

Stable Ramsey's theorem and measure

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Ramsey's theorem

Definition

Fix $n, k \in \mathbb{N}$ and let X be an infinite set.

- 1 $[X]^n = \{Y \subset X : |Y| = n\}.$
- 2 A k -coloring of exponent n is a map $f : [X]^n \rightarrow k = \{0, \dots, k-1\}.$
- 3 A set H is homogeneous for f if it is infinite and $f \upharpoonright [H]^n$ is constant.

Ramsey's theorem

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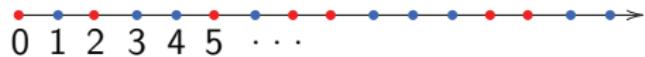
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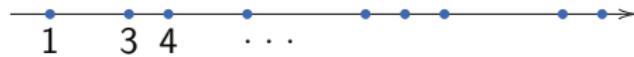
Fix $n, k \in \mathbb{N}$. Every coloring $f : [\mathbb{N}]^n \rightarrow k$ has a homogeneous set.

Ramsey's theorem



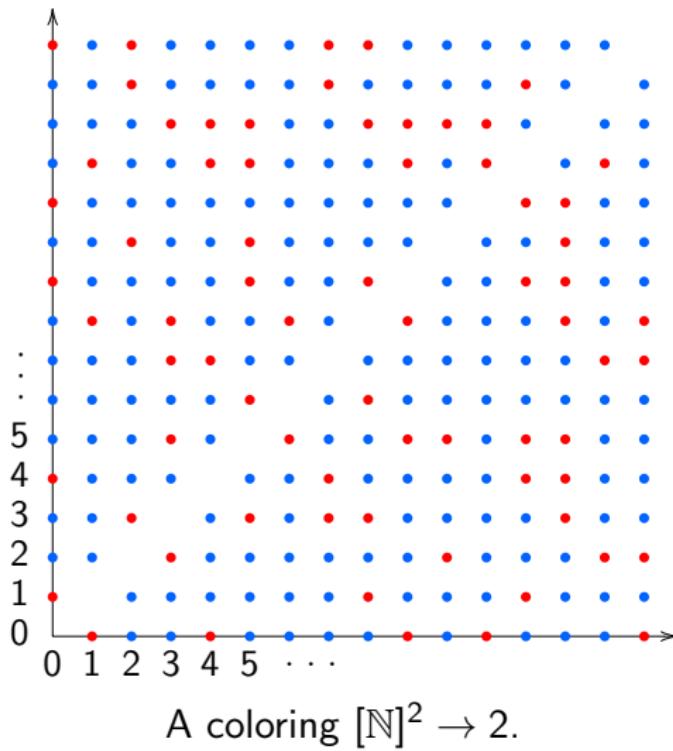
A coloring $[\mathbb{N}]^1 \rightarrow 2$.

Ramsey's theorem

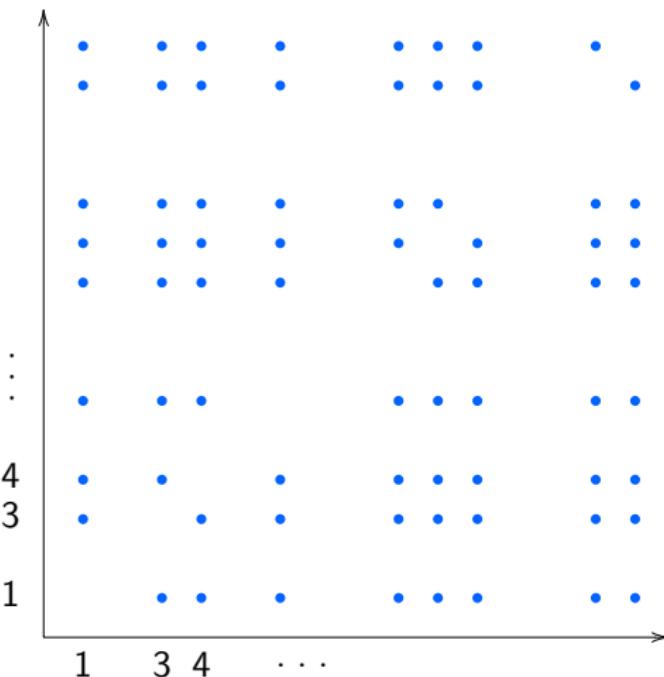


$H = \{1, 3, 4, 6, 9, 10, 11, 14, 15, \dots\}$ is homogeneous.

Ramsey's theorem



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Ramsey's theorem

Corollary (Chain Antichain Condition)

Every partial ordering of \mathbb{N} has an infinite chain or infinite antichain.

Ramsey's theorem

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Proof.

Let \preceq be a partial ordering of \mathbb{N} . Define $f : [\mathbb{N}]^2 \rightarrow 2$ by

$$f(\{x, y\}) = \begin{cases} 0 & \text{if } x \text{ and } y \text{ are comparable under } \preceq \\ 1 & \text{otherwise} \end{cases} .$$

Apply Ramsey's theorem to get a homogeneous set H for f .

If f has value 0 on $[H]^2$ then H is a chain under \preceq , and if f has value 1 on $[H]^2$ then H is an antichain. □

Ramsey's theorem

Moral

Complete disorder is impossible: in any configuration or arrangement of objects, however complicated or disorganized, some amount of structure and regularity is necessary.

Computability theory

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- Computability-theoretic approach: let's try to gauge the complexity, with respect to some hierarchy, of homogeneous sets.
- We restrict to *computable* colorings, i.e., $f : [\mathbb{N}]^n \rightarrow k$ for which there is an algorithm to compute $f(\{x_1, \dots, x_n\})$ given $x_1, \dots, x_n \in \mathbb{N}$.

Computability theory

Theorem (Specker, 1969)

There exists a comp. $f : [\mathbb{N}]^2 \rightarrow 2$ with no computable homogeneous set.

Computability theory

Definition (Turing jump)

Let X be a set.

- 1 The (*Turing*) *jump* of X is the set X' defined as

$$\{(e, x) : \text{the } e\text{th program with access to } X \text{ halts on input } x\}.$$

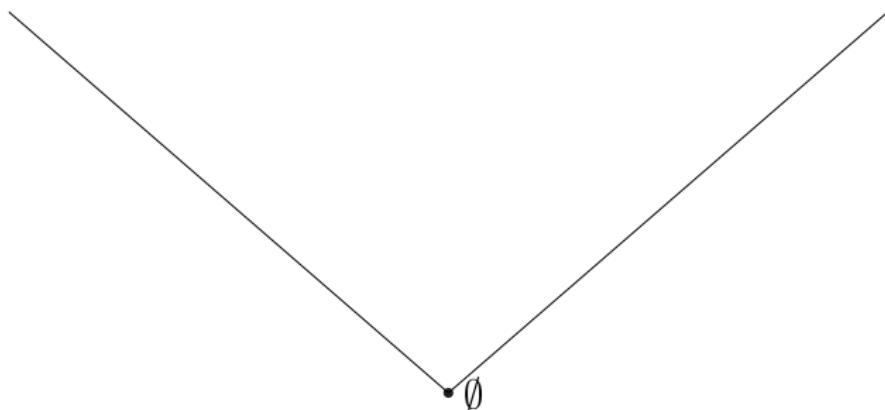
- 2 For $n \geq 1$, the $(n + 1)$ st jump of X is the jump of the n th.

Computability theory

- \emptyset

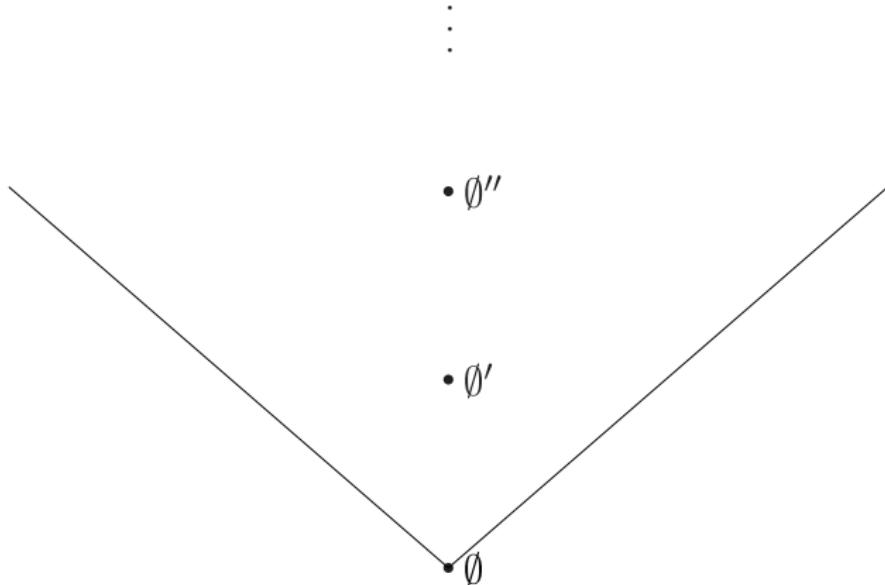
Turing hierarchy (\leq_T)

Computability theory



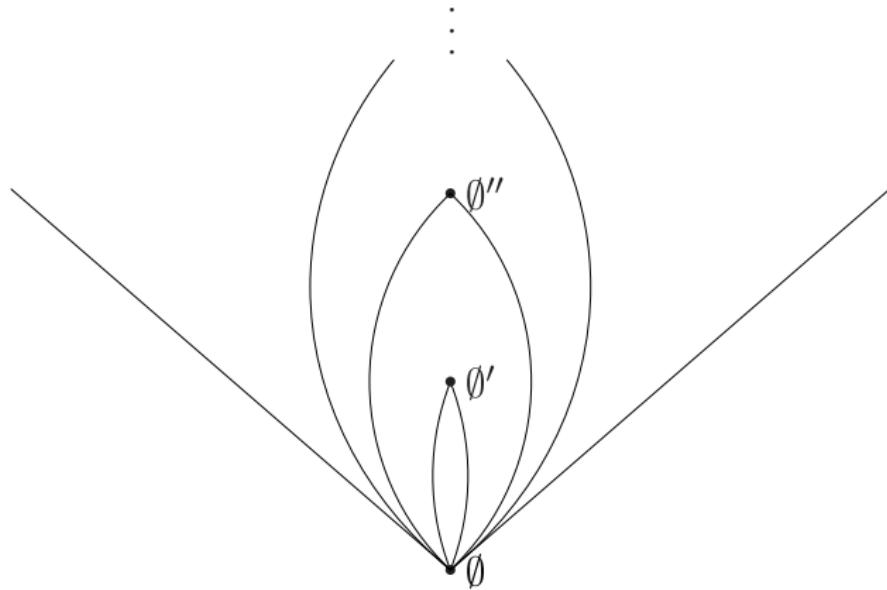
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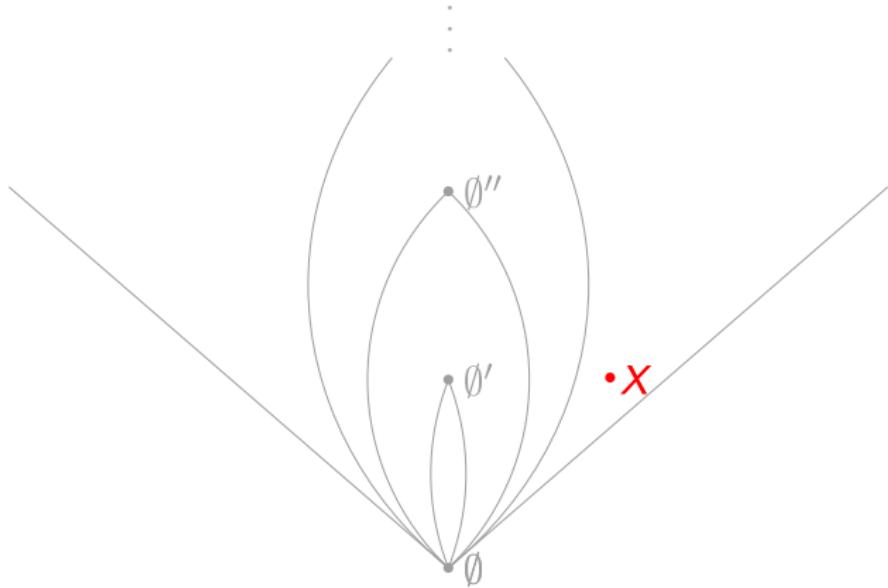
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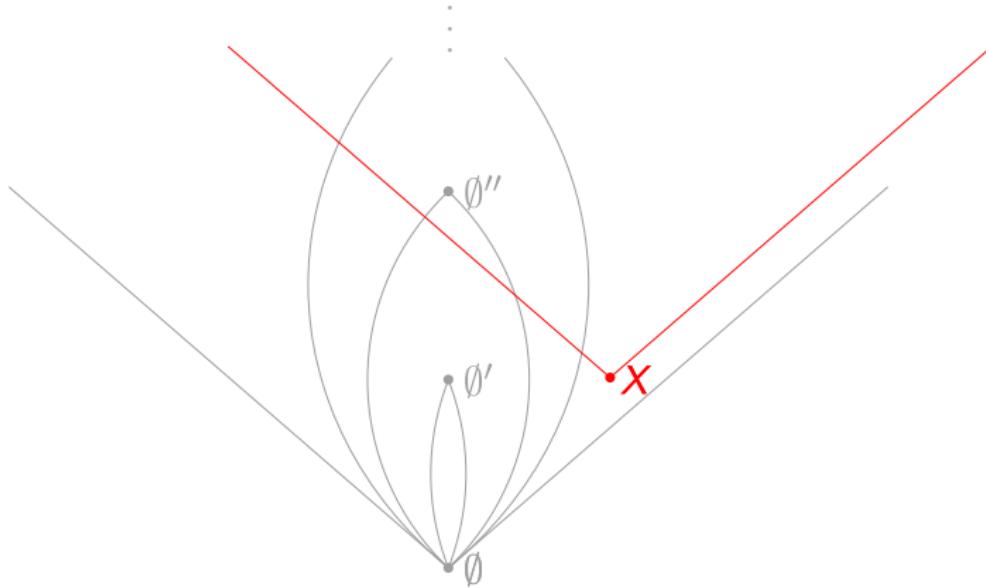
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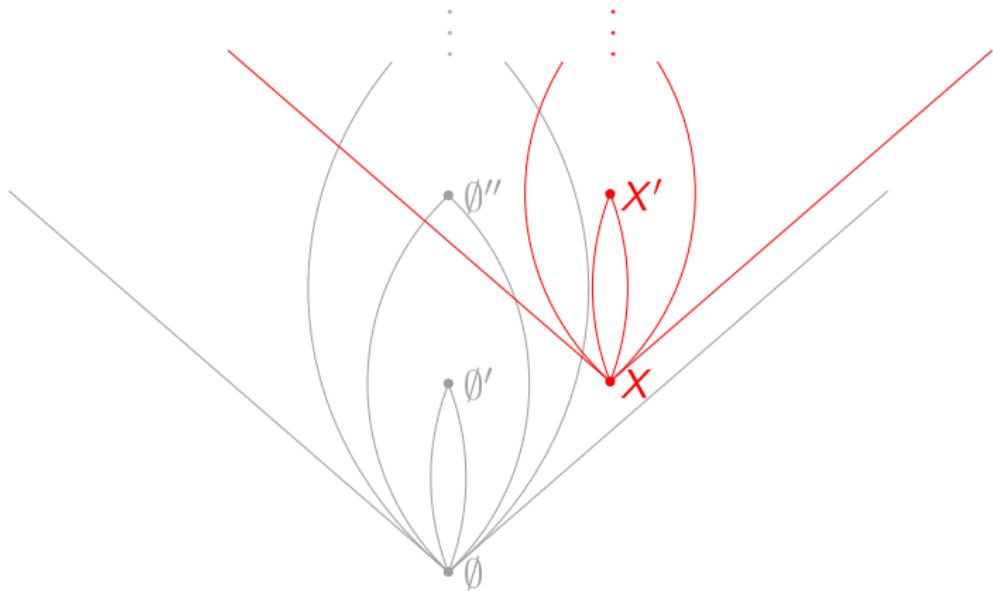
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Computability theory

Theorem (Jockusch, 1972)

For each $n \geq 3$, there exists a computable $f : [\mathbb{N}]^n \rightarrow 2$ such that $H \geq_T \emptyset'$ (in fact, $H \geq_T \emptyset^{(n-2)}$) for every homogeneous set H of f .

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- This reduces the investigation to the $n = 2$ case.

Theorem (Seetapun, 1995)

For every noncomputable set C , every computable $f : [\mathbb{N}]^2 \rightarrow k$ has a homogeneous set $H \not\geq_T C$.

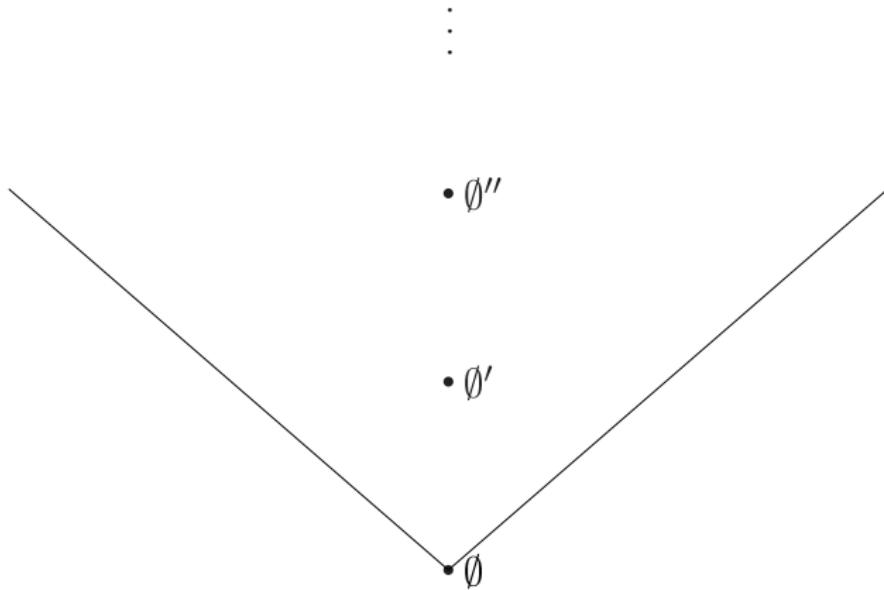
Computability theory

Definition (Low and low_2 sets)

A set X is

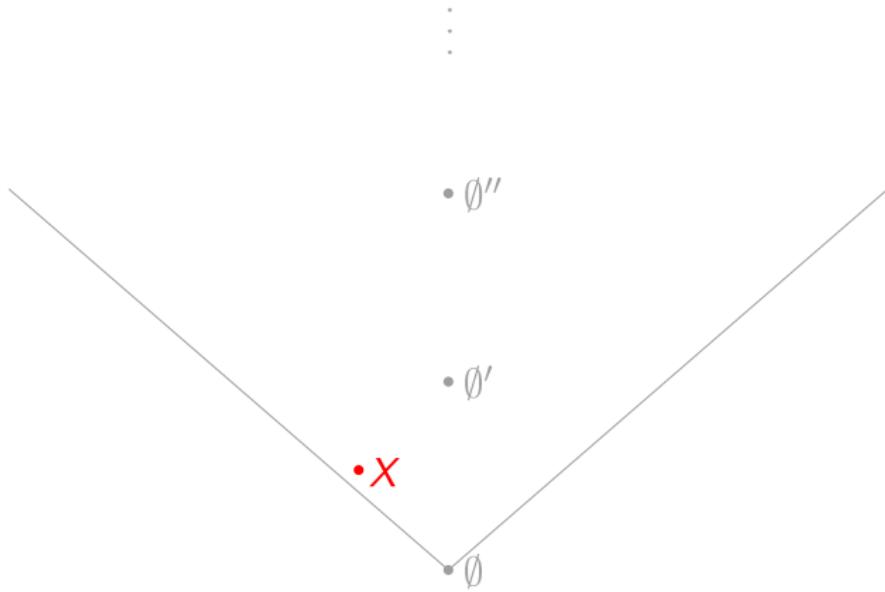
- 1** *low* if $X' \leq_T \emptyset'$;
- 2** *low₂* if $X'' \leq_T \emptyset''$.

Computability theory



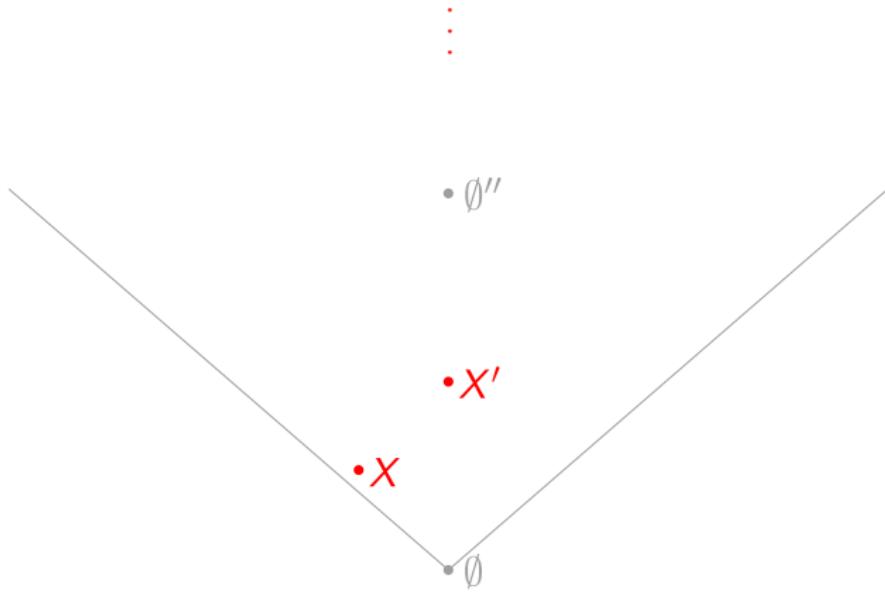
A low set.

Computability theory



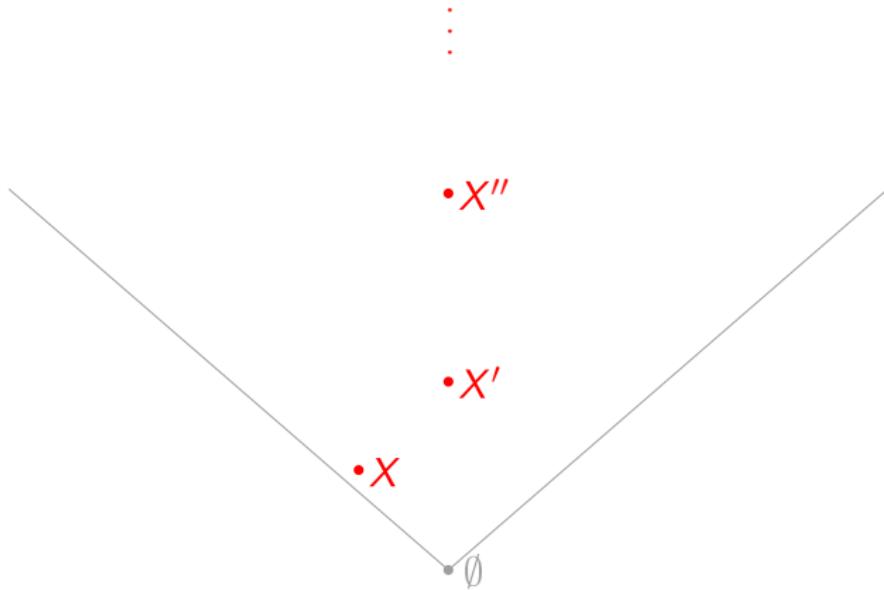
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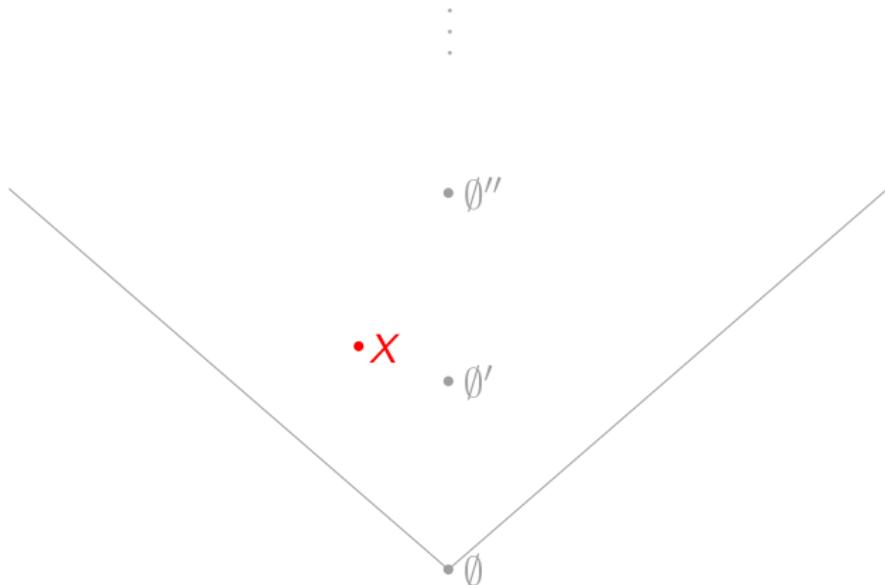
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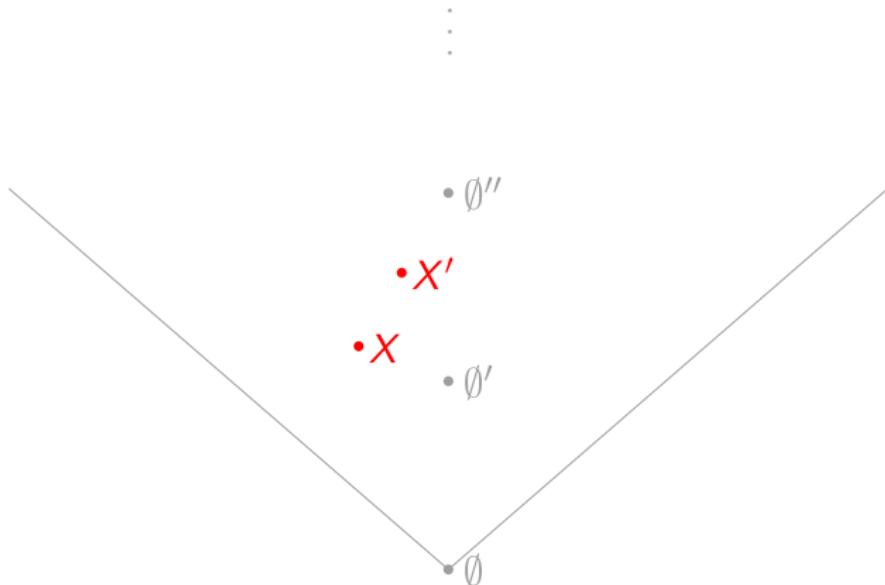
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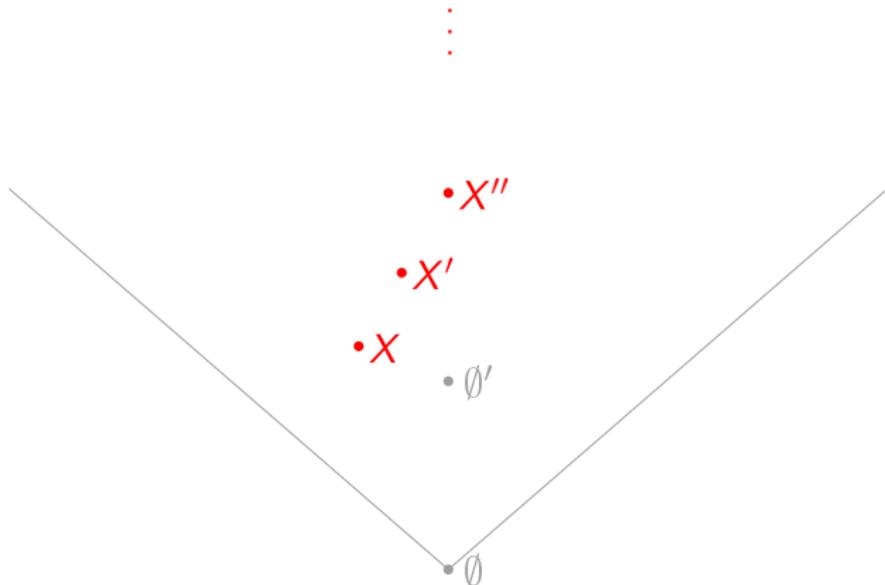
A low_2 set.

Computability theory



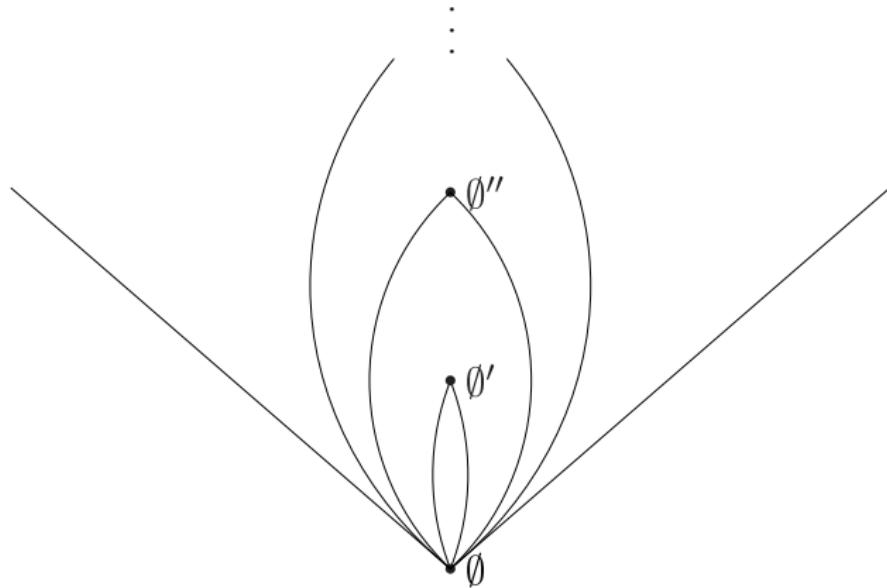
A low_2 set.

Computability theory



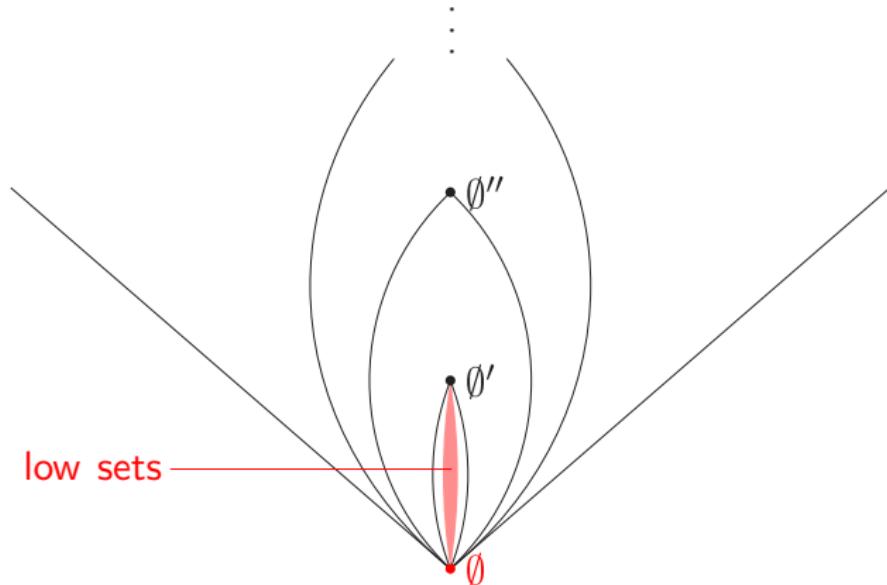
A low₂ set.

Computability theory



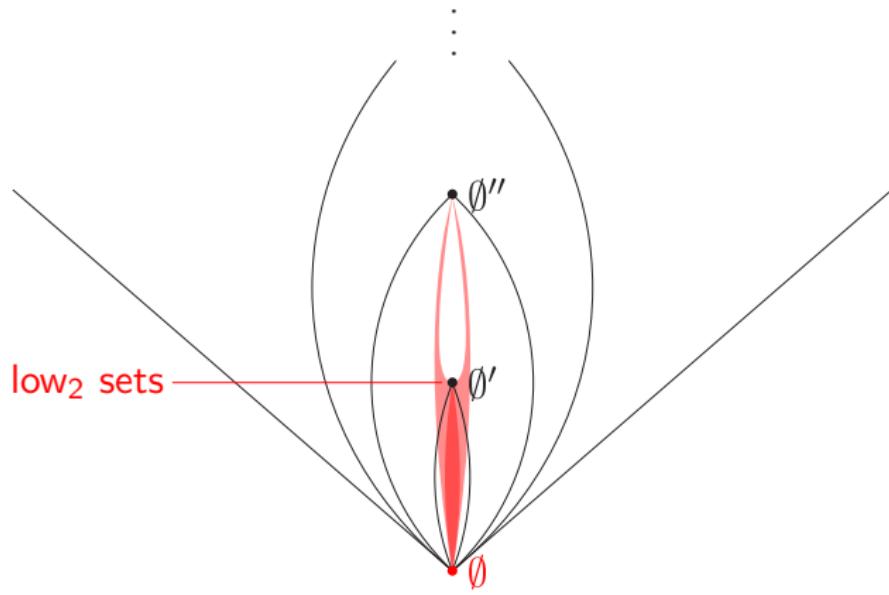
Lowness hierarchy.

Computability theory



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Theorem (Jockusch, 1972)

There exists a comp. $f : [\mathbb{N}]^2 \rightarrow 2$ with no low hom. set (or even $\leq_T \emptyset'$).

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Theorem (Cholak, Jockusch, and Slaman, 2001)

Every computable $f : [\mathbb{N}]^2 \rightarrow k$ has a low_2 homogeneous set H .

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Theorem (Cholak, Jockusch, and Slaman, 2001)

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Question (Cholak, Jockusch, and Slaman, 2001)

Can cone avoidance be added to this result?

Computability theory

Theorem (Dzhafarov and Jockusch, 2009)

For every noncomputable set C , every computable $f : [\mathbb{N}]^2 \rightarrow k$ has a low₂ homogeneous set that does not compute C .

Stable Ramsey's theorem

Definition

A coloring $f : [\mathbb{N}]^2 \rightarrow k$ is *stable* if for all x , $\lim_{y \rightarrow \infty} f(\{x, y\})$ exists.

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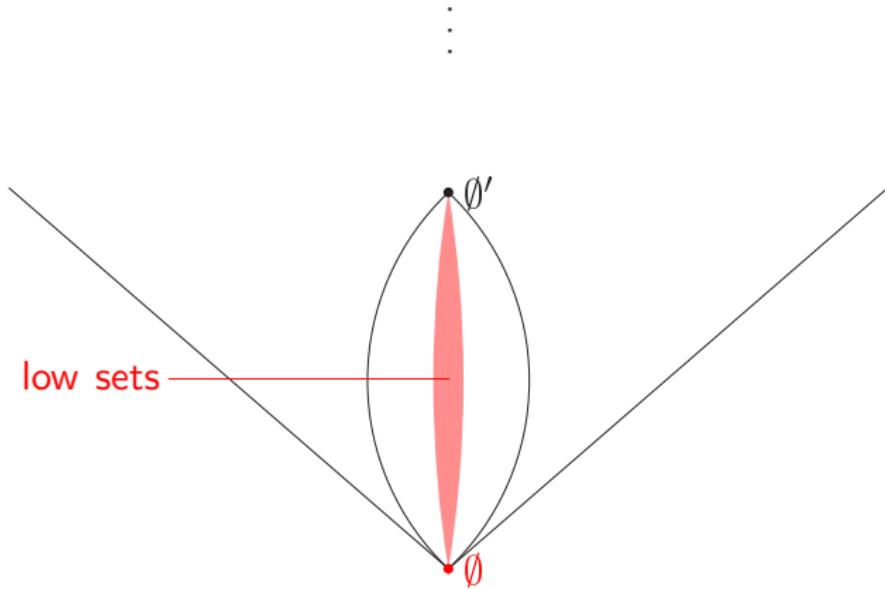
Theorem (Jockusch, 1972)

Every computable stable $f : [\mathbb{N}]^2 \rightarrow k$ has a homogeneous set $H \leq_T \emptyset'$.

Theorem (Downey, Hirschfeldt, Lempp, and Solomon, 2001)

There exists a computable stable $f : [\omega]^2 \rightarrow 2$ with no low hom. set.

Stable Ramsey's theorem



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Theorem (Hirschfeldt, Jockusch, Kjos-Hanssen, Lempp, and Slaman, 2008)

Every computable stable $f : [\omega]^2 \rightarrow 2$ has an homogeneous set $H <_T \emptyset'$.

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Theorem (Hirschfeldt, Jockusch, Kjos-Hanssen, Lempp, and Slaman, 2008)

Every computable stable $f : [\omega]^2 \rightarrow 2$ has an homogeneous set $H <_T \emptyset'$.

Theorem (Mileti, 2005)

There is no single set

- $X <_T \emptyset'$
- or X low₂

that computes a homogeneous set for every comp. stable $f : [\omega]^2 \rightarrow 2$.

Δ_2^0 measure

- Is there a way of telling which of the above results are “typical”?

Δ_2^0 measure

- Is there a way of telling which of the above results are “typical”?
- Measure-theoretic approach: define a notion of measure or nullity on the class of sets computable in \emptyset' .

Δ_2^0 measure

Definition (Martingales)

- 1 A *martingale* is a function $M : \{\text{finite binary strings}\} \rightarrow \mathbb{Q}^{\geq 0}$ such that for every string σ ,

$$M(\sigma) = \frac{M(\sigma0) + M(\sigma1)}{2}.$$

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- 2 M succeeds on a set X if

$$\limsup_{n \rightarrow \infty} M(X \upharpoonright n) = \infty.$$

Δ_2^0 measure

Theorem (Ville, 1939)

A class \mathcal{C} of reals has Lebesgue measure zero if and only if there is a martingale that succeeds on every $X \in \mathcal{C}$.

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Definition (Δ_2^0 nullity)

A class \mathcal{C} of reals has Δ_2^0 measure zero if there is a martingale M computable in \emptyset' that succeeds on every $X \in \mathcal{C}$.

Δ_2^0 measure

- Reasonable notion of measure on the class of sets computable in \emptyset' :
 - The class of all sets computable in \emptyset' does not have Δ_2^0 measure zero.
 - For each set X computable in \emptyset' , the singleton $\{X\}$ has Δ_2^0 measure zero.

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 - The class of all sets computable in \emptyset' does not have Δ_2^0 measure zero.
 - For each set X computable in \emptyset' , the singleton $\{X\}$ has Δ_2^0 measure zero.
- Gives us a notion of measure on the class of computable stable colorings.

Δ_2^0 measure

Theorem (Downey, Hirschfeldt, Lempp, and Solomon, 2001)

There exists a computable stable $f : [\omega]^2 \rightarrow 2$ with no low homogeneous set.

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Theorem (Hirschfeldt and Terwijn, 2008)

The class of computable stable colorings with no low homogeneous set has Δ_2^0 measure zero.

Δ_2^0 measure

Theorem (Mileti, 2005)

There is no single set $X <_T \emptyset'$ that computes a homogeneous set for every computable stable $f : [\omega]^2 \rightarrow 2$.

Δ_2^0 measure

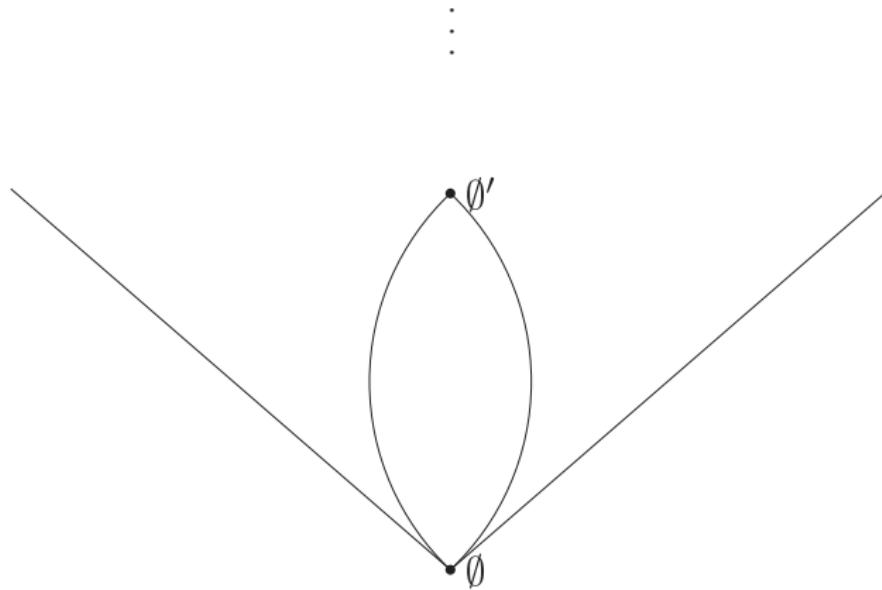
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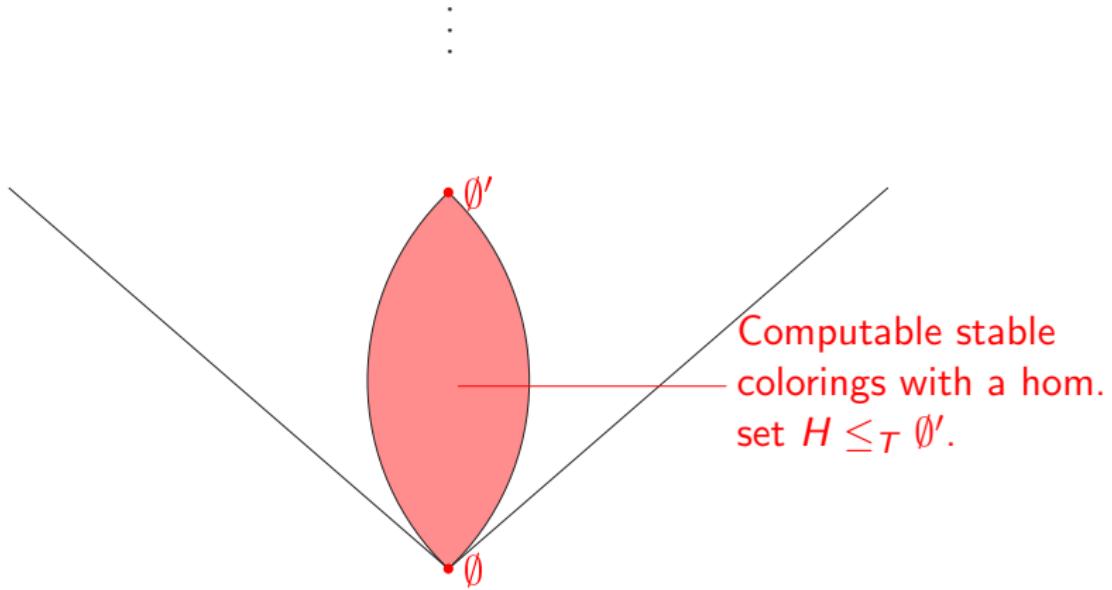
Theorem (Dzhafarov, submitted)

For each set $X <_T \emptyset'$, the class of computable stable colorings having a homogeneous set $H \leq_T X$ has Δ_2^0 measure zero.

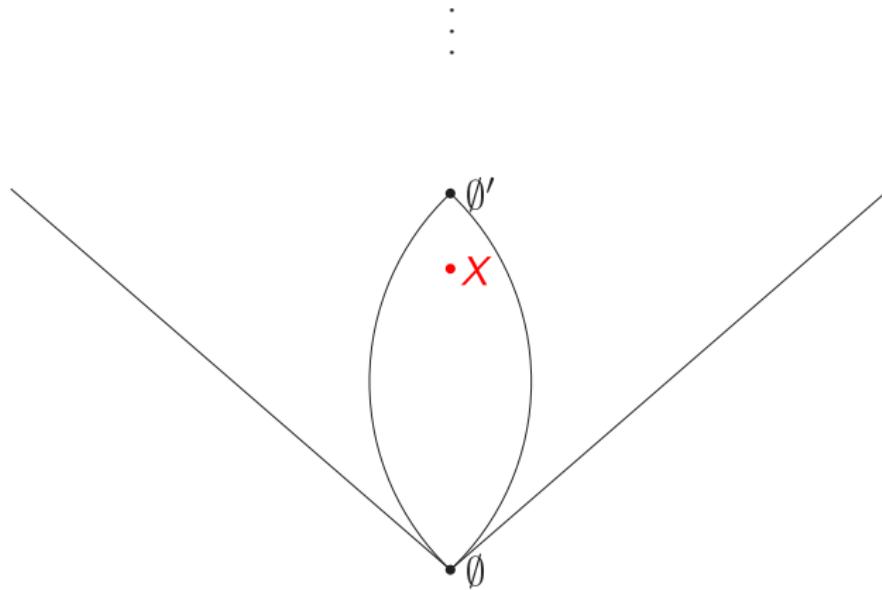
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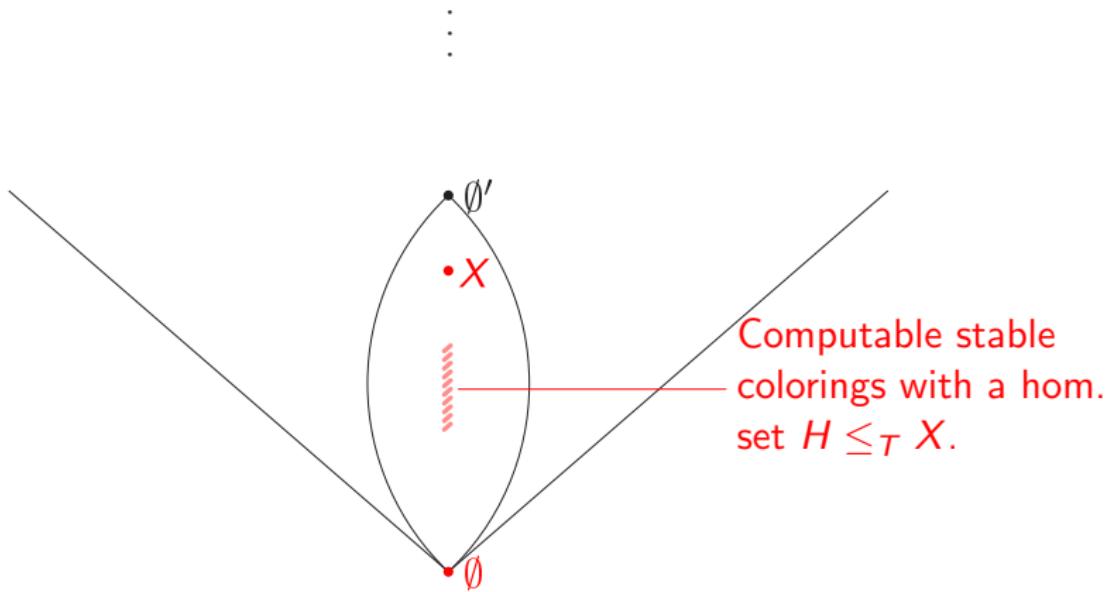
Δ_2^0 measure



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Theorem (Mileti, 2005)

There is no single low_2 set that computes a homogeneous set for every computable stable $f : [\omega]^2 \rightarrow 2$.

Δ_2^0 measure

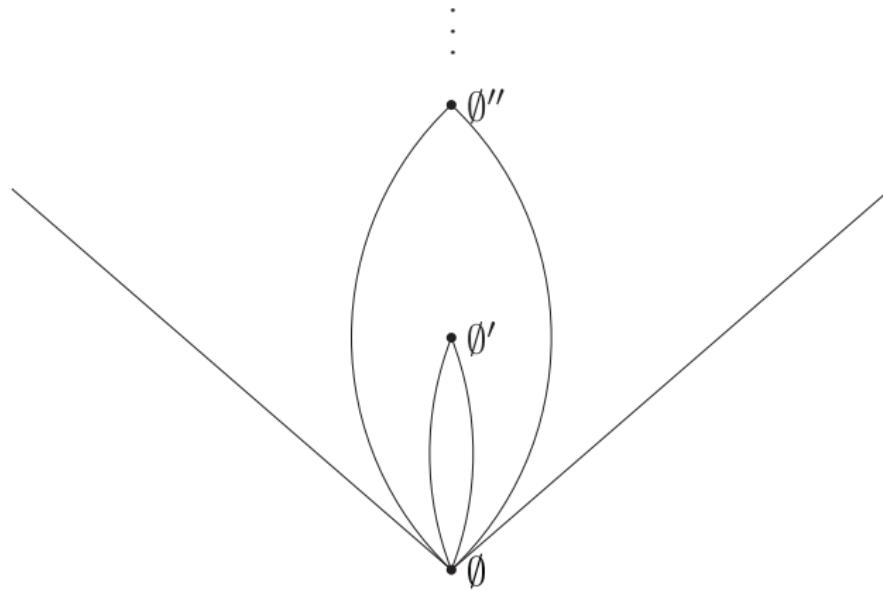
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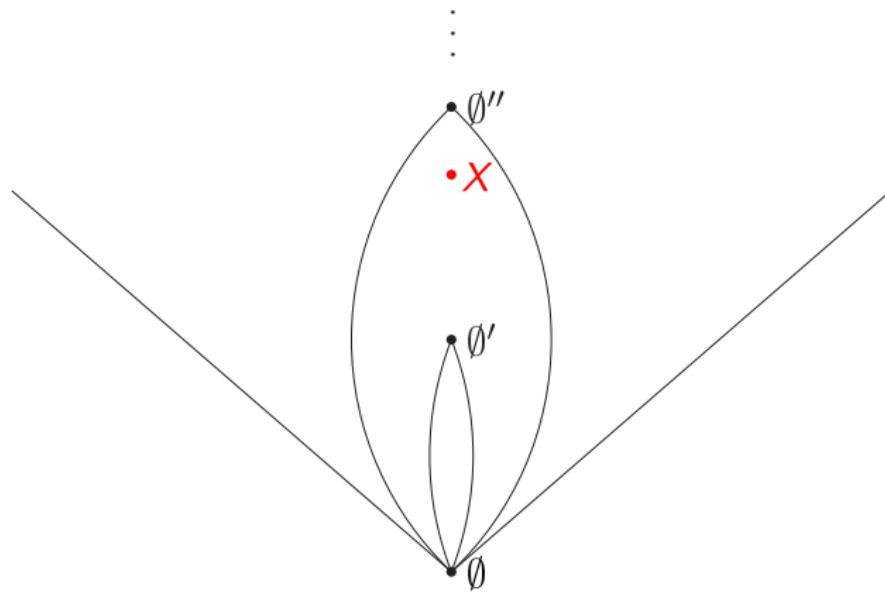
Theorem (Dzhafarov, submitted)

There is a set $X \leq_T 0''$ such that the class of computable stable colorings having a hom. set $H \leq_T X$ does not have Δ_2^0 measure zero but is not equal to the class of all computable stable colorings.

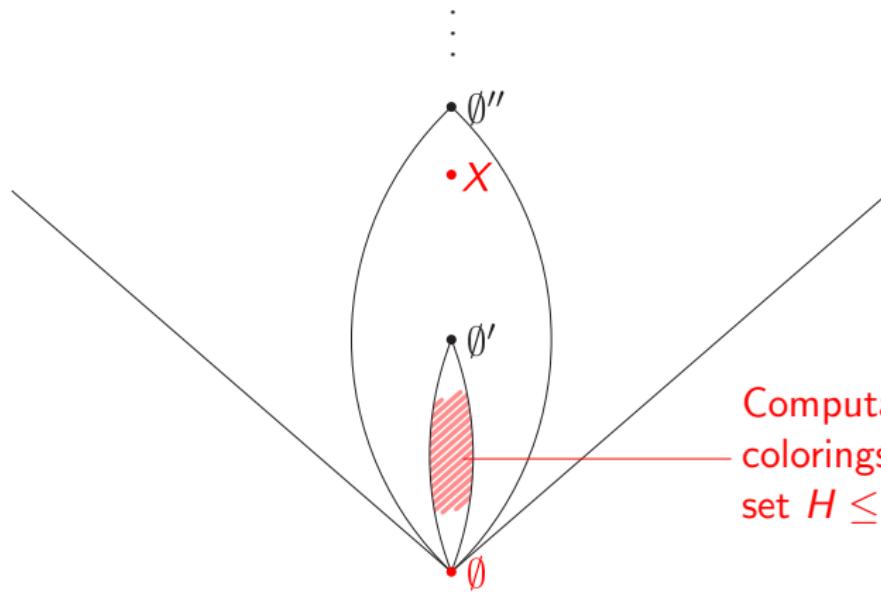
Δ_2^0 measure



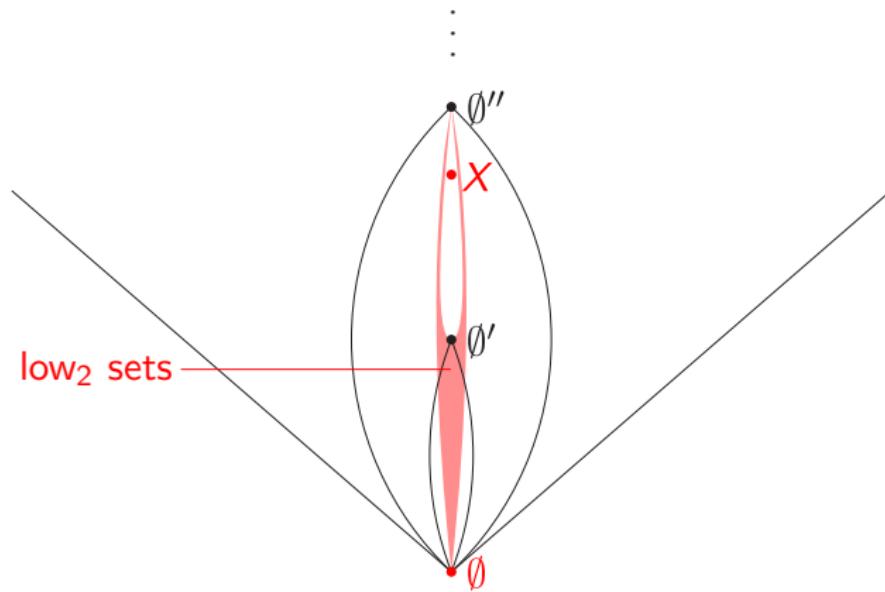
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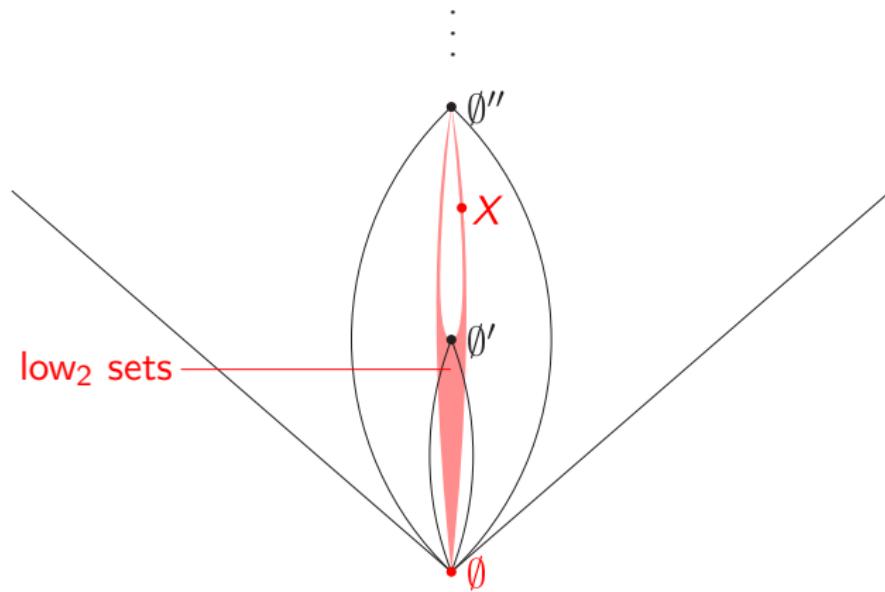


Δ_2^0 measure



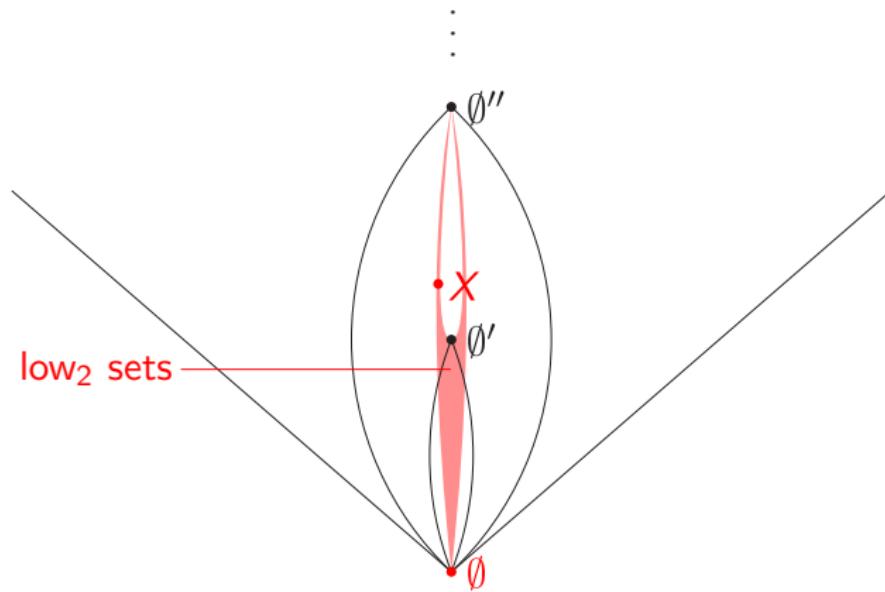
Open Question: Can X be chosen low₂?

Δ_2^0 measure



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Questions

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- What is the situation if instead of Δ_2^0 measure we use an effective notion of dimension?

Thank you for your attention.