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1. Code Templates

1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration. xset r rate 150 100 xmodmap -e "remove Lock = Caps_Lock" -e "keysym Caps_Lock = Escape" -e "add Lock = Caps_Lock" echo "set nocp et sw=4 ts=4 sts=4 si cindent ru noeb showcmd showmode | syn on | colorscheme slate" > ~/ vimrc 1.2. C++ Header. A C++ header. #include <cstdlib>------// 11 #include <cstring>------// d0 6 10 11 12 13 15 #include <utility>-----// d8 17 using namespace std;-----// 7b _____// 7e 19 #define all(o) (o).begin(), (o).end()-----// 03 20 #define allr(o) (o).rbegin(), (o).rend()-----// 31 21 #define pb push_back-----// ec 22 const int INF = 2147483647;-----// 2c 23 const double EPS = 1e-9;-----// 45 24 const double pi = acos(-1);-----// ab 25 typedef long long ll;-----// 92 26 typedef unsigned long long ull;-----// 8a typedef pair<int, int> ii;------// 44 28 typedef vector<int> vi;-----// 20 typedef vector<ii> vii;-----// c9 30 typedef vector<vi>vvi;------// 58 31 typedef vector<vii>vvii;------// d5 32 template <class T> T mod(T a, T b) { return (a % b + b) % b; }------// 6e 33 template <class T> int size(const T &x) { return x.size(); }-----// 75 1.3. **Java Template.** A Java template. import java.util.*;-----// 37 import java.math.*;-----// 89 2 import java.io.*;-----// 28 -----// a3 4 public class Main {-----// 17 5 ----public static void main(String[] args) throws Exception {------// 02 6 -----Scanner in = new Scanner(System.in);------// ef -----PrintWriter out = new PrintWriter(System.out, false);------// 62 -----// code------// e6 9 -----out.flush();-----// 56 10 ----}------// 79 11 }-----// 00 2. Data Structures 2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure. **struct** union_find {------// 42 ----vi parent;--------// 1c ----union_find(int n) { parent.resize(cnt = n); for (int i = 0; i < cnt; i++) parent[i] = i; }-----// 94 ----int find(int i) { return parent[i] == i ? i : (parent[i] = find(parent[i])); }------// df ----bool unite(int i, int j) {------// 8a ------int ip = find(i), jp = find(j);-----// 36 ------parent[ip] = jp; return ip != jp; } };------// f8

```
2.2. Segment Tree. An implementation of a Segment Tree.
  // const int ID = INF;-----// d2
  // int f(int a, int b) { return min(a, b); }-----// 0c
  const int ID = 0; ------// 82
int f(int a, int b) { return a + b; }-----// 6d
4
  struct segment_tree {------// e5
5
  ----int n; vi data;------// 1c
6
  ----segment_tree(const vi &arr) : n(size(arr)), data(4*n) { mk(arr, 0, n-1, 0); }------// 8a
  9
  ------int m = (l + r) / 2;-----// 54
10
  -----return data[i] = f(mk(arr, l, m, 2*i+1), mk(arr, m+1, r, 2*i+2)); }------// 4d
11
  13
  -----if (a <= l && r <= b) return data[i];------// 81
15
  ------int m = (l + r) / 2;-----// 75
16
  -----return f(q(a, b, l, m, 2*i+1), q(a, b, m+1, r, 2*i+2)); }------// 38
17
  18
19
  20
  -----if (l == i && r == i) return data[j] = v;------// 59
21
  ------int m = (l + r) / 2;-------// 9e
22
  ------return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); } };------// 84
  2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It supports adjusting the
  i-th element in O(\log n) time, and computing the sum of numbers in the range i..j in O(\log n) time. It only needs O(n)
  struct fenwick_tree {------// 98
1
  ----int n; vi data;-------// d3
  ----fenwick_tree(<u>int</u> _n) : n(_n), data(vi(n)) { }-----// db
3
  ----void update(int at, int by) { while (at < n) data[at] += by, at |= at + 1; }------// 91
  ----int query(int at) {------// 87
5
  -----int res = 0;-------// 97
6
  -------while (at >= 0) res += data[at], at = (at & (at + 1)) - 1;-------// c6
  ------return res; }------// fb
8
  ----int rsq(int a, int b) { return query(b) - query(a - 1); }------// c8
9
  };-----// c4
10
  struct fenwick_tree_sq {------// 09
11
  ----int n; fenwick_tree x1, x0;--------// 7b
12
  ----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)), x0(fenwick_tree(n)) { }------// 30
13
  ----// insert f(y) = my + c if x <= y------// d3
14
  ----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }-----// 77
15
  ----int query(int x) { return x*x1.query(x) + x0.query(x); }-----// d1 };------// 39
16
17
  void range_update(fenwick_tree_sq &s, int a, int b, int k) {------// cb
18
  ----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 6d
19
  int range_query(fenwick_tree_sq &s, int a, int b) {------// c2
20
  ----return s.query(b) - s.query(a-1); }-----// ec
21
  2.4. Matrix. A Matrix class.
  template <class K> bool eq(K a, K b) { return a == b; }-----// 2a
  template <> bool eq<double a, double b) { return abs(a - b) < EPS; }-----// a7
2
  4
  private:-----// 03
  ----vector<T> data;------// af
6
  ----int cnt;-------// 41
  ----inline T& at(int i, int j) { return data[i * cols + j]; }------// a4
8
  public:----// bc
  ----int rows, cols;-------// a8
10
  ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) { data.assign(cnt, T(0)); }------// 06
11
    matrix(const matrix& other) : rows(other.rows), cols(other.cols), cnt(other.cnt), data(other.data) { }
12
    -Τω operator()(int i, int j) { return at(i, j); }------// θε
13
    void operator += (const matrix \& other) { for (int i=0; i< cnt; i++) data[i] += other.data[i]; }
    void operator -= (const matrix\& other) { for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }
15
    -<mark>void</mark> operator *=(T other) {    for (int i = 0; i < cnt; i++) data[i] *= other; }------// 9e
16
  ----matrix<T> operator +(const matrix& other) { matrix<T> res(*this); res += other; return res; }---// 22
17
  ----matrix<T> operator -(const matrix\& other) { matrix<T> res(*this); res -= other; return res; }---// c\theta
```

18

```
----matrix<T> operator *(T other) { matrix<T> res(*this); res *= other; return res; }------// 8c
19
 ----matrix<T> operator *(const matrix& other) {------// 56
20
 -----matrix<T> res(rows, other.cols);------// 3d
21
 ------for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)------// b\theta
22
 -----for (int k = 0; k < cols; k++) res(i, j) += at(i, k) * other.data[k * other.cols + j];--// cc
23
 -----return res; }------// ae
24
 ----matrix<T> transpose() {-----// e7
 ------matrix<T> res(cols, rows);-------// 74
26
 ------for (int i = 0; i < rows; i++) for (int j = 0; j < cols; j++) res(j, i) = at(i, j);------// 2d
 -----return res; }------// b0
28
 ----matrix<T> pow(int p) {------// ed
 ------matrix<T> res(rows, cols), sq(*this);------// 80
30
 ------for (int i = 0; i < rows; i++) res(i, i) = T(1);------// a0
31
 ------while (p) {------// 6f
32
 ------if (p & 1) res = res * sq;------// b7
33
 -----p >>= 1,------// 88
34
 -----if (p) sq = sq * sq;-----// d3
35
 36
 ----matrix<T> rref() {------// 6d
37
 -----matrix<T> mat(*this);-----// e9
 ------for (int r = 0, c = 0; c < cols; c++) {-------// 48
39
 ------int k = r;-----// 33
40
 -----while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// ef
41
 -----if (k >= rows) continue;------// 78
42
 -----if (k != r) for (int i = 0; i < cols; i++) swap(mat.at(k, i), mat.at(r, i));------// 5f
43
 ------if (!eq<T>(mat(r, c), T(1))) for (int i = cols-1; i >= c; i--) mat(r, i) /= mat(r, c);--// 16
44
 -----for (int i = 0; i < rows; i++) {------// d1
45
 -----T m = mat(i, c);------// 00
46
        if (i != r \& \& !eq<T>(m, T(0))) for (int j = 0; j < cols; j++) mat(i, j) -= m * mat(r, j);
47
 ------} r++;-------// e7
48
 49
 };-----// 0f
50
 2.5. Trie. A Trie class.
 template <class T>------// 82
1
 class trie {------// 9a
2
 private:-----// f4
3
 ----struct node {-------// ae
 ------map<T, node*> children;------// a0
5
 ------int prefixes, words;-------// e2
6
 -----node() { prefixes = words = 0; } };------// 42
 public:----// 88
8
 ----node* root;------// a9
9
 ----trie() : root(new node()) { }-----// 8f
10
 ----template <class I>------// 89
 ----void insert(I begin, I end) {------// 3c
12
 -----node* cur = root;-------// 82
13
 14
 15
 -----if (begin == end) { cur->words++; break; }-----// db
16
 -----else {------// 3e
17
 -----T head = *beqin;-----// fb
  -----typename map<T, node*>::const_iterator it = cur->children.find(head);------// 07
19
        if (it == cur->children.end()) it = cur->children.insert(pair<T, node*>(head, new node())).first;
20
 -----begin++, cur = it->second; } } }-----// e1
21
 ----template<class I>------// 9e
 ----int countMatches(I begin, I end) {------// cf
23
 -----node* cur = root;------// e2
24
 ------while (true) {------// aa
25
 -----if (begin == end) return cur->words;-----// ef
26
 -----else {------// b5
27
 -----T head = *begin;------// 15
28
 -----typename map<T, node*>::const_iterator it = cur->children.find(head);------// 5d
29
 -----if (it == cur->children.end()) return 0;------// 1e
30
 32
 ----int countPrefixes(I begin, I end) {------// f2
 34
 ------while (true) {--------// 6f
```

```
36
 -----else {------// 16
37
 -----T head = *begin;------// b3
38
 -----typename map<T, node*>::const_iterator it = cur->children.find(head);------// 82
39
 -----if (it == cur->children.end()) return 0;------// 84
40
 -----begin++, cur = it->second; } } };-----// 96
41
 2.6. AVL Tree. A fast, easily augmentable, balanced binary search tree.
 #define AVL_MULTISET 0------// b5
1
 2
 3
 class avl_tree {------// ff
4
 public:-----// f6
5
 ----struct node {------// 45
 -----T item; node *p, *l, *r;------// a6
7
 ------<mark>int</mark> size, height;------// 33
    node(\textbf{const} \ T \ \& item, \ node \ *p = NULL) \ : \ item(item), \ p(p), \ l(NULL), \ r(NULL), \ size(1), \ height(0) \ \{ \ \} \ \};
9
 ----avl_tree() : root(NULL) { }-------// 95
10
 11
 ----node* find(const T &item) const {------// ab
12
 -----node *cur = root;------// 94
13
 -------while (cur) {-------// 2b
14
 -----if (cur->item < item) cur = cur->r;------// f2
 16
 -----else break; }-----// db
 -----return cur; }-----// 6f
18
 ----node* insert(const T &item) {------// c9
19
 -----node *prev = NULL, **cur = &root;-----// c7
20
 ------while (*cur) {--------// 60
21
 -----prev = *cur;------// 8c
22
 ------if ((*cur)->item < item) cur = &((*cur)->r);------// 21
23
 #if AVL MULTISET-----// 7f
24
 -----else cur = &((*cur)->l);-----// 95
25
 ------else if (item < (*cur)->item) cur = &((*cur)->l);------// 6b
27
 -----else return *cur;------// 85
28
 29
 30
 -----node *n = new node(item, prev);-----// 73
31
 -----*cur = n, fix(n); return n; }------// b5
32
 ----void erase(const T &item) { erase(find(item)); }-----// 7a
33
 ----void erase(node *n, bool free = true) {------// 75
34
 -----if (!n) return;------// 85
35
 -----if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;------// 95
36
 ------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;-------// 75
 -----else if (n->l && n->r) {------// ea
38
 -----node *s = successor(n);-----// e3
39
 -----erase(s, false);-----// ed
40
 ------s->p = n->p, s->l = n->l, s->r = n->r;------// dc
41
 -----if (n->l) n->l->p = s;-----// ed
42
 43
 -----parent_leg(n) = s, fix(s);-----// cd
44
 ------return;-------// 75
45
 -----} else parent_leg(n) = NULL;-----// 11
46
 -----fix(n->p), n->p = n->l = n->r = NULL; 64
47
 -----if (free) delete n; }------// 03
48
 ----node* successor(node *n) const {-----// 39
49
 -----if (!n) return NULL;------// 99
50
 -----if (n->r) return nth(0, n->r);------// 27
51
 ------node *p = n->p;------// 46
52
 -------while (p && p->r == n) n = p, p = p->p;------// fd
53
 -----return p; }-----// 92
54
 ----node* predecessor(node *n) const {------// 10
55
 56
 -----if (n->l) return nth(n->l->size-1, n->l);------// 39
-----node *p = n->p;------// c8
58
 -------while (p && p->l == n) n = p, p = p->p;------// ce
 -----return p; }------// 76
60
 ----inline int size() const { return sz(root); }------// 19
```

```
----void clear() { delete_tree(root), root = NULL; }------// e5
62
  ----node* nth(int n, node *cur = NULL) const {------// 2c
63
  ------if (!cur) cur = root;------// c1
64
  ------while (cur) {-------// bb
65
  -----if (n < sz(cur->l)) cur = cur->l;------// 69
66
  ------else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-------// fb
67
  69
  private:-----// f6
  ----inline int sz(node *n) const { return n ? n->size : 0; }------// 5c
71
  ----inline int height(node *n) const { return n ? n->height : -1; }-------// 48
  ----inline bool left_heavy(node *n) const { return n && height(n->l) > height(n->r); }-------// ef
73
  ----inline bool right_heavy(node *n) const { return n && height(n->r) > height(n->l); }------// 3d
74
  ----inline bool too_heavy(node *n) const { return n \&\& abs(height(n->l) - height(n->r)) > 1; }-----// 83
75
  ----void delete_tree(node *n) { if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }------// e2
76
  77
78
  ------if (n->p->l == n) return n->p->l;------// dc
79
  ------if (n->p->r == n) return n->p->r;------// ce
80
  -----assert(false); }------// 7c
----void augment(node *n) {------// bd
82
  -----if (!n) return;------// 65
83
  -----n->size = 1 + sz(n->l) + sz(n->r);-------// c5
84
  -----n->height = 1 + max(height(n->l), height(n->r)); }-----// 4b
85
  86
87
  -----l->p = n->p; N-----// a4
88
  -----parent_leg(n) = l; \[ \] ------// f5
89
  ------n->l = l->r; N------// de
90
  91
  -----l->r = n, n->p = l; N------// 86
92
  -----augment(n), augment(l)------// 46
93
  95
  ----void fix(node *n) {-------// 70
-----while (n) { augment(n);------// 1d
97
  ------if (too_heavy(n)) {------// 13
98
  -----if (left_heavy(n) && right_heavy(n->l)) left_rotate(n->l);------// c5
99
  ------else if (right_heavy(n) && left_heavy(n->r)) right_rotate(n->r);------// 11
100
  -----if (left_heavy(n)) right_rotate(n);-------// 7e
-----else left_rotate(n);-----// 00
101
102
  -----n = n->p; }-----// c7
103
  -----n = n->p; } } };------// 2f
104
   Also a very simple wrapper over the AVL tree that implements a map interface.
  #include "avl_tree.cpp"-----// 01
  -----// ba
2
  template <class K, class V>-----// da
  4
  public:-----// 5d
  ----struct node {------// 2f
6
  ------K key; V value;------// 32
  -----node(K k, V v) : key(k), value(v) { }------// 29
  ------bool operator <(const node &other) const { return key < other.key; } };------// 92
9
  ----avl_tree<node> tree;------// b1
----V& operator [](K key) {------// 7c
10
11
  -----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
12
  ------if (!n) n = tree.insert(node(key, V(0)));-------// cb
-----return n->item.value;-----// ec
13
14
  ----}------// 2e
15
  };-----// af
  2.7. Heap. An implementation of a binary heap.
  #define RESIZE------// d0
  #define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
  struct default_int_cmp {------// 8d
  ----default_int_cmp() { }-----// 35
  ----bool operator ()(const int &a, const int &b) { return a < b; } };------// e9
```

```
template <class Compare = default_int_cmp>------// 30
6
 class heap {------// 05
 private:-----// 39
8
 ----<mark>int</mark> len, count, *q, *loc, tmp;------// 0a
9
 ----Compare _cmp;------// 98
10
 ----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }------// a0
11
 ----inline void swp(int i, int j) { SWP(q[i], q[j]), SWP(loc[q[i]], loc[q[j]]); }------// f6
12
 13
 -----int p = (i - 1) / 2;------// 52
15
 17
 ----void sink(int i) {------// 12
18
 ------while (true) {------// 28
19
 -----int l = 2*i + 1, r = l + 1;-----// 0f
20
 -----if (l >= count) break;-----// 7c
21
 ------int m = r >= count || cmp(l, r) ? l : r;-----// ae
22
 -----if (!cmp(m, i)) break;-----// 63
23
 24
 public:-----// 9b
 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {-------// 3e
26
 -----q = new int[len], loc = new int[len];-----// d7
27
 -----memset(loc, 255, len << 2); }-----// 89
28
 ----~heap() { delete[] q; delete[] loc; }------// a3
29
 ----void push(int n, bool fix = true) {--------// 37
30
 -----if (len == count || n >= len) {-------// 6f
31
 #ifdef RESIZE-----// ba
32
 -----int newlen = 2 * len;------// f2
33
 ------while (n >= newlen) newlen *= 2;------// 11
 ------int *newq = new int[newlen], *newloc = new int[newlen];-----// d6
35
 ------for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i];------// f5
 -----memset(newloc + len, 255, (newlen - len) << 2);------// da
37
 -----delete[] q, delete[] loc;-----// 86
38
 -----loc = newloc, q = newq, len = newlen;-----// 48
39
 #else-----// 8c
40
 -----assert(false);-------// 3d
41
 #endif------// 2f
42
 43
 -----assert(loc[n] == -1);------// 71
44
 -----loc[n] = count, q[count++] = n;------// bc
45
 -----if (fix) swim(count-1); }------// d7
46
 ----void pop(bool fix = true) {------// d7
47
 ------assert(count > 0);------// fc
48
 -----loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 5a
49
 ------if (fix) sink(0);-------// fa
50
51
 ----int top() { assert(count > 0); return q[0]; }-----// d5
52
   void heapify() { for (int i = \text{count } -1; i > 0; i - -) if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }
53
 ----void update_key(int n) { assert(loc[n] != -1), swim(loc[n]), sink(loc[n]); }------// b8
54
 ----bool empty() { return count == 0; }------// 17
55
 ----int size() { return count; }-------// f8
 57
 2.8. Skiplist. An implementation of a skiplist.
 #define BP 0.20------// aa
1
 #define MAX_LEVEL 10------// 56
 unsigned int bernoulli(unsigned int MAX) {-----// 7b
3
 ----unsigned int cnt = 0;------// 28
4
 ----while(((float) rand() / RAND_MAX) < BP && cnt < MAX) cnt++;------// d1
5
 ----return cnt; }------// a1
6
 7
 8
 9
 -----int *lens:------// 07
10
 -----node **next;------// 0c
 ------#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))-----// 25
12
 ------node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}------// 7c
 ------node() { free(lens); free(next); }; };------// aa
14
 ----int current_level, _size;-------// 61
```

```
----node *head:-------// b7
16
    ----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };------// 7a
17
    18
19
    20
21
    -----node *x = head; \[\ni_------// 0f\]
22
    -----node *update[MAX_LEVEL + 1]; N-----// 01
23
    ------memset(update, 0, MAX_LEVEL + 1); N------// 38
24
    25
    -----pos[i] = pos[i + 1]; N------// 68
26
    27
    28
    29
    ----int size() { return _size; }------// 69
    ----void clear() { while(head->next && head->next[0]) erase(head->next[0]->item); }------// e9
31
       \verb| node *find(T target) { FIND_UPDATE(x->next[i]->item, target); } \\ \textit{return } x \& \& x->item == target ? x : \texttt{NULL}; } \\ | A target | A tar
32
    ----node *nth(int k) { FIND_UPDATE(pos[i] + x->lens[i], k+1); return x; }-----// 9e
33
    ----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target); return pos[0]; }------// 0d
34
    ----node* insert(T target) {-------// c4
-----FIND_UPDATE(x->next[i]->item, target);-----// 19
35
36
    -----if(x && x->item == target) return x; // SET-----// 44
37
    -------int lvl = bernoulli(MAX_LEVEL);------// 1d
38
    ---------if(lvl > current_level) current_level = lvl;--------// da
------x = new node(lvl, target);------// a2
40
    ------for(int i = 0; i <= lvl; i++) {-------// ac
41
    -----x->next[i] = update[i]->next[i];-----// 63
42
    -----x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];-----// c8
43
    -----update[i]->next[i] = x;-----// 06
44
    -----update[i]->lens[i] = pos[0] + 1 - pos[i];-----// f5
45
    46
    ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// b9
47
    -----size++;------// fe
-----return x; }-----// a6
49
    ----void erase(T target) {-------// 7e
50
    ------FIND_UPDATE(x->next[i]->item, target);------// d7
51
    53
    -----if(update[i]->next[i] == x) {------// ac
54
    -----update[i]->next[i] = x->next[i];-----// 62
55
    -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;------// 6b
56
   57
58
    -----delete x; _size--;------// 79
59
    ------while(current_level > 0 &\( \&\) head->next[current_level] == NULL) current_level--; \ \} \};---// 72
60
```

3. Graphs

3.1. **Breadth-First Search.** An implementation of a breadth-first search that counts the number of edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted graph (which is represented with an adjacency list). Note that it assumes that the two vertices are connected. It runs in O(|V| + |E|) time.

```
int bfs(int start, int end, vvi& adj_list) {------// d7
1
 ----queue<ii>> 0;------// 75
 ----Q.push(ii(start, 0));------// 49
3
 -----// Øb
4
 ----while (true) {------// @a
5
 -----ii cur = Q.front(); Q.pop();-----// e8
6
 ------if (cur.first == end)------// 6f
8
 -----return cur.second;------// 8a
9
 -----// 3c
10
 -----vi& adj = adj_list[cur.first];-----// 3f
 ------for (vi::iterator it = adj.begin(); it != adj.end(); it++)------// bb
12
 -----Q.push(ii(*it, cur.second + 1));-----// b7
 14
 }-----// 7d
15
```

Another implementation that doesn't assume the two vertices are connected. If there is no path from the starting vertex to the ending vertex, a-1 is returned.

```
int bfs(int start, int end, vvi& adj_list) {------// d7
 ----set<<mark>int</mark>> visited;------// b3
----queue<ii>> 0;------// bb
3
 ----Q.push(ii(start, 0));------// 3a
 ----visited.insert(start);------// b2
5
 -----// db
 ----while (!0.emptv()) {------// f7
 -----i cur = Q.front(); Q.pop();-----// 03
8
 -----// 9c
9
 10
 -----return cur.second;------// b9
11
 12
 -----vi& adj = adj_list[cur.first];-----// f9
13
 ------for (vi::iterator it = adj.begin(); it != adj.end(); it++)------// 44
14
 15
16
17
 18
 19
 20
 21
 }------// 03
 3.2. Single-Source Shortest Paths.
 3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|) time.
 struct cmp {------// a5
 ----bool operator()(int a, int b) { return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// 7e
3
 };-----// f4
 pair<int*, int*, dijkstra(int n, int s, vii *adj) {------// 97
6
 ----dist = new int[n];------// b2
 8
 ----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;------// 09
9
 ----set<<mark>int</mark>, cmp> pq;-----// 0d
10
 ----dist[s] = 0, pq.insert(s);-----// 35
11
 ----while (!pq.empty()) {------// f2
12
 ------int cur = *pq.begin(); pq.erase(pq.begin());-----// 99
13
 ------for (int i = 0; i < size(adj[cur]); i++) {-------// 6e
14
 -----int nxt = adj[cur][i].first, ndist = dist[cur] + adj[cur][i].second;------// 92
15
 -----if (ndist < dist[nxt]) pq.erase(nxt), dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);// 73
 17
 18
 -----// 9c
19
 ----return pair<int*, int*>(dist, dad);------// d2
20
 }------// 08
21
 3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths problem in O(|V||E|)
 time. It is slower than Dijkstra's algorithm, but it works on graphs with negative edges and has the ability to detect negative
 cycles, neither of which Dijkstra's algorithm can do.
 int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {-------// cf
 ----has_negative_cycle = false;------// 47
2
 ----int* dist = new int[n];------// 7f
 ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
4
 ----for (int i = 0; i < n - 1; i++)------// a1
5
 -----for (int j = 0; j < n; j++)-----// c4
6
 ------if (dist[j] != INF)-----// 4e
 -----for (int k = 0; k < size(adj[j]); k++)-----// 3f
 -----dist[adj[j][k].first] = min(dist[adj[j][k].first], dist[j] + adj[j][k].second);-// 6e
9
 ----for (int j = 0; j < n; j++)-----// f5
10
 ------for (int k = 0; k < size(adj[j]); k++)-----// b2
11
 -----if (dist[j] + adj[j][k].second < dist[adj[j][k].first])-----// ce
 -----has_negative_cycle = true;-----// 72
13
 ----return dist;-----// 33
14
 }-----// ed
15
```

```
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths problem in O(|V|^3)
 ----for (int k = 0; k < n; k++)------// 49
 -----for (int i = 0; i < n; i++)------// 21
 -----for (int j = 0; j < n; j++)-----// 77
 -----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
  -----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                      -----// 86
 3.4. Strongly Connected Components.
 3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed graph in O(|V| + |E|)
 #include "../data-structures/union_find.cpp"------// 5e
  -----// 11
 vi order;------// 9b
  6
 ----int v; visited[u] = true;------// e3
  ----for (int i = 0; i < size(adj[u]); i++) if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-----// b0
8
 ----order.push_back(u);-----// 16
9
 }-----// 9b
10
 -----// e5
11
 pair<union_find, vi> scc(const vvi &adj) {------// e3
12
  ----int n = size(adj), u, v;------// c3
13
 ----order.clear();-----// fa
14
 ----union_find uf(n);------// c2
15
 ----vi dag;------// 57
 ----vvi rev(n);-------// f7
17
 ----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++) rev[adj[i][j]].push_back(i);-// 48
 ----visited.resize(n), fill(all(visited), false);-----// af
19
 ----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);------// 59
20
 ----fill(all(visited), false);------// d4
21
 ----stack<<u>int</u>> S;-----// 95
22
 ----for (int i = n-1; i >= 0; i--) {------// 96
23
 -----if (visited[order[i]]) continue;-----// ef
24
 -----S.push(order[i]), dag.push_back(order[i]);------// 87
 ------while (!S.empty()) {-------// 3f
26
 -----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);------// 53
 -----for (int i = 0; i < size(adj[u]); i++) if (!visited[v = adj[u][i]]) S.push(v);------// 02
28
 ----}------// 54
30
  ----return pair<union_find, vi>(uf, dag);------// 0a
31
 }------// e8
32
 3.5. Minimum Spanning Tree.
 3.5.1. Kruskal's algorithm.
 #include "../data-structures/union_find.cpp"-----// 5e
  -----// 11
2
 // n is the number of vertices-----// 18
 // edges is a list of edges of the form (weight, (a, b))-----// c6
 // the edges in the minimum spanning tree are returned on the same form-----// 4d
 vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {-------// a7
  ----union_find uf(n);------// 04
 ----sort(all(edges));------// f5
 ----vector<pair<int, ii> > res;------// 9b
 ----for (int i = 0; i < size(edges); i++)-----// fd
10
  -----if (uf.find(edges[i].second.first) != uf.find(edges[i].second.second)) {-------// 01
11
 ----res.push_back(edges[i]);-----// 39
12
 -----uf.unite(edges[i].second.first, edges[i].second.second);------// 83
13
 ----return res;------// 36
15
```

3.6. Topological Sort.

```
3.6.1. Modified Depth-First Search.
 void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res, bool& has_cycle) {------// ba
 ----color[cur] = 1;------// ed
 ----for (int i = 0; i < size(adj[cur]); i++) {-------// af
 ------int nxt = adj[cur][i];-----// 97
4
 -----if (color[nxt] == 0)-----// f7
 -----tsort_dfs(nxt, color, adj, res, has_cycle);-----// 33
6
 -----else if (color[nxt] == 1)-----// 0c
 ------has_cycle = true;------// 0e
8
 -----if (has_cycle) return;-----// b7
9
 ----}-----// 2f
10
 ----color[cur] = 2;------// b6
11
 ----res.push(cur);------// 7c
12
 }------// d4
13
 ,
------// 8d
14
 vi tsort(int n, vvi adj, bool& has_cycle) {------// b0
15
 ----has_cycle = false;-----// c1
 ----stack<int> S;------// be
17
 ---vi res;-----// fc
18
 ----char* color = new char[n];------// 0b
19
 ----memset(color, 0, n);------// 15
20
 ----for (int i = 0; i < n; i++) {-------// eb
21
 ------if (!color[i]) {-------// fe
22
 -----tsort_dfs(i, color, adj, S, has_cycle);-----// 7a
23
 -----if (has_cycle) return res;------// c2
24
 25
 ----}-----// bf
26
 ----while (!S.empty()) res.push_back(S.top()), S.pop();------// 5b
27
 ----return res;-------// 82
28
 }-----// 42
29
 3.7. Bipartite Matching. The alternating paths algorithm solves bipartite matching in O(mn^2) time, where m, n are
 the number of vertices on the left and right side of the bipartite graph, respectively.
 vi* adj;-----// cc
1
 2
 3
 int alternating_path(int left) {-------// da
 ----if (done[left]) return 0;------// 08
5
 ----done[left] = true;------// f2
6
 ----for (int i = 0; i < size(adj[left]); i++) {------// 34
 ------int right = adj[left][i];-----// b6
 -----if (owner[right] == -1 || alternating_path(owner[right])) {------// d2
9
 ------owner[right] = left; return 1;------// 26
-----} }------// 7a
10
11
 ----return 0; }-------// 83
12
 3.8. Maximum Flow.
 3.8.1. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in O(|V||E|^2). It computes
 the maximum flow of a flow network.
 struct mf_edge {-----// b3
1
 ----int u, v, w; mf_edge* rev;------// ab
2
 -----u = _u; v = _v; w = _w; rev = _rev; } };------// b1
4
 pair<int, vector<vector<mf_edge*> > > max_flow(int n, int s, int t, vii* adj) {------// 57
5
 ----int flow = 0, cur, cap;------// ac
6
 ----vector<vector<mf_edge*> > g(n);-----// 07
 ----vector<mf_edge*> back(n);------// 14
 ----mf_edge *ce, *z;------// 09
9
 ----for (int i = 0; i < n; i++) {------// be
10
 ------for (int j = 0; j < size(adj[i]); j++) {------// 21
11
 -----g[i].push_back(ce = new mf_edge(i, adj[i][j].first, adj[i][j].second));-----// cc
12
 -----g[ce->v].push_back(ce->rev = new mf_edge(adj[i][j].first, i, 0, ce)); } }------// 03
13
 ----while (true) {------// 61
 ------back.assign(n, NULL);-----// 59
15
 -----queue<int> Q; Q.push(s);-----// 0a
 -------while (!Q.empty() && (cur = Q.front()) != t) {------// d5
17
 ------Q.pop();-----// da
18
 -----for (int i = 0; i < size(g[cur]); i++) {------// 9d
19
```

20

21

-----if (nxt->v != s && nxt->w > 0 && !back[nxt->v])------// 3a

```
------Q.push((back[nxt->v] = nxt)->v); } }------// 68
22
 -----if (!back[t] || back[t]->w == 0) break;-----// dc
23
  ------for (int i = 0; i < size(g[t]); i++) {------// d2
24
 ------if (!(z = g[t][i]->rev) || (!back[z->u] && z->u != s)) continue;------// 06
25
 -----for (cap = z->w, ce = back[z->u]; ce && cap > 0; ce = back[ce->u])------// 10
26
 ------cap = min(cap, ce->w);------// c8
-----if (cap == 0) continue;----// f8
27
28
 -----assert(cap < INF);-----// 7e
 -----z->w -= cap, z->rev->w += cap;-------// 37
30
 -----for (ce = back[z->u]; ce; ce = back[ce->u])------// 30
 ------ce->w -= cap, ce->rev->w += cap;------// a8
32
  -----flow += cap; } }------// 8e
33
 ----return make_pair(flow, g); }-----// 73
34
 3.9. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modified to find shortest path
 to augment each time (instead of just any path). It computes the maximum flow of a flow network, and when there are
 multiple maximum flows, finds the maximum flow with minimum cost.
 ----int u, v, w, c;-------// a5
 ----mcmf_edge(int _u, int _v, int _w, int _c, mcmf_edge* _rev = NULL) {-------// f7
4
 -----u = _u; v = _v; w = _w; c = _c; rev = _rev;------// b2
  6
 };-----// e4
  -----// 31
 ii min_cost_max_flow(int n, int s, int t, vector<pair<int, ii> >* adj) {------// 4d
9
 ----vector<mcmf_edge*>* g = new vector<mcmf_edge*>[n];------// 0c
----for (int i = 0; i < n; i++) {-------// a7
10
11
 -----for (int j = 0; j < size(adj[i]); j++) {------// a1
12
      mcmf_edge *cur = new mcmf_edge(i, adj[i][j].first, adj[i][j].second.first, adj[i][j].second.second),
13
 *rev = new mcmf_edge(adj[i][j].first, i, 0, -adj[i][j].second.second, cur);---// d6
14
 -----cur->rev = rev;------// 84
15
 -----g[i].push_back(cur);------// ad
 -----g[adj[i][j].first].push_back(rev);-----// 24
17
 18
 ----}------// a8
19
 ----int flow = 0, cost = 0;------// ae
20
 ----mcmf_edge** back = new mcmf_edge*[n];------// 27
21
  ----int* dist = new int[n];------// 60
22
 ----while (true) {------// bd
23
 ------for (int i = 0; i < n; i++) back[i] = NULL, dist[i] = INF;-------// 9a
24
 25
26
 ------for (int j = 0; j < n; j++)------// f2
 -----if (dist[j] != INF)-----// 07
28
 -----for (int k = 0; k < size(g[j]); k++)-----// 1b
29
 30
 -----dist[g[j][k]->v] = dist[j] + g[j][k]->c;-----// d2
31
 -----back[g[j][k]->v] = g[j][k];-----// b9
32
 33
 -----mcmf_edge* cure = back[t];------// 4a
-----if (cure == NULL) break;-----// 79
35
 ------<mark>int</mark> cap = INF;------// 70
 37
 -----if (cure->u == s) break;-----// 41
39
  -----cure = back[cure->u];-----// d6
40
 41
 -----assert(cap > 0 && cap < INF);------// 92
42
 -----cure = back[t];-----// 35
43
  ------while (true) {------// 58
44
 -----cost += cap * cure->c;-----// 6d
45
 -----cure->w -= cap;------// le
46
 ------cure->rev->w += cap; ------// 4e
-----if (cure->u == s) break; -----// 4e
47
48
  -----cure = back[cure->u];-----// d1
49
 50
  ------flow += cap;-------// 40
51
```

```
----}------// db
52
 ----// instead of deleting g, we could also------// f2
53
 ----// use it to get info about the actual flow------// 80
54
 ----for (int i = 0; i < n; i++)------// 99
55
 ------for (int j = 0; j < size(g[i]); j++)------// a5
56
 -----delete g[i][j];-----// 7f
57
 ----delete[] g;------// a0
 ----delete[] back;------// 10
59
 ----delete[] dist;------// 76
 ----return ii(flow, cost);-----// 47
61
 }------// e4
```

3.10. All Pairs Maximum Flow.

3.10.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is constructed using Gusfield's algorithm in $O(|V|^3)$ plus |V|-1 times the time it takes to calculate the maximum flow. If Edmonds Karp's algorithm is used to calculate the max flow the running time is $O(|V|^3 + |V|^2|E|^2)$.

```
#include "edmonds_karps.cpp"-----// 7d
1
 2
 ---vector<int> flows(size(adj)); ----int *back = new int[size(adj)]; -----// 04
3
 ----memset(back, -1, sizeof(int) * size(adj));------// 2e
5
 ----queue<ii>> Q; Q.push(ii(s, INF));-----// d1
 ----while (!Q.empty()) {------// ea
 -----ii bla = Q.front(); Q.pop();-----// 7b
 ------int cur = bla.first;-----// ce
9
 ------for (int i = 0; i < size(adj[cur]); i++) {-------// 85
10
 -----ii nxt = adj[cur][i];-----// 91
11
 -----if (back[nxt.first] == -1 && nxt.first != s) {------// 84
12
 -----flows[nxt.first] = min(bla.second, nxt.second);------// 42
13
 -----back[nxt.first] = cur;-----// 93
14
 -----Q.push(ii(nxt.first, min(bla.second, nxt.second)));-----// fa
15
 16
 ----delete[] back;-----// 52
 ----return flows;-----// 71
18
 }-----// df
19
 pair<set<int>, int> max_flow_min_cut(int n, int s, int t, vii *adj) {------// 14
20
 ----pair<<u>int</u>, vector<vector<mf_edge*> >> mf = max_flow(n, s, t, adj);------// b6
21
 22
23
 ----queue<int> Q; Q.push(s);-----// 4a
24
 25
 ----while (!Q.empty()) {------// 73
26
 ------int cur = Q.front(); Q.pop();------// e8
27
 -----cut.insert(cur);------// 67
 ------for (int i = 0; i < size(mf.second[cur]); i++) {-------// 90
29
 -----mf_edge *nxt = mf.second[cur][i];-----// 06
30
 -----if (back[nxt->v] == -1 && nxt->w > 0) {------// 1a
31
 -----back[nxt->v] = cur;-----// dd
32
 -----Q.push(nxt->v);------// 57
33
 34
 ----delete[] back;------// 58
 ----for(int i = 0; i < size(mf.second); i++)------// f6
36
 ------for(int j = 0; j < size(mf.second[i]); j++)------// 72
37
 38
 }------// d7
40
 vector<vii> gomory_hu_tree(int n, vii *adj) {------// 4f
41
 ----assert(n >= 0);-----// 81
42
 43
 44
 ----if(n < 2) return T;------// df
45
 ----map<ii, set<int> > cuts;------// d1
46
 ----pair<set<int>, int> m = max_flow_min_cut(n, 0, 1, adj);------// 6d
47
 ---cuts[ii(0, 1)] = m.first;-----// a1
---T[0].push_back(ii(1, m.second));-----// ca
49
 ----T[1].push_back(ii(0, m.second));------// 9d
 ----int *back = new int[n];------// 35
51
 ----int t1 = 0, t2 = -1;-------// 8f
```

```
----vector<pair<int, ii>> edges;------// d2
53
 ----for(int k = 2; k < n; ++k) {------// 78
54
 ------while(true) {------// b0
55
 56
 -----memset(back, -1, sizeof(int) * n);------// 4a
57
 -----queue<int> Q; Q.push(t1);-----// ef
58
 ------while (!Q.empty()) {-----// b9
 -----// 7e
60
 61
 -----ii nxt = T[cur][i];------// 4d
62
 -----if(T[cur][i].second < m_arc) {------// a0
 -----m_arc = T[cur][i].second;-----// 3e
64
 -----alpha = cur;------// 26
65
 -----beta = T[cur][i].first;-----// a7
66
 67
 -----if (back[nxt.first] == -1) {------// fb
 69
70
 71
 -----if(alpha > beta) swap(alpha, beta);-----// b2
 -----if(cuts[ii(alpha, beta)].find(k) != cuts[ii(alpha, beta)].end())------// 50
73
 -----t1 = alpha, t2 = beta;-----// 3d
 -----else------// 68
75
 -----t1 = beta, t2 = alpha;-----// 64
76
 -----T[t1].erase(find(T[t1].begin(), T[t1].end(), ii(t2, m_arc)));------// 2f
77
 -----edges.push_back(make_pair(t1, ii(t2, m_arc)));------// e8
78
 79
 -----for(int i = 0; i < edges.size(); ++i)-----// ef
80
 -----T[edges[i].first].push_back(ii(edges[i].second.first, edges[i].second.second));-// 92
81
 ------edges.clear();-----// f3
82
 -----pair<set<<u>int</u>>, <u>int</u>> m = max_flow_min_cut(n, t1, k, adj);------// 03
83
 -----cuts[ii(t1, k)] = m.first;------// 39
84
 -----T[t1].push_back(ii(k, m.second));-----// c0
85
 -----T[k].push_back(ii(t1, m.second));-----// 62
86
 -----break;-----// b3
87
 88
 ----delete[] back;-----// ec
89
 ----return T;------// f9
 }-----// d2
91
                     4. Strings
 4.1. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is also a prefix of S. The Z
 algorithm computes these Z values in O(n) time, where n = |S|. Z values can, for example, be used to find all occurrences
 of a pattern P in a string T in linear time. This is accomplished by computing Z values of S = TP, and looking for all i
 such that Z_i \geq |T|.
 int* z_values(string s) {------// ce
 ----int n = size(s);------// 45
 ----int* z = new int[n];------// 04
3
 ----int l = 0, r = 0;------// e5
 ----for (int i = 1; i < n; i++) {--------// f4
-----z[i] = 0;------// ce
 -----if (i > r) {------// 2c
 -----l = r = i;------// 1a
9
 ------while (r < n && s[r - l] == s[r]) r++;------// 0a
10
 -----z[i] = r - l; r--;-----// a2
11
 ------} else if (z[i - l] < r - i + 1) z[i] = z[i - l];------// a6
12
 -----else {------// 9d
13
 -----l = i;------// 4f
14
 ------while (r < n && s[r - l] == s[r]) r++;-------// d8
15
 -----z[i] = r - l; r--; } }------// d0
16
 ----return z;-----// 2d
 }------// f3
18
                    5. Mathematics
 5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common terms.
 template <class T>------// 82
```

```
private:-----// 8e
  public:----// 0f
5
  6
  ----fraction(T n_, T d_) {------// 03
7
  ------assert(d_ != 0);------// 3d
8
  -----n = n_, d = d_;------// 06
  10
  -----T g = gcd(abs(n), abs(d));-----// fc
11
  -----n /= g, d /= g; }------// a1
12
  ----fraction(T n_) : n(n_), d(1) { }------// 84
  ----fraction(const fraction<T>\& other) : n(other.n), d(other.d) { }------// 01
14
15
   fraction < T> operator + (const fraction < T> \& other) const \{ return fraction < T> (n * other.d + other.n * d, d * other.d); \}
   fraction < T> operator - (const fraction < T> \& other) const \{ return fraction < T> (n * other.d - other.n * d, d * other.d); \}
16
   fraction<T> operator *(const fraction<T>\& other) const { return fraction<T>(n * other.n, d * other.d); }
17
   fraction < T> operator / (const fraction < T> \& other) const { return fraction < T> (n * other.d, d * other.n); } 
18
  ----bool operator <(const fraction<T>& other) const { return n * other.d < other.n * d; }------// 3a
19
 ----bool operator <=(const fraction<T>\& other) const { return !(other < *this); }------// 9a
20
  ----bool operator >(const fraction<T>& other) const { return other < *this; }------// 8b
21
 ----bool operator >=(const fraction<T>\& other) const { return !(*this < other); }------// 34
  ----bool operator ==(const fraction<T>\& other) const { return n == other.n \&\& d == other.d; }-----// 1c
23
  ----bool operator !=(const fraction<T>& other) const { return !(*this == other); }------// 84
 };-----// 4c
25
  5.2. Big Integer. A big integer class.
  class intx-----// 2c
1
  {------// 1f
2
 public:----// 51
3
  ----intx() { normalize(1); } -------// 5c
  ----intx(string n) { init(n); }------// 24
  ----intx(int n) { stringstream ss; ss << n; init(ss.str()); }-----// ae
6
  ----intx(const intx& other) : sign(other.sign), data(other.data) { }------// 01
7
  -----// 6e
8
  ----friend bool operator <(const intx& a, const intx& b);------// d9
  ----friend intx operator +(const intx& a, const intx& b);------// 56
10
  ----friend intx operator -(const intx& a, const intx& b);-----// 43
  ----friend intx operator *(const intx& a, const intx& b);------// d0
12
  ----friend intx operator /(const intx& a, const intx& b);-----// a2
13
  ----friend intx operator %(const intx& a, const intx& b);------// 14
14
  ----friend ostream& operator <<(ostream& outs, const intx& n);-----// cc
15
 -----// 9e
16
 protected:-----// ef
17
  ----int sign;------// 06
18
  ----vector<<mark>unsigned int</mark>> data;------// d5
19
  ----static const unsigned int radix = 1000000000U;------// 90
  ----int size() const { return data.size(); }------// @a
21
  23
  ------if (n.empty()) n = "0";------// ad
24
 ------if (n[0] == '-') res.sign = -1, n = n.substr(1);------// 12
25
 26
27
28
 -----int idx = i - j;------// 65
 30
 -----digit = digit * 10 + (n[idx] - '0');------// 63
 32
  -----res.data.push_back(digit);-----// cc
33
 -----}-------// e1
34
  -----data = res.data;-----// 5a
35
  -----normalize(res.sign);-----// ce
36
  37
  ----void normalize(int nsign) {------// 60
38
  -----if (data.empty()) data.push_back(0);------// c1
39
  ------for (int i = data.size() - 1; i > 0 && data[i] == 0; i--) data.erase(data.begin() + i);-----// 5b
  -----sign = data.size() == 1 && data[0] == 0 ? 1 : nsign;-----// 2a
41
  ----}------// 0b
  43
```

```
------for (int i = 0; i < size(); i++) d[i + n] = data[i];------// 38
45
  -----intx res; res.data = d; res.normalize(sign);-----// 40
46
  -----return res;------// 7c
47
  48
 }:-----// f6
49
 ostream& operator <<(ostream& outs, const intx& n) {-----// 86
50
  51
  ----bool first = true;------// 99
52
  ----for (int i = n.size() - 1; i >= 0; i--) {------// 38
  -----if (first) outs << n.data[i], first = false;-----// d7
54
  -----else {------// b4
  -----unsigned int cur = n.data[i];-----// 07
56
  -----stringstream ss; ss << cur;-----// 27
57
 -----string s = ss.str();------// ba
58
  -----int len = s.size();-----// 12
59
  ------while (len < 9) outs << '0', len++;------// b4
60
  -----outs << s;-----// cd
61
  62
  ----}------// 1c
63
  65
  bool operator <(const intx& a, const intx& b) {------// 00
66
  ----if (a.sign != b.sign) return a.sign < b.sign;------// b6
67
  ----if (a.size() != b.size()) return a.sign == 1 ? a.size() < b.size() : a.size() > b.size();-----// 7a
68
  ----for (int i = a.size() - 1; i >= 0; i--) if (a.data[i] != b.data[i])------// d5
69
  -----return a.sign == 1 ? a.data[i] < b.data[i] : a.data[i] > b.data[i];-----// da
70
  ----return false;-----// 0e
71
 }-----// 3f
72
  intx operator +(const intx& a, const intx& b) {------// 8f
73
  ----if (a.sign != b.sign) return intx() - (intx() - a - b);------// 15
74
  ----intx c; c.data.clear();------// ca
  ----unsigned long long carry = 0;-----// fb
76
  77
  -----carry += (i < a.size() ? a.data[i] : OULL) + (i < b.size() ? b.data[i] : OULL);-----// 3c
78
  -----c.data.push_back(carry % intx::radix);------// e1
79
  -----carry /= intx::radix;-----// 3b
80
  81
  ----c.normalize(a.sign);------// 73
82
  83
  }------// be
84
 intx operator -(const intx& a, const intx& b) {------// b9
85
  ----if (a < b) return intx() - (b - a);------// 71
87
  ----intx c; c.data.clear();------// 3e
88
  ----long long borrow = 0;------// 47
89
  ----for (int i = 0; i < a.size(); i++) {------// a3
  ------borrow = a.data[i] - borrow - (i < b.size() ? b.data[i] : OULL);------// e6
91
  -----c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);------// 4c
92
  ------borrow = borrow < 0 ? 1 : 0;------// a3
93
  94
  ----c.normalize(a.sign);------// ea
  96
  }------// 04
  .
-----// 0a
98
  ----int n = max(a.size(), b.size());------// c0
100
  ----if (n == 1) {-------// f5
101
  -----unsigned long long res = static_cast<unsigned long long>(a.data[0]) * b.data[0];------// e2
102
  -----stringstream ss; ss << res;------// e1
103
  -----intx result(ss.str());------// 8d
  -----result.normalize(a.sign * b.sign);------// ee
105
  -----return result;------// 69
106
  107
  ----if (n & 1) n++;------// 12
108
  100
  110
  ----buff1.reserve(n2); buff2.reserve(n2);------// 0c
111
  ----for (int at = n2 - 1; at >= 0; at--) {------// 1a
112
  ------int idx = n - at - 1;------// @a
```

```
------buff1.push_back(idx < a.size() ? a.data[idx] : 0);-------// cc
114
  ------buff2.push_back(idx < b.size() ? b.data[idx] : θ);-----// da
115
  116
  ----intx i, k;------// 9c
117
  ----i.data = buff1; k.data = buff2;-----// 3c
118
  ----buff1.clear(); buff2.clear();------// 71
119
  ----for (int at = n - 1; at >= n2; at--) {------// 75
120
  ------int idx = n - at - 1;------// ef
121
  ------buffl.push_back(idx < a.size() ? a.data[idx] : 0);------// af
  ------buff2.push_back(idx < b.size() ? b.data[idx] : 0);-----// e8
123
  ---}-----// 73
  ----intx j, l;-------// 5a
125
  ----j.data = buff1; l.data = buff2;-----// 4c
126
  ----intx ik = i * k, jl = j * l;------// 02
127
  ----intx res = ik.mult_radix(n) + ((i + j) * (k + l) - (ik + jl)).mult_radix(n2) + jl;------// 6a
128
  ---res.normalize(a.sign * b.sign);-----// 44
---return res;-----// 20
129
130
  }-----// 3f
131
  intx operator /(const intx& n, const intx& d)-----// ab
132
  {------// cb
  134
  ----intx q, r; q.data.assign(n.size(), 0);------// 5e
  ----for (int i = n.size() - 1; i >= 0; i--) {------// 78
136
  -----r.data.insert(r.data.begin(), 0);-----// 35
137
  -----intx y; y.data[0] = n.data[i];------// f1
-----r = r + y;------// b4
138
139
  140
141
  ...-q.normalize(n.sign * d.sign);------// f7
----return q;------// f9
142
143
  }------
144
  intx operator %(const intx& n, const intx& d) {------// dd
145
  ----intx r;------// d6
146
  ----for (int i = n.size() - 1; i >= 0; i--) {-------// 7b
147
  -----r.data.insert(r.data.begin(), 0);-----// 3d
148
  149
150
151
152
  ----r.normalize(n.sign * d.sign);------// 27
153
  154
  }------// 4c
155
  5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose k items out of a
  int nck(int n, int k) {-----// f6
  ----if (n - k < k) k = n - k;-----// 18
  ----int res = 1;------// cb
3
  ----for (int i = 1; i <= k; i++) res = res * (n - (k - i)) / i;------// bd
  ----return res;------// e4
5
  }------// 03
  5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two integers a, b.
  int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }------// d9
   The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b and also finds two
  integers x, y such that a \times x + b \times y = d.
  int egcd(int a, int b, int& x, int& y) {------// 85
1
  2
3
  ------int d = egcd(b, a % b, x, y);------// 34
  -----x -= a / b * y;------// 4a
5
  -----swap(x, y);------// 26
-----return d;-----// db
  5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is prime.
  bool is_prime(int n) {-------// 6c
  ----if (n < 2) return false;------// c9
```

```
----if (n % 2 == 0 || n % 3 == 0) return false;------// 0f
  5
  ----for (int i = 5; i <= s; i += 6) if (n % i == 0 \mid \mid n \% (i + 2) == 0) return false;-----// 6d
  ----return true; }------// 0b
  5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
  1
  3
  ----memset(prime, 1, mx + 1);-----// 28
  6
  8
  -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
  ------for (int j = sq; j <= mx; j += v) prime[j] = false; }-----// 2e
10
  ----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3);------// 29
  ----delete[] prime; // can be used for O(1) lookup-----// 36
12
                                -----// 72
13
  5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
  #include "eqcd.cpp"------// 55
1
  -----// e8
  5
  ----return x < 0 ? x + m : x;------// 3c
6
                -----// 69
  5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
  T mod_pow(T b, T e, T m) {------// aa
2
  -----if (e & T(1)) res = mod(res * b, m);------// 41
5
  -----b = mod(b * b, m), e >>= T(1); }------// b3
  ----return res:-------// eb
  }-----// c5
  5.9. \  \, \textbf{Chinese Remainder Theorem.} \  \, \textbf{An implementation of the Chinese Remainder Theorem.}
  #include "egcd.cpp"------// 55
  int crt(const vi& as, const vi& ns) {-----// c3
  ----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
  ----for (int i = 0; i < cnt; i++) N *= ns[i]; ------// 88
  ----for (int i = 0; i < cnt; i++) egcd(ns[i], l = N/ns[i], r, s), x += as[i] * s * l;------// d6
  ----return mod(x, N); }------// 31
  5.10. Formulas.
     • Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen
       once: P_k^n = \frac{n!}{(n-k)!}
     • Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen
       multiple times: n^k
      Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type 2, ..., n_k objects
      of type k: \binom{n}{n_1, n_2, \dots, n_k} = \frac{n!}{n_1! \times n_2! \times \dots \times n_k!}
Number of ways to choose k objects from a total of n objects where order does not matter and each item can only
       be chosen once:
       \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^k \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0
     • Number of ways to choose k objects from a total of n objects where order does not matter and each item can be
      chosen multiple times: f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}
Number of integer solutions to x_1 + x_2 + \cdots + x_n = k where x_i \geq 0: f_k^n
      Number of subsets of a set with n elements: 2^n
      |A \cup B| = |A| + |B| - |A \cap B|
      |A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|
      Number of ways to walk from the lower-left corner to the upper-right corner of an n \times m grid by walking only up
       and to the right: \binom{n+m}{m}
```

Number of strings with n sets of brackets such that the brackets are balanced:

 $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$

- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an $n \times n$ lattice which do not rise above the main diagonal: C_n
- Number of permutations of n objects with exactly k ascending sequences or runs: $\binom{n}{k} = \binom{n}{n-k-1} = k \binom{n-1}{k} + (n-k+1) \binom{n-1}{k-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{n-1} = 1$
- Number of permutations of n objects with exactly k cycles: $\begin{bmatrix} n \\ k \end{bmatrix} = \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} + (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix}$
- Number of ways to partition n objects into k sets: $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$

6. Geometry

6.1. Primitives. Geometry primitives.

```
#define P(p) const point &p-----// b8
  #define L(p0, p1) P(p0), P(p1)-----// 30
  -----// 20
  typedef complex<double> point;-----// f8
  typedef vector<point> polygon;-----// 16
  double dot(P(a), P(b)) { return real(conj(a) * b); }-----// 43
  double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// f2
  point\ rotate(P(p),\ P(about),\ \textbf{double}\ radians)\ \{\ \textbf{return}\ (p\ -\ about)\ *\ exp(point(\theta,\ radians))\ +\ about;\ \}
9
  point reflect(P(p), L(about1, about2)) {-------// a4 ----point z = p - about1, w = about2 - about1;------// a1
10
11
  ----return conj(z / w) * w + about1; }-----// e1
12
  bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// cb</pre>
13
  double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// 33
14
  bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }-----// a7
15
  16
  double angle(P(a), P(b), P(c)) { return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 7d
17
  double signed_angle(P(a), P(b), P(c)) { return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }
18
  19
  ----// NOTE: check for parallel/collinear lines before calling this function------// 44
20
  ----point r = b - a, s = q - p;-----// be
21
  ----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;-----// af
22
  ----if (segment && (t < 0-EPS \mid\mid t > 1+EPS \mid\mid u < 0-EPS \mid\mid u > 1+EPS)) return false;-----// 21
23
  24
25
  }------// 80
26
  point closest_point(L(a, b), P(c), bool segment = false) {------// e7
27
  29
  ------if (dot(a - b, c - a) > 0) return a;------// 3e
30
  ---}-----// 27
31
  ----double t = dot(c - a, b - a) / norm(b - a);------// 0b
32
  ----return a + t * (b - a);-----// 05
33
  34
  double polygon_area_signed(polygon p) {------// 88
35
  ----double area = 0; int cnt = size(p);-----// bd
36
  ----for (int i = 1; i + 1 < cnt; i++) area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 96
37
  ----return area / 2;------// b2
38
  double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }-----// fe
40
  // pair<polygon, polygon> cut_polygon(const polygon &poly, point a, point b) {------// ec
41
  //--- polygon left, right;-----// b4
//--- point it(-100, -100);-----// 93
42
43
  //--- for (int i = 0, cnt = poly.size(); i < cnt; i++) {------// fe
44
  //----- int j = i == cnt-1 ? 0 : i + 1;------// 3c
45
  //----- point p = poly[i], q = poly[j];------// 27
//----- if (ccw(a, b, p) <= 0) left.push_back(p);-----// 2c
46
47
  //----- if (ccw(a, b, p) >= 0) right.push_back(p);-----// c0
  //-----// myintersect = intersect where (a,b) is a line, (p,q) is a line segment------// 94
49
  //----- if (myintersect(a, b, p, q, it)) left.push_back(it), right.push_back(it);-----// 14
  //---- }------// 7a
//---- return pair<polygon, polygon>(left, right);------// ea
51
  // }------// 13
53
  6.2. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
  #include "primitives.cpp"------// e0
  point ch_main;------// 38
  bool ch_cmp(P(a), P(b)) {------// 7b
```

```
----if (collinear(ch_main, a, b)) return abs(a - ch_main) < abs(b - ch_main);-------------// 35
 ----return atan2(imag(a) - imag(ch_main), real(a) - real(ch_main)) < atan2(imag(b) - imag(ch_main), -// b4
5
  -----real(b) - real(ch_main)); }------// db
6
 polygon convex_hull(polygon pts, bool add_collinear = false) {------// 53
7
 ----int cnt = size(pts), main = 0, i = 1;------// dd
8
  ----if (cnt <= 3) return pts;------// 88
9
 ----for (int i = 1; i < cnt; i++)-----// f8
10
  -----if (imag(pts[i]) < imag(pts[main]) || abs(imag(pts[i]) - imag(pts[main]) < EPS &&------// 58
11
 -----imag(pts[i]) > imag(pts[main])))-----// a5
 -----main = i;-------// c5
13
 ----swap(pts[0], pts[main]);------// ca
 ----ch_main = pts[0];-----// c1
15
  ----sort(++pts.begin(), pts.end(), ch_cmp);------// 7c
16
 ----point prev, now;------// 1a
17
 ----stack<point> S; S.push(pts[cnt - 1]); S.push(pts[0]);-----// 0b
18
 19
20
 -----if (S.empty()) {------// ac
21
 -----S.push(now);-----// 33
22
 ------S.push(pts[i++]);------// 22
 24
 -----prev = S.top();-----// 8e
 -----S.push(now);------// 8c
26
        \text{if (ccw(prev, now, pts[i])} > 0 \text{ } | \text{ (add\_collinear } \&\& \text{ abs(ccw(prev, now, pts[i]))} \text{ } < \text{EPS)) } \text{ S.push(pts[i++]); } 
27
 -----else S.pop();-----// b3
28
 29
 ----vector<point> res;------// 0f
30
 ----while (!S.empty()) res.push_back(S.top()), S.pop();-----// f6
31
 ----return res;------// 31
 }-----// ce
33
```

- 6.3. Formulas. Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.
 - $a \cdot b = |a||b|\cos\theta$, where θ is the angle between a and b.

3

5

6

9

10

- $a \times b = |a||b|\sin\theta$, where θ is the signed angle between a and b.
- $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b. Half of that is the area of the triangle formed by a and b.

7. Other Algorithms

7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous function f on the interval [a, b], with a maximum error of ε .

```
-----if (clauses[i].first != clauses[i].second)------// 87
8
 -----adj[-clauses[i].second + n].push_back(clauses[i].first + n);------// 93
9
 ----}-----// d8
10
 ----pair<union_find, vi> res = scc(adj);-----// 9f
11
 ----union_find scc = res.first;------// 42
12
 ----vi dag = res.second;-------// 58
13
 ----vi truth(2*n+1, -1);------// 00
 ----for (int i = 2*n; i >= 0; i--) {------// f4
15
 ------int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);------// 5a
16
 -----if (cur == 0) continue;-------// 26
17
 -----if (p == o) return false;-----// 33
 -----if (truth[p] == -1) truth[p] = 1;------// c3
19
 -----truth[cur + n] = truth[p];-----// b3
20
 -----truth[o] = 1 - truth[p];-----// 80
21
 ------if (truth[p] == 1) all_truthy.push_back(cur);------// 5c
22
 23
 ----return true;------// eb
^{24}
 25
 7.3. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \dots, k-1\}.
 vector<int> nth_permutation(int cnt, int n) {------// 78
 ----vector<int> idx(cnt), per(cnt), fac(cnt);-----// 9e
2
 ----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
 ----for (int i = 1; i <= cnt; i++) fac[i - 1] = n % i, n /= i;------// 04
4
   for (int i = cnt - 1; i >= 0; i--) per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);
5
 ----return per;------// fd
6
 }-----// ed
 7.4. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
 ii find_cycle(int x0, int (*f)(int)) {------// a5
1
 ----int t = f(x0), h = f(t), mu = 0, lam = 1;------// 8d
2
 ----while (t != h) t = f(t), h = f(f(h));------// 79
3
 ----h = x0;------// 04
 ----while (t != h) t = f(t), h = f(h), mu++;-----// 9d
5
 ----h = f(t);-----// 00
 ----while (t != h) h = f(h), lam++;-----// 5e
 ----return ii(mu, lam);------// b4
8
 }------// 42
9
 7.5. Dates. Functions to simplify date calculations.
 int intToDay(int jd) { return jd % 7; }------// 89
1
 int dateToInt(int y, int m, int d) {------// 96
 ----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
 -----367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1
4
 -----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
 -----d - 32075;-----// e0
6
 }-----// fa
 8
 ----int x, n, i, j;------// 00
9
 ---x = id + 68569;-----// 11
10
 ---n = 4 * x / 146097;-----// 2f
11
 ----x -= (146097 * n + 3) / 4;-----// 58
12
 ----i = (4000 * (x + 1)) / 1461001;------// 0d
13
 ----x -= 1461 * i / 4 - 31;-----// 09
14
 ----j = 80 * x / 2447;------// 3d
15
 ----d = x - 2447 * j / 80;-----// eb
 ----x = j / 11,-----// b7
17
 ----m = j + 2 - 12 * x;------// 82
18
 ----y = 100 * (n - 49) + i + x;------// 70
19
        .....// af
20
```

8. Useful Information

8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?

- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo m_1, m_2, \ldots, m_k , where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When n = 0, n = -1, n = 1, $n = 2^{31} 1$ or $n = -2^{31}$? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

```
void readn(register int *n) {------// dc
 ----int sign = 1;-------// 32
 ----register char c;-------// a5
3
 ---*n = 0;
 ----while((c = getc_unlocked(stdin)) != '\n') {-------// f3
5
 -----switch(c) {-----// 0c
      -': sign = -1; break;-----// 28
 -----case '
 -----case ' ': goto hell;------// fd
 -----case '\n': goto hell;------// 79
9
 -----default: *n *= 10; *n += c - '0'; break;-----// c0
10
 11
 ----}------// c3
12
 hell:----// ba
13
 ----*n *= siqn;-------// a0
14
 }------// 67
```

8.3. Worst Time Complexity.

n	Worst AC Algorithm	Comment
≤ 10	$O(n!), O(n^6)$	e.g. Enumerating a permutation
≤ 15	$O(2^n \times n^2)$	e.g. DP TSP
≤ 20	$O(2^n), O(n^5)$	e.g. $DP + bitmask technique$
≤ 50	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
$\le 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
$\le 10^{5}$	$O(n \log_2 n)$	e.g. Merge sort, building a Segment tree
$\leq 10^6$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \le 10^6$ (e.g. to read input)

8.4. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.