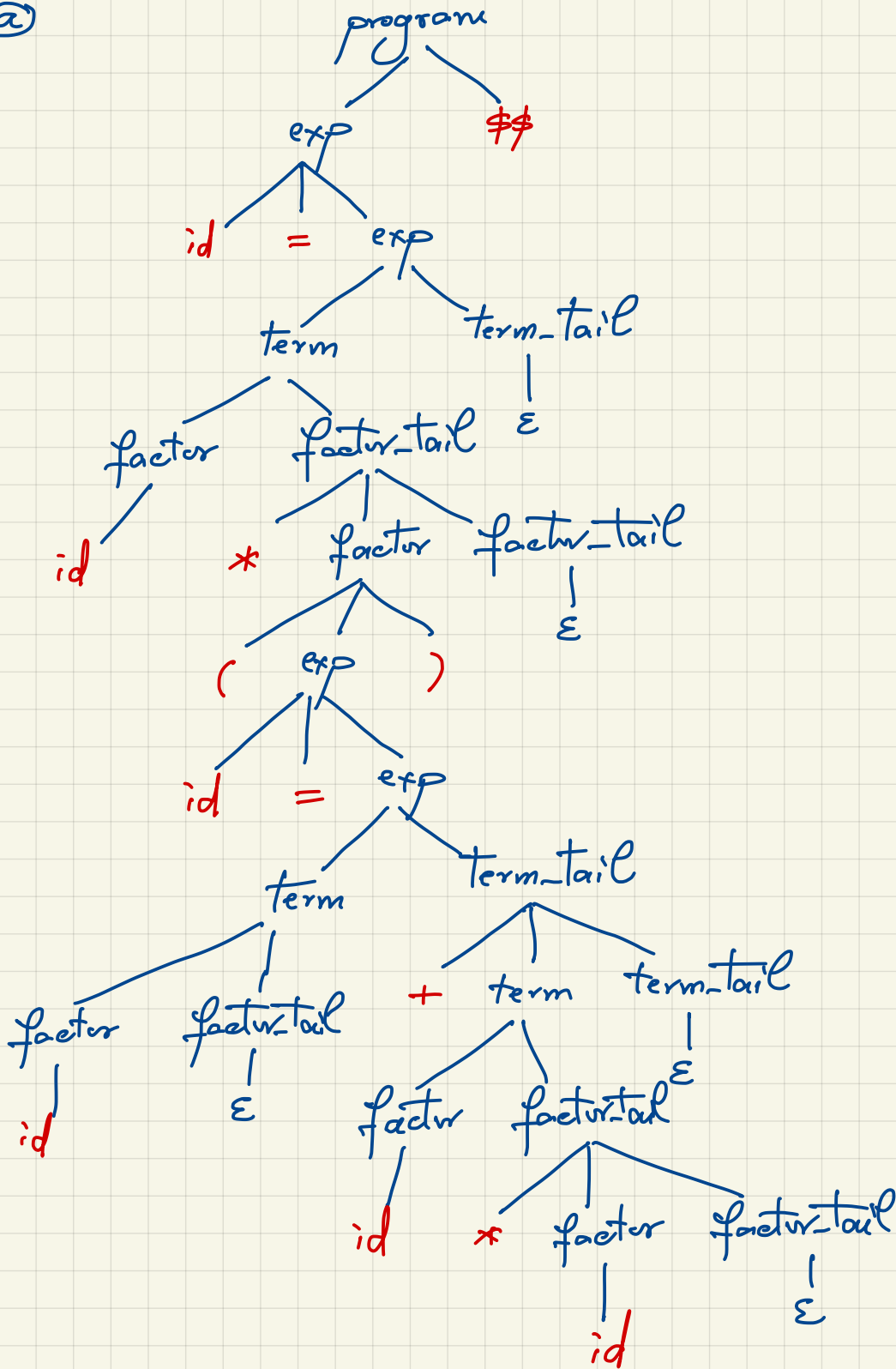


A2-sol

(CS3342 - win. 2024)

Q₁

a)



⑥

α	$FIRST(\alpha)$
$exp \ \&\&$	$id: (1,1), (: (1,2)$
$id = exp$	$id: (0,-)$
$term \ term_tail$	$(, id: (1,5)$
$+ term \ term_tail$	$+: (0,-)$
ϵ	\emptyset
$factor \ factor_tail$	$(: (1,8), id: (1,9)$
$* factor \ factor_tail$	$*: (0,-)$
ϵ	\emptyset
(exp)	$(: (0,-)$
id	$id: (0,-)$

A	$FOLLOW(A)$
$program$	\emptyset
exp	$\&\&: (2,0),): (2,8)$
$term$	$+: (2,2), \&\&,: (3,2)$
$term_tail$	$\&\&,): (3,2)$
$factor$	$*: (2,5), +, \&\&,): (3,5)$
$factor_tail$	$+, \&\&,): (3,5)$

③

i	PREDICT(i)
0	id, (
1	(id
2	(, id
3	+
4	\$\$,)
5	(, id
6	*
7	+, \$\$,)
8	(
9	id

]

④
conflict \Rightarrow not LL(1)

③ There is only one conflict, between productions 1 and 2, on id.

$$2. \text{exp} \rightarrow \text{term term_tail}$$

- can be replaced by:

$$2'. \text{exp} \rightarrow \text{factor factor_tail term_tail}$$

- because there is only one production with LHS term.

- replace again by:

$$2''a: \text{exp} \rightarrow (\text{exp}) \text{ factor_tail term_tail}$$

$$2''b: \text{exp} \rightarrow \text{id factor_tail term_tail} \quad \leftarrow \text{this is causing the conflict with prod. 1}$$

- eliminate the common prefix id between:

$$1. \text{exp} \rightarrow \text{id} = \text{exp}$$

$$2''b: \text{exp} \rightarrow \text{id factor_tail term_tail}$$

gives:

$$2'''a. \text{exp} \rightarrow \text{id exp_tail}$$

$$2'''b. \text{exp_tail} \rightarrow = \text{exp}$$

$$2'''c. \text{exp_tail} \rightarrow \text{factor_tail term_tail}$$

The new grammar is obtained by replacing productions 1 and 2 by $2''a$, $2'''a-c$, see next page.

1. $\text{program} \rightarrow \text{exp } \$$
2. $\text{exp} \rightarrow (\text{exp}) \text{ factor_tail term_tail}$
3. $\text{exp} \rightarrow \text{id exp_tail}$
4. $\text{exp_tail} \rightarrow = \text{exp}$
5. $\text{exp_tail} \rightarrow \text{factor_tail term_tail}$
6. $\text{term_tail} \rightarrow + \text{term term_tail}$
7. $\text{term_tail} \rightarrow \epsilon$
8. $\text{term} \rightarrow \text{factor factor_tail}$
9. $\text{factor_tail} \rightarrow * \text{factor factor_tail}$
10. $\text{factor_tail} \rightarrow \epsilon$
11. $\text{factor} \rightarrow (\text{exp})$
12. $\text{factor} \rightarrow \text{id}$

The only places where conflicts can appear are:

- productions 2, 3 — clearly no conflicts because RHS's start with different tokens

- productions 4, 5 : $\text{FIRST}(\text{factor_tail term_tail}) = \{*, +\}$
does not include '=', therefore no conflict

\Rightarrow new grammar is LL(1).

Q₂

a)

1. $P \rightarrow E \$$
2. $E \rightarrow E \text{ or } T$
3. $E \rightarrow T$
4. $T \rightarrow T \text{ and } F$
5. $T \rightarrow F$
6. $F \rightarrow (E)$
7. $F \rightarrow id$

b)

X	FIRST(X)	FOLLOW(X)
P	(, id	\emptyset
E	(, id	$\$, or,)$
T	(, id	$\$, or,), and$
F	(, id	$\$, or,), and$

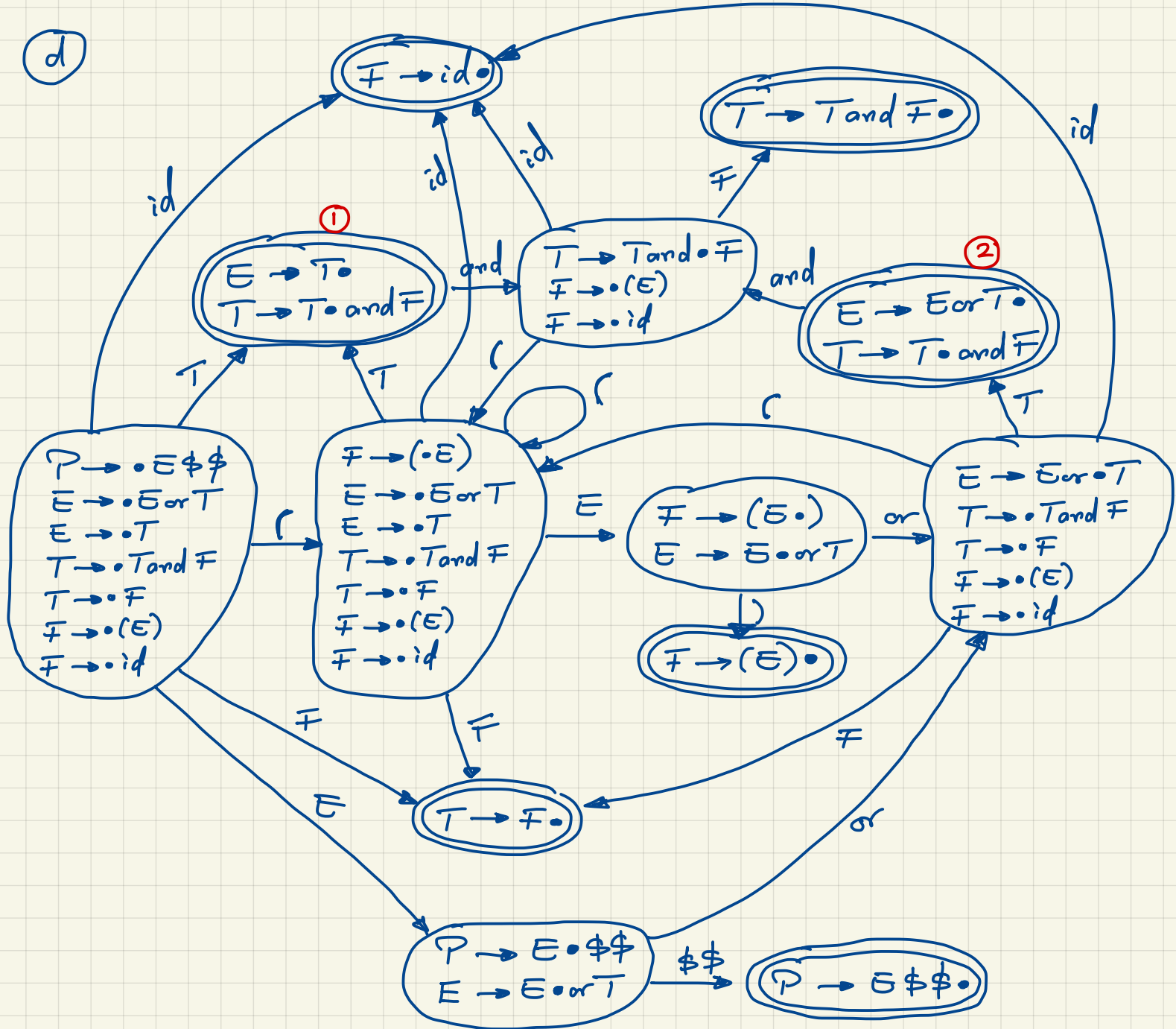
i	PREDICT(i)
1.	(, id
2.	(, id
3.	(, id
4.	(, id
5.	(, id
6.	(
7.	id

c)

conflicts

\Rightarrow not LL(1)

conflicts



There are two states with potential conflicts : ① and ②.

① : $\text{and} \notin \text{FOLLOW}(\epsilon) \Rightarrow \text{no conflict}$

② : same

$\Rightarrow \text{SLR}(1)$

Q3 a) We separate the arms and cases, so that an arm can have several cases. Attribute set stores the values, dup stores duplicates.

```

S      → switch ( expr ) { arm_list }
      ▷ S.dup = arm_list.dup
arm_list → case_list stmt more_arms
      ▷ arm_list.set = case_list.set ∪ more_arms.set
      ▷ arm_list.dup = case_list.dup ∪ more_arms.dup ∪ (case_list.set ∩ more_arms.set)
more_arms → arm_list
      ▷ more_arms.set = arm_list.set
      ▷ more_arms.dup = arm_list.dup
more_arms → ε
      ▷ more_arms.set = ∅
      ▷ more_arms.dup = ∅
case_list → case expr : more_cases
      ▷ case_list.set = {expr.val} ∪ more_cases.set
      ▷ case_list.dup = more_cases.dup ∪ ({expr.val} ∩ more_cases.set)
more_cases → case_list
      ▷ more_cases.set = case_list.set
      ▷ more_cases.dup = case_list.dup
more_cases → ε
      ▷ more_cases.set = ∅
      ▷ more_cases.dup = ∅

```

JFLAP : (q3.jff)

File Input Test Convert Help

Editor Build LL(1) Parse

Do Selected Do Step Do All Next Parse

Table Text Size

S	→	s(e){A}
A	→	CtM
M	→	A
M	→	ε
C	→	ce:N
N	→	C
N	→	ε

	FIRST	FOLLOW
A	{ c }	{ }
C	{ c }	{ t }
M	{ ε, c }	{ }
N	{ ε, c }	{ t }
S	{ s }	{ \$ }

	()	:	c	e	s	t	{	}	\$
A				CtM						
C				ce:N						
M				A					ε	
N				C			ε			
S						s(e){A}				

Table Text Size

no conflict
⇒ LL(1)

