ZAP : An ARM v4T Compatible Soft Processor

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1 Introduction

ZAP is an ARM® v4T compatible soft processor core. The code is fully open source and is released under the MIT license.

1.1 Features

• Fully ARM v4T compatible.

- Executes the 32-bit wide ARM v4 instruction set.
- Executes the 16-bit wide compressed instruction set.

• Deeper pipeline for better clock speed.

- The processor is built around a 9 stage pipeline to achieve a high operating frequency. The pipeline has a data forwarding capability to allow back to back instructions to execute without stalls. Non trivial shifts require their operands a cycle early. Loads have a 3 cycle latency and the pipeline will stall if an attempt is made to access the register within the latency period.

• Supports interrupt and abort signaling

- Features dedicated high level sensitive IRQ, FIQ and memory abort pins.

• Can be interfaced with caches/MMU.

- The CPSR of the processor is exposed as a port allowing for implementation of a virtual memory system.
- Memory stall may be indicates to the core via dedicated ports to allow caches to be connected.

• Coprocessor interface provided.

The coprocessor interface simply exposes internal signals of the core. It is up to the coprocessor
to interpret and process instructions correctly.

• Supports M-variant multiplication instructions

- These instructions are supported:

```
MUL, MLA, SMULL, UMULL, SMLAL, UMLAL
```

• The core is configurable

 The processor may be synthesized without compressed instruction support and/or coprocessor interface support to save area and improve speed.

• Designed for FPGA synthesis

- Most memory structures of the processor map efficiently onto FPGA block RAMs. The register file is overclocked by a 2x clock to allow for 2 write ports.
- The branch predictor memory also efficiently maps to FPGA block RAM.
- No device specific instantiations are made to allow for portability across FPGA vendors.

Faster performance of memory instructions

1 INTRODUCTION 6

- Memory instructions with writeback can be issued as a single instruction since the register file is built to have 2 write ports. This may improve performance.

- The processor core is written entirely in synthesizable Verilog-2001.
- A branch predictor is installed to compensate for the longer pipeline.
 - Branches within a 2KB block of memory can be mapped into the predictor without conflict.
 The predictor basically uses a bimodal prediction algorithm (2-bit saturating counter per branch entry).

· Uses a base restored abort model

Uses a base restored abort model making it easier to write exception handlers. Basically, on a
fault in between a multiple memory transfer, the processor rolls back the base pointer register as
it were before the operation took place.

1.2 About This Manual

The purpose of this manual is to document the processor core's design. This document is very incomplete. I will try my best to update it.

2 Configuring the core and testbench

Throughout, it is assumed that \$ZAP_HOME points to the working directory of the project. Core/testbench configuration may be done using defines. The defines file is located in \$ZAP_HOME/includes/config.vh.

See table 1

Define	Purpose	Required for	Comments		
CORE CONFIGURATION					
THUMB_EN	Enabling compressed instruction support	Core Setup	Enabling this increases core area and reduces performance.		
COPROC_IF_EN	Enabling coprocessor suport. Extra ports get added.	Core Setup	Enabling this increases core area and reduces performance.		
		ONFIGURATION			
IRQ_EN	Generates periodic IRQ pulses.	Testbench	_		
SIM	Generates extra messages.	Testbench	Must be UNDEFINED for correct synthesis of the core in Xilinx since some debugging structures in RTL are removed if this is not defined. ==; MUST BE UNDEFINED FOR SYNTHESIS ;==		
VCD_FILE_PATH	Set the path to the VCD data dump.	Testbench	_		
MEMORY_IMAGE	Path to the memory image Verilog file.	Testbench	-		
MAX_CLOCK_CYCLES	Set the number of cycles the simulation should run before terminating.	Testbench	_		
SEED	Set the testbench seed. The seed influences randomness.	Testbench	_		

Table 1: Configuring the core and testbench

3 Compiling Code and Running Simulation

NOTE: If you want to quickly test the processor with sw/asm/prog.s and sw/asm/prog.c sample programs, see Section 3.1

3.0.1 Generating a binary using GNU tools

You can use the existing GNU toolchain to generate code for the processor. This section will briefly explain the procedure. For the purposes of this discussion, let us assume these are the source files...

```
main.c
fact.c
startup.s
misc.s
linker.ld This is the linker script.
```

Generate a bunch of object files.

```
arm-none-eabi-as -mcpu=arm7tdmi -g startup.s -o startup.o arm-none-eabi-as -mcpu=arm7tdmi -g misc.s -o misc.o arm-none-eabi-gcc -c -mcpu=arm7tdmi -g main.c -o main.o arm-none-eabi-gcc -c -mcpu=arm7tdmi -g fact.c -o fact.o
```

Link them up using a linker script...

```
arm-none-eabi-ld -T linker.ld startup.o misc.o main.o fact.o -o prog.elf
```

```
Finally generate a flat binary... arm-none-eabi-objcopy -0 binary prog.elf prog.bin
```

The .bin file generated is the flat binary.

3.0.2 Generating a Verilog memory map

```
perl $ZAP_HOME/scripts/bin2mem.pl prog.bin prog.v The prog.v file looks like this...
```

```
mem[0] = 8'b00;
mem[1] = 8'b01;
```

3.0.3 Invoking the simulator

Ensure config. vh is set up correctly.

Your command must look like this (It is a single command)...

```
iverilog $ZAP_HOME/rtl/*.v $ZAP_HOME/rtl/*/*.v $ZAP_HOME/testbench/*.v
$ZAP_HOME/models/ram/ram.v -I$ZAP_HOME/includes -DSEED=22
```

The rtl/*.v and rtl/*/*.v collect all of the synthesizable Verilog-2001 files, the testbench/*.v collects all of the testbench (In this situations, the ram.v file is a part of the testbench).

Provide some seed value (22 is used in the example). Ensure you edit the config.vh file before running the simulation to correctly point to the memory map, vcd target output path etc for the simulator to pick up.

3.1 Run Sample Code quickly...

A sample .s and .c file is present in \$ZAP_HOME/sw/s and \$ZAP_HOME/sw/c respectively. To translate them to binary and to a Verilog memory map, you can run the Perl script (See NOTE below)

```
perl $ZAP_HOME/debug/run_sim.pl
```

NOTE: Ensure you set all the variables in the Perl script as per table 2...

Variable	Purpose	
ZAP_HOME Set this to the project working directory.		
LOG_FILE_PATH	Set this to the place where you want the log file to	
	be created.	
ASM_PATH	Set this to the location of your startup assembly file.	
C_PATH	Set this to the location of your C file.	
LINKER_PATH	Set this to the location of the linker script.	
TARGET_BIN_PATH	Set this to the target bin file location.	

VCD_PATH	Set this to location where the VCD is to be created.	
	This must match what is in config.vh	
MEMORY_IMAGE	Set this to the location where the memory image	
	must be created. This must match what is in con-	
	fig.vh	

Table 2: Variables in Perl script

4 IO Ports of the Core

Table 3 lists the ports of the processor core.

Port	Direction	Description	Comments
i_clk	I	Core clock	_
i_clk_2x	I	Register file clock	Must run syn- chronously with the core clock at twice the frequency.
i_reset	I	Active high reset	This runs through a reset filter (using 2 series D flip-flops) that then drives a global reset throughout the design.
i_instruction[31:0]	I	32-bit instruction from the instruction cache.	_
i_valid	I	Indicates that the instruction on the 32-bit bus is valid and ready to be clocked into the instruction register (Fetch stage). If this is 0, the PC does not change.	_
i_instr_abort	I	Indicates that the instruction cache experienced an instruction abort.	When abort is asserted, the i_valid port is treated as a DON'T CARE.
o_read_en	0	Current memory access is a read access.	_
o_write_en	0	Current memory access is a write access.	-
o_address[31:0]	O	Address of the current memory access.	The address is always 32-bit aligned. Specific bytes are addressed using byte enables.
o_mem_translate	О	Current memory access must adopt a user view of memory	
i_data_stall	Ι	Indicates that the data memory has stalled. The entire pipeline freezes when this happens (PC does not change).	_

: data abaset	Т	Indicates that the data	
i_data_abort	I		
		memory access aborted.	
		When this is asserted,	
		the i_data_stall port	
		must be deasserted.	
i_rd_data[31:0]	I	Data received from data	_
		memory is clocked in	
		from this port if there is	
		no data memory stall.	
o_wr_data[31:0]	0	Data to write out to	_
		memory	
o_ben[3:0]	0	Byte enables. [0] deals	_
		with [7:0] of the data	
		and [3] with [31:24].	
i_fiq	I	FIQ level sensitive sig-	_
1-114	1	nal (Active high)	
	I		
i_irq	1	IRQ level sensitive sig-	_
50 (03		nal (Active high)	
o_pc[31:0]	0	Program Counter	Lower bit is ALWAYS
			0.
o_cpsr[31:0]	О	Current PSR	
i_copro_done	I	Coprocessor done indi-	_
		cation.	
o_copro_dav	О	Data on the coprocessor	_
		output ports of the core	
		are valid	
o_copro_word	0	The entire 32-bit co-	_
[31:0]		processor instruction is	
		presented on the port.	
i_copro_reg_en	I	Coprocessor wishes to	_
		take charge of the regis-	
		ter file.	
i copro regur index	Ī	Coprocessor write reg-	
i_copro_reg_wr_index [5:0]	1	ister index.	
	Т		
i_copro_reg_rd_index	1	Coprocessor read regis-	_
[5:0]	T	ter index	
i_copro_reg_wr_data	I	Coprocessor data	_
[31:0]		to write to the reg-	
		ister mentioned in	
		i_copro_reg_wr_index.	

o_copro_reg_rd_data	0	Data read from reg-	_
[31:0]		ister file to be read	
		into coprocessor.	
		This corresponds to	
		register specified on	
		i_copro_reg_rd_index	
		although it is delayed	
		by 1 clock cycle.	

Table 3: IO Ports

NOTE: If COPROC_IF_EN is not defined, the coprocessor ports are not available!

NOTE: The coprocessor must generate translated register numbers. For example, if it wishes to write to R13_FIQ, it must generate a write to R25.

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5 Basic Description

Figure 1 should give a basic overview of the pipeline structure (8 stage pipeline). Table 4 gives an overall description of the pipeline stages. Note that the pipeline is an 8 stage pipeline.

FETCH PRE-DECODE DECODE ISSUE SHIFTER EXECUTE MEMORY ACCESS BRANCH MEMORY BLOCK RAM

Figure 1: High level view of the ZAP pipeline

Stage	Description
FETCH	Clocks in data from the instruction cache. Also
	computes PC+8 in case the instruction refers to
	R15 getting PC from the third stage).
BRANCH PREDICTOR RAM	This consists of block RAM dedicated to hold
	branch state for up to 2KB of instructions.
PREDECODE	Decompresses 16-bit instructions, handles LD-
	M/STM instructions and SWAP/SWAPB. Also
	handles the coprocessor interface. Basically con-
	sists of a series of state machines.
DECODE	Decodes 32-bit instructions into a format that can
	be understood by the downstream logic.
ISSUE	Performs preliminary register read by sniffing out
	register value from the bypass network.
VALUE RESTORE	
SHIFTER	
MAC-UNIT	The value restore unit restores correct values by
	getting them from the ALU. This is useful for
	executing back-to-back instructions that typically
	get read incorrectly in issue. The shifter handles
	all shift operations and the MAC unit performs
	multiply- accumulate.
EXECUTE	The execute unit consists of a 32-bit ALU. This
	unit also generates memory addresses and other
	memory control signals.
MEMORY ACCESS	This stage clocks in data from the data cache.
REGISTER FILE	This stage is the register file and is basically 4
	block RAMs to allow for 4 independent read ports
	(TODO: Can be reduced to 3 since one port is re-
	dundant).

Table 4: Pipeline Stage Description

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5.1 Register Mapping

Table 5 shows how banking is implemented in ZAP. The register map to the register file as shown below...

Program Register	Physical Register
USR_R0	0
USR_R1	1
USR_R2	2
USR_R3	3
USR_R4	4
USR_R5	5
USR_R6	6
USR_R7	7
USR_R8	8
USR_R9	9
USR_R10	10
USR_R11	11
USR_R12	12
USR_R13	13
USR_R14	14
PC	15
RAZ_REGISTER	16
CPSR	17
FIQ_R8	18
FIQ_R9	19
FIQ_R10	20
FIQ_R11	21
FIQ_R12	22
FIQ_R13	23
FIQ_R14	24
IRQ_R13	25
IRQ_R14	26
SVC_R13	27
SVC_R14	28
UND_R13	29
UND_R14	30
ABT_R13	31
ABT_R14	32
DUMMY_REG0	33
DUMMY_REG1	34
FIQ_SPSR	35
IRQ_SPSR	36
SVC_SPSR	37
UND_SPSR	38
ABT_SPSR	39
SWI_SPSR	40

Table 5: Register Mapping

The file that describes this translation is in includes/regs.vh. The function that does the translation is in includes/global_functions.vh