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```
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            1. Code Templates
                               ----public static void main(String[] args) throws Exception {--------// 02
                               -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                               ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                               -----// code-----// e6
setxkbmap -option caps:escape
                               -----out.flush():-----// 56
set -o vi
                               xset r rate 150 100
                               }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                           2. Data Structures
syn on | colorscheme slate
                               2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                               struct union find {-----// 42
#include <cassert>------------------// 65 ----int find(int x) { return p[x] < 0 ? x : p[x] = find(p[x]); }-------// ba
#include <cmath>-----// 7d ----bool unite(int x, int y) {-------// 6c
#include <cstdio>------// 2e ------int xp = find(x), yp = find(y);------// 64
#include <cstdlib>------// 11 ------if (xp == yp) return false;------// 0b
#include <cstring>------// d0 -------if (p[xp] < p[yp]) p[xp] += p[yp], p[yp] = xp;------// 3a
#include <ctime>------// 28 -------| p[xp] += p[xp], p[xp] = yp;------// 3e
#include <map>-----// 02
#include <set>------// d2
#include <sstream>------ min(a, b); }------// 18 // int f(int a, int b) { return min(a, b); }------
#include <stack>------// cf const int ID = 0;------// 82
#include <string>------// a9 int f(int a, int b) { return a + b; }------// 6d
#include <vector>-----// 4f ----int n; vi data, lazy;-----// fa
----/ 7e ----seqment_tree(const vi &arr) : n(size(arr)), data(4*n), lazy(4*n,INF) {-----// 96
----for (typeof((o).begin()) u = (o).begin(); u != (o).end(); ++u)------// 1a ----int mk(const vi &arr, int l, int r, int i) {--------// 75
const int INF = 2147483647;------// be -----if (l == r) return data[i] = arr[l];------// 7c
const double EPS = 1e-9;------// 0f ------int m = (l + r) / 2;-------// e9
typedef long long ll;------// 8f ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// c2
typedef unsigned long long ull;-----// 81 ----int q(int a, int b, int l, int r, int i) {------// 08
template <class T> T mod(T a, T b) { return (a % b + b) % b; }------// 70 ----void update(int i, int v) { u(i, v, 0, n-1, 0); }-------// f1
template <class T> int size(const T &x) { return x.size(); }------// 68 ----int u(int i, int v, int l, int r, int j) {-------// 77
                               -----propagate(l, r, j);-----// θc
1.3. Java Template. A Java template.
                               -----if (r < i || i < l) return data[j];------// cc
import java.io.*;-----return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 96
-----// a3 ----void range_update(int a, int b, int v) { ru(a, b, v, 0, n-1, 0); }------// 65
```

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------propagate(l, r, i);-------// ee template <class T>------// 53
------if (r < a || b < l) return data[i];--------// cc public:--------//
------/ 6b ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// fe
-----ru(a, b, v, m+1, r, 2*i+2));------// 2d -----cnt(other.cnt), data(other.data) { }------// ed
------if (l > r || lazy[i] == INF) return; -------// 08 ------for (int i = 0; i < cnt; i++) data[i] += other.data[i]; }------// e5
------data[i] += lazy[i] * (r - l + 1);-------// 5c ----void operator -=(const matrix& other) {-------// 68
------if (l < r) {-------// f2 ------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
------else lazv[2*i+1] += lazv[i];------// a8 ------for (int i = 0; i < cnt; i++) data[i] *= other; }-----// 40
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];------// 3c ----matrix<T> operator +(const matrix& other) {------// ee
------else lazy[2*i+2] += lazy[i];-------// bb ------matrix<T> res(*this); res += other; return res; }------// 5d
------lazy[i] = INF;------res(*this); res -= other; return res; }------// cf
};-----// e6 -----matrix<T> res(*this); res *= other; return res; }------// 37
                                    ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                                   -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                                   ------for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i... in O(\log n) time. It only needs O(n) space.
                                    ------for (int k = 0; k < cols; k++)------// fc
struct fenwick_tree {------res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n; vi data;------// d3 -----return res; }------// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {---------------// dd
----void update(int at, int by) {-------// 76 -----matrix<T> res(cols, rows);-----// b5
-------while (at < n) data[at] += by, at |= at + 1; }-------// fb -------for (int i = 0; i < rows; i++)-------// 9c
------int res = 0;-------// c3 ------return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);-------// bf
};------while (p) {------// cb
struct fenwick_tree_sq {------// d4 ------if (p & 1) res = res * sq;-----// c1
----int n; fenwick_tree x1, x0;--------// 18 -----p >>= 1;-----------------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;-------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {-------// 89
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }-----// 45 ------matrix<T> mat(*this); det = T(1);------------------// 21
----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73 ------for (int r = 0, c = 0; c < cols; c++) {-------// c4
};-------int k = r;------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
------for (int i = 0; i < cols; i++)-----// ab
2.4. Matrix. A Matrix class.
                                   -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
template <> bool eq<double>(double a, double b) { return abs(a - b) < EPS; }---// a7
```

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------if (!eq<T>(mat(r, c), T(1)))-------// 2c ------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;------// 6b
------for (int i = cols-1; i >= c; i--) mat(r, i) /= mat(r, c);-----// 5d ------else if (n->l && n->r) {---------------------------// 6c
------for (int i = 0; i < rows; i++) {---------// 3d ------node *s = successor(n);-------// e5
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 5a
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;------// 43
-----if (!n) return NULL;------// 37
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                           ------if (n->r) return nth(0, n->r);------// 23
#define AVL_MULTISET 0------// b5 -----node *p = n->p;------// a7
template <class T>-----// 22 -----return p; }-----// c7
class avl_tree {------// ff ----node* predecessor(node *n) const {------// b4
public:-----// f6 ------if (!n) return NULL;-------// dd
----struct node {-------// 45 ------if (n->l) return nth(n->l->size-1, n->l);-------// 10
-----T item; node *p, *l, *r;-------// a6 -----node *p = n->p;-------// ea
------int size, height;-------// 33 -------while (p && p->l == n) n = p, p = p->p;------// 6d
-----node(const T & item, node *_p = NULL) : item(_item), p(_p),------// 4f -----return p; }------
------node *cur = root;-------// b4 --------while (cur) {-------// 29
------while (cur) {-------// 8b ------if (n < sz(cur->l)) cur = cur->l;------// 75
------if (cur->item < item) cur = cur->r;------// 71 ------else if (n > sz(cur->l)) n -= sz(cur->l) + 1, cur = cur->r;-----// cd
------else if (item < cur->item) cur = cur->l;------// cd -----else break;------
------else break; }------// 4f ------} return cur; }------// ed
-----return cur; }------// 84 ----int count_less(node *cur) {------// ec
----node* insert(const T &item) {-------// 4e -----int sum = sz(cur->l);------// bf
------node *prev = NULL, **cur = &root;------// 60 -------while (cur) {-------// 6f
-----prev = *cur;------// f0 -----cur = cur->p;-----// eb
#if AVL_MULTISET-----// 0a private:----// d5
#else-----// ff ----inline int height(node *n) const { return n ? n->height : -1; }------// a6
-----else return *cur;------// 54 -----return n && height(n->l) > height(n->r); }------// a8
-----*cur = n, fix(n); return n; }------// 29 -----return n && abs(height(n->l) - height(n->r)) > 1; }-----// f8
----void erase(const T &item) { erase(find(item)); }------// 67 ----void delete_tree(node *n) {-------// fd
----void erase(node *n, bool free = true) {-------// 58 ------if (n) { delete_tree(n->r); delete n; } }-----// ef
-----if (!n) return;-----// 96
                           ----node*& parent_leg(node *n) {------// 6a
-----if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;------// 12
```

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------if (!n) return;--------// 0e ----int len, count, *q, *loc, tmp;-------// 0a
------n->size = 1 + sz(n->1) + sz(n->r);-------// 93 ----Compare _cmp;-------// 98
------while (i > 0) {------// 1a
-----parent_leg(n) = 1; \[\bar{N}\]------// fc
                              -----int p = (i - 1) / 2;-----// 77
-----n->l = l->r; \\ \[ \] ------// e8
                             ------if (!cmp(i, p)) break;-----// a9
-----augment(n), augment(l)-------// 81 ------while (true) {------------------------// 3c
----void left_rotate(node *n) { rotate(r, l); }------// 45 ------int l = 2*i + 1, r = l + 1;------// b4
------| else if (right_heavy(n) δδ left_heavy(n->r))------// b9 ----heap(int init_len = 128) : count(θ), len(init_len), _cmp(Compare()) {------// 17
------right_rotate(n->r);-------// 08 ------q = new int[len], loc = new int[len];-------// f8
-----if (left_heavy(n)) right_rotate(n);------// 93 ------memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// d5 ----~heap() { delete[] q; delete[] loc; }------// 09
-----n = n->p; }------// 28 ----void push(int n, bool fix = true) {------// b7
-----n = n->p; } } };-------// a2 ------if (len == count || n >= len) {-------// 0f
                              #ifdef RESIZE-----// a9
Also a very simple wrapper over the AVL tree that implements a map interface.
                              -----int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"-----// 01
                              ------while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                              ------int *newq = new int[newlen], *newloc = new int[newlen];------// e3
template <class K, class V>-----// da
                              -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --// 94
class avl_map {-----// 3f
                              -----/ 18 emset(newloc + len, 255, (newlen - len) << 2);------// 18
public:----// 5d
                              -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                              -----loc = newloc, q = newq, len = newlen;-----// 61
------K key; V value;------// 32
                              #else-----// 54
-----node(K k, V v) : key(k), value(v) { }-----// 29
                              -----assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                              #endif------// 64
----avl_tree<node> tree:-----// b1
                              ------}------// 4b
----V& operator [](K key) {------// 7c
                              -----assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(0)));-----// ba
                              -----loc[n] = count, q[count++] = n;-----// 6b
-----if (!n) n = tree.insert(node(key, V(0)));-----// cb
                              -----if (fix) swim(count-1); }-----// bf
-----return n->item.value:-----// ec
                              ----void pop(bool fix = true) {-------// 43
----}------// 2e
                              -----assert(count > 0);-----// eb
};-----// af
                              -----loc[q[0]] = -1, q[0] = q[-count], loc[q[0]] = 0;------// 50
                              -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                              #define RESIZE-----// d0
                             ----int top() { assert(count > 0); return q[0]; }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                             ----void heapify() { for (int i = count - 1; i > 0; i--)-----// 39
struct default_int_cmp {-----// 8d
```

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------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }-------// 0b ------if(lvl > current_level) current_level = lvl;-------// 8a
----void update_kev(int n) {-------------------------// 26 -----x = new node(lyl, target);-------------------// 36
----bool empty() { return count == 0; }-------// f8 ------x->next[i] = update[i]->next[i];------// 46
----int size() { return count; }--------------------------// 86 -------x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];-----------------// bc
----void clear() { count = 0, memset(loc, 255, len << 2); } };-------// 58 ------update[i] = x;------------// 20
                                        -----update[i]->lens[i] = pos[0] + 1 - pos[i];------// 42
2.7. Skiplist. An implementation of a skiplist.
                                        ------for(int i = lvl + 1; i <= MAX_LEVEL; i++) update[i]->lens[i]++;------// 07
#define BP 0.20-----// aa
#define MAX_LEVEL 10-----// 56
                                        -----return x; }-----// c9
unsigned int bernoulli(unsigned int MAX) {-----// 7b
                                        ----void erase(T target) {------// 4d
----unsigned int cnt = 0;-----// 28
                                        ------FIND_UPDATE(x->next[i]->item, target);-------// 6b
----while(((float) rand() / RAND_MAX) < BP \&\& cnt < MAX) cnt++;------// d1
                                        -----if(x && x->item == target) {------// 76
                                        ------for(int i = 0; i <= current_level; i++) {-------// 97
template<class T> struct skiplist {-----// 34
                                        -----if(update[i]->next[i] == x) {------// b1
----struct node {------// 53
-----T item:-----// e3
                                        -----update[i]->next[i] = x->next[i];------// 59
------int *lens;------// 07
                                        -----update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
                                        -----} else update[i]->lens[i] = update[i]->lens[i] - 1;-----// 88
-----node **next:-----// 0c
                                        -----#define CA(v, t) v((t*)calloc(level+1, sizeof(t)))------// 25
                                        -----delete x; _size--;-----// 81
------node(int level, T i) : item(i), CA(lens, int), CA(next, node*) {}-----// 7c
                                        ------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
-----/node() { free(lens); free(next); }; };-----// aa
                                        -----current_level--; } } };-----// 59
----int current_level, _size;------// 61
---node *head;-----// b7
                                        2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
                                        list supporting deletion and restoration of elements.
----~skiplist() { clear(); delete head; head = NULL; }------// aa
                                        template <class T>-----// 82
----#define FIND_UPDATE(cmp, target) \mathbb{N}------// c3
                                        struct dancing_links {-----// 9e
------int pos[MAX_LEVEL + 2]; N------// 18
                                        ----struct node {------// 62
-----memset(pos, 0, sizeof(pos)); \\\------// f2
                                        -----T item:-----// dd
-----node *update[MAX_LEVEL + 1]; N------// 01 -----node(const T &_item, node *_l = NULL, node *_r = NULL)-----// 6d
------memset(update, 0, MAX_LEVEL + 1); \sqrt{\phantom{a}}
                                        -----: item(_item), l(_l), r(_r) {------// 6d
                                        -----if (l) l->r = this;-----// 97
-----for(int i = MAX\_LEVEL; i >= 0; i--) { \sqrt{\phantom{a}}
                                        -----if (r) r->l = this;-----// 81
-----pos[i] = pos[i + 1]; \[\text{N}\]
                                        -----// 2d
----};------// d3
------update[i] = x; N-------// dd ----dancing_links() { front = back = NULL; }-----// 72
----void clear() { while(head->next && head->next[0])------// 91 -----if (!front) front = back;-----// d2
------erase(head->next[0]->item); }-------// e6 ------return back;---------------------------// εθ
------return x && x->item == target ? x : NULL; }-------// 50 ----node *push_front(const T &item) {-------// 4a
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ------if (!back) back = front;---------------------// 10
-----return pos[0]; }------// 19 -----return front;------// cf
------FIND_UPDATE(x->next[i]->item, target);--------// 3a ----void erase(node *n) {---------------------------// a0
------if(x && x->item == target) return x; // SET--------// 07 ------if (!n->l) front = n->r; else n->l->r = n->r;--------// ab
```

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------if (!n->l) front = n; else n->l->r = n; --------// a5 ------if (p.coord[i] < from.coord[i])------// a0
------if (!n->r) back = n; else n->r->l = n;--------// 9d -------sum += pow(from.coord[i] - p.coord[i], 2.0);------// 00
};-----sum += pow(p.coord[i] - to.coord[i], 2.0);------// 8c
                               ------}------------------------// be
2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                              -----return sqrt(sum); }-----// ef
element.
                               ------bb bound(double l, int c, bool left) {------// b6
#define BITS 15------pt nf(from.coord), nt(to.coord);------// 5c
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }------// b0 -----return bb(nf, nt); } };------// 3b
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49 -------pt p; node *l, *r;-------------------------------// 46
----int nth(int n) {-------// 8a ------node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
------int res = 0;-------// a4 ----node *root;------// 30
----}------if (from > to) return NULL;-------// f4
-----nth_element(pts.begin() + from, pts.begin() + mid,-----// f3
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                              -----pts.begin() + to + 1, cmp(c));------// 97
bor queries.
                               -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// dd
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1) - \cdots / 77
                              -----construct(pts, mid + 1, to, INC(c))); }-----// 03
template <int K>-----// cd ----bool contains(const pt &p) { return _con(p, root, 0); }------// f0
class kd_tree {------// 7e ----bool _con(const pt &p, node *n, int c) {------// 74
----struct pt {-------// 78 ------if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));-------// 81
-------double coord[K];-------// d6 ------if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// 95
-----pt() {}------// c1 -----return true; }------// 86
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }------// 4c ----void insert(const pt \&p) { _ins(p, root, 0); }------// f9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {-------// 09
------double sum = 0.0;------// c4 ------if (!n) n = new node(p, NULL, NULL);------// 4d
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }-----// 2c
-----return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }----// 73
----struct cmp {------// 8f ----void _clr(n->l), _clr(n->r), delete n; }-----// 1a
-------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// bf
------cc = i == 0 ? c : i - 1;------// bc ------pt from(cs);------
-----/return false; } };------// 62
----struct bb {-------// 30 ----pair<pt, bool> _nn(------// e3
------bb(pt _from, pt _to) : from(_from), to(_to) {}-------// 57 ------if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// 19
------double dist(const pt &p) {------// 3f
```

```
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------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;------// 9f -------visited.insert(*it);-------------------// cb
------if (found) mn = min(mn, p.dist(resp));-------// 18 ---}------// 0b
-----pair<pt, bool> res =-----// 33
                                          3.2. Single-Source Shortest Paths.
-----nn(p, n1, b.bound(n->p.coord[c], c, l1), mn, INC(c), same);---// 72
------if (res.second && (!found || p.dist(res.first) < p.dist(resp)))----// 76
                                           3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|)
-----resp = res.first, found = true;-----// 3b
                                          time.
-----}------------------------// aa
                                          int *dist, *dad;-----// 46
-----return make_pair(resp, found); } };-----// dd
                                          struct cmp {-----// a5
                                           ----bool operator()(int a, int b) {-----// bb
                  3. Graphs
                                           -----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
3.1. Breadth-First Search. An implementation of a breadth-first search that counts the number of
                                          };-----// 41
edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted
                                          pair<int*, int*> dijkstra(int n, int s, vii *adj) {-----// 53
graph (which is represented with an adjacency list). Note that it assumes that the two vertices are
                                           ----dist = new int[n];-----// 84
connected. It runs in O(|V| + |E|) time.
                                           ----dad = new int[n];-----// 05
int bfs(int start, int end, vvi& adj_list) {------// d7
                                           ----queue<ii>> Q;-----// 75
                                           ----set<<u>int</u>, cmp> pg;-----// 04
----Q.push(ii(start, 0));------// 49
                                           ----dist[s] = 0, pq.insert(s);-----// 1b
                                           ----while (!pq.empty()) {------// 57
----while (true) {------// 0a
                                           ------int cur = *pg.begin(); pg.erase(pg.begin()):-----// 7d
-----ii cur = Q.front(); Q.pop();-----// e8
                                           ------for (int i = 0; i < size(adj[cur]); i++) {------// 9e
-----// 06
                                           -----int nxt = adj[cur][i].first,-----// b8
-----if (cur.first == end)-----// 6f
                                           -----/ndist = dist[cur] + adj[cur][i].second;------// 0c
-----return cur.second:-----// 8a
                                           -----if (ndist < dist[nxt]) pq.erase(nxt),-----// e4
-----// 3c
                                           -----dist[nxt] = ndist, dad[nxt] = cur, pq.insert(nxt);------// \theta f
-----vi& adj = adj_list[cur.first];-----// 3f
                                           -----}-----// 75
------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)-----// bb
                                           ----}-----// e8
-----Q.push(ii(*it, cur.second + 1));-----// b7
                                           ----return pair<int*, int*>(dist, dad);-----// cc
----}------// 93
                                          }-----// af
}-----// 7d
                                          3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
 Another implementation that doesn't assume the two vertices are connected. If there is no path
                                          problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
from the starting vertex to the ending vertex, a-1 is returned.
                                          negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
int bfs(int start, int end, vvi& adi_list) {------// d7
----set<<u>int</u>> visited;------// b3
----queue<ii> 0;-------(int s, vii* adj, boolδ has_negative_cycle) {------// cf
-----// db ----for (int i = 0; i < n; i++) dist[i] = i == s ? 0 : INF;------// 10
-----ii cur = 0.front(); 0.pop();--------// 03 ------for (int j = 0; j < n; j++)-------// c4
  ------if (cur.first == end)-------// 22 ------for (int k = 0; k < size(adj[j]); k++)------// 3f
-----return cur.second;------// b9 ------dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61
-----/<sub>ba</sub> ------dist[i] + adj[j][k].second);------// 47
------vi& adj = adj_list[cur.first];-------// f9 ----for (int j = 0; j < n; j++)------------------------------// 13
------for (vi::iterator it = adj.beqin(); it != adj.end(); it++)-------// 44 ------for (int k = 0; k < size(adj[j]); k++)---------// a0
------if (visited.find(*it) == visited.end()) {------// 8d ------if (dist[j] + adj[j][k].second < dist[adj[j][k].first])-----// ef
------Q.push(ii(*it, cur.second + 1));-------// ab ------has_negative_cycle = true;---------// 2a
```

```
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----return dist;------// 2e ----return pair<union_find, vi>(uf, dag);-------// f2
3.5. Minimum Spanning Tree.
3.3. All-Pairs Shortest Paths.
                                3.5.1. Kruskal's algorithm.
3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths
                                #include "../data-structures/union_find.cpp"------5
problem in O(|V|^3) time.
                                 void floyd_warshall(int** arr, int n) {------// 21
                                // n is the number of vertices-----// 18
----for (int k = 0; k < n; k++)------// 49
                                // edges is a list of edges of the form (weight, (a, b))-----// c6
-----for (int i = 0; i < n; i++)-----// 21
                                // the edges in the minimum spanning tree are returned on the same form-----// 4d
-----for (int j = 0; j < n; j++)-----// 77
                                vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
------if (arr[i][k] != INF && arr[k][j] != INF)------// b1
                                ----union_find uf(n);------// 04
-----arr[i][j] = min(arr[i][j], arr[i][k] + arr[k][j]);------// e1
                                ----sort(edges.begin(), edges.end());-----// 51
}-----// 86
                                ----vector<pair<int, ii> > res;------// 71
                                ----for (int i = 0; i < size(edges); i++)------// ce
3.4. Strongly Connected Components.
                                -----if (uf.find(edges[i].second.first) !=-----// d5
                                -----uf.find(edges[i].second.second)) {------// 8c
3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                -----res.push_back(edges[i]);-----// d1
graph in O(|V| + |E|) time.
                                -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
#include "../data-structures/union_find.cpp"-----------------------// 5e
                                -----}-------// 5b
-----// 11
                                ----return res:-----// 46
vector<br/>bool> visited:-----// 66
                                vi order;-----// 9b
-----// a5
                                3.6. Topological Sort.
void scc_dfs(const vvi &adj, int u) {-----// a1
                                3.6.1. Modified Depth-First Search.
----int v; visited[u] = true;-----// e3
----for (int i = 0; i < size(adj[u]); i++)-------// c5 void tsort_dfs(int cur, char* color, const vviδ adj, stack<int>δ res,------// ca
------if (!visited[v = adj[u][i]]) scc_dfs(adj, v);--------// 6e ------bool& has_cycle) {------------------------// a8
----order.push_back(u);------// 19 ----color[cur] = 1;------// 5b
}------// dc ----for (int i = 0; i < size(adj[cur]); i++) {-------// 96
----int n = size(adj), u, v;--------------------// bd ------tsort_dfs(nxt, color, adj, res, has_cycle);-------// 5b
----order.clear();-------// 22 ------else if (color[nxt] == 1)------// 53
----union_find uf(n);------// 6d ------has_cycle = true;------// c8
----vi dag;-------if (has_cycle) return;-------// 7e
-----rev[adj[i][j]].push_back(i);-------// 77 ----res.push(cur);------// cb
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04 }------// 9e
------S.push(order[i]), dag.push_back(order[i]);--------// 40 ----char* color = new char[n];--------// b1
-----if (!visited[v = adj[u][j]]) S.push(v);------// e7 -----tsort_dfs(i, color, adj, S, has_cycle);------// 40
```

```
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----return res;------// 07
                       3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
}-----// 1f
                       ing. Running time is O(|E|\sqrt{|V|}).
3.7. Euler Path. Finds an euler path (or circuit) in a directed graph, or reports that none exist.
                       #define MAXN 5000-----// f7
#define MAXV 1000------// 2f int dist[MAXN+1], q[MAXN+1];------// b8
vi adj[MAXV];------// ff struct bipartite_graph {------// 2b
ii start_end() {------// 30 ----bipartite_graph(int _N, int _M) : N(_N), M(_M),------// 8d
----int start = -1, end = -1, any = 0, c = 0;------// 74 -----L(new int[N]), R(new int[M]), adj(new vi[N]) {}------// cd
------else if (indeg[i] != outdeg[i]) return ii(-1,-1);-------// fa ------else dist(v) = INF;-------// b3
----}-----dist(-1) = INF;-------// 96
}-------foreach(u, adj[v]) if(dist(R[*u]) == INF)------// 95
---ii se = start_end();------// 45 -----}
------if (s.empty()) break;------// ee -----if(dist(R[*u]) == dist(v) + 1)------// 64
----}------return true;-------------// fa
}------// aa ------dist(v) = INF;------------// 72
                       ------return false;-----// 97
3.8. Bipartite Matching.
                       -----return true;------// c6
3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
                       ----}-----// f7
O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
                       ----void add_edge(int i, int j) { adj[i].push_back(j); }------// 11
graph, respectively.
                       ----int maximum_matching() {------// 2d
vi* adi:-----// cc
                       ------int matching = 0;-----// f5
bool* done;-----// b1
                       ------memset(L, -1, sizeof(int) * N);------// 8f
int* owner:-----// 26
                       -----memset(R, -1, sizeof(int) * M);-----// 39
int alternating_path(int left) {------// da
                       ------while(bfs()) for(int i = 0; i < N; ++i)------// 77
----if (done[left]) return 0;-------// 08
                       ------matching += L[i] == -1 && dfs(i);------// f1
----done[left] = true;-----// f2
                       -----return matching:-----// fc
----for (int i = 0; i < size(adj[left]); i++) {-------// 34
                       ----}-----// le
-----int right = adj[left][i];------// b6
                       1:----// d3
-----if (owner[right] == -1 || alternating_path(owner[right])) {------// d2
-----owner[right] = left; return 1;------------------------// 26 3.9. Maximum Flow.
```

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the maximum flow of a flow network.

```
int q[MAXV], d[MAXV];-----// e6
-----if (d[s] == -1) break;-----// a0
------while ((x = augment(s, t, INF)) != 0) f += x;------// a6
-----if (res) reset():-----// 21
-----return f:-----// b6
----}-----// 1b
}:-----// 3b
```

 $3.9.1.\ Dinic's\ algorithm.$ An implementation of Dinic's algorithm that runs in $O(|V|^2|E|)$. It computes $3.9.2.\ Edmonds\ Karp's\ algorithm.$ An implementation of Edmonds Karp's algorithm that runs in $O(|V||E|^2)$. It computes the maximum flow of a flow network.

```
#define MAXV 2000-----// ba
#define MAXV 2000-----// ba int q[MAXV], p[MAXV];-----// 7b
                     struct flow_network {-----// 5e
struct flow_network {------// 12 ----struct edge {-----// fc
----struct edge {------// le ------// le ------// cb
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// d3 ------memset(head = new int[n], -1, n << 2);------// 58
------head = new int[n], curh = new int[n];------// 6b ----void destroy() { delete[] head; }------// d5
-----memset(head, -1, n * sizeof(int));------// 56 ----void reset() { e = e_store; }------// 1b
----void destroy() { delete[] head; delete[] curh; }-------// f6 ------e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;-------// 4c
----void add_edge(int u, int v, int uv, int vu = 0) {-------// cd ----}------// ef
-----e.push_back(edge(v, uv, head[u])); head[u] = ecnt++;------// c9 ----int max_flow(int s, int t, bool res = true) {-------// 12
-----e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// 89 ------if (s == t) return 0;-------// d6
----}------e_store = e;--------// 9e
------return (e[i].cap -= ret, e[i^1].cap += ret, ret);------// ac -------while (l < r)------------------------// 2c
------return 0:-------for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)------// c6
----int max_flow(int s, int t, bool res = true) {------------------------------(d[v = e[i].v] == -1 || d[u] + 1 < d[v]))-------// 2f
------memset(d, -1, n * sizeof(int));------// a8 -----at = p[t], f += x;-------// 2d
------l = r = 0, d[q[r++] = t] = 0;-------// 0e -------while (at != -1)-------// cd
------if (e[i^1].cap > 0 && d[e[i].v] == -1)-------// 29 ------if (res) reset();--------// 3b
-----d[q[r++] = e[i].v] = d[v]+1;-----// bc
                     ----}------// 05
-----memcpy(curh, head, n * sizeof(int));------// 10 };------// 75
```

3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modified to find shortest path to augment each time (instead of just any path). It computes the maximum flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with minimum cost. Running time is $O(|V|^2|E|\log|V|)$.

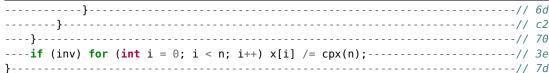
```
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#define MAXV 2000------at = p[t], f += x;-------// 43
------for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];------// ff
----struct edge {-------// 9a ---}------// 11
------int v, cap, cost, nxt;--------// ad };------// ad
-----edge(int _v, int _cap, int _cost, int _nxt)------// ec
                                     3.11. All Pairs Maximum Flow.
-----: v(_v), cap(_cap), cost(_cost), nxt(_nxt) { }------// c4
----}:-----// ad
                                     3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----int n, ecnt, *head;------// 46
                                     structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the max-
----vector<edge> e, e_store;-----// 4b
                                     imum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----flow_network(int _n, int m = -1) : n(_n), ecnt(0) {------// dd
                                     #include "dinic.cpp"-----// 58
-----e.reserve(2 * (m == -1 ? n : m));------// e6
                                      -----// 25
-----memset(head = new int[n], -1, n << 2);------// 6c
                                     bool same[MAXV];-----// 59
----}------// f3
                                     pair<vii, vvi> construct_gh_tree(flow_network \&g) {------// 77
----void destroy() { delete[] head; }------// ac
                                     ----int n = g.n, v;------// 5d
----void reset() { e = e_store; }------// 88
                                     ----vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));-----// 49
----void add_edge(int u, int v, int cost, int uv, int vu=0) {------// b4
                                     ----for (int s = 1; s < n; s++) {------// 9e
-----e.push_back(edge(v, uv, cost, head[u])); head[u] = ecnt++;-----// 43
                                     -----// 9d
-----e.push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;------// 53
                                     -----par[s].second = g.max_flow(s, par[s].first, false);------// 38
----}------// 16
                                     -----memset(d, 0, n * sizeof(int));-----// 79
----ii min_cost_max_flow(int s, int t, bool res = true) {-------// 6d
                                     -----memset(same, 0, n * sizeof(int));-----// b0
-----if (s == t) return ii(0, 0);-----// 34
                                     -----d[q[r++] = s] = 1;------// 8c
-----e_store = e;------// 70
                                     ------while (l < r) {------// 45
-----memset(pot, 0, n << 2);------// 24
                                     -----same[v = q[l++]] = true;-----// c8
------int f = 0, c = 0, v;------// d4
                                     -----for (int i = g.head[v]; i != -1; i = g.e[i].nxt)-----// 33
------while (true) {------// 29
-----memset(d, -1, n << 2);-----// fd
                                     ------if (q.e[i].cap > 0 && d[q.e[i].v] == 0)------// 3f
                                     -----d[q[r++] = g.e[i].v] = 1;-----// f8
-----memset(p, -1, n << 2);-----// b7
                                     -----set<<u>int</u>, cmp> q;-----// d8
-----q.insert(s); d[s] = 0;-----// 1d
                                     ------for (int i = s + 1; i < n; i++)------// 68
------while (!q.empty()) {-----// 04
                                     -----if (par[i].first == par[s].first && same[i]) par[i].first = s;----// ea
                                     -----q.reset();------// 9a
-----'int u = *q.begin();-----// dd
                                     ----}-----// 1e
-----q.erase(q.begin());-----// 20
                                     ----for (int i = 0; i < n; i++) {-------// 2a
-----for (int i = head[u]; i != -1; i = e[i].nxt) {------// 02
                                     -----int mn = INF, cur = i;------// 19
-----if (e[i].cap == 0) continue;-----// 1c
                                     ------while (true) {------// 3a
-----int cd = d[u] + e[i].cost + pot[u] - pot[v = e[i].v];-----// 1d
                                     -----cap[cur][i] = mn;-----// 63
-----if (d[v] == -1 \mid \mid cd < d[v]) {------// d2
                                     -----if (cur == 0) break;-----// 35
-----if (q.find(v) != q.end()) q.erase(q.find(v));------// e2
                                     -----mn = min(mn, par[cur].second), cur = par[cur].first;-----// 28
-----d[v] = cd; p[v] = i;------// f7
                                     -----q.insert(v);-----// 74
                                     ----return make_pair(par, cap);-----// 6b
                                      -----// 99
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {------// 16
-----if (p[t] == -1) break;-----// 09
                                     ---if (s == t) return 0;-----// d4
-----int x = INF, at = p[t];-----// e8
                                     ----int cur = INF, at = s;-----// 65
----while (gh.second[at][t] == -1)-----// ef
```

```
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                                                13
}------node() { prefixes = words = 0; } };------// 42
                        public:----// 88
3.12. Heavy-Light Decomposition.
                        ----node* root:-----// a9
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), below(n) {--// 4f ------cur->prefixes++;------------------------// f1
-----vi tmp(n, ID); values = segment_tree(tmp); }------// a7 ------if (begin == end) { cur->words++; break; }------// db
------it = cur->children.find(head);-------// 77
-----csz(below[u][i]), sz[u] += sz[below[u][i]]; }------// 84 ------pair<T, node*> nw(head, new node());------// cd
------for (int i = 0; i < size(below[u]); i++)----------// a7 ----int countMatches(I begin, I end) {------------------------------// 7f
------if (best == -1 || sz[below[u][i]] > sz[best]) best = below[u][i];--// 19 ------node* cur = root;----------// 32
-------for (int i = 0; i < size(below[u]); i++)-----------// 7d -------if (begin == end) return cur->words;--------// a4
----void build() { int u = curloc = 0;-------// 06 ------T head = *begin;-------// 5c
------it = cur->children.find(head):------// d9
-----u = size(uat) - 1, v = size(vat) - 1:--------// ad -----node* cur = root:-------------------------------// 95
-----res = (loc[uat[u]] < loc[vat[v]] ? uat[u] : vat[v]), u--, v--;----// 13 -------if (begin == end) return cur->prefixes;----------// f5
-----res = f(res, values.query(loc[head[u]], loc[u])), ------// 7c ------it = cur->children.find(head); -------// 43
-----u = parent[head[u]];-------// 4b -------if (it == cur->children.end()) return 0;------// 71
------return f(res, values.query(loc[v] + 1, loc[u])); }-------// 47 -------beqin++, cur = it->second; } } } };-------// 26
----int query(int u, int v) { int l = lca(u, v);------// 04
                        4.2. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
-----return f(query_upto(u, l), query_upto(v, l)); } };-----// 52
                        struct entry { ii nr: int p: }:-----// f9
                        bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77
          4. Strings
                        struct suffix_array {------// 87
4.1. Trie. A Trie class.
                        ----string s; int n; vvi P; vector<entry> L; vi idx;------// b6
class trie {------// 9a ----suffix_array(string _s) : s(_s), n(size(s)) {------// e5
private:------L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);------// 8a
----struct node {------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';------// 8d
```

```
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-----P.push_back(vi(n));-------// 30 ------if (!st) st = qo;-------// e7
------for (int i = 0; i < n; i++)--------// d5 -------s->fail = st->next[a->first];-------// 29
------L[L[i].p = i].nr = ii(P[stp - 1][i],-------// fc ------if (s->fail) {-----------------------// 3b
-----i + cnt < n ? P[stp - 1][i + cnt] : -1);-------// e5 -------if (!s->out) s->out = s->fail->out;-------// 80
-----for (int i = 0; i < n; i++)-------// 85 ------out_node* out = s->out;------// bf
------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;------// fd ------out->next = s->fail->out;------// 65
};-----cur = cur->fail;------// 21
                               -----if (!cur) cur = qo;-----// 2f
4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
                               -----cur = cur->next[*c];-----// 58
state machine from a set of keywords which can be used to search a string for any of the keywords.
                               -----if (!cur) cur = qo;-----// 3f
struct aho_corasick {------// 78
                               ------for (out_node *out = cur->out; out = out->next)-----// eθ
----struct out_node {------// 3e
                               -----res.push_back(out->keyword);----------------------------// 0d
-----string keyword; out_node *next;------// f0
                               -----}-------------// fd
-----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
                               -----return res:-----// c1
----}:------// b9
                               ----struct go_node {------// 40
                               }:-----// 32
-----map<char, go_node*> next;------// 6b
                               4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
-----out_node *out; go_node *fail;-----// 3e
                               also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
-----go_node() { out = NULL; fail = NULL; }-----// Of
                               can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
----qo_node *qo;-----// b8
                               int* z_values(const string &s) {------// 4d
----aho_corasick(vector<string> keywords) {------// 4b
                               ----int n = size(s);-----// 97
-----qo = new qo_node();-----// 77
                               ----int* z = new int[n];-----// c4
------foreach(k, keywords) {-------// e4
-----qo_node *cur = qo;-----// 9d
                               ----int l = 0, r = 0;------// 1c
                               ---z[0] = n;
-----foreach(c, *k)-----// 38
                               ----for (int i = 1; i < n; i++) {------// 7e
-----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d
-----(cur->next[*c] = new go_node());-----// 75 -----z[i] = 0;------
-----cur->out = new out_node(*k, cur->out);-----// 6e -----if (i > r) {-------// 26
-----queue<go_node*> q;------// 8a -------while (r < n && s[r - l] == s[r]) r++;-----// ff
------foreach(a, go->next) g.push(a->second);-------// a3 -----z[i] = r - l; r--;-------// fc
-----go_node *r = q.front(); q.pop();------// 2e -----else {-------
------foreach(a, r->next) {-------// 02
------go_node *s = a->second;------// cb -------while (r < n && s[r - l] == s[r]) r++;------// b3
------q.push(s);------// 76 -----z[i] = r - l; r--; } }------// 8d
-----go_node *st = r->fail;-----// fa ----return z;------// 53
------while (st && st->next.find(a->first) == st->next.end())-----// d7 }------// db
```

```
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------data[i] %= radix;------// bd
------for (int i = n.size() - 1; i >= 0; i--) {--------// la ----int stop = l-1;--------------------// cb
-----r.data.insert(r.data.begin(), 0);------// c7 ----while (stop > 0 && data[stop] == 0) stop--;------// 97
------long long k = 0:-------// cc ---ss << data[stop]:------// 96
------if (d.size() < r.size())--------// b9 ----for (int i = stop - 1; i >= 0; i--)-------// bd
------k = (long long)intx::radix * r.data[d.size()];------// f7 -----ss << setfill('0') << setw(len) << data[i];------// b6
------k /= d.data.back();---------// b7 ----delete[] a; delete[] b;------------// 7e
-----q.data[i] = k;------// d4 }------// d4 }------// d9
5.2. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
-----return pair<intx, intx>(q.normalize(n.sign * d.sign), r);-----// a1
----}-------// 1b
                                   k items out of a total of n items.
----intx operator /(const intx& d) const {-------// a2
                                   int nck(int n, int k) {-----// f6
-----return divmod(*this,d).first; }-----// 1e
                                   ----if (n - k < k) k = n - k:-----// 18
----intx operator %(const intx& d) const {-----------------// 07 ----int res = 1;-----------------------------// cb
-----return divmod(*this,d).second * sign; }------// 5a
                                   ----for (int i = 1; i <= k; i++) res = res * (n - (k - i)) / i;-----// bd
                                   ----return res:-----// e4
                                    -----// 03
5.1.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
#include "intx.cpp"-----// 83
                                   5.3. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
#include "fft.cpp"------// 13
                                   integers a, b.
-----// e0
                                   int gcd(int a, int b) { return b == 0 ? a : gcd(b, a % b); }------// d9
intx fastmul(const intx &an, const intx &bn) {-----// ab
                                    The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
----string as = an.to_string(), bs = bn.to_string();-----// 32
                                   and also finds two integers x, y such that a \times x + b \times y = d.
----int n = size(as), m = size(bs), l = 1,-----// dc
                                   int egcd(int a, int b, int& x, int& y) {------// 85
-----len = 5, radix = 100000,-----// 4f
                                   ----if (b == 0) { x = 1; y = 0; return a; }------// 7b
-----*a = new int[n], alen = 0,-----// b8
-----*b = new int[m], blen = 0;-----// 0a
                                   ----else {------// 00
                                   -----int d = egcd(b, a % b, x, y);-----// 34
----memset(a, 0, n << 2);-----// 1d
                                   -----x -= a / b * y;------// 4a
----memset(b, 0, m << 2);-----// 01
                                   -----swap(x, y);------// 26
----for (int i = n - 1; i >= 0; i -= len, alen++)-----// 6e
------for (int j = min(len - 1, i); j >= 0; j--)-----// 43
                                   -----return d;-----// db
                                   ----}------// 9e
-----a[alen] = a[alen] * 10 + as[i - j] - '0';-----// 14
                                   }-----// 40
----for (int i = m - 1; i >= 0; i -= len, blen++)------// b6
------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
                                   5.4. Trial Division Primality Testing. An optimized trial division to check whether an integer is
-----b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
----while (l < 2*max(alen,blen)) l <<= 1;------// 51
                                   bool is_prime(int n) {------// 6c
----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
                                   ----if (n < 2) return false;------// c9
----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);-----// 35
----ull *data = new ull[l];-----// e7 ----for (int i = 5; i <= s; i += 6)-----// 6c
```

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5.5. Miller-Rabin Primality Test. The Miller-Rabin probabilistic primality test.
                                          ----return res;------// eb
                                         }-----// c5
#include "mod_pow.cpp"-----// c7
bool is_probable_prime(ll n, int k) {------// be
                                          5.9. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
----if (~n & 1) return n == 2;-----// d1
                                          #include "egcd.cpp"-----// 55
----if (n <= 3) return n == 3;-----// 39
                                          int crt(const vi& as, const vi& ns) {-----// c3
----int s = 0; ll d = n - 1;------// 37
                                          ----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
----while (~d & 1) d >>= 1. s++:-----// 35
                                          ----for (int i = 0; i < cnt; i++) N *= ns[i]; ------// 88
----while (k--) {------// c8
                                          ----for (int i = 0; i < cnt; i++)-----// f9
------ll a = (n - 3) * rand() / RAND_MAX + 2;------// 06
                                          -----egcd(ns[i], l = N/ns[i], r, s), x += as[i] * s * l;-----// b0
------ll x = mod_pow(a, d, n); ------// 64
                                          ----return mod(x, N); }-----// 9e
-----if (x == 1 \mid | x == n - 1) continue;-----// 9b
-----bool ok = false;-----// 03
                                         5.10. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
------for (int i = 0: i < s - 1: i++) {------// 6b
#include "eqcd.cpp"-----// 55
------if (x == 1) return false;-----// 4f
                                         vi linear_congruence(int a, int b, int n) {------// c8
-----if (x == n - 1) { ok = true; break; }-----// 74
                                         ----int x, y, d = eqcd(a, n, x, y);------// 7a
------1-------------------------// a9
                                         ----vi res:-----// f5
------if (!ok) return false;------// 00
                                         ----if (b % d != 0) return res;------// 30
----} return true; }-------// bc ----int x0 = mod(b / d * x, n);-------// 48
                                          ----for (int k = 0; k < d; k++) res.push_back(mod(x0 + k * n / d, n));-----// 21
5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
                                          ----return res:-----// 03
vi prime_sieve(int n) {------// 40
                                         }-----// 1c
----int mx = (n - 3) >> 1, sq, v, i = -1;------// 27
                                          5.11. Numeric Integration. Numeric integration using Simpson's rule.
----vi primes;-----// 8f
                                          double integrate(double (*f)(double), double a, double b,-----// 76
----bool* prime = new bool[mx + 1];-----// ef
                                          ------double delta = 1e-6) {------// c0
----memset(prime, 1, mx + 1);------// 28
                                          ----if (abs(a - b) < delta)------// 38
----if (n >= 2) primes.push_back(2);-----// f4
                                          -----return (b-a)/8 *-----// 56
----while (++i <= mx) if (prime[i]) {-----// 73
                                          ------primes.push_back(v = (i << 1) + 3);------// be
                                          ----return integrate(f, a,-----// 64
------if ((sq = i * ((i << 1) + 6) + 3) > mx) break;-----// 2d
                                          -----(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);-----// θε
------for (int j = sq; j <= mx; j += v) prime[j] = false; }------// 2e
                                          -----// 4b
----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3):-----// 29
----delete[] prime; // can be used for O(1) lookup-----// 36
                                          5.12. Fast Fourier Transform. The Cooley-Tukey algorithm for quickly computing the discrete
----return primes; }-----// 72
                                          Fourier transform. Note that this implementation only handles powers of two, make sure to pad with
                                          zeros.
5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
                                          #include <complex>-----// 8e
                                          typedef complex<long double> cpx;-----// 25
-----// e8
                                          void fft(cpx *x, int n, bool inv=false) {------// 23
----for (int i = 0, j = 0; i < n; i++) {------// f2
----int x, y, d = eqcd(a, m, x, y);------// 3e
                                         -----if (i < j) swap(x[i], x[j]);------// 5c
----if (d != 1) return -1;-----// 20
                                         -----int m = n>>1:-----// e5
----return x < 0 ? x + m : x;------// 3c
                                         ------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
}-----// 69
                                         -----i += m:-----// ab
5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
                                          ----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
----while (e) {------------------------// b7 ------cpx t = x[i + mx] * w;-------// f5
```



5.13. Formulas.

- Number of ways to choose k objects from a total of n objects where order matters and each item can only be chosen once: $P_k^n = \frac{n!}{(n-k)!}$
- Number of ways to choose k objects from a total of n objects where order matters and each item can be chosen multiple times: n^k
- Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type 2, ..., n_k objects of type k: $\binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}$
- Number of ways to choose k objects from a total of n objects where order does not matter and each item can only be chosen once:

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^k \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0$$

- Number of ways to choose k objects from a total of n objects where order does not matter and each item can be chosen multiple times: $f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}$
- Number of integer solutions to $x_1 + x_2 + \cdots + x_n = k$ where $x_i \geq 0$: f_k^n
- Number of subsets of a set with n elements: 2^n
- $|A \cup B| = |A| + |B| |A \cap B|$
- $|A \cup B \cup C| = |A| + |B| + |C| |A \cap B| |A \cap C| |B \cap C| + |A \cap B \cap C|$
- Number of ways to walk from the lower-left corner to the upper-right corner of an $n \times m$ grid by walking only up and to the right: $\binom{n+m}{m}$
- Number of strings with n sets of brackets such that the brackets are balanced: $C = \sum_{n=1}^{n-1} C_n C_n = \sum_{n=1}^{n-1} C_n C_n$
- $C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}$
- Number of triangulations of a convex polygon with n points, number of rooted binary trees with n+1 vertices, number of paths across an $n \times n$ lattice which do not rise above the main diagonal: C_n
- ullet Number of permutations of n objects with exactly k ascending sequences or runs:

- Number of permutations of n objects with exactly k cycles: $\begin{bmatrix} n \\ k \end{bmatrix} = \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} + (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix}$
- Number of ways to partition n objects into k sets: $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements): $D_0 = 1, D_1 = 0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points: $\binom{n}{k}D_{n-k}$
- Jacobi symbol: $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- Heron's formula: A triangle with side lengths a, b, c has area $\sqrt{s(s-a)(s-b)(s-c)}$ where $s = \frac{a+b+c}{2}$.
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area $i + \frac{b}{2} 1$.
- Divisor sigma: The sum of divisors of n to the xth power is $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$ where $n = \prod_{i=0}^r p_i^{a_i}$ is the prime factorization.
- Divisor count: A special case of the above is $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$.
- Euler's totient: The number of integers less than n that are comprime to n are $n \prod_{p|n} \left(1 \frac{1}{p}\right)$ where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph, the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover.

• The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.

6. Geometry

6.1. **Primitives.** Geometry primitives.

```
#include <complex>-----// 8e
#define P(p) const point &p-----// b8
#define L(p0, p1) P(p0), P(p1)-----// 30
typedef complex<double> point;-----// e1
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// ff
point rotate(P(p), P(about), double radians) {------// e1
----return (p - about) * exp(point(0, radians)) + about; }-----// cb
point reflect(P(p), L(about1, about2)) {------// c0
----point z = p - about1, w = about2 - about1;-----// 39
----return coni(z / w) * w + about1: }-----// 03
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-----// fc
bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// 6d
double ccw(P(a), P(b), P(c)) { return cross(b - a, c - b); }-----// ca
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }-----// 75
bool collinear(L(a, b), L(p, q)) {-----// 66
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6
double angle(P(a), P(b), P(c)) {------// d0
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }-----// cc
double signed_angle(P(a), P(b), P(c)) {------// fe
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 9e
double progress(P(p), L(a, b)) {-----// d2
----if (abs(real(a) - real(b)) < EPS)------// 9e
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));-----// 35
----else return (real(p) - real(a)) / (real(b) - real(a)); }------// 2c
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) \{-----//d6\}
----// NOTE: check for parallel/collinear lines before calling this function---// 02
----point r = b - a, s = q - p;------// 79
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c; --// a8
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// ae
-----return false:-----// a3
----res = a + t * r:-----// ca
----return true:-----// 17
point closest_point(L(a, b), P(c), bool segment = false) {-----// a1
----if (seament) {-------// c2
------if (dot(b - a, c - b) > 0) return b;------// b5
-----if (dot(a - b, c - a) > 0) return a;-----// cf
----double t = dot(c - a, b - a) / norm(b - a);------// aa
----return a + t * (b - a);-----// 7a
double line_segment_distance(L(a,b), L(c,d)) {------// 99
----double x = INFINITY:-----// 83
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);-----// df
----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true));-----// da
```

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----else {-------------------------// 38 int convex_hull(polygon p) {----------------------// cd
-----x = min(x, abs(b - closest_point(c,d, b, true)));--------// ec ----sort(p.begin(), p.end(), cmp);--------// 3d
}------// b3 ---}-----// d8
                               ----int r = 1:-----// 59
6.2. Polygon. Polygon primitives.
                               ----for (int i = n - 2; i >= 0; i--) {------// 16
#include "primitives.cpp"-----// e0 ------if (p[i] == p[i + 1]) continue;------// c7
double polygon_area_signed(polygon p) {------// 31 -----hull[r++] = p[i];-------// 6d
----double area = 0; int cnt = size(p);-----// a2 ---}
----return area / 2; }-----// e1
double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }-----// 25 6.4. Line Segment Intersection. Computes the intersection between two line segments.
----for (int i = 0, j = n - 1; i < n; j = i++)--------// 77 ------A = B = a; return abs(a - d) < EPS; }-------// ee
------if (collinear(p[i], q, p[j]) &&-------// a5 ----else if (abs(a - b) < EPS) {-------// 03
------0 <= (d = progress(q, p[i], p[j])) && d <= 1)-------// b9 ------A = B = a; double p = progress(a, c,d);-------// c9
-----return 0:------// cc -----return 0.0 <= p && p <= 1.0------// 8a
------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))------// 1f ----else if (abs(c - d) < EPS) {-------// 26
-----in = !in;------// b2 ------A = B = c; double p = progress(c, a,b);------// d9
----return in ? -1 : 1; }------// 77 ------return 0.0 <= p && p <= 1.0------// 8e
//--- polygon left, right;-----// 6b ---else if (collinear(a,b, c,d)) {------// bc
//---- point it(-100, -100);-------// c9 ------double ap = progress(b, c,d);------// a7
//------ int i = i == cnt-1 ? 0: i+1;-------// 8e ---------if (bp < 0.0 || ap > 1.0) return false;-------// 0c
//------ point p = poly[i], q = poly[i];------// 19 ------A = c + max(ap, 0.0) * (d - c);------// f6
//----- if (ccw(a, b, p) \le 0) left.push_back(p);------// 12 ------B = c + min(bp, 1.0) * (d - c);-------// 5c
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// e3 -----return true; }-----
//-----// mvintersect = intersect where------// 24 ----else if (parallel(a.b. c.d)) return false;--------// ca
//-----// (a,b) is a line, (p,q) is a line segment------// f2 ----else if (intersect(a,b, c,d, A, true)) {----------------// 10
//----- if (myintersect(a, b, p, q, it))------// f0 -----B = A; return true; }--------------------// bf
//----------left.push_back(it), right.push_back(it);-------// 21 ----return false;---------------------------// b7
//--- }------// 5e }------// 8b
//--- return pair<polygon, polygon>(left, right);-----// 1d -----// e6
// }-----// 37
                               6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
                               coordinates) on a sphere of radius r.
6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
#include "polygon.cpp"------// 58 double gc_distance(double pLat, double pLong,-----// 7b
point hull[MAXN];------// 43 ----pLat *= pi / 180; pLong *= pi / 180;------// ee
```

```
----qLat *= pi / 180; qLong *= pi / 180;-------------// 75 ----while (low < high) {---------------------------------// a3
-----sin(pLat) * sin(qLat));------// le -----if (f(mid)) high = mid;------// ca
}------// 3f ---}------// 9b
                                             ----assert(f(low));------// 42
6.6. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
                                             ----return low:------// a6
pair of points.
                                            }-----// d3
#include "primitives.cpp"-----// e0
-----// 85
                                            7.2. Ternary Search. Given a function f that is first monotonically increasing and then monotonically
struct cmpx { bool operator ()(const point &a, const point &b) {------// 01
                                            cally decreasing, ternary search finds the x such that f(x) is maximized.
-----return abs(real(a) - real(b)) > EPS ?-----// e9
                                            template <class F>-----// d1
-----real(a) < real(b) : imaq(a) < imaq(b); } };------// 53
                                            double ternary_search_continuous(double lo, double hi, double eps, F f) {-----// e7
struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f ----while (hi - lo > eps) {-------// 3e
----return abs(imag(a) - imag(b)) > EPS ?-----// 0b ------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
-----imag(a) < imag(b) : real(a) < real(b); } };------// a4 -----if (f(m1) < f(m2)) lo = m1;------// 1d
----double mn = INFINITY;-----// f9
-------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);-------// 8b #include "../graph/scc.cpp"--------// c3
-----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));------// fc ------// 63
------jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));-------// 39 bool two_sat(int n, const vii& clauses, vi& all_truthy) {-------// f4
-----cur.insert(pts[i]); }------// 82 ----vvi adj(2*n+1);------// 7b
----return mn; }------// 4c ----for (int i = 0; i < size(clauses); i++) {--------// 9b
                                            -----adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// 17
6.7. Formulas. Let a = (a_x, a_y) and b = (b_x, b_y) be two-dimensional vectors.
                                            -----if (clauses[i].first != clauses[i].second)------// 87
  • a \cdot b = |a||b|\cos\theta, where \theta is the angle between a and b.
                                             -----adj[-clauses[i].second + n].push_back(clauses[i].first + n);-----// 93
  • a \times b = |a||b|\sin\theta, where \theta is the signed angle between a and b.
                                             ----}--------// d8
  • a \times b is equal to the area of the parallelogram with two of its sides formed by a and b. Half
                                            ----pair<union_find, vi> res = scc(adj);------// 9f
                                             ----union_find scc = res.first;------// 42
   of that is the area of the triangle formed by a and b.
                                             ----vi dag = res.second;------// 58
                7. Other Algorithms
                                             ----vi truth(2*n+1, -1);------// 00
                                             ----for (int i = 2*n; i >= 0; i--) {------// f4
7.1. Binary Search. An implementation of binary search that finds a real valued root of the continuous
                                             -----int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n); -// 5a
function f on the interval [a, b], with a maximum error of \varepsilon.
                                             -----if (cur == 0) continue;-----// 26
double binary_search_continuous(double low, double high,------// 8e
                                             -----if (p == o) return false:-----// 33
-----double eps, double (*f)(double)) {-----// c0
                                            -----if (truth[p] == -1) truth[p] = 1;-----// c3
----while (true) {------// 3a
                                            -----truth[cur + n] = truth[p];-----// b3
------double mid = (low + high) / 2, cur = f(mid);-----// 75
                                             -----truth[o] = 1 - truth[p];-----// 80
-----if (abs(cur) < eps) return mid;------// 76
                                             ------if (truth[p] == 1) all_truthy.push_back(cur);------// 5c
-----else if (0 < cur) high = mid;------// e5
                                             ----}-------// d9
-----else low = mid;-----// a7
                                             ----return true;------// eb
----}------// b5
                                            }------// 61
}-----// cb
 Another implementation that takes a binary predicate f, and finds an integer value x on the integer
                                            7.4. Stable Marriage. The Gale-Shapley algorithm for solving the stable marriage problem.
                                            vi stable_marriage(int n, int** m, int** w) {------// e4
interval [a, b] such that f(x) \wedge \neg f(x-1).
----assert(low <= high);-----// 19 ----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
```

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----for (int i = 0; i < n; i++) q.push(i);------// fe ------ptr[i][j]->d = ptr[ni][j];-----// 71
----while (!q.empty()) {------// 55 ------ptr[ni][j]->u = ptr[i][j];-----// c4
------int curm = q.front(); q.pop();--------// ab --------while (true) {-----------------------------------// c6
------int curw = m[curm][i];--------// cf -------if (i == rows || arr[i][nj]) break;------// 8d
------if (eng[curw] == -1) { }-------// 35 -----++nj;-----------------// 1c
-----q.push(enq[curw]);------// 8c -----ptr[i][j]->r = ptr[i][nj];-----// d5
------else continue;------// b4 -------ptr[i][nj]->l = ptr[i][j];--------// 72
}------// 03 ------ptr[rows][0]->l = head;-------// 3b
                             ------head->l = ptr[rows][cols - 1];------// da
7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the
                             -----ptr[rows][cols - 1]->r = head;------// 6b
Exact Cover problem.
                             ------for (int j = 0; j < cols; j++) {------// 97
bool handle_solution(vi rows) { return false; }------// 63
                             -----int cnt = -1:-----// 84
struct exact_cover {-----// 95
                            -----for (int i = 0; i \le rows; i++)------// 96
----struct node {------// 7e ------if (ptr[i][j]->p = ptr[rows][j];-----// cb
------node *l, *r, *u, *d, *p;-------// 19 ------ptr[rows][j]->size = cnt;------// 59
------int row, col, size;------// ae
                            -----node(int _row, int _col) : row(_row), col(_col) {------// c9
                            ------for (int i = 0; i <= rows; i++) delete[] ptr[i];------// bf
-----/c3
                             -----delete[] ptr:-----// 99
----} -----// c1
                             ----}-----// c0
----int rows, cols, *sol;------// 7b
                             ----#define COVER(c, i, j) \sqrt{\phantom{a}}-----// 6a
----bool **arr;------// e6
                             ----node *head:-----// fe
                             ------for (node *i = c->d; i != c; i = i->d) \[ \] ------// a3
----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
                             -----arr = new bool*[rows];-----// cf
                             -----j->d->u = j->u, j->u->d = j->d, j->p->size--;------// 16
-----sol = new int[rows];-----// 5f
                             ----#define UNCOVER(c, i, j) \\-----// d0
------for (int i = 0; i < rows; i++)------// 89
                             ------for (node *i = c->u; i != c; i = i->u) \[ \bigcup_{------// ff} \]
-----/r[i] = new bool[cols], memset(arr[i], 0, cols);-----// 75
----void set_value(int row, int col, bool val = true) { arr[row][col] = val; }-// 03 ------j->p->size++, j->d->u = j->u->d = j; \[ \]------// b6
------node ***ptr = new node**[rows + 1];------// 35 ----bool search(int k = 0) {-------// bb
-----ptr[i] = new node*[cols];------// 0b -----vi res(k);------
-----sort(res.begin(), res.end());------// 87
------else ptr[i][j] = NULL;------// 32 -----return handle_solution(res);-----// 51
------for (int i = 0; i <= rows; i++) {--------// 84 ------node *c = head->r, *tmp = head->r; ------// 8e
-----if (!ptr[i][j]) continue;-----// 35 -----if (c == c->d) return false;-----// b0
-----// b0 ------bool found = false;------// 7f
-----if (ni == rows || arr[ni][j]) break;-----// 19
```

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-----found = search(k + 1);-----// f1
-----for (node *j = r->1; j != r; j = j->1) { UNCOVER(j->p, a, b); j----//ab
                                                                       8. Useful Information
------UNCOVER(c, i, j);------// 3a
                                                    8.1. Tips & Tricks.
-----return found:-----// 80
                                                       • How fast does our algorithm have to be? Can we use brute-force?
                                                       • Does order matter?
                                                       • Is it better to look at the problem in another way? Maybe backwards?
                                                       • Are there subproblems that are recomputed? Can we cache them?
7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \ldots, k-1\}
                                                       • Do we need to remember everything we compute, or just the last few iterations of computation?
1}.
                                                       • Does it help to sort the data?
vector<int> nth_permutation(int cnt, int n) {------// 78
                                                       • Can we speed up lookup by using a map (tree or hash) or an array?
----vector<int> idx(cnt), per(cnt), fac(cnt);------// 9e
                                                       • Can we binary search the answer?
----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
                                                       • Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph
----for (int i = 1; i \le cnt; i++) fac[i - 1] = n % i, n /= i;-----// 04
                                                        into some other kind of a graph (perhaps a DAG, or a flow network)?
----for (int i = cnt - 1; i >= 0; i--)-----// 52
                                                       • Make sure integers are not overflowing.
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// 41
                                                       • Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo
----return per;-----// 84
                                                        m_1, m_2, \ldots, m_k, where m_1, m_2, \ldots, m_k are pairwise coprime integers, and find the real answer
}-----// 97
                                                        using CRT?
                                                       • Are there any edge cases? When n = 0, n = -1, n = 1, n = 2^{31} - 1 or n = -2^{31}? When
7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
                                                        the list is empty, or contains a single element? When the graph is empty, or contains a single
ii find_cycle(int x0, int (*f)(int)) {------// a5
                                                        vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
• Can we use exponentiation by squaring?
----while (t != h) t = f(t), h = f(f(h));-----// 79
----h = x0:
                                                    8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input
----while (t != h) t = f(t), h = f(h), mu++;-----// 9d
                                                    reading/output writing. This can be achieved by reading all input in at once (using fread), and then
----h = f(t):-----// 00
                                                    parsing it manually. Output can also be stored in an output buffer and then dumped once in the end
----while (t != h) h = f(h), lam++:-----// 5e
                                                    (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading
----return ii(mu, lam);-----// b4
                                                    method.
}-----// 42
                                                    void readn(register int *n) {------// dc
                                                    ----int sign = 1;------// 32
7.8. Dates. Functions to simplify date calculations.
                                                    ----register char c:-----// a5
int intToDay(int jd) { return jd % 7; }------// 89
                                                    ---*n = 0;
int dateToInt(int y, int m, int d) {------// 96
                                                    ----while((c = getc_unlocked(stdin)) != '\n') {------// f3
----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +-----// a8
                                                    -----switch(c) {------// 0c
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 -----// d1
                                                    -----case '-': sign = -1; break;-----// 28
-----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
                                                    -----case ' ': goto hell;-----// fd
-----d - 32075;-----// e0
                                                    ------case '\n': goto hell;-----// 79
}-----// fa
                                                    -----default: *n *= 10; *n += c - '0'; break;------// c0
void intToDate(int jd, int &y, int &m, int &d) {------// a1
                                                    -----}-----// 2d
----int x. n. i. i:------// 00
                                                    ----}-----// c3
----x = id + 68569;-----// 11
                                                    hell:-----// ba
----n = 4 * x / 146097;-----// 2f
                                                    ----*n *= sian:------// a0
---x = (146097 * n + 3) / 4;
                                                    ----x -= 1461 * i / 4 - 31:-----// 09
                                                    8.3. 128-bit Integer. GCC has a 128-bit integer data type named __int128. Useful if doing multi-
----j = 80 * x / 2447;-----// 3d
                                                    plication of 64-bit integers, or something needing a little more than 64-bits to represent.
---d = x - 2447 * i / 80:
---x = i / 11:-----// b7
----m = j + 2 - 12 * x;------// 82 8.4. Worst Time Complexity.
```

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|---------------------|---------------------------|---|--|
| n | Worst AC Algorithm | Comment | |
| ≤ 10 | $O(n!), O(n^6)$ | e.g. Enumerating a permutation | |
| ≤ 15 | $O(2^n \times n^2)$ | e.g. DP TSP | |
| ≤ 20 | $O(2^n), O(n^5)$ | e.g. DP + bitmask technique | |
| ≤ 50 | $O(n^4)$ | e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$ | |
| $\leq 10^{2}$ | $O(n^3)$ | e.g. Floyd Warshall's | |
| $\leq 10^{3}$ | $O(n^2)$ | e.g. Bubble/Selection/Insertion sort | |
| $\leq 10^{5}$ | $O(n\log_2 n)$ | e.g. Merge sort, building a Segment tree | |
| $\leq 10^{6}$ | $O(n), O(\log_2 n), O(1)$ | Usually, contest problems have $n \leq 10^6$ (e.g. to read input) | |

8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.