

# Maze Solver using Iterative Deepening Search (IDS) with Binary Search Optimization

## 1. Introduction

This project presents a professional C++ implementation of a maze-solving algorithm based on Iterative Deepening Search (IDS), enhanced with a Binary Search strategy to optimize the depth selection process. The objective is to efficiently find a valid path from a fixed starting cell (0,0) to a user-defined goal position in a two-dimensional maze.

## 2. Maze Representation

The maze is modeled as a two-dimensional grid of size  $N \times M$ , where each cell represents either a traversable path or an obstacle:

- 0: Empty cell (traversable)
- 1: Wall or obstacle (non-traversable)

## 3. Algorithm Overview

The solution combines Depth-Limited Search (DLS) with Binary Search to efficiently determine the minimum depth required to reach the goal.

### 3.1 Depth-Limited Search (DLS)

Depth-Limited Search is a constrained variant of Depth-First Search (DFS). It explores the maze recursively while enforcing a strict upper bound on the maximum search depth. To ensure correctness and prevent infinite loops:

- A *visited* matrix is used to avoid revisiting cells.
- Backtracking is applied to construct the final path once the goal is reached.
- Recursive calls terminate immediately when the depth limit is exceeded.

### 3.2 Binary Search over Depth Limits

Traditional IDS increases the depth limit incrementally (1, 2, 3, ...). In contrast, this implementation applies Binary Search over the depth range [1,  $N \times M$ ] to significantly reduce the number of DLS executions.

## 4. Monotonic Function Justification

The use of Binary Search is mathematically justified by the monotonic nature of the reachability condition in maze traversal.

If a path to the goal exists within a depth limit  $D$ , then the same path (or another valid path) must also exist for any depth limit greater than  $D$ . Conversely, if no path exists within  $D$  steps, it is impossible for any smaller limit.

This behavior forms a step-like boolean function that transitions from False to True exactly once, making it suitable for Binary Search.

## 5. Complexity Analysis

### Time Complexity

A single Depth-Limited Search visits each cell at most once, resulting in a time complexity of  $O(N \times M)$ . Binary Search performs  $\log(N \times M)$  iterations. Therefore, the total time complexity is:

$$O((N \times M) \cdot \log(N \times M))$$

### Space Complexity

The algorithm requires  $O(N \times M)$  space for the visited matrix, recursion stack, and path storage. Hence, the overall space complexity is  $O(N \times M)$ .

## 6. Limitations

While the Binary Search optimized IDS significantly reduces search overhead, it does not guarantee the shortest possible path in all cases, unlike Breadth-First Search (BFS). Additionally, recursive DLS may face stack depth limitations for extremely large mazes.

## 7. Conclusion

This project demonstrates how algorithmic properties such as monotonicity can be leveraged to optimize classical search techniques. The integration of Binary Search with IDS results in a scalable and efficient maze-solving strategy suitable for academic and competitive programming applications.