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```
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          1. Code Templates
                           ----public static void main(String[] args) throws Exception {-------// 02
                           -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                           ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                           -----// code-----// e6
setxkbmap -option caps:escape
                           -----out.flush():-----// 56
set -o vi
                           xset r rate 150 100
                           }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                      2. Data Structures
syn on | colorscheme slate
                           2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                           struct union find {-----// 42
#include <cassert>-----------------// 65 ----int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }-------// ba
#include <cmath>-----// 7d ----bool unite(int x, int y) {-------// 6c
#include <cstdio>------// 2e ------int xp = find(x), yp = find(y);------// 64
#include <cstdlib>------// 11 ------if (xp == yp) return false;------// 0b
#include <cstring>------// d0 -------if (p[xp] < p[yp]) p[xp] += p[yp], p[yp] = xp;------// 3a
#include <ctime>------// 28 -------| p[xp] += p[xp], p[xp] = yp;------// 3e
#include <iostream>------// ec ----int size(int x) { return -p[find(x)]; } };------// 28
#include <map>-----// 02
#include <stack>-----// cf int f(int a, int b) { return min(a, b); }-------// 4f
#include <vector>-----// 4f int f(int a, int b) { return a + b; }-----// dd
using namespace std;------// 7b #endif-----// 7b
-----// 7e struct segment_tree {------------------------// ab
const double pi = acos(-1);------// 49 ----int mk(const vi &arr, int l, int r, int i) {------// 12
typedef unsigned long long ull;------// 81 -----int m = (l + r) / 2;-----// de
typedef vector<vi>vvi;------// 31 ------propagate(l, r, i);-------// 12
typedef vector<vii>vvii;-------// 4b ------if (r < a || b < l) return ID;------// c7
template <class T> T mod(T a, T b) { return (a % b + b) % b; }------// 70 -----if (a <= l && r <= b) return data[i];----------// ce
template <class T> int size(const T &x) { return x.size(); }-----// 68 -----int m = (l + r) / 2;------// 7a
                           -----return f(q(a, b, l, m, 2*i+1), q(a, b, m+1, r, 2*i+2)); }------// 5c
1.3. Java Template. A Java template.
                           ----void update(int i, int v) { u(i, v, 0, n-1, 0); }-----// 90
-----// a3 ------if (l == i && r == i) return data[j] = v;--------// 4a
```

```
2.4. Matrix. A Matrix class.
```

```
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----return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); \}----// 34
----void range_update(int a, int b, int v) { ru(a, b, v, 0, n-1, 0); }------// 71
----int ru(int a, int b, int v, int l, int r, int i) {-------// e0
-----propagate(l, r, i);-----// 19
-----if (l > r) return ID;------// cc
-----if (r < a || b < l) return data[i];-----// d9
-----if (l == r) return data[i] += v;-----// 5f
-----if (a <= l \& r <= b) return (lazy[i] = v) * (r - l + 1) + data[i]:----// 76
-----int m = (l + r) / 2;-----// e7
-----return data[i] = f(ru(a, b, v, l, m, 2*i+1),-----// θe
----}------// 47
----void propagate(int l, int r, int i) {-----// b5
-----if (l > r || lazy[i] == INF) return;-----// 83
-----data[i] += lazy[i] * (r - l + 1);-----// 99
-----if (l < r) {------// dd
------else lazy[2*i+1] += lazy[i];-----// 72
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];-----// dd
------else lazy[2*i+2] += lazy[i];-----// a4
-----lazv[i] = INF:-----// c4
}:-----// 17
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
i...j in O(\log n) time. It only needs O(n) space.
struct fenwick_tree {------// 98
----int n; vi data;------// d3 ------return res; }-----
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {--------// dd
----void update(int at, int by) {------// 76 -----matrix<T> res(cols, rows);-----// b5
-------while (at < n) data[at] += by, at |= at + 1; }-------// fb -------for (int i = 0; i < rows; i++)------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);------// bf
```

```
template <> bool eq<double>(double a, double b) { return abs(a - b) < EPS; }---// a7
                                                template <class T>-----// 53
                                                class matrix {------// 85
                                                public:----// be
                                                ----int rows, cols;------// d3
                                                ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {------// 34
                                                -----data.assign(cnt, T(0)); }-----// d0
                                                ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),-----// fe
                                                -----cnt(other.cnt), data(other.data) { }-----// ed
                                                ----T& operator()(int i, int j) { return at(i, j); }------// e0
                                                ----void operator +=(const matrix& other) {------// c9
                                                ------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
                                                ----void operator -=(const matrix& other) {------// 68
                                                ------for (int i = 0: i < cnt: i++) data[i] -= other.data[i]: }------// 88
                                                ----void operator *=(T other) {------// ba
                                                ------for (int i = 0; i < cnt; i++) data[i] *= other; }------// 40
                                                ----matrix<T> operator +(const matrix& other) {------// ee
                                                ------matrix<T> res(*this); res += other; return res; }------// 5d
                                                ----matrix<T> operator -(const matrix& other) {------// 8f
                                                ------matrix<T> res(*this); res -= other; return res; }------// cf
                                               ----matrix<T> operator *(T other) {------// be
                                                ------matrix<T> res(*this); res *= other; return res; }------// 37
                                                ----matrix<T> operator *(const matrix& other) {------// 95
                                                ------matrix<T> res(rows, other.cols);------// 57
                                                -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
                                                -----for (int k = 0; k < cols; k++)-----// fc
                                               -----res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
};------// 57 -----while (p) {------// cb
struct fenwick_tree_sq {------// d4 -----if (p & 1) res = res * sq;-----// c1
----<mark>int</mark> n; fenwick_tree x1, x0;------// 18 -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {------// 89
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }-----// 45 ------matrix<T> mat(*this); det = T(1);--------// 21
};-----// 13 ------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 ------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f ------if (k >= rows) continue;-------// 3f
----return s.query(b) - s.query(a-1); }-----// f3 ------det *= T(-1);--------------------// 7a
```

```
template <class K> bool eq(K a, K b) { return a == b; }-----// 2a
```

```
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------for (int i = 0; i < cols; i++)-------// ab ----void erase(node *n, bool free = true) {-------// 58
------if (!eq<T>(mat(r, c), T(1)))-------// 2c -----else if (n->l && !n->r) parent_leq(n) = n->l, n->l->p = n->p;-----// 6b
------for (int i = cols-1; i >= c; i--) mat(r, i) /= mat(r, c);-----// 5d ------else if (n->l && n->r) {---------------------------// 6c
------for (int i = 0; i < rows; i++) {----------// 3d ------node *s = successor(n);--------// e5
------T m = mat(i, c);--------// e8 ------erase(s, false);---------// 0a
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 5a
------} return mat; }-------// 8f ------parent_leg(n) = s, fix(s);-------// 82
private:-----// e0 -----return;-------// e5
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;-------// 43
-----if (!n) return NULL;-----// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                           -----if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0------// b5 -----node *p = n->p;------// a7
template <class T>-----// 22 -----return p; }-----// c7
class avl_tree {------// ff ----node* predecessor(node *n) const {------// b4
public:-----// f6 ------if (!n) return NULL;-------// dd
----struct node {-------// 45 -----if (n->l) return nth(n->l->size-1, n->l);------// 10
------int size, height;------// 33 -------while (p && p->l == n) n = p, p = p->p;------// 6d
------l(NULL), r(NULL), size(1), height(0) { } };-------// @d ----inline int size() const { return sz(root); }------// ef
----node *root;------// 91 ----node* nth(int n, node *cur = NULL) const {------// e4
-----node *cur = root;------// b4 ------while (cur) {------// 29
------while (cur) {-------// 8b ------if (n < sz(cur->l)) cur = cur->l;------// 75
-----if (cur->item < item) cur = cur->r;------// 71 -----else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-----// cd
------else if (item < cur->item) cur = cur->l;------// cd -----else break;-----
------else break; }------// 4f ------} return cur; }------// ed
------return cur; }-------// 84 ----int count_less(node *cur) {-------// ec
----node* insert(const T &item) {-------// 4e -----int sum = sz(cur->l);------// bf
-----prev = *cur;-----// f0 -----cur = cur->p;-----// eb
#if AVL_MULTISET-----// 0a private:----// d5
#else------// ff ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
------else return *cur;------// 54 -----return n && height(n->r); }------// a8
-----node *n = new node(item, prev);-------// eb ----inline bool too_heavy(node *n) const {------// @b
-----*cur = n, fix(n); return n; }-----// 29
                           -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }-----// 67
```

```
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------if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }------// ef #define SWP(x,y) tmp = x, x = y, y = tmp------// fb
------if (n->p->l == n) return n->p->l;-------// 83 ----bool operator ()(const int &a, const int &b) { return a < b; } };------// e9
------if (n->p->r == n) return n->p->r;-------// cc template <class Compare = default_int_cmp>------// 30
------if (!n) return;---------// @e ----int len, count, *q, *loc, tmp;--------// @a
------n->height = 1 + max(height(n->l), height(n->r)); }-------// 41 ----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }------// a0
------while (i > 0) {------// 1a
-----parent_leg(n) = l; \[ \]-----// fc
                               -----int p = (i - 1) / 2;-----// 77
-----augment(n), augment(l)-------// 81 ------while (true) {---------------------// 3c
----void fix(node *n) {-------// 0d -------int m = r >= count || cmp(l, r) ? l : r;------// cc
------while (n) { augment(n);-------// 69 ------if (!cmp(m, i)) break;------// 42
------if (too_heavy(n)) {-------// 4c -----swp(m, i), i = m; } }-----// 1d
------if (left_heavy(n) && right_heavy(n->l)) left_rotate(n->l);-----// a9 public;------
------else if (right_heavy(n) && left_heavy(n->r))-------// b9 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {------// b9
-----right_rotate(n->r);------// 08 ------q = new int[len], loc = new int[len];------// f8
------if (left_heavy(n)) right_rotate(n);------// 93 ------memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// d5 ----~heap() { delete[] q; delete[] loc; }------// 09
------n = n->p; }-------// 28 ----void push(int n, bool fix = true) {-------// b7
#ifdef RFSI7F-----// a9
Also a very simple wrapper over the AVL tree that implements a map interface.
                                -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"-----// 01
                                -----while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                                -----int *newq = new int[newlen], *newloc = new int[newlen];-----// e3
template <class K, class V>-----// da
                                -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --//94
class avl_map {------// 3f
                                -----memset(newloc + len, 255, (newlen - len) << 2);-----// 18
public:----// 5d
                                -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                                -----loc = newloc, q = newq, len = newlen;-----// 61
------K key; V value;------// 32
                               #else-----// 54
-----node(K k, V v) : key(k), value(v) { }-----// 29
                                -----assert(false):-----// 84
-----bool operator < (const node &other) const { return key < other.key; } };// 92
                                #endif------// 64
----avl_tree<node> tree:-----// b1
                                ----V& operator [](K key) {------// 7c
                                -----assert(loc[n] == -1);-----// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
                               -----loc[n] = count, q[count++] = n;-----// 6b
------if (!n) n = tree.insert(node(key, V(0)));------// cb
                                ------if (fix) swim(count-1); }------// bf
-----return n->item.value;------// ec
                               ----}------// 2e
                                -----assert(count > 0):-----// eb
}:-----// af
                               ------loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 50
                                -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
```

```
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----}------FIND_UPDATE(x->next[i]->item, target);-------// 3a
----void heapify() { for (int i = count - 1; i > 0; i--)----------// 39 ------int lvl = bernoulli(MAX_LEVEL);----------------------// 7a
------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }--------// 0b ------if(lvl > current_level) current_level = lvl;-----------------------// 8a
----void update_key(int n) {--------------------------// 26 -----x = new node(lvl, target);-------------------// 36
----void clear() { count = 0, memset(loc, 255, len << 2); } };-------// 58 ------update[i] = x;-----------// 20
                                      -----update[i]->lens[i] = pos[0] + 1 - pos[i];------// 42
2.7. Skiplist. An implementation of a skiplist.
                                      ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// 07
#define BP 0.20-----// aa
                                      -----size++;-----// 19
#define MAX_LEVEL 10------// 56
                                      -----return x; }-----// c9
unsigned int bernoulli(unsigned int MAX) {------// 7b
                                      ----void erase(T target) {------// 4d
----unsigned int cnt = 0;-----// 28
----while(((float) rand() / RAND_MAX) < BP \&\& cnt < MAX) cnt++;-----// d1
                                      ------FIND_UPDATE(x->next[i]->item, target);-------// 6b
                                      -----if(x && x->item == target) {-----// 76
----return cnt; }-----// a1
template<class T> struct skiplist {------// 34
                                      ------for(int i = 0; i <= current_level; i++) {-------// 97
                                      -----if(update[i]->next[i] == x) {------// b1
----struct node {------// 53
                                      -----update[i]->next[i] = x->next[i];------// 59
                                      -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
-----int *lens:-----// 07
                                      -----} else update[i]->lens[i] = update[i]->lens[i] - 1;-----// 88
                                      ------#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))-------// 25
                                      -----delete x; _size--;------// 81
-----node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}-----// 7c
                                      ------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
-----/node() { free(lens); free(next); }; };-----// aa
                                      -----current_level--; } } };-----// 59
----int current_level, _size;------// 61
----node *head;------// b7
                                      2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                      list supporting deletion and restoration of elements.
-----skiplist() { clear(); delete head; head = NULL; }------// aa
                                      template <class T>-----// 82
----#define FIND_UPDATE(cmp, target) \|-----// c3
                                      struct dancing_links {-----// 9e
------int pos[MAX_LEVEL + 2]; \[\[\]\------// 18
                                      ----struct node {------// 62
-----memset(pos, 0, sizeof(pos)); \-----// f2
                                      -----T item:-----// dd
-----node *update[MAX_LEVEL + 1]; \| ------// 01
                                      -----node(const T &_item, node *_l = NULL, node *_r = NULL)------// 6d
-----memset(update, 0, MAX_LEVEL + 1); \[\bar{\}\]------// 38
                                      -----: item(_item), l(_l), r(_r) {------// 6d
                                      -----if (l) l->r = this;-----// 97
-----for(int i = MAX\_LEVEL; i >= 0; i--) { \sqrt{\phantom{a}}
                                      -----if (r) r->l = this;-----// 81
-----pos[i] = pos[i + 1]; \\------// 68
                                      ------}-------// 2d
----};------// d3
------update[i] = x; N-----------// dd ----dancing_links() { front = back = NULL; }------// 72
----void clear() { while(head->next && head->next[0])-------// 91 -----if (!front) front = back;-------------// d2
------erase(head->next[0]->item); }-------// e6 ------return back;--------------------------// cθ
------return x && x->item == target ? x : NULL; }-------// 50 ----node *push_front(const T &item) {--------// 4a
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------if (!back) back = front;--------------------// 10
-----return pos[0]; }------// 19 -----return front;------// cf
```

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------if (!n->l) front = n->r; else n->l->r = n->r;-------// ab -------bb(pt _from, pt _to) : from(_from), to(_to) {}------// 57
------if (!n->l) front = n; else n->l->r = n;--------// a5 ------if (p.coord[i] < from.coord[i])------// a0
}:------sum += pow(p.coord[i] - to.coord[i], 2.0);------// 8c
                                   2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                                   -----return sqrt(sum); }-----// ef
element.
                                   ------bb bound(double l, int c, bool left) {------// b6
#define BITS 15------pt nf(from.coord), nt(to.coord);------// 5c
----int cnt[BITS][1<<BITS];------// aa ------else nf.coord[c] = max(nf.coord[c], l);------// 71
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }------// b0 -----return bb(nf, nt); } };------// 3b
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49 ------pt p; node *\, *r;----------------------------// 46
----int nth(int n) {-------// 8a ------node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
------int res = 0;-------// a4 ----node *root;------// 30
------for (int i = BITS-1; i >= 0; i--)------// 99 ----// kd_tree() : root(NULL) { }------// 97
------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4 ----kd_tree(vector<pt> pts) { root = construct(pts, 0, size(pts) - 1, 0); }----// 19
-----return res:------// 3a ----node* construct(vector<pt> &pts, int from, int to, int c) {------// 4e
----}-----if (from > to) return NULL;------// af
}:------// 0a ------int mid = from + (to - from) / 2;-----// 7d
                                   -----nth_element(pts.begin() + from, pts.begin() + mid,-----// d8
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                                   -----pts.begin() + to + 1, cmp(c));-----// 84
bor queries.
                                   -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// f1
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1) - \cdots / 77
                                   ------construct(pts, mid + 1, to, INC(c))); }-----// 50
template <int K>-----// cd ----bool contains(const pt &p) { return _con(p, root, 0); }------// 8a
class kd_tree {------// 7e ----bool _con(const pt &p, node *n, int c) {------// ff
-------double coord[K];-------// d6 ------if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// ae
-----pt() {}------// c1 -----return true; }------// 8e
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }-----// 4c ----void insert(const pt &p) { _ins(p, root, 0); }------// e9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {------// 7d
-----double sum = 0.0;------// c4 -----if (!n) n = new node(p, NULL, NULL);------// 29
------for (int i = 0; i < K; i++)-------// 23 ------else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));------// 13
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }------// f8
------return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }-----// 15
----struct cmp {------// 8f ----void _clr(n->l), _clr(n->r), delete n; }-----// 92
-----cmp(int _c) : c(_c) {}------// a5 -----assert(root);----------// 24
------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// 0d
-----cc = i == 0 ? c : i - 1;------// bc -----pt from(cs);------// af
-----return a.coord[cc] < b.coord[cc];------// b7 -----pt to(cs);------
-----return false; } };------// e2 ____}
----struct bb {------// 30
```

```
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----pair<pt, bool> _nn(-------// cd ----T.erase(T.begin() + i);--------// ca
                                                }-----// 9a
-----const pt &p, node *n, bb b, double &mn, int c, bool same) {------// 1d
-----if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// c5
                                                                      3. Graphs
------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;------// 6d
-----pt resp = n->p;------// 3d
                                                 3.1. Single-Source Shortest Paths.
------if (found) mn = min(mn, p.dist(resp));------// c9
-----node *n1 = n->l, *n2 = n->r;-----// dc
                                                3.1.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|) time.
-----for (int i = 0: i < 2: i++) {------// 74
                                                int *dist. *dad:-----// 46
------if (i == 1 || cmp(c)(n->p, p)) swap(n1, n2), swap(l1, l2);------// ab
                                                 struct cmp {------// a5
-----pair<pt, bool> res =-----// f0
                                                 ----bool operator()(int a, int b) {------// bb
-----nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// ad
                                                 -----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
-----if (res.second && (!found || p.dist(res.first) < p.dist(resp)))----// 17
                                                 }:-----// 41
-----resp = res.first, found = true;-----// 62
                                                 pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
------}------// b7
                                                 ----dist = new int[n];-----// 84
-----return make_pair(resp, found); } };-----// c8
                                                 ----dad = new int[n];-----// 05
                                                 ----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;-------// d6
2.11. Sqrt Decomposition. Design principle that supports many operations in amortized \sqrt{n} per
                                                 ----set<<u>int</u>, cmp> pq;-----// 04
operation.
                                                 ----dist[s] = 0, pq.insert(s);-----// 1b
struct segment {-----// b2
                                                 ----while (!pq.empty()) {------// 57
----vi arr:-----// 8c
                                                 ------int cur = *pq.begin(); pq.erase(pq.begin());-----// 7d
----segment(vi _arr) : arr(_arr) { } };-----// 11
                                                 ------for (int i = 0; i < size(adj[cur]); i++) {------// 9e
vector<segment> T:-----// a1
                                                 -----int nxt = adj[cur][i].first,-----// b8
int K;-----// dc
                                                 -----ndist = dist[cur] + adj[cur][i].second;------// \theta c
void rebuild() {------// 17
                                                 -----if (ndist < dist[nxt]) pq.erase(nxt),------// e4
----int cnt = 0;-----// 14
                                                 -----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);------// 0f
----for (int i = 0; i < size(T); i++)-----// 7d
                                                 -----cnt += size(T[i].arr);-----// 1e
                                                 ----}-----// e8
----K = static_cast<int>(ceil(sqrt(cnt)) + 1e-9);------// 76
                                                 ----return pair<int*, int*>(dist, dad);-----// cc
----vi arr(cnt);------// 41
                                                 }-----// af
----for (int i = 0, at = 0; i < size(T); i++)-----// 24
------for (int j = 0; j < size(T[i].arr); j++)------// 76
                                                3.1.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
-----arr[at++] = T[i].arr[j];-----// 89
                                                 problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
----T.clear();------// b5
                                                negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
----for (int i = 0; i < cnt; i += K)-----// 9f
-----T.push_back(segment(vi(arr.begin()+i, arr.begin()+min(i+K, cnt))));----// 77
                                                 int* bellman_ford(int n. int s. vii* adi. bool& has_negative_cycle) {------// cf
}-----// e5
                                                 ----has_negative_cycle = false;-----// 47
int split(int at) {------// 64
                                                 ----int* dist = new int[n];-----// 7f
----int i = 0;-----// f7
                                                 ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
----while (i < size(T) && at >= size(T[i].arr))------// a7
                                                 ----for (int i = 0; i < n - 1; i++)-----// a1
-----at -= size(T[i].arr), i++;-----// 38
                                                 ------for (int j = 0; j < n; j++)------// c4
----if (i >= size(T)) return size(T);------// 89
                                                 -----if (dist[j] != INF)-----// 4e
----if (at == 0) return i;------// 05
                                                 ------for (int k = 0; k < size(adj[j]); k++)------// 3f
----T.insert(T.begin() + i + 1, segment(vi(T[i].arr.begin() + at, T[i].arr.end())));
                                                 -----dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61
----T[i] = segment(vi(T[i].arr.begin(), T[i].arr.begin() + at));------// 60
                                                 -----dist[j] + adj[j][k].second);------// 47
----return i + 1:-----// c0
                                                 ----for (int j = 0; j < n; j++)-----// 13
}-----// 00
                                                 ------for (int k = 0; k < size(adj[j]); k++)------// a0
void insert(int at, int v) {-----// 87
                                                 ------if (dist[i] + adj[i][k].second < dist[adj[i][k].first])------// ef
----vi arr; arr.push_back(v);------// 30
                                                 -----has_negative_cvcle = true:-----// 2a
----T.insert(T.begin() + split(at), segment(arr));------// 2a
                                                 ----return dist:-----// 2e
}-----// bd
                                                 }-----// c2
void erase(int at) {------// f4
----int i = split(at); split(at + 1);------// 48
                                                3.2. All-Pairs Shortest Paths.
```

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                                                 #define MAXN 5000----// f7
3.2.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                                 int low[MAXN], num[MAXN], curnum;-----// d7
problem in O(|V|^3) time.
                                                 void dfs(const vvi &adj, vi &cp, vii &bri, int u, int p) {------// 22
void floyd_warshall(int** arr, int n) {------// 21
                                                 ----low[u] = num[u] = curnum++;-----// a3
----for (int k = 0; k < n; k++)------// 49
                                                 ----int cnt = 0; bool found = false;-----// 97
------for (int i = 0; i < n; i++)------// 21
                                                 ----for (int i = 0; i < size(adj[u]); i++) {-------// f3
-----for (int j = 0; j < n; j++)-----// 77
                                                 -----int v = adj[u][i];-----// 26
-----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
                                                 -----if (num[v] == -1) {------// f9
-----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                                                 -----dfs(adj, cp, bri, v, u);-----// 7b
}-----// 86
                                                 -----low[u] = min(low[u], low[v]);------// ea
                                                 -----cnt++:-----// 8f
3.3. Strongly Connected Components.
                                                 -----found = found || low[v] >= num[u];-----// fd
                                                 ------if (low[v] > num[u]) bri.push_back(ii(u, v));-------// 52
3.3.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                                 -----} else if (p != v) low[u] = min(low[u], num[v]); }------// c4
graph in O(|V| + |E|) time.
                                                 ----if (found && (p !=-1 \mid \mid cnt > 1)) cp.push_back(u); }-------// dc
#include "../data-structures/union_find.cpp"---------------------// 5e
                                                 pair<vi,vii> cut_points_and_bridges(const vvi &adj) {------// 35
-----// 11
                                                 ----int n = size(adj);-----// 34
vector<br/>bool> visited;------// 66
                                                 ----vi cp; vii bri;------// 63
                                                 ----memset(num, -1, n << 2);------// 4e
-----// a5
                                                 ----curnum = 0:-----// 43
void scc_dfs(const vvi &adj, int u) {-----// a1
                                                 ----for (int i = 0: i < n: i++) if (num[i] == -1) dfs(adj. cp. bri. i. -1):----// e5
----int v; visited[u] = true;-----// e3
                                                 ----return make_pair(cp, bri); }------// 70
----for (int i = 0; i < size(adj[u]); i++)-----// c5
                                                 3.5. Minimum Spanning Tree.
-----if (!visited[v = adj[u][i]]) scc_dfs(adj, v);------// 6e
----order.push_back(u):-----// 19
                                                 3.5.1. Kruskal's algorithm.
}-----// dc
                                                 #include "../data-structures/union_find.cpp"-----// 5e
                                                      -----// 11
pair<union_find, vi> scc(const vvi &adj) {------// 3e
                                                 // n is the number of vertices-----// 18
----int n = size(adj), u, v;-----// bd
                                                 // edges is a list of edges of the form (weight, (a, b))-----// c6
----order.clear():-----// 22
                                                 // the edges in the minimum spanning tree are returned on the same form-----// 4d
----union_find uf(n);------// 6d
                                                 vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
----vi dag:-----// ae
                                                 ----union_find uf(n);------// 04
----vvi rev(n);------// 20
                                                 ----sort(edges.begin(), edges.end());-----// 51
----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++)-----// b9
                                                 ----vector<pair<int, ii> > res;-------// 71
-----rev[adj[i][j]].push_back(i);-----// 77
                                                 ----for (int i = 0; i < size(edges); i++)-----// ce
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04
                                                 ------if (uf.find(edges[i].second.first) !=-----// d5
----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);------// e4
                                                 -----// 8c
----fill(visited.begin(), visited.end(), false);-----// c2
                                                 -----res.push_back(edges[i]);------// d1
----stack<int> S;-----// 04
                                                   -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
----for (int i = n-1; i >= 0; i--) {------// 3f
                                                        -----// 5b
------if (visited[order[i]]) continue;-----// 94
                                                 ----return res;------// 46
-----S.push(order[i]), dag.push_back(order[i]);------// 40
------while (!S.empty()) {------// 03
-----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);-----// 1b
                                                 3.6. Topological Sort.
-----for (int j = 0; j < size(adj[u]); j++)-----// 21
                                                3.6.1. Modified Depth-First Search.
-----if (!visited[v = adj[u][j]]) S.push(v);-----// e7
void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
                                                 ------bool& has_cycle) {------// a8
----return pair<union_find, vi>(uf, dag);-----// f2
                                                ---color[cur] = 1;-----// 5b
}-----// ca
                                                ----for (int i = 0; i < size(adj[cur]); i++) {------// 96
                                                 -----int nxt = adj[cur][i];-----// 53
                                                 -----if (color[nxt] == 0)------// 00
3.4. Cut Points and Bridges.
```

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                                                10
------tsort_dfs(nxt, color, adj, res, has_cycle);-------// 5b ------} else s.push(cur), cur = adj[cur][--outdeg[cur]];-------// d8
------else if (color[nxt] == 1)--------// 53 ---}------// ba
------has_cycle = true;-------// c8 ---return at == 0;----------// c8
3.8. Bipartite Matching.
----color[cur] = 2;-----// 16
----res.push(cur);-----// cb
                         3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
}-----// 9e
                         O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
.
-----// ae
                         graph, respectively.
vi tsort(int n, vvi adj, bool& has_cycle) {------// 37
                         vi* adj;-----// cc
----has_cycle = false;-----// 37
                         bool* done:----// b1
----stack<int> S;-----// 54
                         int* owner;-----// 26
----vi res:-----// d1
                         int alternating_path(int left) {-----// da
----char* color = new char[n];-----// b1
                         ----if (done[left]) return 0;------// 08
----memset(color, 0, n);-----// ce
                         ----done[left] = true;-----// f2
----for (int i = 0; i < n; i++) {------// 96
                         ----for (int i = 0; i < size(adj[left]); i++) {-------// 34
------if (!color[i]) {------// d5
                         ------int right = adj[left][i];------// b6
-----tsort_dfs(i, color, adj, S, has_cycle);-----// 40
                         ------if (owner[right] == -1 || alternating_path(owner[right])) {-------// d2
-----if (has_cycle) return res;------// 6c
                         -----owner[right] = left: return 1:-----// 26
-----}------// 70
                         ------} }-------// 7a
----}------// df
                         ----return 0; }-----// 83
----while (!S.empty()) res.push_back(S.top()), S.pop();------// 94
----return res:-----// 07
                         3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
}-----// 1f
                         ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                         #define MAXN 5000-----// f7
#define MAXE 5000-------// 87 #define dist(v) dist[v == -1 ? MAXN : v]-------// 0f
vi adj[MAXV];------// ff struct bipartite_graph {------// 2b
ii start_end() {------// 30 ---bipartite_graph(int _N, int _M) : N(_N), M(_M),------// 8d
----}-----dist(-1) = INF;-------// 96
}------foreach(u, adj[v]) if(dist(R[*u]) == INF)------// 95
bool euler_path() {------dist(R[*u]) = dist(v) + 1, q[r++] = R[*u];------// 60
---ii se = start_end();------// 45 -------------------// fe
-----cur = s.top(); s.pop();------// d7 ------if(dfs(R[*u])) {-------// d7
```

```
------return true;-------// 1f ------if(s == t) return 0;--------------// 9d
-----}-----memset(d, -1, n * sizeof(int));------// a8
------l = r = 0, d[q[r++] = t] = 0;-------// @e
----void add_edge(int i, int j) { adj[i].push_back(j); }------// 11 -------for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)------// a2
------memset(R, -1, sizeof(int) * M);-------// 39 ------memcpy(curh, head, n * sizeof(int));------// 10
-----return matching;------// fc ------if (res) reset();--------// fc ------// 21
};-----// d3 ---}
                 }:-----// 3b
3.9. Maximum Flow.
                 3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
3.9.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It computes
                 O(|V||E|^2). It computes the maximum flow of a flow network.
the maximum flow of a flow network.
                 #define MAXV 2000-----// ba
-----edge() { }------// 38 ---};------// 38 ---};-------// 38 ---}
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {-------// d3 ------memset(head = new int[n], -1, n << 2);-------// 58
-----e.reserve(2 * (m == -1 ? n : m));------// 3a
------head = new int[n], curh = new int[n];--------// 6b ----void destroy() { delete[] head; }-------// d5
------memset(head, -1, n * sizeof(int));--------// 56 ----void reset() { e = e_store; }------// 1b
----void reset() { e = e_store; }--------// 87 ------e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// bc
----void add_edge(int u, int v, int uv, int vu = 0) {-------// cd ----}------// ef
------for (int &i = curh[v], ret; i != -1; i = e[i].nxt)-------// f9 ------memset(d, -1, n << 2);------------------// 3b
------for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)-----// c6
```

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------(d[v = e[i].v] == -1 || d[u] + 1 < d[v]))------// 2f -----q.insert(s); d[s] = 0;------------// 1d
------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v]; ------// 8a -------for (int i = head[u]; i : != -1; i = e[i].nxt) {-------// \theta 2
------q.insert(v);-------------------// bc
------}------------------// da
3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
                            ------if (p[t] == -1) break:-----// 09
fied to find shortest path to augment each time (instead of just any path). It computes the maximum
                            -----int x = INF, at = p[t];-----// e8
flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with
                            ------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 32
minimum cost. Running time is O(|V|^2|E|\log|V|).
                            -----at = p[t], f += x;-----// 43
#define MAXV 2000------// ba -------while (at != -1)------// 53
int d[MAXV], p[MAXV], pot[MAXV];-----// 80
                           -----e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// 95
----bool operator ()(int i, int j) {-------// 8a -------for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];-----// ff
-----return d[i] == d[j] ? i < j : d[i] < d[j];------// 89 -----}
};-----// cf -----return ii(f, c);------// e7
struct flow_network {------// eb ___}
----struct edge {------// 9a
                           }:-----// d7
------int v, cap, cost, nxt;------// ad
------edge(int _v, int _cap, int _cost, int _nxt)------// ec
                           3.11. All Pairs Maximum Flow.
-----: v(_v), cap(_cap), cost(_cost), nxt(_nxt) { }------// c4
                            3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----}:------// ad
                            structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the
----int n, ecnt, *head;------// 46
                            maximum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----vector<edge> e, e_store;------// 4b
                           #include "dinic.cpp"-----// 58
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// dd
-----e.reserve(2 * (m == -1 ? n : m));------// e6
                            .....// 25
------memset(head = new int[n], -1, n << 2);-------// 6c bool same[MAXV];-------// 59
----void destroy() { delete[] head; }------------------------// ac ----int n = g.n, v;------------------------------// 5d
------push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;-------// 53 ------par[s].second = g.max_flow(s, par[s].first, false);-------// 38
------memset(pot, 0, n << 2);-------// 24 -----same[v = q[l++]] = true;------// c8
------int f = 0, c = 0, v;-------------------// d4 -------for (int i = q.head[v]; i != -1; i = q.e[i].nxt)-------// 33
------d[q[r++] = g.e[i].v] = 1;------// f8
-----set<int, cmp> q;------// d8 ------for (int i = s + 1; i < n; i++)------// 68
```

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                                                               13
-----if (par[i].first == par[s].first && same[i]) par[i].first = s;----// ea ------res = (loc[uat[u]] < loc[vat[v]] ? uat[u] : vat[v]), u--, v--;----// 13
------while (true) {-------// 3a ------u = parent[head[u]];------// 4b
-----cap[curl[i] = mn;------// 63 -----return f(res, values.query(loc[v] + 1, loc[u])); }-----// 47
-----mn = min(mn, par[cur].second), cur = par[cur].first;-----// 28 -----return f(query_upto(u, l), query_upto(v, l)); } };-----// 52
}------// 99 struct tarjan_olca {------// 87
----if (s == t) return 0;------// dd ----vi *adj, answers;-------// dd
----int cur = INF, at = s;------// 65 ----vii *queries;------// 66
------cur = min(cur, qh.first[at].second), at = qh.first[at].first;------// bd ----union_find uf;----------------------------------// 70
----return min(cur, qh.second[at][t]);-------// 6d ----tarjan_olca(int n, vi *_adj) : adj(_adj), uf(n) {--------------// 78
}------// a2 -----colored = new bool[n];------// 8d
                                -----ancestor = new int[n];-----// f2
3.12. Heavy-Light Decomposition.
                                -----queries = new vii[n];-----// 3e
#include "../data-structures/segment_tree.cpp"-----// 16 -----memset(colored, 0, n);------// 6e
struct HLD {-----// 25
                                ----}-----// 6b
----vi sz, head, parent, loc;------// 81 ------queries[x].push_back(ii(y, size(answers)));------// a0
----vvi below; segment_tree values;------// 96 ------queries[v].push_back(ii(x, size(answers)));------// 14
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), below(n) {--// 4f ------answers.push_back(-1);----------------------// ca
-----vi tmp(n, ID); values = segment_tree(tmp); }------// a7 ____}
----void update_cost(int u, int v, int c) {-------// 12 -----ancestor[u] = u;----
------if (parent[v] == u) swap(u, v); assert(parent[u] == v); -------// 9f -------for (int i = 0; i < size(adj[u]); i++) {-------// 2b
-----values.update(loc[u], c); }------// 9a ------// 9a ------int v = adj[u][i];-------// 38
----void csz(int u) { for (int i = 0; i < size(below[u]); i++)------// 63 -----process(v);-----
----void part(int u) {------// 37 ------ancestor[uf.find(u)] = u;-----// d8
------head[u] = curhead; loc[u] = curloc++;-------// 25 ________
------for (int i = 0; i < size(below[u]); i++)-------// a7 ------for (int i = 0; i < size(queries[u]); i++) {------// 34
------if (best == -1 || sz[below[u][i]] > sz[best]) best = below[u][i];--// 19 ------int v = queries[u][i].first;------// 38
-----if (best != -1) part(best);-------// 19 -----if (colored[v]) {------------// c5
------for (int i = 0; i < size(below[u]); i++)------// 7d -----answers[queries[u][i].second] = ancestor[uf.find(v)];-----// 33
------if (below[u][i] != best) part(curhead = below[u][i]); }------// 30 -----}
----void build() { int u = curloc = 0;------// 67
------while (parent[u] != -1) u++;-------// db ____}
-----csz(u); part(curhead = u); }-----// 5e
                                }:-----// 5f
----int lca(int u, int v) {-------// 21
-----vi uat, vat; int res = -1;-----// e2
                                              4. Strings
------while (u != -1) uat.push_back(u), u = parent[head[u]];------// e6
------while (v != -1) vat.push_back(v), v = parent[head[v]];-----// 5b
                                4.1. Trie. A Trie class.
```

```
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----struct node {-------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';-------// 8d
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {------// 46
------for (int i = 0; i < n; i++)-------// d5
public:------L[L[i].p = i].nr = ii(P[stp - 1][i],------// fc
----node* root;------i + cnt < n ? P[stp - 1][i + cnt] : -1);------// e5
----trie() : root(new node()) { }---------// 8f ------sort(L.begin(), L.end()):---------// bc
----template <class I>------(int i = 0; i < n; i++)--------// 85
----void insert(I begin, I end) {-------// 3c -------P[stp][L[i].p] = i > 0 &&------// eb
-------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;------// fd
------while (true) {--------// 67 -----}-----// 73
------else {-------// 3e ----int lcp(int x, int y) {-------// 05
-----typename map<T, node*>::const_iterator it;------// 01 -----if (x == y) return n - x;------------// 7f
-----it = cur->children.find(head);------// 77 ------for (int k = size(P) - 1; k >= 0 && x < n & y < n; k--)------// 07
------pair<T, node*> nw(head, new node());------// cd -----return res;-----
----template<class I>-----// b9
                                  4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
----int countMatches(I begin, I end) {-----// 7f
                                  state machine from a set of keywords which can be used to search a string for any of the keywords.
-----node* cur = root;-----// 32
                                  struct aho_corasick {-----// 78
------while (true) {------// bb
                                  ----struct out_node {------// 3e
-----if (begin == end) return cur->words;-----// a4
                                  -----string keyword; out_node *next;-----// f0
-----else {------// 1e
                                  -----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
-----T head = *begin;-----// 5c
                                  ----}:------// b9
-----typename map<T, node*>::const_iterator it;------// 25
                                  ----struct go_node {------// 40
-----it = cur->children.find(head);-----// d9
-----if (it == cur->children.end()) return 0;-----// 14
                                  -----map<char, qo_node*> next;------// 6b
-----begin++, cur = it->second; } } -----// 7c
                                  -----out_node *out; go_node *fail;-----// 3e
                                  -----go_node() { out = NULL; fail = NULL; }-----// 0f
----template<class I>-----// 9c
                                  ----};------// c0
----int countPrefixes(I begin, I end) {------// 85
                                  ----go_node *go;-----// b8
-----node* cur = root;-----// 95
                                  ----aho_corasick(vector<string> keywords) {------// 4b
------while (true) {------// 3e
                                  -----go = new go_node();-----// 77
-----if (begin == end) return cur->prefixes;-----// f5
                                  ------foreach(k, keywords) {-------// e4
-----else {------// 66
                                  -----qo_node *cur = qo;-----// 9d
-----T head = *begin;-----// 43
                                  -----foreach(c, *k)-----// 38
-----typename map<T, node*>::const_iterator it;------// 7a
                                  -----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d
-----it = cur->children.find(head);------// 43
                                  -----(cur->next[*c] = new qo_node());-----// 75
-----if (it == cur->children.end()) return 0;-----// 71
                                  -----cur->out = new out_node(*k, cur->out);------// 6e
-----begin++, cur = it->second; } } };-----// 26
                                  4.2. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
                                  -----queue<go_node*> q;------// 8a
struct entry { ii nr; int p; };-------// f9 ------foreach(a, go->next) q.push(a->second);------// a3
bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77 -------while (!q.empty()) {-----------------------------// 43
----suffix_array(string _s) : s(_s), n(size(s)) {--------// e5 ------q.push(s);-----------------------// 76
```

```
-----st = st->fail;-----// 3f
------s->fail = st->next[a->first];-------// 29 #define MAXN 100100-------// 29
------out->next = s->fail->out;------// 65 } *st = new state[MAXN+2];-------// 57
-----vector<string> res;------// ef ------char c = s[n++]; int p = last;------// a3
------qo_node *cur = qo;--------// 61 -------while (n - st[p].len - 2 < 0 || c != s[n - st[p].len - 2]) p = st[p].link;
------foreach(c, s) {-------// 6c ------if (!st[p].to[c-BASE]) {------// 05
-----cur = cur->fail;------// 9e -----st[p].to[c-BASE] = q;-----// bb
------if (!cur) cur = qo;------------------// 2f ------st[q].len = st[p].len + 2;-------// 86
-----cur = cur->next[*c];-------// 58 ------do { p = st[p].link;-------// c8
-----res.push_back(out->keyword);------// 0d ------else st[q].link = st[p].to[c-BASE];------// e6
-----return res;------// c1 -----last = st[p].to[c-BASE];------// 30
----}-----return 0; } };------// da
};-----// 32
                             5. Mathematics
4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values 5.1. Big Integer. A big integer class.
can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
                     struct intx {------// cf
accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
                     ----intx() { normalize(1); }------// 6c
int* z_values(const string &s) {------// 4d ----intx(string n) { init(n); }------// b9
----int n = size(s);-------// 97 ----intx(int n) { stringstream ss; ss << n; init(ss.str()); }------// 36
-----z[i] = 0;------// c9 ----static const unsigned int radix = 10000000000U;------// f0
-----l = r = i;------// a7 ----void init(string n) {------// 13
-----z[i] = r - l; r--;--------// fc -----if (n.empty()) n = "0";---------// 99
------} else if (z[i - l] < r - i + 1) z[i] = z[i - l];--------// bf ------if (n[0] == '-') res.sign = -1, n = n.substr(1);------// 3b
-----z[i] = r - l; r--; } }------// 8d --------int idx = i - j;---------// cd
```

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                                                      16
-------if (idx < 0) continue;-------// 52 ------c.data.push_back(carry % intx::radix);------// 86
-----digit = digit * 10 + (n[idx] - '0');------// 1f -----carry /= intx::radix:------// fd
-----res.data.push_back(digit);-------// 07 -----return c.normalize(sign);-------// 20
------data = res.data;-------// 7d ----intx operator -(const intx& b) const {-------// 53
------if (sign > 0 && b.sign < 0) return *this + (-b);-------// 8f
----intx& normalize(int nsign) {---------// 3b ------if (sign < 0 && b.sign < 0) return (-b) - (-*this);------// a1
------if (data.empty()) data.push_back(0);-------// fa ------if (*this < b) return -(b - *this);------// 36
------for (int i = data.size() - 1; i > 0 && data[i] == 0; i--)-------// 27 ------intx c; c.data.clear();--------// 6b
------data.erase(data.begin() + i);-------// 67 ------long long borrow = 0;------// f8
------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
----}------c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
------if (first) outs << n.data[i], first = false;------// 33 ----intx operator *(const intx& b) const {------// bd
-----stringstream ss: ss << cur:------// 8c ------long long carry = 0:------// 20
-----string s = ss.str();-------// 64 -------for (int j = 0; j < b.size() || carry; j++) {-------// c0
-----while (len < intx::dcnt) outs << '0', len++;-------// 0a ------carry += c.data[i + j];----------// 18
------c.data[i + j] = carry % intx::radix;------// 86
------if (sign != b.sign) return sign < b.sign;--------// cf ------assert(!(d.size() == 1 && d.data[0] == 0));-------// e9
------if (size() != b.size())--------// 4d ------intx q, r; q.data.assign(n.size(), 0);-----// ca
------return sign == 1 ? size() < b.size() : size() > b.size();------// 4d -------for (int i = n.size() - 1; i >= 0; i--) {--------// 1a
------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35 ------r.data.insert(r.data.begin(), 0);---------// c7
------return false;-------// ca -------long long k = θ;-------// cc
----intx operator -() const { intx res(*this); res.sign *= -1; return res; }---// 9d -------k = (long long)intx::radix * r.data[d.size()];-------// f7
------if (sign > 0 && b.sign < 0) return *this - (-b):---------// 36 -------r = r - abs(d) * k:------------------// 15
------if (sign < 0 && b.sign > 0) return b - (-*this);--------// 70 --------while (r < 0) r = r + abs(d), k-;--------// 11
-----intx c; c.data.clear();------// 18 -----}
------unsigned long long carry = 0;------// 5c -----return pair<intx, intx>(q.normalize(n.sign * d.sign), r);-----// a1
-----carry += (i < size() ? data[i] : OULL) +------// 91 ---intx operator /(const intx& d) const {------// a2
-----(i < b.size() ? b.data[i] : 0ULL);-------// 0c -----return divmod(*this,d).first; }------// 1e
```

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                                                                                      17
------return divmod(*this,d).second * sign; }-------// 5a ----for (int i = 1; i <= k; i++) res = res * (n - (k - i)) / i;-------// bd
5.1.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
#include "fft.cpp"-----// 13
                                            integers a, b.
                                            int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
intx fastmul(const intx &an, const intx &bn) {------// ab
                                             The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
----string as = an.to_string(), bs = bn.to_string();-----// 32
                                            and also finds two integers x, y such that a \times x + b \times y = d.
----int n = size(as), m = size(bs), l = 1,------// dc
                                            int egcd(int a, int b, int& x, int& y) {-----// 85
-----len = 5, radix = 100000,-----// 4f
                                            ----if (b == 0) { x = 1; y = 0; return a; }------// 7b
-----*a = new int[n], alen = 0,-----// b8
                                            ----else {------// 00
-----*b = new int[m], blen = 0;-----// 0a
                                            ------int d = eqcd(b, a % b, x, y);------// 34
----memset(a, 0, n << 2);------// 1d
                                            -----x -= a / b * y;-----// 4a
----memset(b, 0, m << 2);-----// 01
                                            -----swap(x, y);-----// 26
----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
                                            -----return d;------// db
------for (int j = min(len - 1, i); j >= 0; j--)------// 43
                                            ----}-------// 9e
-----a[alen] = a[alen] * 10 + as[i - j] - '0';-----// 14
----for (int i = m - 1; i >= 0; i -= len, blen++)-----// b6
------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
                                            5.4. Trial Division Primality Testing. An optimized trial division to check whether an integer is
-----b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
                                            prime.
----while (l < 2*max(alen,blen)) l <<= 1;------// 51
                                            bool is_prime(int n) {------// 6c
----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
                                            ----if (n < 2) return false;-----// c9
----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? <math>a[i] : 0, 0);-----// 35
                                            ----if (n < 4) return true;------// d9
----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? b[i] : 0, 0);-----// 66
                                            ----if (n % 2 == 0 || n % 3 == 0) return false;-----// Of
----fft(A, l); fft(B, l);-----// f9
                                            ----if (n < 25) return true;------// ef
----for (int i = 0; i < l; i++) A[i] *= B[i];------// e7
                                            ----int s = static_cast<int>(sqrt(static_cast<double>(n)));------// 64
----fft(A, l, true);-----// d3
                                            ----for (int i = 5; i <= s; i += 6)-----// 6c
----ull *data = new ull[1];-----// e7
                                            ------if (n % i == 0 || n % (i + 2) == 0) return false;------// e9
----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
                                            ----return true; }-----// 43
----for (int i = 0; i < l - 1; i++)-----// 90
-----if (data[i] >= (unsigned int)(radix)) {------// 44
                                            5.5. Miller-Rabin Primality Test. The Miller-Rabin probabilistic primality test.
-----data[i+1] += data[i] / radix:-----// e4
                                            #include "mod_pow.cpp"-----// c7
-----data[i] %= radix;-----// bd
                                            bool is_probable_prime(ll n, int k) {------// be
----if (~n & 1) return n == 2;------// d1
----int stop = l-1;------// cb
                                            ----if (n <= 3) return n == 3;-----// 39
----while (stop > 0 && data[stop] == 0) stop--;------// 97
                                            ----int s = 0; ll d = n - 1;------// 37
----stringstream ss;-----// 42
                                            ----while (~d & 1) d >>= 1, s++;------// 35
----ss << data[stop];-----// 96
                                            ----while (k--) {------// c8
----for (int i = stop - 1; i >= 0; i--)-----// bd
                                            ------ll a = (n - 3) * rand() / RAND_MAX + 2;------// 06
-----ss << setfil('0') << setw(len) << data[i]:-----// b6
                                            -----ll x = mod_pow(a, d, n);------// 64
----delete[] A; delete[] B;-----// f7
----delete[] a; delete[] b;-----// 7e
                                            -----if (x == 1 || x == n - 1) continue;-----// 9b
                                            ------bool ok = false;-----// 03
----delete[] data;-----// 6a
----return intx(ss.str());-----// 38
                                            -----for (int i = 0; i < s - 1; i++) {------// 6b
                                            -----x = (x * x) % n;
}-----// d9
                                            -----if (x == 1) return false;-----// 4f
------}------// a9
k items out of a total of n items.
                                            -----if (!ok) return false;-----// 00
int nck(int n, int k) {-----// f6
----if (n - k < k) k = n - k;-----// 18
                                            ----} return true; }------// bc
```

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5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
                                                      ----return res:------// 03
                                                     }-----// 1c
vi prime_sieve(int n) {-----// 40
----int mx = (n - 3) >> 1, sq, v, i = -1;------// 27
                                                     5.11. Numeric Integration. Numeric integration using Simpson's rule.
----vi primes:-----// 8f
                                                     double integrate(double (*f)(double), double a, double b,-----// 76
----bool* prime = new bool[mx + 1];-----// ef
                                                      ------double delta = 1e-6) {------// c0
----memset(prime, 1, mx + 1);------// 28
                                                      ----if (abs(a - b) < delta)------// 38
----if (n >= 2) primes.push_back(2);-----// f4
                                                      -----return (b-a)/8 *-----// 56
----while (++i <= mx) if (prime[i]) {-----// 73
                                                      -----(f(a) + 3*f((2*a+b)/3) + 3*f((a+2*b)/3) + f(b));-----// e1
-----primes.push_back(v = (i << 1) + 3);-----// be
                                                      ----return integrate(f, a,-----// 64
-----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
                                                      -----(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);------// \theta c
------for (int j = sq; j <= mx; j += v) prime[j] = false; }-----// 2e
                                                        ·----// 4b
----while (++i \le mx) if (prime[i]) primes.push_back((i << 1) + 3);-----// 29
----delete[] prime; // can be used for O(1) lookup-----// 36
                                                     5.12. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete
----return primes; }-----// 72
                                                     Fourier transform. Note that this implementation only handles powers of two, make sure to pad with
5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
                                                      #include <complex>-----// 8e
#include "egcd.cpp"-----// 55
                                                      typedef complex<long double> cpx;-----// 25
-----// e8
                                                      void fft(cpx *x, int n, bool inv=false) {------// 23
int mod_inv(int a, int m) {------// 49
                                                      ----for (int i = 0, j = 0; i < n; i++) {-------// f2
----int x, y, d = eqcd(a, m, x, y);-----// 3e
                                                      ------if (i < j) swap(x[i], x[j]);------// 5c
----if (d != 1) return -1;------// 20
                                                      ------int m = n>>1;------// e5
----return x < 0 ? x + m : x;------// 3c
                                                      ------while (1 <= m && m <= j) j -= m, m >>= 1;-----// fe
                                                      -----j += m:-----// ab
                                                      ----}-----// 1e
5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
                                                      ----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
template <class T>-----// 82
                                                      T mod_pow(T b, T e, T m) {-----// aa
                                                      ------for (int m = 0; m < mx; m++, w *= wp) {------// 40
----T res = T(1);-----// 85
                                                      ------for (int i = m; i < n; i += mx << 1) {------// 33
----while (e) {------// b7
                                                      -----cpx t = x[i + mx] * w;-----// f5
-----if (e & T(1)) res = mod(res * b, m);------// 41
                                                      -----x[i + mx] = x[i] - t;-----// ac
-----b = mod(b * b, m), e >>= T(1); }------// b3
                                                      -----x[i] += t:-----// c7
----return res:-----// eb
                                                      }-----// c5
                                                      ----}-----------// 70
5.9. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
                                                      ----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);------// 3e
#include "eacd.cpp"-----// 55
                                                     }-----// 7d
int crt(const vi& as, const vi& ns) {------// c3
----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
                                                     5.13. Formulas.
----for (int i = 0: i < cnt: i++) N *= ns[i]: ------// 88
                                                         • Number of ways to choose k objects from a total of n objects where order matters and each
----for (int i = 0; i < cnt; i++)-----// f9
                                                          item can only be chosen once: P_k^n = \frac{n!}{(n-k)!}
------egcd(ns[i], l = N/ns[i], r, s), x += as[i] * s * l;-------// b\theta
                                                         • Number of ways to choose k objects from a total of n objects where order matters and each
----return mod(x, N); }-----// 9e
                                                          item can be chosen multiple times: n^k
                                                         • Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type
5.10. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
                                                          2, ..., n_k objects of type k: \binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}
                                                         • Number of ways to choose k objects from a total of n objects where order does not matter
#include "egcd.cpp"-----// 55
                                                          and each item can only be chosen once:
vi linear_congruence(int a, int b, int n) {-----// c8
                                                          \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^k \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0
----int x, y, d = egcd(a, n, x, y);-----// 7a
----vi res;------// f5
                                                         • Number of ways to choose k objects from a total of n objects where order does not matter
                                                          and each item can be chosen multiple times: f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}
----if (b % d != 0) return res:------// 30
----int x\theta = mod(b / d * x, n); ------// 48
                                                         • Number of integer solutions to x_1 + x_2 + \cdots + x_n = k where x_i > 0: f_k^n
```

• Number of subsets of a set with n elements:  $2^n$ 

----for (int k = 0; k < d; k++) res.push\_back(mod(x0 + k \* n / d, n));-----// 21

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- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an  $n \times m$  grid by walking only up and to the right:  $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced:  $C_n = \sum_{k=0}^{n-1} C_k \bar{C}_{n-1-k} = \frac{1}{n+1} {2n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an  $n \times n$  lattice which do not rise above the main diagonal:  $C_n$
- Number of permutations of n objects with exactly k ascending sequences or runs:  $\left\langle \begin{smallmatrix} n \\ k \end{smallmatrix} \right\rangle = \left\langle \begin{smallmatrix} n \\ n-k-1 \end{smallmatrix} \right\rangle = k \left\langle \begin{smallmatrix} n-1 \\ k \end{smallmatrix} \right\rangle + (n-k+1) \left\langle \begin{smallmatrix} n-1 \\ k-1 \end{smallmatrix} \right\rangle = \sum_{i=0}^k (-1)^i \binom{n+1}{i} (k+1-i)^n, \left\langle \begin{smallmatrix} n \\ 0 \end{smallmatrix} \right\rangle = \left\langle \begin{smallmatrix} n \\ n-1 \end{smallmatrix} \right$
- Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1)\binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements):  $D_0 = 1, D_1 =$  $0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points:  $\binom{n}{k}D_{n-k}$
- Jacobi symbol:  $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- Heron's formula: A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} - 1$ .
- Divisor sigma: The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where  $n = \prod_{i=0}^{r} p_i^{a_i}$  is the prime factorization.
- Divisor count: A special case of the above is  $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$ where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- $gcd(2^a 1, 2^b 1) = 2^{gcd(a,b)} 1$

5.14. Numbers and Sequences. Some random prime numbers: 1031, 32771, 1048583, 33554467, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.

## 6. Geometry

## 6.1. **Primitives.** Geometry primitives.

```
#include <complex>-----// 8e
#define P(p) const point &p-----// b8
#define L(p0, p1) P(p0), P(p1)-----// 30
#define C(p0, r) P(p0), double r-----// 08
#define PP(pp) pair<point,point> &pp-----// a1
typedef complex<double> point;-----// 9e
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// 4a -----x = min(x, abs(c - closest_point(a,b, c, true)));-----// 48
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// f3 -----x = min(x, abs(d - closest_point(a,b, d, true)));-----// 75
point rotate(P(p), double radians = pi / 2, P(about) = point(\theta, \theta)) { ------// \theta b
----return (p - about) * exp(point(0, radians)) + about; }-----// f5
point reflect(P(p), L(about1, about2)) {------// 45
```

```
----point z = p - about1, w = about2 - about1;-----// 74
----return conj(z / w) * w + about1; }-----// d1
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// 98
point normalize(P(p), double k = 1.0) { ------// a9
----return abs(p) == 0 ? point(0,0) : p / abs(p) * k; } //TODO: TEST-----// 1c
bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }----// 74
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// ab
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS: }-----// 95
bool collinear(L(a, b), L(p, q)) {-----// de
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// 27
double angle(P(a), P(b), P(c)) {------// 93
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------// a2
double signed_angle(P(a), P(b), P(c)) {------// 46
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 80
double angle(P(p)) { return atan2(imaq(p), real(p)); }-----// c\theta
point perp(P(p)) { return point(-imag(p), real(p)); }-----// 3c
double progress(P(p), L(a, b)) {-----// c7
----if (abs(real(a) - real(b)) < EPS)------// 7d
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));------// b7
----else return (real(p) - real(a)) / (real(b) - real(a)); }------// 6c
----// NOTE: check for parallel/collinear lines before calling this function---// 88
----point r = b - a, s = q - p;------// 54
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// 29
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// 30
-----return false:-----// c0
----res = a + t * r:-----// 88
----return true:-----// 03
point closest_point(L(a, b), P(c), bool segment = false) {------// 06
----if (seament) {-------// 90
-----if (dot(b - a, c - b) > 0) return b;------// 93
-----if (dot(a - b, c - a) > 0) return a;-----// bb
----}------// d5
----double t = dot(c - a, b - a) / norm(b - a);------// 61
----return a + t * (b - a);-----// 4f
}-----// 19
double line_segment_distance(L(a,b), L(c,d)) {-----// f6
----double x = INFINITY:-----// 8c
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// 5f
----else if (abs(a - b) < EPS) \times = abs(a - closest_point(c, d, a, true)); -----// 97
----else if (abs(c - d) < EPS) x = abs(c - closest\_point(a, b, c, true)):-----// 68
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// fa
-----(ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0:-----// bb
----else {------// 5b
-----x = min(x, abs(a - closest_point(c,d, a, true)));
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// 75
---}-----// 60
----return x:-----// 57
```

```
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}------// 8e -----in = !in:-----// b2
int intersect(C(A, rA), C(B, rB), point & res1, point & res2) { ---------// ca ----return in ? -1 : 1; }------------------------------// 77
int intersect(L(A, B), C(0, r), point & res1, point & res2) {-------// ab //----- if (ccw(a, b, p) >= 0) right.push_back(p);-----// e3
---- double h = abs(0 - closest_point(A, B, 0));------// a6 //-----// myintersect = intersect where-----// 24
---- if(r < h - EPS) return 0;------// 52 //------// f2
}-----// 09 //---- return pair<polygon, polygon>(left, right);------// 1d
int tangent(P(A), C(0, r), point & res1, point & res2) {------// f0 // }------// 37
----point v = 0 - A; double d = abs(v):-----// 07
                         6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
----if (d < r - EPS) return 0;-------// b3
----double alpha = asin(r / d), L = sqrt(d*d - r*r);------// 64 #include "polygon.cpp"-----// 58
----v = normalize(v, L);------// 37 #define MAXN 1000-----// 09
----res1 = A + rotate(v, alpha); res2 = A + rotate(v, -alpha);-----// 58
                        point hull[MAXN];-----// 43
----double theta = asin((rB - rA)/abs(A - B));--------------------------------// 3d
----u = normalize(u, rA);------// 58 ------if (i > 0 && p[i] == p[i - 1]) continue;------// b2
----P.first = A + normalize(v, rA); P.second = B + normalize(v, rB);------// 94 -------while (l >= 2 && ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;-----// 20
----Q.first = A + normalize(u, rA); Q.second = B + normalize(u, rB);------// 8e ------hull[l++] = p[i];-------// f7
}-----// e6 ---}-----// d8
                         ----int r = 1:-----// 59
6.2. Polygon. Polygon primitives.
                         ----for (int i = n - 2; i >= 0; i--) {------// 16
#include "primitives.cpp"------// e0 -----if (p[i] == p[i + 1]) continue;------// c7
----double area = 0; int cnt = size(p);------// a2 ---}
----for (int i = 1; i + 1 < cnt; i++)------// d2 ----return l == 1 ? 1 : r - 1;------// 6d
-----area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 7e }------// 79
----return area / 2; }------// e1
double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25 6.4. Line Segment Intersection. Computes the intersection between two line segments.
----for (int i = 0, j = n - 1; i < n; j = i++)------// 77 ------A = B = a; return abs(a - d) < EPS; }------// ee
-----return 0;-----// cc -----return 0.0 <= p && p <= 1.0-----// 8a
------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))-------// 1f ----else if (abs(c - d) < EPS) {-------// 26
```

```
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                                                                                  21
------it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn)):------// fc
----else if (collinear(a,b, c,d)) {-------// bc -------while (it != jt) mn = min(mn, abs(*it - pts[i])), it++;-----// 09
------double ap = progress(a, c,d), bp = progress(b, c,d);-------// a7 ------cur.insert(pts[i]); }-------// 82
------if (ap > bp) swap(ap, bp);--------// b1 ----return mn; }-------// 4c
-----if (bp < 0.0 || ap > 1.0) return false;------// 0c
                                          6.8. 3D Primitives. Three-dimensional geometry primitives.
------A = c + max(ap, 0.0) * (d - c);------// f6
                                          #include <cmath>-----// e5
-----B = c + \min(bp, 1.0) * (d - c); -----// 5c
                                          #define P(p) const point3d &p-----// e5
-----return true; }-----// ab
                                          #define L(p0, p1) P(p0), P(p1)-----// 3c
----else if (parallel(a,b, c,d)) return false;-----// ca
                                          #define PL(p0, p1, p2) P(p0), P(p1), P(p2)-----// 2d
----else if (intersect(a,b, c,d, A, true)) {------// 10
                                          struct point3d {-----// a1
-----B = A; return true; }-----// bf
                                          ----double x, y, z;------// 29
----return false:-----// b7
                                          ----point3d() : x(0), y(0), z(0) {}-----// 8a
}-----// 8b
                                          ----point3d(double _x, double _y, double _z) : x(_x), y(_y), z(_z) {}------// 1c
   ·----// e6
                                          ----point3d operator+(P(p)) const {------// dc
                                          -----return point3d(x + p.x, y + p.y, z + p.z); }-----// d4
6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
coordinates) on a sphere of radius r.
                                          ----point3d operator-(P(p)) const {------// a7
                                         -----return point3d(x - p.x, y - p.y, z - p.z); }------// cc
double gc_distance(double pLat, double pLong,-----// 7b
                                          ----point3d operator-() const {------// 2e
-----// a4
                                          -----return point3d(-x, -y, -z); }------// 77
----pLat *= pi / 180; pLong *= pi / 180;-----// ee
                                         ----qLat *= pi / 180; qLong *= pi / 180;-----// 75
                                         -----return point3d(x * k, y * k, z * k); }-----// 1f
----return r * acos(cos(pLat) * cos(qLat) * cos(pLong - qLong) +-----// e3
-----sin(pLat) * sin(qLat));-----// dc
                                          -----return point3d(x / k, y / k, z / k); }-----// f0
-----// 60
}------// 3f ----double operator%(P(p)) const {-------// 30
                                          -----return x * p.x + y * p.y + z * p.z; }-----// e6
                                         ----point3d operator*(P(p)) const {------// 96
6.6. Triangle Circumcenter. Returns the unique point that is the same distance from all three
points. It is also the center of the unique circle that goes through all three points.
                                          -----return point3d(y*p.z - z*p.v. z*p.v. z*p.x - x*p.z. x*p.v - y*p.x); }------// 02
point circumcenter(point a, point b, point c) {------// 76 -----return sqrt(*this % *this); }------// c9
----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a ------return (*this - p).length(); }-------// 5a
}------// c3 ----double distTo(P(A), P(B)) const {-----------------------------// d8
                                          -----// A and B must be two different points-----// 93
6.7. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
                                          -----return ((*this - A) * (*this - B)).length() / A.distTo(B); }------// 38
pair of points.
                                          ----point3d normalize(double k = 1) const {------// f0
#include "primitives.cpp"------// e0 -----// length() must not return 0--------// b8
-----/<sub>85</sub> ------return (*this) * (k / length()); }-------//<sub>46</sub>
-----real(a) < real(b) : imag(a) < imag(b); } };-------// 53 ------return A + v.normalize((v % (*this - A)) / v.length()); }------// 0c
struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f ----point3d rotate(P(normal)) const {-------// 15
----return abs(imag(a) - imag(b)) > EPS ?-----// 0b -----// normal must have length 1 and be orthogonal to the vector-----// 0b
----sort(pts.beqin(), pts.end(), cmpx());-------// 0c ------return (*this) * cos(alpha) + rotate(normal) * sin(alpha); }------// a8
----set<point, cmpy> cur;-------// bd ----point3d rotatePoint(P(0), P(axe), double alpha) const{-------// f0
----set<point, cmpv>::const_iterator it, jt;------// a6 ------point3d Z = axe.normalize(axe % (*this - 0));-----// 89
----double mn = INFINITY:-----// f9 ------return 0 + Z + (*this - 0 - Z).rotate(alpha, 0); }------// 43
```

```
-----return abs(x) < EPS && abs(y) < EPS && abs(z) < EPS; }-----// 64
                                                6.10. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
----bool isOnLine(L(A, B)) const {------// bc
                                                   • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
-----return ((A - *this) * (B - *this)).isZero(); }------// 8c
                                                   • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
----bool isInSegment(L(A, B)) const {------// e0
                                                   • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
-----return isOnLine(A, B) && ((A - *this) % (B - *this)) < EPS; }------// 52
                                                    of that is the area of the triangle formed by a and b.
----bool isInSegmentStrictly(L(A, B)) const {------// 73
-----return isOnLine(A, B) && ((A - *this) % (B - *this)) < -EPS; }------// 1c
                                                                  7. Other Algorithms
----double getAngle() const {-------// 20
                                                7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
-----return atan2(y, x); }-----// 2a
                                                function f on the interval [a, b], with a maximum error of \varepsilon.
----double getAngle(P(u)) const {------// 19
                                                double binary_search_continuous(double low, double high,-----// 8e
-----return atan2((*this * u).length(), *this % u); }-----// 2f
                                                -----double eps, double (*f)(double)) {------// c0
----bool isOnPlane(PL(A, B, C)) const {------// c8
                                                ----while (true) {------// 3a
-----return abs((A - *this) * (B - *this) % (C - *this)) < EPS; } };------// 16
                                                ------double mid = (low + high) / 2, cur = f(mid);-----// 75
int intersect(L(A, B), L(C, D), point3d &0){-----// 81
                                                -----if (abs(cur) < eps) return mid;------// 76
----if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;-----// 3b
                                                ------else if (0 < cur) high = mid;------// e5
----if (((A - B) * (C - D)).length() < EPS)------// 6c
                                                -----else low = mid:-----// a7
-----return A.isOnLine(C, D) ? 2 : 0;-----// 3d
                                                ----}--------// b5
----point3d normal = ((A - B) * (C - B)).normalize();-----// 9b
----double s1 = (C - A) * (D - A) % normal;-----// 1e
----0 = A + ((B - A) / (s1 + ((D - B) * (C - B) % normal))) * s1;-----// 6e
                                                 Another implementation that takes a binary predicate f, and finds an integer value x on the integer
----return 1: }-----// 24
                                                interval [a,b] such that f(x) \wedge \neg f(x-1).
int intersect(L(A, B), PL(C, D, E), point3d & 0) {-----// ce
                                                int binary_search_discrete(int low, int high, bool (*f)(int)) {------// 51
----double V1 = (C - A) * (D - A) % (E - A):-----// 3c
                                                ----assert(low <= high);-----// 19
----double V2 = (D - B) * (C - B) % (E - B);-----// c8
                                                ----while (low < high) {-----// a3
----0 = A + ((B - A) / (V1 + V2)) * V1;------// 94 -----else low = mid + 1;-----// 03
----return 1; }------// 7a
                                               ----}-----// 9b
bool intersect(P(A), P(nA), P(B), P(nB), point3d &P, point3d &Q) \{-----//24\}
                                               ----assert(f(low));------// 42
---point3d n = nA * nB;-----// d3
                                                ----return low:------// a6
----if (n.isZero()) return false;------// b2
                                                }-----// d3
----point3d v = n * nA;------// c7
----P = A + (n * nA) * ((B - A) % nB / (v % nB));
                                                7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotonically
---0 = P + n;
                                                cally decreasing, ternary search finds the x such that f(x) is maximized.
                                                template <class F>-----// d1
----return true: }-----// 1f
                                                double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
6.9. Polygon Centroid.
                                                ----while (hi - lo > eps) {------// 3e
#include "polygon.cpp"-----// 58
                                                ------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
point polygon_centroid(polygon p) {------// 79
                                                -----if (f(m1) < f(m2)) lo = m1:-----// 1d
----double cx = 0.0, cy = 0.0; -----// d5
                                               -----else hi = m2;-----// b3
----double mnx = 0.0, mny = 0.0;-----// 22
                                               ----}-----// bb
----int n = size(p);------// 2d
                                                ----return hi:-----// fa
----for (int i = 0; i < n; i++)------// 24
                                                }-----// 66
-----mnx = min(mnx, real(p[i])),-----// 6d
                                               7.3. 2SAT. A fast 2SAT solver.
-----mny = min(mny, imag(p[i]));-----// 95
----for (int i = 0; i < n; i++)-----------// df #include "../graph/scc.cpp"------// c3
-----p[i] = point(real(p[i]) - mnx, imag(p[i]) - mny);------// c2 -----// 63
-----cx += (real(p[i]) + real(p[j])) * cross(p[i], p[j]);------// d5 ----vvi adj(2*n+1);------// 7b
----return point(cx, cy) / 6.0 / polygon_area_signed(p) + point(mnx, mny); }---// 2f ------adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// 17
```

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------if (clauses[i].first != clauses[i].second)-------// 87 ----node *head;---------------------------------// fe
-----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93 ----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
----pair<union_find, vi> res = scc(adj);------// 9f -----sol = new int[rows];------// 5f
----vi dag = res.second;------// 58 ------arr[i] = new bool[cols], memset(arr[i], 0, cols);------// 75
-----truth[cur + n] = truth[p]:------// b3 -------for (int j = 0; j < cols; j++)------// f5
-----truth[o] = 1 - truth[p];------// 80 -------if (i == rows || arr[i][j]) ptr[i][j] = new node(i, j);-----// 89
------if (truth[p] == 1) all_truthy.push_back(cur);--------// 5c ------else ptr[i][j] = NULL;-----------------// 32
----}------// d9 -----}------// 98
-----if (!ptr[i][j]) continue;-----// 35
7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
                             -----int ni = i + 1, nj = j + 1;-----// b7
----queue<int> q;-------// f6 -------if (ni == rows + 1) ni = 0;-------// 81
----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));-------// c3 -------if (ni == rows || arr[ni][i]) break;------// 19
----for (int i = 0; i < n; i++) q.push(i);------// fe ------ptr[i][j]->d = ptr[ni][j];-----// 71
----while (!q.empty()) {------// 55 -----ptr[ni][j]:>u = ptr[i][j];-----// c4
------int curm = q.front(); q.pop();------// ab -------while (true) {-------// c6
------for (int &i = at[curm]; i < n; i++) {---------// 9a -------if (nj == cols) nj = 0;-------// e2
------int curw = m[curm][i];-------// cf -------if (i == rows || arr[i][nj]) break;------// 8d
------if (eng[curw] == -1) { }-------// 35
-----q.push(eng[curw]);------// 8c -----ptr[i][j]->r = ptr[i][nj];-----// d5
-----else continue;-----// b4 ------ptr[i][nj]->l = ptr[i][j];------// 72
----}------head = new node(rows, -1);-------// 80
----return res;------head->r = ptr[rows][0];-------// 73
}------ptr[rows][0]->l = head;------// 3b
                             ------head->l = ptr[rows][cols - 1];-----// da
7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the
                             ------ptr[rows][cols - 1]->r = head;------// 6b
Exact Cover problem.
                              ------for (int i = 0: i < cols: i++) {------// 97
bool handle_solution(vi rows) { return false; }------// 63
                             -----int cnt = -1:-----// 84
struct exact_cover {------// 95
                             -----for (int i = 0: i <= rows: i++)------// 96
----struct node {------// 7e ------if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];------// cb
------<mark>int</mark> row, col, size;------// ae _______/ 59
------node(int _row, int _col) : row(_row), col(_col) {-------// c9 ------for (int i = 0; i <= rows; i++) delete[] ptr[i];------// bf
-----size = 0; l = r = u = d = p = NULL; }-----// c3 -----delete[] ptr;-----// 99
----};-------// c1 ...}
----int rows, cols, *sol;------// 7b
                             ----<mark>#</mark>define COVER(c, i, j) \\ \------// 6a
----bool **arr:------// e6
```

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}-----// 42
------for (node *i = c->d; i != c; i = i->d) \sqrt{\phantom{a}}
------for (node *j = i->r; j != i; j = j->r) \sqrt{\phantom{a}}
                                                      7.8. Dates. Functions to simplify date calculations.
-----j->d->u = j->u, j->u->d = j->d, j->p->size--;-----// 16
                                                      int intToDay(int jd) { return jd % 7; }-----// 89
----#define UNCOVER(c, i, j) N-----// d0
                                                      int dateToInt(int y, int m, int d) {-----// 96
------for (node *i = c->u; i != c; i = i->u) \sqrt{\phantom{a}}
                                                      ----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 -----// d1
-----j->p->size++, j->d->u = j->u->d = j; \\ \]
                                                      -----3 * ((v + 4900 + (m - 14) / 12) / 100) / 4 +------// be
------c->r->l = c->l->r = c:------// 91
                                                      -----d - 32075:-----// e0
----bool search(int k = 0) {------// bb
                                                      }-----// fa
-----if (head == head->r) {------// c3
                                                      void intToDate(int jd, int &y, int &m, int &d) {------// a1
-----vi res(k);-----// 9f
                                                      ----int x, n, i, j;------// 00
------for (int i = 0; i < k; i++) res[i] = sol[i];------// 75
                                                      ---x = id + 68569; 11
-----sort(res.begin(), res.end());-----// 87
                                                      ---n = 4 * x / 146097; 2f
-----return handle_solution(res);-----// 51
                                                      ---x = (146097 * n + 3) / 4:
------}-----// f5
                                                      ----i = (4000 * (x + 1)) / 1461001;-----// 0d
-----node *c = head->r, *tmp = head->r;------// 8e
                                                      ----x -= 1461 * i / 4 - 31:-----// 09
-----for ( ; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 00
                                                      ---i = 80 * x / 2447;
-----if (c == c->d) return false;-----// b0
                                                      ---d = x - 2447 * j / 80;
-----COVER(c, i, j);-----// 7a
                                                      ---x = i / 11:-----// b7
------bool found = false:-----// 7f
                                                      ---m = i + 2 - 12 * x;
------for (node *r = c->d; !found && r != c; r = r->d) {-------// 88
                                                      ---v = 100 * (n - 49) + i + x
-----sol[k] = r->row;-----// ef
                                                      }-----// af
------for (node *j = r -> r; j = j -> r) { COVER(j -> p, a, b); }-----// 61
-----found = search(k + 1);-----// f1
                                                                          8. Useful Information
------for (node *j = r->l; j != r; j = j->l) { UNCOVER(j->p, a, b); }----// ab
-------}------// 1a
                                                        Tips & Tricks.
------UNCOVER(c, i, i):-----// 3a
                                                         • How fast does our algorithm have to be? Can we use brute-force?
-----return found:-----// 80
                                                         • Does order matter?
• Is it better to look at the problem in another way? Maybe backwards?
                                                         • Are there subproblems that are recomputed? Can we cache them?
                                                         • Do we need to remember everything we compute, or just the last few iterations of computation?
7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \dots, k-1\}
                                                         • Does it help to sort the data?
1}.
                                                         • Can we speed up lookup by using a map (tree or hash) or an array?
vector<int> nth_permutation(int cnt, int n) {------// 78
                                                         • Can we binary search the answer?
----vector<int> idx(cnt), per(cnt), fac(cnt);-----// 9e
                                                         • Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph
----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
                                                          into some other kind of a graph (perhaps a DAG, or a flow network)?
----for (int i = 1; i <= cnt; i++) fac[i - 1] = n % i, n /= i;------// 04
                                                         • Make sure integers are not overflowing.
----for (int i = cnt - 1; i >= 0; i--)-----// 52
                                                         • Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// 41
                                                          m_1, m_2, \ldots, m_k, where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer
----return per:-----// 84
                                                          using CRT?
}-----// 97
                                                         • Are there any edge cases? When n = 0, n = -1, n = 1, n = 2^{31} - 1 or n = -2^{31}? When
                                                          the list is empty, or contains a single element? When the graph is empty, or contains a single
7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
                                                          vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
ii find_cycle(int x0, int (*f)(int)) {------// a5
                                                         • Can we use exponentiation by squaring?
----int t = f(x0), h = f(t), mu = 0, lam = 1;------// 8d
----while (t != h) t = f(t), h = f(f(h));-----// 79
                                                      8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input
----h = x0:-----// 04
                                                      reading/output writing. This can be achieved by reading all input in at once (using fread), and then
----while (t != h) t = f(t), h = f(h), mu++;-----// 9d
                                                      parsing it manually. Output can also be stored in an output buffer and then dumped once in the end
```

(using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading

----h = f(t):-----// 00

----while (t != h) h = f(h), lam++;-----// 5e

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## 8.3. Worst Time Complexity.

1 0				
	n	Worst AC Algorithm	Comment	
	≤ 10	$O(n!), O(n^6)$	e.g. Enumerating a permutation	
	$\leq 15$	$O(2^n \times n^2)$	e.g. DP TSP	
	$\leq 20$	$O(2^{n}), O(n^{5})$	e.g. $DP + bitmask technique$	
	$\leq 50$	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$	
	$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's	
	$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort	
	$\leq 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree	
	$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)	

## 8.4. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.