Planning Graphs (Pre Lecture)

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Outline

Planning Graphs Construction Analysis

Planning with Planning Graphs
GraphPlan
GraphPlan+SATPlan (BlackBox)



Outline

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Planning with Planning Graphs GraphPlan GraphPlan+SATPlan (BlackBox



Planning Graph Overview

Nodes: literals ∪ actions

Edges: Transition: connects actions with precondition and effect literals,

 $(\ell \times a) \cup (a \times \ell)$

Mutex: conflicts (mutual exclusion) between actions and literals,

 $(\ell \times \ell) \cup (a \times a)$

Levels: Sequences of levels: timesteps







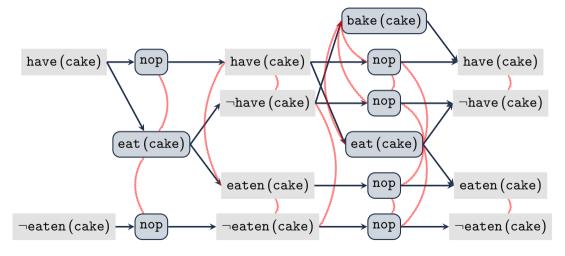
Example: Cake Domain

Operators

Facts



Example: Cake Planning Graph





Plan Graph Construction

- 1. Begin with literals for start state
- 2. Repeatedly add levels:
 - 2.1 Add persistence (nop) actions for each literal
 - 2.2 Add feasible actions
 - 2.3 Mark action mutexes
 - 2.4 Mark literal mutexes

until next level is same as prior level (fixpoint)



Action Mutexes

Conflicting Effect: One action's effect negates the other's effect,

- eff (eat-cake) = eaten-cake $\land \neg$ have-cake
- ▶ eff(nop(have-cake)) = have-cake
- ¬ (eff (eat-cake) ∧ eff (nop (have-cake)))

Conflicting Precondition: One action's precondition is mutexed with the other's precondition,

- pre(eat-cake) = have-cake
- ▶ pre (bake-cake) = ¬have-cake
- ¬ (pre (eat-cake) ∧ pre (bake-cake))

Interference: One action's effect negates the other's precondition,

- ightharpoonup eff (eat-cake) = \neg have-cake
- pre(nop(have-cake)) = have-cake
- ¬ (eff (eat-cake) ∧ nop (have-cake))





Literal Mutexes

Negation: One literal is the negation of the other,

- $ightharpoonup \neg (have-cake \land \neg have-cake)$
- $ightharpoonup \neg (eaten-cake \land \neg eaten-cake)$

Inconsistent Support: Each possible pair of actions to achieve both literals is mutually exclusive

- ► Step 1:
 - ▶ have-cake^[1] \implies nop (have-cake)^[0]
 - ightharpoonup eaten-cake^[1] \Longrightarrow eat-cake^[0]
 - ▶ conflicting effects: $\neg \left(\text{nop} \left(\text{have-cake} \right)^{[0]} \land \text{eat-cake}^{[0]} \right)$
- ► Step 2:
 - $\quad \blacktriangleright \; \mathsf{have-cake}^{[2]} \implies \left(\mathsf{nop}\left(\mathsf{have-cake}\right)^{[1]} \lor \mathsf{bake-cake}^{[1]}\right)$
 - $\hspace{0.2in} \blacktriangleright \hspace{0.2in} \mathtt{eaten-cake}^{[2]} \implies \left(\mathtt{nop} \left(\mathtt{eaten-cake} \right)^{[1]} \lor \mathtt{eat-cake}^{[1]} \right)$
 - ▶ non-conflicting: bake-cake^[1] \land nop (eaten-cake)^[1]



Exercise: Alternate Cake Domain



Exercise: Air Cargo

Operators

Facts

```
(define (problem air)
(define (domain air-cargo)
                                                         (:domain air-cargo)
  (: predicates (plane ?x) (cargo ?x)
                                                         (:objects cargo-0 cargo-1)
               (airport ?x) (at ?x ?y))
                                                                    plane-0 plane-1
  (: action fly : parameters (?p ?x ?y)
                                                                   ATL SFO)
           precondition
                                                         (: init (cargo cargo-0)
           (and (plane ?p) (airport ?x) (airport ?y)
                                                                 cargo cargo-1)
                (at ?p ?x))
                                                                 plane plane-0)
           : effect (and (not (at ?p ?x)) (at ?p ?y)))
                                                                 plane plane-1)
  (: action load : parameters (?c ?p ?a)
                                                                 airport ATL)
           precondition
                                                                 (airport SFO)
           (and (cargo ?c) (plane ?p) (airport ?a)
                                                                 at plane-0 ATL)
                (at ?c ?a) (at ?p ?a))
                                                                 (at plane-1 SFO)
           : effect (and (not (at ?c ?a)) (at ?c ?p)))
                                                                 at cargo = 0 ATL)
  (:action
           unload : parameters (?c ?p ?a)
                                                                 at cargo-1 SFO))
           : precondition
                                                         (: goal (and (at cargo-0 SFO)
           (and (cargo ?c) (plane ?p) (airport ?a)
                                                                      (at cargo-1 ATL)
                (at ?c ?p) (at ?p ?a))
           : effect (and (not (at ?c ?p)) (at ?c ?a))))
                                                                                      MINES
```

Exercise: Air Cargo



Termination of Planning Graph Construction

Theorem

Planning Graphs converge to a fixpoint in a finite number of steps.

Proof Outline

Graph elements increase or decrease monotonically over successive levels:

Literals increase monotonically: Can always persist a literal

Actions increase monotonically: preconditions remain satisfied at successive levels

Mutexes decrease monotonically: mutex at level *i* holds at all levels below *i*

Eventually, can add no more literals or actions and can remove no more mutexes.



Size of Planning Graphs

Theorem

Planning Graphs are polynomial in size of the planning domain.

Proof Outline

- p = |P| propositions, $\ell = 2p$ literals
- ightharpoonup a = |A| actions
- ► Each level:
 - \triangleright $a + \ell$ nodes
 - ▶ max $a * 2\ell$ transition edges (each action to every literal)
 - ▶ max $a^2 + \ell^2$ mutex edges (each action/literal mutex with every other)
- Polynomial number of levels due to monotonically increasing/decreasing elements



Interpreting of Planning Graphs

Feasibility

- ► A literal not in the final level (fixpoint) cannot be achieved of plan graph
- ► Mutexed literals: cannot both hold
 - ► What if goal literals are mutex at end?

Heuristics

- ► Cost to achieve literal: level of the graph
- Cost to achieve conjunction:

Max-level: Maximum cost of arguments

Level-sum: Sum costs of arguments

Set-level: Level where all hold



Analysis

Exercise: Planning Graph Heuristics



Outline

Planning with Planning Graphs GraphPlan GraphPlan+SATPlan (BlackBox)



Overview

- 1. Successively add levels to the planning graph
- 2. At each level.
 - 2.1 If the goals are not mutex, attempt to extract a plan from the graph
 - 2.2 If no plan can be extracted, continue growing the graph



Plan Extraction: GraphPlan Proper

Backwards Search from Final Level:

- 1. Start from last level from graph
- 2. Select conflict free actions from predecessor level to achieve current goal
- 3. New goal is precondition of selected actions
- 4. Repeat



Plan Extraction: GraphPlan + SATPlan (BlackBox Planner)

- 1. Construct planning graph
- 2. Convert planning graph to Boolean formula
- 3. Check SAT
- 4. If UNSAT, repeat

BlackBox vs. SATPlan proper: mutex information



SATPlan Encoding

Transition Function

Operator Encoding: Selected operator's preconditions and effects must hold:

$$o_i^{[k]} \implies \left(\begin{array}{c} \operatorname{precondition \ at \ step \ }^k & \operatorname{effect \ at \ step \ }^{k+1} \\ \operatorname{pre}(o_i)^{[k]} & \wedge & \operatorname{eff}(o_i)^{[k+1]} \end{array} \right)$$

Operator Exclusion: One operator per step:

$$o_i^{[k]} \implies (\neg o_0^{[k]} \wedge \neg o_{(i-1)}^{[k]} \wedge \neg o_{(i+1)}^{[k]} \wedge \neg o_m^{[k]})$$

Frame Axioms: Each proposition *p* is unchanged unless set by an effect:

$$\left(p^{[k]} = p^{[k+1]}\right) ee \left(o_j^{[k]} ee \ldots ee o_{\ell}^{[k]}
ight)$$



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SATPlan vs. Blackbox Encoding

Transition Function

SATPlan Operator Exclusion: One operator per step:

$$o_i^{[k]} \implies \left(\neg o_0^{[k]} \wedge \neg o_{(i-1)}^{[k]} \wedge \neg o_{(i+1)}^{[k]} \wedge \neg o_m^{[k]}\right)$$

BlackBox Operator Exclusion: One operator per step:

$$\left\{o_i^{[k]} \implies \neg o_j^{[k]} \mid o_i^{[k]} \text{ and } o_j^{[k]} \text{ are mutex}\right\}$$



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Blackbox Encoding

State Mutexes

$$\left\{ \neg \left(\ell_i^{[k]} \wedge \ell_j^{[k]} \right) \mid \ell_i^{[k]} \text{ and } \ell_j^{[k]} \text{ are mutex} \right\}$$

Additional constraint restricts search space.



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Summary

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