Monoids and Tigers and Folds, oh my!

Christopher R. Genovese

Department of Statistics & Data Science

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Follow-up Discussion on SWAG

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Folds

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Monoids and Semirings

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Activity

Announcements

- Please fill out Experience survey.
- Please fill out office hours poll.
- Special Office Hour times this week will be posted on Canvas.
- Email subject: [750]
- Please bring your laptop to every class
- Reading:
 - Tuesday: https://36-750.github.io/course-info/shells/, https://git-scm.com/book/en/v2 sections 1.1 and 1.3, and Thinking Languages (almost ready).
- Homework: HW #1 (swag) due Thursday 5 Sep.
- Don't forget to send me your Github account names

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The "Fold" Pattern

In the basic Sliding-Window Aggregation, we repeatedly compute an aggregation over the window with the following pattern (expressed in TL2):

```
accumulator = starting_value
for item in items
accumulator = update_fn(accumulator, item)
```

Line 1 initializes the accumulator. Lines 2-3 successively updates it. Here, update_fn is often called the folding function, reducing function, or update rule. In the SWAG problem, it is just our binary operation.

This is a common pattern called a **fold**.

Examples? Variants?

Folds as "Objects"

We define a *type* to describe folds as a unified object. (We can make this a Object but it need not be, depending on the language.)

A fold is specified by four parameters:

- ① a function that wraps or annotates input values,
- 2 a folding function that does the updates,
- 3 an initial value, and
- 4 a finalizing function that unwraps the final update.

In TL1, this is a type Fold v w r with type variables v, w, and r where

- v is the type of input values;
- w is the type of the wrapped values; and
- r is the type of the final *result* produced.

Here, Fold is a "type constructor," it takes three types and returns a new type. Indeed, it is a function with type Type -> Type -> Type.

In terms of these types, what are the parts of a Fold?

Folds as "Objects" (cont'd)

In TL1, we describe the Fold type as

```
type Fold v w r = forall a.
record Fold where
lift: (v -> w)
step: (a -> w -> a)
init: a
done: (a -> r)
```

Think of this like a function definition where the arguments and return values are types.

The forall a on Line 1 defines a "local" type variable a (for accumulator) that can be anything. Given four appropriate functions for the fields on Lines 3–6, we can *infer* what type a must be, so we don't need to specify it as an independent parameter.

If x has type Fold v w r, then we can use x.lift, x., x.init, and x.done to refer to the fields of x, and we can use Fold.lift, Fold.step, etc. as accessor functions.

Folds as "Objects" (cont'd)

You can represent this type in various ways in various languages. In Python or R, for instance, we might define a class, or we could define a simpler data structure (such as a namedtuple or a list with named components).

For example, in Python:

```
A = TypeVar['A']
class Fold(Generic[V, W, R]):
    def __init__(
        self,
        lift: Callable[[V], W],
        step: Callable[[A, W], A],
        init: A,
        done: Callable[[A], R]
) -> None:
        ...
```

It's also good practice to define function that creates your data structure, if it's not defined through your language (and sometimes even if it is).

Q: What does a Fold for mean look like?

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It is a Fold Num (Num, Int) Num where Num represents numbers (with subtype Int) and (a, b) is a pair with component types a and b.

```
lift x = (x, 1)

step (x, m) (y, n) = (x + y, m + n)

init = 0

done (x, n) = x / n
```

Notice that we made this even more general: the aggregated values can have a different type than the input.

(Ex: counting total number of characters in a sequence of strings.)

Q: What does a Fold for mean look like?

Q: What is the type of swag in terms of the Fold type?

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```
swag : Fold v a b -> Nat -> [v] -> [b]
```

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In general, we have an operation which takes a Fold and a "foldable" container and returns the final result.

```
fold : Foldable c => Fold v w r -> c v -> r
```

Here, c is a type that represents containers containing that can be folded and to the left of the => is a *constraint*. (What's an example of a Foldable besides an array?)

We'll see soon some of what we get from representing fold operations this way.

We can also define related operations like scan (a fold that collects intermediate values), foldRight (a fold from the end, where appropriate), and scanRight.

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Lists

Let \mathcal{A} be a non-empty set. Then $\mathsf{List}(\mathcal{A})$, commonly denoted \mathcal{A}^* , is the set whose elements are *finite*-length tuples of elements from \mathcal{A} .

We can endow the set $\mathsf{List}(\mathcal{A})$ with structure by defining a binary operator $::: \mathsf{List}(\mathcal{A}) \times \mathsf{List}(\mathcal{A}) \longrightarrow \mathsf{List}(\mathcal{A})$ that concatenates two lists. For example, $\langle 0,0 \rangle :: \langle 1,0,1,1 \rangle = \langle 0,0,1,0,1,1 \rangle$. This operator satisfies three **algebraic** laws, which you can confirm for yourself:

$$\begin{split} \langle \rangle &:: \ell = \ell \\ \ell &:: \langle \rangle = \ell \\ \ell_1 &:: (\ell_2 :: \ell_3) = (\ell_1 :: \ell_2) :: \ell_3. \end{split}$$

where $\ell, \ell_1, \ell_2, \ell_3 \in \mathsf{List}(\mathcal{A})$. So, $\langle \rangle$ is an *identity* element for ::, and :: is associative. Notice that :: is *not* commutative.

Monoids

A monoid $\langle \mathcal{M}, \Diamond, e \rangle$ consists of a set \mathcal{M} equipped with a binary operator

- $\lozenge\colon \mathcal{M}\times\mathcal{M} \to \mathcal{M} \text{ and a special element } e\in \mathcal{M} \text{ that satisfy:}$
 - ① \lozenge is associative: $m_1 \lozenge (m_2 \lozenge m_3) = (m_1 \lozenge m_2) \lozenge m_3$ for every $m_1, m_2, m_3 \in \mathcal{M}$, and
- ② e is an identity element: $e \lozenge m = e = m \lozenge e$ for every $m \in \mathcal{M}$.
- \Diamond need not be commutative, but if it is we call this a $\emph{commutative monoid}.$

A special case we will use is an *ordered monoid*, which is a monoid $(\mathcal{M}, \lozenge, e)$ with a partial order \prec such that $x, y \in \mathcal{M}$ with $x \prec y$ implies

$$x \Diamond z \prec y \Diamond z$$
 and $z \Diamond x \prec z \Diamond y$,

for all $z \in \mathcal{M}$.

How might we represent a monoid in a program?

What are some examples of Monoids? (Remember it's not just the set but the identity and operator as well.)

- **1** $(\mathbb{R}, +, 0)$
- $oldsymbol{2}$ ($\mathbb{R}_+,\cdot,1$)

- \bullet ($\mathbb{R}, +, 0$)
- ② $(\mathbb{R}_+,\cdot,1)$
- Booleans with and or or. (What are the units?)

- **1** $(\mathbb{R}, +, 0)$
- **2** $(\mathbb{R}_+,\cdot,1)$
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- 4 Natural numbers with max and ???

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- **7** Functions $\mathcal{X} \longrightarrow \mathcal{X}$ with composition

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- **(1)** ...

Semirings

We say that $\langle \mathcal{S}, \boxplus, \mathbf{0}, \boxdot, \mathbf{1} \rangle$ is a **semiring** when \mathcal{S} is a set with two special elements, denoted by $\mathbf{0}$ and $\mathbf{1}$, and two operators \boxplus and $\boxdot : \mathcal{S} \times \mathcal{S} \longrightarrow \mathcal{S}$ that satisfies

- $(0, \mathbb{H}, \mathbf{0})$ is a *commutative* monoid
- $(2, \mathbb{C}, \mathbf{1})$ is a monoid
- **3 0** annihilates: $x \boxdot \mathbf{0} = \mathbf{0} = \mathbf{0} \boxdot x$

$$a \odot (b \boxplus c) = (a \odot b) \boxplus (a \odot c)$$

 $(b \boxplus c) \odot a = (b \odot a) \boxplus (c \odot a).$

The operator \boxdot need not be commutative; if it is, we have a *commutative* semiring.

Semirings describe sets that are monoids in two separate ways, with some consistency requirements for how the two monoids relate. They generalize the natural numbers and have a huge range of algorithmic applications. We'll see much more of them later.

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Semiring Examples (brief)

- Booleans
- Subsets
- Relations
- 4 Languages
- Matrices
- 6 Polynomials/Sequences
- Tropical (and other numeric) Semirings (min-plus and max-plus)
- 8 ...

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Find the contiguous segment of cells with the largest absolute difference between the number of red and blue cells. Expect that the number of cells can be very large, e.g., $n=10^6$. Can you find a method that scales *linearly* in n?

(h/t Bret Yorgey)



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- How might we think about "red" and "blue" in a way that generalizes the problem?



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- How can we identify the maximizing segment as well?
- How might we think about "red" and "blue" in a way that generalizes the problem?
- Should we be restricted to numbers? What are the algebraic properties that we need from a cell's contents?

(h/t Bret Yorgey)

Kadane's Algorithm (mutable version)

best = current

```
def kadane(cells: list[int]) -> int: # any summable
      "Returns the maximum segment sum of cells w/Kadane's al
      best, current = 0, 0
      for cell in cells:
          current = max(current + cell, 0)
          best = max(best, current)
      return best
Why does this work?
Note we could write the loop contents equivalently as
         current += cell
         if current < 0:
             current = 0
         if current > best:
```

Kadane's Algorithm (towards an immutable version)

Let's express the algorithm in terms of a Fold

```
type State c = record State { total : c, biggest : c }
-- Assume we have agg : c \rightarrow c \rightarrow c that aggregates cells
-- (with analogue of negative reset) and an value c0 : c
-- to start the aggregation
kadane : Fold c (State c) c
kadane = Fold id step (State c0 c0) State.biggest
step : Ordered c => State c -> c -> State c
step (State cur best) value = State next (max best next)
    where
        next = agg cur value
```

Kadane's Algorithm (TL1 version)

The only operations we use in the original were adding, a neutral element, and comparing. We can apply this to cells containing any *ordered Monoid*.

```
bestSoFar : (Ordered c, Foldable t) => Fold a c c -> t a ->
bestSoFar f as = scan f as |> max

kadane : (Monoid a, Ordered a, Foldable t) => t a -> a
kadane = bestSoFar f
   where
        next s a = max munit (s <> a)
        f = Fold {wrap=id, step=next, init=munit, done=id}
```

Here, munit: a is the inferred unit for the Monoid type. It could be a class member or selected by a protocol.

THE END