# 1 Computational Geometa T dis2(const point <T > &p, bool is\_segment

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# 1.1 Geometry

```
59
                                                60
1 const double PI=atan2(0.0,-1.0);
2 template<typename T>
                                                61
3 struct point{
                                                62
    T x,y;
    point(){}
    point(const T&x,const T&y):x(x),y(y){}
    point operator+(const point &b)const{
      return point(x+b.x,y+b.y); }
                                                65
    point operator-(const point &b)const{
      return point(x-b.x,y-b.y); }
    point operator*(const T &b)const{
                                                67
      return point(x*b,y*b); }
                                                68
    point operator/(const T &b)const{
                                                69
      return point(x/b,y/b); }
                                                70
    bool operator == (const point &b)const{
                                                71
      return x==b.x&&y==b.y; }
                                                72
    T dot(const point &b)const{
                                                73
      return x*b.x+y*b.y; }
                                                74
    T cross(const point &b)const{
      return x*b.y-y*b.x; }
                                                76
21
    point normal()const{//求法向量
                                                77
22
      return point(-y,x); }
                                                78
    T abs2()const{//向量長度的平方
                                                79
      return dot(*this); }
                                                80
    T rad(const point &b)const{//兩向量的弧度
                                                81
   return fabs(atan2(fabs(cross(b)),dot(b))); }
                                                82
27
    T getA()const{//對x軸的弧度
                                                83
      T A=atan2(y,x);//超過180度會變負的
                                                84
                                                85
      if(A<=-PI/2)A+=PI*2;</pre>
      return A:
31
32
   template<typename T>
   struct line{
    line(){}
                                                88
    point<T> p1,p2;
    T a,b,c;//ax+by+c=0
    line(const point<T>&x,const point<T>&y):p1
         (x),p2(y){}
    void pton(){//轉成一般式
40
      a=p1.y-p2.y;
      b=p2.x-p1.x;
41
      c=-a*p1.x-b*p1.v:
42
43
    T ori(const point<T> &p)const{//點和有向直
                                                97
          線的關係,>0左邊、=0在線上<0右邊
      return (p2-p1).cross(p-p1);
45
                                                99
46
                                               100
    T btw(const point<T> &p)const{//點投影落在 101
          線段 上 <=0
                                               102
48
      return (p1-p).dot(p2-p);
                                               103
49
    bool point_on_segment(const point<T>&p)
50
                                               104
         const{//點是否在線段上
                                               105
      return ori(p) == 0&&btw(p) <= 0;</pre>
                                               106
                                               107
```

```
=0) const { // 點 跟 直 線 / 線 段 的 距 離 平 方
  point<T> v=p2-p1.v1=p-p1:
                                           109
  if(is_segment){
                                           110
    point<T> v2=p-p2;
                                           111
    if(v.dot(v1)<=0)return v1.abs2();</pre>
                                           112
    if(v.dot(v2)>=0)return v2.abs2();
                                           113
                                           114
 T tmp=v.cross(v1);
  return tmp*tmp/v.abs2();
T seg dis2(const line<T> &1)const{//兩線段 118
  return min({dis2(1.p1,1),dis2(1.p2,1),1. 120
       dis2(p1,1),1.dis2(p2,1)});
                                           121
                                           122
point<T> projection(const point<T> &p)
     const { // 點對直線的投影
                                           123
                                           124
  point<T> n=(p2-p1).normal();
                                           125
 return p-n*(p-p1).dot(n)/n.abs2();
                                           126
point<T> mirror(const point<T> &p)const{
                                          127
  //點對直線的鏡射,要先呼叫pton轉成一般式 128
 noint<T> R:
 T d=a*a+b*b:
 R.x=(b*b*p.x-a*a*p.x-2*a*b*p.y-2*a*c)/d; 130
  R.y=(a*a*p.y-b*b*p.y-2*a*b*p.x-2*b*c)/d; 131
  return R:
                                           133
                                          134
bool equal(const line &1)const{//直線相等
 return ori(1.p1)==0&&ori(1.p2)==0;
                                           136
bool parallel(const line &1)const{
 return (p1-p2).cross(1.p1-1.p2)==0;
                                           137
bool cross seg(const line &1)const{
                                           138
 return (p2-p1).cross(l.p1-p1)*(p2-p1).
       cross(1.p2-p1)<=0;//直線是否交線段
                                          139
int line intersect(const line &l)const{// 140
     直線相交情況,-1無限多點、1交於一點、0141
  return parallel(1)?(ori(1.p1)==0?-1:0)
                                           143
                                           144
                                           145
int seg intersect(const line &1)const{
 T c1=ori(l.p1), c2=ori(l.p2);
 T c3=1.ori(p1), c4=1.ori(p2);
                                           147
  if(c1==0&&c2==0){//共線
    bool b1=btw(1.p1)>=0,b2=btw(1.p2)>=0;
    T a3=1.btw(p1),a4=1.btw(p2);
                                           148
                                           149
    if(b1&&b2&&a3==0&&a4>=0) return 2;
                                           150
    if(b1&&b2&&a3>=0&&a4==0) return 3;
                                           151
   if(b1&&b2&&a3>=0&&a4>=0) return 0;
                                           152
    return -1://無限交點
  }else if(c1*c2<=0&&c3*c4<=0)return 1;</pre>
                                           153
 return 0;//不相交
                                           154
                                           155
point<T> line intersection(const line &l)
                                           156
     const{/*直線交點*/
                                           157
  point<T> a=p2-p1,b=l.p2-l.p1,s=l.p1-p1;
                                           158
  //if(a.cross(b)==0)return INF;
  return p1+a*(s.cross(b)/a.cross(b));
```

```
point<T> seg intersection(const line &1)
          const{//線段交點
                                                  162
       int res=seg intersect(1);
       if(res<=0) assert(0);</pre>
                                                  163
       if(res==2) return p1;
                                                 164
       if(res==3) return p2;
                                                  165
       return line intersection(1);
                                                  166
115 };
                                                  167
116 template<typename T>
   struct polygon{
                                                  168
     polygon(){}
     vector<point<T> > p;//逆時針順序
                                                  169
     T area()const{//面積
                                                  170
       T ans=0;
       for(int i=p.size()-1,j=0;j<(int)p.size()</pre>
                                                  17
            ;i=j++)
                                                 172
          ans+=p[i].cross(p[j]);
                                                 173
       return ans/2;
                                                 174
                                                  175
     point<T> center of mass()const{//重心
                                                  176
       T cx=0, cy=0, w=0;
       for(int i=p.size()-1,j=0;j<(int)p.size()</pre>
             ;i=j++){
                                                  177
         T a=p[i].cross(p[j]);
                                                  178
          cx+=(p[i].x+p[j].x)*a;
          cy+=(p[i].y+p[j].y)*a;
                                                  179
                                                  180
                                                  181
       return point<T>(cx/3/w,cy/3/w);
     char ahas(const point<T>& t)const{//點是否
          在簡單多邊形內,是的話回傳1、在邊上回 183
                                                  184

值 - 1 、 否 則 回 值 a

       bool c=0;
                                                 186
       for(int i=0,j=p.size()-1;i<p.size();j=i</pre>
                                                 188
          if(line<T>(p[i],p[j]).point_on_segment
               (t))return -1;
                                                 190
          else if((p[i].y>t.y)!=(p[j].y>t.y)&&
          t.x<(p[j].x-p[i].x)*(t.y-p[i].y)/(p[j]
                                                  191
              ].y-p[i].y)+p[i].x)
                                                  192
            c=!c;
                                                 193
       return c;
                                                  194
     char point_in_convex(const point<T>&x)
                                                 195
                                                  196
       int l=1,r=(int)p.size()-2;
                                                 197
       while(l<=r){//點是否在凸多邊形內,是的話
                                                  198
             回傳1、在邊上回傳-1、否則回傳0
                                                 199
          int mid=(1+r)/2;
                                                 200
         T a1=(p[mid]-p[0]).cross(x-p[0]);
         T a2=(p[mid+1]-p[0]).cross(x-p[0]);
                                                 202
          if(a1>=0&&a2<=0){
                                                  203
           T res=(p[mid+1]-p[mid]).cross(x-p[
                                                 204
                mid]);
                                                  205
           return res>0?1:(res>=0?-1:0);
          }else if(a1<0)r=mid-1;</pre>
                                                  206
          else l=mid+1:
                                                 207
                                                 208
       return 0;
                                                  209
     vector<T> getA()const{//凸包邊對x軸的夾角
       vector<T>res;//一定是遞增的
```

```
for(size t i=0;i<p.size();++i)</pre>
    res.push back((p[(i+1)\%p.size()]-p[i])
         .getA());
  return res:
bool line intersect(const vector<T>&A,
     const line<T> &1)const{//O(LoaN)
  int f1=upper bound(A.begin(), A.end(),(1.
       p1-l.p2).getA())-A.begin();
  int f2=upper bound(A.begin(), A.end(),(1.
       p2-l.p1).getA())-A.begin();
  return 1.cross seg(line<T>(p[f1],p[f2]))
polygon cut(const line<T> &l)const{//△包
     對 直 線 切 割 , 得 到 直 線 L 左 側 的 凸 包
  polygon ans;
  for(int n=p.size(),i=n-1,j=0;j<n;i=j++){</pre>
    if(l.ori(p[i])>=0){
      ans.p.push back(p[i]);
      if(1.ori(p[j])<0)
        ans.p.push_back(1.
             line intersection(line<T>(p[i
             ],p[j])));
    }else if(l.ori(p[j])>0)
      ans.p.push back(1.line intersection(
           line<T>(p[i],p[j])));
  return ans;
static bool graham cmp(const point<T>& a,
     const point<T>& b){//凸包排序函數
  return (a.x<b.x)||(a.x==b.x&&a.y<b.y);</pre>
void graham(vector<point<T> > &s){//凸包
  sort(s.begin(),s.end(),graham cmp);
  p.resize(s.size()+1);
  for(size t i=0;i<s.size();++i){</pre>
    while (m \ge 2\& (p[m-1]-p[m-2]). cross (s[i
         ]-p[m-2])<=0)--m;
    p[m++]=s[i];
  for(int i=s.size()-2,t=m+1;i>=0;--i){
    while (m>=t&&(p[m-1]-p[m-2]).cross(s[i
         ]-p[m-2])<=0)--m;
    p[m++]=s[i];
  if(s.size()>1)--m;
  p.resize(m);
T diam(){//直徑
  int n=p.size(),t=1;
  T ans=0;p.push_back(p[0]);
  for(int i=0;i<n;i++){</pre>
    point<T> now=p[i+1]-p[i];
    while(now.cross(p[t+1]-p[i])>now.cross
         (p[t]-p[i]))t=(t+1)%n;
    ans=\max(ans,(p[i]-p[t]).abs2());
  return p.pop back(),ans;
T min_cover_rectangle(){//最小覆蓋矩形
  int n=p.size(),t=1,r=1,1;
```

```
if(n<3)return 0;//也可以做最小周長矩形
213
        T ans=1e99; p. push back(p[0]);
                                                      267
        for(int i=0;i<n;i++){</pre>
214
                                                      268
215
          point<T> now=p[i+1]-p[i];
                                                      269
          while(now.cross(p[t+1]-p[i])>now.cross 270 };
216
                (p[t]-p[i]))t=(t+1)%n;
217
          while(now.dot(p[r+1]-p[i])>now.dot(p[r 272 | struct triangle{
                ]-p[i]))r=(r+1)%n;
218
          if(!i)l=r;
          while (now.dot(p[l+1]-p[i]) < =now.dot(p[275])
219
                1]-p[i]))1=(1+1)%n;
220
          T d=now.abs2():
          T tmp=now.cross(p[t]-p[i])*(now.dot(p[ 277
221
               r]-p[i])-now.dot(p[l]-p[i]))/d;
222
          ans=min(ans,tmp);
                                                      279
223
                                                      280
224
        return p.pop back(),ans;
                                                      281
225
                                                      282
      T dis2(polygon &pl){//凸包最近距離平方
226
                                                      283
227
        vector<point<T> > &P=p,&Q=pl.p;
                                                      284
228
        int n=P.size(), m=Q.size(), l=0, r=0;
                                                      285
      for(int i=0;i<n;++i)if(P[i].y<P[l].y)l=i;</pre>
229
                                                     286
      for(int i=0;i<m;++i)if(0[i].y<0[r].y)r=i;</pre>
230
231
        P.push back(P[0]),Q.push back(Q[0]);
                                                      287
232
        T ans=1e99;
                                                      288
        for(int i=0;i<n;++i){</pre>
233
234
          while((P[1]-P[1+1]).cross(Q[r+1]-Q[r]) <sub>289</sub>
                <0)r=(r+1)%m;
235
          ans=min(ans,line\langle T \rangle(P[1],P[1+1]).
                                                      291
                seg_dis2(line<T>(Q[r],Q[r+1])));
236
          l=(1+1)%n;
237
                                                      293
238
        return P.pop_back(),Q.pop_back(),ans;
239
                                                      294
240
      static char sign(const point<T>&t){
                                                      295
241
        return (t.y==0?t.x:t.y)<0;</pre>
                                                      296
242
                                                      297
      static bool angle cmp(const line<T>& A.
243
                                                      298 };
           const line<T>& B){
        point < T > a = A.p2 - A.p1, b = B.p2 - B.p1;
244
                                                      300
245
        return sign(a)<sign(b)||(sign(a)==sign(b</pre>
             )&&a.cross(b)>0);
246
247
      int halfplane_intersection(vector<line<T>
           > &s){//半平面交
        sort(s.begin(),s.end(),angle_cmp);//線段
248
              左 側 為 該 線 段 半 平 面
                                                      306
        int L.R.n=s.size():
249
                                                      307
        vector<point<T> > px(n);
250
                                                      308
251
        vector<line<T> > q(n);
                                                      309
252
        q[L=R=0]=s[0];
                                                      310
253
        for(int i=1;i<n;++i){</pre>
                                                      311
254
          while(L<R&&s[i].ori(px[R-1])<=0)--R;</pre>
255
          while(L<R&&s[i].ori(px[L])<=0)++L;</pre>
                                                      313
256
          a[++R]=s[i];
                                                      314
257
          if(q[R].parallel(q[R-1])){
                                                      315
258
                                                      316
259
            if(q[R].ori(s[i].p1)>0)q[R]=s[i];
                                                      317
260
261
          if(L<R)px[R-1]=q[R-1].</pre>
               line_intersection(q[R]);
                                                      319
262
                                                      320
        while(L<R&&q[L].ori(px[R-1])<=0)--R;</pre>
263
264
        p.clear():
                                                      321
265
        if(R-L<=1)return 0;</pre>
```

```
px[R]=q[R].line intersection(q[L]);
                                                 322 };
        for(int i=L;i<=R;++i)p.push back(px[i]); 323 template<typename T>
       return R-L+1;
                                                 326
271 template<typename T>
                                                 327
     point<T> a,b,c;
                                                 328
      triangle(){}
      triangle(const point<T> &a,const point<T>
          &b, const point<T> &c):a(a),b(b),c(c){}<sub>330</sub>
     T area()const{
       T t=(b-a).cross(c-a)/2;
                                                 332
       return t>0?t:-t:
                                                 333
                                                 334
     point<T> barycenter()const{//重心
                                                 335
       return (a+b+c)/3;
                                                 336
                                                 337
     point<T> circumcenter()const{//外心
       static line<T> u.v:
                                                 339
       u.p1=(a+b)/2;
       u.p2=point<T>(u.p1.x-a.y+b.y,u.p1.y+a.x-340)
            b.x);
                                                 341
       v.p1=(a+c)/2;
                                                 342
       v.p2=point<T>(v.p1.x-a.y+c.y,v.p1.y+a.x-
            c.x);
                                                 343
       return u.line_intersection(v);
                                                 344
     point<T> incenter()const{//内心
                                                 345
       T A=sqrt((b-c).abs2()),B=sqrt((a-c).abs2 346
            ()),C=sqrt((a-b).abs2());
       return point<T>(A*a.x+B*b.x+C*c.x,A*a.y+ 348
            B*b.y+C*c.y)/(A+B+C);
                                                 349
     point<T> perpencenter()const{//垂心
       return barycenter()*3-circumcenter()*2;
                                                 350
                                                 351 };
    template<typename T>
    struct point3D{
     T x, y, z;
     point3D(){}
      point3D(const T&x,const T&y,const T&z):x(x
          ),y(y),z(z){}
      point3D operator+(const point3D &b)const{
        return point3D(x+b.x,y+b.y,z+b.z);}
                                                 358
     point3D operator-(const point3D &b)const{
                                                 359
       return point3D(x-b.x,v-b.v,z-b.z);}
                                                 360
      point3D operator*(const T &b)const{
                                                 361
       return point3D(x*b,y*b,z*b);}
      point3D operator/(const T &b)const{
                                                 362
       return point3D(x/b,y/b,z/b);}
                                                 363
      bool operator==(const point3D &b)const{
                                                 364
       return x==b.x&&y==b.y&&z==b.z;}
     T dot(const point3D &b)const{
                                                 365
       return x*b.x+y*b.y+z*b.z;}
     point3D cross(const point3D &b)const{
       return point3D(y*b.z-z*b.y,z*b.x-x*b.z,x 366
             *b.y-y*b.x);}
     T abs2()const{//向量長度的平方
                                                 368
       return dot(*this);}
     T area2(const point3D &b)const{//和b、原點
           圍成面積的平方
       return cross(b).abs2()/4;}
```

```
371
struct line3D{
                                              372
  point3D<T> p1,p2;
                                              373
  line3D(){}
                                              374
  line3D(const point3D<T> &p1,const point3D< 375
       T> &p2):p1(p1),p2(p2){}
  T dis2(const point3D<T> &p,bool is_segment 377
       =0) const { // 點 跟 直 線 / 線 段 的 距 離 平 方
                                              378
    point3D < T > v = p2 - p1, v1 = p - p1;
    if(is segment){
      point3D<T> v2=p-p2;
      if(v.dot(v1)<=0)return v1.abs2();</pre>
                                              380
      if(v.dot(v2)>=0)return v2.abs2():
                                              381
    point3D<T> tmp=v.cross(v1);
    return tmp.abs2()/v.abs2();
                                              382
  pair<point3D<T>,point3D<T> > closest pair( 383
       const line3D<T> &1)const{
    point3D < T > v1 = (p1 - p2), v2 = (1.p1 - 1.p2);
    point3D<T> N=v1.cross(v2),ab(p1-l.p1);
    //if(N.abs2()==0)return NULL;平行或重合
    T tmp=N.dot(ab),ans=tmp*tmp/N.abs2();//
         最近點對距離
    point3D<T> d1=p2-p1,d2=l.p2-l.p1,D=d1.
         cross(d2),G=1.p1-p1;
    T t1=(G.cross(d2)).dot(D)/D.abs2();
                                              389
    T t2=(G.cross(d1)).dot(D)/D.abs2();
                                              390
    return make pair(p1+d1*t1,1.p1+d2*t2);
                                              391
                                              392
  bool same_side(const point3D<T> &a,const
                                              393
       point3D<T> &b)const{
                                              394
    return (p2-p1).cross(a-p1).dot((p2-p1).
                                             395
         cross(b-p1))>0;
                                              396
                                              397
template<typename T>
                                              398
struct plane{
                                              399
  point3D<T> p0,n;//平面上的點和法向量
                                              400
                                              401
  plane(){}
  plane(const point3D<T> &p0,const point3D<T 402
       > &n):p0(p0),n(n){}
                                              403
                                             404
 T dis2(const point3D<T> &p)const{//點到平
                                              405
       面距離的平方
                                              406
    T tmp=(p-p0).dot(n);
                                              407
    return tmp*tmp/n.abs2();
                                              408
  point3D<T> projection(const point3D<T> &p)
                                              411
    return p-n*(p-p0).dot(n)/n.abs2();
                                              412
  point3D<T> line_intersection(const line3D< ^{413}
       T> &1)const{
    T tmp=n.dot(1.p2-1.p1);//等於0表示平行或 415
         重合該平面
    return 1.p1+(1.p2-1.p1)*(n.dot(p0-1.p1)/
                                              418
                                              419
  line3D<T> plane intersection(const plane &
       pl)const{
    point3D<T> e=n.cross(pl.n),v=n.cross(e);
422
    T tmp=pl.n.dot(v);//等於0表示平行或重合 423
```

```
point3D < T > q = p0 + (v*(pl.n.dot(pl.p0-p0))/
    return line3D<T>(q,q+e);
};
template<typename T>
struct triangle3D{
  point3D<T> a,b,c;
  triangle3D(){}
  triangle3D(const point3D<T> &a,const
       point3D<T> &b, const point3D<T> &c):a(a
       ),b(b),c(c){}
  bool point_in(const point3D<T> &p)const{//
       點在該平面上的投影在三角形中
    return line3D<T>(b,c).same_side(p,a)&&
         line3D<T>(a,c).same side(p,b)&&
         line3D<T>(a,b).same_side(p,c);
};
template<typename T>
struct tetrahedron{//四面體
  point3D<T> a,b,c,d;
  tetrahedron(){}
  tetrahedron(const point3D<T> &a,const
      point3D<T> &b, const point3D<T> &c,
       const point3D<T> &d):a(a),b(b),c(c),d(
      d){}
  T volume6()const{//體積的六倍
    return (d-a).dot((b-a).cross(c-a));
  point3D<T> centroid()const{
    return (a+b+c+d)/4;
  bool point in(const point3D<T> &p)const{
    return triangle3D<T>(a,b,c).point in(p)
        &&triangle3D<T>(c,d,a).point in(p);
};
template<typename T>
struct convexhull3D{
  static const int MAXN=1005;
  struct face{
    int a,b,c;
    face(int a,int b,int c):a(a),b(b),c(c){}
  vector<point3D<T>> pt;
  vector<face> ans:
  int fid[MAXN][MAXN];
  void build(){
    int n=pt.size();
    ans.clear();
    memset(fid,0,sizeof(fid));
    ans.emplace back(0.1.2)://注意不能共線
    ans.emplace back(2,1,0);
    int fton = 0:
    for(int i=3, ftop=1; i<n; ++i,++ftop){</pre>
      vector<face> next;
      for(auto &f:ans){
        T d=(pt[i]-pt[f.a]).dot((pt[f.b]-pt[
             f.a]).cross(pt[f.c]-pt[f.a]));
        if(d<=0) next.push back(f);</pre>
        int ff=0:
        if(d>0) ff=ftop;
        else if(d<0) ff=-ftop;</pre>
```

```
fid[f.a][f.b]=fid[f.b][f.c]=fid[f.c
                                                      1 template < typename IT = point < T > * >
                                                      2 T cloest_pair(_IT L, _IT R){
                 ][f.a]=ff;
                                                          if(R-L <= 1) return INF;</pre>
425
          for(auto &f:ans){
                                                          IT mid = L+(R-L)/2;
426
            if(fid[f.a][f.b]>0 && fid[f.a][f.b
                                                          T x = mid -> x;
427
                 ]!=fid[f.b][f.a])
                                                          T d = min(cloest pair(L,mid),cloest pair(
428
              next.emplace back(f.a,f.b,i);
            if(fid[f.b][f.c]>0 && fid[f.b][f.c
                                                          inplace_merge(L, mid, R, ycmp);
429
                 ]!=fid[f.c][f.b])
                                                          static vector<point> b; b.clear();
              next.emplace back(f.b,f.c,i);
                                                          for(auto u=L;u<R;++u){</pre>
430
                                                            if((u->x-x)*(u->x-x)>=d) continue;
431
            if(fid[f.c][f.a]>0 && fid[f.c][f.a
                                                     10
                 1!=fid[f.a][f.c])
                                                            for(auto v=b.rbegin();v!=b.rend();++v){
                                                     11
              next.emplace_back(f.c,f.a,i);
                                                              T dx=u\rightarrow x-v\rightarrow x, dy=u\rightarrow y-v\rightarrow y;
432
                                                     12
                                                              if(dv*dv>=d) break;
433
                                                     13
434
          ans=next:
                                                     14
                                                              d=min(d,dx*dx+dy*dy);
435
                                                     15
                                                            b.push back(*u);
436
                                                     16
     point3D<T> centroid()const{
                                                     17
437
        point3D<T> res(0,0,0);
438
                                                          return d;
                                                     18
        T vol=0:
439
                                                     19
440
        for(auto &f:ans){
                                                     20
                                                        T closest pair(vector<point<T>> &v){
          T tmp=pt[f.a].dot(pt[f.b].cross(pt[f.c 21
                                                          sort(v.begin(),v.end(),xcmp);
441
                                                     22
                                                          return closest pair(v.begin(),v.end());
          res=res+(pt[f.a]+pt[f.b]+pt[f.c])*tmp;
442
443
          vol+=tmp:
444
        return res/(vol*4);
445
446
                                                             Data Structure
447 };
```

2.1 DLX

#### 1.2 SmallestCircle

```
using PT=point<T>; using CPT=const PT;
2 PT circumcenter(CPT &a,CPT &b,CPT &c){
    PT u=b-a, v=c-a;
    T c1=u.abs2()/2,c2=v.abs2()/2;
    T d=u.cross(v);
    return PT(a.x+(v.y*c1-u.y*c2)/d,a.y+(u.x*
         c2-v.x*c1)/d);
   void solve(PT p[],int n,PT &c,T &r2){
    random shuffle(p,p+n);
    c=p[0]; r2=0; // c, r2 = 圓心, 半徑平方
   for(int i=1;i<n;i++)if((p[i]-c).abs2()>r2){
       c=p[i]; r2=0;
   for(int j=0;j<i;j++)if((p[j]-c).abs2()>r2){
        c.x=(p[i].x+p[j].x)/2;
14
        c.y=(p[i].y+p[j].y)/2;
15
         r2=(p[j]-c).abs2();
   for(int k=0;k<j;k++)if((p[k]-c).abs2()>r2){
          c=circumcenter(p[i],p[j],p[k]);
19
          r2=(p[i]-c).abs2();
20
21
22
```

### 最近點對

```
1 const int MAXN=4100, MAXM=1030, MAXND=16390;
2 struct DLX{
                                                54
     int n,m,sz,ansd;//高是n, 寬是m的稀疏矩陣
                                                55
     int S[MAXM],H[MAXN];
     int row[MAXND], col[MAXND]; //每個節點代表的
          列跟行
     int L[MAXND],R[MAXND],U[MAXND],D[MAXND];
     vector<int> ans,anst;
                                                60
     void init(int n,int m){
                                                61
      n=_n,m=_m;
                                                62
       for(int i=0;i<=m;++i){</pre>
                                                 63
        U[i]=D[i]=i,L[i]=i-1,R[i]=i+1;
                                                64
        S[i]=0;
                                                65
                                                66
       R[m]=0,L[0]=m;
       sz=m, ansd=INT MAX; //ansd 存最優解的個數
       for(int i=1;i<=n;++i)H[i]=-1;</pre>
                                                 69
17
                                                70
     void add(int r,int c){
                                                71
       ++S[col[++sz]=c];
19
                                                 72
       row[sz]=r;
20
       D[sz]=D[c],U[D[c]]=sz,U[sz]=c,D[c]=sz;
21
22
       if(H[r]<0)H[r]=L[sz]=R[sz]=sz;
23
       else R[sz]=R[H[r]], L[R[H[r]]]=sz, L[sz]=H
           [r],R[H[r]]=sz;
24
    #define DFOR(i,A,s) for(int i=A[s];i!=s;i=
25
     void remove(int c){//刪除第c行和所有當前覆
          蓋到第c行的列
```

```
void restore(int c){//恢復第c行和所有當前
    覆蓋到第c行的列,remove的逆操作
  DFOR(i,U,c)DFOR(j,L,i){++S[col[j]],U[D[j
      ]]=j,D[U[j]]=j;}
  L[R[c]]=c,R[L[c]]=c;
void remove2(int nd){//刪除nd所在的行當前
    所有點(包括虛擬節點),只保留nd
  DFOR(i,D,nd)L[R[i]]=L[i],R[L[i]]=R[i];
void restore2(int nd){//刪除nd所在的行當前
    所有點,為remove2的逆操作
  DFOR(i,U,nd)L[R[i]]=R[L[i]]=i;
bool vis[MAXM];
int h(){//估價函數 for IDA*
  int res=0;
  memset(vis,0,sizeof(vis));
  DFOR(i,R,0)if(!vis[i]){
   vis[i]=1;
                                          11
   ++res:
   DFOR(j,D,i)DFOR(k,R,j)vis[col[k]]=1;
                                          13
                                          14
  return res;
                                          15
bool dfs(int d){//for精確覆蓋問題
 if(d+h()>=ansd)return 0;//找最佳解用,找
      任意解可以刪掉
                                          19
  if(!R[0]){ansd=d;return 1;}
  int c=R[0];
  DFOR(i,R,0)if(S[i]<S[c])c=i;</pre>
  remove(c);
                                          23
  DFOR(i,D,c){
                                          24
    ans.push back(row[i]);
   DFOR(j,R,i)remove(col[j]);
   if(dfs(d+1))return 1;
    ans.pop back();
   DFOR(j,L,i)restore(col[j]);
  restore(c):
  return 0;
void dfs2(int d){//for最小重複覆蓋問題
  if(d+h()>=ansd)return;
  if(!R[0]){ansd=d;ans=anst;return;}
                                          34
  int c=R[0];
  DFOR(i,R,0)if(S[i]<S[c])c=i;</pre>
  DFOR(i,D,c){
   anst.push back(row[i]);
    remove2(i);
   DFOR(j,R,i)remove2(j),--S[col[j]];
    dfs2(d+1):
    anst.pop_back();
    DFOR(j,L,i)restore2(j),++S[col[j]];
                                          41
    restore2(i);
                                          42
```

DFOR(i,D,c)DFOR(j,R,i){U[D[j]]=U[j],D[U[

j]]=D[j],--S[col[j]];}

32

33

36

37

38

39

40

43

44

49

50

```
L[R[c]]=L[c],R[L[c]]=R[c];//這裡刪除第c 82
                                     bool exact cover(){//解精確覆蓋問題
                                       return ans.clear(), dfs(0);
    行,若有些行不需要處理可以在開始時呼 83
                                     void min_cover(){//解最小重複覆蓋問題
                                       anst.clear();//暫存用,答案還是存在ans裡
                                       dfs2(0);
                                     #undef DFOR
                                 90 };
```

#### 2.2 Dynamic KD tree

```
1| template<typename T,size_t kd>//有kd個維度
2 struct kd tree{
    struct point{
      T d[kd];
      T dist(const point &x)const{
        T ret=0;
        for(size_t i=0;i<kd;++i)ret+=abs(d[i]-</pre>
             x.d[i]);
        return ret;
      bool operator == (const point &p){
        for(size t i=0;i<kd;++i)</pre>
          if(d[i]!=p.d[i])return 0;
        return 1;
      bool operator<(const point &b)const{</pre>
        return d[0] < b.d[0];</pre>
   };
  private:
    struct node{
      node *1,*r;
      point pid;
      int s:
      node(const point &p):1(0),r(0),pid(p),s
      ~node(){delete l,delete r;}
      void up()\{s=(1?1->s:0)+1+(r?r->s:0);\}
    }*root:
    const double alpha.loga:
    const T INF;//記得要給INF,表示極大值
    int maxn;
    struct __cmp{
      int sort id;
      bool operator()(const node*x,const node*
           y)const{
        return operator()(x->pid,y->pid);
      bool operator()(const point &x,const
           point &y)const{
        if(x.d[sort id]!=y.d[sort id])
          return x.d[sort id]<y.d[sort id];</pre>
        for(size_t i=0;i<kd;++i)</pre>
          if(x.d[i]!=y.d[i])return x.d[i]<y.d[</pre>
               i];
        return 0;
    int size(node *o){return o?o->s:0;}
    vector<node*> A;
```

```
node* build(int k,int l,int r){
                                                         if(erase(cmp(x,u->pid)?u->l:u->r,(k+1)% 165
                                                                                                           return d:
       if(1>r) return 0;
                                                              kd,x)
                                                                                                   166
       if(k==kd) k=0;
48
                                                            return --u->s, 1;
                                                                                                         void rebuild(){
                                                                                                                                                              const point &R){
                                                 107
                                                                                                   167
       int mid=(1+r)/2;
                                                                                                           if(root)rebuild(root,0);
                                                 108
                                                         return 0;
                                                                                                   168
                                                                                                                                                           for(int i=0;i<kd;++i){</pre>
       cmp.sort id = k;
                                                                                                           maxn=root->s;
                                                 109
                                                                                                   169
       nth element(A.begin()+l,A.begin()+mid,A. 110
                                                       T heuristic(const T h[])const{
                                                                                                   170
                                                                                                                                                                  return 0:
            begin()+r+1.cmp):
                                                                                                   171
                                                                                                         T nearest(const point &x.int k){
                                                                                                                                                      37
       node *ret=A[mid];
                                                         for(size t i=0;i<kd;++i)ret+=h[i];</pre>
52
                                                 112
                                                                                                   172
                                                                                                                                                      38
                                                                                                                                                           return 1:
       ret->1 = build(k+1,1,mid-1);
                                                                                                           T mndist=INF,h[kd]={};
53
                                                 113
                                                         return ret;
                                                                                                   173
                                                                                                                                                      39
54
       ret->r = build(k+1,mid+1,r);
                                                                                                   174
                                                                                                           nearest(root,0,x,h,mndist);
                                                 114
55
       ret->up();
                                                 115
                                                       int qM;
                                                                                                   175
                                                                                                           mndist=p0.top().first;
                                                                                                                                                              const point &R){
56
       return ret:
                                                       priority queue<pair<T.point>> p0:
                                                                                                           pQ = priority queue<pair<T,point>>();
                                                 116
                                                                                                                                                           for(int i=0;i<kd;++i){</pre>
                                                       void nearest(node *u,int k,const point &x, 177
57
                                                                                                           return mndist;//回傳離x第k近的點的距離
58
     bool isbad(node*o){
                                                            T *h.T &mndist){
                                                                                                   178
                                                                                                                                                                  1)return 0:
                                                         if(u==0||heuristic(h)>=mndist)return;
59
       return size(o->1)>alpha*o->s||size(o->r) 118
                                                                                                         const vector<point> &range(const point&mi,
                                                                                                   179
            >alpha*o->s;
                                                 119
                                                         T dist=u->pid.dist(x),old=h[k];
                                                                                                              const point&ma){
                                                                                                                                                           return 1;
                                                         /*mndist=std::min(mndist.dist):*/
60
                                                 120
                                                                                                           in range.clear();
                                                                                                   180
                                                                                                                                                      45
     void flatten(node *u, typename vector<node</pre>
                                                         if(dist<mndist){</pre>
61
                                                                                                           range(root,0,mi,ma);
                                                                                                                                                      46 // 單點修改,以單點改值為例
          *>::iterator &it){
                                                           pQ.push(std::make_pair(dist,u->pid));
                                                 122
                                                                                                   182
                                                                                                           return in range;//回傳介於mi到ma之間的點
       if(!u)return:
                                                            if((int)pQ.size()==qM+1)
62
                                                 123
                                                                                                                                                              int k=0){
63
       flatten(u->1,it);
                                                 124
                                                              mndist=p0.top().first,p0.pop();
                                                                                                   183
                                                                                                                                                           if(!u)return;
       *it=u:
64
                                                 125
                                                                                                         int size(){return root?root->s:0;}
                                                                                                   184
                                                                                                                                                           u->down();
65
       flatten(u->r,++it):
                                                 126
                                                         if(x.d[k]<u->pid.d[k]){
                                                                                                   185 };
                                                                                                                                                           if(u->pid==x){
66
                                                 127
                                                           nearest(u \rightarrow 1, (k+1)%kd, x,h,mndist);
                                                                                                                                                             u->data=data;
     void rebuild(node*&u,int k){
                                                           h[k] = abs(x.d[k]-u-pid.d[k]);
67
                                                 128
                                                                                                                                                             u->up2():
68
       if((int)A.size()<u->s)A.resize(u->s);
                                                 129
                                                           nearest(u->r,(k+1)%kd,x,h,mndist);
                                                                                                                                                             return:
69
       auto it=A.begin();
                                                 130
                                                                                                       2.3 kd tree replace segment 54
70
       flatten(u,it);
                                                 131
                                                           nearest(u->r.(k+1)%kd.x.h.mndist);
                                                                                                                                                           cmp.sort id=k;
       u=build(k,0,u->s-1);
                                                           h[k] = abs(x.d[k]-u->pid.d[k]);
71
                                                 132
72
                                                 133
                                                           nearest(u->1,(k+1)%kd,x,h,mndist);
                                                                                                     1 struct node { //kd 樹代 替高維線段樹
                                                                                                                                                                +1)%kd);
73
     bool insert(node*&u,int k,const point &x,
                                                 134
                                                                                                         node *1,*r;
                                                                                                                                                      57
                                                                                                                                                           u->up2():
          int dep){
                                                         h[k]=old;
                                                 135
                                                                                                         point pid, mi, ma;
                                                                                                                                                      58
       if(!u) return u=new node(x), dep<=0;</pre>
74
                                                 136
                                                                                                         int s, data;
                                                                                                                                                      59 //區間修改
75
       ++u->s:
                                                 137
                                                       vector<point>in range;
                                                                                                         node(const point &p,int d):1(0),r(0),pid(p
76
       cmp.sort id=k;
                                                       void range(node *u,int k,const point&mi,
                                                 138
                                                                                                              ),mi(p),ma(p),s(1),data(d),dmin(d),
                                                                                                                                                              point &R, int data){
       if(insert(cmp(x,u->pid)?u->1:u->r,(k+1)%
                                                            const point&ma){
                                                                                                              dmax(d){}
                                                                                                                                                           if(!o)return;
            kd,x,dep-1)){
                                                         if(!u)return;
                                                 139
                                                                                                         void up(){
                                                                                                                                                      62
                                                                                                                                                           o->down();
         if(!isbad(u))return 1;
                                                         bool is=1;
                                                 140
                                                                                                           mi=ma=pid;
                                                                                                                                                           if(range in range(o,L,R)){
79
         rebuild(u,k);
                                                 141
                                                         for(int i=0:i<kd:++i)</pre>
                                                                                                           s=1:
                                                                                                                                                             //區間懶惰標記修改
                                                 142
                                                           if(u->pid.d[i]<mi.d[i]||ma.d[i]<u->pid
80
                                                                                                           if(1){
                                                                                                                                                             o->down();
                                                                .d[i])
                                                                                                                                                      65
81
       return 0;
                                                                                                    10
                                                                                                             for(int i=0;i<kd;++i){</pre>
                                                              { is=0; break; }
                                                                                                                                                      66
                                                                                                                                                             return;
82
                                                 143
                                                                                                               mi.d[i]=min(mi.d[i],1->mi.d[i]);
                                                                                                    11
                                                                                                                                                      67
                                                         if(is) in range.push back(u->pid);
     node *findmin(node*o,int k){
83
                                                 144
                                                                                                               ma.d[i]=max(ma.d[i],1->ma.d[i]);
                                                                                                                                                           if(point_in_range(o,L,R)){
                                                         if(mi.d[k] <= u - > pid.d[k]) range(u - > 1,(k+1))
84
       if(!o)return 0;
                                                 145
       if(cmp.sort_id==k)return o->l?findmin(o
                                                              %kd,mi,ma);
                                                                                                             s+=1->s:
                                                                                                                                                                  一定在區間中
            ->1,(k+1)%kd):o;
                                                 146
                                                         if(ma.d[k])=u-pid.d[k])range(u->r,(k+1)
                                                                                                    15
       node *l=findmin(o->l,(k+1)%kd);
                                                              %kd,mi,ma);
                                                                                                                                                             //單點懶惰標記修改
                                                                                                           if(r){
                                                                                                                                                      70
                                                                                                    16
       node *r=findmin(o->r,(k+1)%kd);
                                                 147
                                                                                                                                                      71
                                                                                                    17
                                                                                                             for(int i=0:i<kd:++i){</pre>
       if(1&&!r)return cmp(1,0)?1:0;
                                                     public:
                                                 148
                                                                                                               mi.d[i]=min(mi.d[i],r->mi.d[i]);
                                                                                                                                                      72
                                                                                                    18
       if(!1&&r)return cmp(r,o)?r:o;
                                                       kd tree(const T &INF, double a=0.75):
                                                                                                               ma.d[i]=max(ma.d[i],r->ma.d[i]);
                                                                                                                                                                ->1,L,R,data);
                                                       root(0),alpha(a),loga(log2(1.0/a)),INF(INF
       if(!1&&!r)return o;
90
                                                                                                    20
       if(cmp(1,r))return cmp(1,o)?1:o;
                                                            ),maxn(1){}
                                                                                                                                                                ->r,L,R,data);
                                                                                                    21
                                                                                                             s+=r->s;
       return cmp(r,o)?r:o;
                                                       ~kd tree(){delete root;}
92
                                                                                                                                                           o->up2();
                                                                                                                                                      74
                                                                                                    22
93
                                                       void clear(){delete root,root=0,maxn=1;}
                                                                                                                                                      75
                                                                                                    23
     bool erase(node *&u,int k,const point &x){ 153
                                                       void build(int n,const point *p){
                                                                                                                                                         //區間查詢,以總和為例
                                                                                                         void up2(){/*其他懶惰標記向上更新*/}
95
       if(!u)return 0;
                                                         delete root, A. resize(maxn=n);
                                                                                                         void down(){/*其他懶惰標記下推*/}
                                                                                                    25
       if(u->pid==x){
                                                 155
                                                         for(int i=0;i<n;++i)A[i]=new node(p[i]);</pre>
                                                                                                                                                               &R){
                                                                                                    26 }*root;
         if(u->r):
                                                         root=build(0,0,n-1);
97
                                                 156
                                                                                                                                                           if(!o)return 0:
                                                                                                    27 //檢查區間包含用的函數
         else if(u->1) u->r=u->1, u->l=0;
                                                 157
                                                                                                                                                           o->down();
                                                                                                       bool range include(node *o,const point &L,
99
         else return delete(u),u=0, 1;
                                                 158
                                                       void insert(const point &x){
                                                                                                            const point &R){
                                                         insert(root,0,x, lg(size(root))/loga);
         --u->s:
                                                 159
                                                                                                                                                      81
                                                                                                                                                           int ans=0:
                                                                                                         for(int i=0:i<kd:++i){</pre>
101
         cmp.sort id=k;
                                                 160
                                                         if(root->s>maxn)maxn=root->s;
                                                                                                           if(L.d[i]>o->ma.d[i]||R.d[i]<o->mi.d[i])
102
         u->pid=findmin(u->r,(k+1)%kd)->pid;
                                                 161
                                                                                                    30
         return erase(u->r,(k+1)%kd,u->pid);
                                                       bool erase(const point &p){
                                                                                                                return 0:
103
                                                 162
                                                                                                                                                                query(o->1,L,R);
                                                                                                         }//(L,R)區間有和o的區間有交集就回傳true
                                                         bool d=erase(root,0,p);
104
                                                 163
       cmp.sort id=k;
                                                         if(root&&root->s<alpha*maxn)rebuild();</pre>
                                                                                                         return 1;
```

```
bool range in range(node *o, const point &L,
   if(L.d[i]>o->mi.d[i]||o->ma.d[i]>R.d[i])
 }//(L,R)區間完全包含o的區間就回傳true
bool point_in_range(node *o,const point &L,
   if(L.d[i]>o->pid.d[i]||R.d[i]<o->pid.d[i
 }//(L,R)區間完全包含o->pid這個點就回傳true
void update(node *u,const point &x,int data,
 update(cmp(x,u->pid)?u->l:u->r,x,data,(k
void update(node *o,const point &L,const
   //這個點在(L.R)區間,但是他的左右子樹不
 if(o->1&&range_include(o->1,L,R))update(o
  if(o->r&&range include(o->r,L,R))update(o
int query(node *o,const point &L,const point
  if(range in range(o,L,R))return o->sum;
 if(point_in_range(o,L,R))ans+=o->data;
 if(o->1&&range include(o->1,L,R))ans+=
```

```
if(o->r&&range include(o->r,L,R))ans+=
     query(o->r,L,R);
return ans;
```

#### 2.4 reference point

```
1 template<typename T>
2 struct _RefC{
    T data:
    int ref;
    _RefC(const T&d=0):data(d),ref(0){}
7 template<typename T>
   struct _rp{
     RefC<T> *p;
    T *operator->(){return &p->data;}
    T & operator*() { return p->data; }
    operator RefC<T>*(){return p;}
    _rp &operator=(const _rp &t){
      if(p&&!--p->ref)delete p;
      p=t.p,p&&++p->ref;
      return *this;
    _rp(_RefC<T> *t=0):p(t){p&&++p->ref;}
    _rp(const _rp &t):p(t.p){p&&++p->ref;}
    ~ rp(){if(p&&!--p->ref)delete p;}
22
   template<typename T>
  inline _rp<T> new_rp(const T&nd){
    return _rp<T>(new _RefC<T>(nd));
```

# 2.5 skew heap

```
1 node *merge(node *a, node *b){
   if(!a||!b) return a?a:b;
   if(b->data<a->data) swap(a,b);
   swap(a->1,a->r);
   a->l=merge(b,a->l);
   return a;
```

# 2.6 undo disjoint set

```
1 struct DisjointSet {
   // save() is like recursive
   // undo() is like return
   int n, fa[MXN], sz[MXN];
   vector<pair<int*,int>> h;
   vector<int> sp;
   void init(int tn) {
      for (int i=0; i<n; i++) sz[fa[i]=i]=1;</pre>
      sp.clear(); h.clear();
```

```
void assign(int *k, int v) {
       h.PB(\{k, *k\});
13
14
15
16
     void save() { sp.PB(SZ(h)); }
     void undo() {
18
       assert(!sp.empty());
       int last=sp.back(); sp.pop_back();
19
20
       while (SZ(h)!=last) {
21
         auto x=h.back(); h.pop_back();
22
         *x.F=x.S;
23
24
25
     int f(int x) {
26
       while (fa[x]!=x) x=fa[x];
27
       return x;
28
29
     void uni(int x, int y) {
30
       x=f(x); y=f(y);
       if (x==y) return ;
31
       if (sz[x]<sz[y]) swap(x, y);</pre>
32
       assign(&sz[x], sz[x]+sz[y]);
34
       assign(&fa[y], x);
35
36 }djs;
```

#### 2.7 整體二分

```
1 | void totBS(int L, int R, vector<Item> M){
   if(Q.empty()) return; //維護全域B陣列
   if(L==R) 整個M的答案=r, return;
   int mid = (L+R)/2:
   vector<Item> mL, mR;
   do modify B with divide(mid, M);
   //讓B陣列在遞迴的時候只會保留[L~mid]的資訊
   undo modify B(mid,M);
   totBS(L,mid,mL);
   totBS(mid+1,R,mR);
```

### Flow

#### 3.1 dinic

```
1 | template < typename T>
2 struct DINIC{
    static const int MAXN=105;
    static const T INF=INT MAX;
    int n, LV[MAXN], cur[MAXN];
    struct edge{
      int v,pre;
      edge(int v,int pre,T cap):v(v),pre(pre),
           cap(cap),r(cap){}
    int g[MAXN];
    vector<edge> e;
```

```
q.push(e[i].v);
        if(e[i].v==t)return 1;
   }
  return 0;
                                              11
                                              12
T dfs(int u,int t,T CF=INF){
                                              13
  if(u==t)return CF;
  T df:
                                              15
  for(int &i=cur[u];~i;i=e[i].pre){
                                              16
    if(LV[e[i].v]==LV[u]+1&&e[i].r){
                                              17
      if(df=dfs(e[i].v,t,min(CF,e[i].r))){ 18
        e[i].r-=df;
        e[i^1].r+=df;
        return df;
                                              21
   }
                                              22
  return LV[u]=0;
                                              23
T dinic(int s,int t,bool clean=true){
                                              25
  if(clean)for(size_t i=0;i<e.size();++i)</pre>
                                              26
    e[i].r=e[i].cap;
                                              27
  T ans=0, f=0;
                                              28
  while(bfs(s,t))while(f=dfs(s,t))ans+=f;
  return ans;
                                              30
```

# Gomory Hu

void init(int n){

directed=false){

g[u]=e.size()-1;

g[v]=e.size()-1;

int bfs(int s,int t){

while(q.size()){

queue<int> q;

q.push(s);

LV[s]=1;

e.clear();

15

16

17

18

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55

56

57

58

60

61

62 };

memset(g,-1,sizeof(int)\*((n=n)+1));

e.push\_back(edge(u,g[v],directed?0:cap))

void add\_edge(int u,int v,T cap,bool

e.push\_back(edge(v,g[u],cap));

memset(LV,0,sizeof(int)\*(n+1));

int u=q.front();q.pop();

memcpy(cur,g,sizeof(int)\*(n+1));

for(int i=g[u];~i;i=e[i].pre){

if(!LV[e[i].v]&&e[i].r){

LV[e[i].v]=LV[u]+1;

```
1 / / 最小割樹+求任兩點間最小割
2 //0-base, root=0
3 LL e[MAXN][MAXN]; //任兩點間最小割
4 int p[MAXN]; //parent
5 ISAP D; // original graph
6 void gomory_hu(){
7 fill(p, p+n, 0);
```

fill(e[0], e[n], INF);

**int** t = p[s];

ISAP F = D;

10

11

12

14

16

31

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33

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38

41

for( int s = 1; s < n; ++s ) {

LL tmp = F.min cut(s, t);

for( int i = 1; i < s; ++i )

for( int i = s+1; i <= n; ++i )</pre>

e[s][i] = e[i][s] = min(tmp, e[t][i]);

if( p[i] == t && F.vis[i] ) p[i] = s;

```
3.3 ISAP with cut
1 | template < typename T>
 struct ISAP{
   static const int MAXN=105;
   static const T INF=INT MAX;
   int n://點數
   int d[MAXN],gap[MAXN],cur[MAXN];
   struct edge{
     int v,pre;
     T cap,r;
     edge(int v,int pre,T cap):v(v),pre(pre),
          cap(cap),r(cap){}
   int g[MAXN];
   vector<edge> e;
   void init(int _n){
     memset(g,-1, sizeof(int)*((n=_n)+1));
     e.clear();
   void add edge(int u,int v,T cap,bool
        directed=false){
     e.push_back(edge(v,g[u],cap));
     g[u]=e.size()-1;
     e.push_back(edge(u,g[v],directed?0:cap))
     g[v]=e.size()-1;
   T dfs(int u, int s, int t, T CF=INF){
     if(u==t)return CF;
     T tf=CF,df;
     for(int &i=cur[u];~i;i=e[i].pre){
       if(e[i].r&&d[u]==d[e[i].v]+1){
         df=dfs(e[i].v,s,t,min(tf,e[i].r));
         e[i].r-=df;
         e[i^1].r+=df;
         if(!(tf-=df)||d[s]==n)return CF-tf;
     for(int i=cur[u]=g[u];~i;i=e[i].pre){
       if(e[i].r&&d[e[i].v]<mh)mh=d[e[i].v];</pre>
     if(!--gap[d[u]])d[s]=n;
     else ++gap[d[u]=++mh];
     return CF-tf;
   T isap(int s,int t,bool clean=true){
     memset(d,0,sizeof(int)*(n+1));
```

memset(gap,0,sizeof(int)\*(n+1));

memcpy(cur,g,sizeof(int)\*(n+1));

```
if(clean) for(size t i=0;i<e.size();++i)</pre>
48
         e[i].r=e[i].cap;
                                                   35
       T MF=0;
49
       for(gap[0]=n;d[s]<n;)MF+=dfs(s,s,t);</pre>
50
                                                   37
       return MF;
51
                                                   38
52
     vector<int> cut_e;//最小割邊集
53
    bool vis[MAXN];
                                                   41
54
                                                   42
55
    void dfs_cut(int u){
       vis[u]=1://表示u屬於source的最小割集
       for(int i=g[u];~i;i=e[i].pre)
         if(e[i].r>0&&!vis[e[i].v])dfs cut(e[i
                                                   45
              1.v);
59
    T min cut(int s,int t){
60
      T ans=isap(s,t);
61
62
       memset(vis,0,sizeof(bool)*(n+1));
63
       dfs cut(s), cut e.clear();
                                                   50
       for(int u=0;u<=n;++u)if(vis[u])</pre>
64
                                                   51
         for(int i=g[u];~i;i=e[i].pre)
                                                   52
           if(!vis[e[i].v])cut e.push back(i);
66
67
       return ans;
                                                   54
68
                                                   55
69 };
```

#### 3.4 MinCostMaxFlow

1 template<typename TP>

```
struct MCMF{
    static const int MAXN=440;
    static const TP INF=999999999;
    struct edge{
      int v,pre;
      TP r,cost;
      edge(int v,int pre,TP r,TP cost):v(v),
           pre(pre),r(r),cost(cost){}
    int n,S,T;
    TP dis[MAXN],PIS,ans;
    bool vis[MAXN];
13
    vector<edge> e;
    int g[MAXN];
    void init(int n){
      memset(g,-1,sizeof(int)*((n=_n)+1));
17
      e.clear();
18
    void add_edge(int u,int v,TP r,TP cost,
19
          bool directed=false){
      e.push_back(edge(v,g[u],r,cost));
      g[u]=e.size()-1;
      e.push back(
      edge(u,g[v],directed?0:r,-cost));
      g[v]=e.size()-1;
24
25
26
    TP augment(int u, TP CF){
      if(u==T||!CF)return ans+=PIS*CF,CF;
       vis[u]=1;
29
      TP r=CF,d;
       for(int i=g[u];~i;i=e[i].pre){
31
         if(e[i].r&&!e[i].cost&&!vis[e[i].v]){
           d=augment(e[i].v,min(r,e[i].r));
32
33
           e[i].r-=d;
```

# Graph

S=s,T=t;

PIS=ans=0:

}return ans;

## 4.1 Augmenting Path

e[i^1].r+=d;

static deque<int>q;

for(int u=0;u<=n;++u)</pre>

TP mincost(int s,int t){

while(modlabel()){

while(augment(S,INF));

while(q.size()){

dis[T]=0,q.push back(T);

return CF-r;

bool modlabel(){

36

39

40

43

44

46

47

48

49

53

56

57

58

59

60

61

62

63

64

65

66

67

68 };

if(!(r-=d))break;

for(int u=0;u<=n;++u)dis[u]=INF;</pre>

int u=q.front();q.pop\_front();

for(int i=g[u];~i;i=e[i].pre){

q.front():S]){

for(int i=g[u];~i;i=e[i].pre)

return PIS+=dis[S], dis[S]<INF;</pre>

q.push\_front(e[i].v);

}else q.push back(e[i].v);

e[i].cost+=dis[e[i].v]-dis[u];

do memset(vis,0,sizeof(bool)\*(n+1));

dis[e[i].v]){

if(e[i^1].r&&(dt=dis[u]-e[i].cost)

if((dis[e[i].v]=dt)<=dis[q.size()?</pre>

```
1 #define MAXN1 505
 2 #define MAXN2 505
 3 int n1, n2; //n1 個點 連向 n2 個點
 4 | int match[MAXN2]; // 屬於 n2的點匹配了哪個點
 5 vector<int > g[MAXN1];//圖 0-base
 6 bool vis[MAXN2];//是否走訪過
   bool dfs(int u){
     for(int v:g[u]){
      if(vis[v]) continue;
      vis[v]=1:
      if(match[v]==-1||dfs(match[v]))
12
        return match[v]=u, 1;
13
14
    return 0;
15 }
16 int max match(){
    memset(match,-1,sizeof(int)*n2);
```

# Augmenting Path multiple

for(int i=0;i<n1;++i){</pre>

if(dfs(i)) ++ans;

return ans;

20

21

22

23

memset(vis,0,sizeof(bool)\*n2);

```
1 #define MAXN1 1005
 2 #define MAXN2 505
 3 int n1, n2; // n1 個點連向 n2 個點,其中 n2 個點可以
        匹配很多邊
 4 vector⟨int⟩ g[MAXN1];//圖 0-base
 5| size t c[MAXN2];//每個屬於n2點最多可以接受幾
        條匹配邊
 6 | vector<int> matchs[MAXN2];//每個屬於n2的點匹
        配了那些點
                                                  36
  bool vis[MAXN2];
                                                  37
  bool dfs(int u){
                                                  38
     for(int v:g[u]){
                                                  39
       if(vis[v])continue;
                                                  40
10
       vis[v] = 1;
11
                                                  41
12
       if(matchs[v].size()<c[v]){</pre>
                                                  42
         return matchs[v].push back(u), 1;
13
                                                  43
       }else for(size t j=0;j<matchs[v].size()</pre>
                                                  44
            ;++j){
                                                  45
15
         if(dfs(matchs[v][j]))
                                                  46
           return matchs[v][j]=u, 1;
16
                                                  47
17
18
                                                  49
19
     return 0;
                                                  50
20
   int max match(){
21
     for(int i=0;i<n2;++i) matchs[i].clear();</pre>
22
23
     int cnt=0:
24
     for(int u=0;u<n1;++u){</pre>
25
       memset(vis,0,sizeof(bool)*n2);
26
       if(dfs(u))++cnt;
27
28
     return cnt;
```

# 4.3 blossom matching

```
1 #define MAXN 505
2 int n; //1-base
3 vector<int> g[MAXN];
4 int MH[MAXN]; //output MH
5 int pa[MAXN],st[MAXN],S[MAXN],v[MAXN],t;
6 int lca(int x,int y){
    for(++t;;swap(x,y)){
      if(!x) continue;
      if(v[x]==t) return x;
      v[x] = t;
      x = st[pa[MH[x]]];
12
13 }
14 #define qpush(x) q.push(x),S[x]=0
```

```
while(q.size()){
    x=q.front(),q.pop();
    for(int y:g[x]){
      if(S[y]==-1){
        pa[y]=x,S[y]=1;
        if(!MH[y]){
          for(int lst;x;y=lst,x=pa[y])
            lst=MH[x],MH[x]=y,MH[y]=x;
          return 1:
        qpush(MH[y]);
      }else if(!S[y]&&st[y]!=st[x]){
        int l=lca(y,x);
        flower(y,x,1,q),flower(x,y,1,q);
 return 0;
int blossom(){
  memset(MH+1,0,sizeof(int)*n);
  int ans=0:
  for(int i=1; i<=n; ++i)</pre>
   if(!MH[i]&&bfs(i)) ++ans;
  return ans;
```

15 | void flower(int x,int y,int l,queue<int>&q){

if(S[y=MH[x]]==1)qpush(y);

memset(S+1,-1,sizeof(int)\*n);

st[x]=st[y]=1, x=pa[y];

iota(st+1, st+n+1, 1);

queue<int>q; qpush(x);

while(st[x]!=1){

pa[x]=y;

bool bfs(int x){

17

19

20

21

22

# 4.4 graphISO

```
1 const int MAXN=1005, K=30; // K要夠大
  const long long A=3,B=11,C=2,D=19,P=0
        xdefaced;
  long long f[K+1][MAXN];
  vector<int> g[MAXN],rg[MAXN];
  int n;
  void init(){
    for(int i=0;i<n;++i){</pre>
      f[0][i]=1;
      g[i].clear(), rg[i].clear();
11
  void add_edge(int u,int v){
    g[u].push_back(v), rg[v].push_back(u);
  long long point_hash(int u){//O(N)
    for(int t=1;t<=K;++t){</pre>
      for(int i=0;i<n;++i){</pre>
         f[t][i]=f[t-1][i]*A%P;
         for(int j:g[i])f[t][i]=(f[t][i]+f[t
              -1][i]*B%P)%P;
         for(int j:rg[i])f[t][i]=(f[t][i]+f[t
              -1][j]*C%P)%P;
```

```
if(i==u)f[t][i]+=D;//如果圖太大的話,
             把這行刪掉,執行一次後f[K]就會是所 46 }
                                                 47 LL KM(){
                                                      memset(My,0,sizeof(int)*(n+1));
        f[t][i]%=P;
                                                      memset(Mx,0,sizeof(int)*(n+1));
                                                 49
23
                                                      memset(ly,0,sizeof(LL)*(n+1));
^{24}
                                                 51
                                                      for(int x=1: x<=n: ++x){</pre>
25
    return f[K][u];
                                                       lx[x] = -INF;
                                                 52
26
                                                 53
                                                        for(int y=1; y<=n; ++y)</pre>
   vector<long long> graph_hash(){
27
                                                          lx[x] = max(lx[x],g[x][y]);
                                                 54
    vector<long long> ans;
                                                 55
    for(int i=0;i<n;++i)ans.push_back(</pre>
                                                 56
                                                      for(int x=1: x<=n: ++x) bfs(x):</pre>
          point_hash(i));//O(N^2)
                                                 57
                                                      LL ans = 0;
    sort(ans.begin(),ans.end());
                                                      for(int y=1; y<=n; ++y) ans+=g[My[y]][y];</pre>
    return ans;
                                                 59
                                                      return ans:
                                                 60 }
  4.5 KM
                                                    4.6 MaximumClique
1 #define MAXN 405
```

```
1 | struct MaxClique{
2 #define INF 0x3f3f3f3f3f3f3f3f3f
                                                        static const int MAXN=105:
3 int n; // 1-base · 0表示沒有匹配
4 LL g[MAXN][MAXN]; //input graph
5 int My[MAXN], Mx[MAXN]; //output match
6 LL lx[MAXN],ly[MAXN],pa[MAXN],Sy[MAXN];
7 bool vx[MAXN], vy[MAXN];
   void augment(int y){
    for(int x, z; y; y = z){
                                                         void init(int n){
       x=pa[y],z=Mx[x];
                                                          N=n;//0-base
       My[y]=x, Mx[x]=y;
12
13 }
                                                    10
   void bfs(int st){
                                                          g[u][v]=g[v][u]=1;
                                                   11
    for(int i=1; i<=n; ++i)</pre>
                                                    12
       Sy[i] = INF, vx[i]=vy[i]=0;
                                                    13
                                                        int dfs(int ns,int dep){
17
     queue<int> q; q.push(st);
                                                    14
                                                          if(!ns){
     for(;;){
                                                    15
                                                             if(dep>ans){
18
19
       while(q.size()){
                                                    16
                                                               ans=dep;
         int x=q.front(); q.pop();
20
                                                    17
21
                                                    18
                                                               return 1;
         for(int y=1; y<=n; ++y) if(!vy[y]){</pre>
                                                             }else return 0;
22
                                                    19
23
           LL t = lx[x]+ly[y]-g[x][y];
                                                   20
24
           if(t==0){
                                                   21
25
             pa[y]=x;
                                                   22
             if(!My[y]){augment(y);return;}
                                                    23
26
             vy[y]=1,q.push(My[y]);
28
           }else if(Sy[y]>t) pa[y]=x,Sy[y]=t;
                                                   25
29
                                                    26
30
                                                   27
       LL cut = INF;
       for(int y=1; y<=n; ++y)</pre>
         if(!vy[y]&&cut>Sy[y]) cut=Sy[y];
                                                    30
       for(int j=1; j<=n; ++j){</pre>
                                                    31
         if(vx[i]) lx[i] -= cut;
                                                   32
                                                          return 0;
         if(vy[j]) ly[j] += cut;
                                                    33
37
         else Sy[j] -= cut;
                                                         int clique(){
38
                                                          int u.v.ns:
39
       for(int y=1; y<=n; ++y){</pre>
         if(!vy[y]&&Sy[y]==0){
                                                   37
           if(!My[y]){augment(y);return;}
           vy[y]=1, q.push(My[y]);
42
43
                                                    40
                                                           return ans;
```

```
int g[MAXN][MAXN], dp[MAXN], stk[MAXN][MAXN
int sol[MAXN],tmp[MAXN];//sol[0~ans-1]為答
  memset(g,0,sizeof(g));
void add_edge(int u,int v){
      memcpy(sol,tmp,sizeof tmp);
  for(int i=0;i<ns;++i){</pre>
    if(dep+ns-i<=ans)return 0;</pre>
    int u=stk[dep][i],cnt=0;
    if(dep+dp[u]<=ans)return 0;</pre>
    for(int j=i+1; j<ns; ++j){</pre>
      int v=stk[dep][j];
      if(g[u][v])stk[dep+1][cnt++]=v;
    if(dfs(cnt,dep+1))return 1;
  for(ans=0,u=N-1;u>=0;--u){
    for(ns=0,tmp[0]=u,v=u+1;v<N;++v)</pre>
      if(g[u][v])stk[1][ns++]=v;
    dfs(ns,1),dp[u]=ans;
```

# MinimumMeanCycle

1 #include < cfloat > //for DBL\_MAX

43 };

```
1 int dp[MAXN][MAXN]; // 1-base, O(NM)
3 vector<tuple<int,int,int>> edge;
  double mmc(int n){//allow negative weight
     const int INF=0x3f3f3f3f;
     for(int t=0;t<n;++t){</pre>
       memset(dp[t+1],0x3f,sizeof(dp[t+1]));
       for(const auto &e:edge){
         int u,v,w;
         tie(u,v,w) = e;
11
         dp[t+1][v]=min(dp[t+1][v],dp[t][u]+w);
12
13
14
     double res = DBL MAX;
     for(int u=1;u<=n;++u){</pre>
       if(dp[n][u]==INF) continue;
       double val = -DBL MAX;
       for(int t=0;t<n;++t)</pre>
         val=max(val,(dp[n][u]-dp[t][u])*1.0/(n
              -t));
       res=min(res,val);
21
22
     return res;
```

#### 4.8 Rectilinear MST

1 / / 平面曼哈頓最小生成樹構造圖(去除非必要邊)

```
2 #define T int
3 #define INF 0x3f3f3f3f
4 struct point{
    T x,y;
    int id;//從0開始編號
     point(){}
    T dist(const point &p)const{
       return abs(x-p.x)+abs(y-p.y);
10
11 };
12 bool cmpx(const point &a,const point &b){
    return a.x<b.x||(a.x==b.x&&a.y<b.y);</pre>
15
  struct edge{
    int u,v;
     edge(int u,int v,T c):u(u),v(v),cost(c){}
    bool operator<(const edge&e)const{</pre>
       return cost<e.cost;</pre>
22 };
23 struct bit node{
    T mi:
    int id;
    bit node(const T&mi=INF,int id=-1):mi(mi),
```

#### 32 33 34 int bit find(int i,int m){ bit node x;

void bit\_update(int i,const T&data,int id){

if(data<bit[i].mi)bit[i]=bit\_node(data,</pre>

```
for(;i<=m;i+=i&(-i)) if(bit[i].mi<x.mi)x=</pre>
          bit[i]:
37
    return x.id;
38
39
  vector<edge> build graph(int n,point p[]){
    vector<edge> e;//edge for MST
    for(int dir=0;dir<4;++dir){//4種座標變換
      if(dir%2) for(int i=0;i<n;++i) swap(p[i</pre>
            ].x,p[i].y);
       else if(dir==2) for(int i=0;i<n;++i) p[i</pre>
           ].x=-p[i].x;
       sort(p,p+n,cmpx);
      vector<T> ga(n), gb;
       for(int i=0;i<n;++i)ga[i]=p[i].y-p[i].x;</pre>
       gb=ga, sort(gb.begin(),gb.end());
       gb.erase(unique(gb.begin(),gb.end()),gb.
            end());
       int m=gb.size();
      bit=vector<bit node>(m+1);
       for(int i=n-1;i>=0;--i){
         int pos=lower_bound(gb.begin(),gb.end
              (),ga[i])-gb.begin()+1;
```

int ans=bit find(pos,m);

if(~ans)e.push\_back(edge(p[i].id,p[ans

].id,p[i].dist(p[ans])));

bit\_update(pos,p[i].x+p[i].y,i);

vector<bit node> bit;

for(;i;i-=i&(-i)){

27 };

# 4.9 treeISO

55

56

58

} 57

return e;

```
1 const int MAXN=100005;
2 const long long X=12327, P=0xdefaced;
3 vector⟨int⟩ g[MAXN];
 4 bool vis[MAXN];
  long long dfs(int u){//hash ver
    vis[u]=1;
    vector<long long> tmp;
    for(auto v:g[u])if(!vis[v])tmp.PB(dfs(v));
    if(tmp.empty())return 177;
    long long ret=4931;
    sort(tmp.begin(),tmp.end());
    for(auto v:tmp)ret=((ret*X)^v)%P;
    return ret;
14
  string dfs(int x,int p){
    vector<string> c;
    for(int y:g[x])
      if(y!=p)c.emplace_back(dfs(y,x));
    sort(c.begin(),c.end());
```

while (stk.size()>=2){

```
string ret("(");
                                                                  int u = stk.back(); stk.pop_back 1 | static const int MAXN = 20;
                                                                                                                                                              int n;// 0-base
     for(auto &s:c)ret+=s;
                                                                       ();
                                                                                                       2 struct Edge{
                                                                                                                                                              vector<int>G[MAXN];
                                                                                                                                                              int rank[MAXN],label[MAXN];
23
    ret+=")";
                                                                  int v = stk.back(); stk.pop_back
                                                                                                           int u, v;
                                                   51
                                                                                                           Edge(int s, int d) : u(s), v(d) {}
                                                                                                                                                              bool mark[MAXN];
24
    return ret;
                                                                  match[u] = v;
                                                                                                                                                              void init(int n){n= n;
                                                   52
                                                                                                         bool isK33(int n, int degree[]){
                                                                                                                                                                for(int i=0;i<n;++i)G[i].clear();</pre>
                                                   53
                                                                  match[v] = u;
                                                   54
                                                                                                           int t = 0, z = 0;
                                                                                                           for(int i=0;i<n;++i){</pre>
                                                                                                                                                              void add_edge(int u,int v){
                                                   55
                                                                                                             if(degree[i] == 3)++t;
                                                   56
                                                                                                                                                         11
                                                                                                                                                                G[u].push back(v);
  4.10 一般圖最小權完美匹配
                                                   57
                                                            if (!found) break;
                                                                                                             else if(degree[i] == 0)++z;
                                                                                                                                                                G[v].push back(u);
                                                                                                                                                         12
                                                   58
                                                                                                      11
                                                                                                             else return false;
                                                                                                                                                         13
                                                   59
                                                          int ret = 0:
                                                                                                      12
                                                                                                                                                              vector<int> MCS(){
                                                                                                                                                         14
1 | struct Graph {
                                                          for (int i=0; i<n; i++)</pre>
                                                                                                      13
                                                                                                           return t == 6 \&\& t + z == n;
                                                                                                                                                                memset(rank,-1,sizeof(int)*n);
                                                   60
                                                                                                                                                         15
     // Minimum General Weiahted Matchina (
                                                   61
                                                            ret += edge[i][match[i]];
                                                                                                      14
                                                                                                                                                                memset(label.0.sizeof(int)*n);
          Perfect Match) 0-base
                                                   62
                                                                                                      15
                                                                                                         bool isK5(int n, int degree[]){
                                                                                                                                                         17
                                                                                                                                                                priority queue<pair<int,int> > pq;
     static const int MXN = 105;
                                                   63
                                                          return ret;
                                                                                                      16
                                                                                                           int f = 0, z = 0;
                                                                                                                                                                for(int i=0;i<n;++i)pq.push(make pair(0,</pre>
     int n, edge[MXN][MXN];
                                                                                                      17
                                                                                                           for(int i=0:i<n:++i){</pre>
                                                   64
     int match[MXN], dis[MXN], onstk[MXN];
                                                                                                             if(degree[i] == 4)++f;
                                                                                                                                                                for(int i=n-1;i>=0;--i)for(;;){
                                                   65 | }graph;
                                                                                                      18
                                                                                                             else if(degree[i] == 0)++z;
                                                                                                                                                                  int u=pq.top().second;pq.pop();
     vector<int> stk;
                                                                                                      19
                                                                                                                                                         20
     void init(int n) {
                                                                                                             else return false:
                                                                                                                                                                  if(~rank[u])continue;
                                                                                                      20
                                                                                                                                                         21
                                                                                                      21
                                                                                                                                                         22
                                                                                                                                                                  rank[u]=i;
                                                     4.11 全局最小割
       for (int i=0; i<n; i++)</pre>
                                                                                                           return f == 5 && f + z == n;
                                                                                                                                                                  for(auto v:G[u])if(rank[v]==-1){
                                                                                                      22
                                                                                                                                                         23
10
         for (int j=0; j<n; j++)</pre>
                                                                                                      23
                                                                                                                                                         24
                                                                                                                                                                    pq.push(make pair(++label[v],v));
           edge[i][j] = 0;
                                                                                                         // it judge a given graph is Homeomorphic
11
                                                                                                                                                         25
12
                                                                                                              with K33 or K5
                                                                                                                                                         26
                                                                                                                                                                  break;
                                                    1 const int INF=0x3f3f3f3f;
13
     void add_edge(int u, int v, int w) {
                                                                                                      25 bool isHomeomorphic(bool G[MAXN][MAXN],
                                                                                                                                                         27
                                                    1 template<typename T>
       edge[u][v] = edge[v][u] = w;
                                                                                                              const int n){
14
                                                                                                                                                         28
                                                                                                                                                                vector<int> res(n);
                                                      struct stoer_wagner{// 0-base
15
                                                                                                           for(;;){
                                                                                                                                                         29
                                                                                                                                                                for(int i=0;i<n;++i)res[rank[i]]=i;</pre>
                                                        static const int MAXN=150;
                                                                                                      26
    bool SPFA(int u){
                                                                                                      27
                                                                                                             int cnt = 0;
                                                                                                                                                         30
                                                                                                                                                                return res;
16
                                                        T g[MAXN][MAXN], dis[MAXN];
17
       if (onstk[u]) return true;
                                                                                                      28
                                                                                                             for(int i=0;i<n;++i){</pre>
                                                                                                                                                         31
                                                        int nd[MAXN],n,s,t;
18
       stk.push back(u);
                                                                                                      29
                                                                                                               vector<Edge> E;
                                                                                                                                                              bool check(vector<int> ord){//弦圖判定
                                                                                                                                                         32
                                                        void init(int n){
19
       onstk[u] = 1;
                                                                                                      30
                                                                                                               for(int j=0;j<n&E.size()<3;++j)</pre>
                                                                                                                                                                for(int i=0;i<n;++i)rank[ord[i]]=i;</pre>
       for (int v=0; v<n; v++){</pre>
                                                                                                                 if(G[i][j] && i != j)
20
                                                                                                      31
                                                                                                                                                                memset(mark,0,sizeof(bool)*n);
                                                          for(int i=0;i<n;++i)</pre>
         if (u != v && match[u] != v && !onstk[
                                                                                                                   E.push back(Edge(i, j));
                                                                                                      32
                                                                                                                                                                for(int i=0;i<n;++i){</pre>
                                                            for(int j=0;j<n;++j)g[i][j]=0;</pre>
                                                                                                      33
                                                                                                               if(E.size() == 1){
              v]){
                                                                                                                                                                  vector<pair<int,int> > tmp;
                                                                                                                 G[i][E[0].v] = G[E[0].v][i] = false;
           int m = match[v];
                                                                                                      34
22
                                                                                                                                                                  for(auto u:G[ord[i]])if(!mark[u])
                                                        void add_edge(int u,int v,T w){
                                                   ^{12}
23
           if (dis[m] > dis[u] - edge[v][m] +
                                                                                                      35
                                                                                                               }else if(E.size() == 2){
                                                                                                                                                                    tmp.push_back(make_pair(rank[u],u));
                                                   13
                                                          g[u][v]=g[v][u]+=w;
                edge[u][v]){
                                                                                                      36
                                                                                                                 G[i][E[0].v] = G[E[0].v][i] = false; 39
                                                                                                                                                                  sort(tmp.begin(),tmp.end());
                                                   14
             dis[m] = dis[u] - edge[v][m] +
                                                                                                      37
                                                                                                                 G[i][E[1].v] = G[E[1].v][i] = false;
                                                                                                                                                                  if(tmp.size()){
                                                        T min cut(){
                                                   15
                                                                                                      38
                                                                                                                 G[E[0].v][E[1].v] = G[E[1].v][E[0].v
                  edge[u][v];
                                                                                                                                                                    int u=tmp[0].second;
                                                         T ans=INF:
             onstk[v] = 1;
                                                                                                                      ] = true;
                                                                                                                                                         42
                                                                                                                                                                    set<int> S;
                                                   17
                                                          for(int i=0;i<n;++i)nd[i]=i;</pre>
             stk.push back(v);
26
                                                                                                      39
                                                                                                                 ++cnt;
                                                                                                                                                                    for(auto v:G[u])S.insert(v);
                                                          for(int ind,tn=n;tn>1;--tn){
                                                                                                                                                         43
             if (SPFA(m)) return true;
27
                                                                                                      40
                                                                                                                                                                    for(size_t j=1;j<tmp.size();++j)</pre>
                                                            for(int i=1;i<tn;++i)dis[nd[i]]=0;</pre>
                                                                                                                                                         44
28
             stk.pop back();
                                                                                                      41
                                                                                                                                                         45
                                                                                                                                                                      if(!S.count(tmp[j].second))return
                                                   20
                                                            for(int i=1;i<tn;++i){</pre>
             onstk[v] = 0;
                                                                                                             if(cnt == 0)break;
29
                                                                                                      42
                                                   21
                                                              ind=i:
30
                                                                                                      43
                                                                                                                                                         46
                                                   22
                                                              for(int j=i;j<tn;++j){</pre>
                                                                                                           static int degree[MAXN];
31
         }
                                                                                                      44
                                                                                                                                                                  mark[ord[i]]=1;
                                                                                                                                                         47
                                                   23
                                                                dis[nd[j]]+=g[nd[i-1]][nd[j]];
                                                                                                           fill(degree, degree + n, 0);
32
                                                                                                                                                         48
                                                   24
                                                                if(dis[nd[ind]]<dis[nd[j]])ind=j;</pre>
                                                                                                           for(int i=0;i<n;++i){</pre>
       onstk[u] = 0;
                                                                                                      46
                                                                                                                                                         49
                                                                                                                                                                return 1;
                                                   25
       stk.pop back();
                                                                                                      47
                                                                                                             for(int j=i+1; j<n; ++j){</pre>
                                                                                                                                                         50
                                                   26
                                                              swap(nd[ind],nd[i]);
       return false;
                                                                                                      48
                                                                                                               if(!G[i][j])continue;
                                                   27
36
                                                                                                      49
                                                                                                               ++degree[i];
                                                            if(ans>dis[nd[ind]])ans=dis[t=nd[ind
                                                   28
     int solve() {
                                                                                                      50
                                                                                                               ++degree[j];
                                                                 ]],s=nd[ind-1];
       // find a match
                                                                                                      51
                                                            for(int i=0;i<tn;++i)</pre>
                                                   29
       for (int i=0; i<n; i+=2){</pre>
                                                                                                      52
                                                                                                          }
                                                                                                                                                            4.14 最小斯坦納樹 DP
                                                              g[nd[ind-1]][nd[i]]=g[nd[i]][nd[ind
                                                   30
         match[i] = i+1, match[i+1] = i;
                                                                                                           return !(isK33(n, degree) || isK5(n,
40
                                                                   -1]]+=g[nd[i]][nd[ind]];
                                                                                                                degree));
41
                                                   31
42
       for(;;){
                                                   32
                                                          return ans;
                                                                                                                                                          1 //n個點,其中r個要構成斯坦納樹
         int found = 0;
                                                                                                                                                          2 //答案在max(dp[(1<<r)-1][k]) k=0~n-1
         for (int i=0; i<n; i++) dis[i] = onstk</pre>
                                                                                                                                                          3 //p表示要構成斯坦納樹的點集
              [i] = 0;
                                                                                                         4.13 弦圖完美消除序列
         for (int i=0; i<n; i++){</pre>
                                                                                                                                                          4 //0( n^3 + n*3^r + n^2*2^r )
           stk.clear();
                                                                                                                                                          5 #define REP(i,n) for(int i=0;i<(int)n;++i)</pre>
           if (!onstk[i] && SPFA(i)){
47
                                                                                                                                                          6 const int MAXN=30, MAXM=8; // 0-base
                                                              平面圖判定
             found = 1;
                                                                                                                                                            const int INF=0x3f3f3f3f;
48
                                                                                                       1 struct chordal{
```

static const int MAXN=1005;

8 int dp[1<<MAXM][MAXN];</pre>

```
9 int g[MAXN][MAXN];// 🗟
  void init(){memset(g,0x3f,sizeof(g));}
   void add edge(int u,int v,int w){
                                                    39
    g[u][v]=g[v][u]=min(g[v][u],w);
                                                    40
13
                                                    41
   void steiner(int n,int r,int *p){
                                                    42
15
     REP(k,n)REP(i,n)REP(j,n)
                                                    43
       g[i][j]=min(g[i][j],g[i][k]+g[k][j]);
16
                                                    44
     REP(i,n)g[i][i]=0;
     REP(i,r)REP(j,n)dp[1<<i][j]=g[p[i]][j];</pre>
                                                    45
19
     for(int i=1;i<(1<<r);++i){</pre>
                                                    46
       if(!(i&(i-1)))continue;
20
       REP(j,n)dp[i][j]=INF;
21
                                                    47
22
       REP(j,n){
                                                    48
23
         int tmp=INF:
                                                    49
24
         for(int s=i&(i-1);s;s=i&(s-1))
                                                    50
           tmp=min(tmp,dp[s][j]+dp[i^s][j]);
25
                                                    51
         REP(k,n)dp[i][k]=min(dp[i][k],g[j][k]+
26
27
                                                    53
28
                                                    54
29
                                                    55
                                                    56
```

# 4.15 最小樹形圖 朱劉

1 template<typename T>

```
2 struct zhu liu{
    static const int MAXN=110,MAXM=10005;
    struct node{
       int u,v;
       T w, tag;
       node *1.*r:
       node(int u=0, int v=0, T w=0): u(u), v(v), w(v)
            w),tag(0),1(0),r(0){}
       void down(){
10
         w+=tag;
         if(1)1->tag+=tag;
12
         if(r)r->tag+=tag;
13
         tag=0;
14
     }mem[MAXM];//靜態記憶體
     node *pq[MAXN*2],*E[MAXN*2];
     int st[MAXN*2],id[MAXN*2],m;
     void init(int n){
19
       for(int i=1;i<=n;++i){</pre>
         pq[i]=E[i]=0, st[i]=id[i]=i;
20
21
       }m=0;
22
     node *merge(node *a, node *b){//skew heap
       if(!a||!b)return a?a:b;
       a->down(),b->down();
       if(b->w<a->w)return merge(b,a);
       swap(a->1,a->r);
       a->l=merge(b,a->l);
       return a:
     void add edge(int u,int v,T w){
       if(u!=v)pq[v]=merge(pq[v],&(mem[m++]=
            node(u,v,w)));
     int find(int x,int *st){
       return st[x]==x?x:st[x]=find(st[x],st);
```

# 4.16 穩定婚姻模板

無解

T build(int root,int n){

while(pq[i]){

T ans=0; int N=n, all=n;

for(int i=1;i<=N;++i){</pre>

continue:

ans+=E[i]->w;

57

58

59

61

62 };

if(i==root||!pq[i])continue;

pq[i]->down(),E[i]=pq[i];

pq[i]=merge(pq[i]->1,pq[i]->r);

if(find(E[i]->u,id)!=find(i,id))

if(find(E[i]->u,id)==find(i,id))

if(find(E[i]->u,st)==find(i,st)){

if(pq[i])pq[i]->tag-=E[i]->w;

find(E[u]->u,id)){

pq[N]=merge(pq[N],pq[u]);

for(int u=find(E[i]->u,id);u!=i;u=

if(pq[u])pq[u]->tag-=E[u]->w;

}else st[find(i,st)]=find(E[i]->u,st)

return all==1?ans:-INT MAX;//圖不連通就

pq[++N]=pq[i];id[N]=N;

id[find(u,id)]=N;

st[N]=find(i.st):

id[find(i,id)]=N;

24 }

16

17

18

19

20

21

22

23

24

25

33

34

51 }

};

```
1 | queue < int > Q;
2 for ( i : 所有考生 ) {
   設定在第0志願;
   0.push(考生i);
6 while(Q.size()){
   當前考生=Q.front();Q.pop();
   while ( 此考生未分發 ) {
     指標移到下一志願;
     if (已經沒有志願 or 超出志願總數)
     計算該考生在該科系加權後的總分;
11
     if (不符合科系需求) continue;
12
     if (目前科系有餘額) {
13
14
      依加權後分數高低順序將考生id加入科系錄
          取名單中;
      break:
15
16
17
     if (目前科系已額滿) {
      if ( 此考生成績比最低分數還高 ) {
18
        依加權後分數高低順序將考生id加入科系
19
           錄取名單:
        Q.push(被踢出的考生);
21
22
```

# 6 Number Theory

for(int i=2;i<MAXPRIME;++i) {</pre>

bool g\_test(const LL &g, const LL &p, const

vector<LL> &v) {

return false:

if(p==2) return 1;

return true;

vector<LL> v;

Factor(p-1,v);

for(int i=0;i<v.size();++i)</pre>

LL primitive root(const LL &p) {

if(modexp(g,(p-1)/v[i],p)==1)

#### 6.1 basic

```
1 template<typename T>
                                                      void gcd(const T &a,const T &b,T &d,T &x,T &
  5.1 simplex
                                                         if(!b) d=a,x=1,v=0;
                                                         else gcd(b,a%b,d,y,x), y-=x*(a/b);
1 /*taraet:
                                                       long long int phi[N+1];
    max \setminus sum_{j=1}^n A_{0,j}*x_j
                                                       void phiTable(){
                                                         for(int i=1;i<=N;i++)phi[i]=i;</pre>
    \sum_{j=1}^n A_{i,j}*x_j <= A_{i,0} | i=1 \sim m
                                                         for(int i=1;i<=N;i++)for(x=i*2;x<=N;x+=i)</pre>
   x_j >= 0 \mid j=1\sim n
```

phi[x]-=phi[i]; VDB = vector<double>\*/ 10 template < class VDB> void all divdown(const LL &n) {// all n/x VDB simplex(int m,int n,vector<VDB> a){ for(LL a=1;a<=n;a=n/(n/(a+1))){</pre> vector<int> left(m+1), up(n+1); 13 // dosomethina: iota(left.begin(), left.end(), n); 14 iota(up.begin(), up.end(), 0); 15 auto pivot = [&](int x, int y){ const int MAXPRIME = 1000000;

swap(left[x], up[y]); int iscom[MAXPRIME], prime[MAXPRIME], **auto** k = a[x][y]; a[x][y] = 1;primecnt: vector<int> pos; int phi[MAXPRIME], mu[MAXPRIME]; for(int j = 0; j <= n; ++j){</pre> void sieve(void){ a[x][j] /= k;memset(iscom,0,sizeof(iscom)); if(a[x][j] != 0) pos.push\_back(j); primecnt = 0; phi[1] = mu[1] = 1;

if(a[i][y]==0 || i == x) continue; if(!iscom[i]) { k = a[i][y], a[i][y] = 0;prime[primecnt++] = i; for(int j : pos) a[i][j] -= k\*a[x][j]; mu[i] = -1; 26 phi[i] = i-1; 27 28 for(int x,y;;){ 29 for(int j=0;j<primecnt;++j) {</pre>

43

44

for(int i=x=1; i <= m; ++i)</pre> 30 int k = i \* prime[j]; if(a[i][0] < a[x][0]) x = i;31 if(k>=MAXPRIME) break; **if**(a[x][0]>=0) **break**; iscom[k] = prime[j]; 32 for(int j=y=1; j <= n; ++j)</pre> 33 if(i%prime[j]==0) { **if**(a[x][j]<a[x][y]) y = j; mu[k] = 0;if(a[x][y]>=0) return VDB();//infeasible phi[k] = phi[i] \* prime[j];

pivot(x, y); 36 break; 37 } else { for(int x,y;;){ for(int j=y=1; j <= n; ++j)</pre> phi[k] = phi[i] \* (prime[j]-1); 39 if(a[0][j] > a[0][y]) y = j;40 **if**(a[0][y]<=0) **break**; 41 42 for(int i=1; i<=m; ++i) if(a[i][v] > 0)

if(x == -1 || a[i][0]/a[i][y] < a[x][0]/a[x][y]) x = i;if(x == -1) return VDB();//unbounded pivot(x, y);

for(int i = 0; i <= m; ++i){</pre>

Linear Programming

VDB ans(n + 1); for(int i = 1: i <= m: ++i)</pre> if(left[i] <= n) ans[left[i]] = a[i][0];</pre> ans[0] = -a[0][0];return ans;

```
v.erase(unique(v.begin(), v.end()), v.end 115 return -1;
          ());
                                                  116 }
     for(LL g=2;g<p;++g)</pre>
                                                      template<typename T>
57
       if(g test(g,p,v))
58
                                                  119 T Euler(T n){
     puts("primitive root NOT FOUND");
                                                  120
                                                        T ans=n:
60
     return -1:
                                                  121
                                                        for(T i=2:i*i<=n:++i){</pre>
                                                          if(n%i==0){
61
                                                  122
   int Legendre(const LL &a, const LL &p) {
                                                             ans=ans/i*(i-1);
                                                  123
        return modexp(a\%p,(p-1)/2,p); }
                                                  124
                                                             while(n%i==0)n/=i;
                                                  125
    LL inv(const LL &a, const LL &n) {
                                                  126
     LL d,x,v;
                                                        if(n>1)ans=ans/n*(n-1);
65
                                                  127
     gcd(a,n,d,x,y);
                                                  128
                                                        return ans:
67
     return d==1 ? (x+n)%n : -1:
                                                  129
68
                                                  130
                                                       //Chinese remainder theorem
69
   int inv[maxN];
                                                      template<typename T>
70
   LL invtable(int n,LL P){
                                                      T pow mod(T n, T k, T m){
                                                  133
72
     inv[1]=1;
                                                  134
                                                        T ans=1:
73
     for(int i=2;i<n;++i)</pre>
                                                  135
                                                        for(n=(n)=m?n\%m:n);k;k>>=1){
       inv[i]=(P-(P/i))*inv[P%i]%P;
                                                          if(k&1)ans=ans*n%m;
74
                                                  136
75
                                                  137
                                                          n=n*n%m;
76
                                                  138
   LL log mod(const LL &a, const LL &b, const
                                                  139
                                                        return ans;
        LL &p) {
                                                  140 3
     // a ^ x = b \pmod{p}
                                                      template<typename T>
     int m=sqrt(p+.5), e=1;
                                                      T crt(vector<T> &m.vector<T> &a){
80
     LL v=inv(modexp(a,m,p), p);
                                                        T M=1,tM,ans=0;
     map<LL,int> x;
                                                  144
                                                         for(int i=0;i<(int)m.size();++i)M*=m[i];</pre>
82
     x[1]=0:
                                                  145
                                                         for(int i=0;i<(int)a.size();++i){</pre>
83
     for(int i=1;i<m;++i) {</pre>
                                                  146
                                                          tM=M/m[i];
84
       e = LLmul(e,a,p);
                                                  147
                                                          ans=(ans+(a[i]*tM%M)*pow mod(tM,Euler(m[
85
       if(!x.count(e)) x[e] = i;
                                                               i])-1,m[i])%M)%M;
86
                                                          /*如果m[i]是質數, Euler(m[i])-1=m[i]-2,
                                                  148
87
     for(int i=0;i<m;++i) {</pre>
                                                                就不用算Euler了*/
88
       if(x.count(b)) return i*m + x[b];
                                                  149
89
       b = LLmul(b,v,p);
                                                  150
                                                        return ans;
90
                                                  151
     return -1;
                                                  152
92
                                                  153 //java code
93
                                                  154 / / 求 sart (N) 的 連 分 數
    LL Tonelli Shanks(const LL &n, const LL &p)
                                                      public static void Pell(int n){
                                                        BigInteger N,p1,p2,q1,q2,a0,a1,a2,g1,g2,h1
     // x^2 = n \pmod{p}
                                                             ,h2,p,q;
     if(n==0) return 0;
                                                        g1=q2=p1=BigInteger.ZERO;
     if(Legendre(n,p)!=1) while(1) { puts("SQRT
                                                        h1=q1=p2=BigInteger.ONE;
           ROOT does not exist"); }
                                                        a0=a1=BigInteger.valueOf((int)Math.sqrt
                                                  159
     int S = 0;
                                                             (1.0*n));
99
     LL Q = p-1;
                                                        BigInteger ans=a0.multiply(a0);
                                                  160
     while( !(Q&1) ) { Q>>=1; ++S; }
                                                        if(ans.equals(BigInteger.valueOf(n))){
                                                  161
     if(S==1) return modexp(n\%p,(p+1)/4,p);
101
                                                          System.out.println("No solution!");
                                                  162
102
     LL z = 2:
                                                  163
                                                          return ;
103
     for(;Legendre(z,p)!=-1;++z)
                                                  164
104
     LL c = modexp(z,Q,p);
                                                         while(true){
                                                  165
     LL R = modexp(n\%p,(Q+1)/2,p), t = modexp(n
105
                                                           g2=a1.multiply(h1).substract(g1);
                                                  166
          p,0,p);
                                                          h2=N.substract(g2.pow(2)).divide(h1);
                                                  167
106
     int M = S;
                                                          a2=g2.add(a0).divide(h2);
                                                  168
107
     while(1) {
                                                          p=a1.multiply(p2).add(p1);
                                                  169
108
       if(t==1) return R:
                                                           q=a1.multiply(q2).add(q1);
                                                  170
109
       LL b = modexp(c,1L << (M-i-1),p);
                                                           if(p.pow(2).substract(N.multiply(q.pow
                                                  171
       R = LLmul(R,b,p);
                                                                (2))).compareTo(BigInteger.ONE)==0)
       t = LLmul( LLmul(b,b,p), t, p);
                                                               break;
       c = LLmul(b,b,p);
                                                           g1=g2;h1=h2;a1=a2;
                                                  172
       M = i;
113
                                                           p1=p2;p2=p;
```

```
174 | q1=q2;q2=q;
175 | }
176 | System.out.println(p+" "+q);
177 | }
```

#### 6.2 bit\_set

```
1 | void sub set(int S){
    int sub=S:
    do{
       //對某集合的子集合的處理
       sub=(sub-1)&S;
    }while(sub!=S);
  void k sub set(int k,int n){
    int comb=(1<<k)-1,S=1<<n;</pre>
    while(comb<S){</pre>
       //對大小為k的子集合的處理
11
12
      int x=comb&-comb,y=comb+x;
13
       comb = ((comb\&\sim y)/x>>1)|y;
14
15 }
```

### 6.3 cantor\_expansion

```
1 int factorial[MAXN];
 2 void init(){
     factorial[0]=1;
     for(int i=1;i<=MAXN;++i)factorial[i]=</pre>
          factorial[i-1]*i;
   int encode(const vector<int> &s){
     int n=s.size(),res=0;
     for(int i=0;i<n;++i){</pre>
       int t=0;
       for(int j=i+1; j<n;++j)</pre>
         if(s[j]<s[i])++t;
12
       res+=t*factorial[n-i-1];
13
14
     return res;
15
   vector<int> decode(int a,int n){
17
     vector<int> res;
     vector<bool> vis(n,0);
18
     for(int i=n-1:i>=0:--i){
19
       int t=a/factorial[i],j;
20
21
       for(j=0;j<n;++j)</pre>
         if(!vis[j]){
22
23
            if(t==0)break;
24
            --t:
25
26
       res.push_back(j);
27
       vis[j]=1;
28
       a%=factorial[i];
29
30
     return res:
31
```

#### 6.4 FFT

```
1 | template < typename T, typename VT = vector <
        complex<T>>>
  struct FFT{
     const T pi;
     FFT(const T pi=acos((T)-1)):pi(pi){}
     unsigned bit reverse(unsigned a,int len){
  a = ((a\&0x555555555) < < 1) | ((a\&0xAAAAAAAAA) > > 1);
  a=((a&0x33333333U)<<2)|((a&0xCCCCCCCU)>>2);
  a = ((a\&0x0F0F0F0FU) << 4) | ((a\&0xF0F0F0F0U) >> 4);
  a=((a&0x00FF00FFU)<<8)|((a&0xFF00FF00U)>>8);
  a=((a\&0x0000FFFFU)<<16)|((a\&0xFFFF0000U)
        >>16);
       return a>>(32-len);
12
     void fft(bool is_inv,VT &in,VT &out,int N)
13
       int bitlen= lg(N), num=is inv?-1:1;
14
       for(int i=0;i<N;++i)out[bit_reverse(i,</pre>
15
            bitlen)]=in[i];
       for(int step=2;step<=N;step<<=1){</pre>
16
17
         const int mh=step>>1;
18
         for(int i=0:i<mh:++i){</pre>
           complex<T> wi=exp(complex<T>(0,i*num
19
                 *pi/mh));
           for(int j=i;j<N;j+=step){</pre>
20
21
             int k=j+mh;
              complex<T> u=out[j],t=wi*out[k];
22
23
              out[i]=u+t;
24
              out[k]=u-t;
25
26
27
       if(is inv)for(int i=0:i<N:++i)out[i]/=N:
28
29
30
```

#### 6.5 find real root

```
1 / / an*x^n + ... + a1x + a0 = 0;
  int sign(double x){
    return x < -eps ? -1 : x > eps;
  double get(const vector<double>&coef, double
    double e = 1, s = 0;
    for(auto i : coef) s += i*e, e *= x;
    return s;
10
11
  double find(const vector<double>&coef, int n
        , double lo, double hi){
    double sign_lo, sign_hi;
    if( !(sign lo = sign(get(coef,lo))) )
         return lo;
    if( !(sign_hi = sign(get(coef,hi))) )
         return hi;
    if(sign lo * sign hi > 0) return INF;
    for(int stp = 0; stp < 100 && hi - lo >
         eps; ++stp){
```

```
double m = (lo+hi)/2.0;
                                                             f[j+k+(1<< i)] = u-v, f[j+k] = u+v;
19
       int sign mid = sign(get(coef,m));
                                                  22
       if(!sign mid) return m;
                                                      if(inverse) for(auto &a:f) a/=f.size();
20
                                                  23
21
       if(sign lo*sign mid < 0) hi = m;</pre>
                                                  24
                                                      return f;
       else lo = m;
22
23
24
    return (lo+hi)/2.0:
25
                                                     6.7 LinearCongruence
26
   vector<double> cal(vector<double>coef, int n
    vector<double>res:
28
                                                   1 | pair<LL,LL> LinearCongruence(LL a[],LL b[],
    if(n == 1){
29
                                                         LL m[], int n) {
30
       if(sign(coef[1])) res.pb(-coef[0]/coef
            [1]);
       return res;
31
32
33
    vector<double>dcoef(n);
34
    for(int i = 0; i < n; ++i) dcoef[i] = coef
                                                         m[i] /= d:
          [i+1]*(i+1);
    vector<double>droot = cal(dcoef, n-1);
35
    droot.insert(droot.begin(), -INF);
36
37
    droot.pb(INF);
38
    for(int i = 0; i+1 < droot.size(); ++i){</pre>
       double tmp = find(coef, n, droot[i],
39
                                                  12
           droot[i+1]);
                                                              (-1LL,0LL);
      if(tmp < INF) res.pb(tmp);</pre>
40
41
                                                             )*m[i];
42
    return res;
                                                  14
43
                                                  15
44
                                                  16
45
   int main () {
                                                  17
    vector<double>ve:
                                                            ,lastm);
    vector<double>ans = cal(ve, n);
    // 視情況把答案 +eps, 避免 -0
48
                                                     6.8 Lucas
  6.6 FWT
```

```
// a[i]*x = b[i] (mod m[i])
for(int i=0;i<n;++i) {</pre>
  LL x, y, d = extgcd(a[i],m[i],x,y);
                                             20
  if(b[i]%d!=0) return make_pair(-1LL,0LL)
  b[i] = LLmul(b[i]/d,x,m[i]);
LL lastb = b[0], lastm = m[0];
for(int i=1;i<n;++i) {</pre>
  LL x, y, d = extgcd(m[i],lastm,x,y);
  if((lastb-b[i])%d!=0) return make_pair
  lastb = LLmul((lastb-b[i])/d,x,(lastm/d)
  lastm = (lastm/d)*m[i];
  lastb = (lastb+b[i])%lastm;
return make pair(lastb<0?lastb+lastm:lastb</pre>
```

11

13

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65

67

```
1 vector<int> F OR T(vector<int> f, bool
        inverse){
     for(int i=0; (2<<i)<=f.size(); ++i)</pre>
       for(int j=0; j<f.size(); j+=2<<i)</pre>
         for(int k=0; k<(1<<i); ++k)</pre>
           f[j+k+(1<< i)] += f[j+k]*(inverse)
                 ?-1:1);
     return f;
   vector<int> rev(vector<int> A) {
     for(int i=0; i<A.size(); i+=2)</pre>
       swap(A[i],A[i^(A.size()-1)]);
     return A;
12 }
   vector<int> F AND T(vector<int> f, bool
     return rev(F_OR_T(rev(f), inverse));
15
   vector<int> F XOR T(vector<int> f, bool
     for(int i=0; (2<<i)<=f.size(); ++i)</pre>
       for(int j=0; j<f.size(); j+=2<<i)</pre>
         for(int k=0; k<(1<<i); ++k){</pre>
19
           int u=f[j+k], v=f[j+k+(1<<i)];</pre>
```

```
1 int mod_fact(int n,int &e){
    e=0:
    if(n==0)return 1;
    int res=mod fact(n/P,e);
    e += n/P:
    if((n/P)%2==0)return res*fact[n%P]%P;
    return res*(P-fact[n%P])%P;
   int Cmod(int n,int m){
    int a1,a2,a3,e1,e2,e3;
    a1=mod fact(n,e1);
     a2=mod fact(m,e2);
     a3=mod fact(n-m,e3);
    if(e1>e2+e3)return 0;
14
    return a1*inv(a2*a3%P,P)%P;
15
```

# 6.9 Matrix

```
1 template < typename T>
2 struct Matrix{
   using rt = std::vector<T>;
```

```
using mt = std::vector<rt>;
using matrix = Matrix<T>;
int r,c;
                                                71
mt m:
Matrix(int r, int c):r(r),c(c),m(r,rt(c))
                                                73
rt& operator[](int i){return m[i];}
                                                74
matrix operator+(const matrix &a){
                                                75
  matrix rev(r,c);
                                                76
  for(int i=0;i<r;++i)</pre>
    for(int j=0;j<c;++j)</pre>
      rev[i][j]=m[i][j]+a.m[i][j];
  return rev:
matrix operator-(const matrix &a){
  matrix rev(r,c);
  for(int i=0;i<r;++i)</pre>
    for(int j=0;j<c;++j)</pre>
      rev[i][j]=m[i][j]-a.m[i][j];
  return rev;
matrix operator*(const matrix &a){
  matrix rev(r.a.c):
  matrix tmp(a.c.a.r):
  for(int i=0;i<a.r;++i)</pre>
    for(int j=0;j<a.c;++j)</pre>
                                                11
      tmp[j][i]=a.m[i][j];
                                                12
  for(int i=0;i<r;++i)</pre>
    for(int j=0;j<a.c;++j)</pre>
      for(int k=0;k<c;++k)</pre>
        rev.m[i][j]+=m[i][k]*tmp[j][k];
  return rev:
                                                17
bool inverse(){
  Matrix t(r,r+c);
  for(int y=0;y<r;y++){</pre>
    t.m[y][c+y] = 1;
    for(int x=0;x<c;++x)</pre>
      t.m[y][x]=m[y][x];
                                                25
  if( !t.gas() )
    return false;
                                                27
  for(int y=0;y<r;y++)</pre>
    for(int x=0;x<c;++x)</pre>
      m[y][x]=t.m[y][c+x]/t.m[y][y];
  return true;
T gas(){
  vector<T> lazy(r,1);
  bool sign=false;
  for(int i=0;i<r;++i){</pre>
    if( m[i][i]==0 ){
      int j=i+1;
      while(j<r&&!m[j][i])j++;
      if(j==r)continue;
      m[i].swap(m[j]);
                                                39
      sign=!sign;
    for(int j=0;j<r;++j){</pre>
      if(i==j)continue;
      lazy[j]=lazy[j]*m[i][i];
      T mx=m[j][i];
      for(int k=0;k<c;++k)</pre>
        m[j][k]=m[j][k]*m[i][i]-m[i][k]*mx
```

```
T det=sign?-1:1;
       for(int i=0;i<r;++i){</pre>
         det = det*m[i][i];
         det = det/lazy[i];
         for(auto &j:m[i])j/=lazy[i];
       return det:
77 };
```

#### 6.10MillerRobin

```
1 | LL LLmul(LL a, LL b, const LL &mod) {
    LL ans=0:
     while(b) {
      if(b&1) {
         if(ans>=mod) ans-=mod;
      a<<=1, b>>=1;
      if(a>=mod) a-=mod;
    return ans;
  LL mod mul(LL a, LL b, LL m){
    a\%=m,b\%=m;/* fast for m < 2^58 */
    LL y=(LL)((double)a*b/m+0.5);
    LL r=(a*b-y*m)%m;
     return r<0?r+m:r;</pre>
   template<typename T>
  T pow(T a,T b,T mod){//a^b\%mod}
    T ans=1:
    for(;b;a=mod_mul(a,a,mod),b>>=1)
      if(b&1)ans=mod mul(ans,a,mod);
     return ans;
26
  int sprp[3]={2,7,61};//int範圍可解
  int llsprp
        [7]={2,325,9375,28178,450775,9780504,
  | 1795265022};//至少unsigned Long Long範圍
  template<typename T>
  bool isprime(T n,int *sprp,int num){
    if(n==2)return 1;
    if(n<2||n%2==0)return 0;
    int t=0;
    T u=n-1;
     for(;u%2==0;++t)u>>=1;
     for(int i=0;i<num;++i){</pre>
      T a=sprp[i]%n;
      if(a==0||a==1||a==n-1)continue;
      T x=pow(a,u,n);
       if(x==1||x==n-1)continue;
       for(int j=0;j<t;++j){</pre>
        x=mod mul(x,x,n);
        if(x==1)return 0;
        if(x==n-1)break;
      if(x==n-1)continue;
      return 0;
49
    return 1;
```

#### 6.11 NTT

```
1 | 2615053605667*(2^18)+1,3
2 15*(2^27)+1,31
3 479*(2^21)+1,3
4 7*17*(2^23)+1,3
5 3*3*211*(2^19)+1,5
6 25*(2^22)+1,3
  template < typename T, typename VT = vector < T > >
   struct NTT{
     const T P,G;
     NTT(T p=(1<<23)*7*17+1,T g=3):P(p),G(g){}
     unsigned bit reverse(unsigned a, int len){
12
       //Look FFT.cpp
13
     T pow mod(T n, T k, T m) \{
15
       T ans=1;
       for(n=(n)=m?n\%m:n);k;k>>=1){
         if(k&1)ans=ans*n%m;
18
         n=n*n%m;
19
20
       return ans;
21
     void ntt(bool is inv,VT &in,VT &out,int N)
       int bitlen= lg(N);
       for(int i=0;i<N;++i)out[bit reverse(i,</pre>
24
             bitlen) |=in[i];
       for(int step=2,id=1;step<=N;step<<=1,++</pre>
26
         T wn=pow_mod(G,(P-1)>>id,P),wi=1,u,t;
27
         const int mh=step>>1;
         for(int i=0;i<mh;++i){</pre>
28
            for(int j=i;j<N;j+=step){</pre>
29
              u=out[j],t=wi*out[j+mh]%P;
30
31
              out[i]=u+t;
32
              out[j+mh]=u-t;
              if(out[j]>=P)out[j]-=P;
33
34
              if(out[j+mh]<0)out[j+mh]+=P;</pre>
35
36
            wi=wi*wn%P;
37
38
39
       if(is inv){
         for(int i=1;i<N/2;++i)swap(out[i],out[</pre>
40
              N-i]);
41
         T invn=pow_mod(N,P-2,P);
42
          for(int i=0;i<N;++i)out[i]=out[i]*invn</pre>
43
44
45 };
```

# 6.12 Simpson

```
double simpson(double a,double b){
   double c=a+(b-a)/2;
   return (F(a)+4*F(c)+F(b))*(b-a)/6;
}
double asr(double a,double b,double eps,
   double A){
   double c=a+(b-a)/2;
```

```
double L=simpson(a,c),R=simpson(c,b);
if( abs(L+R-A)<15*eps )
    return L+R+(L+R-A)/15.0;
return asr(a,c,eps/2,L)+asr(c,b,eps/2,R);
1
}
double asr(double a,double b,double eps){
return asr(a,b,eps,simpson(a,b));
}</pre>
```

### 6.13 外星模運算

 $1 / a[0]^{a[1]^{a[2]^{...}}$ 

bool is prime[maxn+5];

is prime[1]=1;//一不是質數

for(int i=1;i<=maxn;i++)euler[i]=i;</pre>

2 #define maxn 1000000

int euler[maxn+5];

void init euler(){

```
for(int i=2;i<=maxn;i++){</pre>
       if(!is_prime[i]){//是質數
         euler[i]--;
          for(int j=i<<1;j<=maxn;j+=i){</pre>
11
            is_prime[j]=1;
12
            euler[j]=euler[j]/i*(i-1);
13
14
15
16
17
   LL pow(LL a, LL b, LL mod) { //a^b%mod
     LL ans=1:
     for(;b;a=a*a%mod,b>>=1)
       if(b&1)ans=ans*a%mod;
21
22
     return ans;
23
   bool isless(LL *a,int n,int k){
     if(*a==1)return k>1;
     if(--n==0)return *a<k;</pre>
26
27
     int next=0;
     for(LL b=1;b<k;++next)</pre>
28
29
       b*=*a:
     return isless(a+1,n,next);
31 }
   LL high pow(LL *a, int n, LL mod){
     if(*a==1||--n==0)return *a%mod;
     int k=0,r=euler[mod];
     for(LL tma=1; tma!=pow(*a,k+r,mod);++k)
       tma=tma*(*a)%mod;
37
     if(isless(a+1,n,k))return pow(*a,high_pow(
          a+1,n,k),mod);
     int tmd=high_pow(a+1,n,r), t=(tmd-k+r)%r;
     return pow(*a,k+t,mod);
40
41 LL a[1000005];
   int main(){
     init euler();
     scanf("%d",&t);
46
     #define n 4
     while(t--){
48
       for(int i=0;i<n;++i)scanf("%lld",&a[i]);</pre>
       scanf("%d", \&mod);
49
       printf("%lld\n", high pow(a,n,mod));
```

# 6.14 數位統計

return 0;

52

53

```
1 | 11 d[65], dp[65][2];//up區間是不是完整
2 11 dfs(int p,bool is8,bool up){
    if(!p)return 1; // 回傳@是不是答案
    if(!up&&~dp[p][is8])return dp[p][is8];
    int mx = up?d[p]:9;//可以用的有那些
    11 ans=0;
    for(int i=0;i<=mx;++i){</pre>
      if( is8&&i==7 )continue;
      ans += dfs(p-1, i==8, up&&i==mx);
10
    if(!up)dp[p][is8]=ans;
11
    return ans;
12
13
14
   11 f(11 N){
    int k=0;
    while(N){ // 把數字先分解到陣列
      d[++k] = N%10;
      N/=10;
18
19
20
    return dfs(k,false,true);
```

### 6.15 質因數分解

```
1 | LL func(const LL n, const LL mod, const int c)
     return (LLmul(n,n,mod)+c+mod)%mod;
   LL pollorrho(const LL n, const int c) {//循
         環節長度
     LL a=1, b=1;
     a=func(a,n,c)%n;
     b=func(b,n,c)%n; b=func(b,n,c)%n;
     while(gcd(abs(a-b),n)==1) {
       a=func(a,n,c)%n;
       b=func(b,n,c)%n; b=func(b,n,c)%n;
12
     return gcd(abs(a-b),n);
15
   void prefactor(LL &n, vector<LL> &v) {
     for(int i=0;i<12;++i) {</pre>
18
       while(n%prime[i]==0) {
19
          v.push back(prime[i]);
20
          n/=prime[i];
21
22
23
25
   void smallfactor(LL n, vector<LL> &v) {
     if(n<MAXPRIME) {</pre>
       while(isp[(int)n]) {
```

```
v.push back(isp[(int)n]);
29
         n/=isp[(int)n];
30
31
       v.push back(n);
32
     } else {
       for(int i=0;i<primecnt&&prime[i]*prime[i</pre>
            1<=n:++i) {</pre>
         while(n%prime[i]==0) {
34
35
           v.push back(prime[i]);
           n/=prime[i];
36
37
38
39
       if(n!=1) v.push_back(n);
40
41
42
   void comfactor(const LL &n, vector<LL> &v) {
43
44
     if(n<1e9) {
45
       smallfactor(n,v);
46
       return;
47
48
     if(Isprime(n)) {
49
       v.push back(n);
50
       return;
51
52
     LL d:
53
     for(int c=3;;++c) {
       d = pollorrho(n,c);
55
       if(d!=n) break;
56
57
     comfactor(d,v);
58
     comfactor(n/d,v);
59
60
   void Factor(const LL &x, vector<LL> &v) {
     LL n = x;
     if(n==1) { puts("Factor 1"); return; }
     prefactor(n,v);
     if(n==1) return;
     comfactor(n,v);
67
     sort(v.begin(),v.end());
68
   void AllFactor(const LL &n, vector<LL> &v) {
     vector<LL> tmp;
     Factor(n,tmp);
     v.clear();
     v.push back(1);
     int len;
     LL now=1;
     for(int i=0;i<tmp.size();++i) {</pre>
       if(i==0 || tmp[i]!=tmp[i-1]) {
         len = v.size();
         now = 1;
       now*=tmp[i];
       for(int j=0;j<len;++j)</pre>
         v.push_back(v[j]*now);
85
```

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# 7 String

### 7.1 AC 自動機

```
1 template < char L='a', char R='z'>
   class ac automaton{
    struct joe{
       int next[R-L+1], fail, efl, ed, cnt dp, vis;
       joe():ed(0),cnt_dp(0),vis(0){
         for(int i=0;i<=R-L;++i)next[i]=0;</pre>
    };
   public:
    std::vector<joe> S;
    std::vector<int> q;
    int qs,qe,vt;
    ac_automaton():S(1),qs(0),qe(0),vt(0){}
14
    void clear(){
15
       q.clear();
16
       S.resize(1);
       for(int i=0;i<=R-L;++i)S[0].next[i]=0;</pre>
       S[0].cnt dp=S[0].vis=qs=qe=vt=0;
18
19
20
    void insert(const char *s){
21
22
       for(int i=0,id;s[i];++i){
         id=s[i]-L;
23
         if(!S[o].next[id]){
24
           S.push_back(joe());
           S[o].next[id]=S.size()-1;
28
         o=S[o].next[id];
29
30
       ++S[o].ed;
31
32
    void build_fail(){
33
       S[0].fail=S[0].efl=-1;
34
       a.clear();
35
       q.push_back(0);
36
       ++qe;
       while(qs!=qe){
38
         int pa=q[qs++],id,t;
39
         for(int i=0;i<=R-L;++i){</pre>
           t=S[pa].next[i];
           if(!t)continue;
           id=S[pa].fail;
           while(~id&&!S[id].next[i])id=S[id].
                fail;
           S[t].fail=~id?S[id].next[i]:0;
           S[t].efl=S[S[t].fail].ed?S[t].fail:S
                [S[t].fail].efl;
           q.push back(t);
           ++ae;
48
49
    /*DP出每個前綴在字串s出現的次數並傳回所有
          字串被s匹配成功的次數O(N+M)*/
52
    int match 0(const char *s){
       int ans=0,id,p=0,i;
54
       for(i=0;s[i];++i){
55
         id=s[i]-L;
         while(!S[p].next[id]&&p)p=S[p].fail;
```

```
if(!S[p].next[id])continue;
   p=S[p].next[id];
   ++S[p].cnt dp;/*匹配成功則它所有後綴都
        可以被匹配(DP計算)*/
 for(i=qe-1;i>=0;--i){
   ans+=S[q[i]].cnt dp*S[q[i]].ed;
   if(~S[q[i]].fail)S[S[q[i]].fail].
        cnt_dp+=S[q[i]].cnt_dp;
 return ans;
/*多串匹配走efL邊並傳回所有字串被s匹配成功
    的 次 數 O(N*M^1.5)*/
int match 1(const char *s)const{
 int ans=0,id,p=0,t;
 for(int i=0;s[i];++i){
   id=s[i]-L;
   while(!S[p].next[id]&&p)p=S[p].fail;
   if(!S[p].next[id])continue;
   p=S[p].next[id];
   if(S[p].ed)ans+=S[p].ed;
   for(t=S[p].efl;~t;t=S[t].efl){
     ans+=S[t].ed;/*因為都走efL邊所以保證
          匹配成功*/
 return ans;
/*枚舉(s的子字串nA)的所有相異字串各恰一次
    並傳回次數O(N*M^(1/3))*/
int match_2(const char *s){
 int ans=0,id,p=0,t;
 /*把戳記vt+=1,只要vt沒溢位,所有S[p].
      vis==vt就會變成false
  這種利用vt的方法可以0(1)歸零vis陣列*/
 for(int i=0;s[i];++i){
   id=s[i]-L;
   while(!S[p].next[id]&&p)p=S[p].fail;
   if(!S[p].next[id])continue;
   p=S[p].next[id];
   if(S[p].ed&&S[p].vis!=vt){
     S[p].vis=vt;
     ans+=S[p].ed;
   for(t=S[p].efl;~t&&S[t].vis!=vt;t=S[t
       ].efl){
     S[t].vis=vt;
     ans+=S[t].ed;/*因為都走efL邊所以保證
          匹配成功*/
 return ans;
/*把AC自動機變成真的自動機*/
void evolution(){
 for(qs=1;qs!=qe;){
   int p=q[qs++];
   for(int i=0;i<=R-L;++i)</pre>
     if(S[p].next[i]==0)S[p].next[i]=S[S[
         p].fail].next[i];
```

#### 7.2 hash

112 };

```
1 #define MAXN 1000000
 2 #define mod 1073676287
 3 /*mod 必須要是質數*/
 4 typedef long long T;
 5 char s[MAXN+5];
 6 T h[MAXN+5]; /*hash 陣列*/
 7 T h base[MAXN+5];/*h base[n]=(prime^n)%mod*/
 8 void hash init(int len,T prime){
    h base[0]=1:
     for(int i=1;i<=len;++i){</pre>
       h[i]=(h[i-1]*prime+s[i-1])%mod;
11
       h base[i]=(h base[i-1]*prime)%mod;
12
13
14 }
15 | T get_hash(int l,int r){/*閉區間寫法,設編號
        為 @ ~ len-1*/
     return (h[r+1]-(h[1]*h_base[r-1+1])%mod+
         mod)%mod:
17 }
```

#### 7.3 KMP

```
1 /*產生fail function*/
 void kmp fail(char *s,int len,int *fail){
     int id=-1;
     fail[0]=-1;
     for(int i=1;i<len;++i){</pre>
       while(~id&&s[id+1]!=s[i])id=fail[id];
       if(s[id+1]==s[i])++id;
       fail[i]=id;
 | 11 | /*以字串B匹配字串A,傳回匹配成功的數量(用B的
12 int kmp_match(char *A,int lenA,char *B,int
        lenB,int *fail){
     int id=-1.ans=0:
     for(int i=0;i<lenA;++i){</pre>
       while(~id&&B[id+1]!=A[i])id=fail[id];
       if(B[id+1]==A[i])++id;
16
       if(id==lenB-1){/*匹配成功*/
17
         ++ans, id=fail[id];
18
19
20
21
     return ans;
```

#### 7.4 manacher

#### 7.5 minimal string rotation

```
int min_string_rotation(const string &s){
    int n=s.size(),i=0,j=1,k=0;
    while(i<n&&j<n&&k<n){
        int t=s[(i+k)%n]-s[(j+k)%n];
        ++k;
        if(t){
        if(t>0)i+=k;
        else j+=k;
        if(i==j)++j;
        k=0;
    }
}
return min(i,j);//最小循環表示法起始位置
```

#### 7.6 reverseBWT

```
1 const int MAXN = 305, MAXC = 'Z';
  int ranks[MAXN], tots[MAXC], first[MAXC];
  void rankBWT(const string &bw){
    memset(ranks,0,sizeof(int)*bw.size());
    memset(tots,0,sizeof(tots);
    for(size t i=0;i<bw.size();++i)</pre>
      ranks[i] = tots[int(bw[i])]++;
  void firstCol(){
    memset(first,0,sizeof(first));
    int totc = 0;
11
    for(int c='A';c<='Z';++c){</pre>
13
      if(!tots[c]) continue;
14
      first[c] = totc;
15
       totc += tots[c];
16
17
  string reverseBwt(string bw,int begin){
    rankBWT(bw), firstCol();
20
    int i = begin; //原字串最後一個元素的位置
21
    string res;
22
    do{
23
      char c = bw[i];
      res = c + res:
      i = first[int(c)] + ranks[i];
    }while( i != begin );
    return res;
```

#### suffix array lcp 8.2 tnfshb017 2 sat 1 | struct dominator tree{ puts(""); }else puts("0"); static const int MAXN=5005; int n;// 1-base 66 return 0; 1 #define radix\_sort(x,y){\ vector<int> suc[MAXN],pre[MAXN]; 1 | #include < bits / stdc++.h> for(i=0;i<A;++i)c[i]=0;\</pre> int fa[MAXN],dfn[MAXN],id[MAXN],Time; 2 using namespace std; for(i=0;i<n;++i)c[x[y[i]]]++;\</pre> int semi[MAXN],idom[MAXN]; 3 #define MAXN 8001 for(i=1;i<A;++i)c[i]+=c[i-1];\</pre> int anc[MAXN], best[MAXN]; // disjoint set 4 #define MAXN2 MAXN\*4 for(i=n-1;~i;--i)sa[--c[x[y[i]]]]=y[i];\ vector<int> dom[MAXN];//dominator\_tree 5 #define n(X) ((X)+2\*N)void init(int n){ 6 vector<int> v[MAXN2], rv[MAXN2], vis t; 8.3 橋連涌分量 #define AC(r,a,b)\ 7 int N,M; 10 n=\_n; r[a]!=r[b]||a+k>=n||r[a+k]!=r[b+k]for(int i=1;i<=n;++i)suc[i].clear(),pre[</pre> void addedge(int s,int e){ 11 void suffix array(const char \*s,int n,int \* i].clear(); v[s].push back(e); sa,int \*rank,int \*tmp,int \*c){ 12 10 rv[e].push\_back(s); int A = 'z' + 1, i, k, id = 0;13 void add edge(int u,int v){ 11 } 1 #define N 1005 for(i=0;i<n;++i)rank[tmp[i]=i]=s[i];</pre> 14 suc[u].push back(v); 12 int scc[MAXN2]; struct edge{ 12 radix sort(rank,tmp); pre[v].push\_back(u); 13 bool vis[MAXN2]={false}; 15 int u,v; for(k=1;id<n-1;k<<=1){</pre> 13 16 void dfs(vector<int> \*uv,int n,int k=-1){ bool is\_bridge; for(id=0,i=n-k;i<n;++i)tmp[id++]=i;</pre> 14 17 void dfs(int u){ 15 vis[n]=true; edge(int u=0, int v=0):u(u),v(v), is bridge 15 for(i=0:i<n:++i)</pre> dfn[u]=++Time,id[Time]=u; for(int i=0;i<uv[n].size();++i)</pre> 18 16 $(0)\{\}$ if(sa[i]>=k)tmp[id++]=sa[i]-k; 16 19 for(auto v:suc[u]){ 17 if(!vis[uv[n][i]]) 17 radix sort(rank,tmp); 20 if(dfn[v])continue; dfs(uv,uv[n][i],k); 18 vector<edge> E; swap(rank,tmp); dfs(v),fa[dfn[v]]=dfn[u]; if(uv==v)vis\_t.push\_back(n); 18 21 19 vector<int> G[N];// 1-base 19 for(rank[sa[0]]=id=0,i=1;i<n;++i)</pre> 22 20 scc[n]=k; int low[N], vis[N], Time; 20 rank[sa[i]]=id+=AC(tmp,sa[i-1],sa[i]); 21 23 int bcc\_id[N],bridge\_cnt,bcc\_cnt;// 1-base void solve(){ 21 A=id+1; 24 int find(int x){ 22 int st[N],top;//BCC用 for(int i=1;i<=N;++i){</pre> 22 if(x==anc[x])return x; 23 25 void add edge(int u,int v){ 23 int y=find(anc[x]); 26 if(!vis[i])dfs(v,i); G[u].push\_back(E.size()); 27 if(semi[best[x]]>semi[best[anc[x]]])best 25 if(!vis[n(i)])dfs(v,n(i)); //h:高度數組 sa:後綴數組 rank:排名 E.emplace\_back(u,v); void suffix\_array\_lcp(const char \*s,int len, [x]=best[anc[x]]; 26 G[v].push back(E.size()); return anc[x]=y; 27 memset(vis,0,sizeof(vis)); int \*h,int \*sa,int \*rank){ 28 E.emplace back(v,u); 28 **int** c=0; for(int i=0;i<len;++i)rank[sa[i]]=i;</pre> 29 17 void tarjan(int r){ 29 for(int i=vis\_t.size()-1;i>=0;--i) 30 for(int i=0,k=0;i<len;++i){</pre> 18 | void dfs(int u,int re=-1){//u當前點,re為u連 31 Time=0: 30 if(!vis[vis\_t[i]]) if(rank[i]==0)continue; 接前一個點的邊 32 for(int t=1;t<=n;++t){</pre> 31 dfs(rv,vis\_t[i],c++); **if**(k)--k; int v; 32 } 30 while(s[i+k]==s[sa[rank[i]-1]+k])++k; 33 dfn[t]=idom[t]=0: //u=r或 是 u無 法 到 達 r 時 low[u]=vis[u]=++Time; 33 int main(){ h[rank[i]]=k; idom[id[u]]=0 st[top++]=u; int a,b; 32 dom[t].clear(); 34 34 for(int e:G[u]){ 22 h[0]=0;// h[k]=lcp(sa[k],sa[k-1]);scanf("%d%d",&N,&M); 33 anc[t]=best[t]=semi[t]=t; 35 v=E[e].v; 36 for(int i=1;i<=N;++i){</pre> 36 if(!vis[v]){ // (A or B)&(!A & !B) A^B 37 37 dfs(r); 25 dfs(v,e^1);//e^1反向邊 38 a=i\*2-1;38 for(int y=Time;y>=2;--y){ 26 low[u]=min(low[u],low[v]); b=i\*2; int x=fa[y],idy=id[y]; 39 39 27 if(vis[u]<low[v]){</pre> 7.8 Z for(auto z:pre[idy]){ 40 addedge(n(a),b); 40 28 E[e].is bridge=E[e^1].is bridge=1; 41 if(!(z=dfn[z]))continue; addedge(n(b),a); 29 ++bridge\_cnt; addedge(a,n(b)); 42 find(z); 30 addedge(b,n(a)); semi[y]=min(semi[y],semi[best[z]]); 43 1 void z\_alg(char \*s,int len,int \*z){ 31 }else if(vis[v]<vis[u]&&e!=re)</pre> 44 44 int 1=0, r=0; while(M--){ 32 low[u]=min(low[u], vis[v]); 45 dom[semi[y]].push\_back(y); z[0]=len; 33 scanf("%d%d",&a,&b); 46 anc[y]=x; for(int i=1;i<len;++i){</pre> a = a>0?a\*2-1:-a\*2;**if**(vis[u]==low[u]){//處理*BCC* for(auto z:dom[x]){ 47 z[i]=i>r?0:(i-1+z[i-1]< z[1]?z[i-1]:r-ib = b>0?b\*2-1:-b\*2;++bcc cnt;// 1-base find(z); +1); // A or B idom[z]=semi[best[z]]<x?best[z]:x;</pre> 36 do bcc id[v=st[--top]]=bcc cnt;//每個點 while(i+z[i]<len&&s[i+z[i]]==s[z[i]])++z</pre> 50 addedge(n(a),b); 所在的BCC [i]; 51 addedge(n(b),a); 51 dom[x].clear(); while(v!=u); if(i+z[i]-1>r)r=i+z[i]-1,l=i; 37 52 38 solve(); 53 for(int u=2;u<=Time;++u){</pre> 39 bool check=true; 54 if(idom[u]!=semi[u])idom[u]=idom[idom[ 54 void bcc init(int n){ for(int i=1;i<=2\*N;++i)</pre> Time=bcc\_cnt=bridge\_cnt=top=0; 551 dom[id[idom[u]]].push\_back(id[u]); **if**(scc[i]==scc[n(i)]) E.clear(); check=false; 56 for(int i=1;i<=n;++i){</pre> 43 Tarjan if(check){ 57 44 G[i].clear(); 59 printf("% $d \setminus n$ ",N); 58 | }dom; 45 vis[i]=bcc id[i]=0; for(int i=1;i<=2\*N;i+=2){</pre> 46 61 if(scc[i]>scc[i+2\*N]) putchar('+'); dominator tree

62

else putchar('-');

### 8.4 雙連通分量 & 割點

```
1 #define N 1005
vector<int> G[N];// 1-base
3 | vector<int> bcc[N];//存每塊雙連通分量的點
4 int low[N], vis[N], Time;
5 int bcc_id[N],bcc_cnt;// 1-base
6 bool is cut[N];//是否為割點
  int st[N],top;
   void dfs(int u,int pa=-1){//u當前點,pa父親
    int t, child=0;
    low[u]=vis[u]=++Time;
    st[top++]=u;
12
    for(int v:G[u]){
13
      if(!vis[v]){
14
        dfs(v,u),++child;
         low[u]=min(low[u],low[v]);
15
         if(vis[u]<=low[v]){</pre>
16
17
           is cut[u]=1;
           bcc[++bcc cnt].clear();
18
19
             bcc_id[t=st[--top]]=bcc_cnt;
20
21
             bcc[bcc cnt].push back(t);
           }while(t!=v);
22
23
           bcc id[u]=bcc cnt;
           bcc[bcc cnt].push back(u);
24
25
26
      }else if(vis[v]<vis[u]&&v!=pa)//反向邊
        low[u] = min(low[u], vis[v]);
27
    }//u是dfs樹的根要特判
28
29
    if(pa==-1&&child<2)is cut[u]=0;</pre>
30
   void bcc init(int n){
31
    Time=bcc cnt=top=0:
32
    for(int i=1;i<=n;++i){</pre>
33
34
      G[i].clear();
      is cut[i]=vis[i]=bcc id[i]=0;
35
36
37 }
```

# 9 Tree\_problem

### 9.1 HeavyLight

```
if(max son[u]==-1||siz[v]>siz[max son[u
           ]])max son[u]=v;
      siz[u]+=siz[v];
15
16
17
   void build link(int u,int top){
    link[u]=++cnt:
    link_top[u]=top;
20
    if(max son[u]==-1)return;
21
    build_link(max_son[u],top);
     for(auto v:G[u]){
      if(v==max son[u]||v==pa[u])continue;
      build_link(v,v);
25
26
27
  int find_lca(int a,int b){
    //求LCA,可以在過程中對區間進行處理
    int ta=link_top[a],tb=link_top[b];
    while(ta!=tb){
      if(dep[ta]<dep[tb]){</pre>
33
        swap(ta,tb);
34
        swap(a,b);
35
      //這裡可以對a所在的鏈做區間處理
36
      //區間為(Link[ta],Link[a])
37
38
      ta=link top[a=pa[ta]];
39
    //最後a,b會在同一條鏈,若a!=b還要在進行一
         次區間處理
    return dep[a]<dep[b]?a:b;</pre>
42 }
```

#### 9.2 LCA

```
1 | const int MAXN=100000; // 1-base
2 const int MLG=17; //Log2(MAXN)+1;
3 int pa[MLG+2][MAXN+5];
4 int dep[MAXN+5];
   vector<int> G[MAXN+5];
   void dfs(int x,int p=0){//dfs(root);
     pa[0][x]=p;
     for(int i=0;i<=MLG;++i)</pre>
       pa[i+1][x]=pa[i][pa[i][x]];
     for(auto &i:G[x]){
11
      if(i==p)continue;
12
       dep[i]=dep[x]+1;
13
       dfs(i,x);
14
15 }
   inline int jump(int x,int d){
     for(int i=0;i<=MLG;++i)</pre>
      if((d>>i)&1) x=pa[i][x];
19
     return x;
20 }
   inline int find_lca(int a,int b){
    if(dep[a]>dep[b])swap(a,b);
     b=jump(b,dep[b]-dep[a]);
     if(a==b)return a;
     for(int i=MLG;i>=0;--i){
      if(pa[i][a]!=pa[i][b]){
         a=pa[i][a];
27
         b=pa[i][b];
```

```
nd[x].ch[1]=last;
29
30
                                                51
                                                       up(x);
    return pa[0][a];
                                                52
                                                       last=x;
31
32
                                                53
                                                       x=nd[x].pa;
                                                54
                                                     return last;//access後splay tree的根
                                                56
        link cut tree
                                                   void access(int x,bool is=0){//is=0就是一般
                                                        的access
                                                     int last=0;
 1 | struct splay tree{
                                                     while(x){
    int ch[2],pa;//子節點跟父母
                                                       splay(x);
     bool rev; // 反轉的懶惰標記
                                                       if(is&&!nd[x].pa){
                                                         //printf("%d\n",max(nd[last].ma,nd[nd[
     splay tree():pa(0),rev(0){ch[0]=ch[1]=0;}
  };
                                                             x1.ch[1]].ma));
 6 vector<splay_tree> nd;
                                                63
 7 // 有的時候用vector會TLE,要注意
                                                64
                                                       nd[x].ch[1]=last;
                                                65
                                                       up(x);
 8 | // 這邊以node [0] 作為null 節點
                                                       last=x;
 9 bool isroot(int x){//判斷是否為這棵splay
                                                       x=nd[x].pa;
                                                67
       tree的根
                                                68
     return nd[nd[x].pa].ch[0]!=x&&nd[nd[x].pa
         ].ch[1]!=x;
                                                70
                                                   void query_edge(int u,int v){
11 }
                                                71
                                                     access(u);
12
  void down(int x){// 懶 惰 標 記 下 推
                                                     access(v,1);
                                                72
13
     if(nd[x].rev){
                                                73
      if(nd[x].ch[0])nd[nd[x].ch[0]].rev^=1;
                                                   void make root(int x){
                                                74
      if(nd[x].ch[1])nd[nd[x].ch[1]].rev^=1;
15
                                                75
                                                     access(x),splay(x);
16
      swap(nd[x].ch[0],nd[x].ch[1]);
                                                76
                                                     nd[x].rev^=1;
      nd[x].rev=0;
17
                                                77
18
                                                78
                                                   void make root(int x){
19 }
                                                     nd[access(x)].rev^=1;
   void push_down(int x){//所有祖先懶惰標記下推
                                                80
                                                     splay(x);
    if(!isroot(x))push down(nd[x].pa);
                                                81
22
     down(x);
                                                   void cut(int x,int y){
23
                                                     make root(x);
                                                     access(y);
24 | void up(int x){}//將子節點的資訊向上更新
                                                     splay(y);
                                                85
25 | void rotate(int x){//旋轉,會自行判斷轉的方
                                                     nd[y].ch[0]=0;
                                                     nd[x].pa=0;
     int y=nd[x].pa,z=nd[y].pa,d=(nd[y].ch[1]==
         x);
                                                89
                                                   void cut_parents(int x){
27
     nd[x].pa=z;
                                                     access(x);
     if(!isroot(y))nd[z].ch[nd[z].ch[1]==y]=x;
                                                     splay(x);
     nd[y].ch[d]=nd[x].ch[d^1];
                                                92
                                                     nd[nd[x].ch[0]].pa=0;
     nd[nd[y].ch[d]].pa=y;
30
                                                93
                                                     nd[x].ch[0]=0;
    nd[y].pa=x,nd[x].ch[d^1]=y;
31
                                                94
32
    up(y),up(x);
                                                   void link(int x,int y){
33
                                                     make root(x);
  void splay(int x){//將x伸展到splay tree的根
                                                     nd[x].pa=y;
    push down(x);
                                                98
     while(!isroot(x)){
                                                   int find root(int x){
      int y=nd[x].pa;
                                                     x=access(x);
      if(!isroot(y)){
                                                     while(nd[x].ch[0])x=nd[x].ch[0];
39
         int z=nd[y].pa;
                                                     splay(x);
                                               102
         if((nd[z].ch[0]==y)^(nd[y].ch[0]==x))
40
                                                     return x;
                                               103
             rotate(v);
                                               104
         else rotate(x);
                                               105 int query(int u, int v){
42
                                               106 // 傳回uv路徑splav tree的根結點
43
      rotate(x);
                                                   //這種寫法無法求LCA
44
                                               108
                                                     make root(u);
45
                                                     return access(v);
                                               109
  int access(int x){
```

110

int query lca(int u,int v){

int last=0;

splay(x);

while(x){

vis[i]=0;

```
112 // 假 設 求 鏈 上 點 權 的 總 和 , sum 是 子 樹 的 權 重 和
        data 是 節 點 的 權 重
                                                  13 }
                                                     void get dis(vector<int> &dis,int u,int pa,
     access(u);
                                                          int d){
     int lca=access(v);
                                                       dis.push_back(d);
115
     splay(u);
                                                  15
                                                   16
                                                        for(size t i=0;i<g[u].size();++i){</pre>
     if(u==lca){
116
                                                         int v=g[u][i].first,w=g[u][i].second;
       //return nd[lca].data+nd[nd[lca].ch[1]].
117
                                                         if(v!=pa&&!vis[v])get_dis(dis,v,u,d+w);
                                                  18
                                                   19
118
     }else{
                                                  20 }
       //return nd[lca].data+nd[nd[lca].ch[1]].
                                                  21 | vector < int > dis; // 這東西如果放在函數裡會TLE
            sum+nd[u].sum
                                                  22 int cal(int u,int d){
120
                                                       dis.clear();
121
   struct EDGE{
                                                       get dis(dis,u,-1,d);
     int a,b,w;
                                                       sort(dis.begin(),dis.end());
   }e[10005];
                                                  26
                                                       int l=0, r=dis.size()-1, res=0;
   vector<pair<int,int>> G[10005];
                                                         while(1<r&&dis[1]+dis[r]>k)--r;
127 //first表示子節點, second表示邊的編號
                                                         res+=r-(1++);
int pa[10005],edge_node[10005];
                                                   30
                                                       return res;
                                                  31
129 | //pa 是父母節點,暫存用的,edge_node 是每個編
                                                  32 }
        被存在哪個點裡面的陣列
                                                  33 pair<int,int> tree centroid(int u,int pa,
130 void bfs(int root){
                                                          const int sz){
   //在建構的時候把每個點都設成一個splay tree
                                                       size[u]=1;//找樹重心·second是重心
     queue<int > q;
                                                       pair<int,int> res(INT MAX,-1);
     for(int i=1;i<=n;++i)pa[i]=0;</pre>
133
                                                       int ma=0;
     q.push(root);
134
                                                        for(size_t i=0;i<g[u].size();++i){</pre>
135
     while(q.size()){
                                                         int v=g[u][i].first;
                                                  38
136
       int u=q.front();
                                                  39
                                                         if(v==pa||vis[v])continue;
137
       q.pop();
                                                         res=min(res,tree centroid(v,u,sz));
                                                  40
       for(auto P:G[u]){
138
                                                  41
                                                         size[u]+=size[v];
         int v=P.first;
139
                                                  42
                                                         ma=max(ma,size[v]);
140
         if(v!=pa[u]){
                                                  43
           pa[v]=u;
141
                                                  44
                                                       ma=max(ma,sz-size[u]);
142
           nd[v].pa=u;
                                                       return min(res,make_pair(ma,u));
143
           nd[v].data=e[P.second].w;
                                                  46
           edge node[P.second]=v;
144
                                                      int tree_DC(int u,int sz){
                                                  47
145
           up(v);
                                                       int center=tree centroid(u,-1,sz).second;
           q.push(v);
146
                                                       int ans=cal(center,0);
147
                                                  50
                                                       vis[center]=1;
148
                                                       for(size t i=0;i<g[center].size();++i){</pre>
149
                                                         int v=g[center][i].first,w=g[center][i].
150
    void change(int x,int b){
                                                  53
                                                         if(vis[v])continue;
152
     splay(x);
                                                  54
                                                         ans-=cal(v,w);
153
     //nd[x].data=b;
                                                  55
                                                         ans+=tree DC(v,size[v]);
154
     up(x);
                                                  56
155
                                                  57
                                                       return ans;
                                                  58
                                                  59
                                                     int main(){
                                                       while(scanf("%d%d",&n,&k),n||k){
                                                  60
   9.4 POJ tree
                                                  61
                                                         init();
                                                         for(int i=1;i<n;++i){</pre>
                                                  62
                                                  63
                                                           int u,v,w;
 1 | #include < bits / stdc++.h>
                                                            scanf("%d%d%d",&u,&v,&w);
                                                  64
 2 using namespace std;
                                                  65
                                                            g[u].push_back(make_pair(v,w));
 3 #define MAXN 10005
                                                  66
                                                           g[v].push back(make pair(u,w));
                                                  67
 5 vector<pair<int,int> >g[MAXN];
                                                  68
                                                         printf("%d\n",tree_DC(1,n));
 6 int size[MAXN];
                                                  69
 7 bool vis[MAXN];
                                                  70
                                                       return 0;
 8 inline void init(){
                                                  71 }
     for(int i=0;i<=n;++i){</pre>
       g[i].clear();
```

# 10 default

### 10.1 debug

#### 10.2 ext

```
1 | #include < bits / extc++.h>
2 #include < ext/pd_ds/assoc_container.hpp>
3 #include<ext/pd ds/tree policy.hpp>
4 using namespace __gnu_cxx;
5 using namespace __gnu_pbds;
6 template<typename T>
 vsing pbds_set = tree<T,null_type,less<T>,
       rb tree tag,
       tree_order_statistics_node_update>;
 8 template<typename T, typename U>
9 using pbds map = tree<T,U,less<T>,
       rb_tree_tag,
       tree_order_statistics_node_update>;
using heap=__gnu_pbds::priority_queue<int>;
11 //s.find_by_order(1);//0 base
12 //s.order_of_key(1);
```

#### 10.3 IncStack

```
14     rl.rlim_cur=ks;
15     res=setrlimit(RLIMIT_STACK,&rl);
16     }
17 }
```

#### 10.4 input

```
inline int read(){
   int x=0; bool f=0; char c=getchar();
   while(ch<'0'||'9'<ch)f|=ch=='-',ch=getchar();

while('0'<=ch&&ch<='9')x=x*10-'0'+ch,ch=getchar();
   return f?-x:x;

// #!/bin/bash
// g++ -std=c++11 -02 -Wall -Wextra -Wno-unused-result -DDEBUG $1 && ./a.out
/-fsanitize=address -fsanitize=undefined -fsanitize=return</pre>
```

# 11 language

#### 11.1 CNF

```
1 | #define MAXN 55
  struct CNF{
    int s,x,y;//s->xy \mid s->x, if y==-1
    int cost;
    CNF(){}
    CNF(int s, int x, int y, int c):s(s), x(x), y(y)
         ),cost(c){}
7 };
s int state;//規則數量
9 | map < char , int > rule ; // 每個字元對應到的規則
        小寫字母為終端字符
  vector<CNF> cnf;
  void init(){
    state=0;
13
    rule.clear();
14
    cnf.clear();
15
  void add to cnf(char s,const string &p,int
       cost){
    //加入一個s -> 的文法,代價為cost
    if(rule.find(s)==rule.end())rule[s]=state
    for(auto c:p)if(rule.find(c)==rule.end())
         rule[c]=state++;
    if(p.size()==1){
      cnf.push back(CNF(rule[s],rule[p[0]],-1,
           cost));
22
    }else{
23
      int left=rule[s];
      int sz=p.size();
^{24}
      for(int i=0;i<sz-2;++i){</pre>
```

```
cnf.push back(CNF(left,rule[p[i]],
             state,0));
         left=state++;
       cnf.push_back(CNF(left,rule[p[sz-2]],
29
           rule[p[sz-1]],cost));
30
31
32 vector<long long> dp[MAXN][MAXN];
33 | vector<bool> neg_INF[MAXN][MAXN];//如果花費
        是負的可能會有無限小的情形
34 void relax(int l,int r,const CNF &c,long
       long cost,bool neg_c=0){
    if(!neg_INF[1][r][c.s]&&(neg_INF[1][r][c.x
          ]||cost<dp[1][r][c.s])){
       if(neg_c||neg_INF[1][r][c.x]){
         dp[1][r][c.s]=0;
         neg INF[1][r][c.s]=true;
       }else dp[1][r][c.s]=cost;
   void bellman(int l,int r,int n){
    for(int k=1;k<=state;++k)</pre>
      for(auto c:cnf)
         if(c.y==-1)relax(1,r,c,dp[1][r][c.x]+c
             .cost,k==n);
   void cyk(const vector<int> &tok){
    for(int i=0;i<(int)tok.size();++i){</pre>
       for(int j=0;j<(int)tok.size();++j){</pre>
         dp[i][j]=vector<long long>(state+1,
             INT MAX);
         neg_INF[i][j]=vector<bool>(state+1,
             false);
       dp[i][i][tok[i]]=0;
       bellman(i,i,tok.size());
    for(int r=1;r<(int)tok.size();++r){</pre>
       for(int l=r-1;l>=0;--1){
         for(int k=1;k<r;++k)</pre>
           for(auto c:cnf)
             if(~c.y)relax(1,r,c,dp[1][k][c.x]+
                  dp[k+1][r][c.v]+c.cost);
         bellman(l,r,tok.size());
62
```

# 12 other

### 12.1 WhatDay

# 12.2 上下最大正方形

```
1 void solve(int n,int a[],int b[]){// 1-base
    int ans=0;
     deque<int>da,db;
     for(int l=1,r=1;r<=n;++r){</pre>
       while(da.size()&&a[da.back()]>=a[r]){
         da.pop back();
       da.push back(r);
       while(db.size()&&b[db.back()]>=b[r]){
         db.pop back();
       db.push back(r):
       for(int d=a[da.front()]+b[db.front()];r-
         if(da.front()==1)da.pop_front();
15
         if(db.front()==1)db.pop front();
         if(da.size()&&db.size()){
           d=a[da.front()]+b[db.front()];
      ans=max(ans,r-l+1);
    printf("%d\n",ans);
```

# 12.3 最大矩形

```
1 | LL max rectangle(vector<int> s){
     stack<pair<int,int > > st;
     st.push(make_pair(-1,0));
     s.push back(0);
     LL ans=0;
     for(size_t i=0;i<s.size();++i){</pre>
       int h=s[i];
       pair<int,int > now=make_pair(h,i);
       while(h<st.top().first){</pre>
         now=st.top();
11
         st.pop();
         ans=max(ans,(LL)(i-now.second)*now.
              first):
       if(h>st.top().first){
14
         st.push(make_pair(h,now.second));
15
16
17
18
    return ans;
```

# 13 zformula

### 13.1 formula

# 13.1.1 Pick 公式

給定頂點坐標均是整點的簡單多邊形·面積 = 內部格點數 + 邊上格點數/2-1

#### 13.1.2 圖論

- 1. 對於平面圖  $\cdot F = E V + C + 1 \cdot C$  是連通分量 數
- 2. 對於平面圖  $\cdot E \leq 3V 6$
- 3. 對於連通圖 G·最大獨立點集的大小設為 I(G)·最大匹配大小設為 M(G)·最小點覆蓋設為 Cv(G)·最小邊覆蓋設為 Ce(G)。對於任意連通圖:
  - (a) I(G) + Cv(G) = |V|
  - (b) M(G) + Ce(G) = |V|
- 4. 對於連通二分圖:
  - (a) I(G) = Cv(G)
  - (b) M(G) = Ce(G)
- 5. 最大權閉合圖:
  - (a)  $C(u,v) = \infty, (u,v) \in E$
  - (b)  $C(S, v) = W_v, W_v > 0$
  - (c)  $C(v,T) = -W_v, W_v > 0$
  - (d) ans= $\sum_{W_v>0} W_v flow(S, T)$
- 6. 最大密度子圖:
  - (a)  $\vec{x} \max\left(\frac{W_e + W_v}{|V'|}\right), e \in E', v \in V'$
  - (b)  $U = \sum_{v \in V} 2W_v + \sum_{e \in E} W_e$
  - (c)  $C(u, v) = W_{(u,v)}, (u, v) \in E$  · 雙向邊
  - (d)  $C(S, v) = U, v \in V$
  - (e)  $D_u = \sum_{(u,v) \in E} W_{(u,v)}$
  - (f)  $C(v,T) = U + 2g D_v 2W_v, v \in V$
  - (g) 二分搜 g:  $l=0, r=U, eps=1/n^2$ if  $((U\times|V|-flow(S,T))/2>0)$  l=midelse r=mid
  - (h) ans= $min\_cut(S,T)$
  - (i) |E| = 0 要特殊判斷
- 7. 弦圖:
  - (a) 點數大於 3 的環都要有一條弦
  - (b) 完美消除序列從後往前依次給每個點染色,給 每個點染上可以染的最小顏色
  - (c) 最大團大小 = 色數
  - (d) 最大獨立集: 完美消除序列從前往後能選就選
  - (e) 最小團覆蓋: 最大獨立集的點和他延伸的邊構成
  - f) 區間圖是弦圖
  - (g) 區間圖的完美消除序列: 將區間按造又端點由 小到大排序
  - (h) 區間圖染色: 用線段樹做

### 13.1.3 dinic 特殊圖複雜度

- 1. 單位流: $O\left(\min\left(V^{3/2},E^{1/2}\right)E\right)$
- 2. 二分圖: $O\left(V^{1/2}E\right)$

#### 13.1.4 0-1 分數規劃

 $x_i = \{0,1\} \cdot x_i$  可能會有其他限制  $\cdot$  求  $max\left(\sum_{i=1}^{N} S_i x_i\right)$ 

- 1.  $D(i,g) = B_i g \times C_i$
- 2.  $f(g) = \sum D(i, g)x_i$
- 3. f(g) = 0 時 g 為最佳解 f(g) < 0 沒有意義
- 4. 因為 f(g) 單調可以二分搜 g
- 5. 或用 Dinkelbach 通常比較快

```
1| binary search(){
    while(r-l>eps){
     g=(1+r)/2;
     for(i:所有元素)D[i]=B[i]-g*C[i];//D(i,g)
     找出一組合法x[i]使f(g)最大;
     if(f(g)>0) l=g;
     else r=g;
    Ans = r;
  Dinkelbach(){
    g=任意狀態(通常設為0);
14
      for(i:所有元素)D[i]=B[i]-g*C[i];//D(i,g)
      找出一組合法x[i]使f(g)最大;
     p=0,q=0;
      for(i: 所有元素)
      if(x[i])p+=B[i],q+=C[i];
     g=p/q;//更新解,注意q=0的情況
    }while(abs(Ans-g)>EPS);
    return Ans;
```

# 13.1.5 學長公式

- 1.  $\sum_{d|n} \phi(n) = n$
- 2.  $g(n) = \sum_{d|n} f(d) = f(n) = \sum_{d|n} \mu(d) \times g(n/d)$
- 3. Harmonic series  $H_n = \ln(n) + \gamma + 1/(2n) 1/(12n^2) + 1/(120n^4)$
- 4.  $\gamma = 0.5772156649015328606065120900824024310421$
- 5. 格雷碼 =  $n \oplus (n >> 1)$
- 6.  $SG(A+B) = SG(A) \oplus SG(B)$
- 7. 選轉矩陣  $M(\theta) = \begin{pmatrix} \cos\theta & -\sin\theta \\ \sin\theta & \cos\theta \end{pmatrix}$

#### 13.1.6 基本數論

- 1.  $\sum_{d|n} \mu(n) = [n == 1]$
- 2.  $g(m) = \sum_{d|m} f(d) \Leftrightarrow f(m) = \sum_{d|m} \mu(d) \times g(m/d)$
- 3.  $\sum_{i=1}^{n} \sum_{j=1}^{m} \underline{G} [ \underline{y} ] \underline{z} ] = \sum_{i=1}^{m} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor \left\lfloor \frac{m}{d} \right\rfloor$
- 4.  $\sum_{i=1}^{n} \sum_{j=1}^{n} lcm(i,j) = n \sum_{d|n} d \times \phi(d)$

#### 13.1.7 排組公式

- 1. k 卡特蘭  $\frac{C_n^{kn}}{n(k-1)+1} \cdot C_m^n = \frac{n!}{m!(n-m)!}$ 2.  $H(n,m) \cong x_1 + x_2 \dots + x_n = k, num = C_k^{n+k-1}$
- 3. Stirling number of  $2^{nd}$ , n 人分 k 組方法數目
  - (a) S(0,0) = S(n,n) = 1
  - (b) S(n,0) = 0
  - (c) S(n,k) = kS(n-1,k) + S(n-1,k-1)
- 4. Bell number, n 人分任意多組方法數目

  - $\begin{array}{ll} \text{(a)} & B_0 = 1 \\ \text{(b)} & B_n = \sum_{i=0}^n S(n,i) \\ \text{(c)} & B_{n+1} = \sum_{k=0}^n C_k^n B_k \\ \text{(d)} & B_{p+n} \equiv B_n + B_{n+1} mod p \text{, p is prime} \\ \text{(e)} & B_p^{m} +_n \equiv m B_n + B_{n+1} mod p \text{, p is prime} \\ \end{array}$
  - (f) From  $B_0: 1, 1, 2, 5, 15, 52$ ,
  - 203, 877, 4140, 21147, 115975
- 5. Derangement, 錯排, 沒有人在自己位置上
  - (a)  $D_n = n!(1 \frac{1}{1!} + \frac{1}{2!} \frac{1}{3!} \dots + (-1)^n \frac{1}{n!})$ (b)  $D_n = (n-1)(D_{n-1} + D_{n-2}), D_0 =$  $1, D_1 = 0$
  - (c) From  $D_0: 1, 0, 1, 2, 9, 44$ , 265, 1854, 14833, 133496
- 6. Binomial Equality

  - (a)  $\sum_{k} \binom{r}{m+k} \binom{s}{n-k} = \binom{r+s}{m+n}$ (b)  $\sum_{k} \binom{l}{l+k} \binom{s}{n+k} = \binom{l+s}{l-m+n}$ (c)  $\sum_{k} \binom{l}{l+k} \binom{s+k}{n} \binom{s-1}{l-1} = (-1)^{l+m} \binom{s-m}{n-l}$ (d)  $\sum_{k \le l} \binom{l-k}{m} \binom{s}{k-n} (-1)^k = (-1)^{l+m} \binom{s-m-1}{l-n-m}$ (e)  $\sum_{0 \le k \le l} \binom{l-k}{m} \binom{n}{n} = \binom{l+q+1}{m+n+1}$ (f)  $\binom{r}{l} = (-1)^{k} \binom{l+k}{l-r-1}$

  - (f)  $\binom{r}{k} = (-1)^k \binom{k-r-1}{k}$

  - (g)  $\binom{k}{n}\binom{m}{k} = \binom{r}{k}\binom{r-k}{m-k}$ (h)  $\sum_{k \le n} \binom{r+k}{k} = \binom{r+n+1}{n}$ (i)  $\sum_{0 \le k \le n} \binom{k}{m} = \binom{m+1}{m+1}$ (j)  $\sum_{k \le m} \binom{m+r}{k} x^k y^k$  $\sum_{k \le m}^{-} {\binom{-r}{k}} (-x)^k (x+y)^{m-k}$

## 13.1.8 冪次, 冪次和

1.  $a^b \% P = a^{b\% \varphi(p) + \varphi(p)}, b > \varphi(p)$ 2.  $1^3 + 2^3 + 3^3 + \ldots + n^3 = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4}$ 2.  $1^{5} + 2^{c} + 3^{5} + \dots + n^{5} = \frac{n}{4} + \frac{n}{2} + \frac{n}{4}$ 3.  $1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n}{5} + \frac{n^{4}}{2} + \frac{n^{3}}{3} - \frac{n}{30}$ 4.  $1^{5} + 2^{5} + 3^{5} + \dots + n^{5} = \frac{n^{6}}{6} + \frac{n^{5}}{2} + \frac{5n^{4}}{12} - \frac{n^{2}}{12}$ 5.  $0^{k} + 1^{k} + 2^{k} + \dots + n^{k} = P(k), P(k) = \frac{(n+1)^{k+1} - \sum_{k=0}^{k-1} C_{k}^{k+1} P(i)}{k+1}, P(0) = n+1$ 6.  $\sum_{k=0}^{m-1} k^{n} = \frac{1}{n+1} \sum_{k=0}^{n} C_{k}^{n+1} B_{k} m^{n+1-k}$ 7.  $\sum_{j=0}^{m} C_j^{m+1} B_j = 0, B_0 = 1$ 8. 除了  $B_1 = -1/2$ ,剩下的奇數項都是 0 9.  $B_2 = 1/6, B_4 = -1/30, B_6 = 1/42, B_8 = {}^{18}$  $-1/30, B_{10} = 5/66, B_{12} = -691/2730, B_{14} = {}^{19}$  $7/6, B_{16} = -3617/510, B_{18}$  $43867/798, B_{20} = -174611/330,$ 

- 13.1.9 Burnside's lemma
  - 1.  $|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$

  - 3. G 表示有幾種轉法, $X^g$  表示在那種轉法下,有幾種 是會保持對稱的 $\cdot t$  是顏色數 $\cdot c(q)$  是循環節不動的
  - 4. 正立方體塗三顏色,轉 0 有 36 個元素不變,轉 90 有 6 種·每種有 3<sup>3</sup> 不變·180 有 3 × 3<sup>4</sup>· 120(角) 有  $8 \times 3^2 \cdot 180(邊)$  有  $6 \times 3^3 \cdot$  全部  $\frac{1}{24} \left( 3^{6} + 6 \times 3^{3} + 3 \times 3^{4} + 8 \times 3^{2} + 6 \times 3^{3} \right) =$

#### 13.1.10 Count on a tree

- 1. Rooted tree:  $s_{n+1} = \frac{1}{n} \sum_{i=1}^{n} (i \times a_i \times {8 \choose 9})$ ;  $\sum_{j=1}^{\lfloor n/i \rfloor} a_{n+1-i \times j})$
- 2. Unrooted tree:

  - (a) Odd: $a_n \sum_{i=1}^{n/2} a_i a_{n-i}$ (b) Even: $Odd + \frac{1}{2} a_{n/2} (a_{n/2} + 1)$
- 3. Spanning Tree
  - (a) 完全圖  $n^n 2$
  - (b) 一般圖 (Kirchhoff's theorem)M[i][i] = $degree(\hat{V}_i), M[i][j] = -1, \text{if have } E(i, j), 0$ if no edge. delete any one row and col in A, ans = det(A)

#### 13.2 java

#### 13.2.1 文件操作

```
1 import java.io.*;
   2 import java.util.*;
   3 import java.math.*;
= 4 import java.text.*;
   6 public class Main{
       public static void main(String args[]){
            throws FileNotFoundException,
            IOException
         Scanner sc = new Scanner(new FileReader(
              "a.in"));
         PrintWriter pw = new PrintWriter(new
             FileWriter("a.out"));
         n=sc.nextInt()://读入下一个INT
         m=sc.nextInt();
         for(ci=1; ci<=c; ++ci){</pre>
           pw.println("Case #"+ci+": easy for
               output");
         pw.close();// 关闭流并释放,这个很重要
              否则是没有输出的
```

sc.close();// 关闭流并释放

#### 13.2.2 优先队列

```
1 | PriorityQueue queue = new PriorityQueue( 1,
      new Comparator(){
    public int compare( Point a, Point b ){
    if( a.x < b.x || a.x == b.x && a.y < b.y )
    else if( a.x == b.x && a.y == b.y )
     return 0;
    else return 1;
```

#### 13.2.3 Map

```
1 | Map map = new HashMap();
2 map.put("sa","dd");
3 String str = map.get("sa").toString;
 for(Object obj : map.keySet()){
   Object value = map.get(obj );
```

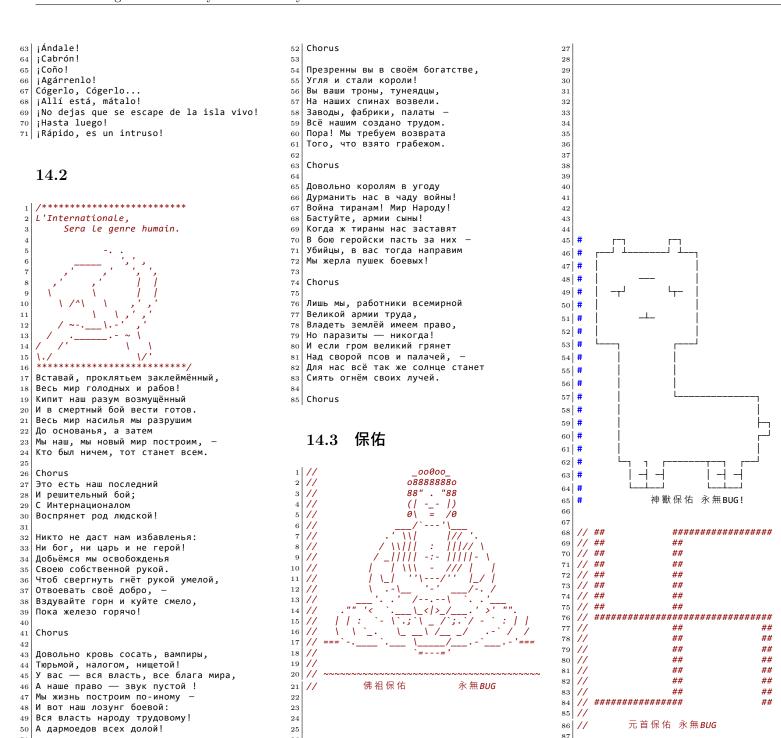
#### 13.2.4 sort

```
1 | static class cmp implements Comparator{
    public int compare(Object o1,Object o2){
    BigInteger b1=(BigInteger)o1;
    BigInteger b2=(BigInteger)o2;
    return b1.compareTo(b2);
  public static void main(String[] args)
       throws IOException{
    Scanner cin = new Scanner(System.in);
    n=cin.nextInt();
    BigInteger[] seg = new BigInteger[n];
    for (int i=0;i<n;i++)</pre>
    seg[i]=cin.nextBigInteger();
15
    Arrays.sort(seg, new cmp());
```

### 14

### 14.1 ganadoQuote

```
1 ¡Allí está!
  ¡Un forastero!
  ¡Agarrenlo!
  ¡Os voy a romper a pedazos!
  ¡Te voy a hacer picadillo!
  ¡Te vov a matar!
  ¡Míralo, está herido!
  ¡Sos cerdo!
  ¿Dónde estás?
11
  ¡Detrás de tí, imbécil!
   ¡No dejes que se escape!
  ¡Basta, hijo de puta!
14
  Lord Saddler...
16
   ¡Mátalo!
  ¡Allí está!
17
18 Morir es vivir.
19 Sííííí, ¡Quiero matar!
20 Muere, muere, muere....
21 Cerebros, cerebros, cerebros...
22 Cógedlo, cógedlo, cógedlo...
23 Lord Saddler...
24
  Dieciséis.
  ¡Va por él!
  ¡Muérete!
27
  ¡Cógelo!
28
  ¡Te voy a matar!
29
  ¡Bloqueale el paso!
31 ¡Te cogí!
32 ¡No dejes que se escape!
  ¿Qué carajo estás haciendo aquí? ¡Lárgate,
       cabrón!
  Hay un rumor de que hay un extranjero entre
       nosotros.
36 Nuestro jefe se encargará de la rata.
  Su "Las Plagas" es mucho mejor que la
       nuestra.
38 Tienes razón, es un hombre.
39 Usa los músculos.
40 Se vuelve loco!
41 ¡Hey, acá!
42 ¡Por aquí!
43 ¡El Gigante!
44 ¡Del Lago!
45 ¡Cógelo!
46 ¡Cógenlo!
47 ¡Allí!
48 ¡Rápido!
49 ¡Empieza a rezar!
50 :Mátenlos!
51 ¡Te voy a romper en pedazos!
52 ¡La campana!
53 Ya es hora de rezar.
54 Tenemos que irnos.
55 ¡Maldita sea, mierda!
56 ¡Ya es hora de aplastar!
58 Puedes correr, pero no te puedes esconder!
59 ¡Sos cerdo!
60 ¡Está en la trampa!
61 Ah, que madre!
62 ¡Vámonos!
```





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