

# COL 351 Lecture 35 2023/04/12

Topic : More Polynomial-time Verification,  
coNP

Announcements:

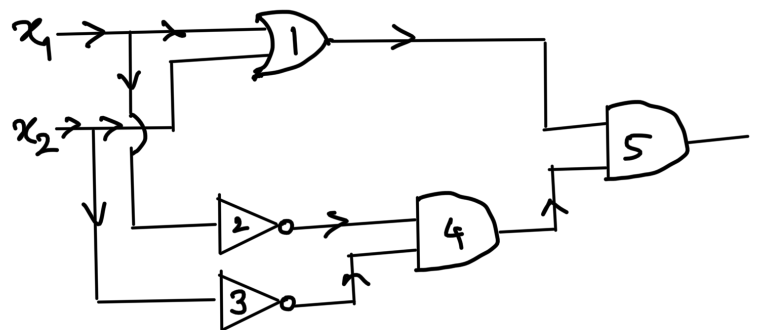
1. Pre-quiz tutorial today 19:30 - onward Bharti-501
2. Quiz 4 tomorrow 12:00 - 13:00, LH 114,
3. Next lecture tomorrow 19:30 - 20:30 Bharti - 501.

## CIRCUIT-SAT

Input: A combinational  
boolean circuit over  $n$   
variables, containing  $m$  gates.

Output: Does there exist  
a boolean input (a.k.a. a  
satisfying assignment)  
to the circuit that makes its  
output TRUE?

$\square$ : AND  $\cup$ : OR  $\triangleright$ : NOT



# vars = 2

# gates = 5

gate1:  $x_1$  OR  $x_2$

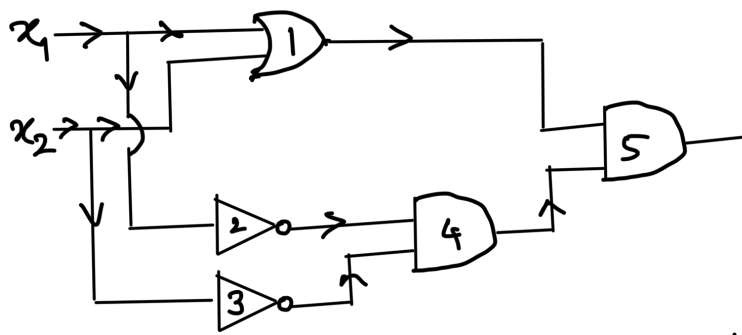
gate2: NOT  $x_1$

gate3: NOT  $x_2$

gate4: gate2 AND gate3

gate5: gate4 AND gate1.

OUTPUT: gate5.



Prover

Verifier

"The circuit has a satisfying assignment"

"Prove it"

"Here is an assignment:

$x_1 = \dots \quad x_2 \dots \dots \quad x_n = \dots$

Runs through the description of the circuit, and checks whether the output is TRUE.

GRAPHISO:

Input: Two graphs,  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ .

Output: Are  $G_1$  and  $G_2$  isomorphic to each other?

i.e.  $\exists$ ? a bijective function  $h: V_1 \rightarrow V_2$

such that  $\forall u_i, v_i \in V_1$

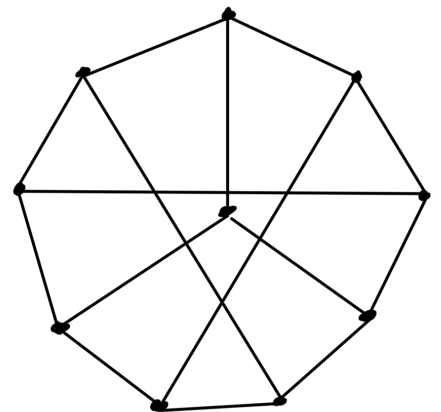
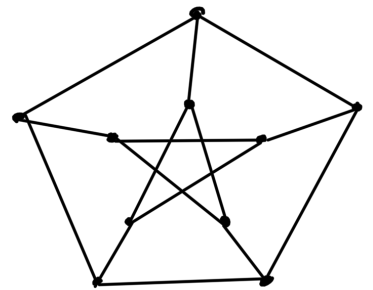
$\{u_i, v_i\} \in E_1 \iff \{h(u_i), h(v_i)\} \in E_2$

Verifier  $(G_1, G_2, h) \quad (h: V_1 \rightarrow V_2)$

- Check whether  $h$  is bijective.

- For each pair  $(u_i, v_i) \in V_1 \times V_1$ :

Check whether  $\{u_i, v_i\} \in E_1 \iff \{h(u_i), h(v_i)\} \in E_2$ .



PRIME:

Input: +ve integer  $m$  (in binary)

Output: Is  $m$  a prime?

Is  $\text{PRIME} \in \text{NP}$ ? — not clear.

$\overline{\text{PRIME}} \in \text{NP}$ .

Verifier for  $\overline{\text{PRIME}}$ : On input  $m, r$ :

If  $r$  divides  $m$  and  $1 < r < m$ , accept

Else reject.

for  $i = 2$  to  $\lfloor \sqrt{m} \rfloor$ :

if  $i$  divides  $m$ :

Return NO

Return YES

Not a polynomial-time algorithm.

Fact [AKS]:  $\text{PRIME} \in \text{P}$   
 $(\therefore \overline{\text{PRIME}} \in \text{P})$  } Well outside the scope of COL351

Claim: If  $Q \in \text{P}$  then  $\overline{Q} \in \text{P}$ .

Definition: A decision problem  $Q$  is said to be in the complexity class  $\text{coNP}$  if  $\overline{Q} \in \text{NP}$ .

Claim:  $\text{P} \subseteq \text{NP} \cap \text{coNP}$ .

Claim: If  $Q \in \text{NP} \cap \text{coNP}$ , then  $\overline{Q} \in \text{NP} \cap \text{coNP}$ .

Open question: Which of the following is true?

