

Read the instructions carefully.

More instructions specific to this exam:

- Answers must be accompanied by rigorous proofs. Unlike the previous exams, vague ideas might not get marks.
 - While describing the behavior of a Turing machine, it is sufficient to give a high level description, that is, a COL106 or COL351 style algorithm.
 - We have assumed or proved the NP-hardness of only the following problems in class and in Homework 5: **CircuitSAT**, **3SAT**, **Strict3SAT**, **IndSet**, **Clique**, **SubgraphIsomorphism**, **VertexCover**, **SetCover**, **HittingSet**, **DominatingSet**, and the crazy COL352 instructor problem. Only these problems are to be assumed to be NP-hard for this exam.
- [1 mark each]** State whether each of the following statements is true or false. Substantiate each answer with a short (ideally, at most 3 sentences of) justification. (Answers with incorrect and insufficient justifications get 0 marks.)
 - The language $\{r \mid r \text{ is a regular expression over } \{0,1\} \text{ and } \mathcal{L}(r) = \{0,1\}^*\}$ is decidable.
 - If L_1 is a context-free language and L_2 is a regular language then $L_1 \setminus L_2$ and $L_2 \setminus L_1$ are both necessarily context-free.
 - The languages $\{w \mid \mathcal{L}(M_w) \text{ is empty}\}$, $\{w \mid \mathcal{L}(M_w) \text{ is finite}\}$, $\{w \mid \mathcal{L}(M_w) \text{ is decidable}\}$, $\{w \mid \mathcal{L}(M_w) \text{ is enumerable}\}$ are all undecidable. (Refer to Homework 4 for the definition of an enumerable language).
 - 4SAT** (defined analogously to **3SAT**) is Karp-reducible to **3SAT**.
 - Consider the following algorithm which takes a grammar $G = (N, \Sigma, R, S)$ as input and either accepts or rejects it. As usual, N is the set of non-terminals, Σ is the alphabet, R is the set of production rules, and $S \in N$ is the initial non-terminal.

Algorithm 1 $\text{Alg}(N, \Sigma, R, S)$

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 $T \leftarrow \emptyset.$ 
while  $R$  contains a production rule of the form  $A \rightarrow \alpha$ , where  $A \notin T$  and  $\alpha \in T^*$  do
     $T \leftarrow T \cup \{A\}.$ 
end while
If  $S \in T$  then Accept, else Reject.

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- [4 marks]** What interesting property of G is this algorithm testing? Prove your answer.
- [6 marks]** Make minimal changes to the above algorithm so that your modified algorithm accepts a grammar G if and only if $\mathcal{L}(G) \neq \emptyset$. Argue why your algorithm is correct.
- [6 marks]** Consider the following language (which might remind you of **InfVisit** from Quiz 4).

$$\text{Visit} = \{(w, x, k) \mid M_w \text{ when run on } x \text{ visits the } k^{\text{th}} \text{ cell of its tape at least once}\}.$$

Is **Visit** recognizable? Is it co-recognizable? Prove your answers. (As usual, M_w denotes the Turing machine whose description is w , and if w doesn't describe any Turing machine legally, then M_w is the Turing machine with two states: $q_{\text{init}} = q_{\text{reject}}$ and q_{accept} , and whose tape alphabet only consists of the input alphabet and the blank.)