

Computer Networks : Protocols and Practice

Part 1 : Introduction

Olivier Bonaventure
<http://inl.info.ucl.ac.be/>

These slides are licensed under the creative commons attribution share-alike license 3.0. You can the detailed information about his license at <http://creativecommons.org/licenses/by-sa/3.0/>

Module 1 : Basics

- Contents

- • Introduction

- Services in computer networks
 - Connectionless service
 - Connection oriented service
- Layered reference models

A network ...

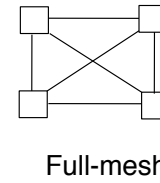
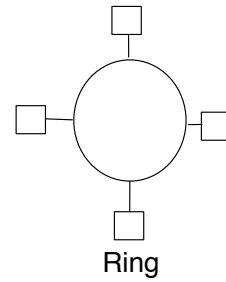
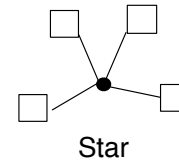
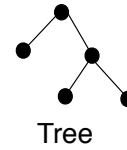
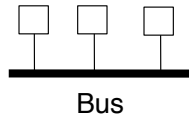
- A network is ...
 - a set of hardware and software that allows to transmit information from one sender to one or more receivers
- Current networks
 - Plain Old Telephone System (POTS)
 - Mobile Telephone
 - Broadcast networks
 - television, radio
 - Computer networks
 - Internet
 - Proprietary networks

Network classification

- Based on their geographical coverage
 - 0.1-1 m : Internal bus/network
 - 10 m - 1 km : Local Area Network (LAN)
 - 1 km - 100 km : Metropolitan Area Network(MAN)
 - 100 km ->... : Wide Area Network (WAN)
 - and more ... : Satellite networks
- Interplanetary network

Network classification (2)

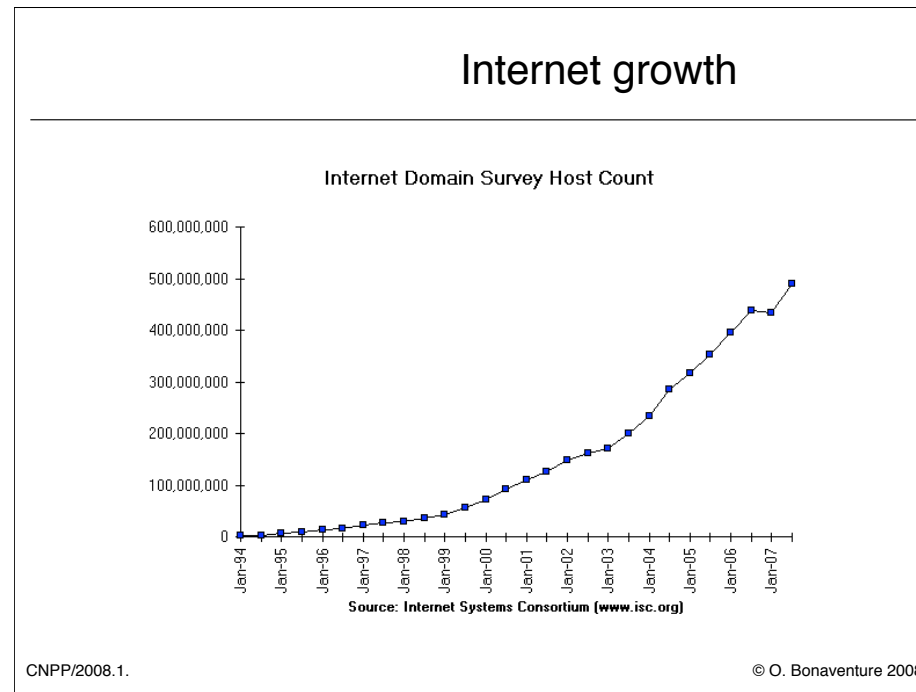
- Based on their topologies



CNPP/2008.1.

© O. Bonaventure 2008

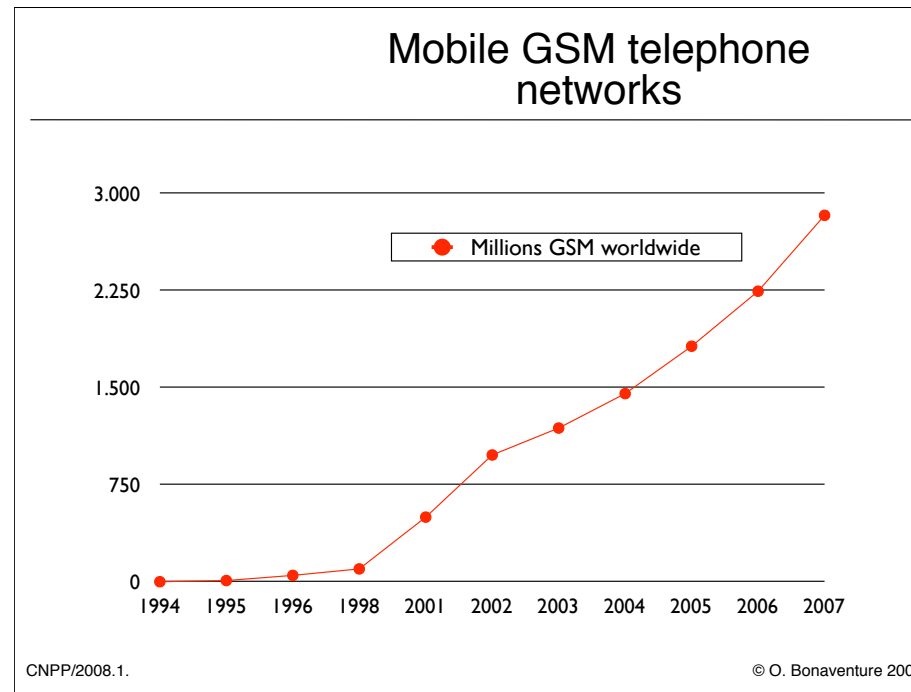
Internet growth



6

Source : <http://www.isc.org>

See also <http://hdr.undp.org/> for information about how Internet is spread across countries

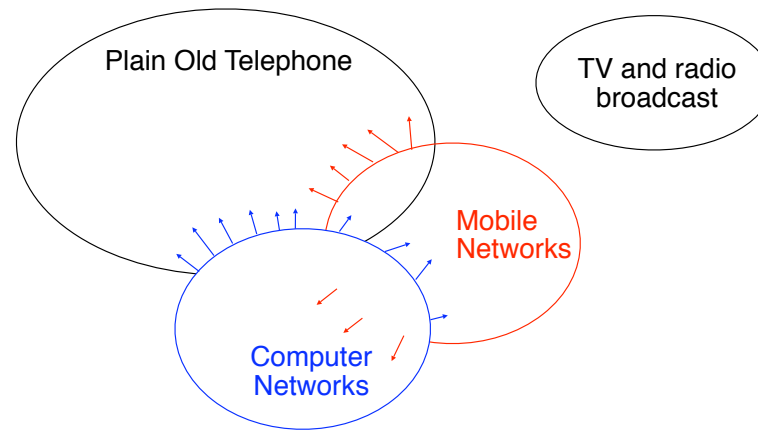


7

Source : <http://www.gsmworld.com/news/statistics/index.shtml> and <http://www.gsmworld.com/about/history.shtml>

The number of fixed telephone lines reached 1 billion worldwide in 2000 and the growth of fixed lines has been rather slow since then to reach 1.2 billions in 2005. See <http://www.itu.int/ITU-D/ict/mdg/storyline/index.html>

Evolution of networks



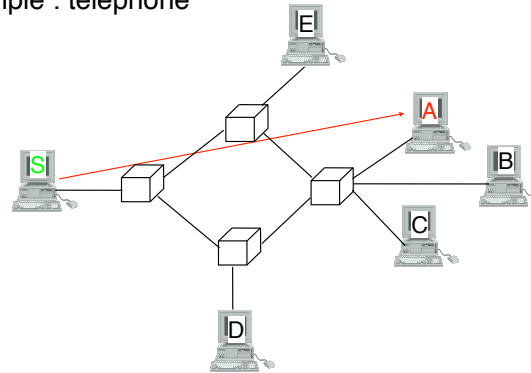
The future

- Most specialists expect
 - A strong convergence between all technologies
 - Triple play
 - Quadruple play
- New services will probably be deployed first (and perhaps exclusively) on data networks
 - Television service provided by telecom operators
 - Mobile data services
 - Mobile television services
 - Voice or video over IP
 - New services

Transmission modes

Unicast

- **Unicast or point-to-point**
 - one **sender**
 - one **receiver**
 - example : telephone



CNPP/2008.1.

© O. Bonaventure 2008

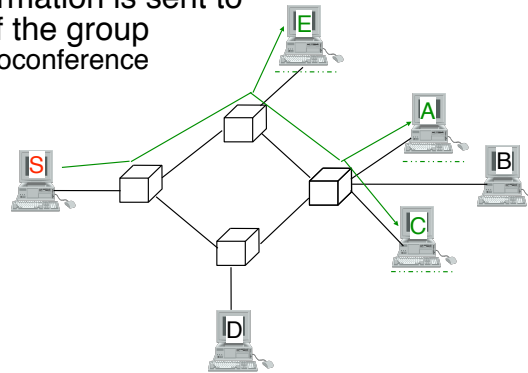
The unicast mode is the most widely used transmission mode. We will study it in details.

Transmission modes

Multicast

□ Multicast or point-to-multipoints

- one **sender**
- a **group of receivers**
- The same information is sent to all members of the group
 - example : videoconference



□ Broadcast

- The same information is sent to everyone

CNPP/2008.1.

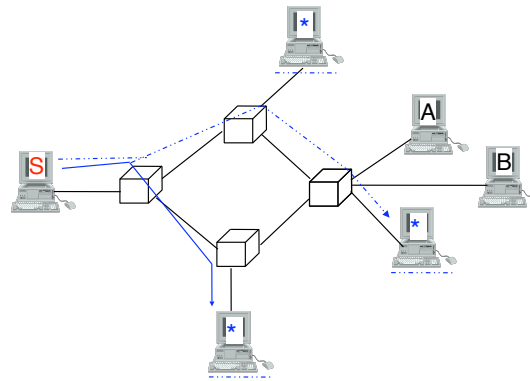
© O. Bonaventure 2008

The multicast transmission mode is described later.

Anycast

□ Anycast

- Information is sent from **one sender** to **one receiver** among a group of possible receivers
 - Example :find server hosting popular content



CNPP/2008.1.

© O. Bonaventure 2008

12

The anycast transmission mode was proposed in

C. Partridge, T. Mendez, W. Milliken, Host Anycasting Service, RFC 1546, 1993

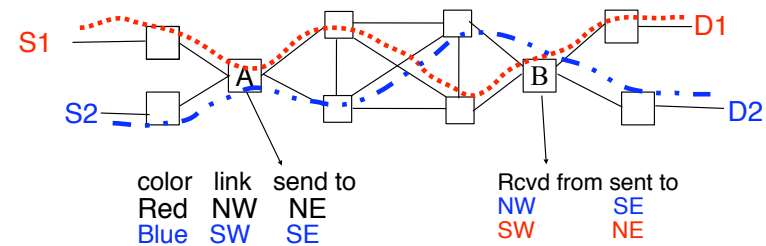
How to carry data through a network ?

□ Circuit switching

□ Principle

- before transmitting data, a circuit is established from the source to the destination hosts
- each intermediate host knows how to forward information received on a circuit that crosses itself

□ Example : POTS



CNPP/2008.1.

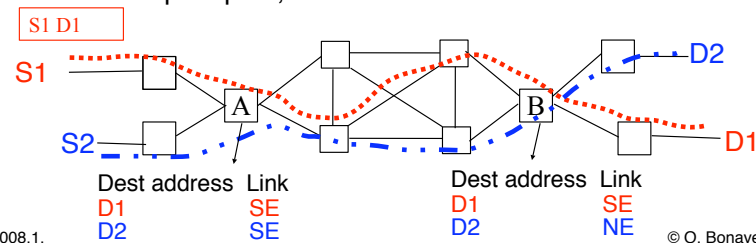
© O. Bonaventure 2008

How to carry data through a network ? (2)

□ Packet switching

□ Principles

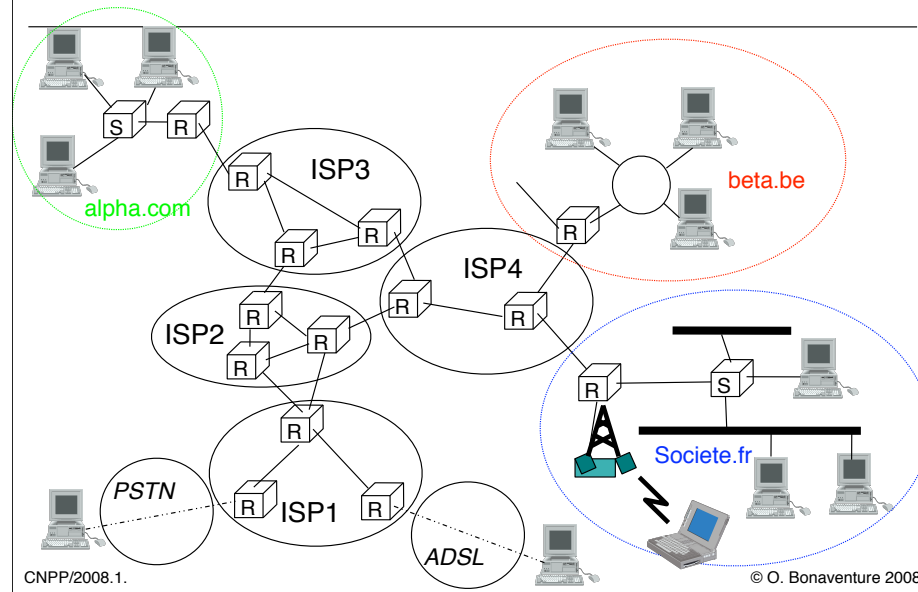
- An address is associated to each host
- data is divided in small **packets**
 - each packet contains
 - the data to be exchange
 - the address of the source host
 - the address of the destination host
- Each intermediate host knows how to reach each destination
- Example : post, Internet



CNPP/2008.1.

© O. Bonaventure 2008

A small Internet



Module 1 : Basics

- Contents

- Introduction

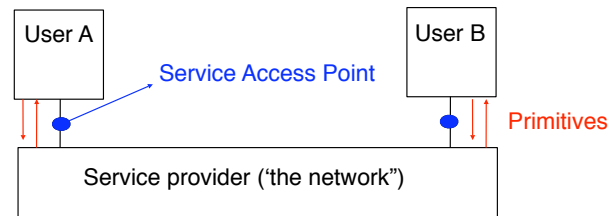
- □ Services in computer networks

- Connectionless service
 - Connection oriented service

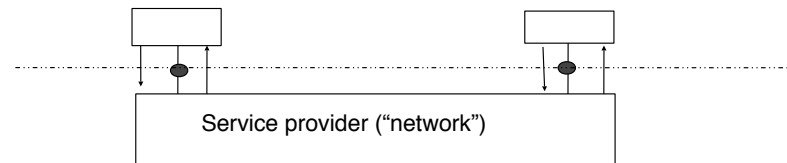
- Layered reference models

Basic concepts

- Abstract model of the network behaviour
 - Network is considered as a black box
 - Users interact with the network by using **primitives** that are exchanged through a **service access point** (SAP)

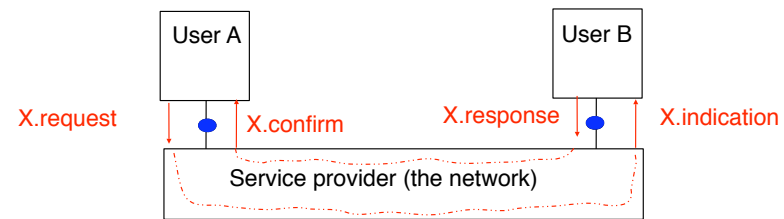


Types of primitives



- Primitive
 - Abstract representation of the interaction between one user and its network provider
 - Can contain parameters such as :
 - source
 - destination
 - message (SDU or Service Data Unit)

Types of primitives (2)



- **X.request**
 - request from a user to a service provider
- **X.indication**
 - primitive generated by the network provider to a user (often related to an earlier and remote X.request primitive)
- **X.response**
 - primitive used to answer to an earlier X.indication primitive
- **X.confirm**
 - primitive generated by the network provider to a user (related to a remote X.response primitive)

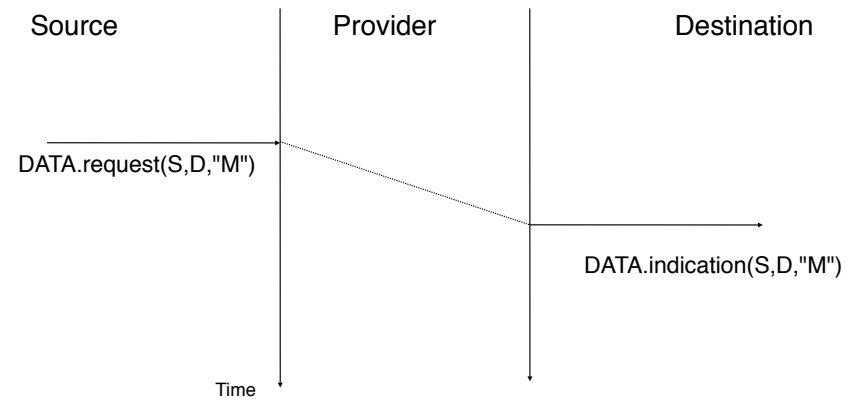
The connectionless service

- Goal
 - Allow a sender to quickly send a message to one receiver
- Principle
 - The sender places the message to be transmitted in a DATA.req primitive and gives it to the network provider
 - The network provider carries the message and delivers it to the receiver by using a DATA.ind primitive
- Utilisation
 - useful to send short-length messages
 - example : post office

Connectionless service

□ Primitives

- DATA.request(source, destination, SDU)
- DATA.indication(source, destination, SDU)



CNPP/2008.1.

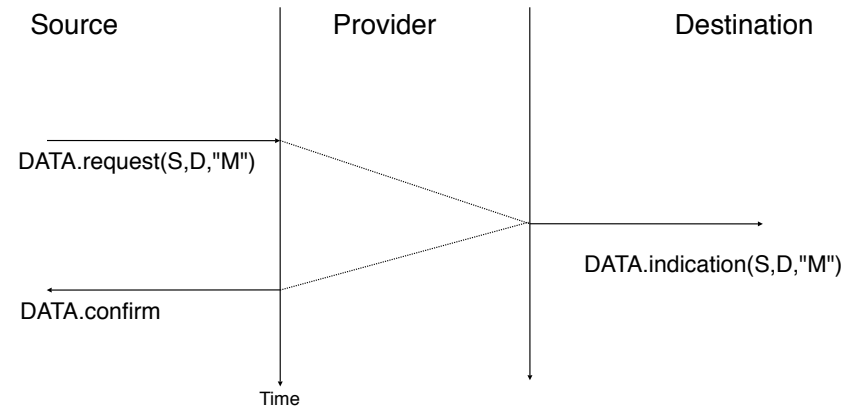
© O. Bonaventure 2008

Connectionless service (2)

- Variants of connectionless service
 - confirmation
 - primitive DATA.confirm delivered by provider to sender to confirm that some message has been delivered to destination
 - reliability
 - reliable connectionless service (no errors)
 - unreliable connectionless service (errors are possible)
 - protection against transmission errors
 - service may or may not detect/correct errors
 - protection against losses
 - the service may or cannot lose messages
 - in sequence delivery
 - not guaranteed
 - in-sequence delivery for all messages sent by one source

Connectionless service (3)

□ Example of acknowledged service



Connection-oriented service

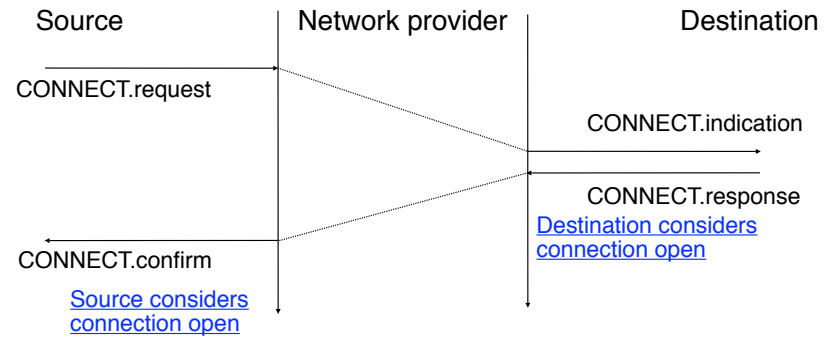
- Goal
 - Create a logical binding (connection) between two users to allow them to efficiently exchange messages
- Main phases of service
 - Connection establishment
 - Data transfer
 - both users can send and receive messages over connection
 - Connection release
- Utilisation
 - useful when the two users either
 - must exchange a large number of messages
 - need a structured exchange
 - example : telephone

Connection oriented service

□ Connection establishment

□ Primitives

- CONNECT.request
- CONNECT.indication
- CONNECT.response
- CONNECT.confirm

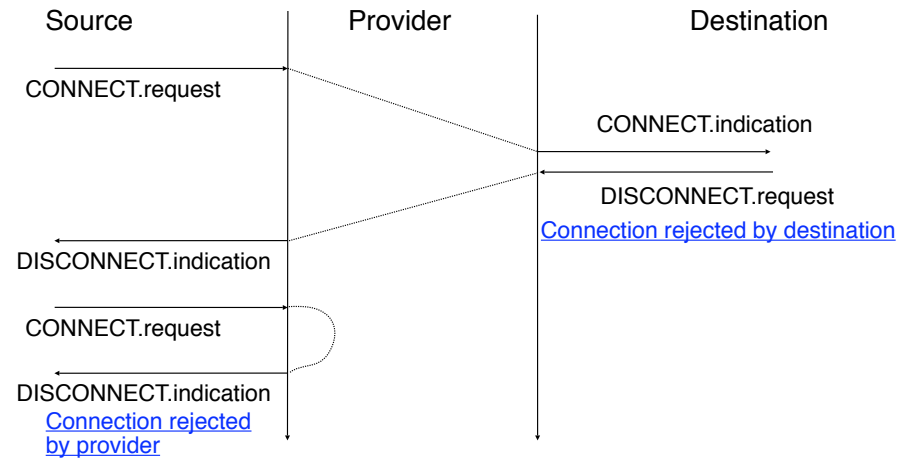


CNPP/2008.1.

© O. Bonaventure 2008

Connection oriented service (2)

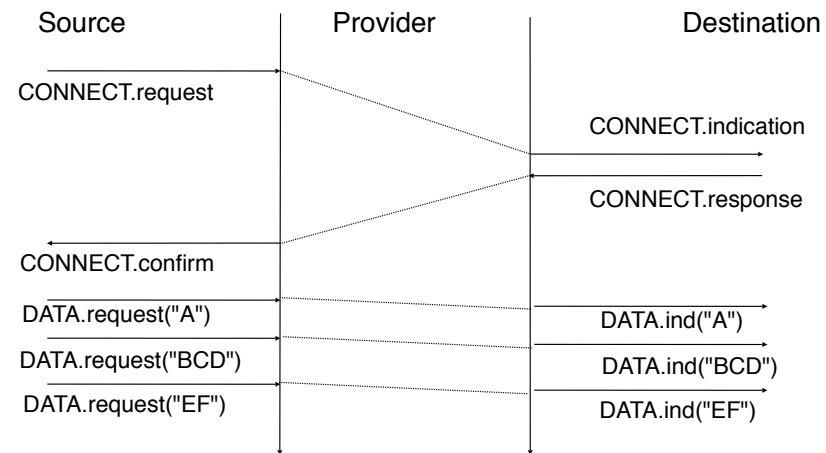
□ Connection can be rejected



CNPP/2008.1.

© O. Bonaventure 2008

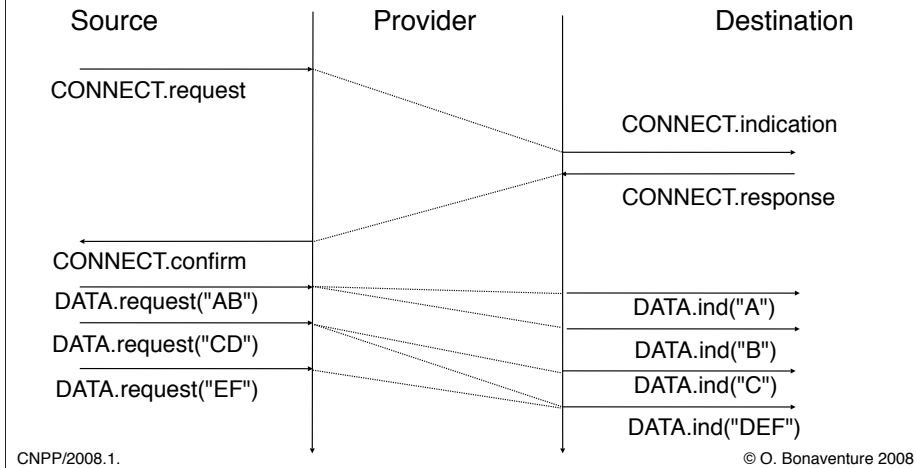
Data transfer : message mode



□ Provider delivers one Data.ind for each Data.req

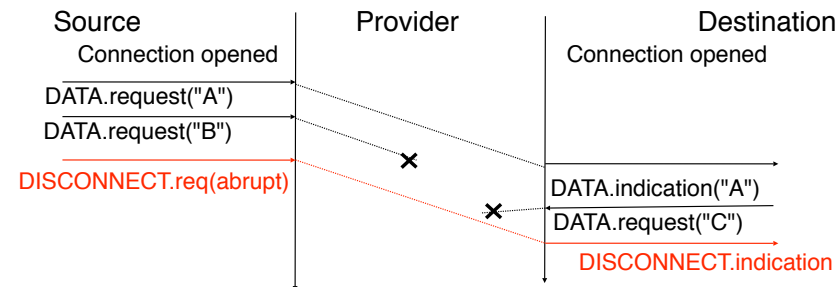
Data transfer : stream mode

- The providers delivers a **stream of characters** from source to destination



Connection release

- Abrupt release
 - SDUs can be lost during connection release

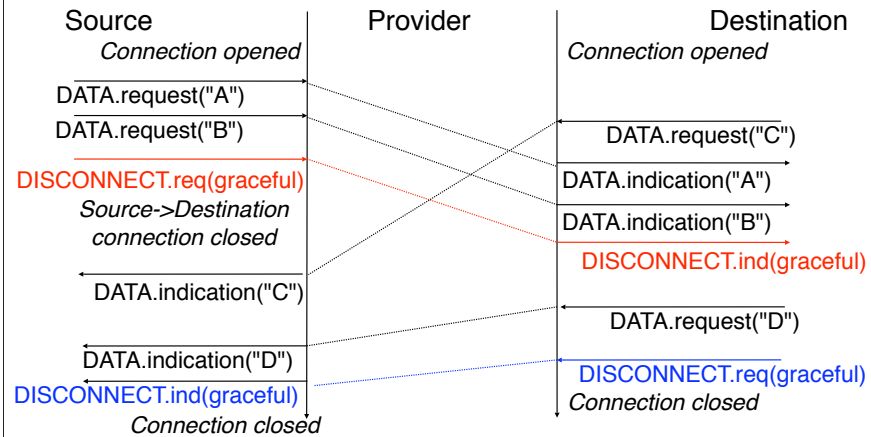


- Such an abrupt connection release can be caused by the network provider or by the users

Connection release (2)

□ Ordered/graceful connection release

- A single direction is closed at a time
- no SDUs can be lost



CNPP/2008.1.

© O. Bonaventure 2008

Characteristics of the connection-oriented service

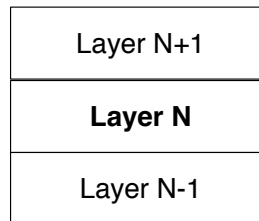
- Possible characteristics
 - bidirectional transmission
 - both users can send and received SDUs
 - reliable delivery
 - All SDUs are delivered in sequence
 - No SDU can be lost
 - No SDU can be corrupted
 - message mode or stream mode
- Connection release
 - Usually abrupt when the provider is forced to release a connection
 - Abrupt or graceful when the users request the end of a connection

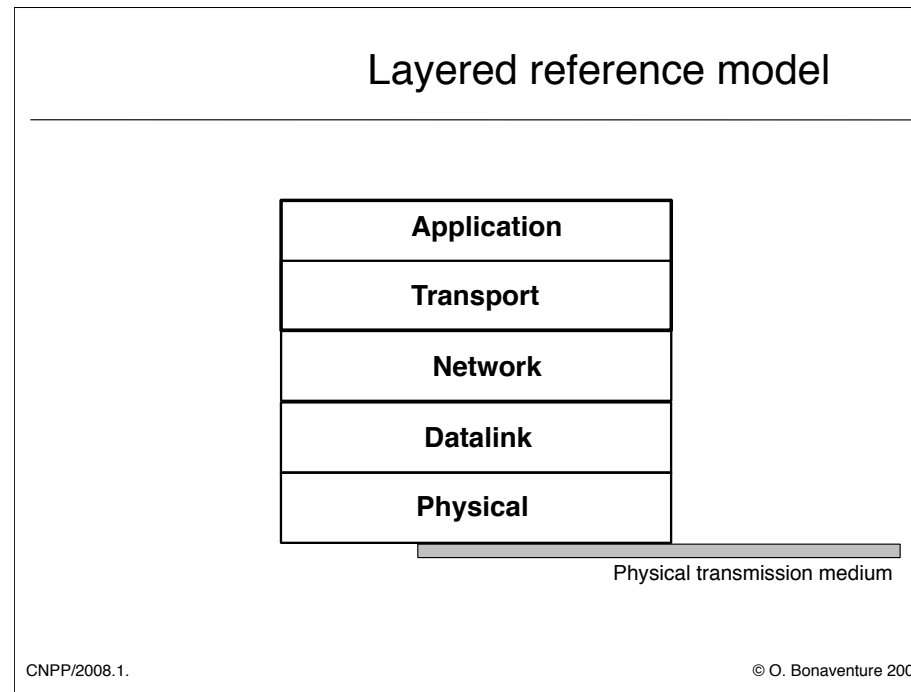
Module 1 : Basics

- Contents
 - Introduction
 - Services in computer networks
 - Connectionless service
 - Connection oriented service
 - • Layered reference models

Layered reference models

- Problem
 - How is it possible to reason about complex systems such as computer networks or the Internet ?
- Solution
 - Divide the network in layers
 - Layer N provides a well defined service to layer N+1 by using the service provided by layer N-1





34

For more information about the design philosophy of the Internet Protocols, see:

D. Clark, The Design Philosophy of the DARPA Internet Protocols, Proceedings of ACM SIGCOMM 1988, August 1998, <http://www.acm.org/sigcomm/ccr/archive/995/jan95/ccr-9501-clark.html>

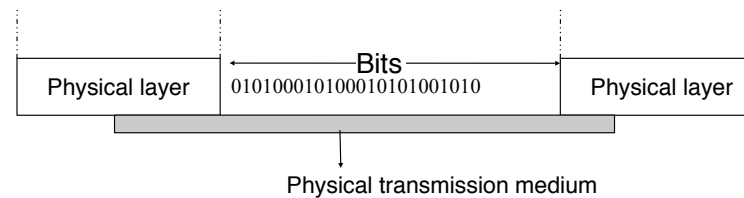
A widely known reference model is the OSI reference model :

H. Zimmerman. Osi reference model - the iso model of architecture for open systems interconnection. IEEE Transactions on Communications, 28(4):425--432, April 1980.

In this reference model, the upper layer is divided in three layers

- application layer
- session layer
- présentation layer

The physical layer



- Goal
 - Transmit bits between two physically connected devices
- Service provided by physical layer
 - bit transmission and reception
 - **unreliable** service
 - The receiver may decode a 1 while the sender sent 0
 - Some transmitted bits may be lost
 - The receiver may decode more bits than the bits that were sent by the sender

Physical layer : an example

- A very simple physical layer operating at one megabit per second
 - One bit is transmitted by sender every microsecond
 - One bit is received by receiver every microsecond
- Sender operation
 - To transmit bit=1, set V=5 Volts during one microsecond
 - To transmit bit=0, set V=-5 Volts during one microsecond
- Receiver operation
 - During each microsecond, measure V
 - If V=5 Volts, a 1 has been decoded
 - If V=-5 Volts, a 0 has been decoded
- Possible problems
 - electromagnetic perturbations
 - clock drift (sender faster than receiver or opposite)

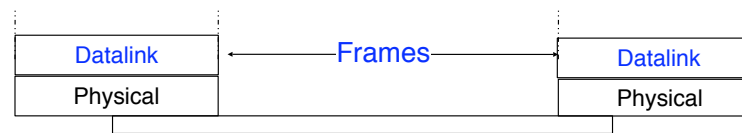
Transmission mediums

- Tapes, CDRoms and DVD
- Twisted pair
 - Telephone networks, ADSL, VDSL, ...
 - bandwidth : from a few megabits to a few 10 Mbps depending on the distance between endpoints
 - Enterprise networks
 - UTP (category 3, category 5)
 - STP (rarely used today)
 - bandwidth : up to 1 Gigabit today
 - new types of cables are being developed to reach 10 Gbps
- Wireless
 - radio
 - optical

Transmission mediums (2)

- ❑ Coaxial cable
 - ❑ Cable TV networks (CATV)
 - ❑ about 1Ghz frequency range
 - ❑ available bandwidth : depends on the split among tv distribution and data transmission
 - ❑ Computer networks
 - ❑ Used a few years ago, but not anymore today
- ❑ Optical fiber
 - ❑ monomode (laser, long distance)
 - ❑ multimode (LED, short distance)
 - ❑ frequency range : up to 100.000 Ghz
 - ❑ available bandwidth
 - ❑ 10 Gbps per wavelength and more
 - ❑ hundreds of wavelength per fiber

The datalink layer



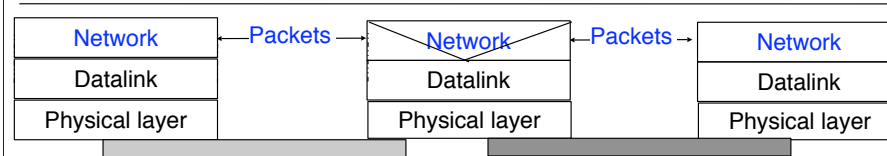
□ Goals

- Provide a service that allows the exchange of frames
 - Frame : structured group of bits
- Support local area networks

□ Services

- Reliable connection-oriented service
- Unreliable connectionless service

The Network Layer



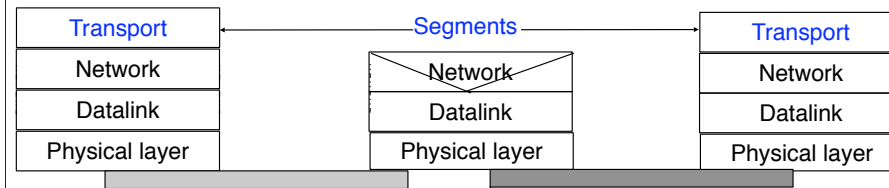
□ Goals

- Allow information to be exchanged between hosts that are not attached to the same physical medium by using relays
- The unit of information in the network layer is called a packet

□ Services

- unreliable connectionless (Internet)
- reliable connection-oriented

The Transport Layer



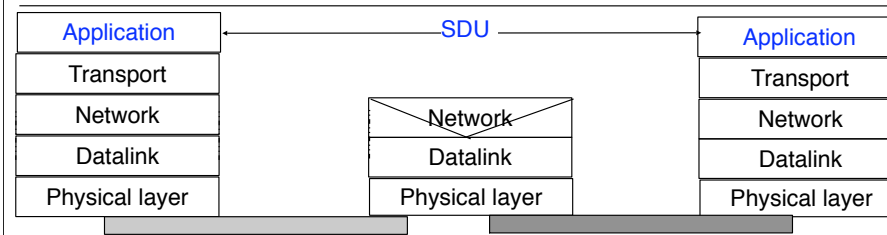
□ Goals

- Ensure a reliable exchange of data between endsystems even if the network layer does not provide a reliable service

□ Services

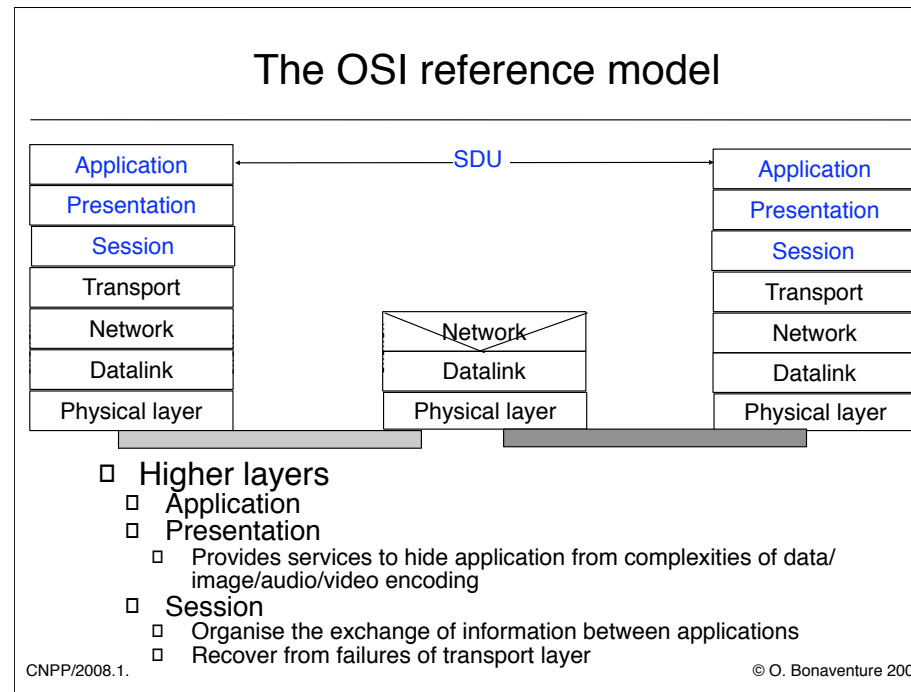
- Unreliable connectionless service
- Reliable connection-oriented service

The application layer



□ Goals

- Exchange useful information between applications by relying on the transport layer that hides the complexity of the network
- Unit of information
 - Service Data Unit, SDU



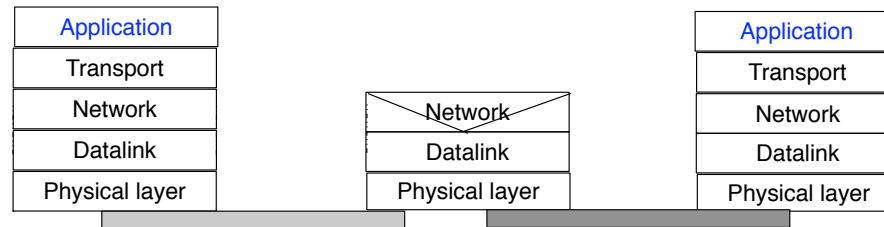
43

A detailed description of the OSI reference model may be found in :

D. Piscitello, L. Chapin, Open Systems Networking : TCP/IP and OSI, Addison Wesley, 1993

Course schedule

- First week : application layer
- Weeks 2-3 : transport layer (key mechanisms)
- Weeks 4-5 : transport layer in Internet (TCP,UDP)
- Weeks 6,7,8 : network layer (IP, RIP, OSPF)
- Weeks 9,10 : interdomain routing (BGP)
- Weeks 11,12 : Datalink layer (Ethernet, 802.11)



Exams and grading

- Oral exam
 - Several questions about theory and exercises
 - 60% of final grade
- Quizzes
 - Short quizzes will be organised randomly during the exercises or the theoretical lessons
 - 5% of final grade for all quizzes
- Project
 - Projects are carried *individually*
 - 35% of final grade for project
- Project can only be done once

Projects

- Each student can choose one project
 - Transport protocol
 - The objective is to implement in Java by using the sockets API a transport protocol that provides a reliable delivery over an unreliable network
 - Routing protocol
 - The objective is to implement in Java by using the sockets API a link state routing protocol that allows routers to efficiently route packets
 - Research project
 - The objective is to perform a detailed analysis of an Internet packet trace to predict the performance of new router mechanisms