

COL 351 Quiz 6B

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TOTAL POINTS

10 / 10

QUESTION 1

1 Q1 10 / 10

$n = k$ in G

+ 0 pts No submission / Incorrect

✓ + 2 pts We first show that this problem is in NP.

The verifier takes as input a graph G and a subset S of size k vertices and checks if S is an independent set or not

✓ + 2 pts We now reduce the independent set problem to the LargeIndSet problem. Let (G, k) be an input to the independent set problem.

We map it to an input (G', k') of the LargeIndSet problem. Let n be the number of vertices in G . We obtain G' as follows: we add a set W of n new vertices to G and there are no edges incident with any vertex in W .

✓ + 1 pts The parameter $k' = k + n$.

✓ + 1 pts Note that if n' is the number of vertices in G' , then

$n' = 2n$ and so, $k' \geq n'/2$

✓ + 2 pts Argue that G has an independent set of size k iff G' has a large independent set of size k' .

Suppose G has an independent S of size k . Then $S \cup W$ is an independent set in G' of size $n + k = k'$

✓ + 2 pts Conversely, suppose G' has an independent set S of size k' . Now at most n of the vertices in S can belong to W . The remaining vertices $S \setminus W$ belong to G and form an independent set in G . Thus, we have an independent set of size at least $k' -$

Quiz 6B (COL 351)

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Give precise arguments. You can use the fact that the following problems are NP-complete: 3-Satisfiability, Clique, Vertex Cover, Independent Set, Subset Sum.

Given an undirected graph G on n vertices, we say that a subset S of vertices in G is a large independent set if S is an independent set AND $|S| \geq n/2$. Prove that the following problem, called LargeIndSet, is NP-complete: given a graph G and a value $k \geq n/2$, does G have a large independent set of size at least k ? Recall that a subset S of vertices is said to be an independent set if there is no edge between any pair of vertices in S .

We will first show that LargeIndSet is in NP.

Given any subset S which is a solution / LargeIndSet, we ~~show~~ can check in polynomial time that $|S| \geq n/2$ AND $|S|$ is independent (check every pair of vertices in $O(|S|^2)$ time)

We will now show that

IS \leq_p LargeIndSet.

G, k

(G, k)

Decision version.

G', k'

(any general instance of IS (G, k) may be converted to a specific instance of LargeIndSet, i.e. it is reducible to LargeIndSet).

If $k \leq n/2$

Construct G' by adding $n - 2k$ disconnected vertices to G , and choose $k' = n - k$. independent $n = |G|$

Else if $k > n/2$,

Choose $G' = G$ and $k' = k$. $G' = G \cup S'$ where S' has $n - 2k$ vertices

Note that this conversion is done in polynomial time (adding $n - 2k$ vertices \rightarrow polynomial)

We will now show that corresponding to every solved instance of $IS(G, k)$, \exists solved instance of $Large\ IndSet(G', k')$ i.e. k vice-versa i.e.

$$IS(G, k) \rightarrow Large\ IndSet(G', k') \quad (1)$$

$$Large\ IndSet(G', k') \rightarrow IS(G, k) \quad (2)$$

(also note that if $k \geq n/2$, $Large\ IndSet(G', k')$ can end as $|G'| \leq k$ if $k < n/2$, " as $|G'| = \frac{2n-2k}{2} \leq k' = \frac{n-k}{2}$)

① Now, if there exists a solution to $IS(G, k)$

- \hookrightarrow if $k \geq n/2 \rightarrow$ it exists in $Large\ IndSet(G', k')$ as $|G'| \leq k$
- if $k < n/2 \rightarrow$ $|G'| = n$

On adding $n-2k$ disconnected vertices to G , we now have $k + n-2k = n-k$ disconnected vertices in G' (disconnected = independent (independent)).

Since $k' = n-k$, we would be able to find a solⁿ for $Large\ IndSet(G', k')$ as $k' = n-k = |G'|$ $k' \geq \frac{|G'|}{2}$

② Now, if \exists a solⁿ to $Large\ IndSet(G', k')$

- \hookrightarrow if $k \geq n/2$ same argument as above ($G' = G$ & $k' = k$)
- if $k < n/2$ Corresponding to solⁿ for $Large\ IndSet(G', k')$ \exists a solⁿ to $IS(G, k)$

Since we added only $n-2k$ independent vertices & $k' = n-k$, \therefore there must exist k independent vertices in the original graph G , & $k' \geq |G'|/2$. Corresponding to those k vertices, we have a solⁿ to $IS(G, k)$.

reduces to

$\therefore IS(G, k) \sim Large\ IndSet(G, k)$ reduces in polynomial time