

**Point Location**

**Kirkpatrick Structure**

**- Performance**

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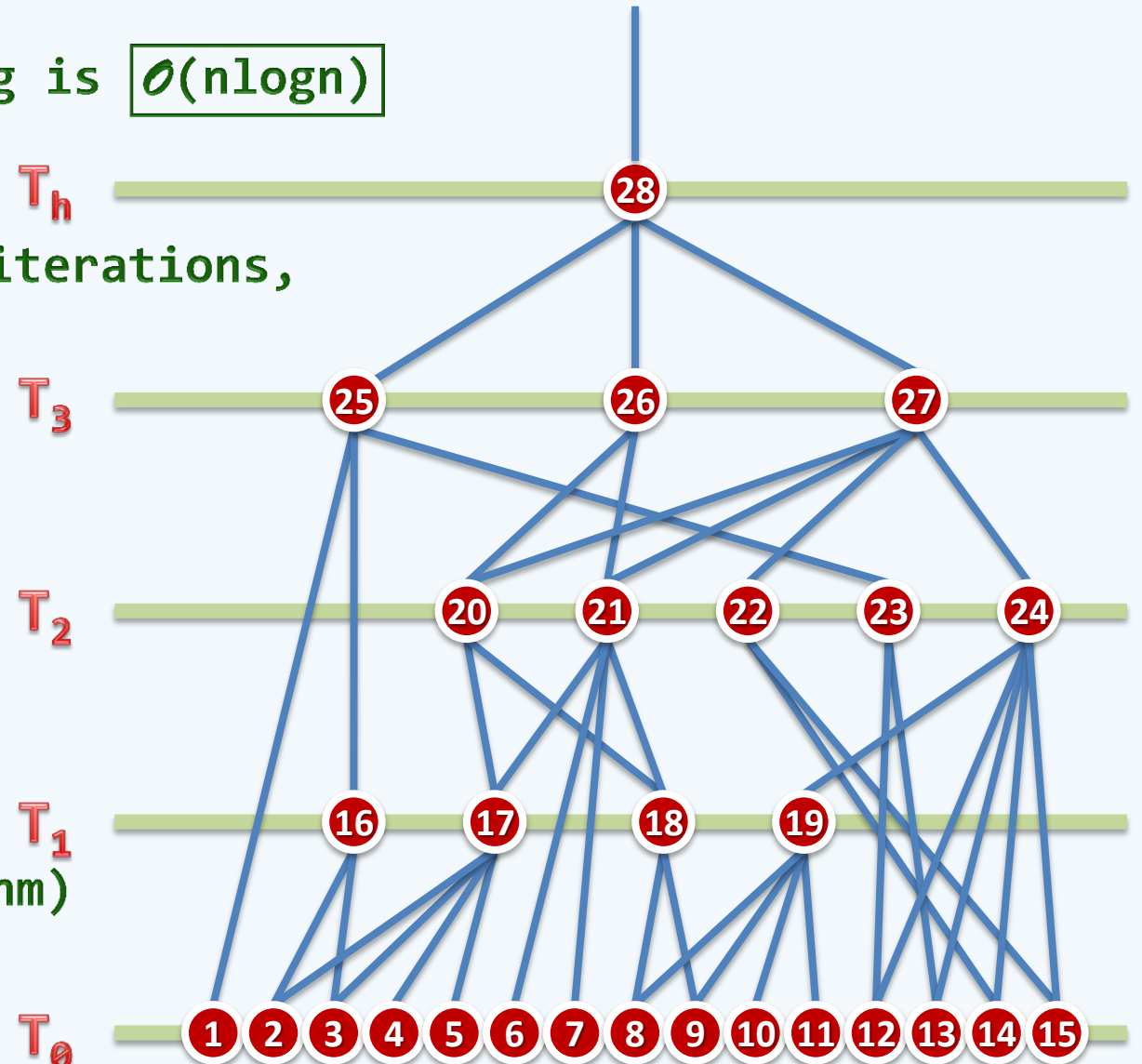
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## Preprocessing

❖ Claim: the time for preprocessing is  $\mathcal{O}(n \log n)$

❖ There are all together  $\mathcal{O}(\log n)$  iterations,  
each of which needs

- $\mathcal{O}(n)$  time  
find an independent set,
- $\mathcal{O}(n)$  time  
to delete vertices,
- $\mathcal{O}(n)$  time to re-triangulate  
(using, say, Chazelle algorithm)
- $\mathcal{O}(8n)$  time  
to test triangle overlapping



# Storage

❖ Claim: Kirkpatrick structure has a **linear** size

❖ The size of a triangulation

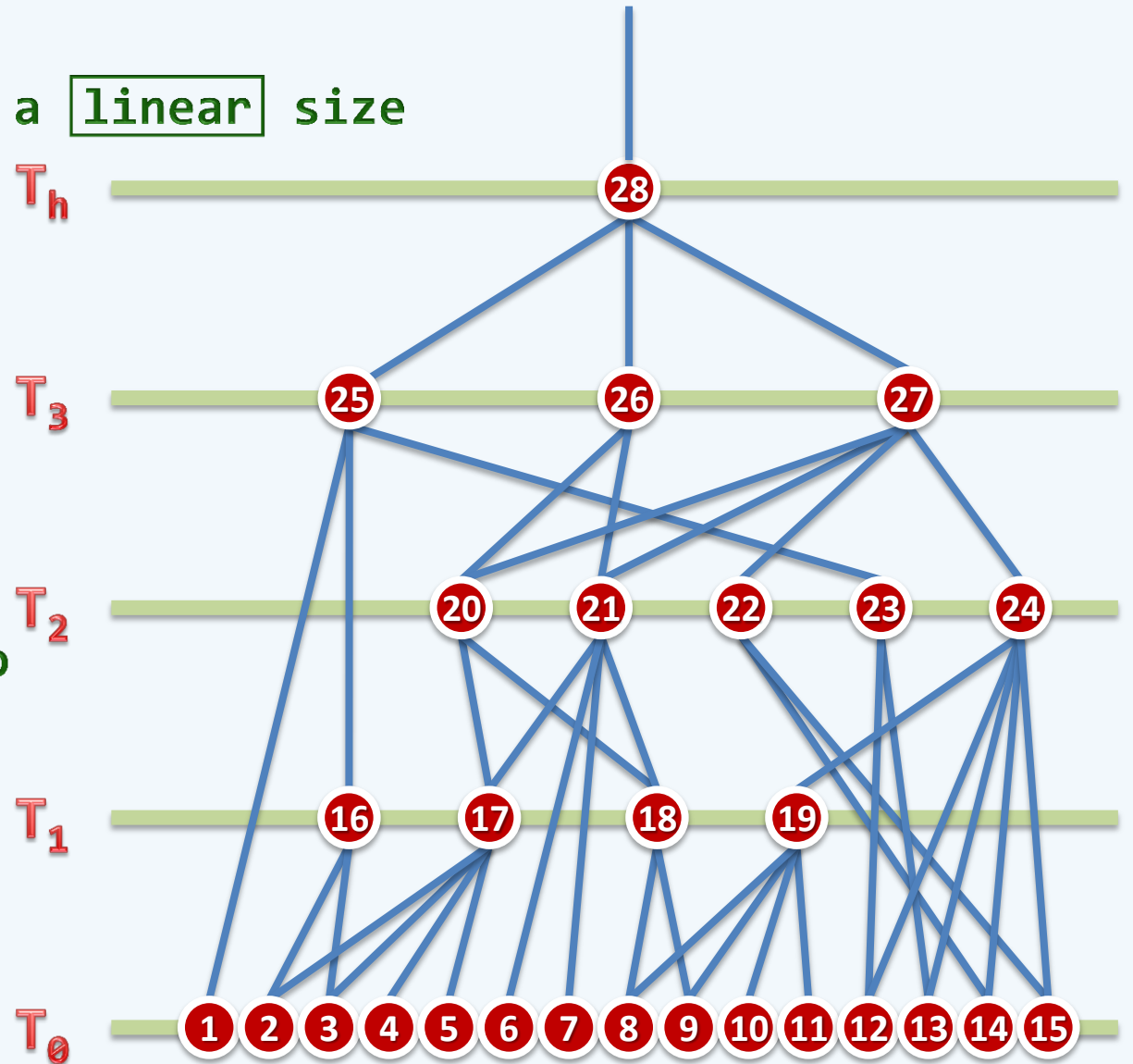
is proportional to

the number of vertices

❖ The total size is proportional to

$$n \times \left[ \left(\frac{17}{18}\right)^0 + \left(\frac{17}{18}\right)^1 + \dots + \left(\frac{17}{18}\right)^{O(\log n)} \right]$$

$$< 18n = \mathcal{O}(n)$$



## Query

- ❖ Claim: each query can be answered in  $\mathcal{O}(\log n)$  time
- ❖ There are  $\mathcal{O}(\log n)$  levels, each of which is a triangulation
- ❖ It takes  $\mathcal{O}(1)$  time to descend one level down
- ❖ In fact, the delay is no more than 8 point-in-triangle tests

