1.2. C++ Header. A C++ header. #include <bits/stdc++.h>-----// 84 using namespace std;-----// 16 typedef long long ll;-------// 47 -----if (l == r) return data[i] = arr[l];------// 5b const int INF = 2147483647;-----// db -----// db ------int m = (l + r) / 2;------// de -----// d8 ------return data[i] = f(mk(arr, l, m, 2*i+1), mk(arr, m+1, r, 2*i+2)); }----// 0a const double EPS = 1e-9;------// 18 ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// f6 const double pi = acos(-1);-----// ec ----int q(int a, int b, int l, int r, int i) {-------// 22 template <class T> T mod(T a, T b) { return (a % b + b) % b; }-----// 3f -----// 3f m = (l + r) / 2;------// 7a 1.3. **Java Template.** A Java template. import java.util.*;-----// 37 import java.math.*;-----// 89 import iava.io.*:-----// 28 _____// a3 public class Main {-----// 17 ----public static void main(String[] args) throws Exception {-------// 02 ------Scanner in = new Scanner(System.in);------// ef

-----PrintWriter out = new PrintWriter(System.out, false):-----// 62

-----// code-----// e6

2. Data Structures

-----propagate(l, r, i);-----// 19 -----out.flush();-----// 56 -----if (l > r) return ID;------// cc -----if (r < a || b < l) return data[i];-----// d9 }-----// 00 -----if (l == r) return data[i] += v;------// 5f -----if (a <= l & r <= b) return (lazy[i] = v) * (r - l + 1) + data[i];----// 76 ------int m = (l + r) / 2;-----// e7 2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.

-----return f(q(a, b, l, m, 2*i+1), q(a, b, m+1, r, 2*i+2)); }-----// 5c

----**void** update(**int** i, **int** v) { u(i, v, 0, n-1, 0); }------// 90

----int u(int i, int v, int l, int r, int j) {-------// 02

-----propagate(l, r, j);-----// ae

-----if (r < i || i < l) return data[j];------// 92

-----if (l == i && r == i) return data[j] = v;------// 4a

------int m = (l + r) / 2;-----// cb

-----return data[j] = $f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----//34$

----**void** range_update(**int** a, **int** b, **int** v) { ru(a, b, v, 0, n-1, 0); }------// 71

----int ru(int a, int b, int v, int l, int r, int i) {------------// e0

```
Reykjavík University
------if (lazy[2*i+1] == INF) lazy[2*i+1] = lazy[i];------// ee ----void operator *= (T other) {--------// ba
------else lazy[2*i+1] += lazy[i];-------// 72 ------for (int i = 0; i < cnt; i++) data[i] *= other; }------// 40
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];------// dd ----matrix<T> operator +(const matrix& other) {-------// ee
------else lazy[2*i+2] += lazy[i];-------// a4 -----matrix<T> res(*this); res += other; return res; }-----// 5d
-----lazy[i] = INF;------res(*this); res -= other; return res; }------// cf
}:-----// 17 -----matrix<T> res(*this); res *= other; return res; }------// 37
                                               ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                                               -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                                               -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i.. j in O(\log n) time. It only needs O(n) space.
                                               -----for (int k = 0; k < cols; k++)-----// fc
struct fenwick_tree {------// 98
                                               -----res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n: vi data:-----// d3
                                               -----return res; }-----// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-----// db
                                               ----matrix<T> transpose() {------// dd
----void update(int at, int by) {------// 76
                                              ------matrix<T> res(cols, rows);------// b5
-------while (at < n) data[at] += by, at |= at + 1; }------// fb
                                               ------for (int i = 0; i < rows; i++)------// 9c
----int query(int at) {------// 71
                                               ------for (int j = 0; j < cols; j++) res(j, i) = at(i, j);-----// a3
-----int res = 0:-----// c3
                                               -----return res; }-----// c3
------while (at >= 0) res += data[at], at = (at & (at + 1)) - 1;------// 37
                                               ----matrix<T> pow(int p) {------// 68
-----return res: }-----// e4
                                               ------matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }------// be
                                               ------for (int i = 0; i < rows; i++) res(i, i) = T(1);------// bf
}:-----// 57
                                              ------while (p) {------// cb
struct fenwick_tree_sq {-----// d4
                                              ------if (p & 1) res = res * sq;------// c1
----int n; fenwick_tree x1, x0;------// 18
                                              -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e
                                              ------if (p) sq = sq * sq;------// 9c
-----x0(fenwick_tree(n)) { }------// 7c
                                              ------} return res; }------// 50
----// insert f(y) = my + c if x <= y------// 17
                                               ----matrix<T> rref(T &det) {------// 89
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }----// 45
                                               -----matrix<T> mat(*this): det = T(1):-----// 21
----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73
                                               ------for (int r = 0, c = 0; c < cols; c++) {-------// c4
};-----// 13
                                               -----int k = r:-----// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {------// 89
                                               -----/while (k < rows \&\& eq<T>(mat(k, c), T(0))) k++;-----//f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f
                                               -----if (k >= rows) continue;-----// 3f
int range_query(fenwick_tree_sq &s, int a, int b) {------// 15
                                               -----if (k != r) {------// a3
----return s.query(b) - s.query(a-1); }------// f3
                                               -----det *= T(-1);-----// 7a
                                               -----for (int i = 0; i < cols; i++)-----// ab
2.4. Matrix. A Matrix class.
                                               -----swap(mat.at(k, i), mat.at(r, i));------// 8d
template <class K> bool eq(K a, K b) { return a == b; }-----// 2a
                                               template <> bool eq<double>(double a, double b) { return abs(a - b) < EPS; }---// a7
                                               -----T d = mat(r,c);------// 34
template <class T>-----// 53
                                               class matrix {-----// 85
                                               ------for (int i = 0; i < rows; i++) {-------// 17
public:----// be
                                               -----T m = mat(i, c):-----// 29
----int rows. cols:-----// d3
                                               -----if (i != r && !eq<T>(m, T(0)))-----// 2b
----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {------// 34
                                               -----for (int j = 0; j < cols; j++) mat(i, j) -= m * mat(r, j);-// \theta 3
-----data.assign(cnt, T(0)); }-----// d0
                                               ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),-----// fe
                                               ------} return mat; }------// 47
-----cnt(other.cnt), data(other.data) { }-----// ed
                                               private:-----// ac
----T& operator()(int i, int j) { return at(i, j); }-----// e0
                                               ----int cnt;------// 69
----void operator +=(const matrix& other) {------// c9
                                               ----vector<T> data:-----// e1
------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
                                               ----inline T& at(int i, int j) { return data[i * cols + j]; }------// 31
----void operator -=(const matrix& other) {------// 68
                                               }:-----// cc
------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
```

```
Reykjavík University
                                         ------if (!n) return NULL;-----// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                                         -----if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0-----// b5
                                          -----node *p = n->p;-----// a7
  -----// 61
                                         template <class T>-----// 22
                                         -----return p; }------// c7
class avl_tree {------// ff
                                         ----node* predecessor(node *n) const {-------// b4
public:----// f6
                                         -----if (!n) return NULL;-----// dd
----struct node {------// 45
                                         -----if (n->l) return nth(n->l->size-1, n->l):------// 10
-----T item; node *p, *l, *r;------// a6
                                         -----node *p = n->p;-----// ea
------int size, height;------// 33
                                         ------while (p && p->l == n) n = p, p = p->p;------// 6d
------node(const T &_item, node *_p = NULL) : item(_item), p(_p),------// 4f
                                         -----return p; }-----// e7
------l(NULL), r(NULL), size(1), height(0) { } };------// 0d
                                         ----inline int size() const { return sz(root); }------// ef
----avl_tree() : root(NULL) { }------// 5d
                                         ----void clear() { delete_tree(root), root = NULL; }-----// 84
----node *root:-----// 91
                                         ----node* nth(int n, node *cur = NULL) const {------// e4
----node* find(const T &item) const {------// 65
                                         -----if (!cur) cur = root:-----// e5
-----node *cur = root:-----// b4
                                         ------while (cur) {------// 29
------while (cur) {-------// 8b
                                         -----if (n < sz(cur->l)) cur = cur->l;------// 75
-----if (cur->item < item) cur = cur->r;------// 71
                                         -------else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;------// cd
------else if (item < cur->item) cur = cur->l;------// cd
                                          ------else break;-----// c0
-----else break; }-----// 4f
                                         -----} return cur; }------// ed
-----return cur; }-----// 84
                                         ----int count_less(node *cur) {-------// ec
----node* insert(const T &item) {------// 4e
                                         ------int sum = sz(cur->l);-----// bf
-----node *prev = NULL, **cur = &root;-----// 60
                                         ------while (cur) {-------// 6f
------while (*cur) {------// aa
                                         -----prev = *cur;-----// f0
                                         -----cur = cur->p;-----// eb
------if ((*cur)->item < item) cur = &((*cur)->r);------// 1b
                                         -----} return sum; }------// a0
#if AVL_MULTISET-----// 0a
                                         private:----// d5
------else cur = &((*cur)->l);-----// eb
                                         ----inline int sz(node *n) const { return n ? n->size : 0; }------// 3f
#else-----// ff
                                         ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
----inline bool left_heavy(node *n) const {-------// a0
------else return *cur;-----// 54
                                         -----return n && height(n->l) > height(n->r); }------// a8
#endif------// af
                                         ----inline bool right_heavy(node *n) const {-------// 27
-----return n && height(n->r) > height(n->l); }------// c8
-----node *n = new node(item, prev);-----// eb
                                         ----inline bool too_heavy(node *n) const {------// 0b
-----*cur = n, fix(n); return n; }-----// 29
                                         -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }-----// 67
                                         ----void delete_tree(node *n) {-------// fd
----void erase(node *n, bool free = true) {------// 58
                                         -----if (n) { delete_tree(n->l), delete_tree(n->r); delete n; } }-----// ef
-----if (!n) return;-----// 96
                                         ----node*& parent_leg(node *n) {------// 6a
-----if (!n->l \&\& n->r) parent_leg(n) = n->r, n->r->p = n->p;------// 12
                                         ------if (!n->p) return root;------// ac
-----else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;------// 6b
                                         -----if (n->p->l == n) return n->p->l;------// 83
-----else if (n->l && n->r) {------// 6c
                                         ------if (n->p->r == n) return n->p->r;------// cc
-----node *s = successor(n);-----// e5
                                         -----assert(false): }-----// 20
-----erase(s. false):-----// 0a
                                         ----void augment(node *n) {------// 72
-----if (!n) return:-----// 0e
------if (n->l) n->l->p = s;------// θb
                                         -----n->size = 1 + sz(n->1) + sz(n->r);------// 93
-----if (n->r) n->r->p = s;-----// ed
                                         -----parent_leg(n) = s, fix(s);------// 82
                                         ----#define rotate(l, r) \------// 62
-----return;-----// e5
                                         -----node *l = n->l; \\\------// 7a
-----} else parent_leg(n) = NULL;------// de
                                         -----l->p = n->p; \\-----// 2b
------fix(n > p), n > p = n > l = n > r = NULL;------// 43
                                         -----parent_leg(n) = l; \|-----// fc
-----if (free) delete n; }------// 23
----node* successor(node *n) const {------// 23
```

```
Reykjavík University
                                           ------if (!cmp(i, p)) break;-----// a9
-----n->l = l->r; \\ ------// e8
                                           -----swp(i, p), i = p; } }-----// 93
-----if (l->r) l->r->p = n; \\------// d3
                                            ----void sink(int i) {------// ce
------l->r = n, n->p = l; \\-------// eb
                                            ------while (true) {------// 3c
-----augment(n), augment(\overline{\mathsf{l}})------// 81
                                            ------int l = 2*i + 1, r = l + 1;-----// b4
----void left_rotate(node *n) { rotate(r, l); }-----// 45
                                            -----if (l >= count) break;-----// d5
----void right_rotate(node *n) { rotate(l, r); }-----// ca
                                            ------int m = r >= count || cmp(l, r) ? l : r;-------// cc
----void fix(node *n) {------// 0d
                                            -----if (!cmp(m, i)) break;-----// 42
------while (n) { augment(n);------// 69
                                            -----swp(m, i), i = m; } }-----// 1d
-----if (too_heavy(n)) {------// 4c
                                           public:----// cd
-----if (left_heavy(n) && right_heavy(n->l)) left_rotate(n->l);----// a9
                                            ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {-----// 17
------else if (right_heavy(n) && left_heavy(n->r))------// b9
                                            -----q = new int[len], loc = new int[len];-----// f8
-----right_rotate(n->r);------// 08
                                            -----memset(loc, 255, len << 2); }-----// f7
-----if (left_heavy(n)) right_rotate(n);-----// 93
                                            ----~heap() { delete[] q; delete[] loc; }-----// 09
-----else left_rotate(n);-----// d5
                                            ----void push(int n, bool fix = true) {------// b7
-----n = n->p; }-----// 28
                                           -----if (len == count || n >= len) {------// 0f
-----n = n->p; } };-----// a2
                                            #ifdef RESIZE-----// a9
 Also a very simple wrapper over the AVL tree that implements a map interface.
                                            -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"-----// 01
                                            -----while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                                           ------int *newq = new int[newlen], *newloc = new int[newlen];------// e3
template <class K. class V>-----// da
                                           -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --//94
class avl_map {-----// 3f
                                           -----memset(newloc + len, 255, (newlen - len) << 2);-----// 18
public:----// 5d
                                           -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                                           -----loc = newloc, q = newq, len = newlen;-----// 61
-----K key; V value;------// 32
                                           #else-----// 54
-----node(K k, V v) : key(k), value(v) { }------// 29
                                           -----assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                                           #endif-----// 64
----avl_tree<node> tree;------// b1
                                           -----}-----// 4b
---- V& operator [](K key) {------// 7c
                                           ------assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
                                           -----loc[n] = count, q[count++] = n;------// 6b
-----if (!n) n = tree.insert(node(key, V(0)));------// cb
                                           ------if (fix) swim(count-1); }------// bf
-----return n->item.value;-----// ec
                                           ----void pop(bool fix = true) {------// 43
----}-----// 2e
                                           -----assert(count > 0);-----// eb
}:----// af
                                           -----loc[q[0]] = -1, q[0] = q[-count], loc[q[0]] = 0;------// 50
                                           -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                                            #define RESIZE-----// d0
                                            ----int top() { assert(count > 0); return q[0]; }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                                            ----void heapify() { for (int i = count - 1; i > 0; i--)------// 39
struct default_int_cmp {------// 8d
                                            -----if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }-----// θb
----default_int_cmp() { }------// 35
                                            ----void update_key(int n) {------// 26
----bool operator ()(const int &a, const int &b) { return a < b; } };------// e9
                                            -----assert(loc[n] != -1), swim(loc[n]), sink(loc[n]); }------// 7d
template <class Compare = default_int_cmp>-----// 30
                                            ----bool empty() { return count == 0; }-----// f8
class heap {------// 05
                                            ----int size() { return count; }------// 86
private:-----// 39
                                            ----void clear() { count = 0, memset(loc, 255, len << 2); } };-----// 58
----int len, count, *q, *loc, tmp;-----// 0a
----Compare _cmp;-----// 98
                                           2.7. Skiplist. An implementation of a skiplist.
----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }------// a0
----inline void swp(int i, int j) {------------------------// 1c #define BP 0.20------------------------------// aa
------SWP(q[i], q[j]), SWP(loc[q[i]], loc[q[j]]); }------// 67 #define MAX_LEVEL 10--------------------------// 56
------while (i > 0) {-------// 1a ----unsigned int cnt = 0;------// 28
```

```
Reykjavík University
----return cnt; }---------------------// al -------FIND_UPDATE(x->next[i]->item, target);-------// bb
-----int *lens:------update[i]->next[i] = x->next[i];------// 59
------update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
---node *head:----// b7
                                -----current_level--; } } };-----// 59
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                 2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----~skiplist() { clear(); delete head; head = NULL; }------// aa
                                 list supporting deletion and restoration of elements.
----#define FIND_UPDATE(cmp, target) N------// c3
                                 template <class T>-----// 82
------int pos[MAX_LEVEL + 2]; \[\[ \]------// 18
                                 struct dancing_links {-----// 9e
-----memset(pos, 0, sizeof(pos)); N-----// f2
                                 ----struct node {------// 62
-----node *x = head; \[\]-----// Of
                                 -----T item;-----// dd
-----node *update[MAX_LEVEL + 1]; \[\scrip_------// 01\]
                                -----node *l, *r;-----// 32
------memset(update, 0, MAX_LEVEL + 1); N------// 38 ------node(const T &_item, node *_l = NULL, node *_r = NULL)------// 6d
                                 -----: item(_item), l(_l), r(_r) {------// 6d
------for(int i = MAX_LEVEL; i >= 0; i--) { \[ \sqrt{--------//87} \]
                                 -----if (l) l->r = this;-----// 97
-----if (r) r->l = this;-----// 81
------}-------// 2d
-----pos[i] += x->lens[i]; x = x->next[i]; } \| \]
                                 ----}:-------// d3
------update[i] = x; N-------// dd ----node *front, *back;------// aa
----void clear() { while(head->next && head->next[0])-------// 91 ------back = new node(item, back, NULL);------// c4
------erase(head->next[0]->item); }------// e6 ------if (!front) front = back;------// d2
----node *find(T target) { FIND_UPDATE(x->next[i]->item, target);------// 36 ------return back;-----
-----return x && x->item == target ? x : NULL; }-------// 50 ...}
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------front = new node(item, NULL, front);------// 47
------return pos[0]; }------// 19 ------if (!back) back = front;------// 10
----node* insert(T target) {--------------// 80 -----return front;------------// cf
------FIND_UPDATE(x->next[i]->item, target);-------// 3a ---}
------int lvl = bernoulli(MAX_LEVEL);--------// 7a ------if (!n->l) front = n->r; else n->l->r = n->r;------// ab
------if(lvl > current_level) current_level = lvl;-------// 8a ------if (!n->r) back = n->l; else n->r->l = n->l;------// 1b
-----x = new node(lvl, target);------// 36 ---}
-----x->next[i] = update[i]->next[i];------// 46 -----if (!n->l) front = n; else n->l->r = n;------// a5
-----x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];------// bc -----if (!n->r) back = n; else n->r->l = n;------// 9d
-----update[i]->next[i] = x;------// 20 ---}
-----update[i]->lens[i] = pos[0] + 1 - pos[i];-----// 42 };------// 5e
-----}-----// fc
                                 2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;-----// 07
-----size++:-----// 19
-----return x; }-----// c9
                                 #define BITS 15-----// 7b
                                 struct misof_tree {------// fe
----void erase(T target) {------// 4d
                                 ----int cnt[BITS][1<<BITS];-----// aa
```

```
Reykjavík University
-----return res;-----(pts, 0, size(pts) - 1, 0); }----// 3a ----kd_tree(vector<pt> pts) { root = construct(pts, 0, size(pts) - 1, 0); }----// 19
------int mid = from + (to - from) / 2;------// 7d
                                         -----nth_element(pts.begin() + from, pts.begin() + mid,------// d8
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                                         -----pts.begin() + to + 1, cmp(c));------// 84
bor queries.
                                         -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),------// f1
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1)-----// 77
                                         -----/construct(pts, mid + 1, to, INC(c))); }------// 50
template <int K>-----// cd
                                         ----bool contains(const pt ωp) { return _con(p, root, θ); }-----// 8a
class kd_tree {------// 7e
                                         ----bool _con(const pt &p, node *n, int c) {------// ff
public:----// c7
                                         ------if (!n) return false;-----// 95
----struct pt {------// 78
                                         -----if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));-------// 09
-----double coord[K];------// d6
                                         -----if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// ae
-----pt() {}-----// c1
                                         -----return true: }------// 8e
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }------// 4c
                                         ----void insert(const pt &p) { _ins(p, root, 0); }-----// e9
-----double dist(const pt &other) const {------// 6c
                                         ----void _ins(const pt &p, node* &n, int c) {------// 7d
-----double sum = 0.0;-----// c4
                                         -----for (int i = 0; i < K; i++)-----// 23
                                         -----else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));------// 13
------sum += pow(coord[i] - other.coord[i], 2.0);-----// 46
                                         ------else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }------// f8
-----return sqrt(sum); } };-----// ad
                                         ----void clear() { _clr(root); root = NULL; }------// 15
----struct cmp {------// 8f
                                         ----void _clr(node *n) { if (n) _clr(n->l), _clr(n->r), delete n; }------// 92
-----int c:-----// f6
                                         ----pt nearest_neighbour(const pt &p, bool allow_same=true) {-------// 1c
-----cmp(int _c) : c(_c) {}-----// a5
                                         -----assert(root):-----// 24
------bool operator ()(const pt &a, const pt &b) {------// 26
                                         -----double mn = INFINITY, cs[K];-----// 0d
-----for (int i = 0, cc; i <= K; i++) {------// f0
                                         ------for (int i = 0; i < K; i++) cs[i] = -INFINITY;------// 58
-----/cc = i == 0 ? c : i - 1;-----// bc
                                         -----pt from(cs);-----// af
-----if (abs(a.coord[cc] - b.coord[cc]) > EPS)------// 28
                                         ------for (int i = 0; i < K; i++) cs[i] = INFINITY;------// a8
------return a.coord[cc] < b.coord[cc];------// b7
                                         -----pt to(cs);------// a0
-----return _nn(p, root, bb(from, to), mn, 0, allow_same).first;------// 41
-----return false: } }:-----// e2
                                         ----}------// a3
----struct bb {------// 30
                                         ----pair<pt, bool> _nn(------// cd
-----pt from, to;------// 2f
                                         -----const pt \&p, node *n, bb b, double \&mn, int c, bool same) \{------//1d\}
-----bb(pt _from, pt _to) : from(_from), to(_to) {}------// 57
                                         -----if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// c5
-----double dist(const pt &p) {-----// 3f
                                         ------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;-----// 6d
-----double sum = 0.0;-----// d9
                                         -----pt resp = n->p:-----// 3d
------for (int i = 0; i < K; i++) {-------// 65
                                         -----if (found) mn = min(mn, p.dist(resp));------// c9
-----if (p.coord[i] < from.coord[i])-----// a0
                                         -----node *n1 = n->l, *n2 = n->r:-----// dc
-----sum += pow(from.coord[i] - p.coord[i], 2.0);-----// 00
                                         -----for (int i = 0; i < 2; i++) {------// 74
-----else if (p.coord[i] > to.coord[i])-----// 83
                                         ------if (i == 1 || cmp(c)(n->p, p)) swap(n1, n2), swap(l1, l2);------// ab
------sum += pow(p.coord[i] - to.coord[i], 2.0);------// 8c
                                         -----pair<pt. bool> res =-----// f0
------}------// be
                                         -----nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// ad
-----return sqrt(sum); }------// ef
                                         -----if (res.second && (!found || p.dist(res.first) < p.dist(resp)))----// 17
-----bb bound(double l, int c, bool left) {------// b6
                                         -----resp = res.first, found = true;-----// 62
-----pt nf(from.coord), nt(to.coord);-----// 5c
-----if (left) nt.coord[c] = min(nt.coord[c], l);-----// ef
```

```
Reykjavík University
------for (int j = 0; j < n; j++)---------// c4 ------if (visited[order[i]]) continue;-------// 94
------if (dist[j] != INF)--------// 4e ------S.push(order[i]), dag.push_back(order[i]);------// 40
------for (int k = 0; k < size(adj[j]); k++)-------// 3f ------while (!S.empty()) {-------------------------// 03
-----dist[adj[j][k].first] = min(dist[adj[j][k].first], -----// b ------visited[u = b.top()] = true, b.pop(), uf.unite(u, order[i]);-----// b
------has_negative_cycle = true;--------// 2a ----return pair<union_find, vi>(uf, dag);-----------// f2
----return dist;------// 2e }-----// ca
}-----// c2
                                    3.4. Cut Points and Bridges.
3.2. All-Pairs Shortest Paths.
                                    #define MAXN 5000-----// f7
                                    int low[MAXN], num[MAXN], curnum;-----// d7
3.2.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                    void dfs(const vvi &adj, vi &cp, vii &bri, int u, int p) {-------// 22
problem in O(|V|^3) time.
                                    ----low[u] = num[u] = curnum++;-----// a3
void floyd_warshall(int** arr, int n) {------// 21
                                    ----int cnt = 0; bool found = false;-----// 97
----for (int k = 0; k < n; k++)-----// 49
                                    ----for (int i = 0; i < size(adj[u]); i++) {-------// f3
-----for (int i = 0; i < n; i++)-----// 21
                                    -----int v = adj[u][i];-----// 26
-----for (int j = 0; j < n; j++)-----// 77
                                    -----if (num[v] == -1) {------// f9
-----if (arr[i][k] != INF && arr[k][j] != INF)-----// b1
                                    -----dfs(adj, cp, bri, v, u);-----// 7b
-----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                                    -----low[u] = min(low[u], low[v]);-----// ea
}-----// 86
                                    -----Cnt++:-----// 8f
                                    -----found = found || low[v] >= num[u];-----// fd
3.3. Strongly Connected Components.
                                    -----if (low[v] > num[u]) bri.push_back(ii(u, v));------// 52
3.3.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                    -----} else if (p != v) low[u] = min(low[u], num[v]); }------// c4
graph in O(|V| + |E|) time.
                                    ----if (found && (p !=-1 \mid \mid cnt > 1)) cp.push_back(u); }-------// dc
#include "../data-structures/union_find.cpp"------5
                                    pair<vi,vii> cut_points_and_bridges(const vvi &adj) {-----// 35
-----// 11
                                    ----int n = size(adi):-----// 34
vector<br/>bool> visited;-----// 66
                                    ----vi cp; vii bri;------// 63
vi order:-----// 9b
                                    ----memset(num, -1, n << 2):-----// 4e
-----// a5
                                    ----curnum = 0:-----// 43
void scc_dfs(const vvi &adj, int u) {------// a1
                                    ----for (int i = 0; i < n; i++) if (num[i] == -1) dfs(adj, cp, bri, i, -1);----// e5
----int v; visited[u] = true;-----// e3
                                    ----return make_pair(cp, bri); }------// 70
----for (int i = 0; i < size(adj[u]); i++)-----// c5
------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);------// 6e
                                    3.5. Minimum Spanning Tree.
----order.push_back(u);-----// 19
                                    3.5.1. Kruskal's algorithm.
}-----// dc
                                    #include "../data-structures/union_find.cpp"-----// 5e
  -----// 96
pair<union_find, vi> scc(const vvi &adj) {------// 3e
                                    -----// 11
----order.clear();------// 22 // edges is a list of edges of the form (weight, (a, b))------// c6
----union_find uf(n);------// 6d // the edges in the minimum spanning tree are returned on the same form-----// 4d
----vi dag;------// ae vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);-------// e4 ------if (uf.find(edges[i].second.first) !=-------// d5
```

```
Reykjavík University
-----res.push_back(edges[i]);------// d1 ---}----// d1 ----
-----uf.unite(edges[i].second.first, edges[i].second.second);------// a2 ----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 6e
}------// 88 }------// 35
                             bool euler_path() {-----// d7
3.6. Topological Sort.
                             ---ii se = start_end();-----// 45
                             ----int cur = se.first, at = m + 1;-----// 8c
3.6.1. Modified Depth-First Search.
                             ----if (cur == -1) return false;------// 45
void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
                             ----stack<int> s;-----// f6
------bool& has_cycle) {------// a8
                             ----while (true) {------// 04
----color[cur] = 1;-----// 5b
                             -----if (outdeg[cur] == 0) {------// 32
----for (int i = 0; i < size(adj[cur]); i++) {------// 96
                             -----res[--at] = cur:-----// a6
------int nxt = adj[cur][i];------// 53
                             -----if (s.empty()) break;-----// ee
-----if (color[nxt] == 0)------// 00
                             -----cur = s.top(); s.pop();-----// d7
-----tsort_dfs(nxt, color, adj, res, has_cycle);------// 5b
                             ------} else s.push(cur), cur = adj[cur][--outdeg[cur]];-------// d8
------else if (color[nxt] == 1)------// 53
                             ----}-----// ba
-----has_cycle = true;-----// c8
                             ----return at == 0;-----// c8
-----if (has_cycle) return;------// 7e
                             }-----// aa
----}-----// 3d
----color[cur] = 2:-----// 16
                            3.8. Bipartite Matching.
----res.push(cur):-----// cb
                            3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
}------// 9e
                            O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
-----// ae
vi tsort(int n, vvi adj, bool& has_cycle) {-----// 37
                            graph, respectively.
----has_cycle = false;-----// 37
                            vi* adj;-----// cc
----memset(color, 0, n);------// ce ----if (done[left]) return 0;-------// 08
-----tsort_dfs(i, color, adj, S, has_cycle);-------// 40 -----int right = adj[left][i];------------------// b6
------}-----owner[right] = left; return 1;--------// 26
----return res:-----// 07
                            3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
}-----// 1f
                            ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                             #define MAXN 5000-----// f7
#define MAXV 1000------// 2f int dist[MAXN+1], q[MAXN+1];---------------------// b8
#define MAXE 5000-------// 87 #define dist(v) dist[v == -1 ? MAXN : v]--------// 0f
vi adj[MAXV];-----// ff struct bipartite_graph {-----// 2b
ii start_end() {------// 30 ----bipartite_graph(int _N, int _M) : N(_N), M(_M),------// 8d
----int start = -1, end = -1, any = 0, c = 0;-------------// 74 ------L(new int[N]), R(new int[M]), adj(new vi[N]) {}-------// cd
```

```
-----dist(-1) = INF;------// 96 ---}-----// 77
-------while(l < r) {-------// 69 ----void destroy() { delete[] head; delete[] curh; }-------// f6
-----int v = q[l++];------// 0c ----void reset() { e = e_store; }------// 87
-----iter(u, adj[v]) if(dist(R[*u]) == INF)-------// 68 -----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// c9
-----return dist(-1) != INF;--------// c2 -----if (v == t) return f;------// 6d
-----if(v != -1) {-------if ((ret = augment(e[i].v, t, min(f, e[i].cap))) > 0)------// 1f
-----iter(u, adi[v])-------// 56 ------return (e[i].cap -= ret, e[i^1].cap += ret, ret);------// ac
------if(dist(R[*u]) == dist(v) + 1)------// 3f -----return 0;-----
-----if(dfs(R[*u])) {-------// d2 ---}-----// d2
-----return true;------// 2d -----if(s == t) return 0;-------// 9d
------}-----memset(d, -1, n * sizeof(int));--------// a8
----}------while (l < r)-------// 7a
----void add_edge(int i, int j) { adj[i].push_back(j); }------// cb ------for (int v = q[l++], i = head[v]; i != -1; i = e[i].nxt)-----// a2
------memset(L, -1, sizeof(int) * N);--------// 9f ------if (d[s] == -1) break;-------// a0
------memset(R, -1, sizeof(int) * M);--------// 0a ------memcpy(curh, head, n * sizeof(int));------// 10
------while(bfs()) for(int i = 0; i < N; ++i)--------// 31 -------while ((x = augment(s, t, INF)) != 0) f += x;------// a6
};-----// b5 ---}
                      };-----// 3b
3.9. Maximum Flow.
3.9.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It computes 3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
                      O(|V||E|^2). It computes the maximum flow of a flow network.
the maximum flow of a flow network.
----struct edge {-------// 1e ----struct edge {------// fc
-----edge() { }-------edge(<u>int</u>_v, <u>int</u>_cap, <u>int</u>_nxt) : v(_v), cap(_cap), nxt(_nxt) { }----// 7a
-----e.reserve(2 * (m == -1 ? n : m));------// 24 -----memset(head = new int[n], -1, n << 2);-----// 58
------head = new int[n], curh = new int[n];-------// 6b ---}------// 3a
------memset(head, -1, n * sizeof(int));--------// 56 ----void destroy() { delete[] head; }-----------------// d5
```

Reykjavík University

----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {-------// dd ----e.reserve(2 * (m == -1 ? n : m));------// e6

#include "dinic.cpp"-----// 58

-----// 25

```
Reykjavík University
-----int v = queries[u][i].first;-------// 38 ------T head = *beqin;-----------------// 43
----}------// ad
                              4.2. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
};-----// 5f
                              struct entry { ii nr; int p; };-----// f9
             4. Strings
                              bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77
                              struct suffix_array {------// 87
4.1. Trie. A Trie class.
                              ----string s; int n; vvi P; vector<entry> L; vi idx;-----// b6
template <class T>-----// 82
                              ----suffix_array(string _s) : s(_s), n(size(s)) {------// a3
class trie {------// 9a
                              -----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);------// 12
private:----// f4
                              ------for (int i = 0; i < n; i++) P[0][i] = s[i];-----// df
----struct node {------// ae
                              ------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {-------// \theta b
-----map<T. node*> children:-----// a0
                              -----P.push_back(vi(n));-----// 99
------int prefixes, words;------// e2
                              ------for (int i = 0; i < n; i++)-----// ad
-----node() { prefixes = words = 0; } };------// 42
                              -----L[L[i].p = i].nr = ii(P[stp - 1][i],-----// 53
public:----// 88
                              -----i + cnt < n ? P[stp - 1][i + cnt] : -1);------// 93
---node* root:-----// a9
                              -----sort(L.begin(), L.end());-----// a7
----trie() : root(new node()) { }------// 8f
                              ------for (int i = 0; i < n; i++)-----// d6
----template <class I>-----// 89
                              ------P[stp][L[i].p] = i > 0 &&-----// 19
----void insert(I begin, I end) {------// 3c
                              -----L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;-----// 58
-----node* cur = root:-----// 82
                              -----}-----// af
------while (true) {------// 67
                              -----cur->prefixes++;-----// f1
                              ----}-----// 5b
-----if (begin == end) { cur->words++; break; }-----// db
                              ----int lcp(int x, int y) {------// 8c
------else {------// 3e
                              -----int res = 0;-----// b8
-----T head = *begin;-----// fb
                              -----if (x == y) return n - x;-----// a9
------typename map<T, node*>::const_iterator it;------// 01
                              ------for (int k = size(P) - 1; k >= 0 && x < n && y < n; k--)-----// 10
-----it = cur->children.find(head);------// 77
                              -----if (P[k][x] == P[k][y]) x += 1 << k, y += 1 << k, res += 1 << k;---// 8d
-----if (it == cur->children.end()) {------// 95
                              -----return res:-----// 96
-----pair<T, node*> nw(head, new node());-----// cd
                              ----}------// e7
-----it = cur->children.insert(nw).first;------// ae
                              }:-----// 11
-----} begin++, cur = it->second; } } }-----// 64
----template<class I>-----// b9
                              4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
------if (begin == end) return cur->words;------// a4 -----string keyword; out_node *next;------// f0
-----T head = *begin;-------// 5c ---};------// b9
-----it = cur->children.find(head);------// d9 -----map<char, qo_node*> next;------// 6b
-----if (it == cur->children.end()) return 0;------// 14 -----out_node *out; qo_node *fail;--------// 3e
-----begin++, cur = it->second; } } }------// 7c ------go_node() { out = NULL; fail = NULL; }------// 0f
```

```
c\theta can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
----ao_node *ao:-----// b8
                            accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
----aho_corasick(vector<string> keywords) {------// 4b
                            -----qo = new qo_node();------// 77
                            ----int n = size(s);------// 97
-----iter(k, keywords) {------// f2
                            ----int* z = new int[n];-----// c4
-----qo_node *cur = qo;-----// a2
                            ----int l = 0, r = 0;-----// 1c
-----iter(c, *k)-----// 6e
                            ---z[0] = n;
-----(cur->next[*c] = new go_node());-----// af ----z[i] = 0;------
-----cur->out = new out_node(*k, cur->out);------// 3f
                            -----if (i > r) {------// 26
-----l = r = i:-----// a7
-----queue<go_node*> q;------// 2c ------while (r < n && s[r - l] == s[r]) r++;-----// ff
-----iter(a, qo->next) q.push(a->second);-----// db
                            -----z[i] = r - l; r--;-----// fc
------while (!q.empty()) {------// 07
                            ------qo_node *r = q.front(); q.pop();------// e0
                            -----else {-----// b5
-----iter(a, r->next) {------// 18
                            -----l = i:-----// 02
-----qo_node *s = a->second;-----// 55
                            ------while (r < n \&\& s[r - l] == s[r]) r++;
-----a.push(s);-----// b5
                            -----z[i] = r - l; r--; } }-----// 8d
-----qo_node *st = r->fail;-----// 53
                            ----return z:-----// 53
}-----// db
-----st = st->fail;-----// b3
-----if (!st) st = qo;-----// 0b
                            4.5. Eertree. Constructs an Eertree in O(n), one character at a time.
------s->fail = st->next[a->first];------// c1 #define MAXN 100100-----// 29
-----if (s->fail) {-----// 98
                            #define SIGMA 26-----// e2
------if (!s->out) s->out = s->fail->out;------// ad #define BASE 'a'------//
------else {------// 5b char *s = new char[MAXN];-----// db
------out_node* out = s->out;-----// b8 struct state {------// b8
------out->next = s->fail->out;------// 62 } *st = new state[MAXN+2];------// 57
----}------st[1].len = st[1].link = 0; }-------// 34
----vector<string> search(string s) {-------// c4 ----int extend() {-------// c2
-----vector<string> res;------// 79 ------char c = s[n++]; int p = last;------// 25
-----iter(c, s) {-------// 57 ------if (!st[p].to[c-BASE]) {------// 82
-----cur = cur->fail;------// b1 -----st[p].to[c-BASE] = q;-----// fc
------st[q].len = st[p].len + 2;------// g2
-----cur = cur->next[*c];------// 97 ------do { p = st[p].link;-----// 04
------for (out_node *out = cur->out; out = out->next)------// d7 -----if (p == -1) st[q].link = 1;------// 77
-----res.push_back(out->keyword);------// 7c -----else st[q].link = st[p].to[c-BASE];-----// 6a
------}-----return 1; }-------// 29
------last = st[p].to[c-BASE];-------// 42
----}-----return 0; } };------// ec
}:-----// de
                                        5. Mathematics
4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
```

also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values 5.1. Big Integer. A big integer class.

```
Reykjavík University
----intx(string n) { init(n); }------// b9 ------return sign == 1 ? size() < b.size() > b.size();-----// 4d
----intx(int n) { stringstream ss; ss << n; init(ss.str()); }--------// 36 -------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35
----intx(const intx\& other) : sign(other.sign), data(other.data) { }-------// 3b ------return sign == 1 ? data[i] < b.data[i] > b.data[i] > b.data[i];--// 27
----int sign;-------// 26 ------return false;------// ca
----static const unsigned int radix = 1000000000U;------// f0 ----friend intx abs(const intx &n) { return n < 0 ? -n : n; }-----// 02
-----intx res; res.data.clear();--------// 4e ------if (sign < 0 && b.sign > 0) return b - (-*this);-------// 70
------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {-------// e7 ------unsigned long long carry = 0;-------// 5c
------for (int j = intx::dcnt - 1; j >= 0; j--) {--------// 72 ------carry += (i < size() ? data[i] : OULL) +------// 91
-----if (idx < 0) continue;------// 52 ------c.data.push_back(carry % intx::radix);------// 86
-----digit = digit * 10 + (n[idx] - '0');------// 1f -----carry /= intx::radix;------
-----res.data.push_back(digit);------// 07 -----return c.normalize(sign);------// 20
------data = res.data;------// 7d ----intx operator -(const intx& b) const {-------// 53
-----if (sign > 0 && b.sign < 0) return *this + (-b);-------// 8f
-----if (data.empty()) data.push_back(θ);-------// fa -----if (*this < b) return -(b - *this);--------// 36
------for (int i = data.size() - 1; i > 0 && data[i] == 0; i--)-------// 27 ------intx c; c.data.clear();-------// 6b
-----data.erase(data.begin() + i);------// 67 -----long long borrow = 0;------// f8
------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
----}------c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
------for (int i = n.size() - 1; i >= 0; i--) {-------// 63 ---}-----// 85
------if (first) outs << n.data[i], first = false;-----// 33 ----intx operator *(const intx& b) const {------// bd
------unsigned int cur = n.data[i];-------// 0f ------for (int i = 0; i < size(); i++) {-------// 7a
-----stringstream ss; ss << cur;------// 8c ------long long carry = 0;------------------// 20
-----string s = ss.str();------// 64 ------for (int j = 0; j < b.size() || carry; j++) {------// c0
------int len = s.size();------// 0d ------if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
------c.data[i + j] = carry % intx::radix;------// 86
-----return outs;-------// cf -----}------// ge
----bool operator <(const intx& b) const {--------// 21 ----friend pair<intx,intx> divmod(const intx& n, const intx& d) {-------// fb
```

```
Reykjavík University
------assert(!(d.size() == 1 && d.data[0] == 0));-------// e^9 -----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));-------// e^9
-----intx q, r; q.data.assign(n.size(), 0);--------// ca ----for (int i = 0; i < l - 1; i++)------------// 90
------for (int i = n.size() - 1; i >= 0; i--) {--------// 1a ------if (data[i] >= (unsigned int)(radix)) {-------// 44
-----r.data.insert(r.data.begin(), 0);-------// c7 ------data[i+1] += data[i] / radix;-------// e4
-----r = r + n.data[i]:--------// e6 ------data[i] %= radix:--------// bd
-----long long k = 0;------// cc -----}-----// cc
-----if (d.size() < r.size())-------// b9 ---int stop = l-1;-------------------// cb
-----k = (long long)intx::radix * r.data[d.size()];------// f7 ----while (stop > 0 && data[stop] == 0) stop--;-------// 97
------k /= d.data.back();--------// b7 ---ss << data[stop];--------// 96
-----//--- intx dd = abs(d) * t;-------// 42 ----delete[] A; delete[] B;--------// f7
-----//--- while (r + dd < 0) r = r + dd, k -= t; }------// b6 ----delete[] a; delete[] b;----------// 7e
-----return pair<intx, intx>(q.normalize(n.sign * d.sign), r);-----// 7b
                                 5.2. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
----}-----// 3f
----intx operator /(const intx& d) const {------// 86
                                 k items out of a total of n items.
-----return divmod(*this.d).first: }-----// ba
                                 int nck(int n, int k) {-----// f6
----intx operator %(const intx& d) const {-------// 75
                                 ----if (n - k < k) k = n - k;------// 18
-----return divmod(*this,d).second * sign; }------// 46
                                 ---int res = 1:-----// cb
----return res;------// e4
5.1.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
#include "intx.cpp"-----// 83
                                 5.3. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
#include "fft.cpp"-----// 13
                                 integers a, b.
-----// e0
                                 int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
intx fastmul(const intx &an, const intx &bn) {------// ab
                                  The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
----string as = an.to_string(), bs = bn.to_string();------// 32
                                 and also finds two integers x, y such that a \times x + b \times y = d.
----int n = size(as), m = size(bs), l = 1,-----// dc
-----len = 5, radix = 100000,-----// 4f
                                 int egcd(int a, int b, int& x, int& y) {-----// 85
                                 ----if (b == 0) { x = 1; y = 0; return a; }------// 7b
-----*a = new int[n], alen = 0,-----// b8
----memset(b, 0, m << 2);-----// 01 -----x -= a / b * y;-----// 4a
------for (int j = min(len - 1, i); j >= 0; j--)--------// 43 -----return d;-----------------------------// db
------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
                                 5.4. Trial Division Primality Testing. An optimized trial division to check whether an integer is
-----b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
----while (l < 2*max(alen.blen)) l <<= 1:-----// 51
----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? b[i] : 0, 0);------// 66 ----if (n < 4) return true;-------------// d9
```

Reykjavík University

```
------for (int i = m; i < n; i += mx << 1) {------// 6a
-----cpx t = x[i + mx] * w; -----// 12
-----x[i + mx] = x[i] - t;
----x[i] += t;-----// 0e
----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);-----// d2
}-----// 73
void czt(cpx *x, int n, bool inv=false) {-----// 7f
----int len = 2*n+1;------// d6
----while (len & (len - 1)) len &= len - 1;-----// 2d
----len <<= 1:-----// cf
----cpx w = \exp(-2.0 * pi / n * cpx(0,1)),-----// b2
-----*c = new cpx[n], *a = new cpx[len],-----// 9f
-----*b = new cpx[len];------// 70
----for (int i = 0; i < n; i++) c[i] = pow(w, (inv ? -1.0 : 1.0)*i*i/2);-----// 7b
----for (int i = 0; i < n; i++) a[i] = x[i] * c[i], b[i] = 1.0/c[i];-----// df
----for (int i = 0; i < n - 1; i++) b[len - n + i + 1] = 1.0/c[n-i-1]; ------// c1
----fft(a, len); fft(b, len);-----// 92
----for (int i = 0; i < len; i++) a[i] *= b[i];-----// 29
----fft(a, len, true);-----// da
----for (int i = 0; i < n; i++) {------// 3c
-----if (inv) x[i] /= cpx(n);------// 89
----}------// 1d
----delete[] a;------// 90
----delete[] b:-----// 36
----delete[] c;------// 91
```

5.13. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once: $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times: n^k
- Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type
- 2, ..., n_k objects of type k: $\binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}$.

 Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$$

- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times: $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to $x_1 + x_2 + \cdots + x_n = k$ where $x_i > 0$: f_k^n
- Number of subsets of a set with n elements: 2^n
- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an $n \times m$ grid by walking only up and to the right: $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced: $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$

- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an $n \times n$ lattice which do not rise above the main diagonal: C_n
- Number of permutations of n objects with exactly k ascending sequences or runs: $\left\langle {n\atop k}\right\rangle = \left\langle {n\atop n-k-1}\right\rangle = k\left\langle {n-1\atop k}\right\rangle + (n-k+1)\left\langle {n-1\atop k-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle =$
- Number of permutations of n objects with exactly k cycles: $\binom{n}{k} = \binom{n-1}{k-1} + (n-1)\binom{n-1}{k}$
- Number of ways to partition n objects into k sets: $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements): $D_0 = 1, D_1 =$ $0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points: $\binom{n}{k}D_{n-k}$
- Jacobi symbol: $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- Heron's formula: A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} - 1$.
- Divisor sigma: The sum of divisors of n to the xth power is $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$ where $n = \prod_{i=0}^{r} p_i^{a_i}$ is the prime factorization.
- Divisor count: A special case of the above is $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$.
- Euler's totient: The number of integers less than n that are comprime to n are $n \prod_{p|n} \left(1 \frac{1}{p}\right)$ where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- $gcd(2^a 1, 2^b 1) = 2^{gcd(a,b)} 1$

5.14. Numbers and Sequences. Some random prime numbers: 1031, 32771, 1048583, 33554467, 1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.

6. Geometry

6.1. **Primitives.** Geometry primitives.

```
#include <complex>-----// 8e
#define P(p) const point &p-----// b8
#define L(p0, p1) P(p0), P(p1)-----// 30
#define C(p0, r) P(p0), double r-----// 08
#define PP(pp) pair<point,point> &pp-----// a1
typedef complex<double> point;-----// 9e
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// 4a
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// f3
point rotate(P(p), double radians = pi / 2, P(about) = point(0,0)) { ------// \theta b
----return (p - about) * exp(point(0, radians)) + about; }-----// f5
point reflect(P(p), L(about1, about2)) {------// 45
----point z = p - about1, w = about2 - about1;-----// 74
----return conj(z / w) * w + about1; }-----// d1
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// 98
point normalize(P(p), double k = 1.0) { ------// a9
----return abs(p) == 0 ? point(0,0) : p / abs(p) * k; } //TODO: TEST-----// 1c
```

```
Reykjavík University
```

```
bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// 74 ----point v = normalize(B - A, a), u = normalize(rotate(B-A), h);---------// da
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-------// ab ----res1 = A + v + u, res2 = A + v - u;-----------// c2
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }-------// 27 int intersect(L(A, B), C(O, r), point & res1, point & res2) {-------// e4
double signed_angle(P(a), P(b), P(c)) {-------// 46 ---- point H = proj(0 - A, B - A) + A, v = normalize((B - A), sqrt(r∗r - h∗h));// a9
point perp(P(p)) { return point(-imag(p), real(p)); }-----// 3c }-----// b8
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) {-------// b4 ----v = normalize(v, L);-----------------------------// 10
----// NOTE: check for parallel/collinear lines before calling this function---// 88 ----res1 = A + rotate(v, alpha); res2 = A + rotate(v, -alpha);-------// 56
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// 30 }-------// 46
}------// 92 ----point v = rotate(B - A, theta + pi/2), u = rotate(B - A, -(theta + pi/2));-// e3
point closest_point(L(a, b), P(c), bool segment = false) {--------------// 06 ----u = normalize(u, rA);------------------------------// 30
------if (dot(b - a, c - b) > 0) return b;-------// 93 ----0.first = A + normalize(u, rA); 0.second = B + normalize(u, rB);------// 2a
----}------// d5
----double t = dot(c - a, b - a) / norm(b - a);
                                     6.2. Polygon. Polygon primitives.
----return a + t * (b - a);-----// 4f
                                      #include "primitives.cpp"-----// e0
}-----// 19
                                     typedef vector<point> polygon;-----// b3
double line_segment_distance(L(a,b), L(c,d)) {-----// f6
                                      double polygon_area_signed(polygon p) {------// 31
----double x = INFINITY;-----// 8c
                                      ----double area = 0; int cnt = size(p);------// a2
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// 5f
                                      ----for (int i = 1; i + 1 < cnt; i++)-----// d2
----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true)); -----// 97
                                      -----area += cross(p[i] - p[0], p[i + 1] - p[0]);-----// 7e
----else if (abs(c - d) < EPS) x = abs(c - closest\_point(a, b, c, true));-----// 68
                                      ----return area / 2; }------// e1
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// fa
                                      double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25
----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;-----// bb
                                      #define CHK(f,a,b,c) (f(a) < f(b) && f(b) <= f(c) && ccw(a,c,b) < 0-----// b2
----else {------// 5b
                                     int point_in_polygon(polygon p, point q) {------// 58
-----x = min(x, abs(a - closest_point(c,d, a, true)));------// 07
                                      ----int n = size(p); bool in = false; double d;------// 06
-----x = min(x, abs(b - closest_point(c,d, b, true)));------// 75
                                      ----for (int i = 0, j = n - 1; i < n; j = i++)------// 77
-----x = min(x, abs(c - closest_point(a,b, c, true)));------// 48
                                      -----if (collinear(p[i], q, p[i]) &&-----// a5
-----x = min(x, abs(d - closest_point(a,b, d, true)));------// 75
                                      -----0 <= (d = progress(q, p[i], p[j])) && d <= 1)-----// b9
----}------// 60
                                      ------return 0;------// cc
----return x;------// 57
                                      ----for (int i = 0, j = n - 1; i < n; j = i++)-----// 6f
}-----// 8e
                                      -----if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))------// 1f
int intersect(C(A, rA), C(B, rB), point & res1, point & res2) { ------// ca
                                      -----in = !in;-----// b2
----double d = abs(B - A);-----// 06
                                      ----return in ? -1 : 1; }-----// 77
----if ((rA + rB) < (d - EPS) || d < abs(rA - rB) - EPS) return 0;------// 5d
                                     // pair<polygon, polygon> cut_polygon(const polygon &poly, point a, point b) {-// 7b
----double a = (rA*rA - rB*rB + d*d) / 2 / d, h = sqrt(rA*rA - a*a); ------// 5e
                                      //---- polygon left, right;------// 6b
```

```
Reykjavík University
//---- point it(-100, -100);-------// c9 -------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
//------ if (ccw(a, b, p) \le 0) left.push_back(p);------// 12 ------B = c + min(bp, 1.0) * (d - c);--------// 5c
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// e3 ------return true; }-----
//-----// myintersect = intersect where-----// 24 ----else if (parallel(a,b, c,d)) return false;-----// ca
//-----// (a,b) is a line, (p,q) is a line segment------// f2 ----else if (intersect(a,b, c,d, A, true)) {----------------// 10
//----- if (myintersect(a, b, p, q, it))------// f0 -----B = A; return true; }------------------// bf
//-----------left.push_back(it), right.push_back(it);----------// 21 ----return false;------------------------// b7
//--- return pair<polygon, polygon>(left, right);-----// 1d -----// e6
// }-----// 37
                                         6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
                                         coordinates) on a sphere of radius r.
#include "polygon.cpp"------// 58 double qc_distance(double pLat, double pLong,------// 7b
#define MAXN 1000-----// 09
                                        -----/ double qLat, double qLong, double r) {------// a4
point hull[MAXN];------// 43 ----pLat *= pi / 180; pLong *= pi / 180;------// ee
bool cmp(const point &a, const point &b) {------// 32 ----qLat *= pi / 180; qLong *= pi / 180;------// 75
----return abs(real(a) - real(b)) > EPS ?-----// 44 ----return r * acos(cos(pLat) * cos(pLong - qLong) +-----// e3
-----real(a) < real(b) : imag(a) < imag(b); }------// 40 -----sin(pLat) * sin(qLat));------// 1e
----int n = size(p), l = 0;-----// 67 }-----// 3f
----sort(p.begin(), p.end(), cmp);-----// 3d
                                         6.6. Triangle Circumcenter. Returns the unique point that is the same distance from all three
----for (int i = 0; i < n; i++) {------// 6f
-----if (i > 0 && p[i] == p[i - 1]) continue;-----// b2
                                         points. It is also the center of the unique circle that goes through all three points.
                                         #include "primitives.cpp"-----// e0
------while (l >= 2 \&\& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--:-----// 20
                                         point circumcenter(point a, point b, point c) {-----// 76
-----hull[l++] = p[i]:-----//
                                         ----b -= a, c -= a;-----// 41
----}------//
                                         ----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a
----int r = l:-----// 59
----for (int i = n - 2; i >= 0; i--) {------// 16
                                         }-----// c3
-----if (p[i] == p[i + 1]) continue;-----// c7
                                         6.7. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
------while (r - l >= 1 \& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 9f
                                         pair of points.
-----hull[r++] = p[i];-----// 6d
                                         #include "primitives.cpp"-----// e0
----}-----------// 74
                                         -----// 85
----return l == 1 ? 1 : r - 1:-----// 6d
                                         struct cmpx { bool operator ()(const point &a, const point &b) {-----// 01
}-----// 79
                                         -----return abs(real(a) - real(b)) > EPS ?------// e9
6.4. Line Segment Intersection. Computes the intersection between two line segments.
                                         -----real(a) < real(b) : imag(a) < imag(b); } };------// 53
#include "primitives.cpp"------// e0 struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
bool line_segment_intersect(L(a, b), L(c, d), point δA, point δB) {------// 6c ----return abs(imag(a) - imag(b)) > EPS ?--------// θb
------A = B = a; return abs(a - d) < EPS; }-------// ee double closest_pair(vector<point> pts) {------// f1
------A = B = a; double p = progress(a, c,d):-------// c9 ----set<point, cmpv> cur:--------------------------// bd
------return 0.0 <= p && p <= 1.0--------// 8a ----set<point, cmpy>::const_iterator it, jt;--------// a6
------A = B = c; double p = progress(c, a,b);-------// d9 ------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);------// 8b
------return 0.0 <= p && p <= 1.0--------// 8e -----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));------// fc
```

```
Reykjavík University
-----cur.insert(pts[i]); }------// 82 -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < EPS; }-----// d9
----return mn; }------// 4c ----bool isInSegmentStrictly(L(A, B)) const {-------// 0e
                                                     -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < -EPS; }------// ba
                                                     ----double getAngle() const {------// Of
6.8. 3D Primitives. Three-dimensional geometry primitives.
                                                     -----return atan2(y, x); }-----// 40
#define P(p) const point3d &p-----// a7
                                                     ----double getAngle(P(u)) const {------// d5
#define L(p0, p1) P(p0), P(p1)-----// Of
                                                     -----return atan2((*this * u).length(), *this % u); }-----// 79
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)-----// 67
                                                     ----bool isOnPlane(PL(A, B, C)) const {------// 8e
struct point3d {-----// 63
                                                     -----return abs((A - *this) * (B - *this) % (C - *this)) < EPS; } };-----// 74
----double x, y, z;-----// e6
                                                     int line_line_intersect(L(A, B), L(C, D), point3d &0){----------------------// dc
----point3d() : x(0), y(0), z(0) {}-----// af
                                                     ----if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;------// 6a
----point3d(double _x, double _y, double _z) : x(_x), y(_y), z(_z) {}------// fc
                                                     ----if (((A - B) * (C - D)).length() < EPS)-----// 79
----point3d operator+(P(p)) const {------// 17
                                                     -----return A.isOnLine(C. D) ? 2 : 0:-----// 09
-----return point3d(x + p.x, y + p.y, z + p.z); }-----// 8e
                                                     ----point3d normal = ((A - B) * (C - B)).normalize();-----// bc
----point3d operator-(P(p)) const {------// fb
                                                     ----double s1 = (C - A) * (D - A) % normal;------// 68
-----return point3d(x - p.x, y - p.y, z - p.z); }-----// 83
                                                     ----0 = A + ((B - A) / (s1 + ((D - B) * (C - B) % normal))) * s1;------// 56
----point3d operator-() const {------// 89
                                                     ----return 1; }-----// a7
-----return point3d(-x, -y, -z); }-----// d4
                                                     int line_plane_intersect(L(A, B), PL(C, D, E), point3d & 0) {------// 09
----point3d operator*(double k) const {------// 4d
                                                     ----double V1 = (C - A) * (D - A) % (E - A);------// c1
-----return point3d(x * k, y * k, z * k); }-----// fd
                                                     ----double V2 = (D - B) * (C - B) % (E - B);------// 29
----point3d operator/(double k) const {------// 95
                                                     ----if (abs(V1 + V2) < EPS)-------// 81
-----return point3d(x / k, y / k, z / k); }-----// 58
                                                     -----return A.isOnPlane(C, D, E) ? 2 : 0;-----// d5
----double operator%(P(p)) const {------// d1
                                                     ---0 = A + ((B - A) / (V1 + V2)) * V1;
-----return x * p.x + y * p.y + z * p.z; }-----// 09
                                                     ----return 1; }-----// ce
----point3d operator*(P(p)) const {------// 4f
                                                     bool plane_plane_intersect(P(A), P(nA), P(B), P(nB), point3d &P, point3d &Q) {-// 5a
-----return point3d(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }-----// ed
                                                     ----point3d n = nA * nB:-----// 49
----double lenath() const {------// 3e
                                                     ----if (n.isZero()) return false;------// 03
-----return sqrt(*this % *this); }-----// 05
                                                     ----point3d v = n * nA:-----// d7
----double distTo(P(p)) const {------// dd
                                                     ----P = A + (n * nA) * ((B - A) % nB / (v % nB));
------return (*this - p).length(); }-----// 57
                                                     ---0 = P + n:------// 9c
----double distTo(P(A), P(B)) const {------// bd
                                                     ----return true; }------// 1a
-----// A and B must be two different points-----// 4e
-----return ((*this - A) * (*this - B)).length() / A.distTo(B); }------// 6e
                                                     6.9. Polygon Centroid.
----point3d normalize(double k = 1) const {------// db
                                                     #include "polygon.cpp"-----// 58
-----// lenath() must not return 0-----// 3c
                                                     point polygon_centroid(polygon p) {------// 79
-----return (*this) * (k / length()); }-----// d4
                                                     ----double cx = 0.0. cy = 0.0:-----// d5
----point3d qetProjection(P(A), P(B)) const {-------// 86
                                                     ----double mnx = 0.0, mny = 0.0;-----// 22
-----point3d v = B - A;-----// 64
                                                     ----int n = size(p):-----// 2d
-----return A + v.normalize((v % (*this - A)) / v.length()); }------// 53
                                                     ----for (int i = 0; i < n; i++)-----// 24
----point3d rotate(P(normal)) const {------// 55
                                                     -\cdots mnx = min(mnx, real(p[i])),.....// 6d
----// normal must have length 1 and be orthogonal to the vector-----// eb
                                                     -----mny = min(mny, imag(p[i]));-----// 95
    return (*this) * normal; }-----// 5c
                                                     ----for (int i = 0; i < n; i++)-----// df
----point3d rotate(double alpha, P(normal)) const {------// 21
                                                     -----p[i] = point(real(p[i]) - mnx, imaq(p[i]) - mny);-----// c2
-----return (*this) * cos(alpha) + rotate(normal) * sin(alpha); }------// 82
                                                     ----for (int i = 0; i < n; i++) {-------// 06
----point3d rotatePoint(P(0), P(axe), double alpha) const{-----------------// 7a
                                                     ------jnt j = (i + 1) % n;-----// d1
-----point3d Z = axe.normalize(axe % (*this - 0));
                                                     -----return 0 + Z + (*this - 0 - Z).rotate(alpha, 0); }-----// 38
                                                     ----bool isZero() const {------// 64
                                                     ----return point(cx, cy) / 6.0 / polygon_area_signed(p) + point(mnx, mny); }---// 2f
-----return abs(x) < EPS && abs(y) < EPS && abs(z) < EPS; }-----// 15
----bool isOnLine(L(A, B)) const {------// 30
                                                     6.10. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
-----return ((A - *this) * (B - *this)).isZero(); }-----// 58
                                                        • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
----bool isInSegment(L(A, B)) const {------// f1
                                                        • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
```

- of that is the area of the triangle formed by a and b.
 - 7. Other Algorithms

```
function f on the interval [a, b], with a maximum error of \varepsilon.
double binary_search_continuous(double low, double high,-----// 8e
```

```
-----double eps, double (*f)(double)) {------// cθ
----while (true) {------// 3a
------double mid = (low + high) / 2, cur = f(mid);-----// 75
-----if (abs(cur) < eps) return mid;------// 76
-----else if (0 < cur) high = mid;------// e5
-----else low = mid:-----// a7
----}-----// b5
```

Another implementation that takes a binary predicate f, and finds an integer value x on the integer interval [a,b] such that $f(x) \wedge \neg f(x-1)$. int binary_search_discrete(int low, int high, bool (*f)(int)) {-------// 51

```
----assert(low <= high);-----// 19
----while (low < high) {------// a3
------int mid = low + (high - low) / 2;-----// 04
-----if (f(mid)) high = mid;-----// ca
-----else low = mid + 1;-----// 03
----assert(f(low));-----// 42
----return low;------// a6
}-----// d3
```

7.2. **Ternary Search.** Given a function f that is first monotonically increasing and then monotonically cally decreasing, ternary search finds the x such that f(x) is maximized.

```
template <class F>------res[eng[curw] = curm] = curw, ++i; break;------// 5e
-----/double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3:-----// e8
-----else hi = m2;-----// b3
----}------// bb
----return hi:-----// fa
}-----// 66
```

7.3. **2SAT.** A fast 2SAT solver.

```
-----// 63 ------<mark>int</mark> row, col, size;---------------// ae
bool two_sat(int n, const vii& clauses, vi& all_truthy) {-------// f4 -----node(int _row, int _col) : row(_row), col(_col) {------// c9
----all_truthy.clear();-------------------------// 31 --------size = 0; l = r = u = d = p = NULL; }-----------------// c3
----vvi adj(2*n+1);-------// 7b ---};-------// c1
------if (clauses[i].first != clauses[i].second)-------// 87 ----node *head;------
-----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93 ----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
```

```
----union_find scc = res.first:-----// 42
                                            ----vi dag = res.second;------// 58
                                            ----vi truth(2*n+1, -1);------// 00
                                            ----for (int i = 2*n; i >= 0; i--) {------// f4
7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
                                            -----int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);-// 5a
                                            -----if (cur == 0) continue;-----// 26
                                            -----if (p == o) return false;-----// 33
                                            -----if (truth[p] == -1) truth[p] = 1;-----// c3
                                            -----truth[cur + n] = truth[p];-----// b3
                                            -----truth[o] = 1 - truth[p];-----// 80
                                            ------if (truth[p] == 1) all_truthy.push_back(cur);------// 5c
                                            ----}-------// d9
                                            ----return true:-----// eb
                                            7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
                                            ----queue<int> q;-----// f6
                                            ----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));-------// c3
                                            ----for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)-----// 05
                                            -----inv[i][w[i][j]] = j;-----// b9
                                            ----for (int i = 0; i < n; i++) q.push(i);-----// fe
                                            ----while (!q.empty()) {-----// 55
                                            ------int curm = q.front(); q.pop();-----// ab
                                            ------for (int &i = at[curm]; i < n; i++) {-------// 9a
                                            -----int curw = m[curm][i];-----// cf
                                            -----if (eng[curw] == -1) { }-----// 35
                                            ------else if (inv[curw][curm] < inv[curw][eng[curw]])------// 10
```

7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the Exact Cover problem. bool handle_solution(vi rows) { return false; }------// 63

----**return** res:-----// 95

-----q.push(eng[curw]);-----// 8c

-----else continue;-----// b4

```
struct exact_cover {------// 95
                                     ----struct node {------// 7e
#include "../graph/scc.cpp"-----// c3 -----node *l, *r, *u, *d, *p;------------// 19
```

Reykjavík University	
7.8. Dates. Functions to simplify date calculations.	
<pre>int intToDay(int jd) { return jd % 7; }</pre>	
<pre>int dateToInt(int y, int m, int d) {</pre>	// 96
return 1461 * (y + 4800 + (m - 14) / 12) / 4 +	// a8
367 * (m - 2 - (m - 14) / 12 * 12) / 12	// d1
3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +	// be
d - 32075;	// e0
}	
<pre>void intToDate(int jd, int &y, int &m, int &d) {</pre>	// a1
int x, n, i, j;	
x = jd + 68569;	
n = 4 * x / 146097;	// 2f
x -= (146097 * n + 3) / 4;	// 58
i = (4000 * (x + 1)) / 1461001;	
x -= 1461 * i / 4 - 31;	
j = 80 * x / 2447;	
d = x - 2447 * j / 80;	
x = j / 11;	
m = j + 2 - 12 * x;	
y = 100 * (n - 49) + i + x;	// 76

8. Useful Information

8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo m_1, m_2, \ldots, m_k , where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When $n = 0, n = -1, n = 1, n = 2^{31} 1$ or $n = -2^{31}$? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

8.3. Worst Time Complexity.

	n	Worst AC Algorithm	Comment
	≤ 10	$O(n!), O(n^6)$	e.g. Enumerating a permutation
	≤ 15	$O(2^n \times n^2)$	e.g. DP TSP
	≤ 20	$O(2^{n}), O(n^{5})$	e.g. DP + bitmask technique
	≤ 50	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
	$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
	$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
	$\leq 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree
	$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

8.4. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- \bullet snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.