# Week 2 Discussion

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#### **Announcements**

- HW1 due tonight (4/8)
- HW2 released, due 4/21
  - Start early! It's harder than HW1
- Homeworks should be submitted on BruinLearn, under Assignments
- Before submitting
  - Make sure your code compiles on SEASnet server
  - Make sure your function signatures are correct
  - Follow all instructions and specifications
  - Do not submit files in a .zip unless told to do so
- Help and starter code from past TAs
  - https://github.com/CS131-TA-team

# **OCaml**

### Option

- Used to express possibly "invalid" values, such as null pointers, divide-by-zero
- Force proper handling of invalid case when one might exist

```
type 'a option =
    | Some of 'a
    | None
```

#### Example

```
let divide a b = match b with
  |0.0 -> None
  | _ -> Some (a /. b)
```

### Recursive Types

- type keyword can also define recursive data structures
  - Helpful in defining tree-like data structures

```
type tree =
    | Leaf of int
    | Node of tree * tree;;
let my_tree = Node (Node (Leaf 1, Leaf 2), Node (Leaf 3, Leaf 4))
```

How to find leftmost leaf in the tree?

### **OCaml Functions**

```
let sum a b = a + b;;
sum 1 2;;
-: int: = 3
```

- No need for parentheses unless you want to change order of evaluation
  - o sum 1 2\*3;; returns 9
    o sum 1 (2 \* 3);; returns 7
- Lambda function: (fun x  $\rightarrow$  x + 1) 5;;
- Equivalent forms
  - o let add\_one x = x + 1;;
    o let add one = fun x -> x + 1;;

### Currying and Partial Application (1 of 2)

- Internally, all functions in OCaml only have one argument
- Multi-argument functions
  - When supplied with one argument, it returns a new function that expects remaining arguments
- The "real behavior" of add with 2 arguments:

```
o let add = fun a -> (fun b -> a + b)
```

- This is called currying: expressing a function with multiple arguments as a series of functions that takes one argument each
  - The nature of functions in OCaml (and many other functional languages)
- Normal way of expressing and using multi-argument functions is provided for our convenience

### Currying and Partial Application (2 of 2)

• Implement add one with add:

```
let add a b = a + b
let add_one x = add 1 x
```

The x argument can be omitted

```
let add one = add 1
```

- This is called partial application
  - Supply the function with less argument it takes, and get a new function that expects remaining arguments
- Function return types
  - o add\_one has type int -> int
  - add takes an int (1), and returns add one (of type int -> int)
- Type of add: int -> (int -> int)
- -> is right-associative, yielding int -> int -> int, OCaml's output

#### List Module Functions

- filter map: A combination of filter and map
- fold\_left/fold\_right: Summarize a list into a single value
- flatten: Concatenate a list of lists into a single list
- split/combine: Conversion between list of tuples and tuple of lists

### List.filter map

• List.filter\_map f l applies f to every element of l, filters out the None elements and returns the list of Some elements.

```
let square_of_even l = List.filter_map (fun x ->
    if x mod 2 = 0 then Some (x * x) else None) l;;
val square_of_even : int list -> int list = <fun>
```

```
square_of_even [1;2;3;4;5];;
```

-: int list = [4; 16]

### List.fold left

- List.fold\_left f a [b1; ...; bn]Equivalent to f (... (f (f a b1) b2) ...) bn.
- Useful to accumulate over a list

```
\circ let list sum l = List.fold left (fun x y -> x + y) 0 l
```

- More concise version
  - o let list sum = List.fold left (+) 0
- Similarly, fold\_right exists
  - Different associativity
- These two operations are called "reduce" in some other languages

### List.flatten, List.split, List.combine

```
List.flatten [["a";"b"];["c";"d"];["e";"f"]];;
- : string list = ["a"; "b"; "c"; "d"; "e"; "f"]
List.split [("A", 1); ("B", 2); ("C", 3)];;
```

```
- : string list * int list = (["A"; "B"; "C"], [1; 2; 3])
```

```
List.combine ["A"; "B"; "C"] [1; 2; 3];;
-: (string * int) list = [("A", 1); ("B", 2); ("C", 3)]
```

### Pattern Matching with function

 function keyword can be used to simplify pattern matching on functions with single parameter

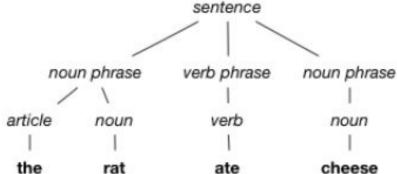
Further simplification if argument is a tuple

```
o let my_fst t = match t with (a, _) -> a
o let my fst (a, ) = a
```

# Homework 2

#### Homework 2

- Problem to solve: how to find the parse tree for a given sentence?
  - Given a context-free grammar
- Finding a parse tree is equivalent to finding the derivation
- Old homework is provided as a hint
  - http://web.cs.ucla.edu/classes/spring20/cs131/hw/hw2-2006-4.html
  - Includes sample code for a similar problem, can be used as a starting point



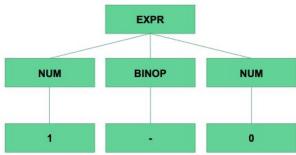
#### HW2 Problem 1

- Grammars are now represented as functions that take a non-terminal and return a list of rules
  - Allows us to access all the rules associated with a particular nonterminal
  - Write a function to convert an old grammar to the new style
- Note that your solution does not need to use pattern matching even though the example does! You need to write a function that takes a non-terminal symbol and returns a list of right-hand sides of rules

#### HW2 Problem 2

- Write a function parse\_tree\_leaves tree that returns a list of leaves in the tree from left to right
  - o In the example below, return ["1"; "-"; "0"]
- Basically, traverse all the leaves in a parse tree
- Tree data structure

```
type ('nonterminal, 'terminal) parse_tree =
   | Node of 'nonterminal * ('nonterminal, 'terminal) parse_tree list
   | Leaf of 'terminal
```



#### HW2 Problem 3 - Matcher

- Write a function make\_matcher gram that returns a matcher for the given grammar
- Matcher tells us if some prefix of the input can be generated using our grammar
  - Using our earlier grammar and input
    ["1"; "+"; "1"; "-"], we can match
    ["1"; "+"; "1"]
- We are not looking for the longest match just the first one that we find using our top-down parsing rules

### HW2 Problem 3 - Acceptor Function

- Given as an argument to the matcher
- Acceptor function takes a suffix and tells us whether it is "acceptable"
  - $\circ$  Some x if acceptable, None if not, where x is usually the suffix itself
- Lets us have additional conditions for matches, such as only accepting full matches
- If acceptor returns None, backtrack
  - If the last rule we tried didn't work

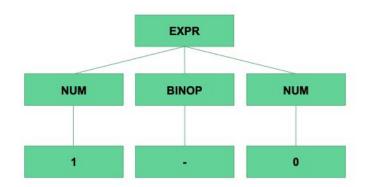
```
let accept_all string = Some string
let accept_empty_suffix = function
    | _::_ -> None
    | x -> Some x
matcher accept all ["9"; "+"; "$"]
```

### HW2 Problem 3 - Suggested Approach

- Expand parse tree using the top-down derivation rules from the earlier slide
- Once you have derived some prefix of the input string, call the acceptor with the unmatched input symbols (suffix)
  - o If the acceptor returns Some value, return that same value
  - o If the acceptor returns None, backtrack and try the next rule

#### Example:

- 1. Input is ["1"; "-"; "0";"+"]
- Use grammar rules until there are no non-terminals and all the terminals match with our input
- 3. Give unmatched suffix ["+"] to acceptor
- Return the acceptor's return value if the suffix was accepted; otherwise, try other derivations



#### HW2 Problem 4 - Parser

- Write a function make\_parser gram that returns a parser for the grammar gram.
- Parser differs from a matcher in two ways:
  - It does not take an acceptor as an input it will only return full matches
  - It returns a parse tree, not a suffix
- Example: parsing ["1"; "-"; "0"; "+"] must return None (non-empty suffix)

#### HW2 Problems 5 - 7

- 5. Write a non-trivial test case for make\_matcher; this should include writing your own grammar
- 6. Using the same grammar, write a test case for make parser
- 7. Write a report
  - Did you use make\_matcher to write make\_parser or vice versa or neither, and why?
  - Explain the weaknesses of your solution, provide test cases if possible
    - There will be some weaknesses, your parser/matcher might fail with a specific type of grammar for example

# Thank You