

# How I Learned to Stop Fuzzing and Find More Bugs

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# Agenda

- **Introduction to fuzzing**
  - What we mean by fuzzing
  - Challenges with fuzzing
- **Introduction to static analysis**
  - How static analysis works
  - Examples of bugs static analysis is good at finding
  - Untapped potential: Customization
- **Experiment**
  - Fuzzing versus static analysis
- **Conclusion**

# What is Fuzzing?

- Encompasses runtime testing that attempts to induce faults in software systems by inputting random or semi-random values
- Introduced by Barton Miller at the University of Wisconsin, Madison in 1990 ([cs.wisc.edu/~bart/fuzz/](http://cs.wisc.edu/~bart/fuzz/))



# Examples of Tools

- We're talking about tools such as:
  - SPIKE  
[www.immunitysec.com/resources-freesoftware.shtml](http://www.immunitysec.com/resources-freesoftware.shtml)
  - Peach  
<http://peachfuzz.sourceforge.net>
  - PROTOS  
<http://www.ee.oulu.fi/research/ouspg/protos/>
  - ... and many more
- But not specialized black box scanning tools:
  - Cenzic
  - SPI Dynamics (except SPI Fuzzer)
  - Watchfire

# The Inventor's Thoughts on Fuzzing

- 1990: “[Fuzzing] is not a substitute for a formal verification or testing procedures, but rather an inexpensive mechanism to identify bugs...”
- 1995: “While [fuzzing] is effective in finding real bugs in real programs, we are not proposing it as a replacement for systematic and formal testing.”
- 2000: “Simple fuzz testing does not replace more extensive formal testing procedures.”

- Barton Miller

## SC-L Digest, Vol 3, Issue 118:

- **"I would assume that "smart" fuzzing could have lots of manipulations of the HH:mm:ss.f format, so this might be findable using black box testing."**

**- Steve Christey**

# Woulda, Coulda, Shoulda

<http://blogs.msdn.com/sdl>

- “It turns out none of the .ANI fuzz templates had a second “anih” record.

This is now addressed, and we are continually enhancing our fuzzing tools to make sure they add manipulations that duplicate arbitrary object elements better.”

- Michael Howard

# How Fuzzing Works

- **Identify sources of input to a program**
- **Permute or generate pseudorandom input**
- **Use an oracle to monitor for failures**
- **Record the input and state that generate faults**



# Input Sources: File Formats

1. **Identify all valid file formats (e.g. JPG, TIFF, PDF, DOC, XLS)**
2. **Collect a library of valid files**
3. **Malform a file**
4. **Cause the program to consume the file and observe its execution for problems**

```
<?xml version="1.0" encoding="utf-8" ?>
<?xml:stylesheet type="text/xsl" href="http://www.w3.org/2001/XMLSchema-instance"
xmlns:xsd="http://www.w3.org/2001/XMLSchema">
  <Words>
    - <Word>
      <Value>House</Value>
      <Definition>A place where people live which provides
        shelter</Definition>
    - <Word>
      <Value>Bus</Value>
      <Definition>A device to transport people</Definition>
    - <Word>
      <Value>Wood</Value>
      <Definition>Part of a tree</Definition>
    - <Word>
      <Value>Serializer</Value>
      <Definition>An object for transforming another object to or from a
        linear sequence of bytes</Definition>
  </Words>
</Dictionary>
```

# Input Sources: Protocols

- **Create bogus messages**  
(e.g. SMTP, TCP/IP, RPC, SOAP, HTTP)
- **Record-fuzz-replay**
  1. Run a sniffer
  2. Collect a few thousand messages
  3. Fuzz the messages
  4. Replay the fuzzed messages



# Dumb Fuzzing



- **Dumb fuzzing: Modify data randomly**
  - Most input will be invalid
  - Makes good error handling test cases
  - Takes a long time to enumerate valid test cases
  - May test the validation logic of high-level protocols instead of the underlying application

# Smart Fuzzing



- **Smart fuzzing: Aware of data structure**
  - Altering content size
  - Replacing null-terminated strings
  - Altering numeric values or flipping signs
  - $0, 2^n +/- 1$
  - Adding invalid headers, altering header values, duplicating headers, ...

# Challenging Questions with Fuzzing

- Microsoft SDL mandates that you run 100,000 iterations per file format/parser.
- If you find a bug, you reset to 0 and start running another 100,000 with a new seed.
- Why? Does this get you what you need?
- How many input sources were missed?
- How much of the program was tested?
- How long did the tests take to run?
- How good were the tests?

# Challenge: Nebulous File Formats / Protocols

- No problem for a standard Web application
- What about proprietary interfaces?
  - Web Service APIs
  - Network servers
  - Thick client software
- Difficult to enumerate input sources to fuzz
- Even harder to generate valid input
- Requires customization
  - Tool must be tuned to specific input sources and formats

# Challenge: Program Semantics / Reachability

- **Example:**

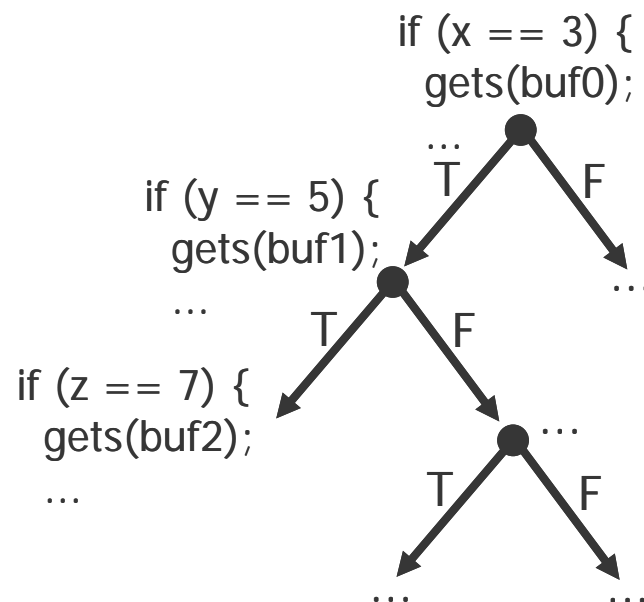
```
if (!strcmp(input1, "static_string") {  
    strcpy(buffer2, input2);  
}
```

- Need to provide value of `input1` equal to "static\_string" and large value of `input2`
- Requires  $N * M$  random inputs to reach bug guarded by two-variable conditions
- May be hard to satisfy some conditionals
- Requires customization
  - Number of input values needed must be narrowed

# Shallow Bugs versus Deep Bugs

## • Fuzzing focuses on shallow bugs

```
foo(int x, int y, int z) {  
  if (x == 3) {          //p = 1/232  
    gets(buf0);  
  
    if (y == 5) {        //p = p * 1/232  
      gets(buf1);  
  
      if (z == 7) {      //p = p * 1/232  
        gets(buf2);  
      }  
    }  
  }  
}
```



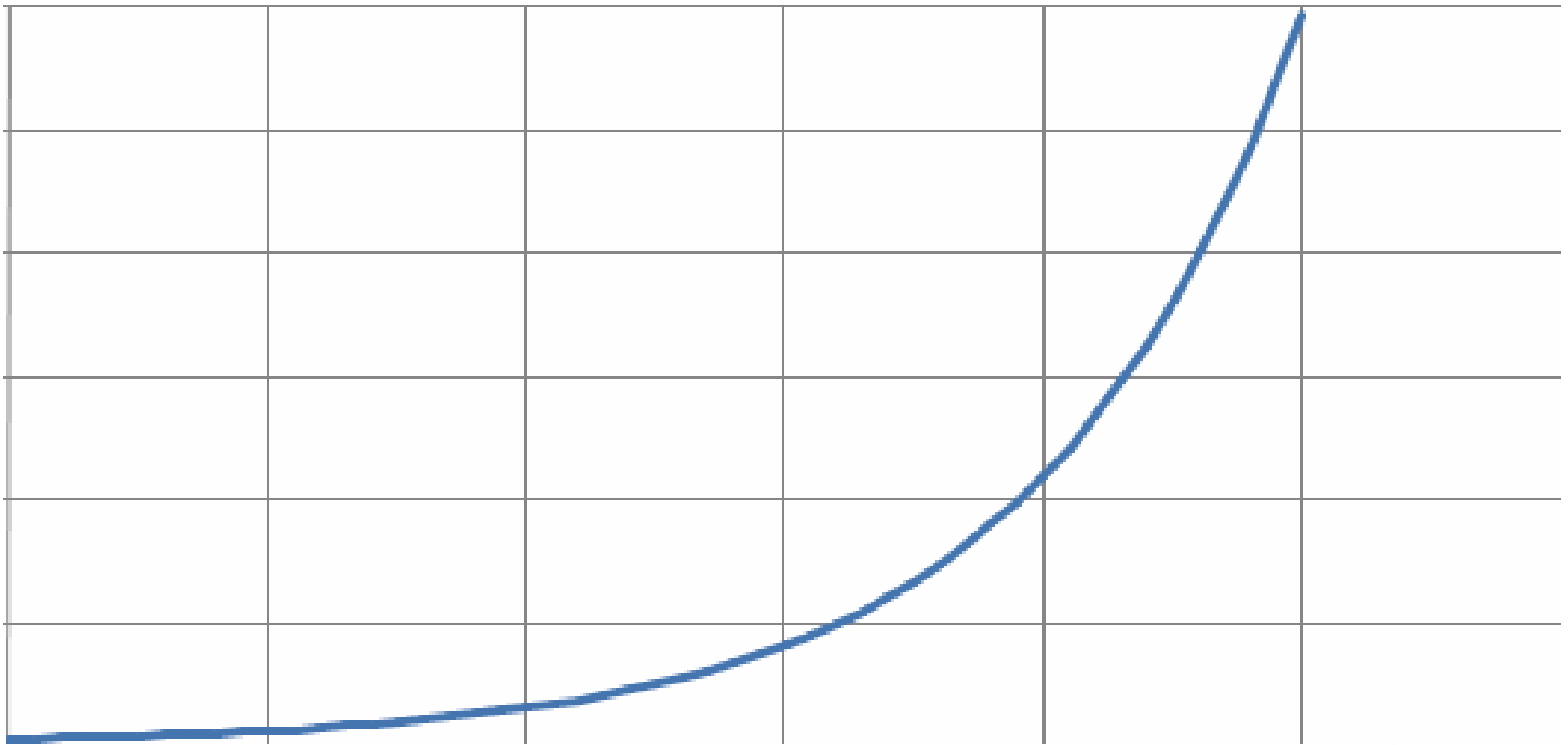
**# of random values of x, y and z to reach each state:**

gets(buf0) =	4,294,967,296
gets(buf1) =	18,446,744,073,709,551,616
gets(buf2) =	79,228,162,514,264,337,593,629,020,928



# Runs Necessary to Reach State Z

- Each conditional adds exponentially to the number of input permutations required to hit a bug
- Running time for the fuzz tests increases accordingly



# Example: Vulnerabilities by Conditional Depth

- **wu-ftpd 2.6.0–Buffer Overflow–extensions.c:**  
`strcpy(curptr->dirname, cwd);`  
Conditional depth: 4
- **wu-ftpd 2.6.0–format string–ftpd.c:**  
`vsnprintf(buf + (n ? 4 : 0), n ?  
sizeof(buf)-4 : sizeof(buf), fmt, ap);`  
Conditional depth: 3
- **OFBiz 1.5–XSS–CommonEvents.java:**  
`out.println(responseString);`  
Conditional depth: 4

# Challenge: Difficult-to-Reach States

- **Airline booking system – overbooked flight**
  - Difficult for a fuzzer to induce
- **Example:**

```
if (flight.seatsAvailable() == 0) {  
    // echo user input to error page - XSS vulnerability  
    ...  
    out.println("Flight "  
                + request.getParameter("flightNumber")  
                + " is overbooked. Please search again.");  
    ...  
}
```

# Challenge: Identifying Errors

- Error reporting conventions differ
- Good design guidelines often require programs to mask errors and error details
- Requires customization
  - Better oracle
  - Binary instrumentation
  - ...



# Finding Bugs with Fuzzing

**Spectrum of ease of detection with fuzzing:**



↖  
**Easy: Shallow cross-site scripting vulnerability (shallowest bugs never leave the client à la JavaScript)**

↗  
**Hard: Many nested conditionals that checks for hard-to-reach states like the overbooked flight**

# Fuzzing Summary

- **Advantages**

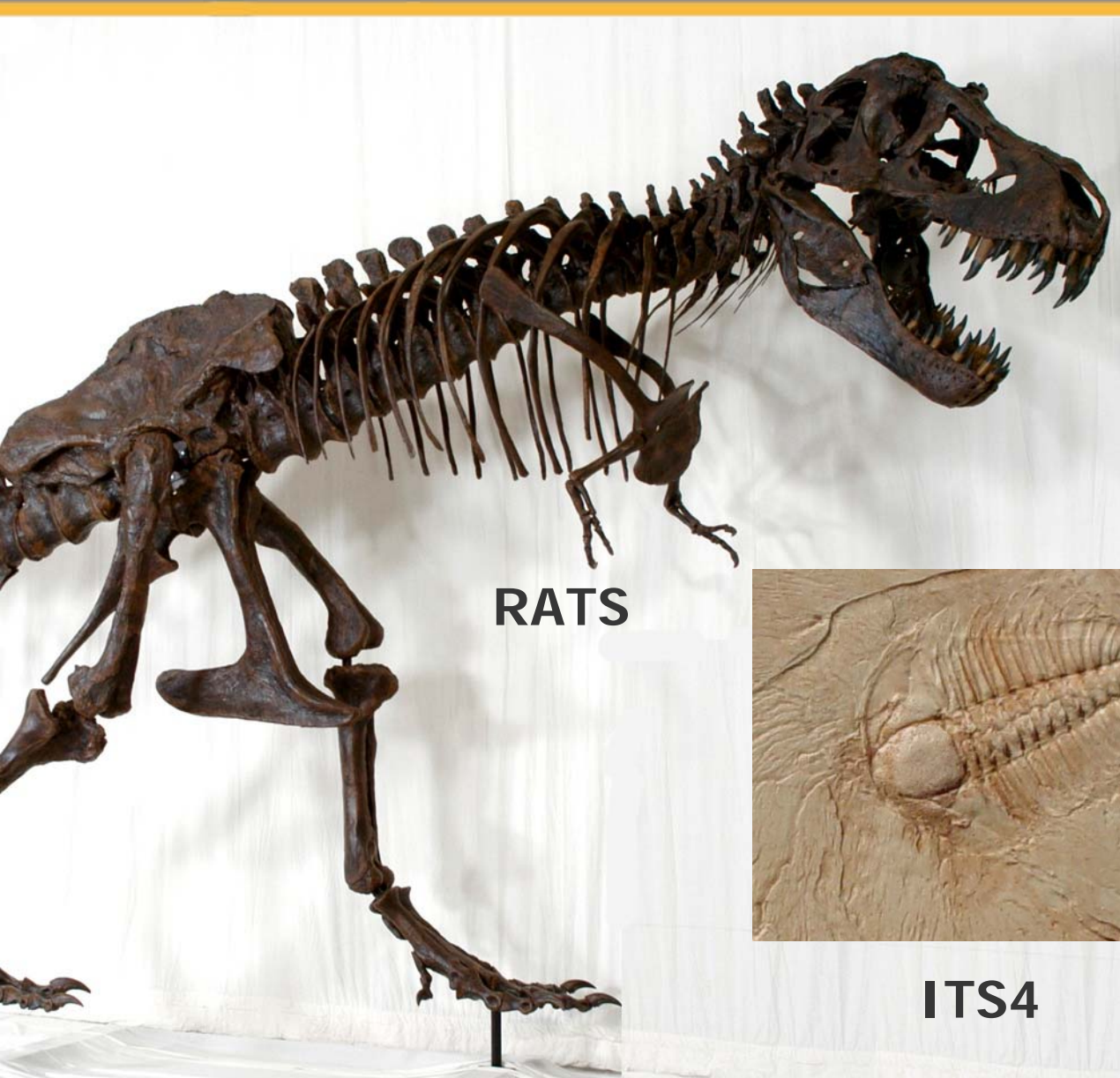
- Requires least effort to find a bug
- Verifiable and reproducible at runtime
- Scalable to programs that use the same file format or protocol

- **Disadvantages**

- Very costly to achieve completeness
- Increasing coverage increases runtime (sometimes exponentially)
- May miss bugs due to inadequate oracle



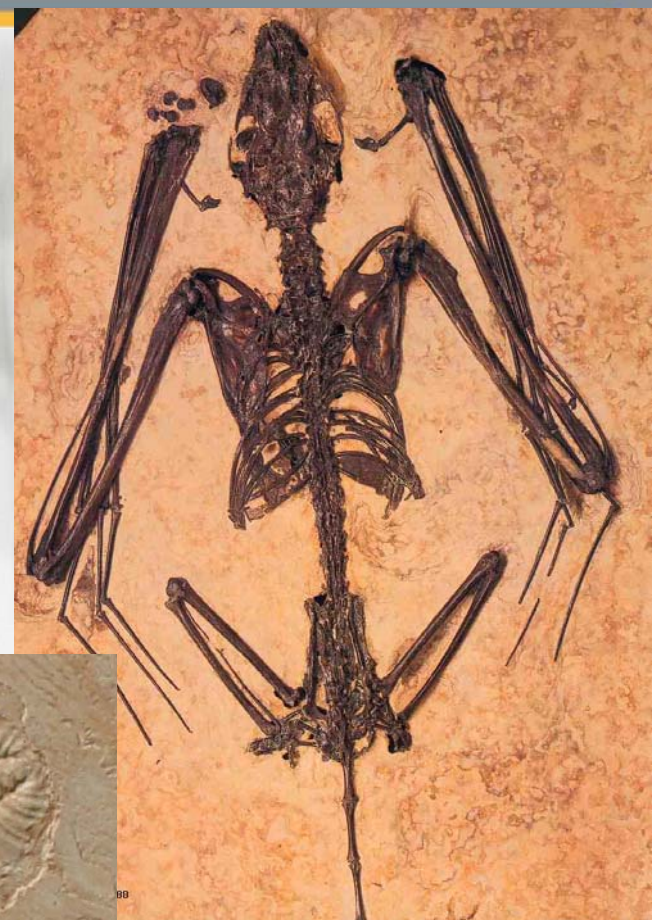
# Prehistoric Static Analysis Tools



**RATS**



**ITS4**



**Flawfinder**

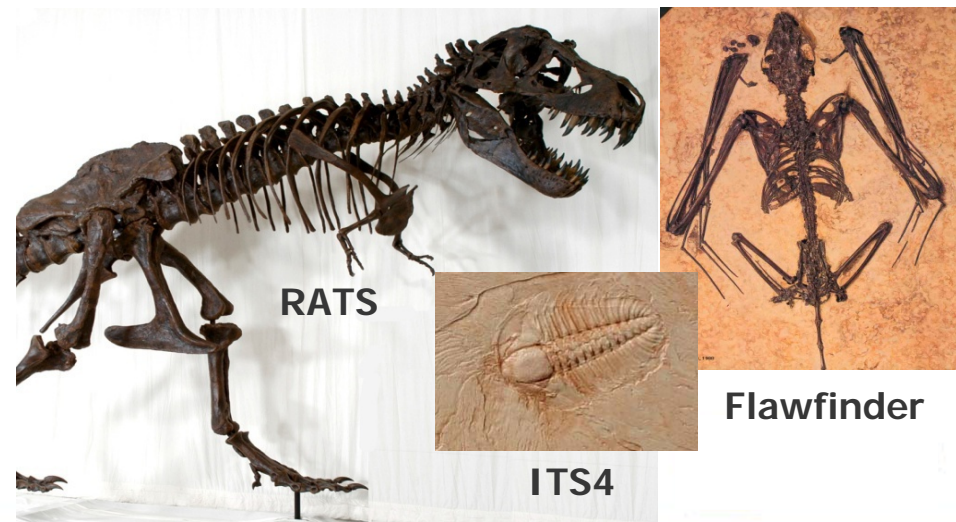
# Prehistoric Static Analysis Tools

## (+) Good

- Help security experts audit code
- Repository for known-bad coding practices

## (-) Bad

- NOT BUG FINDERS
- Not helpful without security expertise





# Misconceptions Prevail

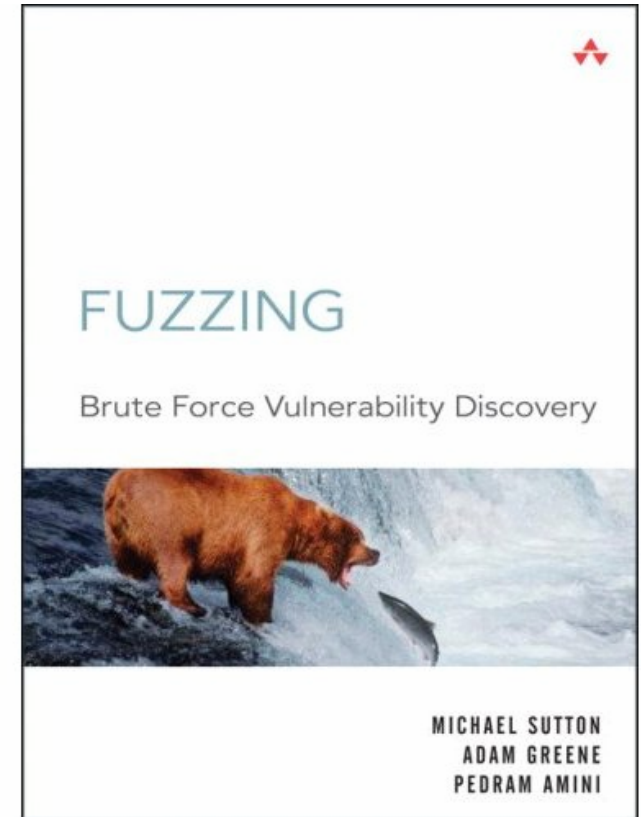
*Fuzzing*, Page 4:

**Low priority**

```
int main(int argc, char** argv)
    char buffer[10];
    strcpy(buf1, "test");
}
```

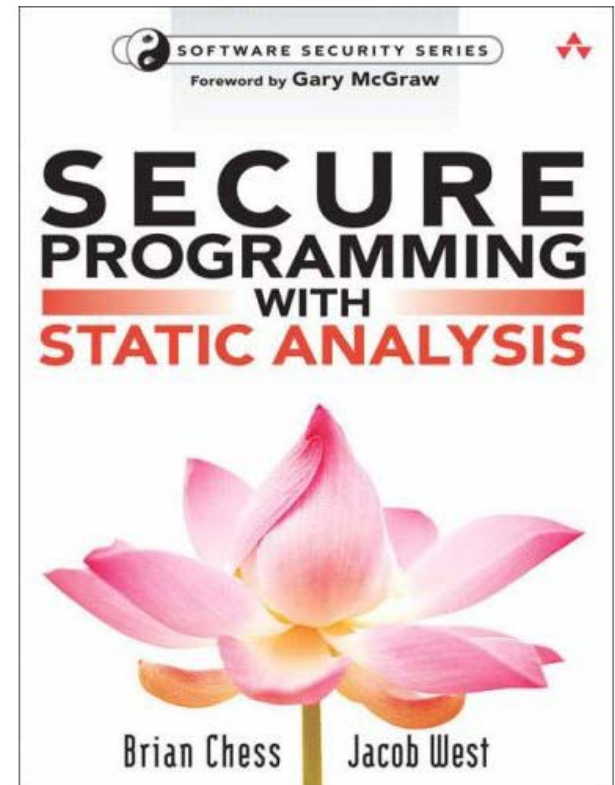
**High priority**

```
int main(int argc, char** argv)
    char buffer[10];
    strcpy(buf1, argv[1]);
}
```



# Static Analysis Is Good For Security

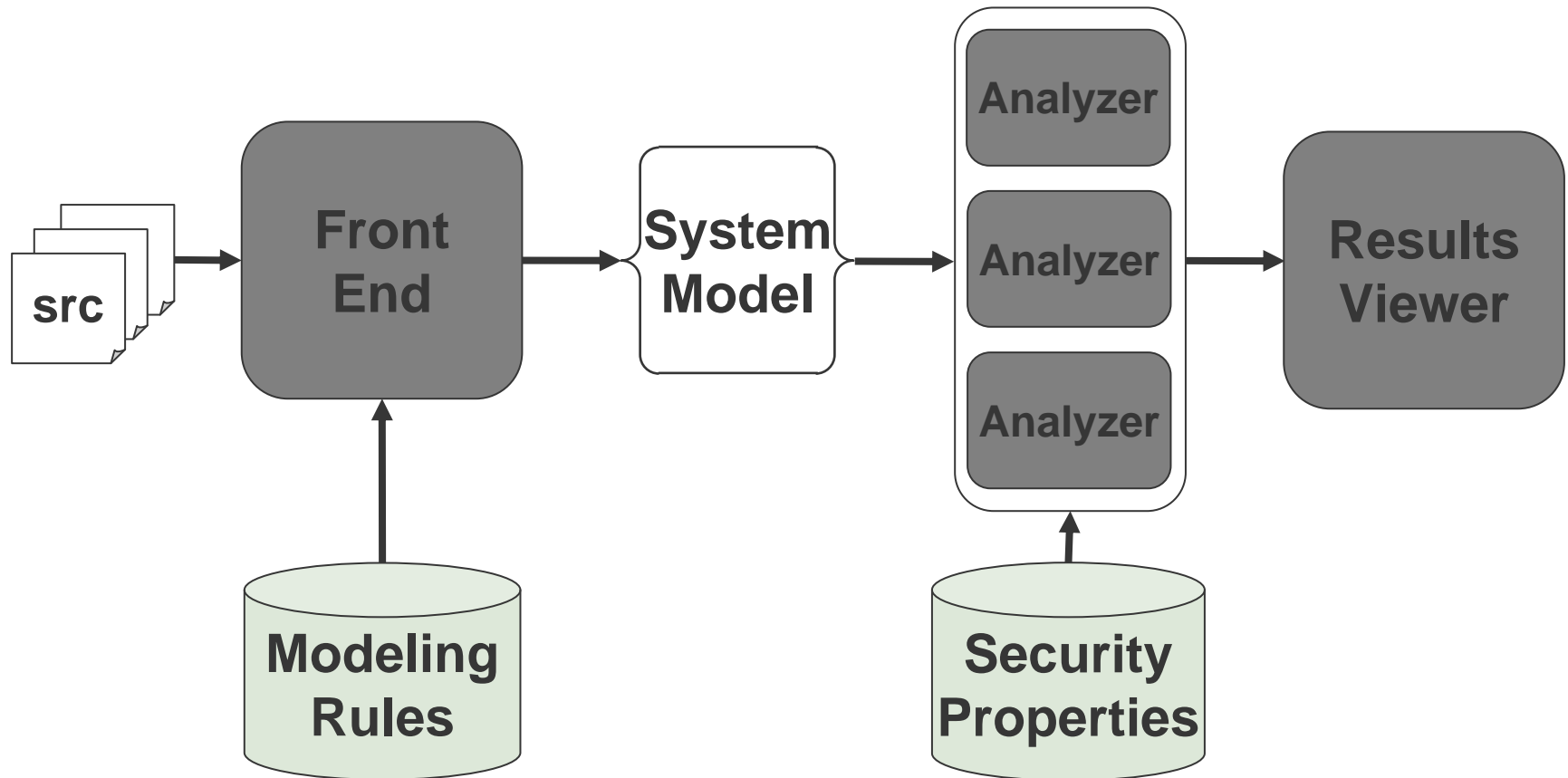
- Fast compared to manual review
- Fast compared to testing
- Complete, consistent coverage
- Brings security knowledge with it
- Makes security review process easier for non-experts
- Useful for all kinds of code, not just Web applications



# Static Analysis: No Silver Bullet

- **Human limitations**
  - Requires access to code
  - User must understand code
- **Tool limitations**
  - Does not understand architecture
  - Does not understand application semantics
  - Does not understand social context

# A Peek Inside a Static Analysis Tool



# Parsing

- **Language support**
  - One language/parser is straightforward
  - Lots of combinations is harder
- **Could analyze compiled code...**
  - Everybody has the binary
  - No need to guess how the compiler works
  - No need for rules
- **...but**
  - Decompilation can be difficult
  - Loss of context hurts
  - Want to report line numbers

# Analysis / Rules: Structural

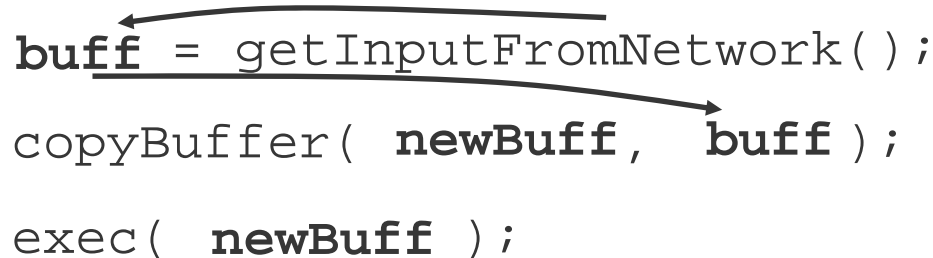
- Identify bugs in the program's structure
- Example: calls to `gets ( )`
- Structural rule:  
`FunctionCall: function is [name == "gets"]`

# Analysis / Rules: Structural

- Identify bugs in the program's structure
- Example: memory leaks caused by `realloc()`  
`buf = realloc(buf, 256);`
- Structural rule:  
FunctionCall c1: (  
    c1.function is [name == "realloc"] and  
    c1 in [AssignmentStatement: rhs is c1 and  
        lhs == c1.arguments[0]  
    ]  
)

# Analysis / Rules: Dataflow Source Rule

- Following interesting values through the program
- Example: Command injection vulnerability



```
buff = getInputFromNetwork();  
copyBuffer( newBuff, buff );  
exec( newBuff );
```

The diagram illustrates the data flow of the example code. A curved arrow originates from the `getInputFromNetwork()` function call in the first line and points to the `buff` variable. Another curved arrow originates from the `buff` variable in the second line and points to the `newBuff` parameter of the `copyBuffer` function.

- Source rule:

Function: `getInputFromNetwork()`

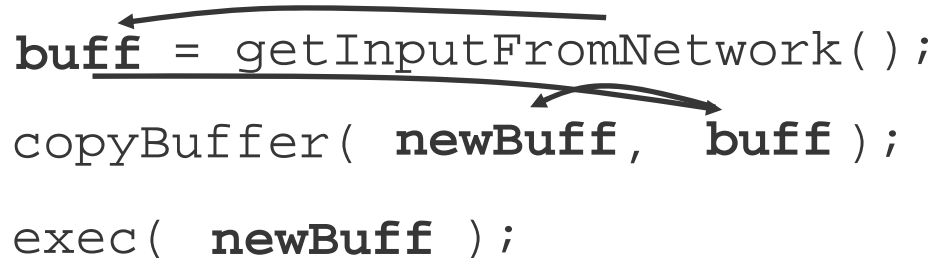
Postcondition: return value is tainted



# Analysis / Rules: Dataflow Pass-Through Rule

- Following interesting values through the program
- Example: Command injection vulnerability

```
buff = getInputFromNetwork();  
copyBuffer( newBuff, buff );  
exec( newBuff );
```



The diagram illustrates the dataflow pass-through rule. It shows three lines of code. The first line, `buff = getInputFromNetwork();`, has a curved arrow pointing from the `buff` variable to the `newBuff` parameter in the second line. The second line, `copyBuffer( newBuff, buff );`, has a straight arrow pointing from the `buff` parameter to the `newBuff` parameter in the third line. The third line, `exec( newBuff );`, shows the `newBuff` variable being passed to the `exec` function. This sequence of arrows demonstrates how the value from `buff` is passed through the `copyBuffer` function to `newBuff`, which is then used in the `exec` function.

- Pass-through rule:

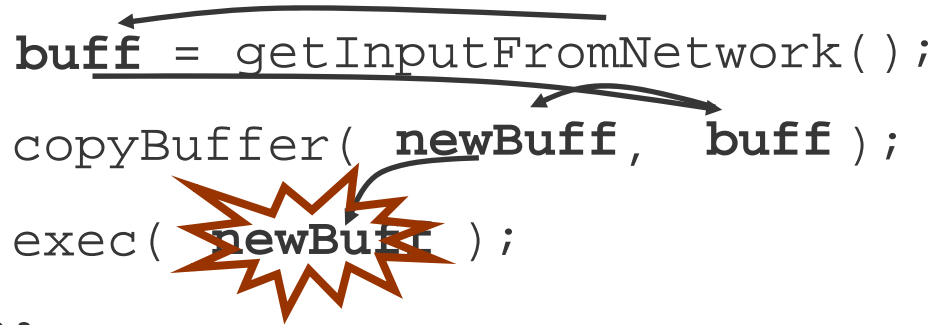
Function: `copyBuffer()`

Postcondition: if the second argument is tainted, then the first argument becomes tainted

# Analysis / Rules: Dataflow Sink Rule

- Following interesting values through the program
- Example: Command injection vulnerability

```
buff = getInputFromNetwork();  
copyBuffer( newBuff, buff );  
exec( newBuff );
```



- Sink rule:

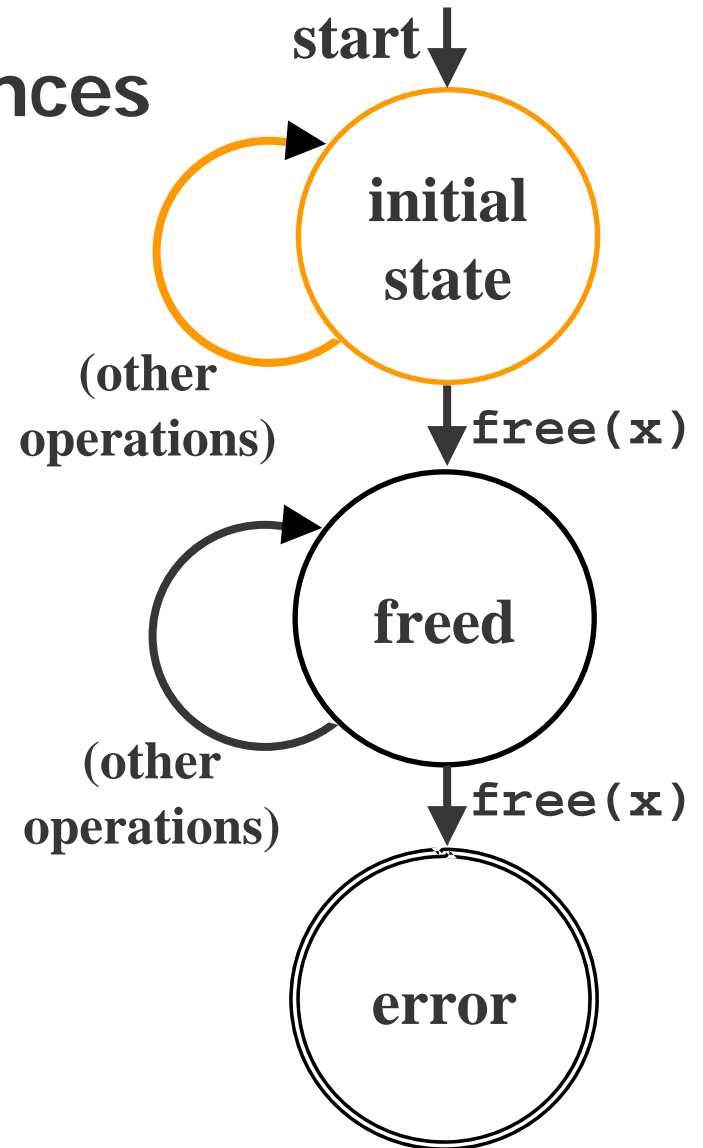
Function: `exec()`

Precondition: the first argument must not be tainted

# Analysis / Rules: Control Flow

- Look for dangerous sequences
- Example: Double-free

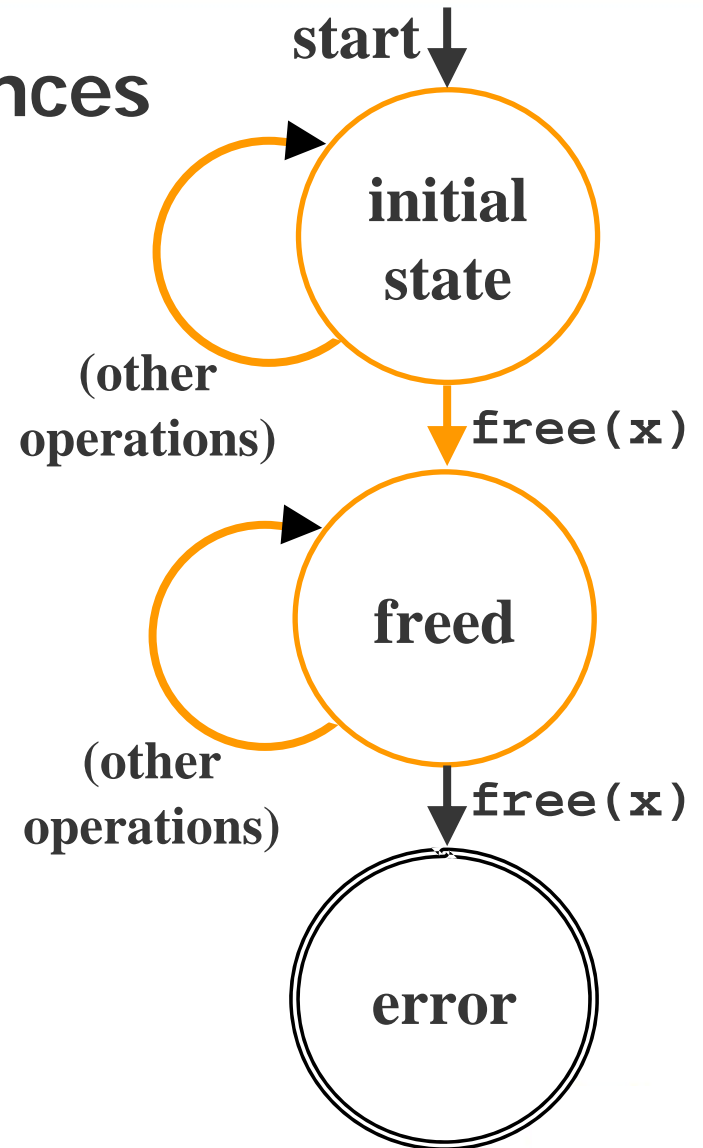
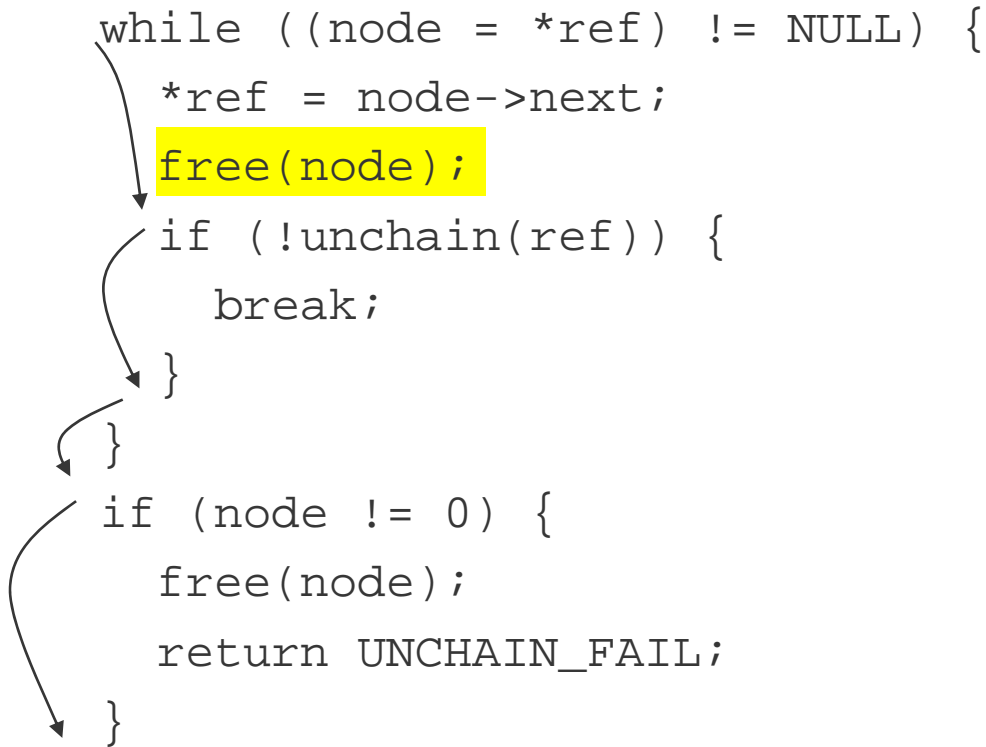
```
while ((node = *ref) != NULL) {  
    *ref = node->next;  
    free(node);  
    if (!unchain(ref)) {  
        break;  
    }  
}  
  
if (node != 0) {  
    free(node);  
    return UNCHAIN_FAIL;  
}
```



# Analysis / Rules: Control Flow

- Look for dangerous sequences
- Example: Double-free

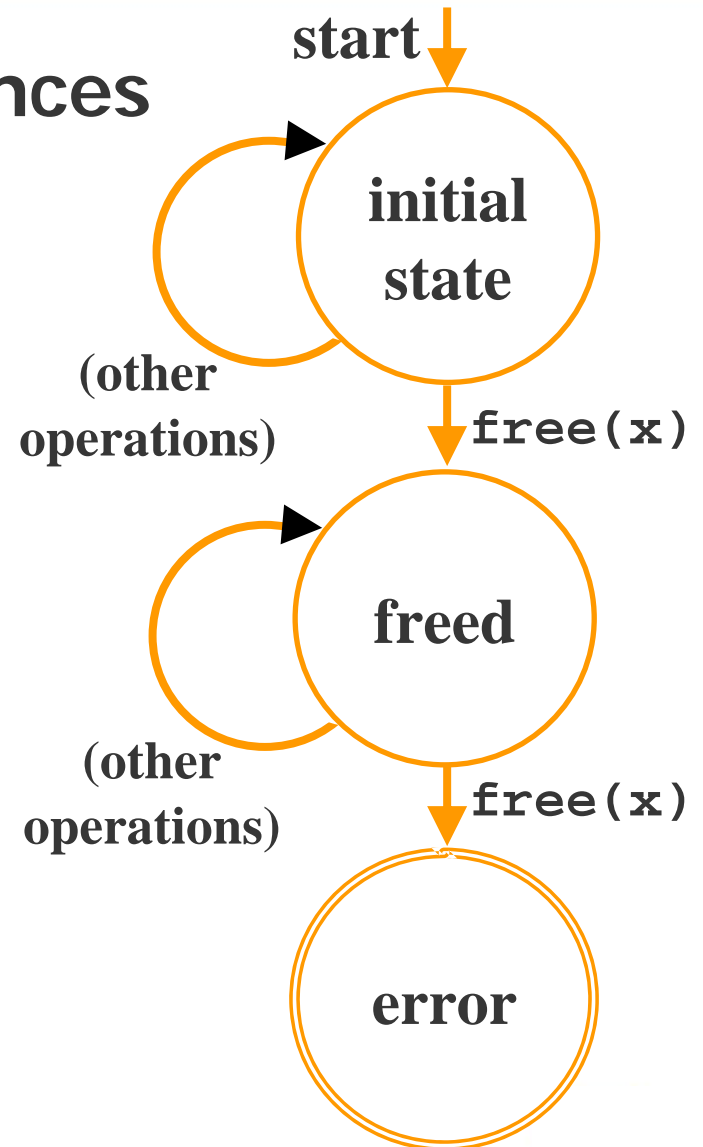
```
while ((node = *ref) != NULL) {  
    *ref = node->next;  
    free(node);  
    if (!unchain(ref)) {  
        break;  
    }  
}  
if (node != 0) {  
    free(node);  
    return UNCHAIN_FAIL;  
}
```



# Analysis / Rules: Control Flow

- Look for dangerous sequences
- Example: Double-free

```
while ((node = *ref) != NULL) {  
    *ref = node->next;  
    free(node);  
    if (!unchain(ref)) {  
        break;  
    }  
}  
  
if (node != 0) {  
    free(node);  
    return UNCHAIN_FAIL;  
}
```



# Only Two Ways to Go Wrong

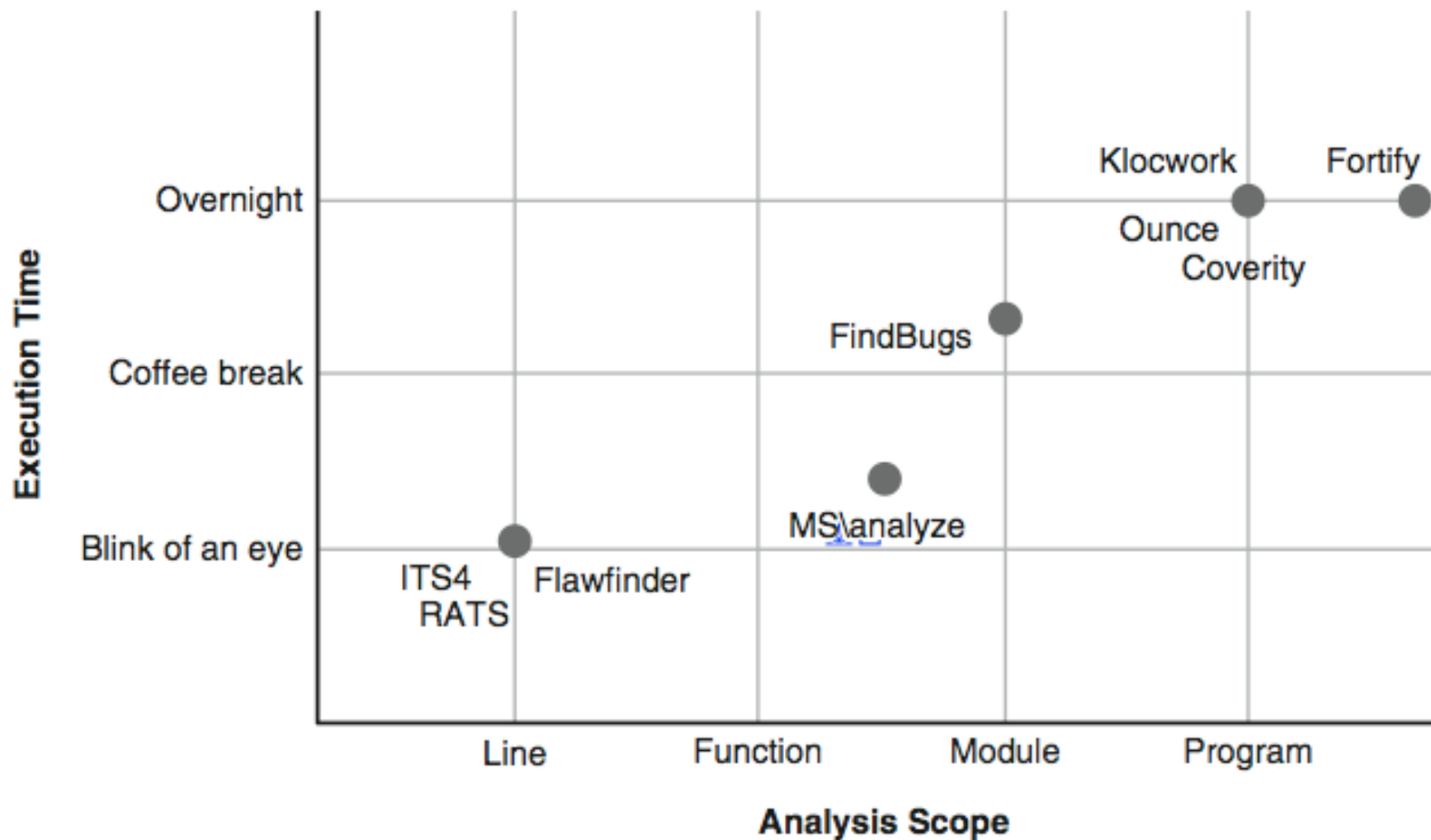
- **False positives**
  - Incomplete/inaccurate model
  - Conservative analysis
  - Missing rules
- **False negatives**
  - Incomplete/inaccurate model
  - “Forgiving” analysis
  - Missing rules



# Untapped Potential: Customization

- **Improve tool understanding of the program**
  - Model the behavior of third-party libraries
  - Describe program semantics
- **Identify program-specific vulnerabilities**
  - Call out targets for manual review
  - Enforce specific coding standards
  - Find vulnerabilities in custom interfaces

# Scope vs. Performance





# Experiment

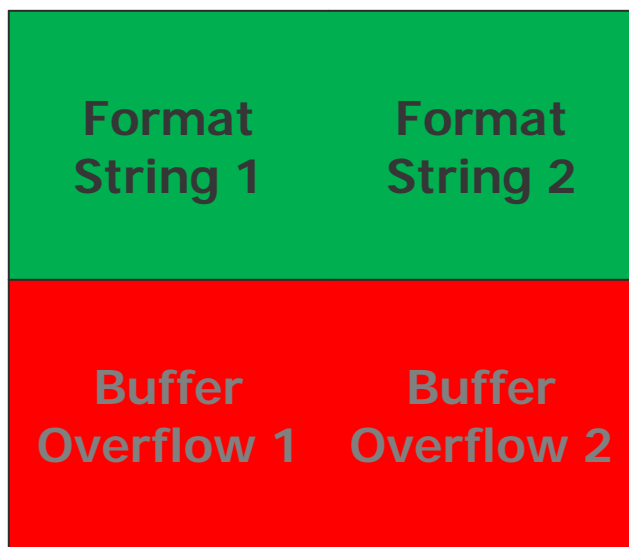
- **Select project: open source mail daemon**
  - qwik-smtpd version .3
  - Contains multiple known vulnerabilities
- **Select tools: fuzzing and static analysis**
  - **Fuzzing: @stake SMTP Fuzz 0.9.16**
    - Customized for SMTP protocol
  - **Static analysis: Fortify**
    - The one we have sitting around
- **Collect data:**
  - Run fuzzing tool on SMTP protocol
  - Run static analysis tool on C source code

# What We Found

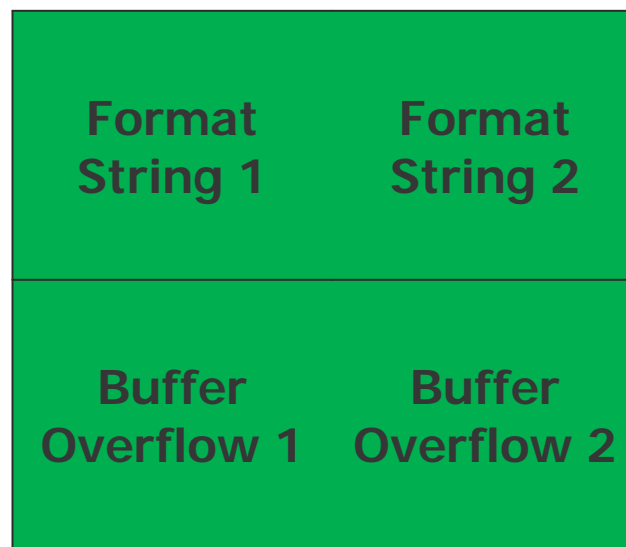
- **Identified four remotely exploitable bugs**
  - Two buffer overflows
  - Two format string vulnerabilities
- **And numerous other locally exploitable vulnerabilities, including:**
  - Buffer overflows
  - Format string vulnerabilities
  - Command injection
  - Memory errors
  - Resource leaks
  - ...

# Results

- Fuzzing found both remotely exploitable format string bugs, but missed both remotely exploitable buffer overflows
- Static analysis: Found all four vulnerabilities



Fuzzing



Static Analysis

# Conclusions

- **Fuzzing...**

- Found exploitable vulnerabilities fast
- Missed critical bugs within its reach
- Missed vulnerabilities from non-SMTP sources
- Would miss bugs behind complex conditions  
(bugs hidden behind multiple header conditions)

# Advantages of Fuzzing over Static Analysis

- **Less involved**
  - Does not require access to or understanding of code
- **Access to context**
  - Does not requires customization to understand program semantics and context
- **The last step**
  - Produces a demonstrable exploit or test case without further human efforts

# Advantages of Static Analysis Over Fuzzing

- **Thoroughness**

- Considers every source of input
- Considers every path through the program

- **Speed**

- Doesn't require running the code
- Customization has almost no impact on performance

- **Visibility**

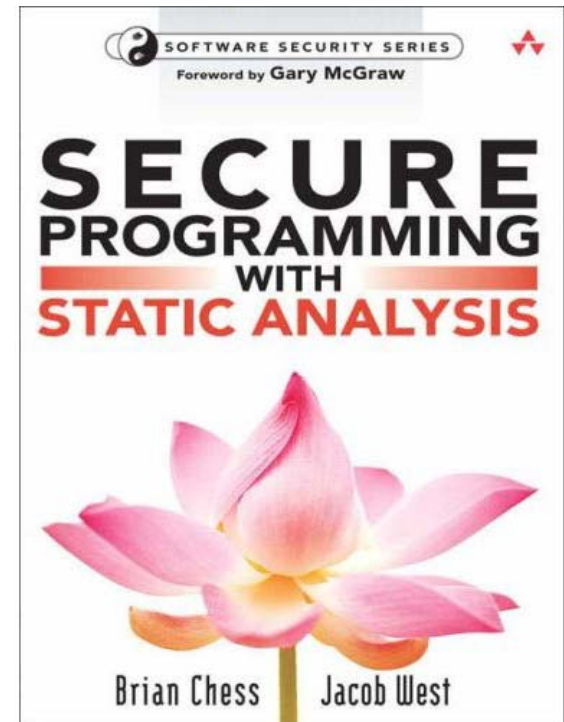
- Identifies vulnerabilities hidden by error handling
- Finds vulnerabilities evidenced through that may be hidden

# Summary

- **Static analysis is spot-on for security**
- **Important attributes**
  - Language support
  - Analysis techniques
  - Rule set
  - Performance
  - Results management
- **Customization**
  - Describe program semantics
  - Model program context

<end>

- PDF of talk will be available here:  
<http://www.fortify.com/presentations>
- Send me email!  
Jacob West <[jacob@fortify.com](mailto:jacob@fortify.com)>



*Secure Programming  
with Static Analysis*