

Triangulation

Beyond  $\Omega(nlogn)$ 

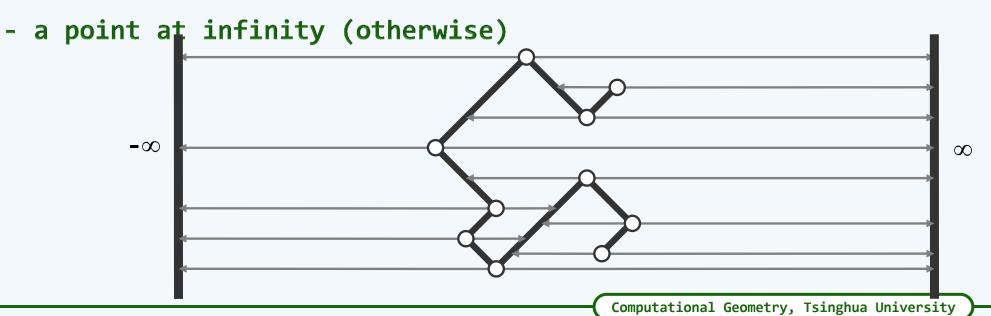
- HVP

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## Horizontal Visibility Partition

- ❖ Trapezoidalization may be extended to simple polygonal chains, where the horizontal lines run to infinity if they meet no obstruction
- ❖ The HVP of a simple polygonal chain is the partition of the plane obtained by adding horizontal edges connecting each vertex to
  - the closest point on the chain on both sides (if exists), or



## Merging HVP's

- $\diamondsuit$  Let  $P_1$  and  $P_2$  be two successively adjacent simple polygonal chains
- ❖ [Chazelle & Incerpi, 1984]

If  $P_1 \cup P_2$  is also a simple polygonal chain, then

 $HVP(P_1 \cup P_2)$  can be obtained from  $HVP(P_1)$  and  $HVP(P_2)$  in linear time

- This result indicates a natural divide-and-conquer algorithm for computing the HVP of a simple polygonal chain
- ❖[Chazelle & Incerpi, 1984]

The HVP of a simple polygonal chain can be computed

in  $|O(n\log n)|$  time, where n is the length of the chain

\*However, the performance of this algorithm is not improved