

# Further linear algebra. Chapter I. Integers.

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Number theory is the theory of  $\mathbb{Z} = \{0, \pm 1, \pm 2, \dots\}$ .

## 1 Euclid's algorithm, Bézout's identity and the greatest common divisor.

- We say that  $a \in \mathbb{Z}$  *divides*  $b \in \mathbb{Z}$  iff there exists  $c \in \mathbb{Z}$  such that  $b = ac$ . We write  $a|b$ .
- A *common divisor* of  $a$  and  $b$  is an integer  $d$  that divides both  $a$  and  $b$ .
- Suppose for simplicity that  $a$  and  $b$  strictly positive. The *greatest common divisor* (sometimes called the *highest common factor*) of  $a$  and  $b$  is a common divisor  $d$  of  $a$  and  $b$  such that any other common divisor is smaller than  $d$ . This is written  $d = \gcd(a, b)$ .

Every  $a \in \mathbb{Z}$  is a divisor of 0, since  $0 = 0 \times a$ , therefore it makes sense to define  $\gcd(a, 0) = a$  if  $a > 0$ . However  $\gcd(0, 0)$  does not exist as any integer is a common factor of 0 and 0.

The following obvious remark : any divisor of  $a \geq 0$  is *smaller or equal* to  $a$ , is often used in the proofs.

Let  $a \geq b > 0$  be two integers. When  $a$  is not divisible by  $b$ , one can still divide with the *remainder*. For example:

$$a = 5, b = 2, a = 2b + 1$$

$$a = 20, b = 3, a = 6b + 2$$

etc..

This leads to the following (fundamental) theorem.

**Theorem 1.1 (Euclidean division)** Let  $a \geq b > 0$  be two integers. There exists a UNIQUE pair of integers  $(q, r)$  satisfying

$$a = qb + r$$

and  $0 \leq r < b$ .

**Proof.** Two things need to be proved : the existence of  $(q, r)$  and its uniqueness.

Let us prove the *existence*.

Consider the set

$$S = \{x, x \text{ integer } \geq 0 : a - xb \geq 0\}$$

The set  $S$  is not empty : 1 belongs to  $S$ . The set  $S$  is bounded : any element  $x$  of  $S$  satisfies  $x \leq \frac{a}{b}$ . Therefore,  $S$  being a bounded set of positive integers,  $S$  is finite and hence contains a maximal element. Let  $q$  be this maximal element and let  $r := a - qb$ .

We need to prove that  $0 \leq r < b$ . By definition  $r \geq 0$  (it belongs to  $S$ ). To prove that  $r < b$ , let us argue by contradiction. Suppose that  $r \geq b$ . As  $r = a - qb$ , we get

$$a - (q + 1)b \geq 0$$

This means that  $q + 1 \in S$  but  $q + 1 > q$ . This contradicts the *maximality* of  $q$ . Therefore  $r < b$  and the existence is proved.

Let us now prove the *uniqueness*.

Again we argue by contradiction. Suppose that there exists a pair  $(q', r')$  satisfying  $a = q'b + r'$  with  $0 \leq r' < b$  and such that  $q' \neq q$ . By subtracting the inequality  $0 \leq r < b$  to this inequality, we get  $-b < r' - r < b$  i.e.

$$|r - r'| < b$$

Now by subtracting  $a = q'b + r'$  to  $a = qb + r$  and taking the modulus, we get

$$|r - r'| = |q - q'|b$$

By assumption  $q \neq q'$ , hence  $|q' - q| \geq 1$  and we get the inequality

$$|r - r'| \geq b$$

The two inequalities satisfied by  $r - r'$  contradict each other, hence  $q = q'$ . Now  $|r - r'| = |q - q'|b = 0$ , hence  $r = r'$ . The uniqueness is proved.  $\square$

**Theorem 1.2** Let  $a \geq b > 0$  be two integers and  $(q, r)$  such that

$$a = bq + r, \quad 0 \leq r < b$$

Then

$$\gcd(a, b) = \gcd(b, r).$$

**Proof.** Let  $A := \gcd(a, b)$  and  $B := \gcd(b, r)$ . As  $r = a - bq$  and  $A$  divides  $a$  and  $b$ ,  $A$  divides  $r$ . Therefore  $A$  is a common factor of  $b$  and  $r$ . As  $B$  is the highest common divisor of  $b$  and  $r$ ,  $A \leq B$ .

In exactly the same way, one proves (left to the reader), that  $B \leq A$  and therefore  $A = B$ .  $\square$

This leads to the following algorithm (the so-called **Euclid's algorithm**).

Let  $a \geq b > 0$  be two integers. We wish to calculate  $\gcd(a, b)$ .

The method is this: Set  $r_1 = a$  and  $r_2 = b$ . We have

$$r_1 = r_2q_1 + r_3$$

with, by the above proposition,  $\gcd(a, b) = \gcd(r_1, r_2) = \gcd(r_2, r_3)$  with  $0 \leq r_3 < r_2$ .

- If  $r_3 = 0$ , then  $\gcd(r_2, r_3) = r_2$  and we are done.
- If  $r_3 \neq 0$ , then divide  $r_2$  by  $r_3$ :

$$r_2 = r_3q_2 + r_4$$

with  $0 \leq r_4 < r_3$ . Again, if  $r_4 = 0$ , then  $\gcd(a, b) = r_3$ , otherwise carry on...

This way one constructs a sequence:

$$r_i = r_{i+1}q_i + r_{i+2}$$

where  $0 \leq r_{i+2} < r_{i+1}$ .

Notice that  $r_{i+2}$  goes **strictly down** hence one **must** at some point find  $r_{i+2} = 0$  and then  $\gcd(a, b) = r_{i+1}$ .

**Remark 1.3** When performing Euclid's algorithm, be very careful not to divide  $q_i$  by  $r_i$ . This is a mistake very easy to make.

**Example 1.4** Take  $a = 27$  and  $b = 7$ . We have

$$\begin{aligned} 27 &= 3 \times 7 + 6r_1 = 27, \quad r_2 = 7, r_3 = 6 \\ 7 &= 1 \times 6 + 1, \quad r_3 = 6, r_4 = 1 \\ 6 &= 6 \times 1 + 0 \quad r_5 = 0 \end{aligned}$$

Therefore

$$\gcd(27, 7) = \gcd(7, 6) = \gcd(6, 1) = \gcd(1, 0) = 1.$$

Another example:

$$\begin{aligned} 555 &= 155 \cdot 3 + 90 \\ 155 &= 90 \cdot 1 + 65 \\ 90 &= 65 \cdot 1 + 25 \\ 65 &= 25 \cdot 2 + 15 \\ 15 &= 10 \cdot 1 + 5 \\ 10 &= 5 \cdot 2 + 0 \\ \gcd(555, 155) &= 5 \end{aligned}$$

Euclid's algorithm is and **algorithm** meaning that no matter what the initial data is, it will yield a gcd in a finite number of steps.

It is easy to implement on a computer. Suppose that you have some standard computer language (Basic, Pascal, Fortran,...) and that it has an instruction  $r := a \bmod b$  which returns the remainder of the Euclidean division of  $a$  by  $b$ .

The implementation of the algorithm would be something like this:

```
Procedure gcd( $a, b$ )
  If  $a < b$  then
    Swap( $a, b$ )
  While  $b \neq 0$ 
    Begin
       $r := a \bmod b$ 
       $a := b$ 
       $b := r$ 
    End
  Return  $a$ 
End
```

The following lemma is very important for what will follow. It is essentially ‘the Euclid’s algorithm’ run backwards.

**Theorem 1.5 (Bézout’s Identity)** *As usual, let  $a \geq b > 0$  be integers. Let  $d = \gcd(a, b)$ . Then there are integers  $h, k \in \mathbb{Z}$  such that*

$$d = ha + kb.$$

Note that in this lemma, the integers  $h$  and  $k$  are not positive, in fact exactly one of them is negative or zero. Prove it !

**Proof.** Consider the sequence given by Euclid’s algorithm:

$$r_i = r_{i+1}q_i + r_{i+2}$$

where  $0 \leq r_{i+2} < r_{i+1}$  with  $r_1 = a, r_2 = b$ .

We will show that each  $r_i$  can be expressed as  $h_i a + k_i b$  with  $h_i, k_i \in \mathbb{Z}$ . In particular, as by Euclid’s algorithm,  $\gcd(a, b)$  is some  $r_i$ , the result will follow.

This is certainly true for  $i = 1, 2$  since  $r_1 = 1 \times a + 0 \times b$  and  $r_2 = 0 \times a + 1 \times b$ .

For the inductive step, assume it is the case for  $r_{i-1}$  and  $r_{i-2}$ , i.e.

$$r_{i-1} = ha + bk, \quad r_{i-2} = h'a + k'b.$$

We have

$$r_{i-2} = q_{i-2}r_{i-1} + r_i.$$

Therefore

$$r_i = h'a + k'b - q_{i-2}(ha + kb) = (h' - q_{i-2}h)a + (k' - q_{i-2}k)b.$$

□

**Example 1.6** *Again we take  $a = 27$  and  $b = 7$ .*

$$27 = 3 \times 7 + 6$$

$$7 = 1 \times 6 + 1$$

$$6 = 6 \times 1 + 0.$$

*Therefore*

$$1 = 7 - 1 \times 6$$

$$= 7 - 1 \times (27 - 3 \times 7)$$

$$= 4 \times 7 - 1 \times 27.$$

*So we take  $h = -1$  and  $k = 4$ .*

Another example

**Example 1.7** Take  $a = 819$  and  $b = 165$ .

$$\begin{aligned}819 &= 165 \times 4 + 159 \\165 &= 159 \times 1 + 6 \\6 &= 3 \times 2\end{aligned}$$

Therefore

$$\begin{aligned}3 &= 159 + (-26)6 = \\(819 + 165(-4)) + (-26)(165 + 159(-1)) &= \\819 + 165(-30) + 159(26) &= \\819 + 165(-30) + (819 + 165(-4))(26) &= \\819(27) + 165(-134)\end{aligned}$$

So we take  $h = 27$  and  $k = -134$ .

**Définition 1.1** Two integers  $a$  and  $b$  are **coprime** if  $\gcd(a, b) = 1$ .

**Proposition 1.8**  $a$  and  $b$  are coprime if and only if there exist integers  $k$  and  $h$  such that  $ha + kb = 1$ .

**Proof.** If  $a$  and  $b$  are coprime, then it's just the Bézout's identity. Suppose that there exist integers  $k$  and  $h$  such that  $ha + kb = 1$ . Let  $d = \gcd(a, b)$ . Then  $d$  divides  $ha + kb$ , hence  $d$  divides 1, hence  $d = 1$ .  $\square$

For example, for any positive integer  $k$ ,  $6 \cdot (7k + 6) + (-7) \cdot (6k + 5) = 1$ , hence  $\gcd(7k + 6, 6k + 5) = 1$ .

One has the following properties of the gcd.

- if  $ha + kb = m$  for some integers  $h, k$ , then  $\gcd(a, b)$  divides  $m$ .
- $\gcd(ca, cb) = c \gcd(a, b)$
- If  $d$  divides  $a$  and  $b$ , then  $\gcd(a/d, b/d) = \gcd(a, b)/d$ . In particular, if  $d = \gcd(a, b)$ , then  $\gcd(a/d, a/b) = 1$ .

The first property is obvious. For the second, write the Bézout's identity  $ha + kb = d$  with  $d = \gcd(a, b)$ . Hence  $cd = hac + kbc$  and hence any common divisor of  $ac$  and  $bc$  divides  $cd$  and hence smaller or equal to  $cd$ . In addition,  $cd$  divides  $ac$  and  $bc$  hence  $\gcd(ca, cb) = cd$ . The next property follows from this.

We now apply the Euclid's algorithm and Bézout's identity to the solution of *linear diophantine equations*.

Let  $a, b, c$  be three positive integers. A linear diophantine equation (in two variables) is the equation

$$ax + by = c$$

A solution is a pair  $(x, y)$  of integers (not necessarily positive) that satisfy this relation.

Such an equation may or may not have solutions. For example, consider  $2x + 4y = 5$ . Quite clearly, if there was a solution, then 2 will divide the right hand side, which is 5. This is not the case, therefore, this equation does not have a solution.

On the other hand, the equation  $2x + 4y = 6$  has many solutions :  $(1, 1), (5, -1), \dots$ . This suggests that the existence of solutions depends on whether or not  $c$  is divisible by the  $\gcd(a, b)$  and that if such is the case, there are many solutions to the equation. This indeed is the case, as shown in the following theorem.

**Theorem 1.9** *Suppose that  $a|bc$  and  $a$  and  $b$  are coprime. Then  $a|c$ .*

*In particular, if  $p$  is a prime and  $p|ab$  then  $p|a$  or  $p|b$ .*

**Proof.**  $a$  and  $b$  are coprime, hence there exist  $h$  and  $k$  such that  $ha + kb = 1$ . Multiply by  $c$  and get  $c = hac + kbc$ .  $a$  divides  $ac$  and  $bc$ , hence  $a$  divides the right hand side. It follows that  $a$  divides  $c$ .

The second statement follows trivially as if  $p$  does not divide  $a$ , then  $a$  and  $p$  are coprime.  $\square$

**Theorem 1.10 (Solution to linear diophantine equations)** *Let  $a, b, c$  be three positive integers, let  $d := \gcd(a, b)$  and consider the equation*

$$ax + by = c$$

1. *This equation has a solution if and only if  $d$  divides  $c$*

2. Suppose that  $d|c$  and let  $(x_0, y_0)$  be a solution. The set of all solutions is  $(x_0 + n\frac{b}{d}, y_0 - n\frac{a}{d})$  where  $n$  runs through the set of all integers (positive and negative).

**Proof.** For the ‘if’ part : Suppose there is a solution  $(x, y)$ . Then  $d$  divides  $ax + by$ . But, as  $ax + by = c$ ,  $d$  divides  $c$ .

For the ‘only if’ part : Suppose that  $d$  divides  $c$  and write  $c = dm$  for some integer  $m$ . By Bézout’s lemma there exist integers  $h, k$  such that

$$d = ha + kb$$

Multiply this relation by  $m$  and get

$$c = dm = (mh)a + (mk)b$$

This shows that  $(x_0 = mh, y_0 = mk)$  is a solution to the equation. That finishes the ‘only if’ part.

Let us now suppose that the equation has a solution (in particular  $d$  divides  $c$ )  $(x_0, y_0)$ . Let  $(x, y)$  be any other solution. Subtract  $ax + by = c$  from  $ax_0 + by_0 = c$  to get

$$a(x_0 - x) + b(y_0 - y) = 0$$

Divide by  $d$  to get

$$\frac{a}{d}(x_0 - x) = -\frac{b}{d}(y_0 - y)$$

This relation shows that  $\frac{a}{d}$  divides  $\frac{b}{d}(y_0 - y)$  but the integers  $\frac{a}{d}$  and  $\frac{b}{d}$  are coprime hence  $\frac{a}{d}$  divides  $y_0 - y$  (by 1.9)

Therefore, there exists an integer  $n$  such that

$$y = y_0 - n\frac{a}{d}$$

Now plug this into the equality  $\frac{a}{d}(x_0 - x) = -\frac{b}{d}(y_0 - y)$  to get that

$$x = x_0 + n\frac{b}{d}$$

□

The proof of this theorem gives a *procedure* for finding solutions, it is as follows:



1. Calculate  $d = \gcd(a, b)$ . If  $d$  does not divide  $c$ , then there are no solutions and you're done. If  $d$  divides  $c$ ,  $c = md$  then there are solutions.
2. Run Euclid's algorithm backwards to find  $h, k$  such that  $d = ha + kb$ . Then  $(x_0 = mh, y_0 = mk)$  is a solution.
3. All solutions are

$$(x_0 + n\frac{b}{d}, y_0 - n\frac{a}{d})$$

where  $n$  runs through all integers.

**Example 1.11** Take  $a = 27, b = 7, c = 5$ . We have found that  $\gcd(a, b) = 1$  (in particular there will be solutions with any  $c$ ) and that  $1 = 4 \times 7 - 1 \times 27$  hence  $h = -1$  and  $k = 4$ .

Our procedure gives a particular solution :  $(-5, 20)$  and the general one  $(-5 + 7n, 20 - 27n)$ .

Take  $a = 666, b = 153, c = 43$ . We have found that  $\gcd(a, b) = 9$ , it does not divide 43, hence no solutions.

Take  $c = 45 = 5 \times 9$ . There will be solutions. We had  $9 = 3 \times 666 - 13 \times 153$ . A particular solution is  $(15, -65)$  and the general one is  $(15 + 17n, -65 - 74n)$ .

(in particular there will be solutions with any  $c$ ) and that  $1 = 4 \times 7 - 1 \times 27$  hence  $h = -1$  and  $k = 4$ .

Our procedure gives a particular solution :  $(-5, 20)$  and the general one is  $(-5 + 7n, 20 - 27n)$ .

## 2 Factorisation into primes.

**Définition 2.1** An integer  $p \geq 2$  is prime iff the only divisors of  $p$  are 1 and  $p$ .

**Lemma 2.1** If  $p|a_1 \cdots a_n$  then there exists  $1 \leq i \leq n$  such that  $p|a_i$ .

**Proof.** One proceeds by induction. True for  $i = 1, 2$  so suppose true for  $n - 1$  and suppose that  $p|a_1 \cdots a_n$ . Let  $A = a_1 \cdots a_{n-1}$  and  $B = a_n$  then  $p|AB$  implies  $p|A$  or  $p|B$ . In the latter case we are done and in the former case the inductive hypothesis implies that  $p|a_i$  for some  $1 \leq i \leq n - 1$ .  $\square$

**Theorem 2.2 (Unique Factorisation Theorem)** *If  $a \geq 2$  is an integer then there are primes  $p_i > 0$  such that*

$$a = p_1 p_2 \cdots p_s.$$

*Moreover this factorisation is unique in the sense that if*

$$a = q_1 q_2 \cdots q_t$$

*for primes  $q_j > 0$  then*

$$s = t$$

*and*

$$\{p_1, \dots, p_s\} = \{q_1, \dots, q_s\}$$

*(equality of sets) In other words, the  $p_i$ s and the  $q_i$ s are the same prime numbers up to reordering.*

**Proof.** For existence suppose the result does not hold. Then there an integer which can not be written as a product of primes. Among all those integers, there is a smallest one (the integers under consideration are greater than two !). Let  $a$  be this smallest integer which is not a product of primes. Certainly  $a$  is not prime so  $a = bc$  with  $1 < b, c < a$ . As  $b$  and  $c$  are strictly smaller than  $a$ , they are products of primes. Write

$$b = p_1 \cdots p_k$$

and

$$c = p_{k+1} \cdots p_l$$

hence

$$a = p_1 \cdots p_l,$$

This contradicts the definition of  $a$  hence the factorisation exists.

For uniqueness suppose that we have an example where there are two distinct factorisations. Again we can choose a *smallest* integer with two different factorisations

$$a = p_1 \cdots p_s = q_1 \cdots q_t.$$

Then  $p_1 | q_1 \cdots q_t$  so by lemma 2.1 we have  $p_1 | q_j$  for some  $1 \leq j \leq t$  then since  $p_1$  and  $q_j$  are primes we have  $p_1 = q_j$ . But then dividing  $a$  by  $p_1$  we have a smaller integer with two distinct factorisations, a contradiction.  $\square$

**Remark 2.3** Of course, the primes in the factorisation  $a = p_1 \cdots p_s$  need not be distinct. For example :  $4 = 2^2$ , here  $p_1 = p_2 = 2$ . Similarly  $8 = 2^3$ ,  $p_1 = p_2 = p_3 = 2$ . Also  $12 = 3 \times 2^2$ ,  $p_1 = 3, p_2 = p_3 = 2$

In fact we have that for any integer  $a \geq 2$ , there exist  $s$  distinct primes  $p_1, \dots, p_s$  and  $t$  integers  $e_i \geq 1$  such that

$$a = p_1^{e_1} \cdots p_s^{e_t}$$

Examples of factorisations:

$$1000 = 2^3 \times 5^3$$

$$144 = 2^4 \times 3^2$$

$$975 = 2^3 \times 5^3$$

Factoring a given integer is hard as there is no procedure like Euclidean algorithm. One usually does it by trial and error. The following trivial lemma helps.

**Lemma 2.4 (Square root test)** Let  $n$  be a composite (not prime) integer. Then  $n$  has a prime divisor  $\leq \sqrt{n}$ .

**Proof.** Write  $n = ab$  with  $1 < a, b < n$ . Suppose that  $a \geq \sqrt{n}$ , then  $n = ab \geq \sqrt{n}b$  hence  $b \leq \sqrt{n}$  and therefore any prime divisor of  $b$  is  $\leq \sqrt{n}$ .  $\square$

For example, suppose you were to factor 3372. Clearly it's divisible by 2 :  $3372 = 2 \times 1686$ . Now, 1686 is again divisible by two :  $1686 = 2 \times 843$  and  $3372 = 2^2 \times 843$ . Now we notice that 3 divides  $843 = 3 \times 281$ . Now the primes  $< \sqrt{281}$  are 2, 3, 5, 7, 11, 13 and 281 is not divisible by any of these. Hence 281 is prime and we get a factorisation:

$$3372 = 2^2 \cdot 3 \cdot 281$$

How many primes there are ? Here is the answer.

**Theorem 2.5 (Euclid's Theorem)** There exist infinitely many primes.

**Proof.** Suppose not, let  $p_1, \dots, p_n$  be all the primes there are. Consider  $Q = p_1 p_2 \cdots p_n + 1$ . Since  $Q$  has a prime factorisation, there is a prime  $p$  that divides  $Q$ . This prime  $p$  has to belong to our list, after reordering we can assume that  $p = p_1$ . Then  $p_1$  divides  $Q - p_1 \cdots p_n = 1$  which is not possible because  $p_1$  is prime.  $\square$

The idea we used here is this : suppose the set of all primes is finite, we *construct* an integer that is not divisible by any of the primes from this set. This is a contradiction.

Can we use the same idea to prove that there are infinitely many primes of a certain form? In some cases yes.

Quite clearly Euclid's theorem shows that there are infinitely many *odd* primes since the only even prime is 2. Put in another way, it shows that there are infinitely many primes of the form  $2k + 1$ .

Let's look at primes of the form  $4k + 3$ . Are there infinitely many of them ?

Suppose there are finitely many and list them  $p_1, \dots, p_r$ . Note that  $p_1 = 3$ . Consider  $Q = 4p_2 \cdots p_r + 3$  (note that we started at  $p_2$  !!!).

The integer  $Q$  is clearly not divisible by 3 (otherwise 3 would divide  $p_2 \cdots p_r$  and  $p_i \neq 3$  for all  $i > 1$ ).

None of the  $p_i$ ,  $i > 2$  divides  $Q$ . Indeed suppose some  $p_i$ ,  $i > 2$  divides  $Q$ . Then

$$4p_2 \cdots p_r + 3 = p_i k$$

which shows that  $p_i$  divides 3 which is not the case.

To get a contradiction, we need to prove that  $Q$  is divisible by a prime of the form  $4k + 3$ , for it will have necessarily be one of the  $p_i$ s and they do not divide  $Q$ .

This is precisely what we are proving.

**Lemma 2.6** *Every integer of the form  $4k + 3$  has a prime factor of the form  $4k + 3$ .*

**Proof.** Let  $N = 4k + 3$ .

The smallest positive integer of this form is 3 which is prime, hence the property holds for 3.

Suppose that the property holds for all integers  $< N$  of the form  $4k + 3$ . If  $N$  is prime, then take for the factor  $N$  itself.

We can and do assume that  $N$  is composite. Write  $N = N_1 N_2$  with  $1 < N_i < N$ . As  $N$  is odd,  $N_1$  and  $N_2$  are odd. Any odd number is of the form  $4k + 1$  or  $4k + 3$ .

Suppose  $N_1 = 4a + 1$  and  $N_2 = 4b + 1$ . Then  $N = N_1 N_2 = (4a + 1)(4b + 1) = 4(4ab + a + b) + 1$  is of the form  $4k + 1$  which contradicts the fact that  $N$  is of the form  $4k + 3$ . Hence one of the  $N_i$ s,  $N_1$  say has a prime factor of the form  $4k + 3$ . As  $N_1 < N$ , by induction assumption,  $N_1$  and hence  $N$  has a prime factor of the form  $4k + 3$ . This finishes the proof.  $\square$

Note that the proof does not work if you try to prove that there are infinitely many primes of the form  $4k + 1$ . This is where it fails. The first prime of this form is  $5 = 4 \times 1 + 1$  but when you try to construct your  $Q$ , you get  $Q = 4 \times 5 + 1 = 21 = 3 \times 7$ . The divisors of  $Q$  are of the form  $4k + 3$ , not  $4k + 1$ ....

In other words, the method fails because the divisors of  $Q$  can have no divisor of the form  $4k + 1$ .

It is however true that there are infinitely many primes of the form  $4k + 1$ , in fact, there is the following spectacular theorem :

**Theorem 2.7 (Dirichlet's theorem on primes in arithmetic progressions)**

*Let  $a$  and  $d$  be two coprime integers. There exist infinitely many primes of the form  $a + kd$ .*

The proof of this theorem is well beyond the scope of this course.

## 2.1 Congruences.

We define  $a \equiv b \pmod{m}$  iff  $m \mid (a - b)$  in other words iff there exists an integer  $k \in \mathbb{Z}$  such that  $a = b + km$ .

We say  $a$  is *congruent* to  $b$  modulo  $m$ .

The *congruency class* of  $a$  is the set of numbers congruent to  $a$  modulo  $m$ . This is written  $[a]$ . In other words

$$[a] = \{a + km : k \in \mathbb{Z}\}$$

Every integer is congruent to one of the numbers  $0, 1, \dots, m - 1$  (can be seen using Euclidean division), so the set of all congruency classes is

$$\mathbb{Z}/m\mathbb{Z} = \{[0], \dots, [m - 1]\}$$

Ex. Take  $m = 3$ , then  $[8] = [5] = [2] = [-1] = [-4] = \dots$

For an integer  $k$ ,  $4k + 1 \equiv 1 \pmod{4}$ ,  $4k + 3 \equiv 3 \pmod{4}$  and  $4k \equiv 0 \pmod{4}$ .

An integer is *even* if and only if it is zero mod 2. An integer is *odd* if and only if it is one mod 2.

Let  $a \geq b$  be two positive integers and let  $(q, r)$  be such that  $a = bq + r$ . Then  $a \equiv r \pmod{b}$ . It may help to think of congruences as the remainders of the Euclidean division.

Another trivial but useful observation is that if  $a \equiv b \pmod{m}$  and  $d$  divides  $a, b$  and  $m$ , then  $\frac{a}{d} \equiv \frac{b}{d} \pmod{\frac{m}{d}}$ .

**Proposition 2.8** *If  $a \equiv b \pmod{m}$  then  $b \equiv a \pmod{m}$ .*

*If  $a \equiv a' \pmod{m}$  and  $b \equiv b' \pmod{m}$  then  $a + b \equiv a' + b' \pmod{m}$  and  $ab \equiv a'b' \pmod{m}$ .*

**Proof.** trivial □

We can rewrite this proposition by simply saying:

$$[a] + [b] = [a + b] \text{ and } [a][b] = [ab]$$

The proposition says that these operations  $+$  and  $\times$  are well defined operations on  $\mathbb{Z}/m\mathbb{Z}$ .

Ex. Write down addition and multiplication tables in  $\mathbb{Z}/3\mathbb{Z}$ ,  $\mathbb{Z}/4\mathbb{Z}$  and  $\mathbb{Z}/6\mathbb{Z}$ .

By an *inverse* of  $a$  modulo  $m$  we mean a number  $c$  such that  $ac \equiv 1 \pmod{m}$ . This is written  $c \equiv a^{-1} \pmod{m}$ .

An element may or may not have an inverse mod  $m$ .

Take  $m = 6$ .  $[5]$  has an inverse in  $\mathbb{Z}/6\mathbb{Z}$  :

$$[5] \times [5] = [25] = [1]$$

While  $[3]$  does not have an inverse : in  $\mathbb{Z}/6\mathbb{Z}$  we have  $[3][2] = [6] = [0]$ . So if  $[3]$  had an inverse, say  $[a]$ , we would have  $[3][a] = [1]$ , and by multiplying by  $[2]$  we would get  $[0] = [2]$  which is not the case.

This suggests that the existence of the inverse of  $a \pmod{m}$  has something to do with common factors of  $a$  and  $m$ . This is indeed the case as shown in the following lemma.

**Lemma 2.9** *An integer  $a$  has an inverse modulo  $m$  if and only if  $a$  and  $m$  are coprime ( $\gcd(a, m) = 1$ ).*

**Proof.** The integer  $a$  has an inverse mod  $m$  if and only if the equation

$$ax + my = 1$$

has a solution. This equation has a solution if and only if  $\gcd(a, m)$  divides 1 which is only possible if  $\gcd(a, m) = 1$ .  $\square$

As usual, the proof of the lemma gives a procedure for finding inverses. Use Euclidean algorithm to calculate  $\gcd(a, m)$ . If it's not one, there is no inverse. If it is one run the algorithm backwards to find  $h$  and  $k$  such that  $ah + mk = 1$  and

$$[a]^{-1} = [h]$$

Notice in particular that if  $p$  is a prime number, then any class  $[x] \neq [0]$  is invertible. The set  $\mathbb{Z}/p\mathbb{Z}$  is a **field** and it is denoted  $\mathbb{F}_p$ .

**Example 2.10** Find  $43^{-1} \bmod 7$ .

*Euclid's algorithm :*

$$43 = 6 \times 7 + 1$$

*They are coprime and  $1 = -6 \times 7 + 1 \times 43$ . Hence  $43^{-1} = 1 \bmod 7$   
Same with  $32^{-1} \bmod 7$ .*

$$32 = 4 \times 7 + 4$$

$$7 = 1 \times 4 + 3$$

$$4 = 1 \times 3 + 1$$

*And*

$$1 = (1 \times 4) + (-1 \times 3) = (-1 \times 7) + (2 \times 4) = (2 \times 32) + (-9 \times 7) = (-9 \times 7) + (2 \times 32)$$

*Hence  $32^{-1} = 2 \bmod 7$ .*

*Same with  $49^{-1} \bmod 15$ .*

$$49 = 3 \times 15 + 4$$

$$15 = 3 \times 4 + 3$$

$$4 = 1 \times 3 + 1$$

And get

$$1 = (1*4)+(-1*3) = (-1*15)+(4*4) = (4*49)+(-13*15) = (-13*15)+(4*49)$$

Hence  $49^{-1} \bmod 15 = 4$ .

**The equation**  $ax \equiv b \pmod{m}$ .

More generally, suppose we want to solve an equation

$$ax \equiv b \pmod{m}$$

By this we mean, find all integers  $x \pmod{m}$  that satisfy the equation.

The equation is equivalent to the existence of an integer  $y$  such that

$$ax + my = b$$

And we know how to solve this !

This equation has a solution if and only if  $d = \gcd(a, m)$  divides  $b$  and we know how to find all the solutions.

In particular, the equation has solutions if and only if  $d$  divides  $b$ .

If this is the case, then to solve the equation, divide it by  $d$ , let  $a' = a/d, m' = m/d, b' = b/d$ . Write  $a'x + m'y = b'$ . Bézout's identity gives  $(h, k)$  such that  $a'h + m'k = 1$ .

By the theorem on solutions of linear diophantine equations, all values of  $x$  are  $\{b'h + nm'\}$  and the solutions of the equation are the  $\{[b'h + nm']\}$ . Notice that there are exactly  $d$  of them.

Let's see a few examples.

$$2x \equiv 4 \pmod{10}.$$

We have  $\gcd(2, 10) = 2$ , it divides 4, there are solutions. Dividing by 2 we get  $x \equiv 2 \pmod{5}$  i.e  $x = 2 + 5n$ .

Now the solutions are  $\{[2], [7]\}$  (classes  $\pmod{10}$ ).

The equation  $2x \equiv 4 \pmod{5}$  has no solutions.

Another example:

$$3x \equiv 6 \pmod{18}$$

$\gcd(3, 18) = 3$  divides 6. We find  $x \equiv 2 \pmod{6}$ . Solutions are  $\{[2], [8], [14]\}$ .

Another:  $10x \equiv 14 \pmod{18}$ .

We have  $\gcd(10, 18) = 2$ , divides 14, we'll find 2 solutions.

Euclid's algorithm gives:

$$18 = 10 + 8$$

$$10 = 8 + 2$$

$$8 = 4 \times 2 + 0$$



and Bézout's identity:

$$2 = 10 + (-1) \times 8 = (-1) \times 18 + 2 \times 10$$

hence

$$14 = -7 \times 18 + 14 \times 10$$

The general solution is  $x = 14 + 9n$ .

The solutions to the congruence are  $\{[14], [5]\}$ .

Notice that when  $a$  and  $m$  are coprime, then there is a unique solution and it is given by  $[a]^{-1}[b]$ .

For example, solve  $99x \equiv 100 \pmod{101}$ .

99 and 101 are coprime, hence there is a unique solution.

Euclid's algorithm gives:

$$101 = 99 + 2$$

$$99 = 2 \times 49 + 1$$

$$2 = 1 \times 2 + 0$$

Bézout's identity:

$$1 = (1 * 99) + (-49 * 2) = (-49 * 101) + (50 * 99)$$

One finds  $[99]^{-1} = [50]$ . The unique solution is  $[50] \times [100] = [5000] = [51]$ .

**Corollary 2.11**  $\mathbb{F}_p^\times = \{[1], [2], \dots, [p-1]\}$  is a group with the operation of multiplication.

**Proof.** A group is a set with a binary operation (in this case multiplication), such that (i) the operation is associative; (ii) there is an identity element; (iii) every element has an inverse. Clearly  $[1]$  is the identity element, and the every element has an inverse because  $1, 2, 3, \dots, p-1$  are coprime with  $p$ .  $\square$

Recall that Lagrange's theorem states that if  $G$  is a finite group and  $H$  is a subgroup, then  $|H|$  divides  $|G|$ . The corollary of this theorem is that if  $a \in G$  and  $k \geq 0$  is the smallest integer such that  $a^k = 1$ , then  $k$  divides  $|G|$ .

**Theorem 2.12 (Fermat's Little Theorem)** If  $p$  is prime and  $a \in \mathbb{Z}$  then

$$a^p \equiv a \pmod{p}.$$

Hence if  $p \nmid a$  then  $a^{p-1} \equiv 1 \pmod{p}$ .

**Proof.** If  $p|a$  then  $a \equiv 0 \pmod{p}$  and  $a^p \equiv 0 \pmod{p}$  so suppose  $p$  does not divide  $a$ , and so  $a \in \mathbb{F}_p^\times$ . Recall that by a corollary to Lagrange's Theorem, the order of an element of a group divides the order of the group. Let  $n$  be the order of  $a$ , so  $a^n \equiv 1$ . But by the corollary to Lagrange's theorem,  $n|p-1$ .  $\square$

Let's look at an example. What is  $33^{22} \pmod{23}$ ? 23 is prime so  $33^{22} \equiv 1 \pmod{23}$ .

How about  $3^{101} \pmod{103}$ ? Well 103 is prime so  $3^{102} \equiv 1 \pmod{103}$ . So  $3^{101} \equiv 3^{-1} \pmod{103}$ . To find  $3^{-1} \pmod{103}$  use Euclid's algorithm.

$$103 = 3 \times 34 + 1.$$

So  $3^{-1} \equiv 34 \pmod{103}$ . Hence  $3^{101} \equiv 34 \pmod{103}$ .

Another example :  $32^6 \pmod{7}$ . We know that  $32^7 \pmod{7}$ . It follows that  $32^7 \equiv 32^{-1} \pmod{7}$ . It suffices to calculate  $32^{-1} \pmod{7}$ . We get

$$32 = 4 * 7 + 4$$

$$7 = 1 * 4 + 3$$

$$4 = 1 * 3 + 1$$

and

$$1 = (1*4) + (-1*3) = (-1*7) + (2*4) = (2*32) + (-9*7) = (-9*7) + (2*32)$$

Hence  $32^{-1} \equiv 2 \pmod{7}$  and  $32^6 \equiv 2 \pmod{7}$ .

Yet another example :  $45^{35} \pmod{13}$ .

We have  $13 \times 2 = 26$  and  $45^{13} \equiv 45 \pmod{13}$ . Hence, as  $35 = 13 \times 2 + 9$ , we have  $45^{35} = 45^2 \times 45^9 = 45^{11} \pmod{13}$ . As  $45^{12} \cong 1 \pmod{13}$ , we have  $45^{11} = 45^{-1} \pmod{13}$ .

We need to calculate  $45^{-1} \pmod{13}$ .

Euclidian algorithm : We get

$$45 = 3 * 13 + 6$$

$$13 = 2 * 6 + 1$$

and

$$1 = (1 * 13) + (-2 * 6) = (-2 * 45) + (7 * 13) = (7 * 13) + (-2 * 45)$$

Hence  $45^{35} \equiv -2 \pmod{13} \equiv 11 \pmod{13}$ .

Let's do  $43^{42} \pmod{13}$ . We have  $43^{39} \equiv 43^3 \pmod{13}$ . Hence  $43^{42} \equiv 43^6 \pmod{13}$ . Now  $43 \pmod{13} \equiv 4 \pmod{13}$ . Hence  $43^{42} \equiv 4^6 \pmod{13}$ . Now  $4^2 = 16 = 3 \pmod{13}$ . Hence  $4^6 = 4^{2^3} = 3^3 = 27 \pmod{13} = 1 \pmod{13}$ . Hence  $43^{42} \pmod{13} \equiv 1 \pmod{13}$ .

And now we get to yet another application of the Bézout's lemma.

We would like to find integers  $z$  that satisfy **two** congruences:  $z \equiv x \pmod{m}$  and  $z \equiv y \pmod{n}$ .

This is not always possible as the example  $z \equiv 3 \pmod{4}$  and  $z \equiv 5 \pmod{8}$  shows. If such a  $z$  existed, one would get  $0 = 1 \pmod{8}$  which is not the case. The reason is that 4 and 8 are not coprime. However, when  $n$  and  $m$  are coprime, we have the following theorem.

**Theorem 2.13 (Chinese Remainder Theorem)** *Suppose  $m$  and  $n$  are coprime; let  $x$  and  $y$  be two integers. Then there is a unique  $[z] \in \mathbb{Z}/nm$  such that  $z \equiv x \pmod{m}$  and  $z \equiv y \pmod{n}$ .*

**Proof.** (existence) By Bezout's Lemma, we can find  $h, k \in \mathbb{Z}$  such that

$$hn + km = 1.$$

Notice that  $hn \equiv 1 \pmod{m}$  and  $km \equiv 1 \pmod{n}$ .

Given  $x, y$  we choose  $z$  by

$$z = hnx + kmy.$$

Clearly  $z \equiv hnx \equiv x \pmod{m}$  and  $z \equiv y \pmod{n}$ .

(uniqueness) For uniqueness, suppose  $z'$  is another solution. Then  $z \equiv z' \pmod{n}$  and  $z \equiv z' \pmod{m}$ . Hence there exist integers  $r, s$  such that

$$z - z' = nr = ms.$$

Since  $hn + km = 1$  we have

$$z - z' = (z - z')hn + (z - z')km = mshn + nrkm = nm(sh + rk).$$

Hence  $z \equiv z' \pmod{nm}$ . □

As usual the proof gives you a procedure to find  $z$ . To find  $z$ , find  $h$  and  $k$  as in the Bézouts lemma (run Euclidean algorithm backwards). Then  $z$  is  $hnx + kmy$ .

Find the unique solution of  $x \equiv 3 \pmod{7}$  and  $x \equiv 9 \pmod{11}$  satisfying  $0 \leq x \leq 76$ .

Solution find  $h, k$  such that  $7h + 11k = 1$  using Euclid:

$$11 = 7 + 4$$

$$7 = 4 + 3$$

$$4 = 3 + 1$$

$$\text{So } 1 = 4 - 3 = 4 - (7 - 4) = 2 \cdot 4 - 7 = 2 \cdot (11 - 7) - 7 = 2 \cdot 11 - 3 \cdot 7.$$

$$\text{Hence let } h = -3 \text{ and } k = 2 \text{ so take } x = -3 \cdot 7 \cdot 9 + 2 \cdot 11 \cdot 3 = -189 + 66 = -123 \equiv 31 \pmod{77}.$$