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------if (l > r) return ID;-------// cc ----return query(segs[id].lid, l, r) + query(segs[id].rid, l, r); }------// ee
-----if (r < a || b < l) return data[i];-----// d9
                                      2.3. Fenwick Tree.
------if (a <= l \& \& r <= b) return (lazy[i] = v) * (r - l + 1) + data[i];----// 06
                                      struct fenwick_tree {-----// 98
-----int m = (l + r) / 2;-----// cc
                                      ----int n; vi data;------// d3
-----return data[i] = f(ru(a, b, v, l, m, 2*i+1),-----// cc
                                      ----fenwick_tree(int _n) : n(_n), data(vi(n)) { }------// db
-----/u(a, b, v, m+1, r, 2*i+2));------// 2b
                                      ----void update(int at, int by) {------// 76
----}------// @b
----void propagate(int l, int r, int i) {-----// a7
                                      ------while (at < n) data[at] += by, at |= at + 1; }------// fb
                                      ----int query(int at) {-------// 71
-----if (l > r || lazy[i] == INF) return;-----// 5f
                                      ------int res = 0;-----// c3
-----data[i] += lazy[i] * (r - l + 1);-----// 44
                                      ------while (at >= 0) res += data[at], at = (at & (at + 1)) - 1;------// 37
-----if (l < r) {------// 28
                                      -----return res; }-----// e4
------if (lazy[2*i+1] == INF) lazy[2*i+1] = lazy[i];------// 4e
                                      ----int rsq(int a, int b) { return query(b) - query(a - 1); }-----// be
------else lazy[2*i+1] += lazy[i];-----// 1e
                                       -----// 57
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];-----// de
                                      struct fenwick_tree_sq {-----// d4
------else lazy[2*i+2] += lazy[i];-----// 74
                                      ----int n; fenwick_tree x1, x0;------// 18
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e
-----lazy[i] = INF;-----// f8
                                      ------x0(fenwick_tree(n)) { }------// 7c
----// insert f(y) = my + c if x <= y------// 17
}:----// ae
                                      ----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }----// 45
                                      ----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73
                                      }:-----// 13
2.2.1. Persistent Segment Tree.
                                      void range_update(fenwick_tree_sq &s, int a, int b, int k) {------// 89
int segcnt = 0;-----// cf
                                      ----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f
struct segment {------// 68
                                      int range_query(fenwick_tree_sq &s, int a, int b) {------// 15
----int l, r, lid, rid, sum;------// fc
                                      ----return s.query(b) - s.query(a-1); }------// f3
} seqs[2000000];-----// dd
int build(int l, int r) {------// 2b
                                      2.4. Matrix.
----int id = seqcnt++;-------------------------// a8 template <> bool eq<double o, double b) { return abs(a - b) < EPS; }---// a7
----else {------// fe ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {------// 56
------data.assign(cnt, T(0)); }------// int m = (l + r) / 2;-------// int m = (l + r) / 2;------// e3
------seqs[id].lid = build(l , m);-------// e3 ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// b5
-----segs[id].rid = build(m + 1, r); }------// 69 ------cnt(other.cnt), data(other.data) { }------// c1
----return id; }--------------------------// c5 ----matrix<T> operator +(const matrix& other) {-------------------------// 33
int update(int idx, int v, int id) {------// b8 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] += other.data[i];-----// f8
----int nid = segcnt++;------matrix<T> res(*this); rep(i,0,cnt) res.data[i] -= other.data[i];-----// 7b
----seqs[nid].l = seqs[id].l;------// 78 -----return res; }-----
----segs[nid].lid = update(idx, v, segs[id].lid);------// 92 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] *= other;------// 05
----seqs[nid].rid = update(idx, v, seqs[id].rid);-------// 06 ------return res; }------
----return nid; }------// e6 ------matrix<T> res(rows, other.cols);------// 4c
----if (r < segs[id].l || segs[id].r < l) return 0;-------// 17 ------res(i, j) += at(i, k) * other.data[k * other.cols + j];------// 17
```

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----matrix<T> pow(int p) {-------------------------// 53 ----r->l = merge(l, r->l); augment(r); return r; }------------------------// cθ
-----if (p) sq = sq * sq;-------// 35 ----return NULL; }------// ae
------} return res; }--------// 22 node* insert(node *t, int x, int y) {--------// 78
------matrix<T> mat(*this); det = T(1), rank = max(rows, cols);------// 7a ----pair<node*, node*> res = split(t, x);--------// ca
-----int k = r;------// 5b node* erase(node *t, int x) {-------// 4d
-----if (k >= rows) { rank-; continue; }------// 1a ---if (t->x < x) t->r = erase(t->r, x);-------// 7c
-----if (k != r) {-------// c4 ----else if (x < t->x) t->l = erase(t->l, x);-------// 48
------swap(mat.at(k, i), mat.at(r, i));-------// 7d int kth(node *t, int k) {---------------------------------// b3
-----rep(i,0,cols) mat(r, i) /= d;------// d1 ----else return kth(t->r, k - tsize(t->l) - 1); }------// f0
-----rep(i,0,rows) {------// f6
                        2.6. Misof Tree.
-----T m = mat(i, c);------// 05
                        #define BITS 15-----// 7b
-----if (i != r && !eq<T>(m, T(0)))------// 1a
                        struct misof_tree {-----// fe
-----rep(j,0,cols) mat(i, j) -= m * mat(r, j);-----// 7b
                        ----int cnt[BITS][1<<BITS];------// aa
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }-----// b0
-----} return mat; }------// b3
                        ----void insert(int x) { for (int i = 0; i < BITS; cnt[i++][x]++, x >>= 1); }--//5a
----matrix<T> transpose() {------// 59
                        ----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49
------matrix<T> res(cols, rows);------// 5b
----rep(i,0,rows) rep(j,0,cols) res(j, i) = at(i, j);-----// 92
                        ----int nth(int n) {-------// 8a
-----return res; } };------// df
                        -----int res = 0;------// a4
                        ------for (int i = BITS-1; i >= 0; i--)------// 99
                        ------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4
2.5. Cartesian Tree.
                        -----return res:-----// 3a
struct node {-----// 36
                        ----}-----// b5
----int x, y, sz;------// e5
                        };-----// 0a
---node *l, *r;-----// 4d
                       2.7. Sqrt Decomposition.
----node(int _x, int _y) : x(_x), y(_y), sz(1), l(NULL), r(NULL) { } };------// 19
pair<node*, node*> split(node *t, int x) {------// 1d ----segment(vi _arr) : arr(_arr) { } };------// 11
----if (t->x < x) {-------------------------// dc
------return make_pair(t, res.second); }-------// e0 ---rep(i,0,size(T))--------// b1
```

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------T.push_back(segment(vi(arr.begin()+i, arr.begin()+min(i+K, cnt))));----// f0 ------return (h[i+1].second-h[i].second)/(h[i].first-h[i+1].first); }------// b9
------if (intersect(n-3) < intersect(n-2)) break;-------// 07
----if (i >= size(T)) return size(T);------// 83 ------swap(h[n-2], h[n-1]);------// bf
----T[i] = segment(vi(T[i].arr.begin(), T[i].arr.begin() + at)); ------//af ------int 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10 = 0, 10
}-----// ea ------int mid = lo + (hi - lo) / 2;---------// 5a
----T.insert(T.begin() + split(at), segment(arr));------// 67 ------return h[res+1].first * x + h[res+1].second; } };------// 84
}-----// cc
                                                         3. Graphs
void erase(int at) {-----// be
----int i = split(at): split(at + 1):-----// da
                                        3.1. Single-Source Shortest Paths.
----T.erase(T.begin() + i);-----// 6b
                                        3.1.1. Dijkstra's algorithm.
}-----// 4b
                                        int *dist. *dad:-----// 46
2.8. Monotonic Queue.
                                        struct cmp {-----// a5
struct min_stack {-------// d8 ----bool operator()(int a, int b) {------------// bb
----stack<int> S. M:------// fe ------return dist[a] != dist[b] ? dist[b] : a < b; }------// e6
------S.push(x);-------// e2 pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
------M.push(M.empty() ? x : min(M.top(), x)); }-------// 92 ----dist = new int[n];-------// 84
struct min_queue {------// b4 -----// b4 -----// 58
----min_stack inp, outp;------// 3d -----rep(i,0,size(adj[cur])) {-----------// a6
----void fix() {-----------------------// 5d --------ndist = dist[cur] + adj[cur][i].second;------// 3a
------dist[nxt] = ndist, dad[nxt] = cur, pg.insert(nxt);------// eb
------if (inp.empty()) return outp.mn();-------// 01 }------// 9b
-----if (outp.empty()) return inp.mn();------// 90
                                        3.1.2. Bellman-Ford algorithm.
-----return min(inp.mn(), outp.mn()); }-----// 97
                                        int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {------// cf
----void pop() { fix(); outp.pop(); }-----// 4f
----bool empty() { return inp.empty() && outp.empty(); }-----// 65
                                        ----has_negative_cycle = false;-----// 47
                                        ----int* dist = new int[n];-----// 7f
}:-----// 60
                                        ----rep(i,0,n) dist[i] = i == s ? 0 : INF;------// df
                                        ----rep(i,0,n-1) rep(j,0,n) if (dist[j] != INF)------// 4d
2.9. Convex Hull Trick.
```

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------dist[adj[j][k].first] = min(dist[adj[j][k].first],------// e1 -------// a5
-----/dist[j] + adj[j][k].second);-------// 18 void scc_dfs(const vvi &adj, int u) {--------------------// a1
------if (dist[j] + adj[j][k].second < dist[adj[j][k].first])--------// 37 ----rep(i,0,size(adj[u]))------------------------------// 2d
------has_negative_cycle = true;-------// f1 ------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-------// a2
}------// a9 }------// 53
                                 -----// 63
3.1.3. IDA^* algorithm.
                                 pair<union_find, vi> scc(const vvi &adj) {------// c2
int n, cur[100], pos;-----// 48 ----int n = size(adj), u, v;------// f8
----int h = 0;-------// 4a ----union_find uf(n);------// a8
---rep(i,0,n) if (cur[i] != 0) h += abs(i - cur[i]);------// 9b ---vi dag;-----
----return h;------// c6 ----vvi rev(n);------// c5
}------// c8 ----rep(i,0,n) rep(j,0,size(adj[i])) rev[adj[i][j]].push_back(i);------// 7e
----if (q + h > d) return g + h;-------// 15 ----fill(visited.begin(), visited.end(), false);------// 59
----if (h == 0) return 0;------// ff ----stack<int> S;------// bb
----int mn = INF;------// 7e ----for (int i = n-1; i >= 0; i--) {-------// 96
----rep(di,-2,3) {-------// 0d ------if (visited[order[i]]) continue;------// db
------if (di == 0) continue; ------// 0a ------S.push(order[i]), dag.push_back(order[i]); ------// 68
------int nxt = pos + di;-------// 76 ------while (!S.empty()) {-------// 9e
------if (nxt == prev) continue; -------// 39 ------visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]); -----// b3
------if (0 <= nxt && nxt < n) {-------// 68 ------rep(j,0,size(adj[u])) if (!visited[v = adj[u][j]]) S.push(v);-----// 1b
-----mn = min(mn, dfs(d, g+1, nxt));------// 22 ----return pair<union_find, vi>(uf, dag);------// 2b
------swap(cur[pos], cur[nxt]);-----// 3b
3.3. Cut Points and Bridges.
-----if (mn == 0) break;-----// 8f
                                 #define MAXN 5000-----// f7
----}------// d3
                                 int low[MAXN], num[MAXN], curnum;-----// d7
----return mn:-----// da
                                 void dfs(const vvi &adj, vi &cp, vii &bri, int u, int p) {------// 22
}-----// f8
                                  ----low[u] = num[u] = curnum++;-----// a3
int idastar() {-----// 22
                                  ----int cnt = 0; bool found = false;-----// 97
----rep(i,0,n) if (cur[i] == 0) pos = i;------// 6b
                                  ----rep(i,0,size(adj[u])) {------// ae
----int d = calch();-----// 38
                                  -----int v = adj[u][i];------// 56
----while (true) {------// 18
                                  -----if (num[v] == -1) {------// 3b
-----int nd = dfs(d, 0, -1):-----// 42
                                 -----dfs(adj, cp, bri, v, u);-----// ba
-----if (nd == 0 || nd == INF) return d;------// b5
                                 -----low[u] = min(low[u], low[v]);-----// be
-----d = nd;-----// f7
                                  -----Cnt++;-----// e0
-----found = found || low[v] >= num[u];-----// 30
}-----// 82
                                  ------if (low[v] > num[u]) bri.push_back(ii(u, v));------// bf
                                  -----} else if (p != v) low[u] = min(low[u], num[v]); }------// 76
3.2. Strongly Connected Components.
                                  ----if (found && (p !=-1 \mid \mid cnt > 1)) cp.push_back(u); }-------// 3e
3.2.1. Kosaraju's algorithm.
                                 pair<vi,vii> cut_points_and_bridges(const vvi &adj) {------// 76
#include "../data-structures/union_find.cpp"---------------------// 5e ----int n = size(adj);-----------------------------------// c8
-----// 11 ----vi cp; vii bri;------------// fb
```

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------e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];------// 95 -------dist[g[j][k]->v]) {-------// 6d
-----rep(i,0,n) if (p[i] != -1) pot[i] += d[i];-------// 86 -------back[g[j][k]->v] = g[j][k];-----// 3d
-----cap = min(cap, cure->w);-----// c3
 A second implementation that is slower but works on negative weights.
                                   -----if (cure->u == s) break;-----// 82
struct flow_network {------// 81
                                   -----cure = back[cure->u];-----// 45
----struct mcmf_edge {------// f6
                                   -----int u, v;-----// e1
                                   -----assert(cap > 0 && cap < INF);-----// ae
-----ll w. c:-----// b4
                                   -----cure = back[t];-----// b9
------mcmf_edge* rev;------// 9d
                                   ------while (true) {------// 2a
-----mcmf_edge(int _u, int _v, ll _w, ll _c, mcmf_edge* _rev = NULL) {-----// ea
                                   -----cost += cap * cure->c;-----// f8
------u = _u: v = _v: w = _w: c = _c: rev = _rev:-----// 83
                                   -----Cure->w -= cap;-----// d1
-----cure->rev->w += cap;-----// cf
----};-------// b9
                                   ------if (cure->u == s) break;------// 8c
----int n;------// b4
                                   -----cure = back[cure->u];------// 60
----vector<pair<int, pair<ll, ll> > * adj;------// 72
----flow_network(int _n) {------// 55
                                   -----n = _n;-----// fa
                                   ------}-------// be
-----adj = new vector<pair<int, pair<ll, ll> > >[n];-----// bb
                                   -----// instead of deleting g, we could also-----// e0
----}------// bd
                                   -----// use it to get info about the actual flow-----// 6c
----void add_edge(int u, int v, ll cost, ll cap) {------// 79
                                   ------for (int i = 0; i < n; i++)------// eb
-----adj[u].push_back(make_pair(v, make_pair(cap, cost)));-----// c8
                                   -----for (int j = 0; j < size(g[i]); j++)------// 82
----}-----// ed
                                   -----delete q[i][j];-----// 06
----pair<ll,ll> min_cost_max_flow(int s, int t) {------// ea
                                   -----delete[] a:-----// 23
-----vector<mcmf_edge*>* q = new vector<mcmf_edge*>[n];-----// ce
                                   -----delete[] back;-----// 5a
------for (int i = 0; i < n; i++) {------// 57
                                   -----delete[] dist;-----// b9
-----for (int j = 0; j < size(adj[i]); j++) {------// 37
                                   -----return make_pair(flow, cost);------// ec
-----mcmf_edge *cur = new mcmf_edge(i, adj[i][j].first,-----// 21
                                   -----adj[i][j].second.first, adj[i][j].second.second),--// 56
                                    -----// bf
-----*rev = new mcmf_edge(adj[i][j].first, i, 0,------// 48
-----/<sub>b1</sub>
                                   3.8. All Pairs Maximum Flow.
-----cur->rev = rev;------// ef
                                   3.8.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is constructed
-----g[i].push_back(cur);-----// 1d
                                   using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the maximum
-----g[adj[i][j].first].push_back(rev);------// 05
                                   flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|). NOTE:
Not sure if it works correctly with disconnected graphs.
------ll flow = 0, cost = 0;------// 68
                                   #include "dinic.cpp"-----// 58
                                   -----// 25
-----mcmf_edge** back = new mcmf_edge*[n];------// e5
------ll* dist = new ll[n];------// 50 bool same[MAXV];------// 59
------for (int i = 0; i < n - 1; i++)-------// be ----rep(s,1,n) {-------// 9e
------for (int j = 0; j < n; j++)-------// 6e -----int l = 0, r = 0;-----------// 08
------if (dist[j] != INF)------// e3 ------par[s].second = g.max_flow(s, par[s].first, false);------// 54
------for (int k = 0; k < size(q[j]); k++)------// 85 -----memset(d, 0, n * sizeof(int));----------// c8
------| (g[j][k]->w > 0 && dist[j] + g[j][k]->c <-----// 7f -----memset(same, 0, n * sizeof(bool));-----------// c9
```

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-d[g[r++] = s] = 1; part(best); part(best); -d[g[r++] = s] = 1;
------while (l < r) {-------// 45 -----rep(i,0.size(adi[u]))-------// 92
-----same[v = q[l++]] = true;------// c5 ------if (adj[u][i] != parent[u] && adj[u][i] != best)-----// e8
-----if (q.e[i].cap > 0 && d[q.e[i].v] == 0)------// 21 ----void build(int r = 0) { curloc = 0, csz(curhead = r), part(r); }------// 78
------while (u ! = -1) uat.push_back(u), u = parent[head[u]]:------// 51
------while (true) {-------// b8 ----int query_upto(int u, int v) { int res = ID;-------// 72
------cap[curl[i] = mn:------// 8d ------while (head[u] != head[v])-------// 69
------if (cur == 0) break;------// fb -----res = f(res, values.query(loc[head[u]], loc[u])),-----// a4
----return make_pair(par, cap);-------// 62 ------return f(query_upto(u, l), query_upto(v, l)); } };------// 5b
}-----// b3
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {-------// 93
                             3.10. Centroid Decomposition.
----if (s == t) return 0;-----// 33
                             #define MAXV 100100-----// 86
----int cur = INF, at = s;-----// e7
                             #define LGMAXV 20-----// aa
----while (gh.second[at][t] == -1)------// 42
                             int imp[MAXV][LGMAXV],....// 6d
-----cur = min(cur, qh.first[at].second), at = qh.first[at].first;-----// 8d
                             ----path[MAXV][LGMAXV],------// 9d
----return min(cur, gh.second[at][t]);-----// 54
                             ----sz[MAXV], seph[MAXV],-----// cf
}------// 46
                             ----shortest[MAXV]:-----// 6b
                             struct centroid_decomposition {------// 99
3.9. Heavy-Light Decomposition.
                             ----int n; vvi adj;-----// e9
#include "../data-structures/segment_tree.cpp"------// 16 ---centroid_decomposition(int _n) : n(_n), adj(n) { }-------// 46
struct HLD {------// 25 ----void add_edge(int a, int b) { adj[a].push_back(b), adj[b].push_back(a); }--// bc
----vvi adj; segment_tree values;---------// 13 ------rep(i,0,size(adj[u])) if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u);--// 78
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), adj(n) {----// 1c ------return sz[u]; }-----
-----vi tmp(n, ID); values = segment_tree(tmp); }-------// f0 ----void makepaths(int sep, int u, int p, int len) {-------// 84
----void add_edge(int u, int v) { adj[u].push_back(v), adj[v].push_back(u); }--// 77 ------jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len;-------------// d9
-----if (parent[v] == u) swap(u, v); assert(parent[u] == v);-------// db -----rep(i,0,size(adj[u])) {--------// f4
-------if (adj[u][i] == p) bad = i;-------// cf
-----sz[u] += csz(adj[parent[adj[u][i]] = u][i]);------// c2 -----if (p == sep) swap(adj[u][bad], adj[u].back()), adj[u].pop_back(); }---// 07
------head[u] = curhead; loc[u] = curloc++;---------// 63 ------down: iter(nxt,adj[sep])---------// 04
-----rep(i,0,size(adj[u]))------// 49 -----sep = *nxt; goto down; }-----// 1a
-----best = adj[u][i];-------// 26 -----rep(i,0,size(adj[sep])) separate(h+1, adj[sep][i]); }------// 90
```

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-----rep(h,0,seph[u]+1)------// c5 ------process(y):------// e8
-----shortest[jmp[u][h]] = min(shortest[jmp[u][h]], path[u][h]); }-----// 11 ------uf.unite(u,v);--------------------------------// 55
----int closest(int u) {-------// 91 ------ancestor[uf.find(u)] = u;------// 1d
-----int mn = INF/2:------// fe -----}-----// fe
-----rep(h,0,seph[u]+1) mn = min(mn, path[u][h] + shortest[jmp[u][h]]);-----// 3e ------colored[u] = true;-------------------------// b9
-----return mn; } };------// 13 -----rep(i,0,size(queries[u])) {-------// d7
                                    -----int v = queries[u][i].first;-----// 89
3.11. Least Common Ancestors, Binary Jumping.
                                    -----if (colored[v]) {------// cb
struct node {-----// 36
                                    -----answers[queries[u][i].second] = ancestor[uf.find(v)];------// 63
---node *p, *imp[20];-----// 24
                                    ----int depth:-----// 10
                                    -----}----// 40
---node(node *_p = NULL) : p(_p) {-----// 78
                                    ----}-------------------------// a9
-----depth = p ? 1 + p->depth : 0;-----// 3b
                                    }:....// le
-----memset(jmp, 0, sizeof(jmp));-----// 64
                                    3.13. Maximum Density Subgraph. Given (weighted) undirected graph G. Binary search density.
-----for (int i = 1; (1<<i) <= depth; i++)------// a8
                                    If g is current density, construct flow network: (S, u, m), (u, T, m + 2g - d_u), (u, v, 1), where m is a
-----jmp[i] = jmp[i-1]->jmp[i-1]; } };-----// 3b
                                    large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has
                                    empty S-component, then maximum density is smaller than q, otherwise it's larger. Distance between
node* st[100000];-----// 65
node* lca(node *a, node *b) {------// 29
                                    valid densities is at least 1/(n(n-1)). Edge case when density is 0. This also works for weighted
                                    graphs by replacing d_u be the weighted degree, and doing more iterations (if weights are not integers).
----if (!a || !b) return NULL;------// cd
----if (a->depth < b->depth) swap(a,b);-----// fe
                                    3.14. Maximum Weighted Independent Set in a Bipartite Graph. This is the same as the
----for (int j = 19; j >= 0; j--)-----// b3
                                    minimum weighted vertex cover. Solve this by constructing a flow network with edges (S, u, w(u))
------while (a->depth - (1<<j) >= b->depth) a = a->jmp[j];-----// cθ
                                    for u \in L, (v, T, w(v)) for v \in R and (u, v, \infty) for (u, v) \in E. The minimum S, T-cut is the answer.
----if (a == b) return a;-----// 08
                                    Vertices adjacent to a cut edge are in the vertex cover.
----for (int j = 19; j >= 0; j--)-----// 11
-------while (a->depth >= (1<<j) && a->jmp[j] != b->imp[j])------// f\theta
                                                    4. Strings
-----a = a->jmp[i], b = b->jmp[i];-----// d\theta
                                    4.1. The Knuth-Morris-Pratt algorithm.
----return a->p; }-----// c5
                                    int* compute_pi(const string &t) {------// a2
3.12. Tarjan's Off-line Lowest Common Ancestors Algorithm.
                                    ----int m = t.size();-----// 8b
#include "../data-structures/union_find.cpp"---------------------// 5e ----int *pit = new int[m + 1];---------
struct tarjan_olca {-------// 87 ---if (0 <= m) pit[0] = 0;-------// 42
----int *ancestor;------// 39 ----if (1 <= m) pit[1] = 0;-------// 34
----vii *queries;------// 66 ------for (int j = pit[i - 1]; ; j = pit[j]) {-------// b5
----union_find uf;-------if (j == 0) { pit[i] = 0; break; }------// 95
-----queries = new vii[n];------// 3e int string_match(const string &s, const string &t) {------// 9e
-----queries[x].push_back(ii(y, size(answers)));-------// a0 -----if (s[i] == t[j]) {--------// 73
-----answers.push_back(-1);-------// ca ------if (j == m) {-------// de
```

```
-----else i++; }-----// b8
                         struct aho_corasick {------// 78
----delete[] pit; return -1; }-------// e3
                         ----struct out_node {------// 3e
                         -----string keyword; out_node *next;-----// f0
4.2. The Z algorithm.
                         -----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
int* z_values(const string &s) {------// 4d
                         ----};------// b9
----int n = size(s):-----// 97
                         ----struct go_node {------// 40
----int* z = new int[n]:-----// c4
                         -----map<char, qo_node*> next;------// 6b
----int l = 0, r = 0:-----// 1c
                         -----out_node *out; go_node *fail;-----// 3e
---z[0] = n;
                         -----qo_node() { out = NULL; fail = NULL; }-----// 0f
---rep(i,1,n) {------// b2
                         ----};-------// c0
----qo_node *qo;------// b8
-----if (i > r) {------// 6d
                         -----l = r = i;-----// 24
                         -----qo = new qo_node();-----// 77
-----iter(k, keywords) {------// f2
-----z[i] = r - l; r--;------// 07
                         -----qo_node *cur = qo;-----// a2
-----} else if (z[i - l] < r - i + 1) z[i] = z[i - l];------// 6f
                         -----iter(c, *k)------// 6e
-----else {------// a8
                         -----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 97
-----l = i:-----// 55
                         -----(cur->next[*c] = new go_node());------// af
-----cur->out = new out_node(*k, cur->out);------// 3f
-z[i] = r - i; r--; \}
                         ----return z;------// 78
                         -----queue<qo_node*> q;------// 2c
                         -----iter(a, go->next) q.push(a->second);------// db
4.3. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
                         ------while (!q.empty()) {------// 07
bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77 -----iter(a, r->next) {-------------------------// 18
-----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);-------// 12 -------while (st &\delta st->next.find(a->first) == st->next.end())------// \thetae
-----rep(i,0,n) P[0][i] = s[i];--------// 5c ------st = st->fail;----------// b3
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {-------// 86 -------if (!st) st = go;------------------------// θb
------P.push_back(vi(n));--------// 53 -------s->fail = st->next[a->first];------// c1
-----rep(i,0,n)-------if (s->fail) {--------// 98
------L[L[i].p = i].nr = ii(P[stp - 1][i],------// e2 -------if (!s->out) s->out = s->fail->out;------// ad
-----sort(L.begin(), L.end());-------// 5f ------out_node* out = s->out;------// b8
-----rep(i,0,n)-------while (out->next) out = out->next;------// b4
------while (cur && cur->next.find(*c) == cur->next.end())-------// df
```

----**void** clear(){ sz = 1; last = len[0] = 0; link[0] = -1; next[0].clear();----// aa

```
-----rep(i,0,size(s)) p[i+1] = (ll)p[i] * b % m;------// 8a
-----rep(i,0,size(s)) h[i+1] = ((ll)h[i] * b + s[i]) % m; }------//10
                                      #include "mod_pow.cpp"-----// c7
----int hash(int l, int r) {------// b2
                                       bool is_probable_prime(ll n, int k) {------// be
-----return (h[r+1] + m - (ll)h[l] * p[r-l+1] % m) % m; } };-----// 26
                                      ----if (~n & 1) return n == 2;------// d1
                                       ----if (n <= 3) return n == 3;------// 39
                                       ----int s = 0: ll d = n - 1:-----// 37
               5. Mathematics
                                       ----while (~d & 1) d >>= 1, s++;------// 35
5.1. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
                                       ----while (k--) {------// c8
k items out of a total of n items. Also contains an implementation of Lucas' theorem for computing
                                       -----ll a = (n - 3) * rand() / RAND_MAX + 2;------// 06
the answer modulo a prime p.
                                       -----ll x = mod_pow(a, d, n);------// 64
                                       -----if (x == 1 \mid | x == n - 1) continue;-----// 9b
int nck(int n, int k) {------// f6
                                       ------bool ok = false;------// 03
----if (n < k) return 0;------// 55
                                       -----rep(i,0,s-1) {------// 13
---k = \min(k, n - k); -----//bd
                                       -----if (x == 1) return false;-----// 5c
----rep(i,1,k+1) res = res * (n - (k - i)) / i;------// 4d
                                       -----if (x == n - 1) { ok = true; break; }-----// a1
                                       -----if (!ok) return false;-----// 37
int nck(int n, int k, int p) {-----// cf
                                      ----} return true; }------// fe
----int res = 1:------// 5c
----while (n || k) {------// e2
                                       5.5. Pollard's \rho algorithm.
----res *= nck(n % p, k % p);-----// cc
                                       // public static int[] seeds = new int[] {2,3,5,7,11,13,1031};-----// 1d
----res %= p, n /= p, k /= p;-----// 0a
                                       // public static BigInteger rho(BigInteger n, BigInteger seed) {-----// 03
----}------// d9
                                       //--- int i = 0,-----// 00
                                       //----- k = 2;-----// 79
                                       //---- BigInteger x = seed,-----// cc
                                       //----y = seed;-----// 31
5.2. Euclidean algorithm.
                                       //--- while (i < 1000000) {-----// 10
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
                                       //----- i++:-----// 8c
                                      //------- x = (x.multiply(x).add(n).subtract(BigInteger.ONE)).mod(n);-----//74
 The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
                                       //----- BigInteger d = y.subtract(x).abs().gcd(n);-----// ce
and also finds two integers x, y such that a \times x + b \times y = d.
                                      //----- if (!d.equals(BigInteger.ONE) && !d.equals(n)) \{------//\ b9\}
int eqcd(int a, int b, int& x, int& y) {------// 85
                                       //----return d;-----// 3b
----if (b == 0) { x = 1; y = 0; return a; }-----// 7b
                                       //------} -------// 7c
                                       //----- if (i == k) {------// 2c
------int d = eqcd(b, a % b, x, y);------// 34
-----x -= a / b * y;-----// 4a
                                       //--- return BigInteger.ONE;-----// 62
5.3. Trial Division Primality Testing.
                                       5.6. Sieve of Eratosthenes.
------if (n % i == 0 || n % (i + 2) == 0) return false;---------// e9 ------primes.push_back(v = (i << 1) + 3);--------// be
```

```
------for (int j = sq; j <= mx; j += v) prime[j] = false; }-------// 2e 5.11. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3);-----// 29
----delete[] prime; // can be used for O(1) lookup-----// 36
                                     #include "eqcd.cpp"-----// 55
----return primes; }------// 72
                                     vi linear_congruence(int a, int b, int n) {-----// c8
                                     ----int x, y, d = egcd(a, n, x, y);------// 7a
5.7. Divisor Sieve.
                                     ----vi res:-----// f5
                                     ----if (b % d != 0) return res;------// 30
vi divisor_sieve(int n) {------// 7f
                                     ----int x\theta = mod(b / d * x, n); ------// 48
----vi minimalDiv(n+1, 2), primes;-----// 37
                                     ----rep(k, 0, d) res.push_back(mod(x0 + k * n / d, n));-----// 7e
----if(n>=2) primes.push_back(2);------// 27
                                     ----return res;------// fe
----minimalDiv[0] = 0;-----// 02
                                      -----// c0
----for(int k=3;k<=n;k+=2) {------// 5d
                                     5.12. Numeric Integration.
------if(minimalDiv[k] == k) primes.push_back(k);-----// 75
                                     double integrate(double (*f)(double), double a, double b,-----// 76
-----rep(i, 1, size(primes))-----// 49
                                     ------double delta = 1e-6) {------// c0
------if(primes[i] > minimalDiv[k] || primes[i]*k > n) break;------// 53
                                     ----if (abs(a - b) < delta)----------------// 38
------else minimalDiv[primes[i]*k] = primes[i];------// 90
                                     -----return (b-a)/8 *-----// 56
----}------// 9d
                                     -----(f(a) + 3*f((2*a+b)/3) + 3*f((a+2*b)/3) + f(b));-----// e1
----return primes; }-----// 93
                                     ----return integrate(f, a,-----// 64
_____// a8
                                     -----(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);-----// θε
                                     }-----// 4b
5.8. Modular Multiplicative Inverse.
                                     5.13. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete
#include "eacd.cpp"-----// 55
                                     Fourier transform. The fft function only supports powers of twos. The czt function implements the
_____// e8
                                     Chirp Z-transform and supports any size, but is slightly slower.
int mod_inv(int a, int m) {------// 49
                                     #include <complex>-----// 8e
----int x, y, d = eqcd(a, m, x, y);------// 3e
                                     typedef complex<long double> cpx;------// 25
----if (d != 1) return -1:-----// 20
                                     // NOTE: n must be a power of two-----// 14
----return x < 0 ? x + m : x;------// 3c
                                     void fft(cpx *x, int n, bool inv=false) {------// 36
}-----// 69
                                     ----for (int i = 0, j = 0; i < n; i++) {------// f9
                                     -----if (i < j) swap(x[i], x[j]);-----// 44
5.9. Primitive Root.
                                     -----int m = n>>1;-----// 9c
#include "mod_pow.cpp"-----// c7
                                     -------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
ll primitive_root(ll m) {------// 8a
                                     -----j += m;-----// 11
----vector<ll> div:-----// f2
                                     ----}--------// d0
----for (ll i = 1; i*i <= m-1; i++) {------// ca
                                     ----for (int mx = 1; mx < n; mx <<= 1) {------// 15
-----if ((m-1) % i == 0) {------// 85
                                     -----if (i < m) div.push_back(i);-----// fd
                                     ------for (int m = 0; m < mx; m++, w *= wp) {------// dc
------if (m/i < m) div.push_back(m/i); } }-----// f2
                                     ------for (int i = m; i < n; i += mx << 1) {------// 6a
----rep(x,2,m) {------// 57
                                     -----cpx t = x[i + mx] * w;-----// 12
------bool ok = true:-----// 17
                                     -----x[i + mx] = x[i] - t;
-----iter(it,div) if (mod_pow(x, *it, m) == 1) { ok = false; break; }-----// a4
                                     -----x[i] += t;-----// 0e
------if (ok) return x; }------// 55
                                     ----return -1; }------// 15
                                     ------}-----// a4
                                     ----}------// bf
5.10. Chinese Remainder Theorem.
                                     ----if (inv) rep(i,0,n) x[i] /= cpx(n);-----// 16
```

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-----*c = new cpx[n], *a = new cpx[len],------// 4e ----if (l == 1) { y[0] = x[0].inv(); return; }-----// 71
-----*b = new cpx[len];------// 30 ---inv(x, y, l>>1);------// bd
-----x[i] = c[i] * a[i];--------// 77 ----if (l == 1) { assert(x[0].x == 1); y[0] = 1; return; }------// 64
------if (inv) x[i] /= cpx(n);---------// b1 ----sqrt(x, y, l>>1);------------// a7
}------// c6 ----rep(i,0,l<<1) T2[i] = T1[i] * T2[i];------// 9f
                           ----ntt(T2, l<<1, true);------// eb
                           5.14. Number-Theoretic Transform.
#include "../mathematics/primitive_root.cpp"------8
                          5.15. Tridiagonal Matrix Algorithm. Solves a tridiagonal system of linear equations a_i x_{i-1} +
int mod = 998244353, q = primitive_root(mod),....// 9c
                          b_i x_i + c_i x_{i+1} = d_i where a_1 = c_n = 0. Beware of numerical instability.
#define MAXN 5000----// f7
#define MAXN (1<<22)-----// 94
                           long double A[MAXN], B[MAXN], C[MAXN], D[MAXN], X[MAXN];------// d8
struct Num {-----// c5
                          void solve(int n) {------// 01
----int x:------// 02
                           ----C[0] /= B[0]; D[0] /= B[0];-----// 94
----Num(ll _x=0) { x = (_x \mod + \mod) \mod; } -----// 1b
                           ----rep(i,1,n-1) C[i] /= B[i] - A[i]*C[i-1];------// 6b
----Num operator +(const Num &b) { return x + b.x; }-----// 08
                           ----rep(i,1,n) D[i] = (D[i] - A[i] * D[i-1]) / (B[i] - A[i] * C[i-1]); ------// 33
----Num operator -(const Num &b) const { return x - b.x; }------// 89
                           ---X[n-1] = D[n-1];
----Num operator *(const Num &b) const { return (ll)x * b.x; }-----// e3
                           ----for (int i = n-2; i>=0; i--) X[i] = D[i] - C[i] * X[i+1]; }------// ad
----Num operator /(const Num &b) const { return (ll)x * b.inv().x; }------// 2a
----Num inv() const { return mod_pow((ll)x, mod-2, mod); }-----// d3
                          5.16. Mertens Function. Mertens function is M(n) = \sum_{i=1}^n \mu(i). Let L \approx (n \log \log n)^{2/3} and the
----Num pow(int p) const { return mod_pow((ll)x, p, mod); }-----// d5
                           algorithm runs in O(n^{2/3}). Can also be easily changed to compute the summatory \Phi.
} T1[MAXN], T2[MAXN];-----// d5
void ntt(Num x[], int n, bool inv = false) {------// 24
                          #define L 9000000-----// 27
                          int mob[L], mer[L];.....// f1
----Num z = inv ? qinv : q;-----// 00
------if (i < j) swap(x[i], x[j]);----------// eb ----if (n < L) return mer[n];----------// 1c
------Num wp = z.pow(p), w = 1;-------// 35 ----return mem[n] = 1 - ans; }------// c2
-----Num t = x[i + mx] * w;-------// 22 ----for (int i = 2; i < L; i++) {--------// 8e
-----x[i] = x[i] + t; } }------// 1e ------mob[i] = -1;-----------// e5
-----Num ni = Num(n).inv();---------// 18 -------mer[j] = 0, mob[j] = (j/i)\%i == 0 ? 0 : -mob[j/i];-----// 26
```

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                                              ----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);------// a5
5.17. Numbers and Sequences. Some random prime numbers: 1031, 32771, 1048583, 33554467,
1073741827,\ 34359738421,\ 1099511627791,\ 35184372088891,\ 1125899906842679,\ 36028797018963971.
                                              ----else if (abs(a - b) < EPS) x = abs(a - closest\_point(c, d, a, true));-----// 23
                                              ----else if (abs(c - d) < EPS) x = abs(c - closest_point(a, b, c, true));-----// 53
                                              6. Geometry
                                              ----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;-----// bf
6.1. Primitives.
                                              ----else {------// e1
#define P(p) const point &p-----// 2e
                                              -----x = min(x, abs(a - closest_point(c,d, a, true)));------// 29
#define L(p0, p1) P(p0), P(p1)-----// cf
                                              -----x = min(x, abs(b - closest_point(c,d, b, true)));------// fe
#define C(p0, r) P(p0), double r-----// f1
                                              -----x = min(x, abs(c - closest_point(a,b, c, true)));-----// 81
#define PP(pp) pair<point, point> &pp-----// e5
                                              -----x = min(x, abs(d - closest_point(a,b, d, true)));------// e4
typedef complex<double> point;-----// 6a
                                              ----}------// c5
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// d2
                                              ----return x:-----// b7
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// 8a
                                                  .....// 27
point rotate(P(p), double radians = pi / 2, P(about) = point(\theta, \theta)) {------// 23
                                              bool intersect(L(a, b), L(p, q), point &res, bool segment = false) {------// d2
----return (p - about) * exp(point(0, radians)) + about; }-----// 25
                                              ----// NOTE: check for parallel/collinear lines before calling this function---// 1b
point reflect(P(p), L(about1, about2)) {------// 50
                                              ----point r = b - a, s = q - p;------// 34
----point z = p - about1, w = about2 - about1;------// 8b
                                              ----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// \theta b
----return conj(z / w) * w + about1; }-----// 83
                                              ----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// e4
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// e7
                                              -----return false;-----// e3
point normalize(P(p), double k = 1.0) {-----// 5f
                                              ----res = a + t * r:-----// 47
----return abs(p) == \theta ? point(\theta,\theta) : p / abs(p) * k; }------// 4a
                                              ----return true:-----// 05
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// 27
                                              1------// 44
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// b3
                                              ,
-----// cc
double angle(P(a), P(b), P(c)) {------// 61
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }-----// c7
                                              6.3. Circles.
double signed_angle(P(a), P(b), P(c)) {------// 4a
                                              #include "primitives.cpp"-----// e0
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }-----// 40
                                              int intersect(C(A, rA), C(B, rB), point & res1, point & res2) {------// 3b
double angle(P(p)) { return atan2(imag(p), real(p)); }-----// e6
                                              ----double d = abs(B - A);-----// 5c
point perp(P(p)) { return point(-imag(p), real(p)); }-----// d9
double progress(P(p), L(a, b)) {-----// b3
                                              ----if ((rA + rB) < (d - EPS) || d < abs(rA - rB) - EPS) return 0;-------// 39
----if (abs(real(a) - real(b)) < EPS)-----// 5e
                                             -----return (imag(p) - imag(a)) / (imag(b) - imag(a));-----// 5e
                                              ----point v = normalize(B - A, a), u = normalize(rotate(B-A), h);
----else return (real(p) - real(a)) / (real(b) - real(a)); }-----// 31
                                              ----res1 = A + v + u, res2 = A + v - u;------// 24
                                              ----if (abs(u) < EPS) return 1; return 2;------// 82
-----// 53
-----// 46 }------// bb
                                              int intersect(L(A, B), C(0, r), point & res1, point & res2) {------// 0e
6.2. Lines.
                                              ---- double h = abs(0 - closest_point(A, B, 0));-----// 24
#include "primitives.cpp"------// e0 --- if(r < h - EPS) return 0;------------// df
-----// 85 ---- point H = proj(0 - A, B - A) + A, v = normalize((B - A), sqrt(r*r - h*h));// 19
point closest_point(L(a, b), P(c), bool segment = false) {------// f2 int tangent(P(A), C(O, r), point & res1, point & res2) {------// aa
------if (dot(b - a, c - b) > 0) return b:-------// 88 ----if (d < r - EPS) return 0:-------// ca
------if (dot(a - b, c - a) > 0) return a;--------// 75 ----double alpha = asin(r / d), L = sqrt(d*d - r*r);--------// 3f
----double t = dot(c - a, b - a) / norm(b - a);-------// 62 ----res1 = A + rotate(v, alpha); res2 = A + rotate(v, -alpha);------// be
----return a + t * (b - a);------// 6e ----if (abs(r - d) < EPS || abs(v) < EPS) return 1;------// bb
}------// 8c ---return 2:------// b9
```

```
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 -----return x * p.x + y * p.y + z * p.z; }------// 60
-----return point3d(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }-----// ed
6.8. Triangle Circumcenter. Returns the unique point that is the same distance from all three
                                                      ----double length() const {------// 3e
points. It is also the center of the unique circle that goes through all three points.
                                                      -----return sqrt(*this % *this); }-----// 05
#include "primitives.cpp"-----// e0
                                                      ----double distTo(P(p)) const {------// dd
point circumcenter(point a, point b, point c) {-----// 76
                                                      -----return (*this - p).length(); }------// 57
----b -= a, c -= a;-----// 41
                                                      ----double distTo(P(A), P(B)) const {------// bd
----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a
                                                      -----// A and B must be two different points-----// 4e
}-----// c3
                                                     -----return ((*this - A) * (*this - B)).length() / A.distTo(B); }------// 6e
                                                      ----point3d normalize(double k = 1) const {------// db
6.9. Closest Pair of Points.
                                                      -----// length() must not return 0-----// 3c
#include "primitives.cpp"-----// e0
                                                      ------return (*this) * (k / length()); }------// d4
-----// 85
                                                      ----point3d getProjection(P(A), P(B)) const {------// 86
struct cmpx { bool operator ()(const point &a, const point &b) {------// 01
                                                      -----point3d v = B - A;-----// 64
------return abs(real(a) - real(b)) > EPS ?-----// e9
                                                      ------return A + v.normalize((v % (*this - A)) / v.length()); }------// 53
-----real(a) < real(b) : imag(a) < imag(b); } };-----// 53
                                                      ----point3d rotate(P(normal)) const {------// 55
struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
                                                      -----// normal must have length 1 and be orthogonal to the vector-----// eb
----return abs(imag(a) - imag(b)) > EPS ?-----// θb
                                                          return (*this) * normal; }-----// 5c
-----imag(a) < imag(b) : real(a) < real(b); } };------// a4
                                                      ----point3d rotate(double alpha, P(normal)) const {------// 21
double closest_pair(vector<point> pts) {------// f1
                                                      -----return (*this) * cos(alpha) + rotate(normal) * sin(alpha); }------// 82
----sort(pts.beqin(), pts.end(), cmpx());-----// 0c
                                                      ----point3d rotatePoint(P(0), P(axe), double alpha) const{-----------------------// 7a
----set<point, cmpy> cur;-----// bd
                                                      -----point3d Z = axe.normalize(axe % (*this - 0));------// ba
----set<point, cmpy>::const_iterator it, jt;------// a6
                                                      -----return 0 + Z + (*this - 0 - Z).rotate(alpha, 0); }------// 38
----double mn = INFINITY;-----// f9
                                                      ----bool isZero() const {------// 64
----for (int i = 0, l = 0; i < size(pts); i++) {-------// ac
                                                      -----return abs(x) < EPS && abs(y) < EPS && abs(z) < EPS; }-----// 15
------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);------// 8b
                                                      ----bool isOnLine(L(A, B)) const {------// 30
-----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));-----// fc
                                                      -----return ((A - *this) * (B - *this)).isZero(); }------// 58
-----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));------// 39
                                                      ----bool isInSegment(L(A, B)) const {------// f1
------while (it != jt) mn = min(mn, abs(*it - pts[i])), it++;------// 09
                                                      -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < EPS; }------// d9
-----cur.insert(pts[i]); }-----// 82
                                                      ----bool isInSegmentStrictly(L(A, B)) const {-------// 0e
----return mn: }-----// 4c
                                                      -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < -EPS; }------// ba
                                                      ----double getAngle() const {------// 0f
6.10. 3D Primitives.
                                                      -----return atan2(y, x); }------// 40
#define P(p) const point3d &p-----// a7
                                                      ----double getAngle(P(u)) const {------// d5
#define L(p0, p1) P(p0), P(p1)-----// Of
                                                      -----return atan2((*this * u).length(), *this % u); }------// 79
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)-----// 67
                                                      ----bool isOnPlane(PL(A, B, C)) const {------// 8e
struct point3d {-----// 63
                                                      -----return abs((A - *this) * (B - *this) * (C - *this)) < EPS; } };-----// 74
----double x, y, z;-----// e6
                                                      int line_line_intersect(L(A, B), L(C, D), point3d &0){-------------------// dc
----point3d() : x(0), y(0), z(0) {}-----// af
                                                      ----if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;------// 6a
----point3d(double _x, double _y, double _z) : x(_x), y(_y), z(_z) {}------// fc
                                                      ----if (((A - B) * (C - D)).length() < EPS)------// 79
----point3d operator+(P(p)) const {------// 17
                                                      -----return A.isOnLine(C, D) ? 2 : 0;-----// 09
-----return point3d(x + p.x, y + p.y, z + p.z); }-----// 8e
                                                      ----point3d normal = ((A - B) * (C - B)).normalize();-----// bc
----point3d operator-(P(p)) const {------// fb
                                                      ----double s1 = (C - A) * (D - A) % normal;------// 68
-----return point3d(x - p.x, y - p.y, z - p.z); }-----// 83
                                                      ----point3d operator-() const {------// 89
                                                      ----return 1; }-----// a7
-----return point3d(-x, -y, -z); }-----// d4
                                                      int line_plane_intersect(L(A, B), PL(C, D, E), point3d & 0) {------------// 09
----point3d operator*(double k) const {------// 4d
                                                      ----double V1 = (C - A) * (D - A) % (E - A);------// c1
-----return point3d(x * k, y * k, z * k); }-----// fd
                                                      ----double V2 = (D - B) * (C - B) % (E - B);------// 29
----point3d operator/(double k) const {------// 95
                                                      ----if (abs(V1 + V2) < EPS)------// 81
-----return point3d(x / k, y / k, z / k); }-----// 58
                                                      -----return A.isOnPlane(C, D, E) ? 2 : 0;-----// d5
----double operator%(P(p)) const {------// d1
```

```
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----0 = A + ((B - A) / (V1 + V2)) * V1;------// 38 // if (h > 1) {------// 26
bool plane_plane_intersect(P(A), P(nA), P(B), P(nB), point3d &P, point3d &Q) \{-\frac{1}{50}\}
----point3d n = nA * nB;------// 49 //--- rep(i,0,h) {-------// 1d
----return true: }------------------------// 1a //--- }-------------------// 1e
                                   //--- caliper A(hull[a], pi/2), B(hull[b], 3*pi/2);-----// 60
6.11. Polygon Centroid.
                                   //--- double done = 0:----// 3c
#include "polygon.cpp"-----// 58
                                   //--- while (true) {-----// 31
point polygon_centroid(polygon p) {------// 79
                                   //---- mx = max(mx, abs(point(hull[a].first,hull[a].second)-----// e3
----double cx = 0.0. cy = 0.0:-----// d5
                                   //------ - point(hull[b].first, hull[b].second)));-----// 24
----double mnx = 0.0, mny = 0.0;-----// 22
                                   //------ double tha = A.angle_to(hull[(a+1)%h]),------// 57
----int n = size(p);------// 2d
                                   //----- thb = B.angle_to(hull[(b+1)%h]);-----// f1
----rep(i,0,n)------// 08
                                   //----- if (tha <= thb) {------// 91
-----mnx = min(mnx, real(p[i])),-----// c6
                                   //----- A.rotate(tha);-----// c9
-----mny = min(mny, imag(p[i]));------// 84
                                   //----- B.rotate(tha);-----// f4
----rep(i,0,n)------// 3f
                                   //---- a = (a+1) % h:----// d4
-----p[i] = point(real(p[i]) - mnx, imag(p[i]) - mny);------// 49
                                   //----- A.move_to(hull[a]):----// b3
----rep(i,0,n) {------// 3c
                                   //-----} else {------// 56
------int j = (i + 1) % n;-----// 5b
                                   //----- A.rotate(thb);-----// 56
//----- B.rotate(thb);-----// 38
----return point(cx, cy) / 6.0 / polygon_area_signed(p) + point(mnx, mny); }---// a1
                                   //------ B.move_to(hull[b1):-----// 38
                                   //-----}
6.12. Rotating Calipers.
                                   //----- done += min(tha, thb);-----// d2
#include "primitives.cpp"-----// e0
                                    //----- if (done > pi) {-----// c2
struct caliper {-----// 8e
                                    //----- break:----// e8
----ii pt;------// 05
                                   //-----}
----double angle;-----// d4
                                   //---- }------// ac
----caliper(ii _pt, double _angle) : pt(_pt), angle(_angle) { }------// 35
----double angle_to(ii pt2) {------// 8b
-----double x = angle - atan2(pt2.second - pt.second, pt2.first - pt.first);// 1e
                                   6.13. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
------while (x \ge pi) x = 2*pi;------// 4a
                                     • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
------while (x <= -pi) x += 2*pi;------// a3
                                     • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
-----return x; }-----// 7d
                                     • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
----void rotate(double by) {------// 57
                                      of that is the area of the triangle formed by a and b.
-----angle -= by;-----// 5d
                                     • Euler's formula: V - E + F = 2
------while (angle < 0) angle += 2*pi;-----// 03
                                     • Side lengths a, b, c can form a triangle iff. a + b > c, b + c > a and a + c > b.
----}-----// 20
                                     • Sum of internal angles of a regular convex n-gon is (n-2)\pi.
----void move_to(ii pt2) { pt = pt2; }-----// 37
----double dist(const caliper &other) {------// 68
                                                 7. Other Algorithms
------point a(pt.first,pt.second),-----// d7
------ c(other.pt.first, other.pt.second);--------// 71 #include "../graph/scc.cpp"------
-----return abs(c - closest_point(a, b, c));------// 58
                                   -----// 63
-----// c5 ----all_truthy.clear();-------// 31
```

```
-----adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// eb ----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
------if (clauses[i].first != clauses[i].second)--------// bc -----arr = new bool*[rows]:--------// cf
----pair<union_find, vi> res = scc(adj);-------// dd ------arr[i] = new bool[cols], memset(arr[i], 0, cols);------// dd
------int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);-// 4f ------rep(i,0,rows+1) {--------// 76
-----if (cur == 0) continue;-------// cd ------ptr[i] = new node*[cols];-------// eb
------if (p == 0) return false;----------// d0 ------rep(j,0,cols)--------// cd
-----truth[cur + n] = truth[p];------// 50 ------else ptr[i][j] = NULL;------// d2
------if (truth[p] == 1) all_truthy.push_back(cur);-------// 55 -----rep(i,0,rows+1) {-----------------------// fc
}--------int ni = i + 1, nj = j + 1;-------// 7a
                                ------while (true) {------// fc
7.2. Stable Marriage.
                                -----if (ni == rows + 1) ni = 0;------// 4c
-----if (ni == rows || arr[ni][j]) break;-----// 8d
----queue<int> a;-----// f6
                                -----+ni:----// 68
----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
                                ----rep(i,0,n) rep(j,0,n) inv[i][w[i][j]] = j;------// f1
                                -----ptr[i][j]->d = ptr[ni][j];-----// 84
----rep(i,0,n) q.push(i);-----// d8
                                -----ptr[ni][j]->u = ptr[i][j];-----// 66
----while (!q.empty()) {------// 68
                                -----while (true) {------// 7f
------int curm = q.front(); q.pop();------// e2
                                -----if (nj == cols) nj = 0;-----// de
------for (int &i = at[curm]; i < n; i++) {-------// 7e
                                -----if (i == rows || arr[i][nj]) break;-----// 4c
-----int curw = m[curm][i];-----// 95
                                -----+ni;-----// c5
------if (eng[curw] == -1) { }-----// f7
                                ------else if (inv[curw][curm] < inv[curw][eng[curw]])------// d6
                                -----ptr[i][j]->r = ptr[i][nj];-----// 60
-----q.push(enq[curw]);------// 2e
                                -----ptr[i][nj]->l = ptr[i][j];-----// 82
-----else continue;-----// 1d
                                -----/<sub>------</sub>/ 0b
-----res[eng[curw] = curm] = curw, ++i; break;-----// a1
                                -----}-----// c4
                                -----head = new node(rows, -1):-----// 66
----}-----// 3d
                                ------head->r = ptr[rows][0];-----// 3e
----return res:------// 42
                                -----ptr[rows][0]->l = head;------// 8c
}-----// bf
                                -----head->l = ptr[rows][cols - 1];------// 6a
                                -----ptr[rows][cols - 1]->r = head;-----// c1
7.3. Algorithm X.
                                -----rep(j,0,cols) {------// 92
bool handle_solution(vi rows) { return false; }------// 63
                                -----int cnt = -1;-----// d4
struct exact_cover {------// 95
                                -----rep(i,0,rows+1)-----// bd
----struct node {------// 7e
                                -----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];------// f3
-----node *l. *r. *u. *d. *p:-----// 19
                                -----ptr[rows][j]->size = cnt;-----// c2
-----int row, col, size;-----// ae
-----node(int _row, int _col) : row(_row), col(_col) {------// c9
                                -----rep(i,0,rows+1) delete[] ptr[i];-----// a5
-----size = 0; l = r = u = d = p = NULL; }-----// c3
                                -----delete[] ptr;-----// 72
----}:-----// c1
                                ----int rows. cols. *sol:-----// 7b
                                ----#define COVER(c, i, j) \sqrt{\phantom{a}}
----bool **arr:------// e6
                                ------c->r->l = c->l, c->l->r = c->r; \\----------------------// 82
----node *head;-----// fe
```

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```
------for (node *i = c->d; i != c; i = i->d) \------// 62 7.6. Dates.
-----j->d->u = j->u, j->u->d = j->d, j->p->size--;------// c1
                               int dateToInt(int y, int m, int d) {------// 96
----#define UNCOVER(c, i, j) \|------// 89
                               ----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +------// a8
------for (node *i = c->u; i != c; i = i->u) \[ \]------// f0
                              -----367 * (m - 2 - (m - 14) / 12 * 12) / 12 -----// d1
------for (node *j = i->l; j != i; j = j->l) \sqrt{\phantom{a}}
                               -----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
                               -----d - 32075;-----// e0
-----j->p->size++, j->d->u = j->u->d = j; \\ \]
void intToDate(int jd, int &y, int &m, int &d) {------// a1
----bool search(int k = 0) {------// f9
                               ----int x, n, i, j;------// 00
-----if (head == head->r) {------// 75
                               ---x = jd + 68569; ------//11
-----vi res(k);-----// 90
                               ---n = 4 * x / 146097;
-----rep(i,0,k) res[i] = sol[i];------// 2a
                               ---x = (146097 * n + 3) / 4;
-----sort(res.begin(), res.end());------// 63
                               -----return handle_solution(res);-----// 11
                               ----x -= 1461 * i / 4 - 31;-----// 09
                               ---j = 80 * x / 2447;-----// 3d
-----node *c = head->r, *tmp = head->r;-----// a3
                               ---d = x - 2447 * j / 80;
-----for (: tmp != head: tmp = tmp->r) if (tmp->size < c->size) c = tmp:---// 41
                               ----x = j / 11;-----// b7
-----if (c == c->d) return false;------// 02
                               ---m = j + 2 - 12 * x;
-----COVER(c, i, j);-----// f6
                               ---y = 100 * (n - 49) + i + x;
------bool found = false;-----// 8d
------for (node *r = c->d; !found && r != c; r = r->d) {------// 78
-----sol[k] = r->row;-----// cθ
                               7.7. Simulated Annealing. An example use of Simulated Annealing to find a permutation of length
-----for (node *j = r - r; j != r; j = j - r) { COVER(j - p, a, b); } -----// f9
                               n that maximizes \sum_{i=1}^{n-1} |p_i - p_{i+1}|.
-----found = search(k + 1):-----// fb
                              double curtime() { return static_cast<double>(clock()) / CLOCKS_PER_SEC; }-----// 9d
------for (node *j = r->l; j != r; j = j->l) { UNCOVER(j->p, a, b); }----// 87
                              int simulated_annealing(int n, double seconds) {------// 54
                               ----default_random_engine rng;------// 67
------UNCOVER(c, i, j);------// a7
                               ----uniform_real_distribution<double> randfloat(0.0, 1.0);------// ed
                              ----uniform_int_distribution<int> randint(0, n - 2);------// bb
};-----// d7
                              ----// random initial solution-----// 01
                               ----vi sol(n);------// 1c
7.4. nth Permutation.
                               ----rep(i,0,n) sol[i] = i + 1;------// 33
----while (true) {------// 46
7.5. Cycle-Finding.
                               -----if (!(iters & ((1 << 4) - 1))) {------// 5d
----h = x0;------// eb
----h = f(t);-----// compute delta for mutation-----// 84
----return ii(mu, lam);------if (a > 0) delta += abs(sol[a+1] - sol[a-1]) - abs(sol[a] - sol[a-1]);-// 94
}------if (a+2 < n) delta += abs(sol[a] - sol[a+2]) - abs(sol[a+1] - sol[a+2]);
```

8. Useful Information

8.1. Fast Input Reading.

```
void readn(register int *n) {
    int sign = 1;
    register char c;
    *n = 0;
    while((c = getc_unlocked(stdin)) != '\n') {
        switch(c) {
            case '-': sign = -1; break;
            case ' ': goto hell;
            case '\n': goto hell;
            default: *n *= 10; *n += c - '0'; break;
hell:
    *n *= sign;
```

8.2. Bit Hacks.

```
dp[i][j] \neq \min_{i < k < j} dp[i][k] + dp[k][j] + C[i][j]
```

9. Misc

9.1. Debugging Tips.

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
 - Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps 1+eps).
 - Rounding negative numbers?
 - Outputting in scientific notation?
- Wrong Answer?
 - Read the problem statement again!
 - Are multiple test cases being handled correctly? Try repeating the same test case many times.
 - Integer overflow?
 - Think very carefully about boundaries of all input parameters
 - Try out possible edge cases:

```
* n = 0, n = -1, n = 1, n = 2^{31} - 1 or n = -2^{31}
```

- * List is empty, or contains a single element
- * n is even, n is odd
- * Graph is empty, or contains a single vertex
- * Graph is a multigraph (loops or multiple edges)
- * Polygon is concave or non-simple
- Is initial condition wrong for small cases? - Are you sure the algorithm is correct?
- Explain your solution to someone.

- Are you using any functions that you don't completely understand? Maybe STL functions?
- Maybe you (or someone else) should rewrite the solution?

9.2. Solution Ideas.

- Dynamic Programming
 - Drop a parameter, recover from others
 - Swap answer and a parameter
 - Parsing CFGs: CYK Algorithm
 - Optimizations
 - * Convex hull optimization

```
\cdot \operatorname{dp}[i] = \min_{j < i} \{ \operatorname{dp}[j] + b[j] \times a[i] \}
b[j] > b[j+1]
```

- · optionally $a[i] \le a[i+1]$
- $O(n^2)$ to O(n)
- * Divide and conquer optimization
 - $dp[i][j] = \min_{k \le i} \{dp[i-1][k] + C[k][j]\}$
 - A[i][j] < A[i][j+1]
 - · $O(kn^2)$ to $O(kn\log n)$

- $A[i][j-1] \le A[i][j] \le A[i+1][j]$
- $O(n^3)$ to $O(n^2)$
- · sufficient: QI and $C[b][c] \le C[a][d], a \le b \le c \le$
- Greedy
- Randomized
- Optimizations
 - Use bitset (/64)
 - Switch order of loops (cache locality)
- Process queries offline
 - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
 - Mo's algorithm
 - Sqrt decomposition
 - Store 2^k jump pointers
- Data structure techniques
 - Sart buckets
 - Store 2^k jump pointers
 - -2^k merging trick
- Counting
 - Inclusion-exclusion principle
 - Generating functions
- Graphs
 - Can we model the problem as a graph?
 - Can we use any properties of the graph?
 - Strongly connected components

- Cycles (or odd cycles)
- Bipartite (no odd cycles)
 - * Bipartite matching
 - * Hall's marriage theorem
 - * Stable Marriage
- Cut vertex/bridge
- Biconnected components
- Degrees of vertices (odd/even)
- - * Heavy-light decomposition
 - * Centroid decomposition
 - * Least common ancestor
- Eulerian path/circuit
- Chinese postman problem
- Topological sort
- (Min-Cost) Max Flow
- Min Cut
 - * Maximum Density Subgraph
- Huffman Coding
- Min-Cost Arborescence
- Steiner Tree
- Kirchoff's matrix tree theorem
- Prüfer sequences
- Lovász Toggle
- Look at the DFS tree (which has no cross-edges)
- Mathematics
 - Is the function multiplicative?
 - Look for a pattern
 - Permutations
 - * Consider the cycles of the permutation
 - Functions
 - * Sum of piecewise-linear functions is a piecewiselinear function
 - * Sum of convex (concave) functions is convex (concave)
 - Modular arithmetic
 - * Chinese Remainder Theorem
 - * Linear Congruence
 - Sieve
 - System of linear equations
 - Values to big to represent?
 - * Compute using the logarithm
 - * Divide everything by some large value
- Logic
 - -2-SAT
 - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half $(\log(n))$
- Strings
 - Trie (maybe over something weird, like bits)
 - Suffix array

- Suffix automaton (+DP?)
- Aho-Corasick
- eerTree
- Work with S + S
- Hashing
- Euler tour, tree to array
- Segment trees
 - Lazy propagation
 - Persistent
 - Implicit
 - Segment tree of X
- Geometry
 - Minkowski sum (of convex sets)
 - Rotating calipers
 - Sweep line (horizontally or vertically?)
 - Sweep angle
 - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)
- Computing a Convolution? Fast Fourier Transform
- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem
 - Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

Formulas

#labeled rooted trees #labeled unrooted trees $\sum_{i=1}^{n} i^2 = n(n+1)(2n+1)/6$ $!n = n \times !(n-1) + (-1)^n$ $\sum_{i=1}^{n} \binom{n}{i} F_i = F_{2n}$ $\sum_{k=0}^{n} \binom{k}{m} = \binom{n+1}{m+1}$ $a \equiv b \pmod{x,y} \Rightarrow a \equiv b \pmod{\operatorname{lcm}(x,y)}$ $ac \equiv bc \pmod{m} \Rightarrow a \equiv b \pmod{\frac{m}{\gcd(c,m)}}$ $p \text{ prime} \Leftrightarrow (p-1)! \equiv -1 \pmod{p}$ $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x} - 1}{p_i^x - 1}$ $\sum_{k=0}^m (-1)^k \binom{n}{k} = (-1)^m \binom{n-1}{m}$ $2^{\omega(n)} = O(\sqrt{n})$

- $\sum_{i} {n-i \choose i} = F_{n+1}$
- $\sum_{d|n} \phi(d) = n$

- Jacobi symbol: $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- Heron's formula: A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s=\frac{a+b+c}{2}$

- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area
- Euler's totient: The number of integers less than n that are comprime to n are $n \prod_{p|n} \left(1 - \frac{1}{p}\right)$ where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph $G = (L \cup R, E)$, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L, and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then $K = (L \setminus Z) \cup (R \cap Z)$ is the minimum vertex cover.
- A minumum Steiner tree for n vertices requires at most n-2additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- Wilson's theorem: $(n-1)! \equiv -1 \pmod{n}$ iff. n is prime
- Lagrange polynomial through points $(x_0, y_0), \ldots, (x_k, y_k)$ is $L(x) = \sum_{j=0}^{k} y_j \prod_{\substack{0 \le m \le k \\ m \ne j}} \frac{x - x_m}{x_j - x_m}$
- Hook length formula: If λ is a Young diagram and $h_{\lambda}(i,j)$ is the hook-length of cell (i, j), then then the number of Young tableux $d_{\lambda} = n! / \prod h_{\lambda}(i, j)$.
- Möbius inversion formula: If $f(n) = \sum_{d|n} g(d)$, then $g(n) = \sum_{d|n} g(d)$ $\sum_{d|n} \mu(d) f(n/d)$
- 9.3. Markov Chains. A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let $P^{(m)} = (p_{ij}^{(m)})$ be the probability matrix of transitioning from state i to state j in m timesteps, and note that $P^{(1)}$ is the adjacency matrix of the graph. Chapman-Kolmogorov: $p_{ij}^{(m+n)} = \sum_{k} p_{ik}^{(m)} p_{kj}^{(n)}$. It follows that $P^{(m+n)} = P^{(m)} P^{(n)}$ and $P^{(m)} = P^m$. If $p^{(0)}$ is the initial probability distribution (a vector), then $p^{(0)}P^{(m)}$ is the probability distribution after m timesteps. $\sum_{i=1}^{n} i^3 = n^2(n+1)^2/4$ The return times of a state i is $R_i = \{m \mid p_{ii}^{(m)} > 0\}$, and i n = (n-1)(!(n-1)+!(n-2)) is aperiodic if $gcd(R_i) = 1$. A MC is aperiodic if any of its vergence n = 1.

is strongly connected. A distribution π is stationary if $\pi P = \pi$. If MC is irreducible $(\sum_{d|n} \sigma_0(d))^2 = \sum_{d|n} \sigma_0(d)^3 = 1/\mathbb{E}[T_i]$, where T_i is the expected time between two visits $\gcd(n^a-1,n^b-1)$ at $n_i^{\gcd(a,b)}/n_i$ is the expected number of visits at j in between two $\sigma_0(n) = \prod_{i=0}^r (a_i + \text{fonsecutive visits at } i. \text{ A MC is } \text{ergodic if } \lim_{m \to \infty} p^{(0)} P^m = \pi. \text{ A}$ MC is ergodic iff. it is irreducible and aperiodic.

tices is aperiodic. A MC is *irreducible* if the corresponding graph

A MC for a random walk in an undirected weighted graph $\sum_{i=1}^{n} 2^{\omega(i)} = O(n \log n)$ weighted graph can be made weighted by adding 1-weights) has $p_{uv} = w_{uv} / \sum_{x} w_{ux}$. If the graph is connected, then $\pi_u =$ $\sum_{x} w_{ux} / \sum_{v} \sum_{x} w_{vx}$. Such a random walk is aperiodic iff. the graph is not bipartite.

An absorbing MC is of the form $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$ $\sum_{m=0}^{\infty} Q^m = (I_t - Q)^{-1}$. Then, if starting in state i, the expected number of steps till absorption is the i-th entry in N1. If starting in state i, the probability of being absorbed in state j is the (i, j)-th

Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

entry of NR.

9.4. Burnside's Lemma. Let G be a finite group that acts on a set X. For each g in G let X^g denote the set of elements in X that are fixed by q. Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

9.5. **Bézout's identity.** If (x, y) is any solution to ax+by=d (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a,b)}, y - k \frac{a}{\gcd(a,b)}\right)$$

9.6. The Twelvefold Way. Putting n balls into k boxes.

r	$_{\mathrm{Balls}}$	same	distinct	same	distinct	
,	Boxes	same	same	distinct	distinct	Remarks
=	-	$p_k(n)$	$\sum_{i=0}^{k} \begin{Bmatrix} n \\ i \end{Bmatrix}$	$\binom{n+k-1}{k-1}$	k^n	$p_k(n)$: #partition
	$\mathrm{size} \geq 1$	p(n,k)	$\left\{ \begin{smallmatrix} n \\ k \end{smallmatrix} \right\}$	$\binom{n-1}{k-1}$	$k! \begin{Bmatrix} n \\ k \end{Bmatrix}$	p(n,k): #partit
	$\mathrm{size} \leq 1$	$[n \le k]$	$[n \le k]$	$\binom{k}{n}$	$n!\binom{k}{n}$	[cond]: 1 if cond

	1 00 1
Catalan	$C_0 = 1, C_n = \frac{1}{n+1} {2n \choose n} = \sum_{i=0}^{n-1} C_i C_{n-i-1} = \frac{4n-1}{n+1}$
	$\begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1, \begin{bmatrix} n \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ n \end{bmatrix} = 0, \begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} +$
Stirling 2nd kind	$\left\{ \left\{ {n\atop 1} \right\} = \left\{ {n\atop n} \right\} = 1, \left\{ {n\atop k} \right\} = k \left\{ {n-1\atop k} \right\} + \left\{ {n-1\atop k-1} \right\}$
Euler	$\left \left\langle {n \atop 0} \right\rangle = \left\langle {n \atop n-1} \right\rangle = 1, \left\langle {n \atop k} \right\rangle = (k+1) \left\langle {n-1 \atop k} \right\rangle + (n-k)$
Euler 2nd Order	$\left \left\langle $
Bell	$B_1 = 1, B_n = \sum_{k=0}^{n-1} B_k \binom{n-1}{k} = \sum_{k=0}^n \binom{n}{k}^n$

PRACTICE CONTEST CHECKLIST

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Is int128 available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is RAND MAX?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++ and Java.
- $\bullet\,$ Try the submit script.
- Try local programs: i?python[23], factor.
- Try submitting with assert(false) and assert(false).
- Return-value from main.
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