

# COL215\_2018 Minor 1

SAKSHAM DHULL

TOTAL POINTS

**7 / 10**

## QUESTION 1

### 1 Question 1 2 / 2

- ✓ + 2 pts Correct proof
- + 0 pts Incorrect proof
- 1 pts Using derived equalities
- 1 pts Missing steps
- + 0 pts Unattempted

## QUESTION 2

### 2 Question 2 0 / 3

- + 2 pts Correct
- + 1 pts Partially correct
- + 0.5 pts Click here to replace this description.
- ✓ + 0 pts Incorrect
- 💬 Placement of literals in K map is wrong. SO entire solution is wrong

## QUESTION 3

### 3 Question 3 2 / 2

- ✓ + 2 pts 8 or more values of f are correct out of 10
- + 1.5 pts 6 or 7 values of f are correct out of 10
- + 1 pts 4 or 5 values of f are correct out of 10
- + 0.5 pts 2 or 3 values of f are correct out of 10
- + 0 pts Less than 2 values of f are correct out of 10

## QUESTION 4

### 4 Question 4 3 / 3

- + 0 pts 3 or less are correct
- + 1 pts 4 to 7 are correct
- + 2 pts 8 to 11 are correct
- ✓ + 3 pts 12 and more are correct

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1. Using basic properties of Boolean Algebra operators '+' and '.' and the definitions of inverse and identity elements, prove the following equality. Do not use any derived properties.

$$(x + y) \cdot (y + z) \cdot (x' + z) = (x + y) \cdot (x' + z)$$

LHS

[2]

$$(x+y) \cdot (y+z) \cdot (x'+z)$$

$$= (x \cdot y + \cancel{x} \cdot y + x \cdot z + y \cdot z) \cdot (x' + z)$$

$$= \cancel{x} \cdot y \cdot 0 + \bar{x} \cdot y + 0 + \bar{x} \cdot y \cdot z + x y z + y z + x z \cdot \cancel{y}$$

$$= \bar{x} y (1+z) + (x+1) y z + x z \cdot \cancel{y}$$

$$= \bar{x} y + x y z + x z$$

$$= 0 + \bar{x} y + (x+y) z$$

$$= x \cdot \bar{x} + \bar{x} \cdot y + (x+y) z$$

$$= (x+y) \cdot (\bar{x}+z)$$



$x_3 x_4$

$x_1 x_2$

1	0	1	0
1	0	1	1
0	1	1	0
0	0	0	0

$x_1 \bar{x}_2 \bar{x}_4 + \bar{x}_1 x_2 x_4$

$00 \times 0$

1	0	1	0
1	0	0	0
0	1	1	0
0	0	0	1

$x_1 x_2 (x_3)$

$(x_3 \bar{x}_4 + \bar{x}_3 x_4) \cdot x_1 \bar{x}_2$

0	0	0	0	1
0	0	0	0	0
0	0	1	0	1
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0
0	0	0	0	0

$10 \times 1$

$x_1 \bar{x}_2 x_4 + \bar{x}_1 x_2 x_4 + \bar{x}_1 x_2 x_3 + \bar{x}_1 x_2 x_4$

$(x_3 + \bar{x}_4) \cdot \bar{x}_1 \bar{x}_2$

$x_1 x_2 (x_3 + x_4)$

$\bar{x}_1 x_2 x_3$

$\bar{x}_1 x_2 (x_3 + \bar{x}_4)$

$0.5 + 0.6$

$x_1 x_2 x_3 + x_1 \bar{x}_4$

$x_1 x_2$

00	01	10	11
1	0	1	0
1	0	1	1
0	1	1	1
0	0	0	0

00	01	10	11
00	01	10	11
01	10	11	00
10	11	00	01

$\bar{x}_1 \bar{x}_3 \bar{x}_4 + x_3 \bar{x}_4 (x_1 + x_2)$

$\bar{x}_3 x_4 \cdot x_1 + x_1 x_2 x_3 x_4$



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2. Find the minimum cost SoP implementations for the following two functions considered separately (only expressions are required, not circuits). If you were allowed to share product terms among the two, what combined minimum cost implementation is possible. Note that in combined minimum implementation, the individual implementations need not be minimum. Here the cost is measured as "literal count", that is, total number of literals present in the expression, where each appearance of a variable in true or complemented form is counted as a literal. For example, the expression  $x_1 \cdot x_2' + x_1 \cdot x_3$  has a literal count of 4.

$$f(x_1, x_2, x_3, x_4) = \sum m(0, 2, 4, 6, 7, 9) + D(10, 11)$$

$$g(x_1, x_2, x_3, x_4) = \sum m(2, 4, 9, 10, 15) + D(0, 13, 14)$$

f.

$x_1$	$x_2$	$x_3$	$x_4$	f	g
0	0	0	0	1	-
0	0	0	1	0	0
0	0	1	0	1	1
0	0	1	1	0	0
0	1	0	0	1	1
0	1	0	1	0	0
0	1	1	0	1	0
0	1	1	1	1	0
0	1	0	0	0	0
1	0	0	1	1	1
1	0	1	0	-	1
1	0	1	1	-	0
1	1	0	0	0	0
1	1	0	1	0	-
1	1	1	0	0	-
1	1	1	1	0	1

Karnaugh Maps

		$x_3 x_4$			
		00	01	10	11
$x_1 x_2$	00	1	0	1	0
	01	1	0	1	1
	10	0	1	-	-
	11	0	0	0	0

[3]

	00	01	10	11
00	-	0	1	0
01	1	0	0	0
10	0	1	1	0
11	0	-	-	1

$$\text{min cost SOP}(f) = x_1 \cdot \bar{x}_2 \cdot x_4 + \bar{x}_1 \cdot \bar{x}_4 + \bar{x}_1 \cdot x_2 \cdot x_3 + \bar{x}_2 \cdot x_1 \cdot x_3$$

$$\text{min cost SOP}(g) = \bar{x}_1 \cdot \bar{x}_2 \cdot \bar{x}_4 + \bar{x}_1 \cdot x_3 \cdot \bar{x}_4 + x_1 \cdot \bar{x}_3 \cdot x_4 + \bar{x}_2 \cdot x_3 \cdot \bar{x}_4 + x_1 x_2 x_3 x_4$$

with sharing.

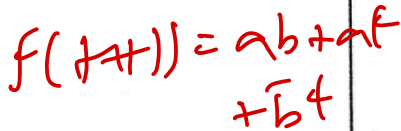
$$\text{min cost}(f) = \bar{x}_1 \bar{x}_2 \bar{x}_4 + x_1 \bar{x}_2 x_3 \bar{x}_4 + x_1 \bar{x}_2 \bar{x}_3 x_4 + \bar{x}_1 x_2 x_3 + \bar{x}_1 \bar{x}_3 \bar{x}_4$$

$$\text{min cost}(g) = \bar{x}_1 \bar{x}_2 \bar{x}_4 + x_1 \bar{x}_2 x_3 \bar{x}_4 + x_1 \bar{x}_2 \bar{x}_3 x_4 + x_1 x_2 x_3 x_4 + \bar{x}_1 \bar{x}_3 \bar{x}_4$$



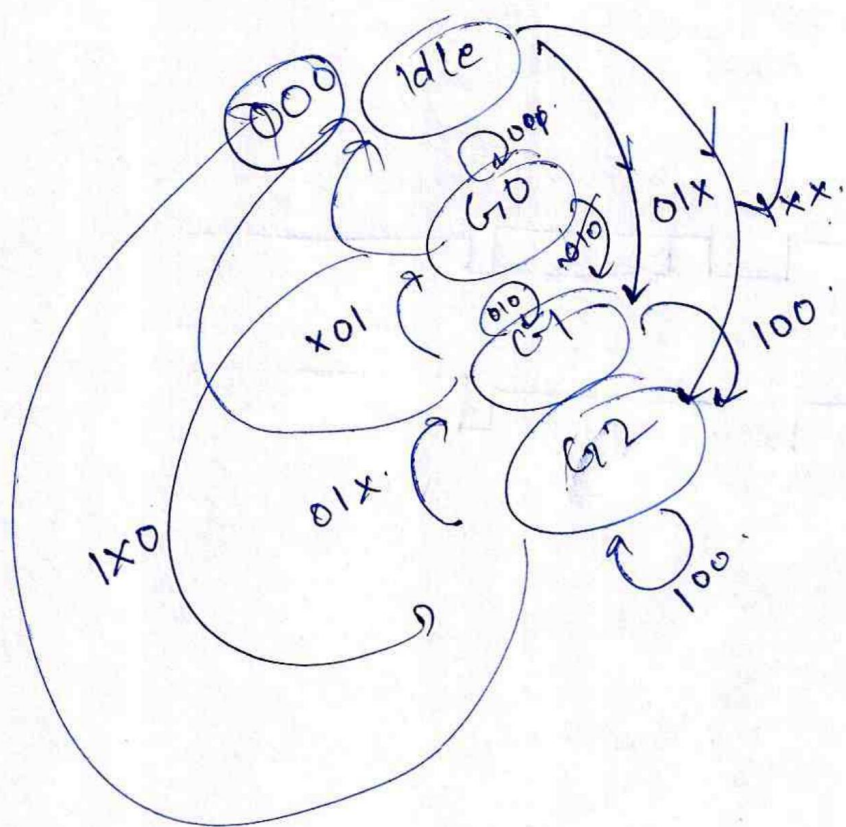
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- $$f(t+n) = a^b + a^f + b^f$$



1, 0, 1





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4. A modified version of the arbiter discussed in class is given below in terms of its Entity and architecture description in VHDL. Draw its state transition diagram. Do not show multiple transitions between any pair of present state and next state.

ENTITY arbiter IS

PORT (clk : IN bit;

req : IN bit\_vector (2 DOWNTO 0);

ack : OUT bit\_vector (2 DOWNTO 0)

);

END arbiter;

ARCHITECTURE FSM OF arbiter IS

TYPE state\_type IS (Idle, G0, G1, G2);

SIGNAL state : state\_type;

BEGIN

PROCESS (clk)

BEGIN

IF clk'EVENT AND clk = '1' THEN ✓

IF req = "000" THEN state &lt;= Idle; ✓

ELSIF req = "100" THEN state &lt;= G2; ✓

ELSIF req = "010" THEN state &lt;= G1; ✓

ELSIF req = "001" THEN state &lt;= G0; ✓

ELSE

CASE state IS

WHEN Idle =&gt;

IF req (2) = '1' THEN state &lt;= G2; ✓

ELSIF req (1) = '1' THEN state &lt;= G1; ✓

END IF;

WHEN G2 =&gt;

IF req (2 DOWNTO 1) = "01" THEN state &lt;= G1; ✓

END IF;

WHEN G1 =&gt;

IF req (1 DOWNTO 0) = "01" THEN state &lt;= G0; ✓

END IF;

WHEN G0 =&gt;

IF req (2) = '1' AND req (0) = '0' THEN state &lt;= G2; ✓

END IF;

END CASE;

END IF;

END IF;

END PROCESS;

WITH state SELECT

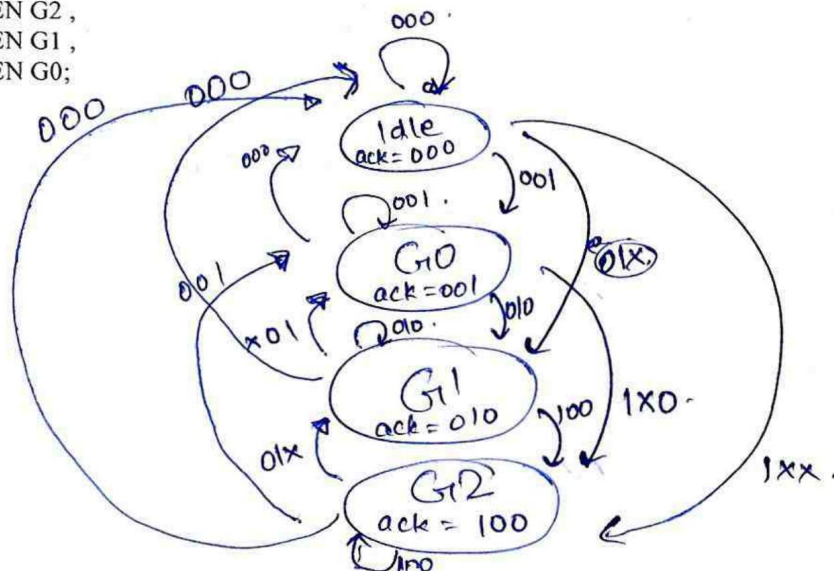
ack &lt;= "000" WHEN Idle ,

"100" WHEN G2 ,

"010" WHEN G1 ,

"001" WHEN G0;

END FSM;



[3]



$$0 + \bar{x}y + xz + yz,$$