

## COL351 Holi2023: Tutorial Problem Set 10

1. Recall the **CircuitSAT** problem discussed in class. The input to the problem is a boolean circuit over variables  $x_1, \dots, x_n$ , where each input of each gate is a variable, a constant (**true** or **false**), or the output of an “earlier” gate. One gate is designated as the output gate, and we are required to determine whether there exists a boolean assignment to the variables that makes the output of the circuit **true** (a.k.a. satisfying assignment). The **CircuitSATSearch** problem has the same input as **CircuitSAT**, but we are required to output a satisfying assignment if such an assignment exists, and otherwise output “NO”. Design a Cook-reduction from **CircuitSATSearch** to **CircuitSAT**.
2. A *dominating set* of a graph is a subset  $D$  of vertices such that every vertex that is not in  $D$  is adjacent to some vertex in  $D$ . The input to the **DominatingSet** problem is a graph  $G$  and an integer  $k$ , and the output is whether  $G$  has a dominating set of size at most  $k$ . Prove that **DominatingSet** is in NP, and that **VertexCover** is Karp-reducible to **DominatingSet**.
3. The input to the **DominatingSetSearch** problem is a graph  $G$  and an integer  $k$ , and the output is a dominating set of  $G$  having at most  $k$  vertices, if such a set exists, and “NO” otherwise. Design a Cook-reduction from **DominatingSetSearch** to **DominatingSet**.
4. Let  $U$  be a finite set and, let  $\mathcal{S} \subseteq 2^U$  be a collection of subsets of  $U$ . We say that a set  $H \subseteq U$  is a *hitting set* of  $\mathcal{S}$  if for all  $A \in \mathcal{S}$ , we have  $H \cap A \neq \emptyset$ , or in other words,  $H$  intersects every set in  $\mathcal{S}$ . In the **HittingSet** problem, we are given  $U$ ,  $\mathcal{S} \subseteq 2^U$ , and an integer  $k$  as input. We are required to determine whether there is a hitting set of  $\mathcal{S}$  having size at most  $k$ . Prove that **HittingSet** is in NP, and that **VertexCover** is Karp-reducible to **HittingSet**.
5. In the **HittingSetOpt** problem, we are given a finite set  $U$  and a collection  $\mathcal{S} \subseteq 2^U$  of subsets of  $U$ , and we are required to output a minimum cardinality hitting set of  $\mathcal{S}$ . Design a Cook-reduction from **HittingSetOpt** to **HittingSet**.
6. Design a Karp-reduction from **HittingSet** to **CircuitSAT**.
7. Recall that in the **GraphIso** problem, we are given two undirected graphs and we wish to determine whether they are isomorphic. Design a Karp-reduction from **GraphIso** to **CircuitSAT**.