# Math 557 August 27

# **Problems**

### Exercise 0.1.

The set of L-terms  $\mathcal{T}^{\mathcal{L}}$  is defined as the *smallest* set (with respect to  $\subseteq$ ) containing variables and constant symbols and that is closed under function symbol application.

- Why do this set exist?
- Does a similar argument work for the set of  $\mathcal{L}$ -formulas?
- Find an example of a (non-empty) family of sets and a property P (applicable to the sets in the family) such that a  $\subseteq$ -smallest set with property P does not exist.

#### Exercise 0.2.

Prove the **Readbility Lemma**: For any term t, exactly one of the following holds.

- 1. There exists a  $v \in \mathcal{V}$  such that  $t \equiv v$ .
- 2. There exists a  $c \in \mathcal{C}^{\mathcal{L}}$  such that  $t \equiv c$ .
- 3. There exists a number  $n \in \mathbb{N}$ , a function symbol  $f \in F_n^{\mathcal{L}}$ , and terms  $t_1, \dots, t_n$  such that  $t \equiv ft_1 \dots t_n$ .

## Exercise 0.3.

Show that no proper initial segment of a term is a term. (Use induction on the length of a term.)

Deduce **unique readability**: The decomposition in (3.) above is unique.

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Prove unique readability for the set of  $\mathcal{L}$ -formulas.

### Exercise 0.4.

The set C is defined as the smallest set of finite strings over the alphabet  $\{\times, \circ\}$  such that

- (C1)  $\times \in C$ ,
- $(C2) \circ \in C$ ,
- (C3) if  $s, t \in C$  and the last symbol in s differs from the first symbol in t, then  $st \in C$  (where st is the concatenation of s and t).

Formulate the property of unique readability for the system (C1)-(C3) and argue that it does not hold for this system.

#### Exercise 0.5.

Let  $\mathcal{L} = (0, <)$  be a language with one constant symbol 0 and one binary relation symbol <. In the following, for sake of (human) readability, we write  $0 < v_1$  instead of  $< 0v_1$ .

Consider the expression

$$\phi \equiv (\exists v_2 \ v_1 < v_2 \land \exists v_1 \ v_2 < 0).$$

- Is this a correct  $\mathcal{L}$ -formula? Characterize every variable occurrence in  $\phi$  as free or bound.