
   * obj.pop();
73
* int param_3 = obj.top();
75
    * int param_4 = obj.getMin();
76
    */
```

```
Success Details >
```

Runtime: 4 ms, faster than 92.81% of Java online submissions for Min Stack.

Memory Usage: 40.9 MB, less than 50.53% of Java online submissions for Min Stack.

Next challenges:

Sliding Window Maximum

Max Stack

Show off your acceptance:







二、数据栈 + 辅助 (最小值) 栈

①、代码图

```
class MinStack {
1
2
       //一、算法思想
3
          //运用:数据栈 + 辅助(最小值)栈 = 双栈模式
4
      //二、算法细节
5
          //1、创建两个栈:
6
                  //①、Deque<Integer> dataStack = new
   LinkedList<>();
8
                  //②、Deque<Integer> minStack = new
   LinkedList<>();
9
          //2、push时:
                 //①、直接入栈: dataStack
10
                  //②、判断minStack是否为空
11
12
                         //如果为空:直接入栈
```

```
//否则:判断入栈值 (与)
13
   minStack.peek()大小
14
                                //如果:入栈值(小于等于),
   则入栈, 因为 = 不入,会导致出栈后,拿peek时空指针
15
                                //否则:不入栈
          //3、pop时:将dataStack直接入栈,并记录值 value,并
16
   用value 和 minStack.peek()比较大小
17
                  //如果:相等时,minStack.pop()也出栈
18
19
20
       //三、测试用例
21
          //push 1, push -2, push 0 --->>> getMin() =
   -2, pop() = 0, getMin() = -2
22
23
      //四、复杂度
          //Time Complexity: O(1) -->>>所有方法操作都是:
24
   0(1)
25
          //Space Complexity: O(n) -->>> 双栈
26
27
       //1、创建:数据栈 + 辅助(最小值)栈
28
       private Deque<Integer> dataStack;
29
       private Deque<Integer> minStack;
30
       /** initialize your data structure here. */
31
32
       public MinStack() {
33
          dataStack = new LinkedList<>();
          minStack = new LinkedList<>();
34
35
       }
36
37
       public void push(int x) {
38
          dataStack.push(x);
          if(minStack.isEmpty() || x <= minStack.peek()){ //</pre>
39
   切记: x == 最小值也要入栈
              minStack.push(x);
40
                                                       //
   因为出栈时: 最小栈可能也会出栈,不入栈导致: 最小栈.peek时空指针
41
          }
42
       }
43
       public void pop() {
44
45
          int value = dataStack.pop();
46
          if(value == minStack.peek()){
47
              minStack.pop();
48
          }
```

```
49
50
        public int top() {
51
            return dataStack.peek();
52
53
        }
54
        public int getMin() {
55
            return minStack.peek();
56
57
        }
58 }
59
```

Success Details >

Runtime: 4~ms, faster than 92.81% of Java online submissions for Min Stack.

Memory Usage: 41.2 MB, less than 22.00% of Java online submissions for Min Stack.

Next challenges:



Max Stack

Show off your acceptance:





三、自定义: 链栈

①、代码图

```
public void push(int x) {
   if(head == null){
     head = new LinkedNode(x, x, head);
   }else{
    head = new LinkedNode(x, Math.min(x, head.minValue), head);
}
                                                        核心在于: 传入当前链表, 通过构造方法, 进行(头插法)操作
public int top() {
    return head.value;
}
public int getMin() {
    return head.minValue;
}
this.minValue = minValue;
this.next = head; 結婚
```

```
1
   class MinStack {
2
      //一、算法思路
3
          //采用: 自己创建(一个链栈) + 头插法
4
      //二、算法细节
 5
          //1、创建:链结点LinkStack(数据结构)
6
7
             //int value: 存放当前入栈值
             //int minvalue: 存放栈最小值
8
9
             //StackNode next: 存放下一个结点
          //2、push时: 若链栈 LinkStack == null --->>> 直接
10
   入栈, 最小值为x
                  //否则: 用Math.min(x, LinkStack.minValue)
11
   找出最小值,
             然后入栈
```

```
//入栈: 采用头插法, 通过head = LinkStack的构造方
12
   法, 巧妙插入
13
           //3、pop时: 直接head = head.next
           //4、top时: 直接head.value
14
15
           //5、getMin时: 直接head.getMin
16
       //三、测试用例
17
18
           //push 1, push -2, push 0 --->>> getMin() = -2,
   pop() = 0, getMin() = -2
19
       //四、复杂度
20
21
           //Time Complexity: 0(1)
22
           //Space Complexity: O(n)
23
24
25
       //1、创建一个: 链栈的头指针
26
       private LinkedNode head;
27
28
       /** initialize your data structure here. */
29
       public MinStack() {
30
       }
31
32
       public void push(int x) {
33
34
           if(head == null){
35
               head = new LinkedNode(x, x, head);
36
           }else{
37
               head = new LinkedNode(x, Math.min(x,
   head.minvalue), head);
           }
38
       }
39
40
41
       public void pop() {
42
           head = head.next;
       }
43
44
45
       public int top() {
           return head.value;
46
47
       }
48
49
       public int getMin() {
           return head.minValue;
50
51
       }
```

```
52
       private class LinkedNode{
53
                             //存放当前入栈值
           private int value;
54
           private int minValue; //存放栈最小值
55
           private LinkedNode next; //存放下一个结点
56
57
           LinkedNode(int value, int minValue, LinkedNode
58
   head){
59
               this.value = value;
60
               this.minValue = minValue;
               this.next = head;
61
62
       }
63
64 }
65
```

Memory Usage: 40.5 MB, less than 85.51% of Java online submissions for Min Stack.

Next challenges:



Max Stack

Show off your acceptance:





20、Valid Parentheses (Easy)

一、栈 + 直接遍历

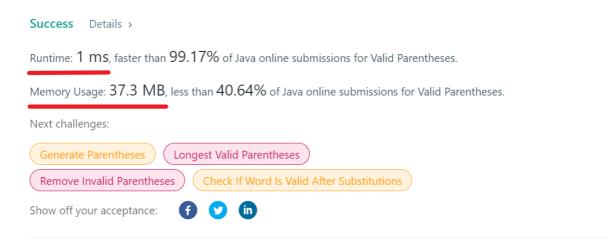
①、代码图

1 可以补充: 前置判断括号个数, if(length % 2 == 1) return false;

```
//2、获取字串长度
int length = s.length();    //因为题目给定: 1 <= s.length <= 10^4 所以不用前置判断了
for(int i = 0; i < length; i++){</pre>
        stack.push(')');
e if(c == '[')
stack.push(']');
e if(c == '{')
                                   精髓在于:
                                       1、入栈相反字符
                                       2、匹配栈顶时为空 || 出栈元素 != 当前字符 -->> false
    return stack.isEmpty();
```

```
class Solution {
1
2
      //一、算法思想
3
         //采用:遍历字符串 + 栈
4
      //二、算法细节
5
6
         //1、创建一个栈: Deque<Character> stack = new
   LinkedList<>()
7
          //2、遍历字符串, 通过if、else if 进行选择
             //@、遇到(,[, { 左括号则入栈它们的(右括号)
8
9
             //②、遇到), ], } 时
                //判断: 栈是否为空, 如果为空 return false
10
                //判断: 右括号是否 == 出栈元素, 如果不等于
11
   return false
12
         //3、return stack.isEmpty() 作为是否是: 有效括号的结
   果
```

```
13
14
15
       //三、测试用例
          //1, ([)] --->>> return false
16
17
          //2、{[]}
                     --->>> return true
18
          //3、[]
                     --->>> return true
19
       //四、复杂度
20
21
          //1 Time Complexity: O(n)
                                          -->> 因为要遍历
   length长度
          //2、Space Complexity: O(n) -->> 用到了: 栈
22
23
     //五、优缺点
24
25
          //1、advantage: 通俗易懂, 代码量少
26
          //2、shortcoming: 拓展性不好, 一旦匹配变多, if将变
   多,而且修改起来困难
27
28
       public boolean isValid(String s) {
29
          //1、创建一个栈
30
          Deque<Character> stack = new LinkedList<>();
31
          //2、获取字串长度
32
          int length = s.length(); //因为题目给定: 1
33
   <= s.length <= 10^4 所以不用前置判断了
34
35
          //3、遍历字符串
          for(int i = 0; i < length; i++){
36
              Character c = s.charAt(i);
37
38
39
              if(c == '(')
                  stack.push(')');
40
              else if(c == '[')
41
42
                  stack.push(']');
              else if(c == '\{'\}
43
                  stack.push('}');
44
45
              else if(stack.isEmpty() || (stack.pop() != c))
46
                  return false:
47
          }
48
49
          return stack.isEmpty();
50
       }
51 }
```



二、遍历字符串 + 栈 + Map

①、代码图

```
public boolean isValid(String s) {
    int length = s.length(); //因为题目给定: 1 <= s.length <= 10^4 所以对此做:判断
    if(length % 2 == 1){    此处用得巧
    map.put('(', ')');
map.put('[', ']');
map.put('{{', '}}');
    for(int i = 0; i < length; i++){
    Character c = s.charAt(i);</pre>
     if(map.containsKey(c))
stack.push(map.get(c));
                                           containsKey()中是有s的, 注意了
         else if(stack.isEmpty() || c != stack.pop())
    return stack.isEmpty();
```

```
class Solution {
2
      //一、算法思想
         //采用: 遍历字符串 + 栈 + Map
3
4
5
      //二、算法细节
         //1、创建一个栈: Deque<Character> stack = new
6
  LinkedList<>()
7
         //2、创建一个: Map<Character, Character> = new
  HashMap<>()
8
         //3、将()、[], {}, 放入map中, 左括号为key, 右括号为
  value
```

```
//4、进行遍历: 当前字符到(Map查找)
10
               //若找到:则将Map的value入栈
11
               //若找不到: 判断当前栈是否为空 || 当前遍历字符
   != 出栈元素
12
                   //栈若为空:说明(右括号)找不到匹配的(左
   括号) -->> 直接return false
13
                   //当前遍历字符!= 出栈元素: 说明匹配不上
   -->> 直接return false
14
         //5、最后用: stack.isEmpty()的原因: 因为若能匹配,则都
   出栈了,不能匹配则(栈不为空)
15
16
      //三、测试用例
17
        //1、([)] --->>> return false
18
         //2、{[]} --->>> return true
19
        //3、[]
                  --->>> return true
20
      //四、复杂度
21
22
         //1、Time Complexity: O(n) -->> 因为要遍历
   length长度
23
         //2 Space Complexity: O(n)
                                     -->> 最坏超过 n
   , 因为用到了Map
24
      //五、优缺点
25
         //1、advantage: 添加了前置判断、拓展性强(增加匹配元
26
   素时)易于添加
        //2、shortcoming: 耗费了空间
27
28
    public boolean isValid(String s) {
29
         //1、获取字串长度
30
31
         int length = s.length(); //因为题目给定: 1
   <= s.length <= 10^4 所以对此做:判断
32
33
         if(length \% 2 == 1){
            return false;
                                 //奇数个括号, 不可
34
   能匹配
35
         }
36
37
         //2、创建一个栈
38
         Deque<Character> stack = new LinkedList<>();
39
40
41
         //3、创建一个Map, 并将成对括号: 放入
```

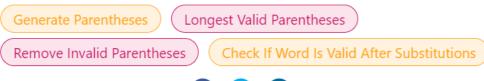
```
Map<Character, Character> map = new HashMap<>();
42
    //直接用map.put()而不用: 匿名内部类, 统一初始化
           map.put('(', ')');
43
    //是因为: map.put()的效率更高
           map.put('[', ']');
44
           map.put('{', '}');
45
46
47
48
           //4、遍历字符串
49
           for(int i = 0; i < length; i++){
50
               Character c = s.charAt(i);
51
               if(map.containsKey(c))
52
                   stack.push(map.get(c));
53
54
               else if(stack.isEmpty() || c != stack.pop())
                   return false;
55
           }
56
57
           return stack.isEmpty();
58
59
       }
60 }
```

Success Details >

Runtime: 1 ms, faster than 99.17% of Java online submissions for Valid Parentheses.

Memory Usage: 37.5 MB, less than 32.48% of Java online submissions for Valid Parentheses.

Next challenges:



Show off your acceptance:







三、哨兵 + stack

①、代码图



```
class Solution {
 1
 2
       //1、创建Map
 3
       private static final Map<Character, Character> map =
   new HashMap<>(){
4
            {
               put('(', ')');
 5
               put('[', ']');
 6
               put('{', '}');
 7
               put('?', '?');
8
                                       //作为哨兵
 9
           }
10
       };
11
12
       public boolean isValid(String s) {
13
           //1、前置欲判断
           if(s.length() % 2 == 1 ||
14
    !map.containsKey(s.charAt(0)))
                return false;
15
16
           //2、创建一个栈
17
18
           Deque<Character> stack = new LinkedList<>();
           stack.push('?');
                                      // (?,?)作为哨兵
19
20
21
           //3、遍历: 判断
           for(char c : s.toCharArray()){
22
23
                if(map.containsKey(c)){
24
                    stack.push(c);
                }else if(map.get(stack.pop()) != c){
25
26
                    return false;
27
                }
28
           }
29
           return stack.size() == 1;
30
31
       }
32 }
```

Success Details >

Runtime: $1\,ms$, faster than 99.17% of Java online submissions for Valid Parentheses.

 ${\sf Memory\ Usage:\ 37.3\ MB,\ less\ than\ 51.26\%\ of\ Java\ online\ submissions\ for\ Valid\ Parentheses.}$

Next challenges:

Generate Parentheses Longest Valid Parentheses

Remove Invalid Parentheses

Check If Word Is Valid After Substitutions

Show off your acceptance:



