

```

1 |----- MODULE TCommit -----|
  | This specification is explained in "Transaction Commit", Lecture 5 of the TLA+ Video Course. |
6 | CONSTANT RM      | The set of participating resource managers |
8 | VARIABLE rmState  | rmState[rm] is the state of resource manager rm. |
9 |-----|
10 | TCTypeOK  $\triangleq$ 
    | The type-correctness invariant |
14 |   rmState  $\in [RM \rightarrow \{\text{"working"}, \text{"prepared"}, \text{"committed"}, \text{"aborted"}\}]$ 
16 | TCInit  $\triangleq$    rmState =  $[r \in RM \mapsto \text{"working"}]$ 
    | The initial predicate. |
21 | canCommit  $\triangleq \forall r \in RM : rmState[r] \in \{\text{"prepared"}, \text{"committed"}\}$ 
    | True iff all RMs are in the "prepared" or "committed" state. |
26 | notCommitted  $\triangleq \forall r \in RM : rmState[r] \neq \text{"committed"}$ 
    | True iff no resource manager has decided to commit. |
30 |-----|
    | We now define the actions that may be performed by the RMs, and then define the complete |
    | next-state action of the specification to be the disjunction of the possible RM actions. |
36 | Prepare(r)  $\triangleq \wedge rmState[r] = \text{"working"}$ 
37 |    $\wedge rmState' = [rmState \text{ EXCEPT } ![r] = \text{"prepared"}]$ 
39 | Decide(r)  $\triangleq \vee \wedge rmState[r] = \text{"prepared"}$ 
40 |    $\wedge canCommit$ 
41 |    $\wedge rmState' = [rmState \text{ EXCEPT } ![r] = \text{"committed"}]$ 
42 |    $\vee \wedge rmState[r] \in \{\text{"working"}, \text{"prepared"}\}$ 
43 |    $\wedge notCommitted$ 
44 |    $\wedge rmState' = [rmState \text{ EXCEPT } ![r] = \text{"aborted"}]$ 
46 | TCNext  $\triangleq \exists r \in RM : Prepare(r) \vee Decide(r)$ 
    | The next-state action. |
50 |-----|
51 | TConsistent  $\triangleq$ 
    | A state predicate asserting that two RMs have not arrived at conflicting decisions. It is an |
    | invariant of the specification. |
56 |  $\forall r1, r2 \in RM : \neg \wedge rmState[r1] = \text{"aborted"}$ 
57 |    $\wedge rmState[r2] = \text{"committed"}$ 
58 |-----|
    | The following part of the spec is not discussed in Video Lecture 5. It will be explained in Video |
    | Lecture 8. |
63 | TCSpec  $\triangleq TCInit \wedge \Box [TCNext]_{rmState}$ 
    | The complete specification of the protocol written as a temporal formula. |

```

69 THEOREM $TCSpec \Rightarrow \Box(TCTypeOK \wedge TCConsistent)$

This theorem asserts the truth of the temporal formula whose meaning is that the state predicate $TCTypeOK \wedge TCInvariant$ is an invariant of the specification $TCSpec$. Invariance of this conjunction is equivalent to invariance of both of the formulas $TCTypeOK$ and $TCConsistent$.

78

\ * Modification History
\ * Last modified Tue May 04 13:49:33 SGT 2021 by nurliyana
\ * Created Tue May 04 12:12:04 SGT 2021 by nurliyana