Mutual Exclusion

Companion slides for The Art of Multiprocessor Programming by Maurice Herlihy & Nir Shavit

Warning

- · You will never use these protocols
 - Get over it
- You are advised to understand them
 - The same issues show up everywhere
 - Except hidden and more complex

Locks (Mutual Exclusion)

```
public interface Lock {
  public void lock();
  public void unlock();
}
```

Locks (Mutual Exclusion)

```
public interface Lock {

public void lock();

public void unlock();
}
```

Locks (Mutual Exclusion)

```
public interface Lock {

public void lock(); acquire lock

public void unlock(); release lock
}
```

Deadlock-Free



- If some thread calls lock()
 - And never returns
 - Then other threads must complete lock() and unlock() calls infinitely often
- System as a whole makes progress
 - Even if individuals starve

Starvation-Free



- If some thread calls lock()
 - It will eventually return
- Individual threads make progress

Two-Thread vs *n* -Thread Solutions

- Two-thread solutions first
 - Illustrate most basic ideas
 - Fits on one slide
- Then n-Thread solutions

Two-Thread Conventions

```
class ... implements Lock {
 // thread-local index, 0 or 1
  public void lock() {
    int i = ThreadID.get();
    int j = 1 - i;
```

Two-Thread Conventions

```
class ... implements Lock {
    ...
    // thread-local index, 0 or 1
    public void lock() {
        int i = ThreadID.get();
        int j = 1 - i;
    ...
}
```

Henceforth: i is current thread, j is other thread

LockOne

LockOne

LockOne

```
class LockOne implements Lock {
private boolean[] flag =
                         new boolean[2];
public void lock() {
 flag[i] = true:
 while (flag[j]) {}
                            Set my flag
                      Wait for other
                      flag to go false
```

Deadlock Freedom

- LockOne Fails deadlock-freedom
 - Concurrent execution can deadlock

```
flag[i] = true; flag[j] = true;
while (flag[j]){} while (flag[i]){}
```

- Sequential executions OK

```
public class LockTwo implements Lock {
  private int victim;
  public void lock() {
    victim = i;
    while (victim == i) {};
  }
  public void unlock() {}
}
```

```
public class LockTwo implements Lock {
  private int victim;
  public void lock() {
    victim = i;
    while (victim == i) {};
  }
  public void unlock() {}
}
```

```
public class LockTwo implements
private int victim;
public void lock() {
    victim = i;
    while (victim == i) {};
}

public void unlock() {}
}
```

```
public class Lock2 implements Lock {
  private int victim;
  public void lock() {
    victim = i;
    while (victim == i) {};
  }
  public void unlock() {}
}
```

LockTwo Claims

- · Satisfies mutual exclusion
 - If thread i in CS
 - Then victim == j
 - Cannot be both 0 and 1
- public void LockTwo() {
 victim = i;
 while (victim == i) {};
 }
- Not deadlock free
 - Sequential execution deadlocks
 - Concurrent execution does not

```
public void lock() {
  flag[i] = true;
  victim = i;
  while (flag[j] && victim == i) {};
}
public void unlock() {
  flag[i] = false;
}
```

Announce I'm

```
public void lock() {
  flag[i] = true;
  victim = i;
  while (flag[j] && victim == i) {};
}
public void unlock() {
  flag[i] = false;
}
```

```
Announce I'm
                          interested
flag[i] = true;
                          Defer to other
victim
while (flag[j] && victim == i) {};
public void unlock() {
flag[i] = false;
```

```
Announce I'm
                           interested
                            Defer to
flag[i] = true;
Victim
                             other
while (flag[j] && victim == i) {}:
                     Wait while other
public void unlock()
                       interested & I'm
flag[i] = false;
                          the victim
```

```
Announce I'm
                          interested
                           Defer to
        = true;
victim
                             other
while (flag[j] && victim == i) {}:
                     Wait while other
           unlock()
flag[i] = false;
                       interested & I'm
                          the victim
        No longer
        interested
```

Mutual Exclusion

```
public void lock() {
  flag[i] = true;
  victim = i;
  while (flag[j] && victim == i) {};
```

- If thread 0 in critical section,
 - flag[0]=true,
 - victim = 1

- If thread 1 in critical section,
 - flag[1]=true,
 - victim = 0

Cannot both be true

Deadlock Free

```
public void lock() {
    ...
    while (flag[j] && victim == i) {};
```

- Thread blocked
 - only at while loop
 - only if it is the victim
- · One or the other must not be the victim

Starvation Free

 Thread i blocked only if j repeatedly reenters so that

```
flag[j] == true and
victim == i
```

- When j re-enters
 - it sets victim toj.
 - So i gets in

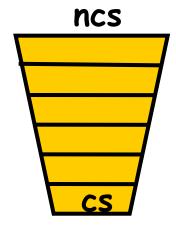
```
public void lock() {
   flag[i] = true;
   victim = i;
   while (flag[j] && victim == i) {};
}

public void unlock() {
   flag[i] = false;
}
```

The Filter Algorithm for *n*Threads

There are n-1 "waiting rooms" called levels

- · At each level
 - At least one enters level
 - At least one blocked if many try



· Only one thread makes it through

```
class Filter implements Lock {
   int[] level; // level[i] for thread i
   int[] victim; // victim[L] for level L
 public Filter(int n) {
                                                     n-1
     level = new int[n];
                             level 0
     victim = new int[n];
     for (int i = 1; i < n; i++) {
         level[i] = 0;
     }}
               Thread 2 at level 4
                                                n-1
                  Art of Multiprocessor
                                                  72
                    Programming
```

```
class Filter implements Lock {
  public void lock(){
    for (int L = 1; L < n; L++) {
      level[i] = L;
      victim[L] = i;
      while ((\exists k != i level[k] >= L) \&\&
             victim[L] == i );
    }}
  public void unlock() {
    level[i] = 0;
  }}
```

```
class Filter implements Lock {
  ...
    for (int L = 1; L < n; L++) {
      level[i]
      victim[L] = i;
      while ((\exists k != i) level[k
             victim[L] == i)
    }}
  public void release(int i) {
    level[i] = 0;
                            One level at a
                                   time
```

```
class Filter implements Lock {
  0.00
  public void lock() {
    for (int L = 1; L < n; L++) {
      <u>level[il</u>
      victim[L]
      while ((\exists k != 1)
                         | evel[k] >= L) &&
              victim[L] ==
                                    Announce
    }}
                                   intention to
  public void release(int i)
    level[i] = 0;
                                  enter level L
```

```
class Filter implements Lock {
                   int level[n];
                   int victim[n];
                    public void lock() {
                                      for (int L = 1; L < n; L++) {
                                                   victim[L] = i:
                                                                                                                                                                                                                   location l
                                                        while ((3 k
                                                                                                                           victim[L
                                   }}
                                                                                                                                                                                                                                                                                       Give priority to
                   public void release(int i)
                                                                                                                                                                                                                                                                                                   anyone but me
                                      level[i] = 0;
                   }}
```

```
Wait as long as someone else is at same
   or higher level, and I'm designated
     while ((\exists k != i) level[k] >= L) \&\&
            victim[L] == i);
 public void release(int i) {
   level[i] = 0;
```

```
class Filter implements Lock {
  int level[n];
  int victim[n];
  public void lock() {
    for (int L = 1; L < n; L++) {
      level[i] = L;
      victim[L] = i;
     while ((\exists k != i) level[k] >= L) \&\&
             victim[L] == i);
Thread enters level L when it completes
                  the loop
```

No Starvation

- · Filter Lock satisfies properties:
 - Just like Peterson Alg at any level
 - So no one starves
- But what about fairness?
 - Threads can be overtaken by others

- Provides First-Come-First-Served
- · How?
 - Take a "number"
 - Wait until lower numbers have been served
- · Lexicographic order
 - -(a,i) > (b,j)
 - If a > b, or a = b and i > j

```
class Bakery implements Lock {
   boolean[] flag;
   Label[] label;
  public Bakery (int n) {
    flag = new boolean[n];
    label = new Label[n];
    for (int i = 0; i < n; i++) {
       flag[i] = false; label[i] = 0;
```

```
class Bakery implements Lock {
   boolean[] flag;
   Label[] label;
  public Bakery (int n) {
    flag = new boolean[n];
    label = new Label[n];
   for (int i = 0; i < n; i++) {
       flag[i] = false; label[i] = 0;
```

Take increasing

```
class Bakery implements Lock { Someone is ... interested public void lock() { flag[i] = true; label[i] = max(label[0], ..., label[n-1])+1; while (∃k flag[k] && (label[i],i) > (label[k],k)); }
```

```
class Bakery implements Lock {
  boolean flag[n];
                        Someone is
  int label[n];
                         interested
 public void lock() {
  flag[i] = true;
              \max(\text{label}[0], ..., \text{label}[n-1])+1;
           && (label[i],i) > (label[k],k));
```

With lower (label,i) in lexicographic order

```
class Bakery implements Lock {
    No longer interested

public void an lock {
    flag[i] = false;
    }
}

labels are always increasing
```

No Deadlock

- There is always one thread with earliest label
- Ties are impossible (why?)

- Suppose A and B in CS together
- Suppose A has earlier label
- When B entered, it must have seen
 - flag[A] is false, or
 - label[A] > label[B]

- · Labels are strictly increasing so
- B must have seen flag[A] == false

- · Labels are strictly increasing so
- B must have seen flag[A] == false
- · Labeling_B \rightarrow read_B(flag[A]) \rightarrow write_A(flag[A]) \rightarrow Labeling_A

- · Labels are strictly increasing so
- B must have seen flag[A] == false
- · Labeling_B \rightarrow read_B(flag[A]) \rightarrow write_A(flag[A]) \rightarrow Labeling_A
- Which contradicts the assumption that A has an earlier label



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