

Algebra

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Lecture 1: Introduction and vector operation review

Contents I

1. Linear algebra: ~ 10-11 lectures

1.1 Review: Vector operations in 2 and 3 dimensional spaces

1.2 System of linear equations (LES)

- ▶ Solution existence of a LES
- ▶ LES and associated augmented matrices
- ▶ Gauss and Gauss-Jordan elimination method

1.3 Matrix Algebra

- ▶ Terminologies
- ▶ Operations with matrices
- ▶ The inverse of a matrix
- ▶ Partitioned matrix and matrix factorization (self-study)

1.4 Determinants

1.5 Vector spaces

1.6 Linear transformations

1.7 Coordinate systems and basic transition matrix

Contents II

2. General algebra: \sim 9-10 lectures

2.1 Propositional logic: Lehman et.al.'s (Chapter 3), or Rosen's (Chapter 1).

2.2 Set theory review: Rosen's textbook

2.3 Relations: Rosen's textbook (chapter 9)

- ▶ Binary relations
- ▶ Popular properties on relations: reflexivity, symmetricity, transitivity, anti-symmetricity.
- ▶ Closures of a relation

2.4 Group theory

- ▶ Basic definitions
- ▶ Generators for a group and cyclic groups
- ▶ Cosets and Lagrange's theorem
- ▶ Homomorphism and isomorphism
- ▶ Quotient groups

References

► Linear algebra

1. D. Lay, Linear Algebra and Its Applications, Pearson New International Edition, Pearson, 2014 (primary).
2. Gilbert Strang, Linear Algebra and Its Applications, Fourth edition, Brooks/Cole Cengage Learning, 2006.
3. Serge Lang, Introduction to Linear Algebra, Second edition, Springer.

► General algebra

1. Kenneth Rosen, Discrete Mathematics and its applications, McGraw Hill education, 2013 ([Logic, relations, group](#)).
2. Eric Lehman, F.T. Leighton, and A. R. Meyer, Mathematics for Computer Science, 2017 ([Logic](#)).
3. Joseph A. Gallian, Contemporary Abstract Algebra, Cengage learning, 2017 ([Groups](#))

Learning outcomes

1. Acquired with essential concepts, structures and methods of propositional algebra, general algebra and linear algebra. Particularly, well-acquired with basic algebraic structures necessary for the comprehension of formal structures in CS;
2. Have the ability to independently develop abstract concepts and to acquire basis techniques of algebra;
3. Acquired for analytical thinking, development of methodological expertise, handling abstract methods, structures and models.

Attendance and exam

1. To enable to pass the final exam, you are recommended to
 - ▶ attend the class at least 70% number of contact hours
 - ▶ do exercises/homework as much as possible
2. To get further point of views as well as to understand applications of the subject in the computer science field, please spend your time for reading textbooks.
3. Exam duration: 90 minutes with 7-12 questions. You are allowed to bring a two-sided A4 written with any contents. Pocket calculator is Not allowed.

Linear Algebra

Lecture 1: Vector operations in 2,3-dimensional spaces

Outline

- ▶ Geometric representation of vectors
- ▶ Coordinate representation of vectors
- ▶ Operations on vectors
- ▶ Inner/dotted product
- ▶ Length/norm of vectors
- ▶ Orthogonal vectors
- ▶ Angle between two vectors

Graphic representation of vectors

Definition (Vector)

A , B are points in the plane (or in the space). The vector \vec{AB} is the directed segment from A to B .

Remark

Some variables, such as length, temperature, time are totally determined by just one number. These variables are so-called scalars. Other variables, such as speed, force need both a magnitude and a direction to be determined completely. These variables are vectors.

Examples

A company produces two products A and B . For one item of a product, the production cost is a 3-tuple vector of material cost, labor cost, and advertisement cost. For instance, the production costs of product A and B are given $c_A = (0.45, 0.25, 0.15)$ and $c_B = (0.4, 0.3, 0.15)$ respectively.

Definition (Equality of vectors)

Two vectors are equal if they have the same direction and the same length.

Notations

- ▶ Vectors are parallel: $\vec{AB} \uparrow\uparrow \vec{CD}$
- ▶ Vectors are antiparallel: $\vec{AB} \uparrow\downarrow \vec{CD}$
- ▶ Vectors: \vec{a}, \vec{b}, \dots
- ▶ Scalars: $\alpha, \beta, \gamma, \dots, \lambda, \mu, \dots$

Definition (Sum of the two vectors \vec{a} and \vec{b})

If we shift \vec{b} in a parallel manner, such that its origin corresponds to the end point of \vec{a} , then $\vec{a} + \vec{b}$ starts at the origin of \vec{a} and ends at the end of \vec{b} .

$$\vec{a} + \vec{b} = \vec{AB} + \vec{BC} = \vec{AC}$$

Theorem

1. *Commutative law:*

$$\vec{a} + \vec{b} = \vec{b} + \vec{a}$$

2. *Associative law:*

$$(\vec{a} + \vec{b}) + \vec{c} = \vec{a} + (\vec{b} + \vec{c})$$

Proof of 1: On the whiteboard

Proof of 2: On the whiteboard

$$\vec{a} + \vec{b} = \vec{AC}$$

$$(\vec{a} + \vec{b}) + \vec{c} = \vec{AC} + \vec{CD} = \vec{AD}$$

$$\vec{b} + \vec{c} = \vec{BD}$$

$$\vec{a} + (\vec{b} + \vec{c}) = \vec{AB} + \vec{BD} = \vec{AD} = (\vec{a} + \vec{b}) + \vec{c}$$

Thus:

$$(\vec{a} + \vec{b}) + \vec{c} = \vec{a} + (\vec{b} + \vec{c})$$

Definition (Special vectors)

1. Zero vector:

$$\vec{0} = \vec{AA}$$

2. Opposite vector to $\vec{a} = \vec{AB}$:

$$-\vec{a} = \vec{BA}$$

Remark

$$\vec{a} + (-\vec{a}) = \vec{AB} + \vec{BA} = \vec{0}$$

Definition

Let \vec{a}, \vec{b} be vectors. Then

$$\vec{a} - \vec{b} = \vec{a} + (-\vec{b})$$

Definition (Multiplication with a scalar)

Given a vector \vec{a} and a scalar λ . Then, the multiplication \vec{a} with a scalar λ is a vector with the following properties:

1. Length $(\lambda \cdot \vec{a}) = |\lambda| \cdot \text{length}(\vec{a})$
2. For $\lambda > 0$ is $\lambda \cdot \vec{a} \uparrow \uparrow \vec{a}$
3. For $\lambda < 0$ is $\lambda \cdot \vec{a} \uparrow \downarrow \vec{a}$
4. For $\lambda = 0$ is $\lambda \cdot \vec{a} = \vec{0}$

Example: On the whiteboard.

Theorem

Let \vec{a}, \vec{b} be vectors and $\lambda, \mu \in \mathbb{R}$. Then $\lambda \cdot \vec{a}$ has the following properties:

1. $(\lambda + \mu) \cdot \vec{a} = \lambda \cdot \vec{a} + \mu \cdot \vec{a}$
2. $(\lambda \cdot \mu) \cdot \vec{a} = \lambda \cdot (\mu \cdot \vec{a})$
3. $\lambda \cdot (\vec{a} + \vec{b}) = \lambda \cdot \vec{a} + \lambda \cdot \vec{b}$

Coordinate representation of vectors

Remark

Even if no direction is required to describe something, it can still be useful to summarize numbers in a vector as the next example will show.

Definition (Vector)

An n -dimensional vector \vec{x} is given by

$$\vec{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$

with $x_1 \in \mathbb{R}, \dots, x_n \in \mathbb{R}$.

Example: Coordinate representation of vectors in \mathbb{R}^3

To describe vectors, we consider a Cartesian coordinate system in the space through the origin $O = (0, 0, 0)$ and the unit points $E_1 = (1, 0, 0)$, $E_2 = (0, 1, 0)$ and $E_3 = (0, 0, 1)$ on the axes. The vectors determined by the unit points are called unit vectors:

$$\begin{aligned}\vec{e}_1 &= \vec{OE}_1 \\ \vec{e}_2 &= \vec{OE}_2 \\ \vec{e}_3 &= \vec{OE}_3\end{aligned}$$

Remark

The vector $\vec{a} = \vec{OA}$ to the point $A = (a_1, a_2, a_3)$ can we then express as a unique linear combination of the unit vectors:

$$\vec{a} = a_1 \cdot \vec{e}_1 + a_2 \cdot \vec{e}_2 + a_3 \cdot \vec{e}_3$$

Remark (Column notation)

$$\vec{a} = a_1 \cdot \vec{e}_1 + a_2 \cdot \vec{e}_2 + a_3 \cdot \vec{e}_3 = \begin{pmatrix} a_1 \\ a_2 \\ a_3 \end{pmatrix}$$

Definition (Arithmetic operations with vectors)

Let

$$\vec{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \text{ and } \vec{y} = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix}$$

► Addition and subtraction

We have

$$\vec{x} \pm \vec{y} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} \pm \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix} = \begin{pmatrix} x_1 \pm y_1 \\ x_2 \pm y_2 \\ \vdots \\ x_n \pm y_n \end{pmatrix}$$

► Multiplication with a scalar

For $\lambda \in \mathbb{R}$ we have

$$\lambda \cdot \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix} = \begin{pmatrix} \lambda \cdot x_1 \\ \lambda \cdot x_2 \\ \vdots \\ \lambda \cdot x_n \end{pmatrix}$$

► Examples (3)

Remark

As we combine different vectors with each other, often linear combinations are created.

Definition (Linear combination)

Let $\vec{x}_1, \dots, \vec{x}_m$ be n -dimensional vectors and $c_1 \in \mathbb{R}, \dots, c_m \in \mathbb{R}$.
Then

$$c_1 \cdot \vec{x}_1 + c_2 \cdot \vec{x}_2 + \dots + c_m \cdot \vec{x}_m = \sum_{i=1}^m c_i \cdot \vec{x}_i$$

is a linear combination of the vectors $\vec{x}_1, \dots, \vec{x}_m$.

Dot/Inner product

Definition

1. In 2-dimensional space, let $\vec{a} = (a_1, a_2)$ and $\vec{b} = (b_1, b_2)$. We define their dot product of \vec{a} and \vec{b} to be

$$\begin{pmatrix} a_1 \\ a_2 \end{pmatrix} \cdot \begin{pmatrix} b_1 \\ b_2 \end{pmatrix} = a_1 b_1 + a_2 b_2$$

2. In 3-dimensional space, their dot or inner product is defined to be

$$\begin{pmatrix} a_1 \\ a_2 \\ a_3 \end{pmatrix} \cdot \begin{pmatrix} b_1 \\ b_2 \\ b_3 \end{pmatrix} = a_1 b_1 + a_2 b_2 + a_3 b_3$$

3. Generally, in n -dimensional space, their dot or inner product is defined to be:

$$\vec{a} \cdot \vec{b} = a_1 b_1 + a_2 b_2 + \cdots + a_n b_n$$

Examples: On the whiteboard.

Remark

We can only compute the dot product of two vectors if they have the same dimensions. For example, it is not possible to compute the dot product as follows:

$$\vec{x} \cdot \vec{y} = \begin{pmatrix} 1 \\ 2 \end{pmatrix} \cdot \begin{pmatrix} 2 \\ 3 \\ 4 \end{pmatrix}.$$

Theorem

Let $\vec{a}, \vec{b}, \vec{c}$ be vectors and $\lambda \in \mathbb{R}$.

1. $(\lambda \cdot \vec{a}) \cdot \vec{b} = \lambda \cdot (\vec{a} \cdot \vec{b})$
2. $\vec{a} \cdot \vec{b} = \vec{b} \cdot \vec{a}$
3. $\vec{a} \cdot (\vec{b} + \vec{c}) = \vec{a} \cdot \vec{b} + \vec{a} \cdot \vec{c}$

► Examples (2)

Definition (Orthogonal vectors)

Two vectors \vec{a} and \vec{b} are called **orthogonal** if $\vec{a} \cdot \vec{b} = 0$.

Examples: On the whiteboard.

The norm or magnitude of a vector

Definition

Let \vec{a} be a vector in \mathbb{R}^n . We define the norm or magnitude of \vec{a} , and denote by $|\vec{a}|$, the number

$$|\vec{a}| = \sqrt{\vec{a} \cdot \vec{a}}.$$

Examples:

1. For

$$\vec{a} = \begin{pmatrix} \lambda_1 \\ \lambda_2 \end{pmatrix},$$

the norm of \vec{a} is

$$|\vec{a}| = \sqrt{\lambda_1^2 + \lambda_2^2}$$

2. For $\vec{a} = (\lambda_1, \lambda_2, \lambda_3)$, the norm of \vec{a} is

$$|\vec{a}| = \sqrt{\lambda_1^2 + \lambda_2^2 + \lambda_3^2}$$

The norm and geometric length

Notes: When $n = 2$ and $n = 3$, then the definition of norm is compatible with the geometric length by using Pythagoras theorem.

Proof 1: Applying the Theorem of Pythagoras

$$a^2 = \lambda_1^2 + \lambda_2^2$$

Proof 2:

$$\begin{aligned} |\vec{a}|^2 &= |\vec{OQ}|^2 + |\vec{QP}|^2 = |\vec{OR}|^2 + |\vec{RQ}|^2 + |\vec{QP}|^2 = \\ &\lambda_1^2 + \lambda_2^2 + \lambda_3^2 \end{aligned}$$

Definition (Distance between two points in the space)

The distance between two points, $A = (a_1, a_2, a_3)$ and $B = (b_1, b_2, b_3)$, in the space is given by:

$$d = |\vec{A} - \vec{B}| = \sqrt{(b_1 - a_1)^2 + (b_2 - a_2)^2 + (b_3 - a_3)^2}$$

► Example

General Pythagoras theorem

Theorem

Let \vec{a} and \vec{b} be vectors in \mathbb{R}^n . Then, vectors \vec{a} and \vec{b} are perpendicular if and only if

$$|\vec{a} + \vec{b}|^2 = |\vec{a}|^2 + |\vec{b}|^2.$$

Angle between two vectors

Remark

How can we calculate the angle φ between two vectors \vec{a} and \vec{b} in term of their coordinates?

► Example

Theorem

Let \vec{a} and \vec{b} be vectors and φ the angle between these vectors.
Then,

$$\vec{a} \cdot \vec{b} = |\vec{a}| \cdot |\vec{b}| \cdot \cos(\varphi)$$

► Example

Remark

Let \vec{a}, \vec{b} be vectors and φ the angle between them. Then

$$\cos(\varphi) = \frac{\vec{a} \cdot \vec{b}}{|\vec{a}| \cdot |\vec{b}|}$$

Linear Algebra

Lecture 2+3: Linear equation systems

Contents

- ▶ Linear equation systems (LES)
- ▶ Matrix representation of a LES
- ▶ Row-echelon form and Gaussian elimination
- ▶ Reduced row-echelon form and Gauss-Jordan elimination

Linear equation systems

In general, if a linear relation exists between the inputs and outputs of a system, we call it a linear system.

Remark (Applications of linear equation systems)

- ▶ Rotations in space
- ▶ Demand calculations for stock holding
- ▶ Modeling of customer streams
- ▶ Flow network
- ▶ Electrical circuits
- ▶ Quantum mechanics

Definition (Linear equation)

A linear equation in n variables:

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n = b,$$

where

a_1, \dots, a_n, b : real numbers

x_1, \dots, x_n : unknown or variables

a_1, \dots, a_n : are called coefficients of the equation.

a_1 : leading coefficient or pivot if non-zero

Examples

► $2x + y - 4z = 0$

► $2x + y - 4z = 3$

Definition (Linear equation system (LES))

A system of m linear equations is called a **linear equation system**:

$$\left\{ \begin{array}{l} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\ \vdots \\ a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n = b_m \end{array} \right. \quad (1)$$

where

$a_{ij}, b_i, i = 1, \dots, m, j = 1, \dots, n$: real numbers;

x_1, \dots, x_n unknowns or variables.

a_{ij} coefficients.

Examples

$$1. \begin{cases} x + y = 3 \\ x - y = -1 \end{cases}$$

$$2. \begin{cases} x + y = 3 \\ 2x + 2y = 6 \end{cases}$$

$$3. \begin{cases} x + y = 3 \\ x + y = 1 \end{cases}$$

Solution set of a LES

- ▶ A **solution** of the linear equation system (1) in n variables is a set of numbers s_1, s_2, \dots, s_n such that for all $i = 1, 2, \dots, n$, then

$$a_{i1}s_1 + a_{i2}s_2 + \dots + a_{in}s_n = b_i.$$

- ▶ Solution set is the set of all solutions of (1).

Consistency of a LES

A system of m linear equations in n variables:

$$\left\{ \begin{array}{l} a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\ \vdots \\ a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n = b_m \end{array} \right.$$

- ▶ **consistent** if it has at least one solution;
- ▶ **inconsistent** if it has no solution;
- ▶ If $b_i = 0$ for all $i = 1, \dots, m$, then the LES is **homogeneous**, otherwise **inhomogeneous**.

Notes: Every system of linear equations has either

1. exactly one solution,
2. infinitely many solution, or
3. no solution.

Equivalent systems

- ▶ Two systems of linear equations are called **equivalent** if they have precisely the same solution set.
- ▶ **Notes:**
 - Eliminations:** We use following operations on a system of linear equations to produce an equivalent system:
 1. **Interchange** two equations
 2. **Multiply** an equation by a nonzero constant
 3. **Add** a multiple of an equation to another equation.

Example

Solve a system of linear equations

$$\begin{cases} x - 2y + 3z &= 9 \\ -x + 3y &= -4 \\ 2x - 5y + 5z &= 17 \end{cases} \quad (2)$$

Solution: Details on the whiteboard.

Coefficient matrix of a LES I

- ▶ The coefficient matrix of the LES 1 is defined by

$$A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \cdots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{pmatrix},$$

- ▶ Notes:

1. Every entry a_{ij} in a matrix is a real number;
2. A matrix with m rows and n columns is said to be of size $m \times n$;
3. If $m = n$, then the matrix is called square of order n ;
4. For a square matrix, the entries $a_{11}, a_{22}, \dots, a_{nn}$ are called the main diagonal entries.

Coefficient matrix of a LES II

- ▶ the right-hand side with the vector

$$\vec{b} = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix}$$

- ▶ Matrix form of the LES 1: $Ax = b$, where

$$A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \cdots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{pmatrix}, \quad x = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}, \quad b = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$

Augmented matrix

- ▶ Augmented matrix for the LES (1)

$$[A|\vec{b}] = \left(\begin{array}{ccc|c} a_{11} & \cdots & a_{1n} & b_1 \\ \vdots & \cdots & \vdots & \vdots \\ a_{m1} & \cdots & a_{mn} & b_m \end{array} \right)$$

- ▶ The n numbers x_1, \dots, x_n of the solution of the LES is called solution vector:

$$\vec{x} = \begin{pmatrix} x_1 \\ \vdots \\ x_m \end{pmatrix}$$

- ▶ Example: The augmented matrix of the LES (2) is

$$[A|b] = \left(\begin{array}{cccc|c} 1 & -2 & 3 & 9 & 9 \\ -1 & 3 & 0 & -4 & -4 \\ 2 & -5 & 5 & 17 & 12 \end{array} \right).$$

Gaussian elimination: Row echelon form

- ▶ It is especially easy to solve a LES, if the coefficient matrix is given in row echelon form.
- ▶ A matrix is in row echelon form if the following three conditions are fulfilled:
 1. All nonzero rows are above any rows of all zeros.
 2. The leading entry (the pivot or the leftmost non-zero entry) in a nonzero row is 1.
 3. Each leading 1 of a row is in a column to the right of the leading 1 of the row above it.
- ▶ Note: All entries in a column below a leading entry are zeros.

Examples I

1. By elimination, the augmented matrix of the LES (2) could be transformed to the row-echelon form:

$$\left(\begin{array}{ccc|c} \textcircled{1} & -2 & 3 & 9 \\ 0 & \textcircled{1} & 3 & 5 \\ 0 & 0 & \textcircled{1} & 2 \end{array} \right)$$

2. The augmented matrix

$$\left(\begin{array}{ccccc|c} 0 & \textcircled{2} & -3 & 4 & 1 & 7 \\ 0 & 0 & 0 & \textcircled{5} & 2 & 4 \\ 0 & 0 & 0 & 0 & \textcircled{-3} & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{array} \right)$$

represents the system of linear equations

Examples II

$$\begin{cases} 2y - 3z + 4w + t &= 7 \\ 5w + 2t &= -4 \\ -3t &= 1 \end{cases}$$

Back-substitution for solving a LES with a coefficient matrix in row echelon form

1. We solve the equations for the leading variables (i.e. the variables corresponding to leading entries, also called **basic variables**).
2. Starting with the last (= bottom) we substitute the calculated variables into the equations above.
3. **Free variables** (i.e. variables not corresponding to leading entries) are assigned arbitrary values.

Remark

If the rows without leading ones are contradictory, then the solution is the empty set. Hence, the LES is inconsistent.

- Examples (3)

Remark

An arbitrary LES is usually not in row echelon form. We must then first transform the system to the row echelon form. The transformation must be done in such way that the solution set remains the same as for the original system.

Remark

Two LES with the same set of solutions are called equivalent.

Remark (Elementary row operations)

- ▶ **Interchange:** Interchange two rows: $r_{ij} : R_i \leftrightarrow R_j$
- ▶ **Scaling:** Multiply all entries in a row by a nonzero constant:
 $r_i^{(k)} : kR_i \rightarrow R_i$
- ▶ **Replacement:** Replace one row by the sum of itself and a multiple of another row: $r_{ij}^{(k)} : kR_i + R_j \rightarrow R_j$.

Theorem

If the augmented matrix $(A^|\vec{b}^*)$ is transformed from $(A|\vec{b})$ by the use of a finite number of elementary row operations, then the solution sets to these two systems are equivalent. We say that the systems are row equivalent.*

Remark (Gauss elimination)

A given augmented matrix is used as a starting point: $\left(A|\vec{b}\right)$.

Using elementary row operations, this augmented matrix is transformed into an equivalent augmented matrix: $\left(A^*|\vec{b}^*\right)$ with A^* in row echelon form.

We proceed as follows:

1. We look for the leftmost column with at least one nonzero element.
2. In this column, we look for the element with the largest absolute value. We interchange this row and the first row. (Pivotsearch to obtain numerical stability).
3. If the uppermost number $a \neq 0$ is, then we multiply the first row by $\frac{1}{a}$ to create a leading 1.
4. We add a multiple of the first row to the other rows, to create zeros below the leading 1.
5. These steps are repeated for a submatrix. The submatrix is obtained, by covering the first row (or the actual first row and all rows above it). The steps are repeated until the whole matrix is in row echelon form.

Examples:

$$1. \begin{bmatrix} 0 & 0 & -2 & 0 & 8 & 12 \\ 2 & 8 & -6 & 4 & 12 & 28 \\ 2 & 4 & -5 & 6 & -5 & 4 \end{bmatrix}$$

$$2. \begin{bmatrix} 1 & -2 & 3 & 9 \\ -1 & 3 & 0 & -4 \\ 2 & -5 & 5 & 17 \end{bmatrix}$$

Gauss-Jordan elimination: The procedure for reducing a matrix to a **reduced row-echelon form** which satisfies the following:

1. it is a row-echelon form;
2. for each non-zero row, the first non-zero entry from the left is equal to **1** and it is called **pivot** or **leading 1**.
3. Every column that has a pivot/leading 1 has zeros in every position above and below its pivot.

Example: On the whiteboard.

Definition (Rank)

Given a matrix A . The number of the nonzero rows of A^* , is called **the rank** of matrix A and denoted by $r(A)$ or simply r in case A is shown.

Remark (Solutions to a LES)

Using the augmented matrix $(A^*|\vec{b}^*)$ in row echelon form, we can identify the kind of solutions that the corresponding LES has (no solution, exactly one solution, infinitely many solutions).

Back-substitution can be used to determine the set of solutions.

Theorem (The existence of solutions)

The LES with the augmented matrix $(A|\vec{b})$ is consistent if after applying the Gaussian elimination algorithm we arrive at the augmented matrix

$(A^*|\vec{b}^*)$ with:

$$b_{r+1}^* = \dots = b_m^* = 0$$

where A is an $m \times n$ matrix.

Then we have:

- ▶ $r = n \Rightarrow$ the LES has exactly one solution
- ▶ $r < n \Rightarrow$ the LES has infinitely many solutions (and $n - r$ of the x_j :s can be taken arbitrary)
- ▶ Examples (5)

Remark (Homogeneous linear equation systems)

A LES with augmented matrix $(A|\vec{0})$ is called a **homogeneous LES**. The right-hand side of a homogeneous LES is always the null vector and has at least the trivial solution $\vec{x} = \vec{0}$.

Theorem

For a homogeneous LES, we have:

- ▶ $r = n \Leftrightarrow$ the LES has only the trivial solution $\vec{0}$
- ▶ $r < n \Leftrightarrow$ the LES has infinitely many solutions and $n - r$ of the x_j :s can be taken arbitrary.

Remark

A homogeneous LES with $n > m$ (with more unknowns than equations) has infinitely many solutions.

Remark

For a homogeneous LES is it sufficient to transform the coefficient matrix A with elementary row operations, since the null vector remains unchanged.

Definition (Square linear equation systems)

A LES with $m = n$ (Number of equations = Number of unknowns) is called a square LES.

Theorem

A square LES has a unique solution if and only if $r = n$.

Remark (Solution of a square LES with a unique solution by Gauss-Jordan elimination)

We can obtain the solution in the following way from the augmented matrix $(A^* | \vec{b}^*)$:

1. We add a multiple of the last row to the other rows, to create zeros above the leading 1.
2. We repeat the first step for the submatrix obtained by covering the last row and repeat this until the first row.

The result is an augmented matrix of the form

$$\left(A^{**} | \vec{b}^{**} \right) = \left(\begin{array}{cccc|c} 1 & 0 & \cdots & 0 & b_1^{**} \\ 0 & 1 & \cdots & 0 & b_2^{**} \\ \vdots & & & \vdots & \vdots \\ 0 & \cdots & & 1 & b_n^{**} \end{array} \right)$$

i.e. we can immediately see the solution:

$$\vec{x} = \begin{pmatrix} b_1^{**} \\ b_2^{**} \\ \vdots \\ b_n^{**} \end{pmatrix}$$

Remark

The Gaussian algorithm together with back-substitution is also called **Gauss-Jordan elimination**.

► Example (2)

Linear Algebra

Lectures 4+5+6: Matrices

Outline

- ▶ Introduction
- ▶ Matrix operations
- ▶ Special matrices
- ▶ The inverse of a matrix
- ▶ Matrix and linear equation system

Introduction

1. Linear transformations like translation, reflection, rotation, can be considered as matrix actions.
2. Our ability to analyze and solve equations will be greatly enhanced when we can perform algebraic operations with matrices.
3. Transformations on equations of a linear system can be written as transformations on rows of a matrix.
4. Matrix algebra provides tools for manipulating matrix equations and creating various useful formulas in ways similar to doing ordinary algebra with real numbers.
5. The inverse of a matrix, if it exists, allows us to treat the matrix as a number which helps to find out the solution of the linear system explicitly.
6. Matrix algebra can be applied in many fields like economics, computer graphics, image processing, etc.

Introduction

Remark

A matrix is a rectangular array, filled with numbers.

- ▶ Example

Definition (Matrix)

A rectangular array A of $m \cdot n$ numbers in m rows and n columns

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix}$$

is called a **matrix of size** $m \times n$.

m is the number of rows of A ;

n is the number of columns of A ;

a_{ij} , for $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$, are called entries of A .

$M_{m \times n}$: the set of all matrices of size $m \times n$.

► Example:

$$\begin{pmatrix} 0 & 1 & -1 \\ 1 & -2 & 3 \end{pmatrix} \in M_{2 \times 3}, \quad \text{and} \quad \begin{pmatrix} 0 & 1 \\ 1 & -1 \\ -1 & 3 \end{pmatrix} \in M_{3 \times 2}$$

Remark

A vector is only a special case of a matrix.

Definition

A matrix consisting of just one column is called a column vector:

$$\begin{pmatrix} a_1 \\ a_2 \\ \vdots \\ a_m \end{pmatrix}$$

Definition

A matrix consisting of just one row is called a **row vector**:

$$\left(\begin{array}{cccc} a_1 & a_2 & \cdots & a_n \end{array} \right)$$

Definition (Equality of matrices)

Two matrices A and B are equal, if they have the same size $m \times n$ and if all elements are equal:

$$a_{ij} = b_{ij} \text{ for all } i = 1, \dots, m, j = 1, \dots, n$$

Matrix Operations: Addition and Subtraction of Matrices

Definition (Addition and Subtraction of Matrices)

Let

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix}$$

and

$$B = \begin{pmatrix} b_{11} & b_{12} & \cdots & b_{1n} \\ b_{21} & b_{22} & \cdots & b_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ b_{m1} & b_{m2} & \cdots & b_{mn} \end{pmatrix}$$

be $m \times n$ matrices.

Then

$$A \pm B = \begin{pmatrix} a_{11} \pm b_{11} & a_{12} \pm b_{12} & \cdots & a_{1n} \pm b_{1n} \\ a_{21} \pm b_{21} & a_{22} \pm b_{22} & \cdots & a_{2n} \pm b_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{m1} \pm b_{m1} & a_{m2} \pm b_{m2} & \cdots & a_{mn} \pm b_{mn} \end{pmatrix},$$

i.e. the matrices are added and subtracted element-wise.

► Example: On the whiteboard

Matrix Operations: Multiplication with a scalar

- Example: On the whiteboard

Definition (Multiplication with a Scalar)

Matrices are element-wise multiplied by a scalar. Let

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix}$$

be an $m \times n$ matrix and let $\lambda \in \mathbb{R}$.

Then

$$\lambda \cdot A = \begin{pmatrix} \lambda \cdot a_{11} & \lambda \cdot a_{12} & \cdots & \lambda \cdot a_{1n} \\ \lambda \cdot a_{21} & \lambda \cdot a_{22} & \cdots & \lambda \cdot a_{2n} \\ \vdots & \vdots & \vdots & \vdots \\ \lambda \cdot a_{m1} & \lambda \cdot a_{m2} & \cdots & \lambda \cdot a_{mn} \end{pmatrix}$$

- Example: On the whiteboard

Theorem (Calculation Rules)

Let A, B, C be $m \times n$ matrices and $\lambda, \mu \in \mathbb{R}$

1. $A + B = B + A$ (Commutative)
2. $(A + B) + C = A + (B + C)$ (Associative)
3. $\lambda \cdot (A + B) = \lambda \cdot A + \lambda \cdot B$ (Distributive)
4. $(\lambda + \mu) \cdot A = \lambda \cdot A + \mu \cdot A$ (Distributive)

► Example: On the whiteboard

Matrix Operations: Matrix Multiplication

► Example

Definition (Matrix multiplication)

For an $m \times n$ matrix A and an $n \times p$ matrix B . The matrix $m \times p$ matrix C with the elements

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj}$$

called the matrix product of A and B , denoted by $C = AB$.

Remark

- ▶ The matrix product AB is only defined if the number of columns of A is equal to the number of rows of B .
- ▶ Write the sizes of A and B beside each other:

$$m \times n \text{ and } q \times p$$

The matrix product is only defined if the inner numbers are equal, i.e. $n = q$. Then the outer numbers tell the size of the matrix product: $m \times p$.

Remark

To determine the element in the i -th row and the j -th column of AB , we multiply the elements of the i -th row of A element-wise with the elements of the j -th column of B and add the generated products:

$$\begin{pmatrix} & \vdots & \\ a_{i1} & \cdots & a_{in} \\ & \vdots & \end{pmatrix} \cdot \begin{pmatrix} & b_{1j} & \\ \cdots & \vdots & \cdots \\ & b_{nj} & \end{pmatrix} = \begin{pmatrix} & \vdots & \\ \cdots & c_{ij} & \cdots \\ & \vdots & \end{pmatrix}$$

- Examples (2): On the whiteboard.

Remark

The product of an $m \times 1$ column vector with a $1 \times m$ row vector is an $m \times m$ matrix:

$$\begin{pmatrix} a_1 \\ \vdots \\ a_m \end{pmatrix} \cdot (b_1 \quad \cdots \quad b_m) = \begin{pmatrix} a_1 b_1 & \cdots & a_1 b_m \\ \vdots & & \vdots \\ a_m b_1 & \cdots & a_m b_m \end{pmatrix}$$

Remark

The product of a $1 \times m$ row vector with an $m \times 1$ column vector is a 1×1 matrix, i.e. a scalar (only a number):

$$\begin{pmatrix} b_1 & \cdots & b_m \end{pmatrix} \cdot \begin{pmatrix} a_1 \\ \vdots \\ a_m \end{pmatrix} = a_1 b_1 + \cdots + a_m b_m$$

This is a dot product written with matrix notation!

► Example

Remark (Missing Calculation Rules)

Many calculation rules for scalars hold also for matrices. But for the multiplication of matrices the commutative law does not hold. In general, is the following NOT fulfilled:

- ▶ Commutative law: $AB = BA$
- ▶ Zero law for products: $AB = 0 \Rightarrow A = 0$ or $B = 0$.
- ▶ Cancellation law: $AC = BC$ and $C \neq 0 \Rightarrow A = B$

Remark (Missing commutative law)

There are some situations, where the commutative law is missing for matrix multiplications:

- ▶ AB is defined, but BA is not defined.
- ▶ AB and BA are of different sizes.
- ▶ It is actually also possible that $AB \neq BA$ is not fulfilled, even if AB and BA have the same size:

$$\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \cdot \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = \begin{pmatrix} 1 & -1 \\ 0 & -1 \end{pmatrix}$$

but

$$\begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \cdot \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 0 & -1 \end{pmatrix}$$

Matrix Operations: Transposing Matrices

Definition (Transposed matrix)

The $n \times m$ matrix that we obtain if we write the rows of the $m \times n$ matrix A as columns, is called the **transposed matrix** of A ,

$$A^T = \begin{pmatrix} a_{11} & \cdots & a_{m1} \\ \vdots & \cdots & \vdots \\ a_{1n} & \cdots & a_{mn} \end{pmatrix}$$

- Example: On the white board.

Properties of transposes

Theorem

Let A and B be matrices. Then

1. $(A^T)^T = A$
2. $(A + B)^T = A^T + B^T$
3. $(cA)^T = cA^T$
4. $(AB)^T = B^T A^T$

Some Special Matrices

Definition (Zero matrix)

$a_{ij} = 0$ for all i, j , i.e.

► Example:

$$\begin{pmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix}$$

Properties of zero matrices

Theorem

Let A be an $m \times n$ matrix and let c be a scalar. Then

1. $A + 0_{m \times n} = A$
2. $A + (-A) = 0_{m \times n}$
3. if $cA = 0_{m \times n}$, then either $c = 0$ or $A = 0_{m \times n}$

Definition (Square matrix)

An $n \times n$ matrix and the elements a_{ii} , $i = 1, \dots, n$ constitute the main diagonal.

► Example

Definition (Diagonal matrix)

A square matrix

$$d_{ij} = 0 \text{ for } i \neq j,$$

i.e.

$$D = \begin{pmatrix} d_{11} & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & d_{nn} \end{pmatrix}$$

► Example

Definition (Identity matrix)

A diagonal matrix with main diagonal elements = 1

$$d_{ij} = 0 \text{ for } i \neq j,$$

and

$$d_{ij} = 1 \text{ for } i = j,$$

is called the identity matrix and denoted by $I_{n \times n}$ or I in case n is shown.

► Example: $I_{2 \times 2} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$, and $I_{3 \times 3} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$,

Definition (Upper triangular matrix)

A square matrix with

$$d_{ij} = 0 \text{ for } i > j,$$

i.e.

$$D = \begin{pmatrix} d_{11} & \cdots & d_{1n} \\ \vdots & & \vdots \\ 0 & \cdots & d_{nn} \end{pmatrix}$$

- Example: On the whiteboard

Definition (Lower triangular matrix)

A square matrix with

$$d_{ij} = 0 \text{ for } i < j,$$

i.e.

$$D = \begin{pmatrix} d_{11} & \cdots & 0 \\ \vdots & & \vdots \\ d_{n1} & \cdots & d_{nn} \end{pmatrix}$$

- Example: On the whiteboard

Definition (Symmetric matrix)

A square matrix A with

$$a_{ij} = a_{ji} \text{ for all } i, j,$$

is called symmetric. In other words, A is symmetric if and only if $A = A^T$.

- ▶ Examples: Transportation matrices, social network matrices, image processing matrices...

- ▶ $\begin{pmatrix} 2 & -1 & -1 \\ -1 & 1 & -2 \\ -1 & -2 & 3 \end{pmatrix}$ is symmetric and $\begin{pmatrix} 2 & -1 & -1 \\ -1 & 1 & -2 \\ -1 & 0 & 3 \end{pmatrix}$ is not symmetric

The inverse of a matrix

Definition (Regular matrices)

An $n \times n$ matrix A is called **regular** if $\text{rank}(A) = n$ and **singular** if $\text{rank}(A) < n$.

Remark

A regular \Leftrightarrow

$$\text{rank}(A) = n \Leftrightarrow$$

Every LES $A\vec{x} = \vec{b}$ has exactly one solution \Leftrightarrow

The homogeneous LES $A\vec{x} = \vec{0}$ has only the trivial solution

Theorem

If the $n \times n$ matrix A is regular, then a uniquely determined matrix $n \times n$ matrix X with $AX = I$ exists.

Definition (Invertible matrix)

An $n \times n$ matrix A is called **invertible** if an $n \times n$ matrix X exists with $AX = XA = I$, where I is the identity matrix of size $n \times n$.

Remark

The matrix X with $AX = I$ is called the inverse to A , and is written $X = A^{-1}$.

Remark

1. A regular $\Leftrightarrow A$ invertible.
2. If A is invertible, then is X uniquely determined.
3. For the inverse A^{-1} , the following holds: A^{-1} is regular and

$$AA^{-1} = A^{-1}A = I$$

with I the identity matrix.

Power of a square matrix

Given a square matrix $A \in M_{n \times n}$. We define

1. $A^0 = I$

2. $A^k = \underbrace{AA \cdots A}_{k \text{ times}} \ (k \in \mathbb{N})$.

3. $A^r \cdot A^s = A^{r+s}$

4. $(A^r)^s = A^{rs}$

5. If $D = \begin{pmatrix} d_1 & 0 & \cdots & 0 \\ 0 & d_2 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & d_n \end{pmatrix}$, then $D^k = \begin{pmatrix} d_1^k & 0 & \cdots & 0 \\ 0 & d_2^k & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & d_n^k \end{pmatrix}$

Theorem (Properties of inverse matrices)

Let A be an invertible matrix, let k be a positive integer, and c be a scalar not equal to zero. Then

1. A^{-1} is invertible and $(A^{-1})^{-1} = A$
2. A^k is invertible and $(A^k)^{-1} = (A^{-1})^k$
3. cA is invertible and $(cA)^{-1} = \frac{1}{c}A^{-1}$
4. A^T is invertible and $(A^T)^{-1} = (A^{-1})^T$
5. If A and B are invertible, then AB is invertible and $(AB)^{-1} = B^{-1}A^{-1}$.

Computation of the inverse

1. Augment A to the right with the identity matrix
2. Perform a Gauss-Jordan elimination

Then we obtain the inverse matrix A^{-1} on the right-hand side:

$$(A|I) = (I|A^{-1})$$

► Notes:

1. If A cannot be row reduced to I , then A is singular.
2. At first it is usually not obvious if the matrix A is invertible at all. If this is not the case, then we would arrive at a zero row at the left-hand side of the augmented matrix. We would then conclude that A is singular and stop the process.

Examples

- ▶ Find the inverse of the following matrix $\begin{pmatrix} 1 & -1 & 0 \\ 1 & 0 & -1 \\ -6 & 2 & 3 \end{pmatrix}$
- ▶ Solution: On the whiteboard.

Notes

From the definitions for matrix operations, a linear system of m equations and n variables x_1, \dots, x_n can be written as matrix operation:

$$Ax = b,$$

where A is the coefficient matrix of the system, x (resp. b) represents column vector of variables (the right hand side) of the system.

$$x = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}, \quad b = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{pmatrix}.$$

Theorem

Let A be an invertible $n \times n$ matrix.

Then the LES $A\vec{x} = \vec{b}$ has the unique solution

$$\vec{x} = A^{-1}\vec{b}$$

► Example

Remark

We can use this theorem if we have to solve many LES with the same coefficient matrix A but different right-hand sides \vec{b} .

Linear Algebra

Lectures 5+6: Determinants and applications

Outlines

- ▶ Determinant of a matrix
- ▶ Properties of determinants
- ▶ Gaussian elimination for calculating determinants
- ▶ Applications in inverse finding and solution solving

Determinants

Determinant is a number assigned to each square matrix. A single number, determinant, can tell only so much about a matrix. Still, it is amazing how much this number can do. We can use determinant for

1. the invertibility existence of a matrix;
2. the solution existence of a linear system;
3. finding the inverse of a matrix using cofactors;
4. solving explicitly a linear system by Cramer's rule;
5. measuring the dependence of $A^{-1}b$ on each element of b . If one parameter is changed in an experiment, or one observation is corrected, the "influence coefficient" in A^{-1} is a ratio of determinants.
6. measuring the amount by which a linear transformation changes the area of a figure. Finding the volume of a box in n -dimensional space.

Determinant of a 2×2 matrix

The number

$$\begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} = a_{11}a_{22} - a_{12}a_{21}$$

is called **the determinant** of the 2×2 matrix

$$\begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}$$

Determinant of a 3×3 matrix

The number

$$\begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} \\ = a_{11}a_{22}a_{33} + a_{12}a_{23}a_{31} + a_{13}a_{21}a_{32} - a_{13}a_{22}a_{31} - a_{11}a_{23}a_{32} - a_{12}a_{21}a_{33} \\ = a_{11} \cdot M_{11} - a_{12} \cdot M_{12} + a_{13} M_{13}$$

is called **the determinant** of the 3×3 matrix

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix}$$

where M_{ij} is **the determinant of the matrix obtained from A by removing its row i and column j** , and called **minor** of the entry a_{ij} .

Remarks

- For the computation of the determinant of a 3×3 matrix, we can use **Sarrus' rule**. Accordingly, its determinant is equal to the difference of the down-diagonal sum and up-diagonal sum of the following matrix:

$$\begin{array}{ccc|cc} a_{11} & a_{12} & a_{13} & a_{11} & a_{12} \\ a_{21} & a_{22} & a_{23} & a_{21} & a_{22} \\ a_{31} & a_{32} & a_{33} & a_{31} & a_{32} \end{array}$$

- We could also compute the determinant by expanding one row or one column and then compute the determinant for the remaining 2×2 matrices. This can also be generalized and then applied to the determinant for an $n \times n$ matrix.

Cofactor expansion

In general, the determinant of a square matrix of size $n \times n$ is a number defined inductively as follow:

1. Pick any one row (column) of A ;
2. For each entry in the row chosen, find its co-factor;
3. Multiply each entry in the row (column) chosen by its co-factor and take the sum that results the determinant of A .

Determinant of an $n \times n$ matrix

- ▶ **Minor** M_{ij} of the entry a_{ij} : the determinant of the $(n-1) \times (n-1)$ matrix obtained from A by removing its row i and column j .
- ▶ **Cofactor** of a_{ij}

$$C_{ij} = (-1)^{i+j} M_{ij}.$$

- ▶ The determinant of A is given by
 1. Cofactor expansion along the row i :

$$\det A = |A| = \sum_{j=1}^n a_{ij} C_{ij} = a_{i1} C_{i1} + a_{i2} C_{i2} + \cdots + a_{in} C_{in}.$$

2. Cofactor expansion along the column j :

$$\det A = |A| = \sum_{i=1}^n a_{ij} C_{ij} = a_{1j} C_{1j} + a_{2j} C_{2j} + \cdots + a_{nj} C_{nj}.$$

Examples

- ▶ **Notes:** The row (or column) containing the **most zeros** often if the best choice for expansion by cofactors.
- ▶ **Examples:** If **A** is an **triangular matrix**, then its determinant is the product of the entries on the main diagonal, where
 1. **upper triangular matrix:** All the entries below the main diagonal are zeros;
 2. **lower triangular matrix:** All the entries above the main diagonal are zeros;
 3. **diagonal matrix:** All the entries above and below the main diagonal are zeros.
- ▶ **Example:** Find the determinant

$$\begin{vmatrix} 2 & 1 & 0 \\ 1 & 1 & 4 \\ -3 & 2 & 5 \end{vmatrix}$$

Solution: On the white board.

Numerical note

By today's standard, a 25×25 matrix is small. Yet it would be impossible to calculate a 25×25 determinant by cofactor expansion. In general, a cofactor expansion requires over $n!$ multiplications, and $25!$ is approximately 1.5×10^{25} , it would spend approximately 500,000 years to compute a 25×25 determinant by this method.

Basic properties

Theorem

- ▶ If A contains a zero row or a zero column, then we have $\det(A) = 0$.
- ▶ $\det(A^T) = \det(A)$
- ▶ If A is a triangular matrix, then we have:

$$\det(A) = a_{11} \cdot a_{22} \cdot \dots \cdot a_{nn}$$

- ▶ $\det(I) = 1$, where I is the identity matrix.
- ▶ $\det(AB) = \det(A) \det(B)$. Consequently, $\det(A^{-1}) = \frac{1}{\det(A)}$.

Remark: In general, $\det(A + B) \neq \det(A) + \det(B)$.

Row operation rules for determinants

Theorem (Determinants and elementary row operations)

Let A be a square matrix

1. If a multiple of one row (column) of A is added to another row (column) to produce a matrix B , then $\det B = \det A$.
2. If two rows of A are interchanged to produce B , then $\det B = -\det A$
3. If one row of A is multiplied by k to produce B , then $\det B = k \det A$.

Proof.

On the white board. □

Remark: Most computer programs that compute $\det A$ for a general matrix A use the elementary row operations to reduce A to triangular matrix.

Examples

Compute $\det(A)$, where $A = \begin{pmatrix} 2 & -8 & 6 & 8 \\ 3 & -9 & 5 & 10 \\ -3 & 0 & 1 & -2 \\ 1 & -4 & 0 & 6 \end{pmatrix}$

Solution: On the whiteboard!

Zero determinants

If A is a square matrix and any of the following conditions is true, then $\det(A) = 0$.

1. An entire row (column) consists of zeros;
2. Two rows (columns) are equal;
3. One row (resp. column) is a multiple of another row (resp. column).
4. One row (resp. column) is a linear combination of other rows (resp. column).

Examples

$$1. \begin{vmatrix} 1 & k & k^2 \\ 1 & k & k^2 \\ 1 & k & k^2 \end{vmatrix} = 0$$

$$2. \begin{vmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \end{vmatrix} = 0$$

$$3. \begin{vmatrix} a_1 & b_1 & a_1 + b_1 \\ a_2 & b_2 & a_2 + b_2 \\ a_3 & b_3 & a_3 + b_3 \end{vmatrix} = 0$$

Theorem

For a square matrix A we have:

$$A \text{ is regular} \Leftrightarrow$$

$$A \text{ is invertible} \Leftrightarrow$$

$$\det(A) \neq 0$$

Remark

If A invertible, then

$$\det(A^{-1}) = \frac{1}{\det(A)}$$

Theorem (Cramer's rule)

If A is an invertible square matrix (i.e. $\det(A) \neq 0$), then the LES $A\vec{x} = \vec{b}$ has the unique solution

$$\vec{x} = \frac{1}{\det(A)} \begin{pmatrix} D_1 \\ \vdots \\ D_n \end{pmatrix}$$

Here D_k is the determinant we would obtain if in $\det(A)$ the k -th column is replaced by \vec{b} .

Proof.

On the whiteboard.



Remarks:

- ▶ Using determinants we can give a very nice and explicit formula for the solution to a linear system with n variables and n equations in case it has unique solution.
- ▶ Cramer's rule is needed in a variety of theoretical calculations. However, the formula is inefficient for hand calculations, except for 2×2 or perhaps 3×3 matrices because of the determinant computing complexity.
- ▶ In general, for solving a linear system, Gaussian and Gaussian-Jordan eliminations are used more often.

Examples

Determine the value of s for which the system has a unique solution, and use Cramer's rule to describe the solution

$$\begin{cases} 3cx - 2y = 4 \\ -6x + cy = 1 \end{cases}$$

Solution:

- View the system as $Ax = b$. Then

$$A = \begin{bmatrix} 3c & -2 \\ -6 & c \end{bmatrix}, \quad D_1 = \begin{bmatrix} 4 & -2 \\ 1 & c \end{bmatrix}, \quad D_2 = \begin{bmatrix} 3c & 4 \\ -6 & 1 \end{bmatrix},$$

- The system has unique solution if and only if $\det(A) \neq 0$:

$$\det(A) = 3c^2 - 12 = 3(c - 2)(c + 2) \neq 0 \Leftrightarrow c \neq \pm 2.$$

- When $c \neq \pm 2$, the solution to system is

$$x = \frac{\det(D_1)}{\det A} = \frac{4c + 2}{3(c - 2)(c + 2)},$$
$$y = \frac{\det(D_2)}{\det A} = \frac{c + 8}{(c - 2)(c + 2)}.$$

Inverse matrix finding: Case of 2×2 matrix

Recall:

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}, \quad A^{-1} = \frac{1}{\det A} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$$

A formula for the inverse matrix

Theorem

Let A be an invertible matrix and let C be cofactor matrix of A , i.e. C_{ij} are cofactors of entries a_{ij} . Then,

$$A^{-1} = \frac{1}{\det A} C^T.$$

Remark: The matrix C^T is called the **adjoint matrix** of A , denoted by $\text{adj}(A)$.

Proof.

On the white board.



Examples

Using cofactors for finding the inverse of

1.

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}, \quad A^{-1} = \frac{1}{ad - bc} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$$

2.

$$A = \begin{pmatrix} 3 & 1 & -2 \\ -1 & 1 & 2 \\ 1 & -2 & 1 \end{pmatrix}.$$

Solution: On the white board.

Linear Algebra

Lectures 7+8+9+10+11: Vector spaces

Introduction

- ▶ Actually, a study of vector spaces is not much different from a study of \mathbb{R}^n itself. We can use our geometric experience with \mathbb{R}^2 , \mathbb{R}^3 to visualize many general concepts.
- ▶ We can use vector space terminology to tie together important facts about rectangular matrices like rank concept.
- ▶ In applications of linear algebra, subspaces of \mathbb{R}^n usually arise in one of two ways: as the set of solutions of a **homogeneous linear system**, or as **the set of all linear combinations** of certain specified vectors. These lead to considering the **null** and **column spaces** of a matrix/linear transformation.

Outline

- ▶ Vector Spaces
- ▶ Subspaces of Vector Spaces
- ▶ Spanning Sets and Linear Independence
- ▶ Basis and Dimension
- ▶ Rank of a Matrix and Systems of Linear Equations
- ▶ Four fundamental subspaces
- ▶ Linear transformations
- ▶ Coordinate systems

Vector space I

Let V be a set on which two operations (**vector addition:** $+: V \times V \rightarrow V$ and **scalar multiplication:** $\cdot: \mathbb{R} \times V \rightarrow V$) are defined. If the following axioms are satisfied for every u , v , and w in V and every scalar (real number) λ and μ , then V is called a *vector space*.

Vector space II

1. $u + v \in V$ (closed with the addition)
2. $u + v = v + u$ (commutative addition)
3. $u + v + w = u + (v + w)$ (associative addition)
4. V has a zero vector 0 s.t. $u + 0 = u$ for all $u \in V$
5. For each u in V , there is an opposite vector in V , denoted by $-u$, s.t. $u + (-u) = 0$
6. $\lambda u \in V$ (closed with scalar multiplication)
7. $\lambda(u + v) = \lambda u + \lambda v$
8. $(\lambda + \mu)u = \lambda u + \mu u$
9. $\lambda\mu u = \lambda(\mu u)$
10. $1.u = u$.

Examples I

1. n -tuple space: \mathbb{R}^n , with the vector addition and the multiplication with scalar.

- ▶ Vector addition

$$(u_1, u_2, \dots, u_n) + (v_1, v_2, \dots, v_n) = (u_1 + v_1, u_2 + v_2, \dots, u_n + v_n).$$

- ▶ Scalar multiplication

$$k(u_1, u_2, \dots, u_n) = (ku_1, ku_2, \dots, ku_n).$$

2. Matrix space: $V = M_{m \times n}$ with the matrix addition and scalar multiplication. Example: $m = n = 2$

- ▶ Matrix addition

$$\begin{pmatrix} u_{11} & u_{12} \\ u_{21} & u_{22} \end{pmatrix} + \begin{pmatrix} v_{11} & v_{12} \\ v_{21} & v_{22} \end{pmatrix} = \begin{pmatrix} u_{11} + v_{11} & u_{12} + v_{12} \\ u_{21} + v_{21} & u_{22} + v_{22} \end{pmatrix}$$

- ▶ Scalar multiplication

$$k \begin{pmatrix} u_{11} & u_{12} \\ u_{21} & u_{22} \end{pmatrix} = \begin{pmatrix} ku_{11} & ku_{12} \\ ku_{21} & ku_{22} \end{pmatrix}$$

Examples II

3. The set of all real **polynomials of degree not exceeding n** : $V = P_n(x)$ together with the polynomial addition and polynomial scalar multiplication forms a vector space. (**verify this please!**)
4. The set of all **solutions of a homogeneous linear equation system** together with vector addition and scalar multiplication forms a vector space. For instance, the solutions of the equation $x + 2y - 4z = 0$. (**verify this please!**)
5. The set of **all continuous functions** on a given domain with the scalar multiplication and function addition.

Notes: To show that a set is not a vector space, you need only find one axiom that is not satisfied.

Examples:

- ▶ The set of all integers is not a vector space since let $\frac{1}{2} \in \mathbb{R}$ and $1 \in V$ then $(\frac{1}{2})(1) = \frac{1}{2} \notin V$.
- ▶ The set of all second-degree polynomials is not a vector space, since let $p(x) = x^2 - 2x - 1$ and $q(x) = -x^2 - 1$, then $p(x) + q(x) = -2x - 2$ which is not a second-degree polynomial.
- ▶ The set of all solutions of a non-homogeneous linear equation system ([verify this please](#)).

Subspaces

Definition

V : vector space

$W \subseteq V$ and $W \neq \emptyset$.

W is called a **subspace** of V if W together with the addition and the multiplication with a scalar inherited from V is a vector space.

Notes:

1. Two trivial subspaces of V : Zero vector space $\{\vec{0}\}$, and V .
2. Test for a subspace: a non-empty set W is a subspace of V if and only if it is closed with the addition and the multiplication with a scalar, i.e.,
 - ▶ If u and v are in W , then $u + v$ is in W
 - ▶ If u is in W and λ is any scalar, then λu is in W .

Examples

- ▶ A set of points on a line through the origin in the plane is a subspace of \mathbb{R}^2 .
- ▶ Let W be the set of all 2×2 symmetric matrices. Then W is a subspace of $M_{2 \times 2}$.
- ▶ The set of singular matrices of size 2×2 is not a subspace of $M_{2 \times 2}$.
- ▶ The set of invertible matrices of size 2×2 is not a subspace of $M_{2 \times 2}$.
- ▶ The set of solutions of a homogeneous system with n variables is a subspace of \mathbb{R}^n .

Lemma

The intersection of two subspaces is a subspace.

Linear combinations I

Definition (Linear combinations)

Let V be a vector space and let $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k$ be vectors in V . Let $\alpha_1, \alpha_2, \dots, \alpha_k$ be numbers. An expression of type

$$\alpha_1 \mathbf{v}_1 + \alpha_2 \mathbf{v}_2 + \cdots + \alpha_k \mathbf{v}_k$$

is called a **linear combination** of $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k$. The numbers $\alpha_1, \alpha_2, \dots, \alpha_n$ are called the **coefficients** of the linear combination.

Lemma

*The set of all linear combinations of vectors $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ in a vector space V is a subspace of V . This subspace is denoted by $\text{span}(S)$ and called **the span** of S .*

Examples

1. All linear combinations of vector $(1, 2)$ in the plane \mathbb{R}^2 are on the line through two points $(1, 2)$ and the origin.
2. All linear combinations of vectors $(1, 1, 0)$ and $(1, 0, 0)$ in \mathbb{R}^3 are on the plane through 3 points $(0, 1, 0)$, $(1, 0, 0)$ and the origin.

Spanning set

1. The set of all linear combinations of k vectors $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ is called **the span of S** and denoted by $\text{span}(S)$. Hence,

$$\text{span}(S) = \{a_1\mathbf{v}_1 + a_2\mathbf{v}_2 + \dots + a_k\mathbf{v}_k : \forall a_1, a_2, \dots, a_k \in \mathbb{R}\}.$$

2. A **spanning set of a vector space**: Given a vector space V and a set $S \subseteq V$. If every vector of V can be written as a linear combination of vectors in S , i.e., $V = \text{span}(S)$, then S is called a **spanning set of V** .

1. **Some terminologies:** Given a set of vectors S . If $\text{span}(S) = V$, then we can say that

- ▶ S **spans (generates)** V ;
- ▶ V **is spanned by** S ;
- ▶ S is a **spanning set** of V .

2. Notes:

- i) $\text{span}(\emptyset) = \{\vec{0}\}$;
- ii) $S \subseteq \text{span}(S)$
- iii) $S_1, S_2 \subseteq V$, and $S_1 \subseteq S_2$, then $\text{span}(S_1) \subseteq \text{span}(S_2)$.

Lemma

Let $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ be vectors in V . Then,

- i) $\text{span}(S)$ is a subspace of V .
- ii) $\text{span}(S)$ is the smallest subspace of V that contains S .

Examples

1. $\mathbf{e}_1 = (1, 0, 0), \mathbf{e}_2 = (0, 1, 0), \mathbf{e}_3 = (0, 0, 1)$ is a spanning set for \mathbb{R}^3 ;
2. $\mathbf{e}_1 = 1, \mathbf{e}_2 = x, \mathbf{e}_3 = x^2$ is a spanning set for $P_2(x)$;
3. Let H be the set of all vectors of the form $(a - 3b, b - a, a, b)$, where a and b are arbitrary scalars. We can write down column vector

$$\begin{bmatrix} a - 3b \\ b - a \\ a \\ b \end{bmatrix} = a \begin{bmatrix} 1 \\ -1 \\ 1 \\ 0 \end{bmatrix} + b \begin{bmatrix} -3 \\ 1 \\ 0 \\ 1 \end{bmatrix}$$

which shows that $H = \text{span}\{\mathbf{v}_1, \mathbf{v}_2\}$, where $\mathbf{v}_1 = (1, -1, 1, 0)$ and $\mathbf{v}_2 = (-3, 1, 0, 1)$.

Details are on the whiteboard!

Linear independence and linear dependence

Definition

Let $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ be a set of vectors in a vector space V . Consider the equation

$$c_1\mathbf{v}_1 + \dots + c_k\mathbf{v}_k = \vec{0}.$$

- i) If the equation has only the trivial solution $c_1 = c_2 = \dots = c_k = 0$, then S is called **linearly independent**.
- ii) If the equation has a nontrivial solution (i.e., not all zeros), then S is called **linearly dependent**.

Notes

1. \emptyset is linearly independent
2. If $\vec{0} \in S$, then S is linearly dependent.
3. If $\mathbf{v} \neq \vec{0}$, then $\{\mathbf{v}\}$ is linearly independent
4. If $S_1 \subseteq S_2$, then
 - ▶ S_1 is linearly dependent $\Rightarrow S_2$ is linearly dependent;
 - ▶ S_2 is linearly independent $\Rightarrow S_1$ is linearly independent;

Examples: Testing for linearly independent

1. Vectors $(1, 1)$ and $(-3, 2)$ are linearly independent in \mathbb{R}^2 .
2. Determine whether the following set of vectors in P_2 is linearly independent or dependent.

$$S = \{1 + x - 2x^2, 2 + 5x - x^2, x + x^2\}.$$

3. Determine whether the following set of vectors in 2×2 matrix space is linearly independent or dependent.

$$S = \left\{ \begin{bmatrix} 2 & 1 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 \\ 2 & 0 \end{bmatrix} \right\}$$

4. Let $p_1(t) = 1$, $p_2(t) = t$ and $p_3(t) = 4 - t$ be polynomials. Then the set $\{p_1, p_2, p_3\}$ are linearly dependent since $p_3 + p_2 - 4p_1 = 0$.

Solution: On the whiteboard!

Remarks

The main difference between linear dependence in \mathbb{R}^n and in a general vector space is that when the vectors are not n -tuples, the homogeneous equation usually cannot be written as a system of n linear equations. That is, the vectors cannot be made into the columns of a matrix A in order to study the equation $Ax = 0$. We must rely instead on the definition of linear dependence.

Theorem

Let V be a vector space. Let $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k$ be linearly independent elements in V . Let $\alpha_1, \dots, \alpha_k$ and β_1, \dots, β_k be numbers such that

$$\alpha_1 \mathbf{v}_1 + \dots + \alpha_k \mathbf{v}_k = \beta_1 \mathbf{v}_1 + \dots + \beta_k \mathbf{v}_k.$$

Then we must have

$$\alpha_1 = \beta_1, \alpha_2 = \beta_2, \dots, \alpha_k = \beta_k.$$

Proof.

On the whiteboard!



Linear dependence and spanning set

Theorem (Spanning set theorem)

Let $S = \{v_1, v_2, \dots, v_k\}$ be a set in a vector space V , and let $H = \text{span}(S)$.

- ▶ If one of the vectors in S , say v_k , is a linear combination of the remaining vectors in S , then the set formed from S by removing v_k still spans H .
- ▶ If $H \neq \{0\}$, then some subset of S is a basis for H .

Proof.

On the whiteboard!



Basis and dimension I

► Definition of a basis: Let V be a vector space and $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n\} \subseteq V$. If

(i) S spans V (i.e., $\text{span}(S) = V$)

(ii) S is linearly independent,

then S is called a **basis** of V .

► Examples:

i) \emptyset is a basis for $\{\vec{0}\}$

ii) the standard basis for \mathbb{R}^3 : $\{\mathbf{e}_1, \mathbf{e}_2, \mathbf{e}_3\}$, where

$$\mathbf{e}_1 = (1, 0, 0), \mathbf{e}_2 = (0, 1, 0), \mathbf{e}_3 = (0, 0, 1).$$

iii) Another basis for \mathbb{R}^3 :

$$\mathbf{v}_1 = (1, 1, 1), \mathbf{v}_2 = (1, 1, 0), \mathbf{v}_3 = (1, 0, 0).$$

Basis and dimension II

iv) the standard basis for $M_{2 \times 2}$:

$$\left\{ \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} \right\}$$

v) The standard basis for polynomials of degree $\leq n$:

$$\{1, x, x^2, \dots, x^n\}.$$

vi) Another basis for polynomials of degree $\leq n$:

$$\{1, (1+x), (1+x)^2, \dots, (1+x)^n\}.$$

Properties of basis

Theorem

Let V be a vector space.

1. If $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n\}$ is a basis for V , then every vector in V can be written in one and only one way as a linear combination of vectors in S .
2. Let S be a basis of V . If $|S| = n$, then every set containing more than n vectors in V is linearly dependent.

Maximal subset of linearly independence

A set of vectors $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ in V is called a **maximal** subset of linearly independent elements if S is linearly independent, and if $S \cup \{\mathbf{v}\}$ is linearly dependent for any $\mathbf{v} \notin S$.

Theorem

1. If $S = \{\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_k\}$ is a spanning set of V and if S is a maximal subset of linearly independent elements, then S is a basis of V .
2. Let V be a vector space and suppose that one basis has n elements, and another basis has m elements. Then $m = n$, i.e., all bases for a finite-dimensional vector space has the same number of vectors.

Remarks

1. When a vector space V has a basis of finite elements, then the number of vectors in a basis for V is called the **dimension** of V , denoted by $\dim(V)$. In other words, if S is a finite set and S is a basis of V , then

$$\dim(V) = |S|.$$

2. A vector space V is called **finite dimensional**, if it has a basis consisting of a finite number of elements.
3. If a vector space V is not finite dimensional, then it is called **infinite dimensional**.

Remarks

1. $\dim(\{\vec{0}\}) = 0$.
2. $\dim(V) = n$ and $S \subseteq V$.
 - ▶ If S is a generating set, then $|S| \geq n$.
 - ▶ If S is a linear independent set, then $|S| \leq n$.
 - ▶ If S is a basis, then $|S| = n$.
3. If $\dim(V) = n$ and if W is a subspace of V , then $\dim(W) \leq n$.

Examples

1. $\dim(\mathbb{R}^b) = n$
2. $\dim(\mathbb{M}_{m \times n}) = mn$
3. $\dim(P_n(x)) = n + 1$, where $P_n(x)$ the vector space of all polynomials with degree $\leq n$.
4. $\dim(P(x)) = \infty$, where $P(x)$ the vector space of all polynomials.
5. Find the dimension of the subspace

$$H = \{(a - 3b + 6c, 5a + 4d, b - 2c - d, 5d) : a, b, c, d \in \mathbb{R}\}.$$

Answer: On the whiteboard.

Null and column spaces

In applications of linear algebra, subspaces of \mathbb{R}^n usually arise in one two ways:

- ▶ as the set of all solutions to a homogeneous linear system
- ▶ as the set of all linear combinations of certain vectors.

These two subspaces are called null and column spaces. In the sense of linear transformation, they also called kernel and range of the linear transformation respectively.

Null and column spaces

Given a matrix A of size $m \times n$.

1. **Null space** of A is the set of all solutions of the homogeneous equation $A\mathbf{x} = \mathbf{0}$. In set notation,

$$\text{Nul}(A) = \{\mathbf{x} \in \mathbb{R}^n : A\mathbf{x} = \mathbf{0}\}$$

2. **Column space** of A is the set of all linear combinations of the columns of A . In the set notation,

$$\text{Col}(A) = \text{span}\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_n\} = \{\mathbf{b} : \mathbf{b} = A\mathbf{x} \text{ for some } \mathbf{x} \in \mathbb{R}^n\},$$

where \mathbf{a}_i are columns of A .

Theorem

Given a matrix A of size $m \times n$. The null space $\text{Null}(A)$ of A is a subspace of \mathbb{R}^n . Equivalently, the set of all solutions to a homogeneous linear system with m equations, n unknown is a subspace of \mathbb{R}^n .

Remarks

- ▶ There is no obvious relation between vectors in $\text{Null}(A)$ and the entries in A .
- ▶ The $\text{Null}(A)$ is defined implicitly, because it is defined by a condition that must be checked.
- ▶ No explicit list or description of the elements in $\text{Null}(A)$ is given. However, solving the equation $Ax = 0$ produces an explicit description of $\text{Null}(A)$.

Finding a basis for $\text{Null}(A)$

1. Find the reduced echelon-form $(A^{**}|0)$ of the augmented matrix $(A|0)$.
2. Write dependent variables in terms of free variables (free variables are variables corresponding to non-pivot columns of A^{**})
3. Find general solution to $Ax = 0$ in term of free variables.
4. Decompose the general solution vector into a linear combination of vectors where the coefficients are the free variables.
5. Vectors in the above step form a basis for $\text{Null}(A)$.
6. Moreover, $\dim(\text{Null}(A))$ is equal to the number of free variables.

Examples I

Find a spanning set and a basis for $\text{Null}(A)$.

1. $A = \begin{pmatrix} 1 & -3 & -2 \\ -5 & 9 & 1 \end{pmatrix}$.

Solution: We reduce the augmented matrix $(A|0)$ to reduced-row echelon form:

$$\begin{aligned} A &= \left(\begin{array}{ccc|c} 1 & -3 & -2 & 0 \\ -5 & 9 & 1 & 0 \end{array} \right) \rightarrow \left(\begin{array}{ccc|c} 1 & -3 & -2 & 0 \\ 0 & -6 & -9 & 0 \end{array} \right) \rightarrow \left(\begin{array}{ccc|c} \frac{1}{2} & 0 & \frac{1}{2} & 0 \\ 0 & -6 & -9 & 0 \end{array} \right) \\ &\rightarrow \left(\begin{array}{ccc|c} 1 & 0 & 1 & 0 \\ 0 & 1 & \frac{3}{2} & 0 \end{array} \right) \end{aligned}$$

Hence, x_1, x_2 are dependent variables and x_3 is free variable.

We get

$$x = (x_1, x_2, x_3) = (-x_3, -\frac{3}{2}x_3, x_3) = x_3(-1, -\frac{3}{2}, 1)$$

and $(-1, -\frac{3}{2}, 1)$ is a basis for $\text{Null}(A)$.

Examples II

2. $A = \begin{pmatrix} -3 & 6 & -1 & 1 & -7 \\ 1 & -2 & 2 & 3 & -1 \\ 2 & -4 & 5 & 8 & -4 \end{pmatrix}.$

Solution: On the whiteboard.

Column space

Theorem

Given a matrix A of size $m \times n$. Assume that A^* is the row-echelon form of A . Then, the column space of A is the space spanned by column of A corresponding to pivot columns in A^* .

Examples: $A = \begin{pmatrix} 1 & -3 & -2 \\ -5 & 9 & 1 \end{pmatrix}$. We have

$$A^* = \begin{pmatrix} 1 & -3 & -2 \\ 0 & -6 & -9 \end{pmatrix},$$

and

$$\{(1, -5), (-3, 9)\}$$

is a basis for $\text{Col}(A)$.

Examples

Find a basis and the dimension for $\text{Col}(A)$ where

$$A = \begin{pmatrix} -3 & 6 & -1 & 1 & -7 \\ 1 & -2 & 2 & 3 & -1 \\ 2 & -4 & 5 & 8 & -4 \end{pmatrix}.$$

Solution: On the whiteboard.

Row space

- ▶ Similarly, we can define the row space of a matrix, denoted by $\text{row}(A)$, which is a subspace generated by its rows. Note that $\dim(\text{row}(A))$ is equal to the number of non-zero rows in the row-echelon form A^* . Furthermore, non-zero rows of A^* forms a basis for $\text{row}(A)$.
- ▶ Given $A \in M_{m \times n}$. It is straightforward that

$$\dim(\text{Col}(A)) = \dim(\text{row}(A)).$$

- ▶ We define the **rank** of a matrix A , denoted by $r(A)$, as the dimension of row (column) space of A .

The following is straightforward from the fact that

number of pivot columns + number of non-pivot columns = number of columns

- ▶ $r(A) + \dim(\text{null}(A)) = n$

- ▶ $r(A) + \dim(\text{null}(A^T)) = m$

Theorem

Let $A \in M_{n \times n}$ and A be invertible. The following are equivalent

- ▶ Columns of A form a basis of \mathbb{R}^n .
- ▶ $\text{col}(A) = \mathbb{R}^n$
- ▶ $\dim(\text{col}(A)) = n$
- ▶ $r(A) = n$
- ▶ $\text{null}(A) = \{\vec{0}\}$
- ▶ $\dim(\text{null}(A)) = 0$

Linear algebra: Linear mappings

Lectures 7+8+9+10+11: Vector spaces

Introduction

- ▶ Among mappings, the linear mappings are the most important.
- ▶ A good deal of mathematics is devoted to reducing questions concerning arbitrary mappings to linear mappings.
- ▶ On the other hand, it is often possible to approximate an arbitrary mapping by a linear one, whose study is much easier than the study of the original mapping. This is done in the calculus of several variables.

Linear mappings

Let V, W be vector spaces on \mathbb{R} . A **linear mapping**

$$T : V \rightarrow W$$

is a mapping which satisfies the following two properties

- ▶ $T(u + v) = T(u) + T(v)$ for any $u, v \in V$
- ▶ $T(cu) = cT(u)$ for any $c \in \mathbb{R}$ and $u \in V$.

Examples I

- ▶ (Projection):

$$\begin{aligned}T : \mathbb{R}^3 &\rightarrow \mathbb{R}^2 \\(x, y, z) &\mapsto (x, y)\end{aligned}$$

- ▶ (Inner product):

$$\begin{aligned}T : \mathbb{R}^3 &\rightarrow \mathbb{R} \\(x, y, z) &\mapsto ax + by + cz,\end{aligned}$$

where a, b, c are given numbers in \mathbb{R} .

- ▶ (Linear mapping given by a matrix):

$$\begin{aligned}T : \mathbb{R}^n &\rightarrow \mathbb{R}^m \\ \mathbf{v} &\mapsto A\mathbf{v},\end{aligned}$$

where A is a given matrix.

Examples II

- ▶ (A linear transformation from $M_{m \times n}$ to $M_{n \times m}$)

$$T : M_{m \times n} \rightarrow M_{n \times m}$$

$$A \mapsto A^T$$

- ▶ (Differentiation): T transforms a differentiable function f to its derivative.

Properties of linear mappings

Given a linear mapping

$$T : V \rightarrow W.$$

Then

- ▶ $T(\vec{0}) = \vec{0}$;
- ▶ $T(-\mathbf{v}) = -T(\mathbf{v})$
- ▶ If $\mathbf{v} = c_1 v_1 + \cdots + c_k v_k$, then

$$T(\mathbf{v}) = c_1 T(v_1) + c_2 T(v_2) + \cdots + c_k T(v_k).$$

Examples

Consider a map T from $M_{m \times n}$ to $M_{n \times m}$:

$$\begin{aligned} T : M_{m \times n} &\rightarrow M_{n \times m} \\ A &\mapsto A^T \end{aligned}$$

Then T is a linear transformation since

$$T(A + B) = (A + B)^T = A^T + B^T = T(A) + T(B),$$

and

$$T(cA) = (cA)^T = cA^T = cT(A).$$

Theorem

Let V and W be vector spaces. Let $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n$ be a basis for V . Let $\mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_n$ be arbitrary elements of W . Then there exists a unique linear mapping $T : V \rightarrow W$ such that

$$T(\mathbf{v}_1) = \mathbf{w}_1, \dots, T(\mathbf{v}_n) = \mathbf{w}_n.$$

In the other words, a linear mapping is well defined by its images of a basis.

The kernel and image of a linear mapping

- **Kernel of a linear map T :** Let $T : V \rightarrow W$ be a linear map. Then the set of all vectors v in V that satisfies $T(v) = 0$ is called the **kernel** of T and is denoted by $\ker(T)$.

$$\ker(T) = \{v : T(v) = 0\}.$$

- **Examples:** Given

$$T(x) = Ax = \begin{bmatrix} 1 & -1 & -2 \\ -1 & 2 & 3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}.$$

Then, $\ker(T)$ coincides with solution set $\text{null}(A)$ of $Ax = 0$. To calculate $\ker(T)$, we use Gauss-Jordan elimination as presented in the previous lecture.

Observation

The kernel of a linear transformation $T : V \rightarrow W$ is an essentially solution set of a homogeneous system of linear equations, and therefore is a subspace of the domain V . For this reason, sometimes T is called the **nullspace** of T .

Some popular linear transformations

On the whiteboard!

1. Reflection through a line
2. Rotation
3. Projection

- ▶ Composition of linear transformation
- ▶ Inverse linear transformation
- ▶ Matrix of a transformation in any bases:

Coordinate systems and basis changes

When a basis \mathcal{B} is chosen for an n -dimensional vector space V , the associated coordinate mapping onto \mathbb{R}^n provides a coordinate system for V .

In some applications, a problem is described initially using a basis \mathcal{B} , but the problem's solution is aided by changing \mathcal{B} to a new basis \mathcal{C} . Each vector is assigned a new \mathcal{C} -coordinate vector. We show in this section how the coordinates of a given vector associated to different bases are related.

Coordinate systems and changes of basis

Given a finite dimensional vector space V such that $\dim V = n$. Let $\mathcal{B} = \{v_1, v_2, \dots, v_n\}$. The coordinates of a vector $v \in V$ relative to a basis \mathcal{B} or \mathcal{B} -coordinates of v , denoted by $[v]_{\mathcal{B}}$, are real numbers $\lambda_1, \lambda_2, \dots, \lambda_n$ such that

$$v = \lambda_1 v_1 + \lambda_2 v_2 + \cdots + \lambda_n v_n.$$

Examples:

1. Let $V = \mathbb{R}^n$ and \mathcal{B} the standard basis of V and $v = (x_1, x_2, \dots, x_n) \in \mathbb{R}^n$, then $[v]_{\mathcal{B}} = (x_1, \dots, x_n)$.
2. Given $V = \mathbb{R}^2$ and $\mathcal{B} = \{(1, 0), (1, 2)\}$. Let $u = (1, 1) \in \mathbb{R}^2$. We have $(1, 1) = \frac{1}{2}(1, 0) + \frac{1}{2}(1, 2)$. Hence, the coordinates of $(1, 1)$ relative to the basis \mathcal{B} is $[(1, 1)]_{\mathcal{B}} = (\frac{1}{2}, \frac{1}{2})$.

The matrix of a linear transformation

Given two finite dimensional spaces V and W of dimension n and m respectively. Assume that $\mathcal{B} = \{v_1, v_2, \dots, v_n\}$ and $\mathcal{C} = \{w_1, w_2, \dots, w_m\}$ are bases for V and W respectively. Let T be a linear transformation from V to W . We define the matrix representation for T associated to bases \mathcal{B} and \mathcal{C} as follows:

$$A = [[T(v_1)]_{\mathcal{C}}, [T(v_2)]_{\mathcal{C}}, \dots, [T(v_n)]_{\mathcal{C}}],$$

where $[T(v_i)]_{\mathcal{C}}$ is the column i of A .

Examples

1. Let $\mathcal{B} = \{v_1, v_2\}$ and $\mathcal{C} = \{w_1, w_2, w_3\}$ and let T be a linear transformation such that $T(v_1) = 3w_1 - 2w_2 + 5w_3$ and $T(v_2) = 4w_1 + 7w_2 - w_3$. Then $A = \begin{pmatrix} 3 & 4 \\ -2 & 7 \\ 5 & -1 \end{pmatrix}$.
2. Let $\mathcal{B} = \{(1, -3), (-2, 4)\}$ and $\mathcal{C} = \{(-7, 9), (-5, 7)\}$ and let T be the identity transformation on \mathbb{R}^2 . Then,

$$T((1, -3)) = (1, -3) = 2(-7, 9) - 3(-5, 7)$$

and

$$T((-2, 4)) = (-2, 4) = \frac{-3}{2}(-7, 9) + \frac{5}{2}(-5, 7).$$

Hence, the matrix representation for T associated to \mathcal{B} and \mathcal{C} is $\begin{pmatrix} 2 & \frac{-3}{2} \\ -3 & \frac{5}{2} \end{pmatrix}$.

Matrix representation for a linear transformation

Matrix representation allows us to find the coordinates of the image of x under T in term of coordinates of x as follows:

Theorem

Given a linear transformation T from V to W . Let A be the matrix representation of T associated to bases B and C on V and W . Then for any $x \in V$, we have

$$[T(x)]_C = A \cdot [x]_B.$$

Change-of-coordinates matrix

Remark:

If \mathcal{B} and \mathcal{C} are bases for the same vector space V , and if T is the identity transformation $T(v) = v$ for any $v \in V$, then the matrix A is called the change-of-coordinate matrix from the basis \mathcal{B} to the basis \mathcal{C} .

Examples

Let $\mathcal{B} = \left\{ \begin{pmatrix} 2 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 1 \end{pmatrix} \right\}$ be a basis for \mathbb{R}^2 . Then,

1. the change-of-coordinate matrix from \mathcal{B} to the standard basis is $A = \begin{pmatrix} 2 & -1 \\ 1 & 1 \end{pmatrix}$.
2. Let $\mathbf{v} = (4, 5) \in \mathbb{R}^2$. Then the coordinates $[\mathbf{v}]_{\mathcal{B}}$ of \mathbf{v} relative to \mathcal{B} are given by

$$[\mathbf{v}]_{\mathcal{B}} = A^{-1} \cdot \begin{pmatrix} 4 \\ 5 \end{pmatrix} = \begin{pmatrix} 3 \\ 2 \end{pmatrix}.$$

In general, if A is the change-of-coordinates matrix from \mathcal{B} to the standard basis in \mathbb{R}^n , then for any vector $\mathbf{v} \in \mathbb{R}^n$ we have

$$[\mathbf{v}]_{\mathcal{B}} = A^{-1} \mathbf{v}.$$

Exercises

1. Let $\mathcal{B} = \{(1, 0, 0), (-3, 4, 0), (3, -6, 3)\}$ and $x = (-8, 2, 3)$.
 - a) Find the change-of-coordinates matrix from \mathcal{B} to the standard basis;
 - b) Find $[x]_{\mathcal{B}}$.
2. Let $P_2(x)$ be the set of polynomials of degree at most 2. Let $\mathcal{B} = \{1 + x, 1 + x^2, x + x^2\}$.
 - a) Show that \mathcal{B} is a basis for $P_2(x)$.
 - b) Find the change-of-coordinates matrix from \mathcal{B} to the standard basis of $P_2(x)$.
 - c) Find the coordinate vector of the polynomial $f(x) = 6 + 3x - x^2$ relative to the basis \mathcal{B} .

The following allows us to calculate the change-of-coordinates matrix between two bases based on the change-of-coordinates matrix from those bases to the standard basis.

Theorem

Let \mathcal{E} be the standard basis on \mathbb{R}^n . Let \mathcal{B} and \mathcal{C} be two bases on \mathbb{R}^n . Assume that B and C are change-of-coordinates matrices from bases \mathcal{B} and \mathcal{C} to \mathcal{E} respectively. Then, the change-of-coordinates matrix from \mathcal{B} to \mathcal{C} is determined by $A = CB^{-1}$.

Part 2: General Algebra

Topics cover:

- ▶ Propositional logic
- ▶ Relations
- ▶ Number theory and its applications in cryptography

Part 2: General Algebra

Lectures 12: Propositional Algebra

Introduction

Propositional algebra is a subject investigating boolean values. It can be applied in designing circuits, specifications of programs, computer architecture.

Statements

- ▶ A **statement** is an assertion that can be determined to be true or false.
- ▶ The **truth value** of a statement is T if it is true and F if it is false.

Example:

- i) The statement " $2 + 3 = 5$ " has truth value T.
- ii) "What time is it now" is not a statement.
- iii) "Sugar is bitter" is a false statement.

Compound Statements

Connectives:

- ▶ and
- ▶ or
- ▶ not
- ▶ if ... then
- ▶ ... if and only if ...

Statements that involve one or more of the connectives are **compound** statements. Otherwise they are **simple** statements.

Examples of Compound Statements

- ▶ "If you finish your homework, then you can play the game."
- ▶ "This is a question if and only if this is the correct answer."
- ▶ "I have read this and I understand the problem."

Connectives and Symbols

Connective	Symbol	Formal name
not	\neg	negation
and	\wedge	conjunction
or	\vee	disjunction
if ... then	\Rightarrow	conditional
... if and only if...	\Leftrightarrow	biconditional

Connective Or

- ▶ Note that the connective Or in logic is used in the inclusive sense, not the exclusive sense as in English.
- ▶ The logical statement

"It is raining or the sun is shining"

means

- ▶ it is raining, or
- ▶ the sun is shining, or
- ▶ it is raining and the sun is shining.

Conditional connective

Note: There are several ways to express the conditional statement

$p \rightarrow q$

- ▶ If p then q
- ▶ q if p
- ▶ p is sufficient for q
- ▶ q is a necessary condition for p
- ▶ p only if q

Translating English Sentences into Propositional Sentences

Let p be the statement "The moped is still driving" and let q be the statement "The traffic light is red".

► $p \vee q$:

"The moped is still driving or the traffic light is red (or both)."

Translating English Sentences into Propositional Sentences

Let p be the statement "The moped is still driving" and let q be the statement "The traffic light is red".

► $q \Rightarrow p$:

"The traffic light is red then the moped is still driving."



Translating English Sentences into Propositional Sentences

Let p be the statement "The moped is still driving" and let q be the statement "The traffic light is red".

► $\neg p \wedge q$:

"The moped isn't driving anymore and the traffic light is red."

Summary of Truth Values

- ▶ The truth value of a compound statement is determined from the truth values of its simple components under certain rules.
- ▶ These rules are summarized in a so-called **truth table**.
Precisely, truth table is a table describing all possible cases of a compound statement in term of truth values of its simple components.

For instance, the truth table for the negation of a simple statement p is given by

p	$\neg p$
T	F
F	T

Truth Value

- ▶ If p and q are statements, then the truth value of the statement $p \vee q$ is T except when both p and q have truth value F.
- ▶ The truth value of $p \wedge q$ is F except if both p and q are true.

Truth Table for Some Connectives

p	q	$p \vee q$	$p \wedge q$	$p \Rightarrow q$	$p \Leftrightarrow q$
T	T	T	T	T	T
T	F	T	F	F	F
F	T	T	F	T	F
F	F	F	F	T	T

Example 1

Complete the truth table!

p	q	$p \wedge \neg q$
T	T	
T	F	
F	T	
F	F	

Example 1

p	q	$p \wedge \neg q$		p	q	$p \wedge \neg q$
T	T	tFf	\Rightarrow	T	T	F
T	F	tTt		T	F	T
F	T	fFf		F	T	F
F	F	fFt		F	F	F

The lower case t and f were used to record truth values in intermediate steps.

Example 2

Complete the truth table!

p	q	r	$(p \Rightarrow q) \Rightarrow (q \vee r)$
T	T	T	
T	T	F	
T	F	T	
T	F	F	
F	T	T	
F	T	F	
F	F	T	
F	F	F	

Example 2 Continued

Continue to complete the truth table!

p	q	r	$(p \Rightarrow q) \Rightarrow (q \vee r)$
T	T	T	t
T	T	F	t
T	F	T	f
T	F	F	f
F	T	T	t
F	T	F	t
F	F	T	t
F	F	F	t

Example 2 Continued

Continue to complete the truth table!

p	q	r	$(p \Rightarrow q)$	\Rightarrow	$(q \vee r)$
T	T	T	t		t
T	T	F	t		t
T	F	T	f		t
T	F	F	f		f
F	T	T	t		t
F	T	F	t		t
F	F	T	t		t
F	F	F	t		f

Example 2 Continued

Continue to complete the truth table!

p	q	r	$(p \Rightarrow q)$	\Rightarrow	$(q \vee r)$
T	T	T	t	T	t
T	T	F	t	T	t
T	F	T	f	T	t
T	F	F	f	T	f
F	T	T	t	T	t
F	T	F	t	T	t
F	F	T	t	T	t
F	F	F	t	F	f

Remark Truth Tables

- ▶ When you're constructing a truth table, you have to consider all possible assignments of True (T) and False (F) to the component statements.
- ▶ Suppose the component statements are P , Q , and R . Each of these statements can be either true or false, so there are $2^3 = 8$ possibilities.

Tautology and Contradiction

- ▶ A statement that is always true is called logically true or a tautology.
- ▶ A statement that is always false is called logically false or a contradiction.

Logical equivalence

- ▶ We cannot construct any more than 16 truth tables involving two statements.
- ▶ This is because such a truth table has 4 rows and the truth value of each row is T or F.
- ▶ However, we can certainly construct more than 16 statements involving two statements.
- ▶ What happens is that many (in fact infinitely many) statements have identical truth tables.
- ▶ We say that the statements r and s are **logically equivalent** if their truth tables are identical.

Logical equivalence

p	q	$\neg p \vee q$
T	T	T
T	F	F
F	T	T
F	F	T

This is equivalent to the truth table of $p \Rightarrow q$.

Logical equivalence

- ▶ The statements r and s are equivalent if and only if $r \Leftrightarrow s$ is a tautology.

Implication

- ▶ We say that r **implies** s if s is true whenever r is true.
- ▶ If r implies s then we write $r \Rightarrow s$.
- ▶ Alternatively, r implies s if and only if the statement $r \Rightarrow s$ is a tautology.
- ▶ We say that s is **logically deducible** from r .
- ▶ For example, in a mathematical theorem the hypothesis implies the conclusion or the conclusion is deducible from the hypothesis.

Order of Precedence

Operator	Precedence
\neg	1
\wedge	2
\vee	3
\Rightarrow	4
\Leftrightarrow	5

Order of Precedence

- \neg has higher precedence than \wedge , i.e.

$$\neg p \wedge q$$

should be read

$$(\neg p) \wedge q$$

and not

$$\neg(p \wedge q)$$

Order of Precedence

- \wedge has higher precedence than \vee , i.e.

$$p \wedge q \vee r$$

should be read

$$(p \wedge q) \vee r$$

and not

$$p \wedge (q \vee r)$$

Order of Precedence

- ▶ \vee has higher precedence than \Rightarrow , i.e.

$$p \vee q \Rightarrow r$$

should be read

$$(p \vee q) \Rightarrow r$$

and not

$$p \vee (q \Rightarrow r)$$

Order of Precedence

- \Rightarrow has higher precedence than \Leftrightarrow , i.e.

$$p \Rightarrow q \Leftrightarrow r$$

should be read

$$(p \Rightarrow q) \Leftrightarrow r$$

and not

$$p \Rightarrow (q \Leftrightarrow r)$$

Commutative Laws

► $p \wedge q \Leftrightarrow q \wedge p$

► $p \vee q \Leftrightarrow q \vee p$

Associative Laws

$$\blacktriangleright p \wedge (q \wedge r) \Leftrightarrow (p \wedge q) \wedge r$$

$$\blacktriangleright p \vee (q \vee r) \Leftrightarrow (p \vee q) \vee r$$

Distributive Laws

- ▶ $p \wedge (q \vee r) \Leftrightarrow p \wedge q \vee p \wedge r$
- ▶ $p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$

Idempotent (or Tautology)

► $p \wedge p \Leftrightarrow p$

► $p \vee p \Leftrightarrow p$

Absorption

► $p \wedge (p \vee q) \Leftrightarrow p$

► $p \vee (p \wedge q) \Leftrightarrow p$

Complementation

$$\blacktriangleright p \wedge \neg p \Leftrightarrow F$$

$$\blacktriangleright p \vee \neg p \Leftrightarrow T$$

Law of Involution (Double Complementation)

$$\blacktriangleright \neg\neg p \Leftrightarrow p$$

Laws of de Morgan

$$\blacktriangleright \neg p \wedge \neg q \Leftrightarrow \neg(p \vee q)$$

$$\blacktriangleright \neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$$

Identity Elements

Disjunction:

$$\blacktriangleright p \vee T \Leftrightarrow T$$

$$\blacktriangleright p \vee F \Leftrightarrow p$$

Conjunction:

$$\blacktriangleright p \wedge T \Leftrightarrow p$$

$$\blacktriangleright p \wedge F \Leftrightarrow F$$

Logic gates

Some popular logic gates

- ▶ Negation gate
- ▶ AND gate
- ▶ OR gate
- ▶ NAND gate
- ▶ NOR gate
- ▶ XOR gate

Details are on the white board.

Rules of Inference

- ▶ An argument is a sequence of statements that end with a **conclusion**. An argument is **valid** if the conclusion follows from the truth of the preceding statements (**premises** or **hypotheses**).
- ▶ In propositional logic, an argument is valid if it is based on a tautology.
- ▶ Arguments that are not based on tautology are called **fallacies**.

Name	Rule of Inference	Tautology
Addition	$\frac{p}{\therefore p \vee q}$	$p \rightarrow (p \vee q)$
Simplification	$\frac{p \wedge q}{\therefore p}$	$(p \wedge q) \rightarrow p$
Modus ponens	$\frac{p \quad p \rightarrow q}{\therefore q}$	$p \wedge (p \rightarrow q) \rightarrow q$
Modus tollens	$\frac{\neg q \quad p \rightarrow q}{\therefore \neg p}$	$(\neg q) \wedge (p \rightarrow q) \rightarrow \neg p$
Hypothetical syllogism	$\frac{p \rightarrow q \quad q \rightarrow r}{\therefore p \rightarrow r}$	$(p \rightarrow q) \wedge (q \rightarrow r) \rightarrow (p \rightarrow r)$
Disjunctive syllogism	$\frac{\neg p \quad p \vee q}{\therefore q}$	$(p \vee q) \wedge (\neg p) \rightarrow q$

Given the hypotheses:

- ▶ "It is not sunny and is cold"
- ▶ "We go swimming only if it is sunny"
- ▶ "If we do not go swimming then we will play soccer"
- ▶ "If we play soccer then we will go home by sunset"

Show that these hypotheses lead to the conclusion: "We will go home by sunset".

Given the hypotheses:

- ▶ "If you send me an email, I will finish writing the program"
- ▶ "If you do not send email then I will go to bed early"
- ▶ "If I go to bed early then I will go jogging tomorrow morning"

Show that these hypotheses lead to the conclusion: "If I do not finish writing the program then I will go jogging tomorrow morning".

Some fallacies

- ▶ Fallacy of affirming the conclusion: $[(p \rightarrow q) \wedge q] \rightarrow p$
- ▶ Fallacy of denying the hypothesis: $[(p \rightarrow q) \wedge \neg p] \rightarrow \neg q$

Part 2: General Algebra

Lecture 14: Set theory review and relations

Sets

- ▶ A set L is a collection of objects, called *elements* of the set. A set can be represented by listing its elements between braces:

Ex.: $L = \{a, b, c, d\}$

- ▶ Notation:
 - ▶ An element belongs to a set: \in
 - ▶ An element does not belong to a set: \notin

Ex.: $a \in L$ and $z \notin L$

Sets

- Recurrence and order is not important:

Example:

$\{\text{red}, \text{blue}, \text{red}\} = \{\text{red}, \text{blue}\}$ *and*

$\{3, 1, 9\} = \{9, 3, 1\} = \{1, 3, 9\}$

- Two sets A and B are equal if and only if they have the same elements.

Ex.: $A = B$

Sets

- ▶ The elements must not have anything in common or be especially similar:

$$\text{Ex.: } L = \{a, 7, \{c, \text{red}\}\}$$

- ▶ A set with no elements is called *empty set*:

$$\emptyset = \{\}$$

- ▶ A set A is a subset of B , if all elements of A are in B :

$$A \subseteq B$$

Sets

- ▶ We call A a proper subset of B if $A \subseteq B$ but $A \neq B$, i.e. $A \subset B$, i.e. there is some element in B which is not in A .
- ▶ $A = B$ if and only if $A \subseteq B$ and $B \subseteq A$
- ▶ $\emptyset \subseteq$ all sets
- ▶ If the set is finite, its number of elements is represented by $|A|$

Characterization of Elements

- ▶ Infinite sets are written with dot notation:

Ex.: The set of natural numbers, $\mathbb{N} = \{0, 1, 2, \dots\}$

- ▶ To specify a certain property for the elements of the set:

$B = \{x \mid x \text{ has property } P_1 \text{ and } P_2 \text{ and } \dots\}$

Ex.: The set of odd natural numbers:

$B = \{x \mid x \in \mathbb{N} \text{ and } x \text{ not divisible by } 2\}$

Complement

► Complement:

Let A be a subset of the universal set U , i.e. $A \subseteq U$.

The complement \overline{A} of A in U is the set of elements that do not belong to A :

$$\overline{A} = \{x \mid x \in U \text{ and } x \notin A\}$$

Union

- Union:

$$A \cup B = \{x \mid x \in A \text{ or } x \in B\},$$

i.e. the set of elements that belong to either of A and B .

$$\text{Ex.: } \{1, 2, 3\} \cup \{3, 4, 5\} = \{1, 2, 3, 4, 5\}$$

Intersection

- Intersection:

$$A \cap B = \{x \mid x \in A \text{ and } x \in B\},$$

i.e. the result is the set containing the common elements of two sets:

$$\text{Ex.: } \{1, 2, 3\} \cap \{3, 4, 5\} = \{3\}$$

Difference (or Relative Complement)

- Difference:

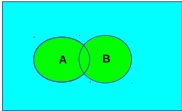
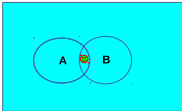
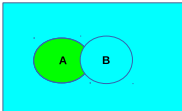
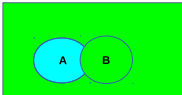
$$A - B = A \setminus B = \{x \mid x \in A \text{ and } x \notin B\},$$

i.e. the result is the set of elements that belong to a set but not to another:

$$\text{Ex.: } \{1, 2, 3\} - \{3, 4, 5\} = \{1, 2\}$$

- The sets A and B are disjoint, if $A \cap B = \emptyset$, i.e. if they have no common elements.

Example: Sets as Venn Diagrams

$A \cup B$	
$A \cap B$	
$A \setminus B$	
	

Cartesian Product

- ▶ The Cartesian product of two sets A and B is the set, which elements are ordered pairs (a, b) , with $a \in A$ and $b \in B$:

$$A \times B = \{(a, b) \mid a \in A \text{ and } b \in B\}$$

- ▶ *Two examples:*

If \mathbb{R} are the real numbers, then the Cartesian product

$\mathbb{R} \times \mathbb{R} = \mathbb{R}^2$ is the real plane.

The set $\mathbb{R} \times \mathbb{R} \times \mathbb{R} = \mathbb{R}^3$ is the real space.

Example: Cartesian Product

Let $A = \{1, 2, 3\}$ and $B = \{a, b\}$. Then is

$$A \times B = \{(1, a), (1, b), (2, a), (2, b), (3, a), (3, b)\}$$

Thus the Cartesian product has $3 \cdot 2 = 6$ elements.

Cartesian Product

- ▶ Notice: If $A \neq B$, then we have: $A \times B \neq B \times A$
- ▶ For finite sets A and B we have:

$$|A \times B| = |A| \cdot |B|,$$

since we can choose the first coordinate in the pair in $|A|$ ways and the second in $|B|$ ways.

The multiplication principle gives that the number of elements of $A \times B$ equals $|A| \cdot |B|$.

Computer representation of sets I

Let U be a universal set. Fix an arbitrary order of elements of U for instance a_1, a_2, \dots, a_n .

If A is a subset of U , represent A with a bit string of length n , where the i th bit is 1 if a_i is in A , and is 0 if a_i is not in A .

Example. Let $U = \{1, 2, 3, 4, 5, 6, 7, 8\}$. Then the subset

- ▶ $A = \{1, 3, 4, 6\}$ is represented as the bit string 10110100
- ▶ $B = \{4, 1, 3, 5\}$ is represented as the bit string 10111000
- ▶ $A \cap B$ is represented as the bit string (please check!):

$$10110100 \wedge 10111000 = 10110000.$$

Hence, $A \cap B = \{1, 3, 4\}$.

- ▶ $A \cup B$ is represented as the bit string (please check!):

$$10110100 \vee 10111000 = 10111100.$$

Hence, $A \cup B = \{1, 3, 4, 5, 6\}$.

Computer representation of sets II

Exercises:

1. Propose an algorithm to find the union, intersection, symmetric difference of the two given sets. What is the running time (complexity) of your algorithm?
2. Using the binary representation of subsets to prove that if $|A| = n$ then the power set $\mathcal{P}(A)$ of A is of cardinality 2^n .

Part 2: General Algebra

Relations

Topics covered:

- ▶ Relations and Their Properties
- ▶ Relation representations
- ▶ Equivalence Relations

What is a relation?

- ▶ Let A and B be sets. A (binary) relation from A to B is a subset of $A \times B$.
- ▶ A subset $A^2 = A \times A$ is called a relation in A .

Examples: Relation

1. Order relation \leq on real numbers: Let A be the set of real numbers. We define

$$R = \{(i, j) | i, j \in \mathbb{R} \text{ and } i \leq j\}.$$

Then R is a relation on \mathbb{R} .

2. Student-Grade relation: Let $A = \{Minh, Binh, Trung\}$ be the set of students and $B = \{1, 2, 3, 4\}$ the set of grade. Then

$$R = \{(Minh, 1), (Binh, 2), (Trung, 4)\}$$

is a relation from student set to grade set.

Examples: Relation

3. Congruence relation \equiv : Let $A = B$ be the set of integer numbers. Then

$$R = \{(a, b) \in \mathbb{Z} \times \mathbb{Z} : a \equiv b \pmod{n}\}$$

is a relation and it is called congruence relation on \mathbb{Z} .

4. "Divide" relation \parallel : Let $A = B = \mathbb{N}$ be the set of natural numbers. Then $R = \{(a, b) \in \mathbb{N} \times \mathbb{N} : a \mid b\}$ is a relation, and it is called "divide" relation.

Properties of Relations

Let R be a relation in set A .

► Reflexive:

$R \subseteq A \times A$ is reflexive, if $(a, a) \in R$ for all $a \in A$

Ex.: \leq

► Symmetric:

$R \subseteq A \times A$ is symmetric, if and only if $(b, a) \in R$ for $(a, b) \in R$

Bsp.: Brothers and sisters

Properties of Relations

Let R be a relation in A .

► Antisymmetric:

R is antisymmetric, if and only if $(b, a) \notin R$, if $(a, b) \in R$
(and a and b differ)

Ex.: Mother and son

► Transitive:

If $(a, b) \in R$ and $(b, c) \in R$, then we also have $(a, c) \in R$

Ex.: Forefathers (Child, Father, Grandfather)

What is an equivalence relation?

- ▶ An equivalence relation is a relation, which is reflexive, symmetric and transitive.
- ▶ Such a relation partitions the set A in disjoint subsets, called equivalence classes.

▶ *Ex.:*

$A =$ the set of different kinds of light bulbs.

One equivalence class: 40 W-light bulbs,
another equivalence class: 60 W-light bulbs.

Example: Equivalence Relation

- ▶ The relation *to be of the same age as somebody else* is an equivalence relation. What are the different equivalence classes?

Example. Check if the following relations are reflexive, symmetric, antisymmetric and transitive

- (1) R = set of pairs of students in a class that have the same birthday
- (2) R = set of pairs of two integers (a, b) with a divisible by b
- (3) R = set of pairs of real numbers (x, y) with $x + y = 0$
- (4) R = set of pairs of negative numbers (a, b) with $a > 2b$

Note. The relation R is transitive if and only if $R^2 \subseteq R$. (Explain!)

Another example for the properties of relations

Consider the set of all production animals on a farm. We define a relation as follows: Two animals have a relation if they are of the same species.

The cow Karin for example is in relation with the steer Olof, but not with the hen Helga.

This is an equivalence relation since it is

- ▶ reflexive: Every animal is of the same species as itself.
- ▶ symmetric: If animal 1 is of the same species as animal 2, then is also animal 2 of the same species as animal 1.
- ▶ transitive: If Karin and Lara are of the same species and Lara and Olof are of the same species, then are also Karin of the same species: All three are cattle.

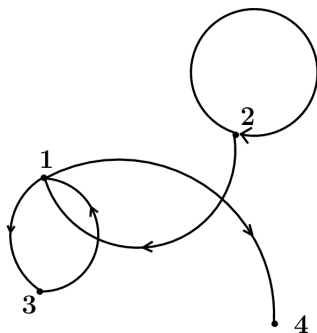
An equivalence class here consists of the animals belonging to a certain species.

Relation representations

Relations via directed graphs.

Let $R = \{(1, 3), (2, 2), (1, 4), (2, 1), (3, 1)\}$ on $A = \{1, 2, 3, 4\}$.

We can use digraphs to represent this relation.



Question. What are the properties of digraphs representing relations which are

- ▶ reflexive
- ▶ symmetric
- ▶ antisymmetric
- ▶ transitive

Operations on Zero-One Matrices I

- ▶ A matrix $A = (a_{ij})_{m \times n}$ is called a **zero-one matrix** if $a_{ij} \in \{0, 1\}$ for all $i = 1, 2, \dots, m$ and $j = 1, 2, \dots, n$.
- ▶ Zero-one matrices are used to represent discrete structure like graphs, relation in computer networks, network flow;
- ▶ Zero-one matrices are used to represent black-white pixels in image processing.

Let A and B be two 0-1 matrices of the same sizes:

$$A = [a_{ij}] = \begin{bmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{bmatrix}, \quad B = [b_{ij}] = \begin{bmatrix} b_{11} & \cdots & b_{1n} \\ \vdots & \ddots & \vdots \\ b_{m1} & \cdots & b_{mn} \end{bmatrix}.$$

Operations on Zero-One Matrices II

- ▶ The meet of A and B is $A \wedge B = [a_{ij} \wedge b_{ij}]$
- ▶ The join of A and B is $A \vee B = [a_{ij} \vee b_{ij}]$
- ▶ The symmetric difference of A and B is $A \oplus B = [a_{ij} \oplus b_{ij}]$
- ▶ The complement of A is $\bar{A} = [\bar{a}_{ij}]$

Example. Let $A = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$ and $B = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 0 \end{bmatrix}$. Then

- ▶ $A \wedge B = \begin{bmatrix} 1 \wedge 0 & 0 \wedge 1 & 1 \wedge 0 \\ 0 \wedge 1 & 1 \wedge 1 & 0 \wedge 0 \end{bmatrix} = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}$
- ▶ $A \vee B = \begin{bmatrix} 1 \vee 0 & 0 \vee 1 & 1 \vee 0 \\ 0 \vee 1 & 1 \vee 1 & 0 \vee 0 \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 0 \end{bmatrix}$
- ▶ $\bar{A} = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$

Boolean product $A \odot B$ I

Let A be an $m \times n$ zero-one matrix, B be an $n \times p$ zero-one matrix. The Boolean product of A and B , denoted by $A \odot B$, is an $m \times p$ zero-one matrix with the entry c_{ij} is defined by

$$c_{ij} = (a_{i1} \wedge b_{1j}) \vee (a_{i2} \wedge b_{2j}) \cdots \vee (a_{in} \wedge b_{nj}).$$

Example.

$$\begin{bmatrix} 0 & 1 & 0 \end{bmatrix} \odot \begin{bmatrix} 1 \\ 0 \\ 1 \end{bmatrix} = [(0 \wedge 1) \vee (1 \wedge 0) \vee (0 \wedge 1)] = [0]$$

and

$$\begin{bmatrix} 0 & 1 & 0 \end{bmatrix} \odot \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix} = [(0 \wedge 0) \vee (1 \wedge 1) \vee (0 \wedge 1)] = [1]$$

Boolean product $A \odot B$ II

Example.

$$\begin{aligned} & \begin{bmatrix} 0 & 1 & 0 \\ 1 & 1 & 0 \end{bmatrix} \odot \begin{bmatrix} 1 & 1 \\ 0 & 1 \\ 1 & 0 \end{bmatrix} \\ &= \begin{bmatrix} (0 \wedge 1) \vee (1 \wedge 0) \vee (0 \wedge 1) & (0 \wedge 1) \vee (1 \wedge 1) \vee (0 \wedge 0) \\ (1 \wedge 1) \vee (1 \wedge 0) \vee (0 \wedge 1) & (1 \wedge 1) \vee (1 \wedge 1) \vee (0 \wedge 0) \end{bmatrix} \\ &= \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \end{aligned}$$

Boolean powers I

Definition

Let A be a square zero-one matrix of size $n \times n$. The r th Boolean power of A , denoted by $A^{[r]}$, is the matrix

$$A^{[r]} = A \odot A \odot \cdots \odot A.$$

By convention, $A^{[0]} = I_n$.

Example: Let $A = \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix}$. Then (please verify)

$$\triangleright A^{[2]} = A \odot A = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{bmatrix}$$

$$\triangleright A^{[3]} = A^{[2]} \odot A = \begin{bmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}$$

Boolean powers II

$$\blacktriangleright A^{[4]} = A^{[3]} \odot A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

$$\blacktriangleright A^{[5]} = A^{[4]} \odot A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix}$$

Combining Relations

Let R and S be relations from A to B .

- ▶ $R \cup S$ = Union relation of R and S
- ▶ $R \cap S$ = Intersection relation of R and S
- ▶ $R \oplus S$ = Exclusive relation of R and S
- ▶ $R - S$ = Difference relation of R and S
- ▶ \overline{R} = Complementary relation of R (in $A \times B$)
- ▶ R^{-1} = Inverse relation of R , which consists of all pairs (b, a) where $(a, b) \in R$.
- ▶ Composite of two relations: Let R and S be relations

$$A \xrightarrow{R} B \xrightarrow{S} C.$$

The composition $S \circ R$ is the set of all pairs (a, c) such that there exists $b \in B$ with $(a, b) \in R$ and $(b, c) \in S$.

Relations as 0 – 1 matrices

Given a relation R from a set of m elements $\{a_1, a_2, \dots, a_m\}$ to a set of n elements $\{b_1, b_2, \dots, b_n\}$. To represent R we use a matrix of size $m \times n$, denoted by $M_R = [a_{ij}]$, whose entries are defined as follows:

$$a_{ij} = \begin{cases} 1 & \text{if } (a_i, b_j) \in R, \\ 0 & \text{if } (a_i, b_j) \notin R \end{cases}$$

Example: Let

$R = \{(apple, sour), (orange, sweet), (apple, sweet), (kiwi, sour)\}$
be a relation from $\{apple, orange, kiwi\}$ to $\{sour, sweet, bitter\}$.
Then, the boolean matrix M_R for R is

$$M_R = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 0 \\ 0 & 1 & 0 \end{pmatrix}.$$

Operations on relations and operations on boolean matrices

Theorem

Let R and S be relations on A with their representing matrices M_R and M_S . Then:

- ▶ $M_{R \cup S} = M_R \vee M_S$
- ▶ $M_{R \cap S} = M_R \wedge M_S$
- ▶ $M_{R \oplus S} = M_R \oplus M_S$
- ▶ $M_{\overline{R}} = \overline{M_R}$
- ▶ $M_{R^{-1}} = (M_R)^T$
- ▶ $M_{S \circ R} = M_R \odot M_S$
- ▶ $M_{R^n} = M_R \odot M_R \cdots \odot M_R =: M_R^{[n]}$

Proof.

On the whiteboard.



Question 1. What are the properties of the representing matrix of a relation which is:

- ▶ reflexive
- ▶ symmetric
- ▶ antisymmetric
- ▶ transitive

Question 2. Which properties: reflexive, symmetric, antisymmetric, or transitive that are possessed by the relation represented by the following matrix:

(a)
$$\begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

(b)
$$\begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

Question 3. Let A be a set of cardinality n . How many relations on A are there that are:

- ▶ reflexive?
- ▶ symmetric?
- ▶ reflexive and symmetric?

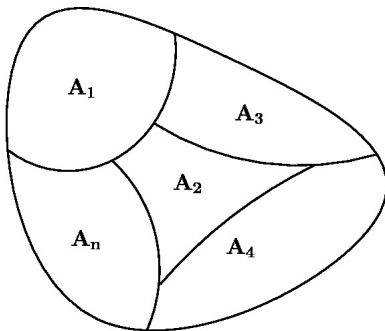
Equivalence relations and partitions of a set

Next, we will prove that an equivalence relation on a set can be considered as a set partition on that set.

Definition (**Partition of a set**)

Let A be a set. The non-empty subsets A_1, A_2, \dots, A_n are called a **partition into n blocks** of the set A if:

- (i) $A_1 \cup A_2 \cup \dots \cup A_n = A$
- (ii) $A_i \cap A_j = \emptyset$ for all $i \neq j$.



Partition of a set

Example.

Which of the following subsets form a partition of the set of all real numbers?

- (a) The subset of positive integers, the subset of negative integers;
- (b) The subset of non-positive integers, the subset of non-negative integers;
- (c) The subset of rational numbers, the subset of irrational numbers;
- (d) The closed intervals $[n, n + 1]$ where n is an integer;
- (e) The intervals $(n, n + 1]$ where n is an integer.

Equivalence Relations

Definition

The relation R on A is called an **equivalence relation** if it is reflexive, symmetric and transitive.

Example. The following relations are equivalence relations:

- (a) The congruence modular m relation on the set of integers;
- (b) The born-in-the-same-city relation in a class;

(c) The relation whose representing matrix is $M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$.

(d) $R = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (4, 4)\}$ on
 $A = \{1, 2, 3, 4\}$.

Equivalence relations and set partitions

Theorem

Let R be an equivalence relation on A . Then

1. The set A is partitioned into disjoint subsets, elements in each subset are mutually related to each other (and to them-self), and any two elements of two different subsets are not related by R .
2. These subsets are called *equivalence classes* of R .
3. Therefore, an equivalence relation on A forms a partition of A consisting of equivalence classes of R .

Conversely, let $P = \{A_1, A_2, \dots, A_n\}$ be a partition of A . Consider the relation

$R = \{(a, b) \in A \times A : \text{both } a, b \text{ are in the same block of } P\}$. Then R is an equivalence class on A . Moreover, the equivalence classes of R are the blocks of the partition P .

Examples

Find all equivalence classes of each equivalence relation.

- (a) The congruence modular m relation of the set of integers
- (b) The born-in-the-same-city relation in a class.

- (c) The relation whose representing matrix is $M_R = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$

- (d) $R = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 3), (4, 4)\}$ on
 $A = \{1, 2, 3, 4\}$

Examples

1. Define an equivalence relation whose equivalence classes are the subsets of the following partitions:
 - 1.1 Partition $\{1, 3, 5\}, \{2\}, \{4\}$ of $\{1, 2, 3, 4, 5\}$.
 - 1.2 Partition of the set of integers consisting of the subset of even numbers and the subset of odd numbers.
2. How many equivalence relations are there on the set $\{1, 2, 3, 4\}$?

Closures of Relations

Let R be a relation on A . The relation obtained from R by adding a minimal number of new pairs so that the new relation is reflexive (symmetric, antisymmetric, transitive) is called the reflexive closure (symmetric, antisymmetric, transitive closure) of R .

Example 1. Find the reflexive closure of:

- ▶ $R = \{(1, 2), (2, 3), (3, 3)\}$ on $A = \{1, 2, 3\}$.
- ▶ $R = \{(a, b) \mid a \neq b\}$ on the set of real numbers.
- ▶ $R = \{(a, b) \mid a \geq b\}$ on the set of real numbers.

Example 2. Find the symmetric closure of:

- ▶ $R = \{(1, 2), (2, 3), (3, 3)\}$ on $A = \{1, 2, 3\}$.
- ▶ $R = \{(a, b) \mid a > b\}$ on the set of real numbers.
- ▶ $R = \{(a, b) \mid a \neq b\}$ on the set of real numbers.

Example 3. Find the antisymmetric closure of
 $R = \{(1, 2), (1, 3), (2, 2), (2, 1)\}$ on $\{1, 2, 3\}$.

Transitive Closure

Let

$$R = \{(1, 3), (1, 4), (2, 1), (3, 2)\}$$

on the set $A = \{1, 2, 3, 4\}$.

The relation R is not transitive.

To get a transitive relation from R , we first need to add the pairs of R^2 .

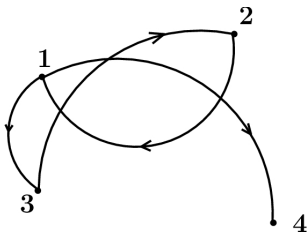
These new pairs together with those of R may form new ones, meaning we might need to add pairs of R^3, \dots

Denote R^* the transitive closure of R . Then R^* is the infinite union

$$R^* = R \cup R^2 \cup R^3 \cup \dots$$

Connectivity Relations

Let $R = \{(1, 3), (1, 4), (2, 1), (3, 2)\}$ on $A = \{1, 2, 3, 4\}$



- ▶ R^n = set of pairs of vertices (u, v) for which there is a path of length n from u to v .
- ▶ $R \cup R^2 \cup \dots \cup R^n$ = set of pairs of vertices (u, v) for which there is a path of length at most n from u to v .
- ▶ $R^* = R \cup R^2 \cup R^3 \cup \dots$ = set of pairs of vertices (u, v) for which there is a path from u to v .

An Algorithm for Computing Transitive Closures

Let R be a relation on a set of n elements. It is showed that to compute R^* we only need to compute a finite union of relations.

Theorem

Let R be a relation on a set of n elements. Then

$$R^* = R \cup R^2 \cup R^3 \cup \dots \cup R^n.$$

Equivalently, the matrix of R^* is determined from the following equation

$$M_{R^*} = M_R \vee M_R^{[2]} \vee M_R^{[3]} \vee \dots \vee M_R^{[n]}$$

Example. Find the transitive closure of

$R = \{(1, 3), (1, 4), (2, 1), (3, 2)\}$ on $A = \{1, 2, 3, 4\}$.

$$M_R = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$M_R^{[2]} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$M_R^{[3]} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$M_R^{[4]} = \begin{bmatrix} 0 & 0 & 1 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\text{Then } M_{R^*} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Warshall's Algorithm for Computing Transitive Closures

Let R be a relation on a set of n elements. To compute the transitive closure of R , Warshall's algorithm constructs a sequence of matrices M_0, M_1, \dots recursively:

- ▶ $M_0 := M_R$
- ▶ $M_k := M_{k-1} \vee$ (Boolean product of k th column and k th row of M_{k-1})

The algorithm terminates when $k = n$, and the matrix M_n is the matrix of the transitive closure of R .

Example. Use Warshall's algorithm to compute the transitive closure of $R = \{(1, 3), (1, 4), (2, 1), (3, 2)\}$ on $A = \{1, 2, 3, 4\}$.

Exercise. Read and understand the proof of correctness of Warshall's algorithm in textbook.

n -ary Relations and Their Applications

An n -ary relation on the sets A_1, A_2, \dots, A_n is a subset of $A_1 \times A_2 \times \dots \times A_n$

Exampe. Given the sets

$A_1 = \{\text{NamNT}, \text{TrungTT}, \text{HuongNTQ}, \text{ThaoNP}, \text{HienPQ}\},$

$A_2 = \{00198, 00011, 00345, 00786, 00321, 00546\},$

$A_3 = \{\text{MAD111}, \text{MAA101}, \text{MAS291}, \text{PFC111}, \text{DSA1}\}$

$A_4 = \{3.5, 4.0, 5.7, 8.0, 9.5, 6.4\}$

A relation R on these sets can be expressed as a database

Name	Code	Subject	Grade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

Name	Code	Subject	Grade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

We say R is a relation on 4 domains (Name, Code, Subject and Grade).

An element of this relation is a record consisting of 4 fields. For example

(TrungTT, 00786, DSA1, 3.5)

Primary key - Composite key

- ▶ A domain of an n -ary relation is a **primary key** if the value from this domain of the n -tuple determines this n -tuple.
- ▶ A set of domains of an n -ary relation is a **composite key** if they determines uniquely n -tuples.

Example. In the 4-ary relation,

Name	Code	Subject	Grade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

- ▶ Each domain **Name** or **Code** can be primary key.
- ▶ Each domain **Subject** or **Grade** is not primary key, but together they form a composite key.

Operations on n -ary Relations

Selection operator S_C . Given relation

Name	Code	Subject	Grade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

- ▶ If C is the condition (Subject = "DSA1") then the selection S_C produces a 4-ary relation with 2 records

Name	Code	Subject	Grade
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

- ▶ If C is the condition (Subject = "DSA1") \wedge (Grade > 5.0) then the selection S_C produces a relation with how many records?

Projection operator P . Given the relation

Name	Code	Subject	Gtade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

- ▶ The projection $P_{2,3,4}$ produces a 3-ary relation with 3 records

Code	Subject	Grade
00345	MAS291	8.0
00786	DSA1	3.5
00321	DSA1	8.0

- ▶ The projection $P_{2,4}$ on the above relation will produce a relation of how many domains and how many records?

Join operator \Join . Given two relations

Name	Code	Subject
HuongNTQ	00345	MAS291
NamNT	00011	MAD121
TrungTT	00786	DSA1
HienPQ	00321	DSA1

Code	Subject	Grade
00345	MAS291	8.0
00786	DSA1	3.5
00546	MAD121	5.7
00321	DSA1	8.0

The join operator \Join_2 used on these relations will produce the relation

Name	Code	Subject	Grade
HuongNTQ	00345	MAS291	8.0
TrungTT	00786	DSA1	3.5
HienPQ	00321	DSA1	8.0

General Algebra

Lectures 17-21: Number theory and cryptography

Covered topics:

1. Divisibility and Modular Arithmetic
2. Greatest common divisors and least common multiples
3. Operations on \mathbb{Z}_n : Addition, multiplication, exponentiation
4. Applications:
 - a) Cryptography
 - b) Hash function: Self-study!
 - c) Check digit

Reference:

1. Rosen's textbook: Chapter 4 (Number Theory and Cryptography)
2. Bogart's textbook (e-book: [DM-for-CS-Bogart et.al.pdf](#)): Chapter 2 (Cryptography and Number Theory)

The Integers and Division

Let a, b be integers with $a \neq 0$. We say the integer a divides b if there is an integer m such that $b = am$.

If a divides b we also write:

- ▶ b is divisible by a
- ▶ b is a multiple of a
- ▶ a is a factor of b
- ▶ $a \mid b$

Theorem

Let a, b, c be integers. Then:

- ▶ If $a \mid b$ and $a \mid c$ then $a \mid (b + c)$
- ▶ If $a \mid b$ then $a \mid bc$ for all c
- ▶ If $a \mid b$ and $b \mid c$ then $a \mid c$

The Division Algorithm I

Theorem (Euclid's division theorem)

Given two arbitrary integers m and n . Then there exist unique integers q and r such that $m = nq + r$ and $0 \leq r < n$.

Definition

1. In the division algorithm, m is called the **dividend**, n is the **divisor**, q is the **quotient** and r is the **remainder**.
2. The remainder r is also called m modulus n , denoted by $m \bmod n$. In other words, $m \bmod n$ is the smallest non-negative integer r such that $m - r$ divides by n .
3. The quotient in the division m by n is denoted by $m \div n$.

The Division Algorithm II

Example. Find the remainder and the quotient of the division:

- (a) -23 is divided by 7
- (b) -125 is divided by 11
- (c) $(-23).(23)$ is divided by 7
- (d) 2^{2016} is divided by 7 (find the remainder only)

Properties:

1. $(a + b) \bmod n = (a \bmod n + b \bmod n) \bmod n$
2. $(a \cdot b) \bmod n = (a \bmod n) \cdot (b \bmod n)$
3. $a^r \bmod n = (a \bmod n)^r \bmod n$

Greatest Common Divisors and Least Common Multiples I

A positive integer p greater than 1 is called a **prime** number if the only prime factors of p are 1 and p . An integer greater than 1 that is not prime is called **composite** number.

The Fundamental Theorem of Arithmetic

Any integer greater than 1 can be written uniquely as a product of powers of distinct primes.

Theorem

There are infinitely many primes.

Greatest Common Divisors and Least Common Multiples II

Examples and exercises

1. Which are prime, which are composite?
 - ▶ 119
 - ▶ $n! - 1$
 - ▶ $2^{2000} + 1$
 - ▶ $n^4 + 4$
 - ▶ $2^k - 1$
2. Show that if $2^n - 1$ is prime, then n is prime.
3. Determine whether the integers in each of these sets are pairwise relatively prime
 - ▶ 7, 8, 9, 11
 - ▶ 14, 15, 21
 - ▶ $2^{35} - 1$ and $2^{2015} - 1$

Greatest Common Divisors and Least Common Multiples III

4. The value of the Euler's totient function ϕ at a positive integer n is defined to be the number of positive integers less than or equal to n that are relatively prime to n . Find the following
- a) $\phi(4), \phi(10), \phi(13)$
 - b) $\phi(p^k)$ where p is a prime number and k is a positive integer
 - c) $\phi(pq)$ where p and q are prime numbers.
 - d*) $\phi(p_1^{a_1} p_2^{a_2} \dots p_n^{a_n})$ where p_i are distinct prime numbers.

GCD and LCM I

- ▶ Let a and b be two integers, not both 0. The greatest integer d that is a divisor of both a and b is called **greatest common divisor** of a and b , denoted by $\gcd(a, b)$.
- ▶ Let a and b be two positive integers. The smallest positive integer d that is divisible by both a and b is called the **least common multiple** of a and b , denoted by $\text{lcm}(a, b)$.

Theorem. Assume that prime factorizations of a and b are given

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}, \text{ and } b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n}.$$

Then:

1. $\gcd(a, b) = p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \cdots p_n^{\min(a_n, b_n)}$
2. $\text{lcm}(a, b) = p_1^{\max(a_1, b_1)} p_2^{\max(a_2, b_2)} \cdots p_n^{\max(a_n, b_n)}$

GCD and LCM II

Example: What are the greatest common divisors and the least common multiple of these pairs of integers?

1. $3^7 \cdot 5^3 \cdot 7^3$ and $3^9 \cdot 5^2 \cdot 11^2$

2. 1000 and 625

3. $3^{13} \cdot 5^{17}$ and $2^{12} \cdot 7^{21}$

4. 1111 and 0.

Theorem

Let a, b be positive integers. Then

$$ab = \gcd(a, b) \cdot \text{lcm}(a, b)$$

Relatively prime

- ▶ Two integers a, b are called **relatively prime** if $\gcd(a, b) = 1$.
- ▶ A set of integers are called **pairwise relatively prime** if any two integers in the set are relatively prime.

Example. Which sets are pairwise relatively prime?

(a) $(13, 24, 49)$

(b) $(14, 23, 35, 61)$

Euclid's Algorithm I

Theorem

Let $a > b$ be positive integers. Put $r = a \bmod b$. Then

$$\gcd(a, b) = \gcd(b, r).$$

Proof.

On the white board



Euclid's Algorithm II

Euclid's idea to find $\gcd(a, b)$: Using an iteration as follows.

Assume that $a \geq b$. Let $r_0 = a$ and $r_1 = b$. Applying the division algorithm successively we obtain

$$r_0 = r_1 q_1 + r_2 \qquad 0 \leq r_2 < r_1 \quad a \text{ is divided by } b$$

$$r_1 = r_2 q_2 + r_3 \qquad 0 \leq r_3 < r_2$$

$$\vdots$$

$$r_{n-1} = r_n q_n \qquad r_{n+1} = 0$$

The procedure stops at the step n where $r_{n+1} = 0$ and $\gcd(a, b) = r_n$.

Euclid's Algorithm III

Remarks:

1. Because of the strict decreasing of the sequence of remainders

$$r_0 > r_1 > \cdots \geq 0,$$

eventually the algorithm stops at a step n such that $r_{n+1} = 0$.

2. Correctness of Euclid's algorithm:

$$\gcd(a, b) = \gcd(r_1, r_2) = \cdots = \gcd(r_{n-1}, r_n) = \gcd(r_n, 0) = r_n.$$

Euclid's Algorithm IV

Example: Find $\gcd(91, 287)$

$$287 = 91 \cdot 3 + 14$$

$$91 = 14 \cdot 6 + 7$$

$$14 = 7 \cdot 2 + 0$$

Hence,

$$\gcd(287, 91) = \gcd(91, 14) = \gcd(14, 7) = \gcd(7, 0) = 7.$$

Euclid's Algorithm V

Procedure GCD(a, b : positive integers)

$x := a$

$y := b$

while $y \neq 0$

$r := x \bmod y$

$x := y$

$y := r$

Print(x)

GCD as a linear combination I

Theorem (Bezout's Theorem)

If a and b are positive integers, then there exists x and y such that

$$\gcd(a, b) = ax + by.$$

Based on Euclid's division, we can find an explicit way to express \gcd as a linear combination of original numbers.

Examples: Express $\gcd(91, 287) = 7$ as a linear combination of 91 and 287. By Euclid's division:

$$287 = 91 \cdot 3 + 14$$

$$91 = 14 \cdot 6 + 7$$

$$14 = 7 \cdot 2 + 0$$

GCD as a linear combination II

Using the next-to-last division, we have

$$7 = 91 - 6 \cdot 14$$

$$14 = 287 - 3 \cdot 91$$

Hence, $7 = 91 - 6 \cdot (287 - 3 \cdot 91) = (-6) \cdot (287) + 10 \cdot 91$.

Exercises.

1. Propose or writing a computer program to find the multiplicative inverse of a number a in \mathbb{Z}_n .
2. Using Euclid's algorithm to find the greatest common divisors of the following pairs

▶ $\gcd(12, 18)$

▶ $\gcd(111, 201)$

▶ $\gcd(210, 126)$

▶ $\gcd(1000, 5040)$

▶ $\gcd(12345, 54321)$

GCD as a linear combination III

- Express the greatest common divisor as a linear combination of original numbers

▶ 21, 44

▶ 36, 48

▶ 34, 55

▶ 117, 213

Arithmetic modulo n

Let \mathbb{Z}_n be the set of all possible remainders when dividing a number by n . Then

$$\mathbb{Z}_n = \{0, 1, \dots, n-1\}.$$

We define the addition and multiplication on \mathbb{Z}_n as follows:

$$i +_n j = i + j \mod n \quad (3)$$

$$i \cdot_n j = i \cdot j \mod n, \quad (4)$$

where the addition and the multiplication on the right-hand sides of these equations are the ordinary addition and multiplication of integers respectively.

Examples: Find

▶ $10 +_{12} 5 = 10 + 5 \mod 12 = 3;$

▶ $10 \cdot_{12} 5 = 10 \cdot 5 \mod 12 = 2;$

▶ $3 \cdot_{12} 2 =$

▶ $3 \cdot_{12} 6 =$

▶ $3 \cdot_{12} 10 =$

▶ a such that $3 \cdot_{12} a = 1$

Groups

Let G be a non-empty set.

- ▶ A binary operation $*$ on G is a map

$$* : G \times G \rightarrow G$$

$$(a, b) \mapsto a * b$$

- ▶ A group is a pair $(G, *)$ of set G together with a binary operation $*$ on G that satisfies the following properties
 1. G is closed under the operation $*$: $a * b \in G, \forall a, b \in G$.
 2. $*$ is associative, i.e., $a * b * c = a * (b * c)$
 3. G has an element e , e is called identity element of G , such that

$$e * a = a * e = a, \quad \text{for all } a \in G.$$

4. For each $a \in G$, there exist an element b such that

$$a * b = b * a = e.$$

Then b is called the inverse of a , and denoted by a^{-1} .

- ▶ The number of element of a group $(G, *)$ is called the order of G . Denoted by $|G|$.

Examples I

1. The set of integer numbers with standard addition $(\mathbb{Z}, +)$ forms a group with the identity 0 and the inverse of a is $-a$.
2. The set $(\mathbb{Z}_n, +_n)$ forms a group, where $e = 0$ and

$$a^{-1} = \begin{cases} n - a, & \text{if } a \neq 0 \\ a, & \text{if } a = 0 \end{cases}.$$

For instance, on $(\mathbb{Z}_7, +_7)$, we have

$$5 +_7 6 = 5 + 6 \mod 7 = 4 \in \mathbb{Z}_7,$$

and $5^{-1} = 2$.

3. The set of invertible matrices of size $n \times n$ together with matrix multiplication forms a group where the identity element is the identity matrix and the inverse of a matrix A is its inverse matrix A^{-1} .
4. The set of non-zero rationals \mathbb{Q}^* with the standard multiplication forms a group with identity $e = 1$ and $\left(\frac{p}{q}\right)^{-1} = \frac{q}{p}$ for any $p, q \in \mathbb{Z}^*$.

Examples II

5. The set (\mathbb{Z}_n, \cdot_n) is not a group since there exist some elements, for instance 0, not invertible by multiplication modulo n .
For example, on (\mathbb{Z}_6, \cdot_6) , we have

$$2 \cdot_6 4 = 2 \cdot 4 \mod 6 = 2,$$

and the elements 0, 2, 3, 4 are not invertible (see the table of multiplication modulo 6 below).

Examples III

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

Table: Multiplication table on \mathbb{Z}_6

Multiplicative group $(\mathbb{Z}_n^\times, \cdot_n)$

Let \mathbb{Z}_n^\times be the set of all invertible elements of \mathbb{Z}_n under the multiplication \cdot_n .

Example:

- ▶ $\mathbb{Z}_6^\times = \{1, 5\}$
- ▶ $\mathbb{Z}_8^\times = \{1, 3, 5, 7\}$
- ▶ $\mathbb{Z}_{12}^\times = \{1, 5, 7, 11\}$

Theorem

Given integer number n . Then

1. an element a of \mathbb{Z}_n is multiplicatively invertible if and only if $\gcd(n, a) = 1$.
2. The set $(\mathbb{Z}_n^\times, \cdot_n)$ is a group of order $\phi(n)$.

Remark: The inverse of an element a in \mathbb{Z}_n^\times can be found using Euclid's algorithm based on Bezout's theorem.

Examples

On the whiteboard!

Cryptography: Basic definitions I

Cryptography is the study of methods to send and receive secret messages and used mainly for military and diplomatic communications. List of terms

1. **Sender** (often called **Alice**) who sends messages
2. **Receiver** (often called **Bob**) who receives messages
3. **Adversary** who wants to steal the information of messages
4. **Plaintext** is the original message
5. **Ciphertext** is the encoded message
6. **Secret code, encode, decode, crack, codebook ...**

Cryptography: Basic definitions II

Two kinds of cryptography:

1. The **private key cryptography**: the sender and receiver agree in advance on a secret code and then send messages using that code;
2. The **public key cryptography**: the encoding method can be published. Each person has a public key used to encrypt messages and a secret key used to decrypt an encrypted message

Caesar cipher: Cryptography using addition mod n

Definition

A **Caesar cipher** is one in which each letter of the alphabet is shifted by a fixed amount.

Example:

- ▶ Alice will encrypt messages by shifting each letter three letter forward in the alphabet.

Plaintext: MEET YOU AT THE PARK

Ciphertext: PHHW BRX LQ WKH SDUN

- ▶ For decrypting to get original message, the receiver (Bob) will shift back each letter three letters in the alphabet, for instance, the original message recovered from the ciphertext RQH LI EB ODQG DQG WZR LI EB VHD is "to be honest, I do not know the answer, please help me".
- ▶ How could the adversary do to decode? In the worst case, just try every ways of shifting. Hence, it could take 26 tests!

Caesar cipher: Cryptography using addition mod n II

Mathematical form of the Caesar cipher is

$$m + a \pmod{26},$$

where m is original message, a is the number of places shifting.

Remark: Even for more complicated encryption and with the support of computer or if the adversary could steal the codebook, he would decode the message easily! This leads to another way for cryptography called Public Key Cryptography!

Cryptography using multiplication mod n I

The encryption by using multiplication mod n is written in mathematical form as a function

$$\begin{aligned} P : \mathbb{Z}_n &\rightarrow \mathbb{Z}_n \\ m &\mapsto m \cdot a \pmod{n}, \end{aligned}$$

where a is a pre-chosen number.

Example 1: Take $n = 7$, $a = 5$.

- ▶ The message $m = 6$ is encrypted by $6 \cdot 5 \pmod{7} = 2$. Hence, Alice will send message $P(m) = 2$ to Bob. The private key in this case is n and a .
- ▶ For decrypting Bob will find a number m (often less than n) such that

$$(m \cdot 5) \pmod{7} = 2.$$

What is the plaintext that Bob can find?

Cryptography using multiplication mod n II

Example 2: Take $n = 12$, $a = 4$. What is the encrypted message that Alice sent if the plaintext is $m = 5$. How will Bob decode if he received the message $m' = 8$?

Remark: In cryptography,

1. encoding function f should be easy to compute;
2. decoding must give the unique solution, easy to compute with the support of private key;
3. cracking (without key) should be "difficult" / take time;
4. decoding in the cryptography using addition mod n is to find the inverse function of the addition mod n . With the key is the number of place shifting a : $f^{-1}(m) = m - a \bmod n$
5. decoding in the cryptography using multiplication mod n is to find a' such that $a \cdot a' \bmod n = 1$.

Public key cryptography

RSA cryptography will be discussed if time is available (without slides).

Multiplicative inverse I

	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Table: Multiplication table on \mathbb{Z}_7

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	1	2	3	4	5
2	0	2	4	0	2	4
3	0	3	0	3	0	3
4	0	4	2	0	4	2
5	0	5	4	3	2	1

Table: Multiplication table on \mathbb{Z}_6

Multiplicative inverse II

Definition

An element a in \mathbb{Z}_n has a **multiplicative inverse** if there exists an element $a' \in \mathbb{Z}_n$ such that $a \cdot_n a' = 1$.

Multiplicative inverse of a (if exists) is denoted by a^{-1} .

Examples:

- ▶ Multiplicative inverse of 5 in \mathbb{Z}_7 is ...
- ▶ Multiplicative inverse of 5 in \mathbb{Z}_6 is ...
- ▶ Multiplicative inverse of 3 in \mathbb{Z}_7 is ...
- ▶ Multiplicative inverse of 3 in \mathbb{Z}_6 is ...
- ▶ The solution of the equation $3 \cdot_7 x = 5$ is ...
- ▶ The solution of the equation $3 \cdot_6 x = 3$ is ...
- ▶ The solution of the equation $3 \cdot_6 x = 2$ is ...

Multiplicative inverse III

Theorem

Given two integer numbers n and a .

1. The multiplicative inverse of a , if it exists, is unique;
2. Suppose that a has a multiplicative inverse a^{-1} in \mathbb{Z}_n . Then for any $b \in \mathbb{Z}_n$, the equation $a \cdot_n x = b$ has the unique solution $x = a^{-1} \cdot_n b$;
3. a has a multiplicative inverse in \mathbb{Z}_n if and only if $\gcd(a, n) = 1$.

Exercises

1. If $a \cdot 133 - m \cdot 277 = 1$, does this guarantee that a has a multiplicative inverse mod m ? If so, what is it? If not, why not?
2. Determine whether every nonzero element of \mathbb{Z}_n has a multiplicative inverse for $n = 10$ and $n = 11$. How many elements $a \in \mathbb{Z}_{10}$ such that $a \cdot_{10} 2 = 1$?
3. Using the inverse-solving in Euclid's division to find multiplicative inverse of

$$\triangleright 5 \in \mathbb{Z}_{11}$$

$$\triangleright 3 \in \mathbb{Z}_{10}$$

$$\triangleright 2 \in \mathbb{Z}_{10}$$

$$\triangleright 16 \in \mathbb{Z}_{103}$$

$$\triangleright 22 \in \mathbb{Z}_{31}$$

4. What is the value of the division $\frac{1}{4} \bmod 9$?
5. Knowing that Alice sent a message to Bob using multiplication mod $n = 103$ with the common key $a = 16$. Assume that Bob receives the message $m = 21$. What is the original message that Alice sent Bob?

Exponentiation mod n

The behind RSA encryption is the **exponentiation** in \mathbb{Z}_n . Let $a \in \mathbb{Z}_n$, we define

$$a^j \bmod n = \underbrace{a \cdot_n a \cdot_n \cdots_n a}_{j \text{ factors}}$$

the exponentiation mod of power j in \mathbb{Z}_n .

Theorem

Let p be a prime number. For any fixed non-zero number $a \in \mathbb{Z}_p$, the function

$$\begin{aligned} f : \mathbb{Z}_p &\rightarrow \mathbb{Z}_p \\ i &\mapsto i \cdot_p a \end{aligned}$$

is a bijection.

Theorem (Ferma's Little Theorem)

Let p be a prime number. Then $a^{p-1} \bmod p = 1$ in \mathbb{Z}_p for each non-zero $a \in \mathbb{Z}_p$.

Examples and exercises. Find

1. $7^{121} \bmod 13$

Answer. By Ferma's Little Theorem, we have $7^{12} \bmod 13 = 1$. Hence, by dividing 121 by 12 and applying the exponentiation rule modulo n we get

$$7^{121} \bmod 13 = 7 \cdot (7^{12})^{10} \bmod 13 = 7 \cdot 1^{10} \bmod 13 = 7.$$

Thus, $7^{121} \bmod 13 = 7$.

2. $2^{340} \bmod 11$

3. $3^{71} \bmod 13$

4. $5^{2003} \bmod 7$

5. $5^{2003} \bmod 11$

6. $5^{2003} \bmod 13$

7. $5^{2003} \bmod 7 \cdot 11 \cdot 13$.

Check digits applied for ISBN

1. All books are identified by an ISBN-10 or ISBN - 13 (International Standard Book Numbers).
2. ISBN-10 is a 10 digit code $x_1x_2 \dots x_{10}$ assigned by the publisher.
3. An ISBN-10 consists of blocks identifying the language, the publisher, the number assigned to the book by its publishing company, and finally, a check digit that is either a digit or the letter X (used to represent 10). This check digit is selected so that

$$\sum_{i=1}^{10} ix_i \mod 11 = 0.$$

Examples:

- ▶ The ISBN-10 of the 7th edition book "Discrete Mathematics and Applications" is 0 – 07 – 289905 – 0
- ▶ Is 0 – 84 – 930149 – X a valid ISBN-10?
- ▶ The first nine digits of the ISBN-10 for the fifth edition of the book "DMA" are 0 – 07 – 119881. What is its check digit?