

Arrangement

Planar Arrangement

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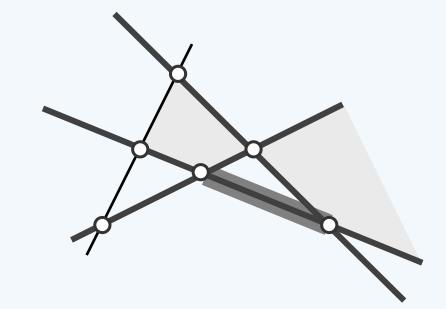
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Arrangement of Lines

❖ Let L be a set of n lines in the plane

 \diamondsuit The arrangement $\mathcal{A}(L)$ is

the plane subdivision induced by L



❖ Complexity: for the arrangement of n lines

#vertices + #edges + #faces = $\Theta(n^2)$

Constructing Planar Arrangements

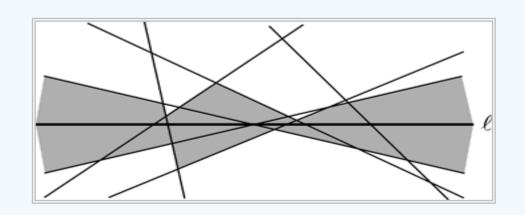
- \diamond Goal: to construct an representation of $\mathcal{A}(L)$ as a |DCEL|
- ❖ Plane sweep needs $|\Theta(n^2\log n)|$ time $//\Theta(n^2)$ events \times $\Theta(\log n)$ time each

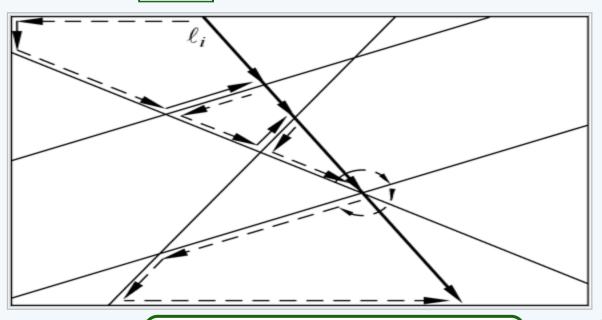
 \Rightarrow RIC algorithm runs in $|O(n^2)|$ time

- //This implies that ...
- \diamond The insertion of each line can be done in |O(n)| time

//how to?

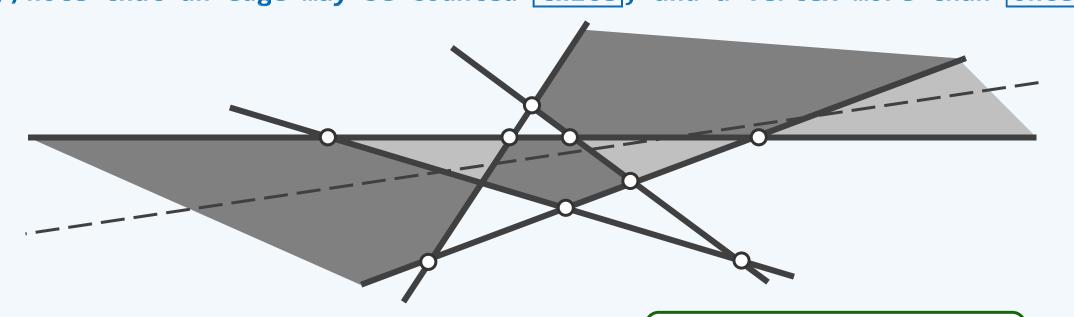
- 1) How many cells are involved?
- 2) What's their total complexity?





Zone Theorem

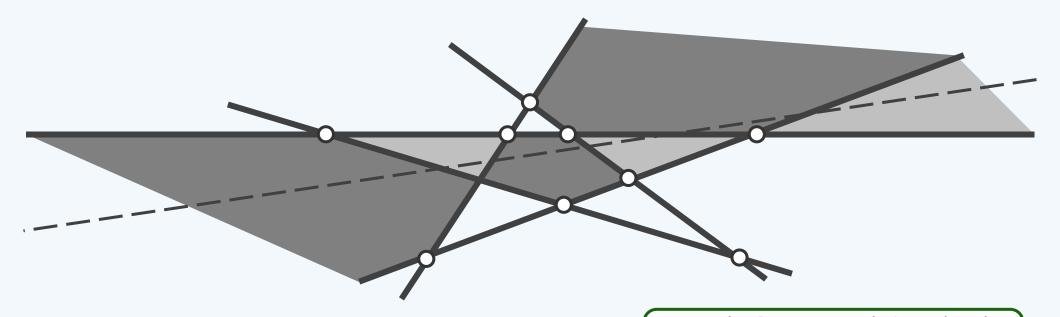
- ❖ The zone of a line 1 in arrangement A(L) is the set of faces of A(L) intersecting 1
- ❖ The complexity of a zone is the total complexity of all its faces //i.e. the total sum of #edges and/or #vertices of these faces //note that an edge may be counted twice, and a vertex more than once



Zone Theorem

❖ [CGAA-8.3]

The complexity of a zone of a line in an arrangement of n lines is O(n)



Complexity of RIC

 \clubsuit The time required to insert l_i is linear in the complexity of l_i 's zone

namely, O(i - 1) by the Zone Theorem

❖ The RIC algorithm

1)constructs the arrangement of n lines in the plane in $o(n^2)$ time and

2) is optimal in terms of the worst case complexity