Gymnasium Bäumlihof

Maturaarbeit

**Theoretical Informatics: Formal languages and Descriptive Complexity**

**A study of the connection of ﬁrst-order logic and context-sensitive languages**

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# Foreword

Some years ago, the Lego Mindstorms [[LEG13](#_bookmark45)] sparked my interest in informatics and programming. By attending some courses at the Phænovum [[Sch09](#_bookmark56)] in Lörrach, I learned how to program using Java, my ﬁrst text-based programming language. At that time, I was planning to work at Boston Dynamics [[Bos92](#_bookmark33)], as I loved being able to physically see what I had achieved. But then came the RoboCup robot [[Rob97](#_bookmark53)]. Me and my colleagues from the Phænovum wanted to program a rolling robot for playing football on a miniature playing ﬁeld. After two years, we were still unable to follow the ball because the Corona pandemic prevented us from working on-site together. Also, and to a greater extent, it didn’t work because the hardware never did what it was meant to, and we spent interminable hours just trying to make it roll forward. As may have become apparent, I got tired of it and decided to move on to something that didn’t include too much hardware and was more abstract. So I chose to participate in the Swiss Olympiad in Informatics [[Swi08](#_bookmark58)], which organizes national programming contests and selects various teams for international Olympiads. There, I got many interesting problems which I loved to solve, found like-minded friends, and was able to participate in several international competitions. Over the years, I began to notice that I enjoy solving tasks theoretically way more than implementing them. This is also something that was reﬂected in my competition scores, where I often came out knowing I solved a lot in theory but failed to get the points.

Some internships at informatics ﬁrms conﬁrmed that actual programming was still too concrete for me.

Thanks to the “Schülerstudium” at the University of Basel, I attended a course on the mathematical background of computer science [[HR22](#_bookmark38)] and one on the theory of computer science [[Rög23](#_bookmark54)]. There, I learned more about abstract concepts such as Complexity theory, computability, logic, and the Chomsky hierarchy. The way in which proofs could be made to hold for every problem with certain properties has fascinated me ever since. Further, logic is a tool that captures mathematical reasoning, and can thus be seen as a formalization of every “logical” thought we have. Also, I don’t have to bother with implementation any more.

While using books to learn more, I found out about the domain of Descriptive Complexity. This domain relates diﬀerent kinds of logic to diﬀerent classes of problems in computer science. I knew I wanted to do my Matura project in this domain, but had diﬃculties ﬁnding an open question which seemed approachable. I asked many people if they had a more exact idea of what I could do. In the end, a friend at the informatics Olympiad asked *ChatGPT* [[Ope22](#_bookmark50)], which told me I could study the connection between the Chomsky hierarchy and Descriptive Complexity. Reﬁning this proposition, I came up with the subject of this project: relating ﬁrst-order logic and context-sensitive languages.

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# Introduction

In our daily lives, we are in contact with various kinds of algorithms at all times. Searching on Google, sending texts, asking questions to [AI](#_bookmark83) chatbots, searching for the fastest route with a [GPS](#_bookmark84); all of these consume resources in energy, storage space, and time. According to a report from 2021 [[Cli21](#_bookmark35)], 3.7% of global carbon emissions come from the [IT](#_bookmark85) domain, with an upward tendency. This footprint is similar to that of the airline industry. At this scale, it is vitally important to understand and ﬁnd out if the resource consumption can be reduced.

To ﬁnd out how to improve, understanding how computers work, what they can compute, and why solving some problems is more diﬃcult than solving others is needed. The ﬁeld of study that investigates this is called Complexity theory. The main tool used in that ﬁeld is the abstraction of computational models and problems. It is often possible to ﬁnd classes of similar problems which allow for generalizations. However, it has proven to be very hard to ﬁnd proofs of either the optimality of algorithms or the separation of complexity classes. One of the emblematic open problems is the [P](#_bookmark78) versus [NP](#_bookmark79) question, which we go further into in appendix [B.2.8](#_bookmark108). Nevertheless, some signiﬁcant results concerning the average complexity of real-world problems and the hardness of cryptographic encryptions have been obtained.

One of the subﬁelds of Complexity theory is Descriptive Complexity. It relates mathematical logic to diﬀerent complexity classes and allows us to get new insights into the underlying structure of problems of certain classes. The other main ﬁeld present in this document, formal languages, is concerned with the abstraction of computational problems as sets, which is often helpful for proofs. This then allows multiple formalisms which describe these sets to be deﬁned. The formal language class this document focuses on, called context-sensitive languages, is part of the “Chomsky hierarchy” introduced by the famous linguist Noam Chomsky in [[Cho59](#_bookmark34)].

Towards the end of the 20th century, many equivalences were proven between fragments and

extensions of various logics and complexity classes, including all classes in the Chomsky hierarchy. A great summary of all the results can be found in Neil Immermans paper “Languages that capture complexity classes” [[Imm87](#_bookmark40)].

In this project, the focus lies on the following question: What are the connections between context- sensitive languages and extensions and restrictions of ﬁrst-order logic? It is known that context- sensitive languages are captured by linear bounded nondeterministic Turing machines and that these in turn are equivalent to a certain extension of second-order logic [[Imm99](#_bookmark42)]. Many complexity classes have multiple formalizations using logic, but this one does not. This is why it is interesting to try to ﬁnd new characterization that can give new insights into the inherent structure of context-sensitive languages. First-order logic has many powerful extensions that allow for such research and thus has become the main focus of this project.

This document starts by introducing the relevant theory in the ﬁrst two chapters. Later, a personal

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study of connections between ﬁrst-order logic and context-sensitive languages is presented.

Both chapter [2](#_bookmark1) and chapter [3](#_bookmark6) present the most important theory of formal languages and De- scriptive Complexity, respectively. Additionally, in section [3.4](#_bookmark14) full proofs concerning equivalences of context-sensitive languages with logics and automata are presented. These proofs are written in a way that shows the motivation behind crucial steps. Some details have been omitted as presenting them would be beyond the scope of this document.

After the introduction to the material, chapter [4](#_bookmark19) describes my personal studies on connections between ﬁrst-order logic and context-sensitive language. Diﬀerent approaches that I tried are presented and investigated. Towards the end of the chapter, the direct connection to context-sensitive languages fades and the objects of study become more related to Savitch’s theorem, which plays a great role in the theory of space-bounded computation. The main result is given in section [4.5](#_bookmark25) and concerns a simulation of alternating Turing machines using iterative logic. This divergence from the original working question is due to the hope that researching other, more loosely connected problems would ultimately help to ﬁnd a solution for the core problem. Also, I aimed to investigate many personal ideas instead of dwelling on approaches that rapidly proved to not develop into successful ways.

Finally, in chapter [5](#_bookmark26), what was accomplished during this project is reﬂected. A review of the working process is presented, and ﬁnally, further possibilities of research are discussed.

For the relevant mathematical background, appendix [A](#_bookmark86) contains the basic deﬁnitions of Set theory, ﬁrst and second-order logic, and Turing machines.

A full collection of descriptions, proofs, techniques, and context information about Descriptive Complexity and languages in the Chomsky hierarchy can be found in appendix [B](#_bookmark92).

Appendix [C](#_bookmark109) contains all the messages and emails exchanged in the course of this project. All notes which were made during this project can be found in appendix [D](#_bookmark117).

In mathematical research, English has traditionally been the language of communication between members of the community, may it be for papers, exchanges or other publications. This document adheres to this tradition and is written in English.

# Formal Languages

## Deﬁnitions

In informatics, we often get an input as a string of characters and want to compute some function on it. Complexity theory mostly focuses on decision problems which ask whether some input fulﬁls some given property. To formalize this, there is the concept of formal languages. The following deﬁnitions are taken from the lecture Theory of Computer Science [[Rög23](#_bookmark54)]. For the mathematical background, refer to appendix [A](#_bookmark86).

**Deﬁnition 2.1** (Alphabet)**.** An alphabet Σ is a ﬁnite set of symbols.

**Deﬁnition 2.2** (Word)**.** A word *w* over some alphabet Σ is a ﬁnite sequence of symbols from Σ. We denote *ε* as the empty word, Σ∗ as the set of all words over Σ, *xy* as the concatenation of the two words *x* and *y*, *xn* as the concatenation of *x* with itself *n* times, and |*x*| as the number of symbols in *x*.

**Deﬁnition 2.3** (Formal language)**.** A formal language is a set of words over some alphabet Σ, or equivalently a subset of Σ∗.

For any computational decision problem, we can reformulate it as the problem of deciding if the input word is contained in the formal language consisting of all words which have the required property.

## Chomsky Hierarchy

One of the multiple ways to categorize formal languages was invented by Avram Noam Chomsky, a modern linguist, in [[Cho59](#_bookmark34)]. It is based on the complexity of deﬁning the language using formal grammars, which are a ﬁnite representation of formal languages (which can be inﬁnite in the general case).

### Grammars

A grammar can informally be seen as a set of rules telling us how to generate all words in a language.

**Deﬁnition 2.4** (Grammar)**.** A grammar is a 4-tuple ⟨*V,* Σ*, R, S*⟩ consisting of

*V* : The set of non-terminal symbols.

Σ: The alphabet of terminal symbols. All words generated by the grammar are in Σ∗.

*R*: The set of rules of the form *a* → *b* with *a* and *b* being words over the alphabet consisting of the union of *V* and Σ. Any rule in *R* must have at least one symbol from *V* on its left-hand side.

*S*: The start symbol from the set *V* .

The non-terminal symbols in *V* are symbols that are not in Σ and exist for the purpose of steering the process of word generation. No non-terminal symbols appear in any of the words generated by the grammar.

To generate the words, there is the concept of derivations.

**Deﬁnition 2.5** (Derivation)**.** First, one derivation step is deﬁned: We say that *u*′ can be derived directly from *u* if

* *u* is of the form *xyz* and *u*′ is of the form *xy*′*z* for some words *x, y, y*′*, z* consisting of symbols in

Σ and *V* .

* there exists a rule *y* → *y*′ in *R*.

We say that a word is in the *generated language* of a grammar if it consists only of symbols in Σ

and can be derived in a ﬁnite number of steps from *S*.

**Example 2.1.** Consider the grammar ⟨{*S*}*,* {*a, b*}*, R, S*⟩ with

*R* = n*S* → *aSb, S* → *ε*o

The generated language of this grammar is {*ε, ab, aabb, . . .* } = {*anbn* | *n* ∈ N0}.

Now that we have a tool to describe some inﬁnite languages using a ﬁnite description, we can further diﬀerentiate the complexity of a language by the minimum required complexity of the rules in any formal grammar describing that language. In the main section, only the context-sensitive languages are presented, as they are important for later chapters. In appendix [B.1](#_bookmark93), explanations for regular languages (appendix [B.1.1](#_bookmark94)), context-free languages (appendix [B.1.2](#_bookmark95)), and recursive languages (appendix [B.1.3](#_bookmark96)) are provided.

### Context-Sensitive Languages

We can deﬁne the most signiﬁcant class of languages for this document by multiple equivalent restric- tions on the grammars.

One restriction is that all rules have to be of the form *αβγ* → *αφγ* with *α* and *γ* being words over Σ ∪ *V* , *β* being a symbol of *V* and *φ* being a nonempty word over Σ ∪ *V* . This means that only the non-terminal symbol is allowed to change. Additionally, if *S* is the start symbol and never occurs on the right-hand side of any rule, we may include the exception *S* → *ε*.

Equivalently, we can require that for any rule *u* → *v* we have |*u*| ≤ |*v*| as shown in [[Par02](#_bookmark51)]. This means that applying any rule makes the result longer. Again, we also allow the exception *S* → *ε* if *S* does not occur on the right-hand side of any rule in *R*. These grammars are called noncontracting.

**Example 2.2.** Consider the grammar ⟨{*S, B*}*,* {*a, b, c*}*, R, S*⟩ with

*R* = *S* → *abc, S* → *aSBc, cB* → *Bc, bB* → *bb*

f \_

It generates the language *anbncn* for *n* ∈ N1 and is noncontracting.

The last restriction, the one that is most useful for proofs, is the Kuroda normal form presented in [[Pet22](#_bookmark52)], where all rules have one of the following structures:

* + - * *A* → *BC*
      * *AB* → *CD*
      * *A* → *a*
      * *S* → *ε* if *S* is the start symbol and does not occur on any right-hand side of a rule.

where *A, B, C, D, S* ∈ *V* and *a* ∈ Σ.

The corresponding formalism for these languages is the linearly bounded nondeterministic Turing machine, which can only write on the tape cells that contained the input word in the beginning. This and an equivalent extension of second-order logic are proven in section [3.4.1](#_bookmark15).

# Descriptive Complexity

## Aims

In mathematics, abstraction is one of the most important tools, enabling us to make general statements and prove them for all concrete instances of a concept. Formal logic takes this even further and makes it possible to abstract mathematical thought itself. That is the reason why Descriptive Complexity uses formal logic to describe computational problems in a way that captures their mathematical essence.

In computer science, we are always interested in the amount of resources needed to compute some function or solve a certain problem, speaking in terms of time and storage space. By focusing on decision problems[1](#_bookmark10), we can deﬁne a logical characterization of a problem as a formula *φ* which is true if and only if a structure satisﬁes the required properties. By looking at the required complexity of formulas that describe some problems in terms of relations, operators, variables, and other metrics, we can often ﬁnd remarkably natural classes of logic that correspond to computational classes of problems.

Using results from this research, many insights into the underlying structure of real-world problems can be made, which in turn can lead to better ways to tackle them. Further, Descriptive Complexity has applications in database theory, computer-aided veriﬁcation, and proofs.

## Tools

Some tools and techniques which are used later for proofs need to be introduced ﬁrst. In this section, only Complexity theory is presented. In appendix [B.2.1](#_bookmark98) and appendix [B.2.2](#_bookmark99), Ehrenfeucht-Fraïssé games and ﬁrst-order reductions are also presented. These tools are not needed for chapter [4](#_bookmark19), but I tried using them extensively during the research phase of this project. Deﬁnitions are again taken from [[Rög23](#_bookmark54)] and [[Imm99](#_bookmark42)].

### Complexity Theory

Complexity theory is the study of the resources, measured mostly in time and space, needed to compute the answers to computational problems[2](#_bookmark11). As diﬀerent computers can handle tasks at various speeds, a notation which only considers the asymptotic use of resources is used.

**Deﬁnition 3.1** (Big-O notation)**.** Let *f, g* be functions *f, g* : N → R+.

1Any problem can be reduced to boolean queries, for example by having a boolean query meaning “the i*th* bit of an encoding of the answer is 1”

2The speciﬁc model of computation is not important as all give almost the same results. We assume Turing machines, which are deﬁned in appendix [A.4](#_bookmark91).

We say that *f* (*n*) ∈ O(*g*(*n*)) if there exists positive integers *n*0*, c* such that for all *n* ≥ *n*0 we have

*f* (*n*) ≤ *c* · *g*(*n*)

Complexity classes can be deﬁned as the set of all problems for which there exists a Turing machine satisfying some bounds that computes the answer to the problem. Now, some common complexity classes are deﬁned.

**Deﬁnition 3.2** ([DTIME](#_bookmark72)[*t*(*n*)])**.** We say that a decision problem is in [DTIME](#_bookmark72)[*t*(*n*)] if there exists a deterministic Turing machine that takes a maximum of *f* (*n*) steps on any input of size *n* and *f* (*n*) ∈ O(*t*(*n*)).

**Deﬁnition 3.3** ([P](#_bookmark78))**.** We say that a decision problem is in [P](#_bookmark78) if there exists a polynomial *q* such that the problem is in [DTIME](#_bookmark72)[*q*(*n*)].

**Deﬁnition 3.4** ([DSPACE](#_bookmark73)[*t*(*n*)])**.** We say that a decision problem is in [DSPACE](#_bookmark73)[*t*(*n*)] if there exists a deterministic Turing machine that visits a maximum of *f* (*n*) tape cells on any input of size *n* and *f* (*n*) ∈ O(*t*(*n*)).

**Deﬁnition 3.5** ([PSPACE](#_bookmark80))**.** We say that a decision problem is in [PSPACE](#_bookmark80) if there exists a polynomial

*q* such that the problem is in [DSPACE](#_bookmark73)[*q*(*n*)].

We can do the same for nondeterministic Turing machines, and get the corresponding complexity classes [NTIME](#_bookmark74), [NP](#_bookmark79), [NSPACE](#_bookmark75), and [NPSPACE](#_bookmark81). There, we always take the maximum of tape cells and steps over any computation branch.

The complexity class [P](#_bookmark78) has a special meaning for computer scientists, as the problems in [P](#_bookmark78) are deemed “feasible” on modern computers.

## Important Results

In this section, only Savitch’s theorem is introduced, which Savitch presented in [[Sav70](#_bookmark55)]. In ap- pendix [B.2.3](#_bookmark100), the space hierarchy theorem is presented.

### Savitch’s Theorem

Savitch’s theorem relates nondeterministic space-bounded Turing machines with deterministic space- bounded Turing machines. As an intermediate step, it introduces alternating Turing machines. These Turing machines are a generalization of nondeterministic Turing machines, which can be seen as machines taking the “or” over all their computation paths.

**Deﬁnition 3.6** (Alternating Turing Machine)**.** An alternating Turing machine is a Turing machine with two types of states: existential and universal. The accepting conditions change compared to a nondeterministic Turing machine and depend on the current state of the alternating Turing machine. If it is in an existential state, the current conﬁguration is accepting if and only if *at least one* of the directly reachable conﬁgurations is accepting. If the machine is in a universal state, the current conﬁguration is accepting if and only if *all* the directly reachable conﬁgurations are accepting.

In addition to taking the “or” over its children, an alternating Turing machine can thus also take the “and”.

We deﬁne [ATIME](#_bookmark76)[*t*(*n*)] and [ASPACE](#_bookmark77)[*t*(*n*)] analogously to [NTIME](#_bookmark74)[*t*(*n*)] and [NSPACE](#_bookmark75)[*t*(*n*)] in section [3.2.1](#_bookmark9).

We can proceed to the proof of Savitch’s theorem using the technique presented in [[Imm99](#_bookmark42)].

**Theorem 3.1** (Savitch’s Theorem)**.** *For all space constructible functions with t*(*n*) ≥ log *n we have*

[NSPACE](#_bookmark75)[*t*(*n*)] ⊆ [ATIME](#_bookmark76)[*t*(*n*)2] ⊆ [DSPACE](#_bookmark73)[*t*(*n*)2]

*Proof.* We start with the ﬁrst inclusion, [NSPACE](#_bookmark75)[*t*(*n*)] ⊆ [ATIME](#_bookmark76)[*t*(*n*)2]. We need to show that any [NSPACE](#_bookmark75)[*t*(*n*)] Turing machine can be simulated by an [ATIME](#_bookmark76)[*t*(*n*)2] alternating Turing machine. Let *N* be a [NSPACE](#_bookmark75)[*t*(*n*)] Turing machine. Without loss of generality, we assume that *N* clears its tape after accepting and goes back to the ﬁrst cell.

Consider *Gw*, the computation graph of *N* on input *w*. This graph consists of conﬁgurations as vertices and directed edges from each conﬁguration to all conﬁgurations that are directly reachable from it. We see that *N* accepts *w* if and only if there is a path from the start conﬁguration *s* to the accepting conﬁguration *t* in *Gw*. We now present a routine *P* (*d, x, y*) which asserts that there is a path of length at most 2*d* from vertex *x* to *y*. Recursively, we can deﬁne *P* as follows:

*P* (*d, x, y*) = (∃*z*)(*P* (*d* − 1*, x, z*) ∧ *P* (*d* − 1*, z, y*))

This formula asserts that there exists a middle vertex *z* for which there is a path of length at most 2*d*−1 from *x* to *z* and from *z* to *y*. The base case for *d* = 0 is to check whether *x* can be reached from *y* by a transition in the transition table of *N* . Using an alternating Turing machine, the formula can be evaluated by using existential states to ﬁnd the middle vertex *z*, and then a universal state choosing which one of the two new paths should be checked.

For the runtime analysis, we proceed as follows. It takes O(*t*(*n*)) time to write down the middle vertex *z* as each conﬁguration includes the tape contents, which have length O(*t*(*n*)). Further, one of the new shorter paths *P* (*d* − 1*, z, y*) is then evaluated. By induction, we ﬁnd that we need O(*d* · *t*(*n*)) time to compute *P* (*d, x, y*). There are only 2O(*t*(*n*)) possible conﬁgurations, so we get that the initial *d* is also in O(*t*(*n*)), and thus our total runtime is O(*t*(*n*) · *t*(*n*)) = O(*t*(*n*)2).

For the second inclusion, by substituting *s*(*n*) for *t*(*n*)2, we need to simulate an [ATIME](#_bookmark76)[*s*(*n*)] Turing machine *A* using a [DSPACE](#_bookmark73)[*s*(*n*)] Turing machine. Again, we consider the computation graph of *A* of on input *w*. This graph has depth O(*s*(*n*)) and size 2O(*s*(*n*)).

A [DSPACE](#_bookmark73)[*s*(*n*)] Turing machine *D* can systematically search this computation graph to simulate

*A*. This is done by keeping a string of choices *c*1*c*2 *. . . cr* of length O(*s*(*n*)) made up to this point. Each choice takes up O(1) place, as the number of transitions from one state is bounded by a constant depending only on *A*. Note that these choices uniquely determine which state *A* is in.

Now, *D* can ﬁnd the answer recursively. If in the simulation, *A* is in a halt state, *D* reports this back to the previous state. In an existential state, *D* simulates the children of that state, and if it gets

a positive result from one of them, *D* also reports a positive result. In a universal state, *D* simulates the children of that state, and if it gets a positive result from all of them, *D* reports a positive result. Overall, *D* only uses O(*s*(*n*)) space to simulate *A*. This shows that we can simulate an

[ATIME](#_bookmark76)[*t*(*n*)2] Turing machine *A* using a [DSPACE](#_bookmark73)[*t*(*n*)2] Turing machine.

Thus, the second part of the theorem follows and by the transitivity of ⊆ we have [NSPACE](#_bookmark75)[*t*(*n*)] ⊆

[DSPACE](#_bookmark73)[*t*(*n*)2].

It is not known if the containment is strict for any of the inclusions. Neither is it known if the quadratic overhead for simulating nondeterministic space is optimal. From this theorem, we also get the following corollary.

**Corollary 3.1.1.** *We have* [PSPACE](#_bookmark80) = [NPSPACE](#_bookmark81)*. Proof.*

[NPSPACE](#_bookmark81) = [NSPACE](#_bookmark75)[*nk*] by deﬁnition

L

*k*∈N

⊆ L

[DSPACE](#_bookmark73)[*n*2*k*] by Savitch’s theorem

*k*∈N

⊆ L

[DSPACE](#_bookmark73)[*nk*]

*k*∈N

= [PSPACE](#_bookmark80) by deﬁnition

L

= [DSPACE](#_bookmark73)[*nk*]

*k*∈N

⊆ L

[NSPACE](#_bookmark75)[*nk*]

*k*∈N

= [NPSPACE](#_bookmark81)

## Results concerning the Chomsky Hierarchy

Now that we have seen most of the required theory, we can start to apply it to the main subject of this document: the Chomsky hierarchy.

The deﬁnitions of ﬁrst- and second-order logic can be found in appendix [A.2](#_bookmark88) and appendix [A.3](#_bookmark90), respectively. For this section, we deﬁne the vocabulary on strings to be *σ* = ⟨*Qa, Qb, . . . , Qz,* ≤

*,* 0*,* 1*,* max⟩. We have a unary predicate *Qa* for each character *a* in Σ, a total ordering on the universe, and the constants 0*,* 1, and max. For a universe |A| = {0*,* 1*, . . . , n* − 1}, we require that max = *n* − 1. Also, for each *a* ∈ Σ we require that *Qa*(*x*) is true if and only if the *xth* character of the string is *a*.

Only the results for context-sensitive languages are discussed in this section. Further results can be found in appendix [B.2](#_bookmark97).

### Context-Sensitive Languages

This is the language class that interests us most, as it has been studied less extensively than other language classes. Nevertheless, there are some known formalisms. One of them is the linear bounded nondeterministic Turing machine. This ﬁrst equivalence is shown in appendix [B.2.6](#_bookmark104).

We deﬁne [MSO](#_bookmark63)([TC](#_bookmark64)) to be second-order logic restricted to quantiﬁcation over unary relations supplemented with the transitive closure operator. Then, the second formalism is given by the logic [MSO](#_bookmark63)([TC](#_bookmark64)). The equivalence was established by Immerman in [[Imm87](#_bookmark40)]. The transitive closure operator takes the transitive closure over some graph and is deﬁned as follows:

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**Deﬁnition 3.7** (Transitive closure)**.** Let *ϕ a, b* be a formula with 2*k* free variables. We can see this formula as a directed edge relation over the graph with vertices *c k*. Then, the transitive closure

(*TCa,bϕ* (*a, b*)) (*u, v*) is true if and only if there is a path in the graph generated by *ϕ* from *u* to *v*.

∈ |A|

**Theorem 3.2.** *A language L is context-sensitive if and only if it can be described by a formula in* [*MSO*](#_bookmark63)([*TC*](#_bookmark64))*.*

*Proof.* By theorem [B.7](#_bookmark105), we can also show the equivalence of [MSO](#_bookmark63)([TC](#_bookmark64)) and [NSPACE](#_bookmark75)[*n*] Turing ma- chines.

First, we show that any formula in [MSO](#_bookmark63)([TC](#_bookmark64)) can be evaluated by a [NSPACE](#_bookmark75)[*n*] Turing machine

*N* . The ﬁrst step is to notice that any relation of [MSO](#_bookmark63) can be represented on the tape in O(*n*) space. For any sub-formula of the form *TCa,bϕ a, b* (*u, v*), *N* writes down *u* and guesses the next vertex on the tape. Then, *N* can check if the transition is valid by evaluating *ϕ*. If it is, *N* replaces *u* with the new vertex and repeats until it reaches *v*. Because Immerman showed that [NSPACE](#_bookmark75) is closed under complementation in [[Imm88](#_bookmark41)], we know that *N* can also compute any sub-formula of the form

(  ( ))

(  ( ))

¬ *TCa,bϕ a, b* (*u, v*). The other parts of the formula can be evaluated easily in linear space, as *N* just needs to write down relations when quantifying and remember in which part of the constant-size formula it is.

To prove the other direction, consider a [NSPACE](#_bookmark75)[*n*] Turing machine *N* . We deﬁne the string

vocabulary for our logic to be over the alphabet Σ which contains the symbol *ba ,q* for all pairs of

*i j*

tape symbols and Turing machine states in addition to all normal symbols *ai*. We use these symbols to completely represent an instantaneous conﬁguration of the computation of *N* . This is done by creating a string from the tape content and replacing the character at the position of the read/write

head by the *ba ,q* which contains the actual character and state of *N* .

*i j*

Consider the tuple *X* = ⟨*Qb*1 *, . . . , Qbr* ⟩ containing a unary relation for each symbol in Σ. This tuple completely represents an instantaneous conﬁguration of the tape and the Turing machine. We can now write a formula *φ*(*X, Y* ) which is true if and only if a transition from state *X* to *Y* is possible in *N* . This can be done by a big disjunction over all rules of *N* . After one step of *N* , the new character *b* at position *x* is only determined by the actual character at position *x* and the two characters left and right of it[3](#_bookmark17). This means that we can write all transitions of *N* in the form ⟨*bk, bl, bm*⟩ → ⟨*bi, bj, bw*⟩. If

3we include markers at the two ends of the tape, which makes this well-deﬁned

*P* is the set of all transitions, we have

*φ*(*X, Y* ) ≡ ∃*i* ∀*j* |*i* − *j*| *>* 1 → /\

*bi*∈Σ

(*YQbi* (*j*) ↔ *XQbi* (*j*)) ∧

⟨*bk,bl,bm*⟩→V⟨*bu,bv,bw*⟩∈*P*

(*XQbk* (*i* − 1) ∧ *XQbl* (*i*) ∧ *XQbm* (*i* + 1) ∧ *YQbu* (*i* − 1) ∧ *YQbv* (*i*) ∧ *YQbw* (*i* + 1))

The ﬁrst line asserts that apart from the indices where the head is, the tape stays unchanged. The second line states that there exists a rule in *P* such that the characters in *Y* at positions *i* − 1*, i*, and *i* + 1 follow from the previous characters in *X*.

Now, we can take the transitive closure over *φ*, starting with the start conﬁguration and ending at an accepting conﬁguration. Without loss of generality, we can assume that *N* clears its tape after accepting, therefore the end position is unique. Thus, our formula holds if and only if there is an accepting path in *N* on the input word *w*.

This concludes the equivalence of [MSO](#_bookmark63)([TC](#_bookmark64)) and context-sensitive languages. An interesting normal form for [MSO](#_bookmark63)([TC](#_bookmark64)) can be derived from this proof.

**Corollary 3.2.1.** *Every formula in* [*MSO*](#_bookmark63)*(*[*TC*](#_bookmark64)*) can be written in the form*

(*TCX,Y φ*(*X, Y* )) (*U, V* )

*Proof.* We have given an explicit way to convert any formula into a Turing machine and back in theorem [3.2](#_bookmark16). By construction, this always gives us a formula of the required form.

Returning to our research question, we now have one characterization of context-sensitive lan- guages, but using second-order logic. Still, we can say that these proofs show that context-sensitive languages are tightly connected to the connectivity of graphs with 2*cn* vertices. This is because the transitive closure answers exactly all questions regarding the existence of paths between two ver- tices in a graph, and corollary [3.2.1](#_bookmark18) combined with theorem [3.2](#_bookmark16) shows that all formulas capturing a context-sensitive language can be written as a formula checking such a path on a graph.

# Personal Contribution

In this chapter, results of searches for new restrictions of ﬁrst-order logic and other related concepts are presented. I did not ﬁnd any fundamentally new result about the characterization of context-sensitive languages using ﬁrst-order logic but managed to prove that various approaches could not lead to an equivalence. Further, I also lowered some upper bounds for simulating alternating Turing machines using iterative logic. Using a restricted form of iterative logic, I could lower the bounds for an iterative simulation of a nondeterministic Turing machine.

## Direct Transformation

Similar to Immerman in [[Imm99](#_bookmark42)], we use extended variables to model second-order variables in ﬁrst- order logic. This then directly gives us a characterization of [NSPACE](#_bookmark75)[*s*(*n*)] without requiring any new insights, as we can just use the same technique as for second-order logic.

**Deﬁnition 4.1** (Extended Variable)**.** The logic [FO-VAR](#_bookmark66)[1*, s*(*n*)] has two types of variables. Domain variables are the ﬁrst type, which we denote by lowercase characters, ranging from 0 to *n*−1. Extended variables are the other type, which we denote by uppercase characters, ranging from 0 to 2*s*(*n*) log(*n*) −1, and thus having *s*(*n*) log(*n*) bits. The extended variables are not allowed to appear as an input to any relation apart from [BIT](#_bookmark65). Thus, we can only query if a speciﬁc bit in the binary representation of an extended variable is on. For extended variables with more than *n* bits, we can extend [BIT](#_bookmark65) to accept a tuple of domain variables encoding the position we want to query. This makes extended variables with a polynomial number of bits possible.

The extended variables in [FO-VAR](#_bookmark66)[1*, nk/* log(*n*)] have exactly the same capabilities as second-order variables of arity *k* in second-order logic. This is true as for every second-order variable *X* we can have an extended variable *X*′ such that bit *x* of *X*′ is set if and only if *X*(*x*) is true. By generalizing the proof of theorem [3.2](#_bookmark16), we have [NSPACE](#_bookmark75)[*nk*] = [SO](#_bookmark62)([TC](#_bookmark64), arity *k*). We now prove that [SO](#_bookmark62)([TC](#_bookmark64), arity *k*)

= [FO-VAR](#_bookmark66)[1*, nk/* log(*n*)]([TC](#_bookmark64)), thus also capturing [NSPACE](#_bookmark75)[*nk*] with an extension of ﬁrst-order logic.

*Proof.* We show by induction on the structures of the formulas that each formula in [SO](#_bookmark62)([TC](#_bookmark64), arity *k*) has an equivalent formula in [FO-VAR](#_bookmark66)[1*, nk/* log(*n*)]([TC](#_bookmark64)) and vice versa.

Any atomic formula in [SO](#_bookmark62)([TC](#_bookmark64), arity *k*) which does not include any second-order variable is trivially writable in [FO-VAR](#_bookmark66)[1*, nk/* log(*n*)]([TC](#_bookmark64)). The same holds for formulas in [FO-VAR](#_bookmark66)[1*, nk/* log(*n*)]([TC](#_bookmark64)) without extended variables. An atomic formula with a second-order variable *Y* (*x*) can be represented as querying the bit *x* of an extended variable *Y* ′. We can do the same transformation backwards for an atomic formula of the form [BIT](#_bookmark65)(*Y* ′*, x*) and get a formula of the form *Y* (*x*) with a second-order variable *Y* .

Using this, we can induct on the structure of the formulas. Taking the conjunction, disjunction, or negation of equivalent formulas generates new formulas which are also equivalent. When a formula quantiﬁes over a second-order or extended variable, we can directly exchange this with a quantiﬁcation over the respective other type of variable. By induction, the sub-formulas without the quantiﬁcation are equivalent. Binding the new second-order or extended variable makes the new formulas equivalent.

A formula that takes the transitive closure can also be reformulated by replacing all variables of the current type with variables of the other type. Again using induction on the sub-formulas, we get that the new formulas are equivalent.

By induction, we then have that every formula in either logic has an equivalent formula in the other logic.

Using this equivalence and the proof in section [3.4.1](#_bookmark15), we get that [FO-VAR](#_bookmark66)[1*, n/* log(*n*)]([TC](#_bookmark64)) de- scribes exactly the context-sensitive languages. This gives us our ﬁrst characterization of context- sensitive languages in ﬁrst-order logic.

## Analogues to Proof for [DSPACE](#_bookmark73)

For [DSPACE](#_bookmark73)[*s*(*n*)], there are multiple other logical characterizations which do not use the transitive closure operator.

First, the logics [FO](#_bookmark61)[*t*(*n*)], which formalizes iterative deﬁnitions, and [VAR](#_bookmark69)[*k*], which restricts the number of variables but allows for unbounded ﬁrst-order iterations, are deﬁned.

**Deﬁnition 4.2** ([FO](#_bookmark61)[*t*(*n*)])**.** Let *Q*1*, . . . , Qn* be a series of quantiﬁers, *s*1*, . . . , sn* a series of variables, and *M*1 *. . . , Mn* a series of quantiﬁer-free formulas. A quantiﬁer block has the form

[QB](#_bookmark70) = (*Q*1*s*1*.Mn*) *. . .* (*Qnsn.Mn*)

For a universal quantiﬁer, (∀*s.M* )*φ* ≡ ∀*s*(*M* → *φ*). For an existential quantiﬁer, (∃*s.M* )*φ* ≡ ∃*s*(*M* ∧*φ*). Both of these equivalences eﬀectively mean that the quantiﬁers are restricted to range only over the elements that satisfy *M* . Then, a formula of [FO](#_bookmark61)[*t*(*n*)] is of the form

([[QB](#_bookmark70)]*t*(*n*)*M*0) (*c/s*)

where *M*0 is a quantiﬁer-free formula, *c* = ⟨*c*1*, . . . , cn*⟩ is a tuple of constants, and *s* = ⟨*s*1*, . . . , sn*⟩ are the variables occurring in the quantiﬁer block. The notation (*c/s*) means that in the beginning, we deﬁne *s*1 := *c*1*, . . . , sn* := *cn*. [[QB](#_bookmark70)]*t*(*n*) stands for [QB](#_bookmark70) literally repeated *t*(*n*) times. The truth values of these formulas for a speciﬁc structure are deﬁned by evaluating the formula obtained by iterating the quantiﬁer block *t*(∥A∥) times.

This gives us a formalism for iterative procedures.

If we restrict the number of variables such that there are at most *k* distinct variables, but allow some additional boolean variables[1](#_bookmark22), we get [FO-VAR](#_bookmark66)[*t*(*n*)*, k*]. In particular, we can reuse the same

1variables which have only two possible values

variable multiple times in the same quantiﬁer block. Additionally, we deﬁne

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[VAR](#_bookmark69)[*k*] = [FO-VAR](#_bookmark66)[2*cnk , k*]

*c*=1

This is the same as saying that we allow unbounded iterations, as after at most 2*cnk* iterations, the truth value of the formula will loop or stay the same. The constant *c* depends only on the number of boolean variables in the formula.

For [DSPACE](#_bookmark73), we have

[DSPACE](#_bookmark73)[*nk*] = [VAR](#_bookmark69)[*k* + 1]

A proof of this can be found in [[Imm99](#_bookmark42)]. In the main part of the proof, a construction is made to simulate a [DSPACE](#_bookmark73)[*nk*] Turing machine using [VAR](#_bookmark69)[*k* + 1]. There, the relation *Ct*(*x, b*) is inductively deﬁned to mean that at time *t*, the character on the tape at position *x* is the one encoded by *b*, where *b* is a tuple of boolean variables. The characters are deﬁned in the same way as in the proof of theorem [3.2](#_bookmark16) to include the machine state as well as the tape symbol. As the Turing machines under consideration are deterministic, the character at position *x* is uniquely determined. Further, it only depends on the characters at positions *x* − 1*, x*, and *x* + 1 at time *t* − 1. Because it can be assumed without loss of generality that the Turing machine returns to the starting tape cell and clears its tape

after accepting, ﬁnding out the value of *C*2*cnk* (0*, b*) is suﬃcient to determine if the machine accepts.

If we want to extend this to nondeterministic Turing machines by deﬁning *Ct*(*x, b*) to mean that

the character at position *x can* be *b* at time *t*, we run into problems. As we do not have the guarantee that the computation is deterministic, we can not say any more that the new state only depends on the possible characters on the tape at time *t* − 1. Doing this would mean that impossible states could be reached. This would happen because the formula does not remember which combinations of characters are possible, which is required as the characters of each tape cell can not be chosen independently. One way to ﬁx this is by remembering the nondeterministic choices made in the previous steps. This would make the values deterministic again, ﬁxing the issue. In the worst case, this would take O(2*cnk* ) additional bits. But we already know by Savitch’s theorem (section [3.3.1](#_bookmark13)) and the equivalence for [DSPACE](#_bookmark73) that we only need log(*n*) · (*k* + 2) bits to represent a [NSPACE](#_bookmark75)[*nk*] computation. This is much less than the 2*cnk* bits needed using this approach.

Let [VAR](#_bookmark69)[*r*(*n*)] be the generalization of [VAR](#_bookmark69)[*k*] which means that a total of *r*(*n*) + O(1) bits of variables are allowed. Any proof which would show that [NSPACE](#_bookmark75)[*nk*] can be expressed in [VAR](#_bookmark69)[*r*(*n*)] with log(*n*) · (*k* + 1) ≤ *r*(*n*) *<* log(*n*) · (*k* + 2) would be an improvement on Savitch’s theorem. That is because the proof for [DSPACE](#_bookmark73)[*nk*] can easily be extended to more general polynomial functions, as Immerman showed in [[IBB99](#_bookmark39)]. If Savitch’s theorem can be improved, we can also describe [NSPACE](#_bookmark75) with less than *k* + 2 variables and unbounded iterations.

## Mixing Iterations and Transitive Closure

One question that came up during my research was: Can we ﬁnd a logic that is more restricted than [FO-VAR](#_bookmark66)[*s*(*n*)*, s*(*n*)*/* log(*n*)] but still captures [NSPACE](#_bookmark75)[*s*(*n*)]? This led to the idea of mixing

up iterative procedures with the transitive closure operator. We now show that a success of this approach would mean that there exists a restriction using only iterative procedures that also contains [NSPACE](#_bookmark75)[*s*(*n*)]. As formulas mixing up multiple operators and procedures are more diﬃcult to analyse, it makes this approach unpractical.

There are two ways of combining the transitive closure with iterated formulas. The ﬁrst is sur- rounding a formula in [FO-VAR](#_bookmark66)[*t*(*n*)*, r*(*n*)*/* log(*n*)] with a [TC](#_bookmark64) operator using variables of size *r*(*n*). For this to contain [NSPACE](#_bookmark75)[*s*(*n*)], for any formula of the form ([TC](#_bookmark64)*a,bφ*)(*c, d*) we need a quantiﬁer block, a formula *ϕ* and some constants *e, f, g* such that

([TC](#_bookmark64)*a,b* ([[QB](#_bookmark70)]*t*(*n*)*ϕ*(*e/s*))) (*f, g*) ≡ ([TC](#_bookmark64)*a,bφ*)(*c, d*)

We know that we can simulate any [TC](#_bookmark64) formula with variables of size *r*(*n*) using *r*(*n*) iterations of a quantiﬁer block. This is done by the technique of halving paths used in Savitch’s theorem. An exact formula achieving this is presented in section [4.4](#_bookmark24). Like this, we ﬁnd an equivalent formula:

([TC](#_bookmark64)*a,b* ([[QB](#_bookmark70)]*t*(*n*)*ϕ*(*e/s*))) (*f, g*) ≡ [[QB](#_bookmark70)1]*r*(*n*)[[QB](#_bookmark70)2]*t*(*n*)*ϕ*(*e/s, f/a, g/b*)

The two quantiﬁer blocks can then be merged into one by additionally maintaining a counter variable of size log(*r*(*n*)) such that the quantiﬁer block acts like [[QB](#_bookmark70)1] during the ﬁrst *r*(*n*) round and then acts like [[QB](#_bookmark70)2].

The other case in which we write the [TC](#_bookmark64) formula after the quantiﬁer block can be treated similarly.

([[QB](#_bookmark70)]*t*(*n*) ([TC](#_bookmark64)*a,bϕ*) (*f, g*)) (*e/s*) ≡ ([TC](#_bookmark64)*a,bφ*)(*c, d*)

If the condition *c* · *r*(*n*) *>* log(*t*(*n*)) is satisﬁed for some ﬁxed *c*, we can again merge the quantiﬁer blocks by using the counting trick. As any *t*(*n*) *>* 2*cr*(*n*) can be seen as having an unbounded amount of iterations, we can assume without loss of generality that the condition holds.

In both cases, we have a formula in [FO-VAR](#_bookmark66)[*r*(*n*) + *t*(*n*)*, r*(*n*)*/* log(*n*)]. Now, there are multiple cases

*r*(*n*) ≤ *t*(*n*)**:** Then, [FO-VAR](#_bookmark66)[*r*(*n*) + *t*(*n*)*, r*(*n*)*/* log(*n*)] = [FO-VAR](#_bookmark66)[*t*(*n*)*, r*(*n*)*/* log(*n*)]. This is only interesting in the case where *r*(*n*) *< s*(*n*)

*r*(*n*) *> t*(*n*)**:** Then, [FO-VAR](#_bookmark66)[*r*(*n*) + *t*(*n*)*, r*(*n*)*/* log(*n*)] = [FO-VAR](#_bookmark66)[*r*(*n*)*, r*(*n*)*/* log(*n*)]. This is again only interesting when *r*(*n*) *< s*(*n*).

We thus get that any improvement of our simulation results happens only when we have less than *s*(*n*) variable bits. One result following from this is that if *t*(*n*)*r*(*n*) *< s*(*n*)2, Savitch’s theorem can be improved as

[FO-VAR](#_bookmark66)[*t*(*n*)*, r*(*n*)*/* log(*n*)] ⊆ [DSPACE](#_bookmark73)[*t*(*n*)*r*(*n*)]

This same method also tells us that *t*(*n*)*r*(*n*) ≥ *s*(*n*) because otherwise, we could simulate a [NSPACE](#_bookmark75)[*s*(*n*)] Turing machine deterministically using less than *s*(*n*) space, which is a contradiction to the space hierarchy theorem described in appendix [B.2.3](#_bookmark100).

## Restricting Universal Quantiﬁcation

Another approach that I attempted was restricting universal quantiﬁcation in the [FO-VAR](#_bookmark66) formulas. I called the resulting logic [FO∃-VAR](#_bookmark67). In [FO∃-VAR](#_bookmark67), universal quantiﬁcation is only allowed over boolean variables. We denote boolean variables by *bj*. Further, the requirement of *Mi* to be quantiﬁer-free is dropped. This step was motivated by the fact that nondeterministic Turing machines in essence capture existential quantiﬁcation.

We use a similar idea to the one used in Savitch’s theorem: guessing the middle of a path and

checking if both shorter paths are connected. A formula in [FO -VAR](#_bookmark67)[*s*(*n*)*, s*(*n*)*/* log(*n*)] which is equivalent to a formula in [FO-VAR](#_bookmark66)[1*, s*(*n*)*/* log(*n*)]([TC](#_bookmark64)) of the form (*TCX,Y φ*) (*C, D*) is

[∃](#_bookmark67)

[QB](#_bookmark70) ≡ (∀*b*1*.M*1)(∃*Z*)(∀*b*2)(∃*A, B.M*2)(∃*X, Y .M*3)

*M*1 ≡ ¬(∀*z*([BIT](#_bookmark65)(*X, z*) ↔ [BIT](#_bookmark65)(*Y , z*)) ∨ *φ*(*X, Y* ))

*M*2 ≡ (*b*2 ∧ ∀*z*([BIT](#_bookmark65)(*X, z*) ↔ [BIT](#_bookmark65)(*A, z*)) ∧ ∀*z*([BIT](#_bookmark65)(*B, z*) ↔ [BIT](#_bookmark65)(*Z, z*)))∨

(¬*b*2 ∧ ∀*z*([BIT](#_bookmark65)(*Z, z*) ↔ [BIT](#_bookmark65)(*A, z*)) ∧ ∀*z*([BIT](#_bookmark65)(*B, z*) ↔ [BIT](#_bookmark65)(*Y , z*)))

*M*3 ≡ ∀*z*([BIT](#_bookmark65)(*X, z*) ↔ [BIT](#_bookmark65)(*A, z*)) ∧ ∀*z*([BIT](#_bookmark65)(*B, z*) ↔ [BIT](#_bookmark65)(*Y , z*))

and ﬁnally

for some constant *c*.

[[QB](#_bookmark70)]*cs*(*n*)(*false*)(*C/X, D/Y* )

In [QB](#_bookmark70), the ﬁrst quantiﬁcation and *M*1 mean that the formula breaks with the actual branch being true whenever *X* and *Y* are connected. When this happens, quantiﬁcation is done universally over the empty set, which is deﬁned as true. If no connection exists yet, a middle conﬁguration *Z* is guessed, and both sides are checked by universally choosing *b*2. In *M*2, the formula then checks that *A* and *B* are the endpoints of the path determined by *b*2*, X, Z*, and *Y* . After this, *A* and *B* are copied to *Y* and *X* in *M*3. When iterating this, a path from *C* to *D* is guessed, and every connection between adjacent states is checked.

We deﬁne [FO∀*,* ∃-VAR](#_bookmark68)[*t*(*n*)*, f* (*n*)*, s*(*n*)] to contain all formulas with *t*(*n*) iterations of a quantiﬁer block, extended variables of *f* (*n*) bits which can be universally quantiﬁed, and extended variables of *s*(*n*) log(*n*) bits which can be used in existential quantiﬁcation. Using this, we can make the more general statement that for any function *r*(*n*) ≤ 2*cs*(*n*) for some constant *c*, we can deﬁne a formula in

[FO∀*,* ∃-VAR](#_bookmark68)[*s*(*n*)*/* log(*r*(*n*))*,* log(*r*(*n*))*, s*(*n*)*r*(*n*)*/* log(*n*)]

which simulates a nondeterministic Turing machine using *s*(*n*) space. These formulas are very similar to the one deﬁned for [FO∃-VAR](#_bookmark67). The only diﬀerence lies in how the paths are split. Instead of splitting the paths into two parts, they are split into *r*(*n*) parts. This is done by guessing *r*(*n*) − 1 middle conﬁgurations at once in an existential extended variable and choosing which part to check

using the universal extended variable. Thus, for a path of length 2*cs*(*n*), the formula needs at most

( )

iterations.

log*r*(*n*)

(2*cs*(*n*)) =

log 2*cs*(*n*) log(*r*(*n*))

= *cs*(*n*)*/* log(*r*(*n*))

As for simulating this formula with a Turing machine, the product of the number of iterations and the size of the extended variables are important, we do not gain any tighter results from this.

## Alternating Bounds

In the proof of Savitch’s theorem, alternating Turing machines were used as an intermediate step for proving that [NSPACE](#_bookmark75)[*s*(*n*)] is a subset of [DSPACE](#_bookmark73)[*s*(*n*)2]. We can also add another interme- diate step using [FO-VAR](#_bookmark66)[*s*(*n*)*, s*(*n*)*/* log(*n*)] with the method shown in section [4.4](#_bookmark24). It is also quite easy to simulate any formula in [FO-VAR](#_bookmark66)[*t*(*n*)*, s*(*n*)*/* log(*n*)] with an alternating Turing machine in [ATSR](#_bookmark82)[*t*(*n*)*s*(*n*)*, s*(*n*)*, t*(*n*)], where in this new class [ATSR](#_bookmark82), we specify time, space and reversals. By reversals, we mean the number of times we switch from universal states to existential states and back. A Turing machine in [ATSR](#_bookmark82)[*t*(*n*)*s*(*n*)*, s*(*n*)*, t*(*n*)] does this by writing the current values of all variables on the tape while iterating and evaluating the quantiﬁer-free ﬁrst-order formulas on the way, which can be done eﬃciently in terms of space.

We now investigate two ways of simulating an [ATSR](#_bookmark82)[*t*(*n*)*, s*(*n*)*, r*(*n*)] Turing machine using [FO-VAR](#_bookmark66).

The ﬁrst one is a straightforward simulation of each step of the computation. For this, our formula encodes the whole state of the Turing machines in a tuple of *s*(*n*) bit extended variables and iteratively guesses the next conﬁguration. We also assume without loss of generality that a predicate *existential*(*X*) exists and that there exists exactly one accepting state, denoted by *E*.

[QB](#_bookmark70) ≡ (∀*b.M*1)(∃*A.φ*(*X, A*))(∀*B.M*2)(∃*X.*∀*z*([BIT](#_bookmark65)(*X, z*) ↔ [BIT](#_bookmark65)(*B, z*))) *M*1 ≡ ¬∀*z*([BIT](#_bookmark65)(*X, z*) ↔ [BIT](#_bookmark65)(*E, z*))

*M*2 ≡ (*existential*(*X*) → ∀*z*([BIT](#_bookmark65)(*A, z*) ↔ [BIT](#_bookmark65)(*B, z*))) ∧ *φ*(*X, B*)

with

([[QB](#_bookmark70)]*t*(*n*)(*false*)) (*S/X*)

This gives us a formula in [FO-VAR](#_bookmark66)[*t*(*n*)*, s*(*n*)*/* log(*n*)] which simulates a computation of an alternating Turing machine. Some of these *Mi* are not quantiﬁer-free. The quantiﬁers could be moved out of the *Mi*, but would make the formula less readable.

This formula works as follows: The formula starts by checking if the accepting conﬁguration has already been reached using *M*1. Next, if the simulation is not done yet, the formula existentially guesses the next conﬁguration, restricting this quantiﬁcation to only those conﬁgurations which are directly connected to the actual one. In *M*2, either this existentially guessed next state is retained if the simulated Turing machine is in an existential state or the existentially quantiﬁed state is ignored

and a universal quantiﬁcation over all states which are connected to the current one is done. The last thing the formula does is copying *B* to *X*, making it the new start of the path.

By combining the above techniques for switching between universal and existential quantiﬁcation with the path-halving method used in Savitch’s theorem, we get an even stronger result.

From now on, we refer to conﬁgurations as vertices to emphasize that we are searching for paths on a graph. On a high level, our plan is to do the following:

1. From the actual starting vertex, ﬁnd all reachable vertices using only vertices which have the same type (existential or universal) as the starting vertex.
   1. For existential states, we can use the normal technique of halving paths naively.
   2. For universal states, we need to quantify over all ending vertices and then consider two cases:
      * We claim that there is no path going to this vertex that uses only universal states. Then, we need to show that for each middle vertex, one of the two new paths still does not exist.
      * We claim that a path exists. In this case, we can try to ﬁnd a path normally.
2. Repeat with all the reached vertices.

Before presenting the formula, we ﬁrst think of how many iterations are needed. Because of the halving trick, log(*a*) iterations are used for a path of length *a*. By summing over all path segments with only one quantiﬁcation type, we can bound the total number of iterations. As we are simulating an alternating Turing machine which makes at most *t*(*n*) steps, the total of all path segments can not exceed *t*(*n*). To get an upper bound on the number of iterations, we can use Jensen’s inequality ([[Mar19](#_bookmark47)], p.28). We get

I:*r*(*n*) log(*aj* )  I:*r*(*n*) *aj* 

*j*=0

*j*=0

*r*(*n*) ≤ log  *r*(*n*) 

≤ log ( *t*(*n*) 1

*r*(*n*)

because of the condition that *r*(*n*) *a t*(*n*). Multiplying this with *r*(*n*), we get that a path can require a maximum of *r*(*n*) log (  *t*(*n*) ) iterations. Thus, the formula is in

*j*=0

I: *j* ≤

*r*(*n*)

[FO-VAR](#_bookmark66) [*r*(*n*) log (*t*(*n*)*/r*(*n*)) *, s*(*n*)*/* log(*n*)].

The following formula simulates an [ATSR](#_bookmark82)[*t*(*n*)*, s*(*n*)*, r*(*n*)] Turing machine:

[QB](#_bookmark70) ≡ (∀*b*0*.M*1)(∃*b*0*.N*1)(∃*b*1*.M*2)(∃*Ze.M*3)(∀*Z.N*2)

(∃*b*2*e*)(∀*b*2*.N*3)(∃*A, B.M*4)(∀*S, E.M*5)(∃*A, B.M*6)(∃*I, X.M*7)(∃*b*0*.N*4)(∃*bpath.N*5)

*M*1 ≡ *bpath* → ¬(*I* = *Y* ∧ (*S* = *E* ∨ *φ*(*S, E*)))

*N*1 ≡ ¬*bpath* → ¬(*S* = *E* ∨ *φ*(*S, E*))

*M*2 ≡ ¬*b*1 ↔ (*S* = *E* ∨ *φ*(*S, E*))

*M*3 ≡ *existential*(*Ze*) ↔ *existential*(*X*)

*N*2 ≡ (*bpath* → *Z* = *Ze*) ∧ (*existential*(*Z*) ↔ *existential*(*X*))

*N*3 ≡ ¬*bpath* → *b*2 = *b*2*e*

*M*4 ≡ (*b*1 ∧ ((¬*b*2 ∧ *A* = *S* ∧ *B* = *Z*) ∨ (*b*2 ∧ *A* = *Z* ∧ *B* = *E*)))∨

(¬*b*1 ∧ *A* = *I* ∧ (*B* = *Y* ∨ (*existential*(*B*) ↔ ¬*existential*(*I*))))

*M*5 ≡ *S* = *A* ∧ (((*b*1 ∨ *existential*(*S*)) ∧ *E* = *B*)∨

(¬(*b*1 ∨ *existential*(*S*)) ∧ (*E* = *Y* ∨ (*existential*(*E*) ↔ ¬*existential*(*S*)))))

*M*6 ≡ (¬*b*1 ∧ *S* = *A* ∧ *E* = *B*) ∨ (*b*1 ∧ *X* = *A* ∧ *I* = *B*) *M*7 ≡ *X* = *A* ∧ *I* = *B*

*N*4 ≡ *b*0 = *bpath*

*N*5 ≡ (*b*1 ∧ *bpath* = *b*0) ∨ (¬*b*1 ∧ (*existential*(*X*) → *bpath*))

and

([[QB](#_bookmark70)]*r*(*n*) log(*t*(*n*)*/r*(*n*))(¬*bpath*)) (true*/bpath, AS/S, AS/X, AS/I, AS/S, AS/E, AE/Y* )

being the ﬁnal formula.

The meanings of the variables are explained below:

*Y* **:** the accepting vertex

*X***:** the starting vertex in the actual component

*I***:** the vertex for which it is being checked whether a path from *X* exists

*S***:** the start of the path segment which is currently checked

*E***:** the end of the path segment which is currently checked

*Z, Ze***:** the guessed middle vertex of a path segment from *S* to *E A, B, b*0**:** variables used for copying or as dummies

*bpath***:** whether the formula claims that there is a path or that there is no path from *S* to *E b*1**:** remembers whether *S* and *E* are connected

*b*2*, b*2*e***:** which half of the *S* → *Z* → *E* path is currently checked

As this formula is quite intricate, it is proposed to read it four times, each time focusing on a diﬀerent stage of the search.

1. First, we look at the formula in the easiest case: It is trying to ﬁnd a path from *X* to *I* using *S* and *E*, which are not directly connected. In this case, both *bpath* and *b*1 are true and [QB](#_bookmark70) works like this: In *M*1, the formula checks and sees that it is not done, as *S* and *E* are not connected. As *bpath* is true, nothing is done in *N*1. In *M*2, *b*1 is set to true. In *M*3, a middle vertex between *S* and *E* is guessed which has the same type as *X*. *N*2 copies *Ze* into *Z*. Next, a side of the path is chosen. As *bpath* is true, in *N*3 the choice made for *b*2*e* is ignored and the side is universally chosen. *M*4 copies the relevant new starting and ending vertices to *A* and *B*. Because *b*1 is true,

*A* and *B* are copied to *S* and *E* in *M*5. In *M*6 and *M*7 combined, *X* and *I* are copied to *A* and

*B* and back again. The same is done with *bpath* using *N*4, *N*5, and *b*0. In the end, the formula gets a new conﬁguration with the guessed middle vertex replacing either the former ending or starting vertex. If no path exists, this continues until in the formula ¬*bpath* this branch is marked as false.

1. The second case we look at is the case where the formula claims that no path exists. Here, *bpath* is false. Because of this, in *M*1 nothing is done. In *N*1, the formula checks if it found a connection, and if it did, it sets this logical branch to false because of the existential quantiﬁcation over the empty set. If *N*1 is true, the formula goes on to *M*2, where it establishes *b*1. As *M*3 is ignored in *N*2, it has no eﬀect. In *N*2, a universal quantiﬁcation over all middle vertices *Z* which could lie on a path between *S* and *E* is done. The next two quantiﬁers quantify existentially which half of *S* → *Z* → *E* is not connected. This half must exist, as otherwise, a path from *S* to *E* exists, which would contradict ¬*bpath*. Because *b*1 is again true, the relevant part of the path is copied to *A* and *B* in *M*4 and back to *S* and *E* in *M*5. In *M*6, *M*7, *N*4, and *N*5 the formula does the same as in the ﬁrst case, which is not changing *X*, *I*, and *bpath*. If no path exists, this continues until the formula ¬*bpath* is hit. At this point, this branch of the formula is marked as true, as eﬀectively, no path from *X* to *I* exists by making these choices.
2. The third possibility we need to consider is that the formula is currently trying to ﬁnd a path from *X* to *I* in a universal component and that it succeeds in ﬁnding a connection between *S* and *E*. In this case, *bpath* is true and *b*1 is false. In *M*1, if *I* = *Y* , this logical branch is done as the simulation reached the accepting vertex. Otherwise, *N*1 can be disregarded, and in *M*2, *b*1 is set to false. The quantiﬁcation for *Z* and *b*2 can be ignored as these variables are not used in this iteration of the formula. In *M*4, *A* and *B* are guessed such that *A* is the new starting vertex corresponding to the old ending vertex *I*, and *B* is either the accepting vertex *E* or a universal vertex. Then, *A* is copied to *S* in *M*5, marking it as the starting vertex. Further, because the quantiﬁcation type changed, a valid *S* must now be existential. Thus, *B* is also copied to *E* directly. After that, in *M*6 and *M*7, the formula copies *S* to *X* and *E* to *I* via *A* and *B*. The last change happens in *N*5, where *bpath* is set to true. Thus, the variables are set for a new block of existential states.
3. The last case is similar to the third one, with the diﬀerence that the current vertex *X* is exist- ential. In this case, *bpath* is true and *b*1 is false. In *M*1, if *I* = *Y* , this logical branch is done as the simulation reached the accepting vertex. Otherwise, *N*1 can be disregarded, and in *M*2, *b*1

is set to false. The quantiﬁcation for *Z* and *b*2 can be ignored as these variables are not used in this iteration of the formula. In *M*4, *A* and *B* are guessed such that *A* is the new starting vertex corresponding to the old ending vertex *I*, and *B* is either the accepting vertex *E* or an existential vertex. Then, *A* is copied to *S* in *M*5, marking it as the starting vertex. Further, because we changed the quantiﬁcation type, a valid *S* must now be universal. Thus, all vertices which are either the accepting vertex or are existential are universally quantiﬁed for *E*. After that, in *M*6 and *M*7, the formula copies *S* to *X* and *E* to *I* via *A* and *B*. The last change happens in *N*5, where the formula quantiﬁes existentially if it claims that a path from *X* to *I* exists or that no such path does. Thus, the variables are set for a new block of universal states.

The complete formula sets the starting values of the variables. By setting *bpath* = true and all of *X, S, E* and *I* to the staring conﬁguration *AS*, either case three or four is triggered. This is what we want as after this, the component including *AS* will be searched. The accepting vertex *Y* is set to the desired ending *AE*.

These results allow us to enclose an [ATSR](#_bookmark82) class between two [FO-VAR](#_bookmark66) classes with only logarithmic overhead. Thus:

[FO-VAR](#_bookmark66) [*a*(*n*)*, s*(*n*)*/* log *n*] ⊆ [ATSR](#_bookmark82)[*a*(*n*)*s*(*n*)*, s*(*n*)*, a*(*n*)] ⊆ [FO-VAR](#_bookmark66) [*a*(*n*) log(*s*(*n*))*, s*(*n*)*/* log *n*]

# Conclusion and Direction

## Conclusion

In this document, multiple logics which have a relation to Savitch’s theorem were investigated. This has led us to a better notion of how these logics are related to [NSPACE](#_bookmark75)[*s*(*n*)] and thus also to linear bounded automata and context-sensitive languages. We found out that context-sensitive languages correspond to reachability querries on certain graphs. Further, section [4.2](#_bookmark21) seems to indicate that they are not local structures as remembering only local transformations is not enough to capture them. In section [4.4](#_bookmark24), we have seen that even very little boolean universal quantiﬁcation is probably too much for the existential nature of [NSPACE](#_bookmark75)[*n*] Turing machines to handle.

In the beginning, the theory of formal languages and the Chomsky hierarchy were introduced. Next, the main tool of Descriptive Complexity used in this document, Complexity theory, was presented. Afterwards, Savitch’s theorem and its proof were explained. From then on, this theorem was used multiple times. Last, the equivalence between context-sensitive languages and monadic second-order logic was shown.

Personally, I investigated multiple ideas I got while reading the theory.

The ﬁrst one was a direct transformation of the results of the equivalence between second-order transitive closure logic and nondeterministic Turing machines with bounded space to ﬁrst-order logic. This approach did not contain any new insights, but it was able to showcase the close relationship between ﬁrst- and second-order logic.

After this, I wanted to explicitly prove that some methods could not work. This branch of my work started with the investigation of the proof that [DSPACE](#_bookmark73) is equivalent to [VAR](#_bookmark69)[*k* + 1]. There, I showed that naively applying the proof to nondeterministic Turing machines gives wrong results and that a natural extension by remembering decisions is suboptimal. A full proof was impossible to make as it is unclear how to deﬁne a generalization of the proof for [DSPACE](#_bookmark73). The diﬃculty of this problem is illustrated by Savitch’s theorem, which has not been improved for a long time and thus makes it seem unlikely that a simulation of [NSPACE](#_bookmark75) Turing machines with subquadratic [DSPACE](#_bookmark73) Turing machines is possible.

Subsequently, I tried was mixing transitive closure logic with iterative logic. This turned out not to be very powerful, as any interesting result in one of the mixed logics would imply one in pure iterative logic. Nevertheless, I proved upper and lower bounds using Savitch’s theorem and the space hierarchy theorem.

In an attempt to make [FO-VAR](#_bookmark66) less powerful, I also tried restricting universal quantiﬁcation. A new formula made this possible. As for all the other restrictions, I was unable to ﬁnd a way to simulate formulas in my restriction of [FO-VAR](#_bookmark66) using [NSPACE](#_bookmark75) Turing machines. The generalization of this idea by adding more bits to the extended variables and decreasing the iteration count gave

a space-time tradeoﬀ but did not solve the inherent problem of simulating this formula by [NSPACE](#_bookmark75) Turing machines.

The most complicated result I managed to ﬁnd does not have any tight connection to context- sensitive languages. It concerns the simulation of alternating Turing machines using [FO-VAR](#_bookmark66). There, I found quite a tight containment, which bounds some alternating Turing machine classes from above and below by a factor of log(*s*(*n*)). I used a combination of diﬀerent techniques used in previous results to get the logical formula. This result could lead to a path for describing alternating space-time classes exactly using an equivalent logic.

### Working Process

Generally, I learned a lot about logic and Descriptive Complexity during this project. It also provided me with insight into how scientiﬁc research is done, and I was able to immerse myself in a domain that fascinates me. I knew that the probability of obtaining signiﬁcant results was low, and indeed I worked on many diﬀerent ideas, but most of the time hit a dead end.

At the beginning of the working process, I spend a lot of time reading multiple books about Descriptive Complexity and logic. I then researched more papers on the topic and began to understand the lower and upper bounds described there. This took more time than expected, but I still was able to start investigating early enough. As this was my ﬁrst scientiﬁc research project, I did not know how to approach it, so I wrote to from the ETH Zürich. His tips were helpful and gave me a starting point for my research. During the months that followed, I investigated multiple ideas and read a lot of papers concerning things I wanted to prove or needed for my proofs.

Early on, and I discussed that in order to help reduce the stress before handing in, I should start writing the theoretical part of this document. I did this during the summer break, experiencing some diﬃculties in the simpliﬁcation of proofs while still maintaining their correctness. Afterwards, I continued doing research. During this time, I found multiple results which are presented in chapter [4](#_bookmark19), but also spent a lot of time working on hierarchies without notable results. Two months before handing in, the International Olympiad in Informatics took place. My time plan did not account for this and all the preparation required very well, so I had to work quite a lot in the last two weeks to write down my personal contribution.

The main diﬃculties I had during this project are the following:

* + - * First, in the past I often felt that my proofs were not explained in a way that was clear and concise. This made it diﬃcult for the reader to follow the arguments and understand the un- derlying thought process. In this document I focused on this point and was able to improve the quality of the proofs.
      * Second, the methods of research were diﬃcult to ﬁnd. I found it diﬃcult to ﬁnd any resources on this matter in the highly abstract ﬁeld of mathematics. I reached out to multiple professors at the University of Basel and the ETH Zürich, but unfortunately only one of the people I contacted had a little time to help me along the right path.
      * Finally, the time management is a topic I need to improve. I spent a lot of time on ideas that

were clearly not working and thus did not have time to investigate everything I wished to study. Also, I started writing down my own research very late, which generated some stress in the end.

## Directions

Multiple topics could be investigated further and would provide an even deeper understanding of the inner workings of Savitch’s theorem.

In all proofs concerning equivalences between transitive closure logics and [NSPACE](#_bookmark75)[*s*(*n*)], it is as- sumed that *s*(*n*) is at most polynomial. This is still a profound limitation, and ﬁnding a generalization could lead to new insights that could help in the polynomial case and in the end also in the case where *s*(*n*) = *n*, capturing exactly the context-sensitive languages.

In the context of Savitch’s theorem, it would also be interesting to investigate the computation graphs of [NSPACE](#_bookmark75) and [DSPACE](#_bookmark73) Turing machines. This has been discussed on the theoretical com- puter science stack exchange in [[Bar10](#_bookmark32)]. A motivation for this is that the number of nodes in both graphs is the same, only the number of edges varies.

Another approach that could be investigated is looking at normal forms. One that seems to have some potential is the Kuroda normal form described in [[Kur64](#_bookmark43)]. This approach could yield a semantic restriction similar to the one for context-free languages described in appendix [B.2.5](#_bookmark102).

Further, one could investigate hierarchies of transitive closure logics, generalized quantiﬁers and alternating time-space complexity classes. “A double arity hierarchy theorem for transitive closure logic”[[GH96](#_bookmark37)] by Grohe shows a strict hierarchy for logics which combine the transitive closure operator with generalized quantiﬁers but is limited to unordered structures. As string structures are inherently ordered, this does not generalize to cases where we use Turing machines. Some steps in the direction of hierarchies for ordered structures were made in the chapter about proving lower bounds in [[Imm99](#_bookmark42)].

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# List of Acronyms

**FO** ﬁrst-order logic

**SO** second-order logic

**MSO** monadic second-order logic

**TC** transitive closure

**BIT** logical relation telling us if the *xth* bit of some variable is 1

**FO-VAR** iterative logic with restricted variable bits

**FO**∃**-VAR** iterative logic with restricted existential variable bits and boolean universal variables

**FO**∀*,* ∃**-VAR** iterative logic with restricted existential variable bits and restricted universal variable bits

**VAR** unrestricted iterative logic with restricted variable bits

**QB** quantiﬁer block

**STRUCT** set of all structures over a vocabulary

**DTIME** deterministic time complexity class

**DSPACE** deterministic space complexity class **NTIME** nondeterministic time complexity class **NSPACE** nondeterministic space complexity class **ATIME** alternating time complexity class

**ASPACE** alternating space complexity class

**P** deterministic polynomial time complexity class

**NP** nondeterministic polynomial time complexity class

**PSPACE** deterministic polynomial space complexity class

**NPSPACE** nondeterministic polynomial space complexity class

**ATSR** alternating time-space-reversals complexity class

**IT** information technologies **GPS** Global Positioning System **AI** artiﬁcial intelligence

# Mathematical Background

The deﬁnitions are taken from the lectures Discrete Mathematics in Computer Science [[HR22](#_bookmark38)] and Theory of Computer Science [[Rög23](#_bookmark54)], as well as from the book Descriptive Complexity [[Imm99](#_bookmark42)].

## Set Theory

**Set** An unordered collection of distinct elements, written with curly braces {}. If an element *x* is in some set *A*, we write *x* ∈ *A*

**Tuple** An ordered collection of elements, written with pointed braces ⟨⟩

**Set operations** There are multiple ways to form new sets from already existing sets:

**Union** denoted as ∪. An element is in *A* ∪ *B* if and only if it is in *A* or *B*

**Intersection** denoted as ∩. An element is in *A* ∩ *B* if and only if it is in *A* and *B*

**Cartesian product** denoted as ×. *A* × *B* is the set of 2-tuples ⟨*a, b*⟩ with an element *a* of *A*

and an element *b* of *B*

**Cartesian power** *Ak* denotes the Cartesian product of *A* with itself repeated *k* times

**Power set** denoted as P(*A*). Contains all subsets of *A*

## First-Order Logic

**Variable** A variable is an element that can have a value from a set. We denote tuples of variables by

*x* = ⟨*x*1*, . . . , xn*⟩

**Universe** The set over which variables and constants can range

**Relation** A relation of arity *k*, *R*(*x*1*, . . . , xk*) can be either true or false for any *k*-tuple of variables. In this document, we always consider an equality relation =, an ordering relation ≤, and [BIT](#_bookmark65)(*x,* 1*y*), which means that the *yth* bit of *x* is 1 in binary notation, to be present

**Vocabulary** A tuple *τ* = ⟨*Ra*1 *, . . . , Rar , c*1*, . . . , cs*⟩ of relations *Ri* with arity *ai* and constants *cj*[1](#_bookmark89)

1

*r*

**Structure** A tuple A = ⟨|A|*, R*A*, . . . , R*A*, c*A*, . . . , c*A⟩ where |A| is the universe, the constants are

1

*r*

1

*s*

assigned a value from |A|, and each relation is assigned a truth value for each *ai*-tuple in |A|*ai* . The set of all structures for a given universe is denoted as [STRUCT](#_bookmark71)[*τ* ]

1We omit functions, which are included in most textbook deﬁnitions, as they can be simulated by a relation in our case

**First-order formula** A ﬁrst-order formula is inductively deﬁned as follows:

**Atoms** Any formula of the form *R*(*x*1*, . . . , xk*) for some relation of arity *k* is called an atomic formula

**Conjunction** If *φ* and *ψ* are formulas, (*φ* ∧ *ψ*) is a formula **Disjunction** If *φ* and *ψ* are formulas, (*φ* ∨ *ψ*) is a formula **Negation** If *φ* is a formula, ¬*φ* is a formula

**Existential quantiﬁcation** If *φ* is a formula, ∃*xφ* is a formula

**Universal quantiﬁcation** If *φ* is a formula, ∀*xφ* is a formula

**Free variables** A variable in a formula which occurs at least once without being bound by a quantiﬁer whose scope surrounds it

**Semantics** For any structure, we can assign a truth value to any formula over the corresponding vocabulary. If the formula contains free variables, these need to be assigned a value from the universe ﬁrst. We say A satisﬁes *ϕ*, denoted as A |= *ϕ*, if and only if *ϕ* is true under the interpretation of the constants and relations in A. This truth value is inductively assigned to all formulas as follows:

**Atoms** For a formula *ϕ* of the form *R*(*x*1*, . . . , xk*), we have A |= *ϕ* if and only if the interpretation of the relation *R*A maps ⟨*x*1*, . . . , xk*⟩ to true

**Conjunction** We have A |= (*φ* ∧ *ψ*) if and only if A |= *φ* and A |= *ψ* **Disjunction** We have A |= (*φ* ∨ *ψ*) if and only if A |= *φ* or A |= *ψ* **Negation** We have A |= ¬*φ* if and only if A |̸= *φ*

**Existential quantiﬁcation** We have A |= ∃*xφ* if and only if there exists a *y* ∈ |A| such that

A |= *φ*(*y/x*), where *φ*(*y/x*) denotes *φ* with any occurrence of *x* replaced by the element *y*

**Universal quantiﬁcation** We have A |= ∀*xφ* if and only if for all *y* ∈ |A| we have A |= *φ*(*y/x*)

**Logical operator** A logical operator can create a new formula from one or more existing formulas.

One example is the conjunction, which combines two existing formulas.

**First-order queries** A ﬁrst-order query is a map from structures over one vocabulary *σ* to structures over another vocabulary *τ* . The mapping is done in such a way that ﬁrst-order formulas deﬁne the universe, which is a subset of |A|*k* for some *k*, the relation symbols, and all the constants. For a more formally thorough deﬁnition see [[Imm99](#_bookmark42)]

**Isomorphism** An isomorphism is a map *I* : |A| → |B| with A*,* B over the same vocabulary which satisﬁes the following properties:

* *I* is bijective
* for every available relation *Ri* of arity *ai* and every *ai*-tuple *e* = ⟨*e*1*, . . . , ea* ⟩ in |A|*ai* , we

*i*

have

*R*A(*e*1*, . . . , ea* ) ⇔ *R*B(*I*(*e*1)*, . . . , I*(*ea* ))

*i*

*i*

*i*

*i*

* for every constant symbol *cj*, we have *I*(*c*A) = *c*B

*j j*

If such an *I* exists for two structures A and B, we write A ∼= B

## Second-Order Logic

In second-order logic, we extend the capabilities of ﬁrst-order logic with the ability to quantify over relations. We thus also need to extend our deﬁnitions.

**Second-order variables** A relation that is not given in the vocabulary and can be substituted with a speciﬁc relation

**Second-order formula** In addition to the inductive rules of the ﬁrst-order formulas, we can quantify over second-order formulas

**Second-order existential quantiﬁcation** If *φ* is a formula, then ∃*V φ* is a formula

**Second-order universal quantiﬁcation** If *φ* is a formula, then ∀*V φ* is a formula

**Second-order semantics** We also need to extend the ﬁrst-order semantics

**Second-order existential quantiﬁcation** We have A |= ∃*V φ* if and only if there exists a relation *U* over |A| such that A |= *φ*(*U/V* ), where *φ*(*U/V* ) denotes *φ* with any occurrence of *V* replaced by *U*

**Second-order universal quantiﬁcation** We have A |= ∀*V φ* if and only if for all relations *U*

over |A| we have A |= *φ*(*U/V* )

## Turing Machines

Turing machines are the most common model of computation.

**Informal deﬁnition** A Turing machine is an automaton with a ﬁnite number of states and an inﬁnite tape. Using a read/write head, which can read one symbol on the tape, modify one symbol on the tape and move left or right, a Turing Machine can compute functions

**Formal deﬁnition** Formally, a Turing machine is a 7-tuple *M* = ⟨*Q,* Σ*,* Γ*, δ, q*0*, qaccept, qreject*⟩, where

*Q* is the set of states

Σ is the alphabet of the input word

Γ is the set of symbols which can be written or read on the tape, which we call the tape alphabet

*δ* is the transition function, with *δ* : Γ × *Q* → Γ × *Q* × {*L, R*}. This means that when a Turing machine is in state *n* and reads *a* on the tape, *δ* tells us to which state we should transition, which symbol we should write and which direction we should move the read/write head

*q*0 the start state

*qaccept* the accepting state

*qreject* the rejecting state

**Turing computation** In the beginning, the Turing machine is in the start state, the input word is written on a consecutive part of the tape and the read/write head is on the ﬁrst character of the input word. In the following steps, the Turing machine state changes according to the transition function. If at some point the Turing machine enters the accepting or the rejecting state, the computation halts, and the Turing machine is said to have accepted / rejected the input. It can happen that the Turing machine continues indeﬁnitely or loops. In that case, we also say that it has rejected its input. In this document, we ignore the tape content after the computation and focus on decision problems.

**Decidability** If a Turing machine halts on all inputs, we say that it decides a problem, as we can always be sure that the machine will accept or reject an input in ﬁnite time.

**Nondeterministic Turing machine** We can extend the transition function *δ* to allow multiple transitions from a given state. Formally, we then have *δ* : Γ × *Q* → P(Γ × *Q* × {*L, R*}). If at some point multiple transitions are possible from the current state, we can take any of them. If there exists any computational path which leads to an accepting state, the nondeterministic Turing machine accepts. This is not analogous to how real sequential computers work, but allows interesting results, and is not more powerful than a normal deterministic Turing machine.

**Space/Time-constructible functions** A function *f* (*n*) is time constructible if there exists a Turing machine which on input 1*n* writes *f* (*n*) in binary on its tape in time *f* (*n*). Space-constructible functions are deﬁned analogously.

**Church-Turing thesis** The Church-Turing thesis states that anything that can be done on a real- world computer can be done using a Turing machine.

# Mathematical Context and Further Proofs

## Formal languages

### Regular Languages

The regular languages have the most restricted type of grammar. Formally, any regular language can be described by a grammar with rules in *V* × (Σ ∪ Σ*V* ∪ *ε*). This means that we only have exactly one non-terminal on the left-hand side and the right-hand side is either a terminal, the empty word, or a terminal symbol followed by a non-terminal symbol.

**Example B.1.** Consider the grammar ⟨{*S, O*}*,* {*a*}*, R, S*⟩ with

*R* = *S* → *aO, S* → *ε, O* → *aS*

f \_

The generated language consists of exactly all words with even length.

The regular languages have been studied quite thoroughly and have multiple equivalent formalisms:

* + - * The language is recognized by a Deterministic ﬁnite automaton, which processes the input word one character at a time
      * The language can be decided by a read-only Turing machine, that is a Turing machine that can not modify its tape
      * The language can be described by a regular expression

For a more in-depth analysis of regular languages and equivalent formalisms refer to appendix [B.2.4](#_bookmark101) and [[Str94](#_bookmark57)].

### Context-Free Languages

Context-free languages extend the regular languages by allowing arbitrary right-hand sides for the rules of the deﬁning grammar. Formally, that gives us rules in *V* × (Σ ∪ *V* )∗. Most valid arithmetic expressions, logical formulas and formally correct code in programming languages are context-free, as we can see the non-terminal symbols as types which are then converted to speciﬁc expressions of that type.

**Example B.2.** Consider the grammar ⟨{**Exp***,* **NumF***,* **Num**}*,* {0*,* 1*,* (*,* )*,* −*,* +}*, R,* **Exp**⟩ with

**Exp** → (**Exp** − **Exp**)*,* **Exp** → (−**Exp**)*,*

*R* = **Num** → 0**Num***,* **Num** → 1**Num***,*

**Exp** → **NumF***,* **Exp** → (**Exp** + **Exp**)*,*

**Num** → *ε,* **NumF** → 0*,*



**NumF** → 1**Num**





This generates the language of all well-formed formulas using addition and subtraction over binary numbers. For clarity, **Exp** denotes an arbitrary expression, **NumF** any number without leading zeroes and **Num** any number (possibly empty or with leading zeroes).

Those languages have less known formalisms, the Push-Down Automaton (again see [[Rög23](#_bookmark54)]) being the most common. For a characterization of the context-free languages using logic, see appendix [B.2.5](#_bookmark102).

### Recursive Languages

The recursive languages are the most general languages in the hierarchy, as they don’t have any restrictions on the rules. It can be shown that this set of languages is equivalent to the languages recognizable by a Turing machine. By the Church-Turing thesis, this means that these are exactly the languages that can be computed by any of our computers and algorithms. Thus, we have a huge number of equivalent formalisms, including a RAM machine, while-programs and lambda calculus.

It is worth noting that there are languages which are not recursive. One of the most important examples of these languages is the set of all (descriptions) of Turing machines which halt on every input, also known as the halting problem. For the characterization using logic, again refer to appendix [B.2.7](#_bookmark106).

## Descriptive Complexity

### Ehrenfeucht-Fraïssé Games

Ehrenfeucht-Fraïssé games are combinatorial games which have a strong connection to ﬁrst-order formulas and their extensions. Using these games, it is often possible to show inexpressibility results for certain problems in some logic L.

As a motivation, we can look at what it means for a formula to hold on some structure. Assume the formula has the form ∀*xφ*. This can be seen as some opponent choosing some element *a* ∈ |A| and us then needing to show that *φ*(*x/a*) holds. The case where the formula has the form ∃*xψ* can be treated similarly, but we can select the element ourselves.

**Deﬁnition B.1** (Ehrenfeucht-Fraïssé Game)**.** The *k-pebble Ehrenfeucht-Fraïssé Game* G*k* is played by two players, the Spoiler and the Duplicator, on a pair of structures A and B using *k* pairs of pebbles. In each move, the Spoiler places one of the remaining pebbles on an element of one of the two structures. Then, the Duplicator tries to match the move by placing the corresponding pebble on an element of the other structure. For a pebbled element *x* in A, we denote by *p*(*x*) the corresponding

pebbled element in B. We say that the Duplicator wins the *k*-pebble Ehrenfeucht-Fraïssé Game on

A*,* B if after the *k* rounds, the map *i* : |A| → |B| deﬁned as

*c*

B *x* = *c*A

*j*

 *j*

*i*(*x*) = *p*(*x*) *x* is pebbled

*undef ined*

 otherwise

is a partial isomorphism. A partial isomorphism is an isomorphism over some subset of the universe with all relations restricted to that subset.

In the Ehrenfeucht-Fraïssé Game, the Spoiler wants to show that A and B are diﬀerent, whereas the Duplicator wants to show that they are equivalent.

As this is a zero-sum game with full information, one of the two players must have a winning strategy. It can be proven that the Duplicator has a winning strategy for the *k*-pebble Ehrenfeucht- Fraïssé on A and B if and only if A and B agree on all formulas with at most *k* nested quantiﬁers. We can use these facts to prove the inexpressibility of some problems in ﬁrst-order logic. Assume we have a decision problem C. If we can exhibit two structures A*k* ∈ C and B*k* ∉ C for each *k* and a winning strategy for the Duplicator on these two structures, we show that no ﬁrst-order formula deﬁning C can exist. This methodology can be extended to other logics by adding new moves or restrictions to the game.

### Reduction and Completeness

A reduction can informally be seen as a method of using a problem we already solved to solve a new problem by converting this new problem into an instance of the old problem. These reductions can be very useful to deﬁne complete problems for complexity classes, which in turn enables us to prove theorems for all problems of a speciﬁc complexity class.

**Deﬁnition B.2** (ﬁrst-order reduction)**.** Let C be a complexity class and *A* and *B* be two problems over vocabularies *σ* and *τ* . Now suppose that there is some ﬁrst-order query *I* : STRUC[*σ*] → STRUC[*τ* ] for which we have the following property:

A ∈ *A* ⇔ *I*(A) ∈ *B*

Then *I* is a ﬁrst-order reduction from *A* to *B*, denoted as *A* ≤*fo B*.

First-order reductions can then be used to show that some problem is also a member in some complexity class, as in most complexity classes, we can compute the ﬁrst-order query, and then we are left with a problem that we already know is in the required class. The converse can also be shown: for some problem *B* which is not in some complexity class C, if we have *B* ≤*fo A*, then *A* is also not in C, as otherwise, *B* would also be in C, which is a contradiction.

Using the reductions, we can deﬁne completeness.

**Deﬁnition B.3** (Completeness via ﬁrst-order reductions for Complexity Class C)**.** We say some problem *A* is complete for C via ≤*fo* if and only if

* *A* ∈ C
* for all *B* ∈ C, we have *B* ≤*fo A*

Informally, a complete problem captures the essence of the complexity class. Further, they have an application in some proofs of equivalences between complexity classes C and logics L. These proofs follow the following steps as in [[Imm99](#_bookmark42)]:

1. Show that L ⊆ C by providing a way to convert any formula *φ* ∈ L into an algorithm in C.
2. Find a complete problem *T* for C via ﬁrst-order reductions.
3. Show that L is closed under ﬁrst-order reductions, that is that any formula can be extended by ﬁrst-order quantiﬁers and boolean connectives and stay in L.
4. Find a formula for *T* in L, which shows *T* ∈ L.

The above steps work, as for any problem *B* in C, there is a ﬁrst-order reduction *I* to *T* , and both L

and C are complete via these reductions, so we also have *B* ∈ L = C.

### Space Hierarchy Theorem

The space hierarchy theorem states that for both nondeterministic and deterministic space, we have problems that can be solved using *t*(*n*) space, but not with any tighter space bound. Formally, we have

[DSPACE](#_bookmark73)[*o*(*t*(*n*))] ⊊ [DSPACE](#_bookmark73)[*t*(*n*)]

where *o*(*t*) is the set of functions *f* such that *f* ∈ O(*t*) but *t* ∉ O(*f* ), that is all functions that grow more slowly than *t*. This holds for all space constructible *t* ≥ log *n*. The same holds for [NSPACE](#_bookmark75).

We will present a proof for deterministic space.

*Proof.* The proof uses a diagonalization argument by presenting some machine *D* that takes a Turing machine *M* and an input size in unary as input and does the opposite of *M* if it halts. We want to show that for all *M* which run in space *f* (*n*) ∈ *o*(*t*(*n*)), we have an input on which *D* and *M* do not agree. This would show that the language computed by *D* is not in [DSPACE](#_bookmark73)[*o*(*t*)], and thus the strict containment.

On input ⟨*M,* 1*k* ⟩ our machine *D* marks of *t*(|⟨*M,* 1*k* ⟩|) tape cells, which are the cells that are allowed for the computation. Further, we also maintain a counter with size |*M* | · 2*t*(|⟨*M,*1*k* ⟩|), which is the maximum amount of diﬀerent conﬁgurations a Turing machine can pass before looping on a binary tape of size *t*(|⟨*M,* 1*k* ⟩|). Then, we simulate *M* on input ⟨*M,* 1*k* ⟩. If we transcend any bound, we reject. For all *M* in [DSPACE](#_bookmark73)[*o*(*t*)], there is a *k* such that *f* (*n*) ≤ *t*(*n*) by deﬁnition. On this input, the simulation ﬁnishes, and we can invert the output.

This directly gives us an input for which *M* and *D* diﬀer and thus proves our claim. Furthermore, *D* runs in [DSPACE](#_bookmark73)[*t*(*n*)] as by construction we ensured that we do not run inﬁnitely and that we stay within the space bound.

### Regular Languages

Here, we will show that the regular languages are captured exactly by second-order logic where we restrict ourselves to quantify only over predicates of arity one and do not include ≤. Further, we also are not allowed to use ≤, but have access to equality *x* = *y* and the successor relation *x* = *y* + 1. We call this class [MSO](#_bookmark63)[+1].

First, we need to present a formal deﬁnition of deterministic ﬁnite automata.

**Deﬁnition B.4** (DFA)**.** A deterministic ﬁnite automaton is a 5-tuple *M* = ⟨*Q,* Σ*, δ, q*0*, F* ⟩ where

*Q* is the set of states

Σ is the alphabet

*δ* is the transition function mapping a state and a symbol to the next state, so formally *δ* : *Q*×Σ →

*Q*.

*q*0 the start state

*F* a subset of *Q* which are the accepting states.

We say that a DFA *D* accepts a word *w* ∈ Σ∗ if when starting at the start state, if we go through *w* and always transition to the next state according to the actual symbol in *w* and the actual state, we end up in an accepting state.

In [[Rög23](#_bookmark54)] and [[Str94](#_bookmark57)] there is a proof of the following fact we will use in our proof for [MSO](#_bookmark63)[+1]:

**Theorem B.1.** *For any alphabet* Σ*, there is a DFA recognizing language L* ⊆ Σ∗ *if and only if it is regular.*

Now we can start to prove our main theorem for regular languages.

**Theorem B.2.** *For any alphabet* Σ*, a language L* ⊆ Σ∗ *is expressible in* [*MSO*](#_bookmark63)*[*+1*] if and only if it is regular.*

*Proof.* First, we show that any regular language can be expressed in [MSO](#_bookmark63)[+1]. Let *L* be regular, and *DL* be a DFA recognizing the language. We assume *L* does not contain the empty word, otherwise, we can recognize the language *L* \ {*ε*} and then add *φ* ∨ ∀*x*(*x* ̸= *x*), which adds the empty string back. Now let *DL* have *k* states. We can existentially quantify unary relations *X*1*, . . . , Xk* to have the meaning that *Xi*(*y*) is true if and only if *DL* is in state *i* after *y* steps. Then, we need to make consistency checks. We present formulas for each of the consistency checks and then can take the

“and” of those to get our ﬁnal formula ∃*X*1*, . . . , Xk*(*φ*1 ∧ *φ*2 ∧ *φ*3).

**The start state is** *qj* We have

*k*

/\ ↔

*φ*1 := (*i* = *j Xj*(0))

*i*=1

**We end in an accepting state** Let *Ti* be the set of all characters which lead from *qi* to an accepting state. Then we have

*k*

/\

*φ*2 :=

*i*=1

*Xi*

(max)

→ *a*V∈*Ti*

*Qa*(max)

**We move according to the transition function** We have

   /\*k* /\ ( )

∀ ∀  → 

*k*

*x y y* = *x* + 1

*i*=1 *a*∈Σ

(*Xi*(*x*) ∧ *Qa*(*x*)) → *Xδ*(*i,a*)(*y*)

/\ /\ V

*k*

(*Xi*(*y*) ∧ *Qa*(*x*)) →

∧

*i*=1 *a*∈Σ

!

(*Xr*(*x*) ∧ *δ*(*r, a*) = *j*)

By induction, we can show that always exactly one *i* satisﬁes *Xi*(*x*) for any *x*. Thus, if the created formula is satisﬁed, we know that *DL* accepts the word, and thus we have described *L* in [MSO](#_bookmark63)[+1].

*r*=1



For the other direction, we need to introduce two new concepts.

One is the nondeterministic ﬁnite automaton, which is analogous to the nondeterministic Turing machine as it can also have multiple transitions going from the same state. As with the nondetermin- istic Turing machine and the Turing machine, both the DFA and the NFA have the same expressive power.

The other concept is that of (V1*,* V2)-structures. These structures are generalizations of our former vocabulary *σ* as they have characters in *A* × P(V1) × P(V2). These structures are useful as we can make V1 to be the set of free ﬁrst-order variables in a formula *φ* and V2 be the set of free second-order variables in the formula. If at a position *i* in our (V1*,* V2)-structures we have *x* in the ﬁrst-order component of its character, we see this as meaning that *x* = *i*. For the second-order variables in the third component, an *X* at position *i* means that *X*(*i*) holds.

Now, we can prove by induction that all formulas in [MSO](#_bookmark63)[+1] with free variables in V1 and V2 are regular. Sentences, the formulas without free variables are the special case where V1*,* V2 = ∅.

First, we need to check that the (V1*,* V2)-structures are consistent, and no ﬁrst-order variable *x* appears more than once. This can be done by a NFA which has one state for each subset of variables, and extends its subset while going over the string. If a variable appears twice, we enter a state that always loops and rejects.

Then, we see that the atomic formulas can be checked, as *x* = *y*, *x* = *y* + 1 and *Qa*(*x*) are easy to check, and checking *X*(*x*) is equivalent to looking if the occurrence of *x* has *X* in the third component. We always need to take the intersection with the NFA which checks if the structure is valid.

All boolean connectives are also valid, as regular languages are closed under complement, inter- section and union as seen in [[Rög23](#_bookmark54)].

The most diﬃcult case is a formula of the form ∃*xφ* (as ∀*xφ* ≡ ¬∃*x*¬*φ*)). If ∃*xφ* is over (V1*,* V2)- structures, then *φ* is over (V1 ∪ {*x*}*,* V2)-structures. By induction, we know that *φ* deﬁnes a regular language and thus there is a NFA *N* which recognizes it. For the new automaton, we duplicate our states, with the meanings “used *x*” and “not used *x*”. If we are in a state where *x* was used, we can not take any transition with *x* in the second set. If we are in a state where *x* was not used, we can take a transition with *x* in the second set and go to the corresponding state with *x* used or take a transition where *x* is not used and go to the corresponding state where *x* was not used.

The remaining case with second-order variables is treated analogously, without the restriction on the number of times the variable is used, so we do not need to duplicate our states.

By induction, we have thus shown the other direction, and we see that [MSO](#_bookmark63)[+1] and the regular languages are equivalent.

### Context-Free Languages

For the context-free languages, we will show that they have an underlying structure that includes matchings.

A matching relation is a binary relation *M* on the universe {0*, . . . , n* − 1} which has the following properties:

**increasing** If *M* (*i, j*), then *i < j*

**uniqueness** Any *k* in the universe appears at most once in the relation

**non-crossing** If we were to draw arcs for the matching, none of them would intersect. Formally, if

*M* (*i, j*) and *M* (*k, l*), then either *j < k* or *l < i*

Visually, we can think of these relations as nested ranges over the universe.

With matchings, we can deﬁne [FO](#_bookmark61)(∃Match) as the ﬁrst-order logic extended with existential quantiﬁcation over matching relations.

**Theorem B.3.** *For any alphabet* Σ*, a language L* ⊆ Σ∗ *is expressible in* [*FO*](#_bookmark61)*(*∃*Match) if and only if it is context-free.*

For this, we ﬁrst introduce a normal form for context-free grammars. The proof that every context- free grammar can be converted to this form can be found in [[LST95](#_bookmark46)].

**Lemma B.4.** *Every context-free language has a grammar which satisﬁes*

* *All rules are of one of the two forms*
  + *S* → *α with α* ∈ Σ
  + *X* → *αuβ with α, β* ∈ Σ *and u* ∈ (Σ ∪ *V* )∗
* *For all production rules with a right-hand side that has at least one non-terminal we deﬁne its* pattern*. The pattern of rule X* → *v*0*X*1*v*2*X*2 *. . . Xsvs is deﬁned to be v*0|*v*1| *. . .* |*vs, where* | *is a new symbol not in* Σ *We require that for any two rules with the same pattern, they have the same left-hand side and thus the source non-terminal can be uniquely identiﬁed by the pattern.*

For a speciﬁc arch ⟨*i, j*⟩ in a matching *M* on some word structure, we can also determine a pattern. For this, we go through all indices from *i* to *j* and add the character at the actual position. If we are at a starting point of some arch ⟨*k, l*⟩, instead of adding the actual character, we add | and continue at *l* + 1.

Now, we can start with the proof of our theorem.

*Proof.* We want to ﬁnd a formula such that for all arches ⟨*i, j*⟩ in *M* , the substring from *i* to *j* can be derived from the non-terminal for which we have a rule with the pattern of ⟨*i, j*⟩.

For this, we say that arch ⟨*i, j*⟩ *corresponds* to a production rule *p* if they have the same pattern. For any string *u*, we have a formula *ϕu*(*i, j*) which means that the substring from *i* to *j* is *u*. This formula is easy to write as we only need to check each character one by one using the successor relation. We can express correspondence to *p* ≡ *X* → *v*0*X*1*v*2*X*2 *. . . Xsvs* (including rules with terminal right-hand side) by the following formula:

X*p*(*x, y*) ≡ ∃*x*1*, y*1*, . . . , xs, ys* ((*x < x*1 ∧ *x*1 *< y*1 ∧ *y*1 *< x*2 ∧ · · · ∧ *ys < y*) ∧ (*ϕv*0 (*x, x*1 − 1) ∧ *ϕv*1 (*y*1 + 1*, x*2 − 1) ∧ · · · ∧ *ϕvs* (*ys* + 1*, y*)) ∧ (*M* (*x*1*, y*1) ∧ · · · ∧ *M* (*xs, ys*) ∧ *M* (*x, y*)) ∧

∀*k, l* (*M* (*k, l*) → ((*x* ≤ *k* ∧ *l* ≤ *y*) ∨ (*x*1 ≤ *k* ∧ *l* ≤ *y*1) ∨ · · · ∨ (*xs* ≤ *k* ∧ *l* ≤ *ys*))))

The ﬁrst line means that the *vi* are in the right order, the second that they correspond, the third that there are arches between the *vi* and the last that there are no other arches.

Now, we want to be more general, and express that the pattern of ⟨*i, j*⟩ corresponds to some production rule with left-hand side *X*. Let X*X*(*x, y*) be the disjunction of all X*p* with left-hand side

-

*X*. Because our normal form says that the pattern uniquely determines the left-hand side, for each arch which has arches underneath, we have a unique corresponding non-terminal.

Now, we want to have a formula which expresses not only correspondence for a production rule *p*, but also that the non-terminals are correct. For this, we supplement our formula X*p* with a new line, expressing that the non-terminals correspond.

X*p*(*x, y*) ≡ ∃*x*1*, y*1*, . . . , xs, ys* ((*x < x*1 ∧ *x*1 *< y*1 ∧ *y*1 *< x*2 ∧ · · · ∧ *ys < y*) ∧ (*ϕv*0 (*x, x*1 − 1) ∧ *ϕv*1 (*y*1 + 1*, x*2 − 1) ∧ · · · ∧ *ϕvs* (*ys* + 1*, y*)) ∧ (*M* (*x*1*, y*1) ∧ · · · ∧ *M* (*xs, ys*) ∧ *M* (*x, y*)) ∧

-

∀*k, l* (*M* (*k, l*) → ((*x* ≤ *k* ∧ *l* ≤ *y*) ∨ (*x*1 ≤ *k* ∧ *l* ≤ *y*1) ∨ · · · ∨ (*xs* ≤ *k* ∧ *l* ≤ *ys*)))

(X*X*1 (*x*1*, y*1) ∧ · · · ∧ X*Xs* (*xs, ys*)))

- -

Finally, with *P* being our set of rules, we have a formula that tells us a word can be derived from the start symbol:

*M*  *S* (0*,* max) *x, y* *M* (*x, y*) *p*(*x, y*)

∃ X- ∧ ∀  → V X- 

*p*∈*P*

Now, by construction, if and only if a word satisﬁes the formula, there is a derivation from *S* for the word, as each arch can be seen as a derivation step, which we check is valid with our formula. We used our assumption from the normal form to show that the X*p* are only satisﬁed when the arches which are non-terminal do not belong to any other non-terminal symbol than the one in the rule. If two terminal rules coincide, we do not care as no further derivation is possible.

-

Now, we come to the other direction. This direction requires some new notation and lemmas. We

will use the notion of tree languages.

**Deﬁnition B.5** (Tree Language)**.** In a tree language, we have a rooted tree, which has an order on its vertices in leftmost depth-ﬁrst way[1](#_bookmark103). Each node has a label, and each label has an arity which corresponds to the out-degree of the node. A tree language is the set of all trees over a ﬁnite label set. We deﬁne the *leaf alphabet* of a tree language to be the set of 0-ary labels. For any tree *T* in a tree language, we deﬁne the *yield* of *T* to be the leaf labels concatenated according to the order relation

from left to right.

The vocabulary of a tree language T is *τ* = ⟨{0*, . . . , n* − 1}*, Qa, Qb, . . . , Qz,* ≤*,* 0*,* 1*,* max*, C*2⟩ In addition to the relations for each label, there is a child relation *C*(*i, j*) which means “node *i* is a child of node *j*”.

Now, we present two lemmas for recognizable tree languages and their relation to context-free languages.

**Lemma B.5** ([[MW67](#_bookmark49)])**.** *A language L* ⊆ Σ∗ *is context-free if and only if there is a recognizable tree language T with leaf alphabet* Σ *for which a word is in L if and only if it is the yield of some tree in T.*

**Lemma B.6** ([[TW68](#_bookmark59)])**.** *A tree language T is recognizable if and only if there is a monadic second-order sentence that recognizes it.*

We present some relations that can be written in [MSO](#_bookmark63) on trees.

**Lf(***i***)** Node *i* is a leaf Lf(*i*) ≡ ∀*x* ̸= *C*(*x, i*)

**Lc(***i, j***)** Node *i* is the leftmost child of *j* Lc(*i, j*) ≡ *C*(*i, j*) ∧ ∀*x*(*C*(*x, j*) → *i < x*) **Rc(***i, j***)** Node *i* is the rightmost child of *j* Rc(*i, j*) ≡ *C*(*i, j*) ∧ ∀*x*(*C*(*x, j*) → *x < i*) **An(***i, j***)** Node *i* is an ancestor of *j*

An(*i, j*) ≡ ∃*U* (*U* (*i*) ∧ *U* (*j*) ∧ ∀*x*(*U* (*x*) →

((*x* ̸= *i* ↔ ∃*y*(*C*(*x, y*) ∧ *U* (*y*))) ∧ (*x* ̸= *j* ↔ ∃*y*(*C*(*y, x*) ∧ *U* (*y*))))))

**Pt(***U, i, j***)** Node *i* is an ancestor of *j*, *j* is a leaf and *U* contains all nodes in the path from *i* to *j*.

Pt(*U, i, j*) ≡ *U* (*i*) ∧ *U* (*j*) ∧ Lf(*j*) ∧ ∀*x*(*U* (*x*) →

((*x* ̸= *i* ↔ ∃*y*(*C*(*x, y*) ∧ *U* (*y*))) ∧ (*x* ̸= *j* ↔ ∃*y*(*C*(*y, x*) ∧ *U* (*y*)))))

Using these two lemmas, we can continue by presenting for any formula *φ* in [FO](#_bookmark61)(∃Match) a formula

*ϕ* in [MSO](#_bookmark63) over trees such that

*w* |= *φ* ⇔ there exists a tree *T* with yield *w* such that *T* |= *ϕ*

1The order which we ﬁnd by doing a depth-ﬁrst search on the tree, entering the leftmost node ﬁrst

For this, we present a class of trees which correspond to a word with a matching. For any word *w* with matching *M* , we can construct a tree over Σ ∪ {⊕2*,* ⊙2}. We do this using an intermediate step. First, we construct a tree with the wrong arity for nodes of type ⊕ by assigning one node of type ⊕ to each arch, with edges to every direct arch underneath it and every character directly underneath it. If we have multiple trees, we add a new ⊙ node on top with an edge to all roots of these trees. To ﬁx the arity issue, we repeat the following procedure until there are no nodes with more than outdegree 2.

Take some node with an outdegree greater than 2. Take the two leftmost children and add a ⊙ node with an edge to both, and an edge from the new node to the former parent. This procedure will eventually terminate as we always decrease the outdegree of some node by one and add a valid node. We can see that this procedure can also be done backwards if and only if for any node of type ⊕,

the leaf on the leftmost and rightmost path of a node are distinct and no ⊕ node occurs on the path between the two. We can express this property of a tree by the following formula.

Υ ≡ ∀*x*(*Q*⊕(*x*) → (∃*y, z, Uy, Uz*(*y* ̸= *z* ∧ Pt(*Uz, x, z*) ∧ Pt(*Uy, x, y*)∧

∀*r*(*Uy*(*r*) → ((*r* = *y* ∨ ∃*w*(*Uy*(*w*) ∧ Rc(*w, r*))) ∧ (*r* = *y* ∨ *r* = *x* ∨ *Q*⊙(*r*))))∧

∀*r*(*Uz*(*r*) → ((*r* = *z* ∨ ∃*w*(*Uz*(*w*) ∧ Lc(*w, r*))) ∧ (*r* = *y* ∨ *r* = *x* ∨ *Q*⊙(*r*)))))))

Here, in the ﬁrst line we assert that for every node with label ⊕, we have two distinct leaves *y, z* with paths *Uy* and *Uz*. The second and third lines assert the same for *x* and *y*, that they are the rightmost / leftmost leaf and that no other ⊕ type node lies on the path from them to *x*.

Now, we want to convert our formula *φ* over strings with a matching to a formula *γ* over trees. Then we can assert that the tree represents a string with matching and the yield satisﬁes *φ* using Υ ∧ *γ*.

For this, we need to restrict any quantiﬁers in *φ* to the leaves by replacing ∃*xϕ* with ∃*x*Lf(*x*) ∧ *ϕ*

and ∀*xϕ* with ∀*x*Lf(*x*) → *ϕ*. Further, we replace *M* (*z, y*) by

*m*(*z, y*) ≡ ∃*x*(*Q*⊕(*x*) ∧ (∃*Uy, Uz*(*y* ̸= *z* ∧ Pt(*Uz, x, z*) ∧ Pt(*Uy, x, y*)∧

∀*r*(*Uy*(*r*) → ((*r* = *y* ∨ ∃*w*(*Uy*(*w*) ∧ Rc(*w, r*))) ∧ (*r* = *y* ∨ *r* = *x* ∨ *Q*⊙(*r*))))∧

∀*r*(*Uz*(*r*) → ((*r* = *z* ∨ ∃*w*(*Uz*(*w*) ∧ Lc(*w, r*))) ∧ (*r* = *y* ∨ *r* = *x* ∨ *Q*⊙(*r*)))))))

which is very similar to Υ.

This is already everything we need to do, and thus we can conclude that

*w* |= *φ* ⇔ there exists a tree *T* with yield *w* such that *T* |= Υ ∧ *γ*

Thus, we have proved both directions and see that the context-free languages are exactly captured by [FO](#_bookmark61)(∃Match)

### Context-Sensitive Languages

Here, we show the equivalence between context-sensitive languages and nondeterministic Turing ma- chines using O(*n*) space. The two directions of the proof were presented separately in [[Kur64](#_bookmark43)] and [[Lan63](#_bookmark44)].

**Theorem B.7.** *The class of context-sensitive languages is exactly the class of languages accepted by a linear bounded nondeterministic Turing machine.*

*Proof.* For the direction from grammar to Turing machine, we only need to show that our Turing machine can simulate a derivation backwards. We know that every context-sensitive language has a noncontracting grammar *G*. Using this fact, we can construct a nondeterministic Turing machine *N* which scans the current tape and whenever it recognizes a pattern of the right-hand side of a production rule in *G*, it decides whether it replaces it or not by the left-hand side of the rule. If some computation branch of *N* ends up with only the start symbol of *G* on the tape, we accept. Essentially, *N* simulates a derivation of *w* from the start symbol backwards. Because we try all possibilities by nondeterminism, we know that if *N* does not accept *w*, there is no derivation ending in *w* from the start symbol of *G*. As we assumed *G* is noncontracting, replacing the right-hand side of a rule with the left-hand side never makes the word longer, and thus we only need O(|*w*|) space.

The other direction works by explicitly deﬁning a grammar which simulates any linear bounded automaton backwards. Without loss of generality, we include an end marker #.

Let *N* = ⟨*Q,* Σ*,* Γ*, δ, q*0*, qaccept, qreject*⟩ be a [NSPACE](#_bookmark75)[*n*] Turing machine. Then, we construct a grammar *G* = ⟨*V,* Σ*, P, S*⟩ with *V* = *qi*∈*Q aj* ∈Γ{*bqi,aj* } ∪ {*S, L, R,* #} ∪ *aw* ∈Γ\Σ{*aw*}. The *bq,a* represent a position on the tape including the actual state and the actual character, in addition, we also have the start state, the end marker and some utility non-terminals.

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We now add the following rules to *P* :

* For each *ai* ∈ Γ, we add a rule *S* → *Lbqaccept,ai R* to *P* . These rules mean that we are in a ﬁnal accept state. To extend the ﬁnal tape, we add again for each *ai* ∈ Γ the rules *L* → *Lai*, *L* → #, *R* → *aiR* and *R* → # to our rule set. These rules allow derivations from *S* to an end tape of

the form #*ai*1 *. . . bqaccept,ai . . . ai* #.

*j*

*k*

* For each rule ⟨*ak, ql, L*⟩ ∈ *δ*(*qi, aj*), we add rule *baw,ql ak* → *awbaj,qi* for every *aw* ∈ Γ. Similarly, for each rule ⟨*ak, ql, R*⟩ ∈ *δ*(*qi, aj*), we add rule *akbaw,ql* → *baj,qi aw* for every *aw* ∈ Γ. We can see that these rules simulate the nondeterministic Turing machine backwards.
* For the start, we include #*bai,q*0 → #*ai* for all *ai* ∈ Σ This rule allows us to say that we are at the start of our computation and “remove” the read-write head to get our initial input word.

As for any word in *G* we can follow back the derivation to an accepting state, we have that both *N*

and *G* deﬁne the same language.

Thus, we are done and have proven the equivalence of context-sensitive languages and linear bounded automata.

### Recursive Languages

The most general case is interesting as it gives us a logical formalism for all problems which are computable at all.

The proof relies on Diophantine sets. Those sets are the sets that correspond to the tuples which have a solution for some Diophantine equation. A Diophantine equation is a polynomial equation *P* with a tuple *x* of parameters and a tuple *y* of variables. A tuple *x* has a solution if there exists a tuple *y* such that *P* (*x, y*) = 0. The famous MRDP theorem states that the Diophantine sets are exactly the computable sets. At the same time, this shows that Hilbert’s 10*th* problem is unsolvable. A full proof of these facts can be found in [[Mat96](#_bookmark48)].

With this new characterization, there is a quite straightforward characterization of computable sets. The logic [FO](#_bookmark61)(∃N) consists of all formulas *ϕ*(*x*) ≡ ∃*y*(*φ*(*x, y*)) with *φ* having only bounded quantiﬁers (of the form ∃*x < y* or ∀*x < y*), addition, multiplication, equality, and any constant natural number [[Ent20](#_bookmark36)]. The existential quantiﬁers at the beginning are allowed to range over all the natural numbers.

These formulas can deﬁne exactly the Diophantine sets, as the formulas we present are exactly those that mean “there is a tuple *y* of natural numbers such that some polynomial is satisﬁed”. Thus, they also deﬁne exactly all computable sets.

### Open questions

The domain of Descriptive Complexity is full of open questions as the proofs of lower bounds seem to be very diﬃcult in most cases. Further, even separation between complexity classes which seem to take an exponential amount of resources compared to another one in practice can not be shown to be diﬀerent.

**[P](#_bookmark78)**=? [**NP**](#_bookmark79)

The [P](#_bookmark78) vs. [NP](#_bookmark79) question is the most emblematic question in Descriptive Complexity theory. In practice, for any [NP](#_bookmark79)-complete problem, only exponential worst-case algorithms are known. This leads to the widely believed conjuncture that [P](#_bookmark78) ̸= [NP](#_bookmark79). The problem is one of the seven Millennium Problems and a solution of equality or inequality is worth 1 Million US dollars.

The consequences of a solution stating that [P](#_bookmark78) = [NP](#_bookmark79) could have many practical advantages if it was constructive and had a low constant, as many important problems in research and logistics could be solved quickly. It would also mean the breakdown of most of modern cryptography, which relies on problems being intractable. On a conceptual level, it would mean that ﬁnding a proof to a problem is not harder than verifying its correctness, which would greatly impact the work of mathematicians. If a proof of the contrary were known, this would focus the research more on the average case complexity of [NP](#_bookmark79) problems, but because of the continued lack of success on the question, this shift has already widely taken place.

[**NSPACE**](#_bookmark75)[*n*]=? [**DSPACE**](#_bookmark73)[*n*]

This problem is known under the name ﬁrst linear bounded automaton problem since its proposal by Kuroda in [[Kur64](#_bookmark43)], and asks if nondeterminism adds power in the context of bounded space. This comes from the fact that a [NSPACE](#_bookmark75)[*n*] Turing machine can be seen as a Turing machine with a linear bound on its space usage. This theorem is of interest as we know that [NSPACE](#_bookmark75)[*n*] is equivalent to the context-sensitive languages by appendix [B.2.6](#_bookmark104).

Since the proposal, there have been two advances. One is the proof that [NSPACE](#_bookmark75) is closed under complement. The contrary would have implied [NSPACE](#_bookmark75)[*n*] ̸= [DSPACE](#_bookmark73)[*n*] as [DSPACE](#_bookmark73) is closed under complement. The second advance is Savitch’s theorem in section [3.3.1](#_bookmark13) which already gives a bound for simulating [NSPACE](#_bookmark75) using [DSPACE](#_bookmark73) machines. It is not known if this theorem is optimal, that is whether the blow-up by a power of 2 is optimal or if we can do better.

An equality would imply that the context-sensitive languages can be recognized by a deterministic linear bounded automaton, which could make recognizing words in context-sensitive languages easier and faster.

# Communication

## Email

Sehr geehrter,

Ich heisse und bin Gymnasiast am Gymnasium Bäumlihof in Basel. Auf Empfehlung meiner Maturarbeitsbetreuungslehrperson schreibe ich Ihnen.

In meiner Maturarbeit beschäftige ich mich mit Finite Model Theory, insbesondere im Bezug auf deren Verbindung mit Formal Languages. Somit will ich insbesondere untersuchen, ob ich eine korrespondierende Logik zu den Context-Sensitive Languages ﬁnde, da ich keinen Artikel gefunden habe, der eine solche erwähnt.

Die Grundlagen der Theoretischen Informatik konnte ich mir im Rahmen des Schülerstudiums an der Uni Basel aneignen. Auch interessiere ich mich stark für Logik, Complexity Theory und Finite Model Theory und habe darüber Bücher gelesen.

Damit Sie meine Fähigkeiten einschätzen können: Auch bei der Schweizer Informatikolympiade, der Schweizer Mathematikolympiade sowie bei mehreren internationalen Olympiaden (MEMO, CEOI, BOI, WEOI) habe ich teilgenommen.

Jetzt zur eigentlichen Frage: Haben sie Tipps, wie ich vorgehen sollte und in welche Richtung ich forschen sollte? Und halten sie mein Unterfangen überhaupt für möglich?

Vielen Dank im Voraus für Ihre Zeit und Hilfe. Viele Grüsse,

## Answer

Lieber

Bitte nehmen Sie es mir nicht übel, dass ich Ihnen noch nicht geantwortet habe. Ich habe Ihre Mail einem sehr guten Kollegen weitergeleitet, der sich sicherlich bei Ihnen melden wird. Momentan fehlen mir leider die zeitlichen Ressourcen, Ihnen persönlich detaillierter zu schreiben, was ich sehr bedauere.

Liebe Grüsse

## Email

Sehr geehrte ,

Ich heisse (Nickname, falls Sie sich daran erinnern) und bin Gymnasiast am Gymnasium Bäumlihof in Basel. Da ich durch Ihre Vorlesung “Theory of computer science” auf mein Maturarbeitsthema gestossen bin schreibe ich Ihnen jetzt.

In meiner Maturarbeit beschäftige ich mich mit Finite Model Theory, insbesondere in Bezug auf deren Verbindung mit Formal Languages. Somit will ich insbesondere untersuchen, ob ich eine kor- respondierende Logik zu den Context-Sensitive Languages ﬁnde, da ich keinen Artikel gefunden habe, der eine solche erwähnt.

Neben Ihren Vorlesungen habe ich auch durch diverse Bücher über Logik, Complexity Theory und Finite Model Theory Einblicke in das Thema gehabt.

Damit Sie meine Fähigkeiten einschätzen können: Auch bei der Schweizer Informatikolympiade, der Schweizer Mathematikolympiade sowie bei mehreren internationalen Olympiaden (MEMO, CEOI, BOI, WEOI) habe ich teilgenommen.

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In meiner Maturarbeit beschäftige ich mich mit Finite Model Theory, insbesondere im Bezug auf deren Verbindung mit Formal Languages. Somit will ich insbesondere untersuchen, ob ich eine korrespondierende Logik zu den Context-Sensitive Languages ﬁnde, da ich keinen Artikel gefunden habe, der eine solche erwähnt.

Die Grundlagen der Theoretischen Informatik konnte ich mir im Rahmen des Schülerstudiums an der Uni Basel aneignen. Auch interessiere ich mich stark für Logik, Complexity Theory und Finite Model Theory und habe darüber Bücher gelesen.

Damit Sie meine Fähigkeiten einschätzen können: Auch bei der Schweizer Informatikolympiade, der Schweizer Mathematikolympiade sowie bei mehreren internationalen Olympiaden (MEMO, CEOI, BOI, WEOI) habe ich teilgenommen.

Jetzt zur eigentlichen Frage: Haben sie Tipps, wie ich vorgehen sollte und in welche Richtung ich forschen sollte? Und halten sie mein Unterfangen überhaupt für möglich?

Ein volles Mentoring suche ich aktuell nicht, eine Antwort in der Form einer Mail, eines Telefonates oder eines kurzen Treﬀens würde mich aber freuen.

Vielen Dank im Voraus für Ihre Zeit und Hilfe. Viele Grüsse,

## Answer

Dear,

asked me to write you some tips about your planned thesis; I am working in his group as senior scientiﬁc assistant. I hope it is ok if I write in English.

As you are likely aware, context-sensitive languages are exactly the languages that can be recog- nized by nondeterministic Turing machines in linear space. This class is often denoted as NSPACE(n)

= nondeterministic space O(n).

I am not really an expert in the ﬁeld of descriptive complexity theory, but I am not aware of logic-based characterization of this class. There are, however, some related results that are known. For example, characterization of DSPACE(nˆk) is known (see e.g. [https://people.cs.umass.edu/](https://people.cs.umass.edu/~immerman/descriptiveComplexity.html)

[~immerman/descriptiveComplexity.html](https://people.cs.umass.edu/~immerman/descriptiveComplexity.html)), where DSPACE(f(n)) denotes the class of languages ac- cepted by deterministic Turing machines using space O(f(n)). NSPACE(n) is trivially a superset of DSPACE(n), and it is also a subset of DSPACE(nˆ2) (due to Savitch’s theorem). This immediately gives two logic-based characterizations that bound NSPACE(n) from below and from above. This observation is already mentioned in [https://people.cs.umass.edu/~immerman/pub/ams\_notices.](https://people.cs.umass.edu/~immerman/pub/ams_notices.pdf) [pdf](https://people.cs.umass.edu/~immerman/pub/ams_notices.pdf) as Theorem 2.

To get an exact logic-based characterization of NSPACE(n) may be not easy, most likely, it is a very ambitious goal. You could try to adapt the results for DSPACE(nˆk) to NSPACE; even if you do not succeed there, explaining the proof ideas and showing why they are not so easy to transfer to NSPACE would make a valuable thesis. Augmenting such results with explanations of the upper and lower bounds of NSPACE(n) would, in my opinion, form a very solid “Maturaarbeit”.

If you have more questions, feel free to ask. I wish you good luck in writing your thesis and I hope you will also enjoy it.

Best regards,

## Further Communication with

Dear,

First of all, thanks a lot for your response and your time!

English is ok for me, in fact I’m writing my Matura project in English.

I eﬀectively wasn’t aware of the characterisation of DSPACE(nˆk), so this is very valuable advice into which I will look. I had found upper / lower bounds using other formal languages, notably [https:](https://link.springer.com/chapter/10.1007/BFb0022257)

[//link.springer.com/chapter/10.1007/BFb0022257](https://link.springer.com/chapter/10.1007/BFb0022257) (existential SO restricted to binary matching predicates) for context-free languages and [FO](#_bookmark61)(∃N) for enumerable languages (As described in the book “Descriptive Complexity” by Neil Immermann), but they had the disadvantage of being less tight.

If (or rather when) I have more questions, I will write you another e-mail, thanks again for giving me this possibility!

Best,

## Question about writing my MA in English

Sehr geehrter, sehr geehrter,

Gerne würde ich meine Maturarbeit über “Theoretische Informatik und Logig” auf Englisch schreiben, da die mathematische Literatur zum Grossteil auf Englisch ist und ich Mathematik lieber in dieser Sprache löse. Ist dies möglich?

Vielen Dank und viele Grüsse,

Sehr geehrter, sehr geehrter,

Da ich gerne bald mit der schriftlichen Arbeit anfangen würde, wäre ich froh um eine Rückmeldung. Viele Grüsse,

Lieber

Die Schulleitung hat Ihren Antrag genehmigt. Sie dürfen die Maturaarbeit auf Englisch schreiben. Auﬂage der Schulleitung: Koreferent\*in muss über entsprechende Englischkennt-

nisse verfügen, um die Arbeit berwerten zu können.

Liebe Grüsse

# Notes

## Grammars

Equivalent: surrounded by context, length left ≤ length right, [Kuroda normal Form](https://en.wikipedia.org/wiki/Kuroda_normal_form)

## Summary

*FO*[*t*(*n*)]: *A* ∈ *S* ⇔ *A* |= ([*QB*]*t*(|*A*|)*M*0) (*c/x*)

*V AR*[*v* + 1] = LJ∞ *FO* − *V AR*[2*cnv , v* + 1]: Unbounded iterations, using at most *v* + 1 vars

*c*=1

*ITER*[arity *k*]: Simultaneous iterations of relations of arity k as in LFP (no monoticity requirement)

*V AR*[*k* + 1*, r*]*/IT ER*[arity *k, r*]: Same plus restricted variables of total bits *r* + O(1)

*DSP ACE*[*s*(*n*)] = *V AR* 1*k* + 1*,* log ( *s*(*n*) )l = *ITER*[arity *k,* log ( *s*(*n*) )] with *k* = ⌊log (*s*(*n*))⌋ Prob-

*nk*

*nk*

*n*

lem for generalisation: For *Ct*(*x, b*) (tape cell *x* at time *t* is *b*), for nondeterminism we could have the same, but other tape cells are not independent.

*FO* − *V AR*[*s*(*n*)*, v*(*n*)]: *V AR*[*s*(*n*)] + constant number of *v*(*n*) log(*n*) bit variables, which can only be tested by *BIT* .

*SO*[*t*(*n*)*,* arity *k*]: *FO*[*t*(*n*)], but with quantiﬁcation over relation variables of arity at most k

*NSP ACE*[*nk*] = *SO*(arity *k*)(*T C*) ⊆ *SO*(arity *k*)[*nk*]

All formulas of *SO*(arity *k*)(*T C*) can be written in the form

(*TCstruct*1*,struct*2 *α*)(*false, true*)

where *α* is quantiﬁer-free and both structs have at most *O*(*nk*) bits

*nk* !

*SO*(arity *k*)[*t*(*n*)] = *FO* − *V AR*

*t*(*n*)*,* log(*n*)

DSPACE[*s*(*n*)] ⊆ NSPACE[*s*(*n*)] ⊆ FO-VAR 1*s*(*n*)*, s*(*n*) l ⊆ *ATIME*[*s*(*n*)2] ⊆ DSPACE[*s*(*n*)2]

log(*n*)

Proof works by looking at alternating time paths of length 2*r* in the *NSP ACE* computation graph.

* 1. **Extending** *Ct*(*x, b*)

Adding an additional variable for non-deterministic choices *Ct*(*c, x, b*), where *c* are the non- deterministic choices, makes the states size grow exponentially in *O*(2*cnk* ) (one for each move of

D. NOTES D.4. Direct transformation of SO(arity *k*)(*T C*) to FO-VAR 11*, nk* l (*TC*)

log *n*

the NTM). So we have way more memory then *ITER*[*k* + 2], which makes it strictly more powerful. We need to have all the choices as we can only inductively deﬁne *Ct* using the previous related choices.

* 1. **Direct transformation of SO**(**arity** *k*)(*T C*) **to**

**FO-VAR** 11*, nk* l (*TC*)

log *n*

*SO*(arity *k*)(*T C*) ⊆ *FO* − *V AR* 11*, nk* l (*TC*): We know that any formula in *SO*(arity *k*)(*T C*) can

log *n*

be written in the form

(*TCstruct*1*,struct*2 *α*)(*false, true*)

where *α* is quantiﬁer-free. Using the equivalence *BIT* (*Y, x*) ⇔ *Y* (*x*), we can write anything in *α*

using *FO* − *V AR* 11*, nk* l. Any relation in *struct*1*, struct*2 can also be rewritten using the extended

log *n*

variables. As we now any Relation with *false* is converted to 0 and any relation with *true* is converted to 1, we have (*false, true*) that can now be represented by (0*, max*).

*SO*(arity *k*)(*T C*) ⊇ *FO* −*V AR* 11*, nk* l (*TC*): We know that any formula in *FO*(*TC*) can be

log *n*

written as (*TCstruct*1*,struct*2 *α*)(0*, max*). We can extend this result to include the extended variables.

For the usage of the extended variables in *BIT* , we have an atomic formula, so this is trivial. For quantiﬁcation, we can handle it the same way as normal quantiﬁcation with the modiﬁcation of the “counting” variable with a tuple of *k* counting variables, and the successor relation on these *k* variables. We haven’t shown yet that negation can be handled by this. In the proof that *FO*(*pos TC*) = *FO*(*TC*),

we only ask for a successor relation on the vertices (consisting of the extended variables and normal

variables), which we assume from the beginning, so this also holds for *FO* − *V AR* 11*, nk* l (*pos TC*) =

log *n*

*FO* − *V AR* 11*, nk* l (*TC*). Now, we can do the same as above for the other direction to convert a

log *n*

formula of the form (*TCstruct*1*,struct*2 *α*)(0*, max*) to (*TCstruct*1*,struct*2 *α*)(*false, true*).

## Can we prove that there is no way to use *TC* with a *FO* language that has less then *O*(*nk*) variables?

We have Arity theorems [Arity hierarchies Martin Grohe](https://pdf.sciencedirectassets.com/271596/1-s2.0-S0168007200X00292/1-s2.0-0168007295000720/main.pdf?X-Amz-Security-Token=IQoJb3JpZ2luX2VjEA0aCXVzLWVhc3QtMSJHMEUCIQCZEfzPpPqb83EtF%2BLgRMsETvrj14M6SqdN3VHcNPEkdAIgWF%2BnohPDt8o%2FXZc7b7P3HqFm72ZMaOoAg4MLrtTMUtYquwUIpv%2F%2F%2F%2F%2F%2F%2F%2F%2F%2FARAFGgwwNTkwMDM1NDY4NjUiDOiLVE0art2rY62EliqPBQbQSd%2FAcBktlD2Axqi2hf20al4TnosyDnMOkEbqwIQAyET8IMr%2B776j5WCIdUO85KSby8IvsXcoMaeJllnCTWiAtcQOCv0tjt1Rome20%2B50%2BMDR2DM6BHnHFA3gOFm3PAoX427r%2BysNLJl2dbF4YAXE2Rb02hYUm47FnV%2FiV6LDZgVk56JvT2yASPDa1TYaBTCUt09gepbCkOkC1Tzac2OEzReNX5RGnSGRJRwZN7QDFmHMEhiE1YO1sYOGJgnrWvjBNlgRGe66ZDXA6H9FlnP3KfXpA%2BRlBp9SfqipL6iEj8iTLlY8Iw4Lx8SLSR4WwHt5ZSopsZj1SZlTKtBkaHqb4%2ByS8iYsy7JUiFML1LvIE4ttxy7fM6qY5XNU0NSNb6NF%2FStuygMbc5kU3Xs5D%2BYtei30HQKStUmol3nJYpBbgTpaUCPZTWF1xfFeF2r0rHtT8SGvWCZNI0N78o6YYf%2FPiFrC%2FLb%2BaGc0DtkdqrLP3L3Rrf2cIXsxqZ4LwlEdj6fEIKZxw71c95o1h1PlnB72qAg62JKOdNtHM3Gpw7yeuU1kAdNEctOU2WMvhqqjDHlcerDbspa39NRMVIgrKI0wHhphT42iSbZC2rVP7cgpqBFADJqrT1C3ROwrQngxGDjCqpaDgR3Lh%2BFCvJ6Hj6aficQkrXnZKYPd2ZxF0LoEw5F0K1eEAoi%2F31m3Wd8MhY%2FzlDt5Owh4zIlIsv9AEG0byN%2F%2FEs30PT469%2BiRlOlCAJ0iYqpvsVk3whvBPEHae8O3LFUG7gWFK50eQ%2FFQ9KVcy0sLMCNpK286rnXOYvU18pdah3GbEt%2BG642xb%2F3y%2BPNG6oMZ%2B485zsKgusLfgILd7%2FU7w4E8yeE8VrHFPFcw3%2B2wswY6sQFACa99Hzo8imVILu9U2u8E7F2wa2uOniA6UidGdcybvYvso0UTqK53Ki%2FnJ5tT6ADTd52VSNtgviQbNSxdnAdsifHTNGD4xTutf9zwfV1AVAWosXcbKwC%2BGd1Cmoh4JMrd4uMlW2FtCpKGbmKW9pcJkmjNRHacw3IN%2FC%2F%2BI4yRBQZyBwoKevOICAIRcr8XsVMxvpSFg4M6Amz3UbgG1Tf5%2FaSW%2BAn2%2BHYGkfJuV5Io8OU%3D&X-Amz-Algorithm=AWS4-HMAC-SHA256&X-Amz-Date=20240614T123128Z&X-Amz-SignedHeaders=host&X-Amz-Expires=300&X-Amz-Credential=ASIAQ3PHCVTYVAGA3OF6%2F20240614%2Fus-east-1%2Fs3%2Faws4_request&X-Amz-Signature=42e03b07238e340e4071f6b533897676385912b01735ff90b4ce8ee3841179f2&hash=a3824b39f0ff05985504a3e08bf09f2572a737ad6204bcb9d1f18f599e123db5&host=68042c943591013ac2b2430a89b270f6af2c76d8dfd086a07176afe7c76c2c61&pii=0168007295000720&tid=spdf-86d73916-b662-4ed3-b6de-753dfc8128db&sid=9e433d8d71f1564cc768be65fe1f9e328634gxrqb&type=client&tsoh=d3d3LnNjaWVuY). Is this applicable because ordered structure? Diagonalisation in logic?

FO[2*n*], SO[n] can both express REACH*dl*

*Claim:* If extended variables are allowed everything, nothing changes.

*Deﬁne* Superrelations, using the extended variables. Does not work because input to big Not needed as we are working on string structures

If we allow extended variables for relations we have undeﬁned for some, and if we check, we could “copy” bits into FO variable Numerical relations ≤*, SUC*(*X, Y* )*,* +*,* × are deﬁnable via BIT anyways. If no further relations are added, we have already that they reﬂect exactly some NSPACE com-

plexity class, so we can not do better.

**Approach 2** We have that the FO[1, s(n)/log(n)](TC) correspond exactly to our complexity classes. The space hierarchy tells us that we can not do better while using these constructions.

Now assume we add iterated quantiﬁer blocks. We have multiple ways of doing this. We can have the TC inside or outside

*TCs*(*n*)([*QB*]*t*(*n*)*φ*) ≡ *TCO*(*n*)(*ϕ*)

In both cases, we can replace the TC operator with a quantiﬁer block iterated *o*(*n*) times. So we have formulas of the form

[*QB*1]*s*(*n*)[*QB*2]*t*(*n*)*ψ*

We can include a counting variable of which counts to *o*(*n*) and then switches block to get only one combined block.

so we then have a formula in FO[s(n) + t(n), s(n)/log(n)], which is also in FO[t(n), s(n)/log(n)](TC) or FO[s(n), s(n)/log(n)](TC) (or TC[FO[s(n), s(n)/log(n)]], but does just not use the TC operator. So if s(n) *o*(*n*) then this is equivalent or more diﬃcult then showing that *SO*(arity *k*)[*nk*] = *SO*(arity *k*)(*T C*) Otherwise, if *t*(*n*) ≥ *s*(*n*) we have to show that *FO t*(*n*)*,*  *s*(*n*) = *SO*(arity *k*)(*T C*). The only problem arises if *s*(*n*) *> t*(*n*), as then the generated new formula is not in the original class.

∉

*log*(*n*)

1 l

What would happen is that then *SO*(*arity k*)(*T C*) = *FO* 1*t*(*n*)*,*  *s*(*n*) l (*TC*) ⊆ *FO* 1*s*(*n*)*,* *s*(*n*) l ⊆

log(*n*)

log *n*

*SO*(arity *k*)[*nk*] We know that *s*(*n*) must be in Ω(√*n*) by Savitch’s theorem. Also if *s*(*n*) ∈ *o*(*n*), we

would have a stronger result then Savitch’s theorem as then *NSP ACE*[*n*] ⊆ *DSP ACE*[*s*(*n*)2], which is an open problem since almost 50 years.

In the end, this proves that combining TC and quantiﬁer iterations is as diﬃcult either as showing an improvement of Savitch’s theorem or is equivalent to having an alternative characterisation using only iterated quantiﬁers.

## Generalised quantiﬁers

By [double arity](https://link.springer.com/article/10.1007/bf01268616) we have FO(TC) = FO(*QT C*) which is the tc quantiﬁer. The tc Quantiﬁer has type

⟨2*,* 1*,* 1⟩ and *k*-vectorisation to make it generalised to vectorisations. There is a theorem which states that no family of quantiﬁers of arity *<* 2*k* suﬃces to express these properties. So We can eﬀectively say that they are the easyest.

This is true for FO-generalised quantiﬁers on unordered structures, but unclear for [SO-gq](https://core.ac.uk/download/pdf/81931656.pdf) Problem is that we can query bits, so we have even more than just an order

Can we port the results directly?

## FO∃[n]

**FO**∃**[s(n)]** A set *S* ⊆ STRUCT[*τ* ] is in FO∃[s(n)] if and only if there exists formulas *Mi, Nj*, 0 ≤ *i* ≤

*k,* 1 ≤ *j* ≤ *r* from L(*τ* ), a tuple *c* of constants and a quantiﬁer block

*QB* = [(∃*x*1*.M*1)(∀*b*1*.N*1) *. . .* (∃*xk.Mk*)(∀*br.Nr*)]

such that for all A ∈ STRUCT[*τ* ], we have

A ∈ *S* ⇔ A |= ([*QB*] *M*0)(*c/x*)

*s*(|A|)

FO∃-VAR[t(n), s(n)] is the same with extended variables We can write TC as

*QB* ≡ (∀*b*1*.M*1)(∃*Z*)(∀*b*2)(∃*A, B.M*2)(∃*X, Y .M*3)

*M*1 ≡ ¬(∀*z*(BIT(*X, z*) ↔ BIT(*Y , z*)) ∨ *φ*(*X, Y* ))

*M*2 ≡ (*b*2 ∧ ∀*z*(BIT(*X, z*) ↔ BIT(*A, z*)) ∧ ∀*z*(BIT(*B, z*) ↔ BIT(*Z, z*)))∨

(¬*b*2 ∧ ∀*z*(BIT(*Z, z*) ↔ BIT(*A, z*)) ∧ ∀*z*(BIT(*B, z*) ↔ BIT(*Y , z*)))

*M*3 ≡ ∀*z*(BIT(*X, z*) ↔ BIT(*A, z*)) ∧ ∀*z*(BIT(*B, z*) ↔ BIT(*Y , z*))

And then [*QB*]*cs*(*n*)(*false*)(*false/X, true/Y* ) ≡ (*TCX,Y φ*(*X, Y* ))(*false, true*)

Thus, we have NSPACE[s(n)] = FO-VAR[1, s(n)/log(n)](TC) ⊆ FO∃-VAR[s(n), s(n)/log(n)]

For the reverse containment, we can simulate the formula with a NSPACE[s(n)] machine by exist- entially guessing the existentially quantiﬁed variables and remembering the universal choices in *s*(*n*) space. Does not work because we need some backtracking stuﬀ, and thus remembering the existential vars.

Also, *FO* − *V AR*[1*, s*(*n*)*/log*(*n*)](*T C*) ⊆ *FO* − *V AR* 1  *s*(*n*) *, s*(*n*)*r*(*n*) l by doing multiple steps at

log(*r*(*n*))

once, and universal quantiﬁcation over *r*(*n*) bit variables only.

## Alternating machines and FO-VAR

*ATIME* − *SP ACE*[*s*(*n*)2*, s*(*n*)] ⊆ *FO* − *V AR* 1*s*(*n*)2*,*  *s*(*n*) l

log(*n*)

log(*n*)

We have a state space of 2O(*s*(*n*)), so we need only *s*(*n*) bit Variables to encode it. We can assume that there exist formulas *existencial*(*X*) which is true if *X* is an existential state, and one *φ*(*X, Y* ) telling us that we can transition from *X* to *Y* .

Then, we have

*QB* ≡ (∀*b.M*1)(∃*A.φ*(*A, X*))(∀*B.M*2)(∃*X.*∀*z*(BIT(*X, z*) ↔ BIT(*B, z*))) *M*1 ≡ ¬∀*z*(BIT(*X, z*) ↔ BIT(*Y , z*))

*M*2 ≡ (*existential*(*X*) → ∀*z*(BIT(*A, z*) ↔ BIT(*B, z*))) ∧ *φ*(*X, B*)

and

([*QB*]*s*(*n*)2 (*false*)) (*false/X, true/Y* )

*ATSR*[*t*(*n*)*, s*(*n*)*, r*(*n*)] ⊆ *FO* − *V AR* 1*r*(*n*) log ( *t*(*n*) 1 *, s*(*n*) l

*r*(*n*)

log *n*

⊆ *ATSR* 1*r*(*n*)*s*(*n*) log ( *t*(*n*) 1 *, s*(*n*)*, r*(*n*) log ( *t*(*n*) 1l

*r*(*n*)

*r*(*n*)

The second containment is a consequence of how we simulate the formula with the alternating Tur- ing machine, as we have *FO* − *V AR* 1*a*(*n*)*,* *s*(*n*) l ⊆ *ATSR*[*a*(*n*)*s*(*n*)*, s*(*n*)*, a*(*n*)] by using the normal

log *n*

construction in Savitch’s theorem.

In the ﬁrst containment, we try to use the method of simulating NSPACE for each alternation.

First, we show that

*r*(*n*)

X

*r*(*n*)

log(*aj*)

*j*=0

under the condition that I:*r*(*n*) *aj* = *t*(*n*).

*j*=0

≤ *r*(*n*

) log ( *t*(*n*) 1

This can be shown using Jensen’s inequality.

I:*r*(*n*) log(*aj* ) I:*r*(*n*) *aj*  ( *t*(*n*) 1

*j*=0

*j*=0

*r*(*n*) ≤ log  *r*(*n*)  = log *r*(*n*)

as log is concave. Multiplying with *r*(*n*) gives the claim.

*Y* accepting *X* start *I* actual intermediate end *S* actual start *E* actual end *A, B* copy vars *bpath* if we are trying to ﬁnd or not ﬁnd a path

*QB* ≡ (∀*b*0*.M*1)(∃*b*0*.N*1)(∃*b*1*.M*2)(∃*Ze.M*3)(∀*Z.N*2)

(∃*b*2*e*)(∀*b*2*.N*3)(∃*A, B.M*4)(∀*S, E.M*5)(∃*A, B.M*6)(∃*I, X.M*7)(∃*b*0*.N*4)(∃*bpath.N*5)

*M*1 ≡ *bpath* → ¬(*I* = *Y* ∧ (*S* = *E* ∨ *φ*(*S, E*)))

*N*1 ≡ ¬*bpath* → ¬(*S* = *E* ∨ *φ*(*S, E*))

*M*2 ≡ ¬*b*1 ↔ (*S* = *E* ∨ *φ*(*S, E*))

*M*3 ≡ *existential*(*Ze*) ↔ *existential*(*X*)

*N*2 ≡ *bpath* → *Z* = *Ze N*3 ≡ ¬*bpath* → *b*2 = *b*2*e*

*M*4 ≡ (*b*1 ∧ ((¬*b*2 ∧ *A* = *S* ∧ *B* = *Z*) ∨ (*b*2 ∧ *A* = *Z* ∧ *B* = *E*)))∨

(¬*b*1 ∧ *A* = *I* ∧ (*B* = *Y* ∨ (*existential*(*B*) ↔ ¬*existential*(*I*))))

*M*5 ≡ *S* = *A* ∧ (((*b*1 ∨ *existential*(*S*)) ∧ *E* = *B*)∨

(¬(*b*1 ∨ *existential*(*S*)) ∧ (*E* = *Y* ∨ (*existential*(*E*) ↔ ¬*existential*(*S*)))))

*M*6 ≡ (¬*b*1 ∧ *S* = *A* ∧ *E* = *B*) ∨ (*b*1 ∧ *X* = *A* ∧ *I* = *B*) *M*7 ≡ *X* = *A* ∧ *I* = *B*

*N*4 ≡ *b*0 = *bpath*

*N*5 ≡ *b*1 → *bpath* = *b*0

In this formula, we have [*QB*]*r*(*n*) log(*t*(*n*)*/r*(*n*))(¬*bpath*) (*true/bpath, AS/S, X, I, S, E, AE/Y* ) is true if and only if the machine accepts.

( )

=⇒ *FO* − *V AR* 1*a*(*n*)*,* *s*(*n*) l ⊆ *ATSR*[*a*(*n*)*s*(*n*)*, s*(*n*)*, a*(*n*)] ⊆ *FO* − *V AR* 1*a*(*n*) log(*s*(*n*))*,* *s*(*n*) l

log *n*

log *n*

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