# Module Title: INFORMATICS 1A Exam Diet (Dec/April/Aug): DECEMBER 2005 Brief notes on answers:

```
1. type Day = Int
  type Month = String
  type Date = (Day, Month)
  months :: [Month]
  months = ["Jan", "Feb", "Mar", "Apr", "May", "Jun",
             "Jul", "Aug", "Sep", "Oct", "Nov", "Dec"]
   (a) index :: Month -> Int
      index m = head [i | (i,m') \leftarrow zip [1..] months, <math>m==m']
      or
      index :: Month -> Int
      index "Jan" = 1
      index "Feb" = 2
      index "Dec" = 12
   (b) sensible :: Date -> Bool
      sensible (d,m) = 1 \le d \&\& d \le 31 \&\& elem m months
   (c) before :: Date -> Date -> Bool
      before (d,m) (d',m') = index m < index m' || index m == index m' && d < d'
   (d) show2 :: Int -> String
      show2 n = [intToDigit (n 'div' 10), intToDigit (n 'mod' 10)]
      showDate :: Date -> String
      showDate (d,m) | sensible (d,m) = show2 d ++ "/" ++ show2 (index m)
```

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2. (a) six :: Int -> Int
      six n = (if n 'mod' 2 == 0 then 2 else 1) * (if n 'mod' 3 == 0
      then 3 else 1)
      or
      six :: Int -> Int
      six n = six' (n 'mod' 6)
      six' 0 = 6
      six' 1 = 1
      six' 2 = 2
      six' 3 = 3
      six' 4 = 2
      six' 5 = 1
   (b) sixes :: [Int] -> [Int]
      sixes xs = [ six x | x \leftarrow xs, x > 0 ]
   (c) sixes :: [Int] -> [Int]
      sixes [] = []
      sixes (x:xs) \mid x > 0 = six x : sixes xs
                  | otherwise = sixes xs
```

```
3. (a) befores :: [Date] -> Bool
    befores xs = and [ before x x' | (x,x') <- zip xs (drop 1 xs) ]

(b) befores :: [Date] -> Bool
    befores [] = True
    befores [x] = True
    befores (x:y:zs) = before x y && befores (y:zs)
```

## Part B COMPUTATION AND LOGIC

- 4. The truth tables and outcomes are as follows:
  - (a) This is contingent.

			E1	E2	E3	
	a	b	not(a)	E1 or b	$b \rightarrow a$	$E2 \leftrightarrow E3$
	t	t	f	t	t	t
	t	f	f	f	t	f
ĺ	f	t	t	t	f	f
	f	f	t	t	t	t

(b) This is a tautology.

		E1	E2	E3	E4		
$\mid a \mid$	b	$a \rightarrow b$	not(a)	not(b)	$E3 \rightarrow E2$	$E1 \rightarrow E4$	
t	t	t	f	f	t	t	
t	f	f	f	t	f	t	
f	t	t	t	f	t	t	
f	f	t	t	t	t	t	

(c) This is an inconsistency.

		E1	E2	E3	E4	
$\mid a \mid$	b	$a \rightarrow b$	not(a)	E2 or b	not(E1)	E4 and E3
t	t	t	f	t	f	f
t	f	f	f	f	t	f
f	t	t	t	t	f	f
f	f	t	t	t	f	f

- 5. For this question the answer requires the circuit to be described as a logical expression; then the expression solved in two different ways.
  - (a) The expression is:

$$not((not(a \ or \ b) \ or \ c) \ and \ (not(c \ and \ d) \ or \ e)) \ or \ e$$

The truth table is:

					E1	E2	E3	E4	E5	E6	E7	E8	
a	b	c	d	e	$a\ or\ b$	not(E1)	$E2\ or\ c$	$c \ and \ d$	not(E4)	$E5\ or\ e$	$E3\ and\ E6$	not(E7)	$E8 \ or \ e$
t	t	t	t	t	t	f	t	t	f	t	t	f	t
t	t	t	t	f	t	f	t	t	f	f	f	t	t
t	t	f	t	t	t	f	f	f	t	t	f	t	t
t	t	f	t	f	t	f	f	f	t	t	f	t	t
t	f	t	t	t	t	f	t	t	f	t	t	f	t
t	f	t	t	f	t	f	t	t	f	f	f	t	t
t	f	f	t	t	t	f	f	f	t	t	f	t	t
t	f	f	t	f	t	f	f	f	t	t	f	t	t

The expression is true (with a and d true) regardless of the truth values of the other inputs. Hence those other inputs are redundant as far as the behaviour of the circuit is concerned.

- (b) An appropriate proof is via the following sequence of proof rules (some details of sub-proofs omitted here the student should supply them):
  - To prove  $[b, (a \ or \ b) \rightarrow c, c \rightarrow e] \vdash e$ :
  - $imp\_elim$  gives  $[b, (a \ or \ b) \rightarrow c, c \rightarrow e] \vdash c$
  - $imp\_elim$  gives  $[b, (a \ or \ b) \rightarrow c, c \rightarrow e] \vdash (a \ or \ b)$
  - or\_intro\_right gives  $[b, (a \text{ or } b) \rightarrow c, c \rightarrow e] \vdash b$

### 6. Answers are:

- (a) A plausible FSM is as follows, where ic is inserting a card; rs is requesting a statement; ds is dispensing a statement; rm is requesting money; re is refusing money; dm is dispensing money; and ec is ejecting a card.
- (b) The transition relation is described as the following table:

	ic	rs	ds	rm	re	dm	ec
1	2						
2		5		3			
3					2	4	
4							1
5			2				

(c) This FSM is deterministic because there is no state at which the same action leads to two different states.

#### 7. Answers are:

- (a) An appropriate FSM is as follows:
- (b) It is NOT possible to draw the FSM for this acceptor because it would need to be able to count arbitrarily large numbers of 0s.

#### 8. Answers are:

- (a) ((0|10)\*)|((0|10)\*1).
- (b) (0(0|1)\*0)|(1(0|1)\*1).