

Project 2: RISC-V Simulator

1. Description

The goal of this project is to write a simulator program that simulates the RISC-V microarchitecture (single cycle execution). You need to model the basic operations of the RISC-V microarchitecture, such as fetching instructions from the instruction memory, decoding instructions, executing the proper arithmetic operations, reading data from the register file or the data memory, and determining the address of the next instructions. Your simulator will read an executable file (e.g., runme.hex) that contains several instructions represented in hexadecimal.

The simulator should have two modes: debug and run mode. In the debug mode, the simulator will show the clock cycles, PC, and register values after executing an instruction every cycle. In the run mode, your simulator will run the given executable file and show the clock cycles, PC, and register values at the end of the simulation.

You can download a zip file (P2_codes.zip) that includes skeleton code (riscv_sim.c) and several executable files (addi_test.hex, add_test.hex, jal_test.hex, jalr_test.hex, beq_test.hex, sd_ld_test.hex, and runme.hex) from the project section of the LMS. You can use the skeleton code to write your simulator program. The files named “(*_test.hex)” will be useful for testing your simulator to ensure that it can execute the individual instruction correctly.

Please upload the “**riscv_sim.c**” file and a “**report**” to the LMS. In the report, you need to explain your code briefly.

2. Instruction you must simulate

add, addi, jal, jalr, ld, sd, beq

3. How to run

```
./riscv_sim ./runme.hex 0 // debug mode
```

```
./riscv_sim ./runme.hex 1 // run mode
```

4. Expected output of your simulator

a. Debug mode

* All numbers are decimal number

```
$. /riscv_sim runme.hex 0
```

```
-----  
Clock cycles = 1  
PC           = ?
```

```
x0    = 0  
x1    = 0  
x2    = ?  
x3    = 0  
x4    = ?  
x5    = 0  
x6    = 0  
x7    = 0  
x8    = ?  
x9    = 0  
x10   = ?  
x11   = 0  
x12   = 0  
x13   = 0  
x14   = 0  
x15   = 0  
x16   = 0  
x17   = 0  
x18   = 0  
x19   = 0  
x20   = 0  
x21   = 0  
x22   = 0  
x23   = 0  
x24   = 0  
x25   = 0  
x26   = 0  
x27   = 0  
x28   = 0  
x29   = ?  
x30   = 0  
x31   = ?
```

```
-----  
Clock cycles = 2  
PC           = ?
```

```
x0    = 0  
x1    = 0  
x2    = ?  
x3    = 0  
x4    = ?  
x5    = 0  
x6    = 0  
x7    = 0  
x8    = ?  
x9    = 0  
x10   = ?
```

```
x11 = 0
x12 = 0
x13 = 0
x14 = 0
x15 = 0
x16 = 0
x17 = 0
x18 = 0
x19 = 0
x20 = 0
x21 = 0
x22 = 0
x23 = 0
x24 = 0
x25 = 0
x26 = 0
x27 = 0
x28 = 0
x29 = ?
x30 = 0
x31 = ?
```

```
-----
Clock cycles = 3
PC           = ?
```

```
x0  = 0
x1  = ?
x2  = ?
x3  = 0
x4  = ?
x5  = 0
x6  = 0
x7  = 0
x8  = ?
x9  = 0
x10 = ?
x11 = 0
x12 = 0
x13 = 0
x14 = 0
x15 = 0
x16 = 0
x17 = 0
x18 = 0
x19 = 0
x20 = 0
x21 = 0
x22 = 0
x23 = 0
x24 = 0
x25 = 0
x26 = 0
x27 = 0
x28 = 0
x29 = ?
x30 = 0
x31 = ?
```

```
..... •
```

b. run mode

* clock cycle is decimal number

```
$/riscv_sim runme.hex 1
```

```
Clock cycles = ?
```

```
PC           = ?
```

```
x0    = 0
```

```
x1    = 12
```

```
x2    = 800
```

```
x3    = 0
```

```
x4    = 0
```

```
x5    = 0
```

```
x6    = 45
```

```
x7    = 0
```

```
x8    = 0
```

```
x9    = 10
```

```
x10   = 55
```

```
x11   = 0
```

```
x12   = 0
```

```
x13   = 0
```

```
x14   = 0
```

```
x15   = 0
```

```
x16   = 0
```

```
x17   = 0
```

```
x18   = 0
```

```
x19   = 0
```

```
x20   = 0
```

```
x21   = 0
```

```
x22   = 0
```

```
x23   = 0
```

```
x24   = 0
```

```
x25   = 0
```

```
x26   = 0
```

```
x27   = 0
```

```
x28   = 0
```

```
x29   = 0
```

```
x30   = 0
```

```
x31   = 0
```

5. Testing

You can find several executable files (“*_test.hex”) that you can use to test your simulator.

- Expected output of your simulator for each executable file
 - a. addi_test.hex

```
$/riscv_sim addi_test.hex 1

Clock cycles = 32
PC           = 128

x0    = 0
x1    = 10
x2    = 10
x3    = 10
x4    = 10
x5    = 10
x6    = 10
x7    = 10
x8    = 10
x9    = 10
x10   = 10
x11   = 10
x12   = 10
x13   = 10
x14   = 10
x15   = 10
x16   = 10
x17   = 10
x18   = 10
x19   = 10
x20   = 10
x21   = 10
x22   = 10
x23   = 10
x24   = 10
x25   = 10
x26   = 10
x27   = 10
x28   = 10
x29   = 10
x30   = 10
x31   = 10
```

b. add_test.hex

```
$/riscv_sim add_test.hex 1
```

```
Clock cycles = 4
```

```
PC          = 16
```

```
x0    = 0
x1    = 0
x2    = 0
x3    = 0
x4    = 0
x5    = 0
x6    = 0
x7    = 0
x8    = 0
x9    = 10
x10   = 0
x11   = 0
x12   = 0
x13   = 0
x14   = 0
x15   = 0
x16   = 10
x17   = 20
x18   = 30
x19   = 0
x20   = 0
x21   = 0
x22   = 0
x23   = 0
x24   = 0
x25   = 0
x26   = 0
x27   = 0
x28   = 0
x29   = 0
x30   = 0
x31   = 0
```

c. jal_test.hex

```
$/riscv_sim jal_test.hex 1
```

```
Clock cycles = 3
```

```
PC          = 16
```

```
x0    = 0
x1    = 4
x2    = 0
x3    = 0
x4    = 0
x5    = 0
x6    = 0
x7    = 0
x8    = 0
x9    = 10
x10   = 0
x11   = 0
x12   = 0
x13   = 0
x14   = 0
x15   = 0
x16   = 0
x17   = 20
x18   = 0
x19   = 0
x20   = 0
x21   = 0
x22   = 0
x23   = 0
x24   = 0
x25   = 0
x26   = 0
x27   = 0
x28   = 0
x29   = 0
x30   = 0
x31   = 0
```

d. jalr_test.hex

```
$/riscv_sim jalr_test.hex 1
```

```
Clock cycles = 7
```

```
PC          = 36
```

```
x0    = 0
x1    = 0
x2    = 0
x3    = 0
x4    = 0
x5    = 0
x6    = 0
x7    = 0
x8    = 0
x9    = 10
x10   = 0
x11   = 0
x12   = 0
x13   = 0
x14   = 0
x15   = 0
x16   = 10
x17   = 20
x18   = 0
x19   = 20
x20   = 0
x21   = 0
x22   = 0
x23   = 0
x24   = 0
x25   = 0
x26   = 0
x27   = 0
x28   = 0
x29   = 0
x30   = 0
x31   = 4
```


e. sd_ld_test.hex

```
$/riscv_sim sd_ld_test.hex 1
```

```
Clock cycles = 5
```

```
PC          = 20
```

```
x0    = 0
x1    = 0
x2    = 0
x3    = 0
x4    = 30
x5    = 0
x6    = 0
x7    = 0
x8    = 0
x9    = 10
x10   = 0
x11   = 0
x12   = 0
x13   = 0
x14   = 0
x15   = 0
x16   = 30
x17   = 0
x18   = 0
x19   = 0
x20   = 0
x21   = 0
x22   = 0
x23   = 0
x24   = 0
x25   = 0
x26   = 0
x27   = 0
x28   = 0
x29   = 0
x30   = 0
x31   = 0
```

f. beq_test.hex

```
$/riscv_sim beq_test.hex 1
```

```
Clock cycles = 4
```

```
PC          = 20
```

```
x0    = 0
x1    = 0
x2    = 0
x3    = 0
x4    = 0
x5    = 0
x6    = 0
x7    = 0
x8    = 0
x9    = 10
x10   = 0
x11   = 0
x12   = 0
x13   = 0
x14   = 0
x15   = 0
x16   = 30
x17   = 30
x18   = 0
x19   = 0
x20   = 0
x21   = 0
x22   = 0
x23   = 0
x24   = 0
x25   = 0
x26   = 0
x27   = 0
x28   = 0
x29   = 0
x30   = 0
x31   = 0
```

6. Evaluation [total: 100 points]

- a. **Unit test** [45 points] : test your simulator to ensure that it can execute individual instruction (add, addi, jal, jalr, ld, sd, beq) correctly.

- b. **Full test** [55 points] : test your simulator to ensure that it can correctly execute the given executable file. **We will use the “runme.hex” file for the full test.**

Late Day Policy

No late submission is allowed.

Plagiarism

No plagiarism will be tolerated. If the instructor determines that there are substantial similarities between your code and others, you will be given 0 points for all project assignments. Code similarity will be checked with a plagiarism detector tool.

Appendix

Category	Instruction	Example	Meaning	Comments
Arithmetic	Add	add x5, x6, x7	$x5 = x6 + x7$	Three register operands; add
	Subtract	sub x5, x6, x7	$x5 = x6 - x7$	Three register operands; subtract
	Add immediate	addi x5, x6, 20	$x5 = x6 + 20$	Used to add constants
Data transfer	Load doubleword	ld x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Doubleword from memory to register
	Store doubleword	sd x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Doubleword from register to memory
	Load word	lw x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Word from memory to register
	Load word, unsigned	lwu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Unsigned word from memory to register
	Store word	sw x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Word from register to memory
	Load halfword	lh x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Halfword from memory to register
	Load halfword, unsigned	lhu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Unsigned halfword from memory to register
	Store halfword	sh x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Halfword from register to memory
	Load byte	lb x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Byte from memory to register
	Load byte, unsigned	lbu x5, 40(x6)	$x5 = \text{Memory}[x6 + 40]$	Byte unsigned from memory to register
	Store byte	sb x5, 40(x6)	$\text{Memory}[x6 + 40] = x5$	Byte from register to memory
	Load reserved	lr.d x5, (x6)	$x5 = \text{Memory}[x6]$	Load; 1st half of atomic swap
	Store conditional	sc.d x7, x5, (x6)	$\text{Memory}[x6] = x5; x7 = 0/1$	Store; 2nd half of atomic swap
Logical	Load upper immediate	lui x5, 0x12345	$x5 = 0x12345000$	Loads 20-bit constant shifted left 12 bits
	And	and x5, x6, x7	$x5 = x6 \& x7$	Three reg. operands; bit-by-bit AND
	Inclusive or	or x5, x6, x8	$x5 = x6 \mid x8$	Three reg. operands; bit-by-bit OR
	Exclusive or	xor x5, x6, x9	$x5 = x6 \wedge x9$	Three reg. operands; bit-by-bit XOR
	And immediate	andi x5, x6, 20	$x5 = x6 \& 20$	Bit-by-bit AND reg. with constant
	Inclusive or immediate	ori x5, x6, 20	$x5 = x6 \mid 20$	Bit-by-bit OR reg. with constant
Shift	Exclusive or immediate	xori x5, x6, 20	$x5 = x6 \wedge 20$	Bit-by-bit XOR reg. with constant
	Shift left logical	sll x5, x6, x7	$x5 = x6 \ll x7$	Shift left by register
	Shift right logical	srl x5, x6, x7	$x5 = x6 \gg x7$	Shift right by register
	Shift right arithmetic	sra x5, x6, x7	$x5 = x6 \gg x7$	Arithmetic shift right by register
	Shift left logical immediate	slli x5, x6, 3	$x5 = x6 \ll 3$	Shift left by immediate
	Shift right logical immediate	srl_i x5, x6, 3	$x5 = x6 \gg 3$	Shift right by immediate
Shift	Shift right arithmetic immediate	srai x5, x6, 3	$x5 = x6 \gg 3$	Arithmetic shift right by immediate
Conditional branch	Branch if equal	beq x5, x6, 100	if $(x5 == x6)$ go to PC+100	PC-relative branch if registers equal
	Branch if not equal	bne x5, x6, 100	if $(x5 != x6)$ go to PC+100	PC-relative branch if registers not equal
	Branch if less than	blt x5, x6, 100	if $(x5 < x6)$ go to PC+100	PC-relative branch if registers less
	Branch if greater or equal	bge x5, x6, 100	if $(x5 \geq x6)$ go to PC+100	PC-relative branch if registers greater or equal
	Branch if less, unsigned	bltu x5, x6, 100	if $(x5 < x6)$ go to PC+100	PC-relative branch if registers less, unsigned
	Branch if greater or equal, unsigned	bgeu x5, x6, 100	if $(x5 \geq x6)$ go to PC+100	PC-relative branch if registers greater or equal, unsigned
Unconditional branch	Jump and link	jal x1, 100	$x1 = \text{PC}+4$; go to PC+100	PC-relative procedure call
	Jump and link register	jalr x1, 100(x5)	$x1 = \text{PC}+4$; go to $x5+100$	Procedure return; indirect call

Name (Field size)	7 bits	5 bits	5 bits	3 bits	5 bits	7 bits	Comments
R-type	funct7	rs2	rs1	funct3	rd	opcode	Arithmetic instruction format
I-type	immediate[11:0]		rs1	funct3	rd	opcode	Loads & immediate arithmetic
S-type	immed[11:5]	rs2	rs1	funct3	immed[4:0]	opcode	Stores
SB-type	immed[12,10:5]	rs2	rs1	funct3	immed[4:1,11]	opcode	Conditional branch format
UJ-type	immediate[20,10:1,11,19:12]				rd	opcode	Unconditional jump format
U-type	immediate[31:12]				rd	opcode	Upper immediate format

Format	Instruction	Opcode	Funct3	Funct6/7
R-type	add	0110011	000	0000000
	sub	0110011	000	0100000
	sll	0110011	001	0000000
	xor	0110011	100	0000000
	srl	0110011	101	0000000
	sra	0110011	101	0100000
	or	0110011	110	0000000
	and	0110011	111	0000000
I-type	lrd	0110011	011	0001000
	scd	0110011	011	0001100
	lb	0000011	000	n.a.
	lh	0000011	001	n.a.
	lw	0000011	010	n.a.
	ld	0000011	011	n.a.
	lbu	0000011	100	n.a.
	lhu	0000011	101	n.a.
	lwu	0000011	110	n.a.
	addi	0010011	000	n.a.
	slli	0010011	001	0000000
	xori	0010011	100	n.a.
	srlr	0010011	101	0000000
	srai	0010011	101	0100000
	ori	0010011	110	n.a.
	andi	0010011	111	n.a.
	jalr	1100111	000	n.a.

Format	Instruction	Opcode	Funct3	Funct6/7
S-type	sb	0100011	000	n.a.
	sh	0100011	001	n.a.
	sw	0100011	010	n.a.
	sd	0100011	011	n.a.
SB-type	beq	1100011	000	n.a.
	bne	1100011	001	n.a.
	blt	1100011	100	n.a.
	bge	1100011	101	n.a.
	bltu	1100011	110	n.a.
	bgeu	1100011	111	n.a.
U-type	lui	0110111	n.a.	n.a.
UJ-type	jal	1101111	n.a.	n.a.