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           1. Code Templates
                             ----public static void main(String[] args) throws Exception {--------// 02
                             -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                             ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                             -----// code-----// e6
setxkbmap -option caps:escape
                             -----out.flush():-----// 56
set -o vi
                             xset r rate 150 100
                             }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                        2. Data Structures
syn on | colorscheme slate
                             2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                             struct union_find {------// 42
#include <cmath>------// 7d ----union_find(int n) { parent.resize(cnt = n);------// 92
#include <cstdio>------[i] = i; }------// 6f
#include <cstdlib>------// 11 ----int find(int i) {--------// a6
#include <cstring>-------[i] = i ? i : (parent[i] = find(parent[i])); }------// @ -------|/ a9
#include <map>-----// 02
#include <set>------// d2
#include <sstream>------ min(a, b); }------// 18 // int f(int a, int b) { return min(a, b); }------
#include <stack>------// cf const int ID = 0;------// 82
#include <string>------// a9 int f(int a, int b) { return a + b; }------// 6d
#include <vector>-----// 4f ----int n; vi data, lazy;-----// fa
------mk(arr, 0, n-1, 0); }------
#define foreach(u, o) \------// ea ----int mk(const vi &arr, int i, int r, int i) {------// 02
const int INF = 2147483647;-----// be -----int m = (l + r) / 2;-----// 0f
const double pi = acos(-1);------// 49 ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// f5
typedef long long ll;------// 8f ----int q(int a, int b, int l, int r, int i) {-------// ad
typedef unsigned long long ull;-----// 81 -----propagate(l, r, i);-----// f7
typedef vector<vii>vvii;------// 4b ----void update(int i, int v) { u(i, v, 0, n-1, 0); }------// 65
template <class T> T mod(T a, T b) { return (a % b + b) % b; }--------// 70 ----int u(int i, int v, int l, int r, int j) {------------// b5
template <class T> int size(const T &x) { return x.size(); }------// 68 -----propagate(l, r, j);-------// 3c
                             -----if (r < i || i < l) return data[j];------// 6a
1.3. Java Template. A Java template.
                             -----if (l == i && r == i) return data[j] = v;------// 74
import java.math.*;------return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 68
-----// a3 ----int ru(int a, int b, int v, int i) {-------------// d7
public class Main {-------// 17 ------propagate(l, r, i);-----------------// 82
```

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------if (r < a || b < l) return data[i]:---------// bb ----int rows, cols:----------------------------------// d3
------int m = (l + r) / 2; ------// 9d ----matrix(const matrix  other) : rows(other.rows), cols(other.cols), ------// fe
-----ceturn data[i] = f(ru(a, b, v, l, m, 2*i+1), ru(a, b, v, m+1, r, 2*i+2)); ------cnt(other.cnt), data(other.data) { }------// ed
------if (l > r || lazy[i] == INF) return; --------// 42 ------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }-------// e5
------if (l < r) {-------// 6e ------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
------else lazv[2*i+2] += lazv[i];------// d1 -----matrix<T> res(*this); res += other; return res; }-----// 5d
-----lazv[i] = INF;------res(*this); res -= other; return res; }------// cf
};-------matrix<T> res(*this); res *= other; return res; }------// 37
                               ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                               -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                               -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i...j in O(\log n) time. It only needs O(n) space.
                               -----for (int k = 0; k < cols; k++)-----// fc
struct fenwick_tree {------res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n; vi data;------// d3 -----return res; }-----// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {--------------// dd
----void update(int at, int by) {--------// 76 ------matrix<T> res(cols, rows);------// b5
------while (at < n) data[at] += by, at |= at + 1; }------// fb -------for (int i = 0; i < rows; i++)------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);-------// bf
};------while (p) {------// cb
struct fenwick_tree_sq {------// d4 ------if (p & 1) res = res * sq;-----// c1
----int n; fenwick_tree x1, x0;-------// 18 -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;-------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {-------// 89
}:-----// 13 -------int k = r;------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f ------if (k >= rows) continue;------// 3f
----return s.query(b) - s.query(a-1); }------// f3 ------det *= T(-1);-------------------// 7a
                               ------for (int i = 0: i < cols: i++)-----// ab
2.4. Matrix. A Matrix class.
                               -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
class matrix {-----// 85
```

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-----for (int i = 0; i < rows; i++) {--------// 3d ------node *s = successor(n);------// e5
-----T m = mat(i, c):--------// e8 ------erase(s, false):--------// 0a
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 5a
-----} return mat; }------// 8f ------parent_leg(n) = s, fix(s);------// 82
private:-----// e0 -----return;-------// e5
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;------// 43
-----if (!n) return NULL;-----// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                            ------if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0------// b5 -----node *p = n->p;------// a7
template <class T>------// 22 -----return p; }------// c7
class avl_tree {------// ff ----node* predecessor(node *n) const {------// b4
public:-----// f6 -----if (!n) return NULL;------// dd
----struct node {-------// 45 -----if (n->l) return nth(n->l->size-1, n->l);------// 10
------T item; node *p, *l, *r;-------// a6 ------node *p = n->p;------// ea
------node(const T \&_item, node *_p = NULL) : item(_item), p(_p),-------// 4f ------return p; }------
------l(NULL), r(NULL), size(1), height(0) { } };-------// @d ----inline int size() const { return sz(root); }-----// ef
----node *root;------// 91 ----node* nth(int n, node *cur = NULL) const {------// e4
------while (cur) {-------// 8b ------if (n < sz(cur->l)) cur = cur->l;------// 75
------else if (item < cur->item) cur = cur->l;------// cd -----else break;-----
------else break; }------// 4f ------} return cur; }------// ed
-----return cur; }------// 84 ----int count_less(node *cur) {-------// ec
----node* insert(const T &item) {-------// 4e -----int sum = sz(cur->l);------// bf
-------while (*cur) {-------// aa -------if (cur->p->r == cur) sum += 1 + sz(cur->p->l);------// 5c
------prev = *cur;------// f0 ------cur = cur->p;-----// eb
#if AVL_MULTISET-----// 0a private:----// d5
#else-----// ff ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
-----else return *cur;------// 54 -----return n && height(n->r); }-----// a8
#endif-----// af ----inline bool right_heavy(node *n) const {------------// 27
-----*cur = n, fix(n); return n; }-----// 29 -----return n && abs(height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }------// 67 ----void delete_tree(node *n) {-------// fd
----void erase(node *n, bool free = true) {-------// 58 ------if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }-----// ef
------if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;-------// 12 -----if (!n->p) return root;-------// ac
------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;------// 6b ------if (n->p->l == n) return n->p->l;------// 83
-----else if (n->l && n->r) {-------// 6c
```

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------if (n->p->r == n) return n->p->r;-------// cc template <class Compare = default_int_cmp>------// 30
------while (i > 0) {------// 1a
-----parent_leg(n) = l; \[ \]-----// fc
                               -----int p = (i - 1) / 2;-----// 77
-----n->l = l->r; N-----// e8
                               -----if (!cmp(i, p)) break;-----// a9
----void right_rotate(node *n) { rotate(l, r); }------// ca -------if (l >= count) break;-------// d5
------if (too_heavy(n)) {-------// 4c -----swp(m, i), i = m; } }-----// 1d
------else if (right_heavy(n) & left_heavy(n->r))------// b9 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {------// b9
-----right_rotate(n->r);------// 08 -----q = new int[len], loc = new int[len];-----// f8
-----if (left_heavy(n)) right_rotate(n);------// 93 ------memset(loc, 255, len << 2); }------// f7
------else left_rotate(n);------// d5 ----~heap() { delete[] loc; }------// 09
-----n = n->p; }------// 28 ----void push(int n, bool fix = true) {-------// b7
-----n = n->p; } };-------// a2 ------if (len == count || n >= len) {--------// 0f
                               #ifdef RESIZE-----// a9
Also a very simple wrapper over the AVL tree that implements a map interface.
                               -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"------// 01
                               ------while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                               -----int *newq = new int[newlen], *newloc = new int[newlen];------// e3
template <class K, class V>-----// da
                               ------for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --// 94
class avl_map {------// 3f
                               -----/memset(newloc + len, 255, (newlen - len) << 2);-----// 18
public:----// 5d
                               -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                               -----loc = newloc, q = newq, len = newlen;-----// 61
-----K key; V value;-----// 32
                               #else-----// 54
-----node(K k, V v) : key(k), value(v) { }------// 29
                               -----assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                               #endif-----// 64
----avl_tree<node> tree;-----// b1
                               ------}------// 4b
----V& operator [](K key) {------// 7c
                               -----assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(\theta)));------// ba
                               -----loc[n] = count, q[count++] = n;------// 6b
-----if (!n) n = tree.insert(node(key, V(0)));-----// cb
                               -----if (fix) swim(count-1): }------// bf
-----return n->item.value:-----// ec
                               ----void pop(bool fix = true) {-------// 43
----}-------// 2e
                               -----assert(count > 0);-----// eb
};-----// af
                               -----loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 50
                               -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                               #define RESIZE-----// d0
                               ----int top() { assert(count > 0); return q[0]; }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                               ----void heapify() { for (int i = count - 1; i > 0; i--)------// 39
struct default_int_cmp {------// 8d
                               ------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }-----// 0b
----default_int_cmp() { }------// 35
                               ----void update_key(int n) {------// 26
----bool operator ()(const int &a, const int &b) { return a < b; } };------// e9
```

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----int size() { return count; }--------------------// 86 -------x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];-------// bc
----void clear() { count = 0, memset(loc, 255, len << 2); } };-------// 58 ------update[i]->next[i] = x;--------// 20
                                      ------update[i]->lens[i] = pos[0] + 1 - pos[i];------// 42
2.7. Skiplist. An implementation of a skiplist.
#define BP 0.20-----// aa
                                      ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// 07
                                      -----size++;-----// 19
#define MAX_LEVEL 10-----// 56
                                      -----return x; }------// c9
unsigned int bernoulli(unsigned int MAX) {-----// 7b
                                      ----void erase(T target) {------// 4d
----unsigned int cnt = 0;-----// 28
                                      ------FIND_UPDATE(x->next[i]->item, target);-------// 6b
----while(((float) rand() / RAND_MAX) < BP \&\& cnt < MAX) cnt++;-----// d1
                                      -----if(x && x->item == target) {------// 76
----return cnt; }-----// a1
                                      ------for(int i = 0; i <= current_level; i++) {-------// 97
template<class T> struct skiplist {------// 34
                                      -----if(update[i]->next[i] == x) {-------// b1
----struct node {------// 53
                                      -----update[i]->next[i] = x->next[i];------// 59
                                      -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
-----int *lens:-----// 07
-----node **next:-----// 0c
                                      -----} else update[i]->lens[i] = update[i]->lens[i] - 1;-----// 88
                                      ------#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))-------// 25
                                       ------delete x; _size--;------// 81
-----node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}-----// 7c
                                      ------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
------node() { free(lens); free(next); }; };------// aa
                                      -----current_level--; } };-----// 59
----int current_level, _size;-----// 61
----node *head;------// b7
                                      2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                      list supporting deletion and restoration of elements.
----~skiplist() { clear(); delete head; head = NULL; }-----// aa
                                      template <class T>-----// 82
----#define FIND_UPDATE(cmp, target) \-----// c3
                                      struct dancing_links {-----// 9e
------int pos[MAX_LEVEL + 2]; \[\bar{\sqrt{-------//18}}
                                      ----struct node {------// 62
-----// f2
                                      -----T item;-----// dd
-----node *x = head; \[\bigcup_{-----// 0f}\]
                                      -----node *l. *r:-----// 32
-----node *update[MAX_LEVEL + 1]; N-----// 01
                                      -----node(const T &_item, node *_l = NULL, node *_r = NULL)------// 6d
                                      ----: item(_item), l(_l), r(_r) {------// 6d
-----if (l) l->r = this;-----// 97
-----for(int i = MAX\_LEVEL; i >= 0; i--) { \sqrt{\phantom{a}}
                                      -----if (r) r->l = this;-----// 81
-----pos[i] = pos[i + 1]; \mathbb{N}------// 68
                                      ----}:-----// d3
------update[i] = x; N-------// dd ----dancing_links() { front = back = NULL; }-----// 72
----void clear() { while(head->next && head->next[0])-------// 91 -----if (!front) front = back;-------------// d2
------erase(head->next[0]->item); }-------// e6 ------return back;--------------------------// cθ
------return x && x->item == target ? x : NULL; }-------// 50 ----node *push_front(const T &item) {-------// 4a
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------if (!back) back = front;--------------------// 10
-----return pos[0]; }------// 19 -----return front;------// cf
------FIND_UPDATE(x->next[i]->item, target);--------// 3a ----void erase(node *n) {--------------------------// a0
------if(x && x->item == target) return x; // SET-------// 07 -----if (!n->l) front = n->r; else n->l->r = n->r; -----// ab
------int lvl = bernoulli(MAX_LEVEL);-------// 7a -----if (!n->r) back = n->l; else n->r->l = n->l;------// 1b
```

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----}------else if (p.coord[i] > to.coord[i])------// 83
};-----sum += pow(p.coord[i] - to.coord[i], 2.0);-----// 8c
                                 2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                                 -----return sqrt(sum); }-----// ef
element.
                                 -----bb bound(double l, int c, bool left) {------// b6
#define BITS 15------pt nf(from.coord), nt(to.coord);------// 5c
----int cnt[BITS][1<<BITS];------// aa ------else nf.coord[c] = max(nf.coord[c], l);------// 71
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }------// b0 -----return bb(nf, nt); } };------// 3b
----int nth(int n) {--------// 8a ------node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
------for (int i = BITS-1; i >= 0; i--)-------// 99 ----kd_tree() : root(NULL) { }-------// 57
------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4 ----kd_tree(vector<pt> pts) { root = construct(pts, 0, size(pts) - 1, 0); }----// 66
----}------if (from > to) return NULL;-------// f4
------nth_element(pts.begin() + from, pts.begin() + mid,------// f3
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                                 -----pts.begin() + to + 1, cmp(c));------// 97
bor queries.
                                 -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// dd
template <int K>-----// cd ----bool contains(const pt &p) { return _con(p, root, 0); }------// f0
class kd_tree {-----// 7e ----bool _con(const pt &p, node *n, int c) {------// 74
public:------if (!n) return false;------// b7
----struct pt {-------// 78 ------if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));-------// 81
-----pt() {}------// c1 -----return true; }------// 86
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }-----// 4c ----void insert(const pt &p) { _ins(p, root, 0); }------// f9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {-------// 09
------double sum = 0.0;------// c4 -----if (!n) n = new node(p, NULL, NULL);------// 4d
------for (int i = 0; i < K; i++)-------// 23 ------else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));------// a0
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }------// 2c
-----return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }-----// 73
----struct cmp {------// 8f -----// 8f -----// 1a
-----cmp(int _c) : c(_c) {}------// a5 -----assert(root);--------// f8
------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// bf
-----cc = i == 0 ? c : i - 1;------// bc -----pt from(cs);------// 5a
------if (abs(a.coord[cc] - b.coord[cc]) > EPS)------// 28 ------for (int i = 0; i < K; i++) cs[i] = INFINITY;-----// 37
-----return a.coord[cc] < b.coord[cc];-----// b7 -----pt to(cs);------
-----return false; } };------// 62
----struct bb {-------// 30 ----pair<pt, bool> _nn(------// e3
------bb(pt _from, pt _to) : from(_from), to(_to) {}-------// 57 ------if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// 19
-------double dist(const pt &p) {-------/ 3f ------/ 3f -----// 3f -----// 9f
------double sum = 0.0;------// d9 -----pt resp = n->p;-----// 6b
-----for (int i = 0; i < K; i++) {------// 65
```

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-----if (found) mn = min(mn, p.dist(resp));--------// 18 ---}------// 0b
-----node *n1 = n->l, *n2 = n->r;------// aa
                                              -----// 63
------for (int i = 0; i < 2; i++) {-------// 50
                                              ----return -1:-----// f5
                                              }-----// 03
------if (i == 1 || cmp(c)(n->p, p)) swap(n1, n2), swap(l1, l2);------// e^2
-----pair<pt, bool> res =-----// 33
                                              3.2. Single-Source Shortest Paths.
-----_nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// 72
------if (res.second \&\& (!found || p.dist(res.first) < p.dist(resp)))----// 76
                                               3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log |V|)
-----resp = res.first, found = true;-----// 3b
                                               time.
-----}-----// aa
                                               int *dist, *dad;-----// 46
-----return make_pair(resp, found); } };-----// dd
                                               struct cmp {-----// a5
                                               ----bool operator()(int a, int b) {-----// bb
                    3. Graphs
                                               -----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }-----// e6
3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
                                               };-----// 41
edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
                                               pair<int*, int*> dijkstra(int n, int s, vii *adj) {----------------// 53
graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
                                               ----dist = new int[n];-----// 84
connected. It runs in O(|V| + |E|) time.
                                               ----dad = new int[n];-----// 05
int bfs(int start, int end, vvi& adj_list) {------// d7
                                               ----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;------// d6
----queue<ii>> 0;------// 75
                                               ----set<int, cmp> pq;-----// 04
----Q.push(ii(start, 0));-----// 49
                                               ----dist[s] = 0, pq.insert(s);-----// 1b
-----// 0b
                                               ----while (!pq.empty()) {------// 57
----while (true) {------// 0a
                                               ------int cur = *pq.begin(); pq.erase(pq.begin());-----// 7d
-----ii cur = Q.front(); Q.pop();-----// e8
                                               ------for (int i = 0; i < size(adj[cur]); i++) {-------// 9e
-----// 06
                                               ------int nxt = adj[cur][i].first,-----// b8
-----if (cur.first == end)------// 6f
                                               ------ndist = dist[cur] + adj[cur][i].second;------// 0c
-----return cur.second:-----// 8a
                                               ------if (ndist < dist[nxt]) pq.erase(nxt),------// e4
-----// 3c
                                               -----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);-----// 0f
-----vi& adi = adi_list[cur.first]:-----// 3f
                                               ------for (vi::iterator it = adj.begin(); it != adj.end(); it++)-----// bb
                                               ----}----// e8
-----0.push(ii(*it, cur.second + 1));-----// b7
                                               ----return pair<int*, int*>(dist, dad);-----// cc
}-----// af
}-----// 7d
                                               3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
 Another implementation that doesn't assume the two vertices are connected. If there is no path
                                               problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
from the starting vertex to the ending vertex, a-1 is returned.
                                               negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
int bfs(int start, int end, vvi& adj_list) {------// d7
----set<int> visited; ------// b3
----queue<ii>> 0:------// bb
                                               int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {------// cf
                                               ----has_negative_cycle = false;------// 47
----Q.push(ii(start, 0));-----// 3a
                                               ----int* dist = new int[n];------// 7f
----visited.insert(start);-----// b2
                                               ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
-----// db
----while (!Q.empty()) {------// f7
                                               ----for (int i = 0; i < n - 1; i++)------// a1
                                               ------for (int j = 0; j < n; j++)------// c4
-----ii cur = Q.front(); Q.pop();-----// 03
                                              -----if (dist[j] != INF)-----// 4e
-----// 9c
                                              ------for (int k = 0; k < size(adj[j]); k++)-----// 3f
-----if (cur.first == end)-----// 22
-----return cur.second:-----// b9
                                               -----dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61
-----// ba
                                               -----dist[j] + adj[j][k].second);------// 47
                                               ----for (int j = 0; j < n; j++)-----// 13
-----vi& adj = adj_list[cur.first];-----// f9
                                               ------for (int k = 0; k < size(adj[j]); k++)------// a0
------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)-----// 44
                                               -----if (dist[j] + adj[j][k].second < dist[adj[j][k].first])------// ef
-----if (visited.find(*it) == visited.end()) {------// 8d
-----Q.push(ii(*it, cur.second + 1));-----// ab
                                               -----has_negative_cycle = true;------// 2a
                                               ----return dist:-----// 2e
-----visited.insert(*it);-----// cb
```

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```
3.3. All-Pairs Shortest Paths.
                                             3.5.1. Kruskal's algorithm.
                                             #include "../data-structures/union_find.cpp"-----// 5e
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                             -----// 11
problem in O(|V|^3) time.
                                             // n is the number of vertices-----// 18
void floyd_warshall(int** arr, int n) {------// 21
                                             // edges is a list of edges of the form (weight, (a, b))-----// c6
----for (int k = 0; k < n; k++)------// 49
                                             // the edges in the minimum spanning tree are returned on the same form-----// 4d
------for (int i = 0; i < n; i++)------// 21
                                             vector<pair<int, ii> > mst(int n, vector<pair<<mark>int</mark>, ii> > edges) {------// a7
-----for (int j = 0; j < n; j++)-----// 77
                                             ----union_find uf(n):------// 04
-----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
                                             ----sort(edges.begin(), edges.end());-----// 51
------arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                                             ----vector<pair<int, ii> > res;------// 71
}-----// 86
                                             ----for (int i = 0; i < size(edges); i++)-----// ce
                                             -----if (uf.find(edges[i].second.first) !=-----// d5
3.4. Strongly Connected Components.
                                             -----uf.find(edges[i].second.second)) {------// 8c
3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                             -----res.push_back(edges[i]);-----// d1
graph in O(|V| + |E|) time.
                                             -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
#include "../data-structures/union_find.cpp"---------------------// 5e
                                             ------}------// 5b
-----//
                                             ----return res;------// 46
vector<br/>bool> visited;-----// 66
vi order:-----// 9b
                                             3.6. Topological Sort.
-----// a5
void scc_dfs(const vvi &adj, int u) {------// a1
                                             3.6.1. Modified Depth-First Search.
----int v; visited[u] = true;------// e3
                                             void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
----for (int i = 0; i < size(adj[u]); i++)-----// c5
                                             ------bool& has_cycle) {-------// a8
------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);------// 6e
                                             ----color[cur] = 1;------// 5b
----order.push_back(u);-----// 19
                                             ----for (int i = 0; i < size(adj[cur]); i++) {-------// 96
}-----// dc
                                             ------int nxt = adj[cur][i];------// 53
  -----// 96
                                             -----if (color[nxt] == 0)------// 00
pair<union_find, vi> scc(const vvi &adj) {------// 3e
                                             ------tsort_dfs(nxt, color, adj, res, has_cycle);-----------// 5b
----int n = size(adj), u, v;-----// bd
                                             -----else if (color[nxt] == 1)------// 53
----order.clear();-----// 22
                                             -----has_cycle = true;-----// c8
----union_find uf(n);------// 6d
                                             -----if (has_cycle) return;------// 7e
----vi dag;------// ae
                                             ----}------// 3d
----vvi rev(n);------// 20
                                             ----color[cur] = 2;-----// 16
----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++)------// b9
                                             ----res.push(cur):-----// cb
-----rev[adj[i][j]].push_back(i);-----// 77
                                             }------// 9e
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04
                                             -----// ae
----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);------// e4
                                             vi tsort(<mark>int</mark> n, vvi adj, <mark>bool</mark>& has_cycle) {----------------------------// 37
----fill(visited.begin(), visited.end(), false);------// c2
                                             ----has_cycle = false;-----// 37
----stack<<u>int</u>> S;-----// 04
                                             ----stack<int> S;-----// 54
----for (int i = n-1; i >= 0; i--) {------// 3f
                                             ----vi res:-----// d1
-----if (visited[order[i]]) continue;------// 94
                                             ----char* color = new char[n];-----// b1
-----S.push(order[i]), dag.push_back(order[i]);------// 40
                                             ----memset(color, 0, n):-----// ce
------while (!S.empty()) {------// 03
                                             ----for (int i = 0; i < n; i++) {-------// 96
-----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);-----// 1b
                                             ------if (!color[i]) {------// d5
-----for (int j = 0; j < size(adj[u]); j++)-----// 21
                                             -----tsort_dfs(i, color, adj, S, has_cycle);-----// 40
------if (!visited[v = adj[u][j]]) S.push(v);------// e7
                                             ----}-------// df
----return pair<union_find, vi>(uf, dag);-----// f2
                                             ----while (!S.empty()) res.push_back(S.top()), S.pop();------// 94
                                             ----return res;------// 07
                                             }-----// 1f
3.5. Minimum Spanning Tree.
```

3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-

ing. Running time is  $O(|E|\sqrt{|V|})$ .

the maximum flow of a flow network.

#define MAXV 2000-----// ba
int q[MAXV], d[MAXV];-----// e6

```
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----struct edge {-------// 1e ------int v, cap, nxt;------// cb
------edge(int _v, int _cap, int _nxt) : v(_v), cap(_cap), nxt(_nxt) { }-----// bc ----int n, ecnt, *head;----------------------------------// 39
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// d3 ------memset(head = new int[n], -1, n << 2);------// 58
-----e.reserve(2 * (m == -1 ? n : m));------// 24 ---}-----// 3a
------head = new int[n], curh = new int[n];------// 6b ----void destroy() { delete[] head; }-----// d5
------memset(head, -1, n * sizeof(int));--------// 56 ----void reset() { e = e_store; }-------// 1b
----void reset() { e = e_store; }------// 87 ------e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// bc
----void add_edge(int u, int v, int uv, int vu = 0) {-------// cd ----}-----// ef
------for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)-----// c6
-----if(s == t) return 0;-------// 9d -----if (p[t] == -1) break;------// c8
-----e_store = e:------// 57 ------<u>int</u> x = INF, at = p[t];-----// f1
------int f = 0, x, l, r;-------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// aa
------while (true) {-------// b5 -----at = p[t], f += x;------// 6f
-----while (l < r)------// 7a -----}-----// 7a
-----if (d[s] == -1) break;-----// a0
                      };-----// 75
-----memcpy(curh, head, n * sizeof(int));-----// 10
                      3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
fied to find shortest path to augment each time (instead of just any path). It computes the maximum
flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with
-----if (res) reset();------// 21
                      minimum cost. Running time is O(|V|^2|E|\log|V|).
-----return f:-----// b6
                      #define MAXV 2000-----// ba
----}------// 1b
                      int d[MAXV], p[MAXV], pot[MAXV];-----// 80
};-----// 3b
                      struct cmp {-----// d1
3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
                      ----bool operator ()(int i, int j) {------// 8a
O(|V||E|^2). It computes the maximum flow of a flow network.
                      -----return d[i] == d[j] ? i < j : d[i] < d[j];------// 89
#define MAXV 2000-----// ba
                      ----}-----// df
struct flow_network {------// 5e struct flow_network {-----// eb
```

```
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};-----// d7
------int v, cap, cost, nxt;-----// ad
------dge(int _v, int _cap, int _cost, int _nxt)------// ec
                                          3.11. All Pairs Maximum Flow.
-----: v(_v), cap(_cap), cost(_cost), nxt(_nxt) { }-----// c4
----}:------// ad
                                          3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----int n, ecnt, *head;-----// 46
                                          structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the max-
----vector<edge> e, e_store;------// 4b
                                          imum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// dd
                                          #include "dinic.cpp"-----// 58
-----e.reserve(2 * (m == -1 ? n : m));-----// e6
                                                       -----// 25
-----memset(head = new int[n], -1, n << 2);------// 6c
pair<vii, vvi> construct_gh_tree(flow_network &g) {------// 77
----void destroy() { delete[] head; }------// ac
                                          ----int n = q.n, v;------// 5d
----void reset() { e = e_store; }------// 88
                                          ----vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));------// 49
----void add_edge(int u, int v, int cost, int uv, int vu=0) {------// b4
                                          ----for (int s = 1; s < n; s++) {------// 9e
-----e.push_back(edge(v, uv, cost, head[u])); head[u] = ecnt++;------// 43
                                          -----int l = 0, r = 0;-----// 9d
------e.push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;------// 53
                                          -----par[s].second = q.max_flow(s, par[s].first, false);------// 38
----}------// 16
                                          -----memset(d, 0, n * sizeof(int));-----// 79
-----memset(same, 0, n * sizeof(int));-----// b0
-----if (s == t) return ii(0, 0);-----// 34
                                          -\cdots -d[q[r++] = s] = 1;
-----e_store = e:-----// 70
                                          -----while (l < r) {------// 45
-----memset(pot. 0. n << 2):-----// 24
------int f = 0, c = 0, v;-----// d4
                                          -----same[v = q[l++]] = true;-----// c8
                                          ------for (int i = q.head[v]; i != -1; i = q.e[i].nxt)------// 33
------while (true) {------// 29
                                          ------if (q.e[i].cap > 0 \& d[q.e[i].v] == 0)-----// 3f
-----memset(d, -1, n << 2);-----// fd
                                          -----d[q[r++] = g.e[i].v] = 1;-----// f8
-----memset(p, -1, n << 2);-----// b7
-----set<<u>int</u>, cmp> q;-----// d8
                                          -----for (int i = s + 1; i < n; i++)-----// 68
-----q.insert(s); d[s] = 0;-----// 1d
                                          -----if (par[i].first == par[s].first && same[i]) par[i].first = s;----// ea
-----while (!q.empty()) {------// 04
                                          -----q.reset();------// 9a
-----int u = *q.begin();-----// dd
                                          ----}-----// 1e
-----q.erase(q.begin());-----// 20
                                          ----for (int i = 0; i < n; i++) {------// 2a
-----for (int i = head[u]; i != -1; i = e[i].nxt) {------// 02
                                          ------int mn = INF, cur = i;-----// 19
-----if (e[i].cap == 0) continue;-----// 1c
                                          ------while (true) {------// 3a
-----int cd = d[u] + e[i].cost + pot[u] - pot[v = e[i].v]; -----// 1d
                                          -----cap[cur][i] = mn;-----// 63
-----if (d[v] == -1 || cd < d[v]) {------// d2
                                          -----if (cur == 0) break;-----// 35
-----if (q.find(v) != q.end()) q.erase(q.find(v));-----// e^2
                                          -----/mn = min(mn, par[cur].second), cur = par[cur].first;-----// 28
-----d[v] = cd; p[v] = i;------// f7
                                          -----q.insert(v);------// 74
                                          ----}------// 4a
----return make_pair(par, cap);-----// 6b
-----// 99
int compute_max_flow(int s, int t, const pair<vii, vvi> &gh) {------// 16
-----if (p[t] == -1) break;-----// 09
                                          ----if (s == t) return 0;-----// d4
-----int x = INF, at = p[t];-----// e8
                                          ----int cur = INF, at = s;------// 65
------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];-----// 32
                                          ----while (gh.second[at][t] == -1)------// ef
-----at = p[t], f += x;------// 43
                                          -----cur = min(cur, gh.first[at].second), at = gh.first[at].first;-----// bd
-----while (at != -1)-----// 53
                                          ----return min(cur, gh.second[at][t]);-----// 6d
-----e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// 95
                                          }-----// a2
-----c += x * (d[t] + pot[t] - pot[s]);------// 44
-----for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];------// ff
                                          3.12. Heavy-Light Decomposition.
struct HLD {-----// 7a
-----if (res) reset();------// 5e
                                          ----int n:------// 0d
-----return ii(f, c);------// e7
```

```
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----vvi below;-------// 06 ---}------// cc
----vi head, parent, loc;-----// 25
----int curhead, curloc;------// 83
                                                             4. Strings
                                          4.1. Trie. A Trie class.
----HLD(int_n): n(n), sz(n, 1), below(n), head(n), parent(n, -1), loc(n) { }// lc
                                          template <class T>-----// 82
-----// 57
                                          class trie {------// 9a
----void add_edge(int u, int v) {------// 29
------below[u].push_back(v);-----// ad
                                           ----struct node {------// ae
-----parent[v] = u;------// 22
                                           -----map<T, node*> children;-----// aθ
                                           ------int prefixes, words;------// e2
                                           -----node() { prefixes = words = 0; } };------// 42
----void csz(int u) {------// 3b
-----for (int i = 0; i < size(below[u]); i++)-----// 14
                                           ----node* root;------// a9
-----csz(below[u][i]), sz[u] += sz[below[u][i]];-----// ab
                                           ----trie() : root(new node()) { }------// 8f
----}-----// 5c
                                           ----template <class I>------// 89
                                           ----void insert(I begin, I end) {------// 3c
----void part(int u) {------// c8
                                           -----node* cur = root;------// 82
-----head[u] = curhead;------// b5
                                           ------while (true) {-------// 67
-----loc[u] = curloc++;------// 01
                                           -----cur->prefixes++;-----// f1
                                           ------if (begin == end) { cur->words++; break; }------// db
-----int best = -1;------// dd
-----for (int i = 0; i < size(below[u]); i++)------// 67
                                             -----T head = *begin;-----// fb
------if (best == -1 || sz[below[u][i]] > sz[best])------// 30
                                             ------typename <code>map<T, node*>::const_iterator it;-------------// 	heta 1</code>
-----best = below[u][i];-----// 63
                                               -----it = cur->children.find(head);------// 77
                                             ------if (it == cur->children.end()) {------// 95
  ---if (best != -1)------// c1
                                             -----pair<T, node*> nw(head, new node());-----// cd
-----part(best);-----// e2
                                             -----// da
                                           -----} begin++, cur = it->second; } } -----// 64
------for (int i = 0; i < size(below[u]); i++)-----// 1d
                                           ----template<class I>-----// b9
-----if (below[u][i] != best)-----// c1
                                           ----int countMatches(I begin, I end) {------// 7f
-----part(curhead = below[u][i]);-----// 4d
                                           -----node* cur = root;------// 32
                                           -----while (true) {------// bb
                                           -----if (begin == end) return cur->words;------// a4
----void build() {------// 38
                                            -----else {------// 1e
------int u = curloc = 0;------// 3a
                                             -----T head = *begin;-----// 5c
------while (parent[u] != -1) u++;------// 85
                                             -----typename map<T, node*>::const_iterator it;------// 25
-----CSZ(u);-----// 90
                                             -----it = cur->children.find(head);-----// d9
-----part(curhead = u);-----// b2
                                             -----if (it == cur->children.end()) return 0;------// 14
                                           -----begin++, cur = it->second; } } }-----// 7c
                                           ----template<class I>-----// 9c
----int lca(int u, int v) {------// 60
                                           ----int countPrefixes(I begin, I end) {------// 85
-----vi uat, vat;-----// b5
                                           -----node* cur = root;------// 95
------while (u != -1) uat.push_back(u), u = parent[head[u]];------// 65
                                           ------while (true) {------// 3e
------while (v != -1) vat.push_back(v), v = parent[head[v]];------// 40
                                           -----if (begin == end) return cur->prefixes;------// f5
-----u = size(uat) - 1, v = size(vat) - 1;------// 52
                                           -----else {------// 66
-----int res = -1;------// 06
                                             -----T head = *begin;-----// 43
------while (u >= 0 \&\& v >= 0 \&\& head[uat[u]] == head[vat[v]])------// 8e
                                             ------typename map<T, node*>::const_iterator it;------// 7a
-----res = (loc[uat[u]] < loc[vat[v]] ? uat[u] : vat[v]), u--, v--;----// 49
                                             -----it = cur->children.find(head);------// 43
-----return res;------// d3
                                           ------if (it == cur->children.end()) return Θ;------// 71
                                           -----begin++, cur = it->second; } } } ;-----// 26
```

```
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                                           15
-----l = r = i;-------// a7 struct intx {-------// cf
-----z[i] = r - l; r--;--------// fc ----intx(string n) { init(n); }-------// b9
-----z[i] = r - l; r--; } }------// 8d ----static const int dcnt = 9;---------// 12
----return z;------// 53 ----static const unsigned int radix = 1000000000U;------// f0
}-----// db ----int size() const { return data.size(); }------// 29
                      ----void init(string n) {------// 13
                      -----intx res; res.data.clear();-----// 4e
        5. Mathematics
                      -----if (n.empty()) n = "0";------// 99
5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common
                      -----if (n[0] == '-') res.sign = -1, n = n.substr(1);------// 3b
                      ------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {------// e7
private: ------int idx = i - j; ------// cd
public:-----digit = digit * 10 + (n[idx] - '0');------// 1f
----fraction(T n_, T d_) {------------------------// 03 ------res.data.push_back(digit);-----------------// 07
-----assert(d_ != 0);-----// 3d ----}----// 7d
-----n = n_, d = d_;------// 06 ------data = res.data;------// 7d
------if (d < T(0)) n = -n, d = -d;------// be ------normalize(res.sign);------// 76
----fraction(T n_) : n(n_), d(1) { }------// 84 -------if (data.empty()) data.push_back(0);------// fa
------return fraction<T>(n * other.d + other.n * d, d * other.d);}------// 3b ------sign = data.size() == 1 \&\& data[0] == 0 ? 1 : nsign;------// ff
----fraction<T> operator *(const fraction<T>& other) const {-------// 38 ----friend ostream& operator <<(ostream& outs, const intx& n) {-------// 0d
------for (int i = n.size() - 1; i >= 0; i--) {-------// 63
------stringstream ss; ss << cur;-------// 8c
------return other < *this; }------// 6e -----int len = s.size();-----// 0d
-----return n == other.n && d == other.d; }------// 14 -----}
----bool operator !=(const fraction<T>& other) const {-------// ec -----return outs;-----
-----return !(*this == other); }------// d1 ---}-----// b9
};------// 12 ----string to_string() const { stringstream ss; ss << *this; return ss.str(); }// fc
                      ----bool operator <(const intx& b) const {------// 21
5.2. Big Integer. A big integer class.
```

```
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-----if (sign != b.sign) return sign < b.sign;-------// cf -----assert(!(d.size() == 1 && d.data[0] == 0));------// e9
------if (size() != b.size())-------// 4d ------intx q, r; q.data.assign(n.size(), 0);------// ca
-----return sign == 1 ? size() < b.size() : size() > b.size();------// 4d -------for (int i = n.size() - 1; i >= 0; i--) {--------// 1a
------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35 ------r.data.insert(r.data.begin(), 0);--------// c7
------return false;-------// ca -------long long k = 0;--------// cc
------unsigned long long carry = 0;-------// 5c ------return pair<intx, intx>(q.normalize(n.sign * d.sign), r);------// a1
-----carry += (i < size() ? data[i] : OULL) +------// 91 ----intx operator /(const intx& d) const {-------// a2
-----c.data.push_back(carry % intx::radix);-------// 86 ----intx operator %(const intx& d) const {--------// 07
-----carry /= intx::radix;------// fd -----return divmod(*this,d).second * sign; }-----// 5a
-----return c.normalize(sign);------// 20
                                      5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
----}------// 70
                                      #include "intx.cpp"-----// 83
----intx operator -(const intx& b) const {--------// 53
                                      #include "fft.cpp"-----// 13
-----if (sign > 0 && b.sign < 0) return *this + (-b);------// 8f
                                      -----// e0
------if (sign < 0 && b.sign > 0) return -(-*this + b);-----// 1b
                                      intx fastmul(const intx &an, const intx &bn) {------// ab
------if (sign < 0 && b.sign < 0) return (-b) - (-*this);------// a1
                                      ----string as = an.to_string(), bs = bn.to_string();-----// 32
-----if (*this < b) return -(b - *this);-----// 36
                                      ----int n = size(as), m = size(bs), l = 1,-----// dc
-----intx c; c.data.clear();------// 6b
                                      -----len = 5, radix = 100000,-----// 4f
-----long long borrow = 0;-----// f8
                                      -----*a = new int[n], alen = 0,-----// b8
------for (int i = 0; i < size(); i++) {-------// a7
                                      -----*b = new int[m], blen = 0;------// 0a
-----borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
                                      ----memset(a, 0, n << 2);-----// 1d
-----c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
                                      ----memset(b, 0, m << 2);-----// 01
------borrow = borrow < 0 ? 1 : 0;-----// 0d
                                      ----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
-----return c.normalize(sign);-----// 35
                                      -----for (int j = min(len - 1, i); j >= 0; j--)------// 43
                                      -----a[alen] = a[alen] * 10 + as[i - j] - '0';------// 14
----intx operator *(const intx& b) const {------// bd
                                      ----for (int i = m - 1; i >= 0; i -= len, blen++)------// b6
-----intx c; c.data.assign(size() + b.size() + 1, 0);-----// d0
                                      ------for (int j = min(len - 1, i); j >= 0; j--)------// ae
------for (int i = 0; i < size(); i++) {------// 7a
                                      -----b[blen] = b[blen] * 10 + bs[i - j] - 0;------// 9b
                                      ----while (l < 2*max(alen,blen)) l <<= 1;------// 51
-----long long carry = 0;-----// 20
                                      ----cpx *A = new cpx[l], *B = new cpx[l];------// \theta d
------for (int j = 0; j < b.size() || carry; j++) {------// c0
                                      ----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);------// 35
-----if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
-----carry += c.data[i + j];-----// 18
                                      ----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? b[i] : 0, 0);------// 66
-----c.data[i + j] = carry % intx::radix;-----// 86
                                      ----fft(A, l); fft(B, l);-----// f9
-----carry /= intx::radix;-----// 05
                                      ----for (int i = 0; i < l; i++) A[i] *= B[i];------// e7
                                      ----fft(A, l, true);------// d3
----ull *data = new ull[l];-----// e7
-----return c.normalize(sign * b.sign);-----// de
                                      ----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
                                      ----for (int i = 0; i < l - 1; i++)------// 90
-----if (data[i] >= (unsigned int)(radix)) {------// 44
----friend pair<intx,intx> divmod(const intx& n, const intx& d) {------// fb
                                      -----data[i+1] += data[i] / radix;-----// e4
```

```
----return res;------// 03
}-----// 1c
```

5.11. **Numeric Integration.** Numeric integration using Simpson's rule.

5.12. **Fast Fourier Transform.** The Cooley-Tukey algorithm for quickly computing the discrete Fourier transform. Note that this implementation only handles powers of two, make sure to pad with zeros.

```
#include <complex>-----// 8e
typedef complex<long double> cpx;-----// 25
void fft(cpx *x, int n, bool inv=false) {------// 23
----for (int i = 0, j = 0; i < n; i++) {------// f2
------if (i < j) swap(x[i], x[j]);------// 5c
-----int m = n>>1:-----// e5
------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
-----j += m;-----// ab
----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
------for (int m = 0; m < mx; m++, w *= wp) {-------// 40
------for (int i = m; i < n; i += mx << 1) {------// 33
-----cpx t = x[i + mx] * w;-----// f5
-----x[i + mx] = x[i] - t;
-----x[i] += t;-----// c7
----}------------// 70
----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);------// 3e
}-----// 7d
```

## 5.13. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once:  $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times:  $n^k$
- Number of permutations of n objects, where there are  $n_1$  objects of type 1,  $n_2$  objects of type 2, ...,  $n_k$  objects of type k:  $\binom{n}{n_1, n_2, \dots, n_k} = \frac{n!}{n_1! \times n_2! \times \dots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:
- $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times:  $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to  $x_1 + x_2 + \cdots + x_n = k$  where  $x_i \ge 0$ :  $f_k^n$
- Number of subsets of a set with n elements:  $2^n$

- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an  $n \times m$  grid by walking only up and to the right:  $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced:  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an  $n \times n$  lattice which do not rise above the main diagonal:  $C_n$
- Number of permutations of n objects with exactly k ascending sequences or runs:  $\binom{n}{k} = \binom{n}{n-k-1} = k \binom{n-1}{k} + (n-k+1) \binom{n-1}{k-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{n-1} = \binom{n}{1}$
- Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements):  $D_0 = 1, D_1 = 0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points:  $\binom{n}{k}D_{n-k}$
- Heron's formula: A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} 1$ .
- **Divisor sigma:** The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where  $n = \prod_{i=0}^r p_i^{a_i}$  is the prime factorization.
- Divisor count: A special case of the above is  $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$  where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.

#### 6. Geometry

6.1. **Primitives.** Geometry primitives.

```
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                                                           19
bool collinear(L(a, b), L(p, q)) {------// 66 #define CHK(f,a,b,c) (f(a) < f(b) & f(b) <= f(c) & ccw(a,c,b) < 0}
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------------// cc -----for (int i = 0, j = n - 1; i < n; j = i++)-------------// 77
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 9e -------0 <= (d = progress(q, p[i], p[j])) && d <= 1)------// b9
double progress(P(p), L(a, b)) {-------// d2 -----return 0;-----
------if (CHK(real, p[i], q, p[j]) || CHK(real, p[i], q, p[i]) ------// 1f
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) {-------// d6 ----return in ? -1 : 1; }------------------------// 77
----// NOTE: check for parallel/collinear lines before calling this function---// 02 // pair<polygon, polygon cut_polygon(const polygon &poly, point a, point b) {-// 7b
----point r = b - a, s = q - p;-------// 79 //---- polygon left, right;-------// 6b
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// a8 //---- point it(-100, -100);-------// c9
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))-------// ae //---- for (int i = 0, cnt = poly.size(); i < cnt; i++) {--------------// 28
----res = a + t * r; -------// ca //------ point p = poly[i], q = poly[i]; -------// 19
}------if (ccw(a, b, p) >= 0) right.push_back(p);------// e3
------if (dot(b - a, c - b) > 0) return b;------// b5 //----- if (myintersect(a, b, p, q, it))-----// f0
------if (dot(a - b, c - a) > 0) return a:------// cf //------ left.push_back(it), right.push_back(it):------// 21
----double t = dot(c - a, b - a) / norm(b - a);------// aa //--- return pair<polygon, polygon>(left, right);------// 1d
}-----// e5
----double x = INFINITY;-----// 83 #include "polygon.cpp"-----// 58
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&------// ee ----return abs(real(a) - real(b)) > EPS ?------// 44
----else {------// 38 int convex_hull(polygon p) {-------// cd
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// ec ----sort(p.begin(), p.end(), cmp);------// 3d
----return x;------hull[l++] = p[i];------// f7
}-----// b3 ---}-----// d8
                              ----int r = 1:-----// 59
6.2. Polygon. Polygon primitives.
                              ----for (int i = n - 2; i >= 0; i--) {------// 16
#include "primitives.cpp"-----// e0
                              -----if (p[i] == p[i + 1]) continue;-----// c7
typedef vector<point> polygon;-----// b3
                              ------while (r - l >= 1 \&\& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 9f
double polygon_area_signed(polygon p) {------// 31
                              -----hull[r++] = p[i];-----// 6d
----double area = 0; int cnt = size(p);------// a2
                              ----}------// 74
----for (int i = 1; i + 1 < cnt; i++)------------------// d2 ----return l == 1 ? 1 : r - 1;-------------------// 6d
-----area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 7e
                              3-----// 79
----return area / 2: }-----// e1
double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------/25 6.4. Line Segment Intersection. Computes the intersection between two line segments.
```

```
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#include "primitives.cpp"------// e0 ----for (int i = 0, l = 0; i < size(pts); i++) {-------// ac
bool line_segment_intersect(L(a, b), L(c, d), point &A, point &B) {-------// 6c -------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);-------// 8b
------A = B = a; return abs(a - d) < EPS; }--------// ee ------jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));-------// 39
-----\&\& (abs(a - c) + abs(d - a) - abs(d - c)) < EPS: }-----// 27
----else if (abs(c - d) < EPS) {------// 26
                                               6.7. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
------A = B = c; double p = progress(c, a,b);------// d9
                                                  • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
------return 0.0 <= p && p <= 1.0------// 8e
                                                  • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
-----\&\& (abs(c - a) + abs(b - c) - abs(b - a)) < EPS; }-----// 4f
                                                  • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
----else if (collinear(a,b, c,d)) {------// bc
                                                    of that is the area of the triangle formed by a and b.
------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
-----if (ap > bp) swap(ap, bp);-----// b1
                                                                  7. Other Algorithms
------if (bp < 0.0 || ap > 1.0) return false;------// 0c
-----A = c + max(ap, 0.0) * (d - c); -----// f6
                                                7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
                                                function f on the interval [a, b], with a maximum error of \varepsilon.
-----B = c + min(bp, 1.0) * (d - c);------// 5c
-----return true; }-----// ab
                                                double binary_search_continuous(double low, double high,-----// 8e
----else if (parallel(a,b, c,d)) return false;-----// ca
                                                -----double eps, double (*f)(double)) {------// c0
----else if (intersect(a.b. c.d. A. true)) {------// 10
                                                ----while (true) {------// 3a
-----B = A; return true; }------// bf ------double mid = (low + high) / 2, cur = f(mid);------// 75
}-------| 8b ------else if (0 < cur) high = mid;------// e5
----}-------// b5
6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
coordinates) on a sphere of radius r.
                                                 Another implementation that takes a binary predicate f, and finds an integer value x on the integer
double gc_distance(double pLat, double pLong,-----// 7b
                                                interval [a,b] such that f(x) \wedge \neg f(x-1).
-----/ double qLat, double qLong, double r) {------// a4
                                                ----pLat *= pi / 180; pLong *= pi / 180;-----// ee
                                                ----assert(low <= high);-----// 19
----qLat *= pi / 180; qLong *= pi / 180;-----// 75
                                                ----while (low < high) {------// a3
----return r * acos(cos(pLat) * cos(qLat) * cos(pLong - qLong) +-----// e3
                                                -----sin(pLat) * sin(qLat));-----// 1e
                                                -----if (f(mid)) high = mid;------// ca
-----// 60
                                                -----else low = mid + 1;-----// 03
                                                ----}-----// 9b
                                                ----assert(f(low));------// 42
6.6. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
pair of points.
                                                ----return low:------// a6
                                                }-----// d3
#include "primitives.cpp"-----// e0
-----// 85
                                                7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotonically
struct cmpx { bool operator ()(const point &a, const point &b) {-----// 01
                                                cally decreasing, ternary search finds the x such that f(x) is maximized.
-----return abs(real(a) - real(b)) > EPS ?-----// e9
                                                template <class F>-----// d1
-----real(a) < real(b) : imag(a) < imag(b); } };-----// 53
struct cmpy { bool operator ()(const point &a, const point &b) {-----// 6f
                                                double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
                                                ----while (hi - lo > eps) {------// 3e
----return abs(imag(a) - imag(b)) > EPS ?-----// θb
                                                -----/double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
-----imag(a) < imag(b) : real(a) < real(b); } };------//
                                                -----if (f(m1) < f(m2)) lo = m1;------// 1d
double closest_pair(vector<point> pts) {------//
                                                -----else hi = m2;-----// b3
----sort(pts.begin(), pts.end(), cmpx());-----// 0c
                                                ----}------// bb
----set<point, cmpy> cur;-----// bd
----set<point, cmpy>::const_iterator it, jt;-----// a6
----double mn = INFINITY;-----// f9
```

```
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------for (int i = 0; i <= rows; i++)------// 96
                                      7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
-----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];-----// cb
-----ptr[rows][j]->size = cnt;-----// 59
------for (int i = 0; i <= rows; i++) delete[] ptr[i];------// bf
-----delete[] ptr;-----// 99
----#define COVER(c, i, j) N------// 6a
}-----// 42
-----for (node *i = c->d; i != c; i = i->d) \\ ------//
-----j->d->u = j->u, j->u->d = j->d, j->p->size--;-----// 16
----#define UNCOVER(c, i, j) \sqrt{\phantom{a}}-----// d\theta
------for (node *i = c->u; i != c; i = i->u) \[ \]------// ff
----bool search(int k = 0) {-------// bb ------d - 32075;------// e0
-----if (head == head->r) {-------// c3 }-----// fa
-----vi res(k);------vi res(k);-------// 9f void intToDate(int jd, int &w, int &d) {------// a1
------for (int i = 0; i < k; i++) res[i] = sol[i];------// 75 ---int x, n, i, j;----------// 00
-----sort(res.begin(), res.end());------// 87 ----x = jd + 68569;------// 11
-----return handle_solution(res);------// 51 ----n = 4 * x / 146097;------// 2f
------/-node *c = head->r, *tmp = head->r;-------// 8e ---i = (4000 * (x + 1)) / 1461001;------// 0d
------for ( ; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 00 ----x -= 1461 * i / 4 - 31;-----------------// 09
------if (c == c->d) return false;------// b0 ----j = 80 * x / 2447;------// 3d
------COVER(c, i, i);-------// 7a ----d = x - 2447 * i / 80;------// eb
------bool found = false;-----// 7f ----x = j / 11;------// b7
-------for (node *r = c->d; !found && r != c; r = r->d) {---------// 88 ----m = \tilde{j} + 2 - 12 * x;-----------// 82
-----sol[k] = r->row;------// ef ----y = 100 * (n - 49) + i + x;------// 70
-----found = search(k + 1);------// f1
-----for (node *j = r > 1; j != r; j = j > 1) { UNCOVER(j > p, a, b); c > 1.
-----UNCOVER(c, i, j);-----// 3a
-----return found:-----// 80
};-----// d9
7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \ldots, k-1\}
1}.
vector<int> nth_permutation(int cnt, int n) {------// 78
----vector<int> idx(cnt), per(cnt), fac(cnt);------// 9e
----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
----for (int i = 1; i \le cnt; i++) fac[i - 1] = n % i, n /= i;-----// 04
----for (int i = cnt - 1; i >= 0; i--)-----// 52
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);------// 41
----return per:-----// 84
```

```
ii find_cycle(int x0, int (*f)(int)) {------// a5
----int t = f(x0), h = f(t), mu = 0, lam = 1;------// 8d
----while (t != h) t = f(t), h = f(f(h));-----// 79
----h = x0;-----// 04
----while (t != h) t = f(t), h = f(h), mu++;-----// 9d
----h = f(t);-----// 00
```

----**while** (t != h) h = f(h), lam++;-----// 5e

----return ii(mu, lam);-----// b4

# 7.8. **Dates.** Functions to simplify date calculations.

```
int intToDay(int jd) { return jd % 7; }-----// 89
int dateToInt(int y, int m, int d) {-----// 96
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1
```

## 8. Useful Information

### 8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo  $m_1, m_2, \ldots, m_k$ , where  $m_1, m_2, \ldots, m_k$  are pairwise coprime integers, and find the real answer using CRT?

- Are there any edge cases? When  $n = 0, n = -1, n = 1, n = 2^{31} 1$  or  $n = -2^{31}$ ? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. **Fast Input Reading.** If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

8.3. **128-bit Integer.** GCC has a 128-bit integer data type named \_\_int128. Useful if doing multiplication of 64-bit integers, or something needing a little more than 64-bits to represent.

## 8.4. Worst Time Complexity.

n	Worst AC Algorithm	Comment
$\leq 10$	$O(n!), O(n^6)$	e.g. Enumerating a permutation
$\leq 15$	$O(2^n \times n^2)$	e.g. DP TSP
$\leq 20$	$O(2^{n}), O(n^{5})$	e.g. $DP + bitmask technique$
$\leq 50$	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
$\leq 10^{5}$	$O(n \log_2 n)$	e.g. Merge sort, building a Segment tree
$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

### 8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.