

Convex Hull

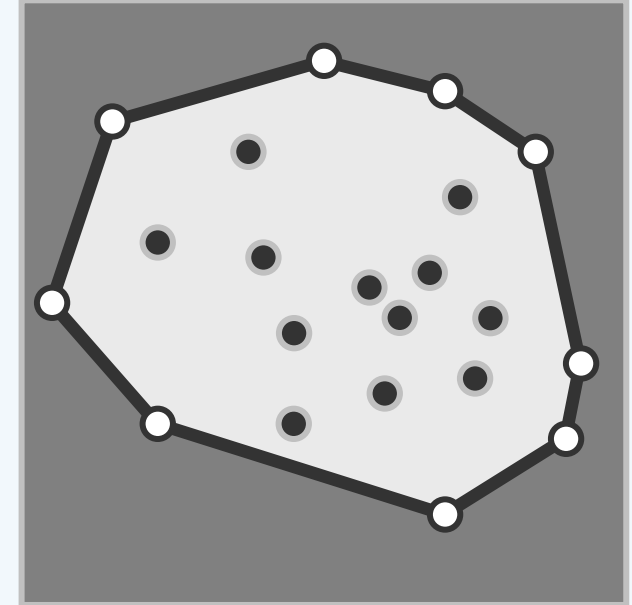
Expected Complexity

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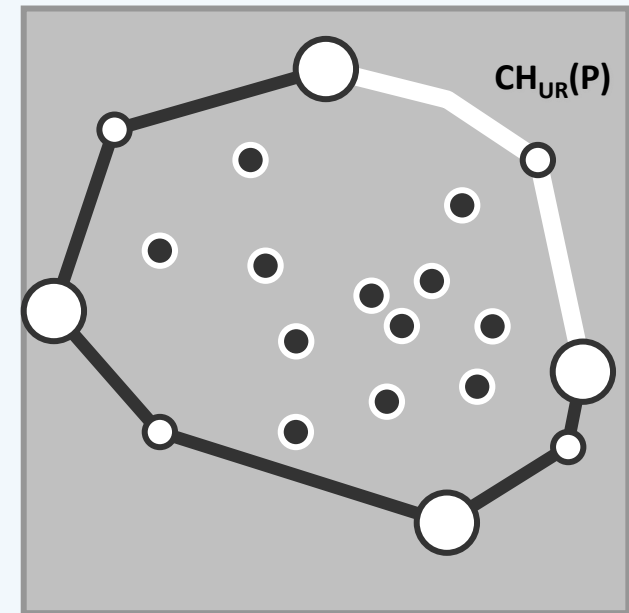
Complexity of Convex Hull

- ❖ Let P be a set of n points in the plane
How many points are expected to lie on $CH(P)$?
- ❖ Point sets with different **distributions** will give quite different answers
- ❖ As an example, this section will discuss the **uniform** and **independent** distribution inside a unit square (or affinely, a box)
- ❖ We will show that the expected number of vertices on $CH(P)$ is **$O(\log n)$**



Quadrant Hulls

- ❖ Actually, instead of the entire hull, we will count only a quarter of $CH(P)$
- ❖ $CH(P)$ is broken by its leftmost, rightmost, highest, and lowest vertices into (no more than) four parts, each of which is called a **quadrant hull**
- ❖ In terms of worst case complexity, to count the points on $CH(P)$, it suffices to count any of the four quadrant hulls
- ❖ W.L.O.G., here we examine the upper-right one, which henceforth will be denoted as $CH_{UR}(P)$



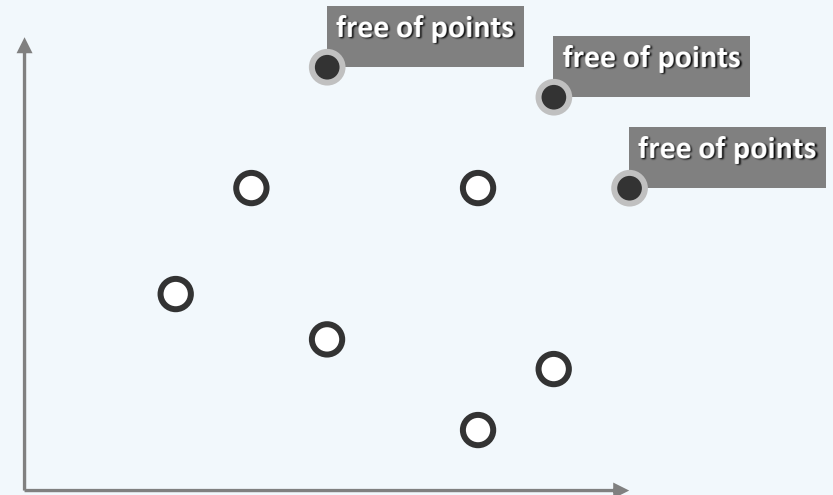
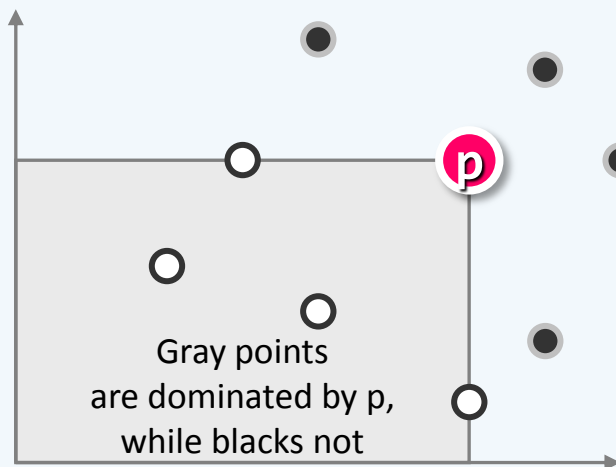
Domination

❖ Given two points $p_1 = (x_1, y_1)$ and $p_2 = (x_2, y_2)$ in the plane

p_1 is said to be **dominated** by p_2 if $x_1 \leq x_2$ and $y_1 \leq y_2$

❖ Ex: how to count dominations for any given P in $O(n \log n)$ time?

Note that there could be up to $\Omega(n^2)$ dominations



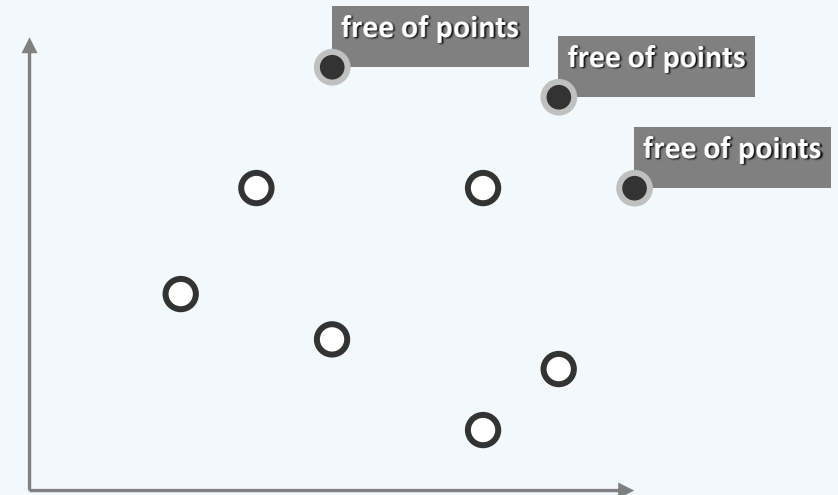
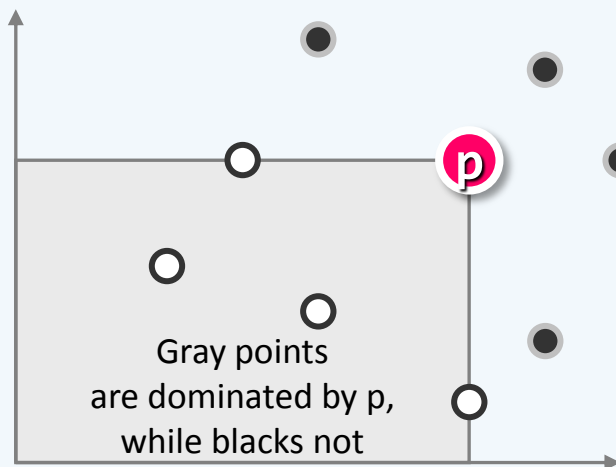
Maximal Point

❖ In fact, we will count a super set of $CH_{UR}(P)$...

❖ Let P be a finite planar point set

A **maximal point** of P is one **not** dominated by any of the other points

❖ The set of all maxima of P is denoted as **MAX(P)**



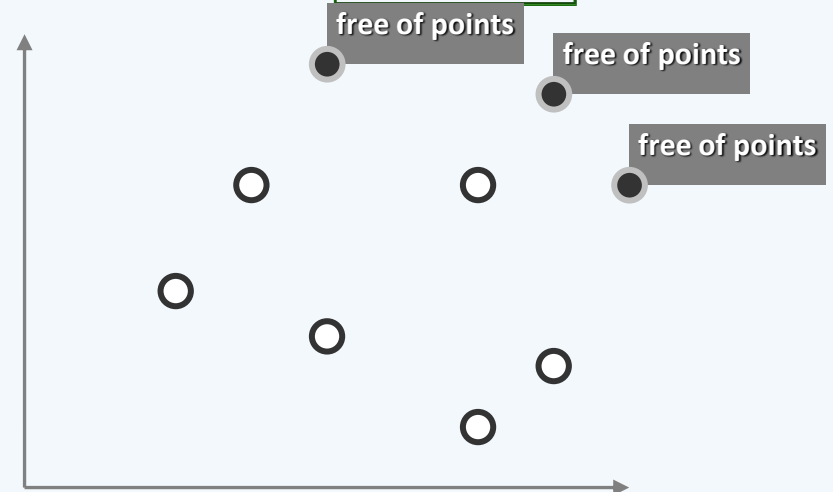
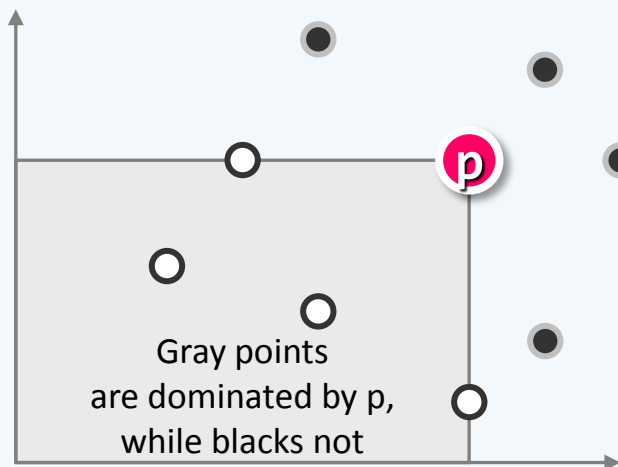
Maxima Set

❖ Note that each vertex of $CH_{UR}(P)$ is a maximal point of P

Therefore, $CH_{UR}(P)$ is a subset of $MAX(P)$, and $|CH_{UR}(P)| \leq |MAX(P)|$ and ...

❖ The expected number of points on $CH_{UR}(P)$ (and hence on $CH(P)$)
is bounded by the expected number of points in $MAX(P)$

❖ We will show next that $|MAX(P)|$ is expected to be $O(\log n)$



expected ($|\text{MAX}(P)|$) = $\mathcal{O}(\log n)$

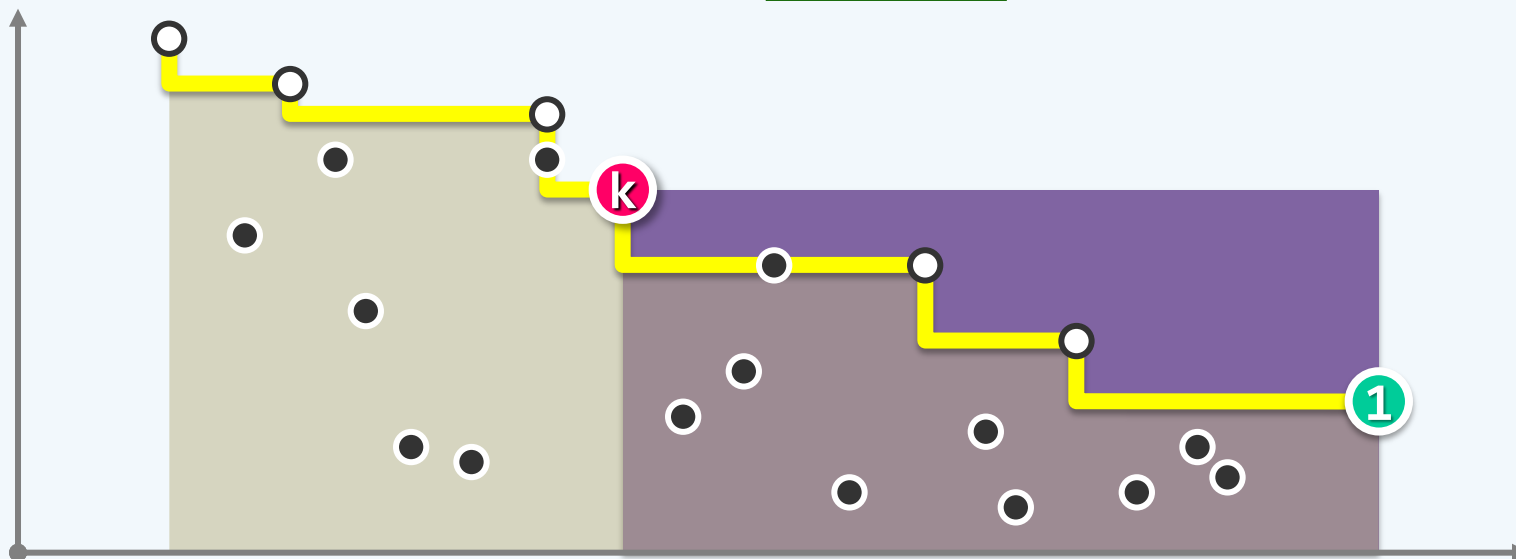
❖ Number all points in P as $\{ p_1, \dots, p_n \}$, from right to left

❖ Consider the point p_k ...

❖ Note that

p_k is a **maximal** point iff

p_k is the **highest** among $\{ p_1, \dots, p_k \}$



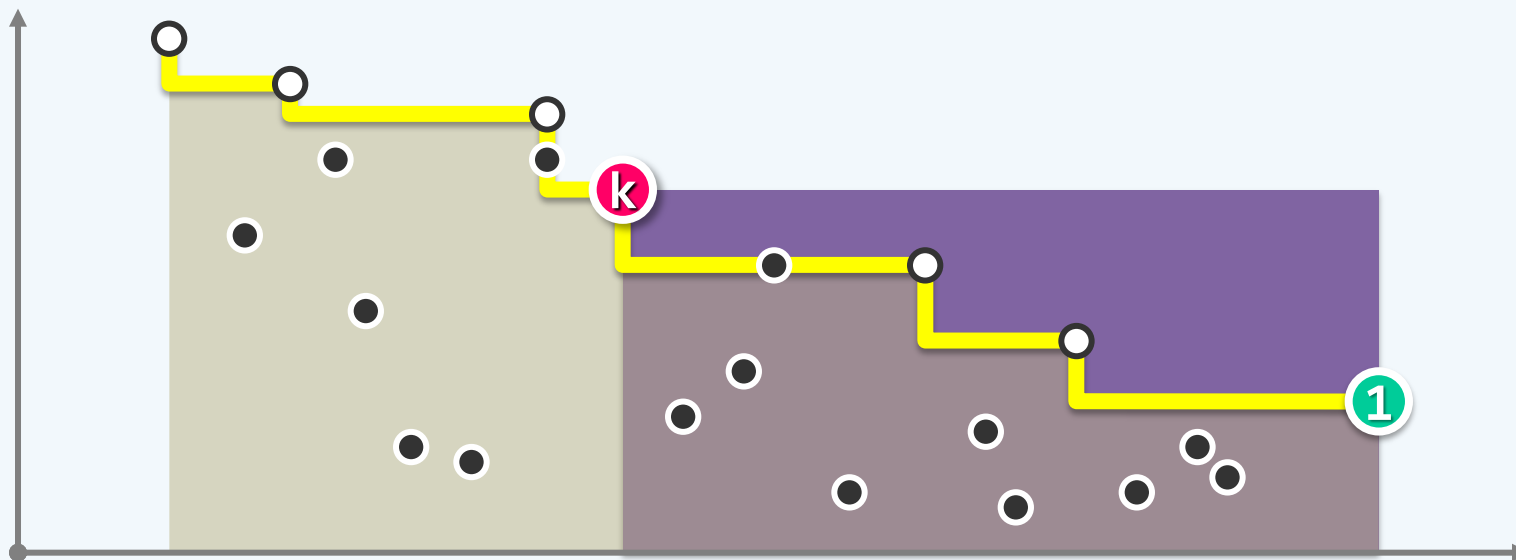
$$\text{expected } (|\text{MAX}(P)|) = \mathcal{O}(\log n)$$

❖ By the assumption of uniform and independent distribution, this is true with a probability of $1/k$

❖ By linearity of expectation

the sum of these expectations gives the expected number of maxima

$$\text{expected } (|\text{MAX}(P)|) = 1/n + 1/(n-1) + \dots + 1/3 + 1/2 + 1 = \mathcal{O}(\log n)$$



Expected Complexity & Expected Area

- ❖ For n points chosen uniformly and independently from the **unit square**,
 - 1) the expected **number of vertices** of the convex hull is $\mathcal{O}(\log n)$
 - 2) the expected **area** of the convex hull is $\mathcal{O}(1 - \log(n)/n)$
- ❖ For n points chosen uniformly and independently from the **unit disk**,
 - 1) the expected **number of vertices** of the convex hull is $\mathcal{O}(n^{1/3})$
 - 2) the expected **area** of the convex hull is $\pi * \mathcal{O}(1 - n^{-2/3})$
- ❖ For n points chosen uniformly and independently from a **triangle**,
the expected **number of vertices** of the convex hull is $\mathcal{O}(\log n)$
- ❖ For n points chosen uniformly and independently from a **convex k -gon**
the expected **number of vertices** of the convex hull is $\mathcal{O}(k \log n)$

Efron's Theorem, 1965

❖ B. Efron, *The convex hull of a random set of points*

Biometrika, 52(3):331-343, 1965

❖ Let C be a compact convex set in the plane

❖ If the expected area of the convex hull of n points,
chosen uniformly and independently from C , is

$$O(1 - f(n)) * \text{Area}(C), \quad \text{where } 1 \geq f(n) \geq 0, \text{ for } n \geq 0,$$

then the expected number of vertices of the convex hull is

$$O(n * f(n/2))$$