Week 1 Discussion

CS 131 Section 1B 1 April 2022 Danning Yu

Administrivia

- About me
 - Danning Yu, danningyu@ucla.edu
 - Did undergrad at UCLA, currently an MS student
 - Took this class W21
- Office Hours: T 9-10am, 1-2pm https://ucla.zoom.us/j/6682857378
- Piazza: https://piazza.com/ucla/spring2022/cs131
- See syllabus for homework deadlines, grading policies, etc.
- Make sure you have a SEASnet account

Assignments

- 6 homeworks + project
 - Project is just another homework, basically
- Make sure everything runs on SEASnet servers
 - Inxsrv[11, 12, 13, or 15].seas.ucla.edu
- Personal opinion: HW2 and HW5 are hardest
- Don't copy solutions from online sources or each other
 - Discussing general ideas and approach is okay, don't share code with each other
 - Submissions will be checked for plagiarism
- HW1 due 4/8 11:55pm

Functional Programming

Functional Programming

- No "side-effects"
 - Calling a function twice with the same arguments gives identical result
 - "Pure function", behaving similar to mathematical functions
 - Side effects present to allow for I/O
- Functions are first class citizens
 - Functions can take other functions as arguments or return them as result (much like other variables)
- Iteration usually implemented with recursion

Why Functional Programming

- Similar ideas can be found in most modern programming languages
 - e.g. Python, C++, Java, Swift, Scala, Kotlin, and many more
 - Hint: if you're struggling with OCaml syntax, try doing the HW in Python first in a functional manner
 - We will see examples later when we cover other languages
- Functional programming makes compiling and testing easier
 - No side effects makes reasoning on the code easier
- Easier to build scalable systems
- Computing can be distributed on multiple machines more easily when there are no shared states or side effects

OCaml

OCaml

- Functional programming language
 - Statically typed: every value (including function) has a type
- Type checking at compile time
 - Catches a lot of errors!
- Type inference
 - In many cases, you don't need to specify types
 - Allows you to deduce if you coded your function correctly
- Garbage collection
- Compiled language
 - But includes an interactive interpreter
- Has object oriented aspect, but we won't use that in this class

OCaml Basics

- Hello, world: print string "Hello, world!\n";;
 - print string is the function, everything that follows are arguments
 - End statement with 2 semicolons when using interpreter (not needed in code files)
 - OCaml will output the return value
- Comments: (* this is a comment *)
- Local bindings ("variables"): let val1 = 5;;
 - Only usable within the scope of the function
- Functions: let <func_name> [args...] = <func_body>

```
# let average a b = # let average a b = (a + b) / 2;; (a +. b) /. 2.0;; val average : int -> int -> int = <fun> val average : float -> float -> float = <fun>
```

Recursive Functions

Recursive functions: add a rec keyword after let

```
o let rec <func_name> [args...] = <func_body>
```

```
# let rec factorial a =
    if a = 1 then 1 else a * factorial (a-1);;
val factorial : int -> int = <fun>
# factorial 5;;
- : int = 120
```

Mutual recursion: use the and keyword

```
let rec even x = 
if x = 0 then true
else odd (x - 1)
and odd x = 
if x = 0 then false
else even (x - 1)
```

Local Variables in Functions

 Add keyword in after the let statement to make the value available in the rest of the function

```
# let average a b =
   let sum = a +. b in
   sum /. 2.0;;
val average : float -> float -> float = <fun>
```

Lambda Functions

- Lambda functions (aka anonymous functions) are defined with keyword fun
 - Useful when supplying a (small) routine as a function argument

```
# (fun x -> x * x) 5;;
- : int = 25
```

The following 2 are equivalent

```
# let square x = x * x; # let square = fun x \rightarrow x * x; val square : int \rightarrow int = \langle fun \rangle
```

Confused? Here's lambda functions in Python

```
o (lambda x, y: x + y) (2, 3)
o high_ord_func = lambda x, func: x + func(x)
o high ord func(2, lambda x: x + 3)
```

OCaml Lists

- **Defining a list**: let x = [1; 2; 3; 4; 5];
- All elements are of same type
- Under the hood, lists can thought of as immutable singly-linked lists
 - Iteration and appending to the end is fast, random access is slow
- Non-empty list consists of a head and a tail
 - Head: the first element: tail: the rest of the list
 - List.hd x;; returns int = 1
 - List.tl x;; returns int list = [2; 3; 4; 5]
- We can add an element to the beginning of the list with cons "operator"::
- 0::[1; 2; 3] gives us [0; 1; 2; 3]
- 0::1::2::[3] also gives us [0; 1; 2; 3]
 - :: is right-associative
- 1::2 or [1]::[2] is not correct
- Lists are immutable: if you want to change a value in the list, need to create new list

Useful List Operations

- List.map: transforms each element with a function
- List.filter: returns a list that contains elements that matches the function
- List.rev: reverses the list
- List.for all: checks if all elements satisfy the specified function
- List.exists: checks if any element in the list satisfies the specified function
- Lambda functions can be useful in these operations

List.filter (fun x -> x < 3) [1; 2; 3; 4; 5];;

```
- : int list = [1; 2]

- : int list = [5; 4; 3; 2; 1]

[# List.for_all (fun x -> x < 3) [1; 2; 3; 4; 5];;
- : bool = false

[# List.for_all (fun x -> x < 6) [1; 2; 3; 4; 5];;
- : bool = true

- : int list = [5; 4; 3; 2; 1]

- : bool = false
```

[# List.rev [1; 2; 3; 4; 5];;

Pattern Matching

- "Better" version of switch statements
 - List possible cases in a cleaner way
- Sequential check from the beginning and find the first rule that matches
- Underscore is often used as a wildcard that matches any value
- Cleaner than conditional when there are many cases
- Compiler will tell you if there exist cases your rules can't match
 - Highly useful in catching bugs!
- Use conditionals with when keyword

```
# let is_zero x = match x with
    0 -> true
    | _ -> false;;
```

```
# let is_zero x =
   if x = 0 then true
   else false;;
```

```
# let first l = match l with
    head :: tail -> head
    | _ -> 0;;
val first : int list -> int = <fun>
# first [1;2;3];;
- : int = 1
```

```
# let rec factorial x = match x with
    x when x < 2 -> 1
    | x -> x * factorial (x - 1);;
val factorial : int -> int = <fun>
```

Tuples

- A collection of heterogeneous values
 - Values separated by commas
 - Typically surrounded by parentheses
- let x = "foo", 3.14, 'c', false;;

 val my tuple : string * float * char * bool
- Accessing tuple elements
 - Tuples with two elements: fst my tuple, snd my tuple
 - Pattern matching

```
# let (x, y) = (3, "foo");;
val x : int = 3
val y : string = "foo"
```

```
# let my_fst (a, b) = a;;
val my_fst : 'a * 'b -> 'a = <fun>
# my_fst (3, "foo");;
- : int = 3
```

Variants

- Another way to have user defined types
- Use as enumeration

```
# type color = Red | Green | Blue;;
type color = Red | Green | Blue
# let my_blue = Blue;;
val my_blue : color = Blue
```

Use as C/C++ like union

```
# let my_print x =
    match x with
    A a -> print_int a
    | B b -> print_string b;;
val my_print : my_type -> unit = <fun>
# my_print (A 8);;
8- : unit = ()
# my_print (B "foobar");;
foobar- : unit = ()
```

Running OCaml From File

Pretend the following is in somecode.ml

- To run: ocaml somecode.ml
- Or: start OCaml interpreter, and copy paste everything in to run
 - Don't forget to add ; ;
- When you submit your code, ocaml <file name> should run without error

Context Free Grammars Homework 1

Context Free Grammars (CFGs)

- A grammar defines a language (programming or otherwise)
 - What sentences are valid for a language
- In programming languages
 - Grammar does not say what instructions in that language mean, it just defines the allowed syntax
 - We can check if print("Hello, World!") is valid, without defining what it does
 - Syntax, not semantics
- There are multiple types of grammars, for this homework, you only need to know context-free grammars
 - Covered in CS 181 as well

CFG Rules

- Symbol:
 - Terminal: symbol that can't be replaced by other symbols (e.g. +)
 - Nonterminal: symbol that can be replaced by other symbols (e.g. BINOP)
- Rule:
 - Defines how a nonterminal symbols can be replaced with other symbols
- Grammar contains a starting symbol (e.g. EXPR), and a set of rules

Sample grammar:

EXPR -> NUM

EXPR -> NUM BINOP NUM

BINOP -> +

BINOP ->
NUM -> 0

NUM -> 1

Possible strings generated by this grammar:

Unreachable Rules

- Sample grammar, now slightly modified
- The nonterminal symbol INCROP can never be reached from EXPR (starting symbol)
 - Thus, rules with INCROP on the left are unreachable rules in this grammar
 - In this homework, you need to remove all the unreachable rules according to grammar and starting point

Sample grammar:

```
EXPR -> NUM

EXPR -> NUM BINOP NUM

BINOP -> +

BINOP -> -

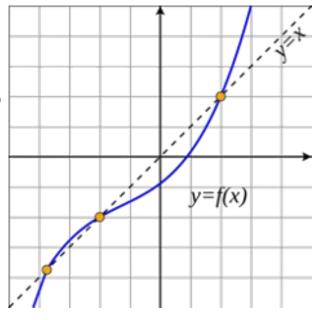
NUM -> 0

NUM -> 1

INCROP -> ++
```

HW1 Question Explanation

- Q1-5: Typical set operations, you can look up the definitions
 - Make sure you watch out for if versus iff
- Q6: Fixed point x, where x = f(x)
 - You can assume a fixed point can be obtained by repeatedly applying the function: f(x), f(f(x)), f(f(f(x))), ...
- Q7: A generalization of fixed point to p > 1
- Q8: Related to questions 7 and 8
- Q9: Copy the following definition into your code



HW1 Tips

- Lists are used to represent sets
 - Sets are a unique collection of elements, but lists can contain duplicates!
- Read through and understand the test cases
 - If problem definition seems unclear, the test cases may help you understand what the expected behavior is
- Read through documentation of List and Stdlib modules
 - Lots of useful functions in there!
 - Problems 1-8 can be written without too much code if you use List and Stdlib modules
- You're allowed to (and encouraged to) reuse functions from earlier problems
- All functions must be free of side effects
- Start early, ask questions on Piazza if you get stuck

Thank You