

# **Lesson 2**

# **Computability**

Math 574 - Topics in Logic  
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Jan Reimann

2-2

# Turing Machines

## Finite Automata:

- model of restricted computational power

$$L = \{0^n 1^n : n \geq 0\}$$

- FA have fixed, finite amount of memory

## Turing Machine (TM)

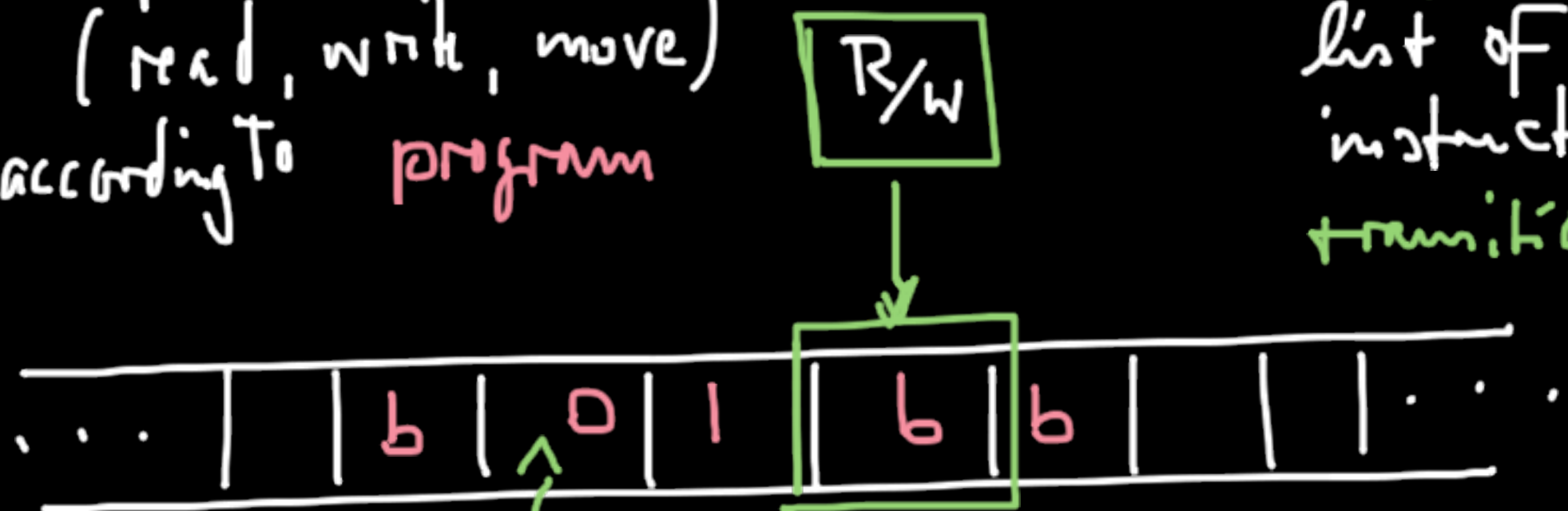
- finitely many states ( $\rightarrow$  program)

+ memory/work tape on which machine can read & write



performs operations  
(read, write, move)  
according to **program**

**program:**  
list of  
instructions  
+ transition function



tape

cell

current cell

holds entry from  
a finite alphabet  $A$

or a special symbol " $b$ "  $\notin A$   
**BLANK**

# Formal Definition of TM

- set of states  $Q$  (finite)

- $q_0 \in Q$  initial state

- $q_f \in Q$  halting state

- Alphabet (finite)  $A$

Special ("blank") symbol " $b$ "  $\notin A$

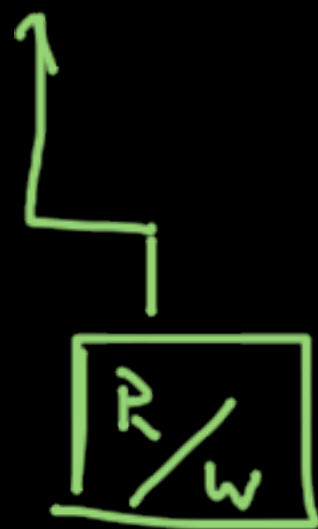
- transition function

$$\delta: Q \times A \cup \{b\} \rightarrow A \cup \{b\} \times \overset{\substack{\text{direction} \\ \swarrow}}{D} \times Q$$

$$D = \{L, R, S\}$$

## Running a TM

• input  $w = w_0 \dots w_{n-1}$ ,  $w_i \in A$



Set the current

state =  $q_0$   
initial state

initial configuration of TM

## Computation Step

- Given: Configuration ("snapshot")

current symbol  
 $x$

current state  
 $q$

- Look up transition function

$$\delta(q, x) = (y, d, p)$$

$\uparrow$   $\uparrow$   $\nwarrow$  new state  
 $\in A \times \{b\}$   $L, R \text{ or } S$





current state = 3

$$\delta(3, 0) = (1, R, 5)$$



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new current state = 5

## Halting & Output

- TM halts if we reach state  $q_F$

- If this happens, TM  
outputs the string (over  $A$ )  
immediately to the right of  
current cell

output:  
011

