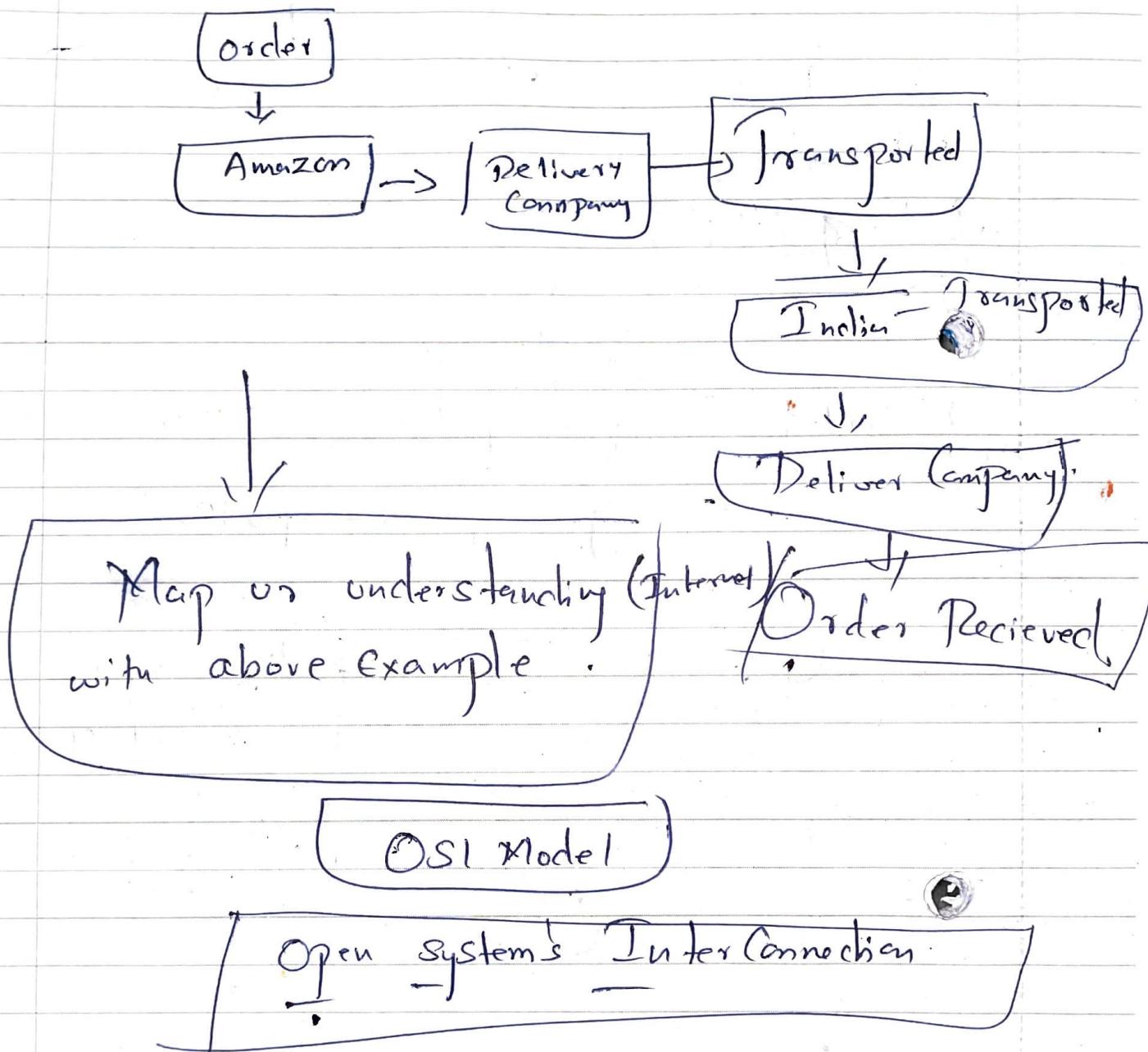
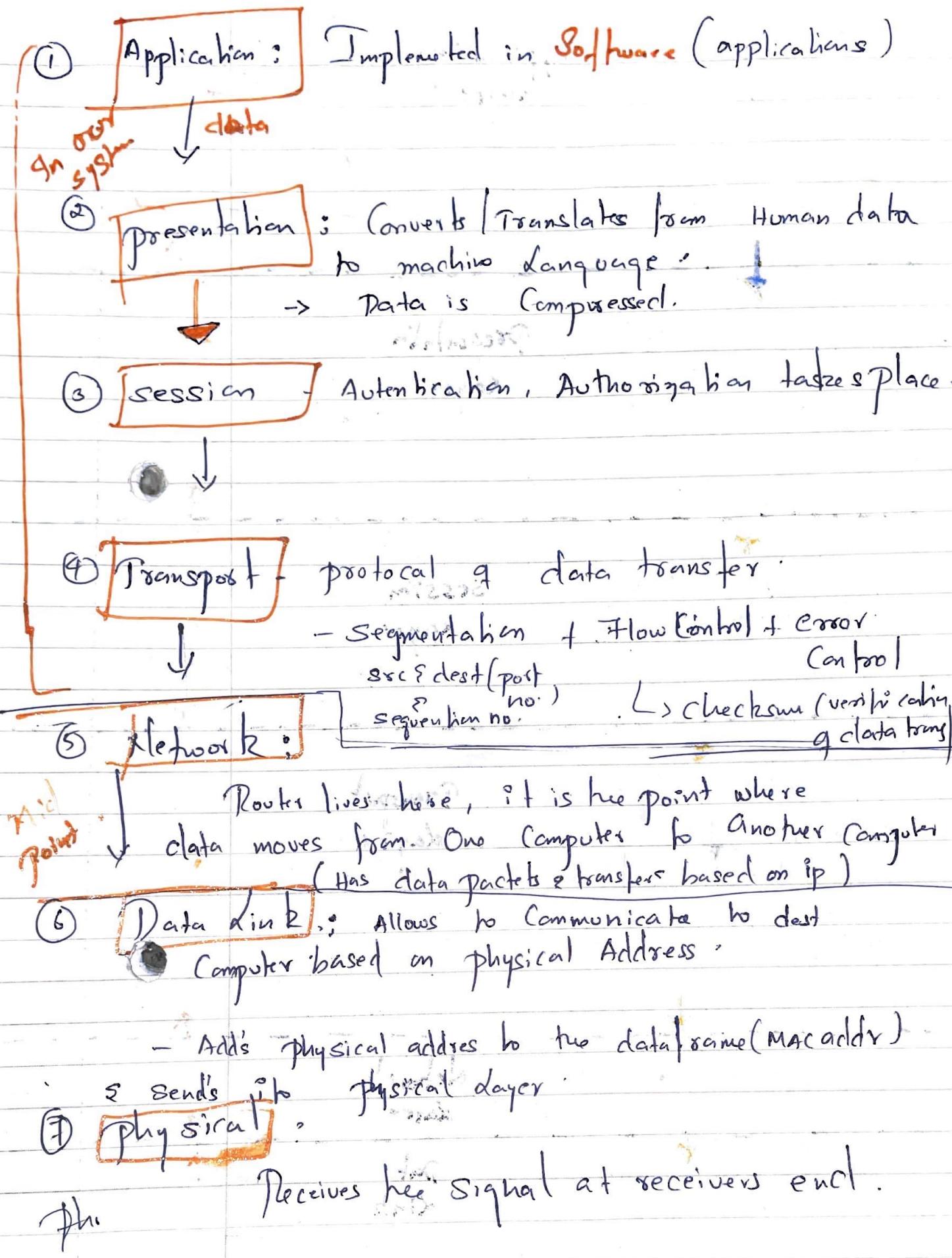


Structure of Network.

Example: Book ~~order~~ Online Ordering



7 Layers.



ME

Application
layer

MONTU

whatup's message is sent

"Hey Montu"

whatup's message is Received
"Hey Montu"

"Hey Montu" → Machine understandable language:
Data compressed.

checks for Authentication
Authorization of user,
data

Set the protocol to the data transfer.

1) Segmentation: SRC | dest · port
Sequence num

2) Flow Control

3) Error Code

A mid point b/w src | dest where transfer of data is made using IP layer

Adds the physical address to data frame & sent's to below layer

Data link

Checks the physical layer & transfers to dest

Physical layer

Hey Montu

Another Model: TCP/IP Model.

The Internet protocol Suite.

5 layers.

More practical.

Application

Transport

Netw

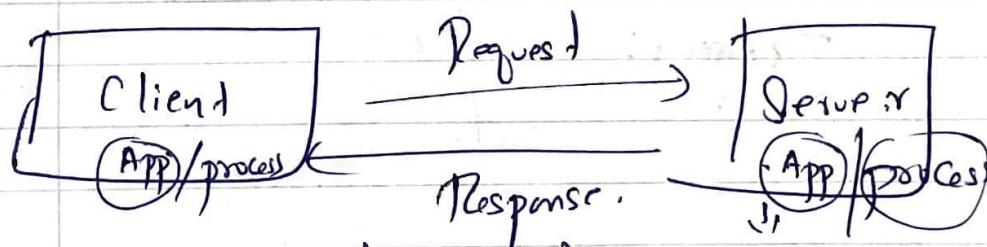
Data link

Physical

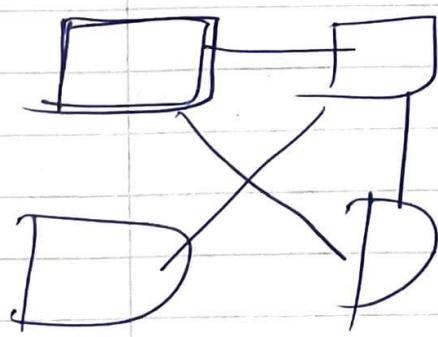
Application Layer:

- * User Interaction
- * For eg. Browser, WhatsApp etc.
- * Lies in our devices
- * Has some protocol
- * Client - Server Architecture.

① Client - Server Arch



② P2P arch. (Peer to peer)



→ Decentralized

→ Every single Computer is
Source / Client

Protocols

web protocols :

TCP | IP :

* HTTP : How HTML page transfer

Dynamic host config * DHCP : Allocates free ip address to connected pp

File transfer, rpm * FTP : files transfer

Simple mail transfer * SMTP : Emails transfer

* POP3 & IMAP : Receive mail

* SSH : Connecting to another system, terminal.

Virtual New Ctrl * VNC : Graphical Ctrl

* Telnet : Port 23

* Connect to free user, using Telnet.

* Not encoded & encrypted.

* UDP :

stateless Connection.

Program: what's app.

process

Send a Message

Record a video

Thread :

(Lighter version
of process)

↓
Set up
page.

↓
Open Camera.

(Single job)

Sockets

: Interface b/w process & Internet

Ports:

Ephemeral ports

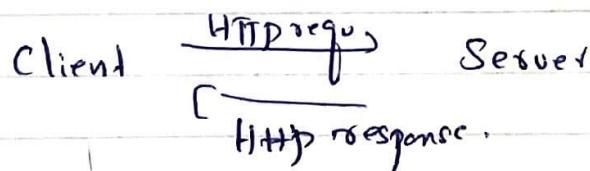
We know that port number determines the "app". But what if user has opened multiple tabs in window? How will know which port to communicate?

That's where Ephemeral ports comes into picture.

Created for Instance &
gets removed after Instance is
closed!

Exists on Client side.

HTTP: Application layer protocol



HTTP uses TCP inside App layer

App layer
↳ stateless protocol

↳ Transfer protocol.
↳ Transport Layer.

Method: Tell's us what to do.

① **GET** : Request the data.

② **POST** : update provide the Input data.

③ **PUT** : puts data at specific location

④ **DELETE** : Delete data

STATUS CODES / Error Code.

200 - Successful

400 - Bad Request

500 - Internal Server error

① 1xx → Information

2xx → Success

3xx → Redirecting

4xx → Client error

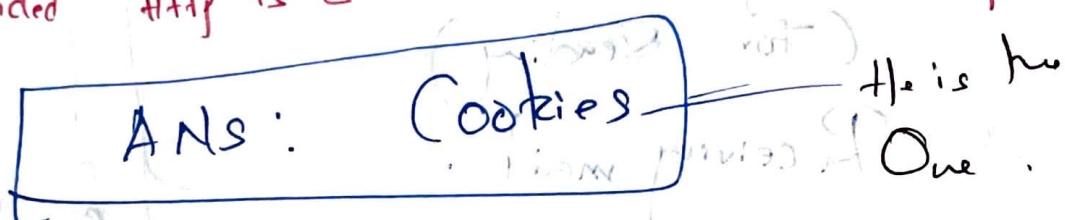
5xx → Server error

NETWORKING:

12/10/1

Scenario:

- Hey, am using my browser to open an Amazon shopping site, logs in and after some days when I try to open amazon site, am still logged in.
- Provided HTTP is stateless, how am I still logged in?



Cookies:

- Unique String → stored in Client browser.
- First time it will save in browser.

↓
Next time, when the request is made to the Server, A Cookies will be sent.

↓
Server: will check and

Auto login

In Inspect page (check for)

- Set Cookies }
- Cookies } In Network Section.

* 3rd party Cookies:

Cookies for websites which we aren't visiting but our cookies from website we visited is used for unvisited website.

How Email works?

SMTP : Simple Mail Transfer Protocol
(for sending)

POP3 : Receiving mail.

forward emails to brother <- point 2) going ON

This application key is how is email transferred

at abc@xyz@yahoo.com. abc@gmail.com
at abc@abc.abc.abc.abc and brother



Sender



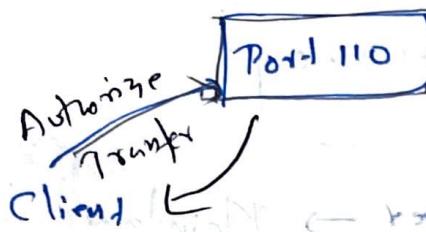
RECEIVER

Command to check name & IP of SMTP server?

nslookup -type=mx gmail.com

mail exchange

POP and Post Office protocol.

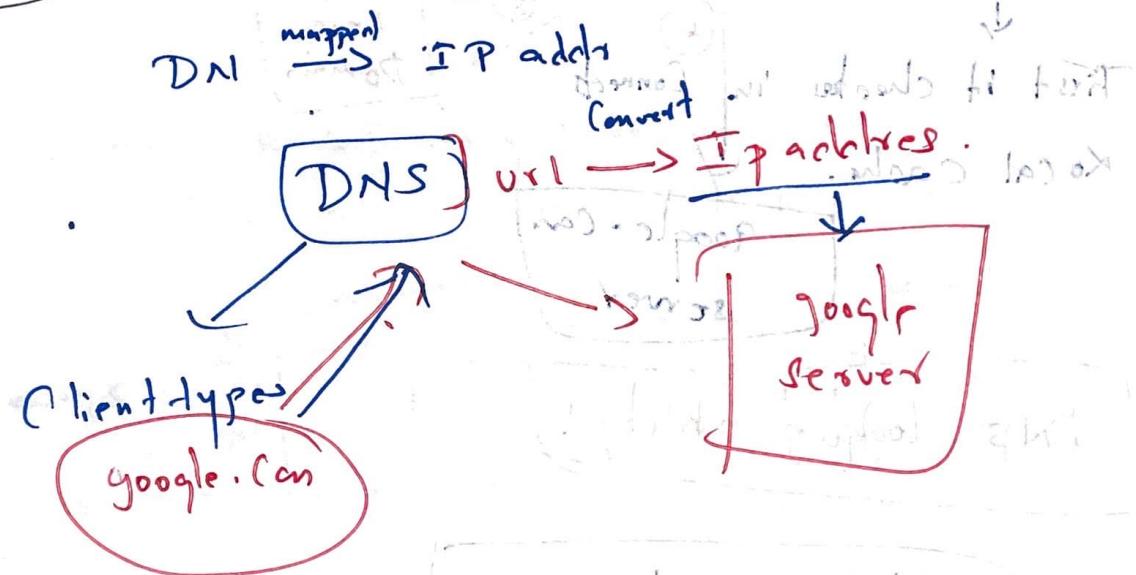


IMAP: Internet Message Access Protocol

↳ (For Receiving emails)

- Allows us to view message in multiple servers.
- Every device are in sync with server.

→ **DNS** - Domain Name System.



DNS → Director

Database which has

Domain names (URL). & ip's

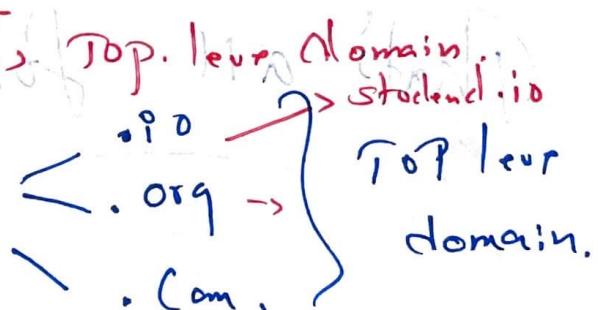
Divided into different class.

Second level Domain

mail.google.com.

Sub domain

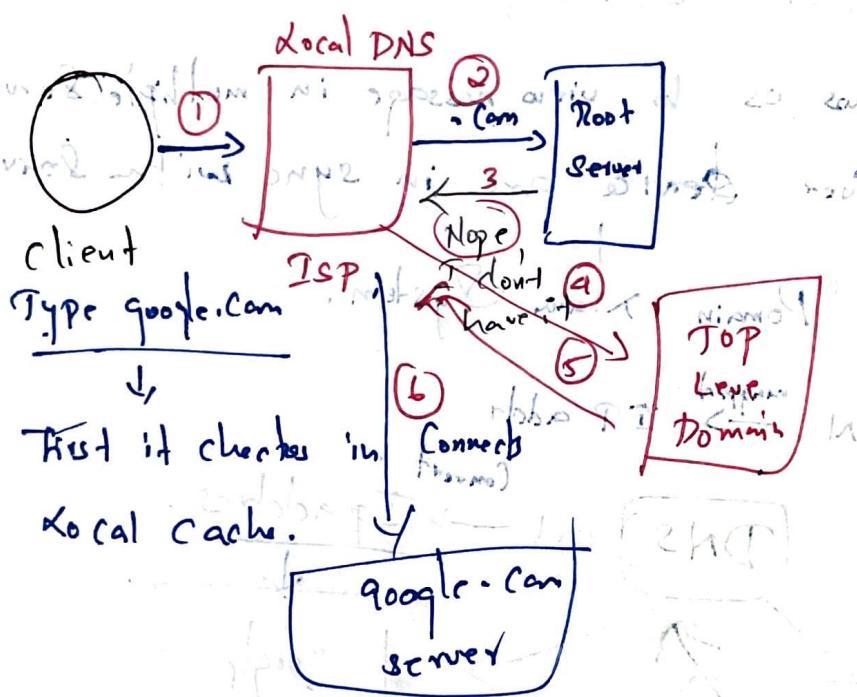
Root DNS Servers



Root DNS Servers? (first) point (Q) Contact.

Chortout : root-server.org

Database for Root DNS Server → Maintained by ICAAN



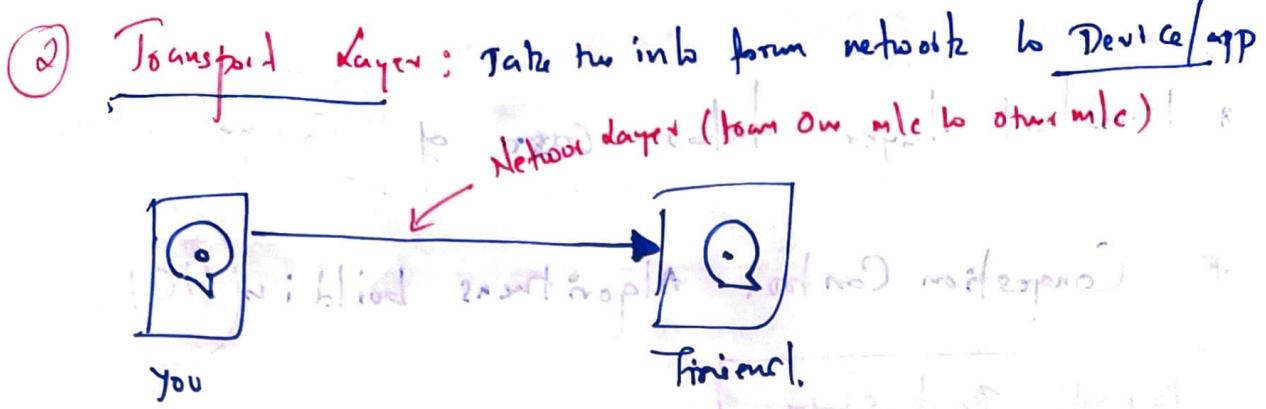
DNS looking up!

Digitized by Google

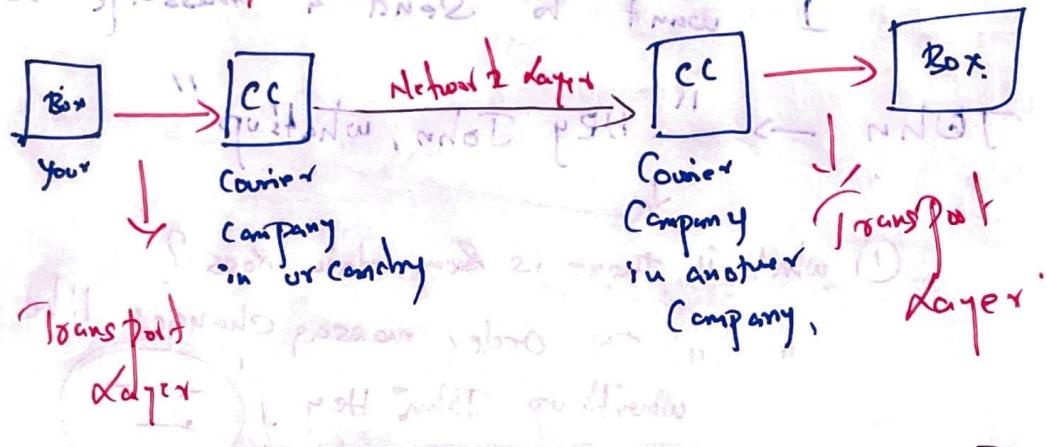
2913. (180) *inca* *viridis*

and traffics are being

That's all folks for Application Layer

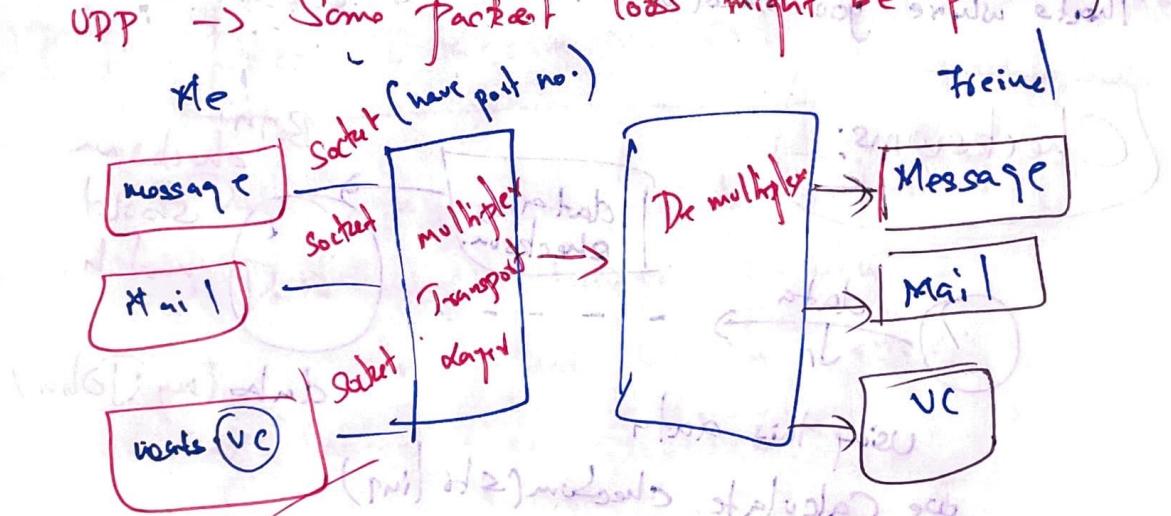


Take an example of Courier Companies

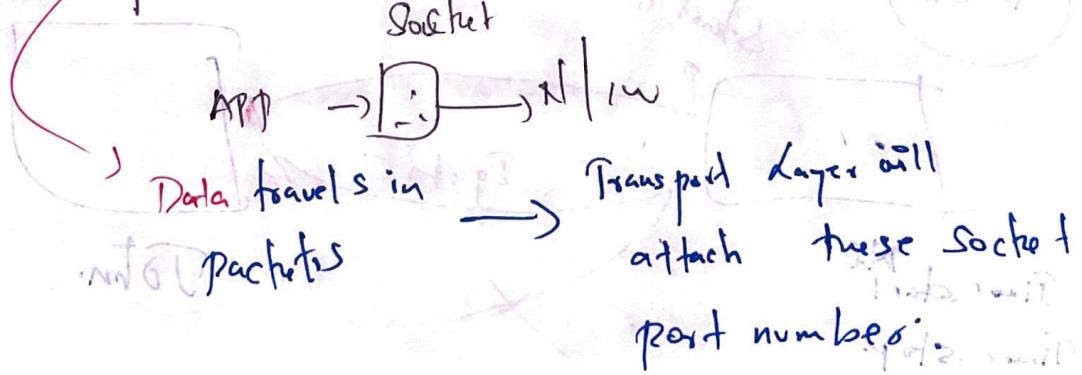


TCP: Ensure 100% of data is transferred whereas *Transport Layer*

UDP → Some packets lost might be present *Transport Layer*



Multiplex: Allows multiple I/P to multiple O/P



* Transport layer makes sure data is sent correctly

* Transport layer takes care of Congestion Control

* Congestion Control: Algorithms built in TCP

Small Back Story

I want to send a message to

John → "Hey John, what's up"

① what if there is some data loss?

the order messag changes like
what's up John Hey,

That's where you need to know Checksum.

Checksums:

data checksum

Both checksum
should match.

data → -----

data (my John)

using this data
we calculate checksum (stop)

TIMERS

Start timer

Stop timer

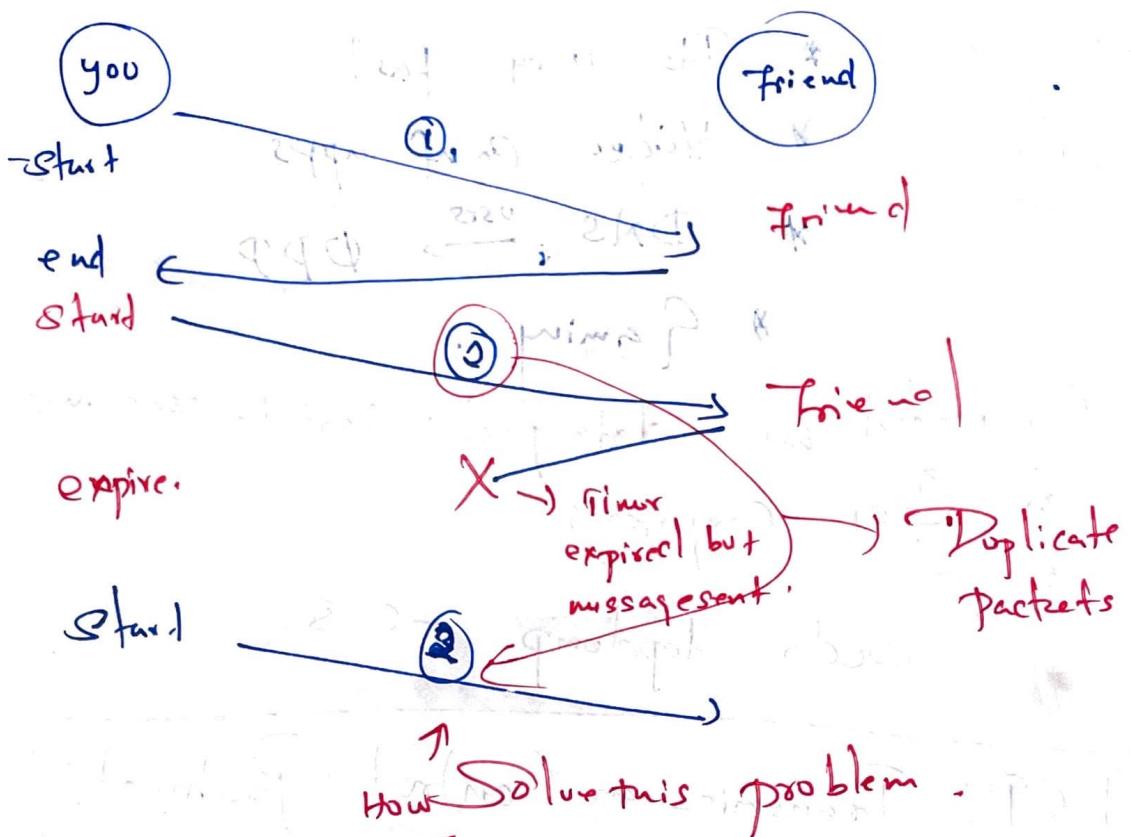
Ignore it

data (John)

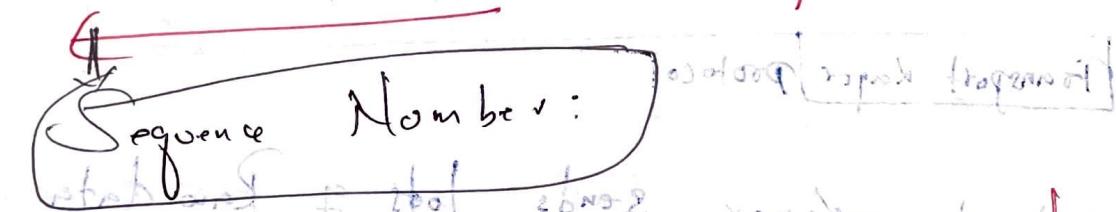
After stop timer

Timer stop

Time expire.



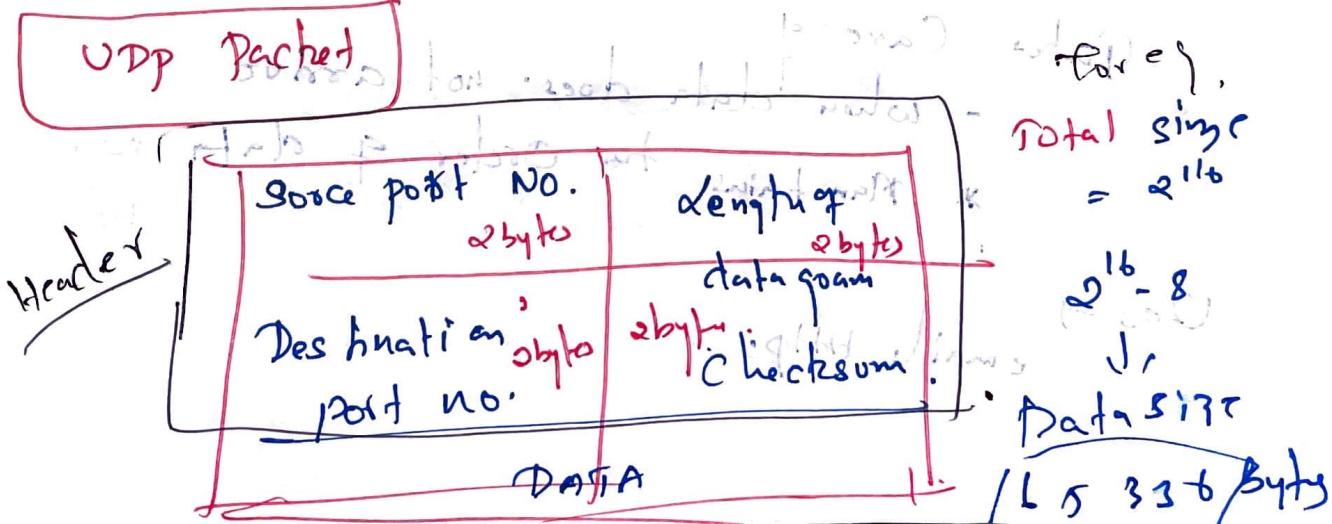
How to solve this problem -



UDP : User Datagram protocol.

- * Data may (n) may not be delivered.
- * Data may/may not change in bw.
- * " " not be in order.
- * Connectionless protocol.
- * UDP uses checksum.

UDP Packet



Use Cases : * Its very fast

or

UDP

* Video Conf apps

* DNS $\xrightarrow{\text{uses}}$ UDP

* Gaming

Command to receive data packets that is received on Our Computer:

sudo tcpdump

TCP: Transmission Control Protocol:

Transport Layer Protocol

* Application layer sends lots of Raw data
→ TCP segments this data -> divide

into data, header headers

what is separate collect the data or

→ it may also collect the data or

network layer

* Congestion Control

Gives care of when data does not arrive

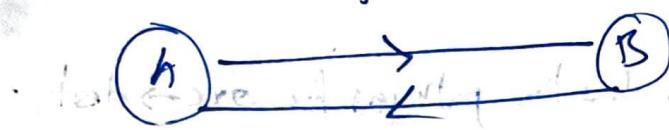
- when data arrives the order of data (seq. no)

(ex →)

email, Http.

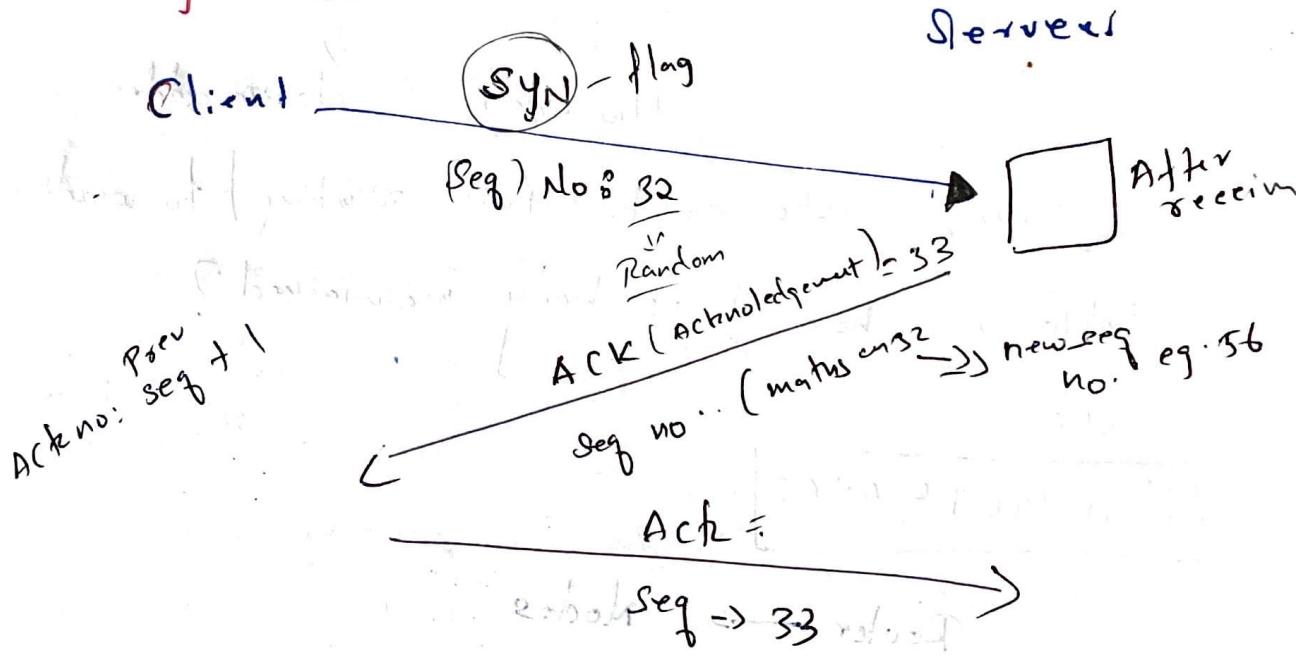
Features :

- * Connection Oriented
- * Error Control
- * Congestion Control
- * Full Duplex (Bidiirectional)

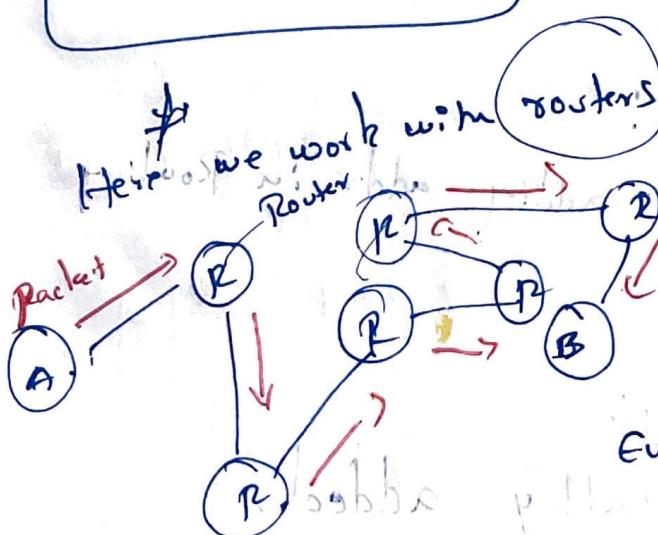


Both A & B can send
info and info
Both files Simultaneous

3-way Handshake: $S \rightarrow D$, $D \rightarrow S$



Network layer



In transport layer

Protocol: Segment

XLLW layer → packets

Data link → frames

Every Router has N/w address.

from A → packet is send with new address
↓

Router Now Router will check the forwarding
Table & then forwards the packet
from A to next Router.

HOP by
HOP

Routing table: It has two paths for src → dest.

Forwarding table: It has only one path

IP address: 192.168.2.30
↓
device addr.
↓
x/w addr device addr.

BTW, who creates this routing & forwarding
table, how it is being maintained?

J,

CONTROL PLANE:

Router → Nodes

Edge → Edges

repeat for every node

2 types of Routing:

① Static Routing: Manually adding address in Routing Table.

Table.

② Dynamic Routing:

- Dynamically added

Internet protocol (IP)

IPv4 : 32 bits are 4 words per host

IPv6 : 128 bits

For e.g.,
format of IPv6 address is 128 bit.

→ 128 bits of IP address is divided into 8 bit. \Rightarrow 32 bit.
 $\frac{0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0}{8 \text{ bit}}$ 1 word of 32 bit which is called as 32 bit word.

$$\frac{192 - 168 \cdot 2}{\text{PSI}} \quad \frac{30}{0.1 \cdot 0.001} \rightarrow$$

device
ID

Subnet IP

e.g., 192.168.1.100

3 types of address

(class of) IP address?

A

0.0.0.0

127.255.255.255

B

128.0.0.0

191.255.255.255

C

192.0.0.0

224.0.0.0

239.255.255.255

D

224.0.0.0

240.0.0.0

255.255.255.255

E

SUBNET MASKING:

For e.g. 2055.192.168.80 : IPv4

Net ID 191 : IPv4

This we can decide.

It is the technique for logically partitioning

a single physical network into multiple

smaller sub-networks (or subnets)

Variable length subnet masking

For eg:

192.0.1.0 /24

24 → 24 bit remains
& remaining is 8 bit (can be variable)

start address
192.0.1.0 + 0.0.0.0
↓ same ↓ same

229.229.229.229 /24

0.0.0.255

Reserved Address

229.229.229.229 /24

0.0.0.255

Ex: Localhost : 127.0.0.1

loopback address, PSS

229.229.229.229 /24

0.0.0.255

Packets: Header is a 20 bytes.

Where it Contains \rightarrow IPV, length, Identification, flags
Protocol, checksum, address, TTL

TTL: Internet Top Level

Router will do TTL = 60 after 60 hops if it
doesn't reach destination
it can drop.

DATA LINK LAYER

IPV6:

IPV6 : $2^{32} \approx 4.3$ billion. (Previously it was sufficient)

IPV6 : $2^{32 \times 4} \approx 2^{128} = \frac{3.4 \times 10^{38}}{\downarrow}$
unique Id.

Cons: * not Backward Compatible.
* So much effort required, IPv4
need to shift, lot of Hardware req.

a:a:a:a:a:a:a:a

↓
Hexadecimal (16 bits)

(.0.0.0.0)

Ex :

(.0.0.0.0)

A B E F

: F 0 0 1 : 3 2 1 0 : 9 1 8 2 : 0 : 0 : 1 : 3

(.0.0.0.0)

MIDDLE BOXES

entity which provides also various services with IP address.

For e.g.: Firewall → Global Internet provider
 → your trusted network

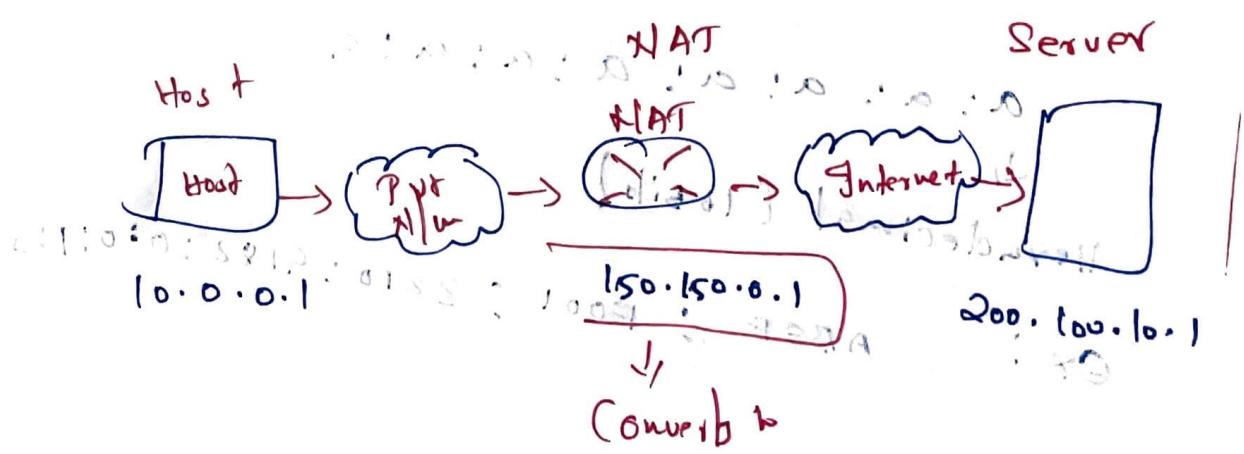
→ Filter good IP packets based on various rules
 Based on ports used in network
 → Address
 → ports and flags
 → Modify packets
 → port no's (Rules can be set)
 → flags
 → protocols.

(for first packet) initial e.g. 192.168.1.1 : HTTP
 & subsequent will be 192.168.1.1 : HTTPS

Stateless → Yes → Stateful firewall,
 by suspending
 to more efficient.

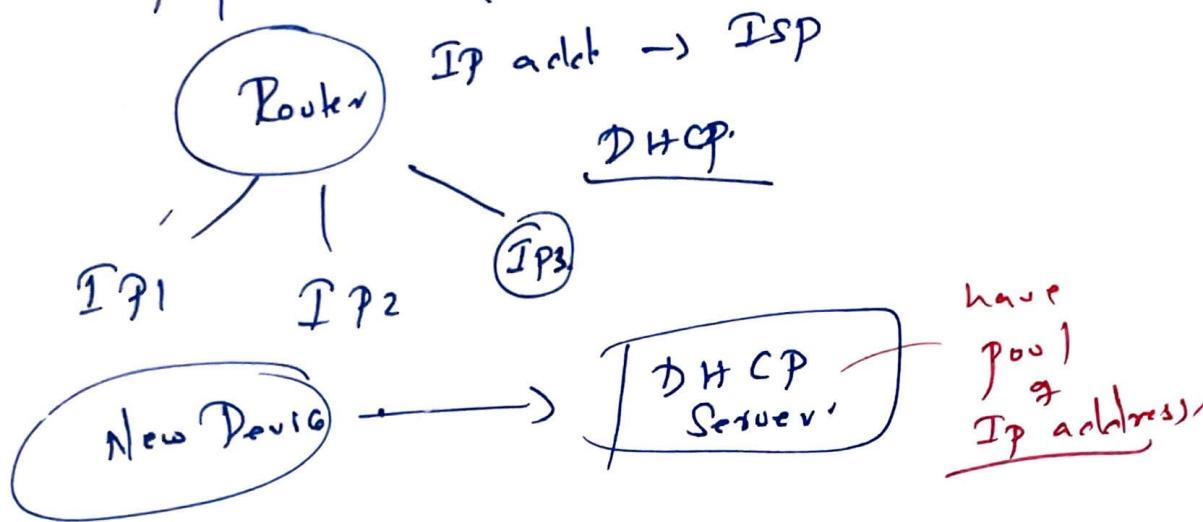
Have this information broadcasted from subsystem

IP membership traffic from or
 ② NAT → Network Address Translation. (Middle Box)

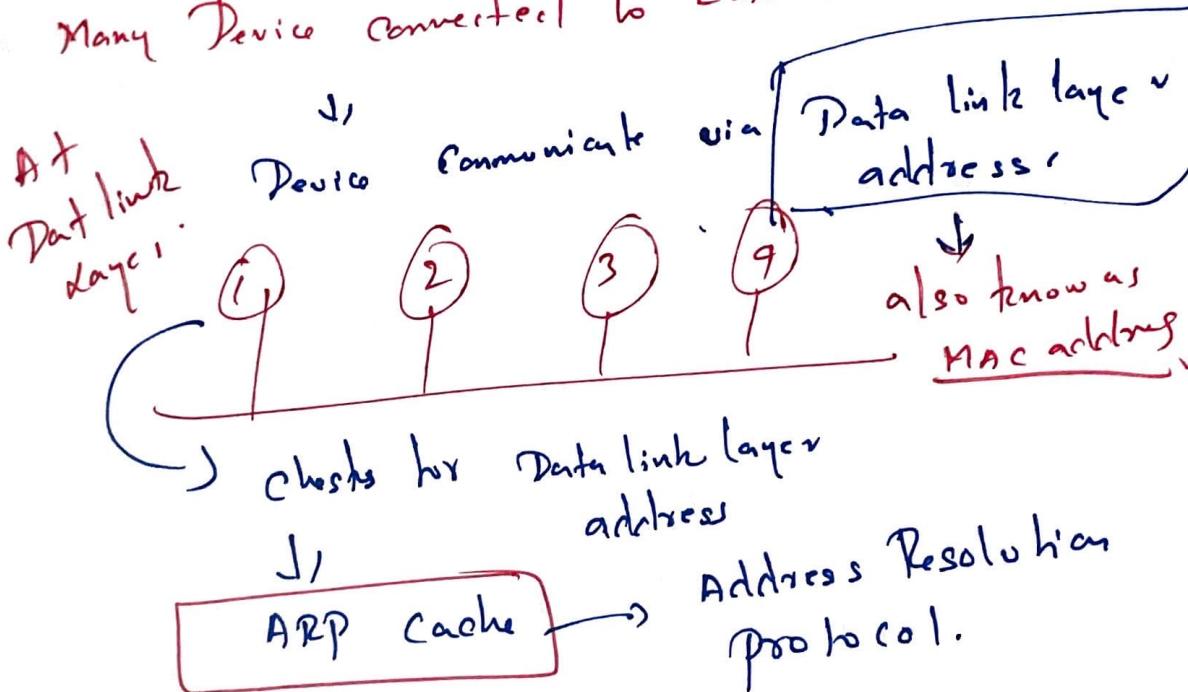


DATA LINK LAYER:

Responsible to sent the packets over a physical layer.



Many Device connected to LAN:



frame: Contains

- ① Data link layer Addr. of Sender.
- ② IP address of destination.