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```
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           1. Code Templates
                             ----public static void main(String[] args) throws Exception {-------// 02
                             -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                             ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                             -----// code-----// e6
setxkbmap -option caps:escape
                             -----out.flush():-----// 56
set -o vi
                             xset r rate 150 100
                             }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                        2. Data Structures
syn on | colorscheme slate
                             2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                             struct union find {-----// 42
#include <cassert>-----------------// 65 ----int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }-------// ba
#include <cmath>-----// 7d ----bool unite(int x, int y) {-------// 6c
#include <cstdio>------// 2e ------int xp = find(x), yp = find(y);------// 64
#include <cstdlib>------// 11 ------if (xp == yp) return false;------// 0b
#include <cstring>------// d0 -------if (p[xp] < p[yp]) p[xp] += p[yp], p[yp] = xp;------// 3a
#include <ctime>------// 28 -------| p[xp] += p[xp], p[xp] = yp;------// 3e
#include <iostream>------// ec ----int size(int x) { return -p[find(x)]; } };------// 28
#include <map>-----// 02
#include <stack>------// cf int f(int a, int b) { return min(a, b); }-------// 4f
#include <vector>-----// 4f int f(int a, int b) { return a + b; }-----// dd
using namespace std;------// 7b #endif-----// 7b #endif------// 7b
-----// 7e struct segment_tree {------------------------// ab
const int INF = 2147483647;------// be ----segment_tree(const vi &arr) : n(size(arr)), data(4*n), lazy(4*n,INF) {-----// f1
const double pi = acos(-1);------// 49 ----int mk(const vi &arr, int l, int r, int i) {------// 12
typedef unsigned long long ull;------// 81 -----int m = (l + r) / 2;-----// de
typedef vector<vi>vvi;------// 31 ------propagate(l, r, i);-------// 12
typedef vector<vii>vvii;-------// 4b ------if (r < a || b < l) return ID;------// c7
template <class T> T mod(T a, T b) { return (a % b + b) % b; }------// 70 -----if (a <= l && r <= b) return data[i];----------// ce
template <class T> int size(const T &x) { return x.size(); }-----// 68 -----int m = (l + r) / 2;------// 7a
                             -----return f(q(a, b, l, m, 2*i+1), q(a, b, m+1, r, 2*i+2)); }------// 5c
1.3. Java Template. A Java template.
                             ----void update(int i, int v) { u(i, v, 0, n-1, 0); }-----// 90
-----// a3 ------if (l == i && r == i) return data[j] = v;--------// 4a
```

```
2.4. Matrix. A Matrix class.
```

```
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----return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); \}----// 34
----void range_update(int a, int b, int v) { ru(a, b, v, 0, n-1, 0); }------// 71
----int ru(int a, int b, int v, int l, int r, int i) {-------// e0
-----propagate(l, r, i);-----// 19
-----if (l > r) return ID;------// cc
-----if (r < a || b < l) return data[i];-----// d9
-----if (l == r) return data[i] += v;-----// 5f
-----if (a <= l \& r <= b) return (lazy[i] = v) * (r - l + 1) + data[i];----// 76
-----int m = (l + r) / 2;-----// e7
-----return data[i] = f(ru(a, b, v, l, m, 2*i+1),-----// θe
----}------// 47
----void propagate(int l, int r, int i) {-----// b5
-----if (l > r || lazy[i] == INF) return;-----// 83
-----data[i] += lazy[i] * (r - l + 1);-----// 99
-----if (l < r) {------// dd
------else lazy[2*i+1] += lazy[i];-----// 72
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];-----// dd
------else lazy[2*i+2] += lazy[i];-----// a4
-----lazv[i] = INF:-----// c4
}:-----// 17
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
i...j in O(\log n) time. It only needs O(n) space.
struct fenwick_tree {------// 98
----int n; vi data;------// d3 ------return res; }-----
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {--------// dd
----void update(int at, int by) {------// 76 -----matrix<T> res(cols, rows);-----// b5
-------while (at < n) data[at] += by, at |= at + 1; }-------// fb -------for (int i = 0; i < rows; i++)------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);------// bf
```

```
template <> bool eq<double>(double a, double b) { return abs(a - b) < EPS; }---// a7
                                                template <class T>-----// 53
                                                class matrix {------// 85
                                                public:----// be
                                                ----int rows, cols;------// d3
                                                ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {------// 34
                                                -----data.assign(cnt, T(0)); }-----// d0
                                                ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),-----// fe
                                                -----cnt(other.cnt), data(other.data) { }-----// ed
                                                ----T& operator()(int i, int j) { return at(i, j); }------// e0
                                                ----void operator +=(const matrix& other) {------// c9
                                                ------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
                                                ----void operator -=(const matrix& other) {------// 68
                                                ------for (int i = 0: i < cnt: i++) data[i] -= other.data[i]: }------// 88
                                                ----void operator *=(T other) {------// ba
                                                ------for (int i = 0; i < cnt; i++) data[i] *= other; }------// 40
                                                ----matrix<T> operator +(const matrix& other) {------// ee
                                                ------matrix<T> res(*this); res += other; return res; }------// 5d
                                                ----matrix<T> operator -(const matrix& other) {------// 8f
                                                ------matrix<T> res(*this); res -= other; return res; }------// cf
                                                ----matrix<T> operator *(T other) {------// be
                                                ------matrix<T> res(*this); res *= other; return res; }------// 37
                                                ----matrix<T> operator *(const matrix& other) {------// 95
                                                ------matrix<T> res(rows, other.cols);------// 57
                                                -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
                                                -----for (int k = 0; k < cols; k++)-----// fc
                                                -----res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
};------// 57 -----while (p) {------// cb
struct fenwick_tree_sq {------// d4 -----if (p & 1) res = res * sq;-----// c1
----<mark>int</mark> n; fenwick_tree x1, x0;-------// 18 -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {------// 89
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }-----// 45 ------matrix<T> mat(*this); det = T(1);------------------// 21
};-----// 13 ------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 ------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f ------if (k >= rows) continue;-------// 3f
----return s.query(b) - s.query(a-1); }-----// f3 ------det *= T(-1);--------------------// 7a
```

```
template <class K> bool eq(K a, K b) { return a == b; }-----// 2a
```

```
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------for (int i = 0; i < cols; i++)-------// ab ----void erase(node *n, bool free = true) {-------// 58
------if (!eq<T>(mat(r, c), T(1)))-------// 2c -----else if (n->l && !n->r) parent_leq(n) = n->l, n->l->p = n->p;-----// 6b
------for (int i = cols-1; i >= c; i--) mat(r, i) /= mat(r, c);-----// 5d ------else if (n->l && n->r) {----------------------------------// 6c
------for (int i = 0; i < rows; i++) {----------// 3d ------node *s = successor(n);--------// e5
------T m = mat(i, c);--------// e8 ------erase(s, false);---------// 0a
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 5a
------} return mat; }-------// 8f ------parent_leg(n) = s, fix(s);-------// 82
private:-----// e0 -----return;-------// e5
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;-------// 43
-----if (!n) return NULL;-----// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                            -----if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0------// b5 -----node *p = n->p;------// a7
template <class T>-----// 22 -----return p; }-----// c7
class avl_tree {------// ff ----node* predecessor(node *n) const {------// b4
public:-----// f6 ------if (!n) return NULL;-------// dd
----struct node {-------// 45 -----if (n->l) return nth(n->l->size-1, n->l);------// 10
------int size, height;------// 33 -------while (p && p->l == n) n = p, p = p->p;------// 6d
------node(const T \&_item, node *_p = NULL) : item(_item), p(_p),-------// 4f ------return p; }------
------l(NULL), r(NULL), size(1), height(0) { } };-------// @d ----inline int size() const { return sz(root); }------// ef
----node *root;------// 91 ----node* nth(int n, node *cur = NULL) const {------// e4
-----node *cur = root;------// b4 ------while (cur) {------// 29
------while (cur) {-------// 8b ------if (n < sz(cur->l)) cur = cur->l;------// 75
-----if (cur->item < item) cur = cur->r;------// 71 -----else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-----// cd
------else if (item < cur->item) cur = cur->l;------// cd -----else break;-----
------else break; }------// 4f ------} return cur; }------// ed
------return cur; }-------// 84 ----int count_less(node *cur) {-------// ec
----node* insert(const T &item) {-------// 4e -----int sum = sz(cur->l);------// bf
-----prev = *cur;-----// f0 -----cur = cur->p;-----// eb
#if AVL_MULTISET-----// 0a private:----// d5
#else------// ff ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
------else return *cur;------// 54 -----return n && height(n->r); }------// a8
-----node *n = new node(item, prev);-------// eb ----inline bool too_heavy(node *n) const {------// @b
-----*cur = n, fix(n); return n; }-----// 29
                            -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }-----// 67
```

```
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------if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }------// ef #define SWP(x,y) tmp = x, x = y, y = tmp------// fb
------if (n->p->l == n) return n->p->l;-------// 83 ----bool operator ()(const int &a, const int &b) { return a < b; } };------// e9
------if (n->p->r == n) return n->p->r;-------// cc template <class Compare = default_int_cmp>------// 30
------n->height = 1 + max(height(n->l), height(n->r)); }-------// 41 ----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }------// a0
------while (i > 0) {------// 1a
-----parent_leg(n) = l; \[ \]-----// fc
                               -----int p = (i - 1) / 2;-----// 77
-----augment(n), augment(l)-------// 81 ------while (true) {---------------------// 3c
----void fix(node *n) {-------// 0d -------int m = r >= count || cmp(l, r) ? l : r;------// cc
------while (n) { augment(n);-------// 69 ------if (!cmp(m, i)) break;------// 42
------if (too_heavy(n)) {-------// 4c -----swp(m, i), i = m; } }-----// 1d
------if (left_heavy(n) && right_heavy(n->l)) left_rotate(n->l);-----// a9 public;------
------else if (right_heavy(n) && left_heavy(n->r))-------// b9 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {------// b9
-----right_rotate(n->r);------// 08 -----q = new int[len], loc = new int[len];------// f8
------if (left_heavy(n)) right_rotate(n);------// 93 ------memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// d5 ----~heap() { delete[] q; delete[] loc; }------// 09
------n = n->p; }-------// 28 ----void push(int n, bool fix = true) {-------// b7
#ifdef RFSI7F-----// a9
Also a very simple wrapper over the AVL tree that implements a map interface.
                               -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"-----// 01
                               -----while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                               -----int *newq = new int[newlen], *newloc = new int[newlen];-----// e3
template <class K, class V>-----// da
                               -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --//94
class avl_map {------// 3f
                               -----memset(newloc + len, 255, (newlen - len) << 2);-----// 18
public:----// 5d
                               -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                               -----loc = newloc, q = newq, len = newlen;-----// 61
------K key; V value;------// 32
                               #else-----// 54
-----node(K k, V v) : key(k), value(v) { }-----// 29
                               -----assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                               #endif------// 64
----avl_tree<node> tree:-----// b1
                               ----V& operator [](K key) {------// 7c
                               -----assert(loc[n] == -1);-----// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
                               -----loc[n] = count, q[count++] = n;-----// 6b
------if (!n) n = tree.insert(node(key, V(0)));------// cb
                               ------if (fix) swim(count-1); }------// bf
-----return n->item.value;------// ec
                               ----}------// 2e
                               -----assert(count > 0):-----// eb
}:-----// af
                               ------loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 50
                               -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
```

```
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----}------FIND_UPDATE(x->next[i]->item, target);-------// 3a
----void heapify() { for (int i = count - 1; i > 0; i--)----------// 39 ------int lvl = bernoulli(MAX_LEVEL);----------------------// 7a
------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }--------// 0b ------if(lvl > current_level) current_level = lvl;-----------------------// 8a
----void update_key(int n) {--------------------------// 26 -----x = new node(lvl, target);-------------------// 36
----void clear() { count = 0, memset(loc, 255, len << 2); } };-------// 58 ------update[i] = x;-----------// 20
                                      -----update[i]->lens[i] = pos[0] + 1 - pos[i];------// 42
2.7. Skiplist. An implementation of a skiplist.
                                      ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// 07
#define BP 0.20-----// aa
                                      -----size++;-----// 19
#define MAX_LEVEL 10------// 56
                                      -----return x; }-----// c9
unsigned int bernoulli(unsigned int MAX) {------// 7b
                                      ----void erase(T target) {------// 4d
----unsigned int cnt = 0;-----// 28
----while(((float) rand() / RAND_MAX) < BP \&\& cnt < MAX) cnt++;-----// d1
                                      ------FIND_UPDATE(x->next[i]->item, target);-------// 6b
                                      -----if(x && x->item == target) {-----// 76
----return cnt; }-----// a1
template<class T> struct skiplist {------// 34
                                      -----for(int i = 0; i <= current_level; i++) {------// 97
                                      -----if(update[i]->next[i] == x) {------// b1
----struct node {------// 53
                                      -----update[i]->next[i] = x->next[i];------// 59
                                      -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
-----int *lens:-----// 07
                                      -----} else update[i]->lens[i] = update[i]->lens[i] - 1;-----// 88
                                      ------#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))-------// 25
                                      -----delete x; _size--;------// 81
-----node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}-----// 7c
                                      ------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
-----/node() { free(lens); free(next); }; };-----// aa
                                      -----current_level--; } } };-----// 59
----int current_level, _size;------// 61
----node *head;------// b7
                                      2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                      list supporting deletion and restoration of elements.
-----skiplist() { clear(); delete head; head = NULL; }------// aa
                                      template <class T>-----// 82
----#define FIND_UPDATE(cmp, target) \|-----// c3
                                      struct dancing_links {-----// 9e
------int pos[MAX_LEVEL + 2]; \[\[\]\------// 18
                                      ----struct node {------// 62
-----memset(pos, 0, sizeof(pos)); \-----// f2
                                      -----T item:-----// dd
-----node *update[MAX_LEVEL + 1]; \| ------// 01
                                      -----node(const T &_item, node *_l = NULL, node *_r = NULL)------// 6d
-----memset(update, 0, MAX_LEVEL + 1); \[\bar{\}\]------// 38
                                      -----: item(_item), l(_l), r(_r) {------// 6d
                                      -----if (l) l->r = this;-----// 97
-----for(int i = MAX\_LEVEL; i >= 0; i--) { \sqrt{\phantom{a}}
                                      -----if (r) r->l = this;-----// 81
-----pos[i] = pos[i + 1]; \\------// 68
                                      ------}-------// 2d
----};------// d3
------update[i] = x; N----------// dd ----dancing_links() { front = back = NULL; }------// 72
----void clear() { while(head->next && head->next[0])-------// 91 -----if (!front) front = back;-------------// d2
------erase(head->next[0]->item); }-------// e6 ------return back;--------------------------// cθ
------return x && x->item == target ? x : NULL; }-------// 50 ----node *push_front(const T &item) {------------------// 4a
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------if (!back) back = front;--------------------// 10
-----return pos[0]; }------// 19 -----return front;------// cf
```

```
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------if (!n->l) front = n->r; else n->l->r = n->r;-------// ab -------bb(pt _from, pt _to) : from(_from), to(_to) {}------// 57
----}-------double sum = 0.0;-------// d9
------if (!n->l) front = n; else n->l->r = n;--------// a5 ------if (p.coord[i] < from.coord[i])------// a0
}:------sum += pow(p.coord[i] - to.coord[i], 2.0);------// 8c
                                    2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                                   -----return sqrt(sum); }-----// ef
element.
                                    ------bb bound(double l, int c, bool left) {------// b6
#define BITS 15------pt nf(from.coord), nt(to.coord);------// 5c
----int cnt[BITS][1<<BITS];-------// aa ------else nf.coord[c] = max(nf.coord[c], l);------// 71
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }------// b0 -----return bb(nf, nt); } };------// 3b
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49 ------pt p; node *\, *r;--------------------------// 46
----int nth(int n) {-------// 8a ------node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
------int res = 0;-------// a4 ----node *root;------// 30
------for (int i = BITS-1; i >= 0; i--)-------// 99 ----kd_tree() : root(NULL) { }-------// 57
------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4 ----kd_tree(vector<pt> pts) { root = construct(pts, 0, size(pts) - 1, 0); }----// 66
------return res:-------// 3a ----node* construct(vector<pt> &pts, int from, int to, int c) {-------// 0b
----}-----if (from > to) return NULL;------// f4
-----nth_element(pts.begin() + from, pts.begin() + mid,------// f3
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                                    -----pts.begin() + to + 1, cmp(c));-----// 97
bor queries.
                                    -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// dd
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1) - \cdots / 77
                                    -----construct(pts, mid + 1, to, INC(c))); }-----// 03
template <int K>-----// cd ----bool contains(const pt &p) { return _con(p, root, 0); }-----// f0
class kd_tree {------// 7e ----bool _con(const pt &p, node *n, int c) {------// 74
-------double coord[K];-------// d6 ------if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// 95
-----pt() {}------// c1 -----return true; }------// 86
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }-----// 4c ----void insert(const pt &p) { _ins(p, root, 0); }-----// f9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {-------// 09
------double sum = 0.0;-------// c4 ------if (!n) n = new node(p, NULL, NULL);------// 4d
------for (int i = 0; i < K; i++)-------// 23 ------else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));------// a0
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }------// 2c
-----return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }----// 73
----struct cmp {------// 8f ----void _clr(n->l), _clr(n->r), delete n; }-----// 1a
-----cmp(int _c) : c(_c) {}------// a5 -----assert(root);------// f8
------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// bf
-----cc = i == 0 ? c : i - 1;------// bc -----pt from(cs);------// 5a
------if (abs(a.coord[cc] - b.coord[cc]) > EPS)------// 28 ------for (int i = 0; i < K; i++) cs[i] = INFINITY;-----// 37
-----return a.coord[cc] < b.coord[cc];-----// b7 -----pt to(cs);------
-----return false; } };------// 6e
----struct bb {------// 30
```

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-----const pt &p, node *n, bb b, double &mn, int c, bool same) {------// aa
                                               }-----// e5
-----if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// 19
------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;------// 9f
                                                                    3. Graphs
-----pt resp = n->p;------// 6b
-----if (found) mn = min(mn, p.dist(resp));------// 18
                                               3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
-----node *n1 = n->l, *n2 = n->r;------// aa
                                               edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
------for (int i = 0: i < 2: i++) {------// 50
                                               graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
------if (i == 1 || cmp(c)(n->p, p)) swap(n1, n2), swap(l1, l2);------// e^2
                                               connected. It runs in O(|V| + |E|) time.
-----pair<pt, bool> res =-----// 33
                                               int bfs(int start, int end, vvi& adj_list) {------// d7
-----nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// 72
                                                ----queue<ii> Q;------// 75
-----if (res.second && (!found || p.dist(res.first) < p.dist(resp)))----// 76
                                                ----Q.push(ii(start, 0));------// 49
-----resp = res.first, found = true;-----// 3b
-----}-----// aa
-----return make_pair(resp, found); } };------// dd
                                                -----ii cur = Q.front(); Q.pop();------// e8
2.11. Sqrt Decomposition. Design principle that supports many operations in amortized \sqrt{n} per
                                                   -----// 06
operation.
                                                ------if (cur.first == end)------// 6f
struct segment {-----// b2
                                               -----return cur.second;------// 8a
----vi arr:-----// 8c
----segment(vi arr) : arr(arr) { } };------// 92
                                               -----vi& adj = adj_list[cur.first];-----// 3f
vector<seament> T:-----// d5
                                               ------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)------// bb
                                               -----Q.push(ii(*it, cur.second + 1));------// b7
int K;-----// 02
                                               ----}-------// 93
                                               }-----// 7d
----int cnt = 0:-----// c1
----for (int i = 0; i < size(T); i++)-----// 71
                                                 Another implementation that doesn't assume the two vertices are connected. If there is no path
-----cnt += size(T[i].arr);------// 97
                                               from the starting vertex to the ending vertex, a-1 is returned.
----K = static_cast<int>(ceil(sqrt(cnt)) + 1e-9);------// 49
                                               int bfs(int start, int end, vvi& adj_list) {------// d7
----vi arr(cnt):-----// da
                                                ----set<<u>int</u>> visited;-----// b3
----for (int i = 0, at = 0; i < size(T); i++)-----// 94
                                                ----queue<ii>> Q;------// bb
------for (int j = 0; j < size(T[i].arr); j++)-----// a9
                                                ----Q.push(ii(start, 0));-----// 3a
-----arr[at++] = T[i].arr[j];-----// 1c
                                                ----visited.insert(start):-----// b2
----T.clear():------// 31
----for (int i = 0; i < cnt; i += K)------// 61
                                                ---while (!0.empty()) {------// f7
-----T.push_back(segment(vi(arr.begin()+i, arr.begin()+min(i+K, cnt))));----// 3e
                                               -----ii cur = Q.front(); Q.pop();-----// 03
int split(int at) {------// da
----int i = 0;------// 82
                                                 ------return cur.second;-----// b9
----while (i < size(T) && at >= size(T[i].arr))------// 2d
-----at -= size(T[i].arr), i++;-----// 16
                                                -----vi& adj = adj_list[cur.first];------// f9
----if (i >= size(T)) return size(T);------// 6e
                                                 -----for (vi::iterator it = adj.begin(); it != adj.end(); it++)------// 44
----if (at == 0) return i;------// 83
                                                ------if (visited.find(*it) == visited.end()) {-------// 8d
----T.insert(T.begin() + i + 1, segment(vi(T[i].arr.begin() + at, T[i].arr.end())));
                                                -----Q.push(ii(*it, cur.second + 1));-------------// ab
----T[i] = segment(vi(T[i].arr.begin(), T[i].arr.begin() + at));------// 05
                                                  -----visited.insert(*it);-----// cb
}-----// 11
void insert(int at, int v) {------// fc
----vi arr; arr.push_back(v);------// 2c
----T.insert(T.begin() + split(at), segment(arr));------// 09
}-----// 1f
void erase(int at) {------// ae
```

```
-----for (int j = 0; j < n; j++)-----// 77
3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|) time.
                                                -----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
int *dist, *dad;-----// 46
                                                   -----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);-----// e1
struct cmp {-----// a5
                                                   -----// 86
----bool operator()(int a, int b) {------// bb
-----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
                                                3.4. Strongly Connected Components.
};-----// 41
pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
                                               3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
----dist = new int[n]:-----// 84
                                                graph in O(|V| + |E|) time.
----dad = new int[n]:-----// 05
                                                #include "../data-structures/union_find.cpp"-----// 5e
----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;-----// d6
                                                -----// 11
----set<int, cmp> pq;-----// 04
                                                vector<br/>bool> visited:-----// 66
----dist[s] = 0, pq.insert(s);------// 1b
                                                vi order:----// 9b
----while (!pq.empty()) {------// 57
                                                -----// a5
------int cur = *pq.beqin(); pq.erase(pq.beqin());-----// 7d
                                                void scc_dfs(const vvi &adj, int u) {------// a1
------for (int i = 0; i < size(adj[cur]); i++) {-------// 9e
                                                ----int v: visited[u] = true:-----// e3
-----int nxt = adj[cur][i].first,-----// b8
                                                ----for (int i = 0; i < size(adj[u]); i++)-----// c5
-----/ndist = dist[cur] + adj[cur][i].second;------// 0c
                                                -----if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-----// 6e
------if (ndist < dist[nxt]) pq.erase(nxt),-----// e4
                                                ----order.push_back(u):-----// 19
-----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);-----// 0f
                                                }-----// dc
------}-----// 75
                                                ----}-----// e8
                                                pair<union_find. vi> scc(const vvi &adi) {------// 3e
----return pair<int*, int*>(dist, dad);------// cc
                                                ----int n = size(adj), u, v;-----// bd
}-----// af
                                                ----order.clear();-----// 22
                                                ----union_find uf(n);-----// 6d
3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
                                                ----vi dag;------// ae
problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
                                                ----vi rev(n):-----// 20
negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
                                                ----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++)------// b9
                                                -----rev[adj[i]]]].push_back(i);-----// 77
int* bellman_ford(int n. int s. vii* adi. bool& has_negative_cvcle) {------// cf
                                                ----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04
----has_negative_cycle = false;------// 47
                                                ----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);------// e4
----int* dist = new int[n];------// 7f
                                                ----fill(visited.begin(), visited.end(), false);-----// c2
----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;-----// 10
                                                ----stack<int> S;-----// 04
----for (int i = 0; i < n - 1; i++)-----// a1
                                                ----for (int i = n-1; i >= 0; i--) {------// 3f
------for (int j = 0; j < n; j++)------// c4
                                                ------if (visited[order[i]]) continue;-----// 94
------if (dist[i] != INF)------// 4e
                                                ------S.push(order[i]), dag.push_back(order[i]);------// 40
-----for (int k = 0; k < size(adj[j]); k++)-----// 3f
                                               ------while (!S.empty()) {------// 03
-----dist[adj[j][k].first] = min(dist[adj[j][k].first], ------//61
                                                -----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);-----// 1b
-----dist[j] + adj[j][k].second);------// 47
                                                -----for (int j = 0; j < size(adj[u]); j++)-----// 21
----for (int j = 0; j < n; j++)-----// 13
                                                -----if (!visited[v = adj[u][j]]) S.push(v);-----// e7
------for (int k = 0; k < size(adj[j]); k++)------// a0
                                                -----}------------------// 98
------if (dist[i] + adj[i][k].second < dist[adj[i][k].first])------// ef
                                                ----}------------// d9
-----has_negative_cycle = true;-----// 2a
                                                ----return pair<union_find. vi>(uf. dag):-----// f2
----return dist;------// 2e
                                                }-----// ca
}-----// c2
                                                3.5. Minimum Spanning Tree.
3.3. All-Pairs Shortest Paths.
                                                3.5.1. Kruskal's algorithm.
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                                #include "../data-structures/union_find.cpp"-----// 5e
problem in O(|V|^3) time.
                                                -----// 11
----for (int k = 0; k < n; k++)-------// 49 // edges is a list of edges of the form (weight, (a, b))------// c6
```

------for (int i = 0; i < n; i++)------// 21 // the edges in the minimum spanning tree are returned on the same form-----// 4d

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                                                       10
-----if (uf.find(edges[i].second.first) !=------// d5 ------else if (indeg[i] == outdeg[i] + 1) end = i, c++;-------// 4f
-----uf.find(edges[i].second.second)) {------// 8c -----else if (indeg[i] != outdeg[i]) return ii(-1,-1);-----// fa
-----uf.unite(edges[i].second.first, edges[i].second.second);-----// a2 ----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 6e
}------// 88 }------// 35
                            bool euler_path() {-----// d7
3.6. Topological Sort.
                            ----ii se = start_end();-----// 45
                            ----int cur = se.first, at = m + 1;-----// 8c
3.6.1. Modified Depth-First Search.
                            ----if (cur == -1) return false;-----// 45
void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
                            ----stack<int> s;------// f6
------bool& has_cycle) {------// a8
                            ----while (true) {------// 04
----color[cur] = 1;------// 5b
                           -----if (outdeg[cur] == 0) {-----// 32
----for (int i = 0; i < size(adj[cur]); i++) {------// 96
                           -----res[--at] = cur;-----// a6
------int nxt = adj[cur][i];------// 53
                            -----if (s.emptv()) break:-----// ee
-----if (color[nxt] == 0)------// 00
                           -----cur = s.top(); s.pop();-----// d7
-----tsort_dfs(nxt, color, adj, res, has_cycle);-----// 5b
                           -----else if (color[nxt] == 1)------// 53
                           ....}.....// ba
-----has_cycle = true;-----// c8
                           ---return at == 0:----// c8
-----if (has_cycle) return;------// 7e
                            }-----// aa
----}------// 3d
----color[cur] = 2;-----// 16
                           3.8. Bipartite Matching.
----res.push(cur):-----// cb
                            3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
}------// 9e
                            O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
-----// ae
                            graph, respectively.
vi tsort(int n, vvi adj, bool& has_cycle) {-----// 37
----has cycle = false:------// 37 vi* adi:------// cc
----stack<int> S;------// 54 bool* done;------// b1
----char* color = new char[n];------// b1 int alternating_path(int left) {------// da
-----tsort_dfs(i, color, adj, S, has_cycle);------// 40 ------int right = adj[left][i];-------// b6
------if (has_cycle) return res;------// 6c -----if (owner[right] == -1 || alternating_path(owner[right])) {------// d2
------}-----owner[right] = left; return 1;--------// 26
----return res;------// 07
ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                            #define MAXN 5000-----// f7
                           int dist[MAXN+1], q[MAXN+1];-----// b8
#define MAXV 1000-----// 2f
#define MAXE 5000------// 87 #define dist(v) dist[v == -1 ? MAXN : v]-------// 0f
vi adj[MAXV];------// ff struct bipartite_graph {------// 2b
```

```
------L(new int[N]), R(new int[M]), adj(new vi[N]) {}------------// cd ----int n, ecnt, *head, *curh;------------------------------// 46
------else dist(v) = INF;---------// b3 ------memset(head, -1, n * sizeof(int));-------// 56
-----dist(-1) = INF:------// 96 ---}-----// 77
------while(l < r) {------// 69 ----void destroy() { delete[] head; delete[] curh; }------// f6
------int v = q[l++];------// 0c ----void reset() { e = e_store; }------// 87
-----if(v != -1) {--------if ((ret = augment(e[i].v, t, min(f, e[i].cap))) > 0)------// 1f
-----foreach(u, adj[v])-------// 43 ------return (e[i].cap -= ret, e[i^1].cap += ret, ret);-----// ac
-----if(dfs(R[*u])) {-------// 75 ---}-----// 75
------return true;------// 1f ------if(s == t) return 0;-------// 9d
-----}-----memset(d, -1, n * sizeof(int));------// a8
----}------while (l < r)------// 7a
----void add_edge(int i, int j) { adj[i].push_back(j); }------// 11 ------for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)-----// a2
-----d[q[r++] = e[i].v] = d[v]+1;------// 28
------memset(L, -1, sizeof(int) * N);--------// 8f ------if (d[s] == -1) break;-------// a0
------memset(R, -1, sizeof(int) * M);-------// 39 ------memcpy(curh, head, n * sizeof(int));------// 10
------while(bfs()) for(int i = 0; i < N; ++i)--------// 77 -------while ((x = augment(s, t, INF)) != 0) f += x;------// a6
};------// d3 ---}
                     }:----// 3b
3.9. Maximum Flow.
                     3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
3.9.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It computes O(|V||E|^2). It computes the maximum flow of a flow network.
the maximum flow of a flow network.
                     #define MAXV 2000-----// ba
#define MAXV 2000-----// ba int q[MAXV], p[MAXV];-----// 7b
----struct edge {--------------------------// le ------int v, cap, nxt;----------------// cb
-----edge() { }------// 38 ---};------// 38 ---};-------// 38 ---}
------edge(int _v, int _cap, int _nxt) : v(_v), cap(_cap), nxt(_nxt) { }-----// bc ----int n, ecnt, *head;------------------------// 39
```

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-----e.reserve(2 * (m == -1 ? n : m));------// 92 ----vector<edge> e, e_store;------// 4b
------memset(head = new int[n], -1, n << 2);-------// 58 ----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {-------// dd
----}-----e.reserve(2 * (m == -1 ? n : m));------// e6
-----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// 4c ----void reset() { e = e_store; }-------// 88
-------e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;--------// bc ----void add_edge(int u, int v, int cost, int uv, int vu=0) {------// b4
----}-----e.push_back(edge(v, uv, cost, head[u])); head[u] = ecnt++;------// 43
------if (s == t) return 0:--------// d6 ---}------// 16
------int f = 0, l, r, v;------------------// 6f ------if (s == t) return ii(0, 0);--------// 34
-----memset(d. -1, n << 2);------// 3b -----memset(pot, 0, n << 2);------// 24
------while (l < r)-----// 2c -----memset(d, -1, n << 2);-----// fd
------for (int u = g[l++], i = head[u]; i != -1; i = e[i].nxt)-----// c6 ------memset(p, -1, n << 2);-----------------------------// b7
-----if (e[i].cap > 0 &&------// 8a -----set<int, cmp> q;-------// d8
------while (at != -1)-------// cd -------int cd = d[u] + e[i].cost + pot[u] - pot[v = e[i].v];-----// 1d
------}-------if (q.find(v) != q.end()) q.erase(q.find(v));------// e2
------if (res) reset();--------d[v] = cd; p[v] = i;-------// f7
------q.insert(v);---------------------// bc
3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
                          -----if (p[t] == -1) break;-----// 09
fied to find shortest path to augment each time (instead of just any path). It computes the maximum
                          -----int x = INF, at = p[t];-----// e8
flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with
                          ------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 32
minimum cost. Running time is O(|V|^2|E|\log|V|).
                          -----at = p[t], f += x;------// 43
#define MAXV 2000-----// ba
                          ------while (at != -1)------// 53
int d[MAXV], pot[MAXV]; .....// 80
                          ------e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// 95
----bool operator ()(int i, int j) {-------// 8a -------for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];-----// ff
-----return d[i] == d[j] ? i < j : d[i] < d[j];------// 89 ____}
----}------// df ------// df ------// 5e
struct flow_network {------// eb ___}
----struct edge {------// 9a }.....// d7
-----int v, Cap, Cost, nxt;-----// ad
------edge(int _v, int _cap, int _cost, int _nxt)------// ec
-----: v(_v), cap(_cap), cost(_cost), nxt(_nxt) { }------// c4 3.11. All Pairs Maximum Flow.
```

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                                                14
------it = cur->children.find(head):------// d9
-----process(v);------// 41 ----template<class I>------// 9c
------int v = queries[u][i].first;-------// 38 -------T head = *beqin;----------// 43
4.2. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
}:-----// 5f
                        struct entry { ii nr; int p; };-----// f9
                        bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }-----// 77</pre>
          4. Strings
                        struct suffix_array {------// 87
4.1. Trie. A Trie class.
                        ----string s; int n; vvi P; vector<entry> L; vi idx;-----// b6
template <class T>-----// 82 ----// REMINDER: Append a large character ('\x7F') to s------// 70
class trie {------// 9a ----suffix_array(string _s) : s(_s), n(size(s)) {------// e5
private:-----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);------// 8a
----struct node {------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';------// 8d
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {------// 46
------<mark>int</mark> prefixes, words;------// e2 ------P.push_back(vi(n));------// 30
------for (int i = 0; i < n; i++)------// d5
public:------L[L[i].p = i].nr = ii(P[stp - 1][i],-----// fc
----node* root;------i + cnt < n ? P[stp - 1][i + cnt] : -1);------// e5
----trie() : root(new node()) { }-------// 8f ------sort(L.begin(), L.end());------// bc
----template <class I>------// 89 -------for (int i = 0; i < n; i++)------// 85
----void insert(I begin, I end) {-------// 3c ------P[stp][L[i].p] = i > 0 &&-----// eb
------while (true) {-------// 67 ----}-----
-----if (begin == end) { cur->words++; break; }------// db ---}
-----else {-------// 3e ----int lcp(int x, int y) {-------// 05
------typename map<T, node*>::const_iterator it;------// 01 -----if (x == y) return n - x;-----// 7f
-----it = cur->children.find(head);------// 77 ------for (int k = size(P) - 1; k >= 0 && x < n && y < n; k--)-----// 07
-----pair<T, node*> nw(head, new node());------// cd -----return res;-----
----template<class I>-----// b9
state machine from a set of keywords which can be used to search a string for any of the keywords.
-----node* cur = root;------// 32
------while (true) {-------// bb struct aho_corasick {-------// 78
-----else {-------// 1e -----string keyword; out_node *next;------// f0
-----T head = *begin;------// 5c -----out_node(string k, out_node *n) : keyword(k), next(n) { }-----// 26
```

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------map<char, go_node*> next;-------// 6b };------// 32
-----out_node *out; go_node *fail;-----// 3e
                                  4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
------qo_node() { out = NULL; fail = NULL; }------// 0f
                                  also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
----qo_node *qo;-----// b8
                                  accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
----aho_corasick(vector<string> keywords) {------// 4b
                                  int* z_values(const string &s) {-----// 4d
-----qo = new qo_node();-----// 77
------foreach(k, keywords) {-------// e4
                                  ----int n = size(s);------// 97
-----go_node *cur = go;-----// 9d ----int* z = new int[n];-----// c4
------foreach(c, *k)------// 38 ----int l = 0, r = 0;-------// 1c
-----(cur->next[*c] = new go_node());-----// 75 ----for (int i = 1; i < n; i++) {-------// 7e
-----cur->out = new out_node(*k, cur->out);------// 6e ----z[i] = 0;------
-----queue<go_node*> q;------// 8a ------l = r = i;-------l = r = i;-------// 87
------foreach(a, go->next) q.push(a->second);-------// a3 -------while (r < n && s[r - l] == s[r]) r++;------// ff
------while (!q.empty()) {--------// 43 ------z[i] = r - l; r--;----------// fc
------go_node *r = q.front(); q.pop();------// 2e -----} else if (z[i - l] < r - i + 1) z[i] = z[i - l];------// bf
-----z[i] = r - l; r--; } }------// 8d
-----st = st->fail;------// db
-----if (!st) st = go;-----// e7
                                  4.5. Palindromic Tree. Constructs a Palindromic Tree in O(n), one character at a time.
-----s--sfail = st->next[a->first];------// 29
                                  #define MAXN 100100-----// 29
-----if (s->fail) {------// 3b
                                  #define SIGMA 26-----// e2
-----if (!s->out) s->out = s->fail->out;-----// 80
                                  #define BASE 'a'-----// a1
-----else {-----// ed
                                  char *s = new char[MAXN];-----// db
-----out_node* out = s->out;-----// bf
                                  struct state {-----// 33
-----// ca
                                  ----int len, link, to[SIGMA];------// 24
-----out->next = s->fail->out;-----// 65
                                  } *st = new state[MAXN+2];-----// 57
struct palindromic_tree {-----// cf
----int last, sz, n;------// 35
                                   ----palindromic_tree() : last(1), sz(2), n(0) {------// b7
-----}-----// e8
                                   -----st[0].len = st[0].link = -1;------// 0e
----vector<string> search(string s) {------// 8d
                                   -----st[1].len = st[1].link = 0; }------// 35
                                  ----int extend() {------// 5d
-----vector<string> res;------// ef
                                   -----char c = s[n++]; int p = last;-----// a3
-----ao_node *cur = ao:-----// 61
                                  -----while (n - st[p].len - 2 < 0 \mid | c != s[n - st[p].len - 2]) p = st[p].link;
-----foreach(c, s) {------// 6c
                                   -----if (!st[p].to[c-BASE]) {------// 05
------while (cur && cur->next.find(*c) == cur->next.end())------// 1f
                                  -----int q = last = sz++;-----// ad
-----cur = cur->fail;-----// 9e
                                   -----st[p].to[c-BASE] = q;-----// bb
-----if (!cur) cur = go;-----// 2f
                                  -----st[q].len = st[p].len + 2;-----// 86
-----cur = cur->next[*c];-----// 58
-----if (!cur) cur = go;-----// 3f
                                  -----do { p = st[p].link;-----// c8
                                  -----} while (p != -1 \&\& (n < st[p].len + 2 || c != s[n - st[p].len - 2]));
-----for (out_node *out = cur->out; out = out->next)------// e0
-----res.push_back(out->keyword);-----// 0d
                                  ------if (p == -1) st[q].link = 1;------// 02
                                  ------else st[q].link = st[p].to[c-BASE];------// e6
-----return 1; }-----// bc
```

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------last = st[p].to[c-BASE];-------// 30 ----static const unsigned int radix = 10000000000U;------// f0
-----return 0: } }:------// da ----int size() const { return data.size(): }------// 29
                                    ----void init(string n) {------// 13
                                    -----intx res; res.data.clear();-----// 4e
              5. Mathematics
                                    -----if (n.empty()) n = "0";------// 99
5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common
                                    -----if (n[0] == '-') res.sign = -1, n = n.substr(1);------// 3b
terms.
                                    ------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {-------// e7
template <class T>------// 82 ------// 82 ------// 98
class fraction {------// cf -----for (int j = intx::dcnt - 1; j >= 0; j--) {-----// 72
private:-----// 8e ------int idx = i - j;-------// cd
public:-----digit = digit * 10 + (n[idx] - '0');-----// 1f
-----assert(d_ != 0);-----// 3d -----// 3d
-----n = n_, d = d_;------// 06 ------data = res.data;-----// 7d
- if (d < T(\theta)) n = -n, d = -d; - normalize(res.sign); - 176
-----T q = gcd(abs(n), abs(d));------// fc ____}
----fraction(T n_) : n(n_), d(1) { }------// 84 ------if (data.empty()) data.push_back(0);------// fa
----fraction(const fraction<T>& other) : n(other.n), d(other.d) { }-------// 01 -------for (int i = data.size() - 1; i > 0 && data[i] == 0; i--)------// 27
----fraction<T> operator +(const fraction<T>& other) const {------// b6 ------data.erase(data.begin() + i);------// 67
-----return fraction<T>(n * other.d + other.n * d, d * other.d);}------// 3b -----sign = data.size() == 1 && data[0] == 0 ? 1 : nsign;------// ff
-----return fraction<T>(n * other.d - other.n * d, d * other.d);}------// 47
----fraction<T> operator *(const fraction<T>& other) const {------// 38 ----friend ostream& operator <<(ostream& outs, const intx& n) {------// 0d
-----return fraction<T>(n * other.n, d * other.d); }------// c5 -----if (n.siqn < 0) outs << '-';--------// c0
----bool operator <(const fraction<T>& other) const {------// 0c -----if (first) outs << n.data[i], first = false;-----// 33
----bool operator <=(const fraction<T>& other) const {------// 48 ------unsigned int cur = n.data[i];------// 0f
------return !(other < *this); }------// 86 -----stringstream ss; ss << cur;-----// 8c
----bool operator >(const fraction<T>& other) const {------// c9 -----string s = ss.str();------
------return other < *this; }------// 6e -------int len = s.size();------// 0d
----bool operator ==(const fraction<T>& other) const {------// 23 _____}
----bool operator !=(const fraction<T>& other) const {------// ec -----return outs;-----
------return !(*this == other); }------// d1 ___}
};-----string to_string() const { stringstream ss; ss << *this; return ss.str(); }// fc
                                    ----bool operator <(const intx& b) const {-------// 21
5.2. Big Integer. A big integer class.
                                    ------if (sign != b.sign) return sign < b.sign;-----// cf
----intx() { normalize(1); }------// 6c ------return sign == 1 ? size() < b.size() > b.size();-----// 4d
----intx(string n) { init(n); }------// b9 -------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35
----intx(int n) { stringstream ss; ss << n; init(ss.str()); }--------// 36 -----return sign == 1 ? data[i] < b.data[i] : data[i] > b.data[i];--// 27
----intx(const intx& other) : sign(other.sign), data(other.data) { }------// 3b -----return false;-----
----vector<unsigned int> data;-----// 19
                                    ----intx operator -() const { intx res(*this); res.sign *= -1; return res; }---// 9d
----static const int dcnt = 9;-----// 12
```

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                                                                            17
------<mark>unsigned long long carry = 0;-------// 5c ------return pair</mark><intx, intx>(q.normalize(n.sign * d.sign), r);------// a1
-----carry += (i < size() ? data[i] : OULL) +------// 91 ----intx operator /(const intx& d) const {-------// a2
-----carry /= intx::radix;------// fd -----return divmod(*this,d).second * sign; }------// 5a
-----return c.normalize(sign);-----// 20
                                      5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
#include "intx.cpp"-----// 83
----intx operator -(const intx& b) const {-------// 53
                                      #include "fft.cpp"-----// 13
------if (sign > 0 && b.sign < 0) return *this + (-b);------// 8f
                                       -----// e0
-----if (sign < 0 && b.sign > 0) return -(-*this + b);------// 1b
                                      intx fastmul(const intx &an, const intx &bn) {------// ab
------if (sign < 0 && b.sign < 0) return (-b) - (-*this);-------// a1
                                       ----string as = an.to_string(), bs = bn.to_string();------// 32
-----if (*this < b) return -(b - *this);------// 36
                                       ----int n = size(as), m = size(bs), l = 1,-----// dc
-----intx c; c.data.clear();------// 6b
-----long long borrow = 0;-----// f8
                                       -----len = 5, radix = 100000,-----// 4f
                                       -----*a = new int[n], alen = 0,------// b8
------for (int i = 0: i < size(): i++) {-------// a7
                                       -----*b = new int[m], blen = 0;------// 0a
-----borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
                                       ----memset(a, 0, n << 2);-----// 1d
-----c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
                                       ----memset(b, 0, m << 2);-----// 01
-----/ od
                                       ----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
------for (int j = min(len - 1, i); j >= 0; j--)------// 43
-----return c.normalize(sign);------// 35
                                      -----a[alen] = a[alen] * 10 + as[i - j] - 0; ------// 14
----}-----------// 85
----intx operator *(const intx& b) const {------// bd
                                       ----for (int i = m - 1; i >= 0; i -= len, blen++)------// b6
                                       ------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
-----intx c; c.data.assign(size() + b.size() + 1, 0);-----// d0
-----for (int i = 0; i < size(); i++) {------// 7a
                                       -----b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
                                       ----while (l < 2*max(alen,blen)) l <<= 1;------// 51
-----long long carry = 0;-----// 20
                                       ----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
------for (int j = 0; j < b.size() || carry; j++) {------// cθ
                                       ----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);------// 35
-----if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
                                       ----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? <math>b[i] : 0, 0);-----// 66
-----carry += c.data[i + j];-----// 18
-----c.data[i + j] = carry % intx::radix;-----// 86
                                      ----fft(A, l); fft(B, l);-----// f9
                                       ----for (int i = 0; i < l; i++) A[i] *= B[i];------// e7
-----carry /= intx::radix;-----// 05
                                       ----fft(A, l, true);------// d3
----ull *data = new ull[l];-----// e7
-----return c.normalize(sign * b.siqn);-----// de
                                       ----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
                                       ----for (int i = 0; i < l - 1; i++)------// 90
----}------// c6
---- friend pair<intx,intx> divmod(const intx\& n, const intx\& d) {------// fb
                                       -----if (data[i] >= (unsigned int)(radix)) {------// 44
                                       -----data[i+1] += data[i] / radix;-----// e4
-----assert(!(d.size() == 1 && d.data[0] == 0));------// e9
                                       -----data[i] %= radix;-----// bd
-----intx q, r; q.data.assign(n.size(), 0);-----// ca
                                       ------for (int i = n.size() - 1; i >= 0; i--) {------// 1a
                                       ----int stop = l-1;------// cb
-----r.data.insert(r.data.begin(), 0);-----// c7
                                       ----while (stop > 0 && data[stop] == 0) stop--;-----// 97
----r = r + n.data[i];-----// e6
                                      ----stringstream ss;-----// 42
-----long long k = 0;-----// cc
                                      ----ss << data[stop];-----// 96
-----if (d.size() < r.size())------// b9
                                      ----for (int i = stop - 1; i >= 0; i--)-----// bd
-----k = (long long)intx::radix * r.data[d.size()];-----// f7
                                       -----ss << setfil('0') << setw(len) << data[i];-----// b6
```

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-----if (x == n - 1) { ok = true; break; }-----// 74
5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
                                           k items out of a total of n items.
                                           -----if (!ok) return false:-----// 00
int nck(int n, int k) {-----// f6
                                           ----} return true; }------// bc
----if (n - k < k) k = n - k;------// 18
----int res = 1:-----// cb
                                           5.7. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
----for (int i = 1; i <= k; i++) res = res * (n - (k - i)) / i:-----// bd
                                           vi prime_sieve(int n) {------// 40
----return res:-----// e4
                                           ----int mx = (n - 3) >> 1, sq, v, i = -1;------// 27
}-----// 03
                                           ----vi primes:------// 8f
                                           ----bool* prime = new bool[mx + 1];-----// ef
5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
                                           ----memset(prime, 1, mx + 1);------// 28
integers a, b.
                                           ----if (n >= 2) primes.push_back(2);------// f4
int qcd(int a, int b) { return b == 0 ? a : qcd(b, a % b); }-----// d9
                                           ----while (++i <= mx) if (prime[i]) {-----// 73
                                           -----primes.push_back(v = (i << 1) + 3);-----// be
 The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
and also finds two integers x, y such that a \times x + b \times y = d.
                                           -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
                                           ------for (int j = sq; j <= mx; j += v) prime[j] = false; }------// 2e
int egcd(int a, int b, int& x, int& y) {------// 85
                                           ----while (++i \le mx) if (prime[i]) primes.push_back((i \le 1) + 3);------// 29
----if (b == 0) { x = 1; y = 0; return a; }------// 7b
                                           ----delete[] prime; // can be used for O(1) lookup-----// 36
----else {------// 00
                                           ----return primes; }-----// 72
------int d = egcd(b, a % b, x, y);------// 34
-----x -= a / b * y;------// 4a
                                           5.8. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
-----swap(x, y);-----// 26
                                           #include "egcd.cpp"-----// 55
-----return d:-----// db
                                              ·----// e8
                                           int mod_inv(int a, int m) {------// 49
}-----// 40
                                           ----int x, y, d = egcd(a, m, x, y);-----// 3e
5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is
                                           ----if (d != 1) return -1;------// 20
                                           ----return x < 0 ? x + m : x;-----// 3c
prime.
bool is_prime(int n) {------// 6c
----if (n < 2) return false;-----// c9
                                           5.9. Modular Exponentiation. A function to perform fast modular exponentiation.
----if (n < 4) return true;------// d9
                                           template <class T>-----// 82
----if (n % 2 == 0 || n % 3 == 0) return false;------// 0f
                                           T mod_pow(T b, T e, T m) {-----// aa
----if (n < 25) return true;------// ef
                                           ----T res = T(1);------// 85
----int s = static_cast<int>(sqrt(static_cast<double>(n))):------// 64
                                           ----while (e) {------// b7
----for (int i = 5; i <= s; i += 6)-----// 6c
                                           -----if (e & T(1)) res = mod(res * b, m);------// 41
------if (n % i == 0 || n % (i + 2) == 0) return false;------// e9
                                           -----b = mod(b * b, m), e >>= T(1);  }-----// b3
----return true; }-----// 43
                                           ----return res:-----// eb
5.6. Miller-Rabin Primality Test. The Miller-Rabin probabilistic primality test.
#include "mod_pow.cpp"-----// c7
                                           5.10. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
bool is_probable_prime(ll n, int k) {------// be
                                           #include "egcd.cpp"-----// 55
----if (~n & 1) return n == 2;-----// d1
                                           int crt(const vi& as, const vi& ns) {------// c3
----if (n <= 3) return n == 3;-----// 39
                                           ----int cnt = size(as), N = 1, x = 0, r, s, l;-----// 55
----int s = 0: ll d = n - 1:-----// 37
                                           ----for (int i = 0; i < cnt; i++) N *= ns[i]; ------// 88
----while (~d & 1) d >>= 1, s++;------// 35
                                           ----for (int i = 0; i < cnt; i++)------// f9
----while (k--) {-------// c8
------ll a = (n - 3) * rand() / RAND_MAX + 2;------// 06
                                           ----return mod(x, N); }------// 9e
-----ll x = mod_pow(a, d, n);-----// 64
```

5.11. Linear Congruence Solver. A function that returns all solutions to  $ax \equiv b \pmod{n}$ , modulo n.

5.12. **Numeric Integration.** Numeric integration using Simpson's rule.

5.13. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete Fourier transform. Note that this implementation only handles powers of two, make sure to pad with zeros.

```
#include <complex>-----// 8e
typedef complex<long double> cpx;------// 25
void fft(cpx *x, int n, bool inv=false) {------// 23
----for (int i = 0, j = 0; i < n; i++) {------// f2
-----if (i < j) swap(x[i], x[i]);-----// 5c
------int m = n>>1;------// e5
-------while (1 <= m && m <= j) j -= m, m >>= 1;-------// fe
-----j += m:-----// ab
----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
-----cpx wp = \exp(\text{cpx}(0, (\text{inv }? -1 : 1) * \text{pi }/\text{mx})), \text{ w} = 1;
-----for (int m = 0; m < mx; m++, w *= wp) {------// 40
-----for (int i = m; i < n; i += mx << 1) {------// 33
-----cpx t = x[i + mx] * w;-----// f5
-----x[i + mx] = x[i] - t;
-----x[i] += t:-----// c7
----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);-----// 3e
```

#### 5.14. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once:  $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times:  $n^k$

- Number of permutations of n objects, where there are  $n_1$  objects of type 1,  $n_2$  objects of type  $2, \ldots, n_k$  objects of type k:  $\binom{n}{n_1, n_2, \ldots, n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$$

- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times:  $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to  $x_1 + x_2 + \cdots + x_n = k$  where  $x_i \geq 0$ :  $f_k^n$
- Number of subsets of a set with n elements:  $2^n$
- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an  $n \times m$  grid by walking only up and to the right:  $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced:  $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} \binom{2n}{n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an  $n \times n$  lattice which do not rise above the main diagonal:  $C_n$
- Number of permutations of n objects with exactly k ascending sequences or runs:  $\binom{n}{k} = \binom{n}{n-k-1} = k \binom{n-1}{k} + (n-k+1) \binom{n-1}{k-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{n-1} = 1$
- Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements):  $D_0 = 1, D_1 = 0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points:  $\binom{n}{k}D_{n-k}$
- Jacobi symbol:  $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- **Heron's formula:** A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} 1$ .
- **Divisor sigma:** The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where  $n = \prod_{i=0}^r p_i^{a_i}$  is the prime factorization.
- Divisor count: A special case of the above is  $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$  where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- $\gcd(2^a 1, 2^b 1) = 2^{\gcd(a,b)} 1$

5.15. **Numbers and Sequences.** Some random prime numbers: 1031, 32771, 1048583, 33554467, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.

## 6. Geometry

6.1. **Primitives.** Geometry primitives.

```
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#define P(p) const point point point point <math>point point point point point point <math>point point point point point point point point point <math>point point poin
double dot(P(a), P(b)) { return real(conj(a) * b); }------// a9 -----x = min(x, abs(d - closest_point(a,b, d, true)));------// cd
point rotate(P(p), P(about), double radians) {-------// e1 ---return x;-------// 9e
----return (p - about) * exp(point(0, radians)) + about; }-----// cb }------// cb
point reflect(P(p), L(about1, about2)) {-----// c0
                                                         6.2. Polygon. Polygon primitives.
----point z = p - about1, w = about2 - about1;-----// 39
----return conj(z / w) * w + about1; }------// 03 #include "primitives.cpp"------// e0
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// fc typedef vector<point> polygon;-----// b3
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }------// ca ----double area = 0; int cnt = size(p);-----// a2
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// 75 ----for (int i = 1; i + 1 < cnt; i++)------// d2
bool collinear(L(a, b), L(p, q)) {------// 66 -----area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 7e
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6 ----return area / 2; }------// e1
double angle(P(a), P(b), P(c)) {-------// d0 double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25
double signed_angle(P(a), P(b), P(c)) {------// fe int point_in_polygon(polygon p, point q) {------// 58
double angle(P(p)) { return atan2(imag(p), real(p)); }------// fc ----for (int i = 0, j = n - 1; i < n; j = i++)------// 77</pre>
point perp(P(p)) { return point(-imag(p), real(p)); }------// 79 -----if (collinear(p[i], q, p[j]) &\dagger--------// a5
double progress(P(p), L(a, b)) {------// 8e -----// 8e -----// 8e d <= 1)-----// b9
----if (abs(real(a) - real(b)) < EPS)--------// bc ------return 0;------
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));-------// 36 ----for (int i = 0, j = n - 1; i < n; j = i++)------// 6f
----else return (real(p) - real(a)) / (real(b) - real(a)); }--------// 58 ------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))-------// 1f
----// NOTE: check for parallel/collinear lines before calling this function---// 79 ----return in ? -1 : 1; }--------------------// 77
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// de //---- polygon left, right;------// 6b
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// d7 //---- point it(-100, -100);-------// c9
-----return false;-----// 00 //---- for (int i = 0, cnt = poly.size(); i < cnt; i++) {-------// 28
------if (dot(b - a, c - b) > 0) return b:-------// 83 //------// 83 //------// f2
-----if (dot(a - b, c - a) > 0) return a;------// d4 //----- if (myintersect(a, b, p, q, it))------// f0
----double t = dot(c - a, b - a) / norm(b - a);------// 22 //----}
----return a + t * (b - a);------// d7 //---- return pair<polygon, polygon>(left, right);------// 1d
}------// 20 // }------// 37
double line_segment_distance(L(a,b), L(c,d)) {-----// da
                                                          6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
----double x = INFINITY;-----// 04
                                                         #include "polygon.cpp"-----// 58
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// 17
                                                         #define MAXN 1000-----// 09
----else if (abs(a - b) < EPS) x = abs(a - closest\_point(c, d, a, true));-----// d9
                                                         point hull[MAXN];-----// 43
----else if (abs(c - d) < EPS) \times = abs(c - closest_point(a, b, c, true)); -----//7f
                                                         bool cmp(const point &a, const point &b) {-----// 32
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&------// c2
                                                          ----return abs(real(a) - real(b)) > EPS ?-----// 44
----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;
                                                          -----real(a) < real(b) : imag(a) < imag(b); }-----// 40
```

```
-----// 60
int convex_hull(polygon p) {-----// cd
----int n = size(p), l = 0;------// 67
                                                }-----// 3f
----sort(p.begin(), p.end(), cmp);-----// 3d
                                                6.6. Triangle Circumcenter. Returns the unique point that is the same distance from all three
----for (int i = 0; i < n; i++) {------// 6f
                                                points. It is also the center of the unique circle that goes through all three points.
------if (i > 0 && p[i] == p[i - 1]) continue;------// b2
                                                #include "primitives.cpp"-----// e0
------while (l >= 2 \&\& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;------// 20
                                                point circumcenter(point a, point b, point c) {-----// 76
-----hull[l++] = p[i];-----//
                                                ----b -= a, c -= a;-----// 41
----}------// d8
                                                ----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a
----int r = 1:-----// 59
----for (int i = n - 2; i >= 0; i--) {------// 16
-----if (p[i] == p[i + 1]) continue;-----// c7
                                                6.7. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
-----while (r - l >= 1 \& \& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 9f
                                                pair of points.
-----hull[r++] = p[i]:-----// 6d
                                                #include "primitives.cpp"-----// e0
----}-----// 74
                                                       .....// 85
----return l == 1 ? 1 : r - 1;------// 6d
                                                struct cmpx { bool operator ()(const point \&a, const point \&b) {------// \theta1
}-----// 79
                                                -----return abs(real(a) - real(b)) > EPS ?-----// e9
                                                -----real(a) < real(b) : imag(a) < imag(b); } };-----// 53
6.4. Line Segment Intersection. Computes the intersection between two line segments.
                                                struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
#include "primitives.cpp"-----// e0
                                                ----return abs(imag(a) - imag(b)) > EPS ?-----// θb
bool line_segment_intersect(L(a, b), L(c, d), point &A, point &B) {------// 6c
                                                -----imag(a) < imag(b) : real(a) < real(b); } };-----// a4
----if (abs(a - b) < EPS &\& abs(c - d) < EPS) {------// db
                                                double closest_pair(vector<point> pts) {------// f1
------A = B = a; return abs(a - d) < EPS; }------// ee
                                                ----sort(pts.begin(), pts.end(), cmpx());-----// 0c
----else if (abs(a - b) < EPS) {------// 03
                                                ----set<point, cmpy> cur;-----// bd
-----A = B = a; double p = progress(a, c,d);------// c9
                                                ----set<point, cmpy>::const_iterator it, jt;-----// a6
-----return 0.0 <= p && p <= 1.0-----// 8a
                                                ----double mn = INFINITY;-----// f9
----for (int i = 0, l = 0; i < size(pts); i++) {------// ac
----else if (abs(c - d) < EPS) {------// 26
                                                ------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);-------// 8b
-----A = B = c; double p = progress(c, a,b);------// d9
                                                -----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));------// fc
-----return 0.0 <= p && p <= 1.0-----// 8e
                                                -----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));------// 39
------while (it != jt) mn = min(mn, abs(*it - pts[i])), it++;------// 09
----else if (collinear(a,b, c,d)) {------// bc
                                                -----cur.insert(pts[i]); }-----// 82
------double ap = progress(a, c,d), bp = progress(b, c,d);-----// a7
                                                ----return mn; }------// 4c
-----if (ap > bp) swap(ap, bp);-----// b1
-----if (bp < 0.0 || ap > 1.0) return false;------// 0c
                                                6.8. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
-----A = c + max(ap, 0.0) * (d - c); ------// f6
                                                   • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
-----B = c + min(bp, 1.0) * (d - c);------// 5c
                                                   • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
-----return true; }-----// ab
                                                   • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
----else if (parallel(a,b, c,d)) return false;-----// ca
                                                    of that is the area of the triangle formed by a and b.
----else if (intersect(a,b, c,d, A, true)) {------// 10
-----B = A; return true; }------// bf
                                                                  7. Other Algorithms
----return false:-----// b7
                                                7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
}-----// 8b
                                                function f on the interval [a, b], with a maximum error of \varepsilon.
      .....// e6
                                                double binary_search_continuous(double low, double high,-----// 8e
6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
                                                -----double eps, double (*f)(double)) {------// c0
coordinates) on a sphere of radius r.
                                                ----while (true) {------// 3a
double gc_distance(double pLat, double pLong,-------// 7b ------double mid = (low + high) / 2, cur = f(mid);------// 75
```

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```
Another implementation that takes a binary predicate f, and finds an integer value x on the integer 7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
interval [a, b] such that f(x) \wedge \neg f(x-1).
                                  ----assert(low <= high);-----// 19 ----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
----while (low < high) {-------// a3 ----for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)------// 05
------int mid = low + (high - low) / 2;------// 04 -----inv[i][w[i][j]] = j;-----// b9
------if (f(mid)) high = mid;------// ca ----for (int i = 0; i < n; i++) q.push(i);-----// fe
-----else low = mid + 1;------// 03 ----while (!q.empty()) {------// 55
----}------int curm = q.front(); q.pop();------// ab
----assert(f(low));------// 42 ------for (int &i = at[curm]; i < n; i++) {-------// 9a
----return low;-------// a6 ---------// cf
}------if (eng[curw] == -1) { }------// 35
                                  ------else if (inv[curw][curm] < inv[curw][enq[curw]])------// 10
7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotoni-
cally decreasing, ternary search finds the x such that f(x) is maximized.
                                  -----else continue:-----// b4
template <class F>-----// d1
                                  -----res[eng[curw] = curm] = curw, ++i; break;------// 5e
                                  -----}-------// 24
double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
                                  ----}------------------// b8
----while (hi - lo > eps) {------// 3e
                                  ----return res:-----// 95
------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
                                   -----if (f(m1) < f(m2)) lo = m1;------// 1d
-----else hi = m2:-----// b3
----}----------// bb
                                  7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the
----return hi:-----// fa
                                  Exact Cover problem.
}-----// 66
                                  bool handle_solution(vi rows) { return false; }------// 63
                                  struct exact_cover {------// 95
7.3. 2SAT. A fast 2SAT solver.
                                  ----struct node {------// 7e
#include "../graph/scc.cpp"-----// c3 -----node *l, *r, *u, *d, *p;------// 19
-----// 63 ------<mark>int</mark> row, col, size;-------// ae
bool two_sat(int n, const vii& clauses, vi& all_truthy) {------// f4 -----node(int _row, int _col) : row(_row), col(_col) {------// c9
----all_truthy.clear();------// 31 ------size = 0; l = r = u = d = p = NULL; }-----// c3
----vvi adi(2*n+1);------// 7b ---}:-----// 7b
------if (clauses[i].first != clauses[i].second)------// 87 ---node *head;-----
-----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93 ----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
----}-----arr = new bool*[rows];------// cf
----pair<union_find, vi> res = scc(adj);------// 9f -----sol = new int[rows];------// 5f
----vi daq = res.second;------// 58 -----arr[i] = new bool[cols], memset(arr[i], 0, cols);------// 75
----vi truth(2*n+1, -1);------// 00 ---}-----// 91
-----if (cur == 0) continue;------// 26 -----node ***ptr = new node**[rows + 1];------// 35
-----if (truth[p] == -1) truth[p] = 1;-------// c3 ------ptr[i] = new node*[cols];------// 0b
-----truth[cur + n] = truth[p];------// b3 -------for (int j = 0; j < cols; j++)-----// f5
------if (truth[p] == 1) all_truthy.push_back(cur);-------// 5c -----else ptr[i][j] = NULL;-------// 32
```

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-----ptr[i][j]->d = ptr[ni][j];--------// 71 -----found = search(k + 1);--------// f1
------ptr[ni][j]->u = ptr[i][j];-------// c4 ------for (node *j = r->l; j != r; j = j->l) { UNCOVER(j->p, a, b); }----// ab
-----/if (nj == cols) nj = 0;--------// e2 ------UNCOVER(c, i, j);------------------// 3a
-----ptr[i][j]->r = ptr[i][nj];-----// d5
                        7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \dots, k-1\}
 -----ptr[i][nj]->l = ptr[i][j];-----// 72
--}------// 92 vector<int> nth_permutation(int cnt, int n) {----------------// 78
-----for (int i = 0; i <= rows; i++)-----// 96
                        7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
 -----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];-----// cb
------ptr[rows][j]->size = cnt;------// 59 ii find_cycle(int x0, int (*f)(int)) {------// a5
                        ------for (int i = 0; i <= rows; i++) delete[] ptr[i];-----// bf
                        ----while (t != h) t = f(t), h = f(f(h));-----// 79
-----delete[] ptr:-----// 99
                        ----h = x0;------// 04
----while (t != h) t = f(t), h = f(h), mu++;------// 9d
----while (t != h) h = f(h), lam++;-----// 5e
------c->r->l = c->l, c->l->r = c->r; \\------// f9
                        ----return ii(mu, lam);-----// b4
-----j->d->u = j->u, j->u->d = j->d, j->\overline{p}->size--;-----// 16
                        7.8. Dates. Functions to simplify date calculations.
----#define UNCOVER(c, i, j) N------// d0
                        int intToDay(int jd) { return jd % 7; }-----// 89
------for (node *i = c->u; i != c; i = i->u) \------// ff
                        int dateToInt(int y, int m, int d) {------// 96
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
-----j->p->size++, j->d->u = j->u->d = j; \\ \]
                        -----367 * (m - 2 - (m - 14) / 12 * 12) / 12 -----// d1
                        -----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
                        -----d - 32075;-----// e0
----bool search(int k = 0) {------// bb
-----if (head == head->r) {------// c3 }-----// fa
------for (int i = 0; i < k; i++) res[i] = sol[i]:-----// 75
                        ----int x, n, i, j;------// 00
-----sort(res.begin(), res.end());------// 87
                        ---x = id + 68569;
-----return handle_solution(res);------// 51
                        ---n = 4 * x / 146097;
                        ----x -= (146097 * n + 3) / 4;-----// 58
-----node *c = head->r, *tmp = head->r;------// 8e
                        ---i = (4000 * (x + 1)) / 1461001;
                        ----x -= 1461 * i / 4 - 31;-----// 09
```

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j = 80 * x / 2447;//	3d
d = x - 2447 * j / 80;//	
x = j / 11;//	b7
m = $j$ + 2 - 12 * $x$ ;//	
y = 100 * (n - 49) + i + x;	70
1//	af

#### 8. Useful Information

## 8.1. Tips & Tricks.

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- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo  $m_1, m_2, \ldots, m_k$ , where  $m_1, m_2, \ldots, m_k$  are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When n = 0, n = -1, n = 1,  $n = 2^{31} 1$  or  $n = -2^{31}$ ? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. **Fast Input Reading.** If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

8.3. **128-bit Integer.** GCC has a 128-bit integer data type named \_\_int128. Useful if doing multiplication of 64-bit integers, or something needing a little more than 64-bits to represent.

# 8.4. Worst Time Complexity.

	1 0	
n	Worst AC Algorithm	Comment
$\leq 10$	$O(n!), O(n^6)$	e.g. Enumerating a permutation
$\leq 15$	$O(2^n \times n^2)$	e.g. DP TSP
$\leq 20$	$O(2^n), O(n^5)$	e.g. $DP + bitmask technique$
$\leq 50$	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
$\le 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree
$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

## 8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.