Problem 5: |15 marks|

Design an iterative version of the quick-sort algorithm using queues. (a) Give a high-level description of your algorithm and explain why it works in a few sentences. (b) Write pseudo code for your iterative algorithm (assume that a queue implementation is already available). (c) In your pseudocode, write down the invariants corresponding to the pivot element.

The algorithm consists of two parts & Partitioning and asort. (a) => Partitioning : " We randomly select a pivot, total supplied supplied with last element of array. Make two pointers i=0 j=n-2 while it; , we check if arr (i) I pivot and arr &ij & pivot. · If not we swap arr [i] & arr [i]. At 'the end, we swap

arr (i) and pivot and return (i) (pivot index) · This ensures all elements betore pivot are smaller and after pivot (b) Partition[1, u, am] are bisser.

19 Sort 8 . ve have two pointers I and a. Initially L=0, u=n-1 · we also have pi (intravioble).

· Pi = partition (L, u, arr.)

· Now we engaged in the order: L, pl-1, pi+1, cl rwhile queue is not empty use por first two eveneutr

assign L= dequequel, u= dequequell and run the loop ogain Thus we keep on doing the same thing for smaller and smaller sub arrays.

queruclinty Q 1=0 u= n-1

pi=partition(1, u, on)

Q.eng(1), Q.eng(pi+1), Q.eng(pi+1), Q.eng(u)

Pushile (Q=!= empty)

l= 9. dog() u=0.deg()

else. Pi= partition(1, u, arr) Q-eng((1),Q.eng(pi-1),Q.eng(pi+1),Qeng(u) End while.

(c) The proof element, after partitioning is always at the final and correct place required after sorting

p= random(1,u)

swap (orrei), amuj)

swap (arr Cpj, arr [u])

it arreiszarries Itt

it an (j) = arrist j -swap (arr (i), arr (j))

After partitioning, elements left to pivot are smaller than it, and those to right are greater thanit.