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```
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          1. Code Templates
                            ----public static void main(String[] args) throws Exception {--------// 02
                            -----Scanner in = new Scanner(System.in):-----// ef
1.1. Basic Configuration. Vim and (Caps Lock = Escape) configuration.
                            ------PrintWriter out = new PrintWriter(System.out, false);------// 62
o.ygtxmal ekrpat # setxkbmap dvorak for dvorak on gwerty
                            -----// code-----// e6
setxkbmap -option caps:escape
                            -----out.flush():-----// 56
set -o vi
                            xset r rate 150 100
                            }-----// 00
cat > ~/.vimrc
set nocp et sw=4 ts=4 sts=4 si cindent hi=1000 nu ru noeb showcmd showmode
                                       2. Data Structures
syn on | colorscheme slate
                            2.1. Union-Find. An implementation of the Union-Find disjoint sets data structure.
1.2. C++ Header. A C++ header.
                            struct union_find {------// 42
#include <cmath>------// 7d ----union_find(int n) { parent.resize(cnt = n);------// 92
#include <cstdio>------[i] = i; }------// 6f
#include <cstdlib>------// 11 ----int find(int i) {--------// a6
#include <cstring>-------[i] = i ? i : (parent[i] = find(parent[i])); }------// @ -------|/ a9
#include <map>-----// 02
#include <set>------// d2
#include <sstream>------ min(a, b); }------// 18 // int f(int a, int b) { return min(a, b); }------
#include <stack>------// cf const int ID = 0;------// 82
#include <string>------// a9 int f(int a, int b) { return a + b; }------// 6d
#include <vector>-----// 4f ----int n; vi data, lazy;-----// fa
------mk(arr, 0, n-1, 0); }------
#define foreach(u, o) \------// ea ----int mk(const vi &arr, int i, int r, int i) {------// 02
const int INF = 2147483647;-----// be -----int m = (l + r) / 2;-----// 0f
const double pi = acos(-1);------// 49 ----int query(int a, int b) { return q(a, b, 0, n-1, 0); }------// f5
typedef long long ll;------// 8f ----int q(int a, int b, int l, int r, int i) {-------// ad
typedef unsigned long long ull;-----// 81 -----propagate(l, r, i);-----// f7
typedef vector<vii>vvii;------// 4b ----void update(int i, int v) { u(i, v, 0, n-1, 0); }------// 65
template <class T> T mod(T a, T b) { return (a % b + b) % b; }--------// 70 ----int u(int i, int v, int l, int r, int j) {------------// b5
template <class T> int size(const T &x) { return x.size(); }------// 68 -----propagate(l, r, j);-------// 3c
                            -----if (r < i || i < l) return data[j];------// 6a
1.3. Java Template. A Java template.
                            -----if (l == i && r == i) return data[j] = v;------// 74
import java.math.*;------return data[j] = f(u(i, v, l, m, 2*j+1), u(i, v, m+1, r, 2*j+2)); }----// 68
-----// a3 ----int ru(int a, int b, int v, int i) {-------------// d7
```

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------int m = (l + r) / 2; ------// 9d ----matrix(const matrix  other) : rows(other.rows), cols(other.cols), ------// fe
-----ceturn data[i] = f(ru(a, b, v, l, m, 2*i+1), ru(a, b, v, m+1, r, 2*i+2)); ------cnt(other.cnt), data(other.data) { }------// ed
------if (l < r) {-------// 6e ------for (int i = 0; i < cnt; i++) data[i] -= other.data[i]; }------// 88
------else lazv[2*i+2] += lazv[i];------// d1 -----matrix<T> res(*this); res += other; return res; }-----// 5d
-----lazv[i] = INF;------res(*this); res -= other; return res; }------// cf
};-------matrix<T> res(*this); res *= other; return res; }------// 37
                             ----matrix<T> operator *(const matrix& other) {------// 95
2.3. Fenwick Tree. A Fenwick Tree is a data structure that represents an array of n numbers. It
                            -----matrix<T> res(rows, other.cols);------// 57
supports adjusting the i-th element in O(\log n) time, and computing the sum of numbers in the range
                            -----for (int i = 0; i < rows; i++) for (int j = 0; j < other.cols; j++)----// 7a
i...j in O(\log n) time. It only needs O(n) space.
                            -----for (int k = 0; k < cols; k++)-----// fc
struct fenwick_tree {------res(i, j) += at(i, k) * other.data[k * other.cols + j];-----// eb
----int n; vi data;------// d3 -----return res; }-----// 70
----fenwick_tree(int _n) : n(_n), data(vi(n)) { }-------// db ----matrix<T> transpose() {--------------// dd
----void update(int at, int by) {--------// 76 ------matrix<T> res(cols, rows);------// b5
------while (at < n) data[at] += by, at |= at + 1; }------// fb -------for (int i = 0; i < rows; i++)------// 9c
------int res = 0;-------// c3 -----return res; }------// c3
-----return res; }------// e4 -----matrix<T> res(rows, cols), sq(*this);------// 4d
----int rsq(int a, int b) { return query(b) - query(a - 1); }-------// be -------for (int i = 0; i < rows; i++) res(i, i) = T(1);-------// bf
};------while (p) {------// cb
struct fenwick_tree_sq {------// d4 ------if (p & 1) res = res * sq;-----// c1
----int n; fenwick_tree x1, x0;-------// 18 -----p >>= 1;------// 68
----fenwick_tree_sq(int _n) : n(_n), x1(fenwick_tree(n)),------// 2e ------if (p) sq = sq * sq;-------// 9c
-----x0(fenwick_tree(n)) { }------// 7c -----} return res; }-------
----// insert f(y) = my + c if x <= y------// 17 ----matrix<T> rref(T &det) {-------// 89
}:-----// 13 -------int k = r;------// e5
void range_update(fenwick_tree_sq &s, int a, int b, int k) {-------// 89 -------while (k < rows && eq<T>(mat(k, c), T(0))) k++;-----// f9
----s.update(a, k, k * (1 - a)); s.update(b+1, -k, k * b); }------// 7f ------if (k >= rows) continue;------// 3f
----return s.query(b) - s.query(a-1); }------// f3 ------det *= T(-1);-------------------// 7a
                             ------for (int i = 0: i < cols: i++)-----// ab
2.4. Matrix. A Matrix class.
                            -----swap(mat.at(k, i), mat.at(r, i));-----// 8d
class matrix {-----// 85
```

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-----for (int i = 0; i < rows; i++) {--------// 3d ------node *s = successor(n);------// 16
-----T m = mat(i, c):--------// e8 ------erase(s, false):-------// 17
------if (i != r && !eq<T>(m, T(0)))-------// 33 ------s->p = n->p, s->l = n->l, s->r = n->r;------// 37
------} return mat; }-------// 8f ------parent_leg(n) = s, fix(s);------// 87
private:-----// e0 -----return;-------// 32
----vector<T> data;------// 41 ------fix(n->p), n->p = n->l = n->r = NULL;------// 70
};-------// b8 ---node* successor(node *n) const {-------// lb
                           -----if (!n) return NULL;-----// b3
2.5. AVL Tree. A fast, easily augmentable, balanced binary search tree.
                           ------if (n->r) return nth(0, n->r);------// 5b
#define AVL_MULTISET 0-----// b5 -----node *p = n->p;-----// 7c
template <class T>------// 22 -----return p; }------// 03
class avl_tree {------// ff ----node* predecessor(node *n) const {------// e6
public:-----// f6 -----if (!n) return NULL;------// 96
----struct node {------// 45 -----if (n->l) return nth(n->l->size-1, n->l);------// 15
-----T item; node *p, *l, *r;--------// a6 -----node *p = n->p;-------// 33
------node(const T &item, node *p = NULL) : item(item), p(p),------// c5 -----return p; }------
------l(NULL), r(NULL), size(1), height(0) { } };--------// e2
----node *root;------// c1 ----node* nth(int n, node *cur = NULL) const {------// f4
----node* find(const T &item) const {-------// d2 -----if (!cur) cur = root;------// 0a
-----node *cur = root;-----// cf ------while (cur) {------// 55
------while (cur) {-------// ad ------if (n < sz(cur->l)) cur = cur->l;------// 25
------else if (item < cur->item) cur = cur->l;------// de -----else break;-----
-------else break; }------// 05 ------} return cur; }-------// 8f
-----return cur; }------// e<sup>7</sup> private:-----// 49
-----node *prev = NULL, **cur = &root;-----// 60
                           ----inline int height(node *n) const { return n ? n->height : -1; }------// e4
------while (*cur) {--------// b0 ----inline bool left_heavy(node *n) const {-------// d7
-----prev = *cur;------// 31 ------return n && height(n->l) > height(n->r); }------// 9d
#endif-----// c6 ----node*& parent_leg(node *n) {-------// 0d
-----*cur = n, fix(n); return n; }------// 86 -----if (n->p->r == n) return n->p->r;-----// 0e
----void erase(const T &item) { erase(find(item)); }------// c0 -----assert(false); }-----
-----if (!n) return;------// 4d -----if (!n) return;-----// 4d
------if (!n->l && n->r) parent_leg(n) = n->r, n->r->p = n->p;-------// f5 -----n->size = 1 + sz(n->l) + sz(n->r);------// 14
------else if (n->l && !n->r) parent_leg(n) = n->l, n->l->p = n->p;-------// 3d -----n->height = 1 + max(height(n->l), height(n->r)); }------// a1
-----else if (n->l && n->r) {------// 1a
```

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-----l->p = n->p; N------// 66 ----void swim(int i) {------// 33
                                     ------while (i > 0) {------// 1a
-----parent_leg(n) = l; \[ \]------// 02
                                     -----int p = (i - 1) / 2;-----// 77
------n->l = l->r; N------// 08
                                     ------if (!cmp(i, p)) break;-----// a9
-----l->r = n, n->p = l; N------// c3 ----void sink(int i) {-------// ca
----void left_rotate(node *n) { rotate(r, l); }------// 43 -------int l = 2*i + 1, r = l + 1;------// b4
------while (n) { augment(n);-------// c9 ------if (!cmp(m, i)) break;------// 42
------if (too_heavy(n)) {-------// a9 -----swp(m, i), i = m; } }-----// 1d
------else if (right_heavy(n) δδ left_heavy(n->r))------// 09 ----heap(int init_len = 128) : count(0), len(init_len), _cmp(Compare()) {------// 17
-----right_rotate(n->r);------// 7c -----q = new int[len], loc = new int[len];------// f8
------if (left_heavy(n)) right_rotate(n);------// 44 -----memset(loc, 255, len << 2); }-----// f7
------else left_rotate(n);------// 02 ----~heap() { delete[] loc; }------// 09
-----n = n->p; }------// af ----void push(int n, bool fix = true) {-------// b7
#ifdef RESIZE-----// a9
 Also a very simple wrapper over the AVL tree that implements a map interface.
                                     ------int newlen = 2 * len;-----// 22
#include "avl_tree.cpp"------// 01
                                     -----while (n >= newlen) newlen *= 2;------// 2f
-----// ba
                                     ------int *newg = new int[newlen], *newloc = new int[newlen];------// e3
template <class K, class V>-----// da
                                     -----for (int i = 0; i < len; i++) newq[i] = q[i], newloc[i] = loc[i]; --// 94
class avl_map {-----// 3f
                                     -----/ 18 emset(newloc + len, 255, (newlen - len) << 2);------// 18
public:----// 5d
                                     -----delete[] q, delete[] loc;-----// 74
----struct node {------// 2f
                                     -----loc = newloc, q = newq, len = newlen;-----// 61
-----K kev: V value:-----// 32
                                     #else-----// 54
-----/ 29 key(k), value(v) { }-----// 29
                                     ------assert(false):-----// 84
-----bool operator <(const node &other) const { return key < other.key; } };// 92
                                     #endif------// 64
----avl_tree<node> tree;------// b1
                                     ----V& operator [](K key) {------// 7c
                                     -----assert(loc[n] == -1);------// 8f
-----typename avl_tree<node>::node *n = tree.find(node(key, V(\theta)));------// ba
                                     -----loc[n] = count, q[count++] = n;-----// 6b
-----if (!n) n = tree.insert(node(key, V(0)));-----// cb
                                     -----if (fix) swim(count-1); }-----// bf
-----return n->item.value;------// ec
                                     ----void pop(bool fix = true) {-------// 43
----}------// 2e
                                     -----assert(count > 0);-----// eb
};-----// af
                                     -----loc[q[0]] = -1, q[0] = q[--count], loc[q[0]] = 0;------// 50
                                     -----if (fix) sink(0);-----// 80
2.6. Heap. An implementation of a binary heap.
                                     ----}-------// 16
#define RESIZE-----// d0
                                     ----int top() { assert(count > 0): return q[0]: }------// ab
#define SWP(x,y) tmp = x, x = y, y = tmp-----// fb
                                     ----void heapify() { for (int i = count - 1; i > 0; i--)-----// 39
struct default_int_cmp {------// 8d
                                     ------if (cmp(i, (i - 1) / 2)) swp(i, (i - 1) / 2); }------// θb
----default_int_cmp() { }------// 35
                                     ----void update_key(int n) {------// 26
----bool operator ()(const int &a, const int &b) { return a < b; } };-----// e9
                                     -----assert(loc[n] != -1), swim(loc[n]), sink(loc[n]); }------// 7d
template <class Compare = default_int_cmp>------// 30
                                     ----bool empty() { return count == 0; }-----// f8
class heap {-----// 05
                                     ----int size() { return count; }-------// 86
private:-----// 39
                                     ----void clear() { count = 0, memset(loc, 255, len << 2); } };------// 58
----int len, count, *q, *loc, tmp;-----// 0a
----Compare _cmp:-----// 98
----inline bool cmp(int i, int j) { return _cmp(q[i], q[j]); }-----// a0 2.7. Skiplist. An implementation of a skiplist.
```

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----unsigned int cnt = 0;------// 28 ----void erase(T target) {-------// 4d
template<class T> struct skiplist {--------// 34 ------for(int i = 0; i <= current_level; i++) {------// 97
-----T item:------update[i]->next[i] = x->next[i];------// 59
------int *lens;------update[i]->lens[i] = update[i]->lens[i] + x->lens[i] - 1;--// b1
------delete x; _size--;-------delete x; _size--;-------// 81
------while(current_level); free(lens); free(next); }; };-------// aa --------while(current_level > 0 && head->next[current_level] == NULL)-----// 7f
----node *head;------// b7
----skiplist() : current_level(0), _size(0), head(new node(MAX_LEVEL, 0)) { };-// 7a
----~skiplist() { clear(); delete head; head = NULL; }-----// aa
                           2.8. Dancing Links. An implementation of Donald Knuth's Dancing Links data structure. A linked
----#define FIND_UPDATE(cmp, target) \[\bar{N}\]------------------------------// c3 list supporting deletion and restoration of elements.
------int pos[MAX_LEVEL + 2]; \[\bar{\cappa}\]-------// 18
                           template <class T>-----// 82
-----memset(pos, 0, sizeof(pos)); \[ \] ------// f2
                           struct dancing_links {-----// 9e
-----node *l, *r;-----// 32
----: item(item), l(l), r(r) {------// 04
-----pos[i] = pos[i + 1]; N-----// 68
                           -----if (l) l->r = this:-----// 1c
-----if (r) r->l = this;-----// 0b
-----pos[i] += x->lens[i]; x = x->next[i]; } \[ \frac{10}{10} \]
----void clear() { while(head->next && head->next[0])-------// 91 ----node *push_back(const T &item) {------// d7
------back = new node(item, back, NULL);-------// 5d
----node *find(T target) { FIND_UPDATE(x->next[i]->item, target);------// 36 -----if (!front) front = back;-----------------------// a2
------return x && x->item == target ? x : NULL; }------// 50 -----return back;------
----int count_less(T target) { FIND_UPDATE(x->next[i]->item, target);------// 80 ----node *push_front(const T &item) {-------------------// ea
------front = new node(item, NULL, front);------// 75
------FIND_UPDATE(x->next[i]->item, target);-------// 3a -----return front;-----
------int lvl = bernoulli(MAX_LEVEL);-------// 7a ----void erase(node *n) {------// 88
------if(lvl > current_level) current_level = lvl;-------// 8a -----if (!n->l) front = n->r; else n->l->r = n->r;------// d5
-----x = new node(lvl, target);-------// 36 ------if (!n->r) back = n->l; else n->r->l = n->l;------// 96
-----x->next[i] = update[i]->next[i];------// 46 ----void restore(node *n) {------// 6d
-----x->lens[i] = pos[i] + update[i]->lens[i] - pos[0];------// bc -----if (!n->l) front = n; else n->l->r = n;------// ab
------update[i]->next[i] = x;-------// 20 ------if (!n->r) back = n; else n->r->l = n;------// 8d
```

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                                           -----}-----// be
2.9. Misof Tree. A simple tree data structure for inserting, erasing, and querying the nth largest
                                            -----return sqrt(sum); }-----// ef
element.
                                           ------bb bound(double l, int c, bool left) {------// b6
#define BITS 15-----// 7h
                                            -----pt nf(from.coord), nt(to.coord);-----// 5c
struct misof_tree {-----// fe
                                            ------if (left) nt.coord[c] = min(nt.coord[c], l);------// ef
----int cnt[BITS][1<<BITS];-----// aa
                                            ------else nf.coord[c] = max(nf.coord[c], l);-----------// 71
----misof_tree() { memset(cnt, 0, sizeof(cnt)); }-----// b0
                                            -----return bb(nf, nt); } };-----// 3b
----void insert(int x) { for (int i = 0; i < BITS; cnt[i++][x]++, x >>= 1); }--// 5a
                                            ----struct node {------// 8d
----void erase(int x) { for (int i = 0; i < BITS; cnt[i++][x]--, x >>= 1); }---// 49
                                            -----pt p; node *l, *r;-----// 46
----int nth(int n) {-------// 8a
                                            -----node(pt _p, node *_l, node *_r) : p(_p), l(_l), r(_r) { } };------// 23
-----int res = 0;------// a4
                                            ----node *root:-----// 30
------for (int i = BITS-1; i >= 0; i--)-----// 99
                                           ----kd_tree() : root(NULL) { }-----// 57
-----if (cnt[i][res <<= 1] <= n) n -= cnt[i][res], res |= 1:-----// f4
                                            ----kd_tree(vector<pt> pts) { root = construct(pts, 0, size(pts) - 1, 0); }----// 66
-----return res:-----// 3a
                                            ----node* construct(vector<pt> &pts, int from, int to, int c) {------// θb
----}------------// b5
                                            -----if (from > to) return NULL;------// f4
};-----// 0a
                                            ------int mid = from + (to - from) / 2;------// 43
                                            -----nth_element(pts.begin() + from, pts.begin() + mid,------// f3
2.10. k-d Tree. A k-dimensional tree supporting fast construction, adding points, and nearest neigh-
                                            -----pts.begin() + to + 1, cmp(c));------// 97
bor queries.
                                            -----return new node(pts[mid], construct(pts, from, mid - 1, INC(c)),-----// dd
#define INC(c) ((c) == K - 1 ? 0 : (c) + 1)-----// 77
                                           -----/ 03
template <int K>-----// cd
                                           ----bool contains(const pt &p) { return _con(p, root, 0); }------// f0
class kd_tree {------// 7e ----bool _con(const pt &p, node *n, int c) {------// 74
public:-----// c7 ------if (!n) return false;------// b7
----struct pt {-------// 78 ------if (cmp(c)(p, n->p)) return _con(p, n->l, INC(c));-------// 81
-------double coord[K];------// d6 ------if (cmp(c)(n->p, p)) return _con(p, n->r, INC(c));------// 95
-----pt() {}------// c1 -----return true; }------// 86
-----pt(double c[K]) { for (int i = 0; i < K; i++) coord[i] = c[i]; }-----// 4c ----void insert(const pt &p) { _ins(p, root, 0); }------// f9
------double dist(const pt &other) const {--------// 6c ----void _ins(const pt &p, node* &n, int c) {-------// 09
-------double sum = 0.0;------// c4 -----if (!n) n = new node(p, NULL, NULL);------// 4d
------for (int i = 0; i < K; i++)-------// 23 ------else if (cmp(c)(p, n->p)) _ins(p, n->l, INC(c));------// a0
-----sum += pow(coord[i] - other.coord[i], 2.0);------// 46 -----else if (cmp(c)(n->p, p)) _ins(p, n->r, INC(c)); }------// 2c
-----return sqrt(sum); } };------// ad ----void clear() { _clr(root); root = NULL; }-----// 73
----struct cmp {------// 8f -----// 8f -----// 1a
-----cmp(int _c) : c(_c) {}------// a5 -----assert(root);----------// f8
------bool operator ()(const pt &a, const pt &b) {-------// 26 -----double mn = INFINITY, cs[K];------// bf
cc = i == 0 ? c : i - 1;------// bc -----pt from(cs);------// 5a
------if (abs(a.coord[cc] - b.coord[cc]) > EPS)------// 28 ------for (int i = 0; i < K; i++) cs[i] = INFINITY;-----// 37
-----return a.coord[cc] < b.coord[cc];-----// b7 -----pt to(cs);------
-----return false; } };------// 6e
----struct bb {-------// 30 ----pair<pt, bool> _nn(------// e3
-----pt from, to;-------const pt &p, node *n, bb b, double &mn, int c, bool same) {------// aa
------bb(pt _from, pt _to) : from(_from), to(_to) {}------// 57 ------if (!n || b.dist(p) > mn) return make_pair(pt(), false);------// 19
------double dist(const pt &p) {-------// 3f -------bool found = same || p.dist(n->p) > EPS, l1 = true, l2 = false;-----// 9f
------double sum = 0.0;------// d9 -----pt resp = n->p;-----// 6b
------if (p.coord[i] < from.coord[i])------// a0 -----node *n1 = n->l, *n2 = n->r;------// aa
-----sum += pow(from.coord[i] - p.coord[i], 2.0);------// 00 ------for (int i = 0; i < 2; i++) {-------// 50
-----sum += pow(p.coord[i] - to.coord[i], 2.0);-----// 8c
```

3.1. **Breadth-First Search.** An implementation of a breadth-first search that counts the number of edges on the shortest path from the starting vertex to the ending vertex in the specified unweighted graph (which is represented with an adjacency list). Note that it assumes that the two vertices are connected. It runs in O(|V| + |E|) time.

Another implementation that doesn't assume the two vertices are connected. If there is no path from the starting vertex to the ending vertex, a -1 is returned.

```
int bfs(int start, int end, vvi& adj_list) {------// d7
----set<<u>int</u>> visited;-----// b3
---queue<ii>> 0;-----// bb
----0.push(ii(start, 0));-----// 3a
----visited.insert(start):-----// b2
-----// db
----while (!Q.empty()) {------// f7
-----ii cur = 0.front(); 0.pop();-----// 03
-----// 9c
-----if (cur.first == end)-----// 22
-----return cur.second;-----// b9
-----// ba
-----vi& adj = adj_list[cur.first];-----// f9
------for (vi::iterator it = adj.begin(); it != adj.end(); it++)------// 44
-----if (visited.find(*it) == visited.end()) {------// 8d
-----Q.push(ii(*it, cur.second + 1));-----// ab
-----visited.insert(*it);-----// cb
----return -1:------// f5
}-----// 03
```

```
3.2. Single-Source Shortest Paths.
```

```
3.2.1. Dijkstra's algorithm. An implementation of Dijkstra's algorithm. It runs in \Theta(|E|\log|V|)
time.
int *dist, *dad;-----// 46
struct cmp {-----// a5
----bool operator()(int a, int b) {------// bb
-----return dist[a] != dist[b] ? dist[a] < dist[b] : a < b; }------// e6
};-----// 41
pair<int*, int*> dijkstra(int n, int s, vii *adj) {-----// 53
----dist = new int[n];-----// 84
----dad = new int[n];-----// 05
----for (int i = 0; i < n; i++) dist[i] = INF, dad[i] = -1;------// d6
----set<<u>int</u>, cmp> pq;-----// 04
----dist[s] = 0, pq.insert(s);-----// 1b
----while (!pq.empty()) {------// 57
------<mark>int</mark> cur = *pq.beqin(); pq.erase(pq.beqin());--------------// 7d
------for (int i = 0; i < size(adj[cur]); i++) {-------// 9e
-----int nxt = adj[cur][i].first,-----// b8
-----/ndist = dist[cur] + adj[cur][i].second;------// θε
-----if (ndist < dist[nxt]) pq.erase(nxt),-----// e4
 -----dist[nxt] = ndist, dad[nxt] = cur, pg.insert(nxt);-----// 0f
----}------// e8
----return pair<int*, int*>(dist, dad);-----// cc
 3.2.2. Bellman-Ford algorithm. The Bellman-Ford algorithm solves the single-source shortest paths
problem in O(|V||E|) time. It is slower than Dijkstra's algorithm, but it works on graphs with
negative edges and has the ability to detect negative cycles, neither of which Dijkstra's algorithm can
int* bellman_ford(int n, int s, vii* adj, bool& has_negative_cycle) {------// cf
----has_negative_cycle = false;------// 47
----int* dist = new int[n];-----// 7f
----for (int i = 0: i < n: i++) dist[i] = i == s ? 0 : INF:------// 10
----for (int i = 0; i < n - 1; i++)-----// a1
```

```
3.3. All-Pairs Shortest Paths.
```

3.3.1. Floyd-Warshall algorithm. The Floyd-Warshall algorithm solves the all-pairs shortest paths problem in  $O(|V|^3)$  time.

-----for (int j = 0; j < n; j++)-----// c4

-----if (dist[j] != INF)-----// 4e

------for (int k = 0; k < size(adj[j]); k++)------// 3f

-----dist[adj[j][k].first] = min(dist[adj[j][k].first],-----// 61

-----dist[j] + adj[j][k].second);------// 47

----for (int j = 0; j < n; j++)-----// 13

------for (int k = 0; k < size(adj[j]); k++)------// aθ

------if (dist[i] + adj[i][k].second < dist[adj[i][k].first])------// ef

-----has\_negative\_cvcle = true:-----// 2a

----return dist:-----// 2e

}-----/<sub>-----</sub>// c2

```
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----for (int k = 0; k < n; k++)-------// 49 // edges is a list of edges of the form (weight, (a, b))------// c6
------for (int i = 0; i < n; i++)------// 21 // the edges in the minimum spanning tree are returned on the same form-----// 4d
------for (int j = 0; j < n; j++)-------// 77 vector<pair<int, ii> > mst(int n, vector<pair<int, ii> > edges) {------// a7
------if (arr[i][k] != INF && arr[k][j] != INF)--------// b1 ----union_find uf(n);------
}------// 86 ----vector<pair<int, ii>> res;-----------------------------------// 71
                                         ----for (int i = 0; i < size(edges); i++)-----// ce
3.4. Strongly Connected Components.
                                         -----if (uf.find(edges[i].second.first) !=-----// d5
                                         -----uf.find(edges[i].second.second)) {-----// 8c
3.4.1. Kosaraju's algorithm. Kosarajus's algorithm finds strongly connected components of a directed
                                         -----res.push_back(edges[i]);-----// d1
graph in O(|V| + |E|) time.
                                         -----uf.unite(edges[i].second.first, edges[i].second.second);------// a2
#include "../data-structures/union_find.cpp"------------------// 5e
                                         -----// 11
                                         ----return res:-----// 46
vector<br/>bool> visited;------// 66
                                         -----// 88
vi order:-----// 9b
-----// a5
void scc_dfs(const vvi &adj, int u) {------// a1
                                        3.6. Topological Sort.
----int v; visited[u] = true;-----// e3
----for (int i = 0; i < size(adj[u]); i++)-----// c5
                                        3.6.1. Modified Depth-First Search.
-----if (!visited[v = adj[u][i]]) scc_dfs(adj, v);------// 6e
                                        void tsort_dfs(int cur, char* color, const vvi& adj, stack<int>& res,-----// ca
----order.push_back(u);-----// 19
                                         ------bool& has_cycle) {------// a8
}-----// dc
                                        ----color[cur] = 1:-----// 5b
   -----// 96
                                        ----for (int i = 0; i < size(adj[cur]); i++) {------// 96
pair<union_find, vi> scc(const vvi &adj) {-----// 3e
                                        -----int nxt = adj[cur][i];-----// 53
----int n = size(adj), u, v;-----// bd
                                        ------if (color[nxt] == 0)------// 00
----order.clear():-----// 22
                                        -----tsort_dfs(nxt, color, adj, res, has_cycle);-----// 5b
----union_find uf(n);------// 6d
                                        -----else if (color[nxt] == 1)------// 53
----vi dag:-----// ae
                                        -----has_cycle = true;-----// c8
----vvi rev(n):------// 20
                                        -----if (has_cycle) return;-----// 7e
----for (int i = 0; i < n; i++) for (int j = 0; j < size(adj[i]); j++)------// b9
                                        ----}------// 3d
-----rev[adj[i]]]].push_back(i);-----// 77
                                        ----color[cur] = 2;-----// 16
----visited.resize(n), fill(visited.begin(), visited.end(), false);------// 04
                                         ----res.push(cur):-----// cb
----for (int i = 0; i < n; i++) if (!visited[i]) scc_dfs(rev, i);------// e4
                                        }-----// 9e
----fill(visited.begin(), visited.end(), false);------// c2
                                        ,
-----// ae
----stack<int> S;------// 04
                                        vi tsort(int n, vvi adj, bool& has_cycle) {------// 37
----for (int i = n-1; i >= 0; i--) {-----// 3f
                                         ----has_cycle = false;-----// 37
-----if (visited[order[i]]) continue;-----// 94
                                        ----stack<<u>int</u>> S;-----// 54
-----S.push(order[i]), dag.push_back(order[i]);------// 40
                                        ----vi res:------// d1
------while (!S.empty()) {------// 03
                                        ----char* color = new char[n];------// b1
-----visited[u = S.top()] = true, S.pop(), uf.unite(u, order[i]);-----// 1b
                                         ----memset(color, 0, n);-----// ce
-----for (int i = 0; i < size(adj[u]); i++)-----// 90
                                        ----for (int i = 0; i < n; i++) {------// 96
-----if (!visited[v = adi[u][i]]) S.push(v):------// 43
                                        ------if (!color[i]) {------// d5
------}------// da
                                        -----tsort_dfs(i, color, adj, S, has_cycle);------// 40
----}------// 7c
                                        -----if (has_cycle) return res;------// 6c
----return pair<union_find, vi>(uf, dag);-----// 94
                                         }-----// 97
                                         ----}-----// df
                                         ----while (!S.empty()) res.push_back(S.top()), S.pop();-----// 94
3.5. Minimum Spanning Tree.
                                         ----return res:-----// 07
                                        }-----// 1f
3.5.1. Kruskal's algorithm.
#include "../data-structures/union_find.cpp"-----------------------// 5e
```

```
10
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vi adj[MAXV];-----// ff #define dist(v) dist[v == -1 ? MAXN : v]------// 0f
------else if (indeq[i] != outdeq[i]) return ii(-1,-1);-------// fa -------for(int v = 0; v < N; ++v) if(L[v] == -1) dist(v) = 0, q[r++] = v;-----// 31
----}------else dist(v) = INF;------// c4
----if ((start == -1) != (end == -1) || (c != 2 && c != 0)) return ii(-1,-1);--// 6e -------dist(-1) = INF;---------------------------// f3
----return ii(start, end);------// 9e -------int v = q[l++];-------// 69
------} else s.push(cur), cur = adj[cur][--outdeg[cur]];--------// d8 ---------if(dfs(R[*u])) {-----------------------// c7
-----dist(v) = INF:-----// d4
                       -----return false:-----// de
3.8. Bipartite Matching.
                       -----return true:------// 7b
3.8.1. Alternating Paths algorithm. The alternating paths algorithm solves bipartite matching in
                       ----}------// 61
O(mn^2) time, where m, n are the number of vertices on the left and right side of the bipartite
                       ----void add_edge(int i, int j) { adj[i].push_back(j); }------// 87
graph, respectively.
                       ----int maximum_matching() {------// ae
vi* adi:-----// cc
                       ------int matching = 0;-----// 7d
bool* done:----// b1
                       -----memset(L, -1, sizeof(int) * N);-----// 16
int* owner;-----// 26
                       -----memset(R, -1, sizeof(int) * M);-----// e4
int alternating_path(int left) {-----// da
                       -------while(bfs()) for(int i = 0; i < N; ++i)------// f6
----if (done[left]) return 0;------// 08
                       -----matching += L[i] == -1 && dfs(i);-----// c9
----done[left] = true;-----// f2
                       -----return matching:-----// 82
----for (int i = 0; i < size(adj[left]); i++) {------// 34
                       ----}------// 86
------int right = adj[left][i];------// b6
                       }:-----// dd
-----if (owner[right] == -1 || alternating_path(owner[right])) {-------// d2
------owner[right] = left; return 1;------// 26
                      3.9. Maximum Flow.
------} }------// 7a
----return 0; }-----// 83
                      3.9.1. Dinic's algorithm. An implementation of Dinic's algorithm that runs in O(|V|^2|E|). It computes
                       the maximum flow of a flow network.
3.8.2. Hopcroft-Karp algorithm. An implementation of Hopcroft-Karp algorithm for bipartite match-
                      #define MAXV 2000-----// ba
```

ing. Running time is  $O(|E|\sqrt{|V|})$ .

int q[MAXV], d[MAXV];-----// e6

```
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----struct edge {-------------------------// le ------int v, cap, nxt;----------------// cb
-----edge() { }------// 38 ---};------// f9
-------edge(int v, int cap, int nxt) : v(v), cap(cap), nxt(nxt) { }-------// f7 ----int n, ecnt, *head;-------------------------------// 00
----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 80 ------memset(head = new int[n], -1, n << 2);------// bc
-----e.reserve(2 * (m == -1 ? n : m));------// 5d ---}-----// 5d ---}
------head = new int[n], curh = new int[n];------// 6d ----void destroy() { delete[] head; }-----// f1
----void reset() { e = e_store; }-------// 60 ------e.push_back(edge(u, vu, head[v])); head[v] = ecnt++;------// 9b
----void add_edge(int u, int v, int uv, int vu = 0) {-------// dd ----}-----// ce
----}------e_store = e;-------// 35
------for (int u = q[l++], i = head[u]; i != -1; i = e[i].nxt)-----// 39
-----if(s == t) return 0;-------// 2e -----if (p[t] == -1) break;------// 84
-----e_store = e:------// 59 ------int x = INF, at = p[t];------// 43
------int f = 0, x, l, r;-------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 70
------while (true) {-------// 36 -----at = p[t], f += x;------// 3c
-----while (l < r)------// b8
-----memcpy(curh, head, n * sizeof(int));-----// 0e
                     3.10. Minimum Cost Maximum Flow. An implementation of Edmonds Karp's algorithm, modi-
fied to find shortest path to augment each time (instead of just any path). It computes the maximum
flow of a flow network, and when there are multiple maximum flows, finds the maximum flow with
-----if (res) reset();------// 71
                     minimum cost. Running time is O(|V|^2|E|\log|V|).
-----return f:-----// 72
                     #define MAXV 2000-----// ba
int d[MAXV], p[MAXV], pot[MAXV];-----// 80
};-----// 97
                     struct cmp {-----// d1
3.9.2. Edmonds Karp's algorithm. An implementation of Edmonds Karp's algorithm that runs in
                     ----bool operator ()(int i, int j) {------// 8a
O(|V||E|^2). It computes the maximum flow of a flow network.
                     -----return d[i] == d[j] ? i < j : d[i] < d[j];------// 89
#define MAXV 2000------// ba ---}-----// df
struct flow_network {------// 5e struct flow_network {-----// eb
```

```
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----struct edge {-------// 9a ---}-----// d1
------int v, cap, cost, nxt;-----// ad
                                          }:-----// 80
-----edge(int v, int cap, int cost, int nxt)------// 01
                                          3.11. All Pairs Maximum Flow.
-----: v(v), cap(cap), cost(cost), nxt(nxt) { }-----// 5d
----}:------// ae
                                          3.11.1. Gomory-Hu Tree. An implementation of the Gomory-Hu Tree. The spanning tree is con-
----int n, ecnt, *head;------// 57
                                          structed using Gusfield's algorithm in O(|V|^2) plus |V|-1 times the time it takes to calculate the max-
----vector<edge> e, e_store;-----// d7
                                          imum flow. If Dinic's algorithm is used to calculate the max flow, the running time is O(|V|^3|E|).
----flow_network(int n, int m = -1) : n(n), ecnt(0) {------// 83
                                          #include "dinic.cpp"-----// 58
-----e.reserve(2 * (m == -1 ? n : m));------// 2c
                                           -----// 25
-----memset(head = new int[n], -1, n << 2);------// e1
                                           ----}------// 64
                                           pair<vii, vvi> construct_gh_tree(flow_network &q) {------// 77
----void destroy() { delete[] head; }-----// 89
                                           ----int n = q.n, v;-----// 5d
----void reset() { e = e_store; }------// 64
                                           ----vii par(n, ii(0, 0)); vvi cap(n, vi(n, -1));------// 49
----void add_edge(int u, int v, int cost, int uv, int vu=0) {------// 30
                                           ----for (int s = 1; s < n; s++) {------// 9e
-----e.push_back(edge(v, uv, cost, head[u])); head[u] = ecnt++;------// 04
                                           -----int l = 0, r = 0;-----// 9d
-----e.push_back(edge(u, vu, -cost, head[v])); head[v] = ecnt++;------// 70
                                           -----par[s].second = q.max_flow(s, par[s].first, false);------// 38
-----memset(d, 0, n * sizeof(int));-----// 79
----ii min_cost_max_flow(int s, int t, bool res = true) {--------// 5e
                                           -----memset(same, 0, n * sizeof(int));-----// b0
-----if (s == t) return ii(0, 0);-----// c4
                                           -----d[a[r++] = s] = 1;------// 8c
-----e_store = e:-----// e8
                                           ------while (l < r) {------// 45
-----/memset(pot. 0. n << 2):-----// 32
                                           -----same[v = q[l++]] = true;-----// c8
------int f = 0, c = 0, v;-----// 9c
                                           -----for (int i = g.head[v]; i != -1; i = g.e[i].nxt)------// 33
------while (true) {------// a7
                                           ------if (q.e[i].cap > 0 \& d[q.e[i].v] == 0)-----// 3f
-----memset(d, -1, n << 2);-----// a2
                                           -----d[q[r++] = g.e[i].v] = 1;------// f8
-----memset(p, -1, n << 2);-----// 71
                                           ------}------// b5
-----set<<u>int</u>, cmp> q;-----// a8
                                           -----for (int i = s + 1; i < n; i++)-----// 68
-----q.insert(s); d[s] = 0;-----// 1d
                                           ------if (par[i].first == par[s].first && same[i]) par[i].first = s;----// ea
-----while (!q.empty()) {------// 6f
                                           -----q.reset();------// 9a
-----int u = *q.begin();-----// e6
                                           ----}-----// le
-----q.erase(q.begin());-----// 0f
                                           ----for (int i = 0; i < n; i++) {------// 2a
------for (int i = head[u]; i != -1; i = e[i].nxt) {------// 9c
                                           -----int mn = INF. cur = i:-----// 19
-----if (e[i].cap == 0) continue;-----// b8
                                           ------while (true) {------// 3a
-----int cd = d[u] + e[i].cost + pot[u] - pot[v = e[i].v];-----// f5
                                           -----cap[curl[i] = mn:-----// 63
-----if (d[v] == -1 \mid \mid cd < d[v]) {------// d2
                                           -----if (cur == 0) break;-----// 35
------if (q.find(v) != q.end()) q.erase(q.find(v));------// 47
                                           -----mn = min(mn, par[cur].second), cur = par[cur].first;-----// 28
-----d[v] = cd; p[v] = i;-----// 52
                                           -----q.insert(v);-----// 0f
                                           ----}------// 4a
----return make_pair(par, cap);------// 6b
}-----// 99
int compute_max_flow(int s, int t, const pair<vii, vvi> &qh) {------// 16
-----if (p[t] == -1) break;-----// 51
                                           ----if (s == t) return 0;-----// d4
-----int x = INF, at = p[t]:-----// fe
                                           ----int cur = INF, at = s;-----// 65
------while (at !=-1) x = min(x, e[at].cap), at = p[e[at^1].v];------// 82
                                           ----while (gh.second[at][t] == -1)------// ef
-----at = p[t], f += x:-----// 30
                                           -----cur = min(cur, qh.first[at].second), at = qh.first[at].first;-----// bd
-----while (at != -1)------// 84
                                           ----return min(cur, gh.second[at][t]);-----// 6d
------[at].cap -= x, e[at^1].cap += x, at = p[e[at^1].v];-----// a5
                                           }-----// a2
-----c += x * (d[t] + pot[t] - pot[s]);
-----for (int i = 0; i < n; i++) if (p[i] != -1) pot[i] += d[i];------// 66
                                                             4. Strings
4.1. Trie. A Trie class.
-----if (res) reset();-----// da
-----return ii(f, c);-----// f4
                                          template <class T>-----// 82
                                           class trie {------// 9a
```

```
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                                                               13
----struct node {-------// ae ------for (int i = 0; i < n; i++) P[0][i] = s[i] - 'a';-------// 69
------for (int stp = 1, cnt = 1; cnt >> 1 < n; stp++, cnt <<= 1) {------// bb
------for (int i = 0; i < n; i++)---------// 50
public:------L[L[i].p = i].nr = ii(P[stp - 1][i],------// 0e
----node* root;------i + cnt < n ? P[stp - 1][i + cnt] : -1);-------// 18
----template <class I>------(int i = 0; i < n; i++)-------// 38
-------L[i].nr == L[i - 1].nr ? P[stp][L[i - 1].p] : i;-------// 61
------while (true) {--------// 67 -----}-----// 39
-----if (begin == end) { cur->words++; break; }------// db ---}------// db ---}
------else {-------// 3e ----int lcp(int x, int y) {-------// 10
-----typename map<T, node*>::const_iterator it;------// 01 -----if (x == y) return n - x;---------// f4
-----it = cur->children.find(head);------// 77 ------for (int k = size(P) - 1; k >= 0 && x < n & y < n; k--)------// 3e
------pair<T, node*> nw(head, new node());------// cd -----return res;-----
-----} begin++, cur = it->second; } } }------// 64 };------// 64
----template<class I>-----// b9
                                4.3. Aho-Corasick Algorithm. An implementation of the Aho-Corasick algorithm. Constructs a
----int countMatches(I begin, I end) {------// 7f
                                state machine from a set of keywords which can be used to search a string for any of the keywords.
-----node* cur = root;-----// 32
                                struct aho_corasick {-----// 78
------while (true) {------// bb
                                ----struct out_node {------// 3e
-----if (begin == end) return cur->words;-----// a4
                                -----string keyword; out_node *next;-----// f0
-----else {------// 1e
                                -----out_node(string k, out_node *n) : keyword(k), next(n) { }------// 26
-----T head = *begin;-----// 5c
                                ----}:------// b9
-----typename map<T, node*>::const_iterator it;------// 25
                                ----struct go_node {------// 40
-----it = cur->children.find(head);-----// d9
                                -----map<char, qo_node*> next;------// 6b
-----if (it == cur->children.end()) return 0;-----// 14
-----begin++, cur = it->second; } } -----// 7c
                                -----out_node *out; go_node *fail;-----// 3e
                                -----go_node() { out = NULL; fail = NULL; }-----// 0f
----template<class I>-----// 9c
----int countPrefixes(I begin, I end) {------// 85
                                ----};------// c0
                                ----go_node *go;-----// b8
-----node* cur = root;-----// 95
                                ----aho_corasick(vector<string> keywords) {------// 4b
------while (true) {------// 3e
                                -----go = new go_node();-----// 77
-----if (begin == end) return cur->prefixes;-----// f5
                                ------foreach(k, keywords) {-------// e4
-----else {------// 66
                                -----qo_node *cur = qo;-----// 9d
-----T head = *begin;-----// 43
                                -----foreach(c, *k)-----// 38
-----typename map<T, node*>::const_iterator it;------// 7a
                                -----cur = cur->next.find(*c) != cur->next.end() ? cur->next[*c] :--// 3d
-----it = cur->children.find(head);------// 43
-----(cur->next[*c] = new go_node());------// 75
                                -----cur->out = new out_node(*k, cur->out);------// 6e
-----begin++, cur = it->second; } } };-----// 26
                                4.2. Suffix Array. An O(n \log n) construction of a Suffix Tree.
                                -----queue<go_node*> q;------// 8a
struct entry { ii nr; int p; };-------// f9 ------foreach(a, go->next) q.push(a->second);------// a3
bool operator <(const entry &a, const entry &b) { return a.nr < b.nr; }------// 77 -------while (!q.empty()) {-----------------------------// 43
```

```
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-----st = st->fail;-----// 3f
-----if (!st) st = go;-----// e7
                                              5. Mathematics
5.1. Fraction. A fraction (rational number) class. Note that numbers are stored in lowest common
-----if (s->fail) {------// 3b
                                 terms.
-----if (!s->out) s->out = s->fail->out;------// 80
                                 template <class T>------// 82
-----else {------// ed
                                 class fraction {------// cf
-----out_node* out = s->out;-----// bf
                                 private:----// 8e
-----while (out->next) out = out->next;-----// ca
                                 ----T gcd(T a, T b) { return b == T(0) ? a : gcd(b, a % b); }------// 86
-----out->next = s->fail->out;-----// 65
                                 public:----// 0f
----T n, d;------// 4b
----fraction(T n_, T d_) {------// 03
-----assert(d_ != 0);-----// 3d
-----}-----// e8
                                 -----n = n_, d = d_;-----// 06
-----if (d < T(0)) n = -n, d = -d;------// be
----vector<string> search(string s) {------// 8d
                                 -----T g = gcd(abs(n), abs(d));-----// fc
-----vector<string> res;------// ef
                                 ----n /= q, d /= q; }------// a1
-----go_node *cur = go;-----// 61
                                 ----fraction(T n_) : n(n_), d(1) { }------// 84
------foreach(c, s) {------// 6c
------while (cur && cur->next.find(*c) == cur->next.end())-----// 1f
                                 ----fraction(const fraction<T>\& other) : n(other.n), d(other.d) { }------// \theta 1
-----cur = cur->fail;-----// 9e
                                 ----fraction<T> operator +(const fraction<T>& other) const {------// b6
-----if (!cur) cur = go;-----// 2f
                                 -----return fraction<T>(n * other.d + other.n * d, d * other.d);}------// 3b
-----cur = cur->next[*c];-----// 58
                                 ----fraction<T> operator -(const fraction<T>& other) const {------// 26
-----if (!cur) cur = go;-----// 3f
                                 -----return fraction<T>(n * other.d - other.n * d, d * other.d);}------// 47
                                 ------for (out_node *out = cur->out; out = out->next)------// e0
                                 -----return fraction<T>(n * other.n, d * other.d); }-----// c5
-----res.push_back(out->keyword);-----// 0d
                                 ----fraction<T> operator /(const fraction<T>& other) const {------// ca
-----return fraction<T>(n * other.d, d * other.n); }------// 35
-----return res:-----// c1
                                 -----return n * other.d < other.n * d; }------// 8c
};-----// 32
                                 ----bool operator <=(const fraction<T>& other) const {-------// 48
4.4. The Z algorithm. Given a string S, Z_i(S) is the longest substring of S starting at i that is
                                -----return !(other < *this); }------// 86
also a prefix of S. The Z algorithm computes these Z values in O(n) time, where n = |S|. Z values
                                 ----bool operator >(const fraction<T>& other) const {------// c9
can, for example, be used to find all occurrences of a pattern P in a string T in linear time. This is
                                 -----return other < *this; }------// 6e
accomplished by computing Z values of S = TP, and looking for all i such that Z_i \geq |T|.
                                 ----bool operator >=(const fraction<T>& other) const {------// 4b
----int* z = new int[n];------// c4 -----return n == other.d; }-----// 14
-----z[i] = 0:-----// c9
-----l = r = i;------// a7 struct intx {------// cf
-----z[i] = r - l; r--;-------// fc ----intx(string n) { init(n); }------// b9
------// b5 ----intx(const intx& other) : sign(other.sign), data(other.data) { }------// 3b
------l = i;-------// 02 ----int sign;-------// 26
-----z[i] = r - l; r--; } }------// 8d ----static const int dcnt = 9;--------// 12
```

```
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                                              15
----static const unsigned int radix = 1000000000U;------// f0 ----friend intx abs(const intx &n) { return n < 0 ? -n : n; }-----// 02
-----intx res; res.data.clear();--------// 4e ------if (sign < 0 && b.sign > 0) return b - (-*this);-------// 70
------for (int i = n.size() - 1; i >= 0; i -= intx::dcnt) {-------// e7 ------unsigned long long carry = 0;------------------------// 5c
------for (int j = intx::dcnt - 1; j >= 0; j--) {--------// 72 ------carry += (i < size() ? data[i] : OULL) +------// 91
-----(i < b.size() ? b.data[i] : OULL);-------// OC
-----if (idx < 0) continue; ------// 52 ------c.data.push_back(carry % intx::radix); ------// 86
-----digit = digit * 10 + (n[idx] - '0');-------// 1f -----carry /= intx::radix;------// fd
-----res.data.push_back(digit);-------// 07 -----return c.normalize(sign);-------// 20
------data = res.data;-------// 7d ----intx operator -(const intx& b) const {-------// 53
-----if (sign > 0 && b.sign < 0) return *this + (-b);-------// 8f
-----if (data.empty()) data.push_back(θ);-------// fa -----if (*this < b) return -(b - *this);--------// 36
-----data.erase(data.begin() + i);------// 67 ------long long borrow = 0;-------// f8
------borrow = data[i] - borrow - (i < b.size() ? b.data[i] : 0ULL);----// a9
----}------c.data.push_back(borrow < 0 ? intx::radix + borrow : borrow);-----// ed
------if (first) outs << n.data[i], first = false;-----// 33 ----intx operator *(const intx& b) const {------// bd
-----string s = ss.str();-------// 64 -------for (int j = 0; j < b.size() || carry; j++) {--------// c0
------int len = s.size();------// 0d -------if (j < b.size()) carry += (long long)data[i] * b.data[j];----// af
------c.data[i + j] = carry % intx::radix;------// 86
-----return outs;-------// cf -----}------// ge
----}-----return c.normalize(sign * b.sign);-------// de
------if (sign != b.sign) return sign < b.sign:-------// cf -----assert(!(d.size() == 1 && d.data[0] == 0)):------// e9
------if (size() != b.size())-------// 4d -----intx q, r; q.data.assign(n.size(), 0);------// ca
------return sign == 1 ? size() < b.size() : size() > b.size(); ------// 4d -------for (int i = n.size() - 1; i >= 0; i--) {--------// 1a
------for (int i = size() - 1; i >= 0; i--) if (data[i] != b.data[i])------// 35 ------r.data.insert(r.data.begin(), 0);-------// c7
------return false;-------// ca -------long long k = θ;-------// cc
```

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                                                                                16
------if (d.size() - 1 < r.size()) k += r.data[d.size() - 1];------// 06 ----delete[] A; delete[] B;-------// f7
------k /= d.data.back();--------// b7 ----delete[] a; delete[] b;---------// 7e
-----}-----// 2f
                                        5.3. Binomial Coefficients. The binomial coefficient \binom{n}{k} = \frac{n!}{k!(n-k)!} is the number of ways to choose
-----return pair<intx, intx>(q.normalize(n.sign * d.sign), r);------// a1
                                         k items out of a total of n items.
----}------// 1b
----intx operator /(const intx& d) const {------// a2
                                        int nck(int n, int k) {-----// f6
                                         ----if (n - k < k) k = n - k;------// 18
-----return divmod(*this,d).first; }-----// 1e
----int res = 1;-----// cb
                                        ----for (int i = 1; i <= k; i++) res = res * (n - (k - i)) / i;-----// bd
-----return divmod(*this,d).second * sign; }------// 5a
                                        ----return res;------// e4
                                         }-----// 03
5.2.1. Fast Multiplication. Fast multiplication for the big integer using Fast Fourier Transform.
                                        5.4. Euclidean algorithm. The Euclidean algorithm computes the greatest common divisor of two
#include "intx.cpp"-----// 83
                                         integers a, b.
#include "fft.cpp"-----// 13
-----// e0
                                        int qcd(int a, int b) { return b == 0 ? a : qcd(b, a % b); }-----// d9
intx fastmul(const intx &an, const intx &bn) {------// ab
                                          The extended Euclidean algorithm computes the greatest common divisor d of two integers a, b
----string as = an.to_string(), bs = bn.to_string();-----// 32
                                         and also finds two integers x, y such that a \times x + b \times y = d.
----int n = size(as), m = size(bs), l = 1,-----// dc
                                         int egcd(int a, int b, int& x, int& y) {------// 85
-----len = 5, radix = 100000,-----// 4f
                                         ----if (b == 0) { x = 1; y = 0; return a; }------// 7b
-----*a = new int[n], alen = 0,-----// b8
                                         ----else {------// 00
-----*b = new int[m], blen = 0;------// 0a
                                         -----int d = eqcd(b, a % b, x, y);-----// 34
----memset(a, 0, n << 2);-----// 1d
                                         -----x -= a / \dot{b} * y;------// 4a
----memset(b, 0, m << 2);-----// 01
                                         ------swap(x, v):-----// 26
----for (int i = n - 1; i >= 0; i -= len, alen++)------// 6e
                                         -----return d:-----// db
------for (int j = min(len - 1, i); j >= 0; j--)------// 43
                                         ----}-----// 9e
-----a[alen] = a[alen] * 10 + as[i - j] - '0';-----// 14
                                        ----for (int i = m - 1; i >= 0; i -= len, blen++)-----// b6
------for (int j = min(len - 1, i); j >= 0; j--)-----// ae
                                        5.5. Trial Division Primality Testing. An optimized trial division to check whether an integer is
------b[blen] = b[blen] * 10 + bs[i - j] - '0';------// 9b
----while (l < 2*max(alen,blen)) l <<= 1;------// 51
                                        bool is_prime(int n) {------// 6c
----cpx *A = new cpx[l], *B = new cpx[l];-----// 0d
                                        ----if (n < 2) return false;-----// c9
----for (int i = 0; i < l; i++) A[i] = cpx(i < alen ? a[i] : 0, 0);-----// 35
                                        ----if (n < 4) return true;------// d9
----for (int i = 0; i < l; i++) B[i] = cpx(i < blen ? b[i] : 0, 0);------// 66 ----if (n % 2 == 0 || n % 3 == 0) return false;------// 0f
----fft(A, l); fft(B, l);------// f9 ----if (n < 25) return true;------// ef
----fft(A, l, true);------// d3 ----for (int i = 5; i <= s; i += 6)------// 6c
----ull *data = new ull[l];------// e7 ------if (n % i == 0 || n % (i + 2) == 0) return false;------// e9
----for (int i = 0; i < l; i++) data[i] = (ull)(round(real(A[i])));------// 06
                                         ----return true: }------// 43
----for (int i = 0; i < l - 1; i++)-----// 90
-----if (data[i] >= (unsigned int)(radix)) {------// 44
                                        5.6. Sieve of Eratosthenes. An optimized implementation of Eratosthenes' Sieve.
------data[i] %= radix;-------// bd ----int mx = (n - 3) >> 1, sq, v, i = -1;--------// 27
----stringstream ss;------// 42 ----if (n >= 2) primes.push_back(2);------// f4
----for (int i = stop - 1; i >= 0; i--)-----// bd ------primes.push_back(v = (i << 1) + 3);-----// be
-----ss << setfill('0') << setw(len) << data[i];-------// b6 -----if ((sq = i * ((i << 1) + 6) + 3) > mx) break;------// 2d
```

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------(a+b)/2, delta) + integrate(f, (a+b)/2, b, delta);------// 0c
----while (++i <= mx) if (prime[i]) primes.push_back((i << 1) + 3):------// 29 }------
----delete[] prime; // can be used for O(1) lookup------// 36
Fourier transform. Note that this implementation only handles powers of two, make sure to pad with
5.7. Modular Multiplicative Inverse. A function to find a modular multiplicative inverse.
#include "egcd.cpp"-----// 55
                                                          #include <complex>-----// 8e
                                                          typedef complex<long double> cpx:-----// 25
int mod_inv(int a, int m) {------// 49
                                                          void fft(cpx *x, int n, bool inv=false) {-----// 23
----int x, y, d = egcd(a, m, x, y);-----// 3e
                                                          ----for (int i = 0, j = 0; i < n; i++) {-------// f2
---if (d != 1) return -1;-----// 20
                                                          ------if (i < j) swap(x[i], x[j]);-----// 5c
----return x < 0 ? x + m : x:-----// 3c
                                                          -----int m = n>>1:-----// e5
                                                          ------while (1 <= m && m <= j) j -= m, m >>= 1;------// fe
                                                           -----i += m:-----// ab
5.8. Modular Exponentiation. A function to perform fast modular exponentiation.
                                                           ----}-------// 1e
template <class T>-----// 82
                                                          ----for (int mx = 1; mx < n; mx <<= 1) {------// 9d
T mod_pow(T b, T e, T m) {------// aa
                                                          -----cpx wp = \exp(\text{cpx}(0, (\text{inv }? -1 : 1) * \text{pi }/\text{mx})), \text{ w} = 1;
----T res = T(1):-----// 85
                                                          ------for (int m = 0; m < mx; m++, w *= wp) {------// 40
----while (e) {------// b7
                                                          ------for (int i = m; i < n; i += mx << 1) {------// 33
-----if (e & T(1)) res = mod(res * b. m):-----// 41
                                                          -----cpx t = x[i + mx] * w:-----// f5
-----b = mod(b * b, m), e >>= T(1); }------// b3
                                                          -----x[i + mx] = x[i] - t:-----// ac
----return res:------// eb
                                                          -----x[i] += t;-----// c7
                                                          ------}------// c2
5.9. Chinese Remainder Theorem. An implementation of the Chinese Remainder Theorem.
                                                           ----}-----// 70
#include "egcd.cpp"-----// 55
                                                           ----if (inv) for (int i = 0: i < n: i++) x[i] /= cpx(n):------// 3e
int crt(const vi& as, const vi& ns) {-----// c3
                                                          }-----// 7d
----int cnt = size(as), N = 1, x = 0, r, s, l;------// 55
----for (int i = 0: i < cnt: i++) N *= ns[i]: ------// 88
                                                          5.13. Formulas.
----for (int i = 0; i < cnt; i++)-----// f9
                                                              • Number of ways to choose k objects from a total of n objects where order matters and each
-----egcd(ns[i], l = N/ns[i], r, s), x += as[i] * s * l;-----// b0
                                                               item can only be chosen once: P_k^n = \frac{n!}{(n-k)!}
----return mod(x, N); }-----// 9e
                                                              \bullet Number of ways to choose k objects from a total of n objects where order matters and each
                                                               item can be chosen multiple times: n^k
5.10. Linear Congruence Solver. A function that returns all solutions to ax \equiv b \pmod{n}, modulo
                                                              • Number of permutations of n objects, where there are n_1 objects of type 1, n_2 objects of type
                                                               2, ..., n_k objects of type k: \binom{n}{n_1, n_2, ..., n_k} = \frac{n!}{n_1! \times n_2! \times \cdots \times n_k!}
#include "egcd.cpp"-----// 55
vi linear_congruence(int a, int b, int n) {-----// c8
                                                              • Number of ways to choose k objects from a total of n objects where order does not matter
                                                               and each item can only be chosen once:
----int x, y, d = egcd(a, n, x, y);-----// 7a
                                                               \binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k} = \binom{n}{n-k} = \prod_{i=1}^{k} \frac{n-(k-i)}{i} = \frac{n!}{k!(n-k)!}, \binom{n}{0} = 1, \binom{0}{k} = 0
----vi res:-----// f5
                                                              • Number of ways to choose k objects from a total of n objects where order does not matter
----if (b % d != 0) return res;-----// 30
                                                               and each item can be chosen multiple times: f_k^n = \binom{n+k-1}{k} = \frac{(n+k-1)!}{k!(n-1)!}
----int x0 = mod(b / d * x, n);------// 48
----for (int k = 0; k < d; k++) res.push_back(mod(x0 + k * n / d, n));------// 21
                                                              • Number of integer solutions to x_1 + x_2 + \cdots + x_n = k where x_i > 0: f_k^n
----return res:-----// 03
                                                              • Number of subsets of a set with n elements: 2^n
}-----// 1c
                                                              • |A \cup B| = |A| + |B| - |A \cap B|
                                                              • |A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|
5.11. Numeric Integration. Numeric integration using Simpson's rule.
                                                              • Number of ways to walk from the lower-left corner to the upper-right corner of an n \times m grid
                                                               by walking only up and to the right: \binom{n+m}{m}
double integrate(double (*f)(double), double a, double b,-----// 76
-----double delta = 1e-6) {------// c0
                                                              • Number of strings with n sets of brackets such that the brackets are balanced:
----if (abs(a - b) < delta)------// 38
                                                               C_n = \sum_{k=0}^{n-1} C_k C_{n-1-k} = \frac{1}{n+1} {2n \choose n}
-----return (b-a)/8 *-----// 56
                                                              • Number of triangulations of a convex polygon with n points, number of rooted binary trees
-----(f(a) + 3*f((2*a+b)/3) + 3*f((a+2*b)/3) + f(b));-----// e1
                                                               with n+1 vertices, number of paths across an n \times n lattice which do not rise above the main
----return integrate(f, a,-----// 64
                                                               diagonal: C_n
```

Reykjavik University ----**else return** (real(p) - real(a)) / (real(b) - real(a)); }------// 2c

- Number of permutations of n objects with exactly k ascending sequences or runs:  $\left\langle {n\atop k}\right\rangle = \left\langle {n\atop n-k-1}\right\rangle = k\left\langle {n-1\atop k}\right\rangle + (n-k+1)\left\langle {n-1\atop k-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle = \sum_{i=0}^k (-1)^i {n+1\choose i} (k+1-i)^n, \left\langle {n\atop 0}\right\rangle = \left\langle {n\atop n-1}\right\rangle =$ • Number of permutations of n objects with exactly k cycles:  $\binom{n}{k} = \binom{n-1}{k-1} + (n-1)\binom{n-1}{k}$
- Number of ways to partition n objects into k sets:  $\binom{n}{k} = k \binom{n-1}{k} + \binom{n-1}{k-1}, \binom{n}{0} = \binom{n}{n} = 1$
- Number of permutations of length n that have no fixed points (derangements):  $D_0 = 1, D_1 =$
- $0, D_n = (n-1)(D_{n-1} + D_{n-2})$
- Number of permutations of length n that have exactly k fixed points:  $\binom{n}{k}D_{n-k}$
- Heron's formula: A triangle with side lengths a, b, c has area  $\sqrt{s(s-a)(s-b)(s-c)}$  where • Pick's theorem: A polygon on an integer grid containing i lattice points and having b lattice
- points on the boundary has area  $i + \frac{b}{2} 1$ . • Divisor sigma: The sum of divisors of n to the xth power is  $\sigma_x(n) = \prod_{i=0}^r \frac{p_i^{(a_i+1)x}-1}{p_i^x-1}$  where
- $n = \prod_{i=0}^{r} p_i^{a_i}$  is the prime factorization. • Divisor count: A special case of the above is  $\sigma_0(n) = \prod_{i=0}^r (a_i + 1)$ .
- Euler's totient: The number of integers less than n that are comprime to n are  $n \prod_{p|n} \left(1 \frac{1}{p}\right)$ where each p is a distinct prime factor of n. • König's theorem: In any bipartite graph, the number of edges in a maximum matching is
- equal to the number of vertices in a minimum vertex cover. • The number of vertices of a graph is equal to its minimum vertex cover number plus the size
- of a maximum independent set.

# 6. Geometry

```
6.1. Primitives. Geometry primitives.
#include <complex>-----// 8e
#define P(p) const point &p-----// b8
#define L(p0, p1) P(p0), P(p1)-----// 30
typedef complex<double> point;-----// e1
double dot(P(a), P(b)) { return real(conj(a) * b); }-----// a9
double cross(P(a), P(b)) { return imag(conj(a) * b); }-----// ff
point rotate(P(p), P(about), double radians) {------// e1
----return (p - about) * exp(point(0, radians)) + about; }-----// cb
point reflect(P(p), L(about1, about2)) {-----// c0
----point z = p - about1, w = about2 - about1;---------------// 39 double polygon_area_signed(polygon p) {-----------------------// 31
point proj(P(u), P(v)) { return dot(u, v) / dot(u, u) * u; }-------// fc ----for (int i = 1; i + 1 < cnt; i++)----------------// d2
bool parallel(L(a, b), L(p, q)) { return abs(cross(b - a, q - p)) < EPS; }-----// 6d ------area += cross(p[i] - p[0], p[i + 1] - p[0]);-------------------------// 7e
bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b, c)) < EPS; }------// 75 double polygon_area(polygon p) { return abs(polygon_area_signed(p)); }------// 25
----return abs(ccw(a, b, p)) < EPS && abs(ccw(a, b, q)) < EPS; }------// d6 int point_in_polygon(polygon p, point q) {-----------------------// 58
----return acos(dot(b - a, c - b) / abs(b - a) / abs(c - b)); }------------// cc -----for (int i = 0, j = n - 1; i < n; j = i++)-------------// 77
----return asin(cross(b - a, c - b) / abs(b - a) / abs(c - b)); }------// 9e -------0 <= (d = progress(q, p[i], p[j])) && d <= 1)------// b9
-----return (imag(p) - imag(a)) / (imag(b) - imag(a));------// 35 ------if (CHK(real, p[i], q, p[j]) || CHK(real, p[j], q, p[i]))------// 1f
```

```
bool intersect(L(a, b), L(p, q), point &res, bool segment = false) \{-----//d6\}
----// NOTE: check for parallel/collinear lines before calling this function---// 02
----point r = b - a, s = q - p;------// 79
----double c = cross(r, s), t = cross(p - a, s) / c, u = cross(p - a, r) / c;--// a8
----if (segment && (t < 0-EPS || t > 1+EPS || u < 0-EPS || u > 1+EPS))------// ae
-----return false:-----// a3
----res = a + t * r:-----// ca
----return true:------// 17
point closest_point(L(a, b), P(c), bool segment = false) {------// a1
----if (segment) {-------// c2
-----if (dot(b - a, c - b) > 0) return b;------// b5
-----if (dot(a - b, c - a) > 0) return a;-----// cf
----}-------// 61
----double t = dot(c - a, b - a) / norm(b - a);------// aa
----return a + t * (b - a);-----// 7a
double line_segment_distance(L(a,b), L(c,d)) {------// 99
----double x = INFINITY:-----// 83
----if (abs(a - b) < EPS && abs(c - d) < EPS) x = abs(a - c):-----// df
----else if (abs(a - b) < EPS) x = abs(a - closest_point(c, d, a, true)); -----// da
----else if (abs(c - d) < EPS) \times = abs(c - closest_point(a, b, c, true)):-----// 52
----else if ((ccw(a, b, c) < 0) != (ccw(a, b, d) < 0) &&------// ee
----- (ccw(c, d, a) < 0) != (ccw(c, d, b) < 0)) x = 0; ------//79
-----x = min(x, abs(a - closest_point(c,d, a, true)));------// f3
-----x = min(x, abs(b - closest_point(c,d, b, true)));-----// ec
-----x = min(x, abs(c - closest_point(a,b, c, true)));
-----x = min(x, abs(d - closest_point(a,b, d, true)));
----}----------// 72
----return x:-----// 0d
}-----// h3
6.2. Polygon. Polygon primitives.
#include "primitives.cpp"-----// e0
typedef vector<point> polygon;-----// b3
```

```
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-----in = !in;------// b2 ------A = B = c; double p = progress(c, a,b);------// d9
----return in ? -1 : 1: }------------------------// 77 -------return 0.0 <= p && p <= 1.0--------------------// 8e
//--- polygon left, right;------// 6b ---else if (collinear(a,b, c,d)) {----------------------// bc
//--- point it(-100, -100);------// c9 ------double ap = progress(a, c,d), bp = progress(b, c,d);------// a7
//------ point p = poly[i], q = poly[i]:------// 19 ------A = c + max(ap, 0.0) * (d - c):------// f6
//------ if (ccw(a, b, p) \le 0) left.push_back(p);------// 12 ------B = c + min(bp, 1.0) * (d - c);-------// 5c
//----- if (ccw(a, b, p) >= 0) right.push_back(p);------// e3 ------return true; }-----
//-----// myintersect = intersect where-----// 24 ----else if (parallel(a,b, c,d)) return false;------// ca
//-----// (a,b) is a line, (p,q) is a line segment------// f2 ----else if (intersect(a,b, c,d, A, true)) {----------------// 10
//----- if (myintersect(a, b, p, q, it))-------// f0 -----B = A; return true; }------
//--- return pair<polygon, polygon>(left, right);------// 1d ------// e6
                                      6.5. Great-Circle Distance. Computes the distance between two points (given as latitude/longitude
6.3. Convex Hull. An algorithm that finds the Convex Hull of a set of points.
                                      coordinates) on a sphere of radius r.
#include "polygon.cpp"-----// 58
#define MAXN 1000-----// 09
                                      double gc_distance(double pLat, double pLong,-----// 7b
point hull[MAXN];-----// 43
                                      ------ double qLat, double qLong, double r) {------// a4
----return abs(real(a) - real(b)) > EPS ?------// 44 ----qLat *= pi / 180; qLong *= pi / 180;-------// 75
                                      ----return r * acos(cos(pLat) * cos(qLat) * cos(pLong - qLong) +------// e3
-----real(a) < real(b) : imag(a) < imag(b): }-----// 40
----for (int i = 0; i < n; i++) {-------// 6f
-----if (i > 0 && p[i] == p[i - 1]) continue;-----// b2
                                      6.6. Closest Pair of Points. A sweep line algorithm for computing the distance between the closest
------while (l >= 2 \& ccw(hull[l - 2], hull[l - 1], p[i]) >= 0) l--;------// 20
                                      pair of points.
-----hull[l++] = p[i];-----// f7
                                      #include "primitives.cpp"-----// e0
----}-----// d8
                                        -----// 85
----int r = l:------// 59
                                      struct cmpx { bool operator ()(const point &a, const point &b) {------// 01
----for (int i = n - 2; i >= 0; i--) {------// 16
                                      -----return abs(real(a) - real(b)) > EPS ?-----// e9
-----if (p[i] == p[i + 1]) continue;-----// c7
                                      -----real(a) < real(b) : imag(a) < imag(b); } };------// 53
------while (r - l >= 1 \&\& ccw(hull[r - 2], hull[r - 1], p[i]) >= 0) r--;----// 9f
                                      struct cmpy { bool operator ()(const point &a, const point &b) {------// 6f
-----hull[r++] = p[i]:-----// 6d
                                      ----return abs(imag(a) - imag(b)) > EPS ?-----// 0b
-----imag(a) < imag(b) : real(a) < real(b); \} };------// a4
----return l == 1 ? 1 : r - 1;------// 6d
                                      double closest_pair(vector<point> pts) {------// f1
}-----// 79
                                      ----sort(all(pts), cmpx());------// 51
6.4. Line Segment Intersection. Computes the intersection between two line segments.
                                      ----set<point, cmpy> cur;-----// a5
                                      ----set<point, cmpy>::const_iterator it, jt;------// 48
#include "primitives.cpp"-----// e0
                                      ----double mn = INFINITY;-----// a4
bool line_segment_intersect(L(a, b), L(c, d), point &A, point &B) {------// 6c
                                      ----for (int i = 0, l = 0; i < size(pts); i++) {-------// 40
----if (abs(a - b) < EPS && abs(c - d) < EPS) {-----// db
------A = B = a; return abs(a - d) < EPS; }-------// ee -------while (real(pts[i]) - real(pts[l]) > mn) cur.erase(pts[l++]);------// ff
------A = B = a; double p = progress(a, c,d);-------// c9 -----jt = cur.upper_bound(point(INFINITY, imag(pts[i]) + mn));------// 58
```

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7. Formulas. Let $a = (a_x, a_y)$ and $b = (b_x, b_y)$ be two-dimensional vectors.	if (clauses[i].first != clauses[i].second)/	
• $a \cdot b =  a  b \cos\theta$ , where $\theta$ is the angle between $a$ and $b$ .	adj[-clauses[i].second + n].push_back(clauses[i].first + n);/	
• $a \times b =  a  b \sin\theta$ , where $\theta$ is the signed angle between $a$ and $b$ .	}}	
• $a \times b$ is equal to the area of the parallelogram with two of its sides formed by a and b. Half	pair <union_find, vi=""> res = scc(adj);/</union_find,>	
of that is the area of the triangle formed by $a$ and $b$ .	union_find scc = res.first;/	
	vi dag = res.second;/	
7. Other Algorithms	vi truth(2*n+1, -1);/	
1. Binary Search. An implementation of binary search that finds a real valued root of the continous	for (int i = 2*n; i >= 0; i) {/	
inction f on the interval $[a, b]$ , with a maximum error of $\varepsilon$ .	int cur = order[i] - n, p = scc.find(cur + n), o = scc.find(-cur + n);-/	
	if (cur == 0) continue;/	// 26
<pre>puble binary_search_continuous(double low, double high,// 8e</pre>	if (p == o) return false;/	// 33
double eps, double (*f)(double)) {// c0	if (truth[p] == -1) truth[p] = 1;/	
while (true) {// 3a	truth[cur + n] = truth[p];/	// b3
double mid = (low + high) / 2, cur = f(mid);// 75	truth[o] = 1 - truth[p];/	// 80
<b>if</b> (abs(cur) < eps) <b>return</b> mid;// /6	if (truth[p] == 1) all_truthy.push_back(cur);/	// 5c
else if (0 < cur) high = mid;// e5	}}	
else low = mid;// a7	<b>return</b> true;/	
}// b5	}/	// 61
// cb	7.4 Ct-11 M The Cele Chember describes for a big at heatable second and beginning	
Another implementation that takes a binary predicate $f$ , and finds an integer value $x$ on the integer	7.4. <b>Stable Marriage.</b> The Gale-Shapley algorithm for solving the stable marriage problem.	
terval $[a,b]$ such that $f(x) \land \neg f(x-1)$ .	<pre>vi stable_marriage(int n, int** m, int** w) {</pre>	'/ e4
<pre>nt binary_search_discrete(int low, int high, bool (*f)(int)) {// 51</pre>	queue< <b>int</b> > q;/	
assert(low <= high);// 19	vi at(n, 0), eng(n, -1), res(n, -1); vvi inv(n, vi(n));/	
while (low < high) {// a3	for (int $i = 0$ ; $i < n$ ; $i++$ ) for (int $j = 0$ ; $j < n$ ; $j++$ )/	
<mark>int</mark> mid = low + (high - low) / 2;// 04	inv[i][w[i][j]] = j;/	// b9
if (f(mid)) high = mid;// ca	for (int i = 0; i < n; i++) q.push(i);/	// fe
else low = mid + 1;// 03	while (!q.empty()) {/	// 55
}// 9b	int curm = q.front(); q.pop();/	
assert(f(low));// 42	for (int &i = at[curm]; i < n; i++) {/	
<b>return</b> low;// a6	int curw = m[curm][i];/	
// d3	if (eng[curw] == -1) { }/	
	else if (inv[curw][curm] < inv[curw][eng[curw]])/	// 10
2. <b>Ternary Search.</b> Given a function $f$ that is first monotonically increasing and then monotonically increasing and the monotonical and the mono		
ally decreasing, ternary search finds the $x$ such that $f(x)$ is maximized.	else continue;/	
emplate <class f="">// d1</class>		
<pre>puble ternary_search_continuous(double lo, double hi, double eps, F f) {// e7</pre>	}	// 24

```
------double m1 = lo + (hi - lo) / 3, m2 = hi - (hi - lo) / 3;-----// e8
-----if (f(m1) < f(m2)) lo = m1;------// 1d
-----else hi = m2;-----// b3
}-----// 66
```

```
7.3. 2SAT. A fast 2SAT solver.
#include "../graph/scc.cpp"-----// c3 -----node *l, *r, *u, *d, *p;------// 19
-----// 63 ------<mark>int</mark> row, col, size;------
bool two_sat(int n, const vii& clauses, vi& all_truthy) {------// f4 -----node(int row, int col) : row(row), col(col) {-------// 68
```

7.5. Algorithm X. An implementation of Knuth's Algorithm X, using dancing links. Solves the Exact Cover problem. bool handle\_solution(vi rows) { return false; }-----// 63

----return res;------// 95

```
struct exact_cover {------// 95
----struct node {------// 7e
```

```
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----exact_cover(int rows, int cols) : rows(rows), cols(cols), head(NULL) {-----// ce
                                          -----arr = new bool*[rows];-----// 15
                                          ------for (node *j = i->r; j != i; j = j->r) \\ --------------// 0e
-----sol = new int[rows];------// 69
                                          -----j->d->u = j->u, j->u->d = j->d, j->\overline{p}->size--;-----// 5a
------for (int i = 0; i < rows; i++)------// c7
                                          ----\#define UNCOVER(c, i, j) \mathbb{N}------// 17
-----arr[i] = new bool[cols], memset(arr[i], 0, cols);------// 68
                                          ------for (node *i = c->u; i != c; i = i->u) \sqrt{\phantom{a}}
----void set_value(int row, int col, bool val = true) { arr[row][col] = val; }-// af
                                          -----j->p->size++, j->d->u = j->u->d = j; \sqrt{\phantom{a}}
----void setup() {------// a8
                                          ------c->r->l = c->l->r = c;------// bb
-----node ***ptr = new node**[rows + 1];-----// da
------for (int i = 0; i <= rows; i++) {------// ce
                                          ----bool search(int k = 0) {------// 4f
                                          -----if (head == head->r) {------// a7
-----ptr[i] = new node*[cols];-----// cc
                                          -----vi res(k);-----// 4f
------for (int j = 0; j < cols; j++)-----// 56
                                          ------for (int i = 0; i < k; i++) res[i] = sol[i];------// c\theta
-----if (i == rows \mid | arr[i][j]) ptr[i][j] = new node(i, j); ------// 95
                                          -----sort(res.begin(), res.end());-----// 3e
-----else ptr[i][j] = NULL;-----// 40
                                          -----return handle_solution(res);-----// dc
                                          -----for (int i = 0; i <= rows; i++) {------// 80
-----for (int j = 0; j < cols; j++) {------// 86
                                          -----node *c = head->r, *tmp = head->r;------// a6
                                          -----for ( ; tmp != head; tmp = tmp->r) if (tmp->size < c->size) c = tmp;---// 1e
-----if (!ptr[i][j]) continue;-----// 76
                                          -----if (c == c->d) return false;-----// 17
-----int ni = i + 1, nj = j + 1;-----// 34
                                          -----COVER(c, i, j);-----// 61
-----while (true) {-----// 7f
-----if (ni == rows + 1) ni = 0;-----// 54
                                          ------bool found = false;-----// 6e
-----if (ni == rows || arr[ni][j]) break;-----// 77
                                          ------for (node *r = c->d; !found && r != c; r = r->d) {------// 1e
-----+ni;-----// c8
                                          -----sol[k] = r->row;------// θb
                                          -----for (node *j = r - r; j != r; j = j - r) { COVER(j - p, a, b); } -----// 3a
-----found = search(k + 1);-----// f4
-----ptr[i][j]->d = ptr[ni][j];-----// a9
                                          -----for (node *j = r > 1; j = j > 1) { UNCOVER(j > p, a, b); }----// 8a
  -----ptr[ni][j]->u = ptr[i][j];-----// c0
                                          ------while (true) {------// 0d
                                          ------UNCOVER(c, i, j);------// 64
-----if (nj == cols) nj = 0;-----// a7
                                          -----return found;------// ff
-----if (i == rows || arr[i][nj]) break;------// e9
                                          -----+nj;-----// a6
                                          };-----// 10
-----ptr[i][j]->r = ptr[i][nj];------// b3
                                          7.6. nth Permutation. A very fast algorithm for computing the nth permutation of the list \{0, 1, \dots, k-1\}
-----ptr[i][nj]->l = ptr[i][j];------// 46
  -----}-----// 83
-----}-----// b4
                                          vector<int> nth_permutation(int cnt, int n) {-----// 78
                                          ----vector<int> idx(cnt), per(cnt), fac(cnt);------// 9e
------head = new node(rows, -1);------// 80
                                          ----for (int i = 0; i < cnt; i++) idx[i] = i;------// 80
------head->r = ptr[rows][0];------// b9
                                          ----for (int i = 1; i <= cnt; i++) fac[i - 1] = n % i, n /= i;-----// 04
-----ptr[rows][0]->l = head;-----// c1
------head->l = ptr[rows][cols - 1];------// 28
                                          ----for (int i = cnt - 1; i >= 0; i--)------// 52
                                          -----per[cnt - i - 1] = idx[fac[i]], idx.erase(idx.begin() + fac[i]);-----// 41
-----ptr[rows][cols - 1]->r = head;------// 83
                                          ----return per;------// 84
------for (int j = 0; j < cols; j++) {-------// 02
                                          }-----// 97
------int cnt = -1;------// 36
-----for (int i = 0; i <= rows; i++)-----// 56
                                          7.7. Cycle-Finding. An implementation of Floyd's Cycle-Finding algorithm.
-----if (ptr[i][j]) cnt++, ptr[i][j]->p = ptr[rows][j];-----// 05
                                          ii find_cycle(int x0, int (*f)(int)) {------// a5
-----ptr[rows][j]->size = cnt;-----// d4
                                          ----int t = f(x0), h = f(t), mu = 0, lam = 1;------// 8d
-----for (int i = 0; i <= rows; i++) delete[] ptr[i];-----// cd
                                          ----while (t != h) t = f(t), h = f(f(h));
-----delete[] ptr;-----// 42
                                          ----h = x0;
                                          ----while (t != h) t = f(t), h = f(h), mu++;------// 9d
----}------// a9
                                          ----h = f(t);-----// 00
----#define COVER(c, i, j) N-----// 23
                                          ----while (t != h) h = f(h), lam++;-----// 5e
```

```
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}------// 42 ----int sign = 1;-------// 32
7.8. Dates. Functions to simplify date calculations.
-----367 * (m - 2 - (m - 14) / 12 * 12) / 12 -------// d1 ------case ' ': goto hell;-------// fd
-----3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +-------// be ------case '\n': goto hell;------// 79
void intToDate(int jd, int &y, int &m, int &d) {------// a1 ---}------// c3
----int x, n, i, j;------// 00
----x = id + 68569;-----// 11
----n = 4 * x / 146097;-----// 2f
---x = (146097 * n + 3) / 4;
----i = (4000 * (x + 1)) / 1461001;-----// 0d
----x -= 1461 * i / 4 - 31;-----// 09
---j = 80 * x / 2447;-----// 3d
----d = x - 2447 * j / 80;-----// eb
---x = i / 11:-----// b7
---m = j + 2 - 12 * x;
---y = 100 * (n - 49) + i + x;
}-----// af
```

### 8. Useful Information

#### 8.1. Tips & Tricks.

- How fast does our algorithm have to be? Can we use brute-force?
- Does order matter?
- Is it better to look at the problem in another way? Maybe backwards?
- Are there subproblems that are recomputed? Can we cache them?
- Do we need to remember everything we compute, or just the last few iterations of computation?
- Does it help to sort the data?
- Can we speed up lookup by using a map (tree or hash) or an array?
- Can we binary search the answer?
- Can we add vertices/edges to the graph to make the problem easier? Can we turn the graph into some other kind of a graph (perhaps a DAG, or a flow network)?
- Make sure integers are not overflowing.
- Is it better to compute the answer modulo n? Perhaps we can compute the answer modulo  $m_1, m_2, \ldots, m_k$ , where  $m_1, m_2, \ldots, m_k$  are pairwise coprime integers, and find the real answer using CRT?
- Are there any edge cases? When  $n=0, n=-1, n=1, n=2^{31}-1$  or  $n=-2^{31}$ ? When the list is empty, or contains a single element? When the graph is empty, or contains a single vertex? When the graph contains self-loops? When the polygon is concave or non-simple?
- Can we use exponentiation by squaring?
- 8.2. Fast Input Reading. If input or output is huge, sometimes it is beneficial to optimize the input reading/output writing. This can be achieved by reading all input in at once (using fread), and then parsing it manually. Output can also be stored in an output buffer and then dumped once in the end (using fwrite). A simpler, but still effective, way to achieve speed is to use the following input reading method.

```
----register char c;-----// a5
----*n = 0:-----// 35
hell:----// ba
---*n *= sign;-----// a0
}-----// 67
```

8.3. 128-bit Integer. GCC has a 128-bit integer data type named \_\_intl28. Useful if doing multiplication of 64-bit integers, or something needing a little more than 64-bits to represent.

## 8.4. Worst Time Complexity.

	n	Worst AC Algorithm	Comment
_	$\leq 10$	$O(n!), O(n^6)$	e.g. Enumerating a permutation
	$\leq 15$	$O(2^n \times n^2)$	e.g. DP TSP
	$\leq 20$	$O(2^n), O(n^5)$	e.g. $DP + bitmask technique$
	$\leq 50$	$O(n^4)$	e.g. DP with 3 dimensions $+ O(n)$ loop, choosing ${}_{n}C_{k} = 4$
	$\leq 10^{2}$	$O(n^3)$	e.g. Floyd Warshall's
	$\leq 10^{3}$	$O(n^2)$	e.g. Bubble/Selection/Insertion sort
	$\leq 10^{5}$	$O(n\log_2 n)$	e.g. Merge sort, building a Segment tree
	$\leq 10^{6}$	$O(n), O(\log_2 n), O(1)$	Usually, contest problems have $n \leq 10^6$ (e.g. to read input)

#### 8.5. Bit Hacks.

- n & -n returns the first set bit in n.
- n & (n 1) is 0 only if n is a power of two.
- snoob(x) returns the next integer that has the same amount of bits set as x. Useful for iterating through subsets of some specified size.

```
----int y = x & -x, z = x + y;-----
----return z | ((x ^ z) >> 2) / y;-----
}------
```