

# Ve203 Discrete Mathematics (Fall 2022)

## Assignment 1

**Date Due:** See canvas

This assignment has a total of **(21 points)**.

**Note:** Unless specified otherwise, you must show the details of your work via logical reasoning for each exercise. Simply writing a final result (whether correct or not) will receive **0 point**. **Explain** (briefly) if you claim something is trivial or straightforward. Provide a counterexample if you are trying to disprove something. It is **NOT OK** to write something like “how do we know that blahblahblah is even true...”

### Exercise 1.1 (6 pts)

Given  $\varphi = (A \rightarrow (B \rightarrow C)) \rightarrow (B \rightarrow (A \rightarrow C))$ ,

- (i) (2 pts) Write the truth table for  $\varphi$ .
- (ii) (2 pts) Write  $\varphi$  in disjunctive normal form.
- (iii) (2 pts) Write  $\varphi$  in conjunctive normal form.

### Exercise 1.2 (6 pts)

The following shows the truth table for all  $2^{2^2} = 16$  different binary logical operators  $\varphi_i, i = 0, \dots, 15$ .

$p$	$q$	$\varphi_0$	$\varphi_1$	$\varphi_2$	$\varphi_3$	$\varphi_4$	$\varphi_5$	$\varphi_6$	$\varphi_7$	$\varphi_8$	$\varphi_9$	$\varphi_{10}$	$\varphi_{11}$	$\varphi_{12}$	$\varphi_{13}$	$\varphi_{14}$	$\varphi_{15}$
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

Using infix notation, for example,  $\varphi_{13}$  can be represented as  $\varphi_{13} = \rightarrow(p, q) = p \rightarrow q$ .

A set  $S$  of logical operators is called *functionally complete* if every compound proposition is logically equivalent to a compound proposition involving only these logical operators in  $S$ . In this exercise, in order to show  $S$  is a functionally complete set, it suffices to verify that for all  $i = 0, \dots, 15$ ,  $\varphi_i$  over logical variables  $p$  and  $q$  can be represented using only operators in  $S$ .

- (i) (1 pt) Show that  $\{\wedge, \vee, \neg\}$  is functionally complete.
- (ii) (1 pt) Show that  $\{\vee, \wedge\}$  is *not* functionally complete.
- (iii) (4 pts) Suppose that the logical variables take on numerical values 0 and 1 as in the table above, and consider  $\varphi_i : \{0, 1\}^2 \rightarrow \{0, 1\}$  given by

$$\varphi_i(p, q) = \begin{cases} 1, & \text{if } a_i p + b_i q + c_i > 0 \\ 0, & \text{otherwise} \end{cases}, \quad a_i, b_i, c_i \in \mathbb{R}$$

Find valid  $a_i, b_i, c_i$  for each  $i = 0, \dots, 15$ .

### Exercise 1.3 (7 pts)

Let  $M$  be a set and let  $X, Y, Z, W \subset M$ . We define the *symmetric difference*:

$$X \triangle Y := (X - Y) \cup (Y - X)$$

- (i) (1 pt) Prove that  $X \triangle Y = (X \cup Y) - (X \cap Y)$ .
- (ii) (1 pt) Prove that  $(M - X) \triangle (M - Y) = X \triangle Y$ .
- (iii) (1 pt) Show that the symmetric difference is associative, i.e.,  $(X \triangle Y) \triangle Z = X \triangle (Y \triangle Z)$ .
- (iv) (1 pt) Prove that  $X \cap (Y \triangle Z) = (X \cap Y) \triangle (X \cap Z)$ .
- (v) (1 pt) Show that  $X \triangle Y = Z \triangle W$  iff  $X \triangle Z = Y \triangle W$ .
- (vi) (1 pt) Indicate the region of  $X \triangle Y \triangle Z$  in a Venn diagram.
- (vii) (1 pt) Sketch a Venn diagram for 4 (distinct) sets using circles (maybe with different radii).

### Exercise 1.4 (2 pts)

Let  $x, y \in \mathbb{R}$ , show that  $\forall x \exists y (xy = 0) \Leftrightarrow \exists y \forall x (xy = 0)$

利用真值表, 求命题的主析取范式和主合取范式. (主析取范式)

No 1. 写命题公式真值表

No 2 对于使A取0的指派, 写出对应的最大项,

No 3 A等值于所有最大项的合取.

$$\neg \varphi_{CNF} = (\neg a \wedge \neg b \wedge \neg c) \vee \dots \vee \dots$$

$$\varphi_{CNF} = \neg(\dots) \vee \neg(\dots) \vee \dots \vee \neg(\dots)$$

De Morgan's Law

$$C - (A \cup B) = (C - A) \cap (C - B)$$

$$C - (A \cap B) = (C - A) \cup (C - B)$$

# Exercise 1.1

$$\varphi = (A \rightarrow (B \rightarrow C)) \rightarrow (B \rightarrow (A \rightarrow C))$$

i) The truth table is as follows:

A	B	C	$B \rightarrow C$	$A \rightarrow C$	$A \rightarrow (B \rightarrow C)$	$B \rightarrow (A \rightarrow C)$	$\varphi$
0	0	0	1	1	1	1	1
0	0	1	1	1	1	1	1
0	1	0	0	1	1	1	1
0	1	1	1	1	1	1	1
1	0	0	1	0	1	1	1
1	0	1	1	1	1	1	1
1	1	0	0	0	0	0	0
1	1	1	1	1	1	1	1

$$\begin{aligned}
 \text{ii) } \varphi &= (A \rightarrow (B \rightarrow C)) \rightarrow (B \rightarrow (A \rightarrow C)) \\
 &= \neg(A \rightarrow (B \rightarrow C)) \vee (B \rightarrow (A \rightarrow C)) \\
 &= \neg(\neg A \vee (B \rightarrow C)) \vee (\neg B \vee (A \rightarrow C)) \\
 &= (A \wedge \neg(B \rightarrow C)) \vee (\neg B) \vee (\neg A) \vee C \\
 &= (A \wedge \neg(\neg B \vee C)) \vee (\neg B) \vee (\neg A) \vee C \\
 &= (A \wedge B \wedge \neg C) \vee (\neg B) \vee (\neg A) \vee C
 \end{aligned}$$

This is the disjunctive normal form of  $\varphi$

$$\begin{aligned}
 \text{iii) } \varphi &= (A \wedge B \wedge \neg C) \vee \neg(B \wedge A) \vee C \\
 &= (A \wedge B) \wedge (\neg C \vee \neg B \vee \neg A) \vee C \\
 &= \neg(\neg(A \wedge B \wedge \neg C) \wedge (B \wedge A)) \vee C \\
 &= \neg(\underbrace{\neg(A \wedge B)}_P \vee C \wedge \underbrace{(B \wedge A)}_Q) \vee C \\
 &= (A \wedge B) \wedge \neg(C \wedge (B \wedge A)) \vee C \\
 &= (A \wedge B) \wedge (\neg C \vee \neg(B \wedge A)) \vee C \\
 &= (A \wedge B) \wedge (\neg C \vee \neg B \vee \neg A) \vee C \\
 &= (C \vee (A \wedge B) \wedge (\neg C \vee \neg B \vee \neg A)) \\
 &= (C \vee (A \wedge B)) \wedge (C \vee \neg C \vee \neg B \vee \neg A) \\
 &= (C \vee A) \wedge (C \vee B) \wedge (C \vee \neg C \vee \neg B \vee \neg A)
 \end{aligned}$$

This is the conjunctive normal form of  $\varphi$ .

$$\begin{aligned}
\text{ii)} \quad \varphi &= A \wedge B \wedge \neg C \vee \neg(B \wedge A) \vee C. \\
&= \neg(\neg(A \wedge B \wedge \neg C) \wedge (B \wedge A)) \vee C. \\
&= \neg(\neg(A \wedge B) \vee C \wedge (B \wedge A)) \vee C. \\
&= (A \wedge B) \wedge \neg(C \wedge B \wedge A) \vee C. \\
&= (A \wedge B) \wedge (\neg C \vee \neg(B \wedge A)) \vee C. \\
&= (A \wedge B) \wedge (\neg C \vee \neg B \vee \neg A) \vee C. \\
&= C \vee ((A \wedge B) \wedge (\neg C \vee \neg B \vee \neg A)). \\
&= (C \vee (A \wedge B)) \wedge (C \vee \neg C \vee \neg B \vee \neg A). \\
&= (C \vee A) \wedge (C \vee B) \wedge (C \vee \neg C \vee \neg B \vee \neg A)
\end{aligned}$$

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## Exercise 2.1

$$i) \quad \varphi_0 = P \wedge \neg P$$

$$\varphi_1 = P \wedge Q$$

$$\varphi_2 = \neg(\neg P \vee Q) = P \wedge (\neg Q)$$

$$\varphi_3 = P$$

$$\varphi_4 = \neg(\neg Q \vee P) = Q \wedge (\neg P)$$

$$\varphi_5 = Q$$

$$\varphi_6 = (\neg P \wedge Q) \vee (P \vee \neg Q)$$

$$\varphi_7 = P \vee Q$$

$$\varphi_8 = \neg(P \vee Q)$$

$$\varphi_9 = (P \wedge Q) \vee (\neg Q \wedge \neg P)$$

$$\varphi_{10} = \neg Q$$

$$\varphi_{11} = \neg Q \vee P$$

$$\varphi_{12} = \neg P$$

$$\varphi_{13} = \neg P \vee Q$$

$$\varphi_{14} = \neg(P \wedge Q)$$

$$\varphi_{15} = P \vee \neg P$$

Therefore,  $\{\wedge, \vee, \neg\}$  is a functionally complete set.

ii) We know that when  $P, Q = 1$ ,  $P \vee Q, P \wedge Q = 1$ , if there doesn't exist  $\neg$ ,  $\varphi_i \neq 0$  when  $P, Q = 1$ .

Therefore, at least  $\varphi_0, \varphi_2, \varphi_4, \dots, \varphi_{14}$  cannot be represented by  $S$ .

Therefore,  $S$  is not functionally complete.

$$\psi_0(p, q) = \begin{cases} 1 & \text{if } a_0 p + b_0 q + c_0 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_0 = -1 \\ b_0 = -1 \\ c_0 = -1 \end{cases} \quad \psi_0(p, q) = \begin{cases} 1 & \text{if } a_0 p + b_0 q + c_0 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_0 = -1 \\ b_0 = -1 \\ c_0 = -1 \end{cases}$$

$$\psi_1(p, q) = \begin{cases} 1 & \text{if } a_1 p + b_1 q + c_1 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_1 = 1 \\ b_1 = 1 \\ c_1 = -1 \end{cases}$$

$$\psi_2(p, q) = \begin{cases} 1 & \text{if } a_2 p + b_2 q + c_2 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_2 = 2 \\ b_2 = -2 \\ c_2 = -1 \end{cases}$$

$$\psi_3(p, q) = \begin{cases} 1 & \text{if } a_3 p + b_3 q + c_3 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_3 = 3 \\ b_3 = -2 \\ c_3 = 0 \end{cases}$$

$$\psi_4(p, q) = \begin{cases} 1 & \text{if } a_4 p + b_4 q + c_4 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_4 = -2 \\ b_4 = 3 \\ c_4 = -1 \end{cases}$$

$$\psi_5(p, q) = \begin{cases} 1 & \text{if } a_5 p + b_5 q + c_5 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_5 = -2 \\ b_5 = 3 \\ c_5 = 0 \end{cases}$$

$$\psi_6(p, q) = \begin{cases} 1 & \text{if } a_6 p + b_6 q + c_6 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\Rightarrow \begin{cases} c \leq 0 \\ a_6 + c > 0 \\ b_6 + c > 0 \\ a_6 + b_6 + c \leq 0 \end{cases}$$

$$0 > a_6 + b_6 + c > -c \geq 0$$

therefore,  $a_6, b_6, c_6$

doesn't exist.

$$\psi_7(p, q) = \begin{cases} 1 & \text{if } a_7 p + b_7 q + c_7 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_7 = 2 \\ b_7 = 2 \\ c_7 = -1 \end{cases}$$

$$\psi_8(p, q) = \begin{cases} 1 & \text{if } a_8 p + b_8 q + c_8 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_8 = -2 \\ b_8 = -2 \\ c_8 = 1 \end{cases}$$

$$\psi_9(p, q) = \begin{cases} 1 & \text{if } a_9 p + b_9 q + c_9 > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} c_9 > 0 \\ a_9 + c_9 \leq 0 \\ b_9 + c_9 \leq 0 \\ a_9 + b_9 + c_9 > 0 \end{cases} \Rightarrow 0 > -c \geq a_9 + b_9 + c > 0$$

therefore,  $a_9, b_9, c_9$  doesn't exist.

$$\psi_{10}(p, q) = \begin{cases} 1 & \text{if } a_{10} p + b_{10} q + c_{10} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{10} = 2 \\ b_{10} = -100 \\ c_{10} = 2 \end{cases}$$

$$\psi_{11}(p, q) = \begin{cases} 1 & \text{if } a_{11} p + b_{11} q + c_{11} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{11} = 4 \\ b_{11} = -2 \\ c_{11} = 1 \end{cases}$$

$$\psi_{12}(p, q) = \begin{cases} 1 & \text{if } a_{12} p + b_{12} q + c_{12} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{12} = -100 \\ b_{12} = 1 \\ c_{12} = 1 \end{cases}$$

$$\psi_{13}(p, q) = \begin{cases} 1 & \text{if } a_{13} p + b_{13} q + c_{13} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{13} = -2 \\ b_{13} = 3 \\ c_{13} = 1 \end{cases}$$

$$\psi_{14}(p, q) = \begin{cases} 1 & \text{if } a_{14} p + b_{14} q + c_{14} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{14} = -1 \\ b_{14} = -1 \\ c_{14} = 3 \end{cases}$$

$$\psi_{15}(p, q) = \begin{cases} 1 & \text{if } a_{15} p + b_{15} q + c_{15} > 0 \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{cases} a_{15} = 1 \\ b_{15} = 1 \\ c_{15} = 1 \end{cases}$$

### Exercise 1.3

Let  $M$  be a set and let  $x, y, z, w \in M$ .  $x \Delta y := (x - y) \cup (y - x)$

(i) for any element  $a \in x \Delta y$

$a$  either belongs to  $x - y$  or  $y - x$

$\Rightarrow$  which means  $a \in x, a \notin y$  or  $a \in y, a \notin x$

for any element  $b \in (x \cup y) - (x \cap y)$

$b$  satisfies that  $b \in x \cup y$  and  $b \notin x \cap y$ .

for any  $a \in x \Delta y$ ,  $a$  satisfies  $a \in x \cup y$  and  $a \notin x \cap y$

for any  $b \in (x \cup y) - (x \cap y)$ ,  $b$  satisfies  $b \in x, b \notin y$  or  $b \in y, b \notin x$

Therefore  $x \Delta y = (x \cup y) - (x \cap y)$

(ii)  $(M - x) \Delta (M - y)$

$$= ((M - x) - (M - y)) \cup ((M - y) - (M - x))$$

for any element that belongs to  $(M - x) - (M - y)$ ,  $a \in M$  and  $a \in y$  and  $a \notin x$

for any element that belongs to  $(M - y) - (M - x)$ ,  $a \in M$  and  $a \in x$  and  $a \notin y$

then for any element  $t \in (M - x) \Delta (M - y)$ ,  $t \in y \setminus x$  or  $t \in x \setminus y$ ,

which is the definition of  $(x \Delta y)$  (since  $x, y \in M$ )

$x \Delta y$ , and that ends the proof.

$$\text{iii) } (x \Delta y) \Delta z = ((x - y) \cup (y - x)) \Delta z = ((x - y) \cup (y - x) - z) \cup$$

$$x \Delta (y \Delta z) = x \Delta ((y - z) \cup (z - y)) = (x - ((y - z) \cup (z - y))) \cup$$

$$((y - z) \cup (z - y) - x)$$

for any  $a \in (x \Delta y) \Delta z$ ,  $a \in x \setminus y \setminus z$  or  $a \notin x \Delta y \setminus z$  or  $a \in z, a \in x, a \in y$  or  $a \in z, a \notin x \setminus y$ .

for any  $b \in X \Delta (Y \Delta Z)$

$b$  satisfies that  $b \in X, b \in Y, b \in Z$  or  $b \in X, b \notin Y, b \notin Z$ .

or  $b \in Y, b \notin Z, b \notin X$  or  $b \in Z, b \notin Y, b \notin X$ .

therefore, for any  $a \in (X \Delta Y) \Delta Z$ , it also satisfies  $a \in X \Delta (Y \Delta Z)$ .

for any  $b \in X \Delta (Y \Delta Z)$ , it also satisfies  $b \in (X \Delta Y) \Delta Z$ .

so  $(X \Delta Y) \Delta Z = X \Delta (Y \Delta Z)$

iv)  $X \cap (Y \Delta Z) = (X \cap Y) \Delta (X \cap Z)$

for any  $a \in X \cap (Y \Delta Z)$ ,  $a \in X$  and  $a \in Y \Delta Z$

for any  $b \in (X \cap Y) \Delta (X \cap Z)$ ,  $b \in (X \cap Y)$  but  $b \notin (X \cap Z)$   
since  $b \in X$ , therefore  $b \notin Z$ .

$\Rightarrow b \in X, b \in Y, b \notin Z$ .

therefore  $a \in (X \cap Y) \Delta (X \cap Z)$ ,  $b \in X \cap (Y \Delta Z)$

proving that  $X \cap (Y \Delta Z) = (X \cap Y) \Delta (X \cap Z)$

v)  $X \Delta Y = Z \Delta W \Leftrightarrow X \Delta Z = Y \Delta W$

$X \Delta Y = Z \Delta W$

$\Leftrightarrow Z \Delta (X \Delta Y) = Z \Delta (Z \Delta W) = (Z \Delta Z) \Delta W = \emptyset \Delta W = W$

$\Leftrightarrow (Z \Delta X \Delta Y) \Delta Y = W \Delta Y$

$\Leftrightarrow Z \Delta X \Delta (Y \Delta Y) = W \Delta Y$

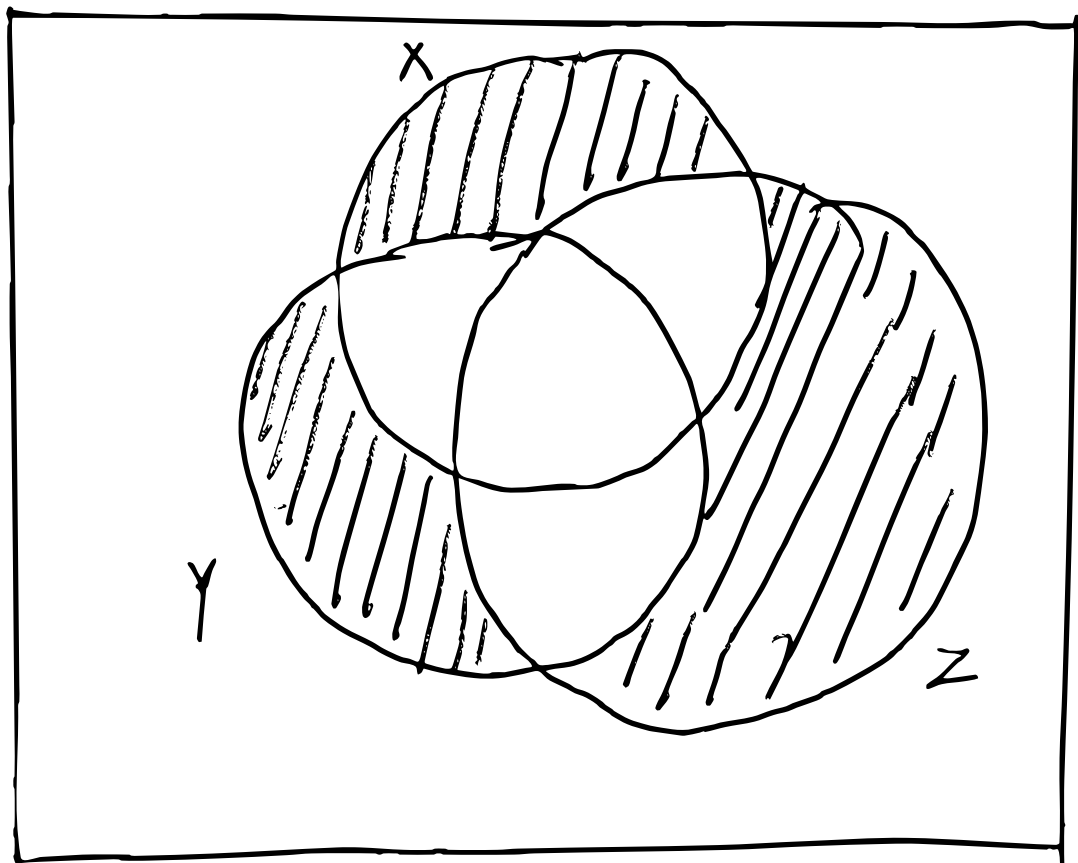
$\Leftrightarrow Z \Delta X = W \Delta Y$

$\Leftrightarrow X \Delta Z = Y \Delta W$

Thus the theorem is proved.

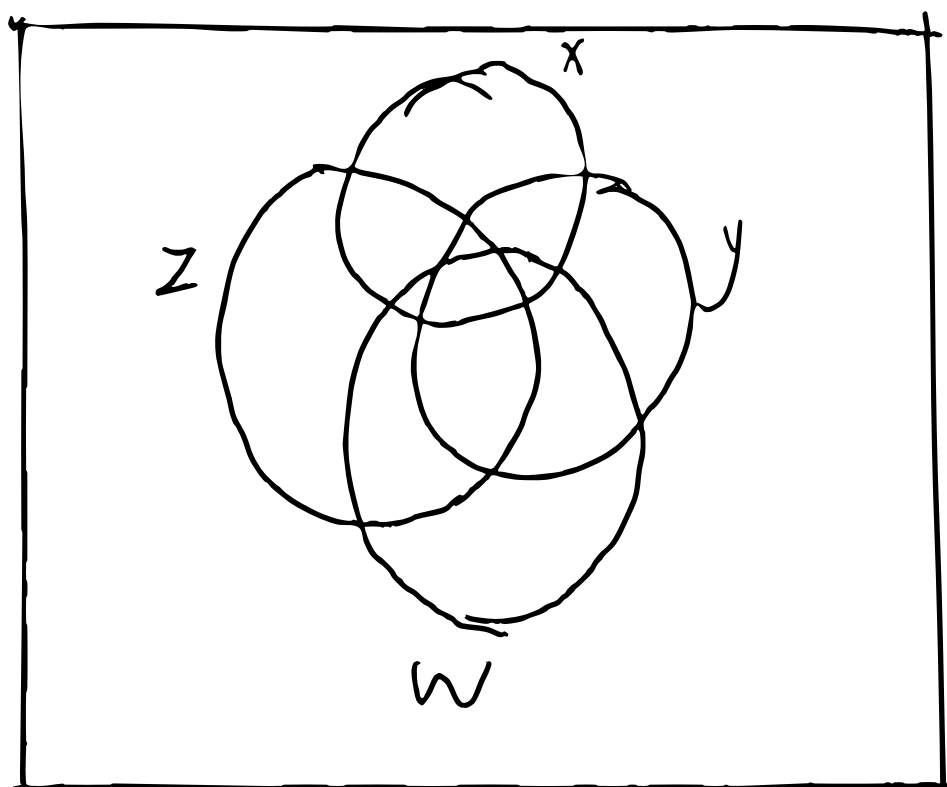


vi)



This is the Venn diagram of  $X \Delta Y \Delta Z$ .

vii)



# Exercise 4

$$\forall x \exists y (xy=0) \Leftrightarrow \exists y \forall x (xy=0)$$

We get  $\exists y \forall x (xy=0) \Rightarrow \forall x \exists y (xy=0)$  for free because.

(ii)

(i)

when we found such a  $y_0$  in (ii), this  $y_0$  can also be replaced in (i)

which can be viewed as a full blown version of  $\exists y (P_{x_1}(y) \wedge P_{x_2}(y)) \Rightarrow \exists y P_{x_1}(y) \wedge \exists y P_{x_2}(y)$ .

$$\exists y \forall x P(x, y) \rightarrow \forall x \exists y P(x, y)$$

$$\Leftrightarrow \exists y \forall x P(x, y) \rightarrow \forall z \exists w P(z, w)$$

$$\Leftrightarrow \forall y \forall z \exists x \exists w (P(x, y) \rightarrow P(z, w))$$

which is indeed true by taking  $x=z$  and  $w=y$ .

And  $\forall x \exists y (xy=0) \Rightarrow \exists y \forall x (xy=0)$  is also a simple proof.

because the we can find a fixed  $y=0$  in  $\forall x \exists y (xy=0)$  that is applicable in  $\exists y \forall x (xy=0)$ .

And that ends the proof.