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```
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           1. Code Templates
                             ----public static void main(String[] args) throws Exception {--------// 02
                             -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                             ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                             -----// code-----// e6
setxkbmap -option caps:escape
                             -----out.flush():-----// 56
set -o vi
                             xset r rate 150 100
                             }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                        2. Data Structures
syn on | colorscheme slate
                             2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                             struct union_find {------// 42
#include <cmath>------// 7d ----union_find(int n) { parent.resize(cnt = n);------// 92
#include <cstdio>------[i] = i; }------// 6f
#include <cstdlib>------// 11 ----int find(int i) {--------// a6
#include <cstring>-------[i] = i ? i : (parent[i] = find(parent[i])); }------// @ -------|/ a9
#include <map>-----// 02
#include <set>------// d2
#include <sstream>------ min(a, b); }------// 18 // int f(int a, int b) { return min(a, b); }------
#include <stack>------// cf const int ID = 0;------// 82
#include <string>------// a9 int f(int a, int b) { return a + b; }------// 6d
#include <vector>-----// 4f ----int n; vi data, lazy;-----// fa
------mk(arr, 0, n-1, 0); }------
#define foreach(u, o) \------// ea ----int mk(const vi &arr, int i, int r, int i) {------// 02
const int INF = 2147483647;-----// be -----int m = (l + r) / 2;-----// 0f
const double pi = acos(-1);------// 49 ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// f5
typedef long long ll;-----// 8f ----int q(int a, int b, int l, int r, int i) {-------// ad
typedef unsigned long long ull;-----// 81 -----propagate(l, r, i);-----// f7
typedef vector<vii>vvii;------// 4b ----void update(int i, int v) { u(i, v, 0, n-1, 0); }------// 65
template <class T> T mod(T a, T b) { return (a % b + b) % b; }--------// 70 ----int u(int i, int v, int l, int r, int j) {-----------// b5
template <class T> int size(const T &x) { return x.size(); }------// 68 -----propagate(l, r, j);-------// 3c
                             -----if (r < i || i < l) return data[j];------// 6a
1.3. Java Template. A Java template.
                             -----if (l == i && r == i) return data[j] = v;------// 74
import java.math.*;------return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 68
-----// a3 ----int ru(int a, int b, int v, int i) {-------------// d7
public class Main {------// 17 -----propagate(l, r, i);-------// 82
```

```
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------int m = (l + r) / 2; ------// 9d ----matrix(const matrix  other) : rows(other.rows), cols(other.cols), ------// fe
-----ceturn data[i] = f(ru(a, b, v, l, m, 2*i+1), ru(a, b, v, m+1, r, 2*i+2)); ------cnt(other.cnt), data(other.data) { }------// ed
------if (l < r) {-------// 6e ------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
------else lazv[2*i+2] += lazv[i];------// d1 -----matrix<T> res(*this); res += other; return res; }-----// 5d
-----lazv[i] = INF;------res(*this); res -= other; return res; }------// cf
};-------matrix<T> res(*this); res *= other; return res; }------// 37
                             ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                            -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                            -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i...j in O(\log n) time. It only needs O(n) space.
                            -----for (int k = 0; k < cols; k++)-----// fc
struct fenwick_tree {------res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n; vi data;------// d3 -----return res; }-----// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {--------------// dd
----void update(int at, int by) {--------// 76 ------matrix<T> res(cols, rows);------// b5
------while (at < n) data[at] += by, at |= at + 1; }------// fb -------for (int i = 0; i < rows; i++)------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);-------// bf
};------while (p) {------// cb
struct fenwick_tree_sq {------// d4 ------if (p & 1) res = res * sq;-----// c1
----int n; fenwick_tree x1, x0;-------// 18 -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;-------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {-------// 89
}:-----// 13 -------int k = r;------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f ------if (k >= rows) continue;------// 3f
----return s.query(b) - s.query(a-1); }------// f3 ------det *= T(-1);-------------------// 7a
                             ------for (int i = 0: i < cols: i++)-----// ab
2.4. Matrix. A Matrix class.
                            -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
class matrix {-----// 85
```

```
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-----for (int i = 0; i < rows; i++) {--------// 3d ------node *s = successor(n);------// 16
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 37
------} return mat; }-------// 8f ------parent_leg(n) = s, fix(s);------// 87
private:-----// e0 -----return;-------// 32
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;------// 70
};-------// b8 ---node* successor(node *n) const {-------// lb
                          -----if (!n) return NULL;-----// b3
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                          ------if (n->r) return nth(0, n->r);------// 5b
#define AVL_MULTISET 0-----// b5 -----node *p = n->p;-----// 7c
template <class T>------// 22 -----return p; }------// 03
class avl_tree {------// ff ----node* predecessor(node *n) const {------// e6
public:-----// f6 -----if (!n) return NULL;------// 96
----struct node {------// 45 -----if (n->l) return nth(n->l->size-1, n->l);------// 15
-----T item; node *p, *l, *r;--------// a6 -----node *p = n->p;-------// 33
------node(const T &item, node *p = NULL) : item(item), p(p),------// c5 -----return p; }------
------l(NULL), r(NULL), size(1), height(0) { } };--------// e2
----node *root;------// c1 ----node* nth(int n, node *cur = NULL) const {------// f4
----node* find(const T &item) const {-------// d2 -----if (!cur) cur = root;------// 0a
-----node *cur = root;-----// cf ------while (cur) {------// 55
------while (cur) {-------// ad ------if (n < sz(cur->l)) cur = cur->l;------// 25
------else if (item < cur->item) cur = cur->l;------// de -----else break;-----
-------else break; }------// 05 ------} return cur; }-------// 8f
-----return cur; }------// e7 private:-----// 49
-----node *prev = NULL, **cur = &root;-----// 60
                          ----inline int height(node *n) const { return n ? n->height : -1; }------// e4
------while (*cur) {--------// b0 ----inline bool left_heavy(node *n) const {-------// d7
-----prev = *cur;------// 31 ------return n && height(n->l) > height(n->r); }------// 9d
#endif-----// c6 ----node*& parent_leg(node *n) {-------// 0d
-----*cur = n, fix(n); return n; }------// 86 -----if (n->p->r == n) return n->p->r;-----// 0e
----void erase(const T &item) { erase(find(item)); }------// c0 -----assert(false); }-----
-----if (!n) return;------// 4d -----if (!n) return;-----// 4d
------if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;-------// f5 -----n->size = 1 + sz(n->l) + sz(n->r);------// 14
------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;-------// 3d -----n->height = 1 + max(height(n->l), height(n->r)); }------// a1
-----else if (n->l && n->r) {------// 1a
```

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-----l->p = n->p; N------// 66 ----void swim(int i) {------// 33
                                     ------while (i > 0) {------// 1a
-----parent_leg(n) = l; \[ \]------// 02
                                     -----int p = (i - 1) / 2;-----// 77
------n->l = l->r; N------// 08
                                     ------if (!cmp(i, p)) break;-----// a9
-----l->r = n, n->p = l; N------// c3 ----void sink(int i) {-------// ca
----void left_rotate(node *n) { rotate(r, l); }------// 43 -------int l = 2*i + 1, r = l + 1;------// b4
------while (n) { augment(n);-------// c9 ------if (!cmp(m, i)) break;------// 42
------if (too_heavy(n)) {-------// a9 -----swp(m, i), i = m; } }-----// 1d
------else if (right_heavy(n) δδ left_heavy(n->r))------// 09 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {------// 17
-----right_rotate(n->r);------// 7c -----q = new int[len], loc = new int[len];------// f8
------if (left_heavy(n)) right_rotate(n);------// 44 -----memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// 02 ----~heap() { delete[] loc; }------// 09
-----n = n->p; }------// af ----void push(int n, bool fix = true) {-------// b7
#ifdef RESIZE-----// a9
 Also a very simple wrapper over the AVL tree that implements a map interface.
                                     ------int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"------// 01
                                     -----while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                                     ------int *newg = new int[newlen], *newloc = new int[newlen];------// e3
template <class K, class V>-----// da
                                     -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i];--// 94
class avl_map {-----// 3f
                                     -----/ 18 emset(newloc + len, 255, (newlen - len) << 2);------// 18
public:----// 5d
                                     -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                                     -----loc = newloc, q = newq, len = newlen;-----// 61
-----K kev: V value:-----// 32
                                     #else-----// 54
-----/ 29 key(k), value(v) { }-----// 29
                                     ------assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                                     #endif------// 64
----avl_tree<node> tree;------// b1
                                     ----V& operator [](K key) {------// 7c
                                     -----assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(\theta)));------// ba
                                     -----loc[n] = count, q[count++] = n;-----// 6b
-----if (!n) n = tree.insert(node(key, V(0)));-----// cb
                                     -----if (fix) swim(count-1); }-----// bf
-----return n->item.value;------// ec
                                     ----void pop(bool fix = true) {-------// 43
----}------// 2e
                                     -----assert(count > 0);-----// eb
};-----// af
                                     -----loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 50
                                     -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                                     ----}-----------// 16
#define RESIZE-----// d0
                                     ----int top() { assert(count > 0): return q[0]: }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                                     ----void heapify() { for (int i = count - 1; i > 0; i--)-----// 39
struct default_int_cmp {------// 8d
                                     ------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }------// θb
----default_int_cmp() { }------// 35
                                     ----void update_key(int n) {------// 26
----bool operator ()(const int &a, const int &b) { return a < b; } };-----// e9
                                     -----assert(loc[n] != -1), swim(loc[n]), sink(loc[n]); }------// 7d
template <class Compare = default_int_cmp>------// 30
                                     ----bool empty() { return count == 0; }-----// f8
class heap {-----// 05
                                     ----int size() { return count; }-------// 86
private:-----// 39
                                     ----void clear() { count = 0, memset(loc, 255, len << 2); } };------// 58
----int len, count, *q, *loc, tmp;-----// 0a
----Compare _cmp:-----// 98
----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }-----// a0 2.7. Skiplist. An implementation of a skiplist.
```

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----unsigned int cnt = 0;------// 28 ----void erase(T target) {-------// 4d
-----T item:------update[i]->next[i] = x->next[i];------// 59
------int *lens;------update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
------delete x; _size--;-------delete x; _size--;-------// 81
------while(current_level); free(lens); free(next); }; };-------// aa --------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
----node *head;------// b7
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
----~skiplist() { clear(); delete head; head = NULL; }-----// aa
                          2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----#define FIND_UPDATE(cmp, target) \[\bar{N}\]------------------------------// c3 list supporting deletion and restoration of elements.
------int pos[MAX_LEVEL + 2]; \[\bar{\cappa}\]-------// 18
                          template <class T>-----// 82
-----memset(pos, 0, sizeof(pos)); \[ \] ------// f2
                          struct dancing_links {-----// 9e
-----node *l, *r;-----// 32
----: item(item), l(l), r(r) {------// 04
-----pos[i] = pos[i + 1]; N-----// 68
                          -----if (l) l->r = this:-----// 1c
-----if (r) r->l = this;-----// 0b
-----pos[i] += x->lens[i]; x = x->next[i]; } \[ \frac{10}{10} \]
----void clear() { while(head->next && head->next[0])-------// 91 ----node *push_back(const T &item) {------// d7
------back = new node(item, back, NULL);-------// 5d
----node *find(T target) { FIND_UPDATE(x->next[i]->item, target);------// 36 -----if (!front) front = back;-----------------------// a2
------return x && x->item == target ? x : NULL; }------// 50 -----return back;------
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ----node *push_front(const T &item) {-------------------// ea
------front = new node(item, NULL, front);------// 75
------FIND_UPDATE(x->next[i]->item, target);-------// 3a -----return front;-----
------int lvl = bernoulli(MAX_LEVEL);-------// 7a ----void erase(node *n) {------// 88
------if(lvl > current_level) current_level = lvl;-------// 8a -----if (!n->l) front = n->r; else n->l->r = n->r;------// d5
-----x = new node(lvl, target);-------// 36 ------if (!n->r) back = n->l; else n->r->l = n->l;------// 96
-----x->next[i] = update[i]->next[i];------// 46 ----void restore(node *n) {------// 6d
-----x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];------// bc -----if (!n->l) front = n; else n->l->r = n;------// ab
------update[i]->next[i] = x;-------// 20 ------if (!n->r) back = n; else n->r->l = n;------// 8d
```

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                                          -----vi& adj = adj_list[cur.first];-----// f9
2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                                          ------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)------// 44
element.
                                          ------if (visited.find(*it) == visited.end()) {------// 8d
#define BITS 15-----// 7b
                                           ------Q.push(ii(*it, cur.second + 1));-------------------// ab
struct misof_tree {-----// fe
----int cnt[BITS][1<<BITS];------// aa
                                          -----visited.insert(*it);-----// cb
                                           ----misof_tree() { memset(cnt, 0, sizeof(cnt)); }-----// b0
                                          ----void insert(int x) { for (int i = 0; i < BITS; cnt[i++][x]++, x >>= 1); }--// 5a
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]---, x >>= 1); }---// 49
                                          ----return -1:------// f5
----int nth(int n) {-------// 8a
                                          }-----// 03
-----int res = 0;------// a4
------for (int i = BITS-1; i >= 0; i--)------// 99
                                          3.2. Single-Source Shortest Paths.
-----if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4
-----return res;------// 3a
                                          3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|)
----}-----// b5
}:-----// @a
                                          int *dist. *dad:-----// 46
                                          struct cmp {-----// a5
                  3. Graphs
                                          ----bool operator()(int a, int b) {-----// bb
                                          -----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
                                          };-----// 41
edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
                                          pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
                                          ----dist = new int[n];-----// 84
connected. It runs in O(|V| + |E|) time.
                                          ----dad = new int[n];------// 05
int bfs(int start, int end, vvi& adj_list) {------// d7
                                          ----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;-------// d6
----queue<ii>> 0:------// 75
                                          ----set<<u>int</u>, cmp> pq;-----// 04
----Q.push(ii(start, 0));-----// 49
                                          ----dist[s] = 0, pq.insert(s);------// 1b
-----// <code>0b</code>
                                          ----while (!pq.empty()) {-----// 57
----while (true) {------// 0a
                                          ------int cur = *pq.begin(); pq.erase(pq.begin());-----// 7d
-----ii cur = Q.front(); Q.pop();-----// e8
                                          ------for (int i = 0; i < size(adj[cur]); i++) {------// 9e
-----// 06
                                          -----int nxt = adj[cur][i].first,-----// b8
-----if (cur.first == end)-----// 6f
                                          -----ndist = dist[cur] + adj[cur][i].second;------// 0c
-----return cur.second;------// 8a
                                          -----// 3c
                                          -----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);------// 0f
-----vi& adj = adj_list[cur.first];-----// 3f
                                          ------for (vi::iterator it = adj.begin(); it != adj.end(); it++)-----// bb
                                          ----}--------// e8
-----0.push(ii(*it, cur.second + 1));-----// b7
                                          ----return pair<int*, int*>(dist, dad);-----// cc
}-----// 7d
 Another implementation that doesn't assume the two vertices are connected. If there is no path
                                          3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
from the starting vertex to the ending vertex, a-1 is returned.
                                          problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
                                          negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
int bfs(int start, int end, vvi& adj_list) {------// d7
----set<<u>int</u>> visited;-----// b3
----Q.push(ii(start, 0));------// 3a ----has_negative_cycle = false;-------// 47
-----// db ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;-------// 10
-----ii cur = Q.front(); Q.pop();-------// 03 ------for (int j = 0; j < n; j++)-------// c4
------if (cur.first == end)-------// 22 ------for (int k = 0; k < size(adj[j]); k++)-----// 3f
```

-----/_{ba} ------dist[j] + adj[j][k].second);-------// 47

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------has_negative_cycle = true;-------// 2a ----return pair<union_find, vi>(uf, dag);-------// 94
----return dist;-------// 2e }------// 97
}-----// c2
                            3.5. Minimum Spanning Tree.
3.3. All-Pairs Shortest Paths.
                            3.5.1. Kruskal's algorithm.
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                            #include "../data-structures/union_find.cpp"-----// 5e
problem in O(|V|^3) time.
                            -----// 11
void floyd_warshall(int** arr, int n) {------// 21 // n is the number of vertices-----// 18
----for (int k = 0; k < n; k++)-------// 49 // edges is a list of edges of the form (weight, (a, b))------// c6
------for (int j = 0; j < n; j++)-------// 77 vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
------arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1 ----sort(edges.begin(), edges.end());-----
}------// 86 ----vector<pair<int, ii>> res;-----------------------------------// 71
                            ----for (int i = 0; i < size(edges); i++)------// ce
3.4. Strongly Connected Components.
                            -----if (uf.find(edges[i].second.first) !=-----// d5
                            ------uf.find(edges[i].second.second)) {------// 8c
3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                            -----res.push_back(edges[i]);-----// d1
graph in O(|V| + |E|) time.
                            -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
#include "../data-structures/union_find.cpp"------// 5e
                            -----// 11
                            ----return res:------// 46
vector<br/>bool> visited;-----// 66
                            }------// 88
vi order;-----// 9b
-----// a5
                            3.6. Topological Sort.
void scc_dfs(const vvi &adj, int u) {-----// a1
----int v; visited[u] = true;------// e3
                            3.6.1. Modified Depth-First Search.
----for (int i = 0; i < size(adj[u]); i++)------// c5 void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-------// 6e ------bool& has_cycle) {----------------------// a8
----order.push_back(u);------// 19 ----color[cur] = 1;------// 5b
}------// dc ----for (int i = 0; i < size(adj[cur]); i++) {-------// 96
-----// 96 ------<mark>int</mark> nxt = adj[cur][i];---------------// 53
----int n = size(adj), u, v;----------// bd -------tsort_dfs(nxt, color, adj, res, has_cycle);------// 5b
----union_find uf(n);------// 6d ------has_cycle = true;------// c8
----vi dag;-------// ae ------if (has_cycle) return;----------------// 7e
-----rev[adj[i][j]].push_back(i);--------// 77 ----res.push(cur);-------// cb
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04 }------// 9e
```

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-----tsort_dfs(i, color, adj, S, has_cycle);-------// 40 -----int right = adj[left][i];-------------// b6
------if (has_cycle) return res;--------// 6c ------if (owner[right] == -1 || alternating_path(owner[right])) {-------// d2
-----}-----owner[right] = left; return 1;--------// 26
----return res;-----// 07
}-----// 1f
                                   3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
                                   ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                                   #define MAXN 5000-----// f7
#define MAXV 1000-----// 2f
                                   int dist[MAXN+1], q[MAXN+1];....// b8
#define MAXE 5000-----// 87
                                   #define dist(v) dist[v == -1 ? MAXN : v]------// 0f
vi adj[MAXV];-----// ff
                                   struct bipartite_graph {------// 2b
----int N, M, *L, *R; vi *adj;------// fc
ii start_end() {------// 30
                                   ----bipartite_graph(int N, int M) : N(N), M(M),------// e7
----int start = -1, end = -1, any = 0, c = 0;-----// 74
                                   -----L(new int[N]), R(new int[M]), adj(new vi[N]) {}------// 46
----for (int i = 0; i < n; i++) {------// 96
-----if (outdeg[i] > 0) any = i;-----// f2
                                   ----~bipartite_graph() { delete[] adj; delete[] L; delete[] R; }------// b9
                                   ----bool bfs() {------// 3e
-----if (indeq[i] + 1 == outdeq[i]) start = i, c++;------// 98
                                   -----int l = 0, r = 0;-----// a4
-----else if (indeg[i] == outdeg[i] + 1) end = i, c++;------// 4f
                                   -----for(int v = 0; v < N; ++v) if(L[v] == -1) dist(v) = 0, q[r++] = v;-----// 31
-----else if (indeg[i] != outdeg[i]) return ii(-1,-1);------//
                                   -----else dist(v) = INF;-----// c4
-----dist(-1) = INF;-----// f3
----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 6e
                                   -----while(l < r) {------// 3f
----if (start == -1) start = end = any;-----// db
----return ii(start, end);-----// 9e
                                   -----int v = q[l++];-----// 69
                                   ------if(dist(v) < dist(-1)) {------// b2
}-----// 35
                                   -----foreach(u, adj[v]) if(dist(R[*u]) == INF)------// 63
bool euler_path() {-----// d7
                                   -----dist(R[*u]) = dist(v) + 1, q[r++] = R[*u];-----// f8
---ii se = start_end();-----// 45
                                   ----int cur = se.first, at = m + 1;------// 8c
                                   ----if (cur == -1) return false:-----// 45
                                   -----return dist(-1) != INF;-----// e4
----stack<int> s;------// f6
                                   ----}------// b5
----while (true) {------// 04
                                   ----bool dfs(int v) {-------// f6
-----if (outdeg[cur] == 0) {------// 32
                                   -----if(v != -1) {------// 6c
-----res[--at] = cur;-----// a6
                                   -----foreach(u, adj[v])-----// 19
-----if (s.empty()) break;-----// ee
                                   -----if(dist(R[*u]) == dist(v) + 1)------// d9
-----cur = s.top(); s.pop();-----// d7
                                   -----if(dfs(R[*u])) {------// c7
-----} else s.push(cur), cur = adj[cur][--outdeq[cur]];------// d8
----}-----------// ba
                                   ------R[*u] = v, L[v] = *u;------// 2e
                                   -----/return true;------// 56
----return at == 0:-----// c8
                                   }-----// aa
                                   -----dist(v) = INF;-----// d4
                                   ------return false;-----// de
3.8. Bipartite Matching.
                                   3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
                                   -----return true:------// 7b
O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
                                   graph, respectively.
                                   ----void add_edge(int i, int j) { adj[i].push_back(j); }------// 87
vi* adj;------// cc ----int maximum_matching() {-----------------------------// ae
int* owner;-----memset(L, -1, sizeof(int) * N);--------------------------------// 16
```

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};-----// dd ----}
3.9. Maximum Flow.
                                          3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
3.9.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It computes
                                          O(|V||E|^2). It computes the maximum flow of a flow network.
the maximum flow of a flow network.
                                          #define MAXV 2000----// ba
#define MAXV 2000-----// ba
                                          int q[MAXV], d[MAXV], p[MAXV];-----// 7b
int q[MAXV], d[MAXV];.....// e6
                                          struct flow_network {------// 5e
struct flow_network {------// 12
                                          ----struct edge {------// fc
----struct edge {------// 1e
                                          ------int v, cap, nxt;-----// cb
-----int v, cap, nxt;-----// ab
                                          ------edge(int v, int cap, int nxt) : v(v), cap(cap), nxt(nxt) { }------// a1
-----edge() { }-----// 38
                                          ----}:-------// f9
------edge(int v, int cap, int nxt) : v(v), cap(cap), nxt(nxt) { }------// f7
                                          ----int n, ecnt, *head;------// 00
----}:-----// 7a
                                          ----vector<edge> e, e_store;-----// 5f
----int n, ecnt, *head, *curh;------// 77
                                          ----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 0d
----vector<edge> e, e_store;------// d0
                                          -----e.reserve(2 * (m == -1 ? n : m));-----// 0d
----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 80
                                          ------memset(head = new int[n], -1, n << 2);------// bc
-----e.reserve(2 * (m == -1 ? n : m));-----// 5d
                                          -----head = new int[n], curh = new int[n];-----// 6d
                                          ----void destroy() { delete[] head; }------// f1
------memset(head, -1, n * sizeof(int));------// f6
                                          ----void reset() { e = e_store; }------// 1c
----void add_edge(int u, int v, int uv, int vu=0) {------// ba
----void destroy() { delete[] head; delete[] curh; }------// 21
                                          -----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// 85
----void reset() { e = e_store; }------// 60
                                          -----e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// 9b
----void add_edge(int u, int v, int uv, int vu = \theta) {------// dd
                                          ----}------// ce
-----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// 7a
                                          ----int max_flow(int s, int t, bool res = true) {-------// b1
-----e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// b2
                                          -----if (s == t) return 0;-----// bb
----}------// 35
                                          -----e_store = e;------// f8
----int augment(int v, int t, int f) {------// a1
                                          -----int f = 0, l, r, v;-----// 62
-----if (v == t) return f;------// 84
                                          ------while (true) {-------// 3a
------for (int &i = curh[v], ret; i != -1; i = e[i].nxt)------// f4
                                          -----memset(d, -1, n << 2);-----// 73
-----if (e[i].cap > 0 && d[e[i].v] + 1 == d[v])------// 8f
                                          -----/memset(p, -1, n \ll 2);-----// \theta d
-----if ((ret = augment(e[i].v, t, min(f, e[i].cap))) > θ)-----// a3
                                          -----return (e[i].cap -= ret, e[i^1].cap += ret, ret);-----// ed
                                          ------while (l < r)-----// af
-----return 0:-----// d3
                                          -----for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)-----// 39
----}------// 48
                                          -----if (e[i].cap > 0 \& (d[v = e[i].v] == -1 \mid | d[u] + 1 < d[v]))
----int max_flow(int s, int t, bool res = true) {------// 87
                                          -----d[v] = d[u] + 1, p[q[r++] = v] = i;------// 14
-----if(s == t) return 0;-----// 2e
                                          -----if (p[t] == -1) break;-----// 84
-----e_store = e;------// 59
                                          -----int x = INF, at = p[t];-----// 43
-----int f = 0, x, l, r;-----// 82
                                          ------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 7\theta
------while (true) {------// 36
                                          -----at = p[t], f += x;------// 3c
------memset(d, -1, n * sizeof(int));------// 6c
-----l = r = 0, d[q[r++] = t] = 0; -----// 47
                                          ------while (at != -1)------// 58
                                          ------e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// 59
------}-------// b8
------for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)-----// cd
                                          -----if (res) reset();------// e7
-----if (e[i^1].cap > 0 && d[e[i].v] == -1)------// 23
                                          -----return f;------// 62
-----d[q[r++] = e[i].v] = d[v]+1;------// c2
                                          ----}------// 65
-----if (d[s] == -1) break;-----// a3
                                          };-----// 83
-----memcpy(curh, head, n * sizeof(int));-----// 0e
-----}-----// ea
                                          3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
```

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flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with x = INF, at p[t]; x = INF, at p[t]; y = INF, at y
minimum cost. Running time is O(|V|^2|E|\log|V|).
                                                     -----at = p[t], f += x;-----// 30
#define MAXV 2000------// ba ------while (at != -1)------// 84
int d[MAXV], pot[MAXV]; .....// 80
                                                    ------e[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// a5
----bool operator ()(int i, int j) {-------// 8a --------for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];-----// 66
-----return d[i] == d[j] ? i < j : d[i] < d[j];------// 89
----}------if (res) reset();------// da
};------return ii(f, c);------// f4
struct flow_network {------// eb ___}
----struct edge {------// 9a
                                                    }:-----// 80
------int v, cap, cost, nxt;-----// ad
-----: v(v), cap(cap), cost(cost), nxt(nxt) { }------// 5d
                                                    3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----}:-----// ae
----int n, ecnt, *head;-----// 57
                                                    structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the max-
                                                    imum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----vector<edge> e, e_store;------// d7
----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 83 #include "dinic.cpp"------// 58
-----e.reserve(2 * (m == -1 ? n : m));------// 2c
                                                    .....// 25
------push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;-------// 70 ------par[s].second = q.max_flow(s, par[s].first, false);-------// 38
----}-----memset(d, 0, n * sizeof(int));-------// 79
------same[v = q[l++]] = true;------// 32
------d[q[r++] = g.e[i].v] = 1;------// f8
-----set<int, cmp> q;-------// a8 ------for (int i = s + 1; i < n; i++)-------// 68
-----q.insert(s); d[s] = 0; d[s] = 0;
------f (q.find(v) != q.end()) q.erase(q.find(v));------// 47 -----mn = min(mn, par[cur].second), cur = par[cur].first;---------// 28
```

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------begin++, cur = it->second; } } } ;------// 26
----return min(cur, gh.second[at][t]);-----// 6d
                          4.2. Suffix Array. An O(n \log n) construction of a Suffix Tree.
}-----// a2
                          struct entry { ii nr; int p; };-----// f9
                          bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77
           4. Strings
                          struct suffix_array {------// 87
4.1. Trie. A Trie class.
                          ----string s; int n; vvi P; vector<entry> L; vi idx;-----// b6
template <class T>-----// 82
                          ----// REMINDER: Append a large character ('\x7F') to s------// 70
class trie {------// 9a ----suffix_array(string s) : s(s), n(size(s)) {------// fc
private: ------L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n); ------// 96
----struct node {------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';------// 69
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {------// bb
------<mark>int</mark> prefixes, words;-------// e2 -------P.push_back(vi(n));-------// e9
------for (int i = 0; i < n; i++)-------// 50
public:------L[L[i].p = i].nr = ii(P[stp - 1][i],-----// 0e
----node* root;------i + cnt < n ? P[stp - 1][i + cnt] : -1);------// 18
----trie() : root(new node()) { }-------// 8f ------sort(L.begin(), L.end());------// 29
----void insert(I begin, I end) {-------// 3c -------P[stp][L[i].p] = i > 0 &&------// 36
------while (true) {-------// 67 ----}
-----if (begin == end) { cur->words++; break; }------// db ---}
------else {-------// 3e ----int lcp(int x, int y) {-------// 10
-----T head = *beqin;------// fb ------int res = 0;------// 62
------typename map<T, node*>::const_iterator it;------// 01 -----if (x == y) return n - x;------// f4
-----it = cur->children.find(head);------// 77 ------for (int k = size(P) - 1; k >= 0 && x < n && y < n; k--)-----// 3e
------pair<T, node*> nw(head, new node());------// cd -----return res;-----
------} begin++, cur = it->second; } } }------// 64 }:-----// 64
----template<class I>-----// b9
-----else {-------// 1e -----string keyword; out_node *next;------// f0
-----T head = *begin;-------------------------// 5c ------out_node(string k, out_node *n) : keyword(k), next(n) { }-------// 26
-----begin++, cur = it->second; } } }------// 7c ------out_node *out; go_node *fail;-------// 3e
-----node* cur = root;------// 95 ---qo_node *qo;------// b8
-------while (true) {---------// 3e ----aho_corasick(vector<string> keywords) {-------// 4b
-----T head = *begin;-------// 43 ------go_node *cur = go;-------// 9d
-----typename map<T, node*>::const_iterator it;-------// 7a -----foreach(c, *k)-------
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-----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d ----for (int i = 1; i < n; i++) {---------// 7e
------foreach(a, qo->next) g.push(a->second);--------// a3 -----z[i] = r - l; r--;----------------// fc
-----st = st->fail;-----// 3f
-----if (!st) st = go;-----// e7
5. Mathematics
-----if (s->fail) {------// 3b
                               5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common
-----if (!s->out) s->out = s->fail->out;------// 80
-----else {------// ed
                               template <class T>-----// 82
-----out_node* out = s->out;-----// bf
                               class fraction {-----// cf
------while (out->next) out = out->next:-----// ca
                               private:----// 8e
-----out->next = s->fail->out;-----// 65
                               ----T qcd(T a, T b) { return b == T(0) ? a : qcd(b, a % b); }------// 86
public:-----// 0f
                               ----T n, d;------// 4b
----fraction(T n_, T d_) {------// 03
-----assert(d_ != 0);-----// 3d
----vector<string> search(string s) {------// 8d
                               -----n = n_, d = d_;------// 06
                               -----if (d < T(0)) n = -n, d = -d; -----// be
-----vector<string> res;------// ef
                               -----T g = gcd(abs(n), abs(d));-----// fc
-----qo_node *cur = qo;-----// 61
                               ------n /= g, d /= g; }------// a1
------foreach(c, s) {------// 6c
------while (cur && cur->next.find(*c) == cur->next.end())-----// 1f
                               ----fraction(T n_) : n(n_), d(1) { }-----// 84
                               ----fraction(const fraction<T>\& other) : n(other.n), d(other.d) { }------// \theta 1
-----cur = cur->fail:-----// 9e
                               ----fraction<T> operator +(const fraction<T>& other) const {------// b6
-----if (!cur) cur = go;-----// 2f
                               -----return fraction<T>(n * other.d + other.n * d, d * other.d);}------// 3b
-----cur = cur->next[*c];-----// 58
-----if (!cur) cur = qo;-----// 3f
                               ----fraction<T> operator -(const fraction<T>& other) const {------// 26
                               -----return fraction<T>(n * other.d - other.n * d, d * other.d);}------// 47
------for (out_node *out = cur->out; out = out->next)------// e0
-----res.push_back(out->keyword);------// 0d
                               ----fraction<T> operator *(const fraction<T>& other) const {------// 38
                               -----return fraction<T>(n * other.n, d * other.d); }-----// c5
                               ----fraction<T> operator /(const fraction<T>& other) const {------// ca
-----return res;-----// c1
                               -----return fraction<T>(n * other.d, d * other.n); }------// 35
----}-----// e4
                               }:-----// 32
                               -----return n * other.d < other.n * d; }------// 8c
4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
                               ----bool operator <=(const fraction<T>& other) const {------// 48
also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
                               -----return !(other < *this); }------// 86
can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
                               ----bool operator >(const fraction<T>& other) const {------// c9
accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
                               -----return other < *this; }-----// 6e
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                                                      14
};------// 12 ----string to_string() const { stringstream ss; ss << *this; return ss.str(); }// fc
                            ----bool operator <(const intx& b) const {------// 21
5.2. Big Integer. A big integer class.
                            ------if (sign != b.sign) return sign < b.sign;------// cf
----intx() { normalize(1); }------// 6c -----return sign == 1 ? size() < b.size() > b.size();-----// 4d
----intx(string n) { init(n); }------// b9 ------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])-----// 35
----intx(int n) { stringstream ss; ss << n; init(ss.str()); }-------// 36 -----return sign == 1 ? data[i] < b.data[i] > b.data[i];--// 27
----intx(const intx& other) : sign(other.sign), data(other.data) { }------// 3b -----return false;-----
----int sign;------------------// 26 ....}
----static const unsigned int radix = 10000000000U;------// f0 ---intx operator +(const intx& b) const {-----// f8
----void init(string n) {-------// 13 ------if (sign < 0 && b.sign > 0) return b - (-*this);------// 70
-----intx res; res.data.clear();------// 4e ------if (sign < 0 && b.sign < 0) return -((-*this) + (-b));------// 59
-----if (n.empty()) n = "0";-------// 99 -----intx c; c.data.clear();------// 18
------if (n[0] == '-') res.sign = -1, n = n.substr(1);------------// 3b ------unsigned long long carry = 0;------------------------// 5c
------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {--------// e7 -------for (int i = 0; i < size() || i < b.size() || carry; i++) {-------// e3
------(i < b.size() ? b.data[i] : OULL);-------// 0c
------c.data.push_back(carry % intx::radix);------// 86
------if (idx < 0) continue;------// 52 -----carry /= intx::radix;-----// fd
-----res.data.push_back(digit);------// 07 ____}
------if (sign < 0 && b.sign > 0) return -(-*this + b);------// 1b
-----if (data.empty()) data.push_back(0);------// fa -----intx c; c.data.clear();------// 6b
-----data.erase(data.begin() + i);------// 67 ------for (int i = 0; i < size(); i++) {------// a7
-----sign = data.size() == 1 && data[0] == 0 ? 1 : nsign;-----// ff ------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
----friend ostream& operator <<(ostream& outs, const intx& n) {------// 0d _____}
------bool first = true;------// 33 ---}
------for (int i = n.size() - 1; i >= 0; i--) {--------// 63 ----intx operator *(const intx& b) const {-------// bd
-----if (first) outs << n.data[i], first = false;----------// 33 ------intx c; c.data.assign(size() + b.size() + 1, 0);------// d0
-----stringstream ss; ss << cur; ------// 8c ------for (int j = 0; j < b.size() || carry; j++) {-------// c0
-----string s = ss.str();------// 64 ------if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
------outs << s;------// 97 ------carry /= intx::radix;-----// 05
-----return outs:-----// cf
```

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                                                            15
------assert(!(d.size() == 1 && d.data[0] == 0));-------// e^9 -------data[i] %= radix;-------------------------// bd
-----long long k = 0;------// cc ---ss << data[stop];------// 96
------if (d.size() < r.size())--------// b9 ----for (int i = stop - 1; i >= 0; i--)-------// bd
------k = (long long)intx::radix * r.data[d.size()];------// f7 -----ss << setfill('0') << setw(len) << data[i];-------// b6
------k /= d.data.back();--------// b7 ----delete[] a; delete[] b;---------// 7e
5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
-----return pair<intx, intx>(q.normalize(n.sign * d.sign), r);------// a1
                              k items out of a total of n items.
----}------// 1b
                              int nck(int n, int k) {------// f6
----intx operator /(const intx& d) const {-------// a2
                               ----if (n - k < k) k = n - k;------// 18
-----return divmod(*this,d).first; }-----// 1e
                              ----int res = 1;-----// cb
----intx operator %(const intx& d) const {------// 07
-----return divmod(*this,d).second * sign; }-----// 5a
                               ----for (int i = 1; i \le k; i++) res = res * (n - (k - i)) / i;------// bd
                              ----return res:------// e4
}:-----// 38
                              }------// 03
5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
                              5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
#include "intx.cpp"-----// 83
#include "fft.cpp"-----// 13
-----// e0
                              int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
intx fastmul(const intx &an. const intx &bn) {------// ab
                               The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
----string as = an.to_string(), bs = bn.to_string();------// 32
                              and also finds two integers x, y such that a \times x + b \times y = d.
----int n = size(as), m = size(bs), l = 1,-----// dc
                              int eqcd(int a, int b, int& x, int& y) {------// 85
-----len = 5, radix = 100000,-----// 4f ----if (b == 0) { x = 1; y = 0; return a; }-----// 7b
-----*b = new int[m], blen = 0;------// 0a -----// 0a -----// 34
----memset(a, 0, n << 2);-----// 1d -----x -= a / b * y;-----// 4a
----memset(b, 0, m << 2);-----// 01 -----swap(x, y);-----// 26
----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e -----return d:-----
------for (int j = min(len - 1, i); j >= 0; j--)------// 43 ---}
----for (int i = m - 1; i >= 0; i -= len, blen++)-----// b6
                              5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is
------for (int j = min(len - 1, i); j >= 0; j--)------// ae
------b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);------// 35 ----if (n < 4) return true;-------// d9
----ull *data = new ull[l];------// e7 ------if (n % i == 0 || n % (i + 2) == 0) return false;------// e9
----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06 ----return true; }-----
```

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5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
                                                     ----return res:------// 03
                                                    }-----// 1c
vi prime_sieve(int n) {-----// 40
----int mx = (n - 3) >> 1, sq, v, i = -1;------// 27
                                                     5.11. Numeric Integration. Numeric integration using Simpson's rule.
----vi primes:-----// 8f
                                                     double integrate(double (*f)(double), double a, double b,-----// 76
----bool* prime = new bool[mx + 1];-----// ef
                                                     ------double delta = 1e-6) {------// c0
----memset(prime, 1, mx + 1);------// 28
                                                     ----if (abs(a - b) < delta)------// 38
----if (n >= 2) primes.push_back(2);-----// f4
                                                     -----return (b-a)/8 *-----// 56
----while (++i <= mx) if (prime[i]) {-----// 73
                                                     -----(f(a) + 3*f((2*a+b)/3) + 3*f((a+2*b)/3) + f(b));-----// e1
-----primes.push_back(v = (i << 1) + 3);-----// be
                                                     ----return integrate(f, a,-----// 64
-----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
                                                     -----(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);------// \theta c
------for (int j = sq; j <= mx; j += v) prime[j] = false; }-----// 2e
                                                       ·----// 4b
----while (++i \le mx) if (prime[i]) primes.push_back((i << 1) + 3);-----// 29
----delete[] prime; // can be used for O(1) lookup-----// 36
                                                    5.12. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete
----return primes; }-----// 72
                                                    Fourier transform. Note that this implementation only handles powers of two, make sure to pad with
5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
                                                     #include <complex>-----// 8e
#include "egcd.cpp"-----// 55
                                                     typedef complex<long double> cpx;-----// 25
-----// e8
                                                     void fft(cpx *x, int n, bool inv=false) {------// 23
int mod_inv(int a, int m) {------// 49
                                                     ----for (int i = 0, j = 0; i < n; i++) {-------// f2
----int x, y, d = eqcd(a, m, x, y);------// 3e
                                                     ------if (i < j) swap(x[i], x[j]);------// 5c
----if (d != 1) return -1;------// 20
                                                     ------int m = n>>1;------// e5
----return x < 0 ? x + m : x;------// 3c
                                                     ------while (1 <= m && m <= j) j -= m, m >>= 1;-----// fe
                                                     -----j += m:-----// ab
                                                     ----}-----// 1e
5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
                                                     ----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
template <class T>-----// 82
                                                     T mod_pow(T b, T e, T m) {-----// aa
                                                     ------for (int m = 0; m < mx; m++, w *= wp) {------// 40
----T res = T(1);-----// 85
                                                     ------for (int i = m; i < n; i += mx << 1) {------// 33
----while (e) {------// b7
                                                     -----cpx t = x[i + mx] * w;-----// f5
-----if (e & T(1)) res = mod(res * b, m);------// 41
                                                     -----x[i + mx] = x[i] - t;-----// ac
-----b = mod(b * b, m), e >>= T(1); }------// b3
                                                     -----x[i] += t:-----// c7
----return res:-----// eb
                                                     }-----// c5
                                                     ----}-----------// 70
5.9. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
                                                     ----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);------// 3e
#include "eacd.cpp"-----// 55
                                                     }-----// 7d
int crt(const vi& as, const vi& ns) {------// c3
----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
                                                    5.13. Formulas.
----for (int i = 0: i < cnt: i++) N *= ns[i]: ------// 88
                                                        • Number of ways to choose k objects from a total of n objects where order matters and each
----for (int i = 0; i < cnt; i++)-----// f9
                                                         item can only be chosen once: P_k^n = \frac{n!}{(n-k)!}
• Number of ways to choose k objects from a total of n objects where order matters and each
----return mod(x, N); }-----// 9e
                                                         item can be chosen multiple times: n^k
                                                        • Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type
5.10. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
                                                         2, ..., n_k objects of type k: \binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}
                                                        • Number of ways to choose k objects from a total of n objects where order does not matter
#include "egcd.cpp"-----// 55
                                                         and each item can only be chosen once:
vi linear_congruence(int a, int b, int n) {-----// c8
                                                         \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^k \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0
----int x, y, d = egcd(a, n, x, y);-----// 7a
----vi res;------// f5
                                                        • Number of ways to choose k objects from a total of n objects where order does not matter
                                                         and each item can be chosen multiple times: f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}
----if (b % d != 0) return res:------// 30
----int x\theta = mod(b / d * x, n); ------// 48
                                                        • Number of integer solutions to x_1 + x_2 + \cdots + x_n = k where x_i > 0: f_k^n
```

• Number of subsets of a set with n elements: 2^n

----for (int k = 0; k < d; k++) res.push_back(mod(x0 + k * n / d, n));-----// 21

• $|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$

- Number of ways to walk from the lower-left corner to the upper-right corner of an n × m grid by walking only up and to the right: (n+m)/m
 Number of strings with n sets of brackets such that the brackets are balanced:
- Number of strings with n sets of brackets such that the brackets are balanced:
 C_n = ∑_{k=0}ⁿ⁻¹ C_kC_{n-1-k} = 1/(n+1) (2n)
 Number of triangulations of a convex polygon with n points, number of rooted binary trees
- with n+1 vertices, number of paths across an n×n lattice which do not rise above the main diagonal: C_n
 Number of permutations of n objects with exactly k ascending sequences or runs:
- Value of permutations of n objects with exactly k ascending sequences of n and n and n and n and n are n and n and n are n and n and n are n are n are n and n are n are n and n are n and n are n are n and n are n are n and n are n are n are n and n are n and n are n are
- Number of permutations of n objects with exactly k cycles: $\binom{n}{k} = \binom{n-1}{k-1} + (n-1)\binom{n-1}{k}$
- Number of ways to partition n objects into k sets: $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$ • Number of permutations of length n that have no fixed points (derangements): $D_0 = 1, D_1 = 1$
- Number of permutations of length n that have exactly k fixed points: $\binom{n}{k}D_{n-k}$
- Heron's formula: A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where
- Fields for initial. A triangle with side lengths a, b, c has area \(\sqrt{s}(s-b)(s-b)(s-b) \) where \(s = \frac{a+b+c}{2} \).
 Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice.
- points on the boundary has area $i + \frac{b}{2} 1$.
- **Divisor sigma:** The sum of divisors of n to the xth power is $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$ where $n = \prod_{i=0}^r p_i^{a_i}$ is the prime factorization.
- Divisor count: A special case of the above is $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$.
- Euler's totient: The number of integers less than n that are comprime to n are n ∏_{p|n} (1 ½) where each p is a distinct prime factor of n.
 König's theorem: In any bipartite graph, the number of edges in a maximum matching is
- equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.

6. Geometry

6.1. **Primitives.** Geometry primitives.

```
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6
double angle(P(a), P(b), P(c)) {-----// d0
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------// cc
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 9e
double progress(P(p), L(a, b)) {------// d2
----if (abs(real(a) - real(b)) < EPS)------// 9e
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));------// 35
----else return (real(p) - real(a)) / (real(b) - real(a)); }------// 2c
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) \{-----//d6\}
----// NOTE: check for parallel/collinear lines before calling this function---// 02
----point r = b - a, s = q - p;------// 79
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// a8
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// ae
-----return false;-----// a3
----res = a + t * r:-----// ca
point closest_point(L(a, b), P(c), bool segment = false) {------// a1
----if (segment) {-------// c2
-----if (dot(b - a, c - b) > 0) return b;------// b5
-----if (dot(a - b, c - a) > 0) return a:-----// cf
----double t = dot(c - a, b - a) / norm(b - a);-----// aa
----return a + t * (b - a);-----// 7a
}-----// e5
double line_segment_distance(L(a,b), L(c,d)) {------// 99
----double x = INFINITY:-----// 83
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// df
----else if (abs(a - b) < EPS) \times = abs(a - closest_point(c, d, a, true)); -----// da
----else if (abs(c - d) < EPS) x = abs(c - closest_point(a, b, c, true));-----// 52
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// ee
----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;-----// 79
----else {------// 38
-----x = min(x, abs(a - closest_point(c,d, a, true)));------// f3
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// ec
-\cdots = min(x, abs(c - closest_point(a,b, c, true)));-----// 36
-----x = min(x, abs(d - closest\_point(a,b, d, true)));
....}-------// 72
----return x:-----// 0d
1.....// b3
6.2. Polygon. Polygon primitives.
#include "primitives.cpp"-----// e0
```

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#define CHK(f,a,b,c) (f(a) < f(b) && f(b) <= f(c) && ccw(a,c,b) < 0)------// b2 #include "primitives.cpp"-----------// e0
int point_in_polygon(polygon p, point g) {-------// 58 bool line_segment_intersect(L(a, b), L(c, d), point &A, point &B) {------// 6c
----for (int i = 0, j = n - 1; i < n; j = i++)------// 77 ------A = B = a; return abs(a - d) < EPS; }------// ee
------if (collinear(p[i], q, p[j]) &&---------// a5 ----else if (abs(a - b) < EPS) {--------------------------// 03
-----return 0;-----// cc -----return 0.0 <= p && p <= 1.0-----// 8a
------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))-------// 1f ----else if (abs(c - d) < EPS) {--------// 26
-----in = !in;------// b2 ------A = B = c; double p = progress(c, a,b);------// d9
----return in ? -1 : 1; }-------------------------// 77 -------return 0.0 <= p && p <= 1.0-------------------// 8e
// pair<polygon, polygon cut_polygon(const polygon &poly, point a, point b) {-// 7b -------& (abs(c - a) + abs(b - c) - abs(b - a)) < EPS; }------// 4f
//--- polygon left, right;-----// 6b ----else if (collinear(a,b, c,d)) {------------// bc
//---- point it(-100, -100);------// c9 -------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
//--- for (int i = 0, cnt = poly.size(); i < cnt; i++) {------// 28 -------if (ap > bp) swap(ap, bp);------// b1
//------ int j = i == cnt-1 ? 0 : i + 1;-------// 8e ---------if (bp < 0.0 || ap > 1.0) return false;-------// 0c
//------ point p = poly[i], q = poly[i];------// 19 ------A = c + max(ap, 0.0) * (d - c);------// f6
//------ if (ccw(a, b, p) \le 0) left.push_back(p);------// 12 ------B = c + min(bp, 1.0) * (d - c);-------// 5c
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// e3 ------return true; }-----
//-----// myintersect = intersect where-----// 24 ----else if (parallel(a,b, c,d)) return false;-------// ca
//-----// (a,b) is a line, (p,q) is a line seament------// f2 ----else if (intersect(a,b, c,d, A, true)) {----------------------// 10
//----- if (myintersect(a, b, p, q, it))------// f0 ------B = A; return true; }------------------// bf
//---------left.push_back(it), right.push_back(it);-------// 21 ----return false:----------------------------// b7
//---- return pair<polygon, polygon>(left, right);------// 1d -----// e6
// }-----// 37
                                                  6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
                                                  coordinates) on a sphere of radius r.
6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
                                                  double gc_distance(double pLat, double pLong,-----// 7b
#include "polygon.cpp"-----// 58
                                                   ----- double qLat, double qLong, double r) {------// a4
#define MAXN 1000-----// 09
                                                   ----pLat *= pi / 180; pLong *= pi / 180;-----// ee
point hull[MAXN];-----// 43
                                                  ----qLat *= pi / 180; qLong *= pi / 180;-----// 75
bool cmp(const point &a, const point &b) {-----// 32
                                                   ----return r * acos(cos(pLat) * cos(pLong - qLong) +------// e3
----return abs(real(a) - real(b)) > EPS ?-----// 44
                                                     -----/:sin(pLat) * sin(qLat));-------------------// 1e
-----real(a) < real(b) : imag(a) < imag(b); }------// 40
                                                             int convex_hull(polygon p) {------// cd
                                                  }-----// 3f
----int n = size(p), l = 0;-----// 67
----sort(p.begin(), p.end(), cmp);-----// 3d
                                                  6.6. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
----for (int i = 0; i < n; i++) {-------// 6f
                                                     • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
------if (i > 0 && p[i] == p[i - 1]) continue;------// b2
                                                     • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
-------while (l \ge 2 \&\& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;------// 20
                                                     • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
-----hull[l++] = p[i];-----// f7
                                                      of that is the area of the triangle formed by a and b.
----}------------// d8
----int r = 1:-----// 59
                                                                     7. Other Algorithms
----for (int i = n - 2; i >= 0; i--) {------// 16
                                                  7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
-----if (p[i] == p[i + 1]) continue;-----// c7
                                                  function f on the interval [a, b], with a maximum error of \varepsilon.
-----while (r - l >= 1 \&\& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--:---// 9f
                                                  double binary_search_continuous(double low, double high,-----// 8e
-----hull[r++] = p[i];-----// 6d
                                                   -----double eps, double (*f)(double)) {-----// c0
----while (true) {------// 3a
----return l == 1 ? 1 : r - 1;------// 6d
                                                  ------double mid = (low + high) / 2, cur = f(mid);-----// 75
}-----// 79
                                                   -----if (abs(cur) < eps) return mid;------// 76
                                                  -----else if (0 < cur) high = mid;------// e5
6.4. Line Segment Intersection. Computes the intersection between two line segments.
```

```
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------else low = mid;-------// a7 ----return true;-------// eb
7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
 Another implementation that takes a binary predicate f, and finds an integer value x on the integer
                                     vi stable_marriage(int n, int** m, int** w) {------// e4
interval [a,b] such that f(x) \wedge \neg f(x-1).
                                      ----queue<int> q;-----// f6
int binary_search_discrete(int low, int high, bool (*f)(int)) {-------// 51
                                      ---vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
----assert(low <= high);-----// 19
                                      ----for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)------// 05
----while (low < high) {------// a3
                                      -----inv[i][w[i][j]] = j;-----// b9
------int mid = low + (high - low) / 2;------// 04
                                      ----for (int i = 0; i < n; i++) q.push(i);-----// fe
-----if (f(mid)) high = mid;-----// ca
                                      ----while (!q.empty()) {------// 55
-----else low = mid + 1;-----// 03
                                      -----int curm = q.front(); q.pop();-----// ab
                                      ------for (int &i = at[curm]; i < n; i++) {-------// 9a
----assert(f(low));------// 42
                                      -----int curw = m[curm][i];-----// cf
                                      ------if (eng[curw] == -1) { }-----// 35
}-----// d3
                                      ------else if (inv[curw][curm] < inv[curw][enq[curw]])------// 10
                                     -----a.push(eng[curw]):-----// 8c
7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotonic
                                      -----else continue;-----// b4
cally decreasing, ternary search finds the x such that f(x) is maximized.
                                     -----res[eng[curw] = curm] = curw, ++i; break;-----// 5e
template <class F>-----// d1
                                      double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
                                      ----}-----------// b8
----while (hi - lo > eps) {------// 3e
                                      ----return res;------// 95
------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
-----if (f(m1) < f(m2)) lo = m1;------// 1d
-----else hi = m2:-----// b3
                                     7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the
----}------// bb
                                     Exact Cover problem.
----return hi:-----// fa
                                     bool handle_solution(vi rows) { return false; }------// 63
}-----// 66
                                     struct exact_cover {------// 95
7.3. 2SAT. A fast 2SAT solver.
                                      ----struct node {------// 7e
bool two_sat(int n, const vii& clauses, vi& all_truthy) {-------// f4 -----node(int row, int col) : row(row), col(col) {-------// 68
----vvi adj(2*n+1);-------// 7b ---};------// 9e
------if (clauses[i].first != clauses[i].second)--------// 87 ----node *head;------------------------------// c2
-----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93 ----exact_cover(int rows, int cols) : rows(rows), cols(cols), head(NULL) {-----// ce
----union_find scc = res.first;--------// 42 ------for (int i = 0; i < rows; i++)-------// c7
----vi dag = res.second;------// 58 -----arr[i] = new bool[cols], memset(arr[i], 0, cols);-----// 68
----vi truth(2*n+1, -1);------// 8b
-----if (cur == 0) continue;-------// 26 -----node ***ptr = new node**[rows + 1];------// da
------if (p == 0) return false; --------// 33 ------for (int i = 0; i <= rows; i++) {--------// ce
------if (truth[p] == -1) truth[p] = 1;--------// c3 ------ptr[i] = new node*[cols];-------// cc
------truth[cur + n] = truth[p];--------// b3 --------for (int j = 0; j < cols; j++)-------// 56
-----truth[o] = 1 - truth[p];-------// 80 -------if (i == rows || arr[i][j]) ptr[i][j] = new node(i, j);-----// 95
------if (truth[p] == 1) all_truthy.push_back(cur);---------// 5c ------else ptr[i][j] = NULL;----------------// 40
```

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-----if (!ptr[i][j]) continue;------// 76 ------for (; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 1e
------if (ni == rows || arr[ni][j]) break;-------// 77 ------for (node *r = c->d; !found && r != c; r = r->d) {-------// 1e
------for (node *j = r->l; j != r; j = j->l) { UNCOVER(j->p, a, b); }----// 8a
------if (nj == cols) nj = 0;------// a7 ------UNCOVER(c, i, j);------// 64
-----ptr[i][j]->r = ptr[i][nj];-----// b3
                         7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \dots, k-1\}
-----ptr[i][nj]->l = ptr[i][j];-----// 46
------head->r = ptr[rows][0];--------// b9 ----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
------head->l = ptr[rows][cols - 1];-------// 28 ----for (int i = cnt - 1; i >= 0; i--)-----// 52
-----int cnt = -1;------// 36 }------// 97
-----for (int i = 0; i <= rows; i++)-----// 56
                         7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
-----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];------// 05
                         ii find_cycle(int x0, int (*f)(int)) {------// a5
-----ptr[rows][j]->size = cnt;-----// d4
                         -----}-----// 8f
------for (int i = 0; i <= rows; i++) delete[] ptr[i];-------// cd ----while (t != h) t = f(t), h = f(f(h));--------// 79
                         ----h = x0;
-----delete[] ptr;-----// 42
                         ----}-----// a9
----while (t != h) h = f(h), lam++;-----// 5e
----return ii(mu, lam);-----// b4
------for (node *i = c->d; i != c; i = i->d) \[\bar{\capacital}\]
-----j->d->u = j->u, j->u->d = j->d, j->p->size--;-----// 5a
                         7.8. Dates. Functions to simplify date calculations.
----#define UNCOVER(c, i, j) \sqrt{\phantom{a}}-----// 17
                         int intToDay(int jd) { return jd % 7; }-----// 89
------for (node *i = c->u; i != c; i = i->u) \[ \bigcup_{------//98}
                         int dateToInt(int y, int m, int d) {------// 96
                         ----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1
-----j->p->size++, j->d->u = j->u->d = j; \\\-------// be
                         -----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
----bool search(int k = 0) {------// 4f
                         -----d - 32075;------// e0
                         }-----// fa
------if (head == head->r) {------// a7
                         -----vi res(k);-----// 4f
                         ----int x, n, i, j;------// 00
------for (int i = 0; i < k; i++) res[i] = sol[i];------// c0
                         ----x = jd + 68569;-----// 11
-----sort(res.begin(), res.end());-----// 3e
                         ----n = 4 * x / 146097;-----// 2f
-----return handle_solution(res);-----// dc
                          ---x = (146097 * n + 3) / 4;
```

i = (4000 * (x + 1)) / 1461001;//	0d
x -= 1461 * i / 4 - 31;//	09
j = 80 * x / 2447;//	3d
d = x - 2447 * j / 80;//	
x = j / 11;//	b7
m = j + 2 - 12 * x;//	82
y = 100 * (n - 49) + i + x;	70
}//	af

8. Useful Information

8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo m_1, m_2, \ldots, m_k , where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When $n = 0, n = -1, n = 1, n = 2^{31} 1$ or $n = -2^{31}$? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

8.3. **128-bit Integer.** GCC has a 128-bit integer data type named __int128. Useful if doing multiplication of 64-bit integers, or something needing a little more than 64-bits to represent.

8.4. Worst Time Complexity.

	n	Worst AC Algorithm	Comment
_	≤ 10	$O(n!), O(n^6)$	e.g. Enumerating a permutation
	≤ 15	$O(2^n \times n^2)$	e.g. DP TSP
	≤ 20	$O(2^n), O(n^5)$	e.g. DP + bitmask technique
	≤ 50	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
	$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
	$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
	$\le 10^{5}$	$O(n \log_2 n)$	e.g. Merge sort, building a Segment tree
	$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.