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            1. Code Templates
                               ----public static void main(String[] args) throws Exception {--------// 02
                               -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                               ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                               -----// code-----// e6
setxkbmap -option caps:escape
                               -----out.flush():-----// 56
set -o vi
                               xset r rate 150 100
                               }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                           2. Data Structures
syn on | colorscheme slate
                               2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                               struct union find {-----// 42
#include <cassert>------------------// 65 ----int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }-------// ba
#include <cmath>-----// 7d ----bool unite(int x, int y) {-------// 6c
#include <cstdio>------// 2e ------int xp = find(x), yp = find(y);------// 64
#include <cstdlib>------// 11 ------if (xp == yp) return false;------// 0b
#include <cstring>------// d0 -------if (p[xp] < p[yp]) p[xp] += p[yp], p[yp] = xp;------// 3a
#include <ctime>------// 28 -------| p[xp] += p[xp], p[xp] = yp;------// 3e
#include <map>-----// 02
#include <set>------// d2
#include <sstream>------ min(a, b); }------// 18 // int f(int a, int b) { return min(a, b); }------
#include <stack>------// cf const int ID = 0;------// 82
#include <string>------// a9 int f(int a, int b) { return a + b; }------// 6d
#include <vector>-----// 4f ----int n; vi data, lazy;-----// fa
----/ 7e ----seqment_tree(const vi &arr) : n(size(arr)), data(4*n), lazy(4*n,INF) {-----// 96
----for (typeof((o).begin()) u = (o).begin(); u != (o).end(); ++u)------// 1a ----int mk(const vi &arr, int l, int r, int i) {--------// 75
const int INF = 2147483647;------// be -----if (l == r) return data[i] = arr[l];------// 7c
const double EPS = 1e-9;------// 0f ------int m = (l + r) / 2;-------// e9
typedef long long ll;------// 8f ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// c2
typedef unsigned long long ull;-----// 81 ----int q(int a, int b, int l, int r, int i) {------// 08
template <class T> T mod(T a, T b) { return (a % b + b) % b; }------// 70 ----void update(int i, int v) { u(i, v, 0, n-1, 0); }------// f1
template <class T> int size(const T &x) { return x.size(); }------// 68 ----int u(int i, int v, int l, int r, int j) {-------// 77
                               -----propagate(l, r, j);-----// θc
1.3. Java Template. A Java template.
                               -----if (r < i || i < l) return data[j];------// cc
import java.io.*;-----return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 96
-----// a3 ----void range_update(int a, int b, int v) { ru(a, b, v, 0, n-1, 0); }------// 65
```

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------propagate(l, r, i);-------// ee template <class T>------// 53
------if (r < a || b < l) return data[i];--------// cc public:--------//
------/ 6b ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// fe
-----ru(a, b, v, m+1, r, 2*i+2));------// 2d -----cnt(other.cnt), data(other.data) { }------// ed
------if (l > r || lazy[i] == INF) return; -------// 08 ------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
------data[i] += lazy[i] * (r - l + 1);-------// 5c ----void operator -=(const matrix& other) {-------// 68
------if (l < r) {-------// f2 ------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
------else lazv[2*i+1] += lazv[i];------// a8 ------for (int i = 0; i < cnt; i++) data[i] *= other; }-----// 40
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];------// 3c ----matrix<T> operator +(const matrix& other) {------// ee
------else lazy[2*i+2] += lazy[i];-------// bb ------matrix<T> res(*this); res += other; return res; }------// 5d
------lazy[i] = INF;------res(*this); res -= other; return res; }------// cf
};-----// e6 -----matrix<T> res(*this); res *= other; return res; }------// 37
                                    ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                                   -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                                   ------for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i... i in O(\log n) time. It only needs O(n) space.
                                    ------for (int k = 0; k < cols; k++)-----// fc
struct fenwick_tree {------res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n; vi data;------// d3 -----return res; }------// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {---------------// dd
----void update(int at, int by) {-------// 76 -----matrix<T> res(cols, rows);-----// b5
-------while (at < n) data[at] += by, at |= at + 1; }-------// fb -------for (int i = 0; i < rows; i++)-------// 9c
------int res = 0;-------// c3 ------return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);-------// bf
};------while (p) {------// cb
struct fenwick_tree_sq {------// d4 ------if (p & 1) res = res * sq;-----// c1
----int n; fenwick_tree x1, x0;--------// 18 -----p >>= 1;-----------------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;-------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {-------// 89
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }-----// 45 ------matrix<T> mat(*this); det = T(1);------------------// 21
----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73 ------for (int r = 0, c = 0; c < cols; c++) {-------// c4
};-------int k = r;------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
------for (int i = 0; i < cols; i++)-----// ab
2.4. Matrix. A Matrix class.
                                   -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
template <> bool eq<double>(double a, double b) { return abs(a - b) < EPS; }---// a7
```

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------if (!eq<T>(mat(r, c), T(1)))-------// 2c ------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;------// 6b
------for (int i = cols-1; i >= c; i--) mat(r, i) /= mat(r, c);-----// 5d ------else if (n->l && n->r) {---------------------------// 6c
------for (int i = 0; i < rows; i++) {---------// 3d ------node *s = successor(n);-------// e5
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 5a
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;------// 43
-----if (!n) return NULL;------// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                           ------if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0------// b5 -----node *p = n->p;------// a7
template <class T>-----// 22 -----return p; }-----// c7
class avl_tree {------// ff ----node* predecessor(node *n) const {------// b4
public:-----// f6 ------if (!n) return NULL;-------// dd
----struct node {-------// 45 ------if (n->l) return nth(n->l->size-1, n->l);-------// 10
-----T item; node *p, *l, *r;-------// a6 -----node *p = n->p;-------// ea
------int size, height;-------// 33 -------while (p && p->l == n) n = p, p = p->p;------// 6d
-----node(const T &_item, node *_p = NULL) : item(_item), p(_p),------// 4f -----return p; }------
------node *cur = root;-------// b4 --------while (cur) {-------// 29
------while (cur) {-------// 8b ------if (n < sz(cur->l)) cur = cur->l;------// 75
------if (cur->item < item) cur = cur->r;------// 71 ------else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-----// cd
------else if (item < cur->item) cur = cur->l;------// cd -----else break;------
------else break; }------// 4f ------} return cur; }------// ed
-----return cur; }------// 84 ----int count_less(node *cur) {------// ec
----node* insert(const T &item) {-------// 4e -----int sum = sz(cur->l);------// bf
------node *prev = NULL, **cur = &root;------// 60 -------while (cur) {-------// 6f
-----prev = *cur;------// f0 -----cur = cur->p;-----// eb
#if AVL_MULTISET-----// 0a private:----// d5
#else-----// ff ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
-----else return *cur;------// 54 -----return n && height(n->l) > height(n->r); }------// a8
-----*cur = n, fix(n); return n; }------// 29 -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }------// 67 ----void delete_tree(node *n) {-------// fd
----void erase(node *n, bool free = true) {-------// 58 ------if (n) { delete_tree(n->r); delete n; } }-----// ef
-----if (!n) return;-----// 96
                           ----node*& parent_leg(node *n) {------// 6a
-----if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;------// 12
```

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------if (!n) return;--------// 0e ----int len, count, *q, *loc, tmp;-------// 0a
------n->size = 1 + sz(n->1) + sz(n->r);-------// 93 ----Compare _cmp;-------// 98
-----l->p = n->p; \|\frac{1}{2} \|\frac{1}{2
                                                  ------while (i > 0) {------// 1a
-----parent_leg(n) = 1; \[ \]\------// fc
                                                  -----int p = (i - 1) / 2;-----// 77
-----n->l = l->r; N-----// e8
                                                  ------if (!cmp(i, p)) break;-----// a9
-----augment(n), augment(l)-------// 81 ------while (true) {-----------------------// 3c
----void left_rotate(node *n) { rotate(r, l); }------// 45 ------int l = 2*i + 1, r = l + 1;------// b4
------| else if (right_heavy(n) δδ left_heavy(n->r))------// b9 ----heap(int init_len = 128) : count(θ), len(init_len), _cmp(Compare()) {------// 17
------right_rotate(n->r);-------// 08 ------q = new int[len], loc = new int[len];-------// f8
-----if (left_heavy(n)) right_rotate(n);------// 93 ------memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// d5 ----~heap() { delete[] q; delete[] loc; }------// 09
-----n = n->p; }------// 28 ----void push(int n, bool fix = true) {------// b7
-----n = n->p; } } };-------// a2 ------if (len == count || n >= len) {-------// 0f
                                                  #ifdef RESIZE-----// a9
 Also a very simple wrapper over the AVL tree that implements a map interface.
                                                  -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"-----// 01
                                                  ------while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                                                  ------int *newq = new int[newlen], *newloc = new int[newlen];------// e3
template <class K, class V>-----// da
                                                  -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --// 94
class avl_map {-----// 3f
                                                  -----/ 18 emset(newloc + len, 255, (newlen - len) << 2);------// 18
public:----// 5d
                                                  -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                                                  -----loc = newloc, q = newq, len = newlen;-----// 61
------K key; V value;------// 32
                                                  #else-----// 54
-----node(K k, V v) : key(k), value(v) { }-----// 29
                                                  -----assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                                                  #endif------// 64
----avl_tree<node> tree:-----// b1
                                                  ------}------// 4b
----V& operator [](K key) {------// 7c
                                                  -----assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
                                                  -----loc[n] = count, q[count++] = n;-----// 6b
-----if (!n) n = tree.insert(node(key, V(0)));-----// cb
                                                  -----if (fix) swim(count-1); }-----// bf
-----return n->item.value:-----// ec
                                                  ----void pop(bool fix = true) {-------// 43
----}------// 2e
                                                  -----assert(count > 0);-----// eb
};-----// af
                                                  -----loc[q[0]] = -1, q[0] = q[-count], loc[q[0]] = 0;------// 50
                                                  -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                                                  #define RESIZE-----// d0
                                                  ----int top() { assert(count > 0); return q[0]; }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                                                  ----void heapify() { for (int i = count - 1; i > 0; i--)-----// 39
struct default_int_cmp {-----// 8d
```

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------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }-------// 0b ------if(lvl > current_level) current_level = lvl;-------// 8a
----void update_kev(int n) {-------------------------// 26 -----x = new node(lyl, target);-------------------// 36
----bool empty() { return count == 0; }-------// f8 ------x->next[i] = update[i]->next[i];------// 46
----int size() { return count; }--------------------------// 86 -------x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];-----------------// bc
----void clear() { count = 0, memset(loc, 255, len << 2); } };-------// 58 ------update[i] = x;------------// 20
                                        -----update[i]->lens[i] = pos[0] + 1 - pos[i];------// 42
2.7. Skiplist. An implementation of a skiplist.
                                        ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// 07
#define BP 0.20-----// aa
                                        -----size++;-----// 19
#define MAX_LEVEL 10-----// 56
                                        -----return x; }-----// c9
unsigned int bernoulli(unsigned int MAX) {-----// 7b
                                        ----void erase(T target) {------// 4d
----unsigned int cnt = 0;------// 28
                                        ------FIND_UPDATE(x->next[i]->item, target);-------// 6b
----while(((float) rand() / RAND_MAX) < BP \&\& cnt < MAX) cnt++;------// d1
                                        -----if(x && x->item == target) {------// 76
                                        ------for(int i = 0; i <= current_level; i++) {-------// 97
template<class T> struct skiplist {-----// 34
                                        -----if(update[i]->next[i] == x) {------// b1
----struct node {------// 53
-----T item:-----// e3
                                        -----update[i]->next[i] = x->next[i];------// 59
------int *lens;------// 07
                                        -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
                                        -----} else update[i]->lens[i] = update[i]->lens[i] - 1;-----// 88
-----node **next:-----// 0c
                                        -----#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))------// 25
                                        -----delete x; _size--;------// 81
------node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}-----// 7c
                                        ------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
-----/node() { free(lens); free(next); }; };-----// aa
                                        -----current_level--; } } };-----// 59
----int current_level, _size;------// 61
---node *head;-----// b7
                                        2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                        list supporting deletion and restoration of elements.
----~skiplist() { clear(); delete head; head = NULL; }------// aa
                                        template <class T>-----// 82
----#define FIND_UPDATE(cmp, target) \sqrt{\phantom{a}}
                                        struct dancing_links {-----// 9e
----struct node {------// 62
-----memset(pos, 0, sizeof(pos)); \\-----// f2
                                        -----T item:-----// dd
-----node *update[MAX_LEVEL + 1]; N------// 01 -----node(const T &_item, node *_l = NULL, node *_r = NULL)-----// 6d
-----memset(update, 0, MAX_LEVEL + 1); \sqrt{\phantom{a}}
                                       -----: item(_item), l(_l), r(_r) {------// 6d
                                        -----if (l) l->r = this;-----// 97
-----for(int i = MAX\_LEVEL; i >= 0; i--) { \sqrt{\phantom{a}}
                                        -----if (r) r->l = this;-----// 81
-----pos[i] = pos[i + 1]; \\ ------// 68
                                        ------}-----// 2d
----};------// d3
------update[i] = x; N-------// dd ----dancing_links() { front = back = NULL; }-----// 72
----void clear() { while(head->next && head->next[0])------// 91 -----if (!front) front = back;-----// d2
------erase(head->next[0]->item); }-------// e6 ------return back;---------------------------// εθ
------return x && x->item == target ? x : NULL; }-------// 50 ----node *push_front(const T &item) {-------// 4a
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------if (!back) back = front;--------------------// 10
-----return pos[0]; }------// 19 -----return front;------// cf
------FIND_UPDATE(x->next[i]->item, target);--------// 3a ----void erase(node *n) {---------------------------// a0
------if(x && x->item == target) return x; // SET--------// 07 ------if (!n->l) front = n->r; else n->l->r = n->r; -------// ab
```

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------if (!n->l) front = n; else n->l->r = n; --------// a5 ------if (p.coord[i] < from.coord[i])------// a0
------if (!n->r) back = n; else n->r->l = n;--------// 9d -------sum += pow(from.coord[i] - p.coord[i], 2.0);------// 00
};-----sum += pow(p.coord[i] - to.coord[i], 2.0);------// 8c
                               ------}------------------------// be
2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                               -----return sqrt(sum); }-----// ef
element.
                                ------bb bound(double l, int c, bool left) {------// b6
#define BITS 15------pt nf(from.coord), nt(to.coord);------// 5c
----int cnt[BITS][1<<BITS];-------// aa ------else nf.coord[c] = max(nf.coord[c], l);------// 71
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }------// b0 -----return bb(nf, nt); } };------// 3b
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49 -------pt p; node *l, *r;-------------------------------// 46
----int nth(int n) {-------// 8a ------node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
------int res = 0;-------// a4 ----node *root;------// 30
----}-----if (from > to) return NULL;-------// f4
-----nth_element(pts.begin() + from, pts.begin() + mid,-----// f3
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                               -----pts.begin() + to + 1, cmp(c));------// 97
bor queries.
                                -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// dd
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1)-----// 77
                               -----construct(pts, mid + 1, to, INC(c))); }-----// 03
template <int K>-----// cd ----bool contains(const pt &p) { return _con(p, root, 0); }------// f0
class kd_tree {------// 7e ----bool _con(const pt &p, node *n, int c) {------// 74
----struct pt {-------// 78 ------if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));-------// 81
-------double coord[K];-------// d6 ------if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// 95
-----pt() {}------// c1 -----return true; }------// 86
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }------// 4c ----void insert(const pt \&p) { _ins(p, root, 0); }------// f9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {-------// 09
------double sum = 0.0;------// c4 ------if (!n) n = new node(p, NULL, NULL);------// 4d
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }-----// 2c
-----return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }----// 73
----struct cmp {------// 8f ----void _clr(n->l), _clr(n->r), delete n; }-----// 1a
-------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// bf
------cc = i == 0 ? c : i - 1;------// bc ------pt from(cs);------
-----/return false; } };------// 62
----struct bb {-------// 30 ----pair<pt, bool> _nn(------// e3
------bb(pt _from, pt _to) : from(_from), to(_to) {}-------// 57 ------if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// 19
------double dist(const pt &p) {------// 3f
```

```
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------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;------// 9f -------visited.insert(*it);---------------------// cb
------if (found) mn = min(mn, p.dist(resp));--------// 18 ---}------// 0b
------for (int i = 0: i < 2: i++) {--------// 50 ---return -1:-------// f5
-----pair<pt, bool> res =-----// 33
-----nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// 72
                                      3.2. Single-Source Shortest Paths.
-----if (res.second & (!found || p.dist(res.first) < p.dist(resp)))----// 76
-----resp = res.first, found = true;-----// 3b
                                      3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|) time.
-----}------------------------// aa
                                      int *dist. *dad:-----// 46
struct cmp {-----// a5
                                       ----bool operator()(int a, int b) {-----// bb
                3. Graphs
                                       -----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
                                      }:-----// 41
edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
                                      pair<int*, int*> dijkstra(int n, int s, vii *adj) {------// 53
graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
                                      ----dist = new int[n]:-----// 84
connected. It runs in O(|V| + |E|) time.
                                       ----dad = new int[n];-----// 05
----Q.push(ii(start, 0));------// 49 ----dist[s] = 0, pq.insert(s);------// 1b
-----ii cur = Q.front(); Q.pop();------// e8 ------for (int i = 0; i < size(adj[cur]); i++) {------// 9e
-----/<sub>06</sub> ------int nxt = adj[cur][i].first,-------//<sub>b8</sub>
-----if (cur.first == end)------// 6f -----ndist = dist[cur] + adj[cur][i].second;------// 0c
------return cur.second;-------// 8a ------if (ndist < dist[nxt]) pg.erase(nxt),------// e4
-----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);------// 0f
-----vi& adj = adj_list[cur.first];-------// 3f -----}
------0.push(ii(*it, cur.second + 1));--------// b7 ----return pair<int*, int*>(dist, dad);------------------// cc
}-----// 7d
 Another implementation that doesn't assume the two vertices are connected. If there is no path 3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
                                       problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
from the starting vertex to the ending vertex, a-1 is returned.
                                      negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
int bfs(int start, int end, vvi& adj_list) {------// d7
----set<<u>int</u>> visited;-----// b3
----queue<ii> 0;------γ bb int* bellman_ford(int n, int s, vii* adj, boolδ has_negative_cycle) {------// cf
-----// db ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
-----ii cur = Q.front(); Q.pop();------// 03 ------for (int j = 0; j < n; j++)------// c4
------if (dist[j] != INF)------// 4e
------if (cur.first == end)-------// 22 ------for (int k = 0; k < size(adj[j]); k++)------// 3f
-----return cur.second;------// b9 ------dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61
-----dist[j] + adj[j][k].second);------// 47
------vi& adj = adj_list[cur.first];-------// f9 ----for (int j = 0; j < n; j++)---------// 13
------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)------// 44 ------for (int k = θ; k < size(adj[j]); k++)--------// aθ
-----if (visited.find(*it) == visited.end()) {-------// 8d -----if (dist[j] + adj[j][k].second < dist[adj[j][k].first])-----// ef
------Q.push(ii(*it, cur.second + 1));-------// ab ------has_negative_cycle = true;---------// 2a
```

```
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----return dist;------// 2e ----return pair<union_find, vi>(uf, dag);-------// f2
3.5. Minimum Spanning Tree.
3.3. All-Pairs Shortest Paths.
                                 3.5.1. Kruskal's algorithm.
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                 #include "../data-structures/union_find.cpp"------5
problem in O(|V|^3) time.
                                  void floyd_warshall(int** arr, int n) {------// 21
                                 // n is the number of vertices-----// 18
----for (int k = 0; k < n; k++)------// 49
                                 // edges is a list of edges of the form (weight, (a, b))-----// c6
-----for (int i = 0; i < n; i++)-----// 21
                                 // the edges in the minimum spanning tree are returned on the same form-----// 4d
-----for (int j = 0; j < n; j++)-----// 77
                                 vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
------if (arr[i][k] != INF && arr[k][j] != INF)------// b1
                                 ----union_find uf(n);------// 04
-----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                                 ----sort(edges.begin(), edges.end());-----// 51
}-----// 86
                                 ----vector<pair<int, ii> > res;------// 71
                                 ----for (int i = 0; i < size(edges); i++)------// ce
3.4. Strongly Connected Components.
                                 -----if (uf.find(edges[i].second.first) !=-----// d5
                                 -----uf.find(edges[i].second.second)) {------// 8c
3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                 -----res.push_back(edges[i]);-----// d1
graph in O(|V| + |E|) time.
                                 -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
#include "../data-structures/union_find.cpp"-----------------------// 5e
                                 -----}-------// 5b
-----// 11
                                 ----return res:-----// 46
vector<br/>bool> visited:-----// 66
                                 vi order;-----// 9b
-----// a5
                                 3.6. Topological Sort.
void scc_dfs(const vvi &adj, int u) {-----// a1
                                 3.6.1. Modified Depth-First Search.
----int v; visited[u] = true;-----// e3
----for (int i = 0; i < size(adj[u]); i++)-------// c5 void tsort_dfs(int cur, char* color, const vviδ adj, stack<int>δ res,------// ca
------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);--------// 6e ------bool& has_cycle) {------------------------// a8
----order.push_back(u);------// 19 ----color[cur] = 1;------// 5b
}------// dc ----for (int i = 0; i < size(adj[cur]); i++) {-------// 96
----int n = size(adj), u, v;--------------------// bd ------tsort_dfs(nxt, color, adj, res, has_cycle);-------// 5b
----order.clear();-------// 22 ------else if (color[nxt] == 1)------// 53
----union_find uf(n);------// 6d ------has_cycle = true;------// c8
----vi dag;-------if (has_cycle) return;-------// 7e
-----rev[adj[i][j]].push_back(i);------// 77 ----res.push(cur);------// cb
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04 }------// 9e
------S.push(order[i]), dag.push_back(order[i]);--------// 40 ----char* color = new char[n];--------// b1
------for (int j = 0; j < size(adj[u]); j++)--------// 21 -----if (!color[i]) {------------------------------// d5
-----if (!visited[v = adj[u][j]]) S.push(v);------// e7 -----tsort_dfs(i, color, adj, S, has_cycle);------// 40
```

```
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----return res;-----// 07
                       3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
}-----// 1f
                       ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                       #define MAXN 5000-----// f7
                       int dist[MAXN+1], q[MAXN+1];-----// b8
#define MAXV 1000-----// 2f
vi adj[MAXV];------// ff struct bipartite_graph {------// 2b
ii start_end() {------// 30 ----bipartite_graph(int _N, int _M) : N(_N), M(_M),------// 8d
----int start = -1, end = -1, any = 0, c = 0;------// 74 -----L(new int[N]), R(new int[M]), adj(new vi[N]) {}------// cd
------else if (indeg[i] != outdeg[i]) return ii(-1,-1);-------// fa ------else dist(v) = INF;-------// b3
----}-----dist(-1) = INF;-------// 96
}-------foreach(u, adj[v]) if(dist(R[*u]) == INF)------// 95
---ii se = start_end();------// 45 -----}
------if (s.empty()) break;------// ee -----if(dist(R[*u]) == dist(v) + 1)------// 64
----}------return true;-------------// fa
}------// aa ------dist(v) = INF;------------// 72
                       ------return false;-----// 97
3.8. Bipartite Matching.
                       -----return true;------// c6
3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
                       ----}-----// f7
O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
                       ----void add_edge(int i, int j) { adj[i].push_back(j); }------// 11
graph, respectively.
                       ----int maximum_matching() {------// 2d
vi* adi:-----// cc
                       ------int matching = 0;-----// f5
bool* done;-----// b1
                       ------memset(L, -1, sizeof(int) * N);------// 8f
int* owner:-----// 26
                       -----memset(R, -1, sizeof(int) * M);-----// 39
int alternating_path(int left) {------// da
                       ------while(bfs()) for(int i = 0; i < N; ++i)------// 77
----if (done[left]) return 0;-------// 08
                       ------matching += L[i] == -1 && dfs(i);------// f1
----done[left] = true;-----// f2
                       -----return matching:-----// fc
----for (int i = 0; i < size(adj[left]); i++) {-------// 34
                       ----}-----// le
-----int right = adj[left][i];------// b6
                       1:----// d3
-----if (owner[right] == -1 || alternating_path(owner[right])) {------// d2
-----owner[right] = left; return 1;------------------------// 26 3.9. Maximum Flow.
```

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the maximum flow of a flow network.

```
int q[MAXV], d[MAXV];-----// e6
-----if (d[s] == -1) break;-----// a0
------while ((x = augment(s, t, INF)) != 0) f += x;------// a6
-----if (res) reset():-----// 21
-----return f:-----// b6
----}-----// 1b
}:-----// 3b
```

 $3.9.1.\ Dinic's\ algorithm.$ An implementation of Dinic's algorithm that runs in $O(|V|^2|E|)$. It computes $3.9.2.\ Edmonds\ Karp's\ algorithm.$ An implementation of Edmonds Karp's algorithm that runs in $O(|V||E|^2)$. It computes the maximum flow of a flow network.

```
#define MAXV 2000-----// ba
#define MAXV 2000-----// ba int q[MAXV], p[MAXV];-----// 7b
                     struct flow_network {-----// 5e
struct flow_network {------// 12 ----struct edge {-----// fc
----struct edge {------// le ------// le ------// cb
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// d3 ------memset(head = new int[n], -1, n << 2);------// 58
------head = new int[n], curh = new int[n];------// 6b ----void destroy() { delete[] head; }------// d5
-----memset(head, -1, n * sizeof(int));------// 56 ----void reset() { e = e_store; }------// 1b
----void destroy() { delete[] head; delete[] curh; }-------// f6 ------e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;-------// 4c
----void add_edge(int u, int v, int uv, int vu = 0) {-------// cd ----}------// ef
-----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// c9 ----int max_flow(int s, int t, bool res = true) {-------// 12
-----e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// 89 ------if (s == t) return 0;-------// d6
----}------e_store = e;--------// 9e
------return (e[i].cap -= ret, e[i^1].cap += ret, ret);------// ac -------while (l < r)------------------------// 2c
------return 0:-------for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)------// c6
----int max_flow(int s, int t, bool res = true) {------------------------------(d[v = e[i].v] == -1 || d[u] + 1 < d[v]))-------// 2f
------memset(d, -1, n * sizeof(int));------// a8 -----at = p[t], f += x;-------// 2d
------l = r = 0, d[q[r++] = t] = 0;-------// 0e -------while (at != -1)-------// cd
------if (e[i^1].cap > 0 && d[e[i].v] == -1)-------// 29 ------if (res) reset();--------// 3b
-----d[q[r++] = e[i].v] = d[v]+1;-----// bc
                     ----}------// 05
-----memcpy(curh, head, n * sizeof(int));------// 10 };------// 75
```

3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modified to find shortest path to augment each time (instead of just any path). It computes the maximum flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with minimum cost. Running time is $O(|V|^2|E|\log|V|)$.

```
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                                                                        12
#define MAXV 2000-------at = p[t], f += x;-------// 43
------for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];------// ff
----struct edge {-------// 9a ---}------// 11
------int v, cap, cost, nxt;--------// ad };------// ad
-----edge(int _v, int _cap, int _cost, int _nxt)------// ec
                                     3.11. All Pairs Maximum Flow.
-----: v(_v), cap(_cap), cost(_cost), nxt(_nxt) { }------// c4
----}:-----// ad
                                     3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----int n, ecnt, *head;------// 46
                                     structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the
----vector<edge> e, e_store;-----// 4b
                                     maximum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// dd
                                     #include "dinic.cpp"-----// 58
-----e.reserve(2 * (m == -1 ? n : m));------// e6
                                      -----// 25
-----memset(head = new int[n], -1, n << 2);------// 6c
                                     bool same[MAXV];-----// 59
----}------// f3
                                     pair<vii, vvi> construct_gh_tree(flow_network \&g) {------// 77
----void destroy() { delete[] head; }------// ac
                                     ----int n = g.n, v;------// 5d
----void reset() { e = e_store; }------// 88
                                     ----vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));-----// 49
----void add_edge(int u, int v, int cost, int uv, int vu=0) {------// b4
                                     ----for (int s = 1; s < n; s++) {------// 9e
-----e.push_back(edge(v, uv, cost, head[u])); head[u] = ecnt++;-----// 43
                                     ------int l = 0, r = 0;------// 9d
-----e.push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;------// 53
                                     -----par[s].second = g.max_flow(s, par[s].first, false);------// 38
----}------// 16
                                     -----memset(d, 0, n * sizeof(int));-----// 79
----ii min_cost_max_flow(int s, int t, bool res = true) {-------// 6d
                                     -----memset(same, 0, n * sizeof(int));-----// b0
-----if (s == t) return ii(0, 0);-----// 34
                                     -----d[q[r++] = s] = 1;------// 8c
-----e_store = e;------// 70
                                     ------while (l < r) {------// 45
-----memset(pot, 0, n << 2);------// 24
                                     -----same[v = q[l++]] = true;-----// c8
------int f = 0, c = 0, v;------// d4
                                     -----for (int i = g.head[v]; i != -1; i = g.e[i].nxt)-----// 33
------while (true) {------// 29
-----memset(d, -1, n << 2);-----// fd
                                     -----if (q.e[i].cap > 0 && d[q.e[i].v] == 0)------// 3f
                                     -----d[q[r++] = g.e[i].v] = 1;-----// f8
-----memset(p, -1, n << 2);-----// b7
                                     -----set<<u>int</u>, cmp> q;-----// d8
-----q.insert(s); d[s] = 0;-----// 1d
                                     ------for (int i = s + 1; i < n; i++)------// 68
------while (!q.empty()) {-----// 04
                                     -----if (par[i].first == par[s].first && same[i]) par[i].first = s;----// ea
                                     -----q.reset();------// 9a
------int u = *q.begin();-----// dd
                                     ----}-----// 1e
-----q.erase(q.begin());-----// 20
                                     ----for (int i = 0; i < n; i++) {-------// 2a
-----for (int i = head[u]; i != -1; i = e[i].nxt) {------// 02
                                     -----int mn = INF, cur = i;------// 19
-----if (e[i].cap == 0) continue;-----// 1c
                                     ------while (true) {------// 3a
-----int cd = d[u] + e[i].cost + pot[v] - pot[v] = e[i].v];
                                     -----cap[cur][i] = mn;-----// 63
-----if (d[v] == -1 || cd < d[v]) {------// d2
                                     -----if (cur == 0) break;-----// 35
-----if (q.find(v) != q.end()) q.erase(q.find(v));-----// e2
                                     -----mn = min(mn, par[cur].second), cur = par[cur].first;-----// 28
-----d[v] = cd; p[v] = i;------// f7
                                     -----q.insert(v);-----// 74
                                     ----return make_pair(par, cap);-----// 6b
                                      -----// 99
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {------// 16
-----if (p[t] == -1) break;-----// 09
                                     ---if (s == t) return 0;-----// d4
-----int x = INF, at = p[t];-----// e8
                                     ----int cur = INF, at = s;-----// 65
----while (gh.second[at][t] == -1)-----// ef
```

```
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                                                    13
}------node() { prefixes = words = 0; } };------// 42
                          public:----// 88
3.12. Heavy-Light Decomposition.
                          ----node* root:-----// a9
struct HLD {------// 25 ----template <class I>---------------------------------// 89
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), below(n) {--// 4f ------cur->prefixes++;------------------------// f1
-----vi tmp(n, ID); values = segment_tree(tmp); }------// a7 ------if (begin == end) { cur->words++; break; }------// db
------it = cur->children.find(head);-------// 77
-----csz(below[u][i]), sz[u] += sz[below[u][i]]; }------// 84 ------pair<T, node*> nw(head, new node());------// cd
------for (int i = 0; i < size(below[u]); i++)----------// a7 ----int countMatches(I begin, I end) {------------------------------// 7f
------if (best == -1 || sz[below[u][i]] > sz[best]) best = below[u][i];--// 19 ------node* cur = root;----------// 32
-------for (int i = 0; i < size(below[u]); i++)-----------// 7d -------if (begin == end) return cur->words;--------// a4
-----if (below[u][i] != best) part(curhead = below[u][i]); }------// 30 -----else {---------------------------------// 1e
----void build() { int u = curloc = 0;-------// 06 ------T head = *begin;-------// 5c
------it = cur->children.find(head):------// d9
-----u = size(uat) - 1, v = size(vat) - 1:--------// ad -----node* cur = root:-------------------------------// 95
-----res = (loc[uat[u]] < loc[vat[v]] ? uat[u] : vat[v]), u--, v--;----// 13 -------if (begin == end) return cur->prefixes;----------// f5
-----res = f(res, values.query(loc[head[u]], loc[u])), ------// 7c ------it = cur->children.find(head); -------// 43
-----u = parent[head[u]];-------// 4b -------if (it == cur->children.end()) return 0;------// 71
------return f(res, values.query(loc[v] + 1, loc[u])); }-------// 47 -------beqin++, cur = it->second; } } } };-------// 26
----int query(int u, int v) { int l = lca(u, v);------// 04
                          4.2. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
-----return f(query_upto(u, l), query_upto(v, l)); } };-----// 52
                          struct entry { ii nr: int p: }:-----// f9
                          bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77
           4. Strings
                          struct suffix_array {------// 87
4.1. Trie. A Trie class.
                          ----string s; int n; vvi P; vector<entry> L; vi idx;------// b6
class trie {------// 9a ----suffix_array(string _s) : s(_s), n(size(s)) {------// e5
private:------L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);------// 8a
----struct node {------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';------// 8d
```

```
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------P.push_back(vi(n)):--------// 30 -------if (!st) st = q_0:-------// e_7
------for (int i = 0; i < n; i++)--------// d5 -------s->fail = st->next[a->first];-------// 29
------L[L[i].p = i].nr = ii(P[stp - 1][i],-------// fc ------if (s->fail) {----------------------// 3b
-----i + cnt < n ? P[stp - 1][i + cnt] : -1);-------// e5 -------if (!s->out) s->out = s->fail->out;-------// 80
-----for (int i = 0; i < n; i++)-------// 85 ------out_node* out = s->out;------// bf
------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;------// fd ------out->next = s->fail->out;------// 65
};------cur = cur->fail;------// 21
                              -----if (!cur) cur = qo;-----// 2f
4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
                              -----cur = cur->next[*c];-----// 58
state machine from a set of keywords which can be used to search a string for any of the keywords.
                              -----if (!cur) cur = qo;-----// 3f
struct aho_corasick {-----// 78
                              -----for (out_node *out = cur->out; out; out = out->next)-----// eθ
----struct out_node {------// 3e
                              -----res.push_back(out->keyword);----------------------------// 0d
-----string keyword; out_node *next;------// f0
                              -----}-------------// fd
-----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
                              -----return res:-----// c1
----}:------// b9
                              ----struct go_node {------// 40
                              }:-----// 32
-----map<char, go_node*> next;------// 6b
                              4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
-----out_node *out; go_node *fail;-----// 3e
                              also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
-----go_node() { out = NULL; fail = NULL; }-----// Of
                              can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
----qo_node *qo;-----// b8
                              int* z_values(const string &s) {------// 4d
----aho_corasick(vector<string> keywords) {------// 4b
                              ----int n = size(s);-----// 97
-----qo = new qo_node();-----// 77
                              ----int* z = new int[n];-----// c4
------foreach(k, keywords) {-------// e4
                              ----int l = 0, r = 0;-----// 1c
-----qo_node *cur = qo;-----// 9d
                              ---z[0] = n;
-----foreach(c, *k)-----// 38
                              ----for (int i = 1; i < n; i++) {------// 7e
-----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d
                              ----z[i] = 0;-----// c9
-----(cur->next[*c] = new qo_node());-----// 75
-----queue<go_node*> q;------// 8a -------while (r < n && s[r - l] == s[r]) r++;-----// ff
------foreach(a, go->next) g.push(a->second);-------// a3 -----z[i] = r - l; r--;-------// fc
-----go_node *r = q.front(); q.pop();------// 2e -----else {-------
------foreach(a, r->next) {-------// 02
------go_node *s = a->second;------// cb -------while (r < n && s[r - l] == s[r]) r++;------// b3
------q.push(s);------// 76 -----z[i] = r - l; r--; } }------// 8d
-----go_node *st = r->fail;-----// fa ----return z;------// 53
------while (st && st->next.find(a->first) == st->next.end())-----// d7 }------// db
```

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                                                               16
------if (first) outs << n.data[i], first = false;------// 33 ----intx operator *(const intx& b) const {-------// bd
------else {------------------------// 1f ------intx c; c.data.assign(size() + b.size() + 1, 0);-------// d0
------unsigned int cur = n.data[i];-------// 0f ------for (int i = 0; i < size(); i++) {-------// 7a
-----stringstream ss; ss << cur; ------// 8c ------long long carry = 0; -------// 20
-----string s = ss.str();------// 64 ------for (int j = 0; j < b.size() || carry; j++) {-------// cθ
------while (len < intx::dcnt) outs << '0', len++;------// 0a ------carry += c.data[i + j];------// 18
------outs << s;------% intx::radix;------// 97 -------c.data[i + j] = carry % intx::radix;------// 86
----}-----return c.normalize(sign * b.sign);------// de
------if (sign != b.sign) return sign < b.sign;--------// cf ------assert(!(d.size() == 1 && d.data[0] == 0));-------// e9
------if (size() != b.size())-------// 4d ------intx q, r; q.data.assign(n.size(), 0);------// ca
-----return sign == 1 ? size() < b.size() : size() > b.size();------// 4d ------for (int i = n.size() - 1; i >= 0; i--) {-------// 1a
------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35 ------r.data.insert(r.data.begin(), 0);--------// c7
------return false;-------// ca -------long long k = θ;--------// cc
-----if (sign > 0 && b.sign < 0) return *this - (-b);---------// 36 -------r = r - abs(d) * k;-------------------------// 15
------<mark>unsigned long long carry = 0;-------// 5c ------return pair</mark><intx, intx>(q.normalize(n.sign * d.sign), r);-----// a1
-----carry += (i < size() ? data[i] : OULL) +-------// 91 ----intx operator /(const intx& d) const {-------// a2
-----(i < b.size() ? b.data[i] : 0ULL);-------// 0c -----return divmod(*this,d).first; }------// 1e
-----carry /= intx::radix;------// fd -----return divmod(*this,d).second * sign; }------// 5a
-----return c.normalize(sign);------// 20
                                5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
----}------------// 70
                                #include "intx.cpp"-----// 83
----intx operator -(const intx& b) const {-------// 53
                                #include "fft.cpp"-----// 13
-----if (sign > 0 && b.sign < 0) return *this + (-b);------// 8f
                                -----// e0
-----if (sign < 0 && b.sign > 0) return -(-*this + b);------// 1b
                                intx fastmul(const intx &an, const intx &bn) {------// ab
------if (sign < 0 && b.sign < 0) return (-b) - (-*this);-----// a1
                                ----string as = an.to_string(), bs = bn.to_string();-----// 32
-----if (*this < b) return -(b - *this);------// 36
                                ----int n = size(as), m = size(bs), l = 1,-----// dc
-----intx c; c.data.clear();-----// 6b
                                -----len = 5, radix = 100000,-----// 4f
-----long long borrow = 0;-----// f8
                                -----*a = new int[n], alen = 0,-----// b8
------for (int i = 0; i < size(); i++) {-------// a7
                                -----*b = new int[m], blen = 0;------// 0a
------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);-----// a9
                                ----memset(a, 0, n << 2);-----// 1d
-----c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
                                ----memset(b, 0, m << 2);-----// 01
-----borrow = borrow < 0 ? 1 : 0;-----// 0d
                                ----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
-----}-----// fa
                                ------for (int j = min(len - 1, i); j >= 0; j--)------// 43
```

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----for (int i = m - 1; i >= 0; i -= len, blen++)-----// b6
                                                }-----// 40
------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
------b[blen] = b[blen] * 10 + bs[i - j] - 0;------// 9b
                                                5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is
----while (l < 2*max(alen,blen)) l <<= 1;------// 51
----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
                                                bool is_prime(int n) {------// 6c
----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);-----// 35
                                                ----if (n < 2) return false;-----// c9
----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? b[i] : 0, 0);-----// 66
                                                ----if (n < 4) return true;-----// d9
----fft(A, l); fft(B, l);-----// f9
----for (int i = 0; i < l; i++) A[i] *= B[i];------// e7
                                                ----if (n % 2 == 0 || n % 3 == 0) return false;-------// 0f
----fft(A, l, true);-----// d3
                                                ----if (n < 25) return true;------// ef
----ull *data = new ull[1];-----// e7
                                                ----int s = static_cast<int>(sqrt(static_cast<double>(n)));------// 64
                                                ----for (int i = 5; i <= s; i += 6)------// 6c
----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
                                                -----if (n % i == 0 || n % (i + 2) == 0) return false;-----// e9
----for (int i = 0; i < l - 1; i++)------// 90
------if (data[i] >= (unsigned int)(radix)) {------// 44
                                                ----return true; }------// 43
-----data[i+1] += data[i] / radix;-----// e4
-----data[i] %= radix;-----// bd
                                                5.6. Miller-Rabin Primality Test. The Miller-Rabin probabilistic primality test.
                                                #include "mod_pow.cpp"-----// c7
----int stop = l-1;-----// cb
                                                bool is_probable_prime(ll n, int k) {------// be
----while (stop > 0 && data[stop] == 0) stop--;------// 97
                                                ----if (~n & 1) return n == 2;------// d1
----stringstream ss;-----// 42
                                                ----if (n <= 3) return n == 3;-----// 39
----ss << data[stop];------// 96
                                                ----int s = 0; ll d = n - 1;------// 37
----for (int i = stop - 1; i >= 0; i--)-----// bd
                                                ----while (~d & 1) d >>= 1, s++;------// 35
-----ss << setfil('0') << setw(len) << data[i];------// b6
                                                ----while (k--) {-------// c8
----delete[] A; delete[] B;-----//
                                                 -----ll a = (n - 3) * rand() / RAND_MAX + 2;------// 06
----delete[] a; delete[] b;-----// 7e
                                                 -----ll x = mod_pow(a, d, n);------// 64
----delete[] data;------// 6a
                                                -----if (x == 1 || x == n - 1) continue;------// 9b
----return intx(ss.str());-----// 38
                                                ------bool ok = false;-----// 03
                                                 ------for (int i = 0; i < s - 1; i++) {-------// 6b
5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
                                                -----if (x == 1) return false;-----// 4f
k items out of a total of n items.
                                                 ------if (x == n - 1) { ok = true; break; }-----// 74
int nck(int n, int k) {-----// f6
                                                 ------------------------------// a9
----if (n - k < k) k = n - k;------// 18
                                                 -----if (!ok) return false;-----// 00
----int res = 1;-----// cb
                                                 ----} return true; }-------// bc
----for (int i = 1; i \le k; i++) res = res * (n - (k - i)) / i;-----// bd
----return res;------// e4
}-----// 03
                                                5.7. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
                                                vi prime_sieve(int n) {-----// 40
5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
                                                 ----int mx = (n - 3) >> 1, sq, v, i = -1;-------// 27
integers a, b.
                                                 ----vi primes:-----// 8f
int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }-----// d9
                                                 ----bool* prime = new bool[mx + 1];------// ef
                                                ----memset(prime, 1, mx + 1);-----// 28
 The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
and also finds two integers x, y such that a \times x + b \times y = d.
                                                ----if (n >= 2) primes.push_back(2);-----// f4
                                                ----while (++i <= mx) if (prime[i]) {------// 73
int egcd(int a, int b, int& x, int& y) {------// 85
----if (b == 0) { x = 1; y = 0; return a; }------// 7b
                                                 -----primes.push_back(v = (i << 1) + 3);-----// be
                                                 -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
----else {------// 00
------int d = egcd(b, a % b, x, y);-----// 34
                                                ------for (int j = sq; j <= mx; j += v) prime[j] = false; }------// 2e
------x = a / b * y;
                                                ----while (++i \le mx) if (prime[i]) primes.push_back((i << 1) + 3);-----// 29
                                                 ----delete[] prime; // can be used for O(1) lookup-----// 36
-----swap(x, y);-----// 26
                                                ----return primes; }-----// 72
-----return d;------// db
```

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 $5.9. \ \, \textbf{Modular Exponentiation.} \ \, \textbf{A function to perform fast modular exponentiation.}$

5.10. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.

5.11. Linear Congruence Solver. A function that returns all solutions to $ax \equiv b \pmod{n}$, modulo

5.12. Numeric Integration. Numeric integration using Simpson's rule.

5.13. **Fast Fourier Transform.** The Cooley-Tukey algorithm for quickly computing the discrete Fourier transform. Note that this implementation only handles powers of two, make sure to pad with zeros.

#include <complex>-----// 8e

```
typedef complex<long double> cpx;------// 25
void fft(cpx *x, int n, bool inv=false) {------// 23
----for (int i = 0, j = 0; i < n; i++) {-------// f2
-----if (i < j) swap(x[i], x[j]);-----// 5c
------int m = n>>1;------// e5
------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
-----i += m:------// ab
----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
-----cpx wp = \exp(cpx(0, (inv ? -1 : 1) * pi / mx)), w = 1;
-----for (int m = 0; m < mx; m++, w *= wp) \{------// 4\theta
-----for (int i = m: i < n: i += mx << 1) {------// 33
-----cpx t = x[i + mx] * w;
-----x[i] += t;-----// c7
-----}-----// c2
----if (inv) for (int i = 0; i < n; i++) x[i] /= cpx(n);------// 3e
```

5.14. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once: $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times: n^k
- Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type 2, ..., n_k objects of type k: $\binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$$

- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times: $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to $x_1 + x_2 + \cdots + x_n = k$ where $x_i \ge 0$: f_k^n
- Number of subsets of a set with n elements: 2^n
- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an $n \times m$ grid by walking only up and to the right: $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced: $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an $n \times n$ lattice which do not rise above the main diagonal: C_n
- Number of permutations of n objects with exactly k ascending sequences or runs: $\binom{n}{k} = \binom{n}{n-k-1} = k \binom{n-1}{k} + (n-k+1) \binom{n-1}{k-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{n-1} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{0} = \sum_{i=0}^{k} (-1)^i \binom{n+1}{i} (k+1-i)^n, \binom{n}{0} = \binom{n}{0}$

Reykjavík University double progress(P(p), L(a, b)) {-----// 8e • Number of permutations of n objects with exactly k cycles: $\binom{n}{k} = \binom{n-1}{k-1} + (n-1) \binom{n-1}{k}$ ----if (abs(real(a) - real(b)) < EPS)------// bc • Number of ways to partition n objects into k sets: $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$ -----return (imag(p) - imag(a)) / (imag(b) - imag(a));------// 36 • Number of permutations of length n that have no fixed points (derangements): $D_0 = 1, D_1 =$ ----**else return** (real(p) - real(a)) / (real(b) - real(a)); }-----// 58 $0, D_n = (n-1)(D_{n-1} + D_{n-2})$ **bool** intersect(L(a, b), L(p, q), point &res, **bool** segment = false) $\{-----//d4\}$ • Number of permutations of length n that have exactly k fixed points: $\binom{n}{k} D_{n-k}$ ----// NOTE: check for parallel/collinear lines before calling this function---// 79 • Jacobi symbol: $\left(\frac{a}{\hbar}\right) = a^{(b-1)/2} \pmod{b}$ ----point r = b - a, s = q - p; -----// θb ----**double** c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c; --// de• Heron's formula: A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where ----**if** (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// d7-----return false:-----// 00 • Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice ----res = a + t * r:-----// c9 points on the boundary has area $i + \frac{b}{2} - 1$. ----**return** true:-----// e7 • Divisor sigma: The sum of divisors of n to the xth power is $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$ where point closest_point(L(a, b), P(c), bool segment = false) {------// 30 $n = \prod_{i=0}^{r} p_i^{a_i}$ is the prime factorization. ----if (segment) {-------// 8f • Divisor count: A special case of the above is $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$. -----if (dot(b - a, c - b) > 0) return b;------// 83 • Euler's totient: The number of integers less than n that are comprime to n are $n \prod_{p|n} \left(1 - \frac{1}{n}\right)$ -----if (dot(a - b, c - a) > 0) return a;-----// d4 where each p is a distinct prime factor of n. -----// 92 • König's theorem: In any bipartite graph, the number of edges in a maximum matching is ----**double** t = dot(c - a, b - a) / norm(b - a);-----// 22 equal to the number of vertices in a minimum vertex cover. ----return a + t * (b - a);-----// d7 • The number of vertices of a graph is equal to its minimum vertex cover number plus the size }-----// 20 of a maximum independent set. double line_segment_distance(L(a,b), L(c,d)) {-----// da • $\gcd(2^a - 1, 2^b - 1) = 2^{\gcd(a,b)} - 1$ ----double x = INFINITY;-----// 04 ----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);------// 17 5.15. Numbers and Sequences. Some random prime numbers: 1031, 32771, 1048583, 33554467, ----else if (abs(a - b) < EPS) $x = abs(a - closest_point(c, d, a, true));$ -----// d9 $1073741827,\ 34359738421,\ 1099511627791,\ 35184372088891,\ 1125899906842679,\ 36028797018963971.$ ----else if $(abs(c - d) < EPS) \times = abs(c - closest_point(a, b, c, true)); -----// 7f$ ----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// c26. Geometry ----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;-----// 93 6.1. **Primitives.** Geometry primitives. ----else {------// 90 #include <complex>-----// 8e -----x = min(x, abs(a - closest_point(c,d, a, true)));-----// b1 #define P(p) const point &p-----// b8 $-----x = min(x, abs(b - closest_point(c,d, b, true)));$ -----x = min(x, abs(c - closest_point(a,b, c, true)));------// 45 -----x = min(x, abs(d - closest_point(a,b, d, true)));-----// cd ----}-------// *30*

```
#define L(p0, p1) P(p0), P(p1)-----// 30
typedef complex<double> point;-----// e1
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// a9
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// ff
point rotate(P(p), P(about), double radians) {-----// dc
----return (p - about) * exp(point(0, radians)) + about; }-----// cb
```

```
point reflect(P(p), L(about1, about2)) {-----// c0
double angle(P(p)) { return atan2(imag(p), real(p)); }------// fc -----if (collinear(p[i], q, p[i]) &⟨---------------// a5
point perp(P(p)) { return point(-imag(p), real(p)); }------// 79 ------0 <= (d = progress(q, p[i], p[j])) && d <= 1)------// b9
```

```
----return x:-----// 9e
                                        6.2. Polygon. Polygon primitives.
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// fc double polygon_area_signed(polygon p) {------// 31
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }------// ca ----for (int i = 1; i + 1 < cnt; i++)-------// d2
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// 75 -----area += cross(p[i] - p[0], p[i + 1] - p[0]);---------// 7e
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6 double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25
double angle(P(a), P(b), P(c)) {-------// d0 #define CHK(f,a,b,c) (f(a) < f(b) & f(b) <= f(c) & ccw(a,c,b) < 0}-----// b2
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }-------// cc int point_in_polygon(polygon p, point q) {-----------// 58
```

```
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----return in ? -1 : 1; }-----// 77 ------return 0.0 <= p && p <= 1.0-----// 8e
// pair<polygon, polygon cut_polygon(const polygon &poly, point a, point b) {-// 7b -------& (abs(c - a) + abs(b - c) - abs(b - a)) < EPS; }-------// 4f
//--- polygon left, right;-----// 6b ----else if (collinear(a,b, c,d)) {-------------// bc
//---- point it(-100, -100);------// c9 -------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
//------ int i = i = cnt-1 ? 0: i+1;-------// 8e -------if (bp < 0.0 \mid | ap > 1.0) return false;------// 0c
//------ point p = poly[i], q = poly[i];------// 19 ------A = c + max(ap, 0.0) * (d - c);------// f6
//------ if (ccw(a, b, p) \le 0) left.push_back(p);------// 12 ------B = c + min(bp, 1.0) * (d - c);-------// 5c
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// e3 ------return true; }-----
//-----// myintersect = intersect where-----// 24 ----else if (parallel(a,b, c,d)) return false;-----// ca
//-----// (a,b) is a line, (p,q) is a line segment------// f2 ----else if (intersect(a,b, c,d, A, true)) {-----------------------// 10
//----- if (myintersect(a, b, p, q, it))------// f0 ------B = A; return true; }------// bf
// }-----// 37
                                    6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
                                    coordinates) on a sphere of radius r.
#define MAXN 1000------ double qLat, double r) {-------// a4
point hull[MAXN];------// 43 ----pLat *= pi / 180; pLong *= pi / 180;-------// ee
bool cmp(const point &a, const point &b) {-------// 32 ----qLat *= pi / 180; qLong *= pi / 180;------// 75
----return abs(real(a) - real(b)) > EPS ?------// 44 ----return r * acos(cos(pLat) * cos(pLong - qLong) +-----// e3
-----real(a) < real(b) : imag(a) < imag(b); }------// 40 -----sin(pLat) * sin(qLat));-----// 1e
----sort(p.begin(), p.end(), cmp);-----// 3d
----for (int i = 0; i < n; i++) {------// 6f
                                    6.6. Triangle Circumcenter. Returns the unique point that is the same distance from all three
-----if (i > 0 && p[i] == p[i - 1]) continue;-----// b2
                                    points. It is also the center of the unique circle that goes through all three points.
                                    #include "primitives.cpp"-----// e0
------while (l >= 2 \&\& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;------// 20
-----hull[l++] = p[i];-----// f7
                                    point circumcenter(point a. point b. point c) {-----// 76
----}-----// d8
                                    ----b -= a, c -= a;-----// 41
----int r = 1:------// 59
                                    }-----// c3
----for (int i = n - 2; i >= 0; i--) {------// 16
-----if (p[i] == p[i + 1]) continue;-----// c7
                                    6.7. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
------while (r - l >= 1 \& \& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 9f
                                    pair of points.
-----hull[r++] = p[i]:-----// 6d
                                    #include "primitives.cpp"-----// e0
----return l == 1 ? 1 : r - 1;------// 6d
                                    -----// 85
                                    struct cmpx { bool operator ()(const point &a, const point &b) {------// 01
}-----// 79
                                     ------return abs(real(a) - real(b)) > EPS ?-----// e9
6.4. Line Segment Intersection. Computes the intersection between two line segments.
                                    -----real(a) < real(b) : imag(a) < imag(b); } };-----// 53
-------A = B = a; return abs(a - d) < EPS; }--------// ee double closest_pair(vector<point> pts) {--------// f1
----else if (abs(a - b) < EPS) {-------// 03 ----sort(pts.begin(), pts.end(), cmpx());------// 0c
------A = B = a; double p = progress(a, c,d);--------// c9 ----set<point, cmpy> cur;------------------------// bd
```

```
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----set<point, cmpy>::const_iterator it, jt;-------// a6 ----return hi;---------------------------// fa
----double mn = INFINITY;-----// f9
                                                                       }-----// 66
----for (int i = 0, l = 0; i < size(pts); i++) {------// ac
------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);------// 8b
                                                                       7.3. 2SAT. A fast 2SAT solver.
-----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));------// fc
                                                                       #include "../graph/scc.cpp"-----// c3
-----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));------// 39
                                                                        -----// 63
-------while (it != jt) mn = min(mn, abs(*it - pts[i])), it++;------// 09
                                                                       bool two_sat(int n, const vii& clauses, vi& all_truthy) {------// f4
-----cur.insert(pts[i]); }-----// 82
                                                                        ----all_truthy.clear();-----// 31
----return mn; }-----// 4c
                                                                        ----vvi adj(2*n+1);------// 7b
                                                                        ----for (int i = 0; i < size(clauses); i++) {-------// 9b
6.8. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
                                                                        -----adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// 17
    • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
                                                                        ------if (clauses[i].first != clauses[i].second)------// 87
    • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
                                                                        -----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93
    • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
                                                                        ----}--------// d8
      of that is the area of the triangle formed by a and b.
                                                                        ----pair<union_find, vi> res = scc(adj);------// 9f
                                                                        ----union_find scc = res.first;------// 42
                          7. Other Algorithms
                                                                        ----vi dag = res.second;-----// 58
                                                                        ----vi truth(2*n+1, -1);-----// 00
7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
                                                                        ----for (int i = 2*n; i >= 0; i--) {------// f4
function f on the interval [a, b], with a maximum error of \varepsilon.
                                                                        -----int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);-// 5a
double binary_search_continuous(double low, double high,-----// 8e
                                                                        -----if (cur == 0) continue;-----// 26
-----double eps, double (*f)(double)) {-----// c0
                                                                        -----if (p == o) return false;-----// 33
----while (true) {------// 3a
                                                                        -----if (truth[p] == -1) truth[p] = 1;------// c3
------double mid = (low + high) / 2, cur = f(mid);-----// 75
                                                                        -----truth[cur + n] = truth[p];-----// b3
-----if (abs(cur) < eps) return mid;-----// 76
                                                                        -----truth[o] = 1 - truth[p];-----// 80
-----else if (0 < cur) high = mid;-----// e5
                                                                        ------if (truth[p] == 1) all_truthy.push_back(cur);-------// 5c
-----else low = mid;-----// a7
                                                                        ----}----------// d9
----}------// b5
                                                                        ----return true;------// eb
}-----// cb
                                                                       }-----// 61
  Another implementation that takes a binary predicate f, and finds an integer value x on the integer
interval [a, b] such that f(x) \wedge \neg f(x-1).
                                                                       7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
int binary_search_discrete(int low, int high, bool (*f)(int)) {------// 51
---assert(low <= high);-----// 19
                                                                       vi stable_marriage(int n, int** m, int** w) {------// e4
----while (low < high) {------// a3
                                                                       ----queue<int> q;------// f6
------int mid = low + (high - low) / 2;-------// 04 ----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
------if (f(mid)) high = mid;-------// ca ----for (int i = 0; i < n; i++) for (int j = 0; j < n; j++)------// 05
-----else low = mid + 1;------// 03 -----inv[i][w[i][j]] = j;-----// b9
----assert(f(low));-------// 42 ----while (!q.empty()) {-------// 55
----return low;------// a6 ------<mark>int</mark> curm = q.front(); q.pop();------// ab
}-------for (int &i = at[curm]; i < n; i++) {-------// 9a
                                                                        -----int curw = m[curm][i];-----// cf
7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotonically increasing and the monotonical and the monotonica
cally decreasing, ternary search finds the x such that f(x) is maximized.
                                                                       ------else if (inv[curw][curm] < inv[curw][eng[curw]])------// 10
template <class F>------q.push(enq[curw]);------// 8c
                                                                       -----else continue;-----// b4
double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
                                                                       -----res[eng[curw] = curm] = curw, ++i; break;------// 5e
----while (hi - lo > eps) {------// 3e
                                                                       -----}-------// 24
------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
                                                                       ----}------// b8
-----if (f(m1) < f(m2)) lo = m1;------// 1d
                                                                       ----return res:-----// 95
-----else hi = m2;-----// b3
                                                                       }-----// 03
----}------// bb
```

```
------else ptr[i][j] = NULL;------// 32 -----return handle_solution(res);-----// 51
------for (int i = 0; i <= rows; i++) {--------// 84 ------node *c = head->r, *tmp = head->r; ------// 8e
------for (int j = 0; j < cols; j++) {---------for (; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 00
-----if (!ptr[i][j]) continue;------// 35 -----if (c == c->d) return false;------// b0
-----/<sub>int</sub> ni = i + 1, nj = j + 1;------// b7 -----COVER(c, i, j);-----// 7a
------while (true) {------// b0 -----<mark>bool</mark> found = false;-----// 7f
-----ptr[ni][j]->u = ptr[i][j];-----// c4 ____}
-----// c6 -----// c6 -----// 3a
------if (i == rows || arr[i][nj]) break;------// 8d ---}
------++ni;------// 1c }
-----ptr[i][nj]->l = ptr[i][j];-----// 72 1}.
------head = new node(rows, -1);-------// 80 ----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
------head->r = ptr[rows][0];-------// 73 ----for (int i = 1; i <= cnt; i++) fac[i - 1] = n % i, n /= i;-----// 04
-----ptr[rows][0]->l = head;------// 3b ----for (int i = cnt - 1; i >= 0; i--)------// 52
```

```
------per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// 41
----return per:-----// 84
}-----// 97
```

7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.

```
ii find_cycle(int x0, int (*f)(int)) {------// a5
----int t = f(x0), h = f(t), mu = 0, lam = 1;------// 8d
----while (t != h) t = f(t), h = f(f(h));------// 79
----h = xθ:-----// 04
----while (t != h) t = f(t), h = f(h), mu++;------// 9d
----h = f(t):-----// 00
----while (t != h) h = f(h), lam++;-----// 5e
----return ii(mu, lam);-----// b4
}-----// 42
7.8. Dates. Functions to simplify date calculations.
```

```
int intToDay(int jd) { return jd % 7; }-----// 89
---x = (146097 * n + 3) / 4;
----i = (4000 * (x + 1)) / 1461001;-----// 0d
----x -= 1461 * i / 4 - 31;-----// 09
---i = 80 * x / 2447:
---d = x - 2447 * j / 80;
----x = i / 11:-----// b7
---m = i + 2 - 12 * x;
---y = 100 * (n - 49) + i + x;
}-----// af
```

8. Useful Information

8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.

- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo m_1, m_2, \ldots, m_k , where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When $n = 0, n = -1, n = 1, n = 2^{31} 1$ or $n = -2^{31}$? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

```
void readn(register int *n) {------// dc
                      ----int sign = 1;------// 32
                      ----register char C:-----// a5
                      ----*n = 0:-----// 35
                      ----while((c = getc_unlocked(stdin)) != '\n') {------// f3
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +------// a8 ------case '-': sign = -1; break;------// 28
------367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1 ------case ' ': goto hell;------// fd
-----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be ------case '\n': goto hell;-------// 79
-----d - 32075:------default: *n *= 10; *n += c - '0'; break;------// c0
```

8.3. 128-bit Integer. GCC has a 128-bit integer data type named __int128. Useful if doing multiplication of 64-bit integers, or something needing a little more than 64-bits to represent.

8.4. Worst Time Complexity.

n	Worst AC Algorithm	Comment	
≤ 10	$O(n!), O(n^6)$	e.g. Enumerating a permutation	
≤ 15	$O(2^n \times n^2)$	$(n \times n^2)$ e.g. DP TSP	
≤ 20	$O(2^n), O(n^5)$	e.g. DP + bitmask technique	
≤ 50	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$	
$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's	
$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort	
$\le 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree	
$< 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)	

8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.

```
----int y = x \& -x, z = x + y;------
----return z | ((x ^ z) >> 2) / y;------
}-----
```