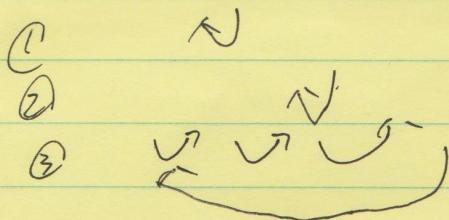
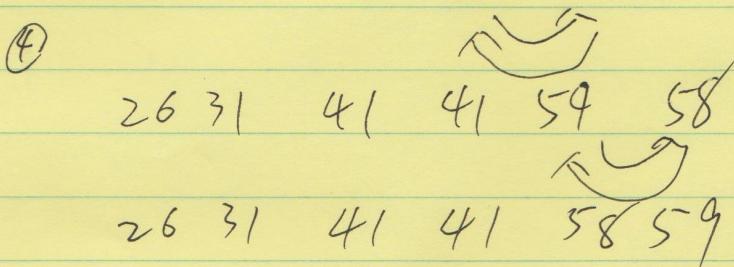


Tarjan

2.1-1 31 41 59 26 41 58



26 31 41 59 41 58



2.1-3 Input $A = \{a_1, a_2, \dots, a_n\}$ and a value v
Output An index i such that $v = A[i]$ or ~~the nil if~~

length = n

① for $i = 1$ to n

② if $A[i] = v$ then
③ return i

④ end if

⑤ end for

⑥ return $n!l$

$$2.2-1 \quad \frac{n^3}{7000} - (cn^2 - cn + 3) = \Theta(n^3)$$

in the ^{search} 2.2-3 . How many elements? Worst case, avg case.
Yes/No ① average is $\frac{n}{2}$. ② worst is n times

$\Theta(n)$

$$2.3 - 1 \quad A = \{3, 41, 152, 26, 38, 57, 9, 49\}$$

~~A~~

$$\begin{array}{cccc}
 3 & 41 & 5 & 2 & 26 & 38 & 57 & 9 & 49 \\
 & \checkmark & & \checkmark & & \checkmark & & \checkmark & \\
 3, 41 & & 26, 52 & & 38, 57 & & 9, 49 \\
 & \backslash & / & & & & / \\
 & 2, 26, 41, 52 & & & & 9, 38, 49, 57 \\
 & & \backslash & & & / \\
 3, 9, 26, 38, 41, 49, 52, 57
 \end{array}$$

2.3 - 3

$$T(n) = \begin{cases} 2 & \text{if } n=2 \\ 2T(n/2) + n & \text{if } n=2^k \text{ for } k \geq 1 \end{cases}$$

$$\therefore T(n) = n g_n$$

$$\textcircled{1} \quad \cancel{T(n) = n g_n \text{ if } n=2}$$

$$\cancel{T(2) = 2g_2 \Rightarrow 2 \cdot 1 = 2} \quad \text{so } \textcircled{1} \text{ right}$$

$$\textcircled{2} \quad \cancel{T(2^k) = 2^k g_{2^k}}$$

$$\therefore \text{for } k \geq 1 \Rightarrow \cancel{T(2^{k+1}) = 2T\left(\frac{2^k}{2}\right) + 2^{k+1}}$$

$$\therefore \cancel{T(2^k) = 2^k g_{2^k}} \Rightarrow 2(T(2^k) + 2^{k+1}) = 2^{k+1}$$

$$= 2^{k+1}$$

base case

①

$$1. n=2 \text{ so } k=1 \quad T(n)=2 = 2 \lg 2$$

$$\begin{aligned} 2. n=2^2 & \text{ so } k=2 \quad T(n)=2T(2^2/2) + 2^2 \\ & = 2T(2) + 4 \\ & = 2 \cdot 2 + 4 = 8 \\ & \text{PER} \quad = 2^2 \lg 2^2 \end{aligned}$$

$$\frac{x^y}{x^{y+1}} = x^{-1}$$

special case

②

$$\therefore \text{if } n=2^t \text{ then } 2^t \lg 2^t$$

$$\begin{aligned} \therefore \text{if } n=2^{t+1} \text{ then } T(2^{t+1}) &= 2T\left(\frac{2^{t+1}}{2}\right) + 2^{t+1} \\ &= 2T(2^t) + 2^{t+1} \\ &= 2^t(2^t \lg 2^t) + 2^{t+1} \\ &= 2^{t+1}(2^t \lg 2^t + 2) \\ &= 2^{t+1}(2^t \lg 2^t + 2^{t+1}) \\ &= 2^{t+1}(2^t \lg 2^{t+1}) \\ &= 2^{t+1}(2^{t+1} \lg 2^{t+1}) \end{aligned}$$

$$3. 1 - \text{Proof Max}(f(n), g(n)) = \Theta(f(n) + g(n))$$

see $\max(f(n), g(n)) = A$, and positive constants c_1, c_2

$$c_1 \leq \max(f(n), g(n)) \leq c_2 (f(n) + g(n))$$

$$\frac{1}{c_1} \leq \frac{\max(f(n), g(n))}{f(n) + g(n)} \leq \frac{1}{c_2}$$

for part ① $\because f(n) \geq 0, g(n) \geq 0$, $\therefore f(n) \geq g(n)$ or $g(n) \geq f(n)$

set $c_1 = 0.1$ $c_2 = 1$ $\left(\frac{f(n) + g(n)}{f(n) + g(n)}\right) \leq \frac{f(n)}{f(n) + g(n)} \leq 0.9$ is right.

② $c_2 = 1$ $A \leq c_2 (f(n) + g(n))$ is right.

$$\therefore ① ② \Rightarrow \max(f(n), g(n)) = \Theta(f(n) + g(n))$$

$$3.1-4 \quad \text{Is } 2^{n+1} = O(2^n) \quad | \quad 2^{2n} = O(2^n)$$

$$0 \leq 2^{n+1} \leq c2^n$$

~~$2 \cdot 2^n \leq c2^n$~~ ~~$\Leftrightarrow c \geq 2$~~ , ~~$c = O(2^n)$~~
 if $c=3$ then $2^{n+1} \leq 32^n$
 $\therefore 2^{n+1} = O(2^n)$ is true

$$0 \leq 2^{2n} \leq c2^n$$

$$2^{2n} \leq c(2^n)^2$$

~~$2^{2n} = O(2^n)$ is right~~

$$2^{2n} \leq c2^n$$

$$2^n \leq c(2^n)$$

~~$n \leq 19c$ can't compare~~
 is FALSE

$$3.1-7 \quad \partial o(g(n)) \cap w(g(n)) \neq \emptyset$$

~~$n > 0$~~ ~~$n \geq n_0$~~ set to γ ~~$n > 0$~~ ~~$n \geq n_0$~~

$$\partial \left(\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \right)$$

$$\partial \left(\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \right)$$

Let $a \in \gamma$, then $a = o(g(n)) = w(g(n))$

\because ① and ② $a \neq \infty$

$\therefore o(g(n)) \cap w(g(n))$ is empty set

3.2-1 if $f(n)$ and $g(n)$ monotonically ↑, then $f(n)+g(n)$ and $f(g(n))$ same.

~~$f(n) + g(n)$ not negative then $f(n) \cdot f(n) \nearrow$~~

$\because m \leq n \quad f(m) \leq f(n), g(m) \leq g(n)$

~~$f(m)+g(m) \leq f(n)+g(n)$ so they are monotonically ↑~~

$f(g(m)) \leq f(g(n))$, $\because f(n)$ is monotonically ↑, $\therefore f(g(n))$ same.

$f(m) \cdot g(m) \leq f(n) \cdot g(m) \leq f(n) \cdot g(n)$, $\therefore f(n) \cdot g(n)$ is monotonically ↑

$$3.1-4 \quad \text{Is } 2^{n+1} = O(2^n) \quad | \quad 2^{2n} = O(2^n)$$

$$\begin{aligned} 0 &\leq 2^{n+1} \leq c2^n \\ 2 \cdot 2^n &\leq c2^n \quad \cancel{\text{if } c \geq 2, \cancel{2^{n+1} = O(2^n)}} \\ \text{if } c=3 &\text{ then } 2^{n+1} \leq 3 \cdot 2^n \\ \therefore 2^{n+1} = O(2^n) &\text{ is true} \end{aligned}$$

$$0 \leq 2^{2n} \leq c \cdot 2^n$$

$$2^{2n} \leq c \cdot 2^n$$

if $2^{2n} = O(2^n)$ is right

$$\begin{aligned} 2^{2n} &\leq c \cdot 2^n \\ 2^n &\leq c \cdot 2^n \\ n &\leq \log c + 1 \\ n &\leq 10c \quad \text{can't compare} \\ \therefore &\text{is false} \end{aligned}$$

$$3.1-7 \quad \text{Is } o(g(n)) \cap w(g(n)) \neq \emptyset$$

set to γ

$$\begin{aligned} o\left(\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0\right) & \quad o\left(\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty\right) \\ n > 0, n \geq n_0 & \quad n > 0, n \geq n_0 \end{aligned}$$

Let $a \in \mathbb{Y}$, then $a = o(g(n)) = w(g(n))$
 $\because \textcircled{1} \text{ and } \textcircled{2} \quad c \neq \infty$
 $\therefore o(g(n)) \cap w(g(n))$ is empty set

3.2-1 if $f(n)$ and $g(n)$ monotonically \nearrow , then $f(n)+g(n)$ and $f(g(n))$ same.

$\cancel{f(n)}$ and $\cancel{g(n)}$ not negative then $f(n)+g(n) \nearrow$

$\therefore \textcircled{1} \quad m < n \quad f(m) \leq f(n), g(m) \leq g(n)$

$f(m) + g(m) \leq f(n) + g(n)$. so they are monotonically \nearrow
 $f(g(m)) \leq f(g(n))$, $\therefore f(n)$ is monotonically \nearrow , $\therefore f(g(n))$ same.

$f(m) \cdot g(m) \leq f(n) \cdot g(n) \leq f(n) \cdot g(n)$, $\therefore f(n) \cdot g(n)$ is monotonically \nearrow

P_{3-2} $\exists i \quad \epsilon > 0$ $c_i \theta$
 $A \Rightarrow O \quad o \quad \cap \quad w \quad \theta$ of B

$@ C^k n$	Yes	No	No	Yes	No	$n^{C>0}$
$\textcircled{1} \quad n^k$	Yes	No	No	No	No	c^n
$\textcircled{2} \quad \ln$	All	No	$\because \sin n \neq 0$	No	Yes	$n^{\sin n}$
$\textcircled{3} \quad 2^n$	Yes	Yes	No	Yes	Yes	$2^{\frac{n}{2}}$
$\textcircled{4} \quad n^{19}$	Yes	No	Yes	No	Yes	c^{19n}
$\textcircled{5} \quad \underline{O(n)}$	Yes	No	Yes	No	Yes	$(O(n))$ $A=B$

$\leq O(O(n))$ $n^k O(n^k)$, $2^n O(2)$,

O and w in the end, because not sure how to do, not on note.

$\Rightarrow \cap$ in $\textcircled{2}$ $A \oplus B$, so other was opposite to O

$= b$ only $\textcircled{2}$ $\textcircled{3}$

$O \quad zh^2 = O(n^2)$ but $zh \neq O(n^2)$, $zh = o(n^2)$, but $zh^2 \neq o(n^2)$
 $f(n) < g(n)$

$n \geq n_0$ need have condition O , so $\textcircled{2}$ is No
 can't same, so $\textcircled{2}$ $\textcircled{3}$ is No

$\textcircled{4}$ if $lg^k n < n^{\epsilon}$ set $k=1 \epsilon=1$ and n' , yet $k=100 \epsilon=1$ $n' \neq n$. No

$\textcircled{5}$ Yes, $O \neq 0 \Rightarrow 0$ is Yes

w generation, \cap need Yes.

$$P3-3 \quad g_1 = \mathcal{R}(g_2)$$

$$f(n) = \Theta(g(n))$$

$$2^{\sqrt{2^{2n}}}, 2^n, 2^{2n+1}$$

- (a)
- 1 2^{4n}
 - 2 $\sqrt[1.414]{(g^2)^n}$, $(\frac{3}{2})^n$, 2^{2n} , $n \cdot 2^n$, $\sqrt[1.5]{2^{gn}}$, e^n , $4^{\frac{gn}{10n}}$
 - 3 $\sqrt[1.414]{(g^2)^n} \cdot n \cdot 2^n$
 - 4 $n! \cdot (19^n)! \cdot (n+1)!$

clean (1) $1, (\sqrt{2})^{19n}, (\frac{3}{2})^n, 2^{\sqrt{2^{19n}}}$, $2^{\sqrt{2^{19n}}}, 2^n, n \cdot 2^n, 2^{2n}, e^n, 4^{\frac{gn}{10n}}$

(1) $1, (\sqrt{2})^{19n}, (\frac{3}{2})^n, 2^{\sqrt{2^{19n}}}, 2^{\sqrt{2^{19n}}}, 2^n, 4^{\frac{19n}{10n}}, 1, 2^n, 2^{2n}, 2^{2n+1}, 2^{2n+2}, n \cdot 2^n, e^n$.

(clean 2) $1 \cdot n, 1 \cdot n \cdot n, \sqrt{19n}, 19 \cdot (19^n), 19 \cdot (19^n), (\sqrt{19n})^{\frac{19n}{10n}}, n \cdot 2^n, e^n, (19^n)^{\frac{19n}{10n}}, n \cdot 2^n, e^n$.

(2) $1 \cdot n, n \cdot n, \sqrt{19n}, (19 \cdot (19^n)), (19 \cdot (19^n)), (19 \cdot (19^n)), n \cdot 2^n, 19 \cdot (19^n), ((19^n))^{\frac{19n}{10n}}, 19^2 \cdot n$

(2-1) n
 (2) $n^{\frac{1}{10n}}, n^{\frac{1}{10n}}, n^2, n^3$
 (4) $((19^n)!)!, n!, (n+1)!$

$g_1 = n g_2 \dots$, in the same cst. $f(n) = \Theta(g(n))$

(b) $f(n) = ((\sqrt{2})^{19n})^{\frac{19n}{10n}}$