

Improvements in randomized benchmarking qiskit_experiments

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Agenda

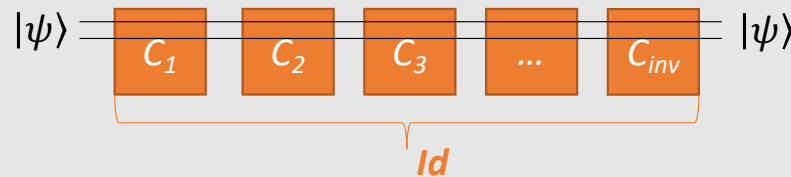
- (Very) short reminder on what randomized benchmarking (RB) is
- The existing algorithm for randomized benchmarking in `qiskit_experiments`, and its bottlenecks
- Main ideas for performance improvements for 1 and 2 qubits
- Main improvements in the code structure
- Benchmarking results

Randomized Benchmarking (RB) – brief reminder

RB is a protocol that provides an estimate of the average error-rate for a set of quantum gate operations

Consists of the following three stages:

- a) Generate RB sequences consisting of random elements from the Clifford group, followed by the Clifford that is the inverse of the random sequence.



- b) Run the RB sequences either on the device or on the simulator (with a noise model) and compare to the initial state
- c) Get the statistics and fit an exponential decaying curve: $Ap^m + B$
Compute error per Clifford (EPC): $r = (1 - p)(d - 1)/d$ ($d = 2^n$)
From this, compute error per gate.

[1] Easwar Magesan, J. M. Gambetta, and Joseph Emerson, *Robust randomized benchmarking of quantum processes*, <https://arxiv.org/pdf/1009.3639>

[2] Easwar Magesan, Jay M. Gambetta, and Joseph Emerson, *Characterizing Quantum Gates via Randomized Benchmarking*, <https://arxiv.org/pdf/1109.6887>
https://github.com/Qiskit/qiskit-tutorials/blob/master/qiskit/advanced/ignis/5a_randomized_benchmarking.ipynb

Current implementation – rough description

```
list_of_circuits = []  
For length in all_lengths  
    circuit = QuantumCircuit  
    current_clifford = Id  
    For i in 1 to length-1  
        generate one random_Clifford  
        append random_Clifford to circuit  
        current_clifford = current_Clifford  $\circ$  random_Clifford  
        append barrier  
    inverse_clifford = inverse(current_clifford)  
    append inverse_clifford to circuit  
    append circuit to list_of_circuits  
  
transpile all circuits in list_of_circuits  
run on backend  
analyze results
```

Current implementation – bottlenecks

```
list_of_circuits = []
```

```
For length in all_lengths
```

```
    circuit = QuantumCircuit
```

```
    current_clifford = Id
```

```
    For i in 1 to length-1
```

```
        generate one random_Clifford
```

```
        append random_Clifford to circuit
```

```
        current_clifford = current_Clifford ◦ random_Clifford
```

```
        append barrier
```

```
    inverse_clifford = inverse(current_clifford)
```

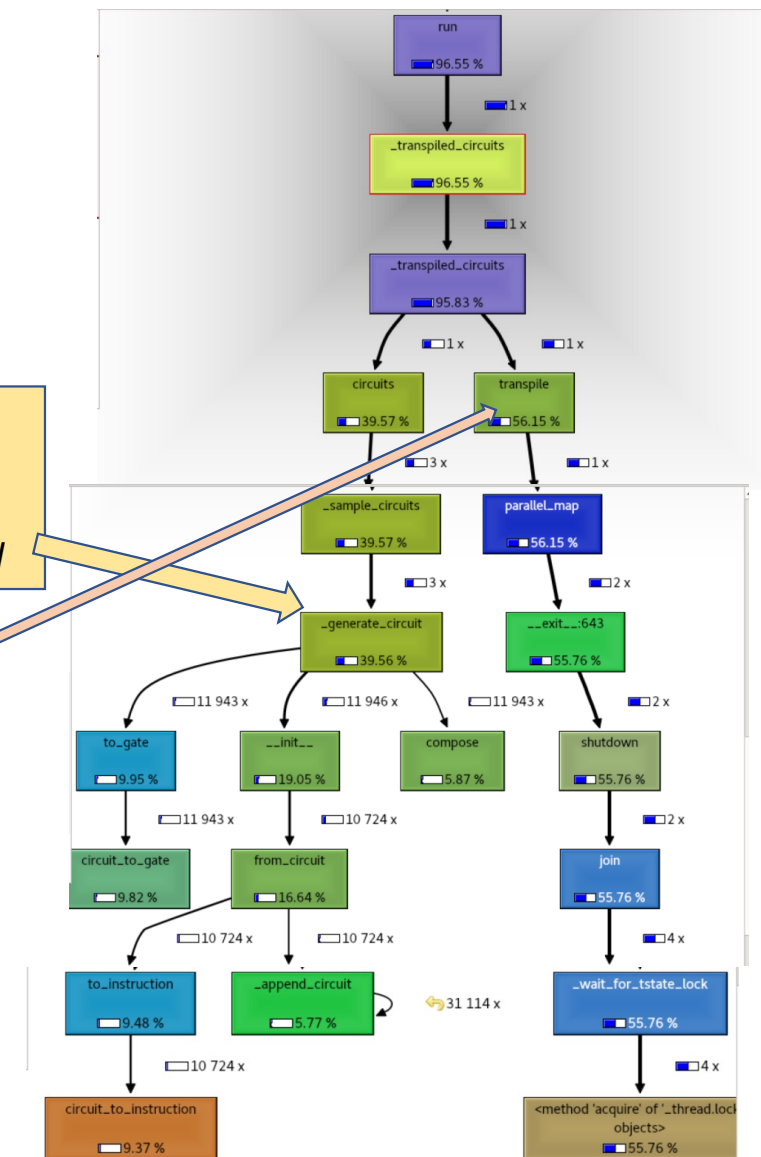
```
    append inverse_clifford to circuit
```

```
    append circuit to list_of_circuits
```

```
transpile all circuits in list_of_circuits
```

```
run on backend
```

```
analyze results
```





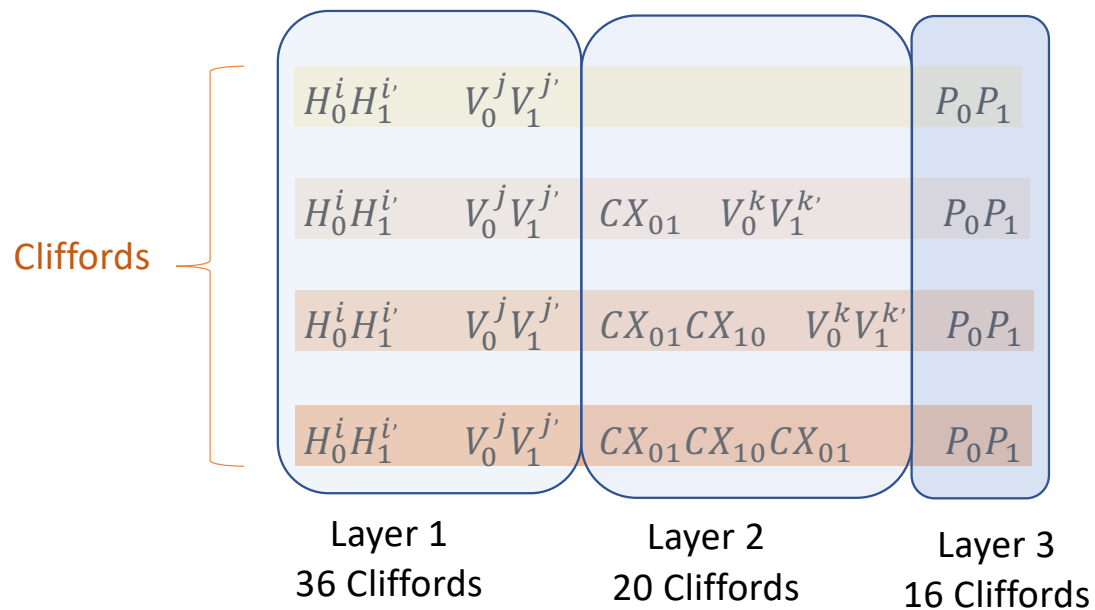
New implementation

Three main ideas:

1. Eliminate transpilation of the entire circuit
2. Speed up operations on Cliffords
3. Reasonable storage space

1. Eliminate transpilation of the entire circuit

- Build the Cliffords in 3 layers
- Transpile these Cliffords in advance

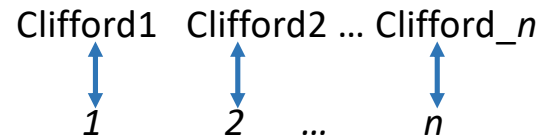


$$i, i' \in \{0,1\}, \quad j, j', k, k' \in \{0,1,2\}, \quad P_0, P_1 \in \{I, X, Y, Z\}$$

$$V = S^\dagger H$$

2. Speed up operations on Cliffords (create, compose and inverse)

Create an isomorphism between the group of Cliffords and a group of integers.



- Clifford_j ◦ Clifford_k = Clifford_m ↔ $j \circ k = m$
- inverse(Clifford_j) = Clifford_k ↔ Inverse(j) = k

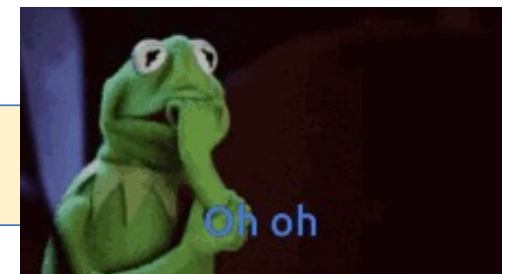
Perform all group operations on the group of integers instead of on the group of Cliffords.

Store a list of all the results of compose. Access the result of $j \circ k$ according to the indices j and k .

Store a list of all the results of inverse. Access the inverse of j according to the index of j .

These lists are constant and are generated once for all RB experiments.

Note – the size of the compose list is n^2 , where n is the size of the Clifford group.
For 2 qubits $11520^2 = 132,710,400$



3. Reduce storage space for compose results

Recall that the Clifford Group for 1 qubit is generated by $clifford_single_gates_1q = \{S, H\}$.
For 2 qubits, it is generated by $clifford_single_gates_2q = \{S(0), S(1), H(0), H(1), CX(0,1), CX(1,0)\}$.

In other words, we can decompose every Clifford into these gates.

So instead of computing $C_1 \circ C_2$, we can decompose C_2 into $C_2 = C_2^0 \circ C_2^1 \circ \dots \circ C_2^k$

Where $C_2^i \in clifford_single_gates$, $0 \leq i \leq k$

And then compute $C_1 \circ C_2^0 \circ C_2^1 \circ \dots \circ C_2^k$.

This means we can reduce the composition table

- For 1 qubit - from 24 X 24 to 24 X 2.
- For 2 qubits – from 11520 X 11520 to 11520 X 6 = 69,120*.

* In practice for convenience, we store the entries for 21 gates, for a total storage of 241,920 entries

New algorithm – rough description

```
transpiled_Cliffords = [transpile(layer1, layer2, layer3)]
```

```
list_of_circuits = []
```

```
For length in all_lengths
```

```
    circuit = QuantumCircuit
```

```
    current_clifford_num = 0
```

```
    For i in 1 to length-1
```

```
        generate a random triplet of integers  $c1, c2, c3$  from a random num in  $[0, \dots, num\_cliffords]$ 
```

```
        rand_cliff = convert_to_Clifford( $c1 \circ c2 \circ c3$ )
```

```
        circuit.append(rand_cliff)
```

```
        circuit.append(barrier)
```

```
        current_clifford_num =
```

```
            current_clifford_num  $\circ c1 \circ c2 \dots \circ c3$  # lookup in list
```

```
    inverse_cliff_num = inverse(current_clifford_num) # lookup in list
```

```
    inverse_cliff = transpiled_Cliffords[inverse_cliff_num]
```

```
    circuit.append(inverse_cliff)
```

```
    append circuit to list_of_circuits
```

```
perform trivial placement on physical qubits
```

```
run on backend
```

```
analyse results
```

Itoko – your turn now!