

1.

local B in

```
thread    // S1
  B=true   // T1
end
thread    // S2
  B=false  // T2
end
if B then  // S3
  skip Browse B // S3.1
end
end
```

S1 T1 S2 T2 S3 ---- displays nothing

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S1 T1 S2 S3 S3.1 T2 ----- displays true

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S1 S2 T1 T2 S3 ---- displays nothing

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S1 S2 T2 S3 T1 ---- displays nothing

// thread2 ran first, and B = false, so B is bound to false. and then thread1 ran after, but B is already bound to false, so B can't be bound to true. so, unification error happens.

S2 S1 T1 S3 S3.1 T2 ----- displays true

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S1 S2 T1 S3 S3.1 T2 ----- displays true

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S2 S1 T2 S3 T1 ----- displays nothing

// thread2 ran first, and B = false, so B is bound to false. and then thread1 ran after, but B is already bound to false, so B can't be bound to true. so, unification error happens.

S1 S2 T2 T1 S3 S3.1 ----- displays true

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S2 T2 S1 T1 S3 S3.1 ----- displays true

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S2 S1 T2 T1 S3 S3.1 ----- displays true

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S2 S1 T1 T2 S3 S3.1 ----- displays true

// thread1 ran first, and B = true, so B is bound to true. and then thread2 ran after, but B is already bound to true, so B can't be bound to false. so, unification error happens.

S2 T2 S1 S3 T1 ----- displays nothing

// thread2 ran first, and B = false, so B is bound to false. and then thread1 ran after, but B is already bound to false, so B can't be bound to true. so, unification error happens.

2.

```
local X Y T in
  thread Y = X end
  X = 3
  skip Browse Y
end
```

```
local T1 T2 in
  T2 = thread 3 end
  T1 = thread (4+3) end
  skip Browse T2
  skip Browse T1
end
```

// In infinity case, the output is Y: unbound, T2: unbound, T1: unbound. The reason is those threads are executed after Browse.

// In finite 1 case, the output is Y:3 T2:3 T1: unbound, the thread runs after "X = 3" so we have Y = X = 3. Then for the "T1 T2" one,

// thread of T2 runs first and then "skip Browse T2", "skip Browse T1", finally, thread of T1. So, we have T2 = 3, T1 = unbound.

3.

```
local Z in
  Z = 3
  thread local X in
    X = 1
    skip Browse X
    skip Browse X
    skip Basic
```

skip Browse X

skip Browse X

skip Basic

skip Browse X

end

end

thread local Y in

Y = 2

skip Browse Y

skip Basic

skip Browse Y

skip Browse Y

skip Browse Y

skip Basic

skip Browse Y

end

end

skip Browse Z

skip Browse Z

skip Browse Z

skip Basic

skip Browse Z

skip Browse Z

end

4. With quantum of 3, B is bounded to true, but with quantum of 5, B become unbound. The reason why is because When it is large enough, thread takes a lower priority, causing the if statement to be executed first, so B becomes unbound.

5.

(a)

function	Maximum X	Time (s)	Memory (bytes)
fib1_sugar	15	37.96	13,435,508,912
fib2_sugar	1800	57.24	20,367,039,712
fib1_thread	13	25.71	3,371,539,608

For fib1_sugar, the

maximum X within one minute is 15. and

for fib2_sugar, the maximum X is 1800, for fib1_thread, the

maximum X is 13.

Why is that result is because, fib1_sugar and fib1_thread

are using recursion without tail call, and there are lots of repetitive calculations

as the progress going through. fib2_sugar is using recursion with tail call and accumulator.

so, there is no repetitive calculations in fib2_sugar. Thus fib2_sugar is more efficient than fib1_sugar and fib1_thread

(b)

num_of_threads n

| n == 0 = 0

| n == 1 = 0

| n == 2 = 1

| otherwise = num_of_threads(n-2)+num_of_threads(n-1)+2

Part 2

1.

```
OddFilter = proc {$ P Out}
  case P
  of nil then Out = nil
  [] '|' (1:X 2: Xr) then
    if ({Mod X 2} == 0) then T in
      Out = (X|T)
      {OddFilter Xr T}
    else
      {OddFilter Xr Out}
    end
  end
end
end
```

2.

```
Consumer = fun {$ P} in
  case P
  of nil then 0
  [] '|' (1:X 2: Xr) then
    (X + {Consumer Xr})
  end
end
end
```

3.

local Producer OddFilter Consumer N L P F S in

thread

Producer = proc {\$ N Limit Out}

if (N<Limit) then T N1 in

Out = (N|T)

N1 = (N + 1)

{Producer N1 Limit T}

else Out = nil

end

end

end

thread

OddFilter = proc {\$ P Out}

case P

of nil then Out = nil

[] '|' (1:X 2: Xr) then

if ({Mod X 2} == 0) then T in

Out = (X|T)

{OddFilter Xr T}

else

{OddFilter Xr Out}

end

end

end

end

thread

Consumer = fun {\$ P} in

case P

of nil then 0

[] ' |(1:X 2:Xr) then

(X + {Consumer Xr})

end

end

end

// Example Testing

N = 0

L = 101

{Producer N L P} // [0 1 2 ... 100]

{OddFilter P F} // [0 2 4 .. 100]

{Consumer P S}

skip Browse F

skip Browse S

end