```
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-----return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 34 int query(int id, int l, int r) {-----------------------// a2
------propagate(l, r, i);-------// 19 ----return query(segs[id].lid, l, r) + query(segs[id].rid, l, r); }------// ee
-----if (l > r) return ID;------// cc
                                  2.3. Fenwick Tree.
-----if (r < a || b < l) return data[i];-----// d9
                                  struct fenwick_tree {------// 98
-----if (a <= l \& x r <= b) return (lazy[i] = v) * (r - l + 1) + data[i];----// 06
                                  ----int n; vi data;------// d3
-----int m = (l + r) / 2;-----// cc
                                  ----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-----// db
-----return data[i] = f(ru(a, b, v, l, m, 2*i+1),-----// cc
                                  ----void update(int at, int by) {------// 76
-----/ ru(a, b, v, m+1, r, 2*i+2));-----// 2b
                                  ------while (at < n) data[at] += by, at |= at + 1; }-----// fb
----}-----// 0b
                                  ----int query(int at) {-------// 71
----void propagate(int l, int r, int i) {-----// a7
                                  ------int res = 0;------// c3
------if (l > r || lazy[i] == INF) return;------// 5f
                                  ------while (at >= 0) res += data[at], at = (at & (at + 1)) - 1;------// 37
-----data[i] += lazy[i] * (r - l + 1);-----// 44
                                  -----return res; }-----// e4
-----if (l < r) {------// 28
                                  ----int rsq(int a, int b) { return query(b) - query(a - 1); }------// be
------if (lazy[2*i+1] == INF) lazy[2*i+1] = lazy[i];------// 4e
                                   -----/-----// 57
------else lazy[2*i+1] += lazy[i];-----// 1e
                                  struct fenwick_tree_sq {-----// d4
------if (lazy[2*i+2] == INF) lazy[2*i+2] = lazy[i];-----// de
                                  ----int n; fenwick_tree x1, x0;------// 18
-----else lazy[2*i+2] += lazy[i];-----// 74
                                  ----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),-----// 2e
-----x0(fenwick_tree(n)) { }------// 7c
-----lazv[i] = INF:-----// f8
                                  ----// insert f(y) = my + c if x <= y------// 17
----void update(int x, int m, int c) { x1.update(x, m); x0.update(x, c); }----// 45
};-----// ae
                                  ----int query(int x) { return x*x1.query(x) + x0.query(x); }------// 73
                                  }:-----// 13
2.2.1. Persistent Segment Tree.
                                  void range_update(fenwick_tree_sq &s, int a, int b, int k) {------// 89
int segcnt = 0;-----// cf
                                  ----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }-----// 7f
struct segment {------// 68
                                  ----int l, r, lid, rid, sum;-----// fc
                                  ----return s.query(b) - s.query(a-1); }------// f3
} seas[2000000]:----// dd
                                  2.4. Matrix.
int build(int l, int r) {-----// 2b
----int id = segcnt++;-------------------------// a8 template <> bool eq<double o, double b) { return abs(a - b) < EPS; }---// a7
----else {------// fe ----matrix(int r, int c) : rows(r), cols(c), cnt(r * c) {-------// 56
------segs[id].lid = build(l , m);-------// e3 ----matrix(const matrix& other) : rows(other.rows), cols(other.cols),------// b5
-----segs[id].rid = build(m + 1, r); }------// 69 -----cnt(other.cnt), data(other.data) { }------// c1
----return id; }------// c5 ----matrix<T> operator +(const matrix& other) {------// 33
int update(int idx, int v, int id) {-------data[i] += other.data[i];-----// b8 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] += other.data[i];-----// f8
----int nid = seqcnt++;-----matrix<T> res(*this); rep(i,0,cnt) res.data[i] -= other.data[i];-----// 7b
----seqs[nid].lid = update(idx, v, seqs[id].lid);------// 92 ------matrix<T> res(*this); rep(i,0,cnt) res.data[i] *= other;------// 05
```

```
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-----rep(i,0,rows) rep(j,0,other.cols) rep(k,0,cols)------// ae node* merge(node *l, node *r) {------// 3c
------if (p & 1) res = res * sq:-------// 62 ------else if (t->x < x) t = t->r:------// da
-----if (p) sq = sq * sq;-------// 35 ----return NULL; }------// ae
------matrix<T> mat(*this); det = T(1), rank = max(rows, cols);------// 7a ----pair<node*, node*> res = split(t, x);-------// ca
------for (int r = 0, c = 0; c < cols; c++) {---------// 8e ----return merge(res.first, merge(new node(x, y), res.second)); }------// 0d
-----int k = r;------// 5b node* erase(node *t, int x) {-------// 4d
-----if (k >= rows) { rank--; continue; }------// 1a ----if (t->x < x) t->r = erase(t->r, x);--------// 7c
------if (k != r) {------// c4 ---else if (x < t->x) t->l = erase(t->l, x);------// 48
-----det *= T(-1);------// 55 ----else { node *old = t; t = merge(t->l, t->r); delete old; }-----// 22
-----rep(i,0,cols) mat(r, i) /= d;------// d1 ----else return kth(t->r, k - tsize(t->l) - 1); }------// f0
-----rep(i,0,rows) {------// f6
                       2.6. Misof Tree.
-----T m = mat(i, c);------// 05
                       #define BITS 15-----// 7b
-----if (i != r && !eq<T>(m, T(0)))-----// 1a
                       struct misof_tree {-----// fe
-----rep(j,0,cols) mat(i, j) -= m * mat(r, j);-----// 7b
                       ----int cnt[BITS][1<<BITS];-----// aa
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }-----// b0
-----} return mat; }------// b3
                       ----void insert(int x) { for (int i = 0; i < BITS; cnt[i++][x]++, x >>= 1); }--// 5a
------matrix<T> res(cols, rows);------// 5b
                       ----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49
-----rep(i,0,rows) rep(j,0,cols) res(j, i) = at(i, j);------// 92
                       ----int nth(int n) {--------// 8a
                       -----int res = 0;------// a4
-----return res; } };-----// df
                       ------for (int i = BITS-1; i >= 0; i--)------// 99
                       -------if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1;------// f4
2.5. Cartesian Tree.
                       -----return res;------// 3a
struct node {-----// 36
                       ----}-----// b5
----int x, y, sz;------// e5
                       l:----// 0a
----node *1, *r;-----// 4d
----node(int _x, int _y) : x(_x), y(_y), sz(1), l(NULL), r(NULL) { } };-----// 19 2.7. Sqrt Decomposition.
----if (t->x < x) {-------------------------// dc
------return make_pair(t, res.second); }-------// e0 ---rep(i,0,size(T))--------// b1
```

```
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-----rep(j,0,size(T[i].arr))-------// a4 };------// a2
-----arr[at++] = T[i].arr[j];-----// f7
----T.clear();------// 4c
                              2.9. Convex Hull Trick.
----for (int i = 0; i < cnt; i += K)-----// 79
                              struct convex_hull_trick {------// 16
-----T.push_back(segment(vi(arr.begin()+i, arr.begin()+min(i+K, cnt))));----// f0
                              }-----// 03
                              ----double intersect(int i) {------// 9b
int split(int at) {------// 71
                              -----return (h[i+1].second-h[i].second)/(h[i].first-h[i+1].first); }-----// b9
----int i = 0:------// 8a
                              ----void add(double m, double b) {------// a4
----while (i < size(T) && at >= size(T[i].arr))------// 6c
                              -----h.push_back(make_pair(m,b));------// f9
-----at -= size(T[i].arr), i++;-----// 9a
                              ------while (size(h) >= 3) {------// f6
----if (i >= size(T)) return size(T);------// 83
                              ------int n = size(h);-----// d8
----if (at == 0) return i;------// 49
                              -----if (intersect(n-3) < intersect(n-2)) break;-----// 07
---T.insert(T.begin() + i + 1, segment(vi(T[i].arr.begin() + at, T[i].arr.end())));
                              -----swap(h[n-2], h[n-1]);-----// bf
----T[i] = segment(vi(T[i].arr.begin(), T[i].arr.begin() + at));------// af
                              -----h.pop_back(); } }-----// 4b
----return i + 1;------// ac
                              ----<mark>double</mark> get_min(<mark>double</mark> x) {------// b0
}-----// ea
                              ------int lo = 0, hi = size(h) - 2, res = -1;-------// 5b
void insert(int at, int v) {------// 5f
                              ------while (lo <= hi) {------// 24
----vi arr; arr.push_back(v);------// 6a
                              -----/int mid = lo + (hi - lo) / 2;-----// 5a
----T.insert(T.begin() + split(at), segment(arr));------// 67
                              ------if (intersect(mid) <= x) res = mid, lo = mid + 1;------// 1d
}-----// cc
                              -----else hi = mid - 1; }-----// b6
void erase(int at) {-----// be
                              ------return h[res+1].first * x + h[res+1].second; } };------// 84
----int i = split(at); split(at + 1);-----// da
----T.erase(T.begin() + i);------// 6b
}-----// 4b
                                           3. Graphs
                              3.1. Single-Source Shortest Paths.
2.8. Monotonic Queue.
                              3.1.1. Dijkstra's algorithm.
struct min_stack {-----// d8
------M.push(M.empty() ? x : min(M.top(), x); }-------// 92 ------return dist[a] != dist[b] ? dist[b] ? dist[b] : a < b; }------// e6
};------// 74 ----rep(i,0,n) dist[i] = INF, dad[i] = -1;-----------------------// 80
----min_stack inp, outp;------// 3d ----dist[s] = 0, pq.insert(s);-------// 1f
-----if (outp.empty()) while (!inp.empty())-------// 3b -----rep(i,0,size(adj[cur])) {------------------------// a6
-----outp.push(inp.top()), inp.pop();-------// 8e ------int nxt = adj[cur][i].first,---------// a4
----}-----ndist = dist[cur] + adj[cur][i].second;------// 3a
-----if (inp.empty()) return outp.mn();------// d2
-----if (outp.empty()) return inp.mn();-------// df
----void pop() { fix(); outp.pop(); }------// 4f }------// 9b
```

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3.1.2. Bellman-Ford algorithm.

```
int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {------// cf
                                 3.2.1. Kosaraju's algorithm.
----has_negative_cycle = false;-----// 47
                                 #include "../data-structures/union_find.cpp"-----// 5e
----int* dist = new int[n];-----// 7f
                                   -----// 11
----rep(i,0,n) dist[i] = i == s ? 0 : INF;-----// df
                                 vector<<mark>bool</mark>> visited;-----// 66
----rep(i,0,n-1) rep(j,0,n) if (dist[j] != INF)------// 4d
                                 vi order:-----// 9b
-----rep(k,0,size(adj[j]))-----// 88
                                 -----// a5
-----dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// e1
                                 void scc_dfs(const vvi &adj, int u) {------// a1
-----dist[j] + adj[j][k].second);-----// 18
                                 ----int v; visited[u] = true;------// e3
----rep(j,0,n) rep(k,0,size(adj[j]))-----// f8
                                 ----rep(i,0,size(adj[u]))-----// 2d
-----if (dist[j] + adj[j][k].second < dist[adj[j][k].first])------// 37
                                 -----if (!visited[v = adj[u][i]]) scc_dfs(adj, v);-----// a2
-----has_negative_cycle = true;-----// f1
                                 ----order.push_back(u);-----// 02
----return dist;-----// 78
                                   -----// 53
}-----// a9
                                   -----// 63
                                 pair<union_find, vi> scc(const vvi &adj) {-----// c2
3.1.3. IDA^* algorithm.
                                 ----int n = size(adj), u, v;-----// f8
int n, cur[100], pos;-----// 48
                                 ----order.clear();-----// 20
int calch() {-----// 88
                                 ----union_find uf(n);-----// a8
----int h = 0:-----// 4a
                                 ----vi dag:-----// 61
----rep(i,0,n) if (cur[i] != 0) h += abs(i - cur[i]);-----// 9b
                                 ----vvi rev(n);------// c5
----return h:-----// c6
                                 ----rep(i,0,n) rep(j,0,size(adj[i])) rev[adj[i][j]].push_back(i);------// 7e
}-----// c8
                                 ----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 80
int dfs(int d, int q, int prev) {------// 12
                                 ----rep(i,0,n) if (!visited[i]) scc_dfs(rev, i);------// 4e
----int h = calch();------// 5d
                                 ----fill(visited.begin(), visited.end(), false);------// 59
----if (q + h > d) return q + h;------// 15
                                 ----stack<<u>int</u>> S;-----// bb
----if (h == 0) return 0;-----// ff
                                 ----for (int i = n-1; i >= 0; i--) {-------// 96
----int mn = INF:-----// 7e
                                 -----if (visited[order[i]]) continue;-----// db
----rep(di,-2.3) {------// 0d
                                 -----if (di == 0) continue:-----// 0a
                                 ------while (!S.empty()) {------// 9e
------int nxt = pos + di;------// 76
                                 -----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);-----// b3
-----if (nxt == prev) continue;-----// 39
                                 -----rep(j,0,size(adj[u])) if (!visited[v = adj[u][j]]) S.push(v);-----// 1b
-----if (0 <= nxt && nxt < n) {------// 68
                                 -----swap(cur[pos], cur[nxt]);-----// 35
                                 -----swap(pos,nxt);-----// 64
                                 ----return pair<union_find. vi>(uf. dag):-----// 2b
-----mn = min(mn, dfs(d, q+1, nxt));------// 22
                                 }-----// 92
-----swap(pos.nxt):-----// 84
----return mn;------// da ----low[u] = num[u] = curnum++;------// a3
}------// f8 ----int cnt = 0; bool found = false;------// 97
int idastar() {------// 22 ----rep(i,0,size(adj[u])) {------// ae
----rep(i,0,n) if (cur[i] == 0) pos = i;------// 6b ------int v = adj[u][i];------// 56
----while (true) {-------------------------// 18 -------dfs(adj, cp, bri, v, u);-------------// ba
-----d = nd;-------found = found || low[v] >= num[u];---------// 30
----}------if (low[v] > num[u]) bri.push_back(ii(u, v));-------// bf
```

3.2. Strongly Connected Components.

```
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pair<vi,vii> cut_points_and_bridges(const vvi &adj) {-------// 76 -----owner[right] = left; return 1;---------// f2
----memset(num. -1. n << 2):-----// 45
                                  3.5.2. Hopcroft-Karp algorithm. Running time is O(|E|\sqrt{|V|}).
#define MAXN 5000-----// f7
----return make_pair(cp, bri); }------// 4c
                                  int dist[MAXN+1], q[MAXN+1];.....// b8
                                  \#define\ dist(v)\ dist[v == -1\ ?\ MAXN\ :\ v]------// 0f
3.4. Euler Path.
                                  struct bipartite_graph {------// 2b
#define MAXV 1000-----// 2f
                                  ----int N, M, *L, *R; vi *adj;------// fc
                                  ----bipartite_graph(int _N, int _M) : N(_N), M(_M),------// 8d
vi adj[MAXV];-----// ff
                                  -----L(new int[N]), R(new int[M]), adj(new vi[N]) {}------// cd
----~bipartite_graph() { delete[] adj; delete[] L; delete[] R; }------// 89
ii start_end() {------// 30
----int start = -1, end = -1, any = 0, c = 0;-----// 74
                                  -----int l = 0, r = 0:-----// 37
---rep(i,0,n) {-----// 20
                                  -----rep(v,0,N) if(L[v] == -1) dist(v) = 0, q[r++] = v;--------// f9
-----if (outdeg[i] > 0) any = i;------// 63
                                  ------else dist(v) = INF;-----// aa
-----if (indeg[i] + 1 == outdeg[i]) start = i, c++;------// 5a
                                  -----dist(-1) = INF;-----// f2
-----else if (indeg[i] == outdeg[i] + 1) end = i, c++;------// 13
                                  ------while(l < r) {------// ba
-----else if (indeg[i] != outdeg[i]) return ii(-1,-1);------// c1
                                  -----int v = q[l++];-----// 50
----}------// ed
                                    ·-----if(dist(v) < dist(-1)) {------// f1
----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 54
                                  -----/jter(u, adj[v]) if(dist(R[*u]) == INF)------// 9b
----if (start == -1) start = end = any;-----// 5e
                                    -----dist(R[*u]) = dist(v) + 1, q[r++] = R[*u];-----// 79
----return ii(start, end);------// a2
                                   }-----// eb
bool euler_path() {------// b4
                                   ----return dist(-1) != INF;-----// 43
----ii se = start_end();------// 8a
----int cur = se.first, at = m + 1;------// b6
                                  ----bool dfs(int v) {-------// 26
----if (cur == -1) return false;-----// ac
                                  -----if(v != -1) {------// d8
----stack<int> s:-----// 1c
                                  -----iter(u, adj[v])------// 99
----while (true) {------// b3
                                    -----if(dist(R[*u]) == dist(v) + 1)------// 74
-----if (outdeg[cur] == 0) {------// 0d
                                    -----if(dfs(R[*u])) {-----// 40
-----res[--at] = cur;-----// bd
                                    -----R[*u] = v, L[v] = *u;------// 47
------if (s.empty()) break;-----// c6
                                      -----// a2
-----cur = s.top(); s.pop();-----// 06
-----} else s.push(cur), cur = adj[cur][--outdeg[cur]];------// 9e
                                   -----dist(v) = INF;-----// 62
----}-------// a4
3.5. Bipartite Matching.
                                  ----void add_edge(int i, int j) { adj[i].push_back(j); }------// 92
3.5.1. Alternating Paths algorithm.
                                  ----int maximum_matching() {------// a2
bool* done;-----memset(L, -1, sizeof(int) * N);-------------------------------// 72
int* owner;-----memset(R, -1, sizeof(int) * M);--------------------------------// bf
----done[left] = true;-------// f2 ------return matching;------// ec
```

```
3.6. Maximum Flow.
               3.7. Minimum Cost Maximum Flow. Running time is O(|V|^2|E|\log|V|). NOTE: Doesn't work
               on negative weights!
3.6.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|).
               #define MAXV 2000-----// ba
#define MAXV 2000------// ba int d[MAXV], pt[MAXV], pt[MAXV];-------// 80
struct flow_network {------// 12 ----bool operator ()(int i, int j) {------// 8a
------int v. cap. nxt;--------// ab ---}-----// ab ---}------// ab
-----edge() { }------// 38 };------// cf
------edge(int _v, int _cap, int _nxt) : v(_v), cap(_cap), nxt(_nxt) { }-----// bc struct flow_network {------------------------------// eb
-----e, reserve(2 * (m == -1 ? n : m)):------// ad
----void reset() { e = e_store; }------// 87 -----memset(head = new int[n], -1, n << 2);------// 6c
------if (v == t) return f;-------// 6d ------e.push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;------// 53
-----if (e[i].cap > 0 && d[e[i].v] + 1 == d[v])------// cc ---ii min_cost_max_flow(int s, int t, bool res = true) {------// 6d
-----if(s == t) return 0;-------// 9d -----memset(d, -1, n << 2);-------// fd
-----e_store = e;------// 57 -----memset(p, -1, n << 2);------// b7
-------q.insert(s); d[s] = 0;-------// 1d
-----while (l < r)------// 7a ------q.erase(q.begin());-------// 20
-----d[q[r++] = e[i].v] = d[v]+1;-----// 28 ------int cd = d[u] + e[i].cost + pot[u] - pot[v = e[i].v];-----// 1d
-----memcpy(curh, head, n * sizeof(int));------// 10 ------if (q.find(v) != q.end()) q.erase(q.find(v));------// e2
```

```
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------int x = INF, at = p[t];-------// e8 -------for (int j = 0; j < n; j++)------// 6e
------while (at != -1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 32 ------if (dist[j] != INF)-------// e3
-----rep(i,0,n) if (p[i] != -1) pot[i] += d[i];-------// 86 -------back[g[j][k]->v] = g[j][k];-----// 3d
------if (res) reset();--------// d7 ------mcmf_edge* cure = back[t];------// b4
------return ii(f, c);-------// 9f -------if (cure == NULL) break;------// ab
-----cap = min(cap, cure->w);-----// c3
A second implementation that is slower but works on negative weights.
                                -----if (cure->u == s) break;-----// 82
struct flow_network {------// 81
                                -----cure = back[cure->u];-----// 45
----struct mcmf_edge {------// f6
                                ------int u, v;------// e1
                                -----assert(cap > 0 && cap < INF);-----// ae
-----ll w, c;-----// b4
                                -----cure = back[t];-----// b9
-----mcmf_edge* rev;-----// 9d
                                -----while (true) {------// 2a
-----mcmf_edge(int _u, int _v, ll _w, ll _c, mcmf_edge* _rev = NULL) {-----// ea
                                -----cost += cap * cure->c;------// f8
------u = _u; v = _v; w = _w; c = _c; rev = _rev;------// 83
                                -----cure->w -= cap:-----// d1
------cure->rev->w += cap;-----// cf
----};------// b9
                                ------if (cure->u == s) break;------// 8c
----int n;------// b4
                                ----vector<pair<int, pair<ll, ll> > * adj;------// 72
                                ----flow_network(int _n) {------// 55
                                -----flow += cap:-----// f2
----n = _n:-----// fa
                                -----adj = new vector<pair<int, pair<ll, ll> > >[n];------// bb
                                -----// instead of deleting g, we could also-----// e0
                                -----// use it to get info about the actual flow-----// 6c
----void add_edge(int u, int v, ll cost, ll cap) {------// 79
                                ------for (int i = 0; i < n; i++)-----// eb
-----adj[u].push_back(make_pair(v, make_pair(cap, cost)));------// c8
                                ------for (int j = 0; j < size(q[i]); j++)------// 82
                                -----delete q[i][j];-----// 06
----pair<ll,ll> min_cost_max_flow(int s, int t) {------// ea
                                -----delete[] q;------// 23
-----vector<mcmf_edge*>* q = new vector<mcmf_edge*>[n];------// ce
                                -----delete[] back;-----// 5a
------for (int i = 0; i < n; i++) {-------// 57
                                -----delete[] dist;-----// b9
-----for (int j = 0; j < size(adj[i]); j++) {------// 37
                                -----return make_pair(flow, cost);------// ec
-----/mcmf_edge *cur = new mcmf_edge(i, adj[i][j].first,-----// 21
                                ----}-------// ad
-----adj[i][j].second.first, adj[i][j].second.second),--// 56
                                 -----// bf
-----*rev = new mcmf_edge(adj[i][j].first, i, 0,------// 48
-----dj[i][j].second.second, cur);-----// b1
                                3.8. All Pairs Maximum Flow.
-----cur->rev = rev;-----// ef
                                3.8.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is constructed
-----g[i].push_back(cur);-----// 1d
                                using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the maximum
-----g[adj[i][j].first].push_back(rev);------// 05
                                flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|). NOTE:
Not sure if it works correctly with disconnected graphs.
#include "dinic.cpp"-----// 58
-----ll flow = 0, cost = 0;-----// 68
-----mcmf_edge** back = new mcmf_edge*[n];-----// e5
                                -----// 25
```

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------int l = 0, r = 0;-------// 08 ------int best = -1;-------// 27
------if (adj[u][i] != parent[u] && (best == -1 || sz[adj[u][i]] > sz[best]))
------memset(same, 0, n * sizeof(bool));------// c9 -------best = adj[u][i];-------// 26
------while (l < r) {------// 45 -----rep(i,0,size(adj[u]))------// 92
------for (int i = q.head[v]; i != -1; i = q.e[i].nxt)------// 66 ------part(curhead = adj[u][i]); }------// 88
------rep(i,s+1,n)--------|/ 71 -------while (u != -1) uat.push_back(u), u = parent[head[u]];------// 51
------if (par[i].first == par[s].first && same[i]) par[i].first = s;-----// 97 -------while (v != -1) vat.push_back(v), v = parent[head[v]];--------------// 6d
------int mn = INF. cur = i;-------// 59 -----return res; }------// 91
------cap[cur][i] = mn;------// 8d ------while (head[u] != head[v])-------// 69
------if (cur == 0) break;------// fb -----res = f(res, values.query(loc[head[u]], loc[u])),-----// a4
----return make_pair(par, cap);-------// 62 ------return f(query_upto(u, l), query_upto(v, l)); } };------// 5b
}-----// b3
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {------// 93
                             3.10. Centroid Decomposition.
---if (s == t) return 0;-----// 33
                             #define MAXV 100100------// 86
----int cur = INF, at = s;-----// e7
                             #define LGMAXV 20-----// aa
----while (qh.second[at][t] == -1)------// 42
                             int imp[MAXV][LGMAXV].....// 6d
-----cur = min(cur, gh.first[at].second), at = qh.first[at].first;-----// 8d
                             ----path[MAXV][LGMAXV],-----// 9d
----return min(cur, gh.second[at][t]);-----// 54
                             ----sz[MAXV], seph[MAXV],-----// cf
}------// 46
                             ----shortest[MAXV]:-----// 6b
                             struct centroid_decomposition {------// 99
3.9. Heavy-Light Decomposition.
                             ----int n: vvi adi:-----// e9
#include "../data-structures/segment_tree.cpp"------// 16 ---centroid_decomposition(int _n) : n(_n), adi(n) { }-------// 46
struct HLD {------// 25 ----void add_edge(int a, int b) { adj[a].push_back(b), adj[b].push_back(a); }--// bc
----vvi adj; segment_tree values;--------// 13 ------rep(i,0,size(adj[u])) if (adj[u][i] != p) sz[u] += dfs(adj[u][i], u);--// 78
----HLD(int _n) : n(_n), sz(n, 1), head(n), parent(n, -1), loc(n), adj(n) {----// 1c ------return sz[u]; }-----
-----vi tmp(n, ID); values = segment_tree(tmp); }-------// f0 ----void makepaths(int sep, int u, int p, int len) {-------// 84
----void add_edge(int u, int v) { adj[u].push_back(v), adj[v].push_back(u); }--// 77 ------jmp[u][seph[sep]] = sep, path[u][seph[sep]] = len;-------------// d9
-----values.update(loc[u], c); }------// 50 ------if (adj[u][i] == p) bad = i;------// cf
-----sz[u] += csz(adj[parent[adj[u][i]] = u][i]);-----// c2 ------if (p == sep) swap(adj[u][bad], adj[u].back()), adj[u].pop_back(); }---// 07
-----return sz[u]; }------// 75 ----void separate(int h=0, int u=0) {------// 03
----void part(int u) {------// c3 ------dfs(u,-1); int sep = u;------// b5
```

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------if (sz[*nxt] < sz[sep] && sz[*nxt] > sz[u]/2) {-------// db ---}-----// db
-----sep = *nxt; goto down; }------// 1a ----void process(int u) {------------------------// 85
------seph[sep] = h, makepaths(sep, sep, -1, 0);--------// ed -----ancestor[u] = u;--------------------------// la
-----rep(i,0,size(adj[sep])) separate(h+1, adj[sep][i]); }------// 90 -----rep(i,0,size(adj[u])) {-------// ce
-----rep(h,0,seph[u]+1)------// c5 -----process(v);------// e8
-----shortest[jmp[u][h]] = min(shortest[jmp[u][h]], path[u][h]); }-----// 11 ------uf.unite(u,v);-------------------------// 55
-----int mn = INF/2:------// fe ----}-----// fe -----}
-----rep(h,0,seph[u]+1) mn = min(mn, path[u][h] + shortest[jmp[u][h]]);-----// 3e ------colored[u] = true;------------------------// b9
-----return mn; } };------// 13 ------rep(i,0,size(queries[u])) {-------// d7
                                       -----int v = queries[u][i].first;-----// 89
3.11. Least Common Ancestors, Binary Jumping.
                                       ------if (colored[v1) {------// cb
struct node {-----// 36
                                       -----answers[queries[u][i].second] = ancestor[uf.find(v)];-----// 63
---node *p. *imp[20]:-----// 24
                                       ----int depth;-----// 10
                                       ----node(node *_p = NULL) : p(_p) {-----// 78
                                       ----}------------------// a9
-----depth = p ? 1 + p->depth : 0;-----// 3b
                                       }:-----// le
-----memset(jmp, 0, sizeof(jmp));-----// 64
                                       3.13. Maximum Density Subgraph. Given (weighted) undirected graph G. Binary search density.
-----jmp[0] = p;-----// 64
                                       If q is current density, construct flow network: (S, u, m), (u, T, m + 2q - d_u), (u, v, 1), where m is a
------for (int i = 1; (1<<i) <= depth; i++)------// a8
                                       large constant (larger than sum of edge weights). Run floating-point max-flow. If minimum cut has
-----jmp[i] = jmp[i-1]->jmp[i-1]; } };-----// 3b
                                       empty S-component, then maximum density is smaller than q, otherwise it's larger. Distance between
node* st[100000]:-----// 65
node* lca(node *a, node *b) {-----// 29
                                       valid densities is at least 1/(n(n-1)). Edge case when density is 0. This also works for weighted
----if (!a || !b) return NULL;------// cd
                                       graphs by replacing d_u be the weighted degree, and doing more iterations (if weights are not integers).
----if (a->depth < b->depth) swap(a,b);-----// fe
                                       3.14. Maximum Weighted Independent Set in a Bipartite Graph. This is the same as the
----for (int j = 19; j >= 0; j--)-----// b3
                                       minimum weighted vertex cover. Solve this by constructing a flow network with edges (S, u, w(u))
------while (a->depth - (1<<j) >= b->depth) a = a->jmp[j];------// c\theta
                                       for u \in L, (v, T, w(v)) for v \in R and (u, v, \infty) for (u, v) \in E. The minimum S, T-cut is the answer.
---if (a == b) return a;-----// 08
                                       Vertices adjacent to a cut edge are in the vertex cover.
----for (int j = 19; j >= 0; j--)-----// 11
-------while (a->depth >= (1 << j) \& a -> jmp[j] != b-> jmp[j])------// f\theta
                                                        4. Strings
-----a = a->jmp[j], b = b->jmp[j];-----// d0
                                       4.1. The Knuth-Morris-Pratt algorithm.
----return a->p; }-----// c5
                                       int* compute_pi(const string &t) {-------// a2
3.12. Tarjan's Off-line Lowest Common Ancestors Algorithm.
                                       ----int m = t.size();-----// 8b
#include "../data-structures/union_find.cpp"-----------------// 5e ----int *pit = new int[m + 1];----------------------------// 8e
struct tarjan_olca {-------// 87 ---if (0 <= m) pit[0] = 0;------// 42
-----ancestor = new int[n];------// f2 ----return pit; }------// e8
-----queries = new vii[n];-------// 3e int string_match(const string &s, const string &t) {------// 9e
-----queries[x].push_back(ii(y, size(answers)));-------// a\theta ------if (s[i] == t[j]) {-------------------------// 73
-----queries[y].push_back(ii(x, size(answers)));-------// 14 -----i++; j++;------------------------------// 7e
-----answers.push_back(-1);------// ca -----if (j == m) {------// de
```

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-----// or j = pit[j];------// ce ------return res;------
------else if (j > 0) j = pit[j];------// 43
-----else i++; }-----// b8
                           4.4. Aho-Corasick Algorithm.
----delete[] pit; return -1; }------// e3
                           struct aho_corasick {------// 78
                           ----struct out_node {------// 3e
4.2. The Z algorithm.
                           -----string keyword; out_node *next;-----// f0
int* z_values(const string &s) {------// 4d
                           -----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
----int n = size(s):-----// 97
                           ----}:-------// b9
----int* z = new int[n];------// c4
                           ----struct qo_node {------// 40
----int l = 0, r = 0;------// 1c
                           -----map<char, go_node*> next;------// 6b
---z[0] = n;
                           -----out_node *out; go_node *fail;-----// 3e
----rep(i,1,n) {------// b2
                           -----go_node() { out = NULL; fail = NULL; }-----// 0f
----};------// c0
-----if (i > r) {------// 6d
                           ----qo_node *qo;------// b8
-----l = r = i;------// 24
                           ----aho_corasick(vector<string> keywords) {-----------------------// 4b
------while (r < n \&\& s[r - l] == s[r]) r++;
                           -----go = new go_node();-----// 77
----z[i] = r - l; r--;-----// 07
                           -----iter(k, keywords) {------// f2
-----} else if (z[i - l] < r - i + 1) z[i] = z[i - l];------// 6f
                           -----qo_node *cur = qo;-----// a2
-----else {------// a8
                           -----iter(c, *k)-----// 6e
-----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 97
-----(cur->next[*c] = new go_node());-----// af
-----z[i] = r - l; r--; } }-----// 13
                           -----cur->out = new out_node(*k, cur->out);-----// 3f
----return z;------// 78
                           }-----// 16
                           -----queue<qo_node*> q;-----// 2c
4.3. Suffix Array. An O(n \log^2 n) construction of a Suffix Tree.
                           -----iter(a, go->next) q.push(a->second);-----// db
----string s; int n; vvi P; vector<entry> L; vi idx;------// b6 ------go_node *s = a->second;-------// 55
-----L = vector<entry>(n), P.push_back(vi(n)), idx = vi(n);-------// 12 ------go_node *st = r->fail;-------------------// 53
------while (st &\& st->next.find(a->first) == st->next.end())------// \theta e
-----P.push_back(vi(n));-------// 53 ------if (!st) st = go;-------// θb
-----rep(i,0,n)-------// 6f ------s->fail = st->next[a->first];------// c1
-----L[L[i].p = i].nr = ii(P[stp - 1][i],-------// e2 ------if (s->fail) {-----------------------// 98
------i + cnt < n ? P[stp - 1][i + cnt] : -1);-------// 43 ------if (!s->out) s->out = s->fail->out;------// ad
------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;------// 55 ------out->next = s->fail->out;--------// 62
------int res = 0;-------// de ---}-----// de
------for (int k = size(P) - 1; k >= 0 && x < n && y < n; k--)-------// fe ------vector<string> res;-------------------------------// 79
```

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-----go_node *cur = go;------// 85 ----int sz, last;---------// 7d
-----iter(c, s) {------// 57 ----string s;-----// f2
------while (cur \&\& cur->next.find(*c) == cur->next.end())-------// df ----suffix_automaton() : len(MAXL*2), link(MAXL*2), occur(MAXL*2), next(MAXL*2),
------if (!cur) cur = go;------// 92 ----void clear(){ sz = 1; last = len[0] = 0; link[0] = -1; next[0].clear();----// aa
-----if (!cur) cur = qo;-------// 01 ----bool issubstr(string other){-------// 3b
------for (out_node *out = cur->out; out; out = out->next)------// d7 ------for(int i = 0, cur = 0; i < size(other); ++i){------// 7f
-----res.push_back(out->keyword);------// 7c -----if(cur == -1) return false; cur = next[cur][other[i]]; }------// 54
------}-----return true; }------// 1a
------return res;-------// 6b ----void extend(char c){ int cur = sz++; len[cur] = len[last] + 1;------// 1d
----}-----next[cur].clear(); isclone[cur] = false; int p = last;------// a9
-----if(p == -1){ link[cur] = 0; }-----// 18
4.5. eerTree.
                                              -----else{ int q = next[p][c];-----// 34
#define MAXN 100100-----// 29
                                              ------if(len[p] + 1 == len[q]){ link[cur] = q; }------// 4d
#define SIGMA 26-----// e2
                                              ------else { int clone = sz++; isclone[clone] = true;-----// 57
#define BASE 'a'-----// a1
                                             -----len[clone] = len[p] + 1;------// 8c
char *s = new char[MAXN];....// db
                                             -----link[clone] = link[q]; next[clone] = next[q];-----// 76
struct state {-----// 33
                                              ------for(; p != -1 && next[p].count(c) && next[p][c] == q; p = link[p]){
----int len, link, to[SIGMA];-------// 24
                                              -----next[p][c] = clone; }-----// 32
} *st = new state[MAXN+2];-----// 57
                                              -----link[q] = link[cur] = clone;-----// 73
struct eertree {------// 78
                                              ------} } last = cur; }-----// b9
----int last, sz, n;------// ba
                                              ----void count(){------// e7
----eertree() : last(1), sz(2), n(0) {------// 83
                                              -----cnt=vi(sz, -1); stack<ii> S; S.push(ii(0,0)); map<char, int>::iterator i;// 56
-----st[0].len = st[0].link = -1;------// 3f
                                              ------while(!S.empty()){------// 4c
-----st[1].len = st[1].link = 0; }-----// 34
                                              -----ii cur = S.top(); S.pop();-----// 67
----int extend() {-------// c2
                                              -----if(cur.second){-----// 78
-----char c = s[n++]; int p = last;-----// 25
                                              ------for(i = next[cur.first].begin();i != next[cur.first].end();++i){
------while (n - st[p].len - 2 < 0 \mid \mid c \mid = s[n - st[p].len - 2]) p = st[p].link;
                                              -----cnt[cur.first] += cnt[(*i).second]; } }-----// da
-----if (!st[p].to[c-BASE]) {------// 82
                                              -----else if(cnt[cur.first] == -1){------// 99
-----int q = last = sz++:-----// 42
                                              ------cnt[cur.first] = 1; S.push(ii(cur.first, 1));-----// bd
------st[p].to[c-BASE] = a:-----// fc
                                              -----for(i = next[cur.first].begin();i != next[cur.first].end();++i){
-----st[q].len = st[p].len + 2;-----// c5
                                              ------S.push(ii((*i).second, 0)); } } } }-----// 61
-----do { p = st[p].link;-----// 04
                                              ----string lexicok(ll k){------// 8b
-----} while (p != -1 \&\& (n < st[p].len + 2 || c != s[n - st[p].len - 2]));
                                              ------int st = 0; string res; map<char,int>::iterator i;------// cf
-----if (p == -1) st[q].link = 1;------// 77
                                              ------while(k){ for(i = next[st].begin(); i != next[st].end(); ++i){------// 69
------else st[q].link = st[p].to[c-BASE];------// 6a
                                              ------if(k <= cnt[(*i).second]){ st = (*i).second; -----// ec
-----return 1; }-----// 29
                                              -----res.push_back((*i).first); k--; break;-----// 63
-----last = st[p].to[c-BASE];-----// 42
                                              -----} else { k -= cnt[(*i).second]; } } }-----// ee
-----return 0: } }:-----// ec
                                              -----return res; }-----// θb
                                              ----void countoccur(){------// ad
4.6. Suffix Automaton. Minimum automata that accepts all suffixes of a string with O(n) construc-
                                              ------for(int i = 0; i < sz; ++i){ occur[i] = 1 - isclone[i]; }-----// 1b
tion. The automata itself is a DAG therefore suitable for DP, examples are counting unique substrings,
                                              -----vii states(sz):-----// dc
occurrences of substrings and suffix.
                                              ------for(int i = 0; i < sz; ++i){ states[i] = ii(len[i],i); }------// 97
// TODO: Add longest common subsring-----// 0e
                                              -----sort(states.begin(), states.end());------// 8d
const int MAXL = 100000;-----// 31
                                              -----for(int i = size(states)-1; i >= 0; --i){ int v = states[i].second; <math>---//a4
struct suffix_automaton {------// e0
                                              ------if(link[v] != -1) { occur[link[v]] += occur[v]; } } }-----// cc
----vi len, link, occur, cnt;-----// 78
                                              }:-----// 32
----vector<map<char,int> > next;------// 90
----vector<bool> isclone:-----// 7b
----ll *occuratleast;-----// f2
```

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-----*c = new cpx[n], *a = new cpx[len],--------// 4e ----if (l == 1) { y[0] = x[0].inv(); return; }------// 71
-----*b = new cpx[len];-------// 30 ---inv(x, y, l>>1);------// bd
----rep(i,0,n) {-------// ff void sqrt(Num x[], Num y[], int l) {-------// 44
------if (inv) x[i] /= cpx(n);------// b1 ----sqrt(x, y, l>>1);-----// a7
}------// c6 ----rep(i,0,l<<1) T2[i] = T1[i] * T2[i];-------// 9f
                       ----ntt(T2, l<<1, true);------// eb
                       5.14. Number-Theoretic Transform.
#include "../mathematics/primitive_root.cpp"--------------// 8c 5.15. Tridiagonal Matrix Algorithm. Solves a tridiagonal system of linear equations a_i x_{i-1} +
int mod = 998244353, q = primitive_root(mod),....// 9c
                       b_i x_i + c_i x_{i+1} = d_i where a_1 = c_n = 0. Beware of numerical instability.
----ginv = mod_pow(g, mod-2, mod), inv2 = mod_pow(2, mod-2, mod); -------//75
                       #define MAXN 5000-----// f7
#define MAXN (1<<22)-----// 94
                       long double A[MAXN], B[MAXN], C[MAXN], D[MAXN], X[MAXN]; ------// d8
struct Num {-----// c5
                       void solve(int n) {------// 01
----int x;------// 02
                       ----C[0] /= B[0]; D[0] /= B[0];-----// 94
----Num(ll _x=0) { x = (_x \mod + \mod) \mod; }-----// 1b
                       ----rep(i.1.n-1) C[i] /= B[i] - A[i]*C[i-1]:-----// 6b
----Num operator +(const Num &b) { return x + b.x; }-----// 08
                       ----rep(i,1,n) D[i] = (D[i] - A[i] * D[i-1]) / (B[i] - A[i] * C[i-1]); ------// 33
----Num operator -(const Num &b) const { return x - b.x; }------// 89
                       ---X[n-1] = D[n-1];
----Num operator *(const Num &b) const { return (ll)x * b.x; }------// e3
                       ----for (int i = n-2; i>=0; i--) X[i] = D[i] - C[i] * X[i+1]; }------// ad
----Num operator /(const Num &b) const { return (ll)x * b.inv().x; }------// 2a
                       5.16. Mertens Function. Mertens function is M(n) = \sum_{i=1}^n \mu(i). Let L \approx (n \log \log n)^{2/3} and the
----Num inv() const { return mod_pow((ll)x, mod-2, mod); }-----// d3
                       algorithm runs in O(n^{2/3}). Can also be easily changed to compute the summatory \Phi.
----Num pow(int p) const { return mod_pow((ll)x, p, mod); }-----// d5
} T1[MAXN], T2[MAXN];-----// d5
                       #define L 9000000-----// 27
                       int mob[L], mer[L];.....// f1
void ntt(Num x[], int n, bool inv = false) {------// 24
------if (i < j) swap(x[i], x[j]);----------// eb ----if (mem.find(n) != mem.end()) return mem[n];-------// 79
-x[i + mx] = x[i] - t;
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                                         double line_segment_distance(L(a,b), L(c,d)) {-----// f3
-----}-----// aa
-----mer[i] = mob[i] + mer[i-1]; } }-----// 3b
                                         ----double x = INFINITY;-----// 64
                                          ----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c);------// a5
5.17. Numbers and Sequences. Some random prime numbers: 1031, 32771, 1048583, 33554467,
                                          ----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true));-----// 23
1073741827, 34359738421, 1099511627791, 35184372088891, 1125899906842679, 36028797018963971.
                                          ----else if (abs(c - d) < EPS) x = abs(c - closest_point(a, b, c, true)); -----// 53
                                          ----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&-----// 6d
                 6. Geometry
                                          -----(ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0;-----// bf
6.1. Primitives.
                                          ----else {------// e1
#define C(p0, r) P(p0), double r-----// f1 -----x = min(x, abs(c - closest_point(a,b, c, true)));-----// 81
#define PP(pp) pair<point, point> &pp-------// e5 -----x = min(x, abs(d - closest_point(a,b, d, true)));------// e4
typedef complex<double> point;-----// 6a ---}-----// 6a
point rotate(P(p), double radians = pi / 2, P(about) = point(\theta,\theta)) {------// 23 bool intersect(L(a, b), L(p, q), point &res, bool segment = false) {-------// d2
----return (p - about) * exp(point(0, radians)) + about; }-----// 25 ----// NOTE: check for parallel/collinear lines before calling this function---// 1b
point reflect(P(p), L(about1, about2)) {------// 50 ---point r = b - a, s = q - p;-----// 34
----return conj(z / w) * w + about1; }------// 83 ----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// e4
point normalize(P(p), double k = 1.0) {------// 5f ----res = a + t * r;------// 47
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// b3
                                         -----// cc
double angle(P(a), P(b), P(c)) {------// 61
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }-----// c7
                                         6.3. Circles.
double signed_angle(P(a), P(b), P(c)) {------// 4a
                                         #include "primitives.cpp"-----// e0
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 40
                                         int intersect(C(A, rA), C(B, rB), point & res1, point & res2) {------// 3b
double angle(P(p)) { return atan2(imag(p), real(p)); }-----// e6
                                          ----double d = abs(B - A);-----// 5c
point perp(P(p)) { return point(-imag(p), real(p)); }-----// d9
                                          ----if ((rA + rB) < (d - EPS) || d < abs(rA - rB) - EPS) return 0;-------// 39
double progress(P(p), L(a, b)) {-----// b3
                                          ----double a = (rA*rA - rB*rB + d*d) / 2 / d, h = sqrt(rA*rA - a*a); ------// 9b
----if (abs(real(a) - real(b)) < EPS)------// 5e
                                          ----point v = \text{normalize}(B - A, a), u = \text{normalize}(\text{rotate}(B-A), h);
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));-----// 5e
                                          ----res1 = A + v + u, res2 = A + v - u;------// 24
----else return (real(p) - real(a)) / (real(b) - real(a)); }-----// 31
                                          ----if (abs(u) < EPS) return 1; return 2;-------// 82
-----// 53
                                            -----// bb
-----// 46
                                         int intersect(L(A, B), C(0, r), point & res1, point & res2) {------// 0e
6.2. Lines.
                                          ---- double h = abs(0 - closest_point(A, B, 0));-----// 24
#include "primitives.cpp"------// e0 ---- if(r < h - EPS) return 0;-----------// df
------// 85 ---- point H = proj(0 - A, B - A) + A, v = normalize((B - A), sqrt(r*r - h*h));// 19
point closest_point(L(a, b), P(c), bool segment = false) {-------// f2 int tangent(P(A), C(0, r), point & res1, point & res2) {-------// aa
-----if (dot(b - a, c - b) > 0) return b;-------// 88 ----if (d < r - EPS) return 0;--------// ca
------if (dot(a - b, c - a) > 0) return a;--------// 75 ----double alpha = asin(r / d), L = sqrt(d*d - r*r);--------// 3f
----}------------------// ce ----v = normalize(v, L);--------// 3f
----double t = dot(c - a, b - a) / norm(b - a);-------// 62 ----res1 = A + rotate(v, alpha); res2 = A + rotate(v, -alpha);------// be
----return a + t * (b - a);------// 6e ----if (abs(r - d) < EPS || abs(v) < EPS) return 1;-----// bb
}------// 8c ----return 2;------// b9
```

```
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----return r * acos(cos(pLat) * cos(pLong - gLong) +------// e3 ------return point3d(x / k, y / k, z / k); }------// 58
-----// 60 ------return x * p.x + y * p.y + z * p.z; }------// 09
}------// 3f ---point3d operator*(P(p)) const {-------// 4f
                                                      -----return point3d(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x); }-----// ed
6.8. Triangle Circumcenter. Returns the unique point that is the same distance from all three
                                                      ----double length() const {------// 3e
points. It is also the center of the unique circle that goes through all three points.
                                                      -----return sqrt(*this % *this); }------// 05
#include "primitives.cpp"-----// e0
                                                      ----double distTo(P(p)) const {------// dd
point circumcenter(point a, point b, point c) {-----// 76
                                                      -----return (*this - p).length(); }-----// 57
----b -= a, c -= a:-----// 41
                                                      ----double distTo(P(A), P(B)) const {------// bd
----return a + perp(b * norm(c) - c * norm(b)) / 2.0 / cross(b, c);------// 7a
                                                      -----// A and B must be two different points-----// 4e
}-----// c3
                                                      -----return ((*this - A) * (*this - B)).length() / A.distTo(B); }------// 6e
                                                      ----point3d normalize(double k = 1) const {------// db
6.9. Closest Pair of Points.
                                                      -----// length() must not return 0-----// 3c
#include "primitives.cpp"-----// e0
                                                      -----return (*this) * (k / length()); }------// d4
-----// 85
                                                      ----point3d getProjection(P(A), P(B)) const {------// 86
struct cmpx { bool operator ()(const point &a. const point &b) {------// 01
                                                      -----point3d v = B - A;-----// 64
-----return abs(real(a) - real(b)) > EPS ?-----// e9
                                                      -----return A + v.normalize((v % (*this - A)) / v.length()); }------// 53
-----real(a) < real(b) : imag(a) < imag(b); } };-----// 53
                                                      ----point3d rotate(P(normal)) const {------// 55
struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
                                                      -----// normal must have length 1 and be orthogonal to the vector-----// eb
----return abs(imag(a) - imag(b)) > EPS ?-----// θb
                                                           return (*this) * normal; }------// 5c
-----imag(a) < imag(b) : real(a) < real(b); } };-----// a4
                                                      ----point3d rotate(double alpha, P(normal)) const {-------// 21
double closest_pair(vector<point> pts) {------// f1
                                                      -----return (*this) * cos(alpha) + rotate(normal) * sin(alpha); }-----// 82
----sort(pts.begin(), pts.end(), cmpx());-----// 0c
                                                      ----point3d rotatePoint(P(0), P(axe), double alpha) const{-------// 7a
----set<point, cmpy> cur;-----// bd
                                                      ------point3d Z = axe.normalize(axe % (*this - 0));------// ba
----set<point, cmpy>::const_iterator it, jt;-----// a6
                                                      -----return 0 + Z + (*this - 0 - Z).rotate(alpha, 0); }------// 38
----double mn = INFINITY:-----// f9
                                                      ----bool isZero() const {------// 64
----for (int i = 0, l = 0; i < size(pts); i++) {------// ac
                                                      -----return abs(x) < EPS && abs(y) < EPS && abs(z) < EPS; }------// 15
------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);------// 8b
                                                      ----bool isOnLine(L(A, B)) const {------// 30
-----it = cur.lower_bound(point(-INFINITY, imag(pts[i]) - mn));-----// fc
                                                      -----return ((A - *this) * (B - *this)).isZero(); }------// 58
-----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));------// 39
                                                      ----bool isInSegment(L(A, B)) const {------// f1
------while (it != jt) mn = min(mn, abs(*it - pts[i])), it++;------// 09
                                                      -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < EPS; }------// d9
-----cur.insert(pts[i]): }-----// 82
                                                      ----bool isInSegmentStrictly(L(A, B)) const {------// θe
----return mn; }-----// 4c
                                                      -----return isOnLine(A, B) && ((A - *this) % (B - *this)) < -EPS: }------// ba
                                                      ----double getAngle() const {------// 0f
6.10. 3D Primitives.
                                                      -----return atan2(y, x); }------// 40
#define P(p) const point3d &p-----// a7
                                                      ----double getAngle(P(u)) const {------// d5
#define L(p0, p1) P(p0), P(p1)-----// Of
                                                      -----return atan2((*this * u).length(), *this % u); }-----// 79
#define PL(p0, p1, p2) P(p0), P(p1), P(p2)-----// 67
                                                      ----bool isOnPlane(PL(A, B, C)) const {------// 8e
struct point3d {-----// 63
                                                      -----return abs((A - *this) * (B - *this) % (C - *this)) < EPS; } };------// 74
----double x, y, z;-----// e6
                                                      <mark>int</mark> line_line_intersect(L(A, B), L(C, D), point3d &0){-----------------// dc
----point3d() : x(0), y(0), z(0) {}------// af
                                                      ----if (abs((B - A) * (C - A) % (D - A)) > EPS) return 0;-------// 6a
----point3d(double _x, double _y, double _z) : x(_x), y(_y), z(_z) {}------// fc
                                                      ----if (((A - B) * (C - D)).length() < EPS)------// 79
----point3d operator+(P(p)) const {------// 17
                                                      -----return A.isOnLine(C, D) ? 2 : 0;-----// 09
-----return point3d(x + p.x, y + p.y, z + p.z); }------// 8e
                                                      ----point3d normal = ((A - B) * (C - B)).normalize();-----// bc
----point3d operator-(P(p)) const {-------// fb
                                                      ----double s1 = (C - A) * (D - A) % normal;------// 68
------return point3d(x - p.x, y - p.y, z - p.z); }------// 83
                                                      ----0 = A + ((B - A) / (s1 + ((D - B) * (C - B) % normal))) * s1;------// 56
----point3d operator-() const {------// 89
                                                      ----return 1; }-----// a7
-----return point3d(-x, -y, -z); }-----// d4
                                                      int line_plane_intersect(L(A, B), PL(C, D, E), point3d & O) {------// 09
----point3d operator*(double k) const {------// 4d
                                                      ----double V1 = (C - A) * (D - A) % (E - A);------// c1
-----return point3d(x * k, y * k, z * k); }-----// fd
                                                      ----double V2 = (D - B) * (C - B) % (E - B);------// 29
----point3d operator/(double k) const {------// 95
```

```
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-----adj[-clauses[i].first + n].push_back(clauses[i].second + n);------// eb ----exact_cover(int _rows, int _cols) : rows(_rows), cols(_cols), head(NULL) {-// b6
------if (clauses[i].first != clauses[i].second)--------// bc -----arr = new bool*[rows];--------// cf
----pair<union_find, vi> res = scc(adj);------// dd
------int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n); -//4f ------rep(i,0,rows+1) {----------------------------------//76
-----if (cur == 0) continue;-------// cd ------ptr[i] = new node*[cols];-------// eb
------if (p == o) return false;---------// d0 ------rep(j,0,cols)--------// cd
-----truth[cur + n] = truth[p];-------// 50 ------else ptr[i][j] = NULL;------// d2
------if (truth[p] == 1) all_truthy.push_back(cur);--------// 55 -----rep(i,0,rows+1) {-----------------------// fc
----}-----rep(j,0,cols) {-----------------// 51
-----while (true) {------// fc
7.2. Stable Marriage.
                               -----if (ni == rows + 1) ni = 0;------// 4c
vi stable_marriage(int n, int** w, int** w) {------// e4
                               -----if (ni == rows || arr[ni][j]) break;-----// 8d
----queue<int> q;------// f6
                               -------++ni;------// 68
----vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));------// c3
                               -----}------------------------// ad
----rep(i,0,n) rep(j,0,n) inv[i][w[i][j]] = j;------// f1
                               -----ptr[i][j]->d = ptr[ni][j];------// 84
----rep(i,0,n) q.push(i);-----// d8
                               -----ptr[ni][j]->u = ptr[i][j];-----// 66
----while (!q.empty()) {------// 68
                               -----while (true) {------// 7f
------int curm = q.front(); q.pop();------// e2
                               -----if (nj == cols) nj = 0;-----// de
------for (int &i = at[curm]; i < n; i++) {-------// 7e
                               ------if (i == rows || arr[i][nj]) break;-----// 4c
-----int curw = m[curm][i];-----// 95
                               ------++nj;-----// c5
-----if (eng[curw] == -1) { }-----// f7
                               ------else if (inv[curw][curm] < inv[curw][eng[curw]])------// d6
                               -----ptr[i][j]->r = ptr[i][nj];-----// 60
-----a.push(eng[curw]):-----// 2e
                               -----ptr[i][nj]->l = ptr[i][j];-----// 82
-----else continue;-----// 1d
                               -----res[eng[curw] = curm] = curw, ++i; break;------// a1
                               ·····}
-----}-----// c4
                               -----head = new node(rows, -1);------// 66
----}-----// 3d
                               ------head->r = ptr[rows][0];------// 3e
----return res:------// 42
                               -----ptr[rows][0]->l = head;------// 8c
}-----// bf
                               -----head->l = ptr[rows][cols - 1];------// 6a
                               -----ptr[rows][cols - 1]->r = head;-----// c1
7.3. Algorithm X.
                               -----rep(j,0,cols) {------// 92
bool handle_solution(vi rows) { return false; }------// 63
                               -----int cnt = -1;-----// d4
struct exact_cover {------// 95
                               -----rep(i,0,rows+1)-----// bd
----struct node {------// 7e
                               -----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];-----// f3
-----node *l, *r, *u, *d, *p;-----// 19
                               -----ptr[rows][j]->size = cnt;-----// c2
------int row, col, size;-----// ae
                               ------}-----// b9
-----node(int _row, int _col) : row(_row), col(_col) {------// c9
                               -----rep(i,0,rows+1) delete[] ptr[i];------// a5
-----size = 0; l = r = u = d = p = NULL; }-----// c3
                               -----delete[] ptr:-----// 72
----}:------// c1
                               ----}------// 19
----int rows, cols, *sol;------// 7b
```

```
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----#define COVER(c, i, j) N------------------------// 91 ----return ii(mu, lam);-----------------------------------// b4
------for (node *j = i->r; j != i; j = j->r) \------// 26
                                       int intToDay(int jd) { return jd % 7; }------// 89
-----j->d->u = j->u, j->u->d = j->d, j->p->size--;-----// c1
                                       int dateToInt(int y, int m, int d) {------// 96
----#define UNCOVER(c, i, j) \sqrt{\phantom{a}}-----// 89
                                       ----return 1461 * (y + 4800 + (m - 14) / 12) / 4 +------// a8
------for (node *i = c->u; i != c; i = i->u) \[ \] ------// f0
                                       ------367 * (m - 2 - (m - 14) / 12 * 12) / 12 ------// d1
                                       -----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +------// be
-----d - 32075;-----// e0
}-----// fa
void intToDate(int jd, int &y, int &m, int &d) {------// a1
----bool search(int k = 0) {------// f9
                                       ----int x, n, i, j;------// 00
-----if (head == head->r) {------// 75
                                       ---x = id + 68569;
-----vi res(k);-----// 90
                                       ----n = 4 * x / 146097;-----// 2f
-----rep(i,0,k) res[i] = sol[i];-----// 2a
                                       ---x = (146097 * n + 3) / 4;
-----sort(res.begin(), res.end());-----// 63
                                       ---i = (4000 * (x + 1)) / 1461001; ------// 0d
-----return handle_solution(res);-----// 11
                                       ----x -= 1461 * i / 4 - 31;-----// 09
-----}-----// 3d
                                       ----j = 80 * x / 2447;
-----node *c = head->r, *tmp = head->r;------// a3
                                       ---d = x - 2447 * j / 80;
-----for (; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp; ---//41
-----if (c == c->d) return false;-----//
                                       ---m = j + 2 - 12 * x;
-----COVER(c, i, j);-----// f6
                                       ---y = 100 * (n - 49) + i + x;
------bool found = false;-----// 8d
------for (node *r = c->d; !found && r != c; r = r->d) {------// 78
-----sol[k] = r->row;------// cθ
                                       7.7. Simulated Annealing. An example use of Simulated Annealing to find a permutation of length
-----for (node *j = r - r; j != r; j = j - r) { COVER(j - p, a, b); } -----// f9
                                       n that maximizes \sum_{i=1}^{n-1} |p_i - p_{i+1}|.
-----found = search(k + 1);-----// fb
                                       double curtime() { return static_cast<double>(clock()) / CLOCKS_PER_SEC; }-----// 9d
------for (node *j = r > l; j = j > l) { UNCOVER(j > p, a, b); }----// 87
                                       int simulated_annealing(int n, double seconds) {------// 54
                                       ----default_random_engine rng;------// 67
-----UNCOVER(c, i, j);-----// a7
                                       ----uniform_real_distribution<double> randfloat(0.0, 1.0);------// ed
----return found;-----// c0
                                       ----uniform_int_distribution<int> randint(0, n - 2);------------------------// bb
                                       ----// random initial solution------// 01
                                       ----vi sol(n);------// 1c
                                       ----rep(i,0,n) sol[i] = i + 1;------// 33
7.4. nth Permutation.
                                       ----random_shuffle(sol.begin(), sol.end());-----// ea
vector<int> nth_permutation(int cnt, int n) {------// 78
                                       ----// initialize score------// 28
----vector<int> idx(cnt), per(cnt), fac(cnt);-----// 9e
                                       ----int score = 0;------// 7d
---rep(i,0,cnt) idx[i] = i;-----// bc
                                       ----rep(i,1,cnt+1) fac[i - 1] = n % i, n /= i;------// 2b
                                       ----int iters = 0;------// 0b
----for (int i = cnt - 1; i >= 0; i--)-----// f9
                                       ----double T0 = 100.0, T1 = 0.001,-----// 5c
-----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// ee
                                        -----/ progress = 0, temp = T0,-----// 3a
----return per;-----// ab
                                       -----/ starttime = curtime();------// d6
                                       ----while (true) {------// 46
7.5. Cycle-Finding.
                                       -----if (!(iters & ((1 << 4) - 1))) {------// 5d
----h = x0;------// eb
----while (t != h) t = f(t), h = f(h), mu++;-----------------// 9d ------int a = randint(rng);-------------------------------// c3
----h = f(t);-----// 00 ------// compute delta for mutation------// 84
```

## 7.8. Fast Input Reading.

```
#labeled rooted trees
 #labeled unrooted trees
                                                                          \sum_{i=1}^{n} i^3 = n^2 (n+1)^2 / 4
\sum_{i=1}^{n} i^2 = n(n+1)(2n+1)/6
!n = n \times !(n-1) + (-1)^n
                                                                          !n = (n-1)(!(n-1)+!(n-2))
\sum_{i=1}^{n} \binom{n}{i} F_i = F_{2n}
                                                                          \sum_{i} \binom{n-i}{i} = F_{n+1}
\sum_{k=0}^{n} {k \choose m} = {n+1 \choose m+1}
                                                                          \sum_{d|n} \phi(d) = n
a \equiv b \pmod{x, y} \Rightarrow a \equiv b \pmod{\operatorname{lcm}(x, y)}
ac \equiv bc \pmod{m} \Rightarrow a \equiv b \pmod{\frac{m}{\gcd(c,m)}}
                                                                         (\sum_{d|n} \sigma_0(d))^2 = \sum_{d|n} \sigma_0(d)^3
                                                                          \gcd(n^a - 1, n^b - 1) = n^{\gcd(a,b)} - 1
p \text{ prime } \Leftrightarrow (p-1)! \equiv -1 \pmod{p}
\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x} - 1}{p_i^x - 1}
                                                                         \sigma_0(n) = \prod_{i=0}^r (a_i + 1)
\sum_{k=0}^{m} (-1)^k \binom{n}{k} = (-1)^m \binom{n-1}{m}
                                                                          \sum_{i=1}^{n} 2^{\omega(i)} = O(n \log n)
2^{\omega(n)} = O(\sqrt{n})
d = v_i t + \frac{1}{2} a t^2
                                                                          d = \frac{v_i + v_f}{2}t
v_f = v_i + at
```

## 7.10. The Twelvefold Way. Putting n balls into k boxes.

Balls	same	distinct	same	distinct	
Boxes	same	same	distinct	distinct	Remarks
-	$p_k(n)$	$\sum_{i=0}^{k} \begin{Bmatrix} n \\ i \end{Bmatrix}$	$\binom{n+k-1}{k-1}$	$k^n$	$p_k(n)$ : #partitions of n into $\leq k$ positive parts
$\mathrm{size} \geq 1$	p(n,k)	$\left\{ {n\atop k} \right\}$	$\binom{n-1}{k-1}$	$k! \begin{Bmatrix} n \\ k \end{Bmatrix}$	p(n,k): #partitions of n into k positive parts
$\mathrm{size} \leq 1$	$[n \leq k]$	$[n \le k]$	$\binom{k}{n}$	$n!\binom{k}{n}$	[ $cond$ ]: 1 if $cond = true$ , else 0

#### 7.9. Bit Hacks.

```
Catalan  C_0 = 1, C_n = \frac{1}{n+1} {2n \choose n} = \sum_{i=0}^{n-1} C_i C_{n-i-1} = \frac{4n-2}{n+1} C_{n-1}  Stirling 1st kind  \begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1, \begin{bmatrix} n \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ n \end{bmatrix} = 0, \begin{bmatrix} n \\ k \end{bmatrix} = (n-1) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}  #perms of the stirling 2nd kind  \begin{cases} n \\ 1 \end{bmatrix} = \begin{cases} n \\ n \end{bmatrix} = 1, \begin{cases} n \\ k \end{bmatrix} = k \begin{cases} n-1 \\ k \end{bmatrix} + \begin{cases} n-1 \\ k-1 \end{bmatrix}  #ways to the stirling 2nd kind  \begin{cases} n \\ 1 \end{bmatrix} = \begin{cases} n \\ n \end{bmatrix} = 1, \begin{cases} n \\ k \end{bmatrix} = (k+1) \begin{pmatrix} n \\ k \end{pmatrix} = (k+1) \begin{pmatrix} n-1 \\ k \end{bmatrix} + (n-k) \begin{pmatrix} n-1 \\ k-1 \end{pmatrix}  #perms of the stirling 2nd Order  \begin{cases} n \\ k \end{bmatrix} = (k+1) \begin{pmatrix} n-1 \\ k \end{bmatrix} + (2n-k-1) \begin{pmatrix} n-1 \\ k-1 \end{pmatrix} + (2n-k-1
```

#perms of n objs with exactly k cycles

#ways to partition n objs into k nonempty sets

#perms of n objs with exactly k ascents

#perms of 1, 1, 2, 2, ..., n, n with exactly k ascents #partitions of 1..n (Stirling 2nd, no limit on k)

## 8. Useful Information

#### 9. Misc

# 9.1. Debugging Tips.

- Stack overflow? Recursive DFS on tree that is actually a long path?
- Floating-point numbers
  - Getting NaN? Make sure acos etc. are not getting values out of their range (perhaps 1+eps).
  - Rounding negative numbers?
  - Outputting in scientific notation?
- Wrong Answer?
  - Read the problem statement again!
  - Are multiple test cases being handled correctly? Try repeating the same test case many times.
  - Integer overflow?
  - Think very carefully about boundaries of all input parameters
  - Try out possible edge cases:
    - \*  $n = 0, n = -\overline{1}, n = 1, n = 2^{31} 1 \text{ or } n = -2^{31}$
    - \* List is empty, or contains a single element
    - \* n is even, n is odd
    - \* Graph is empty, or contains a single vertex
    - \* Graph is a multigraph (loops or multiple edges)
    - \* Polygon is concave or non-simple
  - Is initial condition wrong for small cases?
  - Are you sure the algorithm is correct?
  - Explain your solution to someone.
  - Are you using any functions that you don't completely understand? Maybe STL functions?
  - Maybe you (or someone else) should rewrite the solution?
- Run-Time Error?
  - Is it actually Memory Limit Exceeded?

#### 9.2. Solution Ideas.

- Dynamic Programming
  - Drop a parameter, recover from others
  - Swap answer and a parameter
  - Parsing CFGs: CYK Algorithm
  - Optimizations
    - \* Convex hull optimization
      - $\cdot \operatorname{dp}[i] = \min_{j < i} \{\operatorname{dp}[j] + b[j] \times a[i]\}$
      - $b[j] \geq b[j+1]$
      - · optionally  $a[i] \le a[i+1]$
      - $O(n^2)$  to O(n)
    - \* Divide and conquer optimization
      - $dp[i][j] = \min_{k < j} \{dp[i-1][k] + C[k][j]\}$
      - $\cdot \ A[i][j] \le A[i][j+1]$
      - $O(kn^2)$  to  $O(kn\log n)$
      - · sufficient:  $C[a][c] + C[b][d] \le C[a][d] + C[b][c],$  $a \le b \le c \le d$  (QI)

- \* Knuth optimization
  - $dp[i][j] = \min_{i < k < j} \{dp[i][k] + dp[k][j] + C[i][j]\}$
  - $\cdot \ A[i][j-1] \leq A[i][j] \leq A[i+1][j]$
  - $\cdot O(n^3)$  to  $O(n^2)$
  - · sufficient: QI and  $C[b][c] \le C[a][d], a \le b \le c \le d$
- Greedy
- Randomized
- Optimizations
  - Use bitset (/64)
  - Switch order of loops (cache locality)
- Process queries offline
  - Mo's algorithm
- Square-root decomposition
- Precomputation
- Efficient simulation
  - Mo's algorithm
  - Sqrt decomposition
  - Store  $2^k$  jump pointers
- Data structure techniques
  - Sart buckets
  - Store  $2^k$  jump pointers
  - $-2^k$  merging trick
- Counting
  - Inclusion-exclusion principle
  - Generating functions
- Graphs
  - Can we model the problem as a graph?
  - Can we use any properties of the graph?
  - Strongly connected components
  - Cycles (or odd cycles)
  - Bipartite (no odd cycles)
    - \* Bipartite matching
    - \* Hall's marriage theorem
    - $* \ \, {\rm Stable \ Marriage}$
  - $\ {\rm Cut \ vertex/bridge}$
  - Biconnected components
  - Degrees of vertices (odd/even)
  - Trees
    - \* Heavy-light decomposition
    - \* Centroid decomposition
    - \* Least common ancestor
  - Eulerian path/circuit
  - Chinese postman problem
  - Topological sort
  - (Min-Cost) Max Flow
  - Min Cut
    - \* Maximum Density Subgraph
  - Huffman Coding
  - Min-Cost Arborescence
  - Steiner Tree

- Kirchoff's matrix tree theorem
- Prüfer sequences
- Lovász Toggle
- Look at the DFS tree (which has no cross-edges)
- Mathematics
  - Is the function multiplicative?
  - Look for a pattern
  - Permutations
    - \* Consider the cycles of the permutation
  - Functions
    - \* Sum of piecewise-linear functions is a piecewise-linear function
    - \* Sum of convex (concave) functions is convex (concave)
  - Modular arithmetic
    - \* Chinese Remainder Theorem
    - \* Linear Congruence
  - Sieve
  - System of linear equations
  - Values to big to represent?
    - \* Compute using the logarithm
    - \* Divide everything by some large value
- Logic
  - 2-SAT
  - XOR-SAT (Gauss elimination or Bipartite matching)
- Meet in the middle
- Only work with the smaller half  $(\log(n))$
- Strings
  - Trie (maybe over something weird, like bits)
  - Suffix array
  - Suffix automaton (+DP?)
  - Aho-Corasick
  - eerTree
  - Work with S + S
- Hashing
- Euler tour, tree to array
- Segment trees
  - Lazy propagation
  - Persistent
  - Implicit
  - Segment tree of X
- Geometry
  - Minkowski sum (of convex sets)
  - Rotating calipers
  - Sweep line (horizontally or vertically?)
  - Sweep angle
  - Convex hull
- Fix a parameter (possibly the answer).
- Are there few distinct values?
- Binary search
- Sliding Window (+ Monotonic Queue)

- Computing a Convolution? Fast Fourier Transform
- Exact Cover (+ Algorithm X)
- Cycle-Finding
- What is the smallest set of values that identify the solution? The cycle structure of the permutation? The powers of primes in the factorization?
- Look at the complement problem
  - Minimize something instead of maximizing
- Immediately enforce necessary conditions. (All values greater than 0? Initialize them all to 1)
- Add large constant to negative numbers to make them positive
- Counting/Bucket sort

### 10. Formulas

- Jacobi symbol:  $\left(\frac{a}{b}\right) = a^{(b-1)/2} \pmod{b}$
- **Heron's formula:** A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where  $s = \frac{a+b+c}{2}$ .
- Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice points on the boundary has area  $i + \frac{b}{2} 1$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$  where each p is a distinct prime factor of n.
- König's theorem: In any bipartite graph  $G = (L \cup R, E)$ , the number of edges in a maximum matching is equal to the number of vertices in a minimum vertex cover. Let U be the set of unmatched vertices in L, and Z be the set of vertices that are either in U or are connected to U by an alternating path. Then  $K = (L \setminus Z) \cup (R \cap Z)$  is the minimum vertex cover.
- A minumum Steiner tree for n vertices requires at most n-2 additional Steiner vertices.
- The number of vertices of a graph is equal to its minimum vertex cover number plus the size of a maximum independent set.
- Wilson's theorem:  $(n-1)! \equiv -1 \pmod{n}$  iff. n is prime
- Lagrange polynomial through points  $(x_0, y_0), \ldots, (x_k, y_k)$  is  $L(x) = \sum_{j=0}^k y_j \prod_{\substack{0 \le m \le k \\ m \ne j}} \frac{x x_m}{x_j x_m}$
- Hook length formula: If  $\lambda$  is a Young diagram and  $h_{\lambda}(i,j)$  is the hook-length of cell (i,j), then then the number of Young tableux  $d_{\lambda} = n! / \prod h_{\lambda}(i,j)$ .
- Möbius inversion formula: If  $f(n) = \sum_{d|n} g(d)$ , then  $g(n) = \sum_{d|n} \mu(d) f(n/d)$ . If  $f(n) = \sum_{m=1}^{n} g(\lfloor n/m \rfloor)$ , then  $g(n) = \sum_{m=1}^{n} \mu(m) f(\lfloor \frac{n}{m} \rfloor)$ .
- 10.1. Markov Chains. A Markov Chain can be represented as a weighted directed graph of states, where the weight of an edge represents the probability of transitioning over that edge in one timestep. Let  $P^{(m)}=(p_{ij}^{(m)})$  be the probability matrix of transitioning from state i to state j in m timesteps, and note that  $P^{(1)}$  is the adjacency matrix of the graph. Chapman-Kolmogorov:  $p_{ij}^{(m+n)}=\sum_k p_{ik}^{(m)} p_{kj}^{(n)}$ . It follows that  $P^{(m+n)}=P^{(m)}P^{(n)}$  and

 $P^{(m)} = P^m$ . If  $p^{(0)}$  is the initial probability distribution (a vector), then  $p^{(0)}P^{(m)}$  is the probability distribution after m timesteps.

The return times of a state i is  $R_i = \{m \mid p_{ii}^{(m)} > 0\}$ , and i is aperiodic if  $gcd(R_i) = 1$ . A MC is aperiodic if any of its vertices is aperiodic. A MC is *irreducible* if the corresponding graph is strongly connected.

A distribution  $\pi$  is stationary if  $\pi P = \pi$ . If MC is irreducible then  $\pi_i = 1/\mathbb{E}[T_i]$ , where  $T_i$  is the expected time between two visits at i.  $\pi_j/\pi_i$  is the expected number of visits at j in between two consecutive visits at i. A MC is ergodic if  $\lim_{m\to\infty} p^{(0)}P^m = \pi$ . A MC is ergodic iff. it is irreducible and aperiodic.

A MC for a random walk in an undirected weighted graph (unweighted graph can be made weighted by adding 1-weights) has  $p_{uv} = w_{uv}/\sum_x w_{ux}$ . If the graph is connected, then  $\pi_u = \sum_x w_{ux}/\sum_v \sum_x w_{vx}$ . Such a random walk is aperiodic iff. the graph is not bipartite.

An absorbing MC is of the form  $P = \begin{pmatrix} Q & R \\ 0 & I_r \end{pmatrix}$ . Let  $N = \sum_{m=0}^{\infty} Q^m = (I_t - Q)^{-1}$ . Then, if starting in state i, the expected number of steps till absorption is the i-th entry in N1. If starting in state i, the probability of being absorbed in state j is the (i,j)-th entry of NR.

Many problems on MC can be formulated in terms of a system of recurrence relations, and then solved using Gaussian elimination.

10.2. Burnside's Lemma. Let G be a finite group that acts on a set X. For each g in G let  $X^g$  denote the set of elements in X that are fixed by g. Then the number of orbits

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

10.3. **Bézout's identity.** If (x,y) is any solution to ax+by=d (e.g. found by the Extended Euclidean Algorithm), then all solutions are given by

$$\left(x + k \frac{b}{\gcd(a,b)}, y - k \frac{a}{\gcd(a,b)}\right)$$

10.4. **Misc.** 

10.4.1. Determinants and PM.

$$det(A) = \sum_{\sigma \in S_n} \operatorname{sgn}(\sigma) \prod_{i=1}^{n} a_{i,\sigma(i)}$$

$$perm(A) = \sum_{\sigma \in S_n} \prod_{i=1}^{n} a_{i,\sigma(i)}$$

$$pf(A) = \frac{1}{2^n n!} \sum_{\sigma \in S_{2n}} \operatorname{sgn}(\sigma) \prod_{i=1}^{n} a_{\sigma(2i-1),\sigma(2i)}$$

$$= \sum_{M \in \operatorname{PM}(n)} \operatorname{sgn}(M) \prod_{(i,j) \in M} a_{i,j}$$

10.4.2. BEST Theorem. Number of OST given by Kirchoff's Theorem (remove r/c with root)  $\#\text{OST}(G,r) \cdot \prod_v (d_v - 1)!$ 

10.4.3. Primitive Roots. Only exists when n is  $2,4,p^k,2p^k$ , where p odd prime. Assume n prime. Number of primitive roots  $\phi(\phi(n))$  Let g be primitive root. All primitive roots are of the form  $g^k$  where  $k,\phi(p)$  are coprime.

k-roots:  $g^{i \cdot \phi(n)/k}$  for  $0 \le i < k$ 

10.4.4. Sum of primes. For any multiplicative f:

$$S(n,p) = S(n,p-1) - f(p) \cdot (S(n/p,p-1) - S(p-1,p-1))$$

10.4.5. Floor.

$$\lfloor \lfloor x/y \rfloor / z \rfloor = \lfloor x/(yz) \rfloor$$
$$x\%y = x - y |x/y|$$

## PRACTICE CONTEST CHECKLIST

- How many operations per second? Compare to local machine.
- What is the stack size?
- How to use printf/scanf with long long/long double?
- Is \_\_int128 available?
- Does MLE give RTE or MLE as a verdict? What about stack overflow?
- What is RAND\_MAX?
- How does the judge handle extra spaces (or missing newlines) in the output?
- Look at documentation for programming languages.
- Try different programming languages: C++ and Java.
- Try the submit script.
- Try local programs: i?python[23], factor.
- Try submitting with assert(false) and assert(false).
- Return-value from main.
- Look for directory with sample test cases.
- Remove this page from the notebook.