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```
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----void update(int at, int by) {------------------------// 76 ------matrix<T> res(cols, rows);------------------// b5
-------while (at < n) data[at] += by, at |= at + 1; }-------// fb -------for (int i = 0; i < rows; i++)-------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
------while (at >= 0) res += data[at], at = (at & (at + 1)) - 1;-------// 37 ----matrix<T> pow(int p) {---------// 68
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);-------// 4d
};------// 57 -----while (p) {------// cb
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y-------// 17 ----matrix<T> rref(T &det) {-------// 89
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
------for (int i = 0; i < cols; i++)-----// ab
2.4. Matrix. A Matrix class.
                       -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
public:-----T m = mat(i, c);------// e8
-----data.assign(cnt, T(0)); }-----// d0
                      ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// fe -----} return mat; }-----
-----cnt(other.cnt), data(other.data) { }-----// ed
                       private:----// e0
----T& operator()(int i, int j) { return at(i, j); }------// e0
                       ----int cnt:-----// 6a
----void operator +=(const matrix& other) {------// c9
                       ----vector<T> data;------// 41
------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
                       ----inline T& at(int i, int j) { return data[i * cols + j]; }------// 74
----void operator -=(const matrix& other) {------// 68
                       }:-----// b8
------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
----void operator *=(T other) {------// ba
                       2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
------for (int i = 0; i < cnt; i++) data[i] *= other; }------// 40 #define AVL_MULTISET 0------// b5
------matrix<T> res(*this); res += other; return res; }------// 5d template <class T>-------// 22
------matrix<T> res(*this); res -= other; return res; }------// cf public:------// f6
------matrix<T> res(rows, other.cols);-------// 57 ------node(const T &item, node *p = NULL) : item(item), p(p),------// c5
-----for (int k = 0; k < cols; k++)------// fc ----avl_tree() : root(NULL) { }----------// dc
```

```
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-----node *cur = root;-------// cf ------if (!cur) cur = root;------// 0a
------if (cur->item < item) cur = cur->r;-------// eb ------if (n < sz(cur->l)) cur = cur->l;------// 25
------else if (item < cur->item) cur = cur->l;------// de ------else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-----// 8b
------else break: }-------// 05 ------else break:-------// 4c
-------node *prev = NULL, **cur = &root;----------// 60 ----inline int sz(node *n) const { return n ? n->size : 0; }------// 69
-----prev = *cur;------// 31 ----inline bool left_heavy(node *n) const {-------// d7
-------if ((*cur)->item < item) cur = &((*cur)->r);------// 39 -----return n && height(n->l) > height(n->r); }-----// 9d
#if AVL_MULTISET-----// d1 ----inline bool right_heavy(node *n) const {-------// 91
------else cur = \&((*cur)->l); ------// 3b -----return n \&\& height(n->r) > height(n->l); }-----// 77
#else-----// dc ----inline bool too_heavy(node *n) const {--------// 18
------else if (item < (*cur)->item) cur = &((*cur)->l);------// e5 -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// fb
------else return *cur;------// 19 ----void delete_tree(node *n) {-------// 48
#endif-------if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }------// ea
------node *n = new node(item, prev);-------// 5b -----if (!n->p) return root;-------// af
-----*cur = n, fix(n); return n; }-------// 86 ------if (n->p->l == n) return n->p->l;-------// 95
----void erase(node *n, bool free = true) {-------// 89 ------assert(false); }------
------else if (n->l) && !n->r) parent_leg(n) = n->l, n->l->p = n->p; ------// 3d ------n->size = 1 + sz(n->l) + sz(n->r); -------// 14
-----node *s = successor(n);------// 16 ----#define rotate(l, r) N------// b7
------erase(s, false);-------// 17 -----node *l = n->l; \[\bar{N}\]
-----if (n->l) n->l->p = s;------// 88
                             -----parent_leg(n) = l; \(\bar{\gamma}\)------// 02
-----if (n->r) n->r->p = s;------// 42
                             -----n->l = l->r; \\ \[ \] ------// 08
-----parent_leg(n) = s, fix(s);-----// 87
------} else parent_leg(n) = NULL;---------// 58 ------l->r = n, n->p = l; \[ \bar{N} \]
------if (free) delete n; }------// 99 ----void left_rotate(node *n) { rotate(r, l); }------// 43
-----if (!n) return NULL;------// b3 ----void fix(node *n) {-------// 42
------if (too_heavy(n)) {------// 39
------return p: }--------| 03 ------------| 65e if (right_heavy(n) && left_heavy(n->r))-------// 09
-----node *p = n->p;------node *p = n->p; }-------// af
------while (p && p->l == n) n = p, p = p->p;-------// 03 -----n = n->p; } } };-------// 85
-----return p; }-----// 83
                              Also a very simple wrapper over the AVL tree that implements a map interface.
----inline int size() const { return sz(root); }-----// e2
                             #include "avl_tree.cpp"-----// 01
----void clear() { delete_tree(root), root = NULL; }-----// d4
                              -----// ba
----node* nth(int n, node *cur = NULL) const {------// f4
                             template <class K, class V>-----// da
```

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class avl_map {------// 3f ------memset(newloc + len, 255, (newlen - len) << 2);------// 18
public:-----delete[] q, delete[] loc;------// 74
---avl_tree<node> tree;------// b1 -----}-----// b1 -----}
------typename avl_tree<node>::node *n = tree.find(node(key, V(0)));------// ba ------loc[n] = count, q[count++] = n;------// 6b
------if (!n) n = tree.insert(node(key, V(0)));-------// cb -----if (fix) swim(count-1); }------// bf
----}-----assert(count > 0);-------// eb
------if (fix) sink(0);------// 80
2.6. Heap. An implementation of a binary heap.
                                #define RESIZE-----// d0
                               ----int top() { assert(count > 0); return q[0]; }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                               ----void heapify() { for (int i = count - 1; i > 0; i--)------// 39
struct default_int_cmp {------// 8d
                               -----if (cmp(i, (i-1)/2)) swp(i, (i-1)/2); }------// \theta b
----default_int_cmp() { }------// 35
                               ----void update_key(int n) {------// 26
----bool operator ()(const int &a, const int &b) { return a < b; } };-----// e9
                               -----assert(loc[n] != -1), swim(loc[n]), sink(loc[n]); }------// 7d
template <class Compare = default_int_cmp>------// 30
                               ----bool empty() { return count == 0; }------// f8
class heap {-----// 05
                               ----int size() { return count; }------// 86
private:----// 39
                               ----void clear() { count = 0, memset(loc, 255, len << 2); } };-----// 58
----int len, count, *q, *loc, tmp;-----// 0a
----Compare _cmp;-----// 98
                               2.7. Skiplist. An implementation of a skiplist.
----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }------// a0
----inline void swp(int i, int j) {-------// 1c #define BP 0.20------// aa
------SWP(q[i], q[j]), SWP(loc[q[i]], loc[q[j]]); }------// 67 #define MAX_LEVEL 10--------------------------// 56
----void swim(int i) {------// 33 unsigned int bernoulli(unsigned int MAX) {------// 7b
------while (i > 0) {-------// 1a ----unsigned int cnt = 0;------// 28
------int p = (i - 1) / 2;------// 77 ----while(((float) rand() / RAND_MAX) < BP && cnt < MAX) cnt++;-----// d1
----void sink(int i) {------// ce ----struct node {------// 53
------int l = 2*i + 1, r = l + 1;------// b4 ------int *lens;------// 07
-----<mark>int</mark> m = r >= count || cmp(l, r) ? l : r;------// cc ------#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))------// 25
public:-----// cd ----int current_level, _size;------// 61
-----q = new int[len], loc = new int[len];-------// f8 ----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
------memset(loc, 255, len << 2); }------// f7 -----skiplist() { clear(); delete head; head = NULL; }-----// aa
----~heap() { delete[] q; delete[] loc; }-------// 09 ----#define FIND_UPDATE(cmp, target) \[ \sqrt{\colored}\]-------// c3
-----if (len == count || n >= len) {------// 0f
                               -----memset(pos, 0, sizeof(pos)); \[\sigma_-----// f2\]
#ifdef RESIZE-----// a9
                                -----node *x = head; \[\sqrt{------// 0f}\]
------int newlen = 2 * len;-----// 22
                                -----node *update[MAX_LEVEL + 1]; \[\scrip_------// 01
------while (n >= newlen) newlen *= 2;-----// 2f
                               -----// 38
------int *newq = new int[newlen], *newloc = new int[newlen];------// e3
```

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-----pos[i] = pos[i + 1]; \(\bar{\sqrt{0}}\)-------// 68 -------if (r) r->l = this;------// θb
-----}------// 61
                                        -----pos[i] += x->lens[i]; x = x->next[i]; } \[ \] \[ \] \]
                                        ----node *front, *back;-----// 23
-----update[i] = x; \\ -----// dd
                                        ----dancing_links() { front = back = NULL; }------// 8c
----node *push_back(const T &item) {------// d7
----int size() const { return _size; }------// 9a
                                        ------back = new node(item, back, NULL);------// 5d
----void clear() { while(head->next && head->next[0])-----// 91
                                        -----if (!front) front = back:-----// a2
-----erase(head->next[0]->item); }------// e6
                                        -----return back;-----// b4
----node *find(T target) { FIND_UPDATE(x->next[i]->item, target);------// 36
                                        -----return x && x->item == target ? x : NULL; }-----// 50
                                        ----node *nth(int k) { FIND_UPDATE(pos[i] + x->lens[i], k+1); return x; }-----// b8
                                        ------front = new node(item, NULL, front);------// 75
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80
                                        -----if (!back) back = front;-----// d6
-----return pos[0]; }-----// 19
                                        -----return front;-----// ef
----node* insert(T target) {------// 80
                                        ------FIND_UPDATE(x->next[i]->item, target);------// 3a
                                        ----void erase(node *n) {------// 88
------if(x && x->item == target) return x; // SET------// 07
                                        ------if (!n->l) front = n->r; else n->l->r = n->r; ------// d5
------int lvl = bernoulli(MAX_LEVEL);------// 7a
                                        ------if (!n->r) back = n->l; else n->r->l = n->l;-------// 96
------if(lvl > current_level) current_level = lvl;------// 8a
                                        ----}-------------------------// ae
----x = new node(lvl, target);-----// 36
                                        ----void restore(node *n) {-------// 6d
-----for(int i = 0; i <= lvl; i++) {------// 49
                                        -----if (!n->l) front = n; else n->l->r = n;------// ab
-----x->next[i] = update[i]->next[i];------// 46
                                        ------if (!n->r) back = n; else n->r->l = n;-------------// 8d
-----x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];-----// bc
                                        -----update[i] ->next[i] = x;-----// 20
                                         -----// 4f
-----update[i]->lens[i] = pos[0] + 1 - pos[i];-----// 42
3. Graphs
-----for(int i = lvl + 1: i <= MAX_LEVEL: i++) update[i]->lens[i]++:-----// 07
                                        3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
-----size++;-----// 19
                                        edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
-----return x; }-----// c9
----void erase(T target) {------// 4d
                                        graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
                                        connected. It runs in O(|V| + |E|) time.
------FIND_UPDATE(x->next[i]->item, target);------// 6b
-----if(x && x->item == target) {------// 76
                                        int bfs(int start, int end, vvi& adj_list) {------// d7
                                        ----queue<ii>> Q;------// 75
-----for(int i = 0; i <= current_level; i++) {------// 97
-----update[i]->next[i] = x->next[i];-----// 59 -----// 59
-----current_level--; } } ;-----// 59
                                        -----vi& adj = adj_list[cur.first];-----// 3f
                                        ------for (vi::iterator it = adj.begin(); it != adj.end(); it++)-----// bb
2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
                                        -----Q.push(ii(*it, cur.second + 1));------// b7
list supporting deletion and restoration of elements.
                                        template <class T>-----// 82
                                        }-----// 7d
struct dancing_links {-----// 9e
----struct node {------// 62
                                         Another implementation that doesn't assume the two vertices are connected. If there is no path
                                        from the starting vertex to the ending vertex, a-1 is returned.
-----T item:-----// dd
-----node *l, *r:-----// 32
                                        int bfs(int start, int end, vvi& adj_list) {------// d7
-----node(const T &item, node *l = NULL, node *r = NULL)-----// 88
                                        ----set<<mark>int</mark>> visited;-----// b3
----: item(item), l(l), r(r) {------// 04
                                        ----queue<ii>> 0;------// bb
```

-----if (l) l->r = this;------// 1c ----Q.push(ii(start, 0));------// 3a

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-----// db ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
------if (cur.first == end)-------// 22 ------for (int k = 0; k < size(adj[j]); k++)------// 3f
-----return cur.second;------// b9 ------dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61
-----/<sub>ba</sub> ------dist[j] + adj[j][k].second);-------// 47
-----vi& adj = adj_list[cur.first];-------// f9 ----for (int j = 0; j < n; j++)-------// 13
------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)-------// 44 ------for (int k = 0; k < size(adj[j]); k++)---------// a0
-----if (visited.find(*it) == visited.end()) {-------// 8d -----if (dist[j] + adj[j][k].second < dist[adj[j][k].first])-----// ef
-----visited.insert(*it);-------// cb ---return dist;-----
----}--------// 0b
                                   3.3. All-Pairs Shortest Paths.
-----// 63
----return -1:-----// f5
                                  3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
}-----// 03
                                  problem in O(|V|^3) time.
                                   void floyd_warshall(int** arr, int n) {------// 21
3.2. Single-Source Shortest Paths.
                                   ----for (int k = 0; k < n; k++)-----// 49
                                   ------for (int i = 0; i < n; i++)------// 21
3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|)
                                   -----for (int j = 0; j < n; j++)-----// 77
time.
                                   -----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
int *dist. *dad:-----// 46
                                   -----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);-----// e1
struct cmp {-----// a5
                                  }-----// 86
----bool operator()(int a, int b) {------// bb
-----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
                                  3.4. Strongly Connected Components.
};-----// 41
pair<int*, int*> dijkstra(int n, int s, vii *adj) {-----// 53
                                  3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
----dist = new int[n];-----// 84
                                  graph in O(|V| + |E|) time.
----dad = new int[n];-----// 05
                                  #include "../data-structures/union_find.cpp"-----// 5e
----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;------// d6
                                   -----// 11
                                  vector<br/>bool> visited;------// 66
----set<int. cmp> pg:-----// 04
------int cur = *pq.beqin(); pq.erase(pq.beqin());--------// 7d void scc_dfs(const vvi &adj, int u) {-----------------------------// a1
------int nxt = adj[cur][i].first,-------// b8 ----for (int i = 0; i < size(adj[u]); i++)------// c5
------ndist = dist[cur] + adj[cur][i].second;-------// 0c -----if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-------// 6e
}-----// af ----order.clear();-------// 22
                                   ----union_find uf(n);-----// 6d
3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
                                   ----vi dag;------// ae
problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
                                   ----vvi rev(n):-----// 20
negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
                                   ----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++)------// b9
                                   -----rev[adj[i]]]].push_back(i);-----// 77
int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {-------// cf ----visited.resize(n), fill(visited.begin(), visited.end(), false);-------// 04
```

```
------for (int i = 0; i < size(adi[u]); i++)-------// 90 -----if (!color[i]) {---------------------------// d5
------if (!visited[v = adj[u][i]]) S.push(v);-------// 43 -----tsort_dfs(i, color, adj, S, has_cycle);-------// 40
}-----// 97 ----while (!S.empty()) res.push_back(S.top()), S.pop();------// 94
                                   ----return res:------// 07
3.5. Minimum Spanning Tree.
                                   }-----// 1f
3.5.1. Kruskal's algorithm.
                                   3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
#include "../data-structures/union_find.cpp"-----------------------// 5e
                                   #define MAXV 1000-----// 2f
                                   #define MAXE 5000-----// 87
// n is the number of vertices-----// 18
                                   vi adj[MAXV];-----// ff
// edges is a list of edges of the form (weight, (a, b))-----// c6
                                   // the edges in the minimum spanning tree are returned on the same form-----// 4d
                                   ii start_end() {------// 30
vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
                                   ----int start = -1, end = -1, any = 0, c = 0;------// 74
----union_find uf(n);-----// 04
                                   ----for (int i = 0; i < n; i++) {------// 96
----sort(edges.begin(), edges.end());-----// 51
                                   -----if (outdeg[i] > 0) any = i;-----// f2
-----if (indeg[i] + 1 == outdeg[i]) start = i, c++;------// 98
----for (int i = 0; i < size(edges); i++)-----// ce
                                   ------else if (indeg[i] == outdeg[i] + 1) end = i, c++;------// 4f
-----if (uf.find(edges[i].second.first) !=-----// d5
                                   ------else if (indeg[i] != outdeg[i]) return ii(-1,-1);------// fa
------uf.find(edges[i].second.second)) {------// 8c
                                   ----}------// ef
-----res.push_back(edges[i]);-----// d1
                                   ----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 6e
-----uf.unite(edges[i].second.first, edges[i].second.second);-----// a2
                                   ----if (start == -1) start = end = any;-----// db
-----}-----// 5b
                                   ----return ii(start, end);-----// 9e
----return res;------// 46
                                   }-----// 35
}-----// 88
                                   bool euler_path() {-----// d7
                                   ----ii se = start_end();-----// 45
3.6. Topological Sort.
                                   ----int cur = se.first, at = m + 1;------// 8c
3.6.1. Modified Depth-First Search.
                                   ----if (cur == -1) return false;------// 45
void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
                                   ----stack<int> s;-----// f6
------bool& has_cycle) {------// a8
                                   ----while (true) {------// 04
----color[cur] = 1;------// 5b
                                   -----if (outdeg[cur] == 0) {------// 32
----for (int i = 0; i < size(adj[cur]); i++) {-------// 96
                                   -----res[--at] = cur:-----// a6
------int nxt = adj[cur][i];------// 53
                                   ------if (s.empty()) break;-----// ee
-----if (color[nxt] == 0)------// 00
                                   -----cur = s.top(); s.pop();-----// d7
-----tsort_dfs(nxt, color, adj, res, has_cycle);------// 5b
                                   -----} else s.push(cur), cur = adj[cur][--outdeg[cur]];-----// d8
-----else if (color[nxt] == 1)------// 53
                                   ----}------// ba
-----has_cycle = true;-----// c8
                                   ----return at == 0:-----// c8
-----if (has_cycle) return;-----// 7e
                                   l-----// aa
----}--------// 3d
----color[cur] = 2;-----// 16
                                   3.8. Bipartite Matching.
----res.push(cur):-----// cb
}-----// 9e
                                  3.8.1. Bipartite Matching. The alternating paths algorithm solves bipartite matching in O(mn^2) time,
------// ae where m, n are the number of vertices on the left and right side of the bipartite graph, respectively.
```

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vi* adj;------memset(L, -1, sizeof(int) * N);-------// 16
bool* done:-----memset(R, -1, sizeof(int) * M):-------// e4
----done[left] = true;-------// 86
------int right = adj[left][i];------// b6
                   3.10. Maximum Flow.
------if (owner[right] == -1 || alternating_path(owner[right])) {------// d2
-----owner[right] = left; return 1;-----// 26
                   3.10.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It com-
------} }------// 7a
                   putes the maximum flow of a flow network.
----return 0: }-----// 83
                   #define MAXV 2000-----// ba
3.9. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm.
                   int q[MAXV], d[MAXV];-----// e6
struct bipartite_graph {------// 2b -----edge() { }-----// 38
------L(new int[N]), R(new int[M]), adj(new vi[N]) {}----------// 46 ----int n, ecnt, *head, *curh;------------------------// 77
----~bipartite_graph() { delete[] adj; delete[] L; delete[] R; }-------// b9 ----vector<edge> e, e_store;--------------------------// d0
----bool bfs() {-------// 3e ----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 80
------int l = 0, r = 0; -------// a4 ------e.reserve(2 * (m == -1 ? n : m)); -------// 5d
------else dist(v) = INF;--------// c4 -----memset(head, -1, n * sizeof(int));-------// f6
------while(l < r) {------// 3f ----void destroy() { delete[] head; delete[] curh; }------// 21
------int v = q[l++];------// 69 ----void reset() { e = e_store; }------// 60
------if(dist(v) < dist(-1)) {--------// b2 ----void add_edge(int u, int v, int uv, int vu = 0) {------// dd
-----foreach(u, adj[v]) if(dist(R[*u]) == INF)------// 63 -----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// 7a
------}-----memset(d, -1, n * sizeof(int));--------// 66
----void add_edge(int i, int j) { adj[i].push_back(j); }------// 87 ------for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)-----// 6d
```

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-----memcpv(curh, head, n * sizeof(int)):-------// b6 ---int u, v, w, c:-------
-----if (res) reset();-------// 08 ------u = _u; v = _v; w = _rev;------// b2
}:-----// cf -----// 31
                      ii min_cost_max_flow(int n, int s, int t, vector<pair<int, ii> >* adj) {------// 4d
3.10.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
                     ----vector<mcmf_edqe*>* q = new vector<mcmf_edqe*>[n];------// θε
O(|V||E|^2). It computes the maximum flow of a flow network.
                      ----for (int i = 0; i < n; i++) {------// a7
----int u, v, w; mf_edge* rev;-------------// ab ------mcmf_edge *cur = new mcmf_edge(i, adj[i][j].first,------// 28
----mf_edge(int _u, int _v, int _w, mf_edge* _rev = NULL) {-------// 96 ------adj[i][j].second.first, adj[i][j].second.second),-----// 71
-----ce = new mf_edge(i, adj[i][j].first, adj[i][j].second);-----// ed ----mcmf_edge** back = new mcmf_edge*[n];------// 90
-----g[i].push_back(ce);------// 09 ----int* dist = new int[n];------// 05
------g[ce->v].push_back(ce->rev); } }------// 58 -------for (int i = 0; i < n; i++) back[i] = NULL, dist[i] = INF;------// 41
------back.assign(n, NULL);---------// 4d ------for (int i = 0; i < n - 1; i++)---------// c3
------queue<int> 0; 0.push(s);-------// 18 -------for (int j = 0; j < n; j++)-------// 5e
------while (!Q.empty() && (cur = Q.front()) != t) {------// a7 --------if (dist[j] != INF)------// dd
------mf_edge* nxt = g[cur][i]:------// 86 ------dist[g[j][k]->v]) {-------// ec
------for (int i = 0; i < size(g[t]); i++) {-------// 1e -----mcmf_edge* cure = back[t];------// f8
------if (cap == 0) continue;------// 92 -----cap = min(cap, cure->w);-----// ff
-----assert(cap < INF);--------// fb -------if (cure->u == s) break;-------// ce
-----z->w -= cap, z->rev->w += cap;------// 67 -----cure = back[cure->u];-----// c6
-----ce->w -= cap, ce->rev->w += cap;-------// 9c -----assert(cap > 0 && cap < INF);-------// 72
----return make_pair(flow, q); }-----------------------// f8 -------while (true) {----------------------------------// c9
                      ------cost += cap * cure->c;------// e4
3.11. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
```

minimum cost.

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                                         11
-----cure = back[cure->u];-------// 03 int compute_max_flow(int s, int t, const pair<vii, vvi> &gh) {------// 16
----// instead of deleting q, we could also-------// 5d ------cur = min(cur, qh.first[at].second), at = qh.first[at].first;------// bd
----// use it to get info about the actual flow-------// 5a ----return min(cur, gh.second[at][t]);------// 6d
-----for (int j = 0; j < size(q[i]); j++)-----// 4b
-----delete q[i][j];-----// bb
                              4. Strings
----delete[] q;------// 37
                     4.1. Trie. A Trie class.
----delete[] back;-----// 42
                     template <class T>-----// 82
----delete[] dist;------// 28
                     class trie {-----// 9a
----return ii(flow, cost);------// 32
                     private:----// f4
}-----// 16
                     ----struct node {------// ae
                     -----map<T. node*> children:-----// a0
3.12. All Pairs Maximum Flow.
                     ------int prefixes, words;------// e2
3.12.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
                     -----node() { prefixes = words = 0; } };------// 42
structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the max-
                     public:-----// 88
imum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
                     ----node* root;-----// a9
#include "dinic.cpp"-----// 58
                     ----trie() : root(new node()) { }-----// 8f
-----// 25 ----template <class I>-------// 89
------int l = 0, r = 0;-------// 3e
------memset(d, 0, n * sizeof(int));-------// 79 ------typename map<T, node*>::const_iterator it;------// 01
------memset(same, 0, n * sizeof(int));--------// b0 -----it = cur->children.find(head);-------// 77
------while (l < r) {------// 45 -------pair<T, node*> nw(head, new node());------// cd
-----same[v = g[l++]] = true;------// c8 ------it = cur->children.insert(nw).first;-----// ae
----}------T head = *begin;-------// 5c
-----cap[cur][i] = mn;------// 63 ------begin++, cur = it->second; } } }------// 7c
-----mn = min(mn, par[cur].second), cur = par[cur].first;-------// 28 ----int countPrefixes(I begin, I end) {---------------------------// 85
```

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-----T head = *begin;--------// 43 ------foreach(c, *k)-----------------------// 38
-----typename map<T. node*>::const_iterator it:------// 7a ------cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d
-----it = cur->children.find(head);-------// 43 ------(cur->next[*c] = new go_node());------// 75
------if (it == cur->children.end()) return 0;-------// 71 -----cur->out = new out_node(*k, cur->out);------// 6e
-----begin++, cur = it->second; } } } ;------// 26 -----}-----------------------// 96
                               -----queue<go_node*> q;------// 8a
4.2. Suffix Array. An O(n \log n) construction of a Suffix Tree.
                               ------foreach(a, qo->next) q.push(a->second);------// a3
struct entry { ii nr; int p; };------// f9 ------while (!q.empty()) {------// 43
bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77 ------go_node *r = q.front(); q.pop();------// 2e
struct suffix_array {-------// 87 ------foreach(a, r->next) {------// 25
----string s; int n; vvi P; vector<entry> L; vi idx;------// b6 ------go_node *s = a->second;------// cb
----suffix_array(string s) : s(s), n(size(s)) {-------// 26 -----q.push(s);------------------------// 76
-----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);------// ca ------go_node *st = r->fail;------// fa
------P.push_back(vi(n));------// de ------if (!st) st = go;-----// e7
-----i + cnt < n ? P[stp - 1][i + cnt] : -1);------// e2 ------if (!s->out) s->out = s->fail->out;------// 80
-----sort(L.beqin(), L.end());------// ed
------for (int i = 0; i < n; i++)------// 34 -------out_node* out = s->out;-----// bf
------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;------// 57 ------out->next = s->fail->out;------// 65
----int lcp(int x, int y) {--------// e8
}:------cur = cur->fail;------// 9e
                               ------if (!cur) cur = qo;-----// 2f
4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
                               -----cur = cur->next[*c];------// 58
state machine from a set of keywords which can be used to search a string for any of the keywords.
                               -----if (!cur) cur = go;-----// 3f
struct aho_corasick {------// 78
                               -----for (out_node *out = cur->out; out = out->next)-----// e0
----struct out_node {------// 3e
                               -----/res.push_back(out->keyword);------// 0d
-----string keyword; out_node *next;------// f0
                               -----}-----// fd
-----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
                               return res:----// c1
----};-------// b9
                               ----struct qo_node {------// 40
                               }:-----// 32
-----map<char, qo_node*> next;------// 6b
-----out_node *out; go_node *fail;-----// 3e
                               4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
-----go_node() { out = NULL; fail = NULL; }-----// Of
                               also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
----};------// c0
                               can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
                               accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
----ao_node *ao:-----// b8
-----go_node *cur = go;------// 9d ----int l = 0, r = 0;-------// 1c
```

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---z[0] = n;------// 98 ------return !(*this == other); }------// d1
-z[i] = 0: -c
                          5.2. Big Integer. A big integer class.
------if (i > r) {-------// 26
                          struct intx {-----// cf
-----l = r = i:-----// a7
                          ----intx() { normalize(1); }------// 6c
----intx(string n) { init(n); }-------// b9
----intx(int n) { stringstream ss; ss << n; init(ss.str()); }------// 36
-----} else if (z[i - l] < r - i + 1) z[i] = z[i - l]:----// bf
                          ----intx(const intx& other) : sign(other.sign), data(other.data) { }------// 3b
-----else {------// b5
-----l = i;-----// 02
                          ----int sign;------// 26
                          ----vector<unsigned int> data;------// 19
------while (r < n \&\& s[r - l] == s[r]) r++;
                          ----static const int dcnt = 9;-----// 12
-----z[i] = r - l; r--; } }-----// 8d
                          ----static const unsigned int radix = 1000000000U;-----// f0
----return z;-----// 53
                          ----int size() const { return data.size(); }------// 29
}-----// db
                          ----void init(string n) {------// 13
                          -----intx res; res.data.clear();-----// 4e
          5. Mathematics
                          -----if (n.empty()) n = "0";------// 99
5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common
                          -----if (n[0] == '-') res.sign = -1, n = n.substr(1);------// 3b
                          ------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {-------// e7
public:------digit = digit * 10 + (n[idx] - '0');-------------------------// 1f
-----assert(d_ != 0);-----// 3d -----}----// fb
-----T q = gcd(abs(n), abs(d));--------// fc ---}------// fc
------n /= g, d /= g; }-------// al ----intx& normalize(int nsign) {-------// 3b
----fraction(T n_) : n(n_), d(1) { }-------// 84 ------if (data.empty()) data.push_back(θ);-------// fa
----fraction(const fraction<T>& other) : n(other.n), d(other.d) { }-------// 01 -------for (int i = data.size() - 1; i > 0 && data[i] == 0; i--)-------// 27
------return fraction<T>(n * other.d + other.n * d, d * other.d);}-------// 3b ------sign = data.size() == 1 \& \& data[0] == 0 ? 1 : nsign;-------// ff
------return fraction<T>(n * other.d - other.n * d, d * other.d);}------// 47 ----}
----fraction<T> operator /(const fraction<T>& other) const {-------// ca ------bool first = true;-----------------// 33
------for (int i = n.size() - 1; i >= 0; i--) {-------// 63
-----return n * other.d < other.n * d; }------// 8c -----else {--------|
------stringstream ss; ss << cur;-------// 8c
------return other < *this; }-------// 6e -------int len = s.size();-------// 0d
------return n == other.n && d == other.d: }------// 14 -----}
----bool operator !=(const fraction<T>& other) const {-------// ec -----return outs;-----
```

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------if (sign != b.sign) return sign < b.sign;--------// cf ------assert(!(d.size() == 1 && d.data[0] == 0));-------// e9
------if (size() != b.size())--------// 4d ------intx q, r; q.data.assign(n.size(), θ);-------// ca
-----return sign == 1 ? size() < b.size() : size() > b.size();------// 4d ------for (int i = n.size() - 1; i >= 0; i--) {-------// 1a
------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35 ------r.data.insert(r.data.begin(), 0);--------// c7
------return false;-------// ca -------long long k = θ;--------// cc
------if (sign > 0 && b.sign < 0) return *this - (-b);---------// 36 ------r = r - abs(d) * k;-----------------// 15
-----intx c; c.data.clear();-------// 18 -----}------// 2f
------unsigned long long carry = 0;-------// 5c ------return pair<intx, intx>(q.normalize(n.sign * d.sign), r);------// a1
-----carry += (i < size() ? data[i] : 0ULL) +------// 91 ----intx operator /(const intx& d) const {-------// a2
-----c.data.push_back(carry % intx::radix);-------// 86 ----intx operator %(const intx& d) const {--------// 07
-----carry /= intx::radix;------// fd -----return divmod(*this,d).second * sign; }-----// 5a
-----return c.normalize(sign);------// 20
                                     5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
#include "intx.cpp"-----// 83
----intx operator -(const intx& b) const {--------// 53
                                     #include "fft.cpp"-----// 13
------if (sign > 0 && b.sign < 0) return *this + (-b);------// 8f
                                      -----// e0
-----if (sign < 0 && b.sign > 0) return -(-*this + b);------// 1b
                                     intx fastmul(const intx &an, const intx &bn) {-----// ab
------if (sign < 0 && b.sign < 0) return (-b) - (-*this);------// a1
                                      ----string as = an.to_string(), bs = bn.to_string();-----// 32
-----if (*this < b) return -(b - *this);------// 36
-----intx c; c.data.clear();------// 6b
                                      ----int n = size(as), m = size(bs), l = 1,------// dc
                                      -----len = 5, radix = 100000,-----// 4f
-----long long borrow = 0;-----// f8
------for (int i = 0; i < size(); i++) {------// a7
                                      -----*a = new int[n], alen = 0,-----// b8
                                      -----*b = new int[m], blen = 0;------// 0a
------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);-----// a9
                                      ----memset(a, 0, n << 2);-----// 1d
-----c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
------borrow = borrow < 0 ? 1 : 0;-----// 0d
                                     ----memset(b, 0, m << 2);-----// 01
                                      ----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
------for (int j = min(len - 1, i); j >= 0; j--)-----// 43
-----return c.normalize(sign);------// 35
                                      -----a[alen] = a[alen] * 10 + as[i - j] - '0';------// 14
----}-------// 85
----intx operator *(const intx& b) const {------// bd
                                      ----for (int i = m - 1; i >= 0; i -= len, blen++)------// b6
-----intx c; c.data.assign(size() + b.size() + 1, 0);-----// d0
                                      ------for (int j = min(len - 1, i); j >= 0; j--)------// ae
-----for (int i = 0; i < size(); i++) {------// 7a
                                      -----b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
                                      ----while (l < 2*max(alen,blen)) l <<= 1;------// 51
-----long long carry = 0:-----// 20
                                     ----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
------for (int j = 0; j < b.size() || carry; j++) {------// c0
-----if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
                                      ----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);------// 35
                                      ----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? <math>b[i] : 0, 0);-----// 66
-----carry += c.data[i + j];-----// 18
                                      ----fft(A, l); fft(B, l);-----// f9
-----.c.data[i + j] = carry % intx::radix;------// 86
                                     ----for (int i = 0; i < l; i++) A[i] *= B[i];------// e7
-----carry /= intx::radix;-----// 05
                                     ----fft(A, l, true);------// d3
----ull *data = new ull[l];-----// e7
-----}-----// 9e
                                      ----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
```

```
----for (int i = 0; i < l - 1; i++)-----// 90
                                             5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
-----if (data[i] >= (unsigned int)(radix)) {-------// 44
                                             vi prime_sieve(int n) {-----// 40
-----data[i+1] += data[i] / radix;-----// e4
                                             -----data[i] %= radix;-----// bd
                                             ----vi primes:-----// 8f
                                             ----bool* prime = new bool[mx + 1];------// ef
----int stop = l-1;-----// cb
                                             ----memset(prime, 1, mx + 1);------// 28
----while (stop > 0 && data[stop] == 0) stop--:-----// 97
                                             ----if (n >= 2) primes.push_back(2);-----// f4
----stringstream ss:-----// 42
                                             ----while (++i <= mx) if (prime[i]) {------// 73
----ss << data[stop];------// 96
                                             -----primes.push_back(v = (i << 1) + 3);-----// be
----for (int i = stop - 1; i >= 0; i--)-----// bd
                                             -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
-----ss << setfil('0') << setw(len) << data[i];------// b6
                                             ------for (int i = sq: i <= mx: i += v) prime[i] = false: }------// 2e
----delete[] A; delete[] B;-----// f7
                                             ----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3);------// 29
----delete[] a: delete[] b:-----// 7e
                                             ----delete[] prime; // can be used for O(1) lookup------// 36
----delete[] data;------// 6a
                                             ----return primes; }-----// 72
----return intx(ss.str()):-----// 38
                                             5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
                                             5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
                                               -----// e8
k items out of a total of n items.
                                             int mod_inv(int a, int m) {------// 49
int nck(int n, int k) {------// f6
                                             ----int x, y, d = eqcd(a, m, x, y);------// 3e
----if (n - k < k) k = n - k;------// 18
                                             ----if (d != 1) return -1;------// 20
----int res = 1:------// cb
                                             ----return x < 0 ? x + m : x:-----// 3c
----for (int i = 1; i \le k; i++) res = res * (n - (k - i)) / i;------// bd
}-----// 03
                                             5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
                                             template <class T>-----// 82
5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
                                             T mod_pow(T b, T e, T m) {-----// aa
integers a, b.
                                             ----T res = T(1);-----// 85
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
                                             ----while (e) {------// b7
 The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
                                             -----if (e & T(1)) res = mod(res * b, m);------// 41
and also finds two integers x, y such that a \times x + b \times y = d.
                                             -----b = mod(b * b, m), e >>= T(1); }-----// b3
                                             ----return res:-----// eb
int eqcd(int a, int b, int& x, int& y) {------// 85
----if (b == 0) { x = 1; y = 0; return a; }------// 7b
                                             }-----// c5
----else {------// 06
-----int d = egcd(b, a % b, x, y);-----// 34
                                             5.9. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
-----x -= a / b * y;-----// 4a
                                             #include "egcd.cpp"-----// 55
-----Swap(x, y):-----// 26
                                             int crt(const vi& as, const vi& ns) {-----// c3
                                             ----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
-----return d:-----// db
                                             ----for (int i = 0; i < cnt; i++) N *= ns[i]; ------// 88
----}------------// 9e
                                             ----for (int i = 0; i < cnt; i++)-----// f9
                                             ------egcd(ns[i], l = N/ns[i], r, s), x += as[i] * s * l;-------// b\theta
5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is
                                             ----return mod(x, N); }-----// 9e
prime.
----if (n < 2) return false;------// c9 n.
----if (n % 2 == 0 || n % 3 == 0) return false;------// Of vi linear_congruence(int a, int b, int n) {-------// c8
----if (n < 25) return true;------// ef ----int x, y, d = egcd(a, n, x, y);------// 7a
----int s = static_cast<int>(sqrt(static_cast<double>(n)));------// 64 ----vi res;-----
----for (int i = 5; i <= s; i += 6)------// 6c ----if (b % d != 0) return res;------// 30
----return true; }-----return true; }------// 43 ----for (int k = 0; k < d; k++) res.push_back(mod(x0 + k * n / d, n));------// 21
```

```
----return res;------// 03
}-----// 1c
```

5.11. Numeric Integration. Numeric integration using Simpson's rule.

5.12. **Fast Fourier Transform.** The Cooley-Tukey algorithm for quickly computing the discrete Fourier transform. Note that this implementation only handles powers of two, make sure to pad with zeros.

```
#include <complex>-----// 8e
typedef complex<long double> cpx;-----// 25
void fft(cpx *x, int n, bool inv=false) {------// 23
----for (int i = 0, j = 0; i < n; i++) {------// f2
------if (i < j) swap(x[i], x[j]);------// 5c
-----int m = n>>1:-----// e5
------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
-----j += m;-----// ab
----}-----// 1e
----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
------for (int m = 0; m < mx; m++, w *= wp) {------// 40
------for (int i = m; i < n; i += mx << 1) {------// 33
-----cpx t = x[i + mx] * w;-----// f5
-----x[i + mx] = x[i] - t;
-----x[i] += t:-----// c7
-----}-----// c2
----}------// 70
----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);-----// 3e
}-----// 7d
```

### 5.13. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once:  $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times:  $n^k$
- Number of permutations of n objects, where there are  $n_1$  objects of type 1,  $n_2$  objects of type 2, ...,  $n_k$  objects of type k:  $\binom{n}{n_1, n_2, \dots, n_k} = \frac{n!}{n_1! \times n_2! \times \dots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:
- $\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times:  $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to  $x_1 + x_2 + \cdots + x_n = k$  where  $x_i \ge 0$ :  $f_k^n$
- Number of subsets of a set with n elements:  $2^n$

- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an  $n \times m$  grid by walking only up and to the right:  $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced:  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an  $n \times n$  lattice which do not rise above the main diagonal:  $C_n$
- Number of permutations of n objects with exactly k ascending sequences or runs:  $\binom{n}{k} = \binom{n}{n-k-1} = k \binom{n-1}{k} + (n-k+1) \binom{n-1}{k-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{n-1} = \binom{n}{k} = \binom{n}{k$
- Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Heron's formula: A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} 1$ .
- **Divisor sigma:** The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where  $n = \prod_{i=0}^r p_i^{a_i}$  is the prime factorization.
- Divisor count: A special case of the above is  $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$  where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.

## 6. Geometry

# 6.1. **Primitives.** Geometry primitives.

#include <complex>-----// 8e #define P(p) const point &p-----// b8 #define L(p0, p1) P(p0), P(p1)-----// 30 typedef complex<double> point;------// e1 double dot(P(a), P(b)) { return real(conj(a) \* b); }-----// a9 double cross(P(a), P(b)) { return imag(conj(a) \* b); }-----// ff point rotate(P(p), P(about), double radians) {------// e1 ----**return** (p - about) \* exp(point(0, radians)) + about; }-----// *cb* point reflect(P(p), L(about1, about2)) {------// c0 ----point z = p - about1, w = about2 - about1;-----// 39 ----return conj(z / w) \* w + about1; }-----// 03 point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) \* u; }------// fc**bool** parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// 6d double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// ca **bool** collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }-----// 75 bool collinear(L(a, b), L(p, q)) {------// 66 ----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6 double angle(P(a), P(b), P(c)) {------// d0

```
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----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------// cc ----for (int i = 0, j = n - 1; i < n; j = i++)-----// 77
------if (CHK(real, p[i], q, p[j]) || CHK(real, p[i], q, p[i]) ------// 1f
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) {------// d6 ----return in ? -1 : 1; }--------// 77
----// NOTE: check for parallel/collinear lines before calling this function---// 02 // pair<polygon, polygon cut_polygon (const polygon &poly, point a, point b) {-// 7b
----point r = b - a, s = q - p;-------// 79 //---- polygon left, right;-------// 6b
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// a8 //---- point it(-100, -100);-------// c9
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))-------// ae //---- for (int i = 0, cnt = poly.size(); i < cnt; i++) {-------------------------// 28
------return false:-------// a3 //------ int i = i == cnt-1 ? 0 ; i + 1;-------// 8e
----res = a + t * r;------// ca //------ point p = poly[i], q = poly[j];------// 19
point closest_point(L(a, b), P(c), bool segment = false) {-------// a1 //-----// myintersect = intersect where-----// 24
------if (dot(b - a, c - b) > 0) return b;-------// b5 //----- if (myintersect(a, b, p, q, it))-------// f0
------if (dot(a - b, c - a) > 0) return a;------// cf //------ left.push_back(it), right.push_back(it);------// 21
----double t = dot(c - a, b - a) / norm(b - a);------// aa //--- return pair<polygon, polygon>(left, right);------// 1d
----return a + t * (b - a);------// 7a // }-----// 37
}-----// e5
double line_segment_distance(L(a,b), L(c,d)) {------// 99
                              6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
----double x = INFINITY;------// 83 #include "polygon.cpp"-----// 58
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);------// df #define MAXN 1000------// 09
                              point hull[MAXN];-----// 43
----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true)); -----// da
------ (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;------// 79 -----real(a) < real(b) : imag(a) < imag(b); }------// 40
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// ec ----sort(p.begin(), p.end(), cmp);------// 3d
-----x = min(x, abs(c - closest_point(a,b, c, true)));-------// 36 ----for (int i = 0; i < n; i++) {-------// 6f
-----x = min(x, abs(d - closest_point(a,b, d, true)));-------// e5 ------if (i > 0 && p[i] == p[i - 1]) continue;------// b2
}-----// b3 ---}-----// d8
                              ----int r = l:------// 59
6.2. Polygon. Polygon primitives.
                              ----for (int i = n - 2; i >= 0; i--) {------// 16
#include "primitives.cpp"------// e0 ------if (p[i] == p[i + 1]) continue;------// c7
----double area = 0; int cnt = size(p);-----// a2 ---}
-----area += cross(p[i] - p[0], p[i + 1] - p[0]);------// 7e }-----// 7e
----return area / 2; }------// e1
double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25 6.4. Line Segment Intersection. Computes the intersection between two line segments.
```

```
------A = B = a; return abs(a - d) < EPS; }------// ee
----else if (abs(a - b) < EPS) {------// 03
-----A = B = a; double p = progress(a, c,d);------// c9
-----return 0.0 <= p && p <= 1.0-----// 8a
----else if (abs(c - d) < EPS) {------// 26
------A = B = c; double p = progress(c, a,b);------// d9
-----return 0.0 <= p && p <= 1.0-----// 8e
----else if (collinear(a,b, c,d)) {------// bc
------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
-----if (ap > bp) swap(ap, bp);-----// b1
------if (bp < 0.0 || ap > 1.0) return false;------// 0c
-----A = c + max(ap, 0.0) * (d - c);------// f6
-----B = c + min(bp, 1.0) * (d - c);------// 5c
-----return true; }-----// ab
----else if (parallel(a,b, c,d)) return false;-----// ca
----else if (intersect(a,b, c,d, A, true)) {------// 10
-----B = A; return true; }------// bf
----return false:-----// b7
}-----// 8b
-----// e6
```

6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude coordinates) on a sphere of radius r.

```
double gc_distance(double pLat, double pLong,-----// 7b
-----/ double qLat, double qLong, double r) {-----// a4
----pLat *= pi / 180; pLong *= pi / 180;-----// ee
----qLat *= pi / 180; qLong *= pi / 180;-----// 75
----return r * acos(cos(pLat) * cos(pLong) * cos(qLat) * cos(qLong) +-----// a1
-----cos(pLat) * sin(pLong) * cos(qLat) * sin(qLong) +-----// ea
-----sin(pLat) * sin(qLat)); }-----// 5b
```

- 6.6. Formulas. Let  $a = (a_x, a_y)$  and  $b = (b_x, b_y)$  be two-dimensional vectors.
  - $a \cdot b = |a||b|\cos\theta$ , where  $\theta$  is the angle between a and b.
  - $a \times b = |a||b|\sin\theta$ , where  $\theta$  is the signed angle between a and b.
  - $a \times b$  is equal to the area of the parallelogram with two of its sides formed by a and b. Half of that is the area of the triangle formed by a and b.

#### 7. Other Algorithms

7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous function f on the interval [a, b], with a maximum error of  $\varepsilon$ .

```
double binary_search_continuous(double low, double high,-----// 8e
-----double eps, double (*f)(double)) {------// c0
----while (true) {------// 3a
-----double mid = (low + high) / 2, cur = f(mid);-----// 75
-----if (abs(cur) < eps) return mid;------// 76
-----else if (0 < cur) high = mid;-----// e5
-----else low = mid:-----// a7
----}------// b5
}-----// cb }-----// cb
```

```
Another implementation that takes a binary predicate f, and finds an integer value x on the integer
interval [a,b] such that f(x) \wedge \neg f(x-1).
```

```
----assert(low <= high);-----// 19
----while (low < high) {------// a3
------int mid = low + (high - low) / 2;-----// 04
------if (f(mid)) high = mid;------// ca
-----else low = mid + 1;-----// 03
----}-----// 9b
----assert(f(low));-----// 42
----return low:-----// a6
1-----// d3
7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotoni-
```

cally decreasing, ternary search finds the x such that f(x) is maximized. template <class F>-----// d1

```
double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
----while (hi - lo > eps) {------// 3e
------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
------if (f(m1) < f(m2)) lo = m1;------// 1d
-----else hi = m2:-----// b3
----}--------// bb
----return hi;------// fa
}-----// 66
```

7.3. **2SAT.** A fast 2SAT solver.

```
#include "../graph/scc.cpp"-----// c3
  -----// 63
bool two_sat(int n, const vii& clauses, vi& all_truthy) {------// f4
----all_truthy.clear();------// 31
----vvi adj(2*n+1);------// 7b
----for (int i = 0; i < size(clauses); i++) {-------// 9b
-----adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// 17
-----if (clauses[i].first != clauses[i].second)------// 87
-----adi[-clauses[i].second + nl.push_back(clauses[i].first + n):-----// 93
----}-----------// d8
----pair<union_find, vi> res = scc(adj);------// 9f
----union_find scc = res.first;------// 42
----vi dag = res.second;------// 58
----vi truth(2*n+1, -1);------// 00
----for (int i = 2*n; i >= 0; i--) {------// f4
-----int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n); -\frac{1}{5}
-----if (cur == 0) continue;-----// 26
-----if (p == o) return false;-----// 33
```

-----if (truth[p] == -1) truth[p] = 1;-----// c3

-----truth[cur + n] = truth[p];-----// b3

-----truth[o] = 1 - truth[p];-----// 80

-----if (truth[p] == 1) all\_truthy.push\_back(cur);-------// 5c ----}------// d9

```
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                                           -----if (!ptr[i][j]) continue;-----// 76
7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
                                           ------int ni = i + 1, nj = j + 1;------// 34
vi stable_marriage(int n, int** m, int** w) {------// e4
                                              ----queue<int> q;-----// f6
                                                 -----if (ni == rows + 1) ni = 0;------// 54
----vi at(n, \theta), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
                                           -----/if (ni == rows || arr[ni][j]) break;---------// 77
----for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)------// 05
-----inv[i][w[i][j]] = j;------// b9
                                                     -----// 47
----for (int i = 0; i < n; i++) q.push(i);-----// fe
                                                   -ptr[i][j]->d = ptr[ni][j];-----// a9
----while (!g.empty()) {------// 55
                                                  ---ptr[ni][j]->u = ptr[i][j];-----// c0
-----int curm = q.front(); q.pop();-----// ab
                                             ------while (true) {------// 0d
------for (int &i = at[curm]; i < n; i++) {-------// 9a
                                                 ------if (nj == cols) nj = 0;-----------------// a7
-----int curw = m[curm][i];-----// cf
                                             -----/if (i == rows || arr[i][nj]) break;------------// e9
-----if (eng[curw] == -1) { }------// 35
                                              -----++ni:-----// a6
------else if (inv[curw][curm] < inv[curw][enq[curw]])------// 10
-----q.push(eng[curw]);-----// 8c
                                              -----ptr[i][j]->r = ptr[i][nj];-------------------// b3
-----else continue;-----// b4
                                              -----ptr[i][nj]->l = ptr[i][j];-----// 46
-----res[eng[curw] = curm] = curw, ++i; break;------// 5e
                                           -----head = new node(rows, -1);------// 80
                                           -----head->r = ptr[rows][0];------// b9
                                           -----ptr[rows][0]->l = head:-----// c1
                                           ------head->l = ptr[rows][cols - 1];------// 28
7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the
                                           -----ptr[rows][cols - 1]->r = head;------// 83
Exact Cover problem.
                                           ------for (int j = 0; j < cols; j++) {-------// 02
bool handle_solution(vi rows) { return false; }------// 63
                                           -----// 36
struct exact_cover {------// 95
                                           ------for (int i = 0; i \le rows; i++)------// 56
----struct node {------// 7e
                                           -----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];------// \theta 5
-----node *l, *r, *u, *d, *p;-----// 19
                                           -----/ptr[rows][j]->size = cnt;------// d4
------int row, col, size;------// ae
                                           ------}-----// 8f
-----node(int row, int col) : row(row), col(col) {------// 68
                                           ------for (int i = 0; i <= rows; i++) delete[] ptr[i];-----// cd
-----// 8f
                                           -----delete[] ptr:-----// 42
                                           ----}------// a9
----int rows, cols, *sol;------// 54
                                           ----#define COVER(c, i, j) \\ ------// 23
                                           ------for (node *i = c->d; i != c; i = i->d) \------// 5c
----exact_cover(int rows, int cols) : rows(rows), cols(cols), head(NULL) {-----// ce
                                           ------for (node *j = i->r; j != i; j = j->r) \sqrt{\phantom{a}}
-----arr = new bool*[rows]:-----// 15
                                           -----j->d->u = j->u, j->u->d = j->d, j->p->size--;------// 5a
-----sol = new int[rows];-----// 69
                                           ----<mark>#</mark>define UNCOVER(c, i, j) <mark>\</mark>------// 17
------for (int i = 0; i < rows; i++)-----// c7
                                           ------for (node *i = c->u; i != c; i = i->u) \------// 98
-----arr[i] = new bool[cols], memset(arr[i], 0, cols);------// 68
                                           ----void setup() {-------// a8 -----c->r->l = c->l->r = c;------// bb
-----node ***ptr = new node**[rows + 1];------// da ----bool search(int k = 0) {-------// 4f
------for (int i = 0; i <= rows; i++) {--------// ce -----if (head == head->r) {------// a7
-----ptr[i] = new node*[cols];------// cc -----vi res(k);------
------for (int j = 0; j < cols; j++)------// 56 ------for (int i = 0; i < k; i++) res[i] = sol[i];-----// c0
-----sort(res.begin(), res.end());------// 3e
------else ptr[i][j] = NULL;------// 40 -----return handle_solution(res);-----// dc
------for (int i = 0; i <= rows; i++) {-------// 80
                                           -----node *c = head->r, *tmp = head->r;-----// a6
-----for (int j = 0; j < cols; j++) {------// 86
```

```
------for (; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 1e ----j = 80 * x / 2447;------------------------// 3d
------for (node *r = c->d; !found && r != c; r = r->d) {---------// 1e ----y = 100 * (n - 49) + i + x;--------------// 70
------for (node *j = r - r; j != r; j = j - r) { COVER(j - p, a, b); } -----// 3a
-----found = search(k + 1);-----// f4
                                                                    8. Useful Information
------for (node *j = r > l; j != r; j = j > l) { UNCOVER(j > p, a, b); } ----// 8a
                                                  8.1. Tips & Tricks.
------UNCOVER(c, i, i);------// 64
                                                    • How fast does our algorithm have to be? Can we use brute-force?
-----return found:-----// ff
                                                    • Does order matter?
• Is it better to look at the problem in another way? Maybe backwards?
}:-----// 10
                                                    • Are there subproblems that are recomputed? Can we cache them?
                                                    • Do we need to remember everything we compute, or just the last few iterations of computation?
7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \ldots, k-1\}
                                                    • Does it help to sort the data?
1}.
                                                    • Can we speed up lookup by using a map (tree or hash) or an array?
vector<int> nth_permutation(int cnt, int n) {------// 78
                                                    • Can we binary search the answer?
----vector<int> idx(cnt), per(cnt), fac(cnt);-----// 9e
                                                    • Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph
----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
                                                      into some other kind of a graph (perhaps a DAG, or a flow network)?
----for (int i = 1; i <= cnt; i++) fac[i - 1] = n % i, n /= i;-----// 04

    Make sure integers are not overflowing.

----for (int i = cnt - 1; i >= 0; i--)-----// 52
                                                    • Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// 41
                                                      m_1, m_2, \ldots, m_k, where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer
----return per;------// 84
                                                      using CRT?
}-----// 97
                                                    • Are there any edge cases? When n = 0, n = -1, n = 1, n = 2^{31} - 1 or n = -2^{31}? When
                                                      the list is empty, or contains a single element? When the graph is empty, or contains a single
7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
                                                      vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
ii find_cycle(int x0, int (*f)(int)) {------// a5
                                                    • Can we use exponentiation by squaring?
----int t = f(x\theta), h = f(t), mu = \theta, lam = 1:-----// 8d
----while (t != h) t = f(t), h = f(f(h));-----// 79
                                                 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input
                                                  reading/output writing. This can be achieved by reading all input in at once (using fread), and then
----h = x0;
----while (t != h) t = f(t), h = f(h), mu++;------// 9d
                                                  parsing it manually. Output can also be stored in an output buffer and then dumped once in the end
----h = f(t);-----// 00
                                                  (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading
----while (t != h) h = f(h), lam++:-----// 5e
                                                 method.
----return ii(mu. lam):-----// b4
                                                  void readn(register int *n) {------// dc
}------// 42
                                                  ----int sign = 1;------// 32
                                                  ----register char c;------// a5
7.8. Dates. Functions to simplify date calculations.
                                                  ----*n = 0;-----// 35
int intToDay(int jd) { return jd % 7; }-----// 89
                                                  ----while((c = qetc_unlocked(stdin)) != '\n') {------// f3
int dateToInt(int y, int m, int d) {-----// 96
                                                 -----switch(c) {------// 0c
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-------// a8 ------case '-': sign = -1; break;------// 28
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1 -----case ' ': qoto hell;------
-----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +-------// be ------case '\n': goto hell;-------// 79
-----d - 32075;-------default: *n *= 10; *n += c - '0'; break;------// c0
}------// fa _____}_____// 2d
void intToDate(int jd, int &y, int &m, int &d) {------// a1 ______/ c3
----int x, n, i, j;------// 00
                                                  hell:----// ba
----x = id + 68569;-----// 11
                                                 ----*n *= sign;-----// a0
----n = 4 * x / 146097;-----// 2f ].....// 67
---x = (146097 * n + 3) / 4:
----x -= 1461 * i / 4 - 31;------// 09 plication of 64-bit integers, or something needing a little more than 64-bits to represent.
```

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# 8.4. Worst Time Complexity.

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n	Worst AC Algorithm	Comment
$\leq 10$	$O(n!), O(n^6)$	e.g. Enumerating a permutation
$\leq 15$	$O(2^n \times n^2)$	e.g. DP TSP
$\leq 20$	$O(2^n), O(n^5)$	e.g. $DP + bitmask technique$
$\leq 50$	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
$\le 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree
$\le 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

# 8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- $\bullet$  snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.