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------if (l > r) return ID;-------// cc ----return query(segs[id].lid, l, r) + query(segs[id].rid, l, r); }------// ee
-----if (r < a || b < l) return data[i];-----// d9
                                     2.3. Fenwick Tree.
------if (a <= l \& \& r <= b) return (lazy[i] = v) * (r - l + 1) + data[i];----// 06
                                     struct fenwick_tree {-----// 98
-----int m = (l + r) / 2;-----// cc
                                     ----int n; vi data;------// d3
-----return data[i] = f(ru(a, b, v, l, m, 2*i+1),-----// cc
                                     ----fenwick_tree(int _n) : n(_n), data(vi(n)) { }------// db
-----/u(a, b, v, m+1, r, 2*i+2));-----// 2b
                                     ----void update(int at, int by) {------// 76
----}------// @b
----void propagate(int l, int r, int i) {-----// a7
                                     ------while (at < n) data[at] += by, at |= at + 1; }------// fb
                                     ----int query(int at) {-------// 71
-----if (l > r || lazy[i] == INF) return;-----// 5f
                                     ------int res = 0;-----// c3
-----data[i] += lazy[i] * (r - l + 1);-----// 44
                                     ------while (at >= 0) res += data[at], at = (at \& (at + 1)) - 1;------// 37
-----if (l < r) {------// 28
                                     -----return res; }-----// e4
------if (lazy[2*i+1] == INF) lazy[2*i+1] = lazy[i];------// 4e
                                     ----int rsq(int a, int b) { return query(b) - query(a - 1); }-----// be
------else lazy[2*i+1] += lazy[i];-----// 1e
                                      -----// 57
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];-----// de
                                     struct fenwick_tree_sq {-----// d4
------else lazy[2*i+2] += lazy[i];-----// 74
                                     ----int n; fenwick_tree x1, x0;------// 18
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e
-----lazy[i] = INF;-----// f8
                                     ------x0(fenwick_tree(n)) { }------// 7c
----// insert f(y) = my + c if x <= y------// 17
}:----// ae
                                     ----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }----// 45
                                     ----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73
                                     }:-----// 13
2.2.1. Persistent Segment Tree.
                                     void range_update(fenwick_tree_sq &s, int a, int b, int k) {------// 89
int segcnt = 0;-----// cf
                                     ----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f
struct segment {------// 68
                                     int range_query(fenwick_tree_sq &s, int a, int b) {------// 15
----int l, r, lid, rid, sum;------// fc
                                     ----return s.query(b) - s.query(a-1); }------// f3
} seqs[2000000];-----// dd
int build(int l, int r) {------// 2b
                                     2.4. Matrix.
----int id = seqcnt++;-------------------------// a8 template <> bool eq<double o, double b) { return abs(a - b) < EPS; }---// a7
----else {------// fe ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {------// 56
------data.assign(cnt, T(0)); }------// e3
------seqs[id].lid = build(l , m);-------// e3 ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// b5
-----segs[id].rid = build(m + 1, r); }------// 69 ------cnt(other.cnt), data(other.data) { }------// c1
----return id; }-------------------------// c5 ----matrix<T> operator +(const matrix& other) {-------------------------// 33
int update(int idx, int v, int id) {------// b8 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] += other.data[i];-----// f8
----int nid = segcnt++;------matrix<T> res(*this); rep(i,0,cnt) res.data[i] -= other.data[i];-----// 7b
----seqs[nid].l = seqs[id].l;------// 78 -----return res; }-----
----segs[nid].lid = update(idx, v, segs[id].lid);------// 92 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] *= other;------// 05
----seqs[nid].rid = update(idx, v, seqs[id].rid);-------// 06 ------return res; }------
----return nid; }------// e6 ------matrix<T> res(rows, other.cols);------// 4c
----if (r < segs[id].l || segs[id].r < l) return 0;-------// 17 ------res(i, j) += at(i, k) * other.data[k * other.cols + j];------// 17
```

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----matrix<T> pow(int p) {-------------------------// 53 ----r->l = merge(l, r->l); augment(r); return r; }------------------------// cθ
-----if (p) sq = sq * sq;-------// 35 ----return NULL; }------// ae
------} return res; }--------// 22 node* insert(node *t, int x, int y) {--------// 78
------matrix<T> mat(*this); det = T(1), rank = max(rows, cols);------// 7a ----pair<node*, node*> res = split(t, x);--------// ca
------for (int r = 0, c = 0; c < cols; c++) {---------// 8e ----return merge(res.first, merge(new node(x, y), res.second)); }------// 0d
-----int k = r;------// 5b node* erase(node *t, int x) {-------// 4d
-----if (k >= rows) { rank-; continue; }------// 1a ---if (t->x < x) t->r = erase(t->r, x);-------// 7c
-----if (k != r) {-------// c4 ----else if (x < t->x) t->l = erase(t->l, x);-------// 48
------swap(mat.at(k, i), mat.at(r, i));-------// 7d int kth(node *t, int k) {---------------------------------// b3
-----rep(i,0,cols) mat(r, i) /= d;------// d1 ----else return kth(t->r, k - tsize(t->l) - 1); }------// f0
-----rep(i,0,rows) {------// f6
                         2.6. Misof Tree.
-----T m = mat(i, c);------// 05
                         #define BITS 15-----// 7b
-----if (i != r && !eq<T>(m, T(0)))------// 1a
                         struct misof_tree {-----// fe
-----rep(j,0,cols) mat(i, j) -= m * mat(r, j);-----// 7b
                         ----int cnt[BITS][1<<BITS];------// aa
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }-----// b0
-----} return mat; }------// b3
                         ----void insert(int x) { for (int i = 0; i < BITS; cnt[i++][x]++, x >>= 1); }--//5a
----matrix<T> transpose() {------// 59
                         ----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49
------matrix<T> res(cols, rows);------// 5b
----rep(i,0,rows) rep(j,0,cols) res(j, i) = at(i, j);-----// 92
                         ----int nth(int n) {-------// 8a
-----return res; } };------// df
                         -----int res = 0;------// a4
                         ------for (int i = BITS-1; i >= 0; i--)------// 99
                         ------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4
2.5. Cartesian Tree.
                         -----return res:-----// 3a
struct node {-----// 36
                         ----}-----// b5
----int x, y, sz;------// e5
                         };-----// 0a
---node *l, *r;-----// 4d
                         2.7. Sqrt Decomposition.
----node(int _x, int _y) : x(_x), y(_y), sz(1), l(NULL), r(NULL) { } };------// 19
pair<node*, node*> split(node *t, int x) {------// 1d ----segment(vi _arr) : arr(_arr) { } };------// 11
----if (t->x < x) {-------------------------// dc
------return make_pair(t, res.second); }-------// e0 ---rep(i,0,size(T))--------// b1
```

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------T.push_back(segment(vi(arr.begin()+i, arr.begin()+min(i+K, cnt))));----// f0 ------return (h[i+1].second-h[i].second)/(h[i].first-h[i+1].first); }------// b9
------if (intersect(n-3) < intersect(n-2)) break;-------// 07
----if (i >= size(T)) return size(T);------// 83 ------swap(h[n-2], h[n-1]);------// bf
}-----// ea ------int mid = lo + (hi - lo) / 2;---------// 5a
----T.insert(T.begin() + split(at), segment(arr));------// 67 ------return h[res+1].first * x + h[res+1].second; } };------// 84
}-----// cc
                                3. Graphs
void erase(int at) {-----// be
----int i = split(at): split(at + 1):-----// da
                      3.1. Single-Source Shortest Paths.
----T.erase(T.begin() + i);-----// 6b
                      3.1.1. Dijkstra's algorithm.
}-----// 4b
                      int *dist. *dad:-----// 46
2.8. Monotonic Queue.
                      struct cmp {-----// a5
struct min_stack {-------// d8 ----bool operator()(int a, int b) {------------// bb
----stack<int> S. M:------// fe ------return dist[a] != dist[b] ? dist[b] : a < b; }------// e6
------S.push(x);-------// e2 pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
------M.push(M.empty() ? x : min(M.top(), x)); }-------// 92 ----dist = new int[n];-------// 84
struct min_queue {------// b4 -----// b4 -----// 58
----min_stack inp, outp;------// 3d -----rep(i,0,size(adj[cur])) {-----------// a6
----void fix() {-----------------------// 5d --------ndist = dist[cur] + adj[cur][i].second;------// 3a
------dist[nxt] = ndist, dad[nxt] = cur, pg.insert(nxt);------// eb
-----if (inp.empty()) return outp.mn();------// 01 }------// 9b
-----if (outp.empty()) return inp.mn();------// 90
                      3.1.2. Bellman-Ford algorithm.
-----return min(inp.mn(), outp.mn()); }-----// 97
                      int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {------// cf
----void pop() { fix(); outp.pop(); }-----// 4f
----bool empty() { return inp.empty() && outp.empty(); }-----// 65
                      ----has_negative_cycle = false;-----// 47
                      ----int* dist = new int[n];-----// 7f
}:-----// 60
                      ----rep(i,0,n) dist[i] = i == s ? 0 : INF;------// df
                      ----rep(i,0,n-1) rep(j,0,n) if (dist[j] != INF)------// 4d
2.9. Convex Hull Trick.
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-----rep(k,0,size(adj[j]))----------// 88 vi order;------// 9b
------dist[adj[j][k].first] = min(dist[adj[j][k].first],------// e1 -------// a5
-----/dist[j] + adj[j][k].second);-------// 18 void scc_dfs(const vvi &adj, int u) {--------------------// a1
------if (dist[j] + adj[j][k].second < dist[adj[j][k].first])--------// 37 ----rep(i,0,size(adj[u]))------------------------------// 2d
------has_negative_cycle = true;-------// f1 ------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-------// a2
}------// a9 }------// 53
                                  -----// 63
3.1.3. IDA^* algorithm.
                                  pair<union_find, vi> scc(const vvi &adj) {------// c2
int n, cur[100], pos;-----// 48 ----int n = size(adj), u, v;------// f8
----int h = 0;-------// 4a ----union_find uf(n);------// a8
---rep(i,0,n) if (cur[i] != 0) h += abs(i - cur[i]);------// 9b ---vi dag;-----
----return h;------// c6 ----vvi rev(n);------// c5
}------// c8 ----rep(i,0,n) rep(j,0,size(adj[i])) rev[adj[i][j]].push_back(i);------// 7e
----if (q + h > d) return g + h;-------// 15 ----fill(visited.begin(), visited.end(), false);------// 59
----if (h == 0) return 0;------// ff ----stack<int> S;------// bb
----int mn = INF;------// 7e ----for (int i = n-1; i >= 0; i--) {-------// 96
----rep(di,-2,3) {-------// 0d ------if (visited[order[i]]) continue;------// db
------if (di == 0) continue; ------// 0a ------S.push(order[i]), dag.push_back(order[i]); ------// 68
------int nxt = pos + di;-------// 76 ------while (!S.empty()) {-------// 9e
------if (nxt == prev) continue; -------// 39 ------visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]); -----// b3
------if (0 <= nxt && nxt < n) {-------// 68 ------rep(j,0,size(adj[u])) if (!visited[v = adj[u][j]]) S.push(v);-----// 1b
-----mn = min(mn, dfs(d, g+1, nxt));------// 22 ----return pair<union_find, vi>(uf, dag);------// 2b
------swap(cur[pos], cur[nxt]);-----// 3b
3.3. Cut Points and Bridges.
-----if (mn == 0) break;-----// 8f
                                  #define MAXN 5000-----// f7
----}------// d3
                                  int low[MAXN], num[MAXN], curnum;-----// d7
----return mn:-----// da
                                  void dfs(const vvi &adj, vi &cp, vii &bri, int u, int p) {------// 22
}-----// f8
                                  ----low[u] = num[u] = curnum++;-----// a3
int idastar() {-----// 22
                                  ----int cnt = 0; bool found = false;-----// 97
----rep(i,0,n) if (cur[i] == 0) pos = i;------// 6b
                                  ----rep(i,0,size(adj[u])) {------// ae
----int d = calch();-----// 38
                                  -----int v = adj[u][i];------// 56
----while (true) {------// 18
                                  -----if (num[v] == -1) {------// 3b
-----int nd = dfs(d, 0, -1):-----// 42
                                  -----dfs(adj, cp, bri, v, u);-----// ba
-----if (nd == 0 || nd == INF) return d;------// b5
                                  -----low[u] = min(low[u], low[v]);-----// be
-----d = nd;-----// f7
                                  -----Cnt++;-----// e0
----}------// f9
                                  -----found = found || low[v] >= num[u];-----// 30
}-----// 82
                                  ------if (low[v] > num[u]) bri.push_back(ii(u, v));------// bf
                                  -----} else if (p != v) low[u] = min(low[u], num[v]); }------// 76
3.2. Strongly Connected Components.
                                  ----if (found && (p !=-1 \mid \mid cnt > 1)) cp.push_back(u); }-------// 3e
3.2.1. Kosaraju's algorithm.
                                  pair<vi,vii> cut_points_and_bridges(const vvi &adj) {------// 76
#include "../data-structures/union_find.cpp"---------------------// 5e ----int n = size(adj);-----------------------------------// c8
-----// 11 ----vi cp; vii bri;------------// fb
```

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------e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];------// 95 -------dist[g[j][k]->v]) {-------// 6d
-----rep(i,0,n) if (p[i] != -1) pot[i] += d[i];------// 86 -------back[g[j][k]->v] = g[j][k];-----// 3d
-----cap = min(cap, cure->w);-----// c3
 A second implementation that is slower but works on negative weights.
                                   -----if (cure->u == s) break;-----// 82
struct flow_network {------// 81
                                   -----cure = back[cure->u];-----// 45
----struct mcmf_edge {------// f6
                                   -----int u, v;-----// e1
                                   -----assert(cap > 0 && cap < INF);-----// ae
-----ll w. c:-----// b4
                                   -----cure = back[t];-----// b9
------mcmf_edge* rev;------// 9d
                                   ------while (true) {------// 2a
-----mcmf_edge(int _u, int _v, ll _w, ll _c, mcmf_edge* _rev = NULL) {-----// ea
                                   -----cost += cap * cure->c;-----// f8
------u = _u: v = _v: w = _w: c = _c: rev = _rev:-----// 83
                                   -----Cure->w -= cap;-----// d1
-----cure->rev->w += cap;-----// cf
----};-------// b9
                                   ------if (cure->u == s) break;------// 8c
----int n;------// b4
                                   -----cure = back[cure->u];------// 60
----vector<pair<int, pair<ll, ll> > * adj;------// 72
----flow_network(int _n) {------// 55
                                   -----n = _n;-----// fa
                                   ------}-------// be
-----adj = new vector<pair<int, pair<ll, ll> > >[n];-----// bb
                                   -----// instead of deleting g, we could also-----// e0
----}------// bd
                                   -----// use it to get info about the actual flow-----// 6c
----void add_edge(int u, int v, ll cost, ll cap) {------// 79
                                   ------for (int i = 0; i < n; i++)------// eb
-----adj[u].push_back(make_pair(v, make_pair(cap, cost)));-----// c8
                                   -----for (int j = 0; j < size(g[i]); j++)------// 82
----}-----// ed
                                   -----delete q[i][j];-----// 06
----pair<ll,ll> min_cost_max_flow(int s, int t) {------// ea
                                   -----delete[] a:-----// 23
-----vector<mcmf_edge*>* q = new vector<mcmf_edge*>[n];-----// ce
                                   -----delete[] back;-----// 5a
------for (int i = 0; i < n; i++) {------// 57
                                   -----delete[] dist;-----// b9
-----for (int j = 0; j < size(adj[i]); j++) {------// 37
                                   -----return make_pair(flow, cost);------// ec
-----mcmf_edge *cur = new mcmf_edge(i, adj[i][j].first,-----// 21
                                   -----adj[i][j].second.first, adj[i][j].second.second),--// 56
                                    -----// bf
-----*rev = new mcmf_edge(adj[i][j].first, i, 0,------// 48
-----/<sub>b1</sub>
                                   3.8. All Pairs Maximum Flow.
-----cur->rev = rev;------// ef
                                   3.8.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is constructed
-----g[i].push_back(cur);-----// 1d
                                   using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the maximum
-----g[adj[i][j].first].push_back(rev);------// 05
                                   flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|). NOTE:
Not sure if it works correctly with disconnected graphs.
------ll flow = 0, cost = 0;------// 68
                                   #include "dinic.cpp"-----// 58
                                   -----// 25
-----mcmf_edge** back = new mcmf_edge*[n];------// e5
------ll* dist = new ll[n];------// 50 bool same[MAXV];------// 59
------for (int i = 0; i < n - 1; i++)-------// be ----rep(s,1,n) {-------// 9e
------for (int j = 0; j < n; j++)-------// 6e -----int l = 0, r = 0;-----------// 08
------if (dist[j] != INF)------// e3 ------par[s].second = g.max_flow(s, par[s].first, false);------// 54
------for (int k = 0; k < size(q[j]); k++)------// 85 -----memset(d, 0, n * sizeof(int));----------// c8
------| (g[j][k]->w > 0 && dist[j] + g[j][k]->c <-----// 7f -----memset(same, 0, n * sizeof(bool));-----------// c9
```

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-d[g[r++] = s] = 1; part(best); part(best); -d[g[r++] = s] = 1;
------while (l < r) {-------// 45 -----rep(i,0.size(adi[u]))-------// 92
-----same[v = q[l++]] = true;------// c5 ------if (adj[u][i] != parent[u] && adj[u][i] != best)-----// e8
-----if (q.e[i].cap > 0 && d[q.e[i].v] == 0)------// 21 ----void build(int r = 0) { curloc = 0, csz(curhead = r), part(r); }------// 78
------while (u ! = -1) uat.push_back(u), u = parent[head[u]]:------// 51
------while (true) {-------// b8 ----int query_upto(int u, int v) { int res = ID;-------// 72
------cap[curl[i] = mn:-------// 8d ------while (head[u] != head[v])-------// 69
------if (cur == 0) break;------// fb -----res = f(res, values.query(loc[head[u]], loc[u])),-----// a4
----return make_pair(par, cap);-------// 62 ------return f(query_upto(u, l), query_upto(v, l)); } };------// 5b
}-----// b3
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {-------// 93
                             3.10. Centroid Decomposition.
----if (s == t) return 0;-----// 33
                             #define MAXV 100100-----// 86
----int cur = INF, at = s;-----// e7
                             #define LGMAXV 20-----// aa
----while (gh.second[at][t] == -1)------// 42
                             int imp[MAXV][LGMAXV],....// 6d
-----cur = min(cur, qh.first[at].second), at = qh.first[at].first;-----// 8d
                             ----path[MAXV][LGMAXV],------// 9d
----return min(cur, gh.second[at][t]);-----// 54
                             ----sz[MAXV], seph[MAXV],-----// cf
}------// 46
                             ----shortest[MAXV]:-----// 6b
                             struct centroid_decomposition {------// 99
3.9. Heavy-Light Decomposition.
                             ----int n; vvi adj;-----// e9
#include "../data-structures/segment_tree.cpp"------// 16 ---centroid_decomposition(int _n) : n(_n), adj(n) { }-------// 46
struct HLD {------// 25 ----void add_edge(int a, int b) { adj[a].push_back(b), adj[b].push_back(a); }--// bc
----vvi adj; segment_tree values;---------// 13 ------rep(i,0,size(adj[u])) if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u);--// 78
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), adj(n) {----// 1c ------return sz[u]; }-----
-----vi tmp(n, ID); values = segment_tree(tmp); }-------// f0 ----void makepaths(int sep, int u, int p, int len) {-------// 84
----void add_edge(int u, int v) { adj[u].push_back(v), adj[v].push_back(u); }--// 77 ------jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len;-------------// d9
-----if (parent[v] == u) swap(u, v); assert(parent[u] == v);-------// db -----rep(i,0,size(adj[u])) {--------// f4
-------if (adj[u][i] == p) bad = i;-------// cf
-----sz[u] += csz(adj[parent[adj[u][i]] = u][i]);------// c2 -----if (p == sep) swap(adj[u][bad], adj[u].back()), adj[u].pop_back(); }---// 07
------head[u] = curhead; loc[u] = curloc++;---------// 63 ------down: iter(nxt,adj[sep])--------// 04
-----rep(i,0,size(adj[u]))------// 49 -----sep = *nxt; goto down; }-----// 1a
-----best = adj[u][i];-------// 26 -----rep(i,0,size(adj[sep])) separate(h+1, adj[sep][i]); }------// 90
```

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-----rep(h,0,seph[u]+1)------// c5 ------process(y):------// e8
-----shortest[jmp[u][h]] = min(shortest[jmp[u][h]], path[u][h]); }-----// 11 ------uf.unite(u,v);----------------------------------// 55
----int closest(int u) {-------// 91 ------ancestor[uf.find(u)] = u;------// 1d
-----int mn = INF/2:------// fe -----}-----// fe
-----rep(h,0,seph[u]+1) mn = min(mn, path[u][h] + shortest[jmp[u][h]]);-----// 3e ------colored[u] = true;-------------------------// b9
-----return mn; } };------// 13 -----rep(i,0,size(queries[u])) {-------// d7
                                    -----int v = queries[u][i].first;-----// 89
3.11. Least Common Ancestors, Binary Jumping.
                                    -----if (colored[v]) {------// cb
struct node {-----// 36
                                    -----answers[queries[u][i].second] = ancestor[uf.find(v)];------// 63
---node *p, *imp[20];-----// 24
                                    ----int depth:-----// 10
                                    -----}-----// 40
---node(node *_p = NULL) : p(_p) {-----// 78
                                    ----}-------------------------// a9
-----depth = p ? 1 + p->depth : 0;-----// 3b
                                    }:....// le
-----memset(jmp, 0, sizeof(jmp));-----// 64
                                    3.13. Maximum Density Subgraph. Given (weighted) undirected graph G. Binary search density.
-----for (int i = 1; (1<<i) <= depth; i++)------// a8
                                    If g is current density, construct flow network: (S, u, m), (u, T, m + 2g - d_u), (u, v, 1), where m is a
-----jmp[i] = jmp[i-1]->jmp[i-1]; } };-----// 3b
                                    large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has
                                    empty S-component, then maximum density is smaller than q, otherwise it's larger. Distance between
node* st[100000];-----// 65
node* lca(node *a, node *b) {------// 29
                                    valid densities is at least 1/(n(n-1)). Edge case when density is 0. This also works for weighted
                                    graphs by replacing d_u be the weighted degree, and doing more iterations (if weights are not integers).
----if (!a || !b) return NULL;------// cd
----if (a->depth < b->depth) swap(a,b);-----// fe
                                    3.14. Maximum Weighted Independent Set in a Bipartite Graph. This is the same as the
----for (int j = 19; j >= 0; j--)-----// b3
                                    minimum weighted vertex cover. Solve this by constructing a flow network with edges (S, u, w(u))
------while (a->depth - (1<<j) >= b->depth) a = a->jmp[j];-----// cθ
                                    for u \in L, (v, T, w(v)) for v \in R and (u, v, \infty) for (u, v) \in E. The minimum S, T-cut is the answer.
----if (a == b) return a;-----// 08
                                    Vertices adjacent to a cut edge are in the vertex cover.
----for (int j = 19; j >= 0; j--)-----// 11
------while (a->depth >= (1<<j) && a->jmp[j] != b->imp[j])------// f\theta
                                                    4. Strings
-----a = a->jmp[i], b = b->jmp[i];-----// d\theta
                                    4.1. The Knuth-Morris-Pratt algorithm.
----return a->p; }-----// c5
                                    int* compute_pi(const string &t) {------// a2
3.12. Tarjan's Off-line Lowest Common Ancestors Algorithm.
                                    ----int m = t.size();-----// 8b
#include "../data-structures/union_find.cpp"----------------------// 5e ----int *pit = new int[m + 1];--------
struct tarjan_olca {-------// 87 ---if (0 <= m) pit[0] = 0;-------// 42
----int *ancestor;------// 39 ----if (1 <= m) pit[1] = 0;-------// 34
----vii *queries;------// 66 ------for (int j = pit[i - 1]; ; j = pit[j]) {-------// b5
----union_find uf;-------if (j == 0) { pit[i] = 0; break; }------// 95
-----queries = new vii[n];------// 3e int string_match(const string &s, const string &t) {------// 9e
-----queries[x].push_back(ii(y, size(answers)));-------// a0 -----if (s[i] == t[j]) {--------// 73
-----answers.push_back(-1);-------// ca ------if (j == m) {-------// de
```

```
-----else i++; }-----// b8
                         struct aho_corasick {------// 78
----delete[] pit; return -1; }-------// e3
                         ----struct out_node {------// 3e
                         -----string keyword; out_node *next;-----// f0
4.2. The Z algorithm.
                         -----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
int* z_values(const string &s) {------// 4d
                         ----};------// b9
----int n = size(s):-----// 97
                         ----struct go_node {------// 40
----int* z = new int[n]:-----// c4
                         -----map<char, qo_node*> next;------// 6b
----int l = 0, r = 0:-----// 1c
                         -----out_node *out; go_node *fail;-----// 3e
---z[0] = n;
                         -----qo_node() { out = NULL; fail = NULL; }-----// 0f
---rep(i,1,n) {------// b2
                         ----};-------// c0
----qo_node *qo;------// b8
-----if (i > r) {------// 6d
                         -----l = r = i;-----// 24
                         -----qo = new qo_node();-----// 77
-----iter(k, keywords) {------// f2
-----z[i] = r - l; r--;------// 07
                         -----qo_node *cur = qo;-----// a2
-----} else if (z[i - l] < r - i + 1) z[i] = z[i - l];------// 6f
                         -----iter(c, *k)------// 6e
-----else {------// a8
                         -----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 97
-----l = i:-----// 55
                         -----(cur->next[*c] = new go_node());------// af
-----cur->out = new out_node(*k, cur->out);------// 3f
-z[i] = r - i; r--; \}
                         ----return z;------// 78
                         -----queue<qo_node*> q;------// 2c
                         -----iter(a, go->next) q.push(a->second);------// db
4.3. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
                         ------while (!q.empty()) {------// 07
bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77 -----iter(a, r->next) {-------------------------// 18
-----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);-------// 12 -------while (st &\delta st->next.find(a->first) == st->next.end())------// \thetae
-----rep(i,0,n) P[0][i] = s[i];--------// 5c ------st = st->fail;----------// b3
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {-------// 86 -------if (!st) st = go;------------------------// θb
------P.push_back(vi(n));--------// 53 -------s->fail = st->next[a->first];------// c1
-----rep(i,0,n)-------if (s->fail) {--------// 98
------L[L[i].p = i].nr = ii(P[stp - 1][i],------// e2 -------if (!s->out) s->out = s->fail->out;------// ad
-----sort(L.begin(), L.end());-------// 5f ------out_node* out = s->out;------// b8
-----rep(i,0,n)-------while (out->next) out = out->next;------// b4
------while (cur && cur->next.find(*c) == cur->next.end())------// df
```

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------if (!cur) cur = qo;------// 01 ----bool issubstr(string other){------// 3b
-----res.push_back(out->keyword);------// 7c ------if(cur == -1) return false; cur = next[cur][other[i]]; }------// 54
-----if(p == -1){ link[cur] = 0; }-----// 18
4.5. eerTree.
                                           -----else{ int q = next[p][c];-----// 34
#define MAXN 100100-----// 29
                                           ------if(len[p] + 1 == len[q]){ link[cur] = q; }-----// 4d
#define SIGMA 26-----// e2
                                           ------else { int clone = sz++; isclone[clone] = true;-----// 57
#define BASE 'a'-----// a1
                                           -----len[clone] = len[p] + 1;------// 8c
char *s = new char[MAXN]:----// db
                                           -----link[clone] = link[q]; next[clone] = next[q];-----// 76
struct state {-----// 33
                                           -----for(; p != -1 \& ant[p].count(c) \& ant[p][c] == q; p = link[p]){
-----next[p][c] = clone; }-----// 32
} *st = new state[MAXN+2];-----// 57
                                           -----link[q] = link[cur] = clone;-----// 73
struct eertree {-----// 78
                                           ----int last, sz, n;------// ba
                                           ----void count(){------// e7
----eertree() : last(1), sz(2), n(0) {------// 83
                                           -----cnt=vi(sz, -1); stack<ii> S; S.push(ii(0,0));map<char,int>::iterator i;// 56
-----st[0].len = st[0].link = -1;------// 3f
                                           ------while(!S.emptv()){-------// 4c
-----st[1].len = st[1].link = 0; }------// 34
                                           -----ii cur = S.top(); S.pop();-----// 67
----int extend() {-------// c2
                                           -----if(cur.second){-----// 78
-----char c = s[n++]; int p = last;-----// 25
                                           -----for(i = next[cur.first].begin();i != next[cur.first].end();++i){
-----while (n - st[p].len - 2 < \theta || c != s[n - st[p].len - 2]) p = st[p].link;
                                           -----cnt[cur.first] += cnt[(*i).second]; } }-----// da
------if (!st[p].to[c-BASE]) {------// 82
                                           -----else if(cnt[cur.first] == -1){------// 99
-----int q = last = sz++;-----// 42
                                           -----cnt[cur.first] = 1; S.push(ii(cur.first, 1));-----// bd
-----st[p].to[c-BASE] = q;-----// fc
                                           -----for(i = next[cur.first].begin();i != next[cur.first].end();++i){
-----st[q].len = st[p].len + 2;-----// c5
                                           ------S.push(ii((*i).second, 0)); } } } }-----// 61
-----do { p = st[p].link;-----// 04
                                           ----string lexicok(ll k){------// 8b
-----} while (p != -1 \&\& (n < st[p].len + 2 || c != s[n - st[p].len - 2]));
                                           ------int st = 0; string res; map<char,int>::iterator i;------// cf
-----if (p == -1) st[q].link = 1;------// 77
                                           ------while(k){ for(i = next[st].begin(); i != next[st].end(); ++i){------// 69
------else st[q].link = st[p].to[c-BASE];------// 6a
                                           ------if(k <= cnt[(*i).second]){ st = (*i).second; ------// ec
-----return 1; }-----// 29
                                           -----res.push_back((*i).first); k--; break;-----// 63
-----last = st[p].to[c-BASE];------// 42
                                           -----} else { k -= cnt[(*i).second]; } } -----// ee
-----return 0: } }:-----// ec
                                           -----return res; }-----// 0b
                                           ----void countoccur(){------// ad
4.6. Suffix Automaton. Minimum automata that accepts all suffixes of a string with O(n) construc-
                                           ------for(int i = 0; i < sz; ++i){ occur[i] = 1 - isclone[i]; }-----// 1b
tion. The automata itself is a DAG therefore suitable for DP, examples are counting unique substrings,
                                           -----vii states(sz):-----// dc
occurrences of substrings and suffix.
                                           ------for(int i = 0; i < sz; ++i){ states[i] = ii(len[i],i); }------// 97
// TODO: Add longest common subsring-----/ 0e
                                           -----sort(states.begin(), states.end());------// 8d
const int MAXL = 100000;-----// 31
                                           -----for(int i = size(states)-1; i >= 0; --i){ int v = states[i].second; <math>---//a4
struct suffix_automaton {------// e0
                                           ------if(link[v] != -1) { occur[link[v]] += occur[v]; } } }-----// cc
----vi len, link, occur, cnt;------// 78
                                           }:-----// 32
----vector<map<char,int> > next;------// 90
                                             -----// 56
----vector<bool> isclone;-----// 7b
----ll *occuratleast;------// f2
----int sz. last:-----// 7d
                                                             5. Mathematics
                                           5.1. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
----suffix_automaton() : len(MAXL*2), link(MAXL*2), occur(MAXL*2), next(MAXL*2),
----isclone(MAXL*2) { clear(): }------// a3
                                           k items out of a total of n items. Also contains an implementation of Lucas' theorem for computing
```

the answer modulo a prime p.

----**void** clear(){ sz = 1; last = len[0] = 0; link[0] = -1; next[0].clear();----// aa

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int nck(int n, int k, int p) {-------------------------------// cf ----} return true; }-----------------------------// fe
----int res = 1:------// 5c
                               5.5. Pollard's \rho algorithm.
----while (n || k) {------// e2
-----res *= nck(n % p, k % p);------// 2C // public static int[] seeds = new int[] {2,3,5,7,11,13,1031};------// 1d
                               // public static BigInteger rho(BigInteger n, BigInteger seed) {------// 03
-----res %= p, n /= p, k /= p;-----// 0a
                               //--- int i = 0,-----// 00
----}-----// d9
                               //----- k = 2;-----// 79
----return res;------// 30
}------// Oa //--- BigInteger x = seed,-----// cc
                               //----y = seed;-----// 31
5.2. Euclidean algorithm.
                               //--- while (i < 1000000) {-----// 10
int qcd(int a, int b) { return b == 0 ? a : qcd(b, a % b); }-----// d9
                               //----- i++;-----// 8c
//---- BigInteger d = y.subtract(x).abs().gcd(n);-----// ce
and also finds two integers x, y such that a \times x + b \times y = d.
                               //----- if (!d.equals(BigInteger.ONE) && !d.equals(n)) {------// b9
int egcd(int a, int b, int& x, int& y) {-----// 85
                               //----- return d;-----// 3b
----if (b == 0) { x = 1; y = 0; return a; }------// 7b
                               //------} -------// 7c
----else {------// 00
                               //----- if (i == k) {-----// 2c
------int d = egcd(b, a % b, x, y);-----// 34
                               //----y = x;-----// 89
-----x -= a / b * y;------// 4a
                               -----Swap(x, y):-----// 26
                               //-----}
-----return d:-----// db
                               //---- }------// 96
//--- return BigInteger.ONE;-----// 62
}-----// 40
5.3. Trial Division Primality Testing.
                               5.6. Sieve of Eratosthenes.
bool is_prime(int n) {------// 6c
----if (n < 2) return false;-----// c9
                               vi prime_sieve(int n) {------// 40
                               ----int mx = (n - 3) >> 1, sq, v, i = -1;------// 27
----if (n < 4) return true;------// d9
                               ----vi primes;------// 8f
----if (n % 2 == 0 || n % 3 == 0) return false;-----// Of
----if (n < 25) return true;------// ef
                               ----bool* prime = new bool[mx + 1];------// ef
                               ----memset(prime, 1, mx + 1);-----// 28
----int s = static_cast<int>(sqrt(static_cast<double>(n)));------// 64
----for (int i = 5; i <= s; i += 6)-----// 6c
                               ----if (n >= 2) primes.push_back(2);------// f4
                               ----while (++i <= mx) if (prime[i]) {------// 73
------if (n % i == 0 || n % (i + 2) == 0) return false;------// e9
----return true; }-----// 43
                               -----primes.push_back(v = (i << 1) + 3);-----// be
                               -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
5.4. Miller-Rabin Primality Test.
                               ------for (int j = sq; j <= mx; j += v) prime[j] = false; }------// 2e
#include "mod_pow.cpp"-----// c7 ----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3);-----// 29
bool is_probable_prime(ll n, int k) {------// be
                               ----delete[] prime; // can be used for O(1) lookup-----// 36
----if (~n & 1) return n == 2;------// d1
                               ----return primes; }------// 72
----if (n <= 3) return n == 3;-----// 39
----while (k--) {------------------------// c8 ----vi minimalDiv(n+1, 2), primes;--------// 37
```

```
----for(int k=3;k<=n;k+=2) {------//5d
                                                       5.12. Numeric Integration.
------if(minimalDiv[k] == k) primes.push_back(k);------// 75
                                                       double integrate(double (*f)(double), double a, double b,-----// 76
-----rep(i, 1, size(primes))------// 49
                                                        -----double delta = 1e-6) {------// c0
-----if(primes[i] > minimalDiv[k] || primes[i]*k > n) break;-----// 53
                                                       ----if (abs(a - b) < delta)-------// 38
------else minimalDiv[primes[i]*k] = primes[i];------// 90
                                                       -----return (b-a)/8 *-----// 56
----}-----// 9d
                                                        ----return primes: }-----// 93
                                                        ----return integrate(f, a,-----// 64
                                                       -----(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);------// \theta c
                                                       }-----// 4b
5.8. Modular Multiplicative Inverse.
#include "egcd.cpp"-----// 55
                                                       5.13. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete
-----// e8
                                                       Fourier transform. The fft function only supports powers of twos. The czt function implements the
int mod_inv(int a, int m) {------// 49
                                                       Chirp Z-transform and supports any size, but is slightly slower.
----int x, y, d = eqcd(a, m, x, y);------// 3e
                                                       #include <complex>-----// 8e
----if (d != 1) return -1:------// 20
                                                       typedef complex<long double> cpx;-----// 25
----return x < 0 ? x + m : x;-----// 3c
                                                       // NOTE: n must be a power of two-----// 14
}-----// 69
                                                       void fft(cpx *x, int n, bool inv=false) {------// 36
                                                        ----for (int i = 0, j = 0; i < n; i++) {-------// f9
5.9. Primitive Root.
                                                        -----if (i < j) swap(x[i], x[i]);-----// 44
#include "mod_pow.cpp"-----// c7
                                                       ------int m = n>>1;------// 9c
ll primitive_root(ll m) {------// 8a
                                                       ------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
----vector<ll> div;------// f2
                                                       -----j += m:------// 11
----for (ll i = 1; i*i <= m-1; i++) {------// ca
                                                       ----}------// d0
-----if ((m-1) % i == 0) {------// 85
                                                       ----for (int mx = 1; mx < n; mx <<= 1) {------// 15
------if (i < m) div.push_back(i);-----// fd
                                                       ------if (m/i < m) div.push_back(m/i); } }-----// f2
                                                       -----for (int m = 0; m < mx; m++, w *= wp) {------// dc
----rep(x,2,m) {------// 57
                                                       ------for (int i = m; i < n; i += mx << 1) {------// 6a
-----bool ok = true:-----// 17
                                                       -----cpx t = x[i + mx] * w;-----// 12
-----iter(it.div) if (mod_pow(x, *it. m) == 1) { ok = false: break: }-----// a4
                                                       -----x[i + mx] = x[i] - t:-----// 73
-----if (ok) return x; }------// 55
                                                       -----x[i] += t:-----// 0e
----return -1; }------// 15
                                                       ------}------// a4
5.10. Chinese Remainder Theorem.
                                                        ----}------// bf
#include "eacd.cpp"------// 55
                                                       ----if (inv) rep(i,0,n) x[i] /= cpx(n);------// 16
int crt(const vi& as, const vi& ns) {------// c3 }-----// 1c
----rep(i,0,cnt) N *= ns[i];------// b1 ----int len = 2*n+1;-----// bc
----return mod(x, N); }------// b2 ----len <<= 1;------// b2
                                                        ----cpx w = \exp(-2.0 L * pi / n * cpx(0,1)),-----// 45
5.11. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
                                                       -----*c = new cpx[n], *a = new cpx[len],-----// 4e
                                                        -----*b = new cpx[len];-----// 30
#include "egcd.cpp"------// 55 ----rep(i,0,n) c[i] = pow(w, (inv ? -1.0 : 1.0)*i*i/2);------// 9e
vi linear_congruence(int a, int b, int n) {------// c8 ----rep(i,0,n) a[i] = x[i] * c[i], b[i] = 1.0L/c[i];-----// e9
----int x, y, d = egcd(a, n, x, y); d = egcd(a, x, y); 
}------if (inv) x[i] /= cpx(n);-------// b1
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}------// c6 ----rep(i,0,l<<1) T2[i] = T1[i] * T2[i];-------// 9f
                                                   ----ntt(T2, l<<1, true);------// eb
                                                   5.14. Number-Theoretic Transform.
#include "../mathematics/primitive_root.cpp"------8
                                                   5.15. Tridiagonal Matrix Algorithm. Solves a tridiagonal system of linear equations a_i x_{i-1} +
int mod = 998244353, g = primitive_root(mod),..../9c
                                                   b_i x_i + c_i x_{i+1} = d_i where a_1 = c_n = 0. Beware of numerical instability.
----ginv = mod_pow(g, mod-2, mod), inv2 = mod_pow(2, mod-2, mod);-----// 75
                                                   #define MAXN 5000-----// f7
#define MAXN (1<<22)-----// 94
                                                   long double A[MAXN], B[MAXN], C[MAXN], D[MAXN], X[MAXN];-----// d8
struct Num {-----// c5
                                                   void solve(int n) {------// θ1
                                                   ----C[0] /= B[0]; D[0] /= B[0];-----// 94
----rep(i,1,n-1) C[i] /= B[i] - A[i]*C[i-1];------// 6b
----Num operator +(const Num &b) { return x + b.x; }-----// 08
                                                   ----rep(i,1,n) D[i] = (D[i] - A[i] * D[i-1]) / (B[i] - A[i] * C[i-1]); ------// 33
----Num operator -(const Num &b) const { return x - b.x: }------// 89
----Num operator *(const Num &b) const { return (ll)x * b.x; }------// e3
                                                   ----for (int i = n-2; i>=0; i--) X[i] = D[i] - C[i] * X[i+1]; }------// ad
----Num operator /(const Num &b) const { return (ll)x * b.inv().x; }------// 2a
----Num inv() const { return mod_pow((ll)x, mod-2, mod); }-----// d3
                                                   5.16. Mertens Function. Mertens function is M(n) = \sum_{i=1}^n \mu(i). Let L \approx (n \log \log n)^{2/3} and the
----Num pow(int p) const { return mod_pow((ll)x, p, mod); }-----// d5
                                                   algorithm runs in O(n^{2/3}). Can also be easily changed to compute the summatory \Phi.
} T1[MAXN], T2[MAXN];-----// d5
void ntt(Num x[], int n, bool inv = false) {-----// 24
                                                   #define | 9000000-----// 27
----Num z = inv ? qinv : q;-----// 00
                                                   int mob[L], mer[L];-----// f1
---z = z.pow((mod - 1) / n);
                                                   unordered_map<ll,ll> mem;-----// 30
----for (ll i = 0, j = 0; i < n; i++) {------// fe
                                                   ll M(ll n) {-----// de
-----if (i < j) swap(x[i], x[j]);-----// eb
                                                   ----if (n < L) return mer[n];------// 1c
                                                   ----if (mem.find(n) != mem.end()) return mem[n];------// 79
------while (1 <= k && k <= j) j -= k, k >>= 1;------// d2
                                                   ----ll ans = 0, done = 1;-----// 48
-----i += k: }-----// c4
                                                   ----for (ll i = 2; i*i <= n; i++) ans += M(n/i), done = i; ------// 41
----for (int mx = 1, p = n/2; mx < n; mx <<= 1, p >>= 1) {-------// a2
                                                   ----for (ll i = 1; i*i <= n; i++) ans += mer[i] * (n/i - max(done, n/(i+1))); --// 43
-----Num wp = z.pow(p), w = 1;------// 35
                                                   ----return mem[n] = 1 - ans: }------// c2
------for (int k = 0; k < mx; k++, w = w*wp) {------// 42
                                                   ------for (int i = k; i < n; i += mx << 1) {------// 9a
                                                   ----for (int i = 1; i < L; i++) mer[i] = mob[i] = 1;-----// f7
-----Num t = x[i + mx] * w;-----// 22
                                                   ----for (int i = 2; i < L; i++) {------// 8e
-----x[i + mx] = x[i] - t;
                                                   ------if (mer[i]) {-------// 8b
-----x[i] = x[i] + t; } }-----// 1e
                                                   -----mob[i] = -1:-----// e5
----if (inv) {--------// 6c
                                                   ------for (int j = i+i; j < L; j += i)-----// f0
-----Num ni = Num(n).inv();-----// 18
                                                   -------mer[j] = 0, mob[j] = (j/i)%i == 0 ? 0 : -mob[j/i]; ------// 26
-----rep(i,0,n) { x[i] = x[i] * ni; } } } ------/72
                                                   -----}----------------------// aa
void inv(Num x[], Num y[], int l) {-----// ee
                                                   -----mer[i] = mob[i] + mer[i-1]; } }-----// 3b
----if (l == 1) { y[0] = x[0].inv(); return; }-------// 71
----inv(x, y, l>>1);------// bd
                                                   5.17. Markov Chains. A Markov Chain can be represented as a weighted directed graph of states,
----// NOTE: maybe l<<2 instead of l<<1-----// 24
                                                   where the weight of an edge represents the probability of transitioning over that edge in one timestep.
----rep(i,l>>1,l<<1) T1[i] = y[i] = 0;------// dd
                                                   Let P^{(m)} = (p_{ij}^{(m)}) be the probability matrix of transitioning from state i to state j in m timesteps, and
----rep(i,0,l) T1[i] = x[i];-----// f7
                                                   note that P^{(1)} is the adjacency matrix of the graph. Chapman-Kolmogorov: p_{ij}^{(m+n)} = \sum_k p_{ik}^{(m)} p_{ki}^{(n)}
----ntt(T1, l<<1); ntt(y, l<<1);------// a5
                                                   It follows that P^{(m+n)} = P^{(m)}P^{(n)} and P^{(m)} = P^m. If p^{(0)} is the initial probability distribution (a
----rep(i,0,l<<1) y[i] = y[i]*2 - T1[i] * y[i] * v[i];-----// bc
                                                   vector), then p^{(0)}P^{(m)} is the probability distribution after m timesteps.
----ntt(y, l<<1, true); }------// 40
void sqrt(Num x[], Num y[], int l) {-----// 44
                                                    The return times of a state i is R_i = \{m \mid p_{ii}^{(m)} > 0\}, and i is aperiodic if gcd(R_i) = 1. A MC is
----if (l == 1) { assert(x[0].x == 1); y[0] = 1; return; }------// 64
                                                   aperiodic if any of its vertices is aperiodic. A MC is irreducible if the corresponding graph is strongly
```

----sqrt(x, y, l>>1);-----// a7 connected.

A distribution  $\pi$  is stationary if  $\pi P = \pi$ . If MC is irreducible then  $\pi_i = 1/\mathbb{E}[T_i]$ , where  $T_i$  is the expected time between two visits at i.  $\pi_j/\pi_i$  is the expected number of visits at j in between two consecutive visits at i. A MC is ergodic if  $\lim_{m\to\infty} p^{(0)}P^m = \pi$ . A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (unweighted graph can be made weighted by adding 1-weights) has  $p_{uv} = w_{uv}/\sum_x w_{ux}$ . If the graph is connected, then  $\pi_u = \sum_x w_{ux}/\sum_v \sum_x w_{vx}$ . Such a random walk is aperiodic iff. the graph is not bipartite.

An absorbing MC is of the form  $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$ . Let  $N = \sum_{m=0}^{\infty} Q^m = (I_t - Q)^{-1}$ . Then, if starting

in state i, the expected number of steps till absorption is the i-th entry in N1. If starting in state i, the probability of being absorbed in state j is the (i, j)-th entry of NR.

Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

5.18. Burnside's Lemma. Let G be a finite group that acts on a set X. For each g in G let  $X^g$  denote the set of elements in X that are fixed by g. Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

5.19. **Bézout's identity.** If (x, y) is any solution to ax+by=d (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a,b)}, y - k \frac{a}{\gcd(a,b)}\right)$$

#### 5.20. Formulas.

- Number of permutations of n objects, where there are  $n_1$  objects of type 1,  $n_2$  objects of type 2, ...,  $n_k$  objects of type k:  $\binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times:  $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to  $x_1 + x_2 + \cdots + x_n = k$  where  $x_i > 0$ :  $f_k^n$
- Number of strings with n sets of brackets such that the brackets are balanced:  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an  $n \times n$  lattice which do not rise above the main diagonal:
- Number of trees on n labeled vertices:  $n^{n-2}$
- Number of permutations of n objects with exactly k ascending sequences or runs:

$$\left\langle {n\atop k}\right\rangle = \left\langle {n\atop n-k-1}\right\rangle = k\left\langle {n-1\atop k}\right\rangle + (n-k+1)\left\langle {n-1\atop k-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = 1$$

- Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements):  $D_0 = 1, D_1 = 0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points:  $\binom{n}{k}D_{n-k}$
- Number of trees on n labeled vertices:  $n^{n-2}$
- Jacobi symbol:  $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- Heron's formula: A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} 1$ .

- Divisor sigma: The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where  $n = \prod_{i=0}^r p_i^{a_i}$  is the prime factorization.
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$  where each p is a distinct prime factor of n.
- **König's theorem:** In any bipartite graph  $G = (L \cup R, E)$ , the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L, and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then  $K = (L \setminus Z) \cup (R \cap Z)$  is the minimum vertex cover.
- A minumum Steiner tree for n vertices requires at most n-2 additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a
  maximum independent set.
- $\gcd(n^a 1, n^b 1) = n^{\gcd(a,b)} 1$
- Wilson's theorem:  $(n-1)! \equiv -1 \pmod{n}$  iff. n is prime
- Lagrange polynomial through points  $(x_0, y_0), \ldots, (x_k, y_k)$  is  $L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \le m \le k \\ m \ne j}} \frac{x x_m}{x_j x_m}$
- $\sum_{k=0}^{m} (-1)^k \binom{n}{k} = (-1)^m \binom{n-1}{m}$
- $2^{\omega(n)} = O(\sqrt{n})$ , where  $\omega(n)$  is the number of distinct prime factors
- $\bullet \ \sum_{i=1}^n 2^{\omega(i)} = O(n \log n)$
- Hook length formula: If  $\lambda$  is a Young diagram and  $h_{\lambda}(i,j)$  is the hook-length of cell (i,j), then then the number of Young tableux  $d_{\lambda} = n! / \prod h_{\lambda}(i,j)$ .
- Möbius inversion formula: If  $f(n) = \sum_{d|n} g(d)$ , then  $g(n) = \sum_{d|n} \mu(d) f(n/d)$

5.21. **Numbers and Sequences.** Some random prime numbers: 1031, 32771, 1048583, 33554467, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.

## 6. Geometry

# 6.1. Primitives.

```
#define P(p) const point &p-----// 2e
#define L(p0, p1) P(p0), P(p1)-----// cf
#define C(p0, r) P(p0), double r-----// f1
#define PP(pp) pair<point,point> &pp-----// e5
typedef complex<double> point;------// 6a
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// d2
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// 8a
point rotate(P(p), double radians = pi / 2, P(about) = point(0,0)) \{-----//23\}
----return (p - about) * exp(point(0, radians)) + about; }-----// 25
point reflect(P(p), L(about1, about2)) {-----// 50
----point z = p - about1, w = about2 - about1;------// 8b
----return conj(z / w) * w + about1; }-----// 83
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// e7
point normalize(P(p), double k = 1.0) {------// 5f
----return abs(p) == 0 ? point(0,0) : p / abs(p) * k; }------// 4a
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// 27
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// b3
double angle(P(a), P(b), P(c)) {------// 61
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }-----// c7
double signed_angle(P(a), P(b), P(c)) {------// 4a
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 40
double angle(P(p)) { return atan2(imag(p), real(p)); }-----// e6
point perp(P(p)) {        return point(-imag(p), real(p));    }---------------// d9
```

```
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-----return (imag(p) - imag(a)) / (imag(b) - imag(a));------// 5e ----res1 = A + v + u, res2 = A + v - u;------// 24
-----// 53 }------// bb
-----// 46
                                             -----// ce
                                              int intersect(L(A, B), C(0, r), point & res1, point & res2) {------// 4d
6.2. Lines.
                                              ---- double h = abs(0 - closest_point(A, B, 0));-----// d0
#include "primitives.cpp"-----// e0
                                              ---- if(r < h - EPS) return 0;-----// af
-----// 85
                                              ---- point H = proj(0 - A, B - A) + A, v = normalize((B - A), sqrt(r*r - h*h));// 1a
bool collinear(L(a, b), L(p, q)) {-----// 2f
                                              ---- res1 = H + v; res2 = H - v;-----// fc
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// 3e
                                              ---- if(abs(v) < EPS) return 1; return 2;-----// d1
bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// 8d
                                             }-----// f8
point closest_point(L(a, b), P(c), bool segment = false) {------// f2
                                              .
-----// 74
----if (seament) {-------// f4
                                             int tangent(P(A), C(0, r), point & res1, point & res2) {-----// ec
-----if (dot(b - a, c - b) > 0) return b;-----// 88
                                             ----point v = 0 - A; double d = abs(v);------// 62
-----if (dot(a - b, c - a) > 0) return a;-----// 75
                                              ----if (d < r - EPS) return 0;-----// fc
----}-----// ce
                                             ----double alpha = asin(r / d), L = sqrt(d*d - r*r);------// f0
----double t = dot(c - a, b - a) / norm(b - a);-----// 62
                                              ----v = normalize(v, L);-----// 0f
----return a + t * (b - a);------// 6e
                                              ----res1 = A + rotate(v, alpha); res2 = A + rotate(v, -alpha);------// 9e
}-----// 8c
                                              ----if (abs(r - d) < EPS || abs(v) < EPS) return 1;------// ee
double line_segment_distance(L(a,b), L(c,d)) {------// f3
                                              ----return 2:-----// 4f
----double x = INFINITY:-----// 64
                                             }-----// 3c
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// a5
                                             .....// 29
----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true)); -----// 23
                                              void tangent_outer(point A, double rA, point B, double rB, PP(P), PP(Q)) {-----// 9e
----else if (abs(c - d) < EPS) \times = abs(c - closest_point(a, b, c, true)); -----// 53
                                              ----if (rA - rB > EPS) { swap(rA, rB); swap(A, B); }------// 9c
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// 6d
                                              ----double theta = asin((rB - rA)/abs(A - B));------// 68
----point v = rotate(B - A, theta + pi/2), u = rotate(B - A, -(theta + pi/2)); -// d7
----else {------// e1
                                              ----u = normalize(u, rA):-----// 87
-----x = min(x, abs(a - closest_point(c,d, a, true)));------// 29
                                              ----P.first = A + normalize(v, rA); P.second = B + normalize(v, rB); ------// 2\theta
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// fe
                                             ----Q.first = A + normalize(u, rA); Q.second = B + normalize(u, rB); ------// 52
-----x = min(x, abs(c - closest_point(a,b, c, true)));------// 81
                                             }-----// 63
-----x = min(x, abs(d - closest_point(a,b, d, true)));------// e4
----}------// c5
                                              6.4. Polygon.
----return x:-----// b7
                                             #include "primitives.cpp"-----// e0
}-----// 27
                                             typedef vector<point> polygon;-----// b3
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) \{-----//d2\}
                                              double polygon_area_signed(polygon p) {-----// 31
----// NOTE: check for parallel/collinear lines before calling this function---// 1b
                                              ----double area = 0; int cnt = size(p);-----// a2
----point r = b - a, s = q - p;------// 34
                                              ----rep(i,1,cnt-1) area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 51
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// \theta b
                                              ----return area / 2; }------// 66
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// e4
                                              double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// a4
-----return false;-----// e3
                                             #define CHK(f,a,b,c) (f(a) < f(b) \&\& f(b) <= f(c) \&\& ccw(a,c,b) < 0)-----//8f
----res = a + t * r:-----// 47
                                             int point_in_polygon(polygon p, point q) {------// 5d
----return true:------// 05
                                              ----int n = size(p); bool in = false; double d;------// 69
}-----// 44
                                              ----for (int i = 0, j = n - 1; i < n; j = i++)-----// f3
-----// cc
                                              -----if (collinear(p[i], q, p[j]) &&-----// 9d
6.3. Circles.
                                              ------0 <= (d = progress(q, p[i], p[j])) && d <= 1)------// 4b
#include "primitives.cpp"-----// e0 -----return 0:-----// b3
int intersect(C(A, rA), C(B, rB), point & res1, point & res2) {--------// 3b ----for (int i = 0, j = n - 1; i < n; j = i++)-------// 67</pre>
----double d = abs(B - A);------// 5c ------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))------// b4
```

```
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// pair<polygon, polygon cut_polygon(const polygon &poly, point a, point b) {-// 0d ------A = B = c; double p = progress(c, a,b);-------------// c0
//------ int j = i == cnt-1? 0: i + 1;-------// 02 -------double ap = progress(a, c,d), bp = progress(b, c,d);------// 34
//----- if (ccw(a, b, p) <= 0) left.push_back(p);-----// 8d ------if (bp < 0.0 || ap > 1.0) return false;-----// 14
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// 43 ------A = c + max(ap, 0.0) * (d - c);-------// 14
//----- // myintersect = intersect where------// ba ------B = c + min(bp, 1.0) * (d - c);
//-----// (a,b) is a line, (p,q) is a line segment-----// 7e ------return true;
//----- if (myintersect(a, b, p, q, it))------// 6f ---else if (parallel(a,b, c,d)) return false;-------// fc
//-----------left.push_back(it), right.push_back(it);-------// 8a ----else if (intersect(a,b, c,d, A, true)) {--------// 78
//----} -------B = A; return true; }-------// 7d
//--- return pair<polygon, polygon>(left, right);------// 3d ----return false;-------------------------------// 4b
-----// 83
6.5. Convex Hull. NOTE: Doesn't work on some weird edge cases. (A small case that included
three collinear lines would return the same point on both the upper and lower hull.)
                                      6.7. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
#include "polygon.cpp"-----// 58
                                     coordinates) on a sphere of radius r.
#define MAXN 1000-----// 09 double qc_distance(double pLong,-----// 7b
point hull[MAXN];----- double qLong, double r) {------// a4
bool cmp(const point &a, const point &b) {------// 32 ----pLat *= pi / 180; pLong *= pi / 180;------// ee
----return abs(real(a) - real(b)) > EPS ?------// 44 ----qLat *= pi / 180; qLong *= pi / 180;------// 75
------real(a) < real(b) : imag(a) < imag(b); }-------// 40 ----return r * acos(cos(pLat) * cos(pLong - gLong) +-----// e3
int convex_hull(polygon p) {------// cd ------sin(pLat) * sin(qLat));------// le
----rep(i,0,n) {------// e4
-----if (i > 0 && p[i] == p[i - 1]) continue;-----// c7
                                     6.8. Triangle Circumcenter. Returns the unique point that is the same distance from all three
                                      points. It is also the center of the unique circle that goes through all three points.
------while (l \ge 2 \& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;------// 62
                                      #include "primitives.cpp"-----// e0
------hull[l++] = p[i];-----// bd
                                      point circumcenter(point a, point b, point c) {-----// 76
----}-----// d2
                                     ----b -= a, c -= a;-----// 41
----int r = 1:-----// 30
----for (int i = n - 2; i >= 0; i--) {------// 59
                                      ----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a
-----if (p[i] == p[i + 1]) continue;-----// af
                                      }-----// c3
------while (r - l) = 1 \& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 4d
                                      6.9. Closest Pair of Points.
-----hull[r++] = p[i];-----// f5
                                     #include "primitives.cpp"-----// e0
----}-----------// f6
                                        -----// 85
----return l == 1 ? 1 : r - 1;------// a6
}-----// 6d
                                      struct cmpx { bool operator ()(const point &a, const point &b) {------// 01
                                      -----return abs(real(a) - real(b)) > EPS ?-----// e9
6.6. Line Segment Intersection.
                                      -----real(a) < real(b) : imag(a) < imag(b); } };------// 53
#include "primitives.cpp"------// e0 struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
#include "lines.cpp"------// 54 ----return abs(imag(a) - imag(b)) > EPS ?------// 0b
------A = B = a; return abs(a - d) < EPS; }-------// 37 ----sort(pts.begin(), pts.end(), cmpx());-------// 0c
------A = B = a; double p = progress(a, c,d);---------// 6c ----set<point, cmpy>::const_iterator it, jt;--------// a6
```

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-------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);-------// 8b ------return abs(x) < EPS && abs(y) < EPS; }-------// 15
-----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));-------// 39 -----return ((A - *this) * (B - *this)).isZero(); }-------// 58
------cur.insert(pts[i]); }------// 82 -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < EPS; }-----// d9
-----return isOnLine(A, B) && ((A - *this) % (B - *this)) < -EPS; }------// ba
                                        ----double getAngle() const {------// 0f
6.10. 3D Primitives.
                                        -----return atan2(y, x); }------// 40
#define P(p) const point3d &p-----// a7
                                        ----double getAngle(P(u)) const {------// d5
#define L(p0, p1) P(p0), P(p1)-----// Of
                                        -----return atan2((*this * u).length(), *this % u); }------// 79
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)-----// 67
                                        ----bool isOnPlane(PL(A, B, C)) const {------// 8e
struct point3d {-----// 63
                                        -----return abs((A - *this) * (B - *this) * (C - *this)) < EPS; } };-----// 74
----double x, y, z;-----// e6
                                        int line_line_intersect(L(A, B), L(C, D), point3d &0){------// dc
----point3d() : x(0), y(0), z(0) {}-----// af
                                        ----if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;-------// 6a
----point3d(double _x, double _y, double _z) : x(_x), y(_y), z(_z) {}------// fc
                                        ----if (((A - B) * (C - D)).length() < EPS)-----// 79
----point3d operator+(P(p)) const {------// 17
                                        -----return A.isOnLine(C, D) ? 2 : 0;-----// 09
-----return point3d(x + p.x, y + p.y, z + p.z); }-----// 8e
                                        ----point3d normal = ((A - B) * (C - B)).normalize();-----// bc
----point3d operator-(P(p)) const {------// fb
                                        ----double s1 = (C - A) * (D - A) % normal;------// 68
-----return point3d(x - p.x, y - p.y, z - p.z); }-----// 83
                                        ----0 = A + ((B - A) / (s1 + ((D - B) * (C - B) % normal))) * s1;------// 56
----point3d operator-() const {------// 89
-----return point3d(-x, -y, -z); }-----// d4
                                        int line_plane_intersect(L(A, B), PL(C, D, E), point3d & O) {------// 09
----point3d operator*(double k) const {------// 4d
                                        ----double V1 = (C - A) * (D - A) % (E - A);------// c1
-----return point3d(x * k, y * k, z * k); }------// fd
                                        ----double V2 = (D - B) * (C - B) % (E - B);------// 29
----point3d operator/(double k) const {------// 95
                                        ----if (abs(V1 + V2) < EPS)-------// 81
-----return point3d(x / k, y / k, z / k); }-----// 58
                                        -----return A.isOnPlane(C, D, E) ? 2 : 0;------// d5
----double operator%(P(p)) const {------// d1
                                        ----0 = A + ((B - A) / (V1 + V2)) * V1;-----// 38
-----return x * p.x + y * p.y + z * p.z; }------// 09
                                        ----return 1; }-----// ce
----point3d operator*(P(p)) const {------// 4f
                                        bool plane_plane_intersect(P(A), P(nA), P(B), P(nB), point3d &P, point3d &O) {-// 5a
-----return point3d(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }-----// ed
                                        ----point3d n = nA * nB;-----// 49
----double length() const {------// 3e
                                        ----if (n.isZero()) return false;------// 03
-----return sqrt(*this % *this); }-----// 05
                                        ----point3d v = n * nA;-----// d7
----double distTo(P(p)) const {------// dd
                                        ----P = A + (n * nA) * ((B - A) % nB / (v % nB)); ------// 1a
------return (*this - p).length(); }-----// 57
                                        ---0 = P + n:------// 9c
----double distTo(P(A), P(B)) const {------// bd
                                        ----return true; }------// 1a
-----// A and B must be two different points-----// 4e
-----return ((*this - A) * (*this - B)).length() / A.distTo(B); }------// 6e
----point3d normalize(double k = 1) const {------// db
                                        6.11. Polygon Centroid.
-----// length() must not return 0-----// 3c
------/return (*this) * (k / length()); }--------// d4 #include "polygon.cpp"-------// 58
return (*this) * normal; }------// 5c -----mnx = min(mnx, real(p[i])),------// 6c
------point3d Z = axe.normalize(axe % (*this - 0));-------// ba ----rep(i,0,n) {-----------------------------// 3c
------return 0 + Z + (*this - 0 - Z).rotate(alpha, 0); }-------// 38 ------int j = (i + 1) % n;----------------// 5b
```

```
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----return point(cx, cy) / 6.0 / polygon_area_signed(p) + point(mnx, mny); }---// a1 //------ b = (b+1) % h;---------// 96
                                    //----- B.move_to(hull[b]);-----// 38
                                    //-----}
6.12. Rotating Calipers.
                                    //----- done += min(tha, thb);-----// d2
#include "primitives.cpp"-----// e0
                                    //----- if (done > pi) {-----// c2
struct caliper {-----// 8e
                                    //----- break;-----// e8
---ii pt:-----// 05
                                    //----- }-----// 37
----double angle;-----// d4
                                    //---- }------// ac
----caliper(ii _pt, double _angle) : pt(_pt), angle(_angle) { }------// 35
                                    // }-----// 9c
----double angle_to(ii pt2) {------// 8b
-----double x = angle - atan2(pt2.second - pt.second, pt2.first - pt.first);// 1e
                                    6.13. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
• a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
------while (x <= -pi) x += 2*pi;------// a3
                                      • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
-----return x; }-----// 7d
                                      • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
----void rotate(double by) {------// 57
                                       of that is the area of the triangle formed by a and b.
-----angle -= by;-----// 5d
                                      • Euler's formula: V - E + F = 2
------while (angle < 0) angle += 2*pi;-----// 03
                                      • Side lengths a, b, c can form a triangle iff. a + b > c, b + c > a and a + c > b.
• Sum of internal angles of a regular convex n-gon is (n-2)\pi.
----void move_to(ii pt2) { pt = pt2; }-----// 37
----double dist(const caliper &other) {------// 68
                                                  7. Other Algorithms
------point a(pt.first,pt.second),-----// d7
------ c(other.pt.first, other.pt.second);------// 71 #include "../graph/scc.cpp"-----// c3
-----// c5 ----all_truthy.clear();-------// 31
// int h = convex_hull(pts);-----// 9c ----vvi adj(2*n+1);-----// 7b
// double mx = 0;-----// f1 ----rep(i,0,size(clauses)) {------// 76
// if (h > 1) {------adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// eb
//--- int a = 0,-----// e6 ------if (clauses[i].second)-----// bc
//------b = 0;-------b].second + n].push_back(clauses[i].first + n);-----// f0
//---- rep(i.0,h) {------// ld ---}----// da
//------ if (hull[i].first < hull[a].first)------// ac ----pair<union_find, vi> res = scc(adj);------// 00
//----- if (hull[i].first > hull[b].first)------// 02 ----vi dag = res.second;------// ed
//------b = i;------// 84 ----vi truth(2*n+1, -1);------// c7
//---- caliper A(hull[a], pi/2), B(hull[b], 3*pi/2);------// 60 ------int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);-// 4f
//--- double done = 0;-----// 3c -----if (cur == 0) continue;-----// cd
//--- while (true) {------// 31 ------if (p == o) return false;-----// d0
//----- mx = max(mx, abs(point(hull[a].first,hull[a].second)-----// e3 ------if (truth[p] == -1) truth[p] = 1;------// d3
//-----truth[cur + n] = truth[p];------------// 50
//------ double tha = A.angle_to(hull[(a+1)%h]),------// 57 -----truth[o] = 1 - truth[p];------
//-----thb = B.angle_to(hull[(b+1)%h]);------// f1 ------if (truth[p] == 1) all_truthy.push_back(cur);-----// 55
//----- if (tha <= thb) {------------------------// 91 ...}
//------ A.rotate(tha);------// c9 ----return true;---------------------------// eb
//------- A.move_to(hull[a]);-----------------------// b3 7.2. Stable Marriage.
```

```
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----vi at(n, θ), eng(n, -1), res(n, -1); vvi inv(n, vi(n));--------// c3 ------++ni;-----------+
----while (!q.empty()) {------// 68 ------ptr[ni][j]->u = ptr[i][j];------// 66
------int curm = q.front(); q.pop();--------// e2 -------while (true) {------------------// 7f
------for (int &i = at[curm]; i < n; i++) {--------// de
-----int curw = m[curm][i];-------// 95 ------if (i == rows || arr[i][nj]) break;------// 4c
-----if (eng[curw] == -1) { }-------// f7 ------++nj;------------// c5
-----q.push(eng[curw]);------// 2e ------ptr[i][j]->r = ptr[i][nj];-------// 60
-----else continue;------// 1d -------ptr[i][nj]->l = ptr[i][j];--------// 82
------head->l = ptr[rows][cols - 1];------// 6a
                                  -----ptr[rows][cols - 1]->r = head;-----// c1
7.3. Algorithm X.
                                  -----rep(j,0,cols) {------// 92
bool handle_solution(vi rows) { return false; }------// 63
                                  -----int cnt = -1;------// d4
struct exact_cover {------// 95
                                  -----rep(i,0,rows+1)-----// bd
----struct node {------// 7e
                                  ------if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];------// f3
-----node *l, *r, *u, *d, *p;-----// 19
                                  -----ptr[rows][j]->size = cnt;-----// c2
------int row, col, size;-----// ae
                                  -----node(int _row, int _col) : row(_row), col(_col) {------// c9
                                  -----rep(i,0,rows+1) delete[] ptr[i];-----// a5
-----size = 0; l = r = u = d = p = NULL; }-----// c3
                                  -----delete[] ptr;-----// 72
----}:-----// c1
                                  ----int rows, cols, *sol;------// 7b
                                  ----#define COVER(c, i, j) \------// 91
----bool **arr:-----// e6
                                  ----node *head:-----// fe
                                  ------for (node *i = c->d; i != c; i = i->d) \------// 62
----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
----arr = new bool*[rows];-----//
                                  ----j->d->u = j->u, j->u->d = j->d, j->p->size--;-----// c1
-----sol = new int[rows];-----// 5f
-----rep(i,0,rows)------// 9b
                                 ----#define UNCOVER(c, i, j) \\-----// 89
------arr[i] = new bool[cols], memset(arr[i], 0, cols);------// dd
                                 ------for (node *i = c->u; i != c; i = i->u) \------// f0
                                  ----void set_value(int row, int col, bool val = true) { arr[row][col] = val; }-// 9e
                                  -----j->p->size++, j->d->u = j->u->d = j; \\------// 65
----void setup() {------// a3
                                  ------C->r->l = C->l->r = C:-----// 0e
-----node ***ptr = new node**[rows + 1];------// bd
                                  ----bool search(int k = 0) {------// f9
-----rep(i,0,rows+1) {------// 76
                                  -----if (head == head->r) {------// 75
-----ptr[i] = new node*[cols];-----// eb
                                  -----vi res(k);-----// 90
-----rep(j,0,cols)-----// cd
                                  -----rep(i,0,k) res[i] = sol[i];-----// 2a
-----if (i == rows || arr[i][j]) ptr[i][j] = new node(i, j);-----// 16
                                  -----sort(res.begin(), res.end());-----// 63
-----else ptr[i][j] = NULL;-----// d2
                                  -----return handle_solution(res);-----// 11
                                  -----rep(i,0,rows+1) {------// fc
                                  -----node *c = head->r, *tmp = head->r;-----// a3
-----rep(j,0,cols) {------// 51
                                  -----for (; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 41
-----if (!ptr[i][j]) continue;-----// f7
                                  -----if (c == c->d) return false;-----// 02
-----/int ni = i + 1, nj = j + 1;-----// 7a
                                  -----COVER(c, i, j);-----// f6
-----while (true) {------// fc
                                  ------bool found = false;-----// 8d
------for (node *r = c->d; !found && r != c; r = r->d) {------// 78
-----if (ni == rows || arr[ni][j]) break;------// 8d
```

```
-----sol[k] = r->row;------// c0 7.7. Simulated Annealing. An example use of Simulated Annealing to find a permutation of length
------for (node *j = r->r; j != r; j = j->r) { COVER(j->p, a, b); }-----// f9 n that maximizes \sum_{i=1}^{n-1} |p_i - p_{i+1}|.
-----found = search(k + 1);-----// fb
                                            double curtime() { return static_cast<double>(clock()) / CLOCKS_PER_SEC; }-----// 9d
-----for (node *j = r->l; j != r; j = j->l) { UNCOVER(j->p, a, b); }----// 87
                                            int simulated_annealing(int n, double seconds) {------// 54
------}------// 7c
                                            ----default_random_engine rng;------// 67
------UNCOVER(c, i, j);------// a7
                                            ----uniform_real_distribution<double> randfloat(0.0, 1.0);-----// ed
-----return found:-----// c0
                                            ----uniform_int_distribution<int> randint(0, n - 2);------// bb
----}------// d2
                                            -----// 88
}:-----// d7
                                            ----// random initial solution------// 22
                                             ----vi sol(n);------// 33
7.4. nth Permutation.
                                             ----rep(i,0,n) sol[i] = i + 1;------// ee
vector<int> nth_permutation(int cnt, int n) {------// 78
                                             ----random_shuffle(sol.begin(), sol.end()):-----// 1e
----vector<int> idx(cnt), per(cnt), fac(cnt);-----// 9e
                                             -----// 5b
----rep(i,0,cnt) idx[i] = i;-----// bc
                                             ----// initialize score------// 11
----rep(i,1,cnt+1) fac[i - 1] = n % i, n /= i;------// 2b
                                            ----int score = 0:------// 4d
----for (int i = cnt - 1; i >= 0; i--)-----// f9
                                             ----rep(i,1,n) score += abs(sol[i] - sol[i-1]);------// 74
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// ee
                                             -----// 25
----return per:-----// ab
                                            ----int iters = 0:------// 4d
                                            ----double T0 = 100.0, T1 = 0.001,-----// f4
                                                  progress = 0, temp = T0,-----// 8b
7.5. Cycle-Finding.
                                                  starttime = curtime();-----// a2
ii find_cycle(int x0, int (*f)(int)) {------// a5
                                             ----while (true) {------// db
----int t = f(x\theta), h = f(t), mu = \theta, lam = 1;------// 8d
                                             -----if (!(iters & ((1 << 4) - 1))) {------// e8
----while (t != h) t = f(t), h = f(f(h));-----// 79
                                             -----progress = (curtime() - starttime) / seconds;------// αθ
----h = x0: ------// 04
                                             ------temp = T0 * pow(T1 / T0, progress);-------temp = T0 * pow(T1 / T0, progress);---------------------------// 12
----while (t != h) t = f(t), h = f(h), mu++;-----// 9d
                                             -----if (progress > 1.0) break;-----// 48
----h = f(t):------// 00
----while (t != h) h = f(h), lam++;-----// 5e
                                                   -----// 88
----return ii(mu, lam);-----// b4
                                             -----// random mutation-----// 84
                                             -----int a = randint(rng):-----// f7
7.6. Dates.
                                             -----// compute delta for mutation-----// 4e
int intToDay(int jd) { return jd % 7; }------// 89
                                             -----int delta = 0:-----// 10
int dateToInt(int y, int m, int d) {------// 96
                                             -----if (a > 0) delta += abs(sol[a+1] - sol[a-1]) - abs<math>(sol[a] - sol[a-1]); -//21
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +------// a8
                                             -----if (a+2 < n) delta += abs(sol[a] - sol[a+2]) - abs<math>(sol[a+1] - sol[a+2]);
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1
                                               -----// 22
-----// maybe apply mutation-----// 4d
-----d - 32075;-----// e0
                                             ------if (delta >= 0 || randfloat(rng) < exp(delta / temp)) {-------// a6
}-----// fa
                                             -----swap(sol[a], sol[a+1]);--------------------------------// ce
void intToDate(int jd, int &y, int &m, int &d) {------// a1
                                             -----score += delta;-----// 64
----int x, n, i, j;------// 00
                                             -----//    if (score >= target)    return;----------------------//    a6
---x = jd + 68569;----// 11
                                             ----n = 4 * x / 146097;-----// 2f
                                             -----iters++;-----// 3c
---x = (146097 * n + 3) / 4;
---i = (4000 * (x + 1)) / 1461001;
                                             ----return score;------// d0
----x -= 1461 * i / 4 - 31;-----// 09
---j = 80 * x / 2447;-----// 3d
---d = x - 2447 * j / 80;
                                                             8. Useful Information
----x = j / 11;-----// b7
----m = i + 2 - 12 * x:-----// 82
                                            8.1. Tips & Tricks.
----v = 100 * (n - 49) + i + x
                                               • How fast does our algorithm have to be? Can we use brute-force?
}-----// af
                                               • Does order matter?
```

- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo  $m_1, m_2, \ldots, m_k$ , where  $m_1, m_2, \ldots, m_k$  are pairwise coprime integers, and find the real answer using CRT?
- Can we use exponentiation by squaring?

## 8.2. Fast Input Reading.

```
void readn(register int *n) {------// dc
----int sign = 1;------// 32
----register char c:-----// a5
----*n = 0:-----// 35
----while((c = qetc_unlocked(stdin)) != '\n') {------// f3
-----switch(c) {------// 0c
-----case '-': sign = -1; break;-----// 28
-----case ' ': qoto hell;-----// fd
-----case '\n': goto hell;-----// 79
-----default: *n *= 10; *n += c - '0'; break;-----// c0
-----}------------------------// 2d
----}------// c3
hell:----// ba
----*n *= sign:-----// a0
```

#### 8.3. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.

```
----int y = x \& -x, z = x + y;------// 12
----return z | ((x ^ z) >> 2) / y;-----// 97
```

#### 9. Misc

## 9.1. Debugging Tips.

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
  - Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps 1+eps).
  - Rounding negative numbers?
  - Outputting in scientific notation?
- Wrong Answer?
  - Integer overflow?
  - Think very carefully about boundaries of all input parameters

- Try out possible edge cases:
  - \*  $n = 0, n = -1, n = 1, n = 2^{31} 1$  or  $n = -2^{31}$
  - \* List is empty, or contains a single element
  - \* n is even, n is odd
  - \* Graph is empty, or contains a single vertex
  - \* Graph is a multigraph (loops or multiple edges)
  - \* Polygon is concave or non-simple
- Are multiple test cases being handled correctly? Try repeating the same test case many times.
- Is initial condition wrong for small cases?

### 9.2. Solution Ideas.

- Dynamic Programming
  - Drop a parameter, recover from others
  - Swap answer and a parameter
  - Parsing CFGs: CYK Algorithm
  - Optimizations
    - \* Convex hull optimization

```
\cdot \operatorname{dp}[i] = \min_{i < i} \{ \operatorname{dp}[j] + b[j] \times a[i] \}
```

- b[i] > b[i+1]
- · optionally  $a[i] \le a[i+1]$
- $O(n^2)$  to O(n)
- \* Divide and conquer optimization
  - $\cdot dp[i][j] = \min_{k < i} \{dp[i-1][k] + C[k][j]\}$
  - $A[i][j] \leq A[i][j+1]$
  - $O(kn^2)$  to  $O(kn\log n)$
  - · sufficient:  $C[a][c] + C[b][d] \le C[a][d] + C[b][c], a \le b \le c \le d$  (QI)
- \* Knuth optimization
  - $dp[i][j] = \min_{i < k < j} \{dp[i][k] + dp[k][j] + C[i][j]\}$
  - $A[i][j-1] \le A[i][j] \le A[i+1][j]$
  - $O(n^3)$  to  $O(n^2)$
  - · sufficient: QI and  $C[b][c] \leq C[a][d], a \leq b \leq c \leq d$
- Greedy
- Randomized
- - Use bitset (/64)
  - Switch order of loops (cache locality)
- - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm
  - Sqrt decomposition
  - Store  $2^k$  jump pointers
- Data structure techniques
  - Sart buckets
  - Store  $2^k$  jump pointers
  - $-2^k$  merging trick
- Counting
  - Inclusion-exclusion principle

- Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - \* Bipartite matching
    - \* Hall's marriage theorem
    - \* Stable Marriage
  - Cut vertex/bridge
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - \* Heavy-light decomposition
    - \* Centroid decomposition
    - \* Least common ancestor
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
  - \* Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence
  - Steiner Tree
  - Kirchoff's matrix tree theorem
  - Prüfer sequences
- Mathematics
  - Is the function multiplicative?
  - Look for a pattern
  - Permutations
    - \* Consider the cycles of the permutation
  - Functions
    - \* Sum of piecewise-linear functions is a piecewise-linear function
    - \* Sum of convex (concave) functions is convex (concave)
  - Modular arithmetic
    - \* Chinese Remainder Theorem
    - \* Linear Congruence
  - Sieve
  - System of linear equations
  - Values to big to represent?
    - \* Compute using the logarithm
    - \* Divide everything by some large value
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half  $(\log(n))$
- Strings
  - Trie (maybe over something weird, like bits)

- Suffix array
- Suffix automaton (+DP?)
- Aho-Corasick
- eerTree
- Work with S + S
- Hashing
- Euler tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Geometry
  - Minkowski sum (of convex sets)
  - Rotating calipers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform
- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem
  - Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

Reykjavík University 26

# Practice Contest Checklist.

- How many operations per second?
- What is the stack size?
- $\bullet\,$  How to use printf/scanf with long long/long double?
- Is \_\_int128 available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
  What is RAND\_MAX?