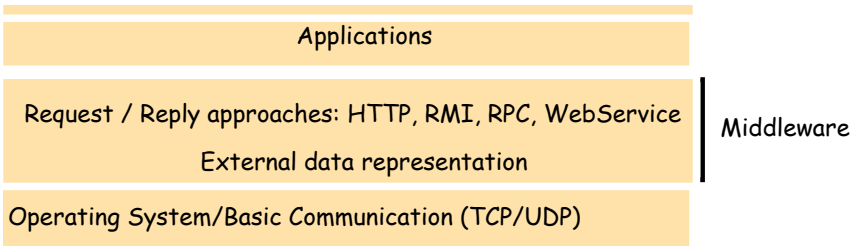
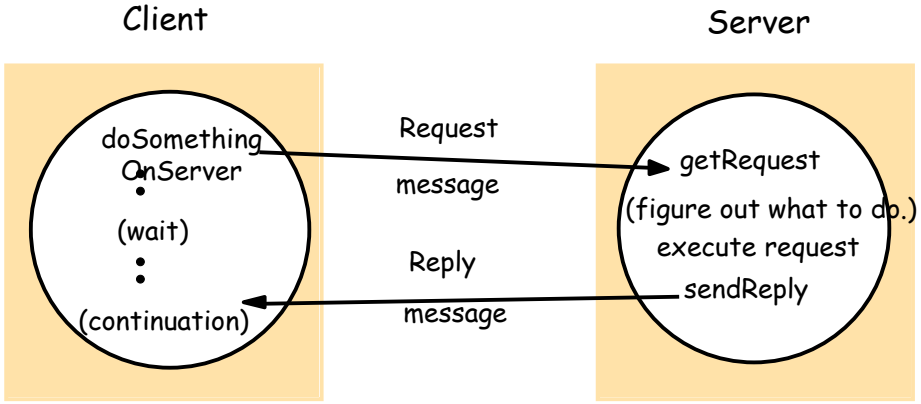
Request Reply
Remote Method Invocation
Remote Service Invocation

Middleware layers



Client-Server communication



Implementation of RR over UDP

- Failure model:
  - ☆ message loss, process crash
- Invocation semantics:
  - ☆ *maybe*
    - provides *best-effort*
    - avoids blocking by *timeout*
      - ▲ give client failure exception after timeout

Implementation of RR over UDP

- Invocation semantics: at-least-once
  - client software
    - resubmits request if it does not receive result within timeout interval (up to N submissions)
    - returns after receiving first reply (ignores further replies it might get)
  - Server might execute request more than once
  - acceptable if service is *idempotent*
  - client resubmits up to N times and then gives up
    - Up to N-1 message losses are accepted by the protocol
    - Client cannot distinguish whether request or reply messages are lost or whether server has crashed
    - After N timeouts it does not know whether server has executed request or not

Implementation of RR over UDP

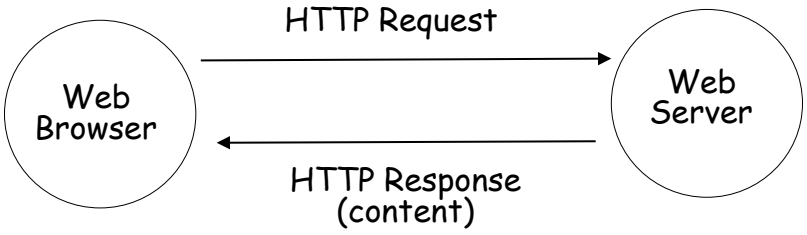
- Invocation semantics: at-most-once (or better exactly-once)
  - client software
    - Same as at-least-once
  - server must detect duplicate requests
    - unique request identifiers
      - typical: IP address of client + process id + sequence id within process
    - server must keep all requests (*history*)
    - once reply is generated, reply is added to corresponding request
    - upon receiving request again,
      - reply is returned without reexecution
      - note: if timeout too small, server might receive second submission of request while executing first: must detect this!
  - garbage collection at server?
    - Client sends ack when it has received response (RRA)
    - Similar optimizations as TCP

Comparison

- Compare at-most-once RR over UDP with TCP
  - ☆ Similar mechanisms but slightly different implementation
    - Exactly-once service execution vs. exactly-once message delivery
- Compare at-most-once RR over UDP with RR using simply TCP
  - ☆ Both provide the same semantics (exactly-once as long as N tries are sufficient)
  - ☆ Performance different
    - Different number of messages sent

HTTP: Hypertext Transfer Protocol

- On top of TCP



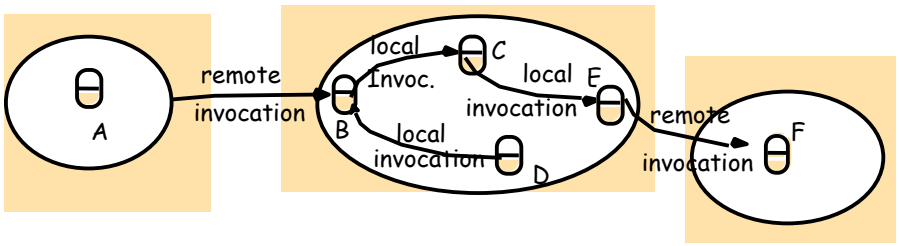
Features

- Request is on a Resource
- Uniform Resource Identifier / Uniform Resource Locator (URL)
  - Location
    - Logical address name plus potentially port
  - Unique Name within location
    - Can be a file (static content)
    - Can be a program
  - http:www.cs.mcgill.ca/~kemme/cs512/index.html
- Most common:
  - Get Request to retrieve static content

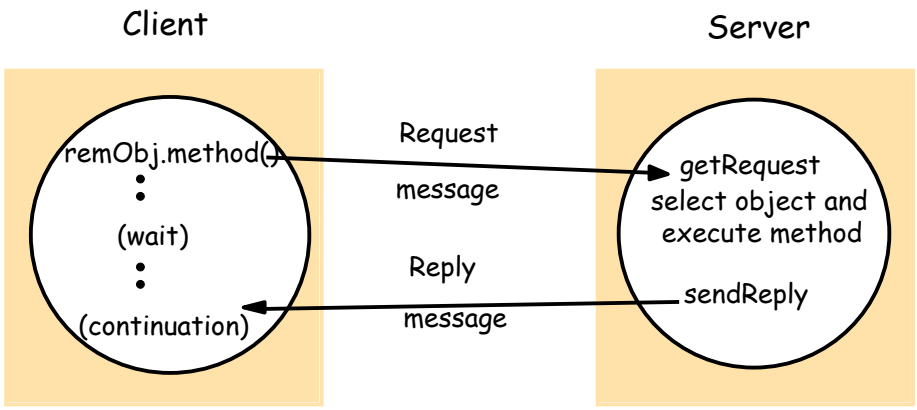
Advanced Requests

- Methods on the resource
  - Get
    - On data: return content
    - On program: run program and return result
  - Put
    - URL already exists: override
    - New URL: create
  - Post
    - Add data to the resource (e.g. a mailing list, database…)
  - Others
  - Parameters: program parameters
    - E.g., data entered in a form…
    - Syntax can get complicated!
- Textual Format

Remote Method Invocation



Client-Server communication



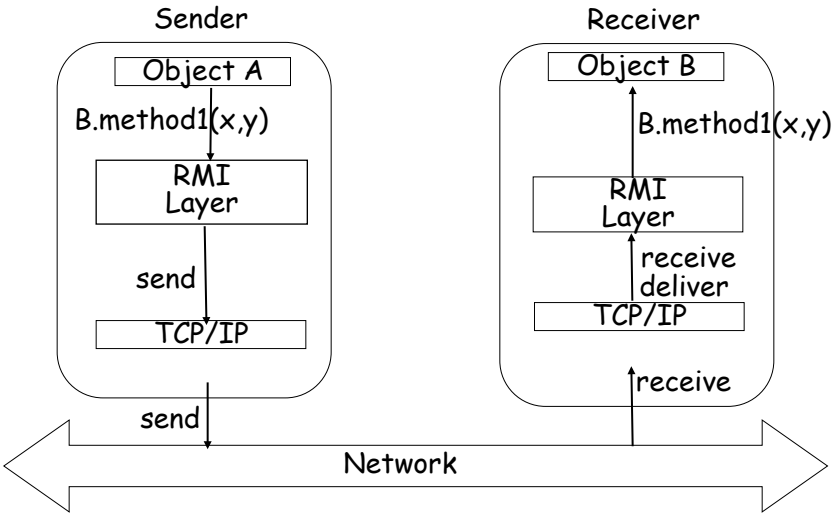
Programming Models for Distributed Applications

- Remote Procedure Call (RPC)
  - Client calls a procedure in a server running on a different machine
  - Idea: call to remote procedure should look like a call to local procedure
  - Theory since 1976, first implementations in the early 80's
- Remote Method Invocation (RMI)
  - Same as RPC in OO-world
  - Object calls method of another object residing on a different machine
  - First: implementations had same architecture than RPC
  - Current developments: more features and variations
  - Trend to truly distributed computing

Review of OO-programming

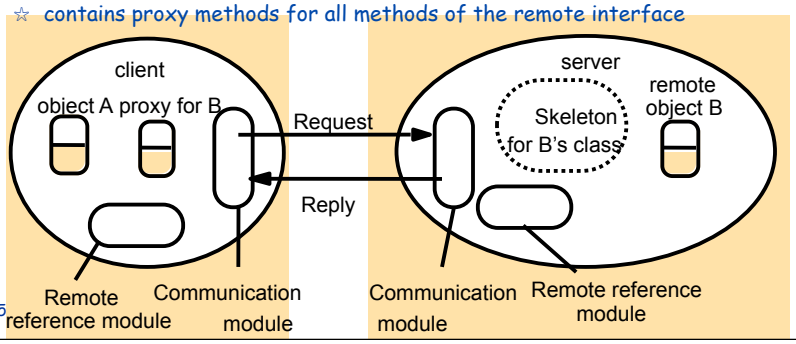
- An object encapsulates both *data* and *methods*
- Objects are accessed via *object references*
- Method Invocation can cause
  - the state of the called object to change
  - further invocations on methods in same or other objects
- Access to variables (direct/indirect)
  - direct (e.g., public): not suitable in distributed world
  - indirect (e.g., via getter/setter methods): encapsulation
  - Assumption from now: variables of an object can only be accessed via method calls
- *Exceptions*
  - are thrown when an error occurs.
  - can be caught (redirection to specific code handling the exception)
  - or delivered to caller.
- *Interfaces*:
  - provide signatures of a set of methods
  - description of input and output parameters of the methods
- *Garbage collection*
  - frees space when objects are no longer needed

RMI I



RMI I

- *Remote Interface File* of an object class
  - Interface file for the set of methods that can be called remotely
  - object class can have further methods that can only be called locally
- Method description in remote interface file
  - input and output parameters are specified
  - a parameter can be an object
- Implementation of local and remote methods
  - only at server
- Proxy object at client (stub)
  - contains proxy methods for all methods of the remote interface



RMI II

- On the client
  - ☆ object A calls method of remote object B
    - it internally calls the proxy-method of the local proxy object.
    - proxy also called *stub*
    - i.e., object A has a object reference to the local proxy object
  - ☆ Within proxy-method and underlying communication module (RR module)
    - marshals *request* message containing
      - ▲ *remote reference* to the remote target object,
      - ▲ its own methodID
      - ▲ and the arguments
    - sends request message to server
    - awaits reply message
    - unmarshals reply
    - returns the results to the invoker

RMI III

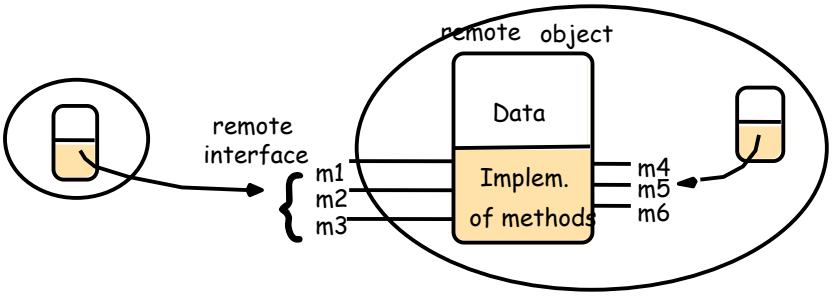
- On the server
  - ☆ RMI communication module receives message
  - ☆ Determines skeleton of object B given the reference and calls skeleton's method based on methodID
  - ☆ Skeleton method unmarshals message and invokes corresponding method on real object B
  - ☆ marshals the result (together with any exception) into result message
  - ☆ Communicatin module sends result message back to RMI layer on caller
  - ☆ Proxies and skeletons are automatically generated with special compiler

Interface File in Java RMI

```
import java.rmi.Remote;
import java.rmi.RemoteException;
public interface NumberStack extends Remote {
    public int getNumber() throws RemoteException;
    public void putNumber(int num) throws RemoteException; }
```

- By extending java.rmi.Remote, this interface' s methods can be called from any virtual machine. Any object that implements this interface becomes a remote object.
- As a member of a remote interface, the getNumber method is a remote method. Therefore the method must be defined as being capable of throwing a java.rmi.RemoteException.
- This exception is thrown by the RMI system during a remote method when communication failure or a protocol error has occurred.
- Any code making a call to a remote method needs to handle this exception by either catching it or declaring it in its throws clause.

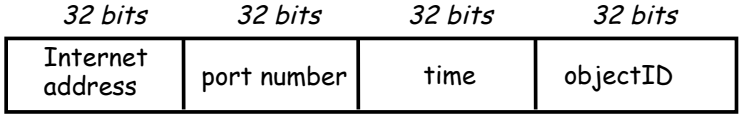
A remote object and its remote interface



Interface Definition Languages

- Java has the concept of interfaces and hence, the remote interface can be described as usual + extending interface to be REMOTE
- In general interfaces are described with Interface Definition Languages
- Examples RMI
  - Corba IDL
  - DCOM IDL for Microsoft's Distributed Common Object Model (DCOM) RMI
- Examples RPC
  - Sun XDR
  - DCE IDL: RPC system of OSF's Distributed Computing Environment DCE

Remote Object Reference

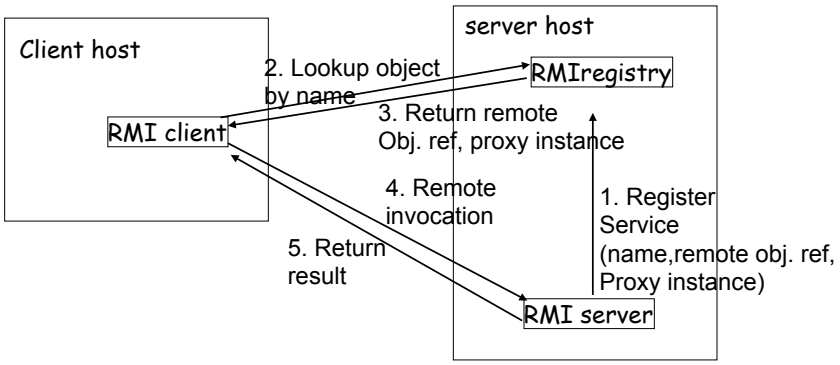


- In non-distributed execution, objects can be accessed via object reference.
- In Java, a variable that appears to hold an object actually holds a reference to that object.
- An object A can invoke the method of a remote object B, if it has access to B's remote object reference.
- Generally: A remote object reference is an identifier that can be used throughout a distributed system to refer to a particular unique remote object.
- Note: In some RMI online information the remote object reference actually refers to a stub/proxy instance
- Conceptually like a URL

Remote reference module

- Maintains remote object table
  - ☆ entry
    - remote object reference and local proxy reference (on client)
    - remote object reference and real object reference (on server)
- Some Tasks
  - ☆ On a server: when a remote object reference arrives in a request message it looks for the corresponding local object reference which points to the real object
  - ☆ On a client: when a remote object reference arrives in a reply message it looks for the proxy object.
  - ☆ When a remote object is to be passed as argument or result for the first time, the remote reference module creates a remote object reference and adds it to the table

A basic Java RMI call



The binder: RMI Registry

- An instance of the Rmiregistry must run on every computer on which a server wants to export a remote interface.
- It maintains a table mapping textual references to remote objects hosted on that computer.
  - //computeName:port/objectName
  - computeName is the host name on which the registry resides and port is the port the rmiregistry listens on

RMI Registry Interface

- Void rebind(String name, Remote obj)
  - Register object by name
  - Override previous registration
- Void bind (String name, Remote obj)
  - Register an object by name
  - If existent throw exception
- Void unbind(String name, remote obj)
- Remote lookup(String name)
  - Used by client
  - Remote object reference is returned
- String[] list()
  - Show all names bound in this registry

Method Parameters

- Input and output parameters can be
  - Primitive types (pass by value)
  - Remote objects (pass the remote object reference)
    - If the recipient of a remote object reference does not have the object proxy, then the object proxy is automatically downloaded
  - Local objects. They are passed by copy, using object serialization (the object must implement the java.io.Serializable Interface).
    - If the recipient of an object passed by value does not already possess the class of the object, its code is automatically downloaded.

Distributed Garbage Collection

- An object must consist as long as there is a local or remote reference held
- Whenever the last references is deleted the object should be garbage collected.
- A classical distributed algorithm problem:
  - Solution: reference counting
- 2-level
  - local proxy garbage collection on each client
    - Whenever a remote object reference enters a process, a proxy is created, when the last reference to this object on the process is deleted, the proxy is deleted
  - object maintenance on the server
    - whenever there is no proxy anymore **and** there is no local reference, you can delete the object

Distributed Garbage Collection

- Algorithm outline:
  - Each server process maintains a list of processes that currently hold a remote object reference
  - Whenever a new proxy is created (a client first requests a reference), the remote reference module on the client first calls a function addRef on the server and then creates the proxy
  - When a proxy on a client is no more needed (detected by the local GC of the client), the remote reference module calls removeRef on the server
  - When the list is empty and there are no local references the remote object is removed
- Pair-wise request-reply between remote reference modules
- Only called when proxies are created/deleted
- Fault-tolerance has to be addressed
  - Idempotent addRef and removeRef
  - removeRef is correct whether addRef worked or not
  - Leases (max. time to live) in case removeRef gets lost

Web Service: Interface

- Remote *Service* Call
  - Set of services / procedures / methods that can be called
- Service Description (all XML)
  - WSDL: web-service definition language
  - Service
    - Described by the format of request / reply messages
  - Bindings
    - Details about protocols
  - Endpoint address
    - Different to RMI!!

Web-Service: Communication

- SOAP: Simple Object Access Protocol
- Messages in XML
  - SOAP messages have a specific format
- Protocol Pattern
  - Request/reply
    - Based on HTTP
    - HTTP contains endpoint address and action for efficiency
  - Other models are offered
    - In-only
    - Out-only
    - Reliable versions

Web-Service with Java

- Java API for developing web-services: JAX-RPC
- Hides all the details of web-services
  - Define your interface using Java interface
    - Some restrictions for parameter types (objects are allowed, but no remote objects)
  - Web-service is a single object that offers the services

Web-Service: Components

- Communication Module: Http
- Container:
  - Dispatcher calls skeleton (according to http header URL)
  - Skeleton (per servlet) translates SOAP message to java servlet call; translates reply back to SOAP