

# Tutorial 7 - Operational Semantics

November 4, 2023

## 1 Grammar

```
 $\langle expr \rangle ::= \langle var \rangle$   
|  $\lambda \langle var \rangle . \langle expr \rangle$   
|  $\langle expr \rangle \langle expr \rangle$   
|  $(\langle expr \rangle, \langle expr \rangle)$   
|  $\text{fst } \langle expr \rangle$   
|  $\text{snd } \langle expr \rangle$   
|  $\text{true}$   
|  $\text{false}$   
|  $\text{and } \langle expr \rangle \langle expr \rangle$   
|  $\text{if } \langle expr \rangle \text{ then } \langle expr \rangle \text{ else } \langle expr \rangle$   
|  $\text{let } \langle var \rangle = \langle expr \rangle \text{ in } \langle expr \rangle$   
|  $(\langle expr \rangle)$ 
```

## 2 Call-By-Name

This is what was presented in tutorial; see `Lambda.hs` for a Haskell implementation.

$$\frac{e_1 \rightarrow e'_1}{e_1 \ e_2 \rightarrow e'_1 \ e_2} \quad \frac{}{(\lambda x. e_1) \ e_2 \rightarrow e_1[e_2/x]}$$
$$\frac{}{\text{fst } (e_1, e_2) \rightarrow e_1} \quad \frac{}{\text{snd } (e_1, e_2) \rightarrow e_2}$$
$$\frac{e_1 \rightarrow e'_1}{\text{and } e_1 \ e_2 \rightarrow \text{and } e'_1 \ e_2} \quad \frac{}{\text{and true } e_2 \rightarrow e_2} \quad \frac{}{\text{and false } e_2 \rightarrow \text{false}}$$
$$\frac{e_1 \rightarrow e'_1}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3 \rightarrow \text{if } e'_1 \text{ then } e_2 \text{ else } e_3} \quad \frac{}{\text{if true then } e_2 \text{ else } e_3 \rightarrow e_2} \quad \frac{}{\text{if false then } e_2 \text{ else } e_3 \rightarrow e_3}$$
$$\frac{}{\text{let } x = e_1 \text{ in } e_2 \rightarrow e_2[e_1/x]}$$

## 3 Call-By-Value

This was not presented in tutorial, and the Haskell implementation is left as an exercise. Recall that  $e$  denotes an arbitrary expression, and  $v$  denotes a normal form; in other words, an expression that cannot be reduced further.

$$\begin{array}{c}
\frac{e_1 \rightarrow e'_1}{e_1 \ e_2 \rightarrow e'_1 \ e_2} \quad \frac{e_2 \rightarrow e'_2}{v_1 \ e_2 \rightarrow v_1 \ e'_2} \quad \frac{}{(\lambda x. e_1) \ v_2 \rightarrow e_1[v_2/x]} \\
\\
\frac{e_1 \rightarrow e'_1}{(e_1, e_2) \rightarrow (e'_1, e_2)} \quad \frac{e_2 \rightarrow e'_2}{(v_1, e_2) \rightarrow (v_1, e'_2)} \quad \frac{}{\text{fst } (v_1, v_2) \rightarrow v_1} \quad \frac{}{\text{snd } (v_1, v_2) \rightarrow v_2} \\
\\
\frac{e_1 \rightarrow e'_1}{\text{and } e_1 \ e_2 \rightarrow \text{and } e'_1 \ e_2} \quad \frac{}{\text{and true } e_2 \rightarrow e_2} \quad \frac{}{\text{and false } e_2 \rightarrow \text{false}} \\
\\
\frac{e_1 \rightarrow e'_1}{\text{if } e_1 \text{ then } e_2 \text{ else } e_3 \rightarrow \text{if } e_1 \text{ then } e_2 \text{ else } e_3} \quad \frac{}{\text{if true then } e_2 \text{ else } e_3 \rightarrow e_2} \quad \frac{}{\text{if false then } e_2 \text{ else } e_3 \rightarrow e_3} \\
\\
\frac{e_1 \rightarrow e'_1}{\text{let } x = e_1 \text{ in } e_2 \rightarrow \text{let } x = e'_1 \text{ in } e_2} \quad \frac{}{\text{let } x = v_1 \text{ in } e_2 \rightarrow e_2[v_1/x]}
\end{array}$$

## 4 Some Useful Exercises

If you are looking for further practice, look no further! Note that no marks will be given for completing these questions; they are just for practice.

- ★ Implement Call-By-Value in Haskell
- ★★ Add 'or' to the language by updating the BNF and operational semantics.
- ★ Implement 'or' in Haskell using your semantics.
- ★★★ Devise a type system for this language; write it out using typing judgements.
- ★★★★★ Prove progress and preservation for your type system.