

COMP 523: Language-based security

Assignment 1 (100 points total)

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Due: Thursday, 20 Sept, 2:35pm

Exercise 1 (20 pts): In the lecture and in Pierce’s book, we define the following operational semantics for a small language of terms. For convenience, we repeat the evaluation rules here.

$$\begin{array}{c}
 \frac{}{\text{pred } (\text{succ } nv) \longrightarrow nv} \text{ E-PRED-SUCC} \quad \frac{}{\text{pred } z \longrightarrow z} \text{ E-PRED-ZERO} \\
 \\
 \frac{t \longrightarrow t'}{\text{succ } t \longrightarrow \text{succ } t'} \text{ E-SUCC} \quad \frac{t \longrightarrow t'}{\text{pred } t \longrightarrow \text{pred } t'} \text{ E-PRED} \\
 \\
 \frac{}{\text{if true then } t_1 \text{ else } t_2 \longrightarrow t_1} \text{ E-IF-TRUE} \quad \frac{}{\text{if false then } t_1 \text{ else } t_2 \longrightarrow t_2} \text{ E-IF-FALSE} \\
 \\
 \frac{t \longrightarrow t'}{\text{if } t \text{ then } t_1 \text{ else } t_2 \longrightarrow \text{if } t' \text{ then } t_1 \text{ else } t_2} \text{ E-IF} \\
 \\
 \frac{}{\text{iszero } z \longrightarrow \text{true}} \text{ E-ISZERO-ZERO} \quad \frac{}{\text{iszero } (\text{succ } nv) \longrightarrow \text{false}} \text{ E-ISZERO-SUCC} \\
 \\
 \frac{t \longrightarrow t'}{\text{iszero } t \longrightarrow \text{iszero } t'} \text{ E-ISZERO}
 \end{array}$$

A friend of yours suggests to replace the evaluation rule E-PRED-SUCC with the rule

$$\frac{}{\text{succ } (\text{pred } t) \longrightarrow t}$$

and the rule E-ISZERO-SUCC with

$$\frac{}{\text{iszero } (\text{succ } t) \longrightarrow \text{false}} \text{ E-ISZERO-SUCC}$$

Is this a good idea? What would you say to her or him? Which basic property discussed in Ch 3 breaks down? – If you think the above rule is good, verify that all the theorems in Ch 3 still hold. If you think the rule is bad, then give a counterexample and explain which theorem does not hold.

Exercise 2 (20pts): Show that for the small-step semantics, we have that all values evaluate to themselves.

If v is a value and $v \longrightarrow^* v'$ then $v = v'$.

Use the following rules for the multi-step relation $t \longrightarrow^* t'$.

$$\frac{}{t \longrightarrow^* t} \text{ refl} \quad \frac{t \longrightarrow^* s \quad s \longrightarrow^* t'}{t \longrightarrow^* t'} \text{ trans} \quad \frac{t \longrightarrow t'}{t \longrightarrow^* t'} \text{ single}$$

Consider only the cases for numerical values.

Exercise 3 (20 pts) An alternative formulation of the multi-step rules

$$\frac{}{t \Longrightarrow^* t} \text{ m-refl} \quad \frac{t \longrightarrow s \quad s \Longrightarrow^* t'}{t \Longrightarrow^* t'} \text{ m-step}$$

Show that the two formulations are equivalent, i.e. $t \longrightarrow^* s$ iff $t \Longrightarrow^* s$. State and prove any additional lemmas you might need.

Exercise 4 (40pts): Extend the language for booleans and arithmetic expressions we have seen in class (see also Ch 3, CH 8 in Pierce) with an expression $\text{leq } t \ t'$ which allows us to check whether t is less than or equal to t' .

1. (5 pts) Define small-step evaluation rules for $\text{leq } t \ t'$.
2. (16 pts) Prove that the rules are deterministic. Justify which cases are impossible and why.
3. (4 pts) Define a typing rule for $\text{leq } t \ t'$.
4. (15 pts) Prove that type preservation holds for this extension.

Exercise 5 (optional) : An alternative style to the small-step semantics seen in class is the *big-step* semantics. The judgment $e \Downarrow v$ describes the complete evaluation of the expression e to some final value v . We concentrate here on the fragment for natural numbers. The rules for big-step evaluation for the small fragment consisting of z , $\text{succ } e$, $\text{pred } e$, and $\text{iszero } e$ are given below.

$$\begin{array}{c} \frac{}{z \Downarrow z} \text{ B-Z} \quad \frac{e \Downarrow v}{\text{succ } e \Downarrow \text{succ } v} \text{ B-SUCC} \\[10pt] \frac{e \Downarrow z}{\text{pred } e \Downarrow z} \text{ B-PRED-ZERO} \quad \frac{e \Downarrow \text{succ } v}{\text{pred } e \Downarrow v} \text{ B-PRED-SUCC} \\[10pt] \frac{e \Downarrow z}{\text{iszero } e \Downarrow \text{true}} \text{ B-ISZERO} \quad \frac{e \Downarrow \text{succ } v}{\text{iszero } e \Downarrow \text{false}} \text{ B-ISSUCC} \end{array}$$

Prove that the small-step and big-step semantics coincide, i.e. $e \Downarrow v$ iff $e \longrightarrow^* v$. In your proofs, you should show the case for handling the predecessor in detail; in particular, state and prove all necessary lemmas. You can sketch the remaining cases for successor and iszero-expression.

Hint: Read Exercise 3.5.17 in TAPL page 42 and the corresponding solution page 498.