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BRIEFINGS

Trace Me If You can: Bypassing Linux Syscall Tracing

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About Rex Guo

- Principal Engineer @ Lacework
 - Behavior anomaly detection (Polygraph)
 - CSPM
- Engineering Manager @ Startups
 - Confluera (XDR)
 - Tetration (CWPP, now part of Cisco)
- Conference speaker at Blackhat, DEFCON,...
- @Xiaofei_REX



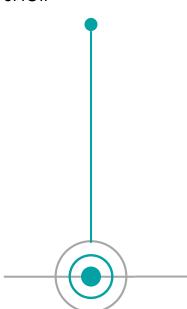


About Junyuan Zeng

- Linkedin
 - Senior Software Engineer: Kubernetes
- JD.com
 - Staff Security Architect/Engineer:
 Cloud native security
- Samsung Research America & FireEye
 - Staff Security Software Engineer/Researcher: Mobile security

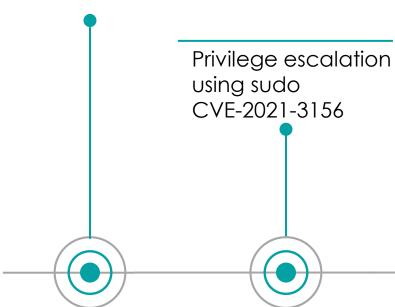


Log4shell RCE on joe-box and executed a reverse shell

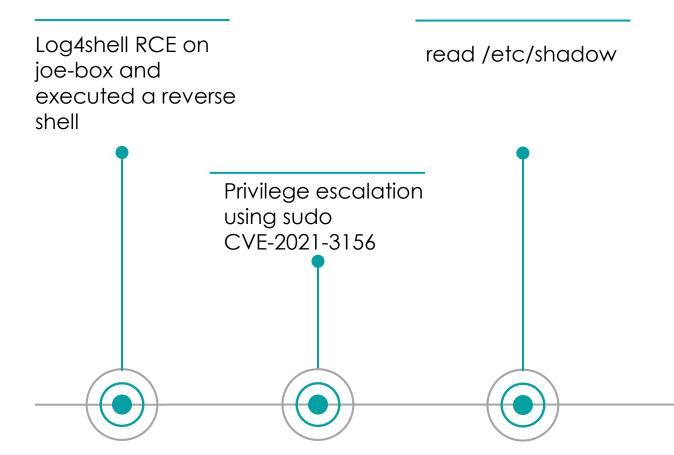




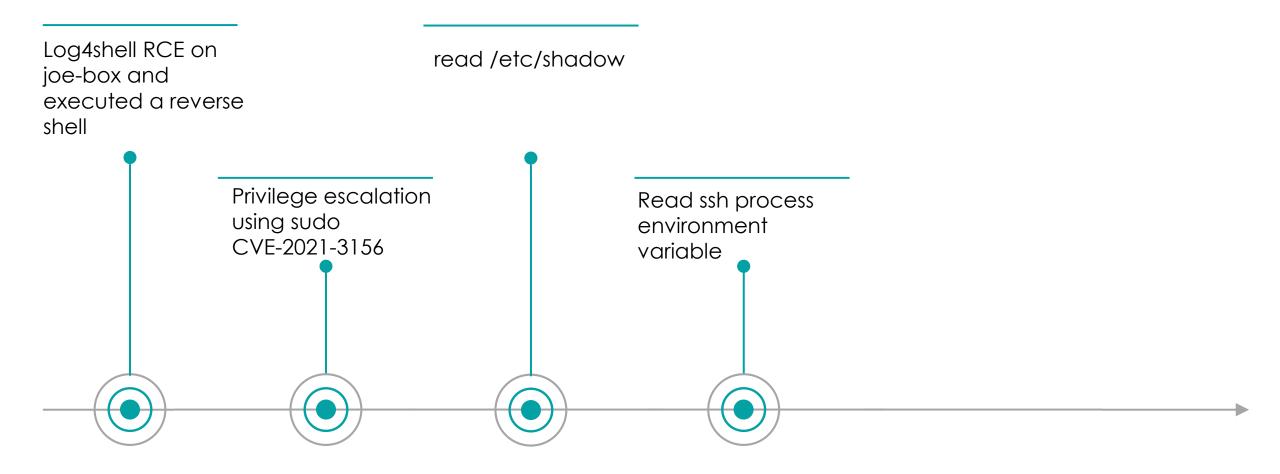
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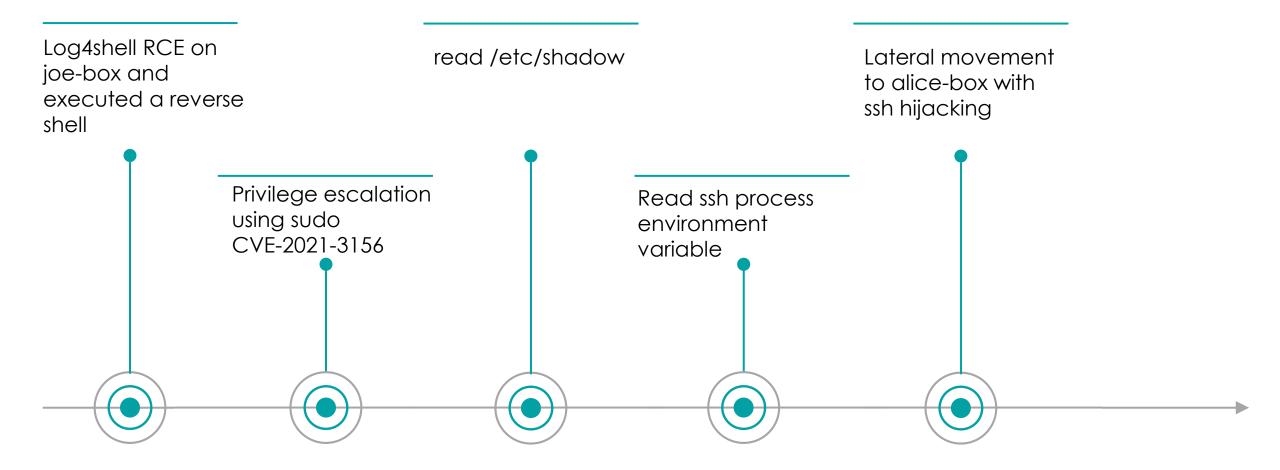




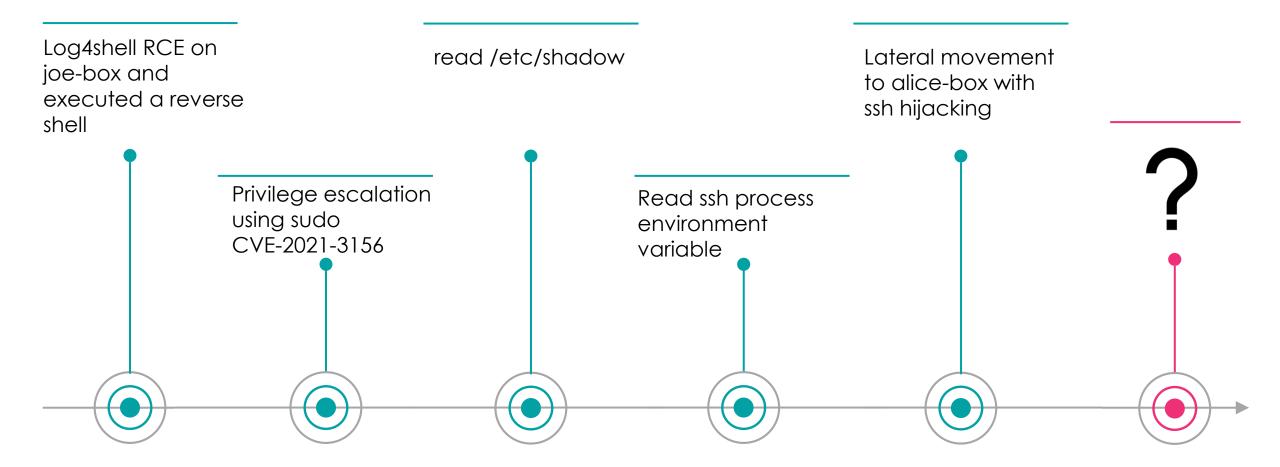






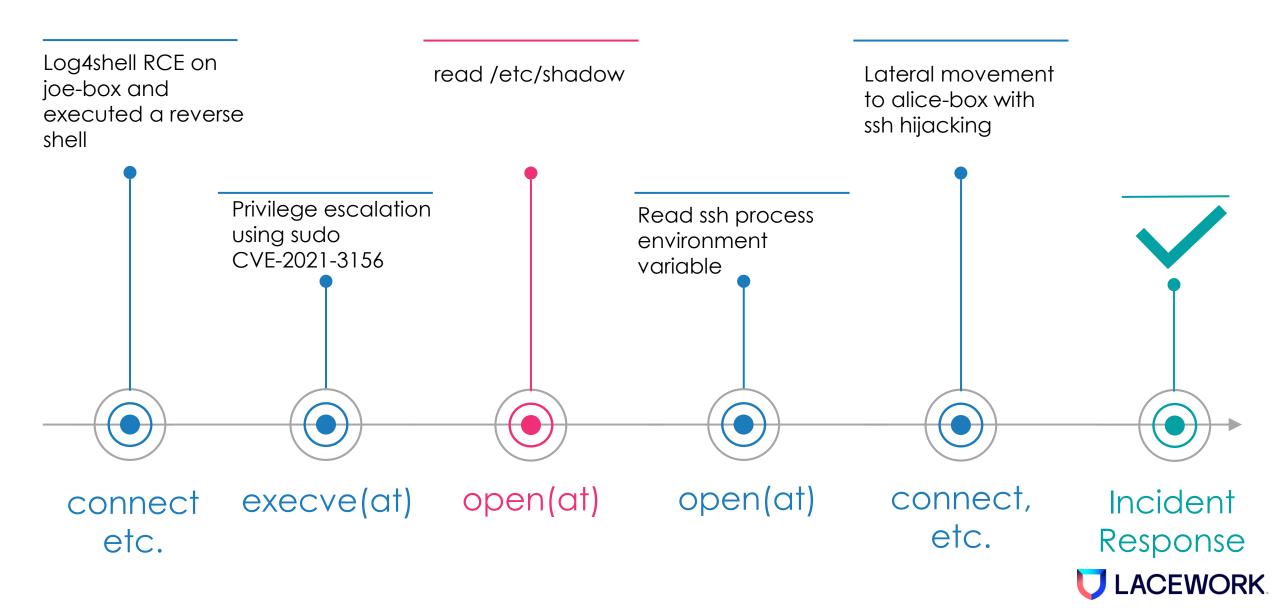








An Incident - A Defender's View



Detection Rule Example

```
rule: untrusted program reads /etc/shadow
condition:
    syscall == open(at)
    and has read permission
    and filename == /etc/shadow
    and program is not in allowlist
```

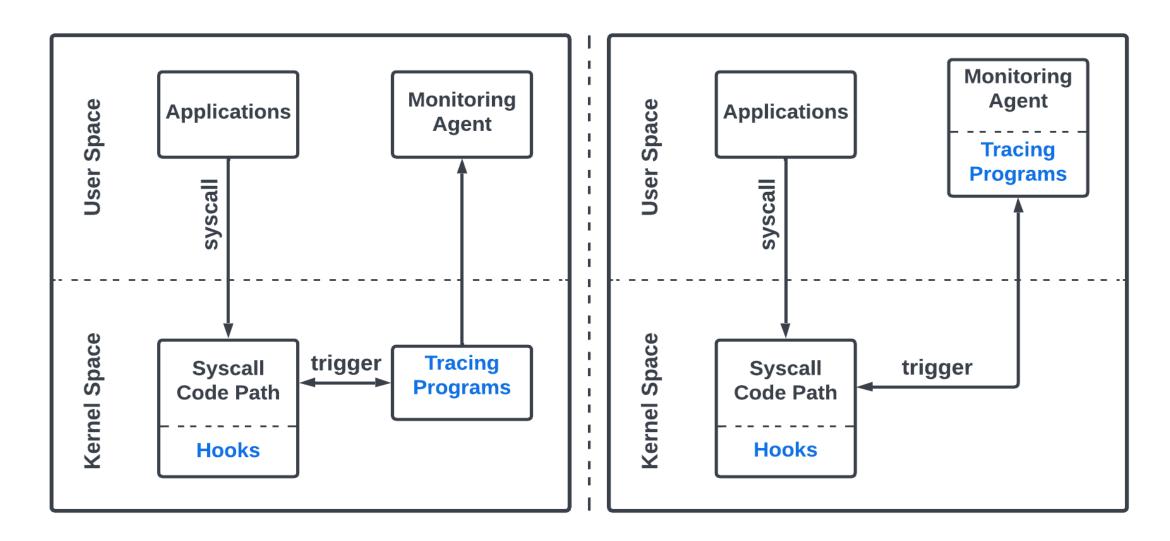


Agenda

- Syscall Tracing
- Vulnerabilities
- Exploitations
- Mitigations
- Takeaways



System Call Tracing





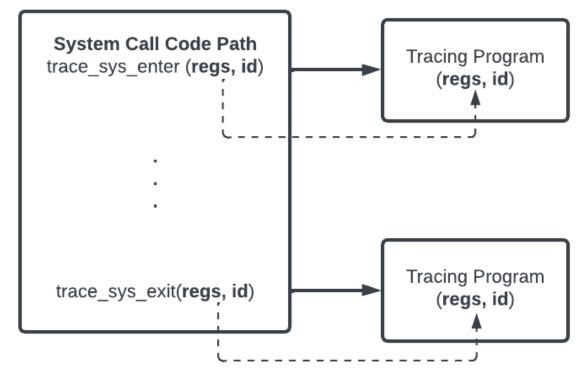
System Call Tracing – Tracing Program

- Tracing programs collect system call data, e.g., arguments
- Tracing programs can "attach" to different hooks
 - tracepoints, kprobe, ptrace etc.
- Tracing programs implementations
 - Linux native mechanisms: ftrace, perf_events etc.
 - Kernel modules, eBPF probe and user space programs



System Call Tracing – tracepoint

- tracepoint
 - Kernel static hook
 - Linux kernel provides sys_enter and sys_exit

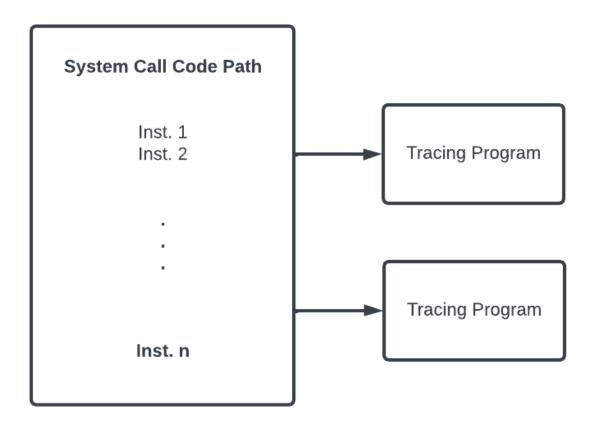


Low overhead but only static interceptions



System Call Tracing – kprobe

- kprobe
 - Dynamic hook in the kernel
 - Register tracing programs on instructions in syscall code path
 - Dynamic but slow compared to tracepoint and need to know exactly how data is placed on the stack and registers





System Call Tracing – ptrace

- ptrace
 - A static hook
 - No Kernel Module/eBPF program are needed
 - Performance overhead is high
 - Can combine with seccomp to reduce overhead
- Others (LD_PRELOAD etc.)



Cloud Workloads

- Virtual machines
 - AWS EC2 instances
 - Google VM instances
- Containers on customer-managed VMs
 - AWS EC2 tasks
 - Standard GKE workloads (e.g. DaemonSet etc.)
 - AKS workloads
- Serverless containers: have no access to the host
 - AWS Fargate tasks
 - GCP Cloud Run services
- Others (AWS Lambda etc.)



System Call Tracing for Cloud Workloads

Workload	System Call Tracing				
VMs	 Hooks: tracepoint, kprobe, ptrace Tracing programs: kernel programs (eBPF, kernel Module), user programs Tools: Falco eBPF/kernel Module, Falco pdig 				
Containers	 Hooks: tracepoint, kprobe, ptrace Tracing programs: kernel programs (eBPF, kernel Module), user programs Tools: Falco eBPF/kernel Module, Falco pdig 				
Serverless Containers	 Hooks: ptrace Tracing programs: user programs Tools: Falco pdig 				



Open Source Projects

- Falco
 - Open source endpoint security monitoring project in CNCF
 - 5K+ github stars
 - Falco supports syscall tracing techniques:
 - tracepoint + kernel module
 - tracepoint + eBPF probe
 - pdig: ptrace + userspace program
- Falco pdig
 - Support syscall tracing of serverless workloads



TOCTOU in Syscall Tracing

- sys_connect(int fd, struct sockaddr __user * uservaddr, int addrlen)
- TOC (Time-Of-Check): tracing programs dereference this user space pointer
- TOU (Time-Of-Use): the kernel dereferences this user space pointer





syscall enter !ptrace_report_syscall(regs, message) ptrace/seccomp/sysenter tracepoint secure computing(struct seccomp data{regs...}) !trace sys enter(regs, regs->orig ax) Syscall Table (x86 64) 42 sys connect ----→long sys connect((int fd, 43 sys accept struct sockaddr user *uservaddr, int addrlen)) 44 sys_sendto Execution Flow struct filename *tmp; ret = move addr to kernel (uservaddr, addrlen, &address); if (!ret) ret = sys connect file (f.file, &address, addrlen, 0); trace sys exit(regs, regs->ax)

ptrace report syscall(regs, message)

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syscall enter ptrace report syscall(regs, message) secure computing(struct seccomp data{regs...}) trace sys enter(regs, regs->orig ax) Syscall Table (x86 64) 42 sys connect →long sys connect((int fd, struct sockaddr __user *uservaddr, int addrlen))! 43 sys accept 44 sys_sendto **Execution Flow** struct filename *tmp; ret = move addr to kernel (uservaddr, addrlen, &address); if (!ret) ret = sys connect file (f.file, &address, addrlen, 0);

trace_sys_exit(regs, regs->ax)
ptrace report syscall(regs, message)



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trace_sys_exit(regs, regs->ax)

ptrace_report_syscall(regs, message)

ptrace/sysexit tracepoint



syscall enter ptrace report syscall(regs, message) Userspace pointer pointing to "socket address" secure computing(struct seccomp data{regs...}) trace sys enter(regs, regs->orig ax) Syscall Table (x86 64) 42 sys connect ----→long sys connect((int fd, 43 sys accept struct sockaddr user *uservaddr, int addrlen)) 44 sys_sendto **Execution Flow** struct filename *tmp; ret = move addr to kernel (uservaddr, addrlen, &address); if (!ret) ret = sys connect file (f.file, &address, addrlen, 0);

trace sys exit(regs, regs->ax)

ptrace report syscall(regs, message)

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syscall enter Syscall Table (x86 64) 42 sys connect 43 sys accept 44 sys_sendto **Execution Flow**

trace sys exit(regs, regs->ax)

ptrace report syscall(regs, message)

```
ptrace report syscall(regs, message)
                                                             Userspace pointer pointing to "socket address"
secure computing(struct seccomp data{regs...})
                                                            Kernel pointer pointing to "socket address"
trace sys enter(regs, regs->orig ax)
                    ---→long sys connect((int fd,
                             struct sockaddr user *uservaddr, int addrlen))
                           struct filename *tmp;
                           ret = move addr to kernel
                                    (uservaddr, addrlen, &address);
                           if (!ret)
                             ret = sys connect file
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```



syscall enter ptrace report syscall(regs, message) secure computing(struct seccomp data{regs...}) trace sys enter(regs, regs->orig ax) Syscall Table (x86 64) 42 sys connect ----→long sys connect((int fd, 43 sys accept struct sockaddr user *uservaddr, int addrlen)) 44 sys_sendto **Execution Flow** struct filename *tmp; ret = move addr to kernel (uservaddr, addrlen, &address); if (!ret) ret = sys connect file (f.file, &address, addrlen, 0); trace sys exit(regs, regs->ax)

ptrace report syscall(regs, message)

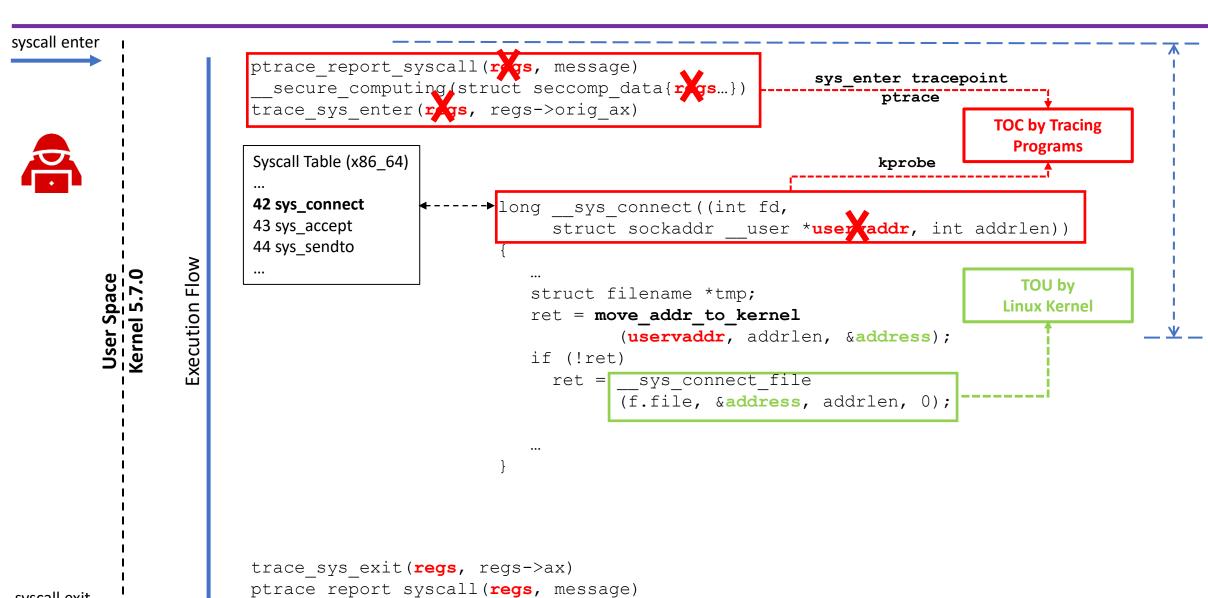
TOU by

Linux Kernel

syscall enter ptrace report syscall(regs, message) secure computing(struct seccomp data{regs...}) trace sys enter(regs, regs->orig ax) Syscall Table (x86 64) 42 sys connect ·→long sys connect((int fd, 43 sys accept struct sockaddr user *uservaddr, int addrlen)) 44 sys_sendto **Execution Flow TOU by** struct filename *tmp; **Linux Kernel** ret = move addr to kernel (uservaddr, addrlen, &address); if (!ret) ret = sys connect file (f.file, &address, addrlen, 0); trace sys exit(regs, regs->ax) ptrace report syscall(regs, message) syscall exit

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syscall exit



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syscall enter



Execution Flow

trace_sys_exit(revs, regs->ax)

ptrace report syscall (ress, message)

```
ptrace report syscall(regs, message)
secure computing(struct seccomp data{regs...})
trace sys enter(regs, regs->orig ax)
Syscall Table (x86 64)
42 sys connect
                      →long sys connect((int fd,
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44 sys sendto
                                                                           TOU by
                           struct filename *tmp;
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                           ret = move addr to kernel
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                           if (!ret)
                             ret = sys connect file
                                    (f.file, &address, addrlen, 0);
```

sys exit tracepoint

ptrace

TOC by Tracing

Programs

```
syscall enter
                        ptrace report syscall(regs, message)
                                                                                                        Falco pdig
                           secure computing(struct seccomp data{regs...})
                        trace sys enter(regs, regs->orig ax)
                         Syscall Table (x86 64)
                         42 sys connect
                                                -→long sys connect((int fd,
                         43 sys accept
                                                        struct sockaddr user *uservaddr, int addrlen))
                         44 sys sendto
                  Execution Flow
                                                                                                         TOU by
                                                      struct filename *tmp;
                                                                                                       Linux Kernel
                                                      ret = move addr to kernel
                                                               (uservaddr, addrlen, &address);
                                                      if (!ret)
                                                        ret = sys connect file
                                                               (f.file, &address, addrlen, 0);
                                                                                                    Falco (<0.31.1) kernel
                         trace sys exit(regs, regs->ax)
                                                                                                       module/eBPF
                         ptrace report syscall (regs, message)
syscall exit
                                                                                                         Falco pdig
```

TOCTOU Windows across Kernels

- TOCTOU windows exist since the initial release of tracepoint/ptrace
- Expected behaviors
- Monitor kernel memory

tracepoint and ptrace have TOCTOU issues!

We knew! They are designed for perf/debug





TOCTOU — Falco

- User space pointers are dereferenced by
 - sys_exit tracepoint (kernel module, eBPF)
 - sys_exit ptrace (pdig)
- Falco older than v0.31.1
 - Check with vendors which commercial versions are affected
- 12/06/2021 Issue reported (CVE-2022-26316)
- 03/11/2022 Mitigation implemented (<u>Advisory</u>)
 - For selected syscalls, compare sys_enter and sys_exit tracepoint data (Falco LKM, eBPF)
 - Compare sys_enter and sys_exit ptrace data (Falco pdig)



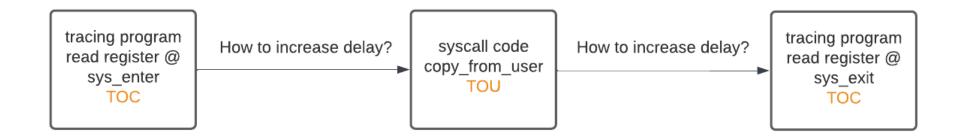
TOCTOU — Falco

• We evaluated the important syscalls in <u>Falco rules</u>.

Syscall	Category	тостои?	Exploitable by blocking condition	Exploitable by DC29 attack
connect	Network	Υ	Υ	Υ
sendto/sendmsg	Network	Υ	N	Υ
open(at)	File	Υ	Υ	Υ
execve	File	N	N*	N*
rename	File	Υ	Υ	Υ
renameat(2)	File	Y	Υ	Υ
mkdir(at)	File	Υ	Υ	Υ
rmdir	File	Υ	Υ	Υ
unlink(at)	File	Υ	Υ	Υ
symlink(at)	File	Υ	Υ	Υ
chmod/fchmod(at)	File	Υ	Υ	Υ
creat	File	Υ	Υ	Υ



Exploit Requirements

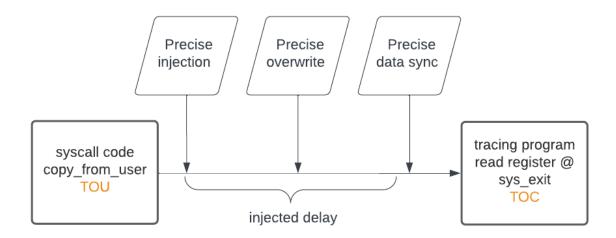


- Exploitation requirements
 - No additional privilege and capabilities
 - Control the time to inject the delay
 - Enough delay for pointer overwrite
 - Reliable



Exploit Strategy 1 (DEFCON 29)

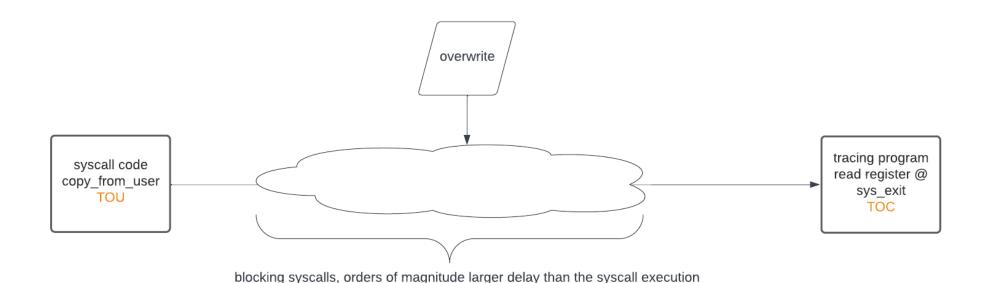
- Injected delay is small
- Requires Userfaultfd syscall for precise injection while pausing the kernel execution
- seccomp can block userfaultfd syscall (e.g., docker default seccomp profile)
- Falco's mitigation was to detect userfaultfd





Exploit Strategy 2

- Injected delay >> the syscall execution time
- No Precise control is required





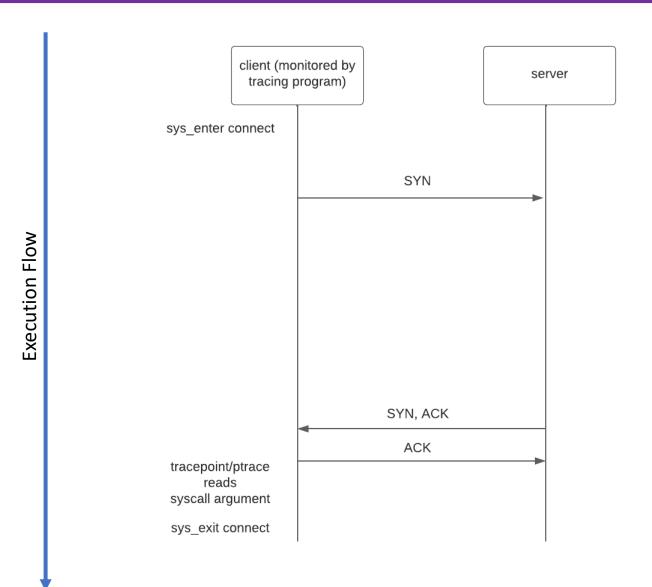
Syscall Built-in Delay

- Attackers can trigger significant syscall delays by introducing:
 - Blocking conditions (attack sys_exit)
 - Seccomp rules (attack sys_enter)
- Syscall can get "blocked"

Categories	Syscalls
Process	fork/exec/exit/wait/
File system	open(at)/symlink(at)/read/write/
Networking	connect/accept/socket/
Security	seccomp/keyctl/
Many others	

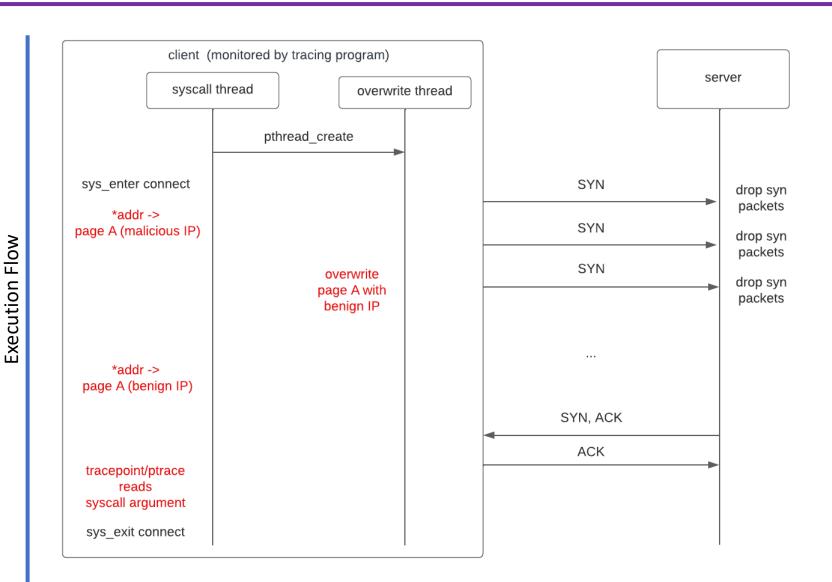


Connect Syscall





Bypassing Connect Syscall Tracing (Demo)





Blocking Syscalls (File Systems)

- File system syscalls are all affected
 - open/openat
 - creat
 - rename/renameat/renameat2
 - mkdir/mkdirat
 - rmdir
 - Other file system syscalls with pointer arguments
- Other syscalls are also affected due to fetching files from file systems.
 - execve/execveat



Filesystem in USErspace - FUSE

- User space filesystem framework
- Used as remote storage FUSE
 - Access the remote files as local ones
 - Faster evolvement and don't panic the kernel etc.
- Remote storage FUSE examples:
 - gcsfuse¹: developed by Google for GCS
 - s3fs-fuse²: Amazon S3
 - BlobFuse³: developed by Azure for Blob storage
 - MezzFS⁴: developed and deployed @ Netflix
 - Many others (sshfs etc.)



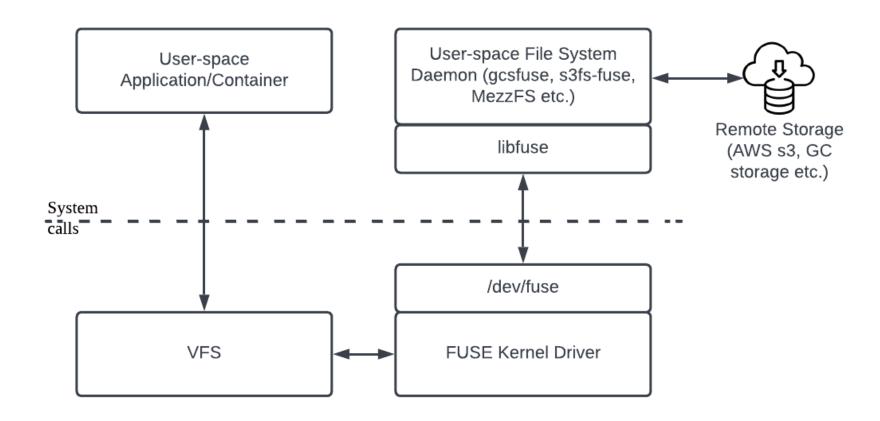
^{1,} https://github.com/GoogleCloudPlatform/gcsfuse

^{2,} https://github.com/s3fs-fuse/s3fs-fuse

^{3,} https://github.com/Azure/azure-storage-fuse

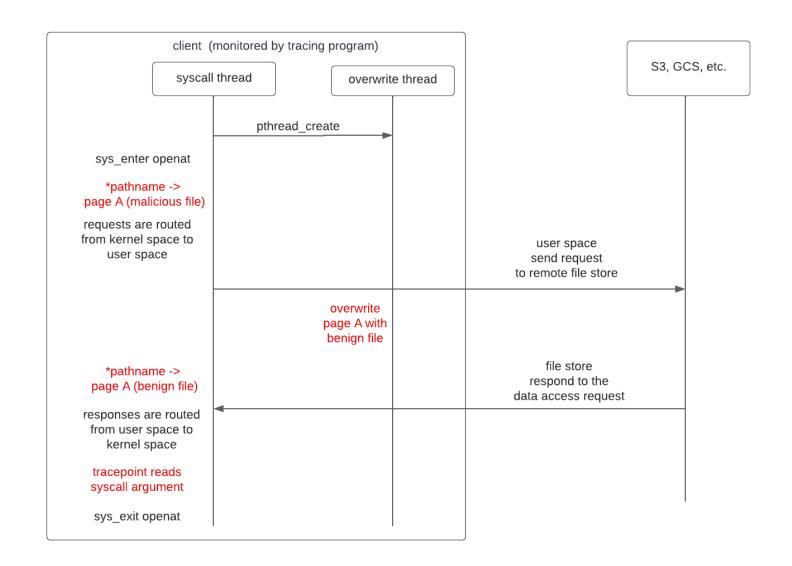
^{4,} https://netflixtechblog.com/mezzfs-mounting-object-storage-in-netflixs-media-processing-platform-cda01c446ba

Remote Storage FUSE - Architecture





Bypassing Openat Tracing (Demo)







TOCTOU – sys_enter (Connect)

```
syscall enter
                        ptrace report syscall(regs, message)
                                                                                                        ptrace
                          secure computing(struct seccomp data{regs...})
                        trace sys enter (regs, regs->orig ax)
                                                                                                       seccomp
                        Syscall Table (x86 64)
                        42 sys connect
                                                ·→long sys connect((int fd,
                        43 sys accept
                                                       struct sockaddr user *uservaddr, int addrlen))
                        44 sys_sendto
                 Execution Flow
                                                                                                       TOU by
                                                     struct filename *tmp;
                                                                                                     Linux Kernel
                                                     ret = move addr to kernel
                                                              (uservaddr, addrlen, &address);
                                                     if (!ret)
                                                       ret = sys connect file
                                                              (f.file, &address, addrlen, 0);
                        trace sys exit(regs, regs->ax)
                        ptrace report syscall(regs, message)
```

syscall exit

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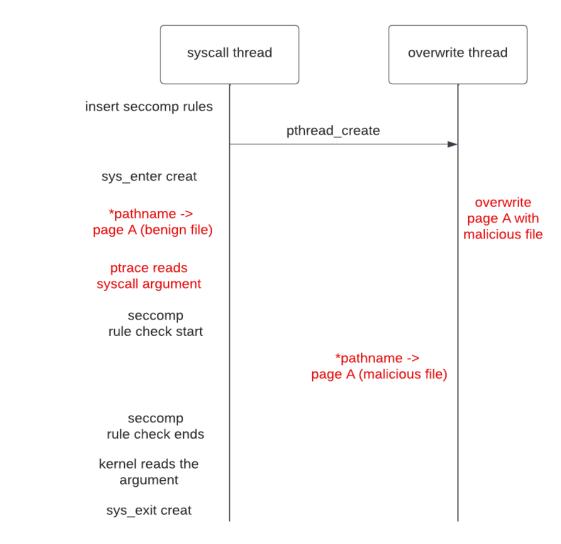
Seccomp Introduction

- Kernel level mechanism to restrict syscalls
- Modern sandboxes heavily relies on seccomp
- Developers can write rules to:
 - allow/block certain syscalls
 - allow/block syscalls based on argument values
- These rules can be quite complex (<u>read more</u>)
 - More rules takes more time to compute
- First inserted rules are evaluated last



Attacking Syscall Enter

Execution Flow

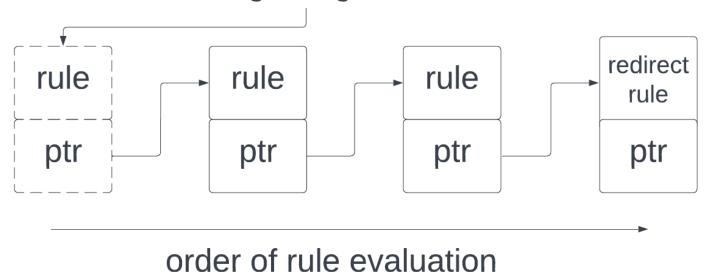




ptrace + seccomp redirect

Tracer starts App

New rules will be inserted at the beginning of the list





Exploitation and Mitigations

Tracing location	TOCTOU Exploitation	Mitigations
ptrace @ sys_enter	Seccomp filter insertion	 ptrace + seccomp redirect to start the app. Inspect seccomp filters already inserted when attaching to a running app
tracepoint @ sys_enter	Unreliable	N/A
tracepoint @ sys_exit	Blocking syscall (This talk)Phantom attack v1 (DEFCON 29)	- Compare tracepoint sys_enter and sys_exit args
ptrace @ sys_exit	Same as above	 Deploy all mitigations for ptrace @ sys_enter Compared the sys_enter and sys_exit syscall args
kprobe @ kernel internal	It depends	Read the kernel copy of the syscall args - LSM (BPF-LSM) - Other interfaces



Key Takeaways

- 1. Linux kernel tracing can be bypassed reliably
 - Check your security tools
- Mitigation is complex (workload type and kernel compatibility)
 - Check your security tools' mitigation claims
- 3. Correlate different data sources
- 4. Know your normal

- Discussing further?
 - @Xiaofei_REX / rex.guo *NOSPAM* lacework DOT com
 - jzeng04 *NOSPAM* gmail DOT com
 - POC: https://github.com/rexguowork/phantom-attack



Acknowledgement

- Joel Schopp (Linux kernel / Security)
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- John Dickson

