

# VE203 Discrete Math

## Spring 2022 — HW7 Solutions

April 19, 2022



### Exercise 7.1

i) For

$$f : B \rightarrow U, \quad B = \{1, 2, 3\}, U = \{1, 2, 3, 4, 5\}$$

we have

Domain	Codomain	Any	Injective	surjective
distinguishable	distinguishable	$5^3$	$\binom{5}{3} \cdot 3!$	0
indistinguishable	distinguishable	$\binom{7}{3}$	$\binom{5}{3}$	0
distinguishable	indistinguishable	5	1	0
indistinguishable	indistinguishable	3	1	0

ii) For

$$f : B \rightarrow U, \quad B = \{1, 2, 3, 4, 5\}, U = \{1, 2, 3\}$$

we have

Domain	Codomain	Any	Injective	Surjective
distinguishable	distinguishable	$3^5$	0	$25 \times 3!$
indistinguishable	distinguishable	$\binom{7}{5}$	0	$\binom{4}{2}$
distinguishable	indistinguishable	41	0	25
indistinguishable	indistinguishable	5	0	2

### Exercise 7.2

By the fundamental theorem of arithmetic there is a unique decomposition for  $n$  in product of primes numbers :  $n = p_1^{a_1} p_2^{a_2} \dots p_k^{a_k}$ . Thus we have :

$$\begin{aligned} \varphi(n) &= \varphi(p_1^{a_1}) \varphi(p_2^{a_2}) \dots \varphi(p_k^{a_k}) \\ \varphi(n) &= p_1^{a_1} p_2^{a_2} \dots p_k^{a_k} \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \dots \left(1 - \frac{1}{p_k}\right) \\ \varphi(n) &= n \left(1 - \frac{1}{p_1}\right) \left(1 - \frac{1}{p_2}\right) \dots \left(1 - \frac{1}{p_k}\right) \end{aligned}$$

We obtain the formula stated before.

$$\begin{aligned}\prod_{i=1}^n (1 - x_i) &= 1 - \sum_{i=1}^n x_i + \sum_{i,j=1}^n x_i x_j - \sum_{i,j,k=1}^n x_i x_j x_k + \cdots + (-1)^n x_1 x_2 \cdots x_n \\ &= \sum_{I \subset \{1,2,\dots,n\}} (-1)^{|I|} \prod_{i \in I} x_i\end{aligned}$$

when you consider the numbers which are multiple of  $p_1$  or  $p_2$ , if you want to count them you have to compute  $\frac{n}{p_1} + \frac{n}{p_2} - \frac{n}{p_1 p_2}$ , you subtract the number of integers which are in the same time a multiple of  $p_1$  and  $p_2$ . Following this reasoning we have :

$$\begin{aligned}\varphi(n) &= n - \sum_{p_i \text{ prime } p_i | n} \frac{n}{p_i} + \sum_{p_i, p_j \text{ prime } p_i, p_j | n} \frac{n}{p_i p_j} - \sum_{p_i, p_j, p_k \text{ prime } p_i, p_j, p_k | n} \frac{n}{p_i p_j p_k} + \cdots \\ &\quad + (-1)^{|Pr|} \frac{n}{p_1 p_2 \cdots p} \\ &= n \left( 1 - \sum \frac{1}{p_i} + \sum \frac{1}{p_i p_j} - \sum \frac{1}{p_i p_j p_k} + \cdots + (-1)^{|Pr|} \frac{1}{p_1 p_2 \cdots p} \right) \\ &= n \prod_{p \in P_r} \left( 1 - \frac{1}{p} \right)\end{aligned}$$

### Exercise 7.3

i) If  $x_i > 0$  and the equality holds. Namely, we have

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 = 100$$

with  $x_i > 0$ , or equivalently,  $x_i \geq 1$ . Then we can divide 100 into

$$\underbrace{1 + 1 + \cdots + 1}_{100 \text{ terms}} = 100$$

Since we have 7 variables, then the number of integer solutions is given by

$$\binom{99}{6}$$

which stands for choosing 6 of + sign out of 99+ sign to form a 7- partition of 100 .

ii) If  $x_i \geq 0$  and the equality holds. Namely, we have

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 = 100$$

with  $x_i \geq 0$ , or equivalently, we have

$$x'_1 + x'_2 + x'_3 + x'_4 + x'_5 + x'_6 + x'_7 = 107$$

with  $x'_i > 0$ , or equivalently,  $x'_i \geq 1$  since  $x'_i := x_i + 1$ . Then by the same argument from (i), we know that the number of integer solutions is given by

$$\binom{106}{6}$$

iii) If  $x_i > 0$  and the equality does not hold. We introduce another variable  $x_8 > 0$  such that

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 + x_8 = 100$$

which is equivalent to the original one. Then by the same argument from (i), we know that the number of integer solutions is given by

$$\binom{99}{7}$$

iv) If  $x_i \geq 0$  and the equality does not hold. We introduce another variable  $x_8 > 0$  such that

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 + x_8 = 100,$$

which is equivalent to the original one. By the same transformation of variables as in (ii), we have

$$x'_1 + x'_2 + x'_3 + x'_4 + x'_5 + x'_6 + x'_7 + x_8 = 107,$$

then now all the variables are strictly greater than 0. By the same argument as in (i), we know that the number of integer solution is given by

$$\binom{106}{7}$$

v) If  $x_i \geq 0$ . Then the number of integer solution is given by the results in (ii) and (iv) combined. Since we know that the solution sets in the case of

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 = 100$$

and the case of

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 < 100$$

are mutually disjoint, hence the number of solution can be combined, namely

$$\binom{106}{6} + \binom{106}{7} = \binom{107}{7}$$

by the recursive identity for binomial coefficient.

## Exercise 7.4

i) From the Master Theorem, we see that the recurrence relation is in the form of

$$T(n) = aT(n/b) + O(n^d)$$

for constants  $a \geq 1, b > 1, d \geq 0$ . Specifically, we see that

$$a = 4, \quad b = 4, \quad d = 1$$

since  $5n \in O(n)$ . Then, we find out

$$\log_b a = \log_4 4 = 1 = d,$$

hence from Master Theorem we conclude that

$$T(n) = O(n^d \log n) = O(n \log n).$$

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since  $4n \in O(n)$ . Then, we find out

$$\log_b a = \log_4 5 > 1 = d,$$

hence from Master Theorem we conclude that

$$T(n) = O(n^{\log_b a}) = O(n^{\log_5 4}).$$

iv)

$$\text{Let } n = 2^m \text{ and } S(m) = T(2^m)$$

$$\Rightarrow S(m) = 4S\left(\frac{m}{2}\right) + m^5$$

$$a = 4, b = 2, d = 5, \Rightarrow \log_b a = 2 < d$$

$$\Rightarrow S(m) = O(m^5)$$

$$\Rightarrow T(n) = S(\log m) = O((\log n)^5)$$

v)

$$\text{Let } n = 2^m \text{ and } S(m) = T(2^m)$$

$$\Rightarrow S(m) = 4S\left(\frac{m}{2}\right) + m^2$$

$$a = 4, b = 2, d = 2, \Rightarrow \log_b a = 2 = d$$

$$\Rightarrow S(m) = O(m^2 \log m)$$

$$\Rightarrow T(n) = S(\log m) = O((\log n)^2 \log \log n)$$

## Exercise 7.5

WIP