## Principles of Database Systems



Intermediate SQL (2)



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### Constituent Parts of SQL (SQL组成部分)

- The SQL language has several parts:
  - Data-definition language (DDL)
  - Data-manipulation language (DML)
  - Integrity (完整性) (included in DDL)
  - View definition (视图定义) (included in DDL)
  - Transaction control (事务控制)
  - Authorization (授权)
  - Embedded SQL and dynamic SQL (嵌入式SQL及动态SQL)









- In some cases, it is not desirable for all users to see the entire logical model.
- Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

**select** ID, name, dept\_name **from** instructor

• Any disadvantages?







• A **view** provides a mechanism to hide certain data from the view of certain users.

• Any relation that is not of the conceptual model but is made visible to a user as a "virtual relation" is called a view.(不是概念模型的一部分,对用户可见的"虚关系")



- SQL provides the **create view** command to create a view, which is stored in the database **in the long run** 
  - create view v as <query expression>
  - where
    - •query expression> is any legal expression
    - the **view** name is represented by v
- View definition is not the same as creating a new relation by evaluating the query expression. (创建视图与 创建关系不同)
  - Rather, a view definition causes the **saving of an expression**; the expression is substituted into queries using the view. (存储的是表达式)

A view of instructors without their salary

create view faculty as
select ID, name, dept\_name
from instructor



• Find all instructors in the Biology department

```
select name
from faculty
where dept_name = 'Biology'
```



Create a view of department salary totals

```
create view
departments_total_salary(dept_name, total_salary)
as
    select dept_name, sum (salary)
    from instructor
    group by dept_name;
```



## **Views Defined Using Other Views**



- Explain the following views:
- create view physics\_fall\_2009 as
   select course.course\_id, sec\_id, building, room\_number
   from course, section
   where course.course\_id = section.course\_id
   and course.dept\_name = 'Physics'
   and section.semester = 'Fall'
   and section.year = '2009';
- create view physics\_fall\_2009\_watson as select course\_id, room\_number from physics\_fall\_2009 where building= 'Watson';



## Views Defined Using Other Views

Expand use of a view in a query/another view

```
create view physics_fall_2009_watson as
(select course_id, room_number
from (select course.course_id, building,
room_number
      from course, section
      where course_id = section.course_id
         and course.dept_name = 'Physics'
          and section.semester = 'Fall'
         and section.year = '2009')
where building= 'Watson';
```



### **Materialized Views**

- Materializing a view(物化视图): create a physical table containing all the tuples in the result of the query defining the view
- If relations used in the query are updated, the materialized view result becomes out of date
  - Need to **maintain** the view(维护视图), by updating the view whenever the underlying relations are updated.



## **Drop View**



• The **Drop View** command deletes the definition the view from the **data dictionary**.

drop view view\_name;

Other views depending on this dropped view should be deleted explicitly.



## Update of a View

 Add a new tuple to faculty view which we defined earlier

insert into faculty values ('30765', 'Green', 'Music');

- Two reasonable approaches:
  - Reject the insertion
  - Insert the tuple

('30765', 'Green', 'Music', null)

into the instructor relation

(必须转化为对实际关系的修改)



# Some Updates cannot be Translated Uniquely

- create view instructor\_info as select ID, name, building from instructor, department where instructor.dept\_name= department.dept\_name;
- **insert into** instructor\_info **values** ('69987', 'White', 'Taylor');
  - which department, if multiple departments in Taylor?
  - what if no department is in Taylor?



# Some Updates cannot be Translated Uniquely



- Most SQL implementations allow updates only on simple views
  - The **from** clause has only **one** database relation.
  - The select clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
  - Any attribute not listed in the **select** clause can be set to null.
  - The query does **not have** a **group** by or **having** clause.



#### And Some Not at All

- create view history\_instructors as select \*
   from instructor
   where dept\_name= 'History';
- What happens if we insert ('25566', 'Brown', 'Biology', 100000) into history\_instructors?
- with check option: if a tuple inserted into the view does not satisfy the view's where clause condition, the insertion is rejected by the database system







- A transaction is a sequence of queries and update statements on DB, executed as a single, and are started implicitly and terminated by one of commit work(提交)or rollback/abort work
  - Commit work: makes the updates performed by the transaction become permanent in the database.
  - Rollback work: undoes all the updates performed by the SQL statements in the transaction.



- Unit of work
- Atomic transaction
  - either fully executed or rolled back as if it never occurred
- Transactions begin implicitly
  - Ended by commit work or rollback work
- But default on most databases: each SQL statement commits automatically
  - Can turn off auto commit for a session (e.g. using API)
  - In SQL:1999, can use: begin atomic .... end
    - Not supported on most databases





## Integrity Constraints



## **Integrity Constraints**

- Integrity constraints guard **against accidental damage** to the database, by ensuring that authorized changes to the database do not result in a loss of data consistency.
  - An instructor name cannot be *null*.
  - No two instructors can have the same instructor ID.
  - The budget of a department must be greater than \$0.00.



## Integrity Constraints on a Single Relation

- primary key
- not null
- Unique
- foreign keys
- check (P), where P is a predicate



### **Not Null**



- Declare name and budget to be not null
  - name varchar(20) not null
  - budget numeric(12,2) not null



## **Unique Constraints**



- unique ( A<sub>1</sub>, A<sub>2</sub>, ..., A<sub>m</sub>)
  - The unique specification states that the attributes  $A_1, A_2, ... A_m$  form a **candidate key**.
  - Candidate keys are permitted to be null (in contrast to primary keys).



#### The check clause



- The **check** clause is applied to relation declaration
  - check (P), where P is a predicate which must be satisfied by every tuple in the relation.
- Example: ensure that the budget of a department must be greater than \$0.00
  - create table department (dept name varchar (20), building varchar (15), budget numeric (12,2), primary key (dept name) check(budget>0));





• Ensures that semester is one of fall, winter, spring or summer:

```
create table section (
course_id varchar (8),
sec_id varchar (8),
semester varchar (6),
year numeric (4,0),
building varchar (15),
room_number varchar (7),
time slot id varchar (4),
primary key (course_id, sec_id, semester, year),
check (semester in ('Fall', 'Winter', 'Spring', 'Summer'))
```

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## **Referential Integrity**

• Ensures that a value that appears in one relation for a given set of attributes **also appears** for a

certain set of attributes in another relation.

Example: If "Biology" is a department name appearing in one of the tuples in the *course* relation, then there exists a tuple in the *department* relation for "Biology". -- foreign key

```
create table course (
    course_id char(5) primary key,
    title varchar(20),
    dept_name varchar(20) references department
)
```



# Cascading Actions in Referential Integrity



 When the DB is modified by Insert, Delete, and Update, the tests must be made in order to preserve the referential integrity constraint.

```
    create table course (
        ...
        dept_name varchar(20),
        foreign key (dept_name) references department
        on delete cascade
```

on update cascade,

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• alternative actions to cascade: set null, set default



# **Integrity Constraint Violation During Transactions**

- E.g. create table person (
  ID char(10) primary key,
  name char(40),
  spouse char(10),
  foreign key spouse references person)
- How to insert a tuple without causing constraint violation?
  - set spouse to null initially, update after inserting all persons (not possible if spouse attribute declared to be **not null**)
  - OR **defer**(延迟) constraint checking



## **Defer Constraint Checking**



- The SQL standard allows a clause initially deferred to be added to a constraint specification.
- For constraints declared as **deferrable**, executing a statement **set constraints** constraint-list **deferred** as part of a transaction causes the checking of the specified constraints to be deferred to the end of that transaction.



# **Complex Check Conditions and Assertions**



check (time\_slot\_idin (select time\_slot\_id from time\_slot))

• Unfortunately: subquery in check clause not supported by pretty much any database.



### **Assertions**



- An **assertion** is a predicate expressing a condition that we wish the database always to satisfy.
  - e.g. domain constraints, referential-integrity constraint

An assertion in SQL takes the form
 create assertion <assertion-name> check
 cpredicate>



#### **Assertions**



• E.g. The value of the attribute tot\_cred for each student must equal the sum of credits of courses that the student has completed successfully.

```
create assertion credits_earned_constraint check
  (not exists (select ID from student
    where tot_cred < > (
```

select sum(credits)

from takes join course

on takes.course\_id= course.course\_id

where student.ID=takes.ID and grade is not

null and grade < > 'F')



#### **Assertions**



- When an assertion is made, the DBMS tests it for validity. Any modification to DB is allowed only if it does not cause that assertion to be violated.
  - This testing may introduce a significant amount of overhead, hence **assertions should be used** with great care.

Not supported by every DBMS.





## Review



### Review



#### Join Expressions

- left outer join, right outer join, full outer join
- inner join = join
- Join types and join conditions

#### Views

- Create view
- Use views in SQL queries
- Update view: with check option

- Atomic
- Commit work, Rollback work



#### Review



- Integrity Constraints
  - Not null
  - unique  $(A_1, A_2, ..., A_m)$ , candidate key, null
  - Check(P)
  - Referential Integrity, foreign key, on delete/update cascade, on delete/update set null, on delete/update set default
  - defer constraint checking
  - create assertion <assertion-name> check
     cpredicate>

