



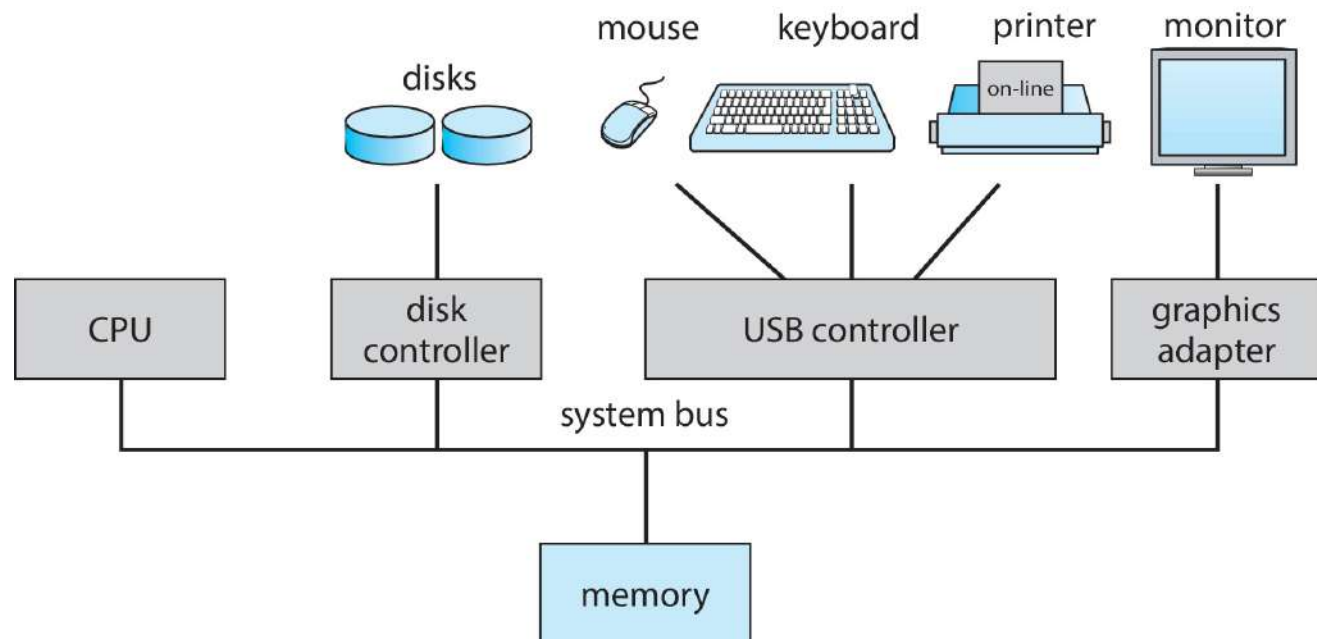
Overview of Computer System Structure





Computer System Organization

- Computer-system operation
 - One or more CPUs, device controllers connect through common **bus** providing access to shared memory
 - Each device controller is in charge of a specific type of device (for example, disk drives, audio devices, or video displays).
 - The CPU and the device controllers can execute in parallel, competing for memory cycles.





Computer-System Operation

- I/O devices and the CPU can execute concurrently
- Each device controller is in charge of a particular device type
- Each device controller has a local buffer
- Each device controller type has an operating system **device driver** to manage it
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller informs CPU that it has finished its operation by causing an **interrupt**





Computer Startup

■ **Booting:**

It is the operation of bringing operating system kernel from the secondary storage and put it in main storage to execute it in CPU.

- There is a program **bootstrap** which is performing this operation when computer is powered up or rebooted
- **Bootstrap program** is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as **firmware**
 - How does Bootstrap work?
 1. Initialize all the aspect of the system, from CPU registers to device controllers to memory contents.
 2. Locate and load the operating system kernel into memory
 3. then the operating system starts executing the first system process, such as “init” in UNIX, and other system processes start.
- When this phase is complete, the system is fully booted, and the system waits for some event to occur.
- The occurrence of an event is usually signaled by an **interrupt** from either the hardware or the software.





Common Functions of Interrupts

- The ***Interrupt*** is either a signal sending by a hardware to CPU, usually during way of the system bus, or a special operation (called a **system call** or a **monitor call**) triggered by a software.
- Interrupt transfers control to the interrupt service routine generally, through the **interrupt vector**, which contains the addresses of all the service routines
- Interrupt architecture must save the address of the interrupted instruction
- A **trap** or **exception** is it is a software-generated interrupt caused either by an error (ex: division by zero or invalid memory access) or by a specific request from a user program (**system call**) that an operating system service be performed.
- Note: An operating system is **interrupt driven**





Operating System Structure

- Multiprogramming (Batch system)
 - Single user cannot always keep CPU and I/O devices busy
 - Multiprogramming organizes jobs (code and data) so CPU always has one to execute
 - A subset of total jobs in system is kept in memory
 - One job selected and run via **job scheduling**
 - When job has to wait (for I/O for example), OS switches to another job
- Multitasking (Timesharing)
 - A logical extension of Batch systems– the CPU switches jobs so frequently that users can interact with each job while it is running, creating **interactive** computing
 - ▶ **Response time** should be < 1 second
 - ▶ Each user has at least one program executing in memory \Rightarrow **process**
 - ▶ If several jobs ready to run at the same time \Rightarrow **CPU scheduling**
 - ▶ If processes don't fit in memory, **swapping** moves them in and out to run





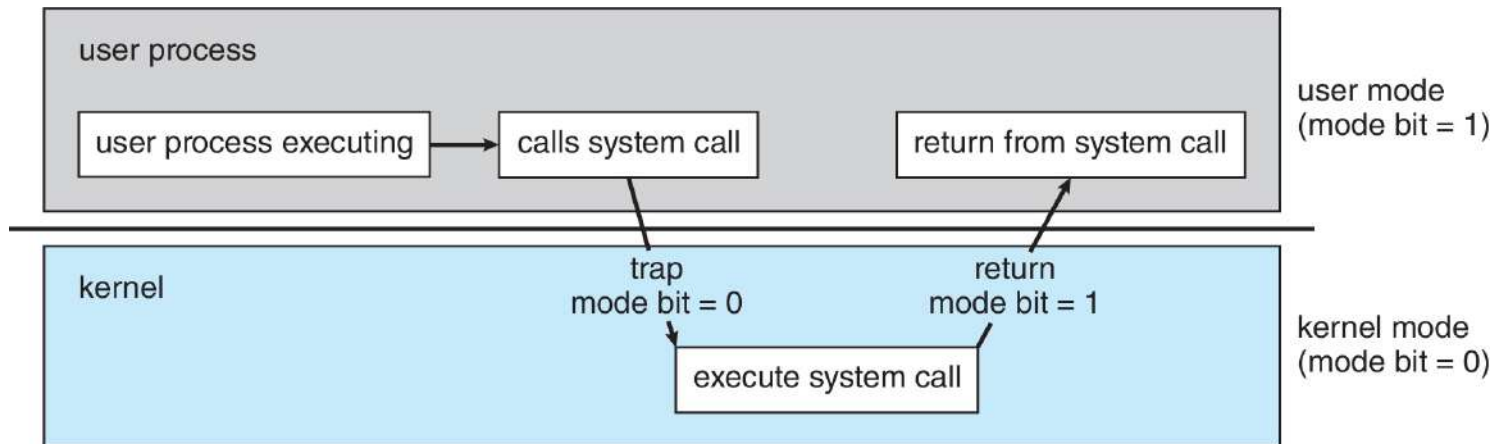
Dual-mode Operation

- **Dual-mode** operation allows OS to protect itself and other system components
 - **User mode** and **kernel mode**
- **Mode bit** provided by hardware
 - Provides ability to distinguish when system is running user code or kernel code.
 - When a user is running \Rightarrow mode bit is “user”
 - When kernel code is executing \Rightarrow mode bit is “kernel”
- How do we guarantee that user does not explicitly set the mode bit to “kernel”?
 - System call changes mode to kernel, return from call resets it to user
- Some instructions designated as **privileged**, only executable in kernel mode





Transition from User to Kernel Mode





Timer

- Timer to prevent infinite loop (or process hogging resources)
 - Timer is set to interrupt the computer after some time period
 - Keep a counter that is decremented by the physical clock
 - Operating system set the counter (privileged instruction)
 - When counter zero generate an interrupt
 - Set up before scheduling process to regain control or terminate program that exceeds allotted time





Computer System Architecture





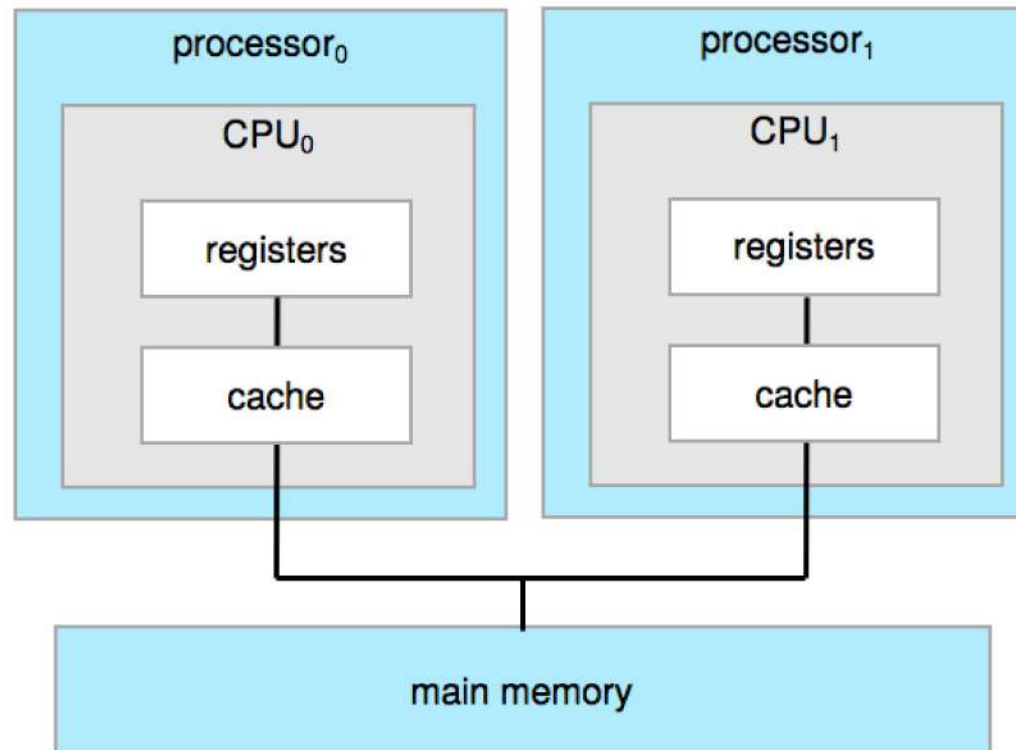
Computer-System Architecture

- Most systems use a single general-purpose processor
 - Most systems have special-purpose processors as well
- **Multiprocessors** systems growing in use and importance
 - Also known as **parallel systems**, **tightly-coupled systems**
 - Advantages include:
 1. **Increased throughput**
 2. **Economy of scale**
 3. **Increased reliability** – graceful degradation or fault tolerance
 - Two types:
 1. **Asymmetric Multiprocessing** – each processor is assigned a specific task.
 2. **Symmetric Multiprocessing** – each processor performs all tasks





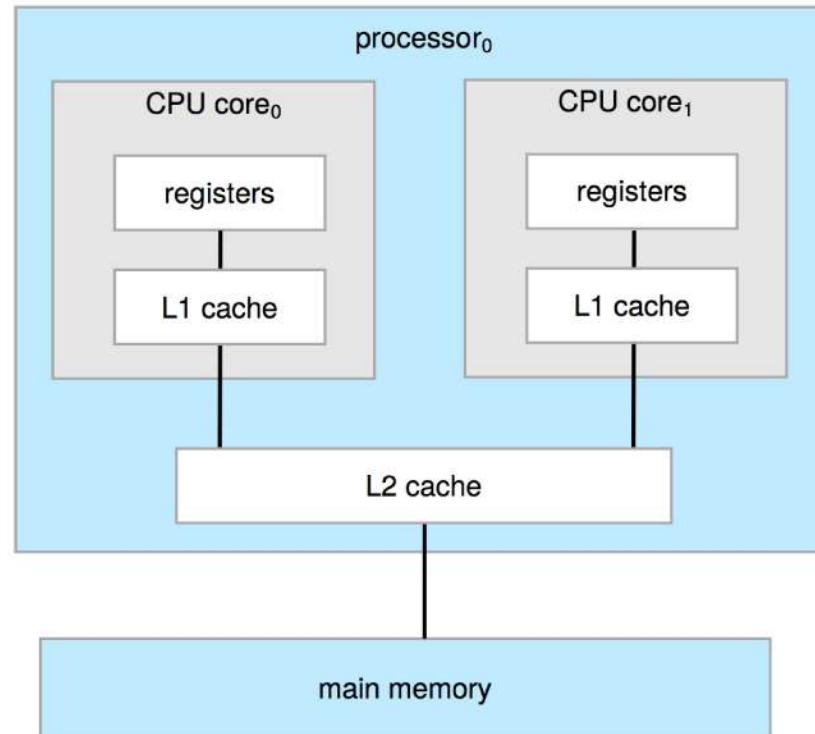
Symmetric Multiprocessing Architecture





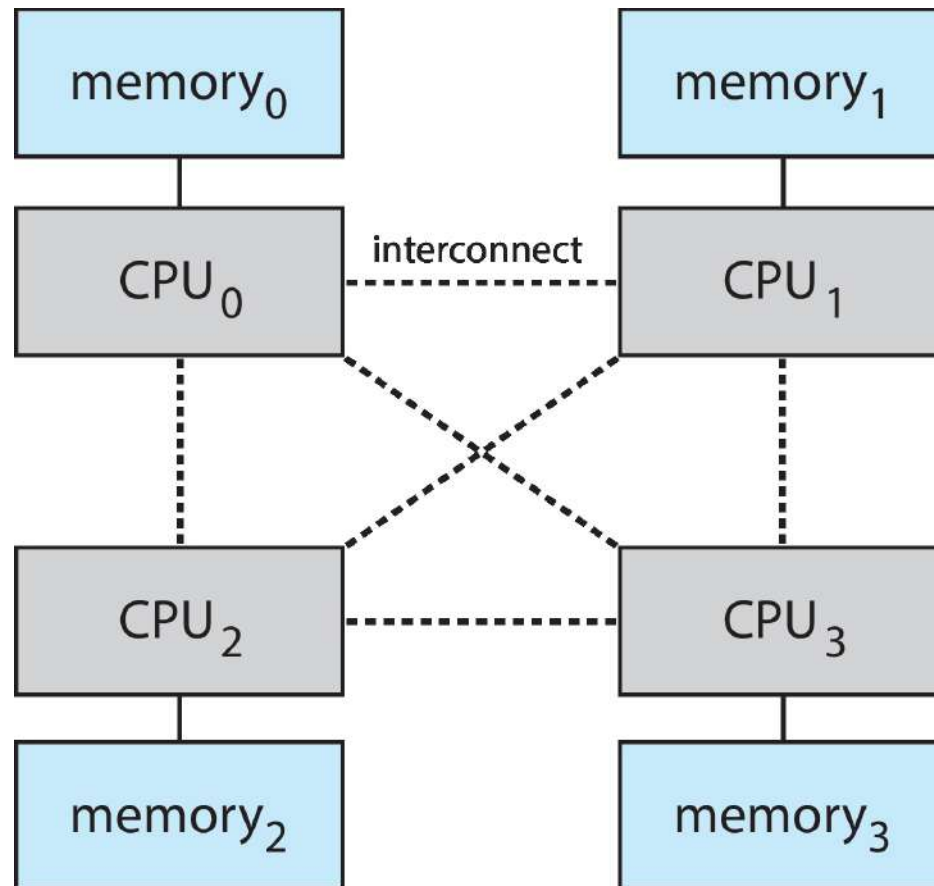
Dual-Core Design

- Multi-chip and **multicore**
- Systems containing all chips
 - Chassis containing multiple separate systems





Non-Uniform Memory Access System





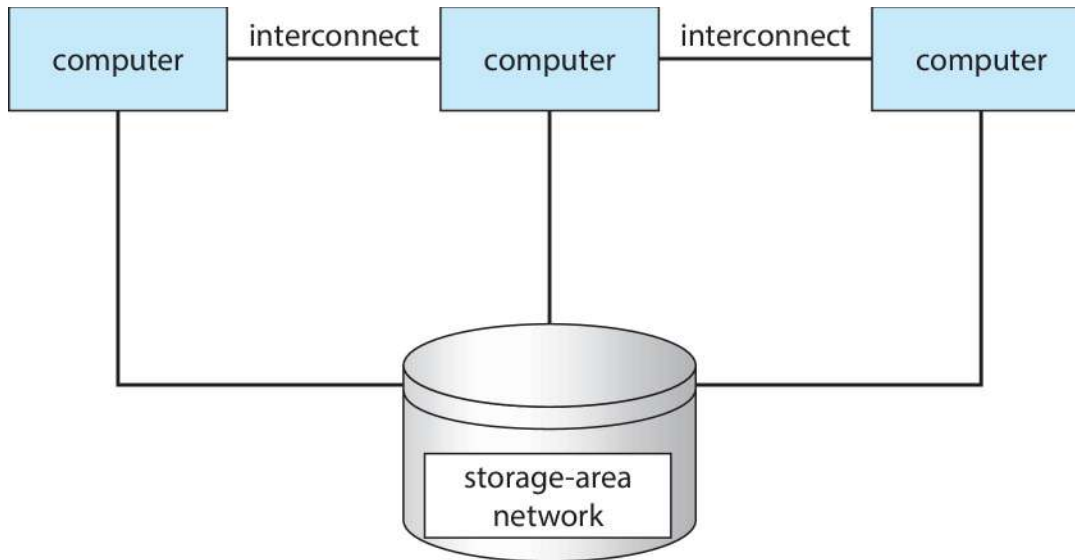
Clustered Systems

- Like multiprocessor systems, but multiple systems working together
 - Usually sharing storage via a **storage-area network (SAN)**
 - Provides a **high-availability** service which survives failures
 - ▶ **Asymmetric clustering** has one machine in hot-standby mode
 - ▶ **Symmetric clustering** has multiple nodes running applications, monitoring each other
 - Some clusters are for **high-performance computing (HPC)**
 - ▶ Applications must be written to use **parallelization**
 - Some have **distributed lock manager (DLM)** to avoid conflicting operations





Clustered Systems





Computer System Environments





Computing Environments

- Traditional
- Mobile
- Distributed System
- Client Server
- Peer-to-Peer
- Cloud computing
- Real-time Embedded





Traditional

- Stand-alone general-purpose machines
- But blurred as most systems interconnect with others (i.e., the Internet)
- **Portals** provide web access to internal systems
- **Network computers (thin clients)** are like Web terminals
- Mobile computers interconnect via **wireless networks**
- Networking becoming ubiquitous – even home systems use **firewalls** to protect home computers from Internet attacks





Mobile

- Handheld smartphones, tablets, etc.
- What is the functional difference between them and a “traditional” laptop?
- Extra feature – more OS features (GPS, gyroscope)
- Allows new types of apps like ***augmented reality***
 - ***augmented reality*** which can be used on several game, is like titling, rotating, and shaking the mobile device
- Use IEEE 802.11 wireless, or cellular data networks for connectivity
- Leaders are **Apple iOS** and **Google Android**





Distributed Systems

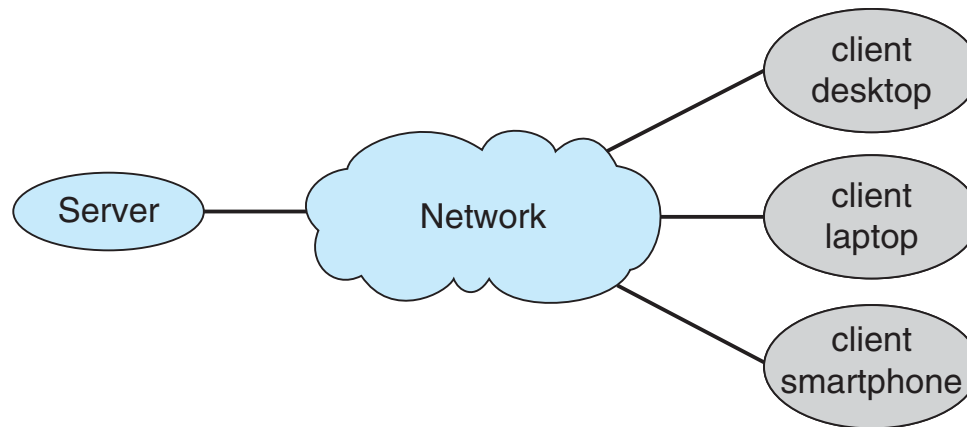
- Collection of separate, possibly heterogeneous, systems networked together
 - **Network** is a communications path, **TCP/IP** most common
 - ▶ **Local Area Network (LAN)**
 - ▶ **Wide Area Network (WAN)**
 - ▶ **Metropolitan Area Network (MAN)**
 - ▶ **Personal Area Network (PAN)**
- **Network Operating System** provides features between systems across network
 - Communication scheme allows systems to exchange messages
 - Illusion of a single system





Client Server

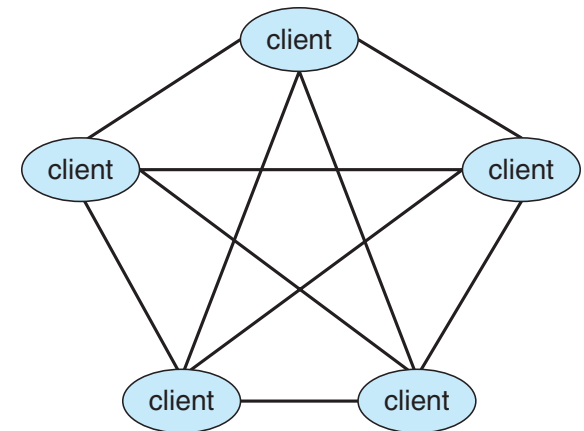
- Client-Server Computing
 - Dumb terminals supplanted by smart PCs
 - Many systems now **servers**, responding to requests generated by **clients**
 - ▶ **Compute-server system** provides an interface to client to request services (i.e., database)
 - ▶ **File-server system** provides interface for clients to store and retrieve files





Peer-to-Peer

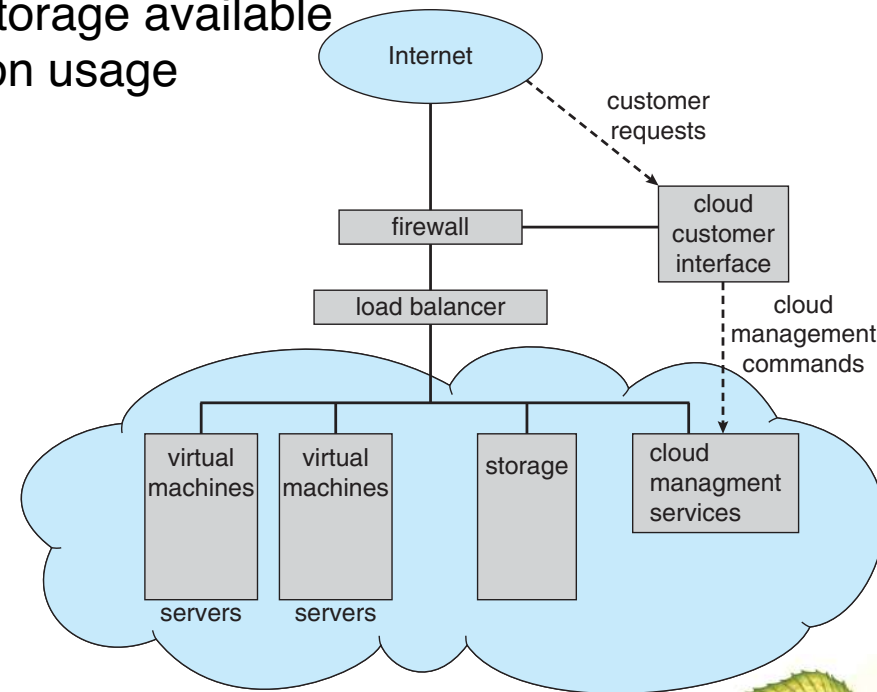
- Another model of distributed system
- P2P does not distinguish clients and servers
 - Instead all nodes are considered peers
 - May each act as client, server or both
 - Node must join P2P network
 - ▶ Registers its service with central lookup service on network, or
 - ▶ Broadcast request for service and respond to requests for service via ***discovery protocol***
 - Examples include Napster and Gnutella, **Voice over IP (VoIP)** such as Skype





Cloud Computing

- Delivers computing, storage, even apps as a service across a network
- Logical extension of virtualization because it uses virtualization as the base for its functionality.
 - Amazon **EC2** has thousands of servers, millions of virtual machines, petabytes of storage available across the Internet, pay based on usage





Real-Time Embedded Systems

- Real-time embedded systems most prevalent form of computers
 - Vary considerable, special purpose, limited purpose OS, **real-time OS**
- Real-time OS has well-defined fixed time constraints
 - Real time system need that the processing must be done within the defined time constraints or the system will fail
 - Processing ***must*** be done within constraint
 - Correct operation only if constraints met





OPERATING SYSTEM SERVICES





Operating System Services

- Operating systems provide an environment for execution of programs and services to programs and users
- One set of operating-system services provides functions that are helpful to the user:
 - **User interface** - Almost all operating systems have a user interface (**UI**).
 - ▶ Varies between **Command-Line (CLI)**, **Graphics User Interface (GUI)**, **touch-screen**, **Batch**
 - **Program execution** - The system must be able to load a program into memory and to run that program, end execution, either normally or abnormally (indicating error)
 - **I/O operations** - A running program may require I/O, which may involve a file or an I/O device
 - **File-system manipulation** - The file system is of particular interest. Programs need to read and write files and directories, create and delete them, search them, list file Information, permission management.





Operating System Services (Cont.)

- One set of operating-system services provides functions that are helpful to the user (Cont.):
 - **Communications** – Processes may exchange information, on the same computer or between computers over a network
 - ▶ Communications may be via shared memory or through message passing (packets moved by the OS)
 - **Error detection** – OS needs to be constantly aware of possible errors
 - ▶ May occur in the CPU and memory hardware, in I/O devices, in user program
 - ▶ For each type of error, OS should take the appropriate action to ensure correct and consistent computing
 - ▶ Debugging facilities can greatly enhance the user's and programmer's abilities to efficiently use the system





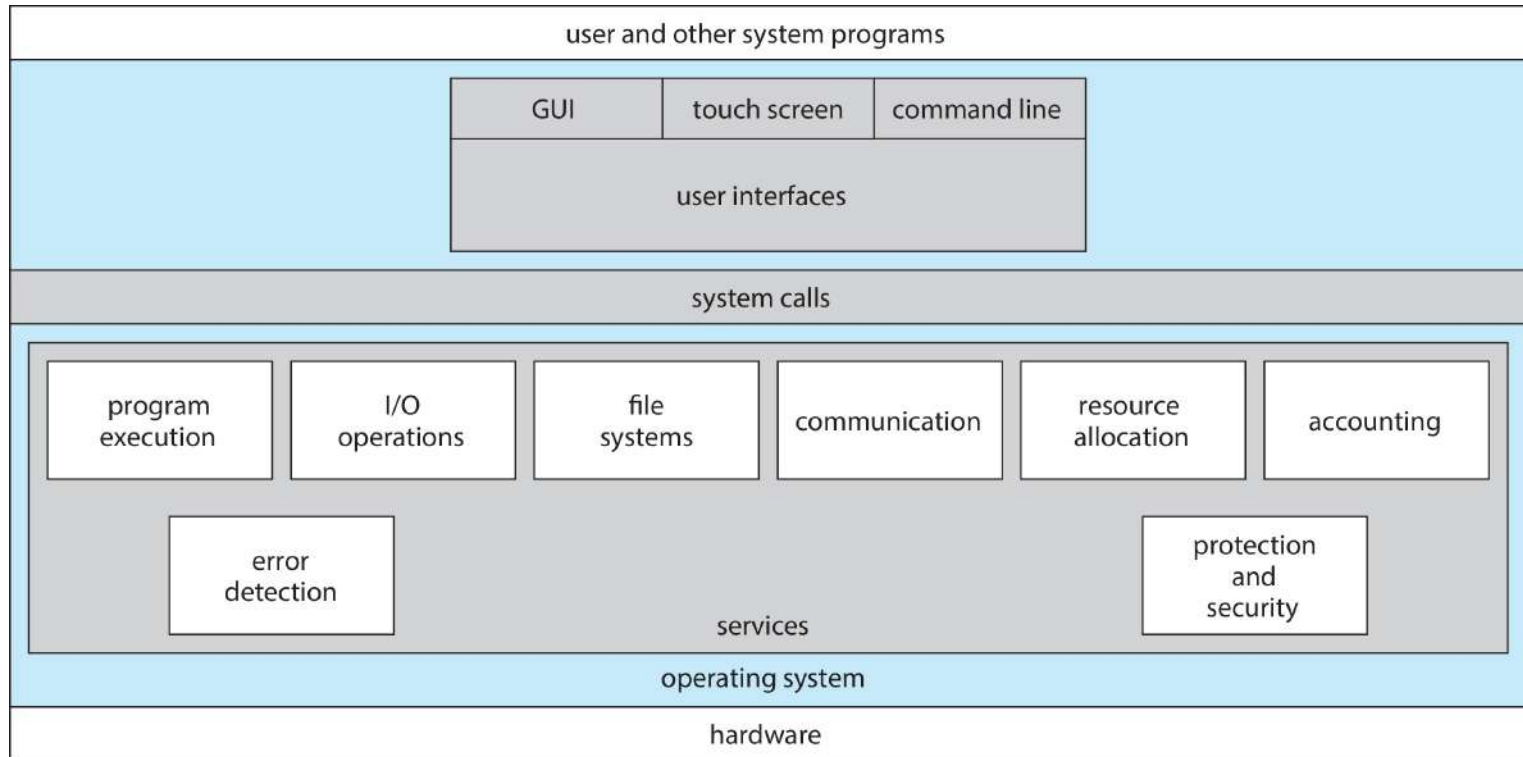
Operating System Services (Cont.)

- Another set of OS functions exists for ensuring the efficient operation of the system itself via resource sharing
 - **Resource allocation** - When multiple users or multiple jobs running concurrently, resources must be allocated to each of them
 - ▶ Many types of resources - CPU cycles, main memory, file storage, I/O devices.
 - **Logging** - To keep track of which users use how much and what kinds of computer resources
 - **Protection and security** - The owners of information stored in a multiuser or networked computer system may want to control use of that information, concurrent processes should not interfere with each other
 - ▶ **Protection** involves ensuring that all access to system resources is controlled
 - ▶ **Security** of the system from outsiders requires user authentication, extends to defending external I/O devices from invalid access attempts





A View of Operating System Services





Command Line interpreter

- CLI allows direct command entry
- Sometimes implemented in kernel, sometimes by systems program
- Sometimes multiple flavors implemented – **shells**
- Primarily fetches a command from user and executes it
- Sometimes commands built-in, sometimes just names of programs
 - If the latter, adding new features doesn't require shell modification





Bourne Shell Command Interpreter

```
1. root@r6181-d5-us01:~ (ssh)
root@r6181-d5-u... %1 ssh %2 root@r6181-d5-us01... %3
Last login: Thu Jul 14 08:47:01 on ttys002
iMacPro:~ pbg$ ssh root@r6181-d5-us01
root@r6181-d5-us01's password:
Last login: Thu Jul 14 06:01:11 2016 from 172.16.16.162
[root@r6181-d5-us01 ~]# uptime
 06:57:48 up 16 days, 10:52,  3 users,  load average: 129.52, 80.33, 56.55
[root@r6181-d5-us01 ~]# df -kh
Filesystem                Size      Used Avail Use% Mounted on
/dev/mapper/vg_ks-lv_root    50G       19G   28G  41% /
tmpfs                      127G      520K   127G   1% /dev/shm
/dev/sda1                   477M       71M   381M  16% /boot
/dev/dssd0000               1.0T      480G   545G  47% /dssd_xfs
tcp://192.168.150.1:3334/orangeifs
                          12T       5.7T   6.4T  47% /mnt/orangeifs
/dev/gpfs-test              23T       1.1T   22T   5% /mnt/gpfs
[root@r6181-d5-us01 ~]#
[root@r6181-d5-us01 ~]# ps aux | sort -nrk 3,3 | head -n 5
root      97653 11.2  6.6 42665344 17520636 ?    S<Ll  Jul13 166:23 /usr/lpp/mmfs/bin/mmfsd
root      69849  6.6  0.0      0      0 ?        S    Jul12 181:54 [vpthread-1-1]
root      69850  6.4  0.0      0      0 ?        S    Jul12 177:42 [vpthread-1-2]
root       3829  3.0  0.0      0      0 ?        S    Jun27 730:04 [rp_thread 7:0]
root       3826  3.0  0.0      0      0 ?        S    Jun27 728:08 [rp_thread 6:0]
[root@r6181-d5-us01 ~]# ls -l /usr/lpp/mmfs/bin/mmfsd
-r-x----- 1 root root 20667161 Jun  3  2015 /usr/lpp/mmfs/bin/mmfsd
[root@r6181-d5-us01 ~]#
```





User Operating System Interface - GUI

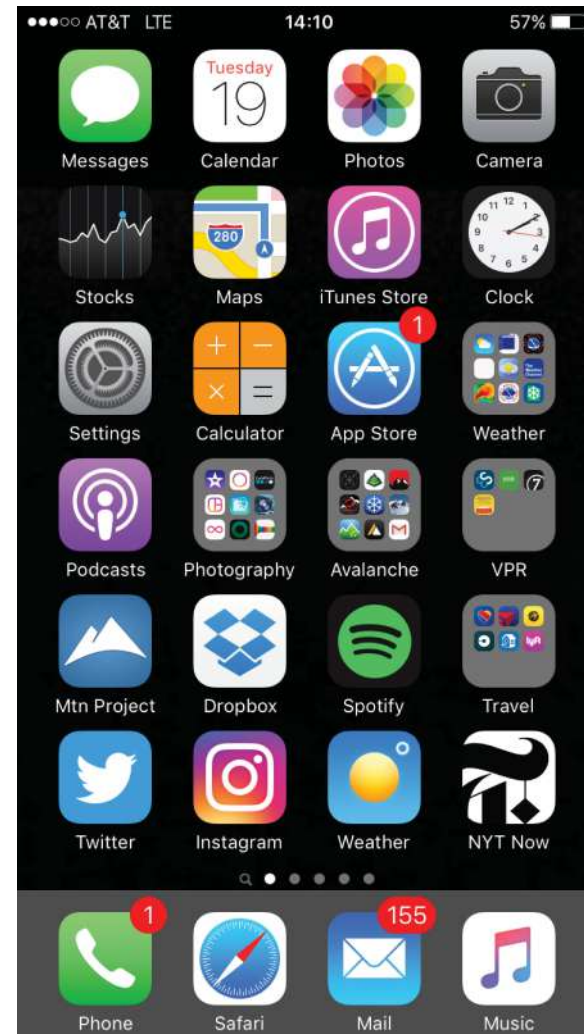
- User-friendly **desktop** metaphor interface
 - Usually mouse, keyboard, and monitor
 - **Icons** represent files, programs, actions, etc
 - Various mouse buttons over objects in the interface cause various actions (provide information, options, execute function, open directory (known as a **folder**))
 - Invented at Xerox PARC
- Many systems now include both CLI and GUI interfaces
 - Microsoft Windows is GUI with CLI “command” shell
 - Apple Mac OS X is “Aqua” GUI interface with UNIX kernel underneath and shells available
 - Unix and Linux have CLI with optional GUI interfaces (CDE, KDE, GNOME)





Touchscreen Interfaces

- Touchscreen devices require new interfaces
 - Mouse not possible or not desired
 - Actions and selection based on gestures
 - Virtual keyboard for text entry
- Voice commands





The Mac OS X GUI





System Calls

- Programming interface to the services provided by the OS
- Typically written in a high-level language (C or C++)
- Mostly accessed by programs via a high-level **Application Programming Interface (API)** rather than direct system call use
- Three most common APIs are Win32 API for Windows, POSIX API for POSIX-based systems (including virtually all versions of UNIX, Linux, and Mac OS X), and Java API for the Java virtual machine (JVM)

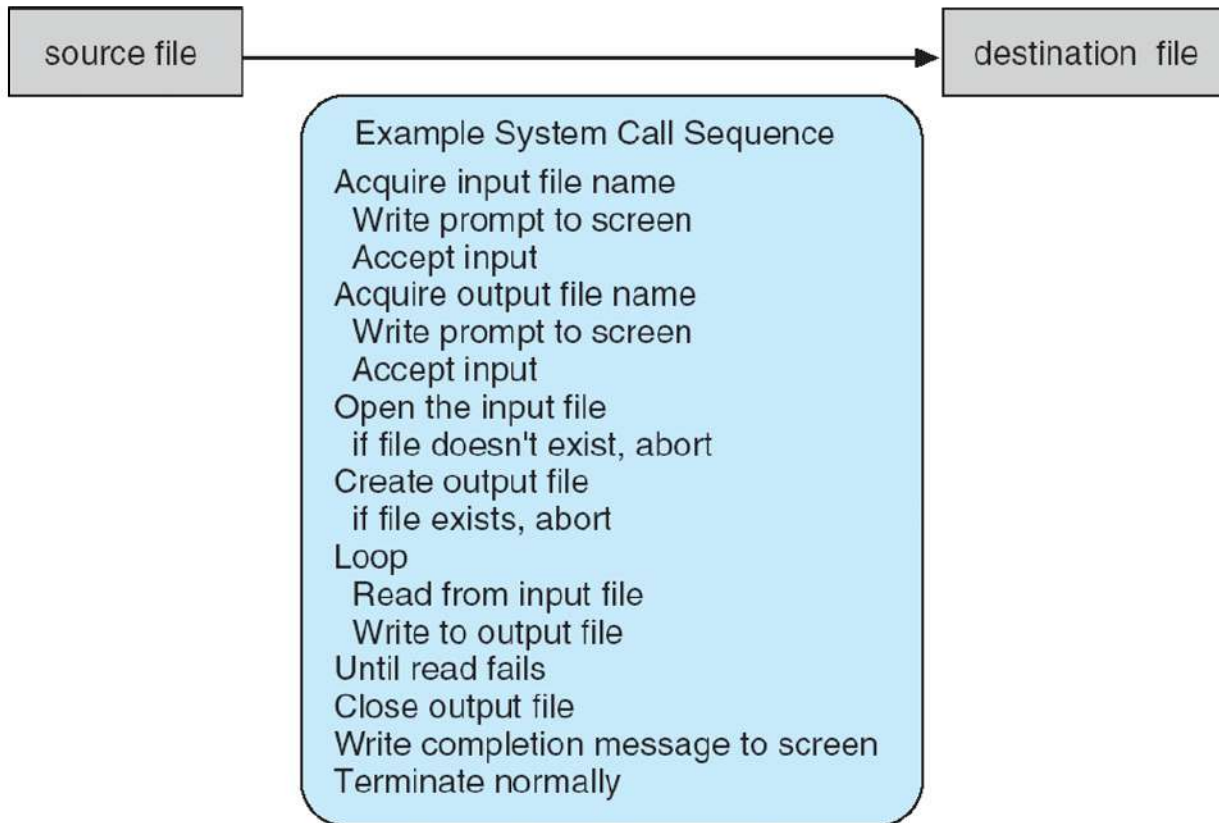
Note that the system-call names used throughout this text are generic





Example of System Calls

- System call sequence to copy the contents of one file to another file





Examples of Windows and Unix System Calls

EXAMPLES OF WINDOWS AND UNIX SYSTEM CALLS

The following illustrates various equivalent system calls for Windows and UNIX operating systems.

	Windows	Unix
Process control	CreateProcess() ExitProcess() WaitForSingleObject()	fork() exit() wait()
File management	CreateFile() ReadFile() WriteFile() CloseHandle()	open() read() write() close()
Device management	SetConsoleMode() ReadConsole() WriteConsole()	ioctl() read() write()
Information maintenance	GetCurrentProcessID() SetTimer() Sleep()	getpid() alarm() sleep()
Communications	CreatePipe() CreateFileMapping() MapViewOfFile()	pipe() shm_open() mmap()
Protection	SetFileSecurity() InitializeSecurityDescriptor() SetSecurityDescriptorGroup()	chmod() umask() chown()





EXAMPLES OF OPERATING SYSTEM STRUCTURE





Monolithic Structure – Original UNIX

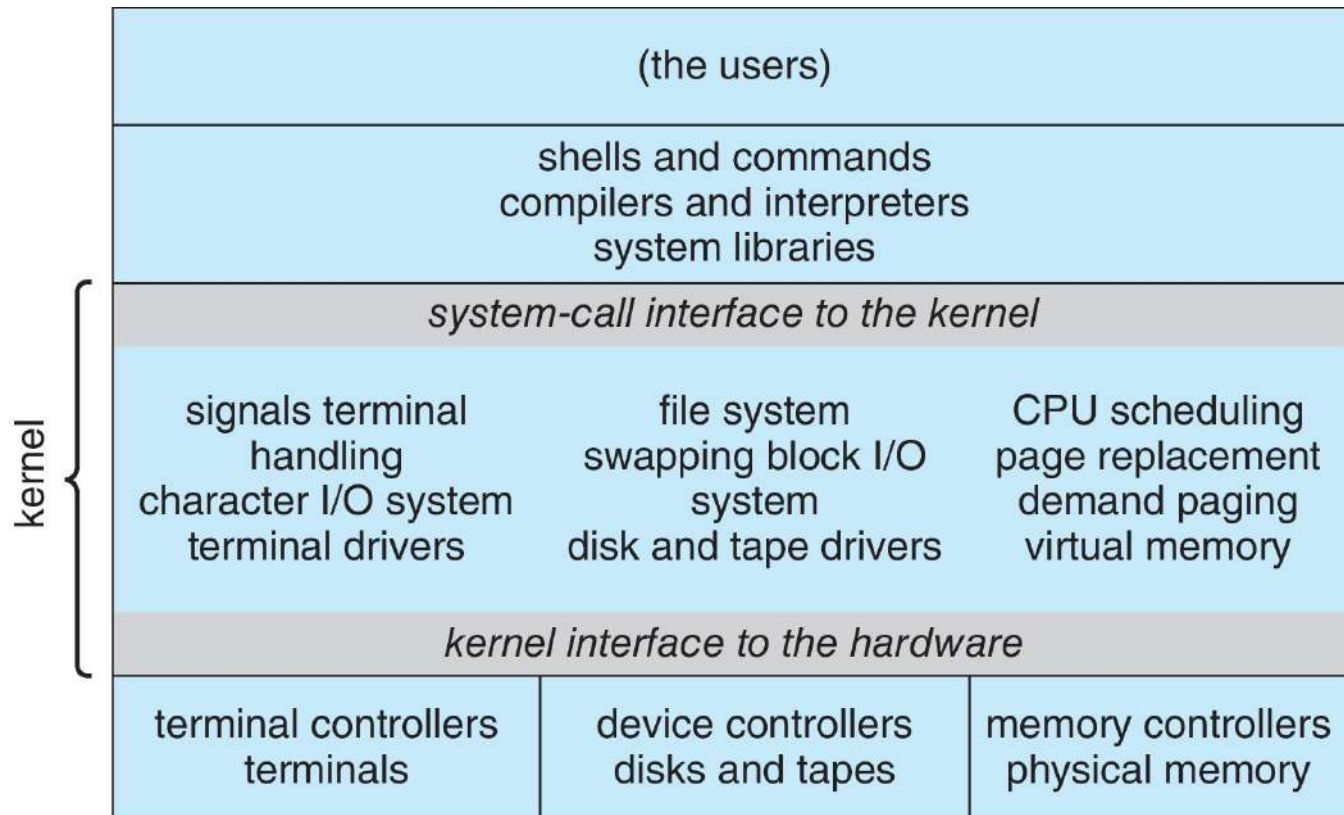
- UNIX – limited by hardware functionality, the original UNIX operating system had limited structuring.
- The UNIX OS consists of two separable parts
 - Systems programs
 - The kernel
 - ▶ Consists of everything below the system-call interface and above the physical hardware
 - ▶ Provides the file system, CPU scheduling, memory management, and other operating-system functions; a large number of functions for one level





Traditional UNIX System Structure

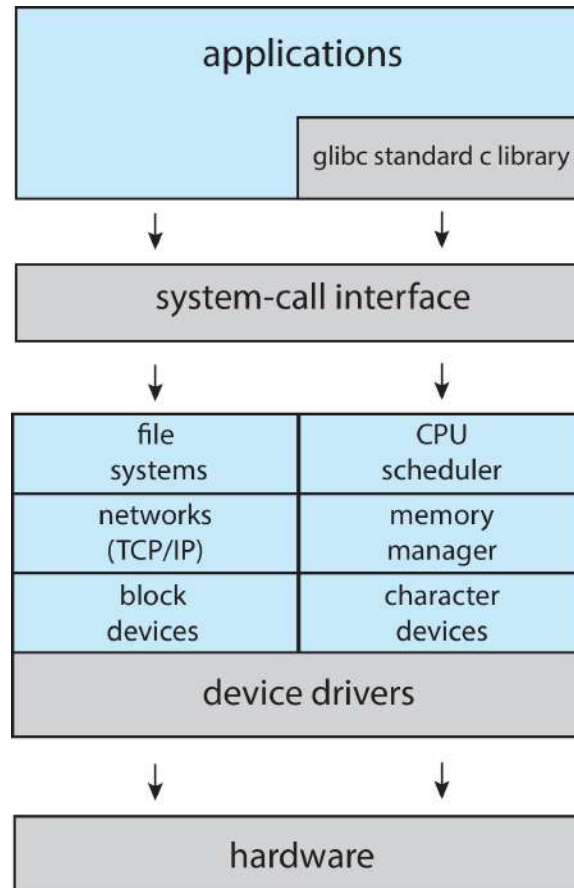
Beyond simple but not fully layered





Linux System Structure

Monolithic plus modular design





iOS

- Apple mobile OS for *iPhone*, *iPad*
 - Structured on Mac OS X, added functionality
 - Does not run OS X applications natively
 - ▶ Also runs on different CPU architecture (ARM vs. Intel)
 - **Cocoa Touch** Objective-C API for developing apps
 - **Media services** layer for graphics, audio, video
 - **Core services** provides cloud computing, databases
 - Core operating system, based on Mac OS X kernel

Cocoa Touch

Media Services

Core Services

Core OS





Android

- Developed by Open Handset Alliance (mostly Google)
 - Open Source
- Similar stack to iOS
- Based on Linux kernel but modified
 - Provides process, memory, device-driver management
 - Adds power management
- Runtime environment includes core set of libraries and Dalvik virtual machine
 - Apps developed in Java plus Android API
 - ▶ Java class files compiled to Java bytecode then translated to executable thnn runs in Dalvik VM
- Libraries include frameworks for web browser (webkit), database (SQLite), multimedia, smaller libc





Android Architecture

