Mars: Accelerating MapReduce on Graphics Processors

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Mars: Accelerating MapReduce on Graphics Processors

What is Mars?

- A MapReduce Programming System, Map + Reduce.
- A Parallel Processing System accelerated by Graphics Processors (GPUs).
- Mars Modules running on:
 - An NVIDIA GPU: MarsCUDA
 - An AMD GPU: MarsBrook
 - A Multi-core CPU: MarsCPU
 - Multi-core CPUs + GPUs: Co-processing
 - Distributed System: MarsHadoop

Mars: Accelerating MapReduce on Graphics Processors

How Good?

- Ease of use. Up to 7 times code saving.
- High performance. An order of magnitude speedup over a state-of-the-art CPU-based MapReduce system.

Agenda

- Why Mars
 - GPGPU
 - MapReduce
- 2 How it works
 - Design
 - Implementation
- Second Second
 - Ease of use
 - High Performance
- Conclusion

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GPU Hardware Trend (1)

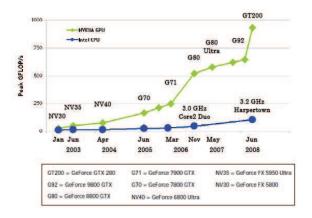


Figure: Floating-Point Operations per Second on NVIDIA GPUs and Intel CPUs.

Source: NVIDA CUDA Programming Guide [4].

GPU Chip

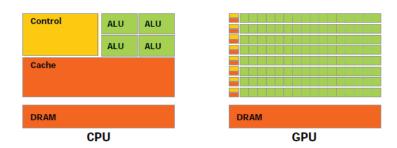


Figure: GPUs devote more transisters to data processing.

Source: NVIDA CUDA Programming Guide [4].

GPU Hardware Trend (2)

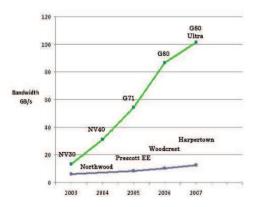
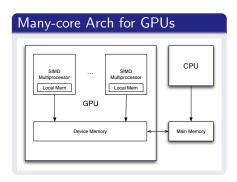


Figure: Bandwidth of NVIDIA GPU memory and CPU memory.

General Purpose GPU Computing



Programability

- NVIDIA CUDA
- AMD Brook+
- OpenCL
- More...

Non-Graphics Workloads on GPUs

Owens et al. [5] A Survey of General-Purpose Computation on Graphics Hardware

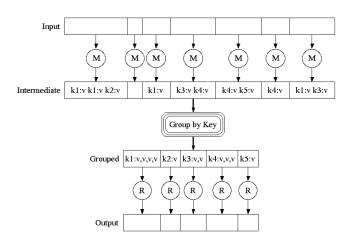
- Linear algebra
- Finance
- Database query
- Machine Learning
- More...
- Data Parallel programs on SIMD multiprocessors.

Map Function and Reduce Function

Jeffrey Dean and Sanjay Ghemawat, MapReduce: Simplified Data Processing on Large Clusters. OSDI'04. [2]

```
Map(void *doc) {
1: for each word w in doc
2: EmitIntermediate(w, 1); // count each word once
}
Reduce(void *word, Iterator values) {
1: int result = 0;
2: for each v in values
3: result += v;
4: Emit(word, result); // output word and its count
}
```

MapReduce Workflow



Source-http://labs.google.com/papers/mapreduce-osdi04-slides/index-auto-0007.html

Implementations of MapReduce

- Distributed Environment
 - Google MapReduce
 - Apache Hadoop (Yahoo, Facebook, ...)
 - MySpace Qizmt
- Multicore CPU
 - Phoenix from Stanford, HPCA'07 [6]/IISWC'09 [7].
- Cell BE
- FPGA
- GPUs
 - From UC-Berkeley, STMCS'08 [1]
 - Merge, from Intel, ASPLOS'08 [3]

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Goals and Challenges

Design Goals

- Programmability. Ease of use.
- Flexibility. Support various multi/many core processors.
- High Performance.

Challenges

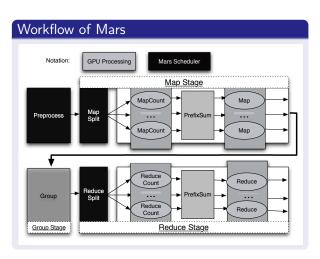
Result output.

- Write conflicts among GPU threads.
- Unknown output buffer size.

Solution

Lock-free scheme

Workflow



Customizing Workflow

- Map Only.
- Map→Group.
- Map→Group \rightarrow Reduce.
 - $Group \rightarrow Reduce.$
- Group.
 Map → Reduce.

Data Structure

Records

Input Records \rightarrow

Map Stage \rightarrow Intermediate Records I \rightarrow **Group Stage** \rightarrow Intermediate Records II \rightarrow **Reduce Stage**

 \rightarrow Output Records

Structure of Arrays

- Key array
- Value array
- Directory index array Variable-sized record
 - Key size, Key offset, Value size, Value offset>
- Chained MapReduce:
 Map1→Group1→Map2→Map3→Map4→Group4

Lock-Free Output

Lock Free

- MapCount
 - Call User defined MapCount function
 - Each function emits intermediate key size and value size
- Prefix sum on intermediate key sizes and value sizes
 - The size of intermediate buffer, allocate at one time
 - The deterministic write position for each Map, lock-free
- Allocate intermediate buffer
- Map
 - Call User defined Map function
 - Output records according to the write position

Lock-Free Output, Example

 $Map1 \rightarrow "123456789"$, $Map2 \rightarrow "abcd"$, $Map3 \rightarrow "ABCDED"$

MapCount

- MapCount1 \rightarrow 9
- MapCount2 → 4
- MapCount3 → 6

Prefix Sum, Allocate buffer, and Map

- 9, 4, 6 size array
- 0, 9, 13 write position array
- 19 output buffer size
- Allocate a buffer of size 19
- "123456789abcdABCDED"

MarsCUDA

Building blocks

- NVIDA CUDA
- Prefix Sum: CUDPP Library, GPU-based Prefix Sum
- Group: GPU-based Bitonic Sort

MarsCUDA – Memory Optimization (1)

Coalesced Access

For a half-warp of threads, simultaneous device memory accesses to consecutive device memory addresses can be coalesced into one transaction. \rightarrow Reduce # of device memory accesses.

Local memory

- Programmable on-chip memory (shared memory in NVIDIA's term).
- Exploit local memory in GPU-based Bitonic Sort for Group Stage.
- Users can explicitly utilize local memory in their Map/Reduce functions.

MarsCUDA – Memory Optimization (2)

Built-in Vector type

- Address Alignment
- float4 and int4
- One load instruction to read data of built-in type, of size up to 16 bytes → Reduce # of memory load instructions, compared with reading scalar type

Page-lock host memory

Prevent OS from paging the locked memory buffer \rightarrow High PCI-E bandwidth

MarsCUDA - Task distribution

Map/Reduce

$\lceil N/B \rceil$ thread blocks

- ullet N: the number of Map or Reduce tasks
- B: the number of GPU threads per thread block, which is practically set to 256
- 1 task per GPU thread

Special case for Reduce

- Communicative and Associative. For example, Integer Addition.
- Parallel reduction for load balanced reduce task distribution.

MarsCPU

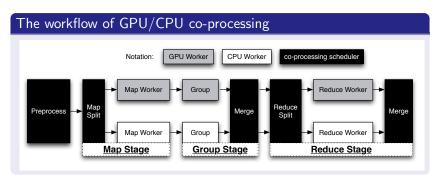
Building blocks

- pthreads
- Group: Parallel Merge Sort

General Mars Design

- Lock Free
- \bullet $\lceil N/T \rceil$ tasks per CPU thread.
 - ullet N: the number of Map or Reduce tasks
 - T: the number of CPU threads
 - ullet N is usually much larger than T

GPU/CPU Co-processing



- ullet I: Total size of input data
- S: Speedup of GPU Worker over CPU Worker
- Workload for GPU Worker: $\frac{SI}{1+S}$ • Workload for CPU Worker: $\frac{I}{1+S}$

MarsHadoop

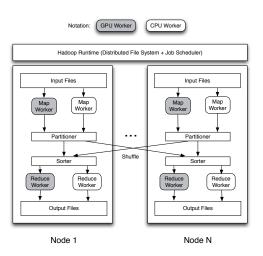


Figure: MarsHadoop. Using Hadoop Streaming.

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Experimental Setup

Machine	PC A	РС В	PC C
GPU	NVIDIA GTX280	NVIDIA 8800GTX	ATI Radeon HD 3870
# GPU core	240	128	320
GPU Core Clock	602 MHz	575 MHz	775 MHz
GPU Memory Clock	1107 MHz	900 MHz	2250 MHz
GPU Memory Band-	141.7 GB/s	86.4 GB/s	72.0 GB/s
width			
GPU Memory size	1024 MB	768 MB	512 MB
CPU	Intel Core2 Quad Q6600	Intel Core2 Quad Q6600	Intel Pentium 4 540
CPU Clock	2400 MHz	2400 MHz	3200 MHz
# CPU core	4	4	2
CPU Memory size	2048 MB	2048 MB	1024 MB
OS	32-bit CentOS Linux	32-bit CentOS Linux	32-bit Windows XP

Applications

Applications	Small	Medium	Large
String Match (SM)	size: 55MB	size: 105MB	size: 160MB
Matrix Multiplication (MM)	256×256	512×512	1024×1024
Black-Scholes (BS)	# option: 1,000,000	# option: 3,000,000	# option: 5,000,000
Similarity Score (SS)	# feature: 128, $#$ docu-	# feature: 128, $#$ docu-	# feature: 128, $#$ docu-
	ments: 512	ments: 1024	ments: 2048
PCA	1000×256	2000×256	4000×256
Monte Carlo (MC)	# option: 500, # samples	# option: 500, # samples	# option: 500, # samples
	per option: 500	per option: 2500	per option: 5000

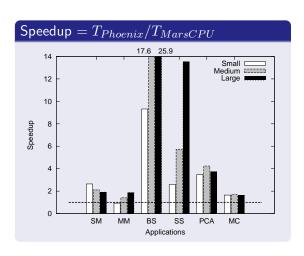
GPU Implementation: MarsCUDA, CUDA CPU Implementation: MarsCPU, Phoenix, pthreads GPUCPU Coprocessing: MarsCUDA + MarsCPU

Code size saving

In lines:

Applications	Phoenix	MarsCUDA/MarsCPU	CUDA
String Match	206	147	157
Matrix Multiplication	178	72	68
Black-Scholes	199	147	721
Similarity Score	125	82	615
Principal component analysis	297	168	583
Monte Carlo	251	203	359

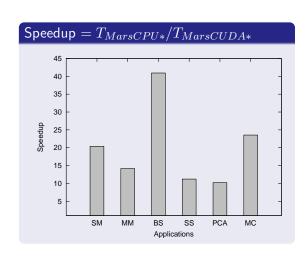
MarsCPU vs Phoenix



Overhead of Phoenix

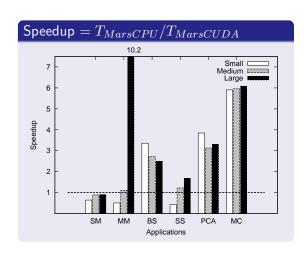
- Always need Reduce stage.
- Lock overhead.
- Re-allocate buffer on the fly.
- Insertion sort on static arrays. Call memmove() frequently.

MarsCUDA vs MarsCPU on Kernel



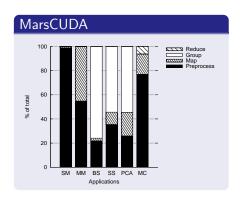
Preprocess + Map + Group + Reduce

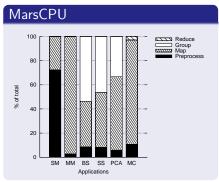
MarsCUDA vs MarsCPU



Preprocess + Map + Group + Reduce

Time Breakdown





Amdahl's Law

Amdahl's Law

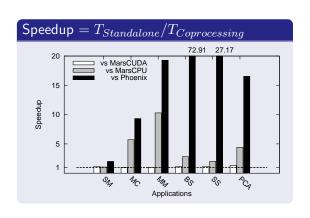
$$\mathsf{Speedup} = \frac{1}{(1-P) + P/S}$$

- *P*: The proportion that is parallelized
- (1 P): The proportion that is not parallelized
- ullet S: Speedup by parallelism

For MarsCUDA

- P: Map + Reduce
- (1 P): Preprocess
- Example: String Match
 - Parallelized: Map stage
 - P = 25%
 - S = 20
 - Speedup $= \frac{1}{(1-25\%)+25\%/20} = 1.3$

GPU/CPU Co-processing

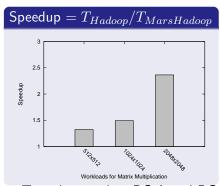


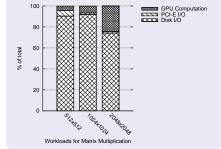
Co-processing over MarsCUDA:

- Speedup = $\frac{S+1}{S}$
- S: Speedup of MarsCUDA over MarsCPU

Time Breakdown

MarsHadoop





Two slave nodes: PC A and PC B

One master node: PC A

Reference



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In IISWC, 2009.

Conclusion

- Mars
 - MarsCUDA for NVIDIA GPU
 - MarsBrook for AMD GPU
 - MarsCPU for multicore CPU
 - GPU/CPU Co-processing
 - MarsHadoop for clusters
- Ease of programming
- High performance

Why Mars How it works Evaluation Conclusion

Thanks! Q&A? http://www.cse.ust.hk/gpuqp/Mars.html