Synchronization: Advanced

Introduction to Computer Systems

Today

- Using semaphores to schedule shared resources
 - Review: Semaphores
 - Producer-consumer problem
 - Readers-writers problem
- Thread safety
- Prethreaded Concurrent Server

Review: Semaphores

- **Semaphore:** non-negative global integer synchronization variable. Manipulated by **P** and **V** operations.
- **■** P(s)
 - If s is nonzero, then decrement s by 1 and return immediately.
 - If s is zero, then suspend thread until s becomes nonzero and the thread is restarted by a V operation.
 - After restarting, the P operation decrements s and returns control to the caller.
- *V(s)*:
 - Increment s by 1.
 - If there are any threads blocked in a P operation waiting for s to become non-zero, then restart exactly one of those threads, which then completes its P operation by decrementing s.
- Semaphore invariant: (s >= 0)

Review: Using semaphores to protect shared resources via mutual exclusion

Basic idea:

- Associate a unique semaphore mutex, initially 1, with each shared variable (or related set of shared variables)
- Surround each access to the shared variable(s) with P(mutex) and V(mutex) operations

```
mutex = 1
P(mutex)
cnt++
V(mutex)
```

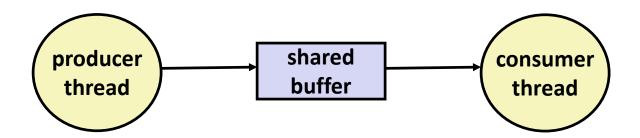
Using Semaphores to Coordinate Access to Shared Resources

- Basic idea: Thread uses a semaphore operation to notify another thread that some condition has become true
 - Use counting semaphores to keep track of resource state and to notify other threads
 - Use mutex to protect access to resource
- Two classic examples:
 - The Producer-Consumer Problem
 - The Readers-Writers Problem

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Producer-Consumer Problem



Common synchronization pattern:

- Producer waits for empty slot, inserts item in buffer, and notifies consumer
- Consumer waits for *item*, removes it from buffer, and notifies producer

Examples

- Multimedia processing:
 - Producer creates MPEG video frames, consumer renders them
- Event-driven graphical user interfaces
 - Producer detects mouse clicks, mouse movements, and keyboard hits and inserts corresponding events in buffer
 - Consumer retrieves events from buffer and paints the display

Producer-Consumer on 1-element Buffer

```
#include "csapp.h"

#define NITERS 5

void *producer(void *arg);
void *consumer(void *arg);

struct {
  int buf; /* shared var */
  sem_t full; /* sems */
  sem_t empty;
} shared;
```

```
int main() {
 pthread t tid producer;
 pthread t tid consumer;
  /* Initialize the semaphores */
  Sem init(&shared.empty, 0, 1);
  Sem init(&shared.full, 0, 0);
  /* Create threads and wait */
 Pthread create (&tid producer, NULL,
                 producer, NULL);
 Pthread create (&tid consumer, NULL,
                 consumer, NULL);
 Pthread join(tid producer, NULL);
 Pthread join(tid consumer, NULL);
 exit(0);
```

Producer-Consumer on 1-element Buffer

Initially: empty==1, full==0

Producer Thread

```
void *producer(void *arg) {
  int i, item;
  for (i=0; i<NITERS; i++) {
    /* Produce item */
    item = i;
    printf("produced %d\n",
            item);
    /* Write item to buf */
    P(&shared.empty);
    shared.buf = item;
    V(&shared.full);
  return NULL;
```

Consumer Thread

```
void *consumer(void *arg) {
  int i, item;
  for (i=0; i<NITERS; i++) {</pre>
    /* Read item from buf */
    P(&shared.full);
    item = shared.buf;
    V(&shared.empty);
    /* Consume item */
    printf("consumed %d\n", item);
  return NULL;
```

Counting with Semaphores

- Remember, it's a non-negative integer
 - So, values greater than 1 are legal
- Lets repeat thing_5() 5 times for every 3 of thing_3()

```
/* thing_5 and thing_3 */
#include "csapp.h"

sem_t five;
sem_t three;

void *five_times(void *arg);
void *three_times(void *arg);
```

```
int main() {
 pthread t tid five, tid three;
 /* initialize the semaphores */
 Sem init(&five, 0, 5);
 Sem init(&three, 0, 3);
 /* create threads and wait */
 Pthread create (&tid five, NULL,
                 five times, NULL);
 Pthread create (&tid three, NULL,
                 three times, NULL);
```

Counting with semaphores (cont)

Initially: five = 5, three = 3

```
/* thing 5() thread */
void *five times(void *arg) {
  int i;
  while (1) {
    for (i=0; i<5; i++) {
      /* wait & thing 5() */
      P(&five);
      thing 5();
    V(&three);
    V(&three);
    V(&three);
  return NULL;
```

```
/* thing 3() thread */
void *three times(void *arg) {
  int i;
  while (1) {
    for (i=0; i<3; i++) {
      /* wait & thing_3() */
     P(&three);
      thing 3();
    V(&five);
   V(&five);
   V(&five);
    V(&five);
    V(&five);
  return NULL;
```

Producer-Consumer on an *n*-element Buffer

- Requires a mutex and two counting semaphores:
 - mutex: enforces mutually exclusive access to the buffer
 - slots: counts the available slots in the buffer
 - i t.ems: counts the available items in the buffer
- Implemented using a shared buffer package called sbuf.

sbuf Package - Declarations

```
#include "csapp.h"
typedef struct {
   int *buf; /* Buffer array */
   int n; /* Maximum number of slots */
   int front;      /* buf[(front+1)%n] is first item */
   int rear; /* buf[rear%n] is last item */
   sem_t mutex; /* Protects accesses to buf */
   sem_t slots; /* Counts available slots */
   sem_t items; /* Counts available items */
} sbuf t:
void sbuf_init(sbuf_t *sp, int n);
void sbuf deinit(sbuf t *sp);
void sbuf insert(sbuf t *sp, int item);
int sbuf remove(sbuf t *sp);
                                                                   sbuf.h
```

sbuf Package - Implementation

Initializing and deinitializing a shared buffer:

```
/* Create an empty, bounded, shared FIFO buffer with n slots */
void sbuf init(sbuf t *sp, int n)
    sp->buf = Calloc(n, sizeof(int));
                         /* Buffer holds max of n items */
    sp->n = n:
    sp->front = sp->rear = 0; /* Empty buffer iff front == rear */
    Sem_init(&sp->mutex, 0, 1); /* Binary semaphore for locking */
    Sem init(&sp->slots, 0, n); /* Initially, buf has n empty slots */
    Sem_init(&sp->items, 0, 0); /* Initially, buf has 0 items */
/* Clean up buffer sp */
void sbuf deinit(sbuf t *sp)
   Free (sp->buf);
```

sbuf Package - Implementation

Inserting an item into a shared buffer:

sbuf Package - Implementation

Removing an item from a shared buffer:

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Readers-Writers Problem

Generalization of the mutual exclusion problem

Problem statement:

- Reader threads only read the object
- Writer threads modify the object
- Writers must have exclusive access to the object
- Unlimited number of readers can access the object

Occurs frequently in real systems, e.g.,

- Online airline reservation system
- Multithreaded caching Web proxy

Variants of Readers-Writers

- First readers-writers problem (favors readers)
 - No reader should be kept waiting unless a writer has already been granted permission to use the object
 - A reader that arrives after a waiting writer gets priority over the writer
- Second readers-writers problem (favors writers)
 - Once a writer is ready to write, it performs its write as soon as possible
 - A reader that arrives after a writer must wait, even if the writer is also waiting
- Starvation (where a thread waits indefinitely) is possible in both cases

Solution to First Readers-Writers Problem

Readers:

```
int readcnt; /* Initially = 0 */
sem_t mutex, w; /* Initially = 1 */
void reader(void)
    while (1) {
         P(&mutex):
         readcnt++:
         if (readcnt == 1) /* First in */
           P(&w):
         V(&mutex);
         /* Critical section */
         /* Reading happens */
         P(&mutex):
         readcnt--;
         if (readcnt == 0) /* Last out */
           V(&w);
         V(&mutex);
```

Writers:

```
void writer(void)
{
    while (1) {
        P(&w);

        /* Critical section */
        /* Writing happens */

        V(&w);
    }
}
```

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Crucial concept: Thread Safety

- Functions called from a thread must be thread-safe
- Def: A function is thread-safe iff it will always produce correct results when called repeatedly from multiple concurrent threads
- Classes of thread-unsafe functions:
 - Class 1: Functions that do not protect shared variables
 - Class 2: Functions that keep state across multiple invocations
 - Class 3: Functions that return a pointer to a static variable
 - Class 4: Functions that call thread-unsafe functions ©

Thread-Unsafe Functions (Class 1)

- Failing to protect shared variables
 - Fix: Use *P* and *V* semaphore operations
 - Example: goodcnt.c
 - Issue: Synchronization operations will slow down code

Thread-Unsafe Functions (Class 2)

- Relying on persistent state across multiple function invocations
 - Example: Random number generator that relies on static state

```
static unsigned int next = 1;
/* rand: return pseudo-random integer on 0..32767 */
int rand(void)
   next = next*1103515245 + 12345;
   return (unsigned int) (next/65536) % 32768;
/* srand: set seed for rand() */
void srand(unsigned int seed)
   next = seed;
```

Thread-Safe Random Number Generator

- Pass state as part of argument
 - and, thereby, eliminate global state

```
/* rand_r - return pseudo-random integer on 0..32767 */
int rand_r(int *nextp)
{
    *nextp = *nextp * 1103515245 + 12345;
    return (unsigned int) (*nextp/65536) % 32768;
}
```

Consequence: programmer using rand_r must maintain seed

Thread-Unsafe Functions (Class 3)

- Returning a pointer to a static variable
- Fix 1. Rewrite function so caller passes address of variable to store result
 - Requires changes in caller and callee
- Fix 2. Lock-and-copy
 - Requires simple changes in caller (and none in callee)
 - However, caller must free memory.

Thread-Unsafe Functions (Class 4)

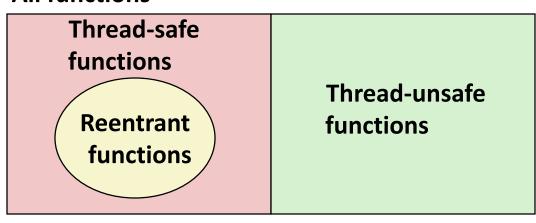
Calling thread-unsafe functions

- Calling one thread-unsafe function makes the entire function that calls it thread-unsafe
- Fix: Modify the function so it calls only thread-safe functions ©

Reentrant Functions

- Def: A function is reentrant iff it accesses no shared variables when called by multiple threads.
 - Important subset of thread-safe functions
 - Require no synchronization operations
 - Only way to make a Class 2 function thread-safe is to make it reetnrant (e.g., rand_r)

All functions



Thread-Safe Library Functions

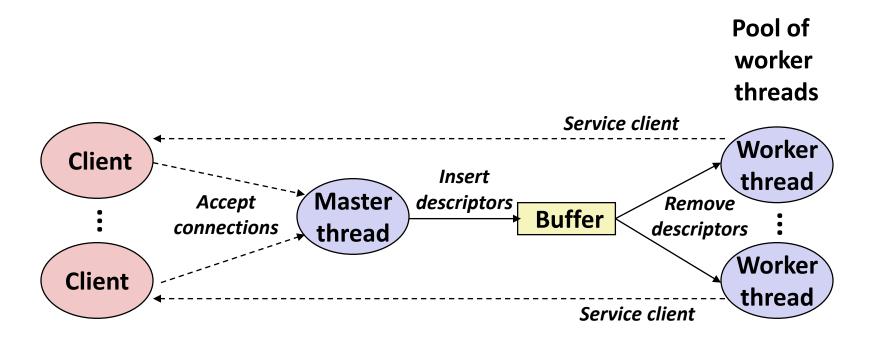
- All functions in the Standard C Library (at the back of your K&R text) are thread-safe
 - Examples: malloc, free, printf, scanf
- Most Unix system calls are thread-safe, with a few exceptions:

Thread-unsafe function	Class	Reentrant version
asctime	3	asctime_r
ctime	3	ctime_r
gethostbyaddr	3	gethostbyaddr_r
gethostbyname	3	gethostbyname_r
inet_ntoa	3	(none)
localtime	3	localtime_r
rand	2	rand_r

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Putting It All Together: Prethreaded Concurrent Server



```
sbuf_t sbuf; /* Shared buffer of connected descriptors */
int main(int argc, char **argv)
{
    int i, listenfd, connfd;
    socklen t clientlen;
    struct sockaddr storage clientaddr;
    pthread t tid:
    listenfd = Open listenfd(argv[1]);
    sbuf init(&sbuf, SBUFSIZE);
    for (i = 0; i < NTHREADS; i++) /* Create worker threads */
        Pthread create (&tid, NULL, thread, NULL);
    while (1) {
         clientlen = sizeof(struct sockaddr storage);
         connfd = Accept(listenfd, (SA *) &clientaddr, &clientlen);
         sbuf insert(&sbuf, connfd); /* Insert connfd in buffer */
                                                               echoservert pre.c
```

Worker thread routine:

```
void *thread(void *vargp)
{
    Pthread_detach(pthread_self());
    while (1) {
        int connfd = sbuf_remove(&sbuf); /* Remove connfd from buf */
        echo_cnt(connfd); /* Service client */
        Close(connfd);
    }
}
echoservert_pre.c
```

echo_cnt initialization routine:

```
static int byte_cnt;  /* Byte counter */
static sem_t mutex;  /* and the mutex that protects it */

static void init_echo_cnt(void)
{
    Sem_init(&mutex, 0, 1);
    byte_cnt = 0;
}
echo_cnt.c
```

Worker thread service routine:

```
void echo cnt(int connfd)
    int n:
    char buf[MAXLINE];
   rio_t rio;
    static pthread once t once = PTHREAD ONCE INIT;
    Pthread once (&once, init echo cnt);
    Rio readinith (&rio, connfd);
    while((n = Rio readlineb(&rio, buf, MAXLINE)) != 0) {
        P(&mutex);
         byte cnt += n;
        printf("thread %d received %d (%d total) bytes on fd %d\n",
                (int) pthread self(), n, byte cnt, connfd);
        V(&mutex):
        Rio writen (connfd, buf, n);
                                                                        echo cnt.o
```