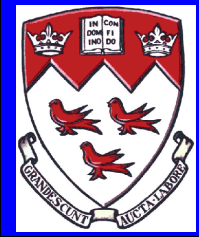




# COMP 273

## Overview of Assembler Programming

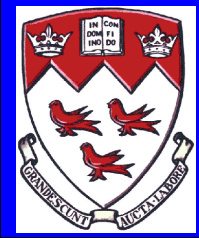
Prof. Joseph Vybihal





# Announcements

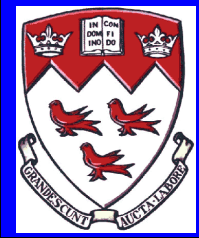
- Project has been posted
- Ass#3 out by end of next week
- MIPS/SPIM/MARS Tutorials
  - TBD





# COMP 273

## Introduction to Computer Systems



# Project



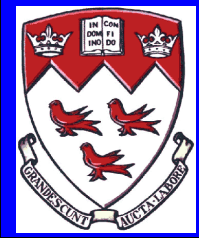
# Development Order

- Now – Teams already picked
- Nov 10 – Select the circuits you will reuse from your assignments and divide the work up between team members, construct the high level architecture (somewhat specified in project)
- Nov 24 – Working basic CPU
- Dec 2 – Working CPU



# Suggestions

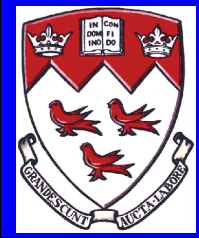
- Divide the CPU into modules and make each one work on its own as a black box.
- Merge working black boxes one at a time.
- Have an overall architectural plan for the CPU early, then
  - Built it in increments
  - Test each increment
- Only think about the bonus if you've got extra time.



# Try This Out at Home

- Download one of the MIPS interpreters and get used to the interface. Load one of the sample programs. Study the code and then get it to run.
  - Load one of the sample programs. Study the code and then get it to run.
  - Write your own code to display the contents of an integer array

# PCSPIM or MARS



# Part 1

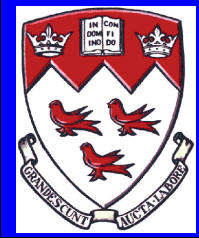
## Assembler Introduction



# Things to know about Assembler Programming

- Just like programming in any language...
- Basic assembler functionality:
  - Math and Logic
  - Data movement
  - Pointers
  - Branching
  - Goto and subroutines
- Functionality not in assembler:
  - No functions (but can be simulated ~subroutines)
  - No objects (but can be simulated)
  - No arrays or data structures (but can be simulated)
  - No variables or constants (in the normal way)
  - Minimal distinction between code and data





# Things to know about Assembler Programming

- Advanced assembler functionality:
  - Access to internal OS routines
  - Access to the control ROMS of all peripherals
- Take care of complexity explosion
  - Keep it simple
  - Modularize with subroutines
  - Use OS functions when useful
- Know how to use the registers and stack
  - Use them over RAM

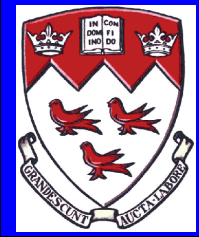
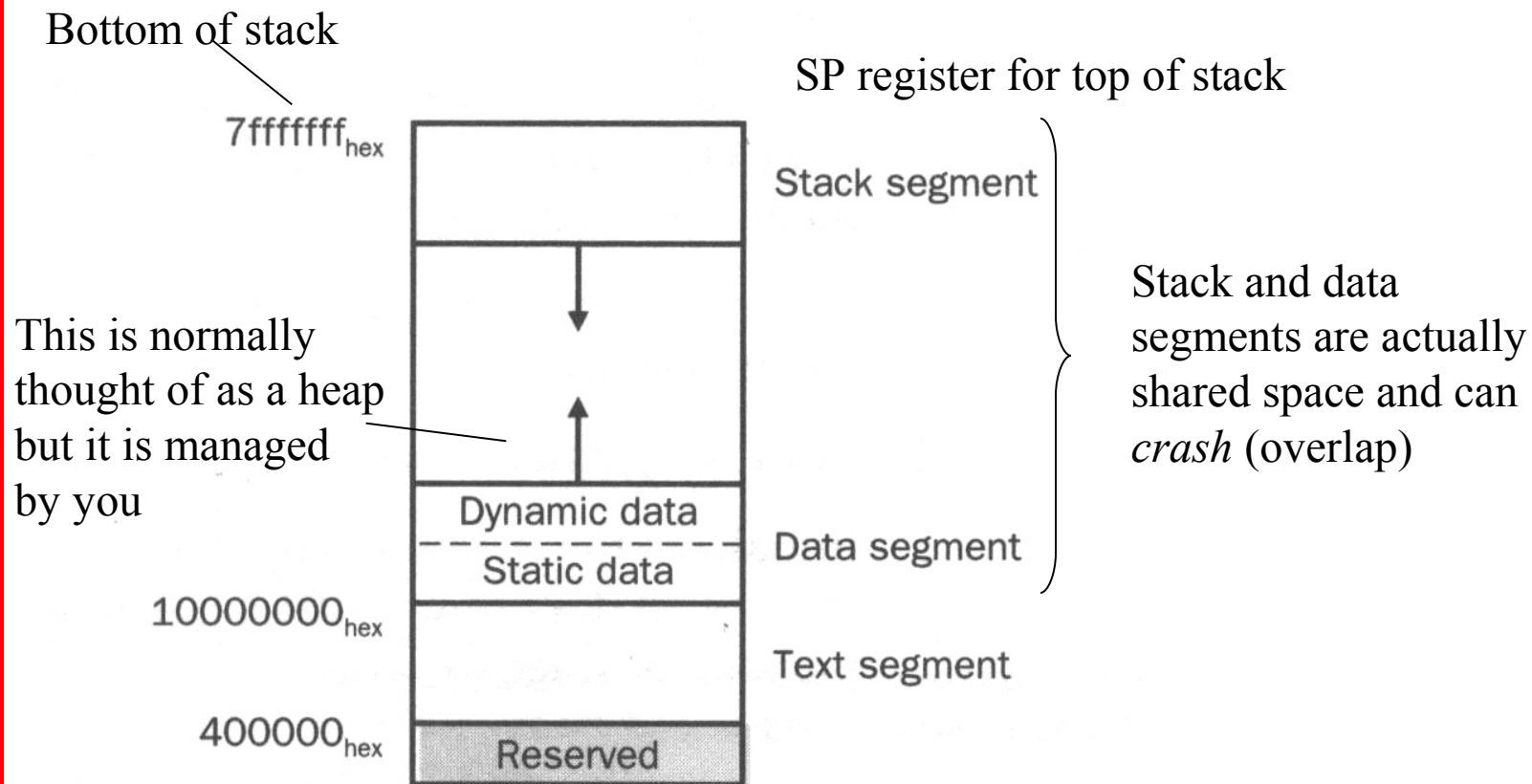


# MIPS Virtual Memory

- $2^{30}$  memory words
  - Memory[0], memory[4], ..., memory[4,294,967,292]
- Accessed only by data transfer instructions
- MIPS uses byte addresses
- Words are 4 bytes
- Memory holds:
  - Data structures
  - Arrays
  - Spilled (saved) registers
  - Assembler instructions
  - Variables and constants



# Virtual Memory Usage



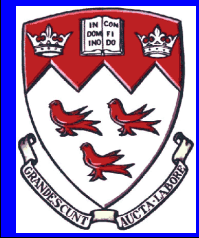


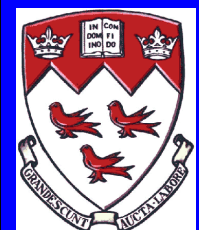
# MIPS Registers

Register	Name	Used For
0	zero	Always returns 0
1	at	Reserved for use by assembler
2-3	v0, v1	Value returned by subroutine
4-7	a0-a3	First few parameters for subroutine
8-15	t0-t7	Temporary: can use without saving
24, 25	t8, t9	Temporary: can use without saving
16-23	s0-s7	If used, <i>must</i> save on stack (or other)
26, 27	k0, k1	Used by interrupt / trap handler
28	gp	A global pointer (extern/static vars...)
29	sp	Stack pointer
30	s8/fp	Frame pointer
31	ra	Subroutine return address

Also:

- Hi
- Lo
- Results of mult & div





Register name	Number	Usage
\$zero	0	constant 0
\$at	1	reserved for assembler
\$v0	2	expression evaluation and results of a function
\$v1	3	expression evaluation and results of a function
\$a0	4	argument 1
\$a1	5	argument 2
\$a2	6	argument 3
\$a3	7	argument 4
\$t0	8	temporary (not preserved across call)
\$t1	9	temporary (not preserved across call)
\$t2	10	temporary (not preserved across call)
\$t3	11	temporary (not preserved across call)
\$t4	12	temporary (not preserved across call)
\$t5	13	temporary (not preserved across call)
\$t6	14	temporary (not preserved across call)
\$t7	15	temporary (not preserved across call)
\$s0	16	saved temporary (preserved across call)
\$s1	17	saved temporary (preserved across call)
\$s2	18	saved temporary (preserved across call)
\$s3	19	saved temporary (preserved across call)
\$s4	20	saved temporary (preserved across call)
\$s5	21	saved temporary (preserved across call)
\$s6	22	saved temporary (preserved across call)
\$s7	23	saved temporary (preserved across call)
\$t8	24	temporary (not preserved across call)
\$t9	25	temporary (not preserved across call)
\$k0	26	reserved for OS kernel
\$k1	27	reserved for OS kernel
\$gp	28	pointer to global area
\$sp	29	stack pointer
\$fp	30	frame pointer
\$ra	31	return address (used by function call)

# MIPS Registers

Return

Parameters

Temporary

Saved

Pointers





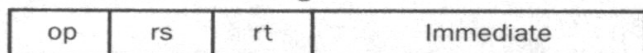
# Addressing Modes

- Register Addressing
  - Operand is a register
  - E.g. `add $s1, $s2, $s3`
- Base or Displacement Addressing
  - Operand is a memory location
  - Register + offset  $\leftarrow$  a constant
  - E.g. `lw $s1, 100($s2)`
- Immediate Addressing
  - Operand is a constant (no addressing)
  - 16-bit constant
  - E.g. `Addi $s1, $s2, 100`
- PC-Relative Addressing
  - Memory location = PC + offset  $\leftarrow$  a constant
  - E.g. `j 2500` or `j label`
- Pseudo-Direct Addressing
  - Memory location = PC (top 6 bits) concat with 26-bit offset
  - Assumes 32-bit addressing
  - E.g. `jal 2500` or `jal label`

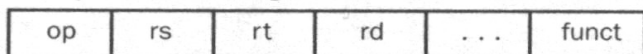
Discuss...



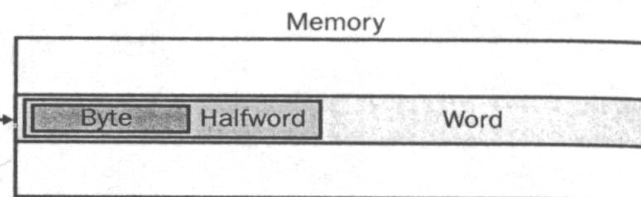
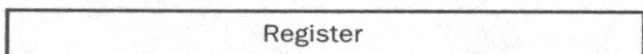
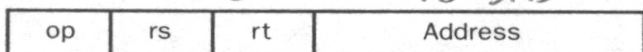
### 1. Immediate addressing *16-bits*



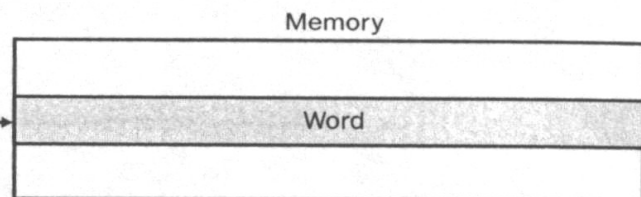
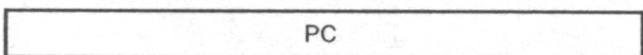
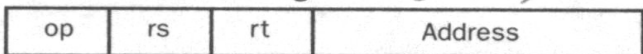
### 2. Register addressing



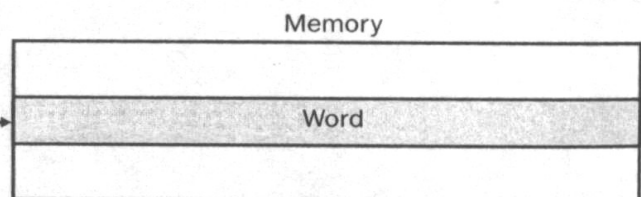
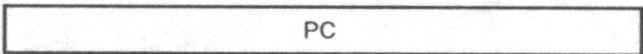
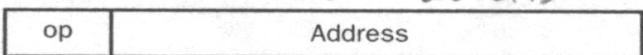
### 3. Base addressing *16-bits*



### 4. PC-relative addressing *16-bits*



### 5. Pseudodirect addressing *26-bits*





# MIPS Instruction Formats

Name	Fields						Comments
Field size	6 bits	5 bits	5 bits	5 bits	5 bits	6 bits	All MIPS instructions 32 bits
R-format	op	rs	rt	rd	shamt	funct	Arithmetic instruction format
I-format	op	rs	rt	address/immediate			Transfer, branch, imm. format
J-format	op	target address					Jump instruction format

R – Register  
I – Immediate  
J - Jump

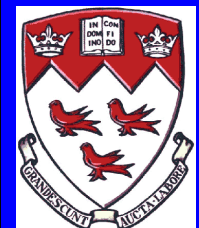
Immediate value  
(like an instruction constant)





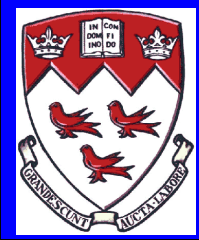
# Introduction to Computer Systems

## COMP 273



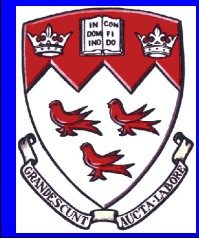
Category	Instruction	Example	Meaning	Comments
Arithmetic	add	add \$s1,\$s2,\$s3	$\$s1 = \$s2 + \$s3$	Three operands; overflow detected
	subtract	sub \$s1,\$s2,\$s3	$\$s1 = \$s2 - \$s3$	Three operands; overflow detected
	add immediate	addi \$s1,\$s2,100	$\$s1 = \$s2 + 100$	+ constant; overflow detected
	add unsigned	addu \$s1,\$s2,\$s3	$\$s1 = \$s2 + \$s3$	Three operands; overflow undetected
	subtract unsigned	subu \$s1,\$s2,\$s3	$\$s1 = \$s2 - \$s3$	Three operands; overflow undetected
	add immediate unsigned	addiu \$s1,\$s2,100	$\$s1 = \$s2 + 100$	+ constant; overflow undetected
	move from coprocessor register	mfc0 \$s1,\$epc	$\$s1 = \$epc$	Used to copy Exception PC plus other special registers
	multiply	mult \$s2,\$s3	Hi, Lo = $\$s2 \times \$s3$	64-bit signed product in Hi, Lo
	multiply unsigned	multu \$s2,\$s3	Hi, Lo = $\$s2 \times \$s3$	64-bit unsigned product in Hi, Lo
	divide	div \$s2,\$s3	Lo = $\$s2 / \$s3$ , Hi = $\$s2 \bmod \$s3$	Lo = quotient, Hi = remainder
	divide unsigned	divu \$s2,\$s3	Lo = $\$s2 / \$s3$ , Hi = $\$s2 \bmod \$s3$	Unsigned quotient and remainder
Logical	move from Hi	mfhi \$s1	$\$s1 = \text{Hi}$	Used to get copy of Hi
	move from Lo	mflo \$s1	$\$s1 = \text{Lo}$	Used to get copy of Lo
	and	and \$s1,\$s2,\$s3	$\$s1 = \$s2 \& \$s3$	Three reg. operands; logical AND
	or	or \$s1,\$s2,\$s3	$\$s1 = \$s2   \$s3$	Three reg. operands; logical OR
	and immediate	andi \$s1,\$s2,100	$\$s1 = \$s2 \& 100$	Logical AND reg, constant
	or immediate	ori \$s1,\$s2,100	$\$s1 = \$s2   100$	Logical OR reg, constant
Data transfer	shift left logical	sll \$s1,\$s2,10	$\$s1 = \$s2 \ll 10$	Shift left by constant
	shift right logical	srl \$s1,\$s2,10	$\$s1 = \$s2 \gg 10$	Shift right by constant
	load word	lw \$s1,100(\$s2)	$\$s1 = \text{Memory}[\$s2+100]$	Word from memory to register
	store word	sw \$s1,100(\$s2)	$\text{Memory}[\$s2 + 100] = \$s1$	Word from register to memory
	load byte unsigned	lbu \$s1,100(\$s2)	$\$s1 = \text{Memory}[\$s2 + 100]$	Byte from memory to register
Conditional branch	store byte	sb \$s1,100(\$s2)	$\text{Memory}[\$s2 + 100] = \$s1$	Byte from register to memory
	load upper immediate	lui \$s1,100	$\$s1 = 100 * 2^{16}$	Loads constant in upper 16 bits
	branch on equal	beq \$s1,\$s2,25	if ( $\$s1 == \$s2$ ) go to PC + 4 + 100	Equal test; PC-relative branch
	branch on not equal	bne \$s1,\$s2,25	if ( $\$s1 != \$s2$ ) go to PC + 4 + 100	Not equal test; PC-relative
	set on less than	slt \$s1,\$s2,\$s3	if ( $\$s2 < \$s3$ ) $\$s1 = 1$ ; else $\$s1 = 0$	Compare less than; two's complement
	set less than immediate	slti \$s1,\$s2,100	if ( $\$s2 < 100$ ) $\$s1 = 1$ ; else $\$s1 = 0$	Compare < constant; two's complement
Unconditional jump	set less than unsigned	sltu \$s1,\$s2,\$s3	if ( $\$s2 < \$s3$ ) $\$s1 = 1$ ; else $\$s1 = 0$	Compare less than; natural numbers
	set less than immediate unsigned	sltiu \$s1,\$s2,100	if ( $\$s2 < 100$ ) $\$s1 = 1$ ; else $\$s1 = 0$	Compare < constant; natural numbers
	jump	j 2500	go to 10000	Jump to target address
	jump register	jr \$ra	go to \$ra	For switch, procedure return
	jump and link	jal 2500	$\$ra = \text{PC} + 4$ ; go to 10000	For procedure call





Name	Format	Example						Comments
		6 bits	5 bits	5 bits	5 bits	5 bits	6 bits	
add	R	0	2	3	1	0	32	add \$1,\$2,\$3
sub	R	0	2	3	1	0	34	sub \$1,\$2,\$3
addi	I	8	2	1	100			addi \$1,\$2,100
addu	R	0	2	3	1	0	33	addu \$1,\$2,\$3
subu	R	0	2	3	1	0	35	subu \$1,\$2,\$3
addiu	I	9	2	1	100			addiu \$1,\$2,100
mfc0	R	16	0	1	14	0	0	mfc0 \$1,\$epc
mult	R	0	2	3	0	0	24	mult \$2,\$3
multu	R	0	2	3	0	0	25	multu \$2,\$3
div	R	0	2	3	0	0	26	div \$2,\$3
divu	R	0	2	3	0	0	27	divu \$2,\$3
mfhi	R	0	0	0	1	0	16	mfhi \$1
mflo	R	0	0	0	1	0	18	mflo \$1
and	R	0	2	3	1	0	36	and \$1,\$2,\$3
or	R	0	2	3	1	0	37	or \$1,\$2,\$3
andi	I	12	2	1	100			andi \$1,\$2,100
ori	I	13	2	1	100			ori \$1,\$2,100
sll	R	0	0	2	1	10	0	sll \$1,\$2,10
srl	R	0	0	2	1	10	2	srl \$1,\$2,10
lw	I	35	2	1	100			lw \$1,100(\$2)
sw	I	43	2	1	100			sw \$1,100(\$2)
lui	I	15	0	1	100			lui \$1,100
beq	I	4	1	2	25			beq \$1,\$2,100
bne	I	5	1	2	25			bne \$1,\$2,100
slt	R	0	2	3	1	0	42	slt \$1,\$2,\$3
slti	I	10	2	1	100			slti \$1,\$2,100
sltu	R	0	2	3	1	0	43	sltu \$1,\$2,\$3
sltiu	I	11	2	1	100			sltiu \$1,\$2,100
j	J	2	2500					j 10000
jr	R	0	31	0	0	0	8	jr \$31
jal	J	3	2500					jal 10000

# Partial MIPS Machine Language



# Part 2

## Coding Examples



# Example 1

C Language:

```
a = b + c;
```

MIPS:

```
lw $t1, b
```

```
lw $t2, c
```

```
add $t0, $t1, $t2
```

```
sw $t0, a
```

Note:

Cannot use variables  
directly in assembler



# Example 2

C Language:

$f = (g + h) - (i + j);$

MIPS:

lw \$s0, g

lw \$s1, h

lw \$s2, i

lw \$s3, j

add \$t0, \$s0, \$s1

add \$t1, \$s2, \$s3

sub \$s4, \$t0, \$t1

sw \$s4, f



# Example 3

## C Language:

```
g = h + A[8];
```

Assume the array A is of size 100 and contains 32-bit integer values.

## MIPS:

Assume: \$s1=g, \$s2=h, \$s3=base address of array

Note: *array* simply means *a contiguous block of memory* (nothing more)

```
lw $t0, 32($s3)
```

```
add $s1, $s2, $t0
```

```
sw $s1, g
```



# Example 4

C Language:

$A[12] = h + A[8];$

MIPS:

Assume similar setup from example 3.

lw \$t0, 32(\$s3)

add \$t0, \$s2, \$t0

sw \$t0, 48(\$s3)

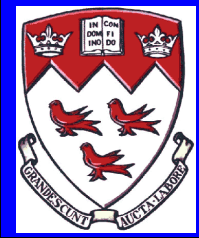


# IF-THEN-ELSE

- `if (i == j) f = g + h; else f = g - h;`

# Assume `f,i,j,g,h` are `$s0,$s3,$s4,$s1,$s2`

```
        bne $s3,$s4,Else    # go to Else if i != j
        add $s0,$s1,$s2     # f=g+h (skipped if i != j)
        j    Exit           # go to Exit
Else:    sub $s0,$s1,$s2     # f=f-h (skipped if i == j)
Exit:
```







# Example 5

## C Language:

```
if (i == A[j])  
{  
    f = g + h;  
} else {  
    f = f - i;  
}
```

## MIPS:

Assume: \$s0=i, \$s1=j, \$s2=f, \$s3=g, \$s4=h  
\$s5=base address of array  
containing 16-bit integer numbers

```
add $t0, $s1, $s1 #calculate offset  
add $t0, $t0, $s5 # base + offset  
lh $t1, 0($t0)  
bne $s0, $t1, L1 # could do beq also  
add $s2, $s3, $s4  
j Exit  
L1: sub $s2, $s2, $s0  
Exit: ...
```



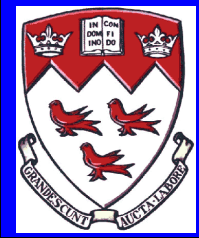
# WHILE LOOP

`while(save[i] == k) i = i + j;`

```
# Assume i,j,k are $s3,$s4,$s5
# Save address in $s6
```

```
Loop: add $t1,$s3,$s3      # Temp reg $t1 = 2 * i
      add $t1,$t1,$t1      # Temp reg $t1 = 4 * i
      add $t1,$t1,$s6      # $t1 = save[i]
      lw  $t0,0($t1)       # $t0 = save[i]
      bne $t0,$s5, Exit    # go to Exit if save != k
      add $s3,$s3,$s4      # i=i+j
      j   Loop
```

```
Exit:
```





# CASE/SWITCH Statement

```
switch(k) {  
    case 0: f = i + j; break;  
    case 1: f = g + h; break;  
}
```

# assume f,g,h,i,j,k are s0,s1,s2,s3,s4,s5

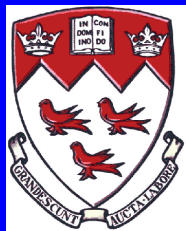
```
switch: addi $t0,$zero, 0          # $t0 = 0  
        beq  $t0,$s5,CASE0        # k == 0?  
        addi $t0,$zero, 1          # $t0 = 1  
        beq  $t0,$s5,CASE1        # k == 1?  
        j    Exit  
CASE0:  add  $s0,$s3,$s4  
        j    Exit  
CASE1:  add  $s0,$s1,$s2  
Exit:
```



# Set Less Than

```
# Assume $s0 has var_a and $s1 has var_b  
# is a < b?
```

```
slt $t0,$s0,$s1      # $t0=1 if a<b otherwise $t0=0  
bne $t0,$zero, Less  # go to Less if $t0 != 0 (if a<b)
```





# Question

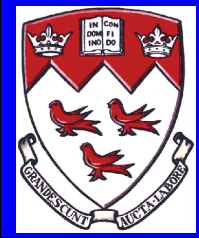
- Write a MIPS program snip that assumes someone's age is in temporary register 1. If the person is an adult, the program will calculate how many days they have been alive (assume 365 for each year) and store that into v0. If the person is underage then the program stores 0 into v0.

Use loop to multiply.



# Question

- Write a code snip that calculates  $x$  to the power of  $y$ .  $X$  and  $Y$  are positive numbers.



# Part 3

## MIPS Assembly Language Formatting



# MIPS Distinctions

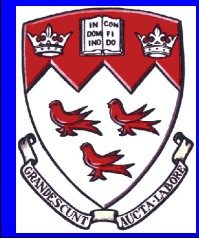
- Bit
  - The fundamental unit in a computer
- Byte
  - The fundamental grouping of units (it has an address)
- Word
  - The common data processing size
  - The size of a register (32 bits)
- Address
  - The built in size of memory addressing
  - The size of the address register  
(may be different from word) (486 24/32 bits)



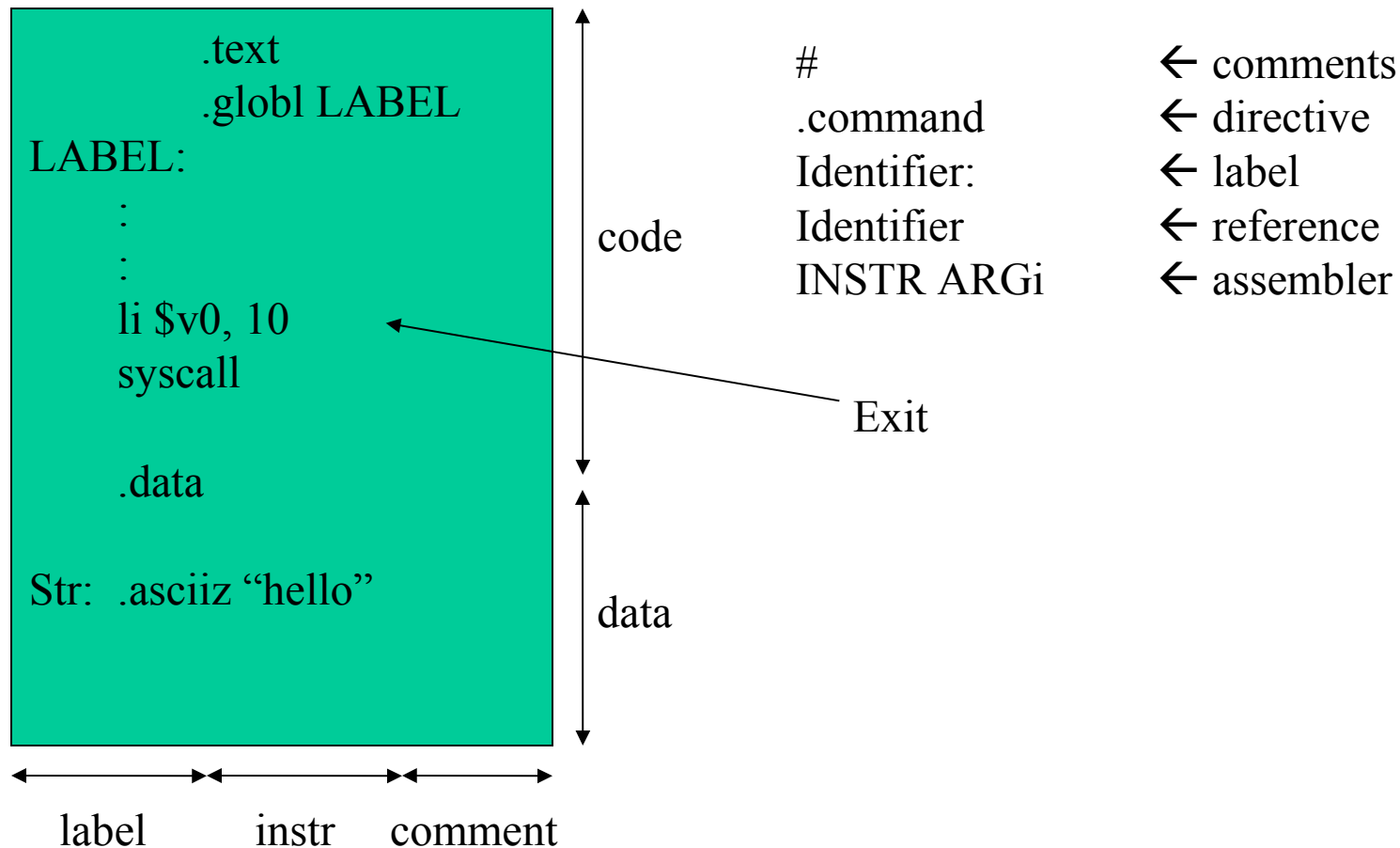


# Assembler Data Directives

- Syntax:
  - LABEL: DIRECTIVE DATA
- Where directives are:
  - .align n .data <addr>
  - .ascii str .extern LABEL SIZE(\$gp)
  - .asciiz str .globl LABEL
  - .space n .text <addr>
  - .byte b1,b2,...,bn ... 8
  - .half h1,h2,...,hn ... 16
  - .word w1,w2,...,wn .... 32
  - .float f1,f2,...,fn ..... 32
  - .double d1,d2,...,dn ..... 64



# The File Structure





# Commenting

(very important in assembler since so cryptic)

```
#
# Program name
# Date created
# Programmer Name
# Description of program
#
# List of major registers used by program
#

##### text segment #####
:
: # comments on almost every line
:
#
# code sectional comments
#

##### data segment #####
:
:

##### end of file: filename #####
```



# Commenting Example

```
##
## length.a - prints out the length of character
## string "str".
##
## t0 - holds each byte from string in turn
## t1 - contains count of characters
## t2 - points to the string
##
```

```
#####
#
# text segment
#
#####
```

```

.text
.globl __start
__start:      # execution starts here
    la $t2,str      # t2 points to the string
    li $t1,0        # t1 holds the count
nextCh: lb $t0,($t2) # get a byte from string
    beqz $t0,strEnd # zero means end of string
    add $t1,$t1,1    # increment count
    add $t2,$t2,1    # move pointer one character
    j nextCh        # go round the loop again

strEnd: la $a0,ans   # system call to print
    li $v0,4         # out a message
    syscall

    move $a0,$t1     # system call to print
    li $v0,1         # out the length worked out
    syscall

    la $a0,endl      # system call to print
    li $v0,4         # out a newline
    syscall

    li $v0,10
    syscall          # au revoir...

```

```
#####
#
# data segment
#
#####
```

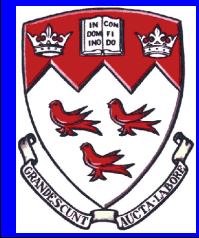
```

.data
str: .ascii "hello world"
ans: .ascii "Length is "
endl: .ascii "\n"

##
## end of file length.a

```

White spaces used to delineate sections but commenting would be good as well.



# Part 4

## More Coding Examples

# Example

```
clear1(int array[], int size)
{
    int i;
    for (i = 0; i < size; i = i + 1)
        array[i] = 0;
}

clear2(int *array, int size)
{
    int *p;
    for (p = &array[0]; p < &array[size]; p = p + 1)
        *p = 0;
}
```





## Array

```

move $t0,$zero    # i = 0
loop1: add $t1,$t0,$t0 # $t1 = i * 2
      add $t1,$t1,$t1 # $t1 = i * 4
      add $t2,$a0,$t1 # $t2 = &array[i]
      sw  $zero, 0($t2) # array[i] = 0
      addi $t0,$t0,1    # i = i + 1
      slt  $t3,$t0,$a1 # $t3 = (i < size)
      bne  $t3,$zero,loop1 # if () go to loop1
  
```

## Pointer

```

move $t0,$a0    # p = & array[0]
add  $t1,$a1,$a1 # $t1 = size * 2
add  $t1,$t1,$t1 # $t1 = size * 4
add  $t2,$a0,$t1 # $t2 = &array[size]
loop2: sw  $zero,0($t0) # Memory[p] = 0
      addi $t0,$t0,4    # p = p + 4
      slt  $t3,$t0,$t2 # $t3=(p<&array[size])
      bne  $t3,$zero,loop2 # if () go to loop2
  
```





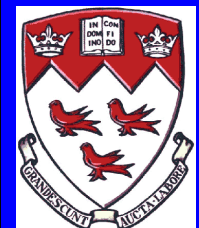


# Another Example

```
sort (int v[], int n)
{
    int i, j;
    for (i = 0; i < n; i = i + 1) {
        for (j = i - 1; j >= 0 && v[j] > v[j + 1]; j = j - 1) { swap(v, j);
        }
    }
}
```







Saving registers			
sort:	addi	\$sp,\$sp, -20	# make room on stack for 5 registers
	sw	\$ra, 16(\$sp)	# save \$ra on stack
	sw	\$s3,12(\$sp)	# save \$s3 on stack
	sw	\$s2, 8(\$sp)	# save \$s2 on stack
	sw	\$s1, 4(\$sp)	# save \$s1 on stack
	sw	\$s0, 0(\$sp)	# save \$s0 on stack
Procedure body			
Move parameters	move	\$s2, \$a0	# copy parameter \$a0 into \$s2 (save \$a0)
	move	\$s3, \$a1	# copy parameter \$a1 into \$s3 (save \$a1)
Outer loop	move	\$s0, \$zero	# i = 0
	for1tst:slt	\$t0, \$s0, \$s3	# reg \$t0 = 0 if \$s0 ≥ \$s3 (i ≥ n)
	beq	\$t0, \$zero, exit1	# go to exit1 if \$s0 ≥ \$s3 (i ≥ n)
Inner loop	addi	\$s1, \$s0, -1	# j = i - 1
	for2tst:slti	\$t0, \$s1, 0	# reg \$t0 = 1 if \$s1 < 0 (j < 0)
	bne	\$t0, \$zero, exit2	# go to exit2 if \$s1 < 0 (j < 0)
	add	\$t1, \$s1, \$s1	# reg \$t1 = j * 2
	add	\$t1, \$t1, \$t1	# reg \$t1 = j * 4
	add	\$t2, \$s2, \$t1	# reg \$t2 = v + (j * 4)
	lw	\$t3, 0(\$t2)	# reg \$t3 = v[j]
	lw	\$t4, 4(\$t2)	# reg \$t4 = v[j + 1]
	slt	\$t0, \$t4, \$t3	# reg \$t0 = 0 if \$t4 ≥ \$t3
	beq	\$t0, \$zero, exit2	# go to exit2 if \$t4 ≥ \$t3
Pass parameters and call	move	\$a0, \$s2	# 1st parameter of swap is v (old \$a0)
	move	\$a1, \$s1	# 2nd parameter of swap is j
	jal	swap	# swap code shown in Figure 3.24
Inner loop	addi	\$s1, \$s1, -1	# j = j - 1
	j	for2tst	# jump to test of inner loop
Outer loop	exit2: addi	\$s0, \$s0, 1	# i = i + 1
	j	for1tst	# jump to test of outer loop
Restoring registers			
exit1:	lw	\$s0, 0(\$sp)	# restore \$s0 from stack
	lw	\$s1, 4(\$sp)	# restore \$s1 from stack
	lw	\$s2, 8(\$sp)	# restore \$s2 from stack
	lw	\$s3,12(\$sp)	# restore \$s3 from stack
	lw	\$ra,16(\$sp)	# restore \$ra from stack
	addi	\$sp,\$sp, 20	# restore stack pointer
Procedure return			
	jr	\$ra	# return to calling routine



# Questions

- Shift the contents of an array left  $n$  cells.