

## Chapter 8 Advanced Counting Techniques

### 8.1 Applications of Recurrence Relations

#### 2. Recurrence Relations

**【Definition】** A *recurrence relation* for the sequence  $\{a_n\}$  is an equation that expresses  $a_n$  in terms of one or more of the previous terms of the sequence, namely,  $a_0, a_1, a_2, \dots, a_{n-1}$ , for all integers  $n$  with  $n \geq n_0$ , where  $n_0$  is a nonnegative integer.

$$a_n = f(a_0, a_1, a_2, \dots, a_{n-1}) \quad n \geq n_0$$

A *solution of a recurrence relation* is a sequence if its terms satisfy the recurrence relation.

Note:

- Normally, there are infinitely many sequences which satisfy a recurrence relation. We distinguish them by the *initial conditions*, the values of  $a_0, a_1, a_2, \dots$  to uniquely identify a sequence.

- The *degree* of a recurrence relation

$$a_n = a_{n-1} + a_{n-8} \quad \text{---- a recurrence relation of degree 8}$$

#### 3. Modeling with Recurrence Relations

#### 8.2 Solving Linear Recurrence Relations

**Linear homogeneous (齐次) recurrence relation of degree  $k$  with constant coefficients**

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} \quad \text{where } c_1, c_2, \dots, c_k \text{ are real numbers, and } c_k \neq 0$$

- *Linear*- linear combination of previous terms
- *constant coefficients*- the coefficients of  $a_i$ s are constants
- *degree  $k$* -  $a_n$  is a function of the previous  $k$  terms of the sequence
- *Homogeneous*- If we put all the  $a_i$ s on the left side of the equation and everything else on the right side, then the right side is 0. Otherwise *nonhomogeneous*.

A sequence satisfying the recurrence relation in the definition is uniquely determined by the recurrence relation and the  $k$  *initial conditions*:

$$a_0 = C_0, a_1 = C_1, \dots, a_{k-1} = C_{k-1}$$

#### 1. Solving Linear Homogeneous Recurrence Relation With Constant Coefficients

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

The basic approach:

To look for solutions of the form  $a_n = r^n$ , where  $r$  is a constant.

Solution Procedure:

- (1) Put all  $a_i$ s on the left side of the equation  $a_n - c_1 a_{n-1} - c_2 a_{n-2} - \dots - c_k a_{n-k} = 0$
- (2) Substitute the solution into the equation, factor out the lowest power of  $r$  and eliminate it.

$$r^n - c_1 r^{n-1} - c_2 r^{n-2} - \dots - c_k r^{n-k} = 0$$

$$r^{n-k}(r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_k) = 0$$

- (3) We obtain the equivalent equation (*Characteristic equation*)

$$r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_k = 0$$

- (4) Find its  $k$  roots  $r_1, r_2, \dots, r_k$  (*Characteristic root*)

These characteristic roots can be used to give an explicit formula for all the solutions of the recurrence relation.

**【Theorem 1】** Let  $c_1, c_2$  be real numbers. Suppose that  $r^2 - c_1 r - c_2 = 0$  has two distinct roots  $r_1, r_2$ . Then the Sequence  $\{a_n\}$  is a solution of the recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2}$  if and only if  $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$  for  $n = 0, 1, 2, \dots$ , where  $\alpha_1, \alpha_2$  are constants.

Find an explicit formula for  $a_n = c_1 a_{n-1} + c_2 a_{n-2}$ :

- (1) Determine the characteristic equation:  $r^2 - c_1 r - c_2 = 0$
- (2) Find its roots:  $r_1, r_2$
- (3) Obtain the general solution:  $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$
- (4) Determine  $\alpha_1, \alpha_2$ :

**【 Theorem 2 】** Let  $c_1, c_2$  be real numbers with  $c_2 \neq 0$ . Suppose that  $r^2 - c_1 r - c_2 = 0$  has only one root  $r_0$ . A sequence  $\{a_n\}$  is a solution of the recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2}$  if and only if  $a_n = \alpha_1 r_0^n + \alpha_2 n r_0^n$  for  $n = 0, 1, 2, \dots$ , where  $\alpha_1, \alpha_2$  are constants.

**【 Theorem 3 】** Let  $c_1, c_2, \dots, c_k$  be real numbers. Suppose that the characteristic equation  $r^k - c_1 r^{k-1} - \dots - c_k = 0$  has  $k$  distinct roots  $r_1, r_2, \dots, r_k$ . Then a sequence  $\{a_n\}$  is a solution of the recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$  if and only if  $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n + \dots + \alpha_k r_k^n$  for  $n = 0, 1, 2, \dots$  where  $\alpha_1, \alpha_2, \dots, \alpha_k$  are constants.

## 2. Linear Nonhomogeneous Recurrence Relation With Constant Coefficients

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n)$$

Where  $c_i (i=1, 2, \dots, k)$  is real numbers,  $F(n)$  is a function not identically zero depending only on  $n$ . the associated *homogeneous* recurrence relation:  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$

Note:

- Solutions to nonhomogeneous case is the sum of solutions to associated homogeneous recurrence system and a particular solution to the nonhomogeneous case.

**【 Theorem 5 】** Let  $\{a_n^{(p)}\}$  be a *particular solution* of the nonhomogeneous linear recurrence relation with constant coefficients  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n)$ . Then every solution is of the form  $\{a_n^{(p)} + a_n^{(h)}\}$ , where  $\{a_n^{(h)}\}$  is a solution of the associated homogeneous recurrence relation  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ .

**【 Theorem 6 】** Assume a linear nonhomogeneous recurrence equation with constant coefficients with the nonlinear part  $F(n)$  of the form

$$F(n) = (b_t n^t + b_{t-1} n^{t-1} + \dots + b_1 n + b_0) s^n$$

If  $s$  is not a root of the characteristic equation of the associated homogeneous recurrence equation, there is a particular solution of the form

$$(p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n$$

If  $s$  is a root of multiplicity  $m$ , a particular solution is of the form

$$n^m (p_t n^t + p_{t-1} n^{t-1} + \dots + p_1 n + p_0) s^n$$

## 8.3 Divide-and-Conquer Algorithms and Recurrence Relations

### 8.4 Generating Functions

#### 1. Generating function for a sequence

**【 Definition 1 】** The *generating function* for the sequence  $a_0, a_1, a_2, \dots, a_k, \dots$  of real numbers is the infinite series

$$G(x) = a_0 + a_1 x + \dots + a_k x^k + \dots = \sum_{k=0}^{\infty} a_k x^k$$

The generating function for a finite sequence of real numbers  $a_0, a_1, a_2, \dots, a_n$  is

$$G(x) = a_0 + a_1x + a_2x^2 + \dots + a_nx^n$$

## 2. Useful Facts About Power Series

**【Theorem 1】** Let  $f(x) = \sum_{k=0}^{\infty} a_k x^k$ ,  $g(x) = \sum_{k=0}^{\infty} b_k x^k$ . Then

$$(1) f(x) + g(x) = \sum_{k=0}^{\infty} (a_k + b_k) x^k$$

$$(2) \alpha \cdot f(x) = \sum_{k=0}^{\infty} \alpha \cdot a_k x^k \quad \alpha \in \mathbb{R}$$

$$(3) x \cdot f'(x) = \sum_{k=0}^{\infty} k \cdot a_k x^k$$

$$(4) f(\alpha x) = \sum_{k=0}^{\infty} \alpha^k \cdot a_k x^k$$

$$(5) f(x)g(x) = \sum_{k=0}^{\infty} \left( \sum_{j=0}^k a_j b_{k-j} \right) x^k$$

### The extended binomial coefficient

**【Definition 2】** Let  $u$  be a real number and  $k$  a nonnegative integer. Then the *extended binomial coefficient* is defined by

$$\binom{u}{k} = \begin{cases} u(u-1)\cdots(u-k+1)/k! & \text{if } k > 0 \\ 1 & \text{if } k = 0 \end{cases}$$

### The extended Binomial Theorem

**【Theorem 2】** Let  $x$  be a real number with  $|x| < 1$  and let  $u$  be a real number. Then

$$(1+x)^u = \sum_{k=0}^{\infty} \binom{u}{k} x^k$$

<b>TABLE 1 Useful Generating Functions.</b>	
$G(x)$	$a_k$
$(1+x)^n = \sum_{k=0}^n C(n, k)x^k$ $= 1 + C(n, 1)x + C(n, 2)x^2 + \cdots + x^n$	$C(n, k)$
$(1+ax)^n = \sum_{k=0}^n C(n, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n, 2)a^2x^2 + \cdots + a^n x^n$	$C(n, k)a^k$
$(1+x^r)^n = \sum_{k=0}^n C(n, k)x^{rk}$ $= 1 + C(n, 1)x^r + C(n, 2)x^{2r} + \cdots + x^{rn}$	$C(n, k/r)$ if $r \mid k$ ; 0 otherwise
$\frac{1-x^{n+1}}{1-x} = \sum_{k=0}^n x^k = 1 + x + x^2 + \cdots + x^n$	1 if $k \leq n$ ; 0 otherwise
$\frac{1}{1-x} = \sum_{k=0}^{\infty} x^k = 1 + x + x^2 + \cdots$	1
$\frac{1}{1-ax} = \sum_{k=0}^{\infty} a^k x^k = 1 + ax + a^2x^2 + \cdots$	$a^k$
$\frac{1}{1-x^r} = \sum_{k=0}^{\infty} x^{rk} = 1 + x^r + x^{2r} + \cdots$	1 if $r \mid k$ ; 0 otherwise
$\frac{1}{(1-x)^2} = \sum_{k=0}^{\infty} (k+1)x^k = 1 + 2x + 3x^2 + \cdots$	$k+1$
$\frac{1}{(1-x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)x^k$ $= 1 + C(n, 1)x + C(n+1, 2)x^2 + \cdots$	$C(n+k-1, k) = C(n+k-1, n-1)$
$\frac{1}{(1+x)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)(-1)^k x^k$ $= 1 - C(n, 1)x + C(n+1, 2)x^2 - \cdots$	$(-1)^k C(n+k-1, k) = (-1)^k C(n+k-1, n-1)$
$\frac{1}{(1-ax)^n} = \sum_{k=0}^{\infty} C(n+k-1, k)a^k x^k$ $= 1 + C(n, 1)ax + C(n+1, 2)a^2x^2 + \cdots$	$C(n+k-1, k)a^k = C(n+k-1, n-1)a^k$
$e^x = \sum_{k=0}^{\infty} \frac{x^k}{k!} = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \cdots$	$1/k!$
$\ln(1+x) = \sum_{k=1}^{\infty} \frac{(-1)^{k+1}}{k} x^k = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \cdots$	$(-1)^{k+1}/k$

#### 4. Counting with Generating Functions

Generating functions can be used to solve a wide variety of counting problems, such as

- Count the number of combinations from a set when repetition is allowed and additional constraints exist.
- Count the number of permutations

## 5. Using Generating Functions to Solve Recurrence Relations

The Methods of Solving Recurrence Relations

- Iterative approach
- Use a systematic way to solve an important class of recurrence relations
- Generating functions    *Method:* (1) Use the recurrence relation to find the generating function of this sequence; (2)  $G(x) \leftrightarrow a_n$

## 6. Proving Identities via Generating Functions

The method of proving combinatorial identities:

- Use combinatorial proofs
- Use generating functions

### 8.5 Inclusion-Exclusion

For the union of three finite sets:

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$$

For the union of  $n$  finite sets:

$$|A_1 \cup A_2 \cup \dots \cup A_n| = \sum_{i=1}^n |A_i| - \sum_{1 \leq i < j \leq n} |A_i \cap A_j| + \sum_{1 \leq i < j < k \leq n} |A_i \cap A_j \cap A_k| - \dots + (-1)^{n+1} |A_1 \cap A_2 \cap \dots \cap A_n|$$

### 8.6 Applications of Inclusion-Exclusion

#### 1. An alternative form of inclusion-exclusion

Problems that ask for the number of elements in a set that have none of  $n$  properties  $P_1, P_2, \dots, P_n$ .

Let  $A_i$  be the subset containing the elements that have property  $P_i$ .

$N(P_1' P_2' \dots P_n')$  ---- The number of elements with none of the properties  $P_1, P_2, \dots, P_n$ .

From the inclusion-exclusion principle, we see that

$$N(P_1' P_2' \dots P_n') = N - |A_1 \cup A_2 \cup \dots \cup A_n| = N - \sum_{1 \leq i \leq n} N(P_i) + \sum_{1 \leq i < j \leq n} N(P_i P_j) - \sum_{1 \leq i < j < k \leq n} N(P_i P_j P_k) + \dots + (-1)^n N(P_1 P_2 \dots P_n)$$

#### 2. The sieve of Eratosthenes

The number of integers not exceeding 100 (and greater than 1) that are divisible by all the primes in a subset of  $\{2, 3, 5, 7\}$  is  $\lfloor 100/N \rfloor$ , where  $N$  is the product of the primes in this subset.

#### 3. The number of onto functions

Theorem:

Let  $m$  and  $n$  be positive integers with  $m \geq n$ . Then, there are

$$n^m - C(n, 1)(n-1)^m + C(n, 2)(n-2)^m - \dots + (-1)^{n-1} C(n, n-1) \cdot 1^m$$

onto functions from a set with  $m$  elements to a set with  $n$  elements.

Applications:

- Assign  $m$  different jobs to  $n$  different employees if every employee is assigned at least one job.
- Distribute  $m$  different toys to  $n$  different children such that each child gets at least one toy.

#### 4. Derangement

A *derangement* is a permutation of objects that leaves no object in its original position.

Theorem: The number of derangements of a set with  $n$  elements is

$$D_n = n! \left[ 1 - \frac{1}{1!} + \frac{1}{2!} - \frac{1}{3!} + \dots + (-1)^n \frac{1}{n!} \right]$$