

# Enhancing Giri: Dynamic Slicing in LLVM

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<http://github.com/liuml07/giri>

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**D**ynamic program slicing has been used in many applications. Giri was a research project from UIUC, which implemented the dynamic backward slicing in LLVM. It was selected as the Google Summer of Code 2013. We achieved several improvements to Giri during GSoC 2013: 1) Update the code to LLVM mainline and make it robust, 2) Reduce the trace size, 3) Make the Giri thread-aware (pthread only), 4) Improve the performance of run-time.

## 1 Introduction

Program slice contains all statements in a program that directly or indirectly affect the value of a variable or instruction [7]. We can further narrow down the notion of the *slice* to *dynamic program slice* [1]. Unlike traditional slicing named static slicing which computes all statements that may have affected the value of a variable at a program point for any arbitrary execution of the program, the dynamic program slicing computes all statements that actually affect the value of a variable/instruction at a program point for a particular execution of the program with given input. There are many applications from research and industry organizations which use (or could benefit from) dynamic slicing. For example, it has long been used in software debugging [2, 4] and testing [3]. However, to the best of our knowledge, there is no publicly available dynamic slicing tool in either GCC or Open64.

Sahoo et. al. use dynamic program slicing to filter/remove false positive candidate root causes for automated software fault localization [5]. They implemented the dynamic program slicing code, named Giri, in LLVM compiler infrastructure (version 3.1) for research purpose. It collects the trace of a user program execution and finally reports the dynamic slice of a given statement. It also maps LLVM IR statements to source code using the debug meta-data to report the dynamic slice in terms of source line numbers. To make the Giri up-to-date and robust for general use, we enhanced the code in several ways.

The report is organized as following. Section 2 illustrates the dynamic slicing and shows the internal architecture of Giri. Section 3 lists the progress we made during the GSoC 2013. Section 4 lists the future work to improve Giri further. Section 5 concludes our work.

## 2 Overview

### 2.1 Dynamic Program Slicing

The following is an example program.

---

```
1 #include <math.h>
2 #include <stdio.h>
3 #include <stdlib.h>
4
5 typedef struct result
6 {
7     int result;
8     int count;
9 } result_t;
10
```

```

11 void calc(int x, result_t *y, result_t *z)
12 {
13     if (x < 0)
14     {
15         y->result = sqrt(x);
16         y->count++;
17         z->result = pow(2, x);
18         z->count++;
19     } else {
20         if (x == 0)
21         {
22             y->result = sqrt(x * 2);
23             y->count++;
24             z->result = pow(3, x);
25             z->count++;
26         } else {
27             y->result = sqrt(x * 3);
28             y->count++;
29             z->result = pow(4, x);
30             z->count++;
31         }
32     }
33 }
34
35 int main(int argc, char *argv[])
36 {
37     int x, ret;
38     result_t y = {0, 0}, z = {0, 0};
39
40     x = atoi(argv[1]);
41
42     calc(x, &y, &z);
43
44     ret = printf("%d\n", y.result);
45     printf("%d\n", z.result);
46
47     return ret;
48 }

```

Given the input 10 to this program from command line, the source line numbers affect the return value are: `<13, 20, 27, 40, 42, 44, 47>`.

The default start point of slicing is the `return` instruction at the `main` function. Note that there is only one `return` when the `main` executes. Giri also supports two more ways to specify the slicing criterion, which are source code line number and LLVM instruction number respectively. See the `test/UnitTests/test4` for more information.

## 2.2 The Design of Giri

Giri handles both data-flow and control-flow dependences (optional) when computing the dynamic backwards slice. It has two phases. In the first pass named *tracing* pass, it instruments the code to record various runtime information in a trace file at run-time for all threads. In the second pass named *slicing* pass, it uses the execution trace to create a program dependence graph for computing dynamic slice. Giri

takes advantage of the LLVM intermediate representation (IR), which is static single assignment (SSA) form, to reduce the size of the trace file.

The tracing pass instruments code to record three different pieces of information: 1) basic block exits; 2) memory accesses and their addresses; and 3) function calls and returns. The slicing pass initializes a work list with the start instruction (also called *criterion*) of slicing. For each dynamic value in the work list, its parent was retrieved if it is not processed yet, which is a basic block. The slicing pass uses static control-dependence analysis to determine which value forced the execution of that basic block. The dynamic value is added to the work list if not processed. Next, it uses the data-dependence analysis to find the most recent execution of the instruction generating that value and add it to the work list too. The slicing pass does not stop until all dynamic values in the work list are processed. In the end, it reports the dynamic slice of the specific program execution as well as the source code line number information.

Since most instructions operate on SSA scalar values, we only refer to dynamic trace whenever needed. For example, we can get all operands of a `call` instruction in the SSA form without accessing the trace file. However, when we find all `store` instructions for a `load` instruction, we need to scan the trace file backward since the data read by the `load` instruction may be written by many possible `store` instructions. For programs with multiple threads, it is possible that the `store` instructions locate at different basic blocks or functions.

To make the slicing pass thread aware, the runtime library writes the thread id (`thread_t`) to each record indicating current thread performing that operation. The slicing pass checks the thread id when scanning the trace file for a given dynamic value or basic block.

## 3 Progress

Our goal of GSoC was to make the Giri code up-to-date with the latest LLVM version, improve its runtime performance, and/or reduce the tracing overhead. There are several things we did in this summer:

### 1. Update the code to LLVM mainline.

- Clean up version 3.1 which works with LLVM 3.1
- Make the code update to the latest LLVM 3.4 (r191529)

2. **Make the Giri run-time library thread safe and the slicing pass thread aware.** The trace records of all threads are written to a single file. In addition to other information, trace records also include the thread id (`pthread_t`), which indicates the thread performing that particular operation. We use locks when we write trace records to the trace file in order to avoid race conditions.
3. **Improve the Giri run-time performance.** We checked every call to function `mmap` as well as the parameters, and eliminated the useless ones to improve performance. For example, at the end of the tracing, it is not needed to `rmmap` once all trace records are synchronized to disk file. Giri now dynamically computes the cache size of trace records to hold in memory before flushing to disk.
4. **Write dozens of unit tests and try more real programs.** Giri was able to handle few large real programs like Squid, Apache, MySQL before GSoC. We added few more real programs to this set. There is also a simple test framework which runs all the unit tests at the top level of `giri/test/` and report the results. In the future development, we shall make sure all unit tests pass locally before submitting a patch to the code repository.
5. **Reduce the trace size.** We truncate the trace file at the end of tracing according to its real size.
6. **Write documents for the project.** There are several sample pages:
  - The Wiki Page Home.
  - How to Compile Giri.
  - Hello World! Example.
  - Command Line Usage With An Example.

## 4 Future Work

The following are the list of future work we are going to address in the next few months. An updated list is maintained online at the wiki page of Giri at [github.com](https://github.com).

1. Improve the performance of locking mechanism of the Giri runtime.
2. Write more unit tests with complicated direct/indirect recursive function calls.

3. Double check the `TraceFile.cpp` to make sure the thread id be checked correctly.
4. Add `tool/Tracer.cpp` to the make list. It was removed when upgrading to LLVM 3.4.
5. Try more real world programs and add them to the `test/` directory.
6. Try large test programs, e.g. *nginx*, *squid*, etc.
7. Make Giri code useful for other platforms besides Linux, e.g. Mac OS X, Cygwin, and FreeBSD.
8. Handle more special function calls, which are intrinsic LLVM instructions (like *libc* functions `printf`, `scanf`, `sprintf` etc.) in `TracingNoGiri::visitSpecialCall()`.
9. Parallelize the code which writes the entry cache containing trace records to trace file and adds trace records to the entry cache.

## 5 Conclusion

As the first publicly available dynamic slicing tool, Giri was improved significantly during the GSoC. We updated the code to LLVM mainline, reduced the trace size, made it thread-aware, and improved the performance of its run-time. We have published our code at <https://github.com/liuml07/giri> [6].

The code is still under active development. Dr. Swarup will direct Mingliang LIU in future to make the Giri code better. There are other researchers from Tsinghua University, Xi'an Jiaotong University and University of Miami, who contacted us for permission to use and contribute to Giri project. For more information, please visit the homepage above.

## References

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